

Andrew Rutherford

CSCI 3104

CPU: 2.8 GHz Intel Core i7

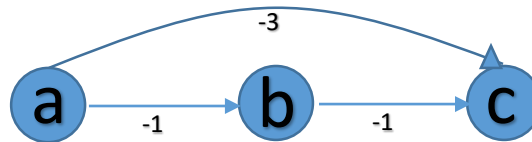
Ram: 16 GB 1600 MHz DDR3

OSX Yosemite

### Homework #5

On my honor, as a University of Colorado at Boulder student, I have neither given nor received any unauthorized help.

1. Adding a large constant to each edge so that all negative edges become positive will not work.



If you were to add a constant of 4 to all the nodes, the distance from a-b-c would be 6, and the distance from a-b would be 1. In reality the shortest distance from a to c would be a-b-c.

2. Running Dijkstra's Algorithm once from city s to city t and again from city t to city s in graph G will give you all shortest path distances between cities s and t. The time complexity of doing this is  $O(|V|^2)$ . In constant time, you can compute the length of the shortest path from city s to city t going through  $e'$  for any  $e' \in E'$ . Which of these paths is shortest give the best edge to add and that length will be the maximum decrease between the two fixed cities. The overall time complexity is  $O(|V|^2 + |E'|)$
3. `[['ME', 'EME', 'MAES', 'MENSA', 'UNSEAM', 'SURNAME', 'ANEURYSM', 'ANEURYSMS', 'NURSERYMAN', 'MENSTRUALLY']]`

This was the first chain output by the algorithm after approximately 45 minutes of run time. A



[illegible]

The programming of the algorithm was a collaborative effort between myself and David Olson.