

# EECS 368

## Programming Language Paradigms

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Is this a valid scheme command?

> (1 + (2 \* 3))

Is this a valid scheme command?

> (1 + (2 \* 3))

No.

expected procedure, given: 2; arguments were: #<procedure:\*> 3

Is this a valid scheme command?

> ' (1 + (2 \* 3))

Is this a valid scheme command?

> ' (1 + (2 \* 3))

Yes! - it is a quoted S-Expression

(1 + (2 \* 3))

This looks like an arithmetical expression, but is actually Scheme data.

We want to write a scheme  
function, value

```
> (value '(1 + (2 * 3)))
```

7

```
> (value '(2 * 3))
```

6

value takes a (quoted) expression,  
and returns a number.

Assumptions: the input is well formed

$$\begin{aligned} E ::= & \text{ number} \\ & | ( E + E ) \\ & | ( E * E ) \end{aligned}$$

What questions do we need to ask?

# What questions do we need to ask?

- Are you at atom?  
(must be a number!)
- Are you an addition?
- Are you a multiplication?



```
(define atom?  
  (lambda (x)  
    (and (not (pair? x)) (not (null? x)))))
```

```
(define value  
  (lambda (nexp)  
    (cond  
      ((atom? nexp) nexp)  
      ((...) ...)   
      ((...) ...)   
    ) ... )
```

```
> (value '4)
```

```
4
```

```
> (value '(2 + 3))
```

```
...
```

We can test for  $(E + E)$  using `'+`

```
(eq? (car (cdr nexp)) '+)
```

```
(define value
  (lambda (nexp)
    (cond
      ((atom? nexp) nexp)
      ((eq? (car (cdr nexp)) '+)
       (+ (value (car nexp))
           (value (car (cdr (cdr nexp))))))
      ((eq? (car (cdr nexp)) '*)
       (* (value (car nexp))
           (value (car (cdr (cdr nexp))))))
    )
  )
)
```

```
> (value '(4 + 3))
```

```
7
```

# We have taken a quoted S-Expression, and given it meaning!

```
> (value '((1 + 2) * (4 + 3)))
```

```
21
```

```
> (value '4)
```

```
4
```

```
> (value '(4))
```

```
car: contract violation
```

```
  expected: pair?
```

```
  given: '()
```

How do we fix the last example?

# Homework!

Write a function which multiplies the height from the root to each element in an s-expression.

```
> (timesdepth* 1 '(1 2 3))  
(1 2 3)  
  
> (timesdepth* 1 '((1) (2) (3)))  
((2) (4) (6))  
  
> (timesdepth* 1 '((1) 2 ((3))))  
((2) 2 ((9)))
```

Test out your function with Racket. If you are stuck, write `timesdepth`, that works on a list of atoms. You can assume that the atoms are numbers.

# Lambda the Ultimate

# eq? and equal? and =

			eq?	equal?	=
'a	'a	⇒	#t	#t	error
'(a)	'(a)	⇒	#f	#t	error
'(a b)	(a b)	⇒	#f	#t	error
'(b a)	'(a b)	⇒	#f	#f	error
'(a (b) c)	'(a (b) c)	⇒	#f	#t	error
1	2	⇒	#f	#f	#f
1.0	1.0	⇒	#t	#t	#t
1	1	⇒	#t	#t	#t
100000	100000	⇒	#t	#t	#t
100000000000	100000000000	⇒	#t	#t	#t

# Keyword-like things and built-in things in scheme

```
define lambda cond else
```

```
#t #f
```

```
if let begin set!
```



# Function-like things in scheme

eq? equal? null? zero? cons?

+ - \* /

= < > <= >=

and or not

cons car cdr list

display

```
(define findindex  
  (lambda (a l)  
    (findindex2 a l 0)))
```

```
(define findindex2
  (lambda (a l ix)
    (cond
      ((null? l) -1)
      (else
       (cond
         ((= a (car l)) ix)
         (else (findindex2 a (cdr l) (+ 1 ix))))))))
```

## passing functions as arguments

```
(define findindex2
  (lambda (cmp a l ix)
    (cond
      ((null? l) -1)
      (else
       (cond
         ((cmp a (car l)) ix)
         (else (findindex2 cmp a (cdr l) (+ 1 ix)))
       )))))
```

```
(define findindex-using
  (lambda (cmp a l)
    (findindex2 cmp a l 0)))
```

# map

```
(define double  
  (lambda (a) (+ a a)))
```

```
> (map double '(1 2 3 4))
```

```
(define double  
  (lambda (a) (+ a a)))
```

```
> (map double '(1 2 3 4))
```

```
(2 4 6 8)
```

# map

```
(define double  
  (lambda (a) (+ a a)))
```

```
> (map double '(1 2 3 4))
```

```
(2 4 6 8)
```

```
> (map  
  (lambda (a) (* a a))  
  '(1 2 3 4))
```

```
(define double  
  (lambda (a) (+ a a)))
```

```
> (map double '(1 2 3 4))
```

```
(2 4 6 8)
```

```
> (map  
  (lambda (a) (* a a))  
  '(1 2 3 4))
```

```
(1 4 9 16)
```



# foldr

```
(define foldr
  (lambda (f z l)
    (cond
      ((null? l) z)
      (else (f (car l) (foldr f z (cdr l)))))))
```

```
> (foldr + 0 '(1 2 3 4 5))
```

# foldr

```
(define foldr
  (lambda (f z l)
    (cond
      ((null? l) z)
      (else (f (car l) (foldr f z (cdr l)))))))
```

```
> (foldr + 0 '(1 2 3 4 5))
```

15

# foldr

```
(define foldr
  (lambda (f z l)
    (cond
      ((null? l) z)
      (else (f (car l) (foldr f z (cdr l)))))))
```

```
> (foldr + 0 '(1 2 3 4 5))
```

15

```
> (foldr cons '() '(1 2 3 4 5))
```

# foldr

```
(define foldr
  (lambda (f z l)
    (cond
      ((null? l) z)
      (else (f (car l) (foldr f z (cdr l)))))))
```

> (foldr + 0 '(1 2 3 4 5))

15

> (foldr cons '() '(1 2 3 4 5))

(1 2 3 4 5)

# Real World Scheme

So far, we can evaluate small functions, but do no real work.

In this class we will step back, and see some other “features” of scheme.

- Special Forms – (cond ...) (define ...) etc.
- Assignment – (set! ...)
- Scoping Rules
- Input and Output

We have met several special forms already

```
(define xyz ...)
```

```
(lambda (a b c) ...)
```

```
(define add1  
  (lambda (x) (+ x 1)))
```

```
(define double (lambda (a) (* a 2)))  
(double 4)
```



```
(define double (lambda (a) (* a 2)))  
(double 4)
```

```
(      (lambda (a) (* a 2))  4  )
```

```
(define double (lambda (a) (* a 2)))  
(double 4)
```

```
(      (lambda (a) (* a 2))  4  )
```

```
(                * 4 2          )
```

```
(let ((v 4)
      (w 9))
      (* v w))
```

```
(let ( (double (lambda (a) (* a 2)))  
      (v 4)  
    )  
  (double v)  
)
```

```
( (lambda (x) (* x x)) 9 )
```

```
(let ( (x 9)
      )
    (* x x)
)
```

```
(define x 1)
```

```
(+ x 1)           ;; ==> 2
```

```
(set! x 4)
```

```
(+ x 1)           ;; ==> 5
```

define creates a binding for x, while set changes it.

```
(define v ..)
```

has essentially the same effect as this assignment expression, if variable is bound:

```
(set! v ...)
```

but is about *namespace*, not assignment!

```
(define x 1)
(define f
  (lambda (y)
    (begin
      (set! x y)
      (+ y y))))
(f 99)
```



```
(define x 1)
(define g (lambda (y) (+ x y)))
(define f (lambda (x) (g 2)))
(f 5)
```

```
(define x 1)
(cond ((= x 1) 99)
      ((= x 9) 100)
      (else 121)
)
```

```
(define x 1)  
(if (= x 22) 99 100)
```

```
> (display "Hello, World")  
Hello, World  
> (write "Hello, World")  
"Hello, World"  
>
```

```
> (set! x (read))
```

```
Hello
```

```
>
```

This is actually reading a scheme s-expression!

```
> (let ((x ""))  
    (begin (set! x (read))  
            (write x))  
    )  
Hello World
```

```
> (let ((x ""))  
    (begin (set! x (read))  
            (write x))  
    )
```

Hello World

Hello