

CPE166 Advanced Logic Design

Lab 4 Report: Simplified Microprocessor Design

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Introduction

A microprocessor is a computer processor that incorporates the functions of a central processing unit on a single integrated circuit (IC). Microprocessors contain both combinational logic and sequential digital logic. For this lab we needed to build a simplified microprocessor in which it would take some inputs and would output a value. The design needed to implement the following logic equation: $R2 = M0 + (\text{not } M1) + \text{Cin}$. The lab will be divided into three parts to reduce the amount of errors at the end. The first part of this lab will be to create the datapath. The datapath contains the ALU, mux2to1, mux4to1, and D-flip-flop. For the second part, we will create the FSM to determine the internal inputs for datapath. The third part of this lab will be to create a top file where we combine the datapath and FSM with three added inputs: m0, m1, and m2.

Part 1: Microprocessor Data Path Design

Design Purpose:

This part of the lab will be to create the datapath circuit as shown in the lab manual. It consists of 3 mux2to1 circuits, 4 D-flip-flops, 1 mux4to1, and the ALU. The ALU will compute the functions as shown in the lab manual depending on the inputs 'a' and 'b'. To test the datapath, we will create a testbench file giving values to m0, m1, m2, ce, w, s, sel, and cin. Each of these internal signals will determine how data will be passed through each of these circuits and that is how we will be able to track if our output is correct.

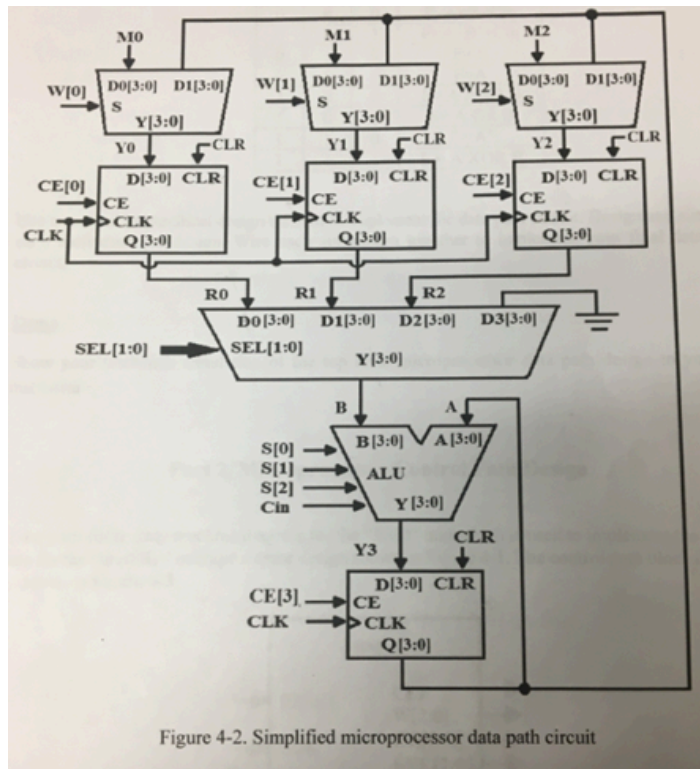


Table 4-1. ALU truth table

| S[2] | S[1] | S[0] | ALU Output F |
|------|------|------|------------------------|
| 0 | 0 | 0 | $F = A + B + C_{in}$ |
| 0 | 0 | 1 | $F = A + B' + C_{in}$ |
| 0 | 1 | 0 | $F = B$ |
| 0 | 1 | 1 | $F = A$ |
| 1 | 0 | 0 | $F = A \text{ AND } B$ |
| 1 | 0 | 1 | $F = A \text{ OR } B$ |
| 1 | 1 | 0 | $F = A'$ |
| 1 | 1 | 1 | $F = A \text{ XOR } B$ |

Design Source Code:

Multiplexer 2-to-1:

```
1 module mux2to1(d0, d1, s, y);
2     input[3:0] d0, d1;
3     input s;
4     output y;
5     reg[3:0] y;
6
7     always@(d0 or d1 or s)
8     begin
9         if(s)
10             y = d1;
11         else
12             y = d0;
13     end
14
15 endmodule
16
```

```
1 module mux2to1_tb;
2
3     reg[3:0] d0, d1;
4     reg s;
5     wire[3:0] y;
6
7     mux2to1 u1(d0, d1, s, y);
8
9     initial begin
10         d1 = 4'b0000; d0 = 4'b0000; s = 1'b0;
11         #10 d1 = 4'b0001; d0 = 4'b0010; s = 1'b1;
12         #10 d1 = 4'b0010; d0 = 4'b1000; s = 1'b0;
13         #10 d1 = 4'b1001; d0 = 4'b0100; s = 1'b1;
14         #10 $stop;
15     end
16 endmodule
17
```

D flip-flop:

```

dff.v
1  module dff (clk, clr, d, q, ce);
2  input[3:0] d;
3  input clk, ce, clr;
4  output q;
5  reg[3:0] q;
6
7  always@(posedge clr or posedge clk)
8  begin
9      if(clr)
10         q <= 0;
11     else if(ce)
12         q <= d;
13 end
14 endmodule
15

```

```

dff_tb.v
1  module dff_tb;
2
3  reg[3:0] d;
4  reg clk, clr, ce;
5
6  wire[3:0] q;
7
8  dff u1(clk, clr, d, q, ce);
9
10 initial clk = 0;
11 always #10 clk = ~clk;
12
13 initial begin
14     clr = 1; d = 4'b0001; ce = 1;
15     #20
16     clr = 0; d = 4'b1000; ce = 0;
17     #20
18     clr = 0; d = 4'b0001; ce = 1;
19     #20
20     clr = 0; d = 4'b1110; ce = 1;
21     #40
22     $stop;
23 end
24
25 endmodule
26

```

Multiplexer 4-to-1:

```

mux4to1.v
1  module mux4to1(d0, d1, d2, d3, sel, y);
2      input[3:0] d0,d1,d2,d3;
3      input[1:0] sel;
4      output y;
5      reg[3:0] y;
6
7      always @(d0 or d1 or d2 or d3 or sel) begin
8          case (sel)
9              2'b00: y = d0;
10             2'b01: y = d1;
11             2'b10: y = d2;
12             2'b11: y = d3;
13         endcase
14     end
15 endmodule
16

```

```

mux4to1_tb.v
1  module mux4to1_tb;
2
3      reg[3:0] d0, d1, d2, d3;
4      reg[1:0] sel;
5      wire[3:0] y;
6
7      mux4to1 u1(d0, d1, d2, d3, sel, y);
8
9      initial begin
10         d0 = 4'b1100; d1 = 4'b1010; d2 = 4'b0001; d3 = 4'b0101; sel = 2'b00;
11         #10
12         d0 = 4'b1100; d1 = 4'b1010; d2 = 4'b0001; d3 = 4'b0101; sel = 2'b01;
13         #10
14         d0 = 4'b1100; d1 = 4'b1010; d2 = 4'b0001; d3 = 4'b0101; sel = 2'b10;
15         #10
16         d0 = 4'b1100; d1 = 4'b1010; d2 = 4'b0001; d3 = 4'b0101; sel = 2'b11;
17         #10
18         $stop;
19     end
20 endmodule
21

```

ALU:

```

1 module alu(a, b, s, cin, y);
2   input[3:0] a, b;
3   input[2:0] s;
4   input cin;
5   output y;
6   reg[3:0] y;
7
8   always@(a or b or s or cin)
9   begin
10      case(s)
11         3'b000: y = a + b + cin;
12         3'b001: y = a + ~b + cin;
13         3'b010: y = b;
14         3'b011: y = a;
15         3'b100: y = a & b;
16         3'b101: y = a | b;
17         3'b110: y = ~a;
18         3'b111: y = a ^ b;
19      endcase
20   end
21 endmodule
22

```

```

1 module alu_tb;
2
3   reg[3:0] a, b;
4   reg[2:0] s;
5   reg cin;
6
7   wire[3:0] y;
8
9   alu u1(a, b, s, cin, y);
10
11   initial begin
12     s = 3'b000; a = 4'b1010; b = 4'b0101; cin = 1'b0;
13     #10
14     s = 3'b001; a = 4'b1000; b = 4'b1111; cin = 1'b1;
15     #10
16     s = 3'b010; a = 4'b0001; b = 4'b0011; cin = 1'b0;
17     #10
18     s = 3'b011; a = 4'b1000; b = 4'b0000; cin = 1'b0;
19     #10
20     s = 3'b100; a = 4'b1111; b = 4'b0000; cin = 1'b0;
21     #10
22     s = 3'b101; a = 4'b1010; b = 4'b0101; cin = 1'b0;
23     #10
24     s = 3'b110; a = 4'b0000; b = 4'b0000; cin = 1'b0;
25     #10
26     s = 3'b111; a = 4'b1111; b = 4'b0000; cin = 1'b0;
27     #10
28     $stop;
29   end
30
31 endmodule
32

```


Datapath:

```

1 module datapath(m0, m1, m2, cin, clr, w, ce, sel, s, clk, r0, r1, r2);
2   input clk, clr, cin;
3   input[3:0] m0, m1, m2;
4   input[3:0] ce;
5   input[2:0] w;
6   input[2:0] s;
7   input[1:0] sel;
8
9   output[3:0] r0, r1, r2;
10
11  wire [3:0] y0, y1, y2, y3, a, b;
12
13  mux2to1 g1(.d0(m0), .d1(a), .s(w[0]), .y(y0));
14  mux2to1 g2(.d0(m1), .d1(a), .s(w[1]), .y(y1));
15  mux2to1 g3(.d0(m2), .d1(a), .s(w[2]), .y(y2));
16
17  dff d1(.clk(clk), .clr(clr), .ce(ce[0]), .d(y0), .q(r0));
18  dff d2(.clk(clk), .clr(clr), .ce(ce[1]), .d(y1), .q(r1));
19  dff d3(.clk(clk), .clr(clr), .ce(ce[2]), .d(y2), .q(r2));
20
21  mux4to1 u1(.d0(r0), .d1(r1), .d2(r2), .d3(4'b0000), .sel(sel), .y(b));
22
23  alu a1(.a(a), .b(b), .s(s), .cin(cin), .y(y3));
24
25  dff d4(.clk(clk), .clr(clr), .ce(ce[3]), .d(y3), .q(a));
26
27  endmodule
28

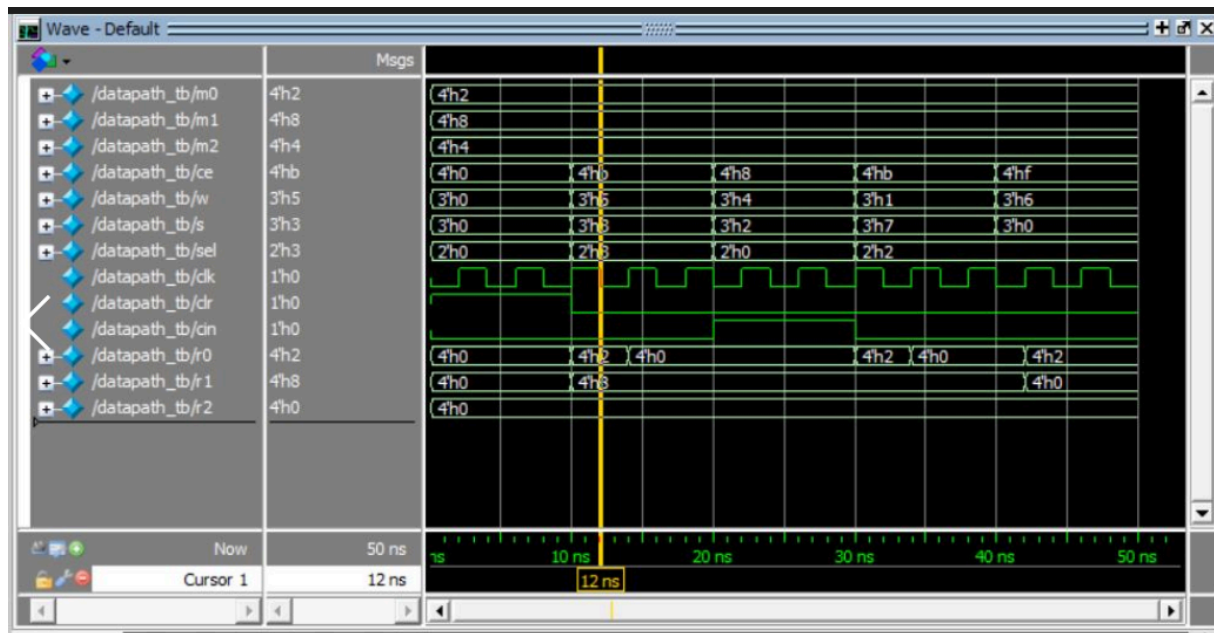
```

```

1 module datapath_tb;
2   reg[3:0] m0, m1, m2, ce;
3   reg[2:0] w, s;
4   reg[1:0] sel;
5   reg clk, clr, cin;
6   wire [3:0] r0, r1, r2;
7
8   datapath uut(m0, m1, m2, cin, clr, w, ce, sel, s, clk, r0, r1, r2);
9
10  initial
11  begin
12    clr = 1'b1; w = 3'b000; ce = 4'b0000; sel = 2'b00; s = 3'b000; clk = 0;
13    m2 = 4'b0100; m0 = 4'b0010; m1 = 4'b1000; cin = 0;
14  end
15
16  always
17  begin
18    #2
19    clk = ~clk;
20  end
21
22  initial
23  begin
24    #10;
25    clr = 1'b0; ce = 4'b1011; sel = 2'b11; s = 3'b011; w = 3'b101; cin = 0;
26    #10;
27    clr = 1'b0; ce = 4'b1000; sel = 2'b00; s = 3'b010; w = 3'b100; cin = 1;
28    #10;
29    clr = 1'b0; ce = 4'b1011; sel = 2'b10; s = 3'b111; w = 3'b001; cin = 0;
30    #10;
31    clr = 1'b0; ce = 4'b1111; sel = 2'b10; s = 3'b000; w = 3'b110; cin = 0;
32    #10;
33    $stop;
34  end
35
36  endmodule
37

```

Simulation Waveform:



In the waveform, we see that: $r0 = 'a'$ because $w[0]$ is '1' and $ce[0]$ is '1'. $r1 = m2$ because $w[1]$ is '0' therefore the data passed through is M1 instead of 'a'. $r2 = 0$ because $ce[2]$ is '0' so neither M2 or a is passed to $r2$. The 4-to-1 mux passes D3 which is '0000' because $sel = '111'$ therefore b is '0000' the alu function performed is $F = A$ because $'s' = '011'$.

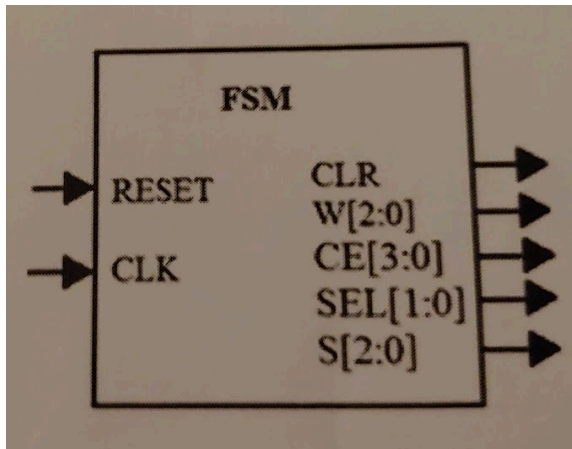
Result Discussion:

This part of the lab is a lot of work but it was manageable. The MUX selects properly, the flip flops store data properly, and the ALU does all of the calculations properly. The data path waveform shows the simplified processor doing the $R2 = M0 + (\text{not } M1) + \text{Cin}$ calculation. The hard part of this lab was creating the testbench for the datapath and telling whether or not it was correct because I didn't know what 'a' was.

Part 2: Microprocessor Control Path Design

Design Purpose:

Once the data path module is built, we can proceed to building the finite state machine. The FSM will determine the internal signals for the datapath. It will also have a reset and a clk. The FSM will have 6 states each with different random data to confirm that each state transitions correctly.



Design Source Code:

```
1 module FSM(clk, reset, clr, w, ce, sel, s);
2
3 input clk, reset;
4
5 output clr;
6 output[2:0] w, s;
7 output[1:0] sel;
8 output[3:0] ce;
9
10 reg clr;
11 reg [2:0] w, s;
12 reg [1:0] sel;
13 reg [3:0] ce;
14
15 reg [2:0] cs, ns;
16
17 parameter s0 = 0, s1 = 1, s2 = 2, s3 = 3, s4 = 4, s5 = 5;
18
19 always@(posedge clk or posedge reset)
20 begin
21     if(reset)
22         cs <= s0;
23     else
24         cs <= ns;
25 end
26
27 always @ (cs)
28 begin
29     case(cs)
30         s0: ns = s1;
31         s1: ns = s2;
32         s2: ns = s3;
33         s3: ns = s4;
34         s4: ns = s5;
35         s5: ns = s5;
36         default: ns = s0;
37     endcase
38 end
39
```

```

40  always@(cs)
41  begin
42      case(cs)
43      s0:  begin
44              clr = 1'b1;
45              w = 3'b000;
46              s = 3'b000;
47              sel = 2'b00;
48              ce = 4'b0000;
49          end
50
51      s1:  begin
52              clr = 1'b0;
53              w = 3'b101;
54              s = 3'b011;
55              sel = 2'b11;
56              ce = 4'b1011;
57          end
58
59      s2:  begin
60              clr = 1'b0;
61              w = 3'b100;
62              s = 3'b010;
63              sel = 2'b00;
64              ce = 4'b1000;
65          end
66
67      s3:  begin
68              clr = 1'b0;
69              w = 3'b001;
70              s = 3'b111;
71              sel = 2'b10;
72              ce = 4'b1011;
73          end
74
75      s4:  begin
76              clr = 1'b0;
77              w = 3'b110;
78              s = 3'b000;
79              sel = 2'b10;
80              ce = 4'b1111;
81          end
82
83      s5:  begin
84              clr = 1'b1;
85              w = 3'b111;
86              s = 3'b111;
87              sel = 2'b11;
88              ce = 4'b1111;
89          end
90      endcase
91  end
92
93  endmodule
94
95

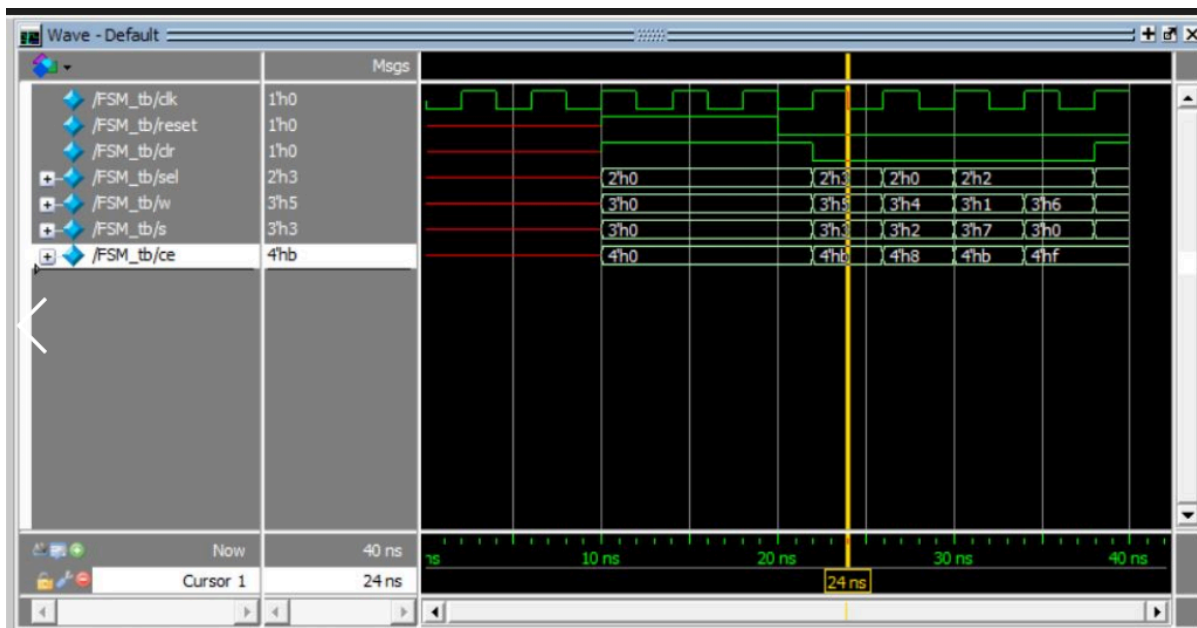
```

Testbench:

```

1 module FSM_tb();
2   reg clk, reset;
3   wire clr;
4   wire [1:0] sel;
5   wire [2:0] w, s;
6   wire [3:0] ce;
7
8   FSM uut(clk, reset, clr, w, ce, sel, s);
9
10  always
11  begin
12    #2
13    clk = ~clk;
14  end
15
16  initial
17  begin
18    clk = 0;
19    #10
20    reset = 1;
21    #10
22    reset = 0;
23    #20
24    $stop;
25  end
26 endmodule
27

```

Simulation Waveform:

Here we can see that the FSM shows the correct data for each state accordingly from the testbench.

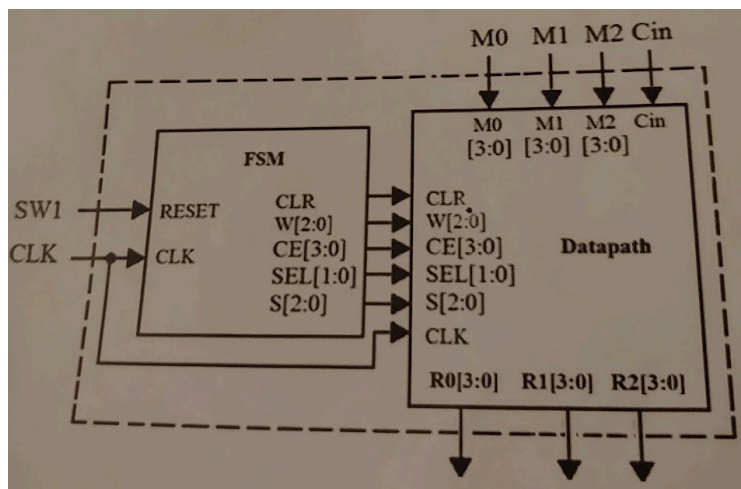
Result Discussion:

The point of the FSM is to control the internal signals to the data path. Since I chose random numbers, all I had to do was create a simple FSM so I can put it together with the datapath circuit for the next part.

Part 3: Final Simplified Microprocessor Design

Design Purpose:

This part of the lab is to put the FSM together with the datapath with a “top” file. The top will determine the inputs m_0 , m_1 , m_2 , and C_{in} . The output should display r_0 , r_1 , and r_2 . The main focus of this last part is to learn how to integrate the data path and control path to design a simplified microprocessor.



Design Source Code:

```

1 module TOP( clk, sw1, m0, m1, m2, cin, r0, r1, r2);
2 input clk, cin, sw1;
3 input [3:0] m0, m1, m2;
4
5 output [3:0] r0, r1, r2;
6
7 wire [2:0] w, s;
8 wire [1:0] sel;
9 wire [3:0] ce;
10 wire clr;
11
12
13 FSM g1(.clk(clk), .reset(sw1), .clr(clr), .w(w), .ce(ce), .sel(sel), .s(s));
14
15 datapath g2(.clk(clk), .clr(clr), .w(w), .ce(ce), .sel(sel), .s(s), .m0(m0), .m1(m1), .m2(m2), .cin(cin)
16
17 endmodule
18

```

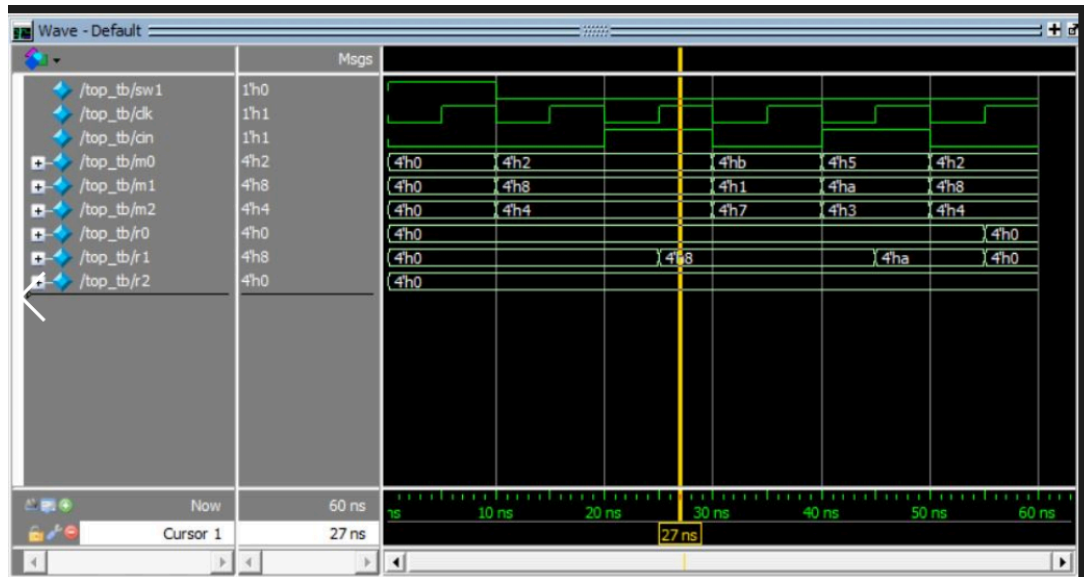
Testbench:

```

1 module top_tb;
2 reg sw1, clk, cin;
3 reg [3:0] m0, m1, m2;
4
5 wire [3:0] r0, r1, r2;
6
7 TOP uut(.clk(clk), .sw1(sw1), .m0(m0), .m1(m1), .m2(m2), .cin(cin), .r0(r0), .r1(r1), .r2(r2));
8
9 initial clk = 0;
10 always
11 #5 clk = ~clk;
12
13 initial begin
14 sw1 = 1; m0 = 4'b0000; m1 = 4'b0000; m2 = 4'b0000; cin = 0;
15 #10;
16 sw1 = 0; m0 = 4'b0010; m1 = 4'b1000; m2 = 4'b0100; cin = 0;
17 #10;
18 sw1 = 0; m0 = 4'b0010; m1 = 4'b1000; m2 = 4'b0100; cin = 1;
19 #10;
20 sw1 = 0; m0 = 4'b1011; m1 = 4'b0001; m2 = 4'b0111; cin = 0;
21 #10;
22 sw1 = 0; m0 = 4'b0101; m1 = 4'b1010; m2 = 4'b0011; cin = 1;
23 #10;
24 sw1 = 0; m0 = 4'b0010; m1 = 4'b1000; m2 = 4'b0100; cin = 0;
25 #10;
26 $stop;
27 end
28
29 endmodule
30

```


Simulation Waveform:



The top inserts m0, m1, and m2 into the datapath and the FSM determines the internal wiring.

Result Discussion:

This part wasn't hard to finish. The only thing we had to do is create a top file to connect the datapath and FSM. Then we just had to create a testbench to show the waveform. It is hard to see if the waveform is correct but if we look at the internal signals that the FSM is passing and what the inputs for the top file is, then we can compare the answer with the equation given in the lab manual to test if it is correct.

Conclusion

The main objective of this lab was to design a simple 4 bit microprocessor that implemented the logic equation $R2 = M0 + \text{not}(M1) + \text{Cin}$. We learned how to design data paths, control paths, and how to integrate both of them. The use of hierarchical design was used to complete this lab and reduce many errors. In all, this lab gave me great experience with combinational and sequential circuits. It took me a while to understand how the data path passed around data but as I looked at the D-flip-flop, multiplexers, and ALU separately it started to make sense. I learned how a simple microprocessor functions from this lab.