

# Goal

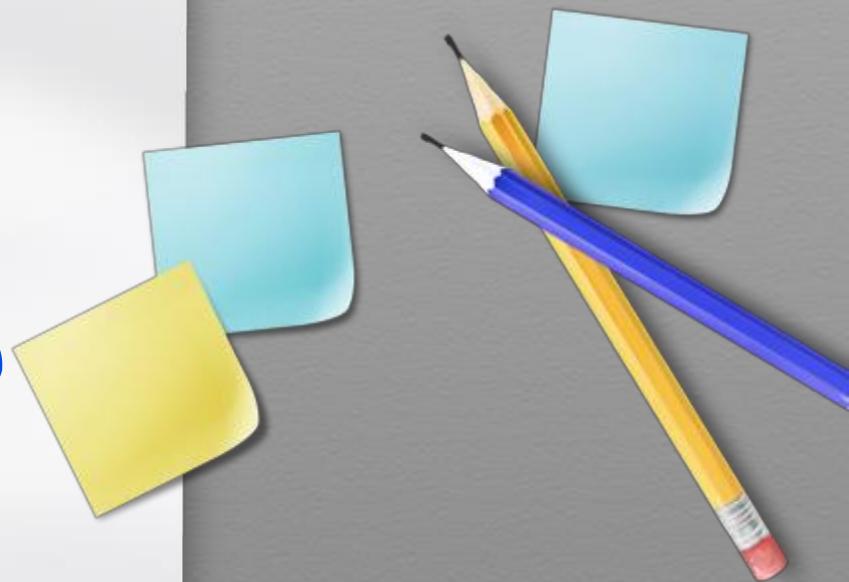
- How to implement the multiplicative inverse of  $GF(2^8)$  of AES

Method using Fermat's Little Theorem

Itoh-Tsujii method

Using Extended Euclidian Algorithm

Reduction to subfield inversion (HW)



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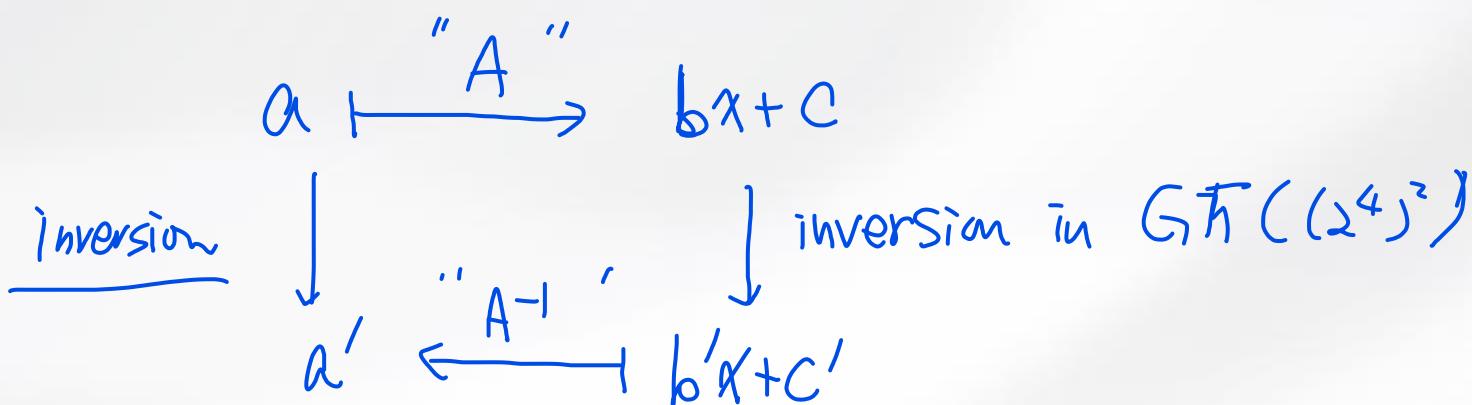


## Multiplicative Inversion in $GF(2^8)$

(4) Reduction to subfield inversion

$$GF(2^8) \xrightarrow{\text{iso} \sim A} GF((2^4)^2), \underline{x^2 + x + \omega}$$

$$\begin{matrix} \omega \\ a \end{matrix} \xrightarrow{\quad} b\chi + c, \quad b, c \in GF(2^4)$$



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## Multiplicative Inversion in $GF(2^8)$

### (4) Reduction to subfield inversion

- Two Approaches to find the inversion in  $GF((2^4)^2)$

✓ (4-1) Direct method: Find  $px + q \in GF((2^4)^2)$  such that  $(bx + c)(px + q) = 1 \Rightarrow (bx + c)^{-1}$

✓ (4-2) Using the following fact: for  $p \in GF((2^n)^m)$ ,

$$p^{-1} = (p^r)^{-1} p^{r-1} \quad p^r \in GF(2^n), r = \frac{2^m - 1}{2^n - 1}$$

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## Multiplicative Inversion in $GF(2^8)$

### (4) Reduction to subfield inversion

Consider  $GF((2^n)^2)$  and defining polynomial  $\underline{x^2 + x + B}$

(4-1) For a given  $\underline{bx + c} \in \boxed{GF((2^n)^2)}$ ,

Find  $\boxed{(bx + c)(px + q) = 1}$  Goal

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## Multiplicative Inversion in $GF(2^8)$

defining  $X^2 + X + B$ .

### (4-1) Reduction to subfield inversion

$$(bx+c)(px+g) = 1$$

$$\Rightarrow bP X^2 + (bg + cP)X + cg = 1$$

Since  $X^2 = x+B$  in  $GF((2^8)^2)$

$$bP(X+B) + (bg + cP)X + cg = 1.$$

$$\Rightarrow \underline{bP + bg + cP}X + \underline{bPB + cg} = 1.$$

$$\therefore \begin{cases} (b+C)P + bg = 0 \\ bBP + cg = 1. \end{cases}$$

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## Multiplicative Inversion in $GF(2^8)$

(4-1) Reduction to subfield inversion

$$(b+c)p + bg = 0 \quad \leftarrow \times c$$

$$bB P + cg = 1 \quad \leftarrow \times b.$$

$$(bc+c^2)p + \cancel{bcg} = 0$$

$$+ \underbrace{b^2Bp + \cancel{bcg}}_{} = b$$

Let  $d := \underbrace{(b^2B + bc + c^2)}_{\text{underlined}} p = b$

$$dP = b \Rightarrow P = d^{-1}b \quad //$$

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## Multiplicative Inversion in $GF(2^8)$

### (4-1) Reduction to subfield inversion

Now, applying  $P = d + b$  to  $bB\beta + cg = 1$ .

$$\Rightarrow b^2Bd^{-1} + cg = 1.$$

$$\Rightarrow cg = b^2Bd^{-1} + 1 = b^2Bd^{-1} + d^{-1}d$$

$$= (b^2B + d)d^{-1}$$

$$= (\cancel{b^2B} + \cancel{b^2B} + bd + c^2)d^{-1}$$

$$\Rightarrow g = (b+c)d^{-1}$$

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## Multiplicative Inversion in $GF(2^8)$

### (4-1) Reduction to subfield inversion

Hence,

$$(b\alpha + c)^{-1} = px + g, \text{ where}$$

$$\begin{cases} p = d+b \\ g = d + (b+c) \end{cases}, \quad \underline{d = b^2B + bc + c^2}$$

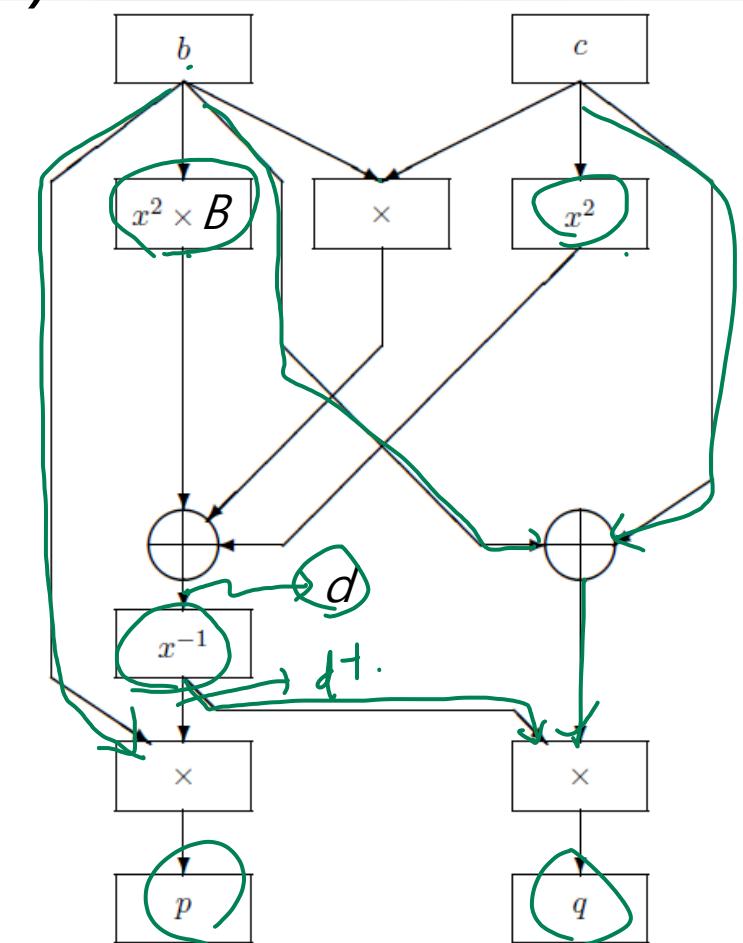


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## Multiplicative Inversion in $GF(2^8)$

$$b^{-1} + c$$

### (4-1) Reduction to subfield inversion



$$d = b^2B + bc + c^2$$

$$P = d + b$$

$$g = d^{-1}(b+c)$$

$$GF(2^8) \approx GF((2^4)^2)$$

Efficient Implementation of the Rijndael S-box, Vincent Rijmen, 2000

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## Multiplicative Inversion in $GF(2^8)$

### (4-2) Reduction to subfield inversion

*Itoh-Tsuji inversion .*

- Fact: For  $P \in GF((2^n)^m)$   $P^{-1} = (P^r)^{-1} P^{r-1}$  ✓

where  $P^r \in GF(2^n)$ . In fact,  $r = \frac{2^{nm} - 1}{2^n - 1}$

$GF(2^8)$   
 $GF(2^4)^*$

$$\text{In our case, } r = \frac{2^{4 \cdot 2} - 1}{2^4 - 1} = \frac{255}{15} = 17$$

$GF((2^2)^2)$

Step 1) compute  $P^{r-1} = P^{16}$  ✓

Step 2) compute  $P^r = (P^{16}) \cdot P$

Step 3) compute  $(P^r)^{-1}$  in  $GF(2^4)$

$p^{-1}$

Step 4) compute  $(P^r)^{-1} \cdot P^{r-1}$  using  $GF((2^2)^2)$  arithmetic

$p^{-1}$

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## Multiplicative Inversion in $GF(2^8)$

### (4-2) Reduction to subfield inversion

Consider the following tower field structure.

$$GF(2^8), \underbrace{z^4 + z + 1}_{\sim} \rightsquigarrow z^4 = z + 1$$



$$GF((2^4)^2), \underbrace{x^2 + x + \phi}, \\ \phi = z^3 + 1 = (1001)$$

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## Multiplicative Inversion in $GF(2^8)$

(4-2) Reduction to subfield inversion

Now, let  $P = P_H X + P_L \in GF(2^8)$ .

$$P_H, P_L \in GF(2^4)$$

$$\begin{aligned} P^{16} &= (P_H X + P_L)^{16} = X^{16} \\ &= P_H^{2^4} X^{2^4} + P_L^{2^4} = P_H X^{2^4} + P_L. \end{aligned}$$

$$\text{Since } A^{2^8} = A \quad \text{in } GF(2^8)$$

$$X^{2^4} = ?$$

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## Multiplicative Inversion in $GF(2^8)$

$$x^2 + x + \phi$$

### (4-2) Reduction to subfield inversion

(Fact)  $x^{2^4} = x + 1$  ✓  $\phi = z^3 + 1$

$$\begin{aligned}
 x^{2^4} &= (x^2)^{2^3} = (x + \phi)^{2^3} = (x + z^3 + 1)^{2^3} \\
 \\ 
 &= (x^2 + z^6 + 1)^{2^2} \\
 &= (x + z^3 + 1 + z^6 + 1)^{2^2} \\
 &= (x^2 + z^6 + z^{12})^2 \\
 &= (x + z^3 + z^6 + z^{12} + 1)^2
 \end{aligned}
 \quad \left| \begin{array}{l}
 = x^2 + z^6 + z^{12} + z^{24} + 1. \\
 = x + z^3 + 1 + z^6 + z^{12} + z^{24} + 1 \\
 = x + z^3 + z^6 + z^{12} + z^{24} \\
 \\ 
 \left( \begin{array}{l}
 z^6 = 2^4 \cdot z^2 \\
 = (z+1)z^2 \\
 = z^3 + z^2
 \end{array} \right)
 \end{array} \right.$$

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## Multiplicative Inversion in $GF(2^8)$

### (4-2) Reduction to subfield inversion

(Fact)  $x^{2^4} = x + 1$

$$z^6 = z^3 + z^2 .$$

$$\begin{aligned}
 &= x + z^3 + z^6 + z^{12} + \cancel{(z^{24})} \cdot (z^6)^4 \\
 &= x + \cancel{z^3} + \cancel{z^3} + z^2 + z^{12} + (z^3 + z^2)^4 \\
 &= x + z^2 + \cancel{z^{12}} + \cancel{z^{12}} + \cancel{(z^8)} = (z^4)^2 . \\
 &= x + z^2 + (z+1)^2 . \\
 &= x + z^2 + z^2 + 1 . \quad \therefore \underline{x^{2^4} = x+1} // \\
 &= x + 1
 \end{aligned}$$

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## Multiplicative Inversion in $GF(2^8)$

### (4-2) Reduction to subfield inversion

Therefore,

$$\text{Step 1) } P^{16} = (P_H X + P_L)^{16} = P_H X^{16} + P_L.$$

$$= P_H (X+1) + P_L.$$

$$= P_H X + (P_H + P_L).$$

$X+1$   
//

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## Multiplicative Inversion in $GF(2^8)$

(4-2) Reduction to subfield inversion

$$\begin{aligned} \text{(Step 2)} \quad P^{19} &= P^8 \cdot P \quad X^2 = X + \phi \\ &= (P_H X + (P_H + P_L))(P_H X + P_L). \\ &= P_H^2 X^2 + \cancel{P_H P_L X} + (P_H^2 + \cancel{P_H P_L})X + P_H P_L + \cancel{P_L^2}. \\ &= \cancel{P_H^2}(X + \phi) + \cancel{P_H^2}X + P_H P_L + \cancel{P_L^2}. \\ &= \underline{\cancel{P_H^2} \phi + P_H P_L + \cancel{P_L^2}}. \end{aligned}$$

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# Multiplicative Inversion in $GF(2^8)$

## (4-2) Reduction to subfield inversion

$$P \rightarrow P^{16} \rightarrow \underline{P^{16} \cdot P} \xrightarrow{\chi^{-1}} (P^{17})^{-1} \rightarrow \underline{(P^{17})^{-1} \cdot P^{16}}.$$

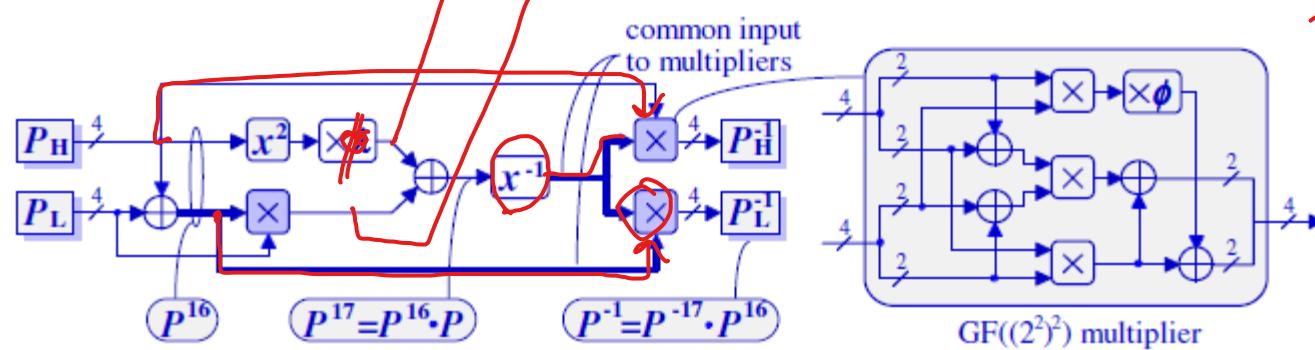
$$P_H X + (P_H + P_L)$$

$$P_H^2 \phi + P_H P_L + P_L^{2a}$$

$$\underline{P_H^2 \phi} + \underline{(P_H + P_L) P_L}.$$

$$d^{-1}(P_H \wedge (P_H \# B))$$

$$\frac{(d - P_H)}{d + (P_H + P_L)}$$



# Q&A



KEEP  
CALM  
AND  
STUDY  
ON

