Announcements:

Quiz today!

Milterm 3: Next Wed. 11/19 7:00-8:30pm Noyes 217

Recall: The chromatic polynomial of G is $\chi(G;k) := number of proper k-colorings of G$

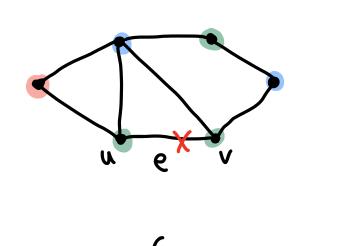
There is a method to compute $\chi(G;k)$ recursively using heletion-contraction, allowing for a computation of $\chi(G;k)$, and thus $\chi(G)$, for any (in dividual) graph G.

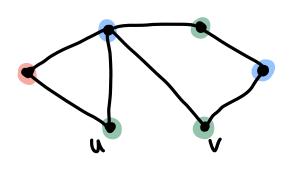
Thm 5.3.6: Let G be a simple graph and $e \in E(G)$. Then,

 $\chi(G;k) = \chi(G \cdot e;k) - \chi(G \cdot e;k)$

Pf: Every proper coloring of G is a proper

coloring of Gre, and a proper coloring of Gre
is a proper coloring of G iff it give distinct
colors to the endpoints u and v of e

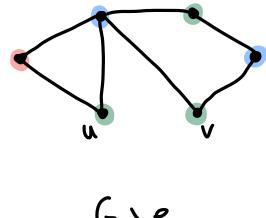




Gre

So
$$\chi(G;k) = \chi(G \cdot e;k) - | \{proper k - colorings of G \cdot e\} | \{where u, v have same color\} |$$

and



G·e

can remove mult, edges

Example 5.3.7:

G.e
$$\stackrel{\sim}{=}$$
 k_3 $\chi(k_3; k) = k(k-1)(k-2)$

$$e^{-6} = k^3 = k^3$$

So,

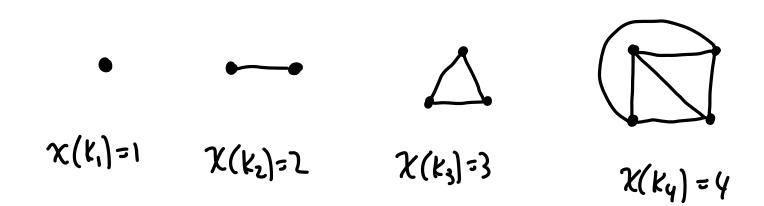
$$\chi(G;k) = \chi(P_4;k) - \chi(K_3;k) = k(k-1)(k^2-3k+3)$$

Chapter 6: Planar Graphs

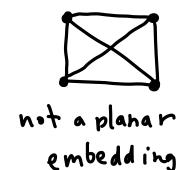
Goal: Find possible values of X(G) for

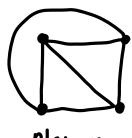
planar graphs G

can be drawn on a piece of paper w/out crossings



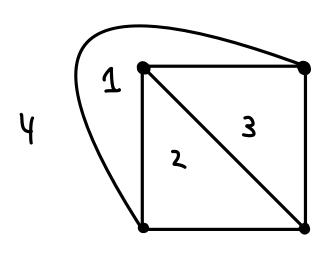
Def 6.1.4: A graph G is planar if it has a drawing w/out crossings, called a planar embedding or a plane graph





Planar embedding

so ky is planar Def: The faces of planar embedding are the maximal conn. regions of the plane not intersecting vertices and edges



Remark: It is surprisingly difficult to make some of these ideas rigorous. Need topology and the "Jordan Curve Theorem"

Prop 6.1.2: Ks and K3,3 are not planar

