

Derivadas de lenguajes regulares

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20 de agosto de 2015



Derivadas

Preliminares

Definición

La derivada de un lenguaje $\mathcal{L} \subseteq \Sigma^*$ con respecto a una cadena $u \in \Sigma^*$ se define como

$$\partial_u \mathcal{L} = \{v | u \cdot v \in \mathcal{L}\}$$

Ejemplo

$$\mathcal{L} = \{ab, abb, abbb, abbbb, abbbbb, abbbbbbb, \dots\} = \mathcal{L}[(ab) \cdot b^*]$$

$$\partial_{ab} \mathcal{L} = \{\varepsilon, b, bb, bbb, bbbb, bbbbb, \dots\} = \mathcal{L}[b^*]$$

Teorema

Si $\mathcal{L} \subseteq \Sigma^*$ es regular, entonces $\partial_u \mathcal{L}$ también es regular.

Derivadas

Problema

Problema

Dada una expresión regular r y una cadena $u \in \Sigma^*$, determinar si $u \in L[\![r]\!]$.

Resolver el problema es equivalente a resolver:

$$u \in L[\![r]\!] \Leftrightarrow \varepsilon \in L[\!\partial_u r\!]$$

Lo cual sucede si y sólo si

$$\nu(\partial_u r) = \varepsilon$$

Derivadas

Decimos que una expresión regular r es anulable, si el lenguaje definido contiene a la cadena vacía, esto es si $\varepsilon \in \mathcal{L}[\![r]\!]$ la función $\nu : RE \rightarrow RE$ tiene la siguiente propiedad:

$$\nu(r) = \begin{cases} \varepsilon & \text{Si } r \text{ es anulable} \\ \emptyset & \text{otro caso} \end{cases}$$

Y está definida como sigue

$$\nu(\varepsilon) = \varepsilon$$

$$\nu(\emptyset) = \emptyset$$

$$\nu(a) = \emptyset$$

$$\nu(r \cdot s) = \nu(r) \cap \nu(s)$$

$$\nu(r + s) = \nu(r) + \nu(s)$$

$$\nu(r^*) = \varepsilon$$

Derivadas

Problema

Usando lo anterior y la definición de ∂_u , generamos un algoritmo que verifique si $u \in L[[r]]$.

Definición

Expresandolo en términos de la relación $r \sim u$ (u caza con r).

$$r \sim \varepsilon \Leftrightarrow v(r) = \varepsilon$$

$$r \sim a \cdot \omega \Leftrightarrow \partial_a r \sim \omega$$

Teorema

$$r \sim u \Leftrightarrow u \in L[[r]]$$

Derivadas

Expresiones regulares

Reglas de Brzozowski para expresiones regulares con respecto a un símbolo $a \in \Sigma^*$.

$$\partial_a \epsilon = \emptyset$$

$$\partial_a a = \epsilon$$

$$\partial_a b = \emptyset \text{ si } b \neq a$$

$$\partial_a \emptyset = \emptyset$$

$$\partial_a(r \cdot s) = (\partial_a r) \cdot s + v(r) \cdot \partial_a s$$

$$\partial_a(r + s) = \partial_a r + \partial_a s$$

$$\partial_a(r^*) = \partial_a r \cdot r^*$$

La reglas se puede extender a cadenas como:

$$\partial_\epsilon r = r$$

$$\partial_{u \cdot a} r = \partial_a(\partial_u r)$$

Derivadas

Ejemplos

$\dot{c}abb \in \mathcal{L}[a \cdot b^*]?$	$\dot{c}aba \in \mathcal{L}[a \cdot b^*] ?$
$a \cdot b^* \sim abb$	$a \cdot b^* \sim aba$

Derivadas

Ejemplos

$\dot{c}abb \in \mathcal{L}[[a \cdot b^*]]?$	$\dot{c}aba \in \mathcal{L}[[a \cdot b^*]] ?$
$a \cdot b^* \sim abb$ $\Leftrightarrow \partial_a(a \cdot b^*) \sim bb$	$a \cdot b^* \sim aba$ $\Leftrightarrow \partial_a(a \cdot b^*) \sim ba$

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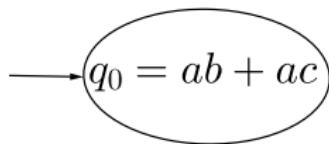
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\Sigma = \{a, b, c\}$

1 $q_0 = \partial_\varepsilon(a \cdot b + a \cdot c) = a \cdot b + a \cdot c$



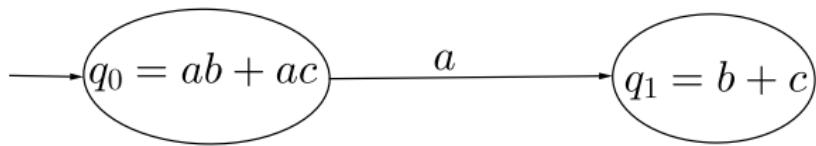
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

- ② $\partial_a q_0 = \partial_a(a \cdot b + a \cdot c) = b + c$, el cual es un estado nuevo, lo llamamos q_1



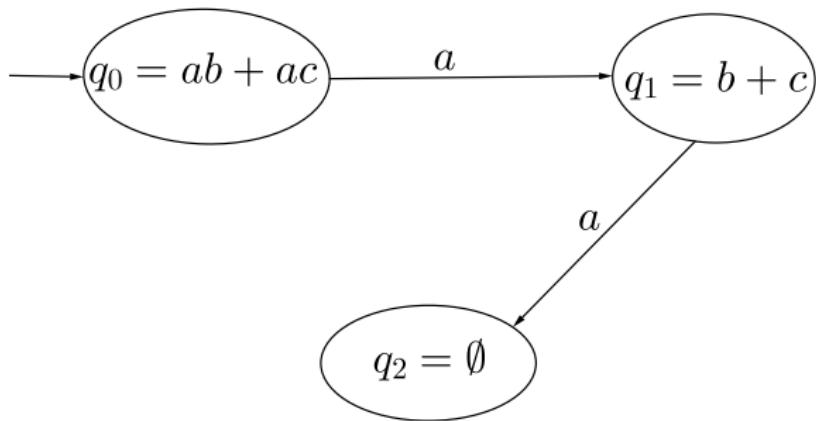
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- ③ $\partial_a q_1 = \partial_a(b + c) = \emptyset$, el cual es un nuevo estado y lo llamaremos q_2



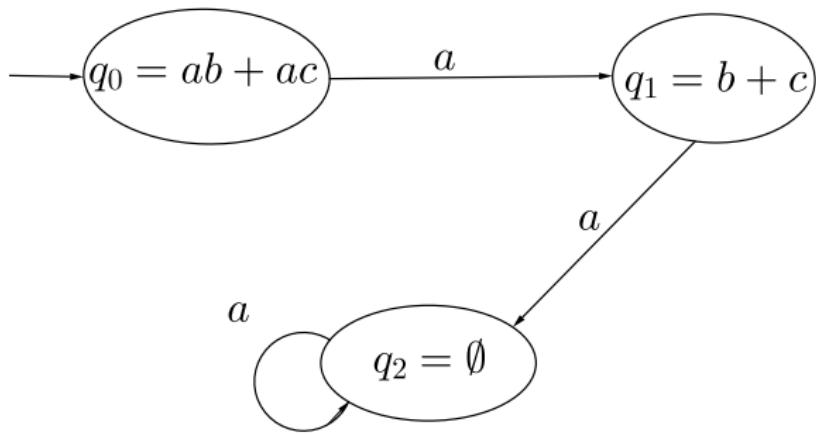
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Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

- ④ $\partial_a q_2 = \partial_a(\emptyset) = \emptyset = q_2$, pues el estado ya existe



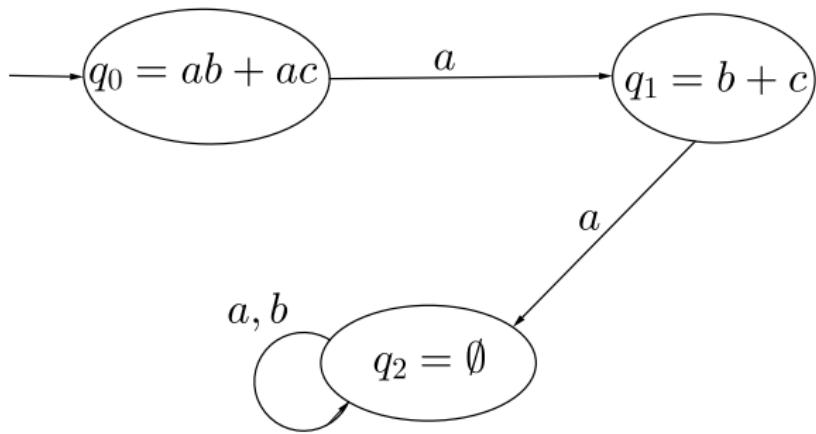
Algoritmo

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Ejemplo

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5 $\partial_b q_2 = \partial_b(\emptyset) = \emptyset = q_2$



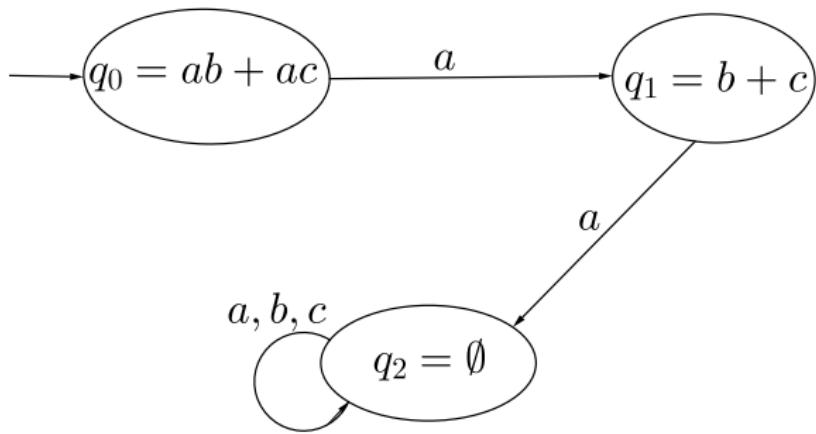
Algoritmo

Construcción de DFA

Ejemplo

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⑥ $\partial_c q_2 = \partial_c(\emptyset) = \emptyset = q_2$



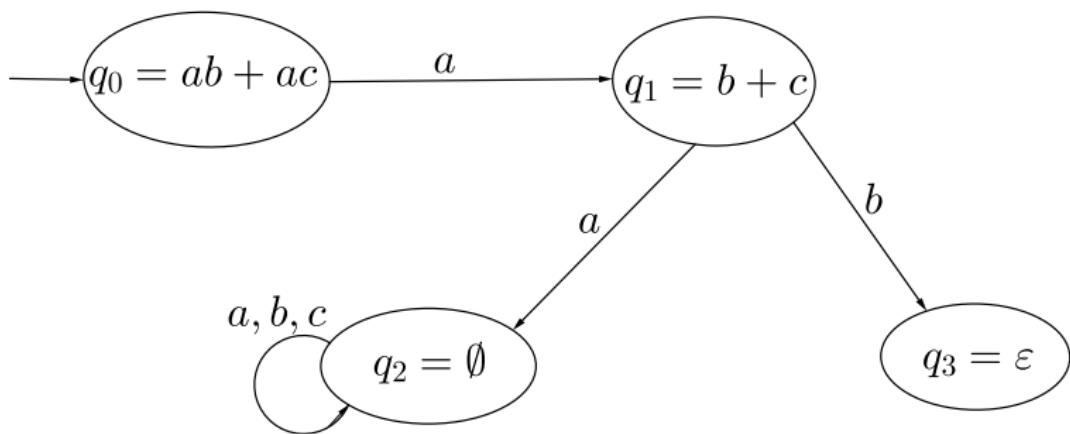
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Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

- ⑦ $\partial_b q_1 = \partial_b(b + c) = (\varepsilon + \emptyset) \equiv \varepsilon$, es un nuevo estado llamado q_3



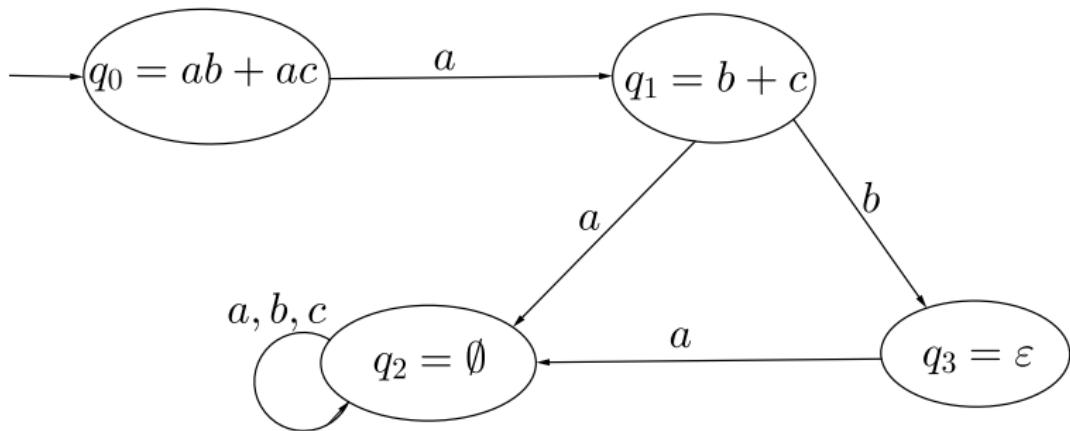
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⑧ $\partial_a q_3 = \partial_a(\varepsilon) = \emptyset = q_2$



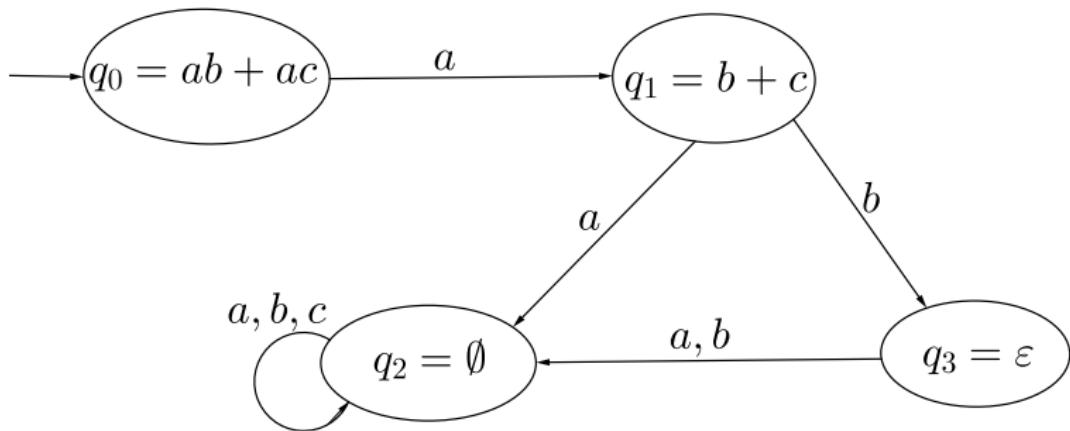
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Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

9 $\partial_b q_3 = \partial_b(\varepsilon) = \emptyset = q_2$



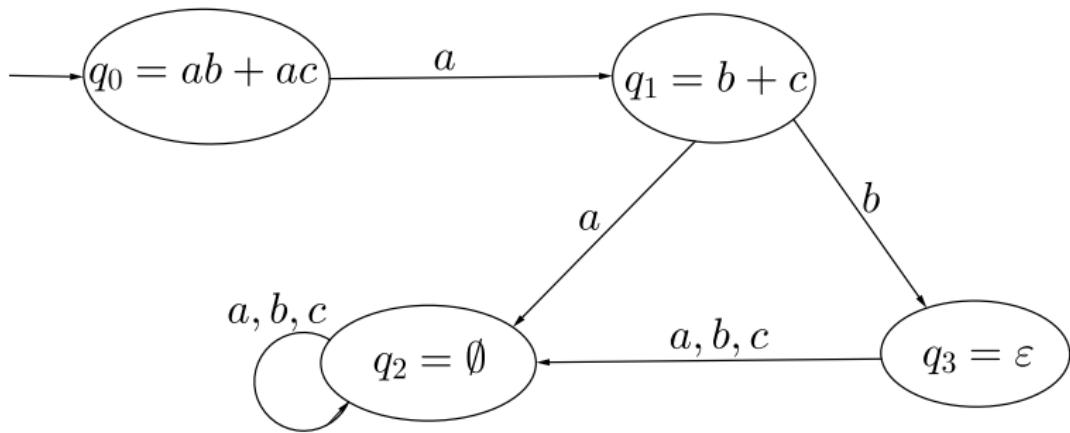
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

10 $\partial_c q_3 = \partial_c(\varepsilon) = \emptyset = q_2$



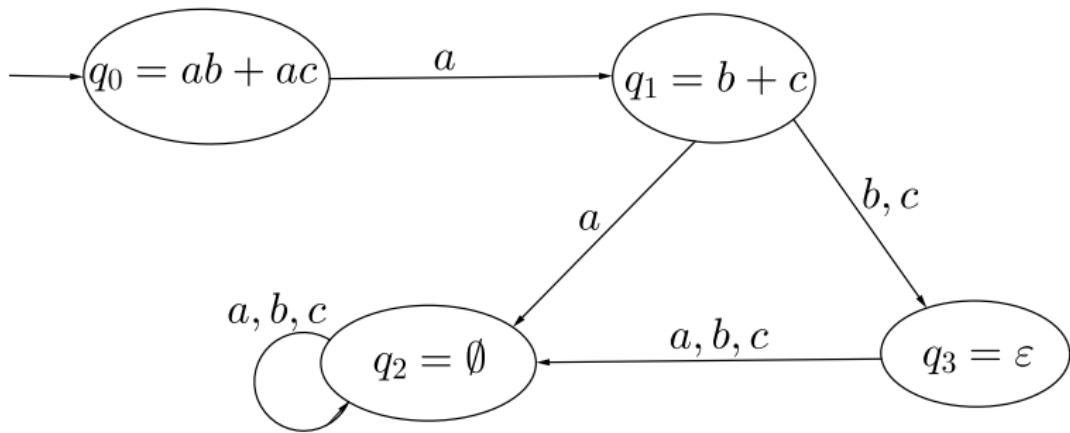
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

① $\partial_c q_1 = \partial_c(b + c) = (\emptyset + \varepsilon) \equiv \varepsilon = q_3$



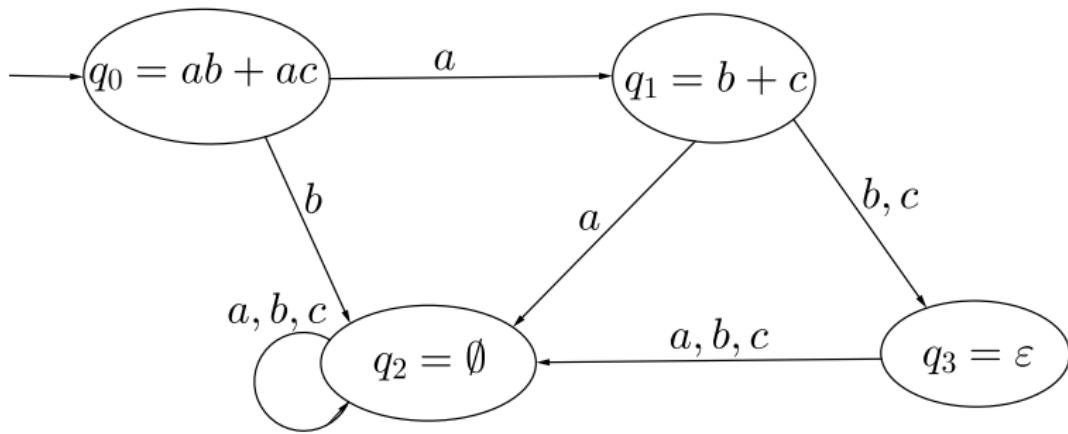
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

⑫ $\partial_b q_0 = \partial_b(a \cdot b + a \cdot c) = \emptyset = q_2$



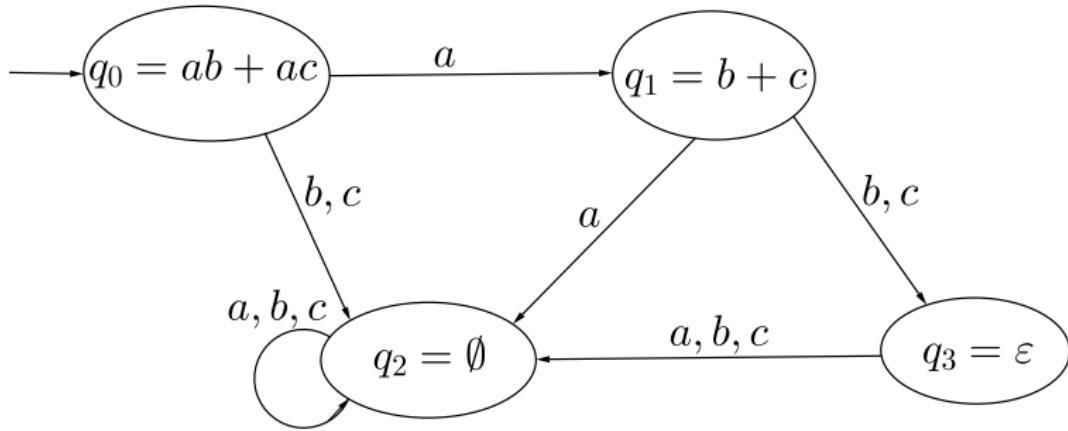
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

⑬ $\partial_c q_0 = \partial_c(a \cdot b + a \cdot c) = \emptyset = q_2$



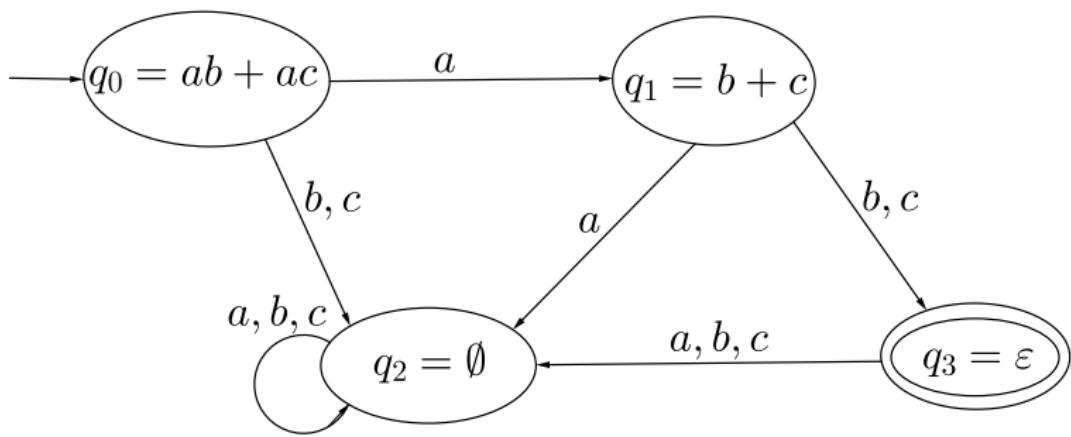
Algoritmo

Construcción de DFA

Ejemplo

Consideremos la expresión regular $r = a \cdot b + a \cdot c$ sobre el alfabeto $\{a, b, c\}$

Los estados finales son todos aquellos estados que sean anulables, en este caso únicamente es anulable $v(q_3) = \varepsilon$



Algoritmo

fun goto $q(c, (Q, \delta)) =$

let $q_c = \partial_c q$

in

if $\exists q' \in Q$ tal que $q' \equiv q_c$

then $(Q, \delta \cup \{(q, c) \rightarrow q'\})$

else

let $Q' = Q \cup \{q_c\}$

let $\delta' = \delta \cup \{(q, c) \rightarrow q_c\}$

in explora (Q', δ', q_c)

and explora $(Q, \delta, q) = \text{fold}(\text{goto } q)(Q, \delta) \Sigma$

fun mkDFA $r =$

let $q_0 = \partial_\varepsilon r$

let $(Q, \delta) = \text{explora}(\{q_0\}, \{\}, q_0)$

let $F = \{q | q \in Q \text{ y } v(q) = \varepsilon\}$

in $(Q, \Sigma, q_0, F, \delta)$