

Announcements:

- HW9 due tomorrow (Thurs. 10/30) at 9am (office hour today)

- Exam 3 graded

Problem Scores:

Mean: 62.9 } out of

Median: 64.5 } 95

Std. dev.: 15.0

Q_1 : 82%

Q_3 : 56%

Q_2 : 50%

Q_4 : 75%

- Plan for rest of semester (rough!)

Wed 11/29: §6.1, §6.3 if time

Fri 12/1: §6.3

Mon 12/4: §6.3 (cont.) and Quiz 4

Wed 12/6: Final exam review

(Some sort of review session + office hours)

Thurs 12/14, 8:00-11:00am: Final exam!

132 Berier Hall (not one of our usual rooms!)

Recall:

Euler's Formula: Let G be a connected plane graph w/
 n vertices, e edges, and f faces. Then,

$$n - e + f = 2$$

Last time: used this to study regular polyhedra

Today: a bunch of corollaries

Remark 6.1.22:

- a) Since n and e don't depend on the planar embedding, neither does f .
- b) Recall that the dual graphs of two different planar embeddings of G can be nonisomorphic. However, if the dual graph has n^* vertices, e^* edges, and f^* faces, $n^* = f$, $e^* = e$, $f^* = n$, so these numbers are independent of planar embedding.
- c) For a graph w/ k conn. components, we have

$$n - e + f = k + 1$$

Thm 6.1.23:

a) If G is a simple planar graph w/ ≥ 3 vertices,
then

$$e(G) \leq 3n(G) - 6$$

b) If G is also \triangle -free, then

$$e(G) \leq 2n(G) - 4$$

Pf:

Corollary (6.1.24) : K_5 and $K_{3,3}$ are non planar
(already proved this)

Pf: Class activity!

Def 6.1.25:

- a) A maximal planar graph is a simple planar graph that is not a spanning subgraph of another planar graph (i.e. adding edges makes the graph nonplanar)
- b) A triangulation is a simple plane graph where every face boundary is a 3 cycle.

Prop 6.1.26: Let G be a simple n -vertex plane graph. The following are equivalent:

A) $e(G) = 3n - 6$

B) G is a triangulation

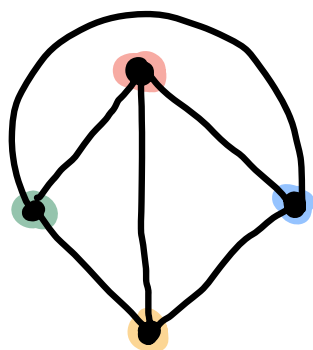
C) G is (an embedding of) a maximal planar graph.

Pf:

Recall our main question for this section: what is the maximum number of colors needed to give any (loopless / simple) planar graph a proper coloring?

k -Color Theorem: Every planar graph is k -colorable.

There is no 3-color theorem since $\chi(K_4) = 4$ and K_4 is planar



Six-Color Theorem (Exercise 6.3.2): Every planar graph is 6-colorable.

Pf: Induction on $n(G)$.

What about 5 colors? 4 colors?

Next time.