

# Bits, Bytes and Integers – Part 1

15-213/18-213/15-513: Introduction to Computer Systems  
2<sup>nd</sup> Lecture, Aug. 31, 2017

**Today's Instructor:**

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# 数据的表示和运算(chapter2,课件2-4)

- 分三个部分介绍
  - 非数值数据的表示、数据的存储
    - 数据宽度单位
    - 硬件特征：大端/小端、对齐存放
  - 数值数据的表示
    - 定点数的编码表示
    - 整数的表示
      - 无符号整数、带符号整数
    - 浮点数的表示
    - C语言程序的整数类型和浮点数类型

# 数据的表示和运算

- 分三个部分介绍（续）

- 数据的运算

- 按位运算和逻辑运算
    - 移位运算
    - 位扩展和位截断运算
    - 无符号和带符号整数的加减运算
    - 无符号和带符号整数的乘除运算
    - 变量与常数之间的乘除运算
    - 浮点数的加减乘除运算

围绕C语言中的运算，解释其在底层机器级的实现方法

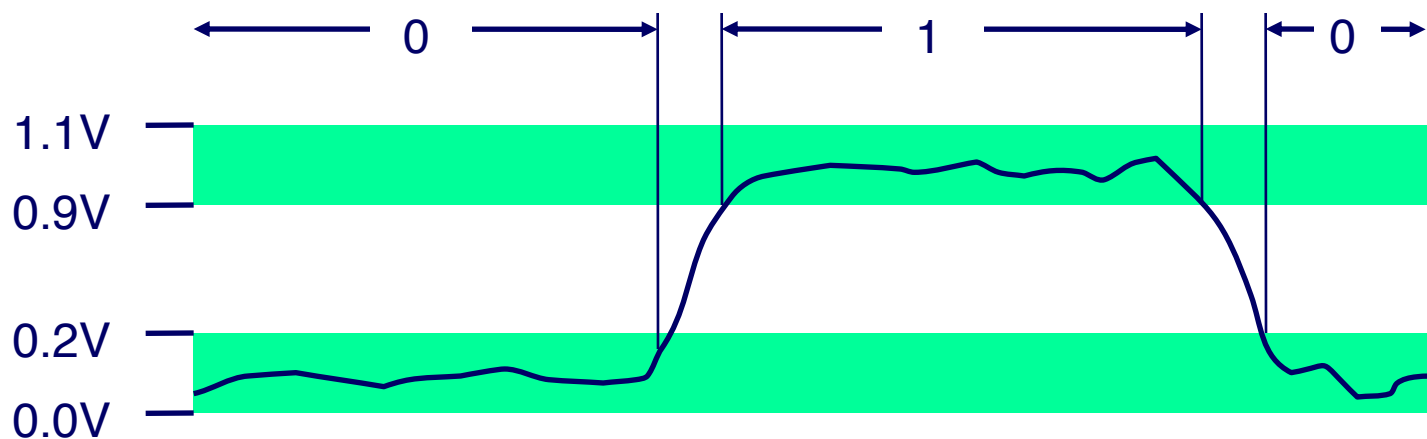
从高级语言程序中的表达式出发，用机器数在具体电路中的执行过程，来解释表达式的执行结果

# Today: Bits, Bytes, and Integers

- **Representing information as bits**
- **Bit-level manipulations**
- **Integers**
  - Representation: unsigned and signed
  - Conversion, casting
  - Expanding, truncating
  - Addition, negation, multiplication, shifting
  - Summary
- **Representations in memory, pointers, strings**

# Everything is bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
  - Computers determine what to do (instructions)
  - ... and represent and manipulate numbers, sets, strings, etc...
- **Why bits? Electronic Implementation**
  - Easy to store with bistable elements
  - Reliably transmitted on noisy and inaccurate wires



# For example, can count in binary

## ■ Base 2 Number Representation

- Represent  $15213_{10}$  as  $11101101101101_2$
- Represent  $1.20_{10}$  as  $1.0011001100110011[0011]..._2$
- Represent  $1.5213 \times 10^4$  as  $1.1101101101101_2 \times 2^{13}$

# Encoding Byte Values

## ■ Byte = 8 bits

- Binary  $00000000_2$  to  $11111111_2$
- Decimal:  $0_{10}$  to  $255_{10}$
- Hexadecimal  $00_{16}$  to  $FF_{16}$ 
  - Base 16 number representation
  - Use characters '0' to '9' and 'A' to 'F'
  - Write  $FA1D37B_{16}$  in C as
    - $0xFA1D37B$
    - $0xfa1d37b$

| Hex | Decimal | Binary |
|-----|---------|--------|
| 0   | 0       | 0000   |
| 1   | 1       | 0001   |
| 2   | 2       | 0010   |
| 3   | 3       | 0011   |
| 4   | 4       | 0100   |
| 5   | 5       | 0101   |
| 6   | 6       | 0110   |
| 7   | 7       | 0111   |
| 8   | 8       | 1000   |
| 9   | 9       | 1001   |
| A   | 10      | 1010   |
| B   | 11      | 1011   |
| C   | 12      | 1100   |
| D   | 13      | 1101   |
| E   | 14      | 1110   |
| F   | 15      | 1111   |

15213: 0011 1011 0110 1101  
           3      B      6      D

# 数据量的度量单位

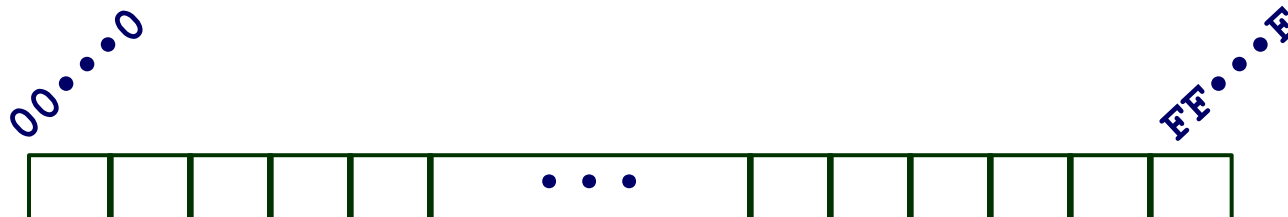
- 存储二进制信息时的度量单位要比字节或字大得多
- 容量经常使用的单位有：
  - “千字节” (KB),  $1\text{KB}=2^{10}\text{字节}=1024\text{B}$
  - “兆字节” (MB),  $1\text{MB}=2^{20}\text{字节}=1024\text{KB}$
  - “千兆字节” (GB),  $1\text{GB}=2^{30}\text{字节}=1024\text{MB}$
  - “兆兆字节” (TB),  $1\text{TB}=2^{40}\text{字节}=1024\text{GB}$
- 通信中的带宽使用的单位有：
  - “千比特/秒” (kb/s),  $1\text{kbps}=10^3\text{ b/s}=1000\text{ bps}$
  - “兆比特/秒” (Mb/s),  $1\text{Mbps}=10^6\text{ b/s}=1000\text{ kbps}$
  - “千兆比特/秒” (Gb/s),  $1\text{Gbps}=10^9\text{ b/s}=1000\text{ Mbps}$
  - “兆兆比特/秒” (Tb/s),  $1\text{Tbps}=10^{12}\text{ b/s}=1000\text{ Gbps}$

如果把b换成B，则表示字节而不是比特（位）

例如，10MBps表示 10兆字节/秒



# Byte-Oriented Memory Organization



## ■ Programs refer to data by address

- Conceptually, envision it as a very large array of bytes
  - In reality, it's not, but can think of it that way
- An address is like an index into that array
  - and, a pointer variable stores an address

## ■ Note: system provides private address spaces to each “process”

- Think of a process as a program being executed
- So, a program can clobber its own data, but not that of others

# 数据的基本宽度

- “字” 和 “字长” 的概念不同

- “字长” 指数据通路的宽度。

(数据通路指CPU内部数据流经的路径以及路径上的部件，主要是CPU内部进行数据运算、存储和传送的部件，这些部件的宽度基本上要一致，才能相互匹配。因此，“字长”等于CPU内部总线的宽度、运算器的位数、通用寄存器的宽度等。)

- “字” 表示被处理信息的单位，用来度量数据类型的宽度。

- 字和字长的宽度可以一样，也可不同。

例如，x86体系结构定义“字”的宽度为16位，但从386开始字长就是32位了。

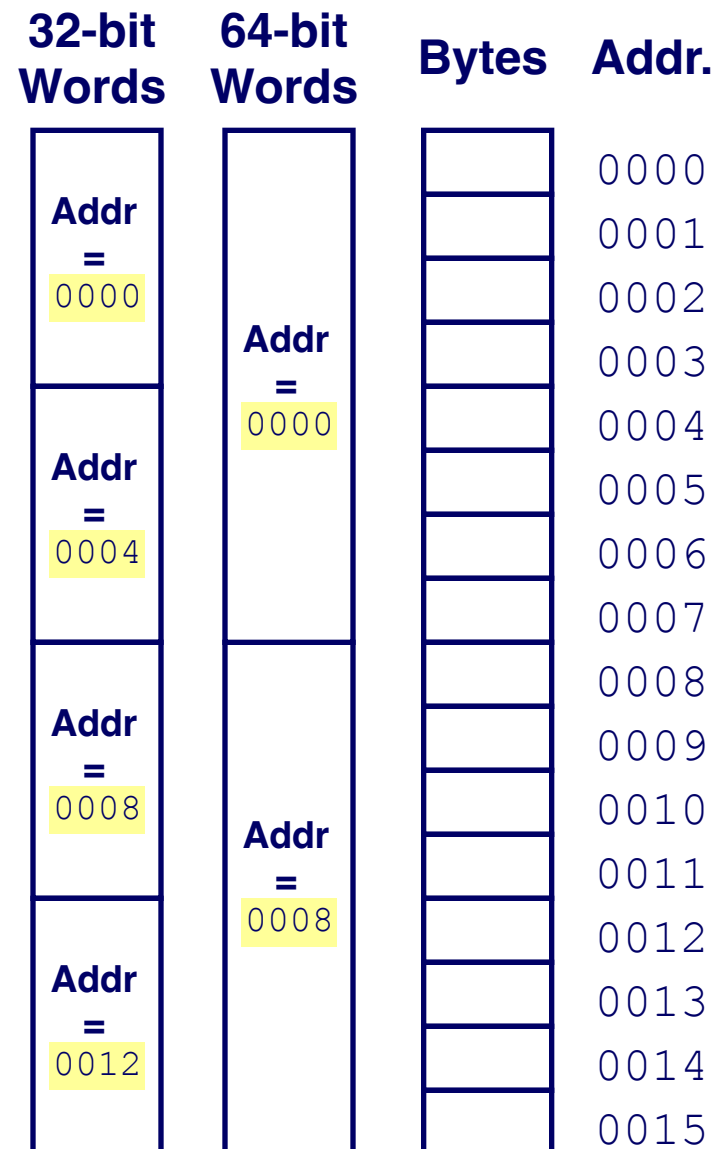
# Machine Words

- **Any given computer has a “Word Size”**
  - Nominal size of integer-valued data
    - and of addresses
  - Until recently, most machines used 32 bits (4 bytes) as word size
    - Limits addresses to 4GB ( $2^{32}$  bytes)
  - Increasingly, machines have 64-bit word size
    - Potentially, could have 18 EB (exabytes) of addressable memory
    - That's  $18.4 \times 10^{18}$
  - Machines still support multiple data formats
    - Fractions or multiples of word size
    - Always integral number of bytes

# Word-Oriented Memory Organization

## ■ Addresses Specify Byte Locations

- Address of first byte in word
- Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)



# Example Data Representations

| C Data Type          | Typical 32-bit | Typical 64-bit | x86-64 |
|----------------------|----------------|----------------|--------|
| <code>char</code>    | 1              | 1              | 1      |
| <code>short</code>   | 2              | 2              | 2      |
| <code>int</code>     | 4              | 4              | 4      |
| <code>long</code>    | 4              | 8              | 8      |
| <code>float</code>   | 4              | 4              | 4      |
| <code>double</code>  | 8              | 8              | 8      |
| <code>pointer</code> | 4              | 8              | 8      |

- So, how are the bytes within a multi-byte word ordered in memory?

# 数据的存储和排列顺序

- 80年代开始，几乎所有机器都用**字节编址**
- ISA设计时要考虑的两个问题：
  - 如何根据一个字节地址取到一个32位的字？ - **字的存放问题**
  - 一个字能否存放在任何字节边界？ - **字的边界对齐问题**

例如，若  $\text{int } i = -65535$ ，存放在内存100号单元（即占100# ~ 103#），则用“取数”指令访问100号单元取出  $i$  时，必须清楚  $i$  的4个字节是如何存放的。

$$65535 = 2^{16} - 1$$

$$[-65535]_{\text{补}} = \text{FFFF0001H}$$

Word:

|            |           |           |                  |
|------------|-----------|-----------|------------------|
| FF<br>103  | FF<br>102 | 00<br>101 | 01<br><b>100</b> |
| msb        |           |           | lsb              |
| <b>100</b> | 101       | 102       | 103              |

little endian word 100#

big endian word 100#

**大端方式 (Big Endian) : MSB所在的地址是数的地址**

e.g. IBM 360/370, Motorola 68k, MIPS, Sparc, HP PA

**小端方式 (Little Endian) : LSB所在的地址是数的地址**

e.g. Intel 80x86, DEC VAX

有些机器两种方式都支持，可通过特定控制位来设定采用哪种方式。

# Representing Integers

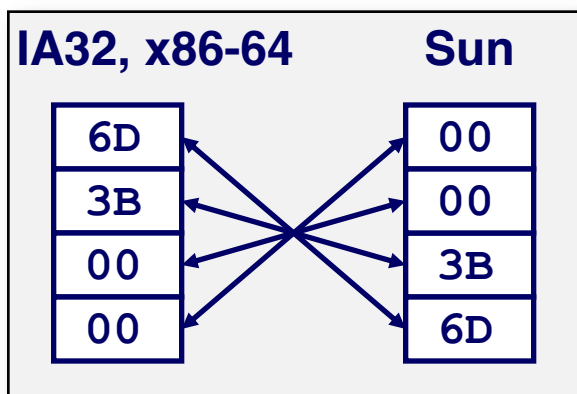
Decimal: 15213

Binary: 0011 1011 0110 1101

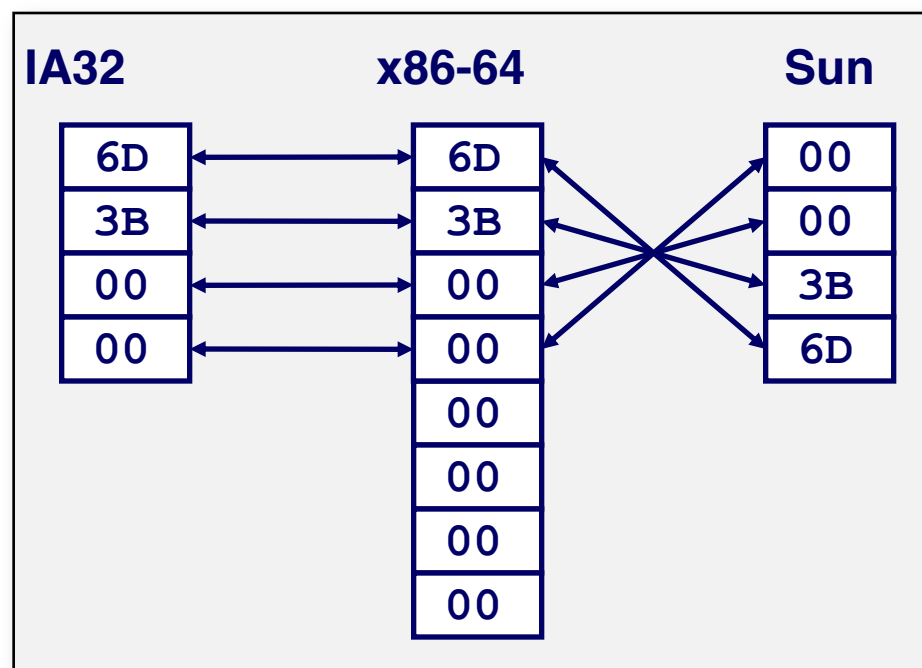
Hex: 3 B 6 D

`int A = 15213;`

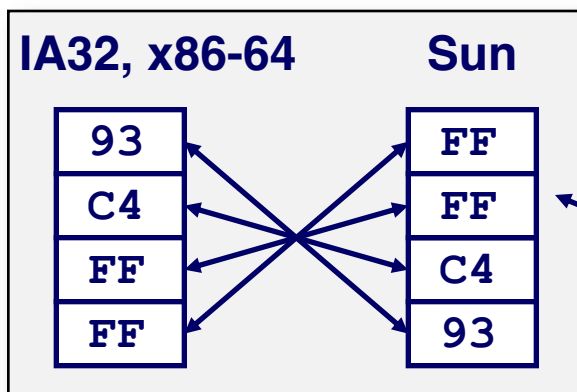
Increasing addresses  
↓



`long int C = 15213;`



`int B = -15213;`



Two's complement representation

# BIG Endian versus Little Endian

Ex3: Memory layout of a instruction located in 1000

假定小端机器中指令: `mov AX, 0x12345(BX)`

其中操作码mov为40H, 寄存器AX和BX的编号分别为0001B和0010B, 立即数占32位, 则存放顺序为:



若在大端机器上, 则存放顺序如何?



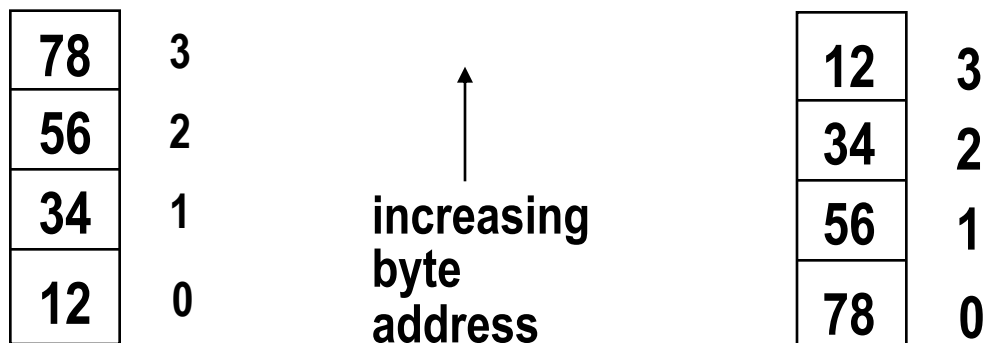
|    |      |    |
|----|------|----|
| 00 | 1005 | 45 |
| 01 | 1004 | 23 |
| 23 | 1003 | 01 |
| 45 | 1002 | 00 |
| 12 | 1001 | 12 |
| 40 | 1000 | 40 |

地址

只需要考虑指令中立即数的顺序!



# Byte Swap Problem (字节交换问题)



Big Endian

Little Endian

上述存放在0号单元的数据（字）是什么？**12345678H?** **78563412H?**

存放方式不同的机器间程序移植或数据通信时，会发生什么问题？

- ◆ 每个系统内部是一致的，但在系统间通信时可能会发生问题！
- ◆ 因为顺序不同，需要进行顺序转换

音、视频和图像等文件格式或处理程序都涉及到字节顺序问题

**ex. Little endian: GIF, PC Paintbrush, Microsoft RTF, etc**

**Big endian: Adobe Photoshop, JPEG, MacPaint, etc**

```

C:\WINDOWS\system32\cmd.exe - debug p160.exe

18E4:0130  31 39 39 34 20 28 F0 46-00 20 99 3B 20 30 01 20 1994 <.F. .; 0.
18E4:0140  31 39 39 35 20 08 24 5B-00 20 88 45 20 4F 01 20 1995 $.[. .E 0.
18E4:0150  B8 D6 18 8E D8 B8 E4 18-8E C0 BB 00 00 BF 00 00 .....
18E4:0160  BE 00 00 B9 15 00 8B 05-26 89 07 8B 45 02 26 89 .....&...E.&.
18E4:0170  47 02 8B 84 A8 00 26 89-47 0A 8B 45 54 26 89 47 G.....&.G..ET&.G
-d es:0
18E4:0000  31 39 37 35 20 10 00 00-00 20 03 00 20 05 00 20 1975 .... .. ..
18E4:0010  31 39 37 36 20 16 00 00-00 20 07 00 20 03 00 20 1976 .... .. ..
18E4:0020  31 39 37 37 20 7E 01 00-00 20 09 00 20 2A 00 20 1977 ~... .. *.
18E4:0030  31 39 37 38 20 4C 05 00-00 20 0D 00 20 68 00 20 1978 L... .. h.
18E4:0040  31 39 37 39 20 56 09 00-00 20 1C 00 20 55 00 20 1979 U... .. U.
18E4:0050  31 39 38 30 20 40 1F 00-00 20 26 00 20 D2 00 20 1980 e... &. ..
18E4:0060  31 39 38 31 20 40 06 00-00 20 82 00 20 0C 00 20 1981 e... .. ..
18E4:0070  31 39 38 32 20 A6 5F 00-00 20 DC 00 20 6F 00 20 1982 _... .. o.
-
-d
18E4:0080  31 39 38 33 20 91 C3 00-00 20 DC 01 20 69 00 20 1983 .... .. i.
18E4:0090  31 39 38 34 20 C7 7C 01-00 20 0A 03 20 7D 00 20 1984 .!.. .. >.
18E4:00A0  31 39 38 35 20 81 24 02-00 20 E9 03 20 8C 00 20 1985 $. .. ..
18E4:00B0  31 39 38 36 20 8A 03 03-00 20 A2 05 20 88 00 20 1986 .... .. ..
18E4:00C0  31 39 38 37 20 7C 47 05-00 20 D2 08 20 99 00 20 1987 !G.. .. ..
18E4:00D0  31 39 38 38 20 EB 03 09-00 20 E9 0A 20 D3 00 20 1988 .... .. ..
18E4:00E0  31 39 38 39 20 CA 42 0C-00 20 C5 0F 20 C7 00 20 1989 ..B... .. ..
18E4:00F0  31 39 39 30 20 18 0D 12-00 20 03 16 20 D1 00 20 1990 .. .. ..

```

# Alignment(对齐)

Alignment: 要求数据的地址是相应的边界地址

- 目前机器字长一般为32位或64位，而存储器地址按字节编址
- 指令系统支持对字节、半字、字及双字的运算，也有位处理指令
- 各种不同长度的数据存放时，有两种处理方式：

- 按边界对齐 （假定存储字的宽度为32位，按字节编址）

- 字地址：4的倍数（低两位为0）

- 半字地址：2的倍数（低位为0）

- 字节地址：任意

每4个字节可同时读写



- 不按边界对齐

坏处：可能会增加访存次数！

（学了存储器组织后会更明白！）

# Alignment(对齐)

如: `int i, short k, double x, char c, short j,.....`

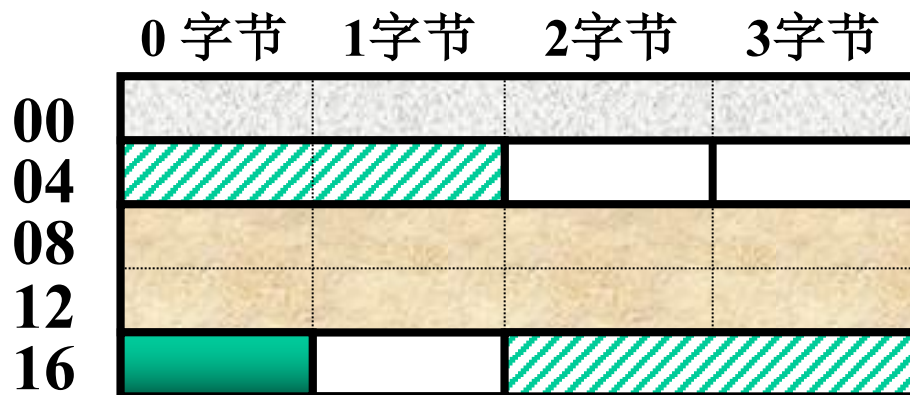
存储器按字节  
编址

每次只能读写  
某个字地址开  
始的4个单元中  
连续的1个、2  
个、3个或4个  
字节

按边界对齐

x: 2个周期

j: 1个周期



则: `&i=0; &k=4; &x=8; &c=16; &j=18;.....`

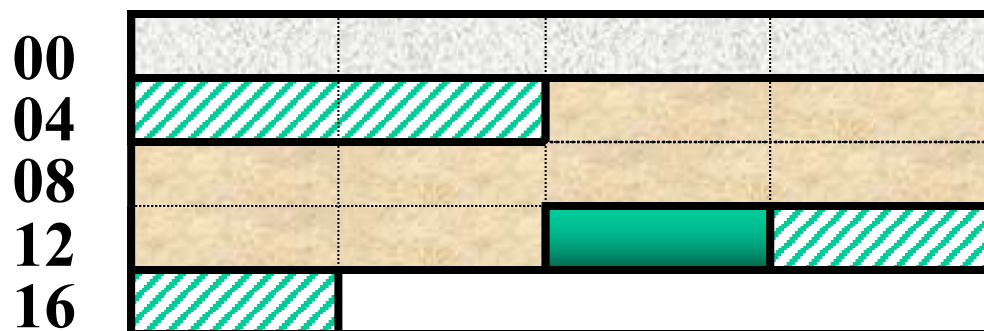
字节0 字节1 字节2 字节3

虽节省了空间，  
但增加了访存次  
数！

边界不对齐

x: 3个周期

j: 2个周期



则: `&i=0; &k=4; &x=6; &c=14; &j=15;.....`

需要权衡，目前  
来看，浪费一点  
存储空间没有关  
系！

# Alignment(对齐) 举例

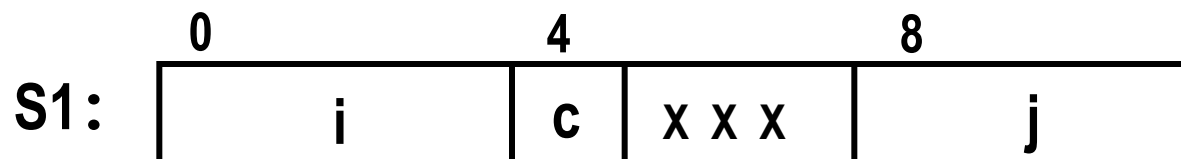
例如，考虑下列两个结构声明：

```
struct S1 {
    int    i;
    char   c;
    int    j;
};
```

```
struct S2 {
    int    i;
    int    j;
    char   c;
};
```

在要求对齐的情况下，哪种结构声明更好？

**S2比S1好**

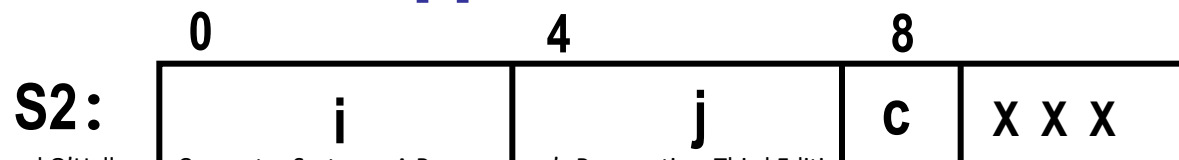


**需要12个字节**



**只需要9个字节**

对于 “struct S2 d[4]”，只分配9个字节能否满足对齐要求？ **不能！**



**也需要12个字节**

# Today: Bits, Bytes, and Integers

- Representing information as bits
- **Bit-level manipulations**
- **Integers**
  - Representation: unsigned and signed
  - Conversion, casting
  - Expanding, truncating
  - Addition, negation, multiplication, shifting
  - Summary
- Representations in memory, pointers, strings

# 逻辑数据的编码表示

- 表示

- 用一位表示。例如，真：1 / 假：0
- N位二进制数可表示N个逻辑数据，或一个位串

- 运算

- 按位进行
- 如:按位与 / 按位或 / 逻辑左移 / 逻辑右移 等

- 识别

- 逻辑数据和数值数据在形式上并无差别，也是一串0/1序列，机器靠指令来识别。

- 位串

- 用来表示若干个状态位或控制位（OS中使用较多）

例如，x86的标志寄存器含义如下：

|  |  |  |  |    |    |    |    |    |    |  |    |  |    |  |    |
|--|--|--|--|----|----|----|----|----|----|--|----|--|----|--|----|
|  |  |  |  | OF | DF | IF | TF | SF | ZF |  | AF |  | PF |  | CF |
|--|--|--|--|----|----|----|----|----|----|--|----|--|----|--|----|

# Boolean Algebra

## ■ Developed by George Boole in 19th Century

- Algebraic representation of logic
  - Encode “True” as 1 and “False” as 0

And

- $A \& B = 1$  when both  $A=1$  and  $B=1$

| $\&$ | 0 | 1 |
|------|---|---|
| 0    | 0 | 0 |
| 1    | 0 | 1 |

Or

- $A | B = 1$  when either  $A=1$  or  $B=1$

| $ $ | 0 | 1 |
|-----|---|---|
| 0   | 0 | 1 |
| 1   | 1 | 1 |

Not

- $\sim A = 1$  when  $A=0$

| $\sim$ |   |
|--------|---|
| 0      | 1 |
| 1      | 0 |

Exclusive-Or (Xor)

- $A \wedge B = 1$  when either  $A=1$  or  $B=1$ , but not both

| $\wedge$ | 0 | 1 |
|----------|---|---|
| 0        | 0 | 1 |
| 1        | 1 | 0 |



# General Boolean Algebras

## ■ Operate on Bit Vectors

- Operations applied bitwise

|                   |                   |                   |                   |
|-------------------|-------------------|-------------------|-------------------|
| 01101001          | 01101001          | 01101001          |                   |
| & 01010101        | 01010101          | ^ 01010101        | ~ 01010101        |
| <u>          </u> | <u>          </u> | <u>          </u> | <u>          </u> |
| 01000001          | 01111101          | 00111100          | 10101010          |

## ■ All of the Properties of Boolean Algebra Apply

# Example: Representing & Manipulating Sets

## ■ Representation

- Width  $w$  bit vector represents subsets of  $\{0, \dots, w-1\}$
- $a_j = 1$  if  $j \in A$

- 01101001       $\{0, 3, 5, 6\}$

- 76543210

- 01010101       $\{0, 2, 4, 6\}$

- 76543210

## ■ Operations

- |                             |          |                        |
|-----------------------------|----------|------------------------|
| ▪ &    Intersection         | 01000001 | $\{0, 6\}$             |
| ▪      Union                | 01111101 | $\{0, 2, 3, 4, 5, 6\}$ |
| ▪ ^    Symmetric difference | 00111100 | $\{2, 3, 4, 5\}$       |
| ▪ ~    Complement           | 10101010 | $\{1, 3, 5, 7\}$       |

# Bit-Level Operations in C

## ■ Operations $\&$ , $|$ , $\sim$ , $\wedge$ Available in C

- Apply to any “integral” data type
  - long, int, short, char, unsigned
- View arguments as bit vectors
- Arguments applied bit-wise

## ■ Examples (Char data type)

- $\sim 0x41 \rightarrow$
- $\sim 0x00 \rightarrow$
- $0x69 \& 0x55 \rightarrow$
- $0x69 | 0x55 \rightarrow$

| Hex | Decimal | Binary |
|-----|---------|--------|
| 0   | 0       | 0000   |
| 1   | 1       | 0001   |
| 2   | 2       | 0010   |
| 3   | 3       | 0011   |
| 4   | 4       | 0100   |
| 5   | 5       | 0101   |
| 6   | 6       | 0110   |
| 7   | 7       | 0111   |
| 8   | 8       | 1000   |
| 9   | 9       | 1001   |
| A   | 10      | 1010   |
| B   | 11      | 1011   |
| C   | 12      | 1100   |
| D   | 13      | 1101   |
| E   | 14      | 1110   |
| F   | 15      | 1111   |

# Bit-Level Operations in C

## ■ Operations $\&$ , $|$ , $\sim$ , $\wedge$ Available in C

- Apply to any “integral” data type
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- View arguments as bit vectors
- Arguments applied bit-wise

## ■ Examples (Char data type)

- $\sim 0x41 \rightarrow 0xBE$ 
  - $\sim 0100\ 0001_2 \rightarrow 1011\ 1110_2$
- $\sim 0x00 \rightarrow 0xFF$ 
  - $\sim 0000\ 0000_2 \rightarrow 1111\ 1111_2$
- $0x69 \& 0x55 \rightarrow 0x41$ 
  - $0110\ 1001_2 \& 0101\ 0101_2 \rightarrow 0100\ 0001_2$
- $0x69 | 0x55 \rightarrow 0x7D$ 
  - $0110\ 1001_2 | 0101\ 0101_2 \rightarrow 0111\ 1101_2$

| Hex | Decimal | Binary |
|-----|---------|--------|
| 0   | 0       | 0000   |
| 1   | 1       | 0001   |
| 2   | 2       | 0010   |
| 3   | 3       | 0011   |
| 4   | 4       | 0100   |
| 5   | 5       | 0101   |
| 6   | 6       | 0110   |
| 7   | 7       | 0111   |
| 8   | 8       | 1000   |
| 9   | 9       | 1001   |
| A   | 10      | 1010   |
| B   | 11      | 1011   |
| C   | 12      | 1100   |
| D   | 13      | 1101   |
| E   | 14      | 1110   |
| F   | 15      | 1111   |

# Contrast: Logic Operations in C

## ■ Contrast to Bit-Level Operators

### ■ Logic Operations: `&`, `||`, `!`

- View 0 as “False”

- Anything nonzero

- Always

- Early

## ■ Example

- `!0x41`

- `!0x00` →

- `!!0x41` → `0x01`

- `0x69 && 0x55` → `0x01`

- `0x69 || 0x55` → `0x01`

- `p && *p` (avoids null pointer access)

**Watch out for `&&` vs. `&` (and `||` vs. `|`)...  
one of the more common oopsies in  
C programming**

# Shift Operations

## ■ Left Shift: $x \ll y$

- Shift bit-vector  $x$  left  $y$  positions
  - Throw away extra bits on left
  - Fill with 0's on right

## ■ Right Shift: $x \gg y$

- Shift bit-vector  $x$  right  $y$  positions
  - Throw away extra bits on right
- Logical shift**
  - Fill with 0's on left
- Arithmetic shift**
  - Replicate most significant bit on left

## ■ Undefined Behavior

- Shift amount  $< 0$  or  $\geq$  word size

|                |          |
|----------------|----------|
| Argument $x$   | 01100010 |
| $\ll 3$        | 00010000 |
| Log. $\gg 2$   | 00011000 |
| Arith. $\gg 2$ | 00011000 |

|                |          |
|----------------|----------|
| Argument $x$   | 10100010 |
| $\ll 3$        | 00010000 |
| Log. $\gg 2$   | 00101000 |
| Arith. $\gg 2$ | 11101000 |

# 西文字符的编码表示

## • 特点

- 是一种拼音文字，用有限几个字母可拼写出所有单词
- 只对有限个字母和数学符号、标点符号等辅助字符编码
- 所有字符总数不超过256个，使用7或8个二进位可表示

## • 表示（常用编码为7位ASCII码）

- 十进制数字：0/1/2.../9（30H）
  - 英文字母：A/B/.../Z/a/b/.../z（41H、61H）
  - 专用符号：+/-/%/\*/&/.....
  - 控制字符（不可打印或显示）（回车：0DH；换行：0AH）
- } 必须熟悉对应的ASCII码！

## • 操作

- 字符串操作，如：传送/比较 等

| ASCII码 | 控制字符 | ASCII码 | 字符      | ASCII码 | 字符 | ASCII码 | 字符  |
|--------|------|--------|---------|--------|----|--------|-----|
| 0      | NUL  | 32     | (space) | 64     | @  | 96     | `   |
| 1      | SOH  | 33     | !       | 65     | A  | 97     | a   |
| 2      | STX  | 34     | "       | 66     | B  | 98     | b   |
| 3      | ETX  | 35     | #       | 67     | C  | 99     | c   |
| 4      | EOT  | 36     | %       | 68     | D  | 100    | d   |
| 5      | ENQ  | 37     | %       | 69     | E  | 101    | e   |
| 6      | ACK  | 38     | &       | 70     | F  | 102    | f   |
| 7      | BEL  | 39     | '       | 71     | G  | 103    | g   |
| 8      | BS   | 40     | (       | 72     | H  | 104    | h   |
| 9      | HT   | 41     | )       | 73     | I  | 105    | i   |
| 10     | LF   | 42     | *       | 74     | J  | 106    | j   |
| 11     | VT   | 43     | +       | 75     | K  | 107    | k   |
| 12     | FF   | 44     | ,       | 76     | L  | 108    | l   |
| 13     | CR   | 45     | -       | 77     | M  | 109    | m   |
| 14     | SO   | 46     | .       | 78     | N  | 110    | n   |
| 15     | SI   | 47     | /       | 79     | O  | 111    | o   |
| 16     | DLE  | 48     | 0       | 80     | P  | 112    | p   |
| 17     | DC1  | 49     | 1       | 81     | Q  | 113    | q   |
| 18     | DC2  | 50     | 2       | 82     | R  | 114    | r   |
| 19     | DC3  | 51     | 3       | 83     | S  | 115    | s   |
| 20     | DC4  | 52     | 4       | 84     | T  | 116    | t   |
| 21     | NAK  | 53     | 5       | 85     | U  | 117    | u   |
| 22     | SYN  | 54     | 6       | 86     | V  | 118    | v   |
| 23     | ETB  | 55     | 7       | 87     | W  | 119    | w   |
| 24     | CAN  | 56     | 8       | 88     | X  | 120    | x   |
| 25     | EM   | 57     | 9       | 89     | Y  | 121    | y   |
| 26     | SUB  | 58     | :       | 90     | Z  | 122    | z   |
| 27     | ESC  | 59     | ;       | 91     | [  | 123    | {   |
| 28     | FS   | 60     | <       | 92     | \  | 124    |     |
| 29     | GS   | 61     | =       | 93     | ]  | 125    | }   |
| 30     | RS   | 62     | >       | 94     | ^  | 126    | ~   |
| 31     | US   | 63     | ?       | 95     | -  | 127    | DEL |



# 汉字及国际字符的编码表示

## ■ 特点

- 汉字是表意文字，一个字就是一个方块图形。
- 汉字数量巨大，总数超过6万字，给汉字在计算机内部的表示、汉字的传输与交换、汉字的输入和输出等带来了一系列问题。

## ■ 编码形式

- 有以下几种汉字代码：
  - 输入码：对汉字用相应按键进行编码表示，用于输入
  - 内码：用于在系统中进行存储、查找、传送等处理
  - 字模点阵或轮廓描述：描述汉字字模点阵或轮廓，用于显示/打印

# Today: Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- **Integers**
  - **Representation: unsigned and signed**
  - Conversion, casting
  - Expanding, truncating
  - Addition, negation, multiplication, shifting
  - Summary
- Representations in memory, pointers, strings
- Summary

# Encoding Integers

## Unsigned

$$D_4 W^*Z+ = \sum_{K=2}^{Y-3} Z_K \cdot 4^K$$

## Two's Complement

$$D_4 V^*Z+ = -Z_{Y-3} \cdot 4^{Y-3} + \sum_{K=2}^{Y-4} Z_K \cdot 4^K$$

```
short int x = 15213;
short int y = -15213;
```

Sign Bit



## ■ C short 2 bytes long

|          | Decimal | Hex   | Binary            |
|----------|---------|-------|-------------------|
| <b>x</b> | 15213   | 3B 6D | 00111011 01101101 |
| <b>y</b> | -15213  | C4 93 | 11000100 10010011 |

## ■ Sign Bit

- For 2's complement, most significant bit indicates sign
  - 0 for nonnegative
  - 1 for negative

# Two-complement: Simple Example

|      | -16 | 8 | 4 | 2 | 1 |            |
|------|-----|---|---|---|---|------------|
| 10 = | 0   | 1 | 0 | 1 | 0 | $8+2 = 10$ |

|       | -16 | 8 | 4 | 2 | 1 |                 |
|-------|-----|---|---|---|---|-----------------|
| -10 = | 1   | 0 | 1 | 1 | 0 | $-16+4+2 = -10$ |

# Two-complement Encoding Example (Cont.)

**x =**            15213: 00111011 01101101  
**y =**            -15213: 11000100 10010011

| Weight     | 15213        |      | -15213        |        |
|------------|--------------|------|---------------|--------|
| 1          | 1            | 1    | 1             | 1      |
| 2          | 0            | 0    | 1             | 2      |
| 4          | 1            | 4    | 0             | 0      |
| 8          | 1            | 8    | 0             | 0      |
| 16         | 0            | 0    | 1             | 16     |
| 32         | 1            | 32   | 0             | 0      |
| 64         | 1            | 64   | 0             | 0      |
| 128        | 0            | 0    | 1             | 128    |
| 256        | 1            | 256  | 0             | 0      |
| 512        | 1            | 512  | 0             | 0      |
| 1024       | 0            | 0    | 1             | 1024   |
| 2048       | 1            | 2048 | 0             | 0      |
| 4096       | 1            | 4096 | 0             | 0      |
| 8192       | 1            | 8192 | 0             | 0      |
| 16384      | 0            | 0    | 1             | 16384  |
| -32768     | 0            | 0    | 1             | -32768 |
| <b>Sum</b> | <b>15213</b> |      | <b>-15213</b> |        |

# Numeric Ranges

## ■ Unsigned Values

- $UMin = 0$   
000...0
- $UMax = 2^w - 1$   
111...1

## ■ Two's Complement Values

- $TMin = -2^{w-1}$   
100...0
- $TMax = 2^{w-1} - 1$   
011...1
- Minus 1  
111...1

### Values for $W = 16$

|             | Decimal       | Hex          | Binary                   |
|-------------|---------------|--------------|--------------------------|
| <b>UMax</b> | <b>65535</b>  | <b>FF FF</b> | <b>11111111 11111111</b> |
| <b>TMax</b> | <b>32767</b>  | <b>7F FF</b> | <b>01111111 11111111</b> |
| <b>TMin</b> | <b>-32768</b> | <b>80 00</b> | <b>10000000 00000000</b> |
| <b>-1</b>   | <b>-1</b>     | <b>FF FF</b> | <b>11111111 11111111</b> |
| <b>0</b>    | <b>0</b>      | <b>00 00</b> | <b>00000000 00000000</b> |

"

# Values for Different Word Sizes

|      | W    |         |                |                            |
|------|------|---------|----------------|----------------------------|
|      | 8    | 16      | 32             | 64                         |
| UMax | 255  | 65,535  | 4,294,967,295  | 18,446,744,073,709,551,615 |
| TMax | 127  | 32,767  | 2,147,483,647  | 9,223,372,036,854,775,807  |
| TMin | -128 | -32,768 | -2,147,483,648 | -9,223,372,036,854,775,808 |

## ■ Observations

- $|TMin| = TMax + 1$ 
  - Asymmetric range
- $UMax = 2 * TMax + 1$

## ■ C Programming

- `#include <limits.h>`
- Declares constants, e.g.,
  - `ULONG_MAX`
  - `LONG_MAX`
  - `LONG_MIN`
- Values platform specific

# Unsigned & Signed Numeric Values

| $X$  | $B2U(X)$ | $B2T(X)$ |
|------|----------|----------|
| 0000 | 0        | 0        |
| 0001 | 1        | 1        |
| 0010 | 2        | 2        |
| 0011 | 3        | 3        |
| 0100 | 4        | 4        |
| 0101 | 5        | 5        |
| 0110 | 6        | 6        |
| 0111 | 7        | 7        |
| 1000 | 8        | -8       |
| 1001 | 9        | -7       |
| 1010 | 10       | -6       |
| 1011 | 11       | -5       |
| 1100 | 12       | -4       |
| 1101 | 13       | -3       |
| 1110 | 14       | -2       |
| 1111 | 15       | -1       |

## ■ Equivalence

- Same encodings for nonnegative values

## ■ Uniqueness

- Every bit pattern represents unique integer value
- Each representable integer has unique bit encoding

## ■ $\Rightarrow$ Can Invert Mappings

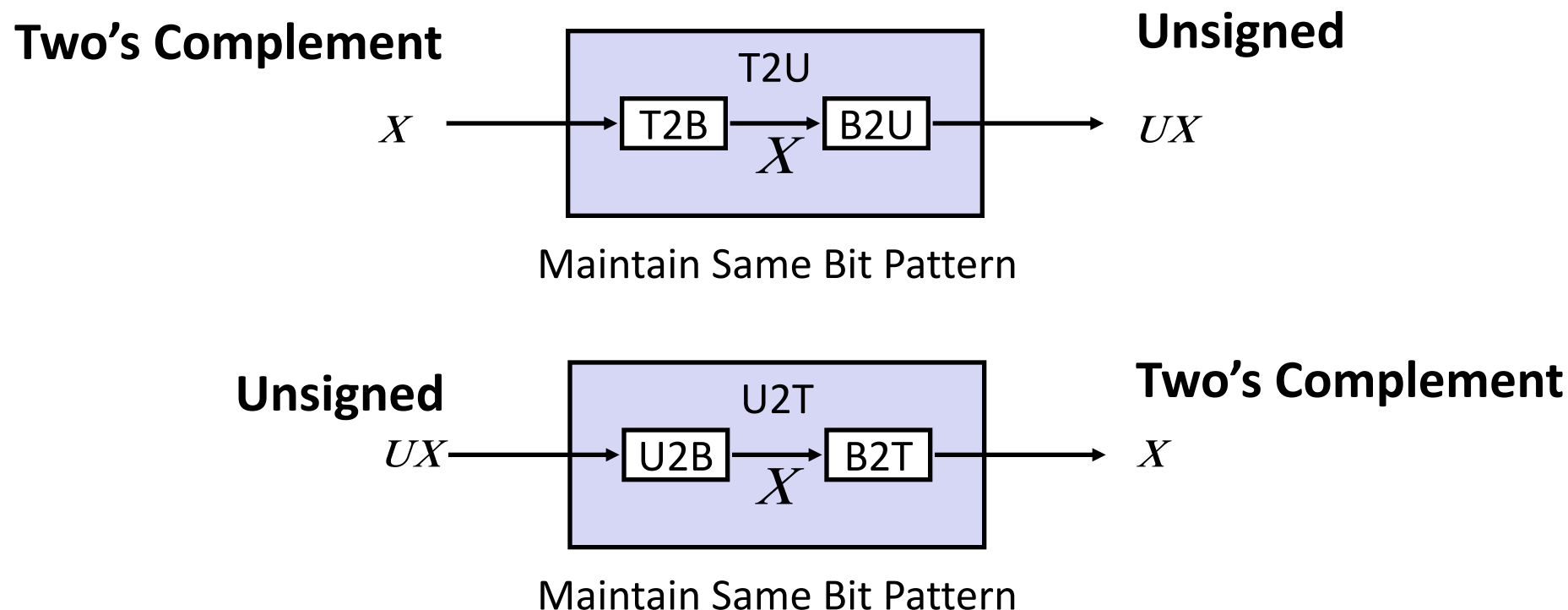
- $U2B(x) = B2U^{-1}(x)$ 
  - Bit pattern for unsigned integer
- $T2B(x) = B2T^{-1}(x)$ 
  - Bit pattern for two's comp integer



# Today: Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- **Integers**
  - Representation: unsigned and signed
  - **Conversion, casting**
  - Expanding, truncating
  - Addition, negation, multiplication, shifting
  - Summary
- Representations in memory, pointers, strings

# Mapping Between Signed & Unsigned



- Mappings between unsigned and two's complement numbers:  
**Keep bit representations and reinterpret**

# Mapping Signed $\leftrightarrow$ Unsigned

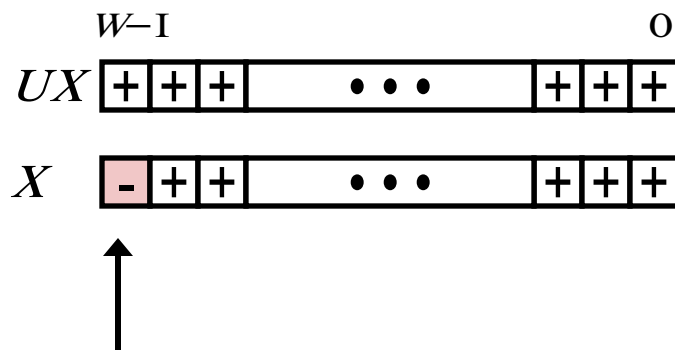
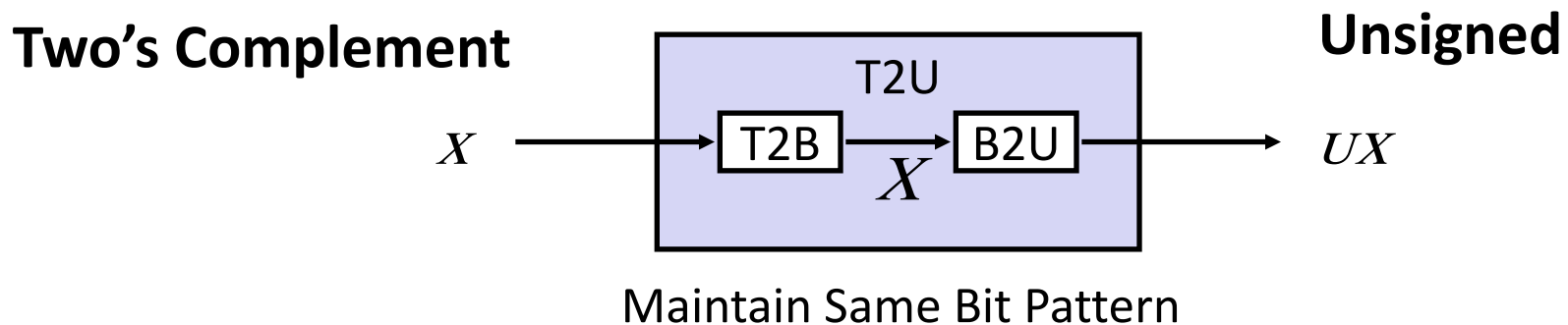
| Bits | Signed |  | Unsigned |
|------|--------|--|----------|
| 0000 | 0      |  | 0        |
| 0001 | 1      |  | 1        |
| 0010 | 2      |  | 2        |
| 0011 | 3      |  | 3        |
| 0100 | 4      |  | 4        |
| 0101 | 5      |  | 5        |
| 0110 | 6      |  | 6        |
| 0111 | 7      |  | 7        |
| 1000 | -8     |  | 8        |
| 1001 | -7     |  | 9        |
| 1010 | -6     |  | 10       |
| 1011 | -5     |  | 11       |
| 1100 | -4     |  | 12       |
| 1101 | -3     |  | 13       |
| 1110 | -2     |  | 14       |
| 1111 | -1     |  | 15       |

$\rightarrow$  **T2U**  $\rightarrow$   
 $\leftarrow$  **U2T**  $\leftarrow$

# Mapping Signed $\leftrightarrow$ Unsigned

| Bits | Signed |          | Unsigned |
|------|--------|----------|----------|
| 0000 | 0      | $=$      | 0        |
| 0001 | 1      |          | 1        |
| 0010 | 2      |          | 2        |
| 0011 | 3      |          | 3        |
| 0100 | 4      |          | 4        |
| 0101 | 5      |          | 5        |
| 0110 | 6      |          | 6        |
| 0111 | 7      |          | 7        |
| 1000 | -8     | $\pm 16$ | 8        |
| 1001 | -7     |          | 9        |
| 1010 | -6     |          | 10       |
| 1011 | -5     |          | 11       |
| 1100 | -4     |          | 12       |
| 1101 | -3     |          | 13       |
| 1110 | -2     |          | 14       |
| 1111 | -1     |          | 15       |

# Relation between Signed & Unsigned

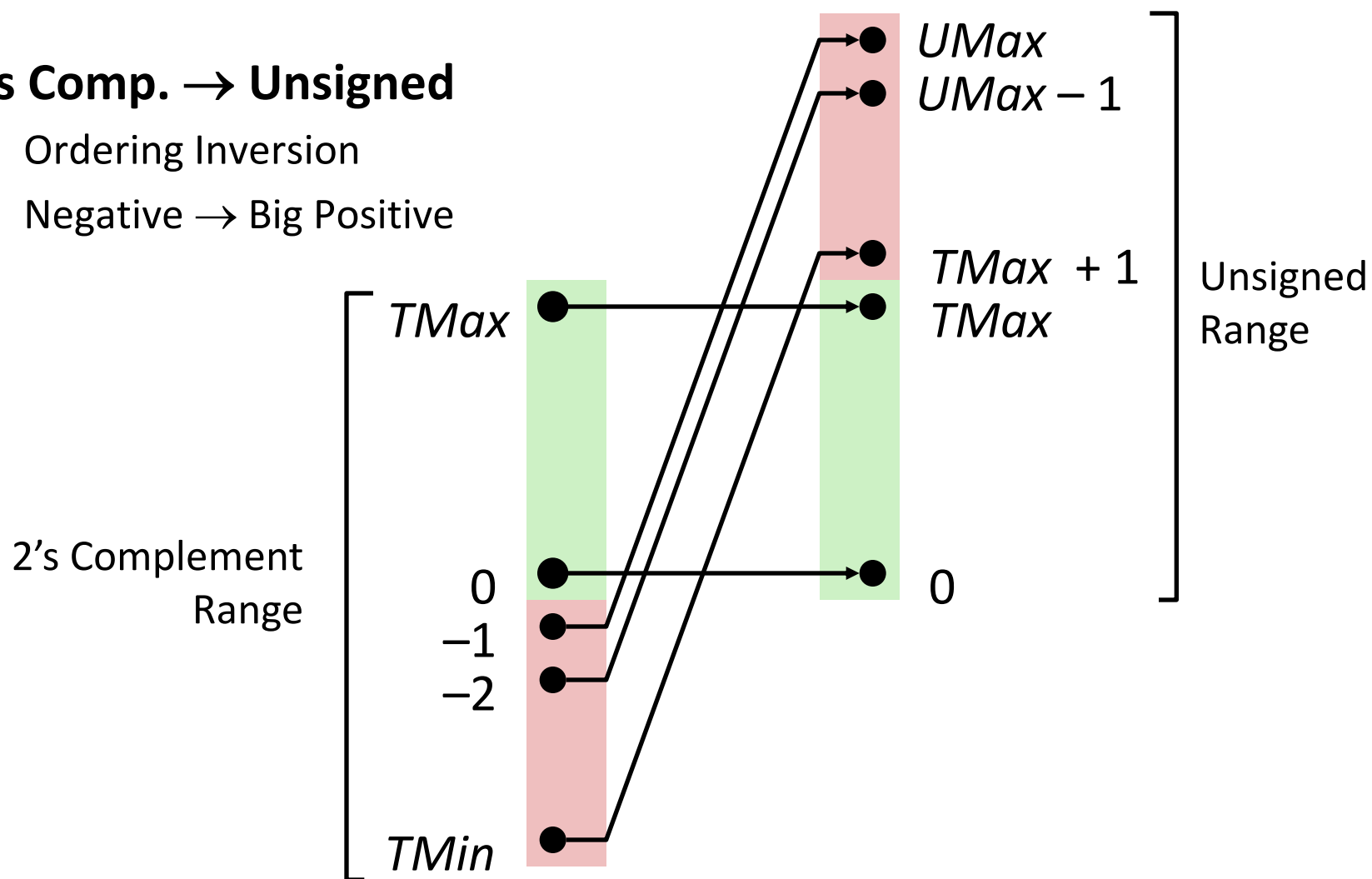


Large negative weight  
*becomes*  
 Large positive weight

# Conversion Visualized

## ■ 2's Comp. → Unsigned

- Ordering Inversion
- Negative → Big Positive



# 补码 $\leftrightarrow$ 无符号数

- 对于满足  $TMin_w \leq x \leq TMax_w$  的  $x$  (补码) 有

$$T2U_w(x) = \begin{cases} x + 2^w, & x < 0 \\ x, & x \geq 0 \end{cases}$$

- 对于满足  $0 \leq u \leq UMax_w$  的  $u$  (无符号数) 有

$$U2T_w(u) = \begin{cases} u - 2^w, & u > TMax_w \\ u, & u \leq TMax_w \end{cases}$$

# Signed vs. Unsigned in C

## ■ Constants

- By default are considered to be **signed integers**
- Unsigned if have “**U**” as suffix

0U, 4294967259U

## ■ Casting

- **Explicit casting** between signed & unsigned same as U2T and T2U

```
int tx, ty;
unsigned ux, uy;
tx = (int) ux;
uy = (unsigned) ty;
```

- **Implicit casting** also occurs via assignments and procedure calls

```
tx = ux;                int fun(unsigned u);
uy = ty;                uy = fun(tx);
```



# Casting Surprises

## ■ Expression Evaluation

- If there is a mix of unsigned and signed in single expression,  
*signed values implicitly cast to unsigned*
- Including comparison operations `<`, `>`, `==`, `<=`, `>=`
- Examples for  $W = 32$ : **TMIN = -2,147,483,648**, **TMAX = 2,147,483,647**

| 关系表达式                          | 运算类型 | 结果 | 说明 |
|--------------------------------|------|----|----|
| 0 == 0U                        |      |    |    |
| -1 < 0                         |      |    |    |
| -1 < 0U                        |      |    |    |
| 2147483647 > -2147483647-1     |      |    |    |
| 2147483647U > -2147483647-1    |      |    |    |
| 2147483647 > (int) 2147483648U |      |    |    |
| -1 > -2                        |      |    |    |
| (unsigned) -1 > -2             |      |    |    |

# C语言程序中的整数

| 关系<br>表达式                        | 类型 | 结果 | 说明   |
|----------------------------------|----|----|--|
| $0 == 0U$                        | 无  | 1  | $00...0B = 00...0B$                        |
| $-1 < 0$                         | 带  | 1  | $11...1B (-1) < 00...0B (0)$               |
| $-1 < 0U$                        | 无  | 0* | $11...1B (2^{32}-1) > 00...0B(0)$          |
| $2147483647 > -2147483647 - 1$   | 带  | 1  | $011...1B (2^{31}-1) > 100...0B (-2^{31})$ |
| $2147483647U > -2147483647 - 1$  | 无  | 0* | $011...1B (2^{31}-1) < 100...0B(2^{31})$   |
| $2147483647 > (int) 2147483648U$ | 带  | 1* | $011...1B (2^{31}-1) > 100...0B (-2^{31})$ |
| $-1 > -2$                        | 带  | 1  | $11...1B (-1) > 11...10B (-2)$             |
| $(unsigned) -1 > -2$             | 无  | 1  | $11...1B (2^{32}-1) > 11...10B (2^{32}-2)$ |

**带\*的结果与常规预想的相反！**

# C语言程序中的整数

例如，考虑以下C代码：

```
1 int x = -1;
2 unsigned u = 2147483648;
3
4 printf ( "x = %u = %d\n" , x, x);
5 printf ( "u = %u = %d\n" , u, u);
```

在32位机器上运行上述代码时，它的输出结果是什么？为什么？

**x = 4294967295 = -1**

**u = 2147483648 = -2147483648**

- ◆ 因为-1的补码整数表示为“11...1”，作为32位无符号数解释时，其值为 $2^{32}-1 = 4\ 294\ 967\ 296-1 = 4\ 294\ 967\ 295$ 。
- ◆  $2^{31}$ 的无符号数表示为“100...0”，被解释为32位带符号整数时，其值为最小负数： $-2^{32-1} = -2^{31} = -2\ 147\ 483\ 648$ 。

# C语言程序中的整数

- 1) 在有些32位系统上, C表达式  $-2147483648 < 2147483647$  的执行结果为false。Why?
- 2) 若定义变量 `int i=-2147483648;`, 则 `i < 2147483647` 的执行结果为true。Why?
- 3) 如果将表达式写成 `-2147483647-1 < 2147483647`, 则结果会怎样呢? Why?

1) 在ISO C90标准下, 2147483648为unsigned类型, 因此  
 $-2147483648 < 2147483647$  按无符号数比较,  
10.....0B比01.....1大, 结果为false。

在ISO C99标准下, 2147483648为int类型, 因此  
 $-2147483648 < 2147483647$  按带符号整数比较,  
10.....0B比01.....1小, 结果为true。

- 2) `i < 2147483647` 按int型数比较, 结果为true。
- 3) `-2147483647-1 < 2147483647` 按int型比较, 结果为true。

# Unsigned vs. Signed: Easy to Make Mistakes

```
unsigned i;  
for (i = cnt-2; i >= 0; i--)  
    a[i] += a[i+1];
```

- Can be very subtle

```
#define DELTA sizeof(int)  
int i;  
for (i = CNT; i-DELTA >= 0; i-= DELTA)  
    . . .
```

# Summary

## Casting Signed $\leftrightarrow$ Unsigned: Basic Rules

- Bit pattern is maintained
- But reinterpreted
- Can have unexpected effects: adding or subtracting  $2^w$
- Expression containing signed and unsigned int
  - `int` is cast to `unsigned`!!

# Today: Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- **Integers**
  - Representation: unsigned and signed
  - Conversion, casting
  - **Expanding, truncating**
  - Addition, negation, multiplication, shifting
  - Summary
- Representations in memory, pointers, strings

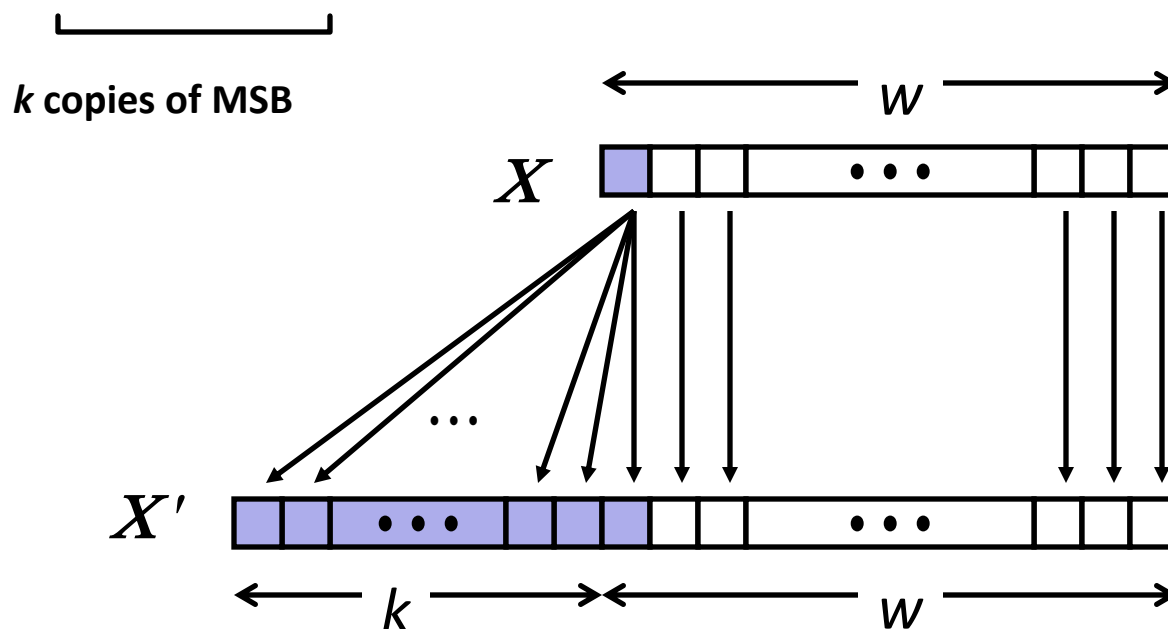
# Sign Extension

## ■ Task:

- Given  $w$ -bit signed integer  $x$
- Convert it to  $w+k$ -bit integer with same value

## ■ Rule:

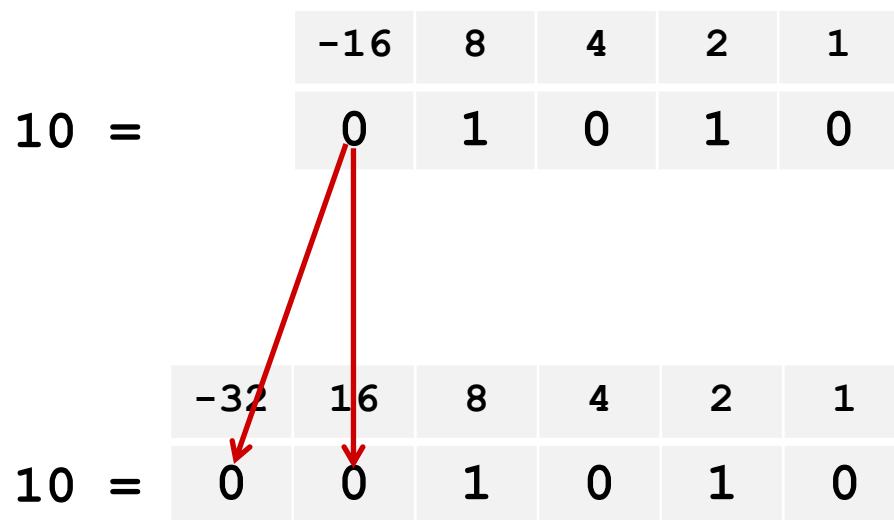
- Make  $k$  copies of sign bit:
- $X' = \underbrace{x_{w-1}, \dots, x_{w-1}}_{k \text{ copies of MSB}}, x_{w-1}, x_{w-2}, \dots, x_0$



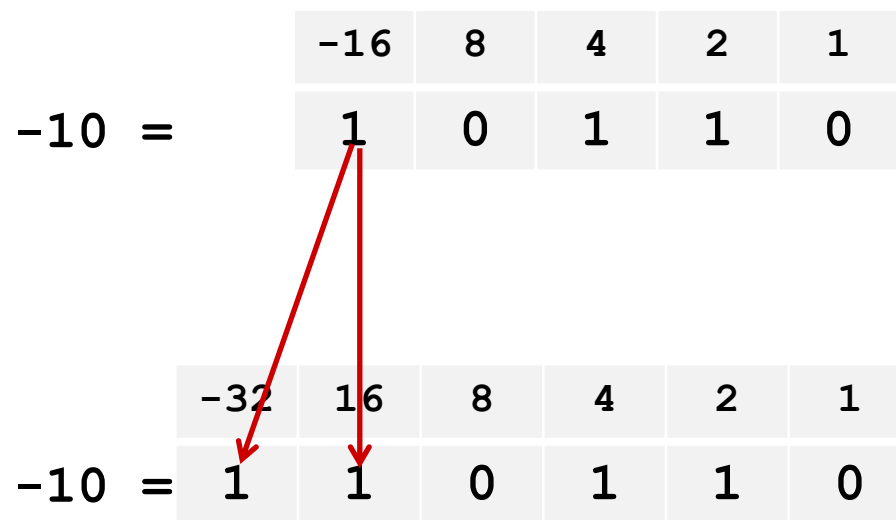


# Sign Extension: Simple Example

## Positive number



## Negative number



# Larger Sign Extension Example

```
short int x = 15213;
int      ix = (int) x;
short int y = -15213;
int      iy = (int) y;
```

|           | Decimal | Hex         | Binary                              |
|-----------|---------|-------------|-------------------------------------|
| <b>x</b>  | 15213   | 3B 6D       | 00111011 01101101                   |
| <b>ix</b> | 15213   | 00 00 3B 6D | 00000000 00000000 00111011 01101101 |
| <b>y</b>  | -15213  | C4 93       | 11000100 10010011                   |
| <b>iy</b> | -15213  | FF FF C4 93 | 11111111 11111111 11000100 10010011 |

- Converting from smaller to larger integer data type
- C automatically performs sign extension

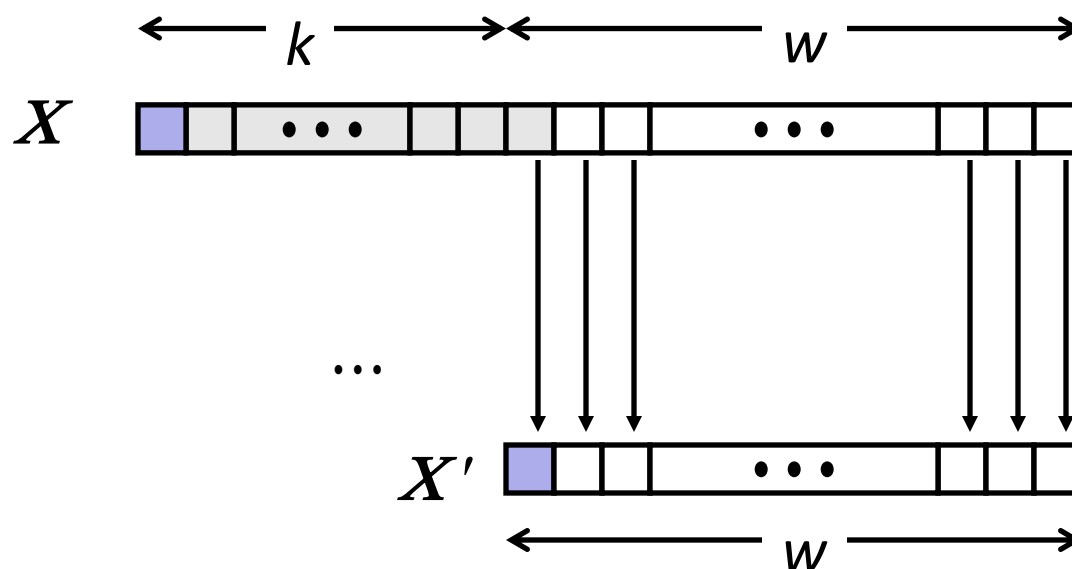
# Truncation

## ■ Task:

- Given  $k+w$ -bit signed or unsigned integer  $X$
- Convert it to  $w$ -bit integer  $X'$  with same value for “small enough”  $X$

## ■ Rule:

- Drop top  $k$  bits:
- $X' = x_{w-1}, x_{w-2}, \dots, x_0$



# Truncation: Simple Example

## No sign change

|     |     |   |   |   |   |
|-----|-----|---|---|---|---|
|     | -16 | 8 | 4 | 2 | 1 |
| 2 = | 0   | 0 | 0 | 1 | 0 |

|     |    |   |   |   |
|-----|----|---|---|---|
|     | -8 | 4 | 2 | 1 |
| 2 = | 0  | 0 | 1 | 0 |

$$2 \bmod 16 = 2$$

|      |     |   |   |   |   |
|------|-----|---|---|---|---|
|      | -16 | 8 | 4 | 2 | 1 |
| -6 = | 1   | 1 | 0 | 1 | 0 |

|      |    |   |   |   |
|------|----|---|---|---|
|      | -8 | 4 | 2 | 1 |
| -6 = | 1  | 0 | 1 | 0 |

$$-6 \bmod 16 = 26 \text{U} \bmod 16 = 10 \text{U} = -6$$

## Sign change

|      |     |   |   |   |   |
|------|-----|---|---|---|---|
|      | -16 | 8 | 4 | 2 | 1 |
| 10 = | 0   | 1 | 0 | 1 | 0 |

|      |    |   |   |   |
|------|----|---|---|---|
|      | -8 | 4 | 2 | 1 |
| -6 = | 1  | 0 | 1 | 0 |

$$10 \bmod 16 = 10 \text{U} \bmod 16 = 10 \text{U} = -6$$

|       |     |   |   |   |   |
|-------|-----|---|---|---|---|
|       | -16 | 8 | 4 | 2 | 1 |
| -10 = | 1   | 0 | 1 | 1 | 0 |

|     |    |   |   |   |
|-----|----|---|---|---|
|     | -8 | 4 | 2 | 1 |
| 6 = | 0  | 1 | 1 | 0 |

$$-10 \bmod 16 = 22 \text{U} \bmod 16 = 6 \text{U} = 6$$

# 举例

**例1（扩展操作）：** 在大端机上输出  
si, usi, i, ui的十进制和十六进制值  
是什么？

```
short si = -32768;
unsigned short usi = si;
int i = si;
unsigned ui = usi;
```

```
si = -32768    80 00
usi = 32768    80 00
i = -32768     FF FF 80 00
ui = 32768     00 00 80 00
```

**例2（截断操作）：** i和j是否相等？

```
int i = 32768;
short si = (short) i;
int j = si;
```

不相等！

```
i = 32768    00 00 80 00
si = -32768   80 00
j = -32768   FF FF 80 00
```

**原因：** 对i截断时发生了“溢出”，即：32768截断为16位数时，因其超出16位能表示的最大值，故无法截断为正确的16位数！

# Summary:

## Expanding, Truncating: Basic Rules

- **Expanding (e.g., short int to int)**
  - Unsigned: zeros added
  - Signed: sign extension
  - Both yield expected result
  
- **Truncating (e.g., unsigned to unsigned short)**
  - Unsigned/signed: bits are truncated
  - Result reinterpreted
  - Unsigned: mod operation
  - Signed: similar to mod
  - For small (in magnitude) numbers yields expected behavior

# 小结

- 非数值数据的表示
  - 逻辑数据用来表示真/假或N位位串，按位运算
  - 西文字符：用ASCII码表示
  - 汉字编码
- 数据的宽度
  - 位、字节、字（不一定等于字长）
  - k / K / M / G / T / P / E / Z / Y 有不同的含义
- 数据的存储排列
  - 数据的地址：连续若干单元中最小的地址，即：从小地址开始存放数据
    - 问题：若一个short型数据si存放在单元0x08000100和0x08000101中，那么si的地址是什么？
  - 大端方式：用MSB存放的地址表示数据的地址
  - 小端方式：用LSB存放的地址表示数据的地址
  - 按边界对齐可减少访存次数

# 小结

- 在机器内部编码后的数称为机器数，其值称为真值
- 定义数值数据有三个要素：进制、定点/浮点、编码
- 整数的表示
  - 无符号数：正整数，用来表示地址等；带符号整数：用补码表示
- C语言中的整数
  - 无符号数：unsigned int ( short / long)；带符号数：int ( short / long)
- 十进制数的表示：用ASCII码或BCD码表示
- 整形运算
  - 算术运算：无符号整数、有符号整数
  - 按位运算：“|”、“&”、“~”、“^”
  - 逻辑运算：“||”、“&&”、“!”
  - 移位运算：“<<”、“>>”（逻辑移位、算术移位）
  - 位扩展和位截断：无符号和有符号