

# Physical optimization (select, sort, projection)

# Knowledge Objectives

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1. Explain how several indexes can be used to solve a complex selection predicate over one table
2. Enumerate three ways to have a table sorted
3. Enumerate seven operations that involve sorting in query processing
4. Explain the merge-sort algorithm
5. Explain how duplicates can be removed from a table depending on the data structure it has (i.e., plain file, B+, Hash, Clustered index)

# Understanding Objectives

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1. Find which indexes would be used to solve a multi-clause selection predicate over one table
2. Calculate the approximate cost of a merge-sort operation, given the available memory, the number of tuples in the table, and the number of tuples that fit in a disk block

# Selection algorithm

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1. Put the predicate in CNF
  1. Remove negations of parenthesis
    - $\text{NOT (A OR B)} = \text{NOT A AND NOT B}$
    - $\text{NOT (A AND B)} = \text{NOT A OR NOT B}$
  2. Move disjunctions into the parenthesis
    - $(\text{A AND B}) \text{ OR C} = (\text{A OR C}) \text{ AND } (\text{B OR C})$
    - $(\text{A AND B}) \text{ OR } (\text{C AND D}) = (\text{A OR C}) \text{ AND } (\text{A OR D}) \text{ AND } (\text{B OR C}) \text{ AND } (\text{B OR D})$
    - $(\text{A AND B}) \text{ OR } (\text{C AND B}) = (\text{A OR C}) \text{ AND B}$
2. Remove disjunctions if possible
  - For each parenthesis, if indexes can be used for all conditions in it, unite the RID lists resulting from accessing those indexes
3. Remove conjunctions if possible
  - For all parenthesis resolved in the previous step, intersect the RID lists produced
4. For each RID (if any) obtained from previous step, go to the table (by following the RID) to check the remaining predicates

# Considerations on the selection algorithm

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Let's suppose that we have a predicate in CNF with disjunctions:

$(A_1 \text{ op}_1 v_1 \text{ OR } A_2 \text{ op}_2 v_2) \text{ AND } \dots \text{ AND } (A_p \text{ op}_p v_p)$

being  $A_i$  attributes of the same relation and  $\text{op}_i$  a comparison operator

- If one of the conditions inside the parenthesis does not allow an index to be used, we cannot use any other index
- If no index can be used at all, we should perform a table scan
- Some kinds of indexes are not useful for some conditions
  - B+ useless for "different from"
  - Hash useless for "different from" and inequalities

# Example of selection

---

- We have tree indexes over attributes A, B and C
- We want to select those tuples in R that fulfill:

$(A=k_1 \text{ AND } B=k_2) \text{ OR } (C=k_3 \text{ AND } D=k_4)$

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$(A \text{ OR } C) \text{ AND } (A \text{ OR } D) \text{ AND } (B \text{ OR } C) \text{ AND } (B \text{ OR } D)$

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- We want to select those tuples in R that fulfill:

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U



# Example of selection

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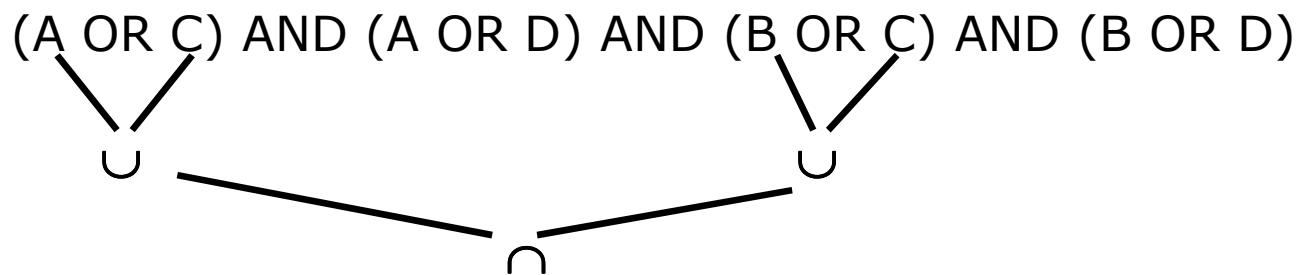
$(A \text{ OR } C) \text{ AND } (A \text{ OR } D) \text{ AND } (B \text{ OR } C) \text{ AND } (B \text{ OR } D)$

Below the expression, there are two 'U' symbols. The first 'U' is positioned under the sub-expression  $(A \text{ OR } C)$  and the second 'U' is positioned under the sub-expression  $(B \text{ OR } C)$ . Lines connect each 'U' to its corresponding sub-expression, indicating a join operation.

## Example of selection

- We have tree indexes over attributes A, B and C
- We want to select those tuples in R that fulfill:

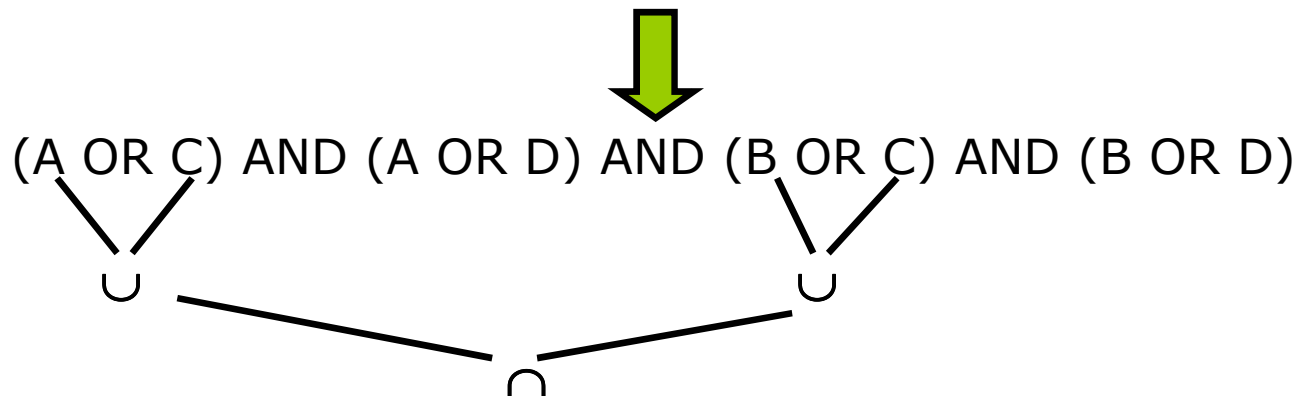
$(A=k_1 \text{ AND } B=k_2) \text{ OR } (C=k_3 \text{ AND } D=k_4)$



# Example of selection

- We have tree indexes over attributes A, B and C
- We want to select those tuples in R that fulfill:

$$(A=k_1 \text{ AND } B=k_2) \text{ OR } (C=k_3 \text{ AND } D=k_4)$$



For each entry in the list of RID resulting from the intersection  
Go to the table's file and Check "(A OR D) AND (B OR D)"

# External sorting algorithms

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- No index, with  $M+1$  memory pages
  - $2B \cdot \lceil \log_M B \rceil - B$
- B+
  - $\lceil |T|/u \rceil + |T|$
- Clustered
  - $\lceil 1.5B \rceil$
- Hash
  - Useless

$$u = \%carga \cdot 2d = (2/3) \cdot 2d$$

# External Merge Sort (I)

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## □ Function *sort()*

Assumption: Data is already in memory

Result: Writes the memory pages (blocks) sorted

Disk accesses:  $B_T$

## □ Function *scan(T)*

Assumption: We have  $B_T$  memory pages

Result:  $T$  has been read into memory

Disk accesses:  $B_T$

# External Merge Sort (II)

## □ Function $merge(T_1, \dots, T_M)$

Assumption:  $T_i$  are sorted and we have  $M+1$  memory pages

Result: Writes the sorted union of all  $T_i$

```

 $t_1 := first(T_1); t_2 := first(T_2); \dots t_M := first(T_M);$ 
while not (end( $T_1$ ) and end( $T_2$ ) and ... and end( $T_M$ ))
     $T^{ord} += t_{min};$ 
     $t_{min} := next(T_{min});$ 
endWhile
  
```

Disk accesses:  $2(B_{T_1} + B_{T_2} + \dots + B_{T_M})$

$min$  = index (1..M) of the  $T_i$  with the minimum current value

*first*  $\Rightarrow$  reads the first block into memory

*next*  $\Rightarrow$  reads a new block if the memory page is empty

$+=$   $\Rightarrow$  writes a block if buffer is full

# External Merge Sort (III)

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## □ Function *mergeSort(T)*

Assumption: We have  $M+1$  memory pages

Result: Sorted  $T$

```
if  $B_T \leq M$  then
  scan( $T$ );  $T^{\text{ord}} := \text{sort}()$ ;
else
   $T_1^{\text{ord}} := \text{mergeSort}(T_1)$ ; ...  $T_M^{\text{ord}} := \text{mergeSort}(T_M)$ ;
   $T^{\text{ord}} := \text{merge}(T_1^{\text{ord}}, \dots, T_M^{\text{ord}})$ ;
endif
```

# External Merge Sort (III)

## □ Function *mergeSort*(*T*)

Assumption: We have  $M+1$  memory pages

Result: Sorted *T*

```

if  $B_T \leq M$  then
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else
   $T_1^{ord} := \text{mergeSort}(T_1)$ ; ...  $T_M^{ord} := \text{mergeSort}(T_M)$ ;
   $T^{ord} := \text{merge}(T_1^{ord}, \dots, T_M^{ord})$ ;
endif

```

Disk accesses:  $2B_T \cdot \lceil \log_M B_T \rceil$

$$\begin{array}{ll}
 B \leq M \rightarrow 2B = & 1 * 2B \\
 M < B \leq M^2 \rightarrow 2B + 2B = & 2 * 2B \\
 M^2 < B \leq M^3 \rightarrow 4B + 2B = & 3 * 2B \\
 M^3 < B \leq M^4 \rightarrow 6B + 2B = & 4 * 2B
 \end{array}$$



# External Merge Sort (III)

## □ Function *mergeSort(T)*

Assumption: We have  $M+1$  memory pages

Result: Sorted  $T$

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if  $B_T \leq M$  then
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   $T^{\text{ord}} := \text{merge}(T_1^{\text{ord}}, \dots, T_M^{\text{ord}})$ ;
endif

```

Disk accesses:  $2B_T \cdot \lceil \log_M B_T \rceil$

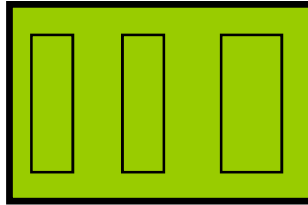
$B_T$  is the real number of blocks (taking into account possible cluster). For the sake of simplicity, we don't distinguish whether output has empty spaces or not

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 B \leq M \rightarrow 2B = & 1 * 2B \\
 M < B \leq M^2 \rightarrow 2B + 2B = & 2 * 2B \\
 M^2 < B \leq M^3 \rightarrow 4B + 2B = & 3 * 2B \\
 M^3 < B \leq M^4 \rightarrow 6B + 2B = & 4 * 2B
 \end{array}$$

# External Merge Sort (IV)

## Accesses

$M=2$



5	3	2	9
3	1	8	4

if  $B_T \leq M$  then  
 $\text{scan}(T); T^{\text{ord}} := \text{sort}();$

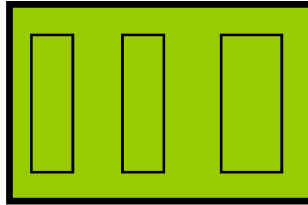
else

$T_1^{\text{ord}} := \text{mergeSort}(T_1);$   
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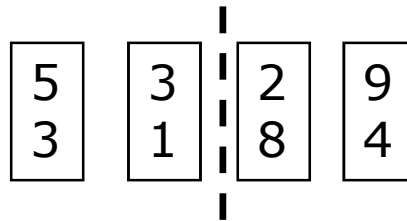
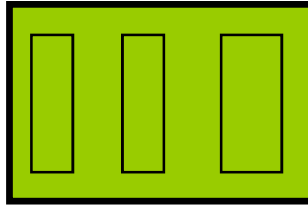
else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
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# External Merge Sort (IV)

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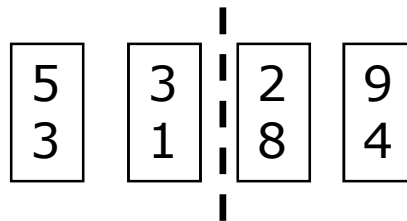
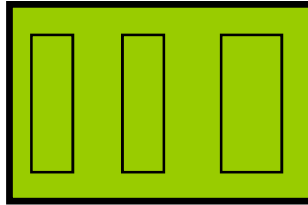
else

$T_1^{\text{ord}} := \text{mergeSort}(T_1);$   
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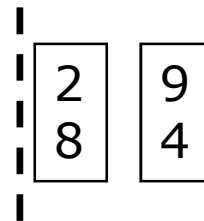
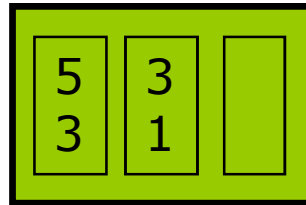
$T_1^{\text{ord}} := \text{mergeSort}(T_1);$   
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# External Merge Sort (IV)

## Accesses

R  
R

M=2



if  $B_T \leq M$  then  
`scan(T); Tord := sort();`

else

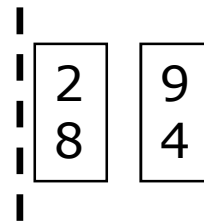
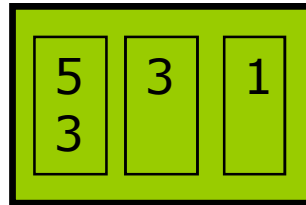
`T1ord := mergeSort(T1);`  
`T2ord := mergeSort(T2);`  
`Tord := merge(T1ord, T2ord);`

# External Merge Sort (IV)

## Accesses

R  
R

M=2



if  $B_T \leq M$  then  
 scan(T);  $T^{ord} := \text{sort}();$

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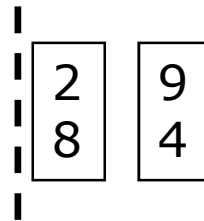
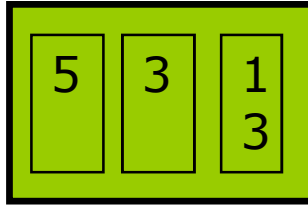
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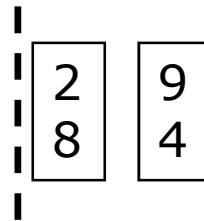
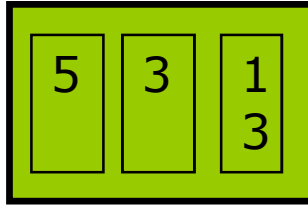


# External Merge Sort (IV)

## Accesses

R  
R  
W

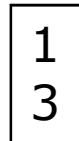
M=2



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 scan(T);  $T^{ord} := \text{sort}();$

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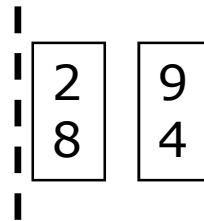
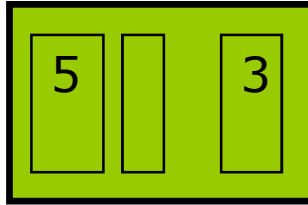


# External Merge Sort (IV)

## Accesses

R  
R  
W

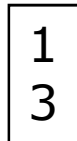
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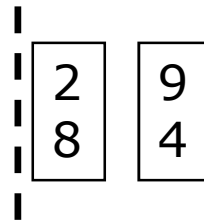
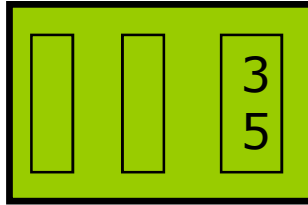


# External Merge Sort (IV)

## Accesses

R  
R  
W

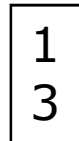
M=2



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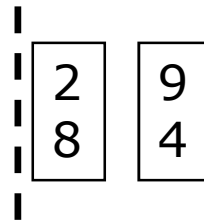
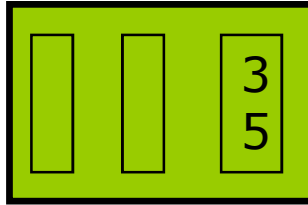


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W

M=2



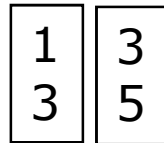
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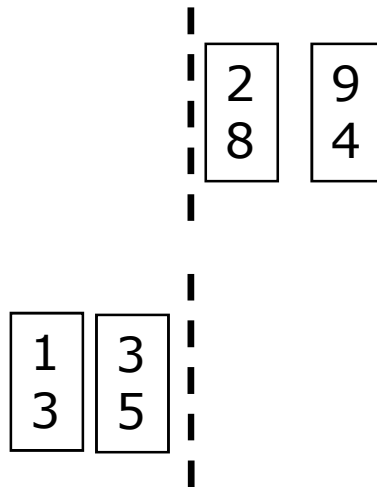
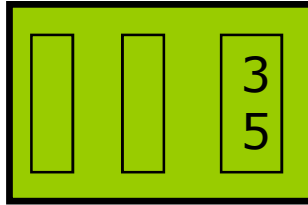


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W

M=2



if  $B_T \leq M$  then  
`scan(T); Tord := sort();`

else

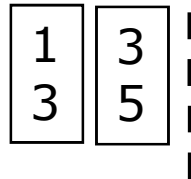
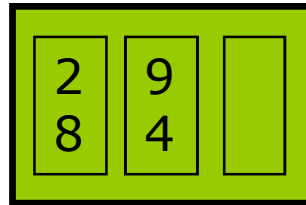
$T_1^{\text{ord}} := \text{mergeSort}(T_1);$   
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# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R

M=2



```
if  $B_T \leq M$  then
    scan(T);  $T^{ord} := \text{sort}()$ ;
```

```
else
```

```
     $T_1^{ord} := \text{mergeSort}(T_1)$ ;
```

```
     $T_2^{ord} := \text{mergeSort}(T_2)$ ;
```

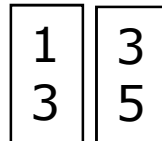
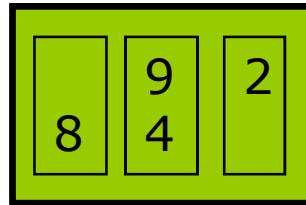
```
     $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;
```

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R

M=2



```

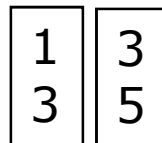
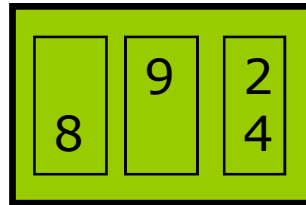
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```

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R

M=2



```

if  $B_T \leq M$  then
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```

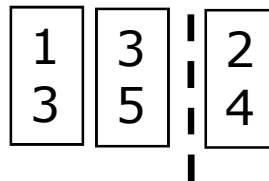
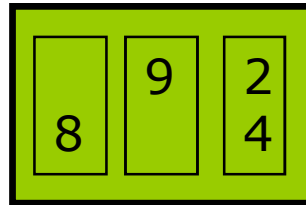


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W

M=2



```

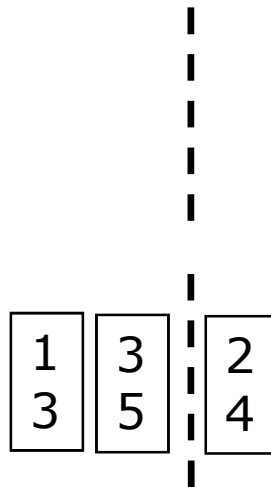
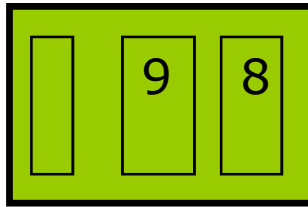
if  $B_T \leq M$  then
    scan(T);  $T^{ord} := \text{sort}();$ 
else
     $T_1^{ord} := \text{mergeSort}(T_1);$ 
     $T_2^{ord} := \text{mergeSort}(T_2);$ 
     $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord});$ 
  
```

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W

M=2



```

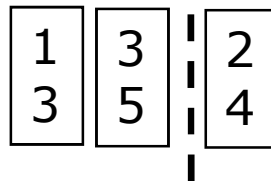
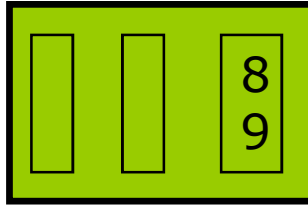
if  $B_T \leq M$  then
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else
     $T_1^{ord} := \text{mergeSort}(T_1)$ ;
     $T_2^{ord} := \text{mergeSort}(T_2)$ ;
     $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;
  
```

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W

M=2



```

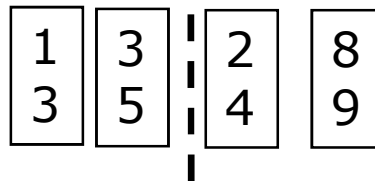
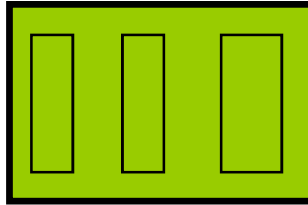
if  $B_T \leq M$  then
    scan(T);  $T^{ord} := \text{sort}()$ ;
else
     $T_1^{ord} := \text{mergeSort}(T_1)$ ;
     $T_2^{ord} := \text{mergeSort}(T_2)$ ;
     $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;
  
```

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W

M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

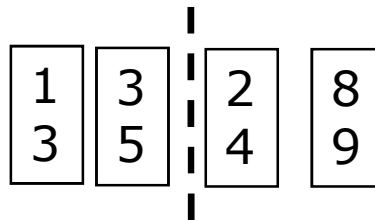
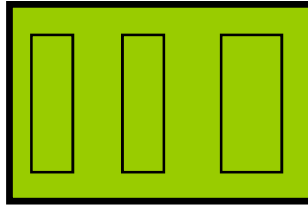
$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W

M=2



```
if  $B_T \leq M$  then
    scan(T);  $T^{ord} := \text{sort}()$ ;
```

```
else
```

```

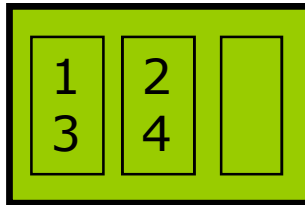
 $T_1^{ord} := \text{mergeSort}(T_1)$ ;
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;
```

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R

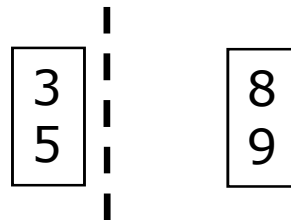
M=2



if  $B_T \leq M$  then  
     scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1);$   
 $T_2^{ord} := \text{mergeSort}(T_2);$   
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord});$

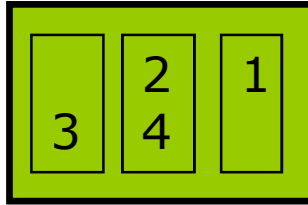


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R

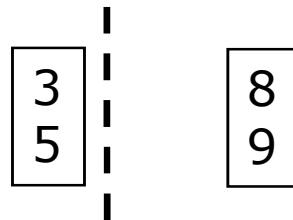
M=2



if  $B_T \leq M$  then  
     scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
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 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

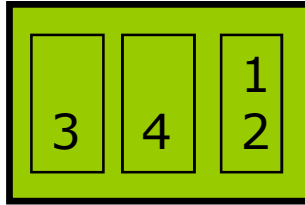


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R

M=2



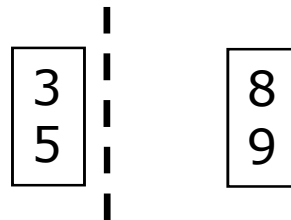
if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;

$T_2^{ord} := \text{mergeSort}(T_2)$ ;

$T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;



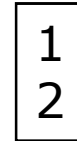
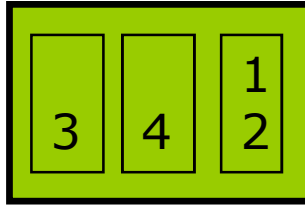


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R

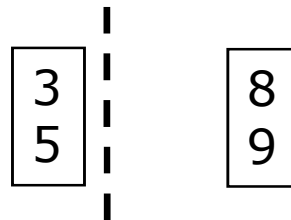
M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

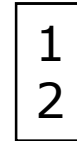
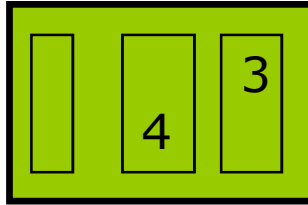


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R

M=2



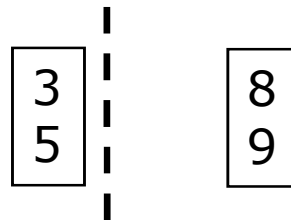
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    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;

$T_2^{ord} := \text{mergeSort}(T_2)$ ;

$T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

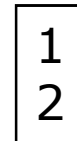
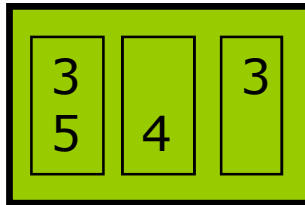


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R

M=2

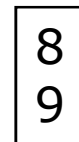


if  $B_T \leq M$  then  
     scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

|  
|  
|  
|  
|

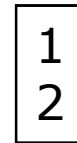
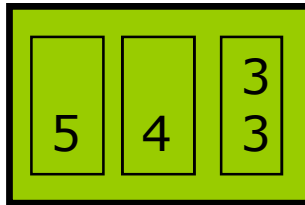


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R

M=2

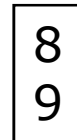


if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

---

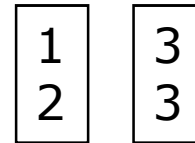
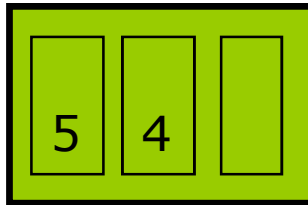


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R

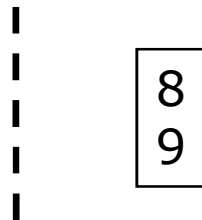
M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

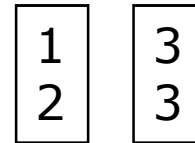
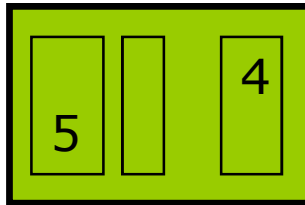


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R

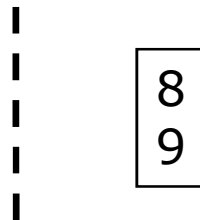
M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
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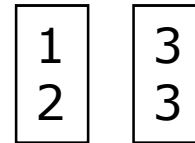
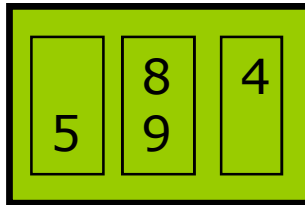


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



```
if  $B_T \leq M$  then
    scan(T);  $T^{ord} := \text{sort}()$ ;
```

```
else
```

```

 $T_1^{ord} := \text{mergeSort}(T_1)$ ;
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;
```

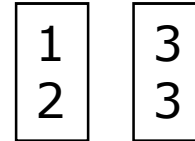
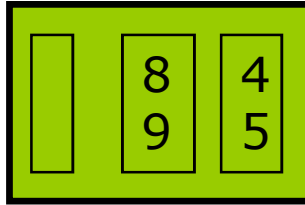
⋮

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



if  $B_T \leq M$  then  
     scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
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 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

⋮

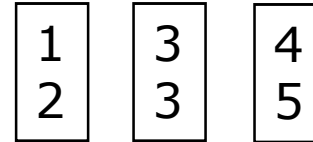
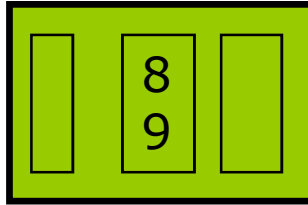


# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;

$T_2^{ord} := \text{mergeSort}(T_2)$ ;

$T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

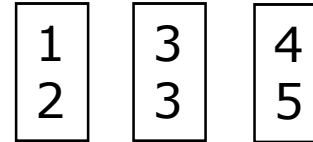
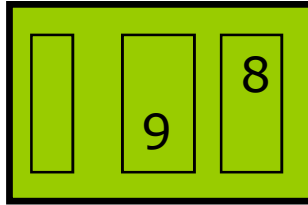
⋮

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

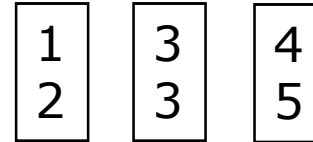
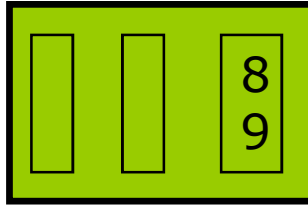
⋮

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
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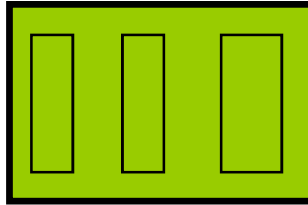
⋮

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



1	3	4	8
2	3	5	9

if  $B_T \leq M$  then  
    scan(T);  $T^{ord} := \text{sort}()$ ;

else

$T_1^{ord} := \text{mergeSort}(T_1)$ ;  
 $T_2^{ord} := \text{mergeSort}(T_2)$ ;  
 $T^{ord} := \text{merge}(T_1^{ord}, T_2^{ord})$ ;

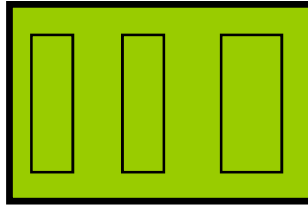
⋮

# External Merge Sort (IV)

## Accesses

R  
R  
W  
W  
R  
R  
W  
W  
R  
R  
R  
R

M=2



1	3	4	8
2	3	5	9

```

if BT ≤ M then
    scan(T); Tord := sort();
else
    T1ord := mergeSort(T1);
    T2ord := mergeSort(T2);
    Tord := merge(T1ord, T2ord);
  
```

# Usefulness of sorting algorithms

---

- ORDER BY
- Duplicates removal
  - DISTINCT
  - UNION
- GROUP BY
- Join
- Difference (anti-join)
- Massive index load

# Projection algorithms

---

## 1. Attribute removal

- a) There is another operation
  - 0
- b) There is no other operation
  - B

## 2. Duplicate removal

- a) No index, with  $M+1$  memory pages
  - $2B \cdot \lceil \log_M B \rceil - B$
- b) B+, useful if M and R are small with regard to B
  - $\lceil |T|/u \rceil$  (probably  $+|T|$ )
- c) Clustered
  - $\lceil 1.5B \rceil$
- d) Hash, useful if M and R are small with regard to B
  - $\lceil 1.25(|T|/2d) \rceil$  (probably  $+|T|$ )

$$u = \%load \cdot 2d = (2/3) \cdot 2d$$

# Summary

---

- Selection algorithms
- External sort algorithms
- Projection algorithms



# Bibliography

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