

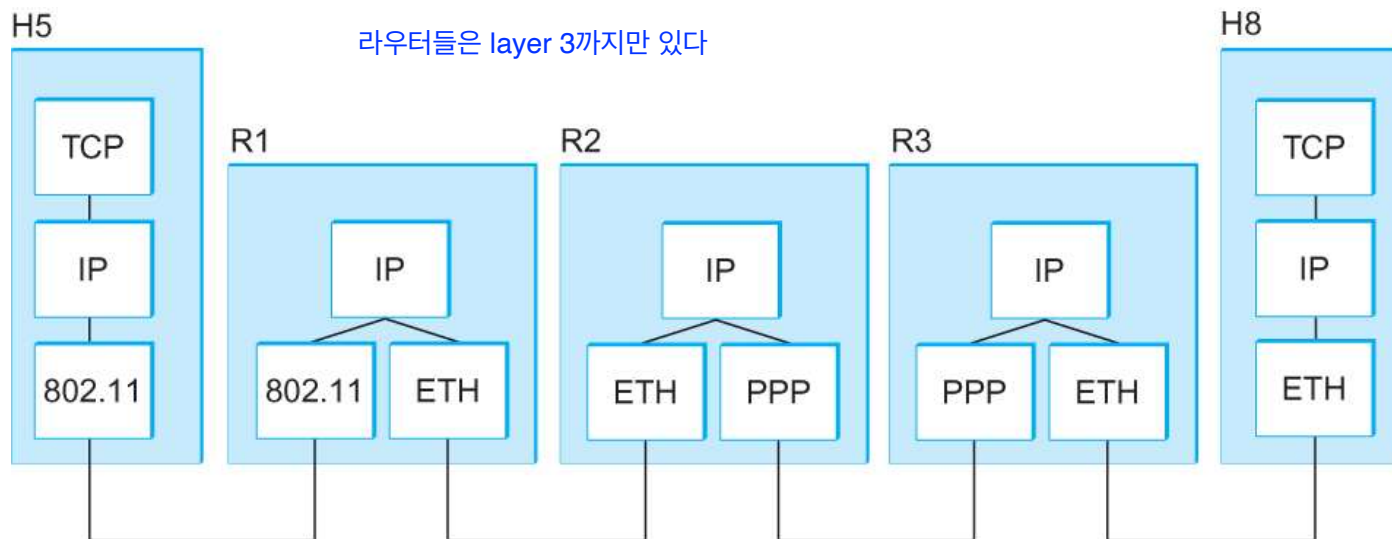
# Internet Protocol (IP)

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Dept. of Computer Science and Engineering  
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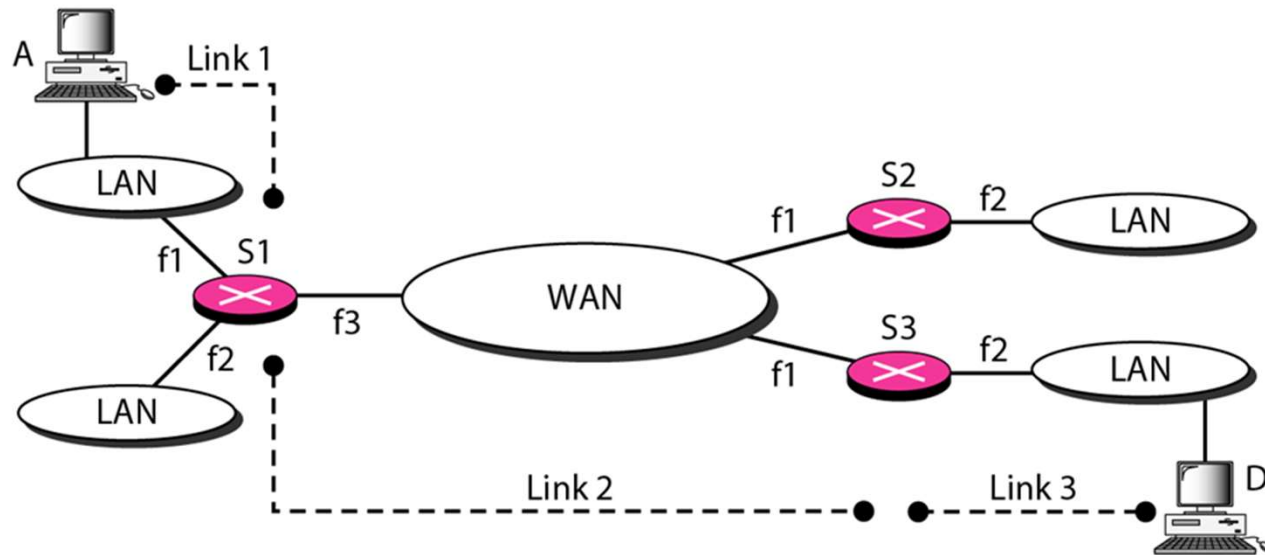
# Internet Protocol (IP)

- A Network layer protocol
- Allows delivery of packets across heterogeneous systems



# IP: A Network Layer Protocol

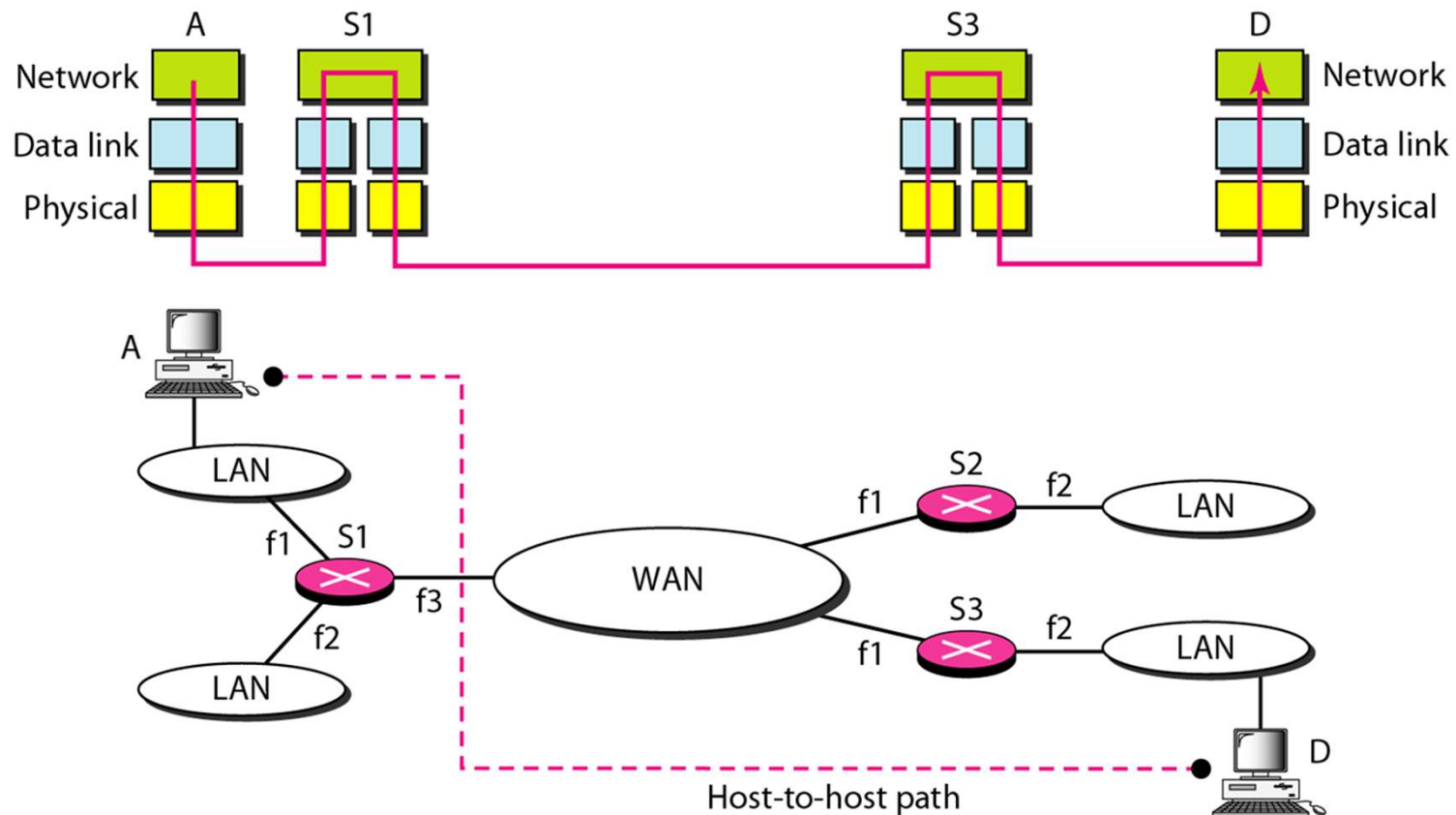
- Packet flow without Network layer



Communication only possible between physically connected hosts

# IP: A Network Layer Protocol

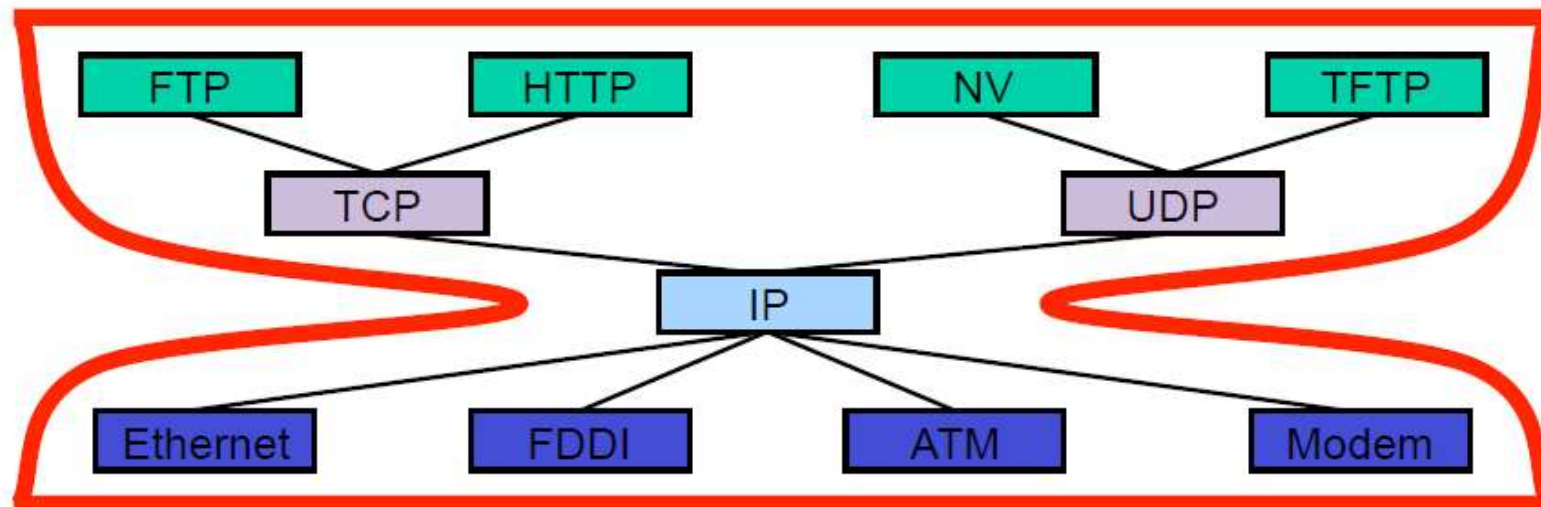
- Packet flow with Network layer



“Multi-hop” communication possible using routers

# Hourglass model

- IP is the unified protocol required for all routers to implement
- Networks may implement different transport, data link, and physical layer protocols, but they all must implement IP



# IP: Service Model

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- Packet delivery model

- Connectionless

최선을 다하겠다 - 보장못한단 뜻

- Best-effort (unreliable)

- Packets may get lost
    - Packets can be delivered out-of-order
    - Duplicate copies may be delivered
    - Packets can be delayed for a long time

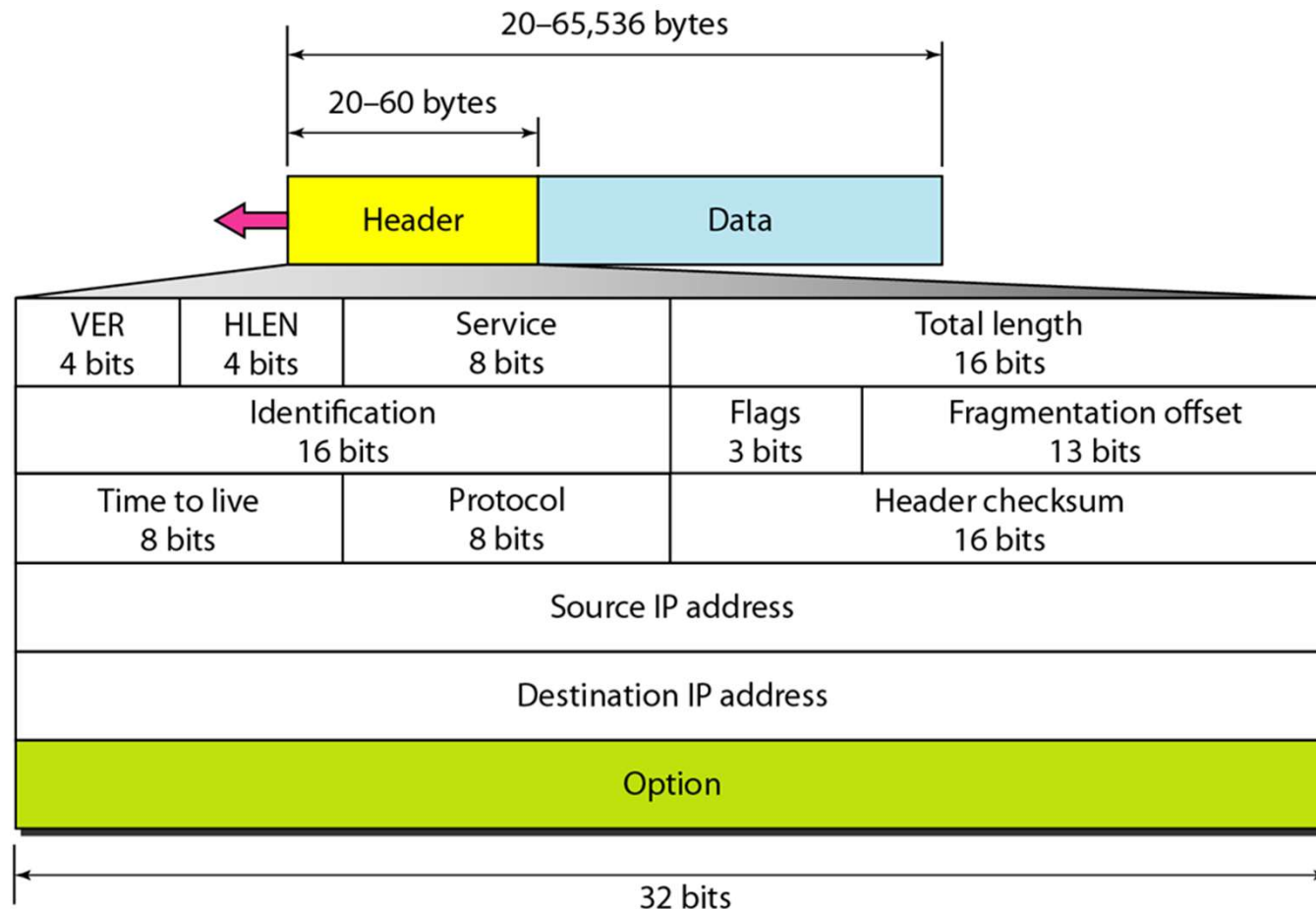
connection oriented - 전화  
connectionless - 우표

- Global Addressing Scheme

- Provides a way to identify all hosts in the network

# IP: Packet Format

- IP Header is attached in front of the packet
- 20 bytes with no option, up to 60 bytes with option



# IP header

- VER: IPv4 or IPv6
- HLEN: length of the header
  - unit: word = 4 bytes
  - If the header is 20 bytes, then HLEN = 5

시험때 이 그림이랑 헤더를 16진수로 줘서  
어느 부분이 src, dest ip 인지 아는 문제 나옴

hlen 값에 4를 곱해야 헤더 크기가 나옴  
ex. hlen : 1111 - 15\*4 byte

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits		Flags 3 bits	Fragmentation offset 13 bits	
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				



# IP header

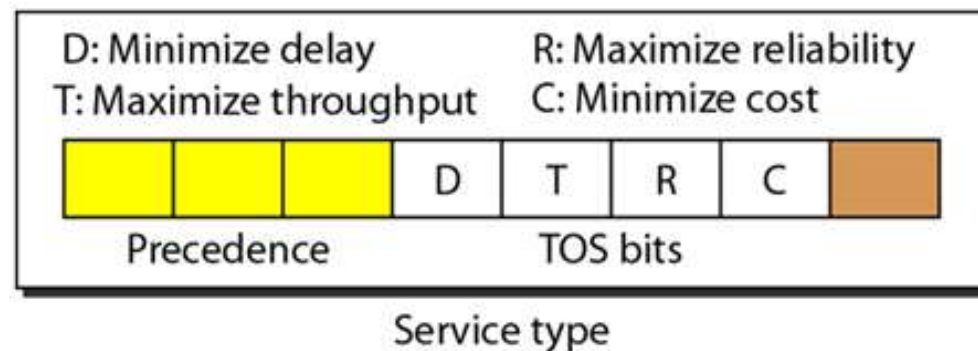
- Service: service class of the packet
  - The packet is processed at the router based on the service class

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

# Service class

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- Precedence: priority (0-7)
  - 7: the highest priority
  - When the router needs to drop a packet, it first drops the packet with the lowest priority
- TOS bits: Type of Service
  - None or only one of the bit can be 1



# Service class

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- Types of service

<i>TOS Bits</i>	<i>Description</i>
0000	Normal (default)
0001	Minimize cost
0010	Maximize reliability
0100	Maximize throughput
1000	Minimize delay

# Service class

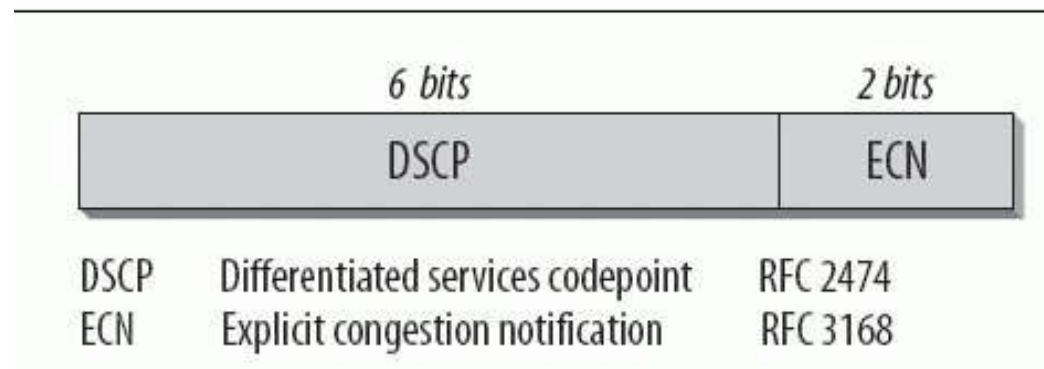
- Packet types and their service classes

<i>Protocol</i>	<i>TOS Bits</i>	<i>Description</i>
ICMP	0000	Normal
BOOTP	0000	Normal
NNTP	0001	Minimize cost
IGP	0010	Maximize reliability
SNMP	0010	Maximize reliability
TELNET	1000	Minimize delay
FTP (data)	0100	Maximize throughput
FTP (control)	1000	Minimize delay
TFTP	1000	Minimize delay
SMTP (command)	1000	Minimize delay
SMTP (data)	0100	Maximize throughput
DNS (UDP query)	1000	Minimize delay
DNS (TCP query)	0000	Normal
DNS (zone)	0100	Maximize throughput

# Service class (Revised)

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- The previous service class field is modified to indicate DSCP (Differentiated Service Code Point) and ECN (Explicit Congestion Notification)



# Service class (Revised)

- DSCP Table

Decimal	DSCP	Description
0	Default	Best Effort
8	CS1	Class 1 (CS1)
10	AF11	Class 1, Gold (AF11)
12	AF12	Class 1, Silver (AF12)
14	AF13	Class 1, Bronze (AF13)
16	CS2	Class 2 (CS2)
18	AF21	Class 2, Gold (AF21)
20	AF22	Class 2, Silver (AF22)
22	AF23	Class 2, Bronze (AF23)
24	CS3	Class 3 (CS3)
26	AF31	Class 3, Gold (AF31)
28	AF32	Class 3, Silver (AF32)
30	AF33	Class 3, Bronze (AF33)
32	CS4	Class 4 (CS4)
34	AF41	Class 4, Gold (AF41)
36	AF42	Class 4, Silver (AF42)
38	AF43	Class 4, Bronze (AF43)
40	CS5	Class 5 (CS5)
46	EF	Expedited Forwarding (EF)
48	CS6	Control (CS6)
56	CS7	Control (CS7)

# IP header

- Total length: the total packet length (header + data)
  - unit: byte

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

# IP header

- Identification, Flags, Fragmentation offset
- Related to fragmentation and assembly
  - Explained later

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				



# IP header

- Time to live: the life of the packet
  - unit: number of hops
  - decremented after going through each router
  - when TTL becomes 0, the packet is no longer forwarded
    - If it did not reach the destination, it is dropped

수명  
라우터를 거칠 때마다  
1씩 떨어짐  
범위 제한의 목적 or  
계속 라우터를 뱅뱅이 도는 걸 방지

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

# IP header

- Protocol: protocol of the upper layer (transport)

<i>Value</i>	<i>Protocol</i>
1	ICMP
2	IGMP
6	TCP
17	UDP
89	OSPF

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits	Protocol 8 bits		Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

# IP header

- Header checksum: used for error detection

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits	Protocol 8 bits		Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

# IP header

- Source IP address (32 bits)
- Destination IP address (32 bits)

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

# IP header

- Option: optional, may not be present
  - Other 20 bytes are mandatory fields
  - Thus, the minimum HLEN is 5
  - When option fields are present, HLEN is larger than 5

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits		4byte
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits	4byte
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits		4byte
Source IP address					4byte
Destination IP address					4byte
Option					

# IP header: exercise

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- The initial 8 bits of an IPv4 packet was the following

ip4                      4byte \* 2  
01000010

헤더는 최소 20 byte인데, 지금 HLEN 값은 2 -> 4byte라서 에러

- The router drops the packet since it is an erroneous packet. Why?

# IP header: exercise

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- HLEN value of an IPv4 packet is 1000 (in binary).  
What is the length of option fields in its IP header?

$(1000)_2 \times 4 \text{ byte} = 32 \text{ byte}$   
 $32 - 20 = 12 \text{ byte option}$

# IP header: exercise

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- HLEN of a packet is 5, and the total length field has the value 0x0028. What is the length of data in this IP packet?

20 byte

$2 \times 16 + 8 = 40\text{byte total}$   
 $40\text{byte} - 5 \times 4\text{byte} = 20\text{byte}$



# IP header: exercise

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- The initial part of an IPv4 packet is like the following:

0x45000028000100000102

4byte                      4byte                      2byte

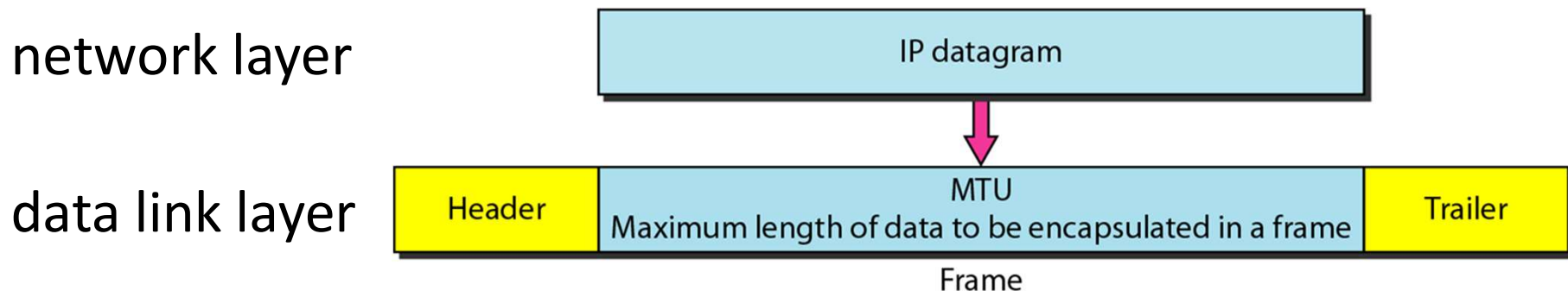
- What is the time-to-live of this packet?

1

# MTU (Maximum Transfer Unit)

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- Every data link layer protocol has an upper limit of IP packet



# MTU (Maximum Transfer Unit)

링크 레이어에서 한번에 보낼 수 있는 프레임 제한

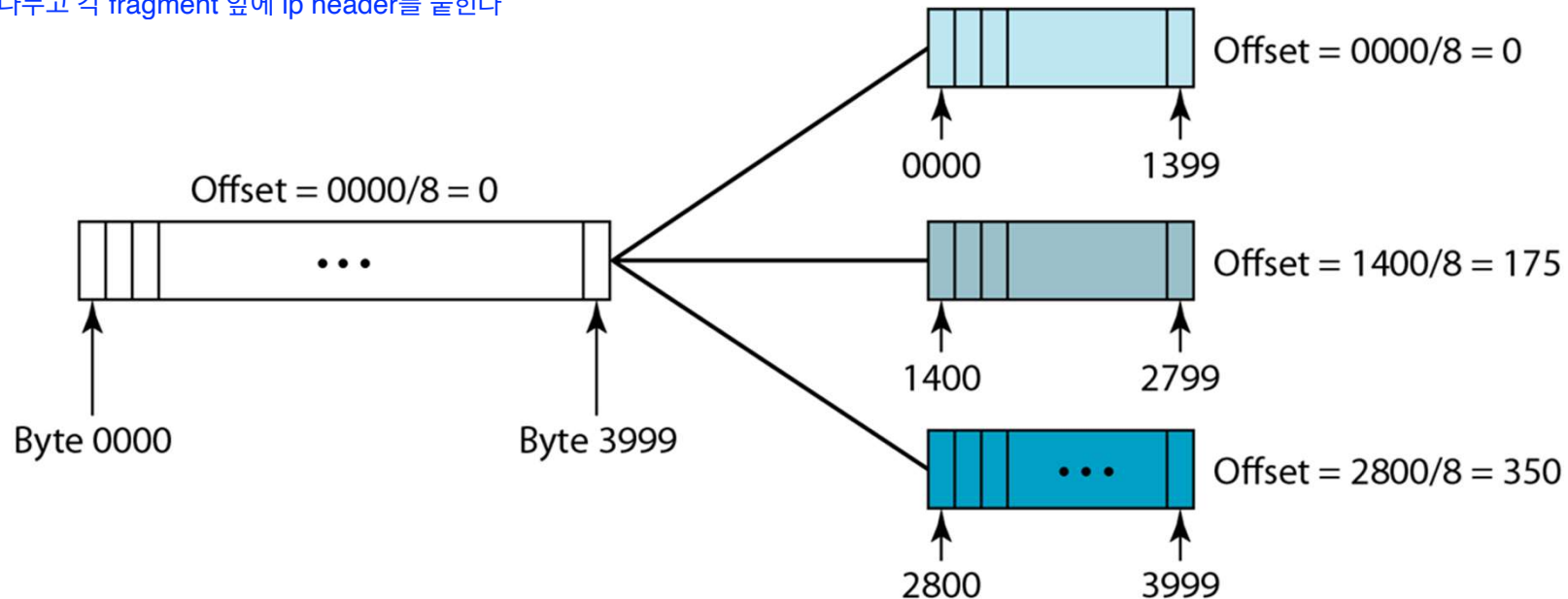
- Different protocols have different MTUs
  - MTU depends on forwarding and error detection methods

<i>Protocol</i>	<i>MTU</i>
Hyperchannel	65,535
Token Ring (16 Mbps)	17,914
Token Ring (4 Mbps)	4,464
FDDI	4,352
Ethernet	1,500
X.25	576
PPP	296

# Fragmentation

- If size of an IP packet is larger than MTU, then the packet needs to be fragmented into multiple packets
  - The IP header is attached to each fragment

fragment로 나누고 각 fragment 앞에 ip header를 붙인다



# Fields related to Fragmentation

- Identification: packet identifier
  - All fragments of the same packet has the same identification number

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
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Source IP address				
Destination IP address				
Option				

# Fields related to Fragmentation

- Flag

- If the bit 'D' is 1: this packet should not be fragmented

- If packet size is larger than MTU, packet is discarded
- D - fragment 못하게 함  
만약 D가 1인데 mtu 보다 크면 버림

- If the bit 'M' is 1: there are more fragments after this

1이면 뒤에 fragment가 더 있다  
0이면 뒤에 fragment가 없다

- If 'M' is 0, this is the last fragment, or this is the only fragment



VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits
Identification 16 bits		Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits
Source IP address			
Destination IP address			
Option			

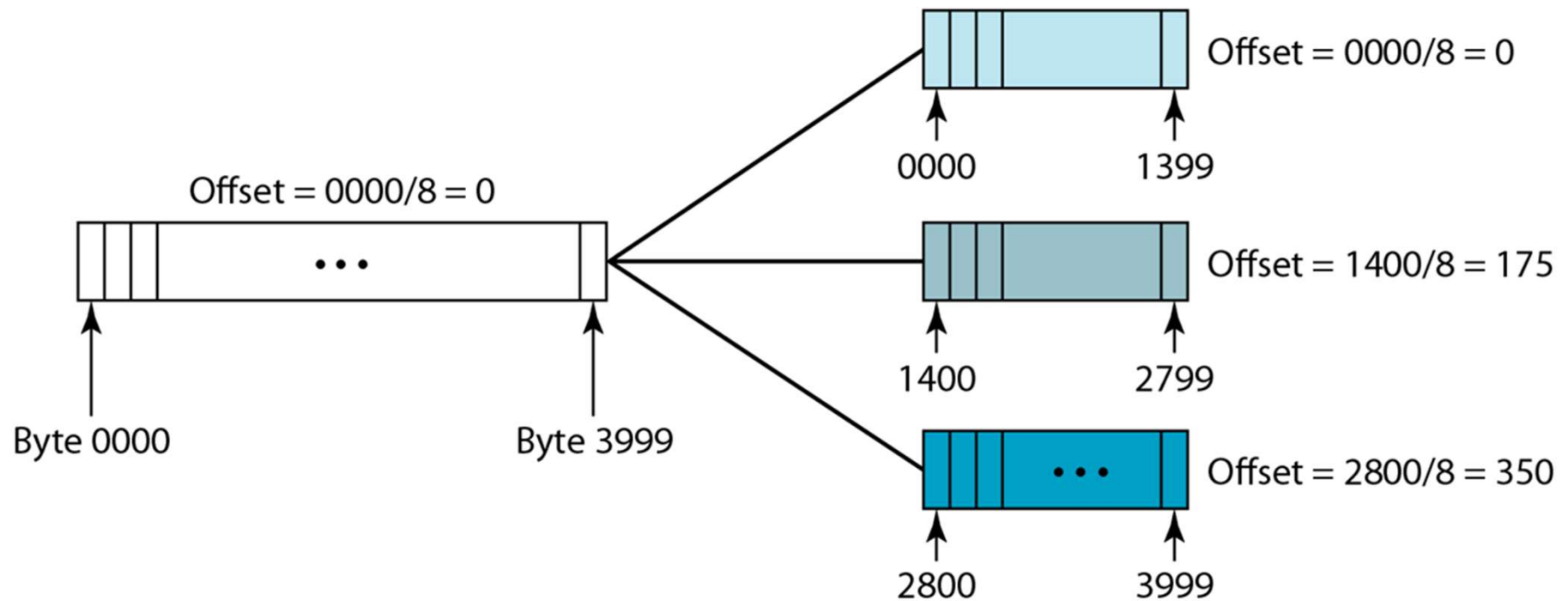
# Fields related to Fragmentation

- Fragmentation offset
  - The position of this fragment in the original packet
    - unit: 8 bytes
    - only considers data bytes (refer to the figures in the next slides)

VER 4 bits	HLEN 4 bits	Service 8 bits	Total length 16 bits	
Identification 16 bits			Flags 3 bits	Fragmentation offset 13 bits
Time to live 8 bits		Protocol 8 bits	Header checksum 16 bits	
Source IP address				
Destination IP address				
Option				

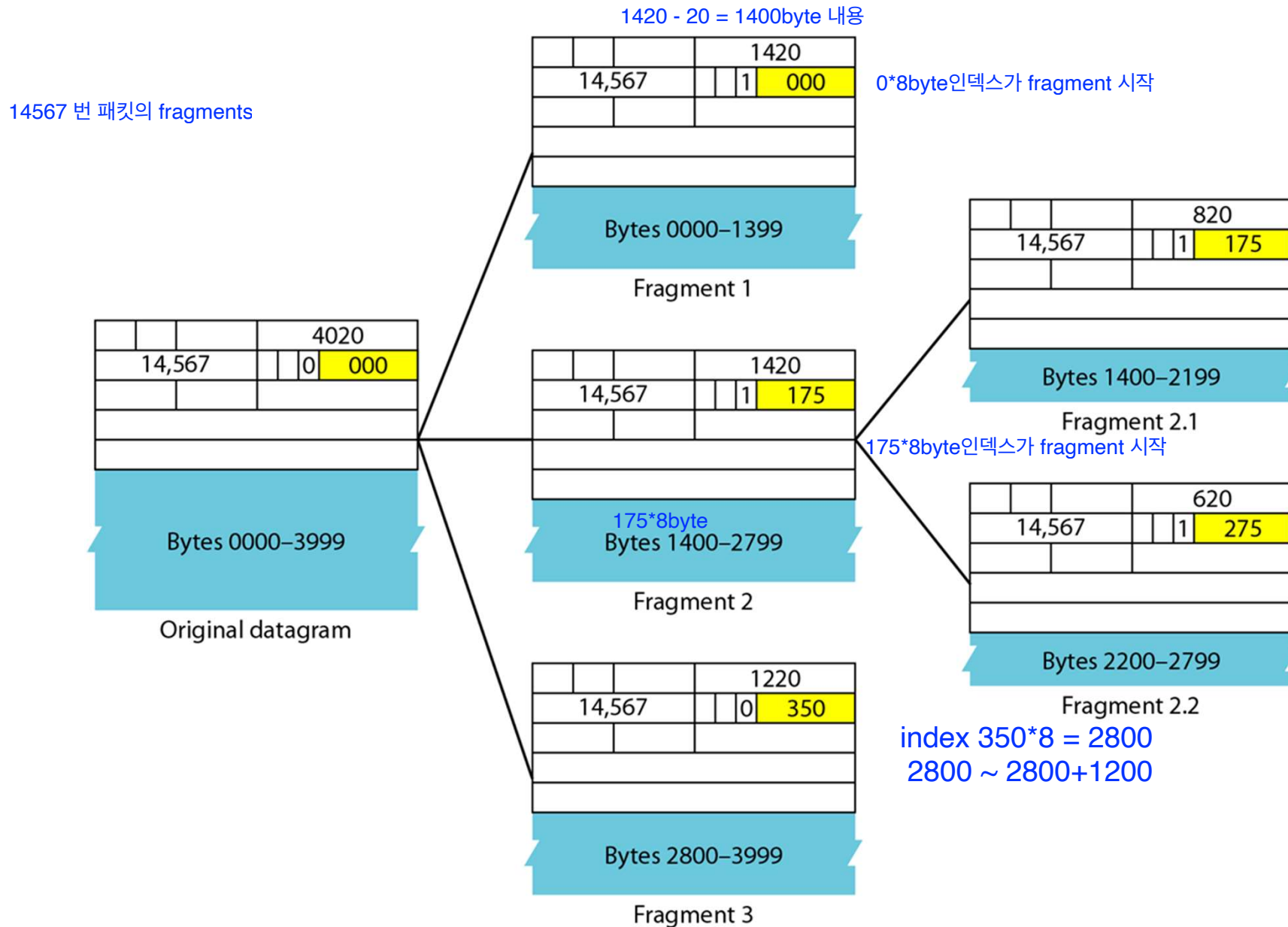
전체에서 이 fragment가 어느 위치인지  
8byte 단위

# Fragmentation offset



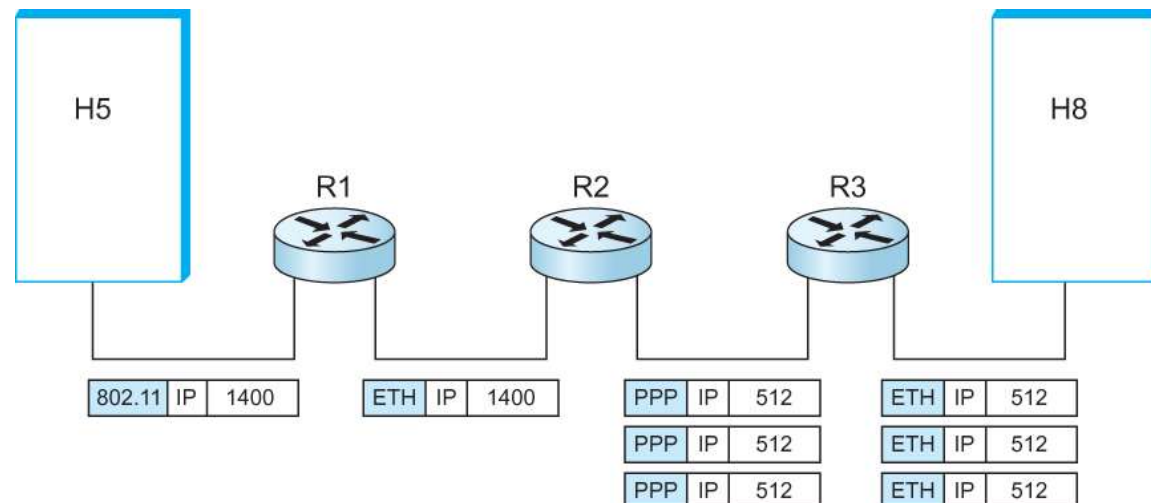


# IP fragmentation: example



# Fragmentation & reassembly

- Fragmentation can take place at any node
- Reassembly only occurs at the final destination
  - If the packet is reassembled at an intermediate node, it may need to be fragmented again
- The whole packet is discarded if one of the fragment does not arrive at the destination



# Fragmentation & reassembly: exercise

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- A packet has arrived, and its 'M' bit is 0. What is the meaning of this?
- Answer: either of the following
  - This packet is not fragmented
  - This packet is fragmented, and this fragment is the last one

# Fragmentation & reassembly: exercise

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- A packet has arrived, and its 'M' bit is 1. What is the meaning of this?
- Answer
  - This packet is fragmented, and this fragment is not the last one

# Fragmentation & reassembly: exercise

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- A packet has arrived, and its 'M' bit is 1, and the fragmentation offset is 0. What is the meaning of this?
- Answer
  - This packet is fragmented and this is the first fragment (there are more fragments after me)

# Fragmentation & reassembly: exercise

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- A packet has arrived, and its fragmentation offset is 100. What is the meaning of this?
- Answer
  - This packet is fragmented, and the position of this fragment is 800 bytes (only counting data bytes)

# Fragmentation & reassembly: exercise

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- A packet has arrived. Its fragmentation offset is 100, HLEN is 5, and total length is 100. What is the first and last position of this fragment in the original packet? (the first and the last bytes)
- the first byte:  $100 \times 8 = \text{byte \#800}$
- the last byte
  - total length: 100 bytes
  - header length:  $5 \times 4 = 20$  bytes
  - data length:  $100 - 20 = 80$  bytes
  - Thus, the last byte is the byte #879

# Fragmentation & reassembly: exercise

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- Assume IP header is always 20 bytes
- An IP packet has a size of 5140 bytes (including header)
- If data link layer MTU is 1500 bytes, then how should the packet be fragmented?



# Fragmentation & reassembly: exercise

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- The packet should be fragmented as follows.
  - MTU = 1500

