


Algorithmics	Student information	Date	Number of session
	UO: 282276	01/03/2022	5
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## Activity 1. Basic recursive models

### Class: Division 1

The complexity of this algorithm is obtained taking the table that says:

$O(n^k)$  if  $a < b^k$   
 $O(n^k * \log n)$  if  $a = b^k$   
 $O(n^{\log_b n})$  if  $a > b^k$

Therefore, the complexity of this algorithm would be  $O(n^k)$  as  $1 < 3$  that is  $O(n)$ .

### Class: Division 2

The complexity of this algorithm is obtained taking the table that says:

$O(n^k)$  if  $a < b^k$   
 $O(n^k * \log n)$  if  $a = b^k$   
 $O(n^{\log_b n})$  if  $a > b^k$

Therefore, the complexity of this algorithm would be  $O(n^{k * \log n})$  as  $2=2$  that is  $O(n \log n)$ .

### Class: Division 3

The complexity of this algorithm is obtained taking the table that says:

$O(n^k)$  if  $a < b^k$   
 $O(n^k * \log n)$  if  $a = b^k$   
 $O(n^{\log_b n})$  if  $a > b^k$

Therefore, the complexity of this algorithm would be  $O(n^{\log_b n})$  as  $2 > 1$  that is  $O(n^{\log n})$

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## Class: Subtraction 1

The different variables of D&C are:

$$A = 1$$

$$B = 1$$

$$K = 0$$

Taking the table:

$O(n^k)$  if  $a < 1$  (It wont ever happen)  
 $O(n^{k+1})$  if  $a = 1$   
 $O(a^{n/b})$  if  $a > 1$

As  $A = 1$  we can agree that the complexity is  $O(n)$

## Class: Subtraction 2

The different variables of D&C are:

$$A = 1$$

$$B = 1$$

$$K = 1$$

Taking the table:

$O(n^k)$  if  $a < 1$  (It wont ever happen)  
 $O(n^{k+1})$  if  $a = 1$   
 $O(a^{n/b})$  if  $a > 1$

As  $A = 1$  we can agree that the complexity is  $O(n^2)$

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## Class: Subtraction 3

The different variables of D&C are:

$$A = 2$$

$$B = 1$$

$$K = 0$$

Taking the table:

$O(n^k)$  if  $a < 1$  (It wont ever happen)  
 $O(n^{k+1})$  if  $a = 1$   
 $O(a^{n/b})$  if  $a > 1$

As  $A > 1$  we can agree that the complexity is  $O(2^n)$

## Class: Division 4

The different variables of D&C are:

$$A = 4$$

$$B = 2$$

$$K = 1$$

Taking the table:

$O(n^k)$  if  $a < b^k$   
 $O(n^k * \log n)$  if  $a = b^k$   
 $O(n^{\log_b a})$  if  $a > b^k$

As  $A > 2^1$  we can agree that the complexity is  $O(n^{\log_2 4}) == O(n^2)$

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## Class: Subtraction 4

The different variables of D&C are:

$$A = 3$$

$$B = 2$$

$$K = 0$$

Taking the table:

$O(n^k)$  if  $a < 1$  (It wont ever happen)  
 $O(n^{k+1})$  if  $a = 1$   
 $O(a^{n/b})$  if  $a > 1$

As  $A > 1$  we can agree that the complexity is  $O(3^{(n/2)})$