Announcements HW9 posted (or will be today) - due sun 11/10 Quiz 7 Wed. in class

§10.1: Graphs

Def: A graph (=(V, E) (onsists of V: a nonempty set of vertices, and

E: a set of edges

Each edge has two vertices as endpoints. If they are the same, the edge is called a loop

Def: A digraph D has the same defin except that edges are directed start end

Def: A graph or digraph is called <u>simple</u> if

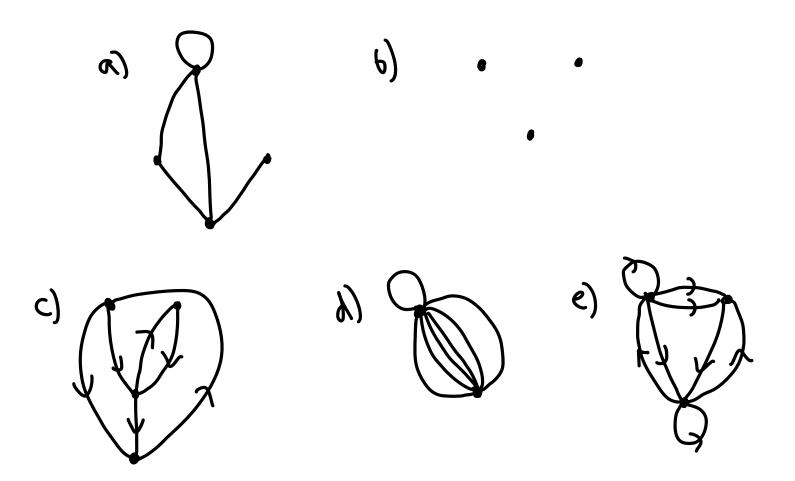
a) it has no loops ("loopless")

b) it has no multiple edges

edges w/ same tail/head (digraph)

(Di)graphs w/ multiple edges are called multi(di) graphs

Class activity: Graph or digraph? Simple? Malti-?



Rosen has many examples of how graphs/digraphs can be used to represent real-world data (Ex 1-13, also on HW)

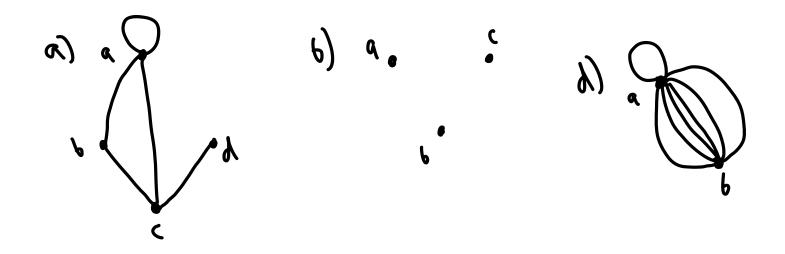
§10.2: Graph terminology, and special types of graph Def: Let G = (V, E) be a graph

a) $u_1v \in V$ $(u \neq v)$ are <u>adjacent</u> or <u>neighbors</u> if there is an edge $e \in E$ with endpoints u and v. e is <u>incident</u> to its endpoints.

b) The <u>neighborhood</u> of $V \in V$ is the set N(V) of all neighbors of V. If $A = \{v_1, v_2, ..., v_k\}$, then $N(A) = N(V_1) \cup N(V_2) \cup ... \cup N(V_k)$, the set of all vertices adjacent to any vertex in A

c) The degree of veV is the number d(v) of edges incident to v (counting loops twice).

Class activity: Find all neighborhoods and Agrees



Handshake theorem: For a graph G with m edges,

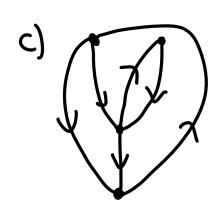
In particular, the number of vertices of odd-degree is always even!

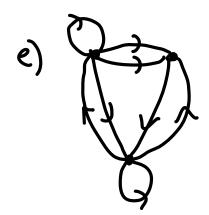
Def: Let D= (V, E) be a digraph, V e V.

The in-legree deg-(v) of v is the number of edges w/ end/head v.

The out-legree degt(v) of v is the number of edges w/ start/tail v.

Class activity: Find all in / out-degrees

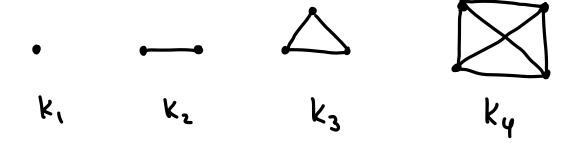


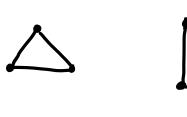


Handshake theorem: For a digraph D with m edges,

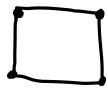
Special (undirected, simple) graphs

a) Complete graph kn: all pairs of vertices are adjacent

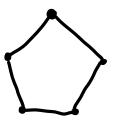




C3



Cy



Cs

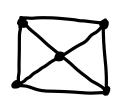


Also Cs (doesn't matter how you place the vertices)

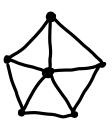
C) Wheel Wn: Cn with a hub



W3



Wy

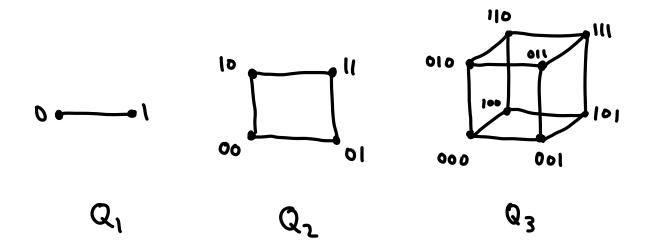


Ws

d) Hypercube Qn

V = { binary strings of length n}

N(v) = { all strings off by one digit from v}



Def: G is bipartite if there is a set partition $V=V_1 UV_2$ such that every edge has one endpoint in disjoint V_1 and the other in V_2

Class activity: Of the above graphs, which are bipartite? (if time)

$$j+1=7$$
 $j=6$

$$\frac{1}{\binom{6}{2} = \binom{3}{k}}$$

$$\binom{10}{2} = \binom{1-3}{k}$$

$$N+1=17 n=16$$

 $2k+1=5$
 $k=2$