

Announcements

Midterm 1 Wed in-class (50 minutes)

Reference sheet allowed (one A4 sheet, both sides)

Sections covered: 1.1-1.4, 2.1-3, 3.1-2, 5.1-2, 6.1 (no 6.2)

Practice problems partial sol's posted

See policy email for more

Today: review

After class: problem session (review)

Tomorrow 10:30-11:30 am: office hour (via Zoom @ snow day link)

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Midterm 1 review:

(Partial) list of topics:

Logic

Propositions

English vs. symbols

and, or, not, implies, for all, there exists

Truth tables

Sets

Roster notation vs. set builder notation

Special sets (\mathbb{N} , \mathbb{R} , \emptyset , etc.)

Venn diagrams

Subset, power set, Cartesian product

Cardinality

Set operations: union, intersection, set-minus, complement

Set identities (1-10)

Proof techniques: elt. chasing, membership tables

Functions

Definition

Domain, codomain, range/image, preimage

Injective/surjective/bijective (& pf. techniques)

Composition

Inverses + invertibility

Algorithms

Definition

Properties (describe and check)

Perform an algorithm

Write an algorithm

Searching / sorting / greedy change

Big - O

Precise def'n of O , Ω , Θ ; proof techniques

Tricks & heuristics ($1 < \log x < x < x^2 < \dots < e^x < \dots$)

Induction

Mathematical vs. strong

Base case, inductive step

Critique proofs

Various examples from class & H/W

Counting

Sum/product/subtraction/division rules

Combining the rules (examples from class, HW)

Examples:

1) Prove that $A \cup (A \cap B) = A$ for all sets A, B

Pf: We prove the result by showing that

a) $A \cup (A \cap B) \subseteq A$ and b) $A \subseteq A \cup (A \cap B)$.

a) Let $x \in A \cup (A \cap B)$. Then either $x \in A$ or $x \in A \cap B$.

In the former case, clearly $x \in A$, and in the latter case, $x \in A$ and $x \in B$, so $x \in A$. Thus, $A \cup (A \cap B) \subseteq A$.

b) Let $x \in A$. Then $x \in A$, so $x \in A \cup (A \cap B)$. Thus, $A \subseteq A \cup (A \cap B)$. □

- 2) How many license plates can be made using either
- three digits followed by three uppercase letters or
 - three uppercase letters followed by three digits

a) We have 10 choices for each of digits 1, 2, 3 and 26 choices for each of digits 4, 5, 6, so by the product rule, there are

$$10 \cdot 10 \cdot 10 \cdot 26 \cdot 26 \cdot 26 = 17576000 \text{ such license plates}$$

b) We have 26 choices for each of digits 1, 2, 3 and 10 choices for each of digits 4, 5, 6, so by the product rule, there are

$$26 \cdot 26 \cdot 26 \cdot 10 \cdot 10 \cdot 10 = 17576000 \text{ such license plates}$$

There is no overlap between the two cases, so in total there are

$$17576000 + 17576000 = 35152000 \text{ valid license plates}$$

3) Consider $f: \mathbb{Z} \rightarrow \mathbb{Z}_{\geq 0}$ where $f(x) = |x| + 1$. Determine, with proof, whether f is inj., surj., bij., or none.

We claim that f is neither injective nor surjective.

Pf: First we consider injectivity. f is injective if and only if whenever $f(x) = f(y)$, $x = y$. However, $f(1) = 2 = f(-1)$, so f is not injective.

Now, f is surjective if and only if $f(\mathbb{Z}) = \mathbb{Z}_{\geq 0}$; in other words, if and only if for every $y \in \mathbb{Z}_{\geq 0}$, there exists $x \in \mathbb{Z}$ s.t. $f(x) = y$. We claim that $0 \notin f(\mathbb{Z})$. To see this, notice that for all $x \in \mathbb{Z}$, $f(x) = |x| + 1 \geq 0 + 1 = 1 > 0$. So $0 \notin f(\mathbb{Z})$, and f is not surjective. \square

4) Name 3 of the 7 properties that algorithms should have, give a short description, and write some pseudocode which fails at only this property.

Properties: input, output, definiteness, correctness, finiteness, effectiveness, generality

Finiteness: For any input, the algorithm should produce the desired output after a finite number of steps

int-sqrt (pos integer x): (find \sqrt{x} if $\sqrt{x} \in \mathbb{Z}$;
 $i := 0$ otherwise, return -1)
while ($i^2 \neq x$)
 $i := i + 1$
return i

Correctness: The algorithm should produce the correct output to the desired problem

int-sqrt (pos integer x):
 $i := 0$
while ($x^2 \neq i$)
 $i := i + 1$
return i

Effectiveness: It must be possible to perform each step exactly, and in a finite amount of time

int-sqrt (pos. integer x):
 $A := \{(i, i^2) \mid i \in \mathbb{Z}_+\}$
If $(i, x) \in A$ for some $i \in \mathbb{Z}$
return i
Otherwise, return -1

5) Find and prove a simple big- Θ estimate for

$$f(x) = (7n^n + n2^n + 3^n)(n! + 3^n)$$

(Heuristic: Choose fastest-growing term of each factor)

Let $g(x) = n^n \cdot n!$. We claim that $f(x)$ is $\Theta(g(x))$.

Pf: We show that f is $O(g)$ and g is $O(f)$

$$\text{Let } C = 18 \quad k = 10$$

Then $n! \geq 3^n$ for all $x > k$ (this needs pf)

and $n^n \leq n2^n$, $n^n \geq 3^n$ for all $x > k$ (this needs pf)

So for all $x > k$,

$$\begin{aligned} |f(x)| &= (7n^n + n2^n + 3^n)(n! + 3^n) \\ &\leq (7n^n + n^n + n^n)(n! + n!) \\ &= 18n^n n! = C|g(x)|, \end{aligned}$$

so f is $O(g)$.

(Other direction: $C = k = 1$ works)