

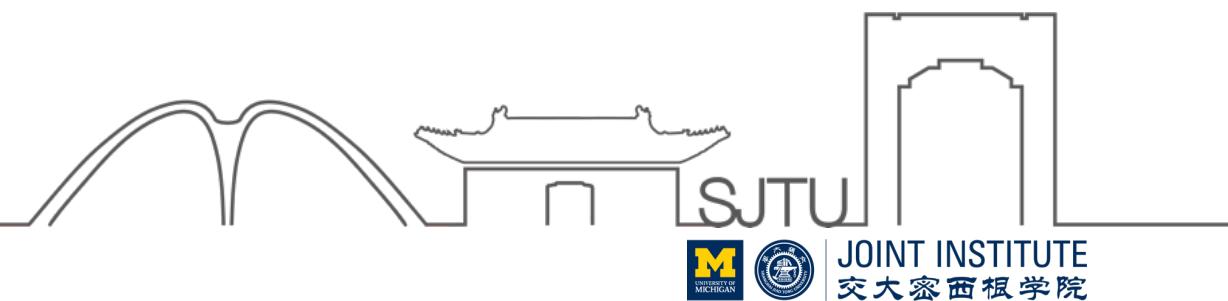


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ECE2700J SU23 RC2

Boolean Algebra & Logic Optimization

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# Boolean Algebra

## Terminology

Take  $F = A' \cdot B \cdot C + A \cdot (B + C')$  as an example:

- Variable: represent a value  
A, B, C
- Literal: appearance of a variable or its complement  
A', B, C, A, B, C' (Repetition of a literal is allowed)
- Product term & Sum term : a product of literals & a sum of literals  
 $A' \cdot B \cdot C$        $A + B$ ,
- Sum-of-products (SOP) form: a sum of product terms  
 $F = A' \cdot B \cdot C + A \cdot B + A \cdot C'$
- Product-of-sums (POS) form: a product of sum terms  
 $F = (A + B) \cdot (A + C) \cdot (B + C')$



# Boolean Algebra

## Theorems

- |                                    |                            |                       |
|------------------------------------|----------------------------|-----------------------|
| • (a) $x + 0 = x;$                 | (b) $x \cdot 0 = 0;$       | <b>(theorem 1)</b>    |
| • (a) $x + x' = 1;$                | (b) $x \cdot x' = 0;$      | <b>(theorem 2)</b>    |
| • (a) $x + x = x;$                 | (b) $x \cdot x = x;$       | <b>(theorem 3)</b>    |
| • (a) $x + 1 = 1;$                 | (b) $x \cdot 1 = x;$       | <b>(theorem 4)</b>    |
| • $(x')' = x;$                     |                            | <b>(involution)</b>   |
|                                    |                            |                       |
| • (a) $x + y = y + x;$             | (b) $xy = yx;$             | <b>(commutative)</b>  |
| • (a) $x + (y + z) = (x + y) + z;$ | (b) $x(yz) = (xy)z;$       | <b>(associative)</b>  |
| • (a) $x(y + z) = xy + xz;$        | (b) $x + yz = (x+y)(x+z);$ | <b>(distributive)</b> |
| • (a) $x + xy = x;$                | (b) $x(x + y) = x;$        | <b>(absorption)</b>   |
| • (a) $xy + xy' = x;$              | (b) $(x + y)(x + y') = x$  | <b>(theorem 5)</b>    |
| • (a) $x + x'y = x + y$            | (b) $x(x' + y) = xy$       | <b>(theorem 6)</b>    |



# Boolean Algebra

## Important Theorems

- DeMorgan's Law

$$(a) (x+y)' = x'y'$$

$$(b) (xy)' = x' + y'$$

$$(x+y)' = x'y'$$

$$(xyz)' = (xy)' + z'$$

$$(xy+z)' = x'+y'+z'$$

$$(D')' = D$$

- Exercise

$$\begin{aligned} & \text{SOP} \\ & ((AB' + C)D' + E)' = A'C'E' + B'C'E' \\ & = ((AB' + C) \cdot D') \cdot E' + DE' \\ & = [(AB' + C) + D] \cdot E' \\ & = [(AB')' \cdot C' + D] \cdot E' \\ & = [(A'B) \cdot C' + D] \cdot E' \end{aligned}$$



# Boolean Algebra

## Minterm and Maxterm

- Minterm is a product of  $n$  literals in which each literal appears exactly once in either true or complemented form, but not both.
  - Minterm is represented by  $m_i$
- Maxterm is a sum of  $n$  literals in which each literal appears exactly once in either true or complemented form, but not both.
  - Maxterm is represented by  $M_i$

A B C .

$\bar{A}B\bar{C}$

$\bar{A}\cdot\bar{B}\cdot\bar{C}$  .



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## Minterm and Maxterm

| x<br>y<br>z | Minterms |             | Maxterms   |             |
|-------------|----------|-------------|------------|-------------|
|             | Term     | Designation | Term       | Designation |
| 0 0 0       | $x'y'z'$ | $m_0$       | $x+y+z$    | $M_0$       |
| 0 0 1       | $x'y'z$  | $m_1$       | $x+y+z'$   | $M_1$       |
| 0 1 0       | $x'yz'$  | $m_2$       | $x+y'+z$   | $M_2$       |
| 0 1 1       | $x'yz$   | $m_3$       | $x+y'+z'$  | $M_3$       |
| 1 0 0       | $xy'z'$  | $m_4$       | $x'+y+z$   | $M_4$       |
| 1 0 1       | $xy'z$   | $m_5$       | $x'+y+z'$  | $M_5$       |
| 1 1 0       | $xyz'$   | $m_6$       | $x'+y'+z$  | $M_6$       |
| 1 1 1       | $xyz$    | $m_7$       | $x'+y'+z'$ | $M_7$       |



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## Minterm in Truth Table

|       | x    | y    | z    | F      |       |
|-------|------|------|------|--------|-------|
|       | con1 | con2 | con3 | result |       |
| 0 0 0 | 0    | 0    | 0    | 0      | $m_0$ |
| 0 0 1 | 0    | 0    | 1    | 1      | $m_1$ |
| 0 1 0 | 0    | 1    | 0    | 0      | $m_2$ |
| 0 1 1 | 0    | 1    | 1    | 1      | $m_3$ |
| 1 0 0 | 1    | 0    | 0    | 1      | $m_4$ |
| 1 0 1 | 1    | 0    | 1    | 1      | $m_5$ |
| 1 1 0 | 1    | 1    | 0    | 0      | $m_6$ |
| 1 1 1 | 1    | 1    | 1    | 0      | $m_7$ |

$$F \Leftrightarrow$$

$$x'y'z$$

$$\begin{aligned}
 F &= \cancel{x'y'z} + \cancel{x'yz} + \cancel{xy'z'} + \cancel{xy'z} \\
 &= m_1 + m_3 + m_4 + m_5 \\
 &= \Sigma m(1, 3, 4, 5)
 \end{aligned}$$



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## Exercise

- Q1. • Find minterm logic equation from these truth table **Sum**.

| x | y | z | F       |
|---|---|---|---------|
| 0 | 0 | 0 | 1 $m_0$ |
| 0 | 0 | 1 | 0 $m_1$ |
| 0 | 1 | 0 | 0 $m_2$ |
| 0 | 1 | 1 | 1 $m_3$ |
| 1 | 0 | 0 | 0 $m_4$ |
| 1 | 0 | 1 | 1 $m_5$ |
| 1 | 1 | 0 | 1 $m_6$ |
| 1 | 1 | 1 | 0 $m_7$ |

$$\begin{aligned} F &= \\ &\bar{x}\bar{y}\bar{z} \\ &+ x'\bar{y}\bar{z} \\ &+ \bar{x}y\bar{z}' \\ &+ x\bar{y}z' \\ &+ xy\bar{z}. \\ \therefore F &= m_0 + m_3 + m_5 + m_6. \\ \therefore F &= \sum m(0, 3, 5, 6) \end{aligned}$$

| w | x | y | z | F |          |
|---|---|---|---|---|----------|
| 0 | 0 | 0 | 0 | 1 | $m_0$    |
| 0 | 0 | 0 | 1 | 0 | $m_1$    |
| 0 | 0 | 1 | 0 | 0 | $m_2$    |
| 0 | 0 | 1 | 1 | 1 | $m_3$    |
| 0 | 1 | 0 | 0 | 0 | $m_4$    |
| 0 | 1 | 0 | 1 | 0 | $m_5$    |
| 0 | 1 | 1 | 0 | 0 | $m_6$    |
| 0 | 1 | 1 | 1 | 1 | $m_7$    |
| 1 | 0 | 0 | 0 | 1 | $m_8$    |
| 1 | 0 | 0 | 1 | 0 | $m_9$    |
| 1 | 0 | 1 | 0 | 0 | $m_{10}$ |
| 1 | 0 | 1 | 1 | 0 | $m_{11}$ |
| 1 | 1 | 0 | 0 | 0 | $m_{12}$ |
| 1 | 1 | 0 | 1 | 0 | $m_{13}$ |
| 1 | 1 | 1 | 0 | 0 | $m_{14}$ |
| 1 | 1 | 1 | 1 | 1 | $m_{15}$ |

$$\begin{aligned} F &= m_0 + m_3 + m_5 + m_6 \\ &+ m_7. \end{aligned}$$



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$$\bar{x}' \cdot y \cdot \bar{z}'$$

## Incompletely Specified Functions

- In a circuit, some input conditions may never happen, then the output is not completely specified
- The corresponding output is designated as “x”, called **don't care**
- A don't care output could be either 0 or 1

- $F = \sum m(1, 3, 4) \text{ with } d(2, 5)$

| x | y | z | F |
|---|---|---|---|
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | X |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | X |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 0 |



# Boolean Algebra

# Incompletely Specified Functions

**Example:** Bell rings,  $F = 1$ . A B C D F

There is a bell with 3 switch A, B and C. Switch A is to make the bell ring, and switch B and C are for the fuses to make sure safety. We know that switch B and C can't be off together otherwise the system will be dangerous. If switch A is off, the bell F will make sound. If all 3 switches are on, energy will be wasted and the bell make sound to remind people not wasting so much energy. Complete the truth table and create logic equations describing the desired behavior for F.

| $(A'B'C + A'BC' + A'BC) \div A$ |   |   |   |
|---------------------------------|---|---|---|
| A                               | B | C | E |
| 0                               | 0 | 0 | 1 |
| 0                               | 0 | 1 | 1 |
| 0                               | 1 | 0 | 1 |
| 0                               | 1 | 1 | 1 |
| 1                               | 0 | 0 | 0 |
| 1                               | 0 | 1 | 0 |
| 1                               | 1 | 0 | 0 |
| 1                               | 1 | 1 | 1 |

$$\sum m(1,2,3,7) + d(0,4). \quad \begin{array}{l} \text{Set} \\ m_0 = 1 \\ m_4 = 0. \end{array}$$
$$= A' + ABC.$$