Morphological epenthesis in Serbo-Croatian noun inflection

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Introduction

- Class I:

Class I:	zavod institute.NOM.SG.M	zavod-a institute.GEN.SG.M
	sel-o village.NOM.SG.N	sel-a village.GEN.SG.N
	tele calf.NOM.SG.N	tele**t**-a calf.GEN.SG.N
Class II:	ʒen-a woman.NOM.SG.F	ʒen-e woman.GEN.SG.F
Class III:	retʃ word.NOM.SG.F	retʃ-i word.GEN.SG.F

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• Class III:	reʧ word.NOM.SG.F	reʧ-i word.GEN.SG.F

morphologically conditioned consonant insertion
–
repairing vowel-final stems

Introduction

- part of a wider algorithmic process of noun inflection
  - Boolean Monadic Recursive Schemes (BMRSs)
  - analyzed as logical transductions on strings
    - stem + [Case Num] → inflected word ~ realization rules
  - hierarchical ordering of more specific and less specific blocking and licensing structures
    - ~ realization rules intrinsically ordered by Pāṇini’s principle (cf. Stump 2001)
  - possibility to compose BMRSs into a complete inflectional system
    - algorithm of inflectional class assignment

Morphologically conditioned consonant insertion

- SC stems are normally C-final; those that are not are repaired via consonant insertion

  → **if another suffix follows immediately**

otherwise the word is vowel-final (bare stem) and considered well-formed (as the nom.sg. suffix is *-o*/*-e*)

	tele	‘calf.NOM.SG’
	tele*t*-a	‘calf.GEN.SG.
/tele + ji/	[teletɕi]	‘calf-like(ADJ.)’

Morphologically conditioned consonant insertion

- *t* is not the general epenthetic consonant in this language

/zaova/	[zaova]	‘sister-in-law’
/violina/	[vijolina]	‘violin’

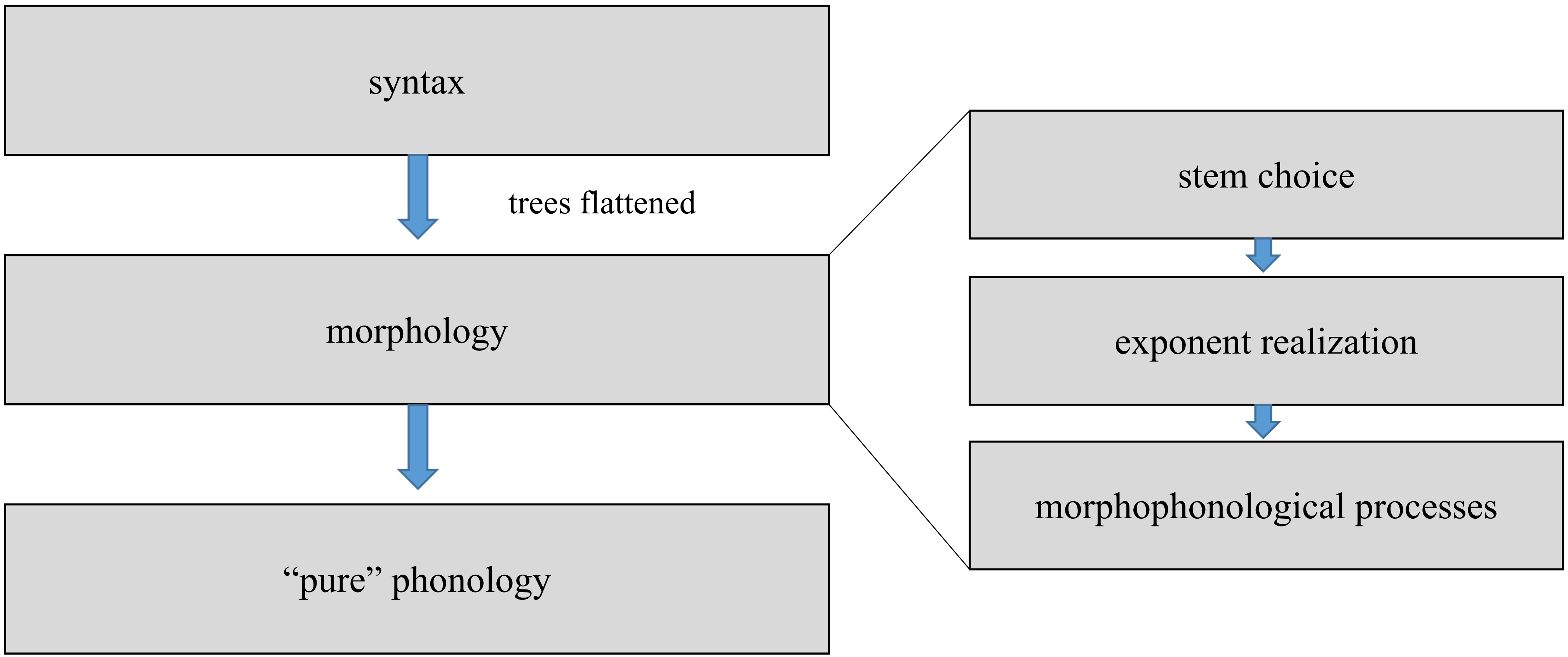
- *t* epenthesized only to repair vowel-final stems, does not depend on vowel hiatus

/tele + ji/	[teletɕi] ‘calf-like’	vs.	/koz + ji/	[kozji] ‘goat-like’
/sirtɕe + ni/	[sirtɕetni] ‘acetic’	vs.	/plod + ni/	[plodni] ‘fertile’

The nature of morphological processes

- morphological processes are regular (Karttunen et al. 1992), most often subregular (Chandlee 2017)
- DM is a tree-based theory, but can be formalized over strings (Ermolaeva & Edmiston 2018)
  - yield function flattens trees, retaining necessary boundary symbols
- In terms of computational power, W&P morphology is equivalent to a collection of regular relations (Karttunen 2003)

Proposed model



Analysis

- *Boolean Monadic Recursive Schemes* (BMRS)
    - incorporates the observation that phonology is regular (Johnson 1972; Kaplan & Kay 1994; Heinz 2018)
    - morphology, too, is regular (Karttunen et al. 1992), most often subregular (Chandlee 2017)
    - unlike FSTs, captures linguistically significant generalizations – we can use representations that are common in linguistics
        - phonological, morphosyntactic features
        - Pāṇini's Principle (= Elsewhere condition) has no obvious finite-state implementation (Karttunen 2003), whereas it is directly captured with BMRS's 'if...then...else' syntax

Analysis

- primitives:
  - *boolean values* ⊤ and ⊥
  - *monadic predicates* $P(t)$ – take a *single* argument *t,* and return ⊤ or ⊥

$\mathbb{V}$ = {a, e, i, o, u}
$\mathbb{C}$ = {b, ʦ, ʧ, ʨ, d, ʤ, ʥ, f, g, h, j, k, l, ʎ, m, n, ɲ, p, r, s, ʃ, t, v, z, ʒ}
$\mathbb{M}$ = {[nom], [gen], [dat], [acc], [voc], [ins], [loc], [sg], [pl]}

$\Sigma = \mathbb{V} \cup \mathbb{C} \cup \mathbb{M}$ $\Sigma_{\bowtie} = \Sigma \cup \{\rtimes, \ltimes, +\}$

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$\mathbb{M}$ = {[nom], [gen], [dat], [acc], [voc], [ins], [loc], [sg], [pl]}

$\Sigma = \mathbb{V} \cup \mathbb{C} \cup \mathbb{M}$ $\Sigma_{\bowtie} = \Sigma \cup \{\rtimes, \ltimes, +\}$

$\mathcal{I}$ = {a(t), …, ʒ(t), [nom](t), …, [pl](t), ⋊(t), ⋉(t), +(t)}
$\mathcal{O}$ = {$\text{a}_{\text{o}}(t)$, …, $\text{ʒ}_{\text{o}}(t)$, $\text{[nom]}_{\text{o}}(t)$, …, $\text{[pl]}_{\text{o}}(t)$, $\rtimes_{\text{o}}(t)$, $\ltimes_{\text{o}}(t)$, $+_{\text{o}}(t)$}

Analysis

- no right word boundary symbol → stems are (contiguous) strings of characters occupying indices belonging to a *countably infinite* set of indices

$\mathbf{S} = \langle D; \sigma_1, \sigma_2, \ldots, p, s \rangle$

$D = \{1, 2, \ldots\}$

- the right word boundary is assigned at the end of the word only after the inflectional exponents are realized (look out for $\boldsymbol{\ltimes(x)}$ in the following slides!)
- (I do this to avoid using copy sets, which might not be a great idea.)

Analysis

- In BMRS, we can refer to both input and output predicates, and define new ones

input								
	1	2	3	4	5	6	7	…
	⋊	t	e	l	e	+	[dat sg]	…
output								
		t	e	l	e	t	u	

- Examples of defining new input predicates: case/number feature combinations, syncretism, defining the class of FCs

Analysis

- necessary combinations of case and number features available for neuter nouns:

`[nom/acc/voc sg]`(x) = if `[nom/acc/voc]`(x) then `[sg]`(x) else $\perp$
`[dat/loc sg]`(x) = if `[dat/loc]`(x) then `[sg]`(x) else $\perp$
`[ins sg]`(x) = if `[ins]`(x) then `[sg]`(x) else $\perp$
`[nom/acc/voc pl]`(x) = if `[nom/acc/voc]`(x) then `[pl]`(x) else $\perp$
`[gen pl]`(x) = if `[gen]`(x) then `[pl]`(x) else $\perp$
`[dat/loc pl]`(x) = if `[dat/loc]`(x) then `[pl]`(x) else $\perp$
`[ins pl]`(x) = if `[ins]`(x) then `[pl]`(x) else $\perp$

(one would expect 14 – 7 cases × 2 numbers – but we do not have 14 distinct phonological forms of overt suffixes)

Analysis

- (non-directional) syncretism

`[dat/loc]`(x) = if `[dat]`(x) then ⊤ else `[loc]`(x)
`[nom/acc/voc]`(x) = if `[nom]`(x) then ⊤ else
if `[acc]`(x) then ⊤ else `[voc]`
`[gen/ nom/acc/voc pl]`(x) = if `[gen]`(x) then ⊤
else `[nom/acc/voc pl]`(x)
`[dat/loc pl/ins pl]`(x) = if `[dat/loc pl]`(x) then ⊤
else `[ins pl]`(x)

Analysis

- effect of fronting consonants (i.e. fronting of o to e after stem-final FCs)

(FC = C[cor, -ant] + ʦ)

[FC](x) = if ʦ(x) then ⊤ else
if ʧ(x) then ⊤ else
if ʨ(x) then ⊤ else
if ʤ(x) then ⊤ else
if ʥ(x) then ⊤ else
if j(x) then ⊤ else
if ʎ(x) then ⊤ else
if ɲ(x) then ⊤ else
if ʃ(x) then ⊤ else ʒ(x)

Analysis

- *t*-insertion

$\mathtt{outseg}(x)$ = if $\mathtt{a}_o(x)$ then $\top$ else
if $\mathtt{b}_o(x)$ then $\top$ else
if $\mathtt{c}_o(x)$ then $\top$ else
(...)
if $\mathtt{z}_o(x)$ then $\top$ else $\mathtt{ʒ}_o(x)$

$\mathtt{t}_o(x)$ =if $+(x)$ then
if $\mathtt{e}(p(x))$ then $\mathtt{outseg}(s(x))$ else $\perp$
else $\mathtt{t}(x)$

Analysis

- Output predicates: (~ realization rules)

$\mathtt{u}_o(x)$ = if $\mathtt{[dat/loc\ sg]}(x)$ then $\top$ else $\mathtt{u}(x)$
$\mathtt{a}_o(x)$ = if $\mathtt{[gen/\ nom/acc/voc\ pl]}(x)$ then $\top$ else
    if $\mathtt{[gen\ pl]}(p(x))$ then $\top$ else
    if $\mathtt{[ins\ pl]}(p(p(x)))$ then $\top$ else $\mathtt{a}(x)$
$\mathtt{o}_o(x)$ = if $\mathtt{[nom/acc/voc\ sg]}(x)$ then $\top$ else
    if $\mathtt{[ins\ sg]}(x)$ then $\top$ else $\mathtt{o}(x)$
$\mathtt{m}_o(x)$ = if $\mathtt{[ins]}(p(x)))$ then $\top$ else $\mathtt{m}(x)$
$\mathtt{i}_o(x)$ = if $\mathtt{[ins\ pl]}(x)$ then $\top$ else $\mathtt{i}(x)$
$\mathtt{e}_o(x)$ = if $\mathtt{o}_o(x)$ then
        if $+(p(x))$ then $\mathtt{[pal]}(p(p(x)))$ else $\perp$
    else $\mathtt{e}(x)$

$\mathtt{l}_o(x) = \mathtt{l}(x)$
$\mathtt{b}_o(x) = \mathtt{b}(x)$
(etc., for {ʦ, ʧ, ʨ, d, ʤ, ʥ, f, g, h, j, k, l, ʎ, n, ɲ, p, r, s, ʃ, v, z, ʒ})

Analysis

	1	2	3	4	5	6	7	8	9
input	⋊	t	e	l	e	+	[ins sg]		…
output		t	e	l	e	t	o	m	

- Output predicates:

$\texttt{u}_o(x)$ = if $\texttt{[dat/loc sg]}(x)$ then $\top$ else $\texttt{u}(x)$
$\texttt{a}_o(x)$ = if $\texttt{[gen/ nom/acc/voc pl]}(x)$ then $\top$ else
if $\texttt{[gen pl]}(p(x))$ then $\top$ else
if $\texttt{[ins pl]}(p(p(x)))$ then $\top$ else $\texttt{a}(x)$
$\texttt{o}_o(x)$ = if $\texttt{[nom/acc/voc sg]}(x)$ then $\top$ else
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$\texttt{i}_o(x)$ = if $\texttt{[ins pl]}(x)$ then $\top$ else $\texttt{i}(x)$
$\texttt{e}_o(x)$ = if $\texttt{o}_o(x)$ then
if $+(p(x))$ then $\texttt{[pal]}(p(p(x)))$ else $\perp$
else $\texttt{e}(x)$

$\texttt{l}_o(x) = \texttt{l}(x)$
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(etc., for {ʦ, ʧ, ʨ, d, ʤ, ʥ, f, g, h, j, k, l, ʎ, n, ɲ, p, r, s, ʃ, v, z, ʒ})

Analysis

$\ltimes(x)$ = if `[nom sg]`(x) then `e`$(p(x))$ else
if `[nom/acc/voc sg]`$(p(x))$ then $\top$ else
if `[dat/loc sg]`$(p(x))$ then $\top$ else
if `[gen/ nom/acc/voc pl]`(x) then $\top$ else
if `[ins sg]`$(p(p(x)))$ then $\top$ else
if `[dat/loc pl/ins pl]`$(p(p(p(x))))$ then $\top$ else $\bot$

out(x) = if $+(x)$ then $\mathtt{t}_o(x)$ else
if $\rtimes(x)$ then $\bot$ else
if $\ltimes(x)$ then $\bot$ else $\top$

Analysis

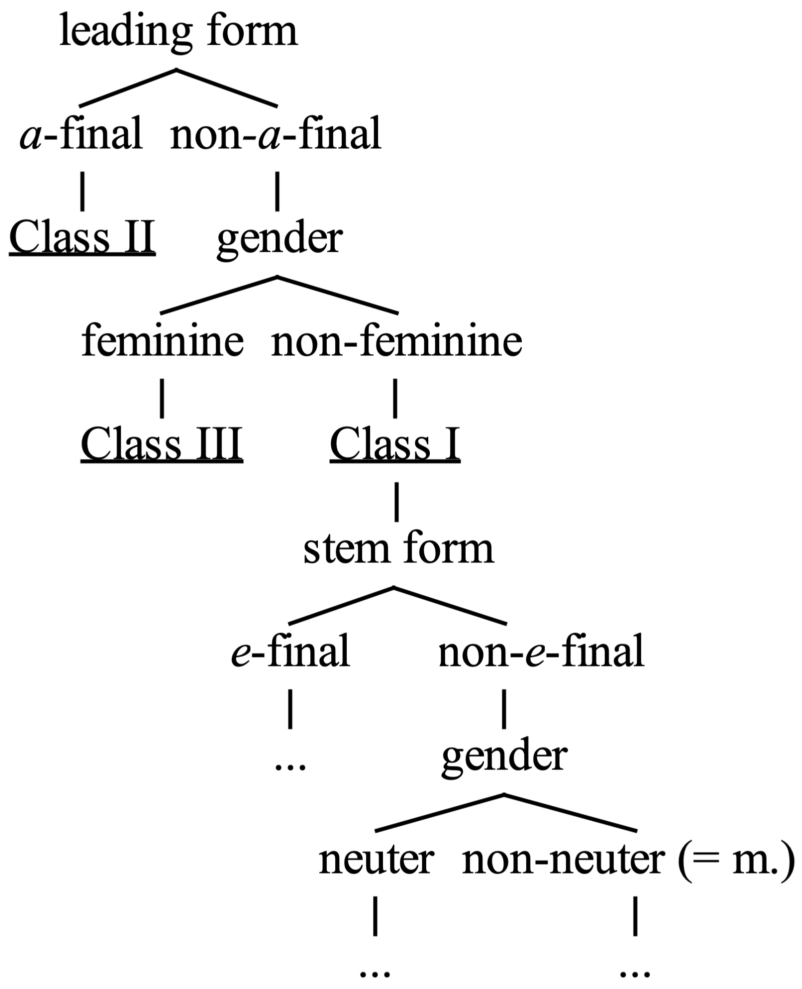
- no copy sets = no epenthesis/insertion of material longer that 1 symbol ☹

- is there a smart way to do this (+ more → non-concatenative morphology) without copy sets? The non-computationally minded often find them counterintuitive

Class membership assignment in SC

- Class membership assignment in SC can be modeled as a predictable process if we take into account a *leading form*/*base* (Wurzel 1990, McCarthy 2005, Albright 2008)
- This I argue to be the nominative singular: it is the most frequent form (subject case), and it ranks highest in the Case Hierarchy (Blake 2001).
- loanwords do not (necessarily) have an oblique form – they are incorporated into the declension system on the basis of their input form (which is borrowed as nominative singular)
- so, lexical entry = {stem, leading form, morphosyntactic info}, but the algorithm should work with a deficient lexical entry too

Class membership assignment in SC



Conclusions

- Inserting *t* repairs aberrant, V-final stems (whenever an overt suffix is attached)
- This happens regardless of whether that suffix is vowel- or consonant-initial
- A means to avoid stipulating listed stems – assuming unpredictable stem allomorphs would basically reduce the phenomenon to an accident
- Assuming morphological epenthesis enables us to express the generalizations explicitly and overtly

Conclusions

- BMRSs can directly capture morphological and phonological generalizations, retaining the computationally restrictive nature of such processes
- Intuitive, easily implementable, extendable to a wider range of phenomena
- Future work: noun class membership assignment base on 'if…then…else' syntax; class membership does not have to be listed information

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