

UC Berkeley EECS
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Computational Structures in Data Science

Lecture 6: Environment Diagrams, Recursion Review, Midterm Review

March 4, 2019 <http://inst.eecs.berkeley.edu/~cs88>

On Computer Science Exams

In computer science exams, we try to assess the student's understanding of concepts and his or her ability to practically apply these.

- In CS, we do not:
 - require extensive memorization (e.g. we allow cheat sheet)
 - require a lot of reading
 - require essay writing skills
- In CS, we do:
 - require the ability to translate a given textual problem into programming code
 - require you to be able to read other people's code
 - value solutions that are almost right over no solution
 - accept solutions we did not think about if they work
 - prioritize math (logic) and science (experiment) over opinion or authority

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How to prepare for a CS exam

- Explain the content of the computational concepts toolbox to somebody else
 - Describe the concept
 - What is an example of using it?
 - When does it not work? Corner cases?
 - Why does it exist?
- Practice programming:
 - Play around with the examples from lecture, lab, homework
 - Think about your own similar examples
- In the exam:
 - Make sure you understand the question: What is the given input? What is the required output?
 - Think of easy cases first (e.g. $n=1$).
 - What is the iteration/recursion doing (e.g. $i=i+1$)?
 - What are corner cases that need explicit handling (e.g. division by zero, negative numbers, empty list)?

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Computational Concepts Toolbox

- Data type: values, literals, operations,
 - e.g., int, float, string
- Expressions, Call expression
- Variables
- Assignment Statement
- Sequences: tuple, list
 - indexing
- Data structures
- Tuple assignment
- Call Expressions
- Function Definition Statement
- Conditional Statement
- Iteration:
 - data-driven (list comprehension)
 - control-driven (for statement)
 - while statement
- Higher Order Functions
 - Functions as Values
 - Functions with functions as argument
 - Assignment of function values
- Recursion
- Environment Diagrams

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Recursion Key concepts – by example

```
def sum_of_squares(n):
    if n < 1:
        return 0
    else:
        return sum_of_squares(n-1) + n**2
```

1. Test for simple "base" case

2. Solution in simple "base" case

3. Assume recursive solution to simpler problem

4. Transform soln of simpler problem into full soln

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In words

- The sum of no numbers is zero
- The sum of 1^2 through n^2 is the
 - sum of 1^2 through $(n-1)^2$
 - plus n^2

```
def sum_of_squares(n):
    if n < 1:
        return 0
    else:
        return sum_of_squares(n-1) + n**2
```

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How does it work?

- Each recursive call gets its own local variables
 - Just like any other function call
- Computes its result (possibly using additional calls)
 - Just like any other function call
- Returns its result and returns control to its caller
 - Just like any other function call
- The function that is called happens to be itself
 - Called on a simpler problem
 - Eventually bottoms out on the simple base case
- Reason about correctness “by induction”
 - Solve a base case
 - Assuming a solution to a smaller problem, extend it

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Local variables

```
def sum_of_squares(n):
    n_squared = n**2
    if n < 1:
        return 0
    else:
        return n_squared + sum_of_squares(n-1)
```

- Each call has its own “frame” of local variables
- What about globals?
- Let’s see the environment diagrams

<https://goo.gl/CiFaUJ>

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Environments Example

Python 3.3

```
def sum_of_squares(n):
    n_squared = n**2
    if n < 1:
        return 1
    else:
        return n_squared + sum_of_squares(n-1)
sum_of_squares(3)
```

Global frame: sum_of_squares → func sum_of_squares(n) (parent=Global)

f1: sum_of_squares (parent=Global): n=3, n_squared=9

python.tutor.com

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Environments Example

Python 3.3

```
def sum_of_squares(n):
    n_squared = n**2
    if n < 1:
        return 1
    else:
        return n_squared + sum_of_squares(n-1)
sum_of_squares(3)
```

Global frame: sum_of_squares → func sum_of_squares(n) (parent=Global)

f1: sum_of_squares (parent=Global): n=3, n_squared=9

f2: sum_of_squares (parent=Global): n=2, n_squared=4

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Environments Example

Python 3.3

```
def sum_of_squares(n):
    n_squared = n**2
    if n < 1:
        return 1
    else:
        return n_squared + sum_of_squares(n-1)
sum_of_squares(3)
```

Global frame: sum_of_squares → func sum_of_squares(n) (parent=Global)

f1: sum_of_squares (parent=Global): n=3, n_squared=9

f2: sum_of_squares (parent=Global): n=2, n_squared=4

f3: sum_of_squares (parent=Global): n=1, n_squared=1

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Environments Example

Python 3.3

```
def sum_of_squares(n):
    n_squared = n**2
    if n < 1:
        return 1
    else:
        return n_squared + sum_of_squares(n-1)
sum_of_squares(3)
```

Global frame: sum_of_squares → func sum_of_squares(n) (parent=Global)

f1: sum_of_squares (parent=Global): n=3, n_squared=9

f2: sum_of_squares (parent=Global): n=2, n_squared=4

f3: sum_of_squares (parent=Global): n=1, n_squared=1

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Environments Example

Python 3.3

```

1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
4         return 1
5     else:
6         return n_squared + sum_of_squares(n-1)
7 sum_of_squares(3)

```

Global frame

sum_of_squares (parent=Global)

F1: sum_of_squares (parent=Global)

n 3
n_squared 9

F2: sum_of_squares (parent=Global)

n 2
n_squared 4

F3: sum_of_squares (parent=Global)

n 1
n_squared 1

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Environments Example

Python 3.3

```

1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
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```

Global frame

sum_of_squares (parent=Global)

F1: sum_of_squares (parent=Global)

n 3
n_squared 9

F2: sum_of_squares (parent=Global)

n 2
n_squared 4

F3: sum_of_squares (parent=Global)

n 1
n_squared 1

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Environments Example

Python 3.3

```

1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
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```

Global frame

sum_of_squares (parent=Global)

F1: sum_of_squares (parent=Global)

n 3
n_squared 9

F2: sum_of_squares (parent=Global)

n 2
n_squared 4

F3: sum_of_squares (parent=Global)

n 1
n_squared 1

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Environments Example

Python 3.3

```

1 def sum_of_squares(n):
2     n_squared = n**2
3     if n == 1:
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```

Global frame

sum_of_squares (parent=Global)

F1: sum_of_squares (parent=Global)

n 3
n_squared 9

F2: sum_of_squares (parent=Global)

n 2
n_squared 4

F3: sum_of_squares (parent=Global)

n 1
n_squared 1

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How much ???

- Time is required to compute `sum_of_squares(n)`?
 - Recursively?
 - Iteratively?
- Space is required to compute `sum_of_squares(n)`?
 - Recursively?
 - Iteratively?
- Count the frames...
- Recursive is linear, iterative is constant!

Linear proportional to cn for some c

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Tail Recursion

- All the work happens on the way down the recursion
- On the way back up, just return

```

def sum_up_squares(i, n, accum):
    """Sum the squares from i to n in incr. order"""
    if i > n:
        Base Case
    else:
        Tail Recursive Case

>>> sum_up_squares(1, 3, 0)
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```

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Tree Recursion

- Break the problem into multiple smaller sub-problems, and solve them recursively

```
def split(x, s):
    return [i for i in s if i <= x], [i for i in s if i > x]

def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""
    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split(pivot, rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

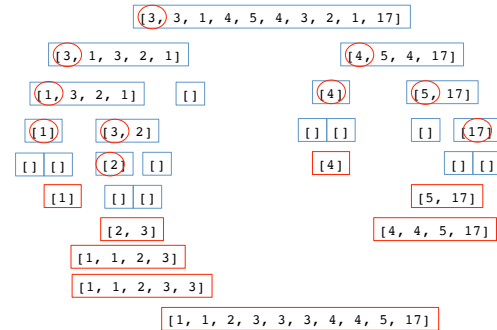
>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 4, 4, 5, 17]
```

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QuickSort Example



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Tree Recursion with HOF

```
def qsort(s):
    """Sort a sequence - split it by the first element,
    sort both parts and put them back together."""
    if not s:
        return []
    else:
        pivot = first(s)
        lessor, more = split_fun(leg_maker(pivot), rest(s))
        return qsort(lessor) + [pivot] + qsort(more)

>>> qsort([3,3,1,4,5,4,3,2,1,17])
[1, 1, 2, 3, 3, 3, 4, 4, 5, 17]
```

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Answers for the Wandering Mind

The computer chooses a random element x of the list generated by `range(0,n)`. What is the smallest amount of iteration/recursion steps the best algorithm needs to guess x ?

$\log_2 n$

How would the algorithm look like?

Guess the binary digits of x starting with the highest significant digit. This is, ask questions of the form "smaller than 2^{n-1} ?" (yes => 0...), "smaller than 2^{n-2} ?" (no => 0 1...), "smaller than $2^{n-2}+2^{n-3}$?", ...

This method is also called: binary search

Quantum physics: Allow less than $\log_2 n$ guesses.

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