

BramDeSerializer

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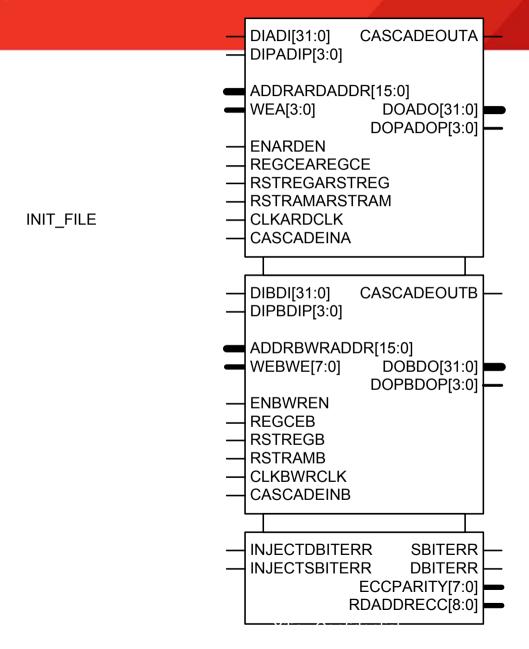
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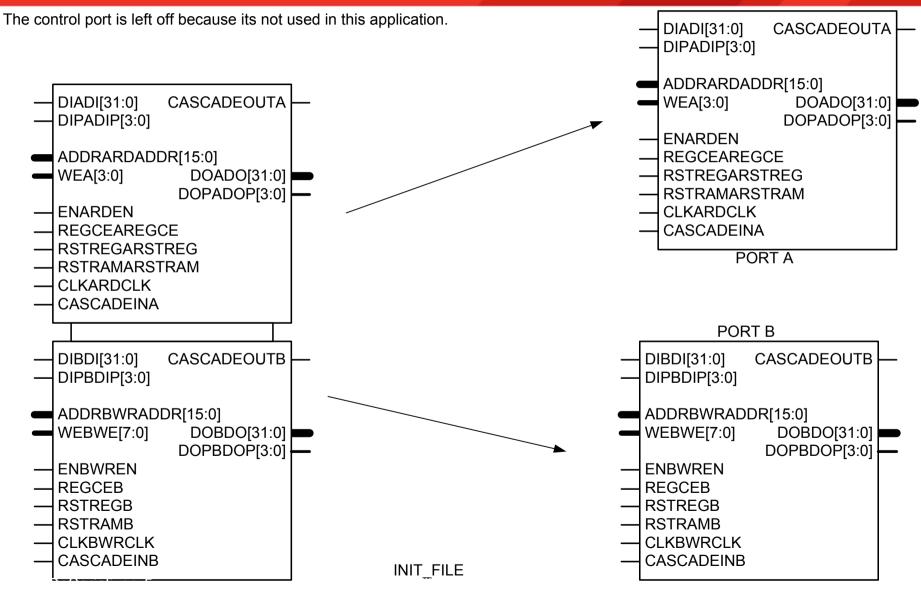


RAMB36E1





RAMB36E1 Used in this document

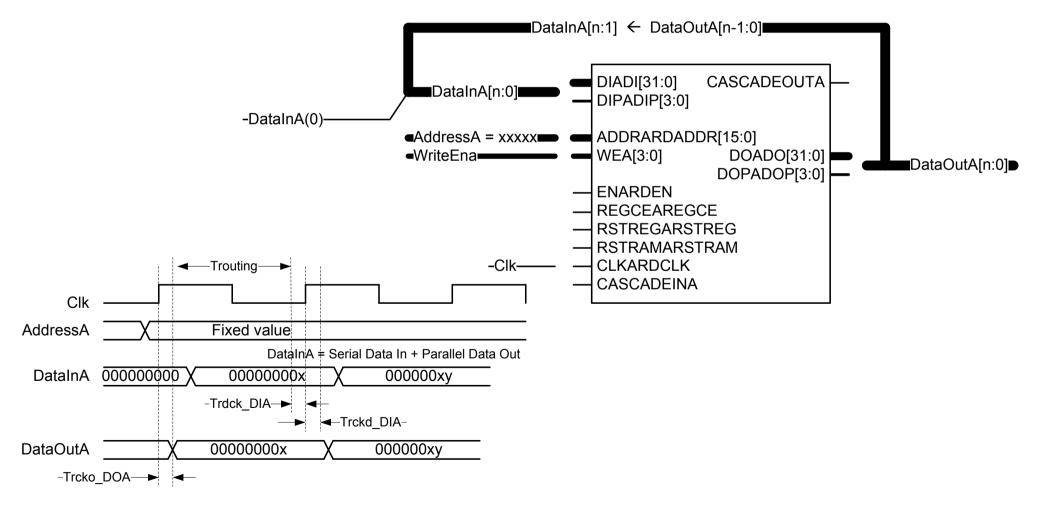




DeSerializer

This sheet shows a 36-bit serial-to-parallel register. Any possible setup of the RAMB36E1 can be turned into a serial-to-parallel register. This is a waste of a RAMB36 because only one address is used to make a serila-to-parallel register.

Hint: With less data bits needed it might be possible to include the address counter into the Block-RAM itself.



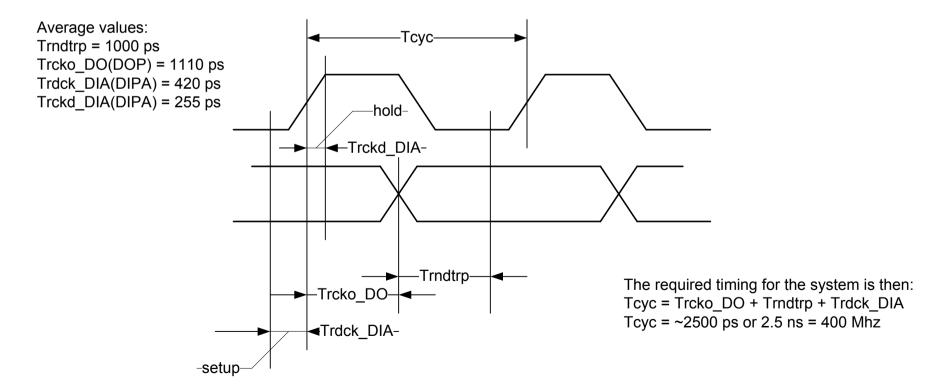


DeSerializer Timing

The most important timing for this application is the routing from the output of the Block-RAM back to the input of the Block-RAM. Let's call this timing Trndtrp (T_{Round Trip}).

This routing must thus be placed under strict timing control when the highest possible speed must be obtained from the de-serializer.

The total round truip timing and thus the maximal operation frequency also depents on the parameters of the RAMB36E1 block. The figure shows a calculation example for a XC6VLX240T–2 component.

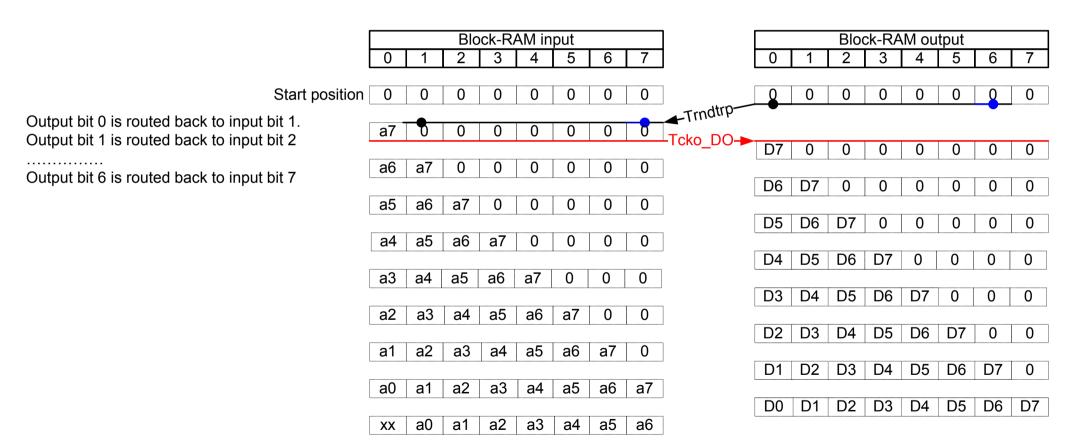




DeSerializer Data Flow

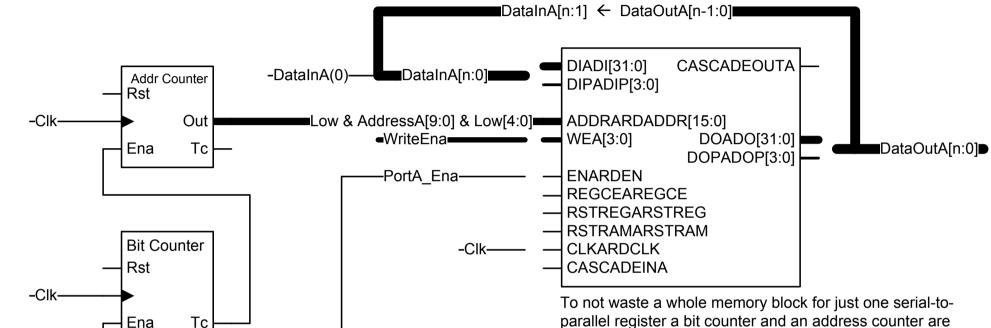
Just to show how the data flows in and out of the Block-Ram a eight bit example is given here

The serial input data is received MSB first and is named: a7, a6, down to a0. The output data of the Block-RAM is named D7, D6, D5 down to D0.





DeSerializer with a Twist



The "Bit counter" must count as far as the data depth of the Block-RAM. Example when the Block-RAM is set to 32-bit+4-bit and all 36bit are used as shift register, the counter must count from 0 to 35 and issue a terminal count at bit 36.

When all address are written. Port A must be set in read mode in order to be able to read all memory contents via the data out bus. parallel register a bit counter and an address counter are added to the design,

The bit counter is enabled at the same time as the Block-RAM. it counts the exact number of bits shifted into a specific memory address.

Except counting the written bits, so far nothing has changed since the previous example.

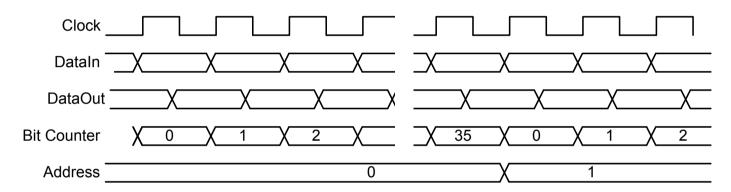
When the bit counter enables at it's terminal count an address counter to increase the address by one suddenly the design becomes interesting.

The Block-RAM is now changed into a pile of serial-to-parallel registers or otherwise written the Block-RAM becomes a serialto-parallel register with a build in memory/FIFO.



-PortA Ena-

DeSerializer with a Twist (waveforms)



Before rising clock edge.

DataIn = New Serial input bit at DataIn(0) and looped back data from DataOut(n-1:0)

After rising clock edge.

DataOut = New Serial data and looped back data from DataOut(n-1:0)

When the bit counter has seen 36 new serial bit shifted into the Block-RAM then its terminal count becomes active and the address counter is enabled to point to the next address.



DeSerializer with a Twist and a Turn

For the examples on following slides the Bloc-RAM is reduced to it's important, tor the application, pins.

The first example adds the dual port capability of the Block-RAM to the design. This way it becomes possible to serially shift data into the Block-RAM, fill the Block-RAM as memory and read the contents in parallel format from port B.

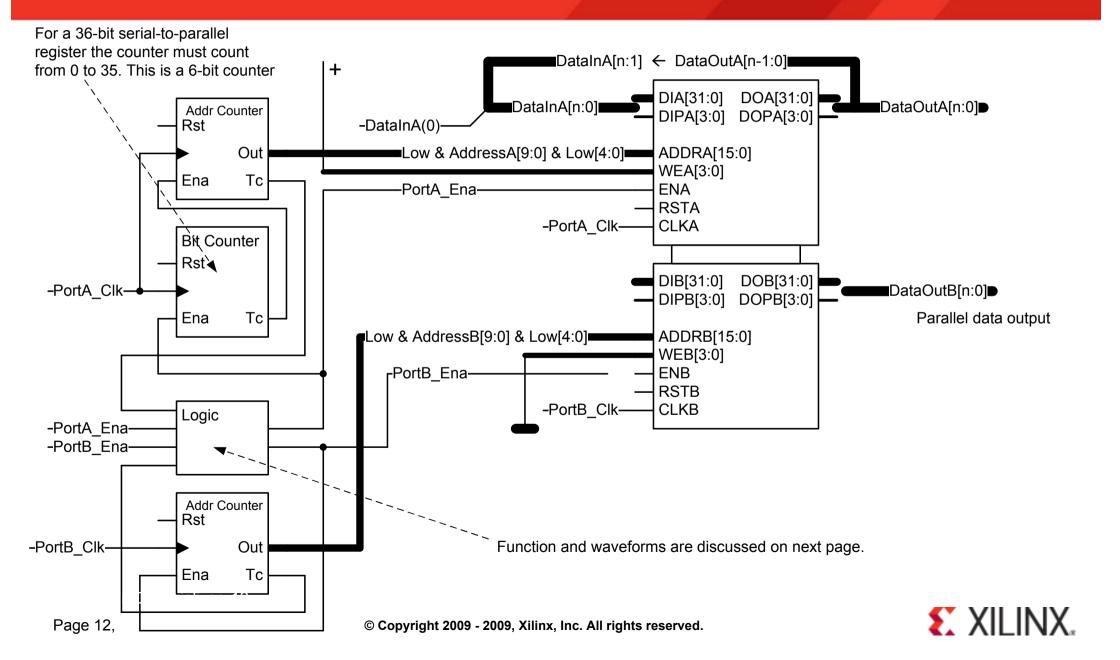
The memory functionality is transformed into a kind of FIFO.

The memory is filled via serial-to-parallel conversion and when a certain address or address range is reached the memory content must be read via port B.

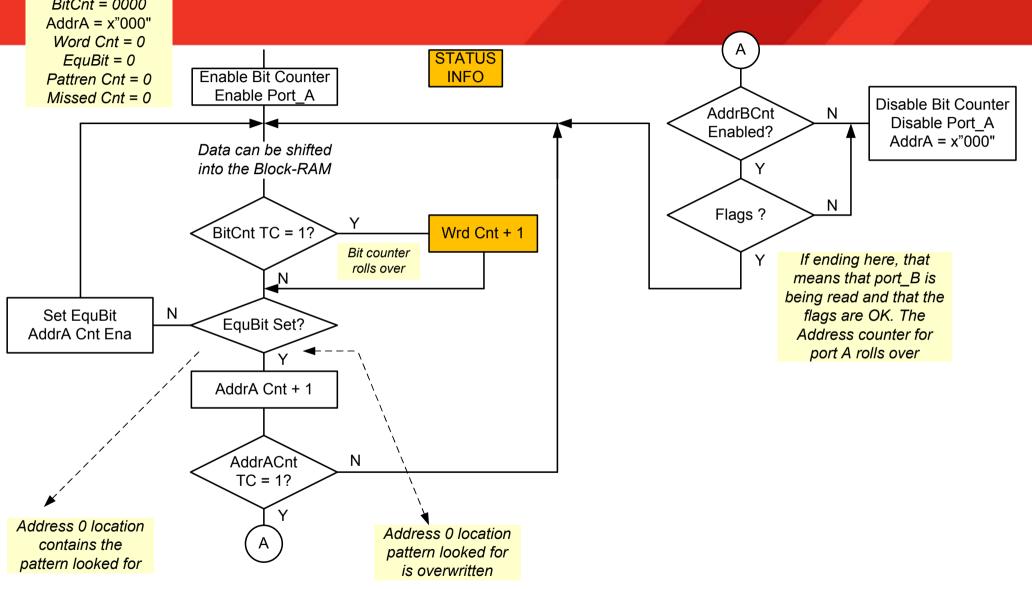
The second example adds something extra to the first example. A pattern comparator. This way it is possible to search a serial bit stream for a certain pattern before starting real reception of serial data. When the pattern is detected the address counter is incremented. Thus only good data is stored into the Block-RAM. As in the previous example it is read throught the B port of the Block-RAM



Example One

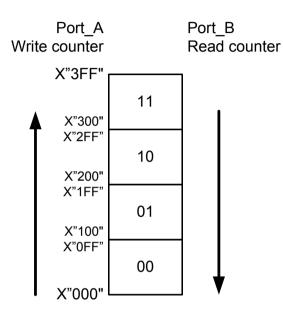


Example One (waveforms)





Example One (Block-RAM address counters & flags)



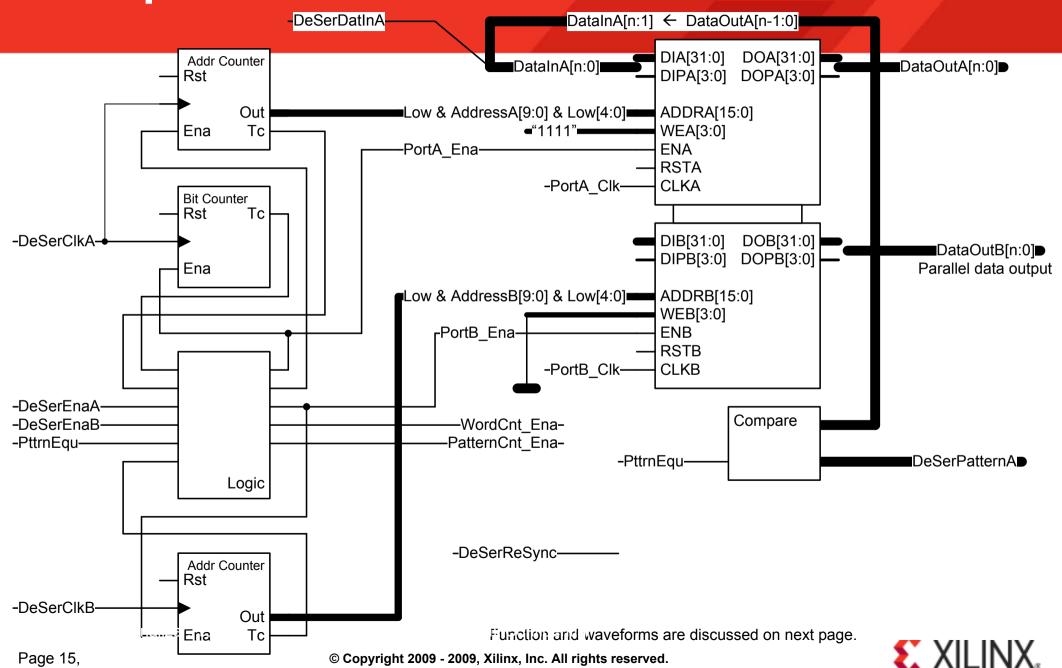
When the read pointer is right behind the write pointer, the FIFO is as good as empty. When the write pointer is right behind the read pointer, the FIFO is nearly full.

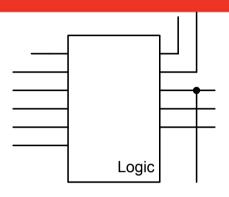
EMPTY		Lo-MID		Hi_M	Hi_MID		FULL	
11	11	00	11	01	11	10	11	
10	_ 10	11 _	10	00 _	10	01	_ 10	
01	= ¹⁰ 01	10 =	01	11	01	00	= 01	
00	00	01	00	10	00	11	00	
Wr	Rd	Wr	Rd	Wr	Rd	Wr	Rd	

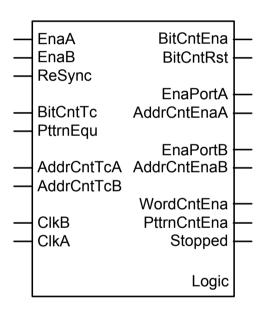
When both pointers are at the same counter value the buffer is considered to be empty, the flags indicate that the buffer is EMPTY.

Assume that one starts writing. Then the write pointer will shift one level up and the flags indicate that the buffer is filled up to the Lo-MID range. When the buffer continues to be written the flag will change to the Hi-MID indication and finally when no read operation is started the FULL flag will be raised.

Example Two









Example Two (waveforms)

