## Announcements:

Quiz 3 this Wed.

Midterm 1 Wed 9/2s in-class (50 minutes)

§ 5.2: Strong Induction and hell-Ordering

Recall:

Principle of Mathematical Induction: P(n) is true for all n if and only if

- · P(I) is true (base case)
- If we assume P(k) is true then P(k+1) is true (induction step)

Strong induction:

P(n) is true for all n if and only if

- · P(I) is true (base case)
- · If we assume P(1), P(2),..., P(k) are all true, then P(k+1) is true (induction step)

Only difference: now we get to assume P(1), --, P(k) in the induction step instead of just P(k)

No downside to using strong induction

Sometimes just assuming P(k) is enough

Other times, it helps to assume P(1),..., P(k)

Ex 2: Show that if nEZ, nZZ, then n can be written as a product of one or more primes

Pf: We prove this by induction. Let P(n) be the statement in can be written as a product of primes"

Base case: 2 is prime, so P(2) is true

Inductive step! Assure that P(1), ..., P(k) are true. If k+1 is prime, then P(k+1) is true. Otherwise,

k+1=ab for integers 25a,6<k+1. By the inductive

hypothesis, a and b can be written as products of prines:

a= P1P2 -- Pn, b= 9,92 -- 9m.

Then

k+1 = P1 .. Png1 .. gm

can also be written as a product of primes, so P(L+1) is true, and P(n) is true for all n=2 by induction

Ex 4: Prove that every amount of postage of 12 cents or more can be formed using just 4 cent and 5 cent stamps

Pf: We prove this via strong induction. Let P(n) be the statement in cents can be formed using just 4-cent 2 5-cent stamps

Base cases: 12 = 4 + 4 + 9 13 = 4 + 4 + 5 15 = 5 + 5 + 5

So P(12), P(13), P(14), and P(15) are all true.

Inductive step: Let k > 15, and assume that P(12), ..., P(k) are all true. We want to show that P(k+1) is true. Since k > 15, k+1 > 16, so k+1-4 > 12, so P(k+1-4) is true. Therefore, we can make k+1-4 cents using 4-cent and 5-cent stamps, so adding a 4-cent stamp gives k+1 cents. Thus,

P(k+1) is true, so P(n) is true for all n=12 by strong induction [

Ex 3: Consider the following game: Two piles of a matches



The players take turns removing \$1 matches from one of the piles. The players who takes the last match wins. Show that Player 2 can always guarantee a win. Class activity: play this game, and try to figure out a strategy.

Pf: We use strong induction. Let P(n) be "Player 2 can win whenever there are initially n matches in each pile"

Base case: If n=1, Player I must remove the I match from one of the piles. Player 2 takes the match from the other pile and wine.

Inductive step: Suppose k=1 and P(1), -, P(k) are true. For k+1 metches per pile, suppose Player I takes r matches from the first pile. Then Player 2 can take r matches from the other pile. If r= k+1, Player 2 wins. If 1 < r < k+1, then each pile has k+1-r

matches remaining, and it is Player 1's turn again. Since  $1 \le k+1-r \le k$ , P(k+1-r) is true, so Player 2 can now guarantee a win. Thus, P(k+1) is true, so by strong in duction, P(n) is true for all n.

## If the:

Well-Ordering Property: Every nonempty subset A of M has a smallest element.

Note: not true for subsets of 72 e.g.  $\{..., -8, -6, -4, ...\}$  or for subsets of  $R_+$  e.g.  $\{x \in R \mid 0 < x < 1\}$ 

Ex 5: Use the well-ordering property to prove the division algorithm!

If  $a \in 7L$ ,  $d \in 7L_+$ , then there are [unique] integers q, r with  $0 \le r < d$  such that a = qd + r

e.g. a=31, d=3

$$31 = \underbrace{10.3 + 1}_{6}$$
 (division w/ remainder)

Pf: Let S be the set of nonnegative integers of the form a-dq, where q is an integer.

S is nonempty since taking  $q \ll 0$  makes  $\alpha - dq \geq 0$ . By the well-ordering property, S has a smallest element  $Y = \alpha - dq_0$ .  $Y \in M$  since  $Y \in S$ . In addition, if  $Y \geq d$ , then  $Y_1 := \alpha - d(q_0 + 1) \in S$  and  $Y_1 < Y_2$ , but this can't happen since  $Y_1 := \alpha + d(q_0 + 1) \in S$  and  $Y_1 < Y_2$ , but this can't happen since  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_1 < Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_1 < Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_1 < Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_1 < Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_1 < Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_2 := \alpha + d(q_0 + 1) \in S$  and  $Y_3 := \alpha + d(q_0 + 1)$