

# Announcement:

Midterm 2 tonight!

7:00pm - 8:30pm in 217 Noyes Lab. (ref. sheet allowed)

Be early!

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Exam covers: Ch 1-3 (focus on Ch. 2, 3),  
plus circuit application of matrix tree thm.

Most focus: topics that appeared in lecture or homework

Some focus: topics in relevant subsections of textbook

Low/no focus: topics in subsections we didn't cover at all

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Types of graphs: (dis.)conn., bipartite, paths, cycles, trees, forests, complete (bipartite) graphs, digraphs, weighted graphs

Walks, trails, circuits

Things graphs have:

Eulerian circuits (Euler Thm.)

Perfect matching (Hall's Thm., Tutte's Thm.)

Trees:

Equiv. def'n's

Prüfer code & Cayley's formula

Spanning subgraphs & spanning trees

Matrix tree thm.

Kirchoff's Laws and Kirchoff's Thm.

Algorithms:

Kruskal (min. wt. spanning tree)

Dijkstra (distances)

Gale-Shapley (stable matching)

Algorithmic thinking

Matchings: general concept

Perfect vs. maximum vs. maximal

M-alt. paths & M-aug. paths

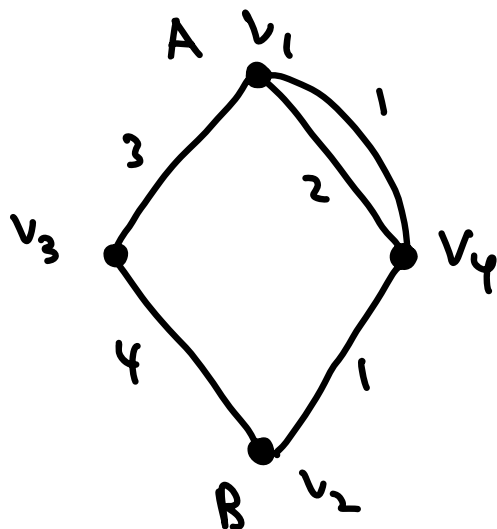
Theorems: Berge, Hall, Tutte, Berge-Tutte,  
Petersen x2

Relationships btwn. matchings, vertex/edge covers, and  
indep. sets

k-factors

Examples:

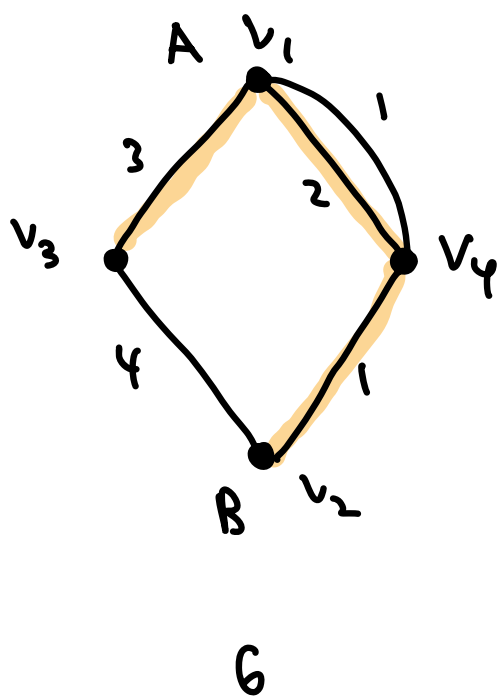
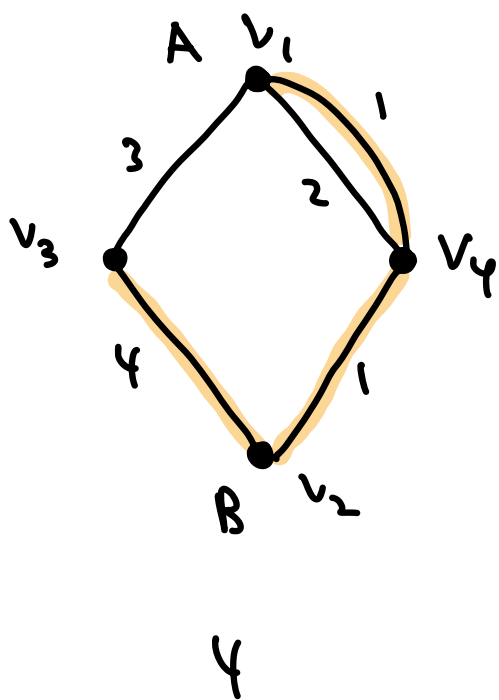
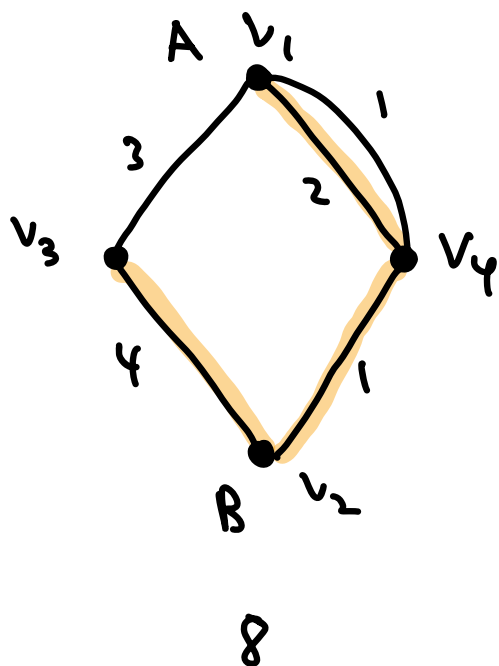
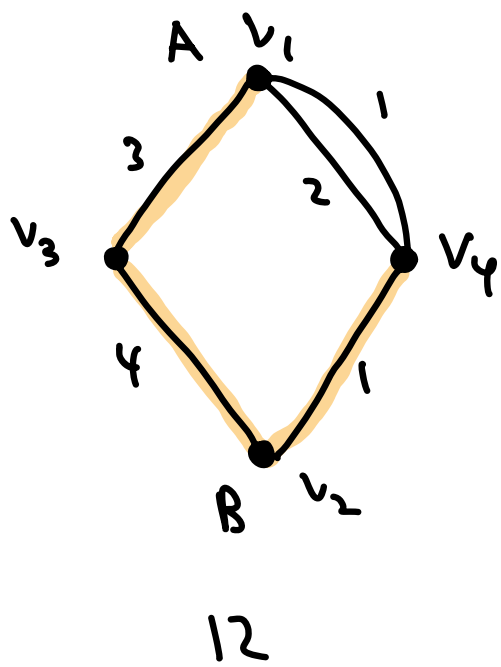
1) Consider the weighted graph  $G$ :

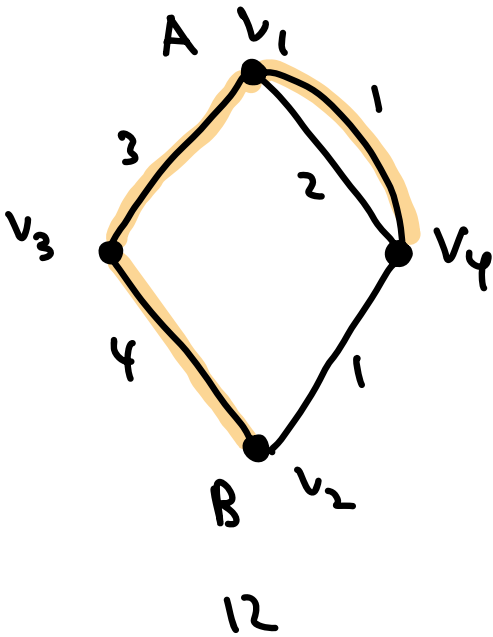
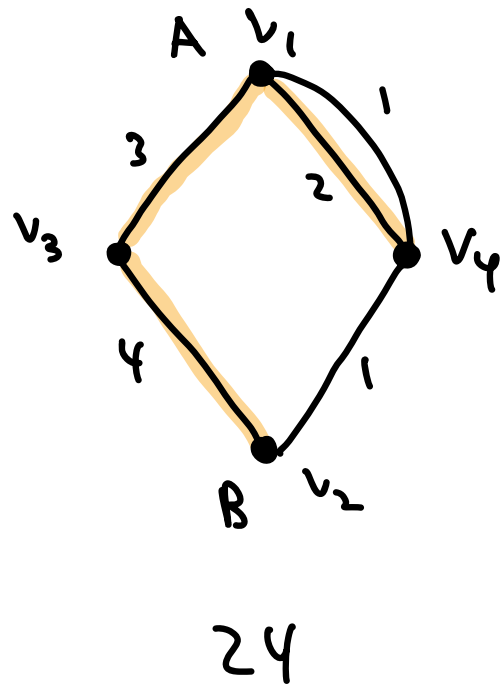
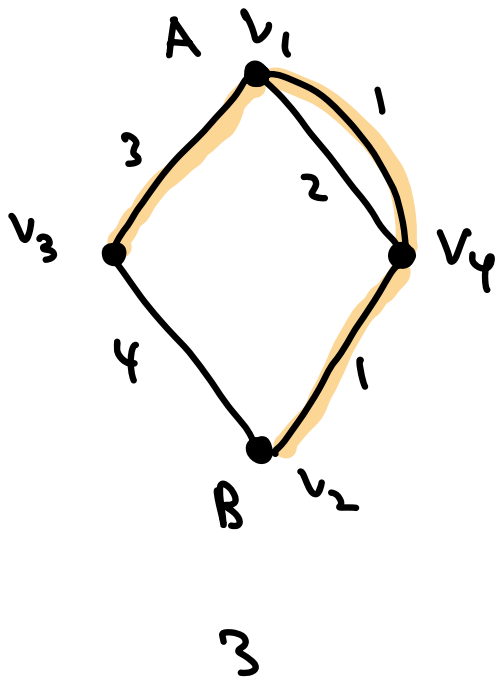


Find the effective resistance from  $A$  to  $B$

Sol'n: Step 1: compute  $\tau(G)$

OR

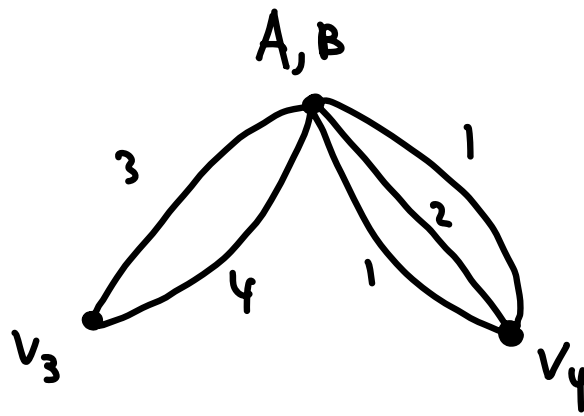




$$\tau(G) = 12 + 8 + 4 + 6 + 3 + 24 + 12$$

Step 2:

$(G \sqcup AB) \cdot AB :$




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2) Let  $G$  be a simple graph s.t.  $\delta(G) \geq k$   
and  $n(G) \geq 2k$ . Prove that  $G$  has a matching  
of size  $\geq k$ .



3) Compute  $\tau(k_{i,m})$