## **Lecture 5 and Lab5 notes**

## **Performance**

1. definition

$$Performance = \frac{1}{Execution \ time}$$

2. Performance ratio

what does "x is n times faster than y" mean?

$$Performance \ ratio = rac{Performance_x}{Performance_y} = rac{Execution \ time_y}{Execution \ time_x}$$

- 3. Measuring execution time
  - Elapsed time: total responce time, including all aspects processing, IO, OS overhead, idle time
  - 2. CPU time: time spent processing a specific task

Minus IO time, other jobs' shared

Includes **user CPU time**(time that the CPU spends executing code in the user space of a program)

and **system CPU time**(includes the time spent executing system calls, handling hardware interrupts, and managing system resources such as memory and devices)

- 4. CPU clocking
  - Operation of digital hardware governed by a constant-rate clock
  - Notions: Clock cycle, Clock period, Clock rate/frequency
  - Clock period (Clock cycle time): duration of a clock cycle
  - Clock frequency (rate): cycles per second
  - $Clock\ cycle\ time = 1/Clock\ rate$
- 5. CPU time

$$CPU\ time = Clock\ cycles\ counts*Clock\ period$$

$$= \frac{Clock\ cycles\ counts}{Clock\ rate}$$

performance can be improved by:

- reducing the number of cycle counts
- · Increasing clock rate
- 6. IC and CPI
  - IC: instruction count for a program, determined by program, ISA, compiler, have nothing to do with the computer itself

- CPI: cycles per instrucion. Generally, CPI means average CPI, because each kind of instruction have different cost. Determined by CPU hardware
- Thus

$$egin{aligned} Clock \ cycles &= IC*CPI \ CPU \ time &= IC*CPI*T_c \ &= rac{IC*CPI}{f} \end{aligned}$$

7. Considering different instruction types

$$egin{aligned} Clock \ cycles &= \sum_{i=1}^{n} (CPI_i * IC_i) \ Avg \ CPI &= rac{\sum_{i=1}^{n} (CPI_i * IC_i)}{\sum_{i=1}^{n} IC_i} \end{aligned}$$

Compiler optimization:

- minimize execution time: e.g. GNU gcc -O3
- minimize code size: e.g. GNU gcc -Os
- 8. Factors

|                        | Instruction Count | СРІ       | Clock Rate |
|------------------------|-------------------|-----------|------------|
| Algorithm              | $\checkmark$      | $\sqrt{}$ |            |
| Language               | √                 | V         |            |
| Compiler               | V                 | √         |            |
| ISA                    | √                 | <b>√</b>  | <b>√</b>   |
| Organization(hardware) |                   | <b>√</b>  | V          |

## **Energy dissipation**

- 1. Energy consumption = dynamic energy + static energy
  - Dynamic energy: (energy spent when transistors switch from  $0\rightarrow 1$   $1\rightarrow 0$ ) is primary
  - Static energy: leakage
- 2. Energy and power

$$Energy \propto Capacitive\ load * Voltage^2$$
 
$$\propto CV^2$$
 
$$Power \propto \frac{1}{2} * Capacitive\ load * Voltage^2 * Frequency\ switched$$

 $\propto \frac{1}{2}CV^2f$ 

- 1. a pitfall: Improving an aspect of a computer and expecting a proportional improvement in overall performance.
- 2. Amdahl's law
  - Calculate how much a computation can be sped up by running part of it in parallel

$$T_{improved} = rac{T_{affected}}{improvement\ factor} + T_{unaffected}$$

## **Benchmark**

- 1. Select a set of programs that represent the workload
- 2. Run these programs on each computer
- 3. Compare the average execution time (**geometric mean**)

$$\sqrt[n]{\prod_{i=1}^n Execution\ time_i}$$

4. SPEC ratio

 $excution\ time_{reference\ machine} \ excution\ time_{measure\ machine}$