MORE SCHEME AND INTERPRETERS

COMPUTER SCIENCE MENTORS

November 9, 2020 - November 12, 2020

Call expressions follow prefix notation, i.e. (<operator> <operand1> <operand2> ... <operandN>)

Evaluating a call expressions closely mirrors Python:

- Evaluate the operator, yielding a procedure p
- Evaluate each operand, each yielding a value argi
- Apply the procedure p with arguments arg1, arg2, ..., argN

Special forms *look* like call expressions but aren't – they implement Scheme language **\$** features and follow special evaluation rules (e.g., short-circuiting).

(Aside: Note that you're free to use a special form name as a variable name, but the name will be looked up *only* in a non-operator position; when used as an operator, it will always refer to the original special form.)

Notable Special Forms:

behavior	syntax
variable assignment	(define <variable-name> <value>)</value></variable-name>
function defining	(define (<function> <op1><opn>) <body>)</body></opn></op1></function>
if / else	(if <condition> <true-expr> <else-expr>)</else-expr></true-expr></condition>
if / elif / else	<pre>(cond (<cond1> <expr1>) (else <else-expr>))</else-expr></expr1></cond1></pre>
and	(and <operand1> <operandn>)</operandn></operand1>
or	(or <operand1> <operandn>)</operandn></operand1>
quote	(quote <operand1>)</operand1>
begin	(begin <expr1> <expr2> <exprn>)</exprn></expr2></expr1>
lambdas	(lambda (<operand1> <operandn>) <body>)</body></operandn></operand1>
let / execute many lines	(let ((<var1> <val1>) (<varn> <valn>)) body)</valn></varn></val1></var1>

737 (Jethne a U)

G
777 a

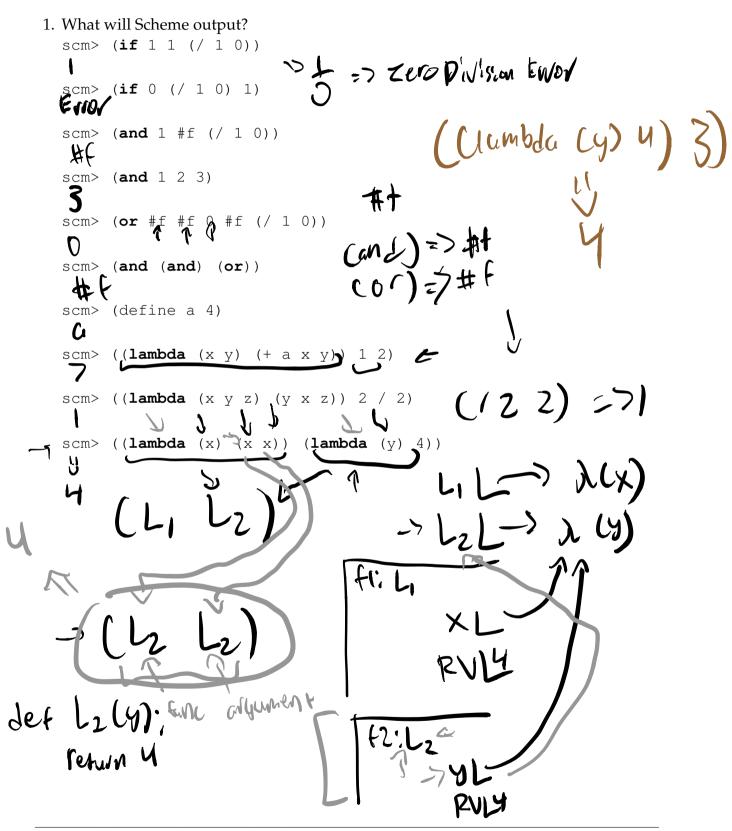
777 a

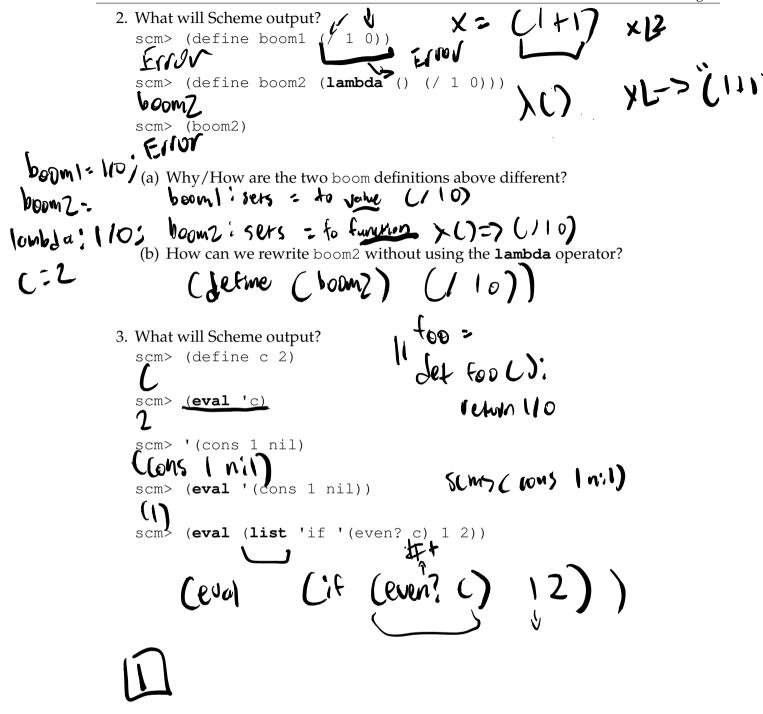
777 a

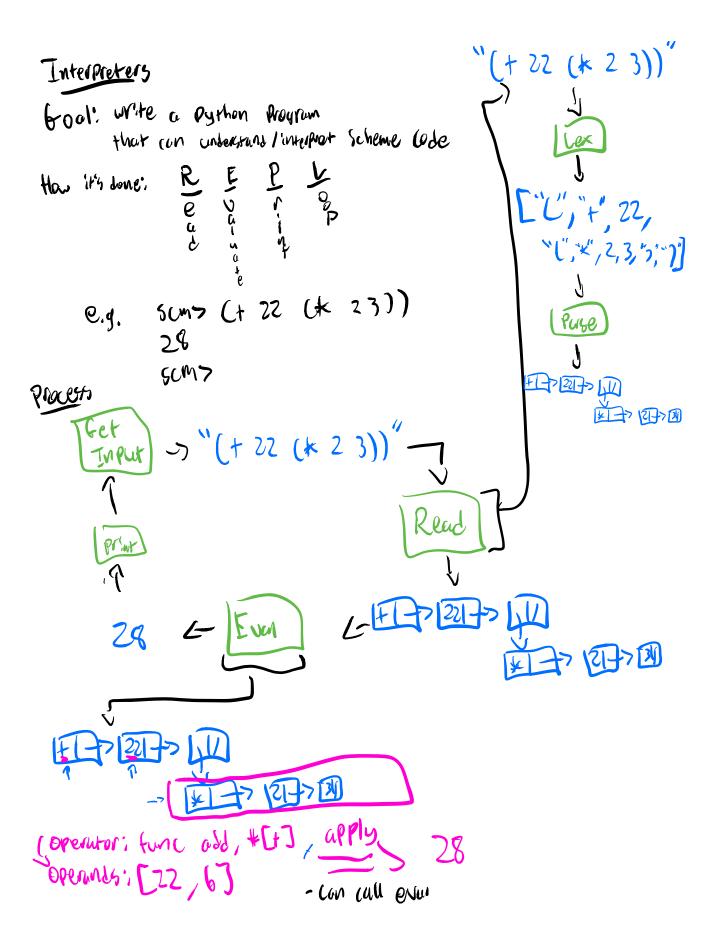
777 a

777 (eval 'a)

1 What Would Scheme Print?







2 Interpreters

The following questions refer to the Scheme interpreter. Assume we're using the implementation seen in lecture and in the Scheme project.

1. What's the purpose of the read stage in a Read-Eval-Print Loop? For our Scheme interpreter, what does it take in, and what does it return?

- 'un', Str'my - out', Vinced l'ist (Pu'ur)

2. What are the two components of the read stage? What do they do?

1. Lex: gets each individual token from impur string

2. Russer: turns tokens => duta smuture (Parr)

3. Write out the constructor for the Pair object the read stage creates with the input string

Point ("sexine", Print avint ("600", Paint")), Paint ("4", Paint")

("x", Print (1))))

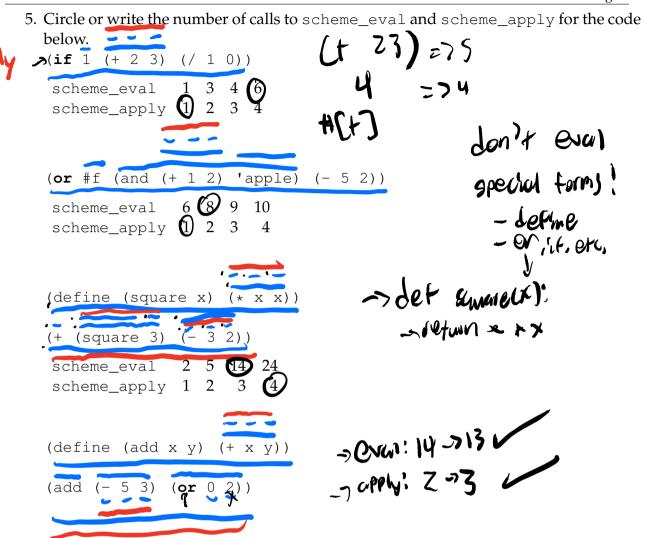
Lower (1)

4. For the previous example, imagine we saved that Pair object to the variable p. How could we check that the expression is a define special form? How would we access the name of the function and the body of the function?

Check (cor ?)

- P. first 22 "define"

- wener P. Second. First. First -body: Process, sevone. First



Code Writing

1. Define **is**-prefix, which takes in a list p and a list lst and determines if p is a prefix of lst. That is, it determines if lst starts with all the elements in p.

```
; Doctests:
 scm> (is-prefix '() '())
 → scm> (is-prefix '() '(1 2))
    scm> (is-prefix '(1) '(1 2))
    #t.
    scm> (is-prefix '(2) '(1 2))
    ; Note here p is longer than 1st
scm> (is-prefix '(1 2) '(1))
#f
    (define (is-prefix p lst)
          PISE (if (= (cor P) (cor 15+))

(is-prefix (cdr P) (cdr 15+))
               (and (= carp) (cor 14t)) (is prefix ...
    )
```

2. Define **apply**-multiple which takes in a single argument function f, a nonnegative integer n, and a value x and returns the result of applying f to x a total of n times.

```
;doctests
scm> (apply-multiple (lambda (x) (* x x)) 3 2)
256
scm> (apply-multiple (lambda (x) (+ x 1)) 10 1)
11
scm> (apply-multiple (lambda (x) (* 1000 x)) 0 5)
5
(define (apply-multiple f n x)
```

)

3. Finish the functions max and max-depth. max takes in two numbers and returns the larger. Function max-depth takes in a list lst and returns the maximum depth of the list. In a nested scheme list, we define the depth as the number of scheme lists a sublist is nested within. A scheme list with no nested lists has a max-depth of 0.