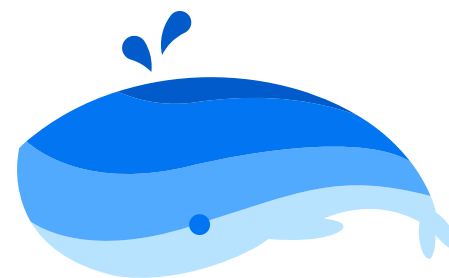


Implement SQL parser and type system in Databend

Databend Cloud is an affordable cloud **data warehouse** developed in Rust. In this exploration, we will dive into Databend's internal, specifically examining its SQL parser and type system that are meticulously crafted for Rust.

by @AndyLok



Who am I

骆迪安

<https://github.com/andylokandy>

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Databend Planner Team
Distributed transaction

Rust contributor
Idris-lang contributor

Databend

<https://databend.rs>

Feature-Rich

Instant Elasticity

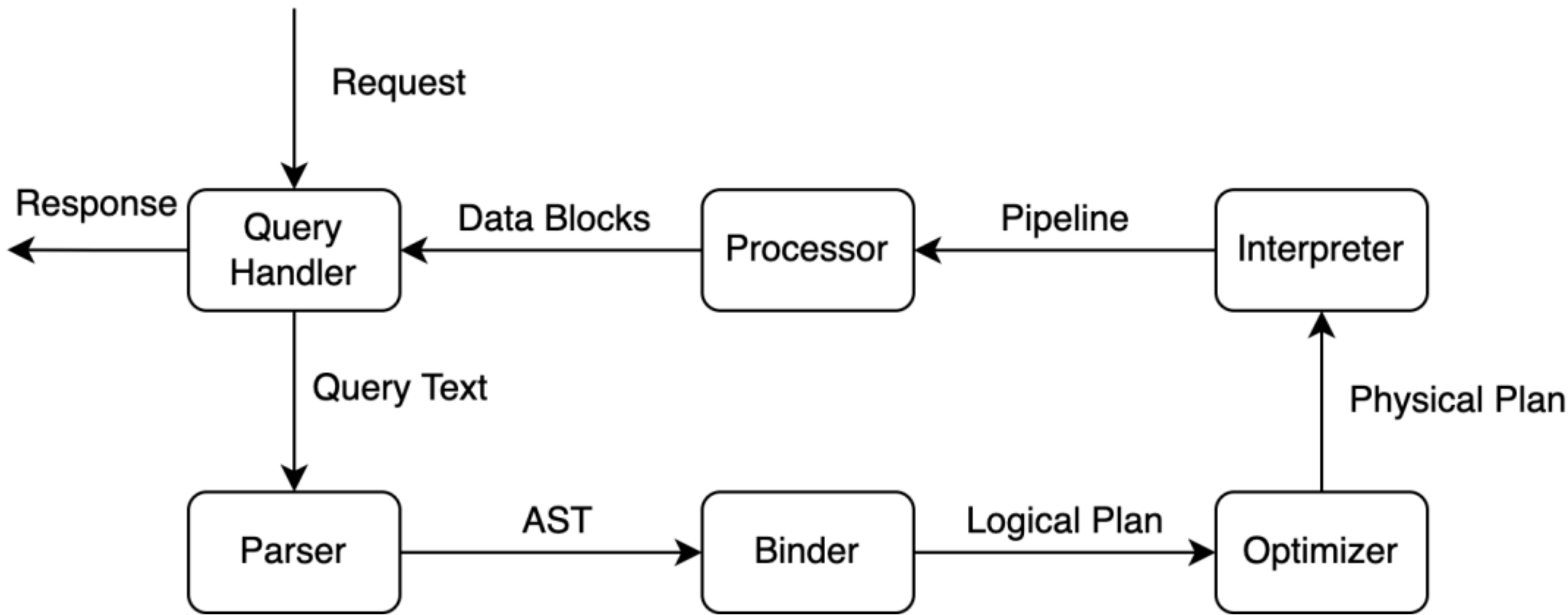
Support for Semi-Structured Data

MySQL/ClickHouse Compatible

Low Cost

Easy To Use

Cloud Native (S3, Azure Blob, Google Cloud Storage,
Alibaba Cloud OSS, etc)



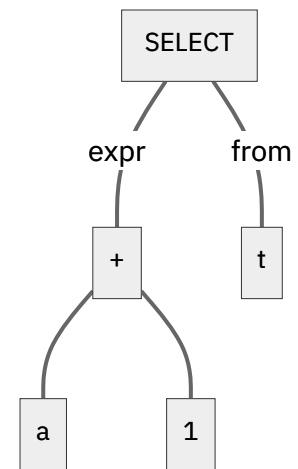
SQL Parser

Check for syntax violation and construct **Abstract Syntax Tree (AST)**

Input

SELECT a + 1 **FROM** t

Output



Tokenize

Convert **String** to **Token Stream**

Parse

Convert **Token Stream** to **AST**

Tokenize

Recognise **tokens** by regular expression

Regular expression for each token kind

```
Ident = [_a-zA-Z][_a-zA-Z0-9]*  
Number = [0-9]+  
Plus = \+  
SELECT = SELECT  
FROM = FROM
```

Tokenize

Recognise **tokens** by regular expression

Input

```
SELECT a + 1 FROM t
```

Output

```
[  
  Keyword(SELECT),  
  Ident(a),  
  BinOp(+),  
  Number(1),  
  Keyword(FROM),  
  Ident(t),  
]
```


Parse

Define **grammar rule** and construct **AST**

Backus–Naur form (BNF) grammars

```
<select_statement> ::= SELECT <expr> FROM <ident>  
<expr> ::= <number>  
          | <ident>  
          | <expr> <op> <expr>  
<op> ::= + | - | * | /
```

Parse

Define **grammar rule** and construct **AST**

Input

```
[  
  Keyword(SELECT),  
  Ident(a),  
  BinOp(+),  
  Number(1),  
  Keyword(FROM),  
  Ident(t),  
]
```

Output

```
SelectStatement {  
  projection: Expr::BinaryOp {  
    op: Op::Plus,  
    args: [  
      Expr::ColumnRef("a"),  
      Expr::Constant(Scalar::Int(1))  
    ]  
  }  
  from: "t",  
}
```

Choosing SQL Parser

sqlparser-rs

ANTLR4

LALRPOP

nom

Tokenizer

<https://github.com/maciejhirsz/logos>

```
#[derive(Logos)]
pub enum TokenKind {
    #[regex(r"[ \t\r\n\f]+", logos::skip)]
    Whitespace,

    #[regex(r#"[_a-zA-Z][_a-zA-Z0-9]*"#)]
    Ident,
    #[regex(r"[0-9]+")]
    Number,

    #[token("+")]
    Plus,

    #[token("SELECT", ignore(ascii_case))]
    SELECT,
    #[token("FROM", ignore(ascii_case))]
    FROM,
}
```

Tokenizer

<https://github.com/maciejhirsz/logos>

Input

```
SELECT a + 1 FROM t
```

Output

```
[  
  Token { kind: TokenKind::Select, text: "SELECT", span: 0..6 },  
  Token { kind: TokenKind::Ident, text: "a", span: 7..8 },  
  Token { kind: TokenKind::Plus, text: "+", span: 9..10 },  
  Token { kind: TokenKind::Number, text: "1", span: 11..12 },  
  Token { kind: TokenKind::From, text: "from", span: 13..17 },  
  Token { kind: TokenKind::Ident, text: "t", span: 18..19 },  
]
```

Tokenizer

<https://github.com/maciejhirsz/logos>

Input

```
SELECT a + 1 FROM t
```

Output

```
[  
  Token { kind: TokenKind::Select, text: "SELECT", span: 0..6 },  
  Token { kind: TokenKind::Ident, text: "a", span: 7..8 },  
  Token { kind: TokenKind::Plus, text: "+", span: 9..10 },  
  Token { kind: TokenKind::Number, text: "1", span: 11..12 },  
  Token { kind: TokenKind::From, text: "from", span: 13..17 },  
  Token { kind: TokenKind::Ident, text: "t", span: 18..19 },  
]
```

Span

SELECT	a	+	1	FROM	t
0..6	7..8	9..10	11..12	13..17	18..19

Error Report

Pretty print the error report thanks to the **span** information

```
error:
--> SQL:1:19
|
1 | create table a (c varchar)
| -----          - ^^^^^ expected `BOOLEAN`, `BOOL`, `UINT8`, `TINYINT`, `UINT16`, `SMALLINT`, or 33 more ...
| |               |
| |               | while parsing `<column name> <type> [DEFAULT <default value>] [COMMENT '<comment>']`
| | while parsing `CREATE TABLE [IF NOT EXISTS] [<database>.]<table> [<source>] [<table_options>]`
```

Parser

<https://github.com/rust-bakery/nom>

Terminal

Recognize only one **token** from **token stream**

Combinator

Combine **terminals** and other small parsers into a larger parser

Terminal

Recognize only one **token** from **token stream**

Recognize a token that has the exactly **text**

```
fn match_text(text: &str)
    -> impl FnMut(&[Token]) -> IResult<&[Token], Token>
{
    satisfy(|token: &Token| token.text == text)(i)
}
```

Recognize a token that is of the **token kind**

```
fn match_token(kind: TokenKind)
    -> impl FnMut(&[Token]) -> IResult<&[Token], Token>
{
    satisfy(|token: &Token| token.kind == kind)(i)
}
```

Combinator

Combine **terminals** and other small parsers into a larger parser

`tuple(a, b, c)`

`alt(a, b, c)`

`many0(a)`

`many1(a)`

`opt(a)`

BNF

Formally defined **grammar rule**

```
<select_statement> ::=  
    SELECT <expr> FROM <ident>
```

Code

Practical Rust code using **nom**

```
tuple((  
    match_token(SELECT),  
    expr,  
    match_token(FROM),  
    match_token(Ident),  
))
```

BNF

Formally defined **grammar rule**

```
<select_statement> ::=  
    SELECT <expr> [FROM <ident>]
```

Code

Practical Rust code using **nom**

```
tuple((  
    match_token(SELECT),  
    expr,  
    opt(tuple((  
        match_token(FROM),  
        match_token(Ident),  
    )))  
))
```

BNF

Formally defined **grammar rule**

```
<select_statement> ::=  
    SELECT <expr>+ [FROM <ident>]
```

Code

Practical Rust code using **nom**

```
tuple((  
    match_token(SELECT),  
    many1(expr),  
    opt(tuple((  
        match_token(FROM),  
        match_token(Ident),  
    )))  
))
```

BNF

Formally defined **grammar rule**

```
<select_statement> ::=
    SELECT (<expr> AS <ident> | <expr>)+
    [FROM <ident>]
```

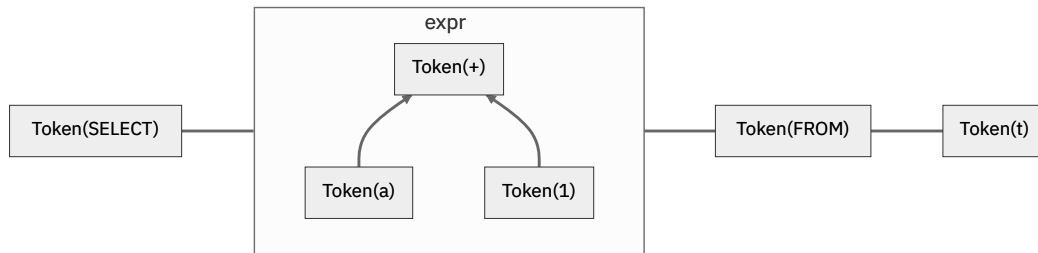
Code

Practical Rust code using **nom**

```
tuple((
    match_token(SELECT),
    many1(alt((
        tuple((
            expr,
            match_token(AS),
            expr,
        )),
        expr,
    ))),
    opt(tuple((
        match_token(FROM),
        match_token(Ident),
    )))
))
```

Parse Tree

The **parse tree** of `SELECT a + 1`



AST

The **AST** of `SELECT a + 1`

```
SelectStatement {  
  projection: Expr::BinaryOp {  
    op: Op::Plus,  
    args: [  
      Expr::ColumnRef("a"),  
      Expr::Constant(Scalar::Int(1))  
    ]  
  }  
  from: "t",  
}
```


map()

Use **map()** to convert **Parse Tree** into **AST**

```
fn select_statement(input: &[Token])
-> IResult< &[Token], SelectStatement >
{
    map(
        tuple((
            match_token(SELECT),
            many1(alt((
                tuple((
                    expr,
                    match_token(AS),
                    expr,
                )),
                expr,
            ))),
            opt(tuple((
                match_token(FROM),
                match_token(Ident),
            )))
        )),
        |(_, projections, _, opt_from)| SelectStatement {
            projections: projections
                .map(|(expr, _, alias)| {
                    Projection::Aliased(expr, alias)
                })
                .collect(),
            from: opt_from.map(|(_, from)| from),
        }
    )(input)
}
```

nom-rule

<https://github.com/andylokandy/nom-rule>

Simplify nom parser using **BNF**-like grammar

Syntax definition using **nom-rule**

```
nom_rule! {  
  SELECT  
  ~ (#expr ~ AS ~ Ident | #expr)+  
  ~ (FROM ~ Ident)?  
}
```

Generated **nom** parser

```
tuple((  
  match_token(SELECT),  
  many1(alt((  
    tuple((  
      expr,  
      match_token(AS),  
      expr,  
    )),  
    expr,  
  )),  
  opt(tuple((  
    match_token(FROM),  
    match_token(Ident),  
  ))  
))
```

nom-rule

<https://github.com/andylokandy/nom-rule>

Simplify nom parser using **BNF**-like grammar

Update the example using **nom-rule**

```
fn select_statement(input: &[Token])
-> IResult< &[Token], SelectStatement >
{
    map(
        nom_rule! {
            SELECT
            ~ (#expr ~ AS ~ Ident | #expr)+
            ~ (FROM ~ Ident)?
        },
        |(_, projections, _, opt_from)| SelectStatement {
            projections: projections
                .map(|(expr, _, alias)| {
                    Projection::Aliased(expr, alias)
                })
                .collect(),
            from: opt_from.map(|(_, from)| from),
        }
    )(input)
}
```

nom-rule

<https://github.com/andylokandy/nom-rule>
Simplify nom parser using **BNF**-like grammar

nom-rule cheatsheet

nom-rule	Translated
TOKEN	match_token(TOKEN)
"+"	match_text("+")
a ~ b ~ c	tuple((a, b, c))
(...)*	many0(...)
(...)+	many1(...)
(...)?	opt(...)
a b c	alt((a, b, c))

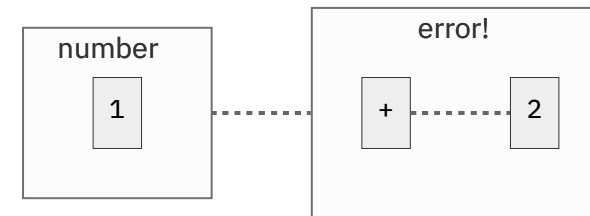
Left Recursion

A classical **problem** every parser will meet when parsing expression

Try to define the syntax for expression like `1 + 2`

```
<expr> ::= <number>  
        | <expr> + <expr>
```

oops! The second rule will **never apply**



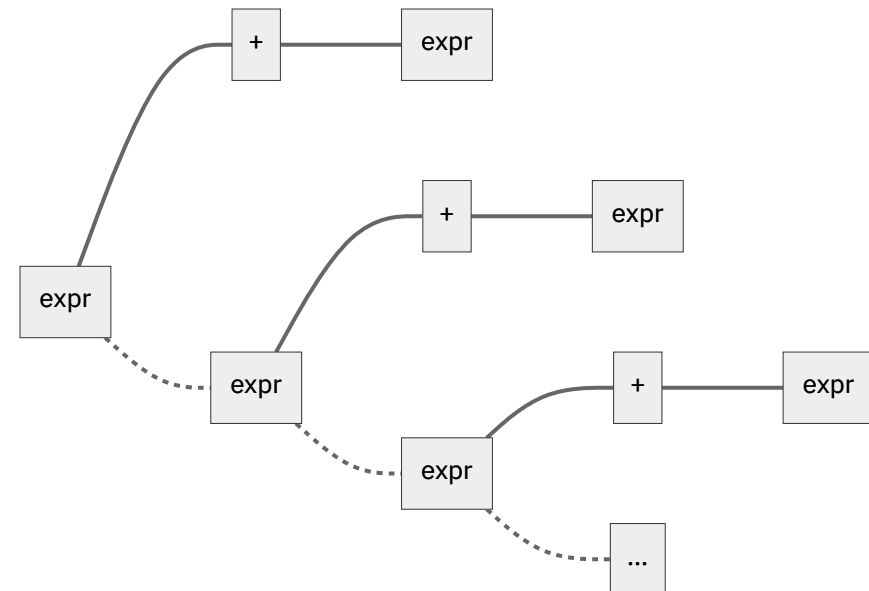
Left Recursion

A classical **problem** every parser will meet when parsing expression

Try again! **Reverse the order!**

```
<expr> ::= <expr> + <expr>  
         | <number>
```

but not gonna work... it'll **never stop**



Pratt Parser

<https://github.com/segeljakt/pratt/>

A decent solution to the problem of **Left Recursion**

Instead of parsing to a tree, parse the flatten elements

`<expr_element> ::= + | <number>`

Pratt Parser

<https://github.com/segeljakt/pratt/>

A decent solution to the problem of **Left Recursion**

Construct the AST using **Pratt Parser**

```
use pratt::PrattParser;

impl<I: Iterator<Item = ExprElement> PrattParser<I> for ExprParser {
    type Output = Expr;

    fn query(&mut self, elem: &ExprElement) -> Affix {
        match elem {
            ExprElement::Plus =>
                Affix::Infix(Precedence(20), Associativity::Left),
            ExprElement::Number(_) => Affix::Nilfix,
        }
    }

    fn primary(&mut self, elem: ExprElement) -> Expr {
        match elem {
            ExprElement::Number(n) => Expr::Number(n),
            _ => unreachable!(),
        }
    }

    fn infix(
        &mut self,
        lhs: Expr,
        elem: ExprElement,
        rhs: Expr,
    ) -> Expr {
        match elem {
            ExprElement::Plus(n) => Expr::Plus(lhs, rhs),
            _ => unreachable!(),
        }
    }
}
```


Type Check

Validate the **semantic** of the SQL

1. Given an AST

```
1 + 'foo'
```

2. **Desugar** AST into function calls

```
plus(1, 'foo')
```

3. Infer type of **literals**

```
1 :: Int8  
'foo' :: String
```

Type Check

Validate the **sematic** of the SQL

4. So we desire to find a function overload tha matches

```
plus(Int8, String)
```

5. Query function overloads for `plus()`

```
plus(Int8, Int8) -> Int8
plus(Int16, Int16) -> Int16
plus(Timestamp, Timestamp) -> Timestamp
plus(Date, Date) -> Date
```

6. However, no matching overload is found, thus typechecker reports a **type error**

```
1 + 'foo'
^ function `plus` has no overload for parameters `(Int8, String)`
```

available overloads:

```
plus(Int8, Int8) -> Int8
plus(Int16, Int16) -> Int16
plus(Timestamp, Timestamp) -> Timestamp
plus(Date, Date) -> Date
```

Type Judgement

Formal rules to **prove type** of a expression

1. Rule to assume the type of **boolean** literal

$\vdash \text{TRUE} : \text{Boolean}$

2. Also, the same for **numbers** and **string**

$\vdash 1 : \text{Int8}$
 $\vdash \text{"foo"} : \text{String}$

3. Rule to prove the type of **function call**

$$\frac{\Gamma \vdash e1 : \text{Int8} \quad \Gamma \vdash e2 : \text{Int8}}{\Gamma \vdash \text{plus}(e1, e2) : \text{Int8}}$$

* $\Gamma : \text{Type Environment}$

Type Environment

The meaning of the mystery `Γ`

What is the type of variable `a`?

```
SELECT 1 + a
```

a. determined by querying the **table metadata**

```
SELECT 1 + a from t
```

b. determined by type checking the **subquery**

```
SELECT 1 + a from (  
    SELECT number as a from numbers(100)  
)
```

Subtype

1. Given an AST

```
1 + 256
```

2. **Desugar** AST to function call

```
plus(Int8, Int16)
```

3. oops! There is no matching overloads

```
plus(Int8, Int8) -> Int8  
plus(Int16, Int16) -> Int16
```

Subtype

4. Because `Int8` is subtype of `Int16`, `Int8` can be cast to `Int16`

```
Int8 <: Int16
```

5. Update the type rule of function call to accept **subtype**

$$\frac{\Gamma \vdash e1 : T1 \quad \Gamma \vdash e2 : T2 \quad T1 <: \text{Int16} \quad T2 <: \text{Int16}}{\Gamma \vdash \text{plus}(e1, e2) : \text{Int16}}$$

6. Finally we prove

```
⊢ plus(1, 256) : Int16
```

Generic

1. Type of **array**

```
⊢ [1, 2, 3] : Array<Int8>
```

2. oops! It's impossible to enumerate all possible types

```
get(Array<Int8>, Int64) -> Int8  
get(Array<Array<Int8>>, Int64) -> Array<Int8>  
get(Array<Array<Array<Int8>>>, Int64) -> Array<Array<Int8>>  
...
```

3. Use **generic** in the function signature

```
get<T>(Array<T>, Int64) -> T
```

Generic

4. Assume we are checking this expression

```
get([1, 2, 3], 0)
```

5. Type check its arguments first

```
get(Array<Int8>, Int8)
```

6. We have got the **substitution** after **unification**, then apply the substitution to the signature

```
{T=Int8}  
// then apply the substitution to the signature  
get<Int8>(Array<Int8>, Int64) -> Int8
```

7. Finally we prove

```
⊢ get([1, 2, 3], 0) : Int8
```


Evaluation

Typechecker produces expression that is ready to be evaluated

```
pub enum Expr {  
  Constant {  
    scalar: Scalar,  
    data_type: DataType,  
  },  
  ColumnRef {  
    id: String,  
    data_type: DataType,  
  },  
  Cast {  
    is_try: bool,  
    expr: Box<Expr>,  
    dest_type: DataType,  
  },  
  FunctionCall {  
    name: String,  
    args: Vec<Expr>,  
    return_type: DataType,  
    eval: Box<Fn([Value]) -> Value>,  
  },  
}
```

Evaluation

Expression of `1 + a`

```
Expr::FunctionCall {  
  name: "plus",  
  args: [  
    Expr::Constant {  
      scalar: Scalar::Int8(1),  
      data_type: DataType::int8,  
    },  
    Expr::ColumnRef {  
      id: "a",  
      data_type: DataType::Int8,  
    },  
  ],  
  return_type: DataType::Int8,  
  eval: Box<Fn([Value]) -> Value>,  
}
```

Evaluation

The place where actual work takes place

```
fn plus_eval(args: &[amp;Value]) -> Value {  
    ...  
}
```

Vectorized input which can either be a **column** or **scalar** (all values are same)

```
enum Value {  
    Scalar(Scalar),  
    Column(Column),  
}  
  
enum Scalar {  
    Int8(i8),  
    Int16(i16),  
    Boolean(bool),  
    ...  
}  
  
enum Column {  
    Int8(Vec<i8>),  
    Int16(Vec<i16>),  
    Boolean(Vec<bool>),  
    ...  
}
```

Evaluation

A working example for `plus(Int8, Int8) -> Int8`

```
fn plus_eval(args: &[amp;Value]) -> Value {
  match (&args[0], &args[1]) {
    (Value::Scalar(Scalar::Int8(lhs)), Value::Scalar(Scalar::Int8(rhs))) => {
      Value::Scalar(Scalar::Int8(lhs + rhs))
    },
    (Value::Column(Column::Int8(lhs)), Value::Scalar(Scalar::Int8(rhs))) => {
      let result: Vec<i8> = lhs
        .iter()
        .map(|lhs| *lhs + rhs)
        .collect();
      Value::Column(Column::Int8(result))
    },
    (Value::Scalar(Scalar::Int8(lhs)), Value::Column(Column::Int8(rhs))) => {
      let result: Vec<i8> = rhs
        .iter()
        .map(|rhs| lhs + *rhs)
        .collect();
      Value::Column(Column::Int8(result))
    },
    (Value::Column(Column::Int8(lhs)), Value::Column(Column::Int8(rhs))) => {
      let result: Vec<i8> = lhs
        .iter()
        .zip(rhs.iter())
        .map(|(lhs, rhs)| *lhs + *rhs)
        .collect();
      Value::Column(Column::Int8(result))
    },
    _ => unreachable!()
  }
}
```

Evaluation

Extract the pattern for **vectorization**

```
fn register_2_arg_int8(&mut self, name: String, eval: impl Fn(i8, i8) -> i8) {
    self.register_function(Function {
        signature: FunctionSignature {
            name: "plus",
            arg: [DataType::Int8, DataType::Int8],
            return_type: DataType::Int8,
        },
        eval: |args: &[Value]| -> Value {
            match (&args[0], &args[1]) {
                (Value::Scalar(Scalar::Int8(lhs)), Value::Scalar(Scalar::Int8(rhs))) => {
                    Value::Scalar(Scalar::Int8(eval(lhs, rhs)))
                },
                (Value::Column(Column::Int8(lhs)), Value::Scalar(Scalar::Int8(rhs))) => {
                    let result: Vec<i8> = lhs
                        .iter()
                        .map(eval(*lhs, rhs))
                        .collect();
                    Value::Column(Column::Int8(result))
                },
                (Value::Scalar(Scalar::Int8(lhs)), Value::Column(Column::Int8(rhs))) => {
                    let result: Vec<i8> = rhs
                        .iter()
                        .map(eval(lhs, *rhs))
                        .collect();
                    Value::Column(Column::Int8(result))
                },
                (Value::Column(Column::Int8(lhs)), Value::Column(Column::Int8(rhs))) => {
                    let result: Vec<i8> = lhs
                        .iter()
                        .zip(rhs.iter())
                        .map(|(lhs, rhs)| eval(*lhs, *rhs))
                        .collect();
                    Value::Column(Column::Int8(result))
                },
                _ => unreachable!()
            }
        },
    });
}
```

Evaluation

As a result, it's easy to register overload for any **vectorized** input

```
registry.register_2_arg_int8("plus", |lhs: i8, rhs: i8| lhs + rhs);  
registry.register_2_arg_int8("minus", |lhs: i8, rhs: i8| lhs - rhs);
```

But still many **duplicated types** to register

```
registry.register_2_arg_int16("plus", |lhs: i16, rhs: i16| lhs + rhs);  
registry.register_2_arg_int16("minus", |lhs: i16, rhs: i16| lhs - rhs);  
registry.register_2_arg_int32("plus", |lhs: i32, rhs: i32| lhs + rhs);  
registry.register_2_arg_int32("minus", |lhs: i32, rhs: i32| lhs - rhs);  
registry.register_2_arg_int64("plus", |lhs: i64, rhs: i64| lhs + rhs);  
registry.register_2_arg_int64("minus", |lhs: i64, rhs: i64| lhs - rhs);
```

Evaluation

```
// Marker type for `Int8`
struct Int8Type;

// Define all methods needed to process data of the type
trait ValueType {
    type Scalar;

    fn data_type() -> DataType;
    fn downcast_scalar(Scalar) -> Self::Scalar;
    fn upcast_scalar(Self::Scalar) -> Scalar;
    fn iter_column(Column) -> impl Iterator<Item = Self::Scalar>;
    fn collect_iter(impl Iterator<Item = Self::Scalar>) -> Column;
}
```

```
impl ValueType for Int8Type {
    type Scalar = i8;

    fn data_type() -> DataType {
        DataType::Int8
    }
    fn downcast_scalar(scalar: Scalar) -> Self::Scalar {
        match scalar {
            Scalar::Int8(val) => val,
            _ => unreachable!(),
        }
    }
    fn upcast_scalar(scalar: Self::Scalar) -> Scalar {
        Scalar::Int8(scalar)
    }
    fn iter_column(col: Column) -> impl Iterator<Item = Self::Scalar> {
        match col {
            Column::Int8(col) => col.iter().cloned(),
            _ => unreachable!(),
        }
    }
    fn collect_iter(iter: impl Iterator<Item = Self::Scalar>) -> Column {
        let col = iter.collect();
        Column::Int8(col)
    }
}
```

Evaluation

Extract the data type using **marker type**

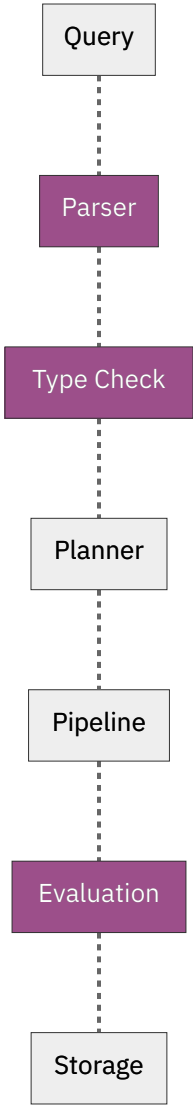
```
fn register_2_arg<I1: ValueType, I2: ValueType, Output: ValueType>(
    &mut self,
    name: String,
    eval: impl Fn(I1::Scalar, I2::Scalar) -> Output::Scalar
) {
    self.register_function(Function {
        signature: FunctionSignature {
            name: "plus",
            arg: [I1::data_type(), I2::data_type()],
            return_type: Output::data_type(),
        },
        eval: |args: &[Value]| -> Value {
            match (&args[0], &args[1]) {
                (Value::Scalar(lhs), Value::Scalar(rhs)) => {
                    let lhs: I1::Scalar = I1::downcast_scalar(lhs);
                    let rhs: I2::Scalar = I2::downcast_scalar(rhs);
                    let res: Output::Scalar = eval(lhs, rhs);
                    Output::upcast_scalar(0::upcast_scalar(res))
                },
                ...
            }
        }
    });
}
```


Evaluation

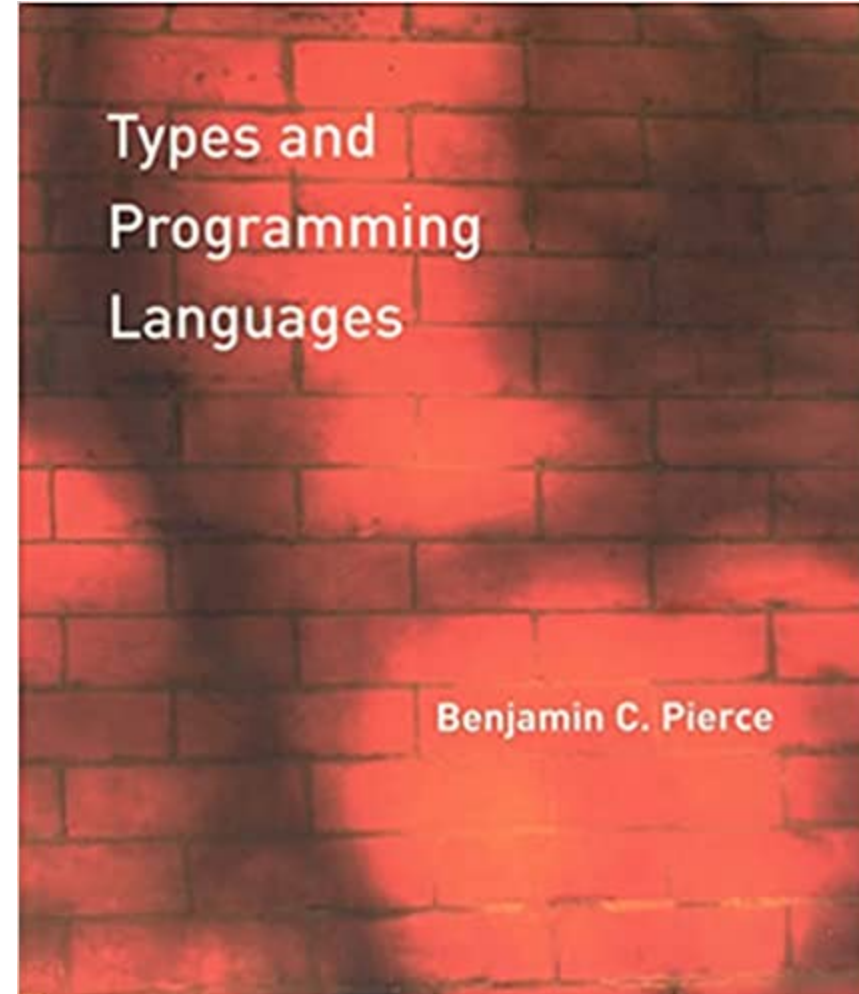
As a result, its easy to register overload for input of **any data type**

```
registry.register_2_arg:<Int8Type, Int8Type, Int8Type>(  
    "plus", |lhs: i8, rhs: i8| lhs + rhs  
);  
registry.register_2_arg:<Int8Type, Int8Type, Int8Type>(  
    "minus", |lhs: i8, rhs: i8| lhs - rhs  
);
```

Conclusion



Conclusion



Types and Programming Languages