Announcements: HWS posted Exam 1 graded

Median: 75/15

Mean: 71.4/15

5th. Lev.: 17.3

Regrade request deadline: next Web. (10/4)

Recall: Matrix tree thm .: For any loopless

graph G, and for any i,

T(G) = det Li(G), Laplacian

where L(G) = D(G) - A(G) is the Laplacian matrix of G.

Pf (Godsil-Royle, Algebraic Graph Theory):

Induction on | E(G)|, using Prop. 2.2.8:

T(G)=T(G/e)+T(G/e)

Base case: no edges:

$$T(G) = \begin{cases} 1, & n=1 \\ 0, & n>1 \end{cases} = det L^{i}(G). \sqrt{\frac{1}{2}}$$

Inductive step:

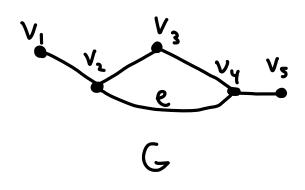
Let e be an edge with endpoints vi and vi.

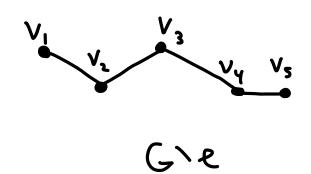
Let
$$E = i \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix}$$

$$E' = E'$$
remove now/col i

Then,
$$L(G) = L(G \setminus e) + E$$
, so $L^{i}(G) = L^{i}(G \setminus e) + E^{i}$

Since E' just has a single hunzero entry, expanding along row i gives





$$\begin{bmatrix}
1 & -1 & 0 & 0 & 0 \\
-1 & 3 & -1 & -1 & 0 \\
0 & -1 & 2 & -1 & 0 \\
0 & 0 & 0 & -1 & 1
\end{bmatrix}$$

$$\begin{bmatrix}
1 & -1 & 0 & 0 & 0 \\
-1 & 2 & -1 & 0 & 0 \\
0 & 0 & -1 & 2 & -1 \\
0 & 0 & 0 & -1 & 1
\end{bmatrix}$$

$$\begin{bmatrix}
1 & -1 & 0 & 0 & 0 \\
-1 & 2 & -1 & 0 & 0 \\
0 & 0 & 0 & -1 & 1
\end{bmatrix}$$

det L2(G) = det L2(G/e) + 1. det L2,4(G/e)

Return to proof:

When forming G-e, consider the combined vertex to be labelled Ui. All edges incident to Ui and/or Vi in G are represented by row/col i in L(G-e).

Therefore, deleting row/col. i from L(G.e) is equivalent to deleting rows/cols. i and i from L(G) i.e. $L^{i}(G-e) = L^{i}(G)$,

so (*) becomes

det Li(G) = det Li(G/e) + det Li(G·e) By the inductive hyp.,

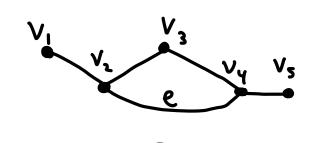
det L'(6,e) = 7(6,e)

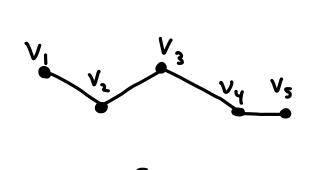
det L'(6-e) = T(6·e),

so by Prop. 2.2.8,

det Li(6) = T(6,e) + T(6,e) = T(6)

Return to example:





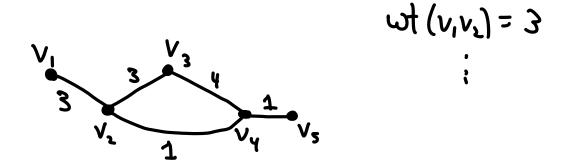
L'(6)

$$\det \begin{bmatrix} 0 & 0 & 0 \\ 0 & 2 & -1 & 0 \\ 0 & -1 & 3 & -1 \\ 0 & 0 & -1 & 1 \end{bmatrix} = 464 \begin{bmatrix} 0 & 0 & -1 & 1 \\ 0 & 2 & -1 & 0 \\ 0 & 2 & -1 & 0 \\ 0 & 0 & -1 & 1 \end{bmatrix} + 464 \begin{bmatrix} 0 & 0 & 1 \\ 0 & 2 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

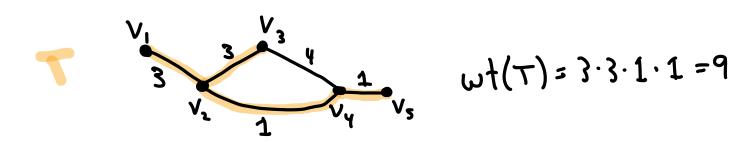
There are many generalizations of the Matrix Tree Theorem. Here's one:

Def:

a) A weighted graph G is a graph together with a function $wt: E(G) \rightarrow \mathbb{R}$



b) If T is a spanning tree of G, the weight of T is $\omega t(T) := TT \omega t(e)$ SEF(T)

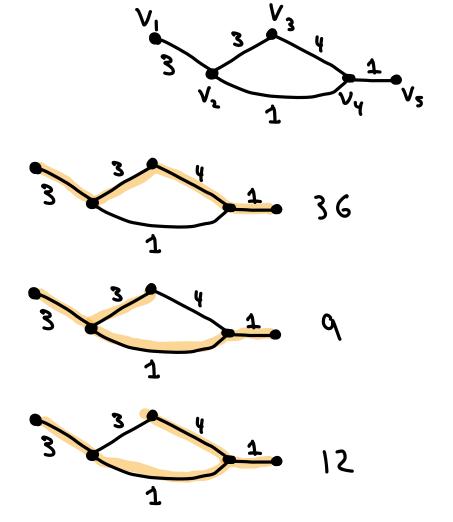


c) The tree sum T(G) of G is $T(G) = \sum_{T \text{ sp. tree of } G} \text{wt}(T)$

d) The (weighted) Laplacian matrix L(G) of G is given by:

Class activity:

Find t(G) and L(G) for:



T(G) = 57

Weighted Matrix Tree Theorem: For any loopless weighted graph G and any i,

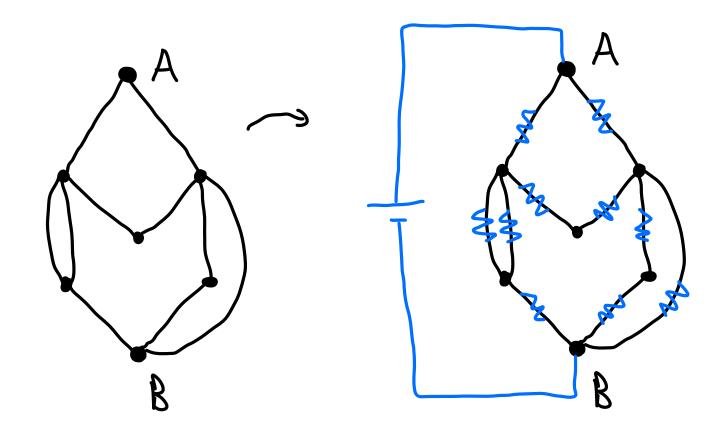
Pf: Homework!

Application / motivation:

Kirchoff's laws for electrical circuits Source: Postnikov lecture notes (link on 412 course website)

Let G be a (loopless) graph, and consider edges of G to represent resistors.

Choose vertices A and B to be connected to a source of electricity



Choose any orientation D of G (doesn't matter which)

Quantities associated to each edge e:

- · Current Ie through e
- · Voltage (or potential difference) Ve across e
- · Resistance Re of e (Re>0)
- Conductance $C_e! = \frac{1}{R_e}$

Three laws:

KI: At any vertex v, the sum of the in-currents equals the sum of the out-currents:

K2: For any cycle (in G, the (signed) sum of voltages is 0: $V_e = 2 + 3 + 1 - 6 = 0$ $V_e = 3 + 3 + 1 - 6 = 0$ $V_e = 3 + 3 + 1 - 6 = 0$ $V_e = 3 + 3 + 1 - 6 = 0$ $V_e = 3 + 3 + 1 - 6 = 0$ $V_e = 3 + 3 + 1 - 6 = 0$

where we traverse C in either direction, and the term involving be is positive iff we traverse e in the way it's oriented in D.

Ohm's Law: Ye & E(0),

Prop: Kz is equivalent to the following condition:

K2: There exists a (unique) function

 $U: V(G) \rightarrow \mathbb{R},$

called the potential function, s.t.

b)
$$V(B) = 0$$

Pf: Homework!