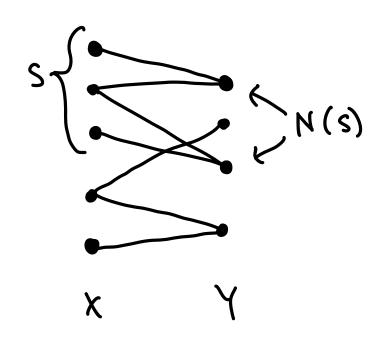
Hall's (Marriage) Thm (3.1.11): Let G be a bipartite graph w/ parts X and Y. Then,

G has a matching \Longrightarrow $|N(s)| \ge |s|$ that saturates X for all $S \le X$

Pf: \Rightarrow If G has such a matching M, the vertices in S are matched to |s| vertices, all of which must be in |N(s)|



Need a defin first

Def 3.1.6: Let McG be a matching.

a) An M-alternating path is a path PCG which alternates btwn. edges in M and edges not in M





b) An M-augmenting path is an M-alternating path whose endpoints are unsaturated





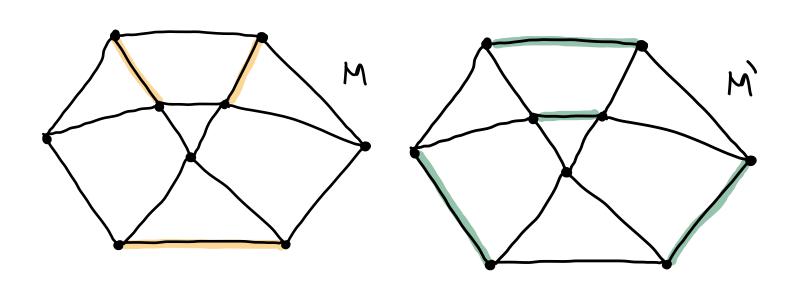
Idea: given an M-augmenting path, swap the edges and non-edges

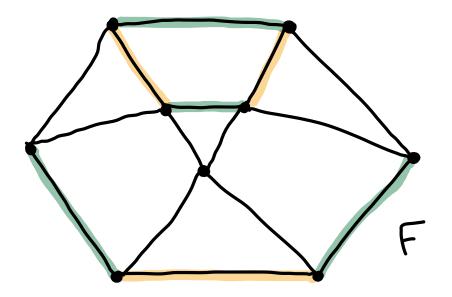




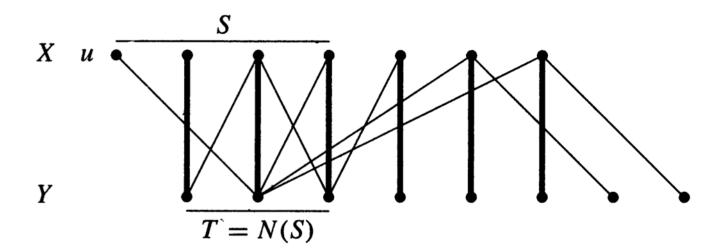
Always gives a larger matching

Thm 3.1.10: Let MSG be a matching. Then, Mis maximum (a) Ghas no M-angmenting path PF:





Back to pf of Hall's Thm!



Def (3.1.14/3.1.19): Let G be a graph

- a) Q S V(G) is a <u>vertex cover</u> of G if every edge in E(G) has ≥ 1 endpoint in Q
- b) L S F(G) is an <u>edge cover</u> of G if every edge in V(G) is incident to ≥ 1 edge in L
- C) d(G) := maximum size of independent set d'(G) := maximum size of matching P(G) := minimum size of vertex cover P'(G) := minimum size of edge cover

Class activity: compute &(G), &'(G), B(G), B'(G) for the graphs below

