## Announ cements

HWY posted (due Wed. 9/27)

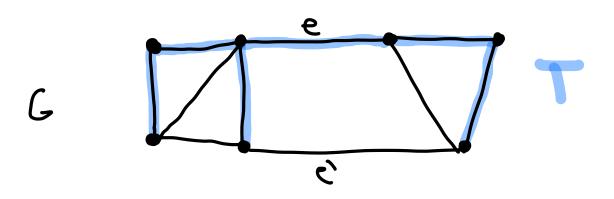
Prop (2.1.6/2.1.7): Let G be a graph w/ spanning trees T, T'.

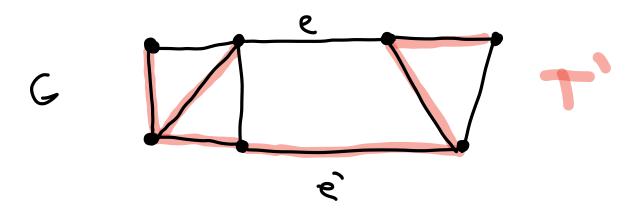
a) For all e ∈ E(T), ∃ e' ∈ E(T') s.t. (T v e') le is a spanning tree of G.

b) For all  $e' \in E(T')$ ,  $\exists e \in E(T)$  s.t.  $(Tue') \setminus e$  is a spanning tree of G.

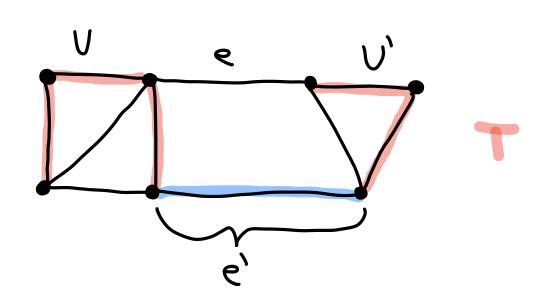
If you tell me which edge to remove, I'll tell you which edge to add

If you tell me which edge to add, I'll tell you which edge to remove





Pf: a) By defin (e) (2.1.4) of a tree, every edge of T is a cut-edge. So let U and U' be the two components of T e. Since T' is connected,  $\exists e' \in E(T')$  with endpoints in U and U', so  $(T \cdot e) \cup e'$  is connected, has vertex set V(G) and has |E(T)| = n(G) - 1 edges, so by  $def^{in}(c)$ ,  $(T \cdot e) \cup e'$  is a spanning tree of G.



b) If T=T', let e=e'. Otherwise, Tue' has n(G) edges, it is not a tree, by defin (c), so by defin (n), it has a cycle C. C is the unique cycle since (Tue') e'=T is a tree and hence acyclic. In addition, C contains ≥2 edges so let e F E(C) s.t. e + e'. By Thm. 1.2.14, e is not a cut edge since it belongs to a cycle, so (Tre) ve is a conn. graph w/ n(6)-1 edges and vertex set V(G); hence a spanning tree.

Note the choice for e

Def 5.1.9/12: G: Graph

- a) The distance dlu,v) from u to v is the shortest length of a u,v-path (so if no path)
- b) The diameter diam G is the maximum distance btwn. any two vertices in G (a if disconn.)

  diam G = max u,v ev(G) d(u,v)
  - c) The eccentricity of a vertex u is  $E(u) := \max_{v \in V(G)} d(u, v)$

(i.e. dian G = maxuev(G) & (u))

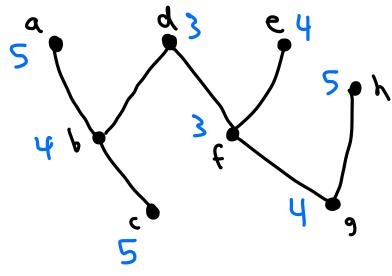
- d) The radius of G is

  rad  $G := \min_{u \in V(G)} E(u)$
- e) The center of G is the induced subgraph

  G[{ueV(6)|E(u)=rad(u)}]

Class activity:

Find dian 6, rad 6, the eccentricity of each vertex, and the center



rad G = 3

diam G = 5center G = 5

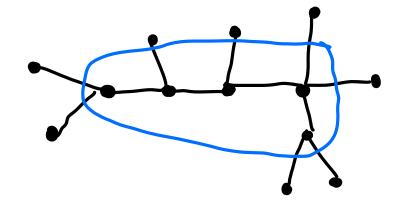
Jordan Tree Theorem: The center of a tree

Pf sketch: induction on n:= n(T)

n < 2: the center is the entire tree • or •

ns2: Let T':= T[non-leaves]. By defin C

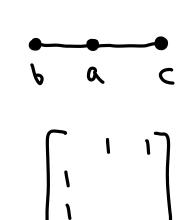
(2.1.4), T' is a tree.

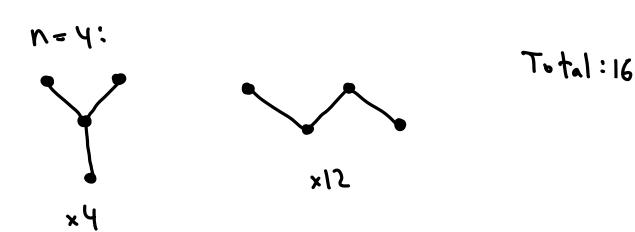


If  $u \in V(T)$ ,  $\mathcal{E}_{T}(u) = \mathcal{E}_{T'}(u) + 1$ , if u is a leaf of T, it's not in the center of T. Thus, the center of T equals the center of T.  $\Pi$ 

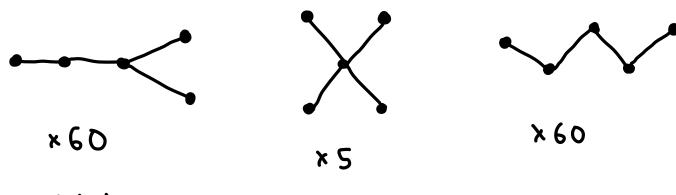
How many (labelled) n-vertex trees are there?

$$h=3:$$





N= 5:



Total: 125

Pattern?

Cayley's Formula (Thm 2.2.3): There are nn-2 labelled trees with n vertices

[For technical reasons, latel set is some SCN]

Pf idea!

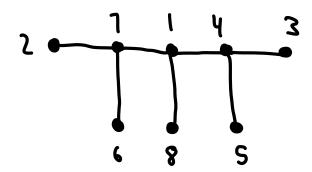
n-vertex trees \implies length n-2 strings of elts. of S

Def: The Prüfer (ode Prn(T) =  $(a_1, ..., a_{n-2})$  of T is given by the following algorithm:

At step i:

- delete the leaf w/ the smallest label
- a; is the label for the lunique) neighbor of the leaf

Ex:



Can go backwards:

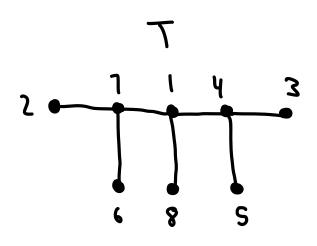
$$Pru(T) = 744171$$

Pf of Cayley's Formula: h=1 good

We prove that for n ? 2

$$T \longleftrightarrow Prn(T)$$

is a bijection.



Cor 2.2.4: Let  $d_{1,-}, d_{n} \in \mathbb{Z}_{\geq 1}$  s.t.  $d_{1} + \cdots + d_{n} = 2n-2$ . Then the number of trees w/ label set  $\{1,-,n\}$  s.t. vertex i has degree  $d_{i}$  is  $\frac{(n-2)!}{\prod (d_{i}-1)!}$