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# Relational translation - 1

# Knowledge objectives

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1. Name the two meanings of the NULL value

# Understanding Objectives

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1. Explain the storage space problem generated by NULL values, and how DBMSs solve it
2. Explain the access time problem generated by NULL values, and how we should take it into account in the DB design
3. Explain the general way to translate classes and associations into relational model

# Application Objectives

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1. Translate a DB query from natural language into SQL, given a DB schema, taking into account the presence of NULL values

# Which rows do we get?

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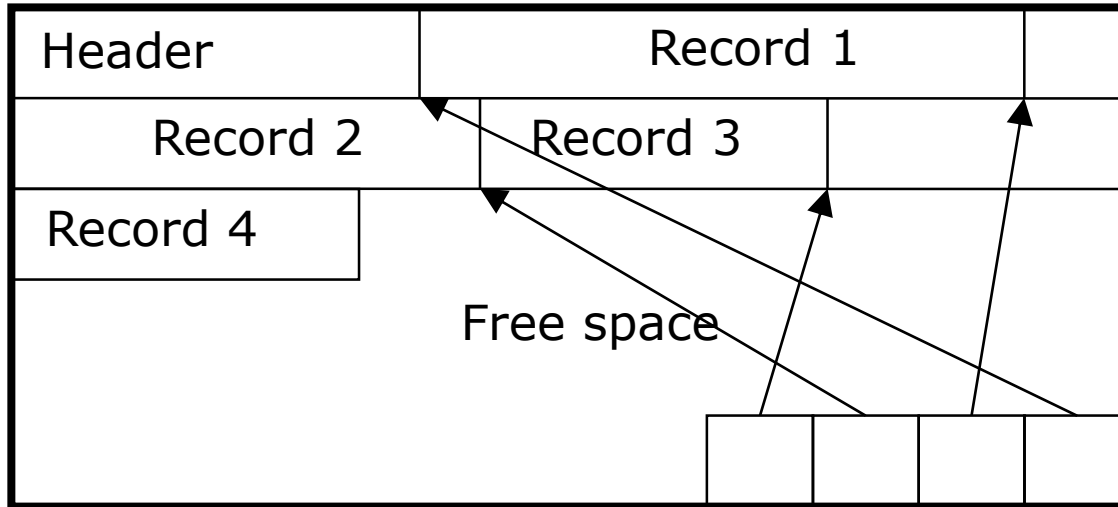
```
SELECT *  
FROM R  
WHERE A=10 OR A<>10;
```

# Null values

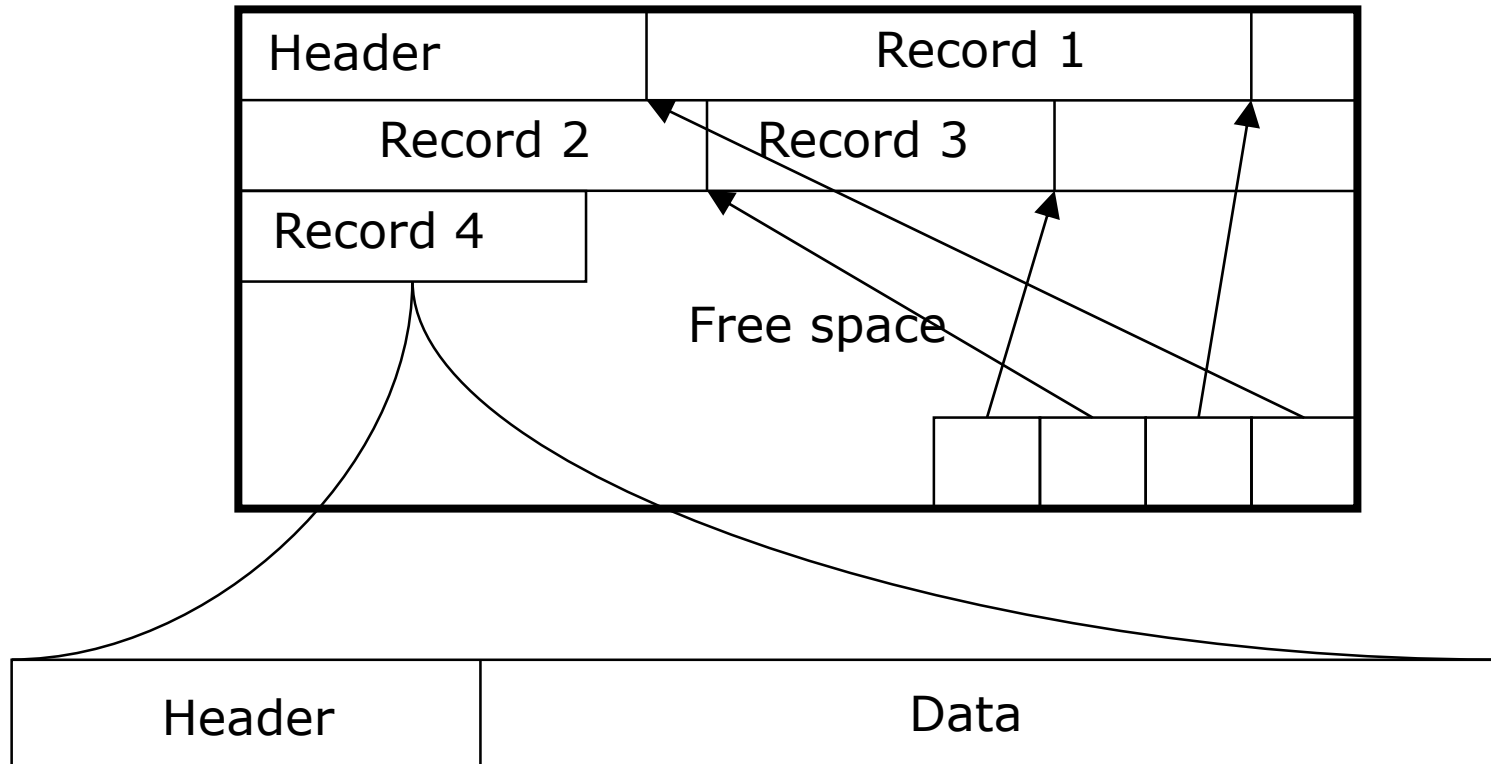
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- ❑ Two meanings
- ❑ Reasons to use them:
  - Inserting a tuple with an unknown value
  - Adding a new attribute to a non-empty relation
  - Special aggregation cases
  - Avoiding exceptions in aggregations with unknown values
- ❑ Representation:
  - Different from any non-null value

# Null values storage

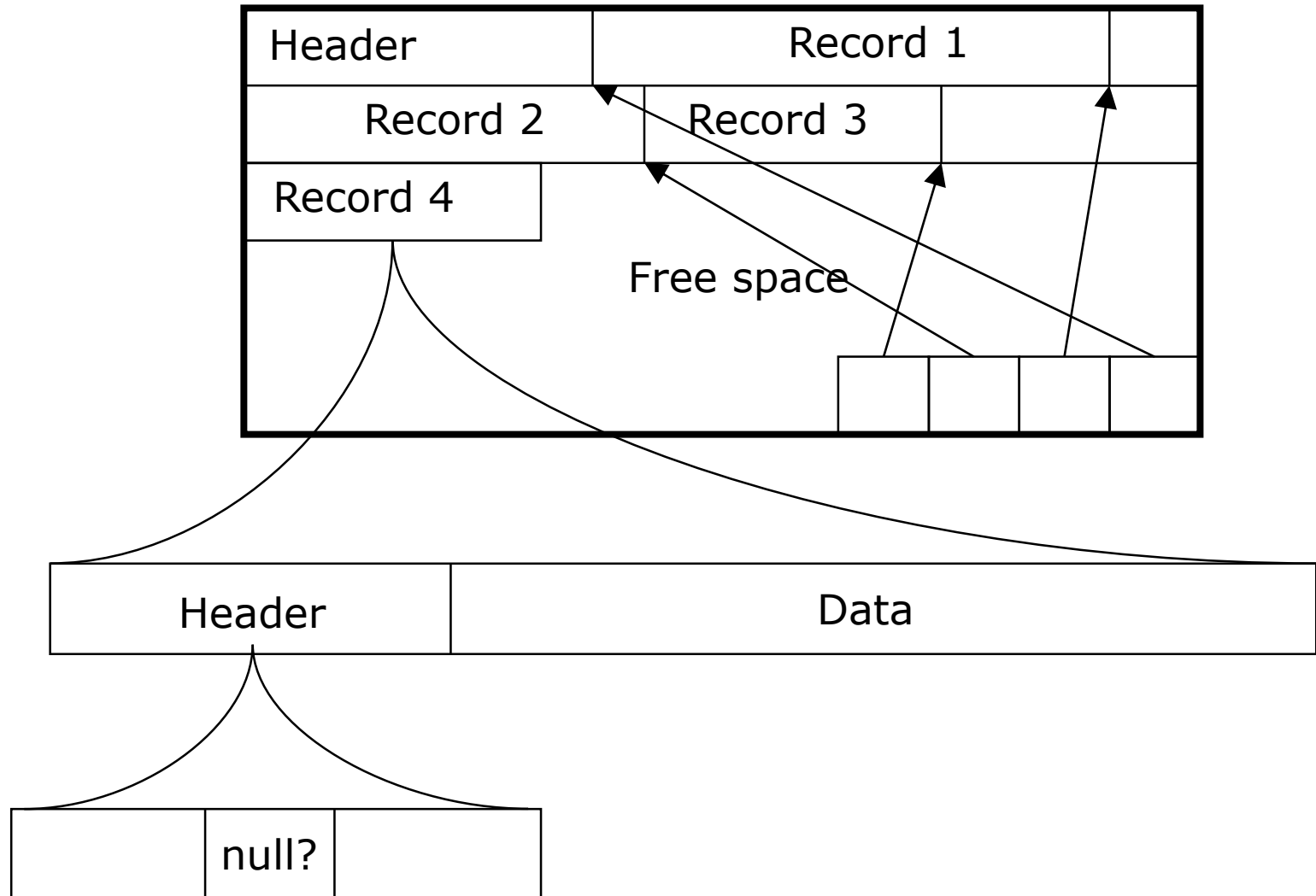


# Null values storage





# Null values storage



# Null values query

a) SELECT name  
FROM People  
WHERE #labour>0;

People


b) SELECT name  
FROM Women  
WHERE #labour>0;

Men


Women


# Ternary logic for null values

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NULL  $\theta$  X  $\rightarrow$  UNKNOWN

NULL=NULL  $\rightarrow$  UNKNOWN

# Ternary logic for null values

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NULL  $\theta$  X  $\rightarrow$  UNKNOWN  
NULL=NULL  $\rightarrow$  UNKNOWN

NOT	
T	F
U	U
F	T

# Ternary logic for null values

NULL  $\theta$  X  $\rightarrow$  UNKNOWN  
NULL=NULL  $\rightarrow$  UNKNOWN

NOT	
T	F
U	U
F	T

AND	T	U	F
T	T	U	F
U	U	U	F
F	F	F	F

# Ternary logic for null values

$\text{NULL} \neq \text{X} \rightarrow \text{UNKNOWN}$   
 $\text{NULL} = \text{NULL} \rightarrow \text{UNKNOWN}$

<b>NOT</b>	
T	F
U	U
F	T

<b>AND</b>	T	U	F
T	T	U	F
U	U	U	F
F	F	F	F

<b>OR</b>	T	U	F
T	T	T	T
U	T	U	U
F	T	U	F

# Consequences of ternary logic

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## □ Queries

- Return rows when the predicate is true
- Do not return those evaluating unknown

## □ Constraints

- Raise an exception when the predicate is false
- Do not raise anything when evaluates unknown

# Effect of null values in aggregates

- COUNT
  - With "\*", counts all tuples
  - With "a", counts those with non-null value for the attribute
- SUM
  - Adds only non-null values
  - Returns null if there are not non-null values
- MIN/MAX
  - Returns null if there are not non-null values
- AVG
  - Its result always coincide with  $SUM(a)/COUNT(a)$

Content	SUM(a)	COUNT(*)	COUNT(a)	AVG(a)	SUM(a) / COUNT(a)	SUM(a) / COUNT(*)	MIN(a)
empty	null	0	0	null	null	null	null
null	null	1	0	null	null	null	null
null 0 1	1	3	2	0.5	0.5	0.33333	0



# Effect of null values in usual other cases

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- $V \text{ IN } (X_1, X_2, \dots)$  is the same as  $V = X_1 \text{ OR } V = X_2 \text{ OR } \dots$ 
  - $\text{NULL IN } (X_1, X_2, \dots) \rightarrow \text{UNKNOWN}$
- GROUP BY
  - Null values are put onto a single group
- UNIQUE
  - Unique constraints treat null as different to everything (different to null and different to other values)
  - This makes every null a different null so that unique constraints accept multiple null values

# Specific comparison for nulls

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SELECT id FROM T WHERE a IS NULL;

VS

SELECT id FROM T WHERE a=NULL;

# Specific comparison for nulls

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SELECT id FROM T WHERE a IS NULL;

VS

~~SELECT id FROM T WHERE a=NULL;~~



# Query examples with nulls

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*teachers living in a city where no student lives*

```
SELECT id  
FROM teachers  
WHERE city NOT IN (SELECT city FROM students);
```

VS

```
SELECT id  
FROM teachers  
WHERE NOT EXISTS (  
    SELECT *  
    FROM students  
    WHERE teachers.city =  
           students.city);
```

# Algebraic operations with nulls

R(	A,	B)
	?	?
	a	?
	a	1
	?	1

S(	A,	B)
	?	?
	a	?
	a	1

T(	A,	B,	C)
	?	?	z
	a	?	y
	?	1	x
	a	?	w
	?	1	v

$R \cup S$	(A,	B)
	?	?
	a	?
	a	1
	?	1

$R \cap S$	(A,	B)
	?	?
	a	?
	a	1

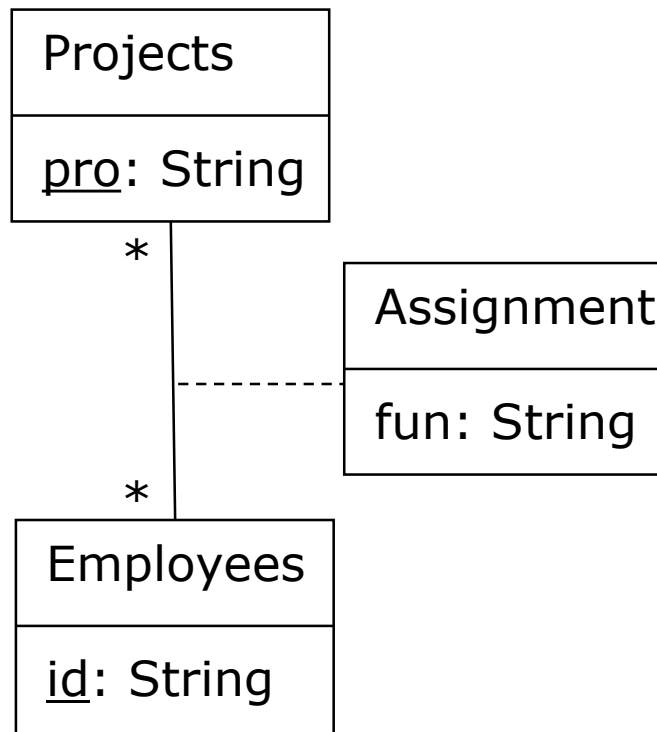
$R - S$	(A,	B)
	?	1

$T[A,B]$	(A,	B)
	?	?
	a	?
	?	1

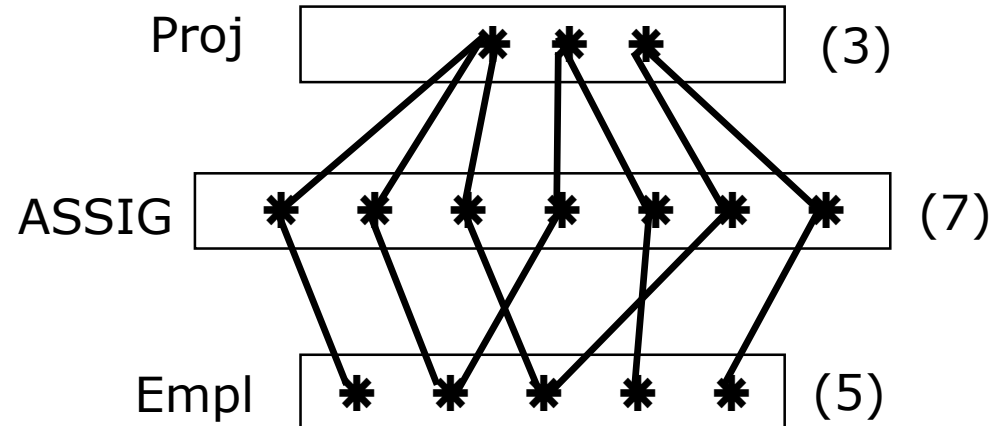
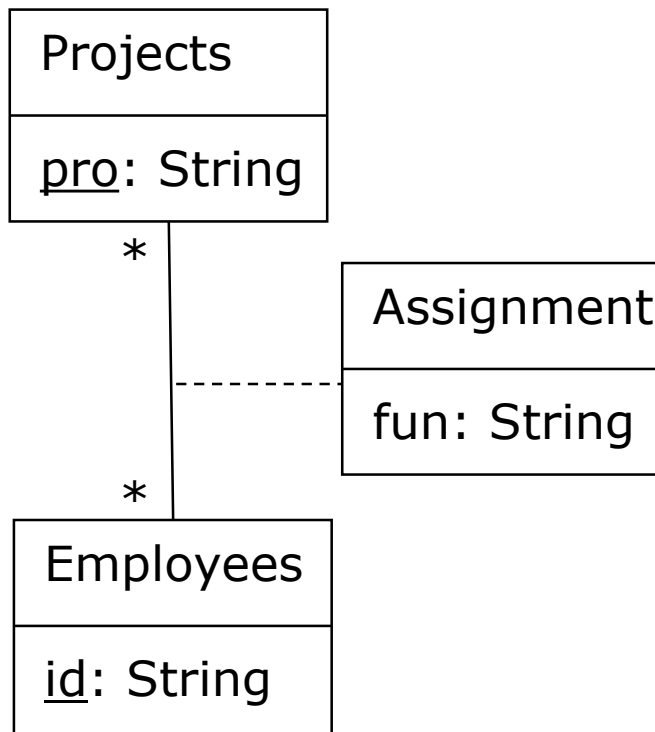
$T[A]$	(A)
	?
	a

$R(B = B)T$	(A,	B,	A',	B',	C)
	a	1	?	1	x
	?	1	?	1	x
	a	1	?	1	v
	?	1	?	1	v

# UML model and instantiation



# UML model and instantiation



# Relational model representation

## □ Long

```
CREATE TABLE Projects (pro CHAR(25), ...);  
CREATE TABLE Employees (id CHAR(9), ...);  
CREATE TABLE Assignments (...);
```

## □ Short

Projects(pro, ...)

Assignments(empl, pro, function)

Employees(id, ...)

ONLY CORRECT  
INSTANTIATIONS OF  
UML DIAGRAM MUST  
BE POSSIBLE TO  
REPRESENT IN THE  
RELATIONAL SCHEMA

**WITHOUT OID!!!**



# Summary

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- Meaning of NULL values
- Consequences of NULL values
  - Space
  - Time
  - Writing queries
- General way to translate classes and associations

# Bibliography

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- ❑ Jaume Sistac et al. *Disseny de bases de dades*. Editorial UOC, 2002. Col·lecció Manuals, number 43
- ❑ J. Melton and A. Simon. *SQL 1999*. Morgan Kaufmann, 2002
- ❑ P. Gultzan and T. Pelzer. *SQL-99 Complete, really*. R&D Books, 1999