

# EECS 368

## Programming Language Paradigms

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# Syntax for Languages

## C++

```
#include <iostream.h>
main()
{
    cout << "Hello World!";
    return 0;
}
```

## Java

```
public class HelloWorld {
    public static void main (String[] args) {
        System.out.println ("Hello World!");
    }
}
```

## JavaScript

```
console.log("Hello, World");
```

## Scheme

```
(display "Hello World!")
(newline)
```

## Haskell

```
main :: IO ()
main = putStrLn "Hello World!"
```

# Putting it together

```
java ::= public class name { methods }
```

# Putting it together

```
java ::= public class name { methods }
```

```
methods ::= method methods  
          |
```

# Putting it together

```
java ::= public class name { methods }
```

```
methods ::= method methods  
          |
```

```
method ::= public static void name ( method_args ) { statement ; }
```

# Putting it together

```
java ::= public class name { methods }
```

```
methods ::= method methods  
          |
```

```
method ::= public static void name ( method_args ) { statement ; }
```

```
method_args ::= non_empty_method_args  
              |
```

```
non_empty_method_args ::= method_arg  
                        | method_arg , non_empty_method_args
```

# Putting it together

`java ::= public class name { methods }`

`methods ::= method methods`  
`|`

`method ::= public static void name ( method_args ) { statement ; }`

`method_args ::= non_empty_method_args`  
`|`

`non_empty_method_args ::= method_arg`  
`| method_arg , non_empty_method_args`

`method_arg ::= ...`

**BNF is a way of expressing valid programs.**

**BNF can express arbitrary sized sequences.**

**Real languages can have large BNF's.**



# Examples of Comma-Separated Sequences

## One or More

```
seq ::= number  
      | seq , seq
```

---

```
seq ::= number  
      | number , seq
```

---

```
seq ::= number  
      | seq , number
```

## Zero or More

```
seq ::= m_seq  
      |  
m_seq ::= number  
         | m_seq , m_seq
```

---

```
seq ::= m_seq  
      |  
m_seq ::= number  
         | number , m_seq
```

---

```
seq ::= m_seq  
      |  
m_seq ::= number  
         | m_seq , number
```

# Another Comma-Separated Sequence

There are other ways of representing sequences

In the last class (right leaning)

```
seq ::= number  
      | number , seq
```

An alternative right leaning

```
seq ::= numbers number  
  
numbers ::= number , numbers  
          |
```

What is the common pattern? How many more are there?

We should be able to think in terms of “*What does this BNF accept?*”

There are some understood extensions, called EBNF.

One or more of something

$$X^+$$

Zero or more of something

$$X^*$$

Optional elements

$$X^?$$

Grouping

$$( \dots )$$

# ? Optional

Optional elements

**X?**

**X** can be a terminal, non-terminal, or a parenthesized sequence of terminals and non-terminals.

Sometimes written [ **X** ]

BNF

```
abc ::= xyz vwx www  
      | xyz www
```

EBNF

```
abc ::= xyz vwx? www
```

# ? More Optional

## BNF

```
seq ::= m_seq  
      |
```

## EBNF

```
seq := m_seq?
```

## BNF

```
seq ::= m_seq  
      |  
m_seq ::= number  
        | number , m_seq
```

## EBNF

```
seq ::= m_seq?  
m_seq ::= number ( , m_seq )?
```

# \* Zero or More

Zero or More

**X**\*

**X** can be a terminal, non-terminal, or a parenthesized sequence of terminals and non-terminals.

Sometimes written { **X** }

BNF

```
seq ::= numbers number
```

```
numbers ::= number , numbers  
           |
```

EBNF

```
seq ::= numbers number
```

```
numbers ::= ( number , )*
```

EBNF

```
seq ::= ( number , )* number
```

## + One or More

One or More

$X_+$

**X** can be a terminal, non-terminal, or a parenthesized sequence of terminals and non-terminals.

BNF (no commas)

```
seq ::= number seq  
      | number
```

EBNF

```
seq ::= number+
```

Linked List EBNF

```
seq ::= number seq?
```

One or more of something

$X_+$

Zero or more of something

$X^*$

Optional elements

$X?$

Grouping

$( \dots )$



# New EBNF Grammar for Mini-Java

`java ::= public class name { method* }`

`method ::= (public|private)  
          static* type name ( method_args ) { statements }`

`method_args ::= method_arg , method_args  
              |  
method_arg ::= type name`

`type ::= type [ ] | name | "int" | "float" | "boolean"`

`statements ::= ( statement ; )*  
statement ::= full_name ( fun_args )`

`full_name ::= ( name . )* name  
fun_args ::= expr , fun_args`

`expr ::= string | int | expr + expr | expr * expr | ( expr )`

$S ::= A S \mid A$   
 $A ::= A S \mid B$   
 $B ::= C S D \mid E$

- What are the terminals and non-terminals for this grammar? (In class)
- Show that this grammar is ambiguous by giving a token sequence that has two possible concrete syntax trees (two derivations)
- Construct an unambiguous grammar to describe the same language, using BNF.
- Construct another unambiguous grammar to describe the same language, using EBNF.

(Adapted from Fundamental Structures of Computer Science, Wulf, Shaw, Hilfinger, Flon, pp 370)