

Announcements

- No class this Friday (10/27)
- No H/w this week (HW8 will be due Wed. 11/8)

- Exam 2 graded

Mean: 63.6

Median: 63.5

Std. dev.: 7.45

Problem scores:

Q_1 : 93%

Q_2 : 28%

Q_3 : 59%

Q_4 : 54%

Q_5 : 93%

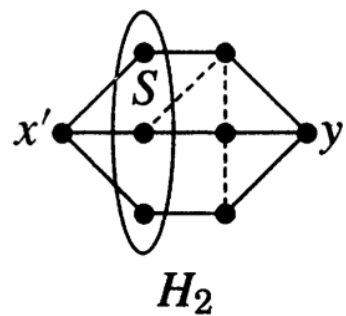
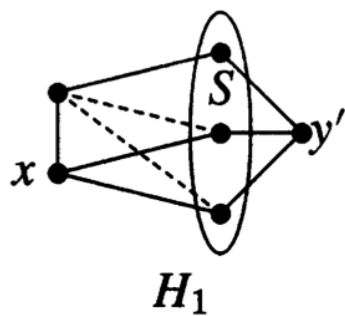
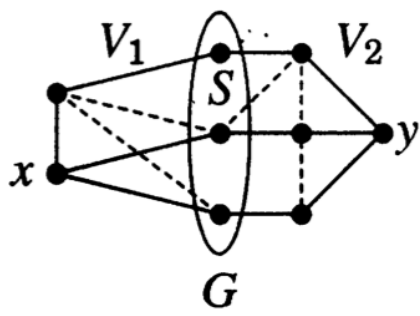
Menger's Theorem: If $x \neq y \in V(G)$ and $xy \notin E(G)$,
then $K(x, y) = \lambda(x, y)$

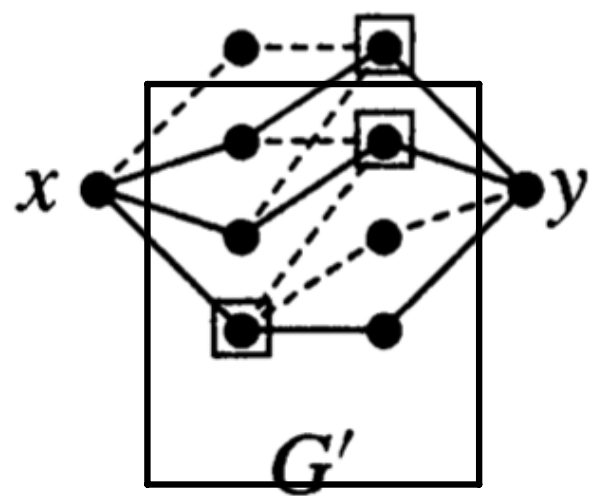
Pf: \geq) An x, y -cut must contain an internal vertex from each path in a set of pairwise internally-disjoint x, y -paths, so taking such a set of size $\lambda(x, y)$ gives $K(x, y) \geq \lambda(x, y)$.

\leq) Induction on $n := n(G)$.

Base case: $n=2$. If $xy \notin G$, then there is no x, y -path

$$K(x, y) = \lambda(x, y) = 0.$$





Similar results hold for directed graphs and for edge cuts

Def 4.2.11: Let D be a digraph

a) A vertex cut of D is a set $S \subseteq D$ s.t.

$D \setminus S$ is not **strongly** connected

b) If $S, T \subseteq V(D)$, $[S, T]$ denotes the set of edges w/ **tail** in S and **head** in T . An edge cut of D is $[S, \bar{S}]$ for some nonempty $S \subseteq V(D)$.

c) (Edge)-connectivity, $K(D)$, $K'(D)$, $K(x, y)$ defined the same w.r.t. vertex/edge cuts.

d) $\lambda(x, y)$ is still the largest # of internally-disjoint x, y -paths

Def: Let G be a graph or digraph.

a) $K'(x, y) = \min.$ size of $F \subseteq E(G)$ s.t. $G \setminus F$ has no x, y -path

b) $\lambda'(x, y) = \max.$ size of set of edge-disjoint x, y -paths

Thm: Let G be a graph or digraph.

a) Let $x \neq y \in V(G)$ with no edge from x to y .

i) IF G is a graph, then $K(x,y) = \lambda(x,y)$ (Menger)

ii) IF G is a digraph, then $K(x,y) = \lambda(x,y)$

b) Let $x \neq y \in V(G)$

i) IF G is a graph, then $K'(x,y) = \lambda'(x,y)$

ii) IF G is a digraph, then $K'(x,y) = \lambda'(x,y)$