

Analysis of Sequential Decoding Complexity Using the Berry-Esseen Inequality*

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Abstract

This study presents a novel technique to estimate the computational complexity of sequential decoding using the Berry-Esseen theorem. Unlike the theoretical bounds determined by the conventional central limit theorem argument, which often holds only for sufficiently large codeword length, the new bound obtained from the Berry-Esseen theorem is valid for any blocklength. The accuracy of the new bound is then examined for two sequential decoding algorithms, an ordering-free variant of the generalized Dijkstra's algorithm (GDA)(or simplified GDA) and the maximum-likelihood sequential decoding algorithm (MLSDA). Empirically investigating codes of small blocklength reveals that the theoretical upper bound for the simplified GDA almost matches the simulation results as the signal-to-noise ratio (SNR) per information bit (γ_b) is greater than or equal to 8 dB. However, the theoretical bound may become markedly higher than the simulated average complexity when γ_b is small. For the MLSDA, the theoretical upper bound is quite close to the simulation results for both high SNR ($\gamma_b \geq 6$ dB) and low SNR ($\gamma_b \leq 2$ dB). Even for moderate SNR, the simulation results and the theoretical bound differ by at most 0.579966 on a \log_{10} scale.

Index Terms

Coding, Decoding, Large Deviations, Convolutional Codes, Maximum-Likelihood, Soft-Decision, Sequential Decoding

I. INTRODUCTION

The Berry-Esseen theorem [6, sec.XVI. 5] states that the distribution of the sum of independent zero-mean random variables $\{X_i\}_{i=1}^n$, normalized by the standard deviation of the sum, differs from the unit Gaussian distribution by no more than $C r_n / s_n^3$, where s_n^2 and r_n are, respectively, the sums of the marginal variances and the marginal absolute third moments, and the Berry-Esseen coefficient, C , is an absolute constant. Specifically, for every $a \in \mathbb{R}$,

$$\left| \Pr \left\{ \frac{1}{s_n} (X_1 + \cdots + X_n) \leq a \right\} - \Phi(a) \right| \leq C \frac{r_n}{s_n^3}, \quad (1)$$

where $\Phi(\cdot)$ represents the unit Gaussian cumulative distribution function (cdf). The remarkable aspect of this theorem is that the upper bound depends only on the variance and the absolute third moment, and therefore, can provide a good probability estimate through the first three

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moments. A typical estimate of the absolute constant is six [6, sec.XVI. 5, Thm. 2]. When $\{X_n\}_{i=1}^n$ are identically distributed, in addition to independent, the absolute constant can be reduced to three, and has been reported to be improved down to 2.05 [6, sec.XVI. 5, Thm. 1]. In 1972, Beek sharpened the constant to 0.7975 [2]. Later, Shiganov further improved the constant down to 0.7915 for an independent sample sum, and, 0.7655, if these samples are also identically distributed [23]. Shiganov's result is generally considered to be the best result yet obtained thus far [22].

In applying this inequality to analyze the computational complexity of sequential decoding algorithms, the original analytical problem is first transformed into one that concerns the asymptotic probability mass of the sum of independent random samples. Inequality (1) can therefore be applied. The complexities of two sequential maximum-likelihood decoding algorithms are then analyzed. One is an ordering-free variant of the generalized Dijkstra's algorithm (GDA) [11] operated over a code tree of linear block codes, and the other is the maximum-likelihood sequential decoding algorithm (MLSDA) [13] that searches for the codeword over a trellis of binary convolutional codes.

The computational effort required by sequential decoding is conventionally determined using a random coding technique, which averages the computational effort over the ensemble of random tree codes [14], [16], [21]. Branching process analysis on sequential decoding complexity has been recently proposed [9], [17], [18]; the results, however, were still derived by averaging over semi-random tree codes. Chevillat and Costello proposed to analyze the computational effort of sequential decoding in terms of the column distance function of a specific time-invariant code [4]; but, the analysis only applied to a situation in which the code was transmitted via binary symmetric channels.

In light of the Berry-Esseen inequality and the large deviations technique, this work presents an alternative approach to derive the theoretical upper bounds on the computational effort of the simplified GDA and the MLSDA for binary codes antipodally transmitted through an additive white Gaussian noise (AWGN) channel. Unlike the bounds established in terms of the conventional central limit theorem argument, which often holds only for sufficiently large codeword length, the new bound is valid for any blocklength. Empirically investigating codes of small blocklength shows that for the trellis-based MLSDA, the theoretical upper bound is quite close to the simulation results for both high SNR and low SNR; even for moderate SNR, the

theoretical upper bound and the simulation results differ by no more than 0.579966 on a \log_{10} scale. For the tree-based ordering-free GDA, the theoretical bound coincides with the simulation results at high SNR; however, the bound tends to be substantially larger than the simulation results at very low SNR. The possible cause of the inaccuracy of the bound at low SNR for the tree-based ordering-free GDA is addressed at the end of this study.

The rest of this paper is organized as follows. Section II derives a probability bound for use of analyzing the sequential decoding complexity due to the Berry-Esseen inequality. Section III presents an analysis of the average computational complexity of the GDA. Section IV briefly introduces the MLSDA, and then analyzes its complexity upper bound. Conclusions are finally drawn in Section V.

Throughout this article, $\Phi(\cdot)$ denotes the unit Gaussian cdf.

II. BERRY-ESSEEN THEOREM AND PROBABILITY BOUND

This section derives an upper probability bound for the sum of independent random samples using the Berry-Esseen inequality. This bound is essential to the analysis of the computational effort of sequential decoding algorithms.

The approach used here is the *large deviations* technique, which is generally applied to compute the *exponent* of an exponentially decaying probability mass. The Berry-Esseen inequality is also applied to evaluate the *subexponential* detail of the concerned probability. With these two techniques, an upper bound of the concerned probability can be established.

Lemma 1: Let $Y_n = \sum_{i=1}^n X_i$ be the sum of i.i.d. random variables. Define

$$\rho(\theta) \triangleq E \left[\left| X_1^{(\theta)} - E[X_1^{(\theta)}] \right|^3 \right], \quad \sigma(\theta) \triangleq \sqrt{\text{Var}[X_1^{(\theta)}]} \text{ and } M(\theta) \triangleq E[e^{\theta X_1}].$$

Then

$$\Pr \{Y_n \leq -nx\} \leq \begin{cases} A_n(\theta) e^{n\theta x} M^n(\theta), & \text{when } E[X_1] > -x; \\ 1, & \text{otherwise,} \end{cases}$$

for some $\theta < 0$, where

$$A_n(\theta) \triangleq \min \left(e^{\theta^2 \cdot n \cdot \sigma^2(\theta)/2} \Phi(\theta \cdot \sqrt{n} \cdot \sigma(\theta)) + 2C \frac{\rho(\theta)}{\sqrt{n} \sigma^3(\theta)}, 1 \right),$$

and $C = 0.7655$.

Proof: The bound subject to the condition of $E[X_1] \leq -x$ is trivial; therefore, only the validity of the upper bound subject to $E[X_1] > -x$ is proved here.

Form a new random variable, $\tilde{Y}_n = \sum_{i=1}^n \tilde{X}_i$, where $\tilde{X}_i = (X_i + x)$. Then $E[e^{\theta \tilde{X}_1}] = E[e^{\theta(X_1+x)}] = e^{\theta x} M(\theta)$. Denote the distribution of \tilde{Y}_n by $F_n(\cdot)$. The *twisted* distribution with parameter θ , corresponding to $F_n(\cdot)$, is defined by

$$dF_n^{(\theta)}(y) \triangleq \frac{\exp\{\theta y\} dF_n(y)}{\int_y \exp\{\theta \hat{y}\} dF_n(\hat{y})} = \frac{\exp\{\theta y\} dF_n(y)}{\exp\{\theta n x\} M^n(\theta)}. \quad (2)$$

Denote the random variable whose probability distribution is $F_n^{(\theta)}(\cdot)$ by $\tilde{Y}_n^{(\theta)}$. A *twisted* counterpart of \tilde{X}_i , denoted by $\tilde{X}_i^{(\theta)}$, is similarly defined. A fundamental large deviation result is that $\tilde{Y}_n^{(\theta)} = \sum_{i=1}^n \tilde{X}_i^{(\theta)}$ almost surely, and $\{\tilde{X}_i^{(\theta)}\}_{i=1}^n$ also forms an independent and identically distributed sequence. In addition, $\tilde{X}_i^{(\theta)} = X_i^{(\theta)} + x$, $\text{Var}[\tilde{X}_1^{(\theta)}] = \sigma^2(\theta)$, and $E[|\tilde{X}_1^{(\theta)} - E[\tilde{X}_1^{(\theta)}]|^3] = \rho(\theta)$. Denote the distribution of $(\tilde{Y}_n^{(\theta)} - E[\tilde{Y}_n^{(\theta)}]) / [\sqrt{n} \sigma(\theta)]$ by $H_n(\cdot)$. Then from equation (2),

$$\begin{aligned} \Pr(\tilde{Y}_n \leq 0) &= \int_{-\infty}^0 dF_n(y) \\ &= e^{n\theta x} M^n(\theta) \int_{-\infty}^0 \exp\{-\theta y\} dF_n^{(\theta)}(y) \\ &= e^{n\theta x} M^n(\theta) \int_{-\infty}^0 \exp\{-\theta \cdot \sqrt{n} \cdot \sigma(\theta) \cdot y\} dH_n(y). \end{aligned} \quad (3)$$

Integrating by parts on (3) with $\lambda(dy) \triangleq -\theta \cdot \sqrt{n} \cdot \sigma(\theta) \cdot \exp\{-\theta \cdot \sqrt{n} \cdot \sigma(\theta) \cdot y\} dy$, and then applying equation (1) yields

$$\begin{aligned} \int_{-\infty}^0 e^{-\theta \sqrt{n} \sigma(\theta) y} dH_n(y) &= \int_{-\infty}^0 [H_n(0) - H_n(y)] \lambda(dy) \\ &\leq \int_{-\infty}^0 \left[\Phi(0) - \Phi(y) + 2C \frac{\rho(\theta)}{\sqrt{n} \sigma^3(\theta)} \right] \lambda(dy) \\ &= \int_{-\infty}^0 [\Phi(0) - \Phi(y)] \lambda(dy) + 2C \frac{\rho(\theta)}{\sqrt{n} \sigma^3(\theta)} \\ &= \exp\left\{ \frac{\theta^2 \cdot n \cdot \sigma^2(\theta)}{2} \right\} \Phi(\theta \cdot \sqrt{n} \cdot \sigma(\theta)) + 2C \frac{\rho(\theta)}{\sqrt{n} \sigma^3(\theta)}, \end{aligned} \quad (4)$$

where (4) holds by, again, applying integration by part. On the other hand,

$$\int_{-\infty}^0 \exp\{-\theta \cdot \sqrt{n} \cdot \sigma(\theta) \cdot y\} dH_n(y) \leq 1$$

can be established by observing that

$$\begin{aligned}
e^{n\theta x} M^n(\theta) \int_{-\infty}^0 \exp\{-\theta \cdot \sqrt{n} \cdot \sigma(\theta) \cdot y\} dH_n(y) &= \Pr\{\tilde{Y}_n \leq 0\} \\
&= \Pr\left\{e^{\theta \tilde{Y}_n} \geq 1\right\} \\
&\leq E[e^{\theta \tilde{Y}_n}] \\
&= e^{n\theta x} M^n(\theta).
\end{aligned}$$

We then conclude that

$$\begin{aligned}
&\Pr(Y_n \leq -nx) \\
&= \Pr(\tilde{Y}_n \leq 0) \\
&\leq \min\left(e^{\theta^2 \cdot n \cdot \sigma^2(\theta)/2} \Phi(\theta \cdot \sqrt{n} \cdot \sigma(\theta)) + 2C \frac{\rho(\theta)}{\sqrt{n} \sigma^3(\theta)}, 1\right) e^{n\theta x} M^n(\theta).
\end{aligned}$$

■

The independent samples with which our decoding problems are concerned actually consist of two i.i.d. sequences, one of which is Gaussian distributed and the other is non-Gaussian distributed. Thus, an upper probability bound concerning such a compound i.i.d. sequence is established in the following lemma.

Lemma 2: Let $Y_n = \sum_{i=1}^n X_i$ be the sum of independent random variables $\{X_i\}_{i=1}^n$, among which $\{X_i\}_{i=1}^d$ are identically Gaussian distributed with positive mean μ and non-zero variance σ^2 , and $\{X_i\}_{i=d+1}^n$ have common marginal distribution as $\min\{X_1, 0\}$. Let $\gamma \triangleq (1/2)(\mu^2/\sigma^2)$. Then

$$\Pr\{Y_n \leq 0\} \leq \mathcal{B}(d, n-d, \gamma),$$

where

$$\mathcal{B}(d, n-d, \gamma) = \begin{cases} \Phi(-\sqrt{2\gamma n}), & \text{if } d = n; \\ \Phi\left(-\frac{(n-d)\hat{\mu} + d\sqrt{2\gamma}}{\sqrt{d}}\right) + \tilde{A}_{n,d}(\lambda) \\ \quad \times \left[\Phi(-\lambda)e^{-\gamma}e^{\lambda^2/2} + \Phi(\sqrt{2\gamma})\right]^{n-d} \\ \quad \times e^{d(-\gamma + \lambda^2/2)} \Phi\left(\frac{(n-d)\hat{\mu} + \lambda d}{\sqrt{d}}\right), & \text{if } 1 > \frac{d}{n} \geq 1 - \frac{\sqrt{4\pi\gamma}e^\gamma}{1 + \sqrt{4\pi\gamma}e^\gamma\Phi(\sqrt{2\gamma})}; \\ 1, & \text{otherwise,} \end{cases}$$

$$\begin{aligned}
\tilde{A}_{n,d}(\lambda) &\triangleq \min \left(e^{(n-d)(\lambda-\sqrt{2\gamma})^2 \tilde{\sigma}^2(\lambda)/2} \Phi \left((\lambda - \sqrt{2\gamma}) \tilde{\sigma}(\lambda) \sqrt{n-d} \right) + \frac{1.531 \cdot \tilde{\rho}(\lambda)}{\tilde{\sigma}^3(\lambda) \sqrt{n-d}}, 1 \right), \\
\hat{\mu} &\triangleq E[X_{d+1}] = -(1/\sqrt{2\pi})e^{-\gamma} + \sqrt{2\gamma}\Phi(-\sqrt{2\gamma}), \\
\tilde{\sigma}^2(\lambda) &\triangleq -\frac{d}{n-d} - \frac{nd}{(n-d)^2}\lambda^2 + \frac{n}{n-d} \frac{1}{1 + \sqrt{2\pi}\lambda e^\gamma \Phi(\sqrt{2\gamma})}, \\
\tilde{\rho}(\lambda) &\triangleq \frac{n}{(n-d)} \frac{\lambda}{[1 + \sqrt{2\pi}\lambda e^\gamma \Phi(\sqrt{2\gamma})]} \left\{ 1 - \frac{d(n+d)}{(n-d)^2}\lambda^2 \right. \\
&\quad + 2 \left[\frac{n^2}{(n-d)^2}\lambda^2 + 2 \right] e^{-d(2n-d)\lambda^2/[2(n-d)^2]} \\
&\quad - \frac{d}{n-d} \left[\frac{n+d}{n-d}\lambda^2 + 3 \right] \sqrt{2\pi}\lambda e^\gamma \Phi(\sqrt{2\gamma}) \\
&\quad \left. - \frac{2n}{n-d} \left[\frac{n^2}{(n-d)^2}\lambda^2 + 3 \right] \sqrt{2\pi}\lambda e^{\lambda^2/2} \Phi \left(-\frac{n}{n-d}\lambda \right) \right\},
\end{aligned}$$

and λ is the unique solution (in $[0, \sqrt{2\gamma})$) of

$$\lambda e^{(1/2)\lambda^2} \Phi(-\lambda) = \frac{1}{\sqrt{2\pi}} \left(1 - \frac{d}{n} \right) - \frac{d}{n} e^\gamma \Phi(\sqrt{2\gamma})\lambda.$$

Proof: Only the bound for $d < n$ is proved since the case of $d = n$ can be easily substantiated.

By noting that $(\mu/\sigma) = \sqrt{2\gamma}$, $\Pr(Y_n \leq 0)$ can be bounded by

$$\begin{aligned}
&\Pr(Y_n \leq 0) \\
&= \Pr\{X_1 + \dots + X_d + X_{d+1} + \dots + X_n \leq 0\} \\
&= \int_{-\infty}^{\infty} \Pr\{X_{d+1} + \dots + X_n \leq -x\} \frac{1}{\sqrt{2\pi d\sigma^2}} e^{-\frac{(x-d\mu)^2}{2d\sigma^2}} dx, \quad (x \rightarrow \sigma x') \\
&= \int_{-\infty}^{\infty} \Pr\{X_{d+1} + \dots + X_n \leq -\sigma x'\} \frac{1}{\sqrt{2\pi d}} e^{-\frac{(x'-d\sqrt{2\gamma})^2}{2d}} dx', \quad (x' \rightarrow (n-d)x'') \\
&= \int_{-\infty}^{\infty} \Pr\{X_{d+1} + \dots + X_n \leq -\sigma(n-d)x''\} \frac{1}{\sqrt{2\pi d/(n-d)^2}} e^{-\frac{(x''-d\sqrt{2\gamma}/(n-d))^2}{2d/(n-d)^2}} dx'' \\
&= \int_{-\infty}^{-\hat{\mu}} \Pr\{X_{d+1} + \dots + X_n \leq -\sigma(n-d)x''\} \frac{1}{\sqrt{2\pi d/(n-d)^2}} e^{-\frac{(x''-d\sqrt{2\gamma}/(n-d))^2}{2d/(n-d)^2}} dx'' \\
&\quad + \int_{-\hat{\mu}}^{\infty} \Pr\{X_{d+1} + \dots + X_n \leq -\sigma(n-d)x''\} \frac{1}{\sqrt{2\pi d/(n-d)^2}} e^{-\frac{(x''-d\sqrt{2\gamma}/(n-d))^2}{2d/(n-d)^2}} dx'' \\
&\leq \int_{-\infty}^{-\hat{\mu}} \frac{1}{\sqrt{2\pi d/(n-d)^2}} e^{-\frac{(x''-d\sqrt{2\gamma}/(n-d))^2}{2d/(n-d)^2}} dx'' \\
&\quad + \int_{-\hat{\mu}}^{\infty} A_{n-d}(\theta) e^{(n-d)\theta\sigma x''} M^{n-d}(\theta) \frac{1}{\sqrt{2\pi d/(n-d)^2}} e^{-\frac{(x''-d\sqrt{2\gamma}/(n-d))^2}{2d/(n-d)^2}} dx''
\end{aligned}$$

$$\begin{aligned}
&= \Phi \left(-\frac{(n-d)\hat{\mu} + d\sqrt{2\gamma}}{\sqrt{d}} \right) \\
&\quad + A_{n-d}(\theta) M^{n-d}(\theta) e^{d(\theta\sigma\sqrt{2\gamma} + \theta^2\sigma^2/2)} \Phi \left(\frac{(n-d)\hat{\mu} + d\sqrt{2\gamma}}{\sqrt{d}} + \theta\sigma\sqrt{d} \right),
\end{aligned}$$

where $M(\theta) = E[e^{\theta X_{d+1}}]$, and the last inequality follows from Lemma 1. We then note that

$$M^{n-d}(\theta) e^{d(\theta\sigma\sqrt{2\gamma} + \theta^2\sigma^2/2)}$$

is exactly the moment generating function of $Y_n = \sum_{i=1}^n X_i$; hence, if $E[Y_n] = d\mu + (n-d)\sigma\hat{\mu} > 0$, then the solution θ of $\partial E[e^{\theta Y_n}]/\partial \theta = 0$ is definitely negative.

For notational convenience, we let $\lambda = (\mu/\sigma) + \sigma\theta = \sqrt{2\gamma} + \sigma\theta$, and yield that

$$M(\theta) = \Phi(-\lambda) e^{-\gamma} e^{\lambda^2/2} + \Phi(\sqrt{2\gamma}) \quad \text{and} \quad e^{\theta\sigma\sqrt{2\gamma} + \theta^2\sigma^2/2} = e^{-\gamma} e^{\lambda^2/2}.$$

Accordingly, the chosen $\lambda = \sqrt{2\gamma} + \sigma\theta$ should satisfy

$$\frac{\partial \left(\left[\Phi(-\lambda) e^{-\gamma} e^{\lambda^2/2} + \Phi(\sqrt{2\gamma}) \right]^{n-d} e^{d(-\gamma + \lambda^2/2)} \right)}{\partial \lambda} = 0,$$

or equivalently,

$$e^{(1/2)\lambda^2} \Phi(-\lambda) = \frac{1}{\sqrt{2\pi}\lambda} \left(1 - \frac{d}{n} \right) - \frac{d}{n} e^{\gamma} \Phi(\sqrt{2\gamma}). \quad (5)$$

As it turns out, the solution $\lambda = \lambda(\gamma)$ of the above equation depends only on γ . Now, by replacing $e^{(1/2)\lambda^2} \Phi(-\lambda)$ with $(1 - d/n) / (\sqrt{2\pi}\lambda) - (d/n) e^{\gamma} \Phi(\sqrt{2\gamma})$, we obtain

$$\begin{aligned}
\tilde{\sigma}^2(\lambda) &\triangleq \frac{\text{Var} \left[X_{d+1}^{(\theta)} \right]}{\sigma^2} \Bigg|_{\theta = (\lambda - \sqrt{2\gamma})/\sigma} \\
&= -\frac{d}{n-d} - \frac{nd}{(n-d)^2} \lambda^2 + \frac{n}{n-d} \frac{1}{1 + \sqrt{2\pi}\lambda e^{\gamma} \Phi(\sqrt{2\gamma})},
\end{aligned}$$

and

$$\begin{aligned}
\tilde{\rho}(\lambda) &\triangleq \frac{E \left[\left| X_{d+1}^{(\theta)} - \hat{\mu} \right|^3 \right]}{\sigma^3} \Bigg|_{\theta=(\lambda-\sqrt{2\gamma})/\sigma} \\
&= \frac{n}{(n-d)} \frac{\lambda}{[1 + \sqrt{2\pi}\lambda e^\gamma \Phi(\sqrt{2\gamma})]} \left\{ 1 - \frac{d(n+d)}{(n-d)^2} \lambda^2 \right. \\
&\quad + 2 \left[\frac{n^2}{(n-d)^2} \lambda^2 + 2 \right] e^{-d(2n-d)\lambda^2/[2(n-d)^2]} \\
&\quad - \frac{d}{n-d} \left[\frac{n+d}{n-d} \lambda^2 + 3 \right] \sqrt{2\pi}\lambda e^\gamma \Phi(\sqrt{2\gamma}) \\
&\quad \left. - \frac{2n}{n-d} \left[\frac{n^2}{(n-d)^2} \lambda^2 + 3 \right] \sqrt{2\pi}\lambda e^{\lambda^2/2} \Phi\left(-\frac{n}{n-d}\lambda\right) \right\}
\end{aligned}$$

Hence, the previously obtained upper bound for $\Pr(Y_n \leq 0)$ can be reformulated as

$$\begin{aligned}
&\Phi\left(-\frac{(n-d)\hat{\mu} + d\sqrt{2\gamma}}{\sqrt{d}}\right) \\
&+ \tilde{A}_{n,d}(\lambda) \left[\Phi(-\lambda)e^{-\gamma}e^{\lambda^2/2} + \Phi(\sqrt{2\gamma}) \right]^{n-d} e^{d(-\gamma+\lambda^2/2)} \Phi\left(\frac{(n-d)\hat{\mu} + \lambda d}{\sqrt{d}}\right),
\end{aligned}$$

where

$$\begin{aligned}
&\tilde{A}_{n,d}(\lambda) \\
&\triangleq A_{n-d}\left(\frac{\lambda - \sqrt{2\gamma}}{\sigma}\right) \\
&= \min\left(e^{(n-d)(\lambda-\sqrt{2\gamma})^2\tilde{\sigma}^2(\lambda)/2} \Phi\left((\lambda - \sqrt{2\gamma})\tilde{\sigma}(\lambda)\sqrt{n-d}\right) + 1.531 \frac{\tilde{\rho}(\lambda)}{\tilde{\sigma}^3(\lambda)\sqrt{n-d}}, 1\right).
\end{aligned}$$

Finally, a simple derivation yields

$$\begin{aligned}
E[Y_n] &= dE[X_1] + (n-d)E[X_{d+1}] \\
&= \sigma \left(d\sqrt{2\gamma} + (n-d) \left[-(1/\sqrt{2\pi})e^{-\gamma} + \sqrt{2\gamma}\Phi(-\sqrt{2\gamma}) \right] \right),
\end{aligned}$$

and hence, the condition of $E[Y_n] > 0$ can be equivalently replaced by

$$\frac{d}{n} \geq 1 - \frac{\sqrt{4\pi\gamma}e^\gamma}{1 + \sqrt{4\pi\gamma}e^\gamma\Phi(\sqrt{2\gamma})}.$$

■

The third argument γ in function $\mathcal{B}(\cdot, \cdot, \cdot)$ is in fact the signal-to-noise ratio for the decoding problems considered in this paper, and is reasonably assumed to be always positive.

III. ANALYSIS OF THE COMPUTATIONAL EFFORT OF THE SIMPLIFIED GENERALIZED DIJKSTRA'S ALGORITHM

In 1993, a novel and fast maximum-likelihood soft-decision decoding algorithm for linear block codes was proposed in [11], and was called the generalized Dijkstra's algorithm (GDA). Computer simulations have shown that the algorithm is highly efficient (that is, with small average computational effort) for certain number of linear block codes [5], [11]. Improvements of the GDA have been subsequently reported [1], [5], [8], [10], [12], [19], [24].

The authors of [12] proposed an upper bound on the average computational effort of an ordering-free variant of the GDA for linear block codes antipodally transmitted via the AWGN channel; however, the bound is valid only for codes with sufficiently large codeword length. In terms of the Berry-Esseen inequality, an alternative upper bound that holds for *any* (thus including, *small*) codeword length can be given.

A. Notations and definitions

Let \mathcal{C} be an (n, k) binary linear block code with codeword length n and dimension k , and let $R \triangleq k/n$ be the *code rate* of \mathcal{C} . Denote the codeword of \mathcal{C} by $\mathbf{x} \triangleq (x_0, x_1, \dots, x_{n-1})$. Also, denote by $\mathbf{r} = (r_0, r_1, \dots, r_{n-1})$ the received vector due to a codeword of \mathcal{C} is transmitted via a time-discrete memoryless channel.

From [3] (also [25], [26]), the *maximum-likelihood* (ML) estimate $\hat{\mathbf{x}} = (\hat{x}_0, \hat{x}_1, \dots, \hat{x}_{n-1})$ for a time-discrete memoryless channel, upon the receipt of \mathbf{r} , satisfies

$$\sum_{j=0}^{n-1} (\phi_j - (-1)^{\hat{x}_j})^2 \leq \sum_{j=0}^{n-1} (\phi_j - (-1)^{x_j})^2 \quad \text{for all } \mathbf{x} \in \mathcal{C}, \quad (6)$$

where $\phi_j \triangleq \ln[\Pr(r_j|0)/\Pr(r_j|1)]$. An immediate implication of equation (6) is that using the log-likelihood ratio vector $\boldsymbol{\phi} = (\phi_0, \phi_1, \dots, \phi_{n-1})$ rather than the received vector \mathbf{r} is sufficient in ML decoding.

When the linear block code is antipodally transmitted through the AWGN channel, the relationship between the binary codeword \mathbf{x} and the received vector \mathbf{r} can be characterized by

$$r_j = (-1)^{x_j} \sqrt{E} + e_j \quad \text{for } 0 \leq j \leq n-1, \quad (7)$$

where E is the signal energy per channel bit, and e_j represents a noise sample of a Gaussian process with single-sided noise power per hertz N_0 . The signal-to-noise ratio for the channel is

therefore $\gamma \triangleq E/N_0$. In order to account for the code redundancy for different code rates, the SNR per information bit $\gamma_b = \gamma/R$ is used instead of γ in the following discussions.

A *code tree* of an (n, k) binary linear block code is formed by representing every codeword as a *code path* on a binary tree of $(n+1)$ levels. A *code path* is a particular *path* that begins at the *start node* at level 0, and ends at one of the *leaf nodes* at level n . There are two branches, respectively labeled by 0 and 1, that leave each node at the first k levels. The remaining nodes at levels k through $(n-1)$ consist of only a single leaving branch. The 2^k rightmost nodes at level n are referred to as *goal nodes*. In notation, $\mathbf{x}_{[\ell]}$ is used to denote a path labeled by $(x_0, x_1, \dots, x_{\ell-1})$. For notational convenience, the subscript “ $[n]$ ” is dropped for the label sequence of a code path, namely $\mathbf{x}_{[n]}$ is briefed by \mathbf{x} . The same notational convention is adopted for other notation including the received vector \mathbf{r} and the log-likelihood ratio vector ϕ .

B. Brief description of the GDA

For completeness, we brief the GDA decoding algorithm in [11] in this subsection.

After obtaining the log-likelihood ratio vector $\phi = (\phi_0, \phi_1, \dots, \phi_{n-1})$, the GDA algorithm first permutes the positions of codeword components such that the codeword component that corresponds to larger absolute value of log-likelihood ratio appears earlier in its position whenever possible, and still the first k positions uniquely determine a code path. The post-permutation codewords thereby result in a new code tree \mathcal{C}^* . Let $\phi^* \triangleq (\phi_0^*, \phi_1^*, \dots, \phi_{n-1}^*)$ be the new log-likelihood ratio vector after permutation, and define the path metric of a path $\mathbf{x}_{[\ell]}$ (over the new code tree \mathcal{C}^*) as $\sum_{j=0}^{\ell-1} (\phi_j^* - (-1)^{x_j})^2$. The path metric of a code path \mathbf{x} is thus given by $\sum_{j=0}^{n-1} (\phi_j^* - (-1)^{x_j})^2$. The algorithm then searches for the code path with the minimum path metric over \mathcal{C}^* , which, from equation (6), is exactly the code path labeled by the permuted ML codeword. As expected, the final step of the algorithm is to output the de-permuted version of the labels of the minimum-metric code path.

The search process of the GDA algorithm is guided by an evaluation function $f(\cdot)$, defined for all paths of a code tree. A simple evaluation function [10] that guarantees the ultimate finding of the minimum-metric code path is

$$f(\mathbf{x}_{[\ell]}|\phi^*) = \sum_{j=0}^{\ell-1} (\phi_j^* - (-1)^{x_j})^2 + \sum_{j=\ell}^{n-1} (|\phi_j^*| - 1)^2.$$

Hence, when a path $\mathbf{x}_{[\ell]}$ is extended to its immediate successor path $\mathbf{x}_{[\ell+1]}$, the evaluation function value is updated by adding the branch metric, $(\phi_\ell^* - (-1)^{x_\ell})^2 - (|\phi_\ell^*| - 1)^2$, to its original value. The algorithm begins the search from the path that contains only the start node. It then extends, among the paths that have been visited, the path with the smallest f -function value. Once the algorithm chooses to extend a path that ends at a goal node, the search process terminates. Notably, any path that ends at level k has already uniquely determined a code path. Hence, once a length- k path is visited and the f -function value associated with its respective code path does not exceed the associated f -function value of any of the later top paths in the stack, the algorithm can ensure that this code path is the targeted one with the minimum code path metric. This indicates that the computational complexity of the GDA is dominantly contributed by those paths up to level k . This justifies our later analysis of the decoding complexity of the GDA, where only the computations due to those paths up to level k are considered.

The *simplified GDA algorithm* is a simplified variant of the GDA algorithm that does not perform permutation. In other words, its codeword search is operated over the unpermuted original code tree \mathcal{C} . Although both algorithms yield the same output, the simplified one was demonstrated to involve a larger branch metric computational load [12].

C. Analysis of the computational effort of the simplified GDA

The computational effort of the simplified GDA can now be analyzed.

Theorem 1 (Complexity of the simplified GDA): Consider an (n, k) binary linear block code antipodally transmitted via an AWGN channel. The average number of branch metric computations evaluated by the simplified GDA, denoted by $L_{\text{SGDA}}(\gamma_b)$, is upper-bounded by

$$L_{\text{SGDA}}(\gamma_b) \leq 2 \sum_{\ell=0}^{k-1} \sum_{d=0}^{\ell} \binom{\ell}{d} \mathcal{B}(d, n - \ell, k\gamma_b/n), \quad (8)$$

where function $\mathcal{B}(\cdot, \cdot, \cdot)$ is defined in Lemma 2.

Proof: Assume without loss of generality that the all-zero codeword $\mathbf{0}$ is transmitted.

Let \mathbf{x}^* label the minimum-metric code path for a given log-likelihood ratio vector ϕ . Then we quote from [12] that for any path $\mathbf{x}_{[\ell]}$ selected for extension by the simplified GDA,

$$f(\mathbf{x}_{[\ell]}|\phi) \leq \sum_{j=0}^{n-1} (\phi_j - (-1)^{x_j^*})^2,$$

which implies that for $\ell < k$,

$$\begin{aligned}
& \Pr [\text{path } \mathbf{x}_{[\ell]} \text{ is extended by the simplified GDA}] \\
& \leq \Pr \left[f(\mathbf{x}_{[\ell]} | \boldsymbol{\phi}) \leq \sum_{j=0}^{n-1} (\phi_j - (-1)^{x_j^*})^2 \right] \\
& \leq \Pr \left[f(\mathbf{x}_{[\ell]} | \boldsymbol{\phi}) \leq \sum_{j=0}^{n-1} (\phi_j - (-1)^0)^2 \right], \tag{9}
\end{aligned}$$

$$= \Pr \left[\sum_{j=0}^{\ell-1} (\phi_j - (-1)^{x_j})^2 + \sum_{j=\ell}^{n-1} (|\phi_j| - 1)^2 \leq \sum_{j=0}^{n-1} (\phi_j - 1)^2 \right], \tag{10}$$

where (9) follows from the assumption that the path metric of the \mathbf{x}^* -labeled code path is the smallest with respect to $\boldsymbol{\phi}$, and hence, does not exceed that of the $\mathbf{0}$ -labeled code path.

Now denote by $\mathcal{J} = \mathcal{J}(\mathbf{x}_{[\ell]})$ the set of index j , where $0 \leq j \leq \ell - 1$, for which $x_j = 1$. Then (10) can be rewritten as

$$\begin{aligned}
& \Pr [\text{path } \mathbf{x}_{[\ell]} \text{ is extended by the simplified GDA}] \\
& \leq \Pr \left[\sum_{j \in \mathcal{J}} \phi_j + \sum_{j=\ell}^{n-1} \min(\phi_j, 0) \leq 0 \right] \\
& = \Pr \left[\sum_{j \in \mathcal{J}} r_j + \sum_{j=\ell}^{n-1} \min(r_j, 0) \leq 0 \right] \tag{11}
\end{aligned}$$

where (11) holds since for the AWGN channel specified in (7), $\phi_j = 4\sqrt{E}r_j/N_0$. As the all-zero codeword is assumed to be transmitted, r_j is Gaussian distributed with mean \sqrt{E} and variance $N_0/2$. Hence, Lemma 2 can be applied to obtain

$$\Pr [\text{path } \mathbf{x}_{[\ell]} \text{ is extended by the simplified GDA}] \leq \mathcal{B}(d, n - \ell, R\gamma_b),$$

where $d = |\mathcal{J}|$ is the Hamming weight of $\mathbf{x}_{[\ell]}$.

Observe that the extension of each path that ends at level ℓ , where $\ell < k$, causes two branch metric computations. Therefore, the expectation value of the number of branch metric evaluations satisfies

$$L_{\text{SGDA}}(\gamma_b) \leq 2 \sum_{\ell=0}^{k-1} \sum_{d=0}^{\ell} \binom{\ell}{d} \mathcal{B}(d, n - \ell, R\gamma_b).$$

■

D. Numerical and simulation results

The accuracy of the previously derived theoretical upper bound for the average computational effort of the simplified GDA is now empirically studied. Two linear block codes are considered — one is a $(24, 12)$ binary extended Golay code, and the other is a $(48, 24)$ binary extended quadratic residue code.

Figures 1 and 2 illustrate the deviation between the simulated results and the theoretical upper bound in Theorem 1. According to these figures, the theoretical upper bound is quite close to the simulation results for high γ_b (above 8 dB). In such a case, the computational complexity of the simplified GDA reduces to its minimum possible values, 24 and 48, for $(24, 12)$ and $(48, 24)$ codes, respectively. As γ_b reaches 1 dB, the theoretical bound for $(48, 24)$ code is around 12 times higher than the simulated average complexity. However, for the $(24, 12)$ code, the theoretical bound and the simulation results differ only by 0.671638 on a \log_{10} scale at $\gamma_b = 1$ dB, and it is when $\gamma_b \leq -8$ dB that the upper bound becomes ten times larger than the simulated complexity. The conclusion section will address the possible cause of the inaccuracy of the theoretical bounds at low SNR.

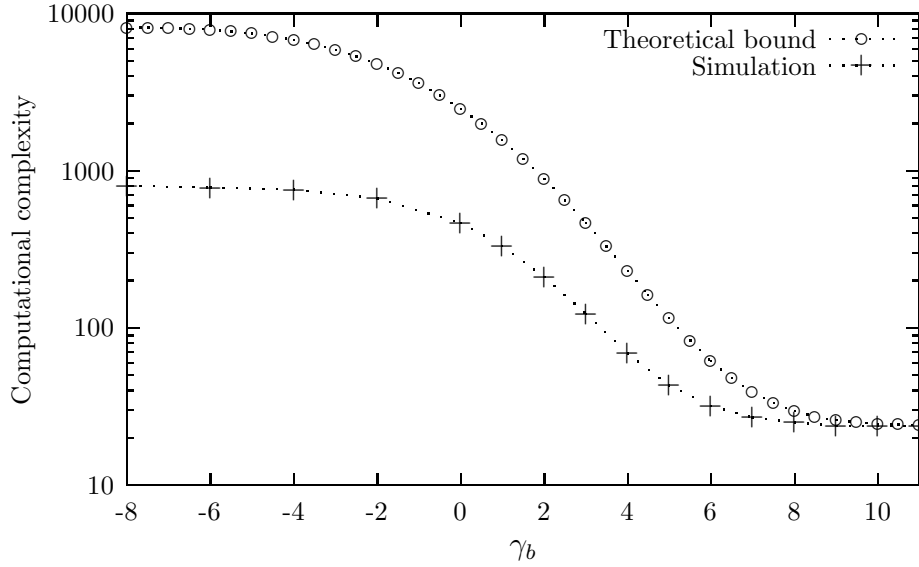


Fig. 1. Average computational complexity of the simplified GDA for $(24, 12)$ binary extended Golay code.

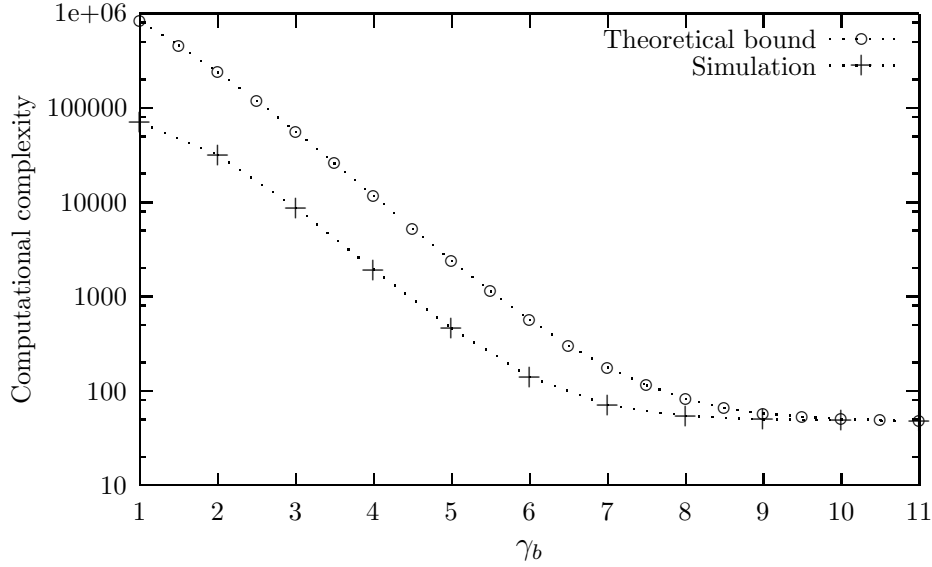


Fig. 2. Average computational complexity of the simplified GDA for (48, 24) binary extended quadratic residue code.

IV. ANALYSIS OF THE COMPUTATIONAL EFFORT OF THE MLSDA

Based on the probability bound established in Lemma 2, the computational complexity of the maximum-likelihood sequential decoding algorithm (MLSDA) proposed in [13] is analyzed for convolutional codes antipodally transmitted via the AWGN channel.

A. Notation and definitions

Let \mathcal{C} be an (n, k, m) binary convolutional code, where k is the number of encoder inputs, n is the number of encoder outputs, and m is its *memory order* defined as the maximum number of shift register stages from an encoder input to an encoder output. Let $R \triangleq k/n$ and $N \triangleq n(L + m)$ be the *code rate* and the *code length* of \mathcal{C} , respectively, where L represents the length of applied information sequence. Denote the codeword of \mathcal{C} by $\mathbf{x} \triangleq (x_0, x_1, \dots, x_{N-1})$. Also denote the left portion of codeword \mathbf{x} by $\mathbf{x}_{(b)} \triangleq (x_0, x_1, \dots, x_b)$. Assume that antipodal signaling is used in the codeword transmission such that the relationship between binary channel codeword \mathbf{x} and received vector $\mathbf{r} \triangleq (r_0, r_1, \dots, r_{N-1})$ is

$$r_j = (-1)^{x_j} \sqrt{E} + e_j, \quad 0 \leq j \leq N-1, \quad (12)$$

where E is the signal energy per channel bit, and e_j is a noise sample of a Gaussian process with single-sided noise power per hertz N_0 . The signal-to-noise ratio per information bit $\gamma_b = (EN)/(N_0kL)$ is again used to account for the code redundancy for various code rates.

A *trellis*, as depicted in Fig. 3 in terms of a specific example, can be obtained from a code tree by combining nodes with the same *state*. States are characterized by the content of the shift-register stages in a convolutional encoder. For convenience, the leftmost node (at level 0) and the rightmost node (at level $L + m$) of a trellis are named the *start node* and the *goal node*, respectively. A path on a trellis from the single start node to the single goal node is called a *code path*. Each branch in the trellis is labeled by an appropriate encoder output of length n .

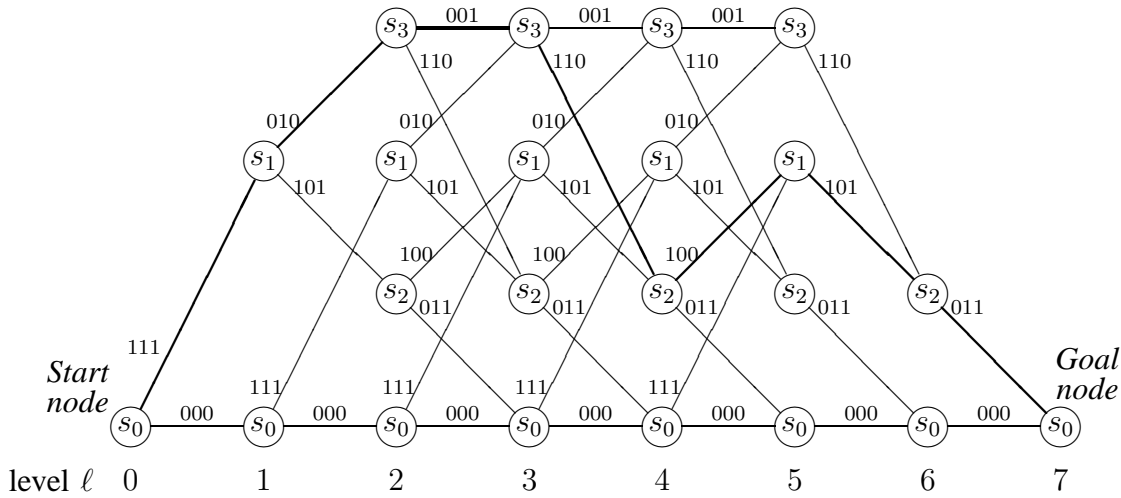


Fig. 3. Trellis for a $(3, 1, 2)$ binary convolutional code with information length $L = 5$. In this case, the code rate $R = 1/3$ and the codeword length $N = 3(5 + 2) = 21$. The code path indicated by the thick line is labeled by 111, 010, 001, 110, 100, 101 and 011, thus its corresponding codeword is $\mathbf{x} = (111010001110100101011)$.

B. Maximum-likelihood soft-decision sequential decoding algorithm (MLSDA)

In [13], a trellis-based sequential decoding algorithm specifically for binary convolutional codes is proposed. The same paper proves that the algorithm performs maximum-likelihood decoding, and is thus named the *maximum-likelihood sequential decoding algorithm* (MLSDA). Unlike the conventional sequential decoding algorithm [7], [15], [20], [27] which requires only

a single stack, the trellis-based MLSDA maintains two stacks — an *Open Stack* and a *Closed Stack*. For completeness, the algorithm is quoted below.

- Step 1. Put the path that contains only the start node into the Open Stack, and assign its path metric as zero.*
- Step 2. Compute the path metric for each of the successor paths of the top path in the Open Stack by adding the branch metric of the extended branch to the path metric of the top path. Put into the Closed Stack both the state and level of the end node of the top path in the Open Stack. Delete the top path from the Open Stack.*
- Step 3. Discard any successor path that ends at a node that has the same state and level as any entry in the Closed Stack. If any successor path merges¹ with a path already in the Open Stack, eliminate the path with higher path metric.*
- Step 4. Insert the remaining successor paths into the Open Stack in order of ascending path metrics.*
- Step 5. If the top path in the Open Stack ends at the single goal node, the algorithm stops; otherwise go to Step 2.*

In the MLSDA, the path metric for a path labeled by $\mathbf{x}_{(\ell_{n-1})} = (x_0, x_1, \dots, x_{\ell_{n-1}})$, upon receipt of $\phi_{(\ell_{n-1})}$, is given by

$$M(\mathbf{x}_{(\ell_{n-1})} | \phi_{(\ell_{n-1})}) \triangleq \sum_{j=0}^{\ell_{n-1}} (y_j \oplus x_j) \times |\phi_j|, \quad (13)$$

where $\phi_j = \log[\Pr(r_j|0)/\Pr(r_j|1)]$ is the j th received log-likelihood ratio, r_j is the j th received scalar, and $y_j = 1$ if $\phi_j < 0$ and $y_j = 0$, otherwise.

C. Analysis of the computational efforts of the MLSDA

Since the nodes at levels L through $(L+m-1)$ have only one branch leaving them, and L is typically much larger than m , the contribution of these nodes to the computational complexity due to path extensions can be reasonably neglected. Hence, the analysis in the following theorem only considers those branch metric computations applied up to level L of the trellis.

Notations that will be used in the next theorem are first introduced. Denote by $s_j(\ell)$ the node that is located at level ℓ and corresponds to state index j . Let $\mathcal{S}_j(\ell)$ be the set of paths

¹“Merging” of two paths means that the two paths end at the same node.

that end at node $s_j(\ell)$. Also let $\mathcal{H}_j(\ell)$ be the set of the Hamming weights of the paths in $\mathcal{S}_j(\ell)$. Denote the minimum Hamming weight in $\mathcal{H}_j(\ell)$ by $d_j^*(\ell)$. As an example, $\mathcal{S}_3(3)$ equals $\{111010001, 000111010\}$ in Fig. 3, which results in $\mathcal{H}_3(3) = \{5, 4\}$ and $d_3^*(3) = 4$.

Theorem 2 (Complexity of the MLSDA): Consider an (n, k, m) binary convolutional code transmitted via an AWGN channel. The average number of branch metric computations evaluated by the MLSDA, denoted by $L_{\text{MLSDA}}(\gamma_b)$, is upper-bounded by

$$L_{\text{MLSDA}}(\gamma_b) \leq 2^k \sum_{\ell=0}^{L-1} \sum_{j=0}^{2^m-1} \mathcal{B}\left(d_j^*(\ell), N - \ell n, \frac{kL}{N} \gamma_b\right),$$

where if $\mathcal{H}_j(\ell)$ is empty, implying the non-existence of state j at level ℓ , then $\mathcal{B}(d_j^*(\ell), N - \ell n, kL\gamma_b/N) = 0$.

Proof: Assume without loss of generality that the all-zero codeword $\mathbf{0}$ is transmitted.

First, observe that for any two paths that end at a common node, only one of them will survive in the Open Stack. In other words, one of the two paths will be discarded either due to a larger path metric or because its end node has the same state and level as an entry in the Closed Stack. In the latter case, the surviving path has clearly reached the common end node earlier, and has already been extended by the MLSDA at some previous time (so that the state and level of its end node has already been stored in the Closed Stack). Accordingly, unlike the code tree search in the GDA, the branch metric computations that follow these two paths will only be performed once. It therefore suffices to derive the computational complexity of the MLSDA based on the nodes that have been extended rather than the paths that have been extended.

Let \mathbf{x}^* label the minimum-metric code path for a given log-likelihood ratio ϕ . Then we claim that if a node $s_j(\ell)$ is extended by the MLSDA, given that $\mathbf{x}_{(\ell n-1)}$ is the only surviving path (in the Open Stack) that ends at this node at the time this node is extended, then

$$M(\mathbf{x}_{(\ell n-1)} | \phi_{(\ell n-1)}) \leq M(\mathbf{x}^* | \phi) \quad (14)$$

The validity of the above claim can be simply proved by contradiction. Suppose $M(\mathbf{x}_{(\ell n-1)} | \phi_{(\ell n-1)}) > M(\mathbf{x}^* | \phi)$. Then the non-negativity of the individual metric $(y_j \oplus x_j) | \phi_j |$, which implies $M(\mathbf{x}^* | \phi) \geq M(\mathbf{x}_{(b)}^* | \phi_{(b)})$ for every $0 \leq b \leq N - 1$, immediately gives $M(\mathbf{x}_{(\ell n-1)} | \phi_{(\ell n-1)}) > M(\mathbf{x}_{(b)}^* | \phi_{(b)})$ for every $0 \leq b \leq N - 1$. Therefore, $\mathbf{x}_{(\ell n-1)}$ cannot be on top of the Open Stack (because some $\mathbf{x}_{(b)}^*$ always exists in the Open Stack), and hence violates the assumption that $s_j(\ell)$ is extended by the MLSDA.

For notational convenience, denote by $\mathcal{A}(s_j(\ell), \mathbf{x}_{(\ell n-1)})$ the event that “ $\mathbf{x}_{(\ell n-1)}$ is the only path in the intersection of $\mathcal{S}_j(\ell)$ and the Open Stack at the time node $s_j(\ell)$ is extended.” Notably,

$$\{\mathcal{A}(s_j(\ell), \mathbf{x}_{(\ell n-1)})\}_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)}$$

are disjoint, and

$$\sum_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \{ \mathcal{A}(s_j(\ell), \mathbf{x}_{(\ell n-1)}) \} = 1.$$

Then according to the above claim,

$$\begin{aligned} & \Pr \{ \text{node } s_j(\ell) \text{ is extended by the MLSDA} \} \\ &= \sum_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \{ \mathcal{A}(s_j(\ell), \mathbf{x}_{(\ell n-1)}) \} \Pr \left\{ \begin{array}{l} \text{node } s_j(\ell) \text{ is extended} \\ \text{by the MLSDA} \end{array} \middle| \mathcal{A}(s_j(\ell), \mathbf{x}_{(\ell n-1)}) \right\} \\ &\leq \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \left\{ \begin{array}{l} \text{node } s_j(\ell) \text{ is extended} \\ \text{by the MLSDA} \end{array} \middle| \mathcal{A}(s_j(\ell), \mathbf{x}_{(\ell n-1)}) \right\} \\ &\leq \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \{ M(\mathbf{x}_{(\ell n-1)} | \phi_{(\ell n-1)}) \leq M(\mathbf{x}^* | \phi) \} \\ &\leq \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \{ M(\mathbf{x}_{(\ell n-1)} | \phi_{(\ell n-1)}) \leq M(\mathbf{0} | \phi) \} \\ &= \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \left\{ \sum_{j=0}^{\ell n-1} (y_j \oplus x_j) |\phi_j| \leq \sum_{j=0}^{N-1} (y_j \oplus 0) |\phi_j| \right\}, \end{aligned}$$

where the replacement of \mathbf{x}^* by the all-zero codeword $\mathbf{0}$ follows from $M(\mathbf{x}^* | \phi) \leq M(\mathbf{0} | \phi)$. We then observe that for the AWGN channel defined through (12), $\phi_j = 4\sqrt{E}r_j/N_0$; hence, y_j can be determined by

$$y_j = \begin{cases} 1, & \text{if } r_j < 0; \\ 0, & \text{otherwise.} \end{cases}$$

This observation, together with the fact that $2(y_j \oplus x_j)|r_j| = r_j[(-1)^{y_j} - (-1)^{x_j}]$, gives

$$\begin{aligned} & \Pr \{ \text{node } s_j(\ell) \text{ is extended by the MLSDA} \} \\ &\leq \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \left\{ \sum_{j=0}^{\ell n-1} (y_j \oplus x_j) |r_j| \leq \sum_{j=0}^{N-1} (y_j \oplus 0) |r_j| \right\}, \\ &= \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \left\{ \sum_{j=0}^{\ell n-1} r_j [(-1)^{y_j} - (-1)^{x_j}] \leq \sum_{j=0}^{N-1} r_j [(-1)^{y_j} - (-1)^0] \right\} \\ &= \max_{\mathbf{x}_{(\ell n-1)} \in \mathcal{S}_j(\ell)} \Pr \left\{ \sum_{j \in \mathcal{J}(\mathbf{x}_{(\ell n-1)})} r_j + \sum_{j=\ell n}^{N-1} \min(r_j, 0) \leq 0 \right\}, \end{aligned}$$

where $\mathcal{J}(\mathbf{x}_{(\ell n-1)})$ is the set of index j , where $0 \leq j \leq \ell n - 1$, for which $x_j = 1$. As r_j is Gaussian distributed with mean \sqrt{E} and variance $N_0/2$ due to the transmission of the all-zero codeword, Proposition 1 (in the Appendix) and Lemma 2 can be applied to obtain

$$\begin{aligned}
& \Pr \{ \text{node } s_j(\ell) \text{ is extended by the MLSDA} \} \\
& \leq \max_{d \in \mathcal{H}_j(\ell)} \Pr \left\{ r_1 + \cdots + r_d + \sum_{j=\ell n}^{N-1} \min(r_j, 0) \leq 0 \right\} \\
& = \Pr \left\{ r_1 + \cdots + r_{d_j^*(\ell)} + \sum_{j=\ell n}^{N-1} \min(r_j, 0) \leq 0 \right\} \\
& \leq \mathcal{B} \left(d_j^*(\ell), N - \ell n, \frac{kL}{N} \gamma_b \right).
\end{aligned}$$

Consequently,

$$L_{\text{MLSDA}}(\gamma_b) \leq 2^k \sum_{\ell=0}^{L-1} \sum_{j=0}^{2^m-1} \mathcal{B} \left(d_j^*(\ell), N - \ell n, \frac{kL}{N} \gamma_b \right),$$

where the multiplication of 2^k is due to the fact that whenever a node is extended, 2^k branch metric computations will follow. ■

D. Numerical and simulation results

The accuracy of the previously derived theoretical upper bound for the computational effort of the MLSDA is now empirically examined using two types of convolutional codes. One is a $(2, 1, 6)$ code with generators 634, 564 (octal); the other is a $(2, 1, 16)$ code with generators 1632044, 1145734 (octal). The lengths of the applied information bits are 60 and 100.

Figures 4–7 present the deviation between the simulated results and the theoretical upper bound on the computational complexity of the MLSDA. According to these figures, the theoretical upper bound is fairly close to the simulation results for both high γ_b (above 6 dB) and low γ_b (below 2 dB). Even for moderate γ_b , they only differ by no more than 0.355, 0.359, 0.573 and 0.580 respectively for Figs. 4–7 on a \log_{10} scale.

A side observation from these figures is that the codes with longer constraint length, although having a lower bit error rate, require more computations. However, such a tradeoff on constraint length and bit error rate can be moderately eased at high SNR. Notably, when $\gamma_b > 6$ dB, the average computational effort of the MLSDA in all four figures is reduced to approximately $2^k L$ in spite of the constraint length.

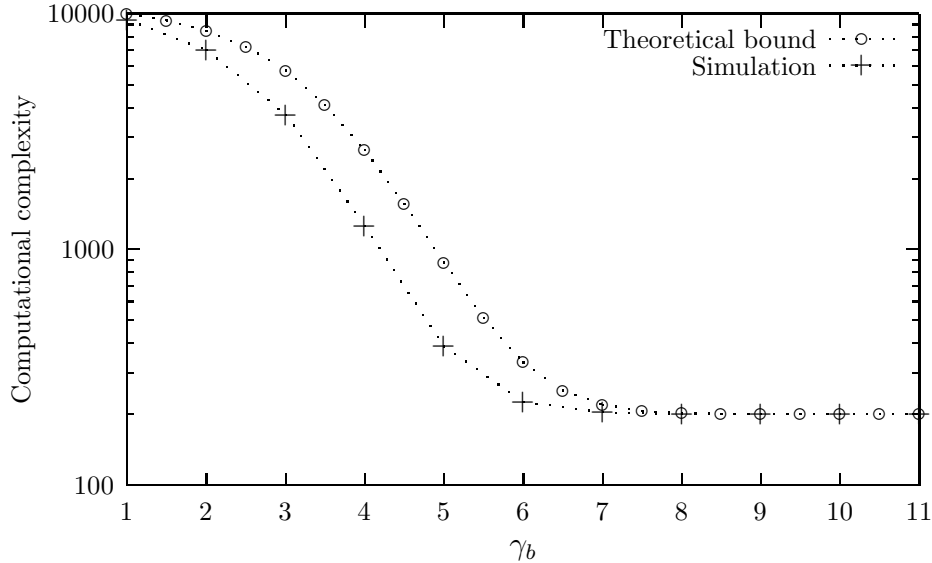


Fig. 4. Average computational complexity of the MLSDA for $(2, 1, 6)$ convolutional code with generators 634, 564 (octal) and information length $L = 100$.

V. CONCLUSIONS

In terms of the Berry-Esseen theorem, this study established theoretical upper bounds on the computational effort of the simplified GDA and the MLSDA for AWGN channels. Empirically investigating their accuracy for codes of small blocklengths shows that the theoretical bound is fairly close to the simulated results for the trellis-based MLSDA. However, for the tree-based simplified GDA, the upper complexity bound tends to be markedly larger than the simulated results at low SNR. Specifically, Figs. 4–7 reveal that adopting the Berry-Esseen technique yields a good upper bound on the decoding complexity of the MLSDA at low SNR; nevertheless, the two curves in Fig. 1, which depict the theoretical and simulated GDA decoding complexities for $(24, 12)$ code, do not seem to be close to each other even at very low SNR. This finding suggests that taking the expectation value of the computational efforts with respect to the path-extending probability bound in (11) over-estimates the true decoding complexity, and the path-extending probability may be much smaller than the bound in (11) due to a high error rate at low SNR. Yet, such an over-estimation can be moderately compensated for, if the codeword search is performed over a trellis, where “path-merging” significantly reduces the difference between the trellis node-extending probability and the theoretical bound.

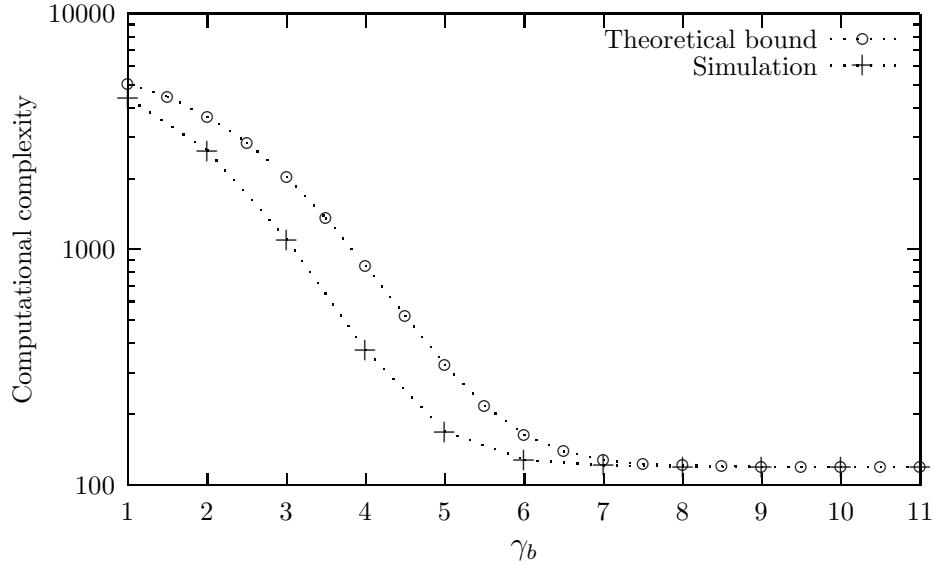


Fig. 5. Average computational complexity of the MLSDA for $(2, 1, 6)$ convolutional code with generators 634, 564 (octal) and information length $L = 60$.

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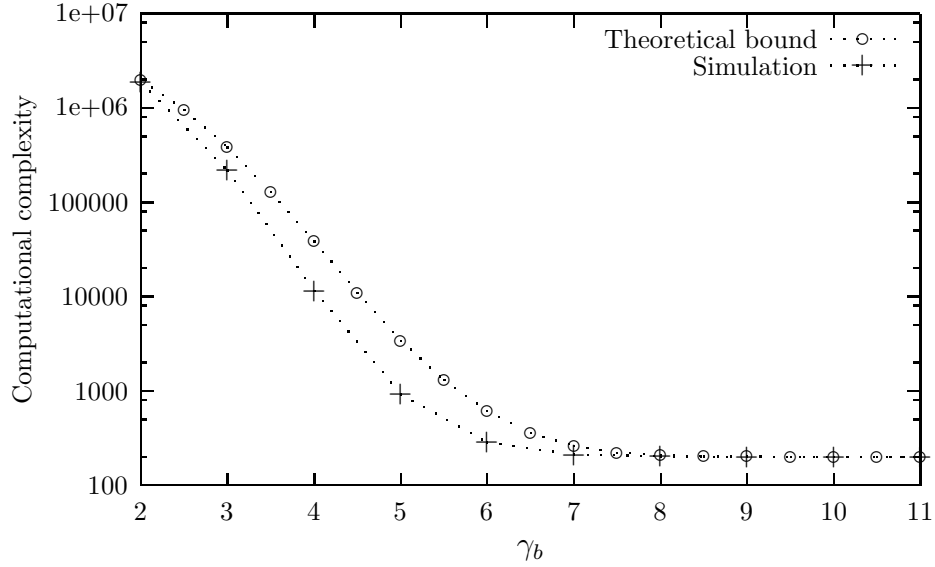


Fig. 6. Average computational complexity of the MLSDA for $(2, 1, 16)$ convolutional code with generators 1632044, 1145734 (octal) and information length $L = 100$.

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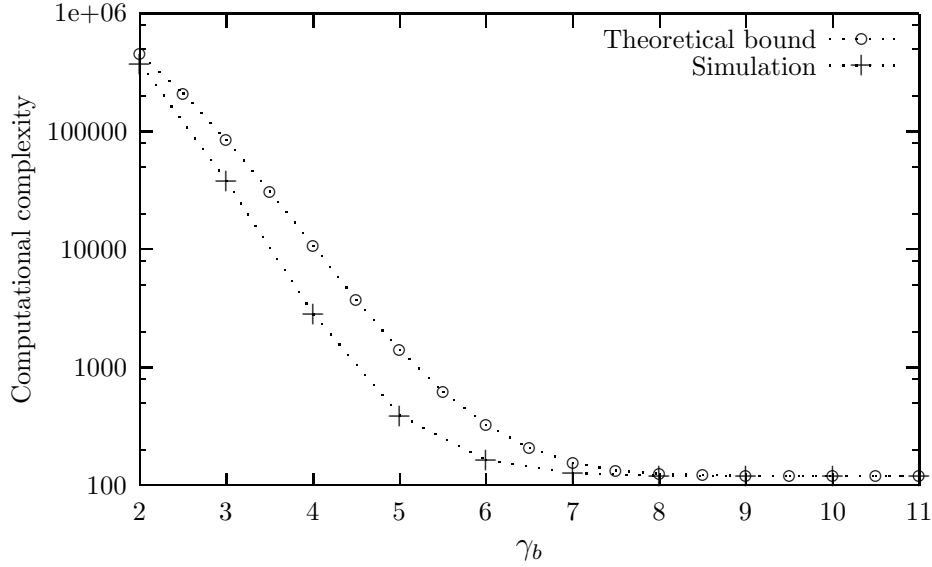


Fig. 7. Average computational complexity of the MLSDA for $(2, 1, 16)$ convolutional code with generators 1632044, 1145734 (octal) and information length $L = 60$.

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APPENDIX

Proposition 1: For a fixed non-negative integer k , the probability mass of

$$\Pr \{r_1 + \dots + r_d + \min(w_1, 0) + \dots + \min(w_k, 0) \leq 0\}$$

is a decreasing function for non-negative integer d , provided that $r_1, r_2, \dots, r_d, w_1, w_2, \dots, w_k$ are i.i.d. with a Gaussian marginal distribution of positive mean μ and variance σ^2 .

Proof: Assume without loss of generality that $\sigma^2 = 1$. Also, assume $k \geq 1$ since the proposition is trivially valid for $k = 0$.

Let $\Omega_d \triangleq r_1 + \dots + r_d$. Denote the probability density function of w_1 by $f(\cdot)$. Then putting $\nu \triangleq \Pr\{w_j = 0\}$ yields

$$\begin{aligned}
& \Pr\{\Omega_d + w_1 + w_2 + \dots + w_k \leq 0\} \\
&= \sum_{j=0}^k \Pr\{\text{exactly } (k-j) \text{ zeros in } (w_1, w_2, \dots, w_k)\} \\
& \quad \Pr\{\Omega_d + w_1 + w_2 + \dots + w_k \leq 0 \mid \text{exactly } (k-j) \text{ zeros in } (w_1, w_2, \dots, w_k)\} \\
&= \binom{k}{0} \nu^k \Pr\{\Omega_d \leq 0\} + \binom{k}{1} \nu^{k-1} (1-\nu) \int_{-\infty}^0 f(x) \Pr\{\Omega_d \leq -x\} dx \\
& \quad + \binom{k}{2} \nu^{k-2} (1-\nu)^2 \int_{-\infty}^0 \int_{-\infty}^0 f(x_1) f(x_2) \Pr\{\Omega_d \leq -(x_1 + x_2)\} dx_1 dx_2 \\
& \quad + \dots \\
& \quad + \binom{k}{k} (1-\nu)^k \int_{-\infty}^0 \dots \int_{-\infty}^0 f(x_1) \dots f(x_k) \Pr\{\Omega_d \leq -(x_1 + \dots + x_k)\} dx_1 \dots dx_k.
\end{aligned}$$

Accordingly, if each of the above $(k+1)$ terms is non-increasing in d , so is their sum. Let

$$\begin{aligned}
q_j(d) &\triangleq \int_{-\infty}^0 \dots \int_{-\infty}^0 f(x_1) \dots f(x_j) \Pr\{\Omega_d \leq -(x_1 + \dots + x_j)\} dx_1 \dots dx_j \\
&= \int_{-\infty}^0 \dots \int_{-\infty}^0 f(x_1) \dots f(x_j) \Phi\left(-\frac{x_1 + \dots + x_j}{\sqrt{d}} - \sqrt{d}\mu\right) dx_1 \dots dx_j.
\end{aligned}$$

Then

$$\begin{aligned}
\frac{\partial q_j(d)}{\partial(\sqrt{d})} &= \int_{-\infty}^0 \dots \int_{-\infty}^0 f(x_1) \dots f(x_j) \\
& \quad \times \left(\frac{x_1 + \dots + x_j}{d} - \mu\right) \frac{1}{\sqrt{2\pi}} e^{-(x_1 + \dots + x_j + d\mu)^2/(2d)} dx_1 \dots dx_j \\
&\leq -\frac{\mu}{\sqrt{2\pi}} \int_{-\infty}^0 \dots \int_{-\infty}^0 f(x_1) \dots f(x_j) e^{-(x_1 + \dots + x_j + d\mu)^2/(2d)} dx_1 \dots dx_j \quad (15) \\
&< 0,
\end{aligned}$$

where (15) follows from $x_i \leq 0$ (according to the range of integration) for $1 \leq i \leq j$. Consequently, $q_j(d)$ is decreasing in d for d positive and every $1 \leq j \leq k$. The proof is completed by noting that the first term, $\Pr\{\Omega_d \leq 0\} = \Phi(-\sqrt{d}\mu)$, is also decreasing in d . ■