Computer Arithmetic Algorithms, 2nd Edition Israel Koren

List of Corrections for the first printing (2002)

- 1. Page 33, reference [6], replace "(2001)." by "(November 2001), 1267-1278."
- 2. Page 36, in the equation on line 3, replace the last line

$$= \left(\sum_{j=0}^{n-2} x_j \ 2^{-(n-1-j)} \right) \cdot A = 2^{(n-1)} \left(\sum_{j=0}^{n-2} x_j \ 2^j \right) \cdot A$$

by

$$= \left(\sum_{j=0}^{n-2} x_j \ 2^{-(n-1-j)} \right) \cdot A = 2^{-(n-1)} \left(\sum_{j=0}^{n-2} x_j \ 2^j \right) \cdot A$$

3. Page 49, line 4 (2nd line of the equation for r_m), replace

$$r_m = 2r_{m-1} - q_m(2Q_{m-1} + q_m 2^{-m})$$

= $2^2r_{m-2} - 2q_{m-1}(2Q_{m-2} + q_{m-1}) - q_m(2Q_{m-1} + q_m 2^{-m})$

by

$$r_m = 2r_{m-1} - q_m(2Q_{m-1} + q_m 2^{-m})$$

= $2^2 r_{m-2} - 2q_{m-1}(2Q_{m-2} + q_{m-1} 2^{-(m-1)}) - q_m(2Q_{m-1} + q_m 2^{-m})$

- 4. Page 68, line 15, replace the sentence "A similar effect is achieved by using a significand of 1.f instead of 0.f, since this adds 1 to the exponent." by "A similar effect is achieved by using a significand of 1.f instead of 0.1f, since this adds 1 to the exponent."
- 5. Page 82, last paragraph, replace the sentence "In the *CLOSE* case we first predict the exponent difference, based on the two least significant bits of the operands, ..." by "In the *CLOSE* case we first predict the exponent difference, based on the two least significant bits of the two exponents, ..."
- 6. Page 83, Table 4.10, Step 1 for the *CLOSE* case, replace "Predict exponent" by "Predict exponent difference"
- 7. Page 110, last sentence should be "Note that unlike Figure ??, the Ladner-Fischer adder employs fundamental carry operators with a fan-out value higher than 2. Such an implementation with fan-outs of up to n/2 requires buffers which add to the overall delay.

- 8. Page 110, Figure 5.10, remove the right portion of the figure.
- 9. Page 129, 1st line (below Figure 5.25) should say: k primary inputs
- 10. Page 159, last sentence should be: The implementation in Figure 6.12 corresponds to the setting a = b = d = 0 and c = e = f = 1.
- 11. Page 167, last sentence in Section 6.5, replace "In this case, the rounding step for $(A \times B + C)$ is performed at the same time as the multiplication by D, by adding the partial product $Incr. \times D$ to the CSA tree." by "In this case, the rounding step for $(X \times Y + Z)$ is performed at the same time as the multiplication by B, by adding the partial product $Incr. \times B$ to the CSA tree."
- 12. Page 180, reference [14], replace "600 floating-point unit" by "6000 floating-point unit"
- 13. Page 206, line 5 from the bottom, replace "where Q_{i-1} is the partially calculated root at step (i-1), i.e., $Q_{i-1} = 0.q_1q_2 \cdot q_{i-1}$." by "where Q_{i-1} is the partially calculated root at step (i-1), i.e., $Q_{i-1} = 0.q_1q_2 \cdots q_{i-1}$."
- 14. Page 208, equation (7.22), replace

$$\sqrt{1.f \cdot 2^{E-1023}} = \begin{cases} \sqrt{0.1f} \cdot 2^{(E-1)/2-1023} & \text{if } E \text{ is odd} \\ \sqrt{0.01f} \cdot 2^{E/2-1-1023} & \text{if } E \text{ is even} \end{cases}$$
 (7.22)

by

$$\sqrt{1.f \cdot 2^{E-1023}} = \begin{cases} \sqrt{0.1f} \cdot 2^{(E+1)/2-1023} & \text{if } E \text{ is odd} \\ \sqrt{0.01f} \cdot 2^{E/2+1-1023} & \text{if } E \text{ is even} \end{cases}$$
 (7.22)

15. Page 241 equation (9.53), replace

$$y_m = y_i \prod_{k=i}^{m-1} (1 + s_k 2^{-k}) \approx y_i \left(1 + s_i 2^{-i} + s_{i+1} 2^{-(i+1)} + s_{i+1} 2^{-(i+1)} + \cdots \right)$$
 (9.53)

by

$$y_m = y_i \prod_{k=i}^{m-1} (1 + s_k 2^{-k}) \approx y_i \left(1 + s_i 2^{-i} + s_{i+1} 2^{-(i+1)} + s_{i+2} 2^{-(i+2)} + \cdots \right)$$
(9.53)

16. Page 275, Exercise 11.9, replace $A - 2^a - 1$ by $A = 2^a - 1$.