

Unit 4

Regular Expressions

Reading: Sipser, chapter 1

Overview

- Regular expressions
- Equivalence of RE and RL

Regular Expressions

- Regular languages (RL) are often described by means of algebraic expressions called *regular expressions (RE)*.
- In arithmetic we use the +, * operations to construct expressions: $(2+3)^*5$
- The *value* of the arithmetic expression is the number 25.
- The *value of a regular expression* is a *regular language*.

Regular operations

- In *RE* we use *regular operations* to construct expressions describing *regular languages*:

$$(0 + 1)^* \circ 0$$

where :

- $r+s$ means r OR s
- r^* means Kleene star of r
- $r \circ s$ (or rs) means concatenation of r and s

Formal definition

- A set of regular expressions over an alphabet Σ is defined inductively as follows:

Basis:

ε , \emptyset , and σ (for all $\sigma \in \Sigma$) are regular expressions.

Induction:

If r and s are RE then the following expressions are also RE:

- (r)
- $r+s$
- $r \circ s$
- r^*

Examples over $\Sigma=\{a,b\}$

ϵ , a, a+b, b*, (a+b)b, ab*, a*+b*

- To avoid using many parentheses, the operations have the following priority hierarchy:
 1. * - highest (do it first)
 2. \circ
 3. + - lowest (do it last)

- Example: $(b+(a \circ (b^*))) = b + ab^*$

(Notations: The symbol \circ can be dropped. Σ means $(\sigma_1+\sigma_2+\sigma_3\dots)$, r^+ means rr^*).

Regular expressions and regular languages

We associate each regular expression r with a regular language $L(r)$ as follows:

- $L(\emptyset) = \emptyset$,
- $L(\varepsilon) = \{\varepsilon\}$,
- $L(\sigma) = \{\sigma\}$ for each $\sigma \in \Sigma$,
- $L(r+s) = L(r) \cup L(s)$,
- $L(r \circ s) = L(r) \circ L(s)$,
- $L(r^*) = (L(r))^*$.

Examples over $\Sigma=\{0,1\}$

In Class: Describe each language as a regular expression.

$$L_1 = \{ w \mid w \text{ has a single } 1 \}$$

$$L_2 = \{ w \mid w \text{ has at least one } 1 \}$$

$$L_3 = \{ w \mid w \text{ contains the string } 110 \}$$

$$L_4 = \{ w \mid |w| \bmod 2 = 0 \}$$

$$L_5 = \{ w \mid w \text{ starts with } 1 \}$$

$$L_6 = \{ w \mid w \text{ ends with } 00 \}$$

Examples over $\Sigma=\{0,1\}$, cont'

$L_7 = \{ w \mid w \text{ starts with } 0 \text{ and ends with } 10 \}$

$L_8 = \{ w \mid w \text{ contains the string } 010 \text{ or the string } 101 \}$

$L_9 = \{ w \mid w \text{ starts and ends with the same letter} \}$

$L_{10} = \{0, 101\}$

$L_{11} = \{w \mid w \text{ does not contain } 11 \text{ as a substring}\}$

$L_{12} = \{w \mid \#1(w) \text{ is even}\}$

$L_{13} = \{w \mid w \text{ does not contain } 101\}$

Properties of regular expressions 1

Useful properties of regular expressions:

- $r+s=s+r$
- $r+\emptyset=\emptyset+r=r$
- $r+r=r$
- $r\emptyset=\emptyset r=\emptyset$
- $rr^*=r^+$
- $r\varepsilon=\varepsilon r=r$

Properties of regular expressions 2

- $r(s+t) = rs+rt$
- $r+(s+t) = (r+s)+t$
- $r(st) = (rs)t$
- $r^* = (r^*)^* = r^*r^* = r^* + r$
- $r^* + r^+ = r^*$

Equivalence of regular expressions

- To prove that two regular expressions r and s are equivalent we need to show that

$$L[r] \subseteq L[s] \text{ and } L[s] \subseteq L[r].$$

- To show that two regular expressions are **not** equivalent we have to find a word that belongs to one expression and does not belong to the other.

Example

Let $r = \varepsilon + (0+1)^*1$

Let $s = (0^*1)^*$

Are r and s equivalent?

Answer: Yes. We will prove

$$L[s] \subseteq L[r] \text{ and } L[r] \subseteq L[s]$$

$$L[r] \supseteq L[s]$$

$$r = \varepsilon + (0+1)^* 1$$

$$s = (0^* 1)^*$$

- Let $w \in L[s] = (0^* 1)^*$.
- $w = \varepsilon$ or $w = x_1 x_2 \dots x_n$, $n > 0$ such that $x_i \in L[0^* 1]$
- If $w = \varepsilon$ then $w \in L[r]$.
- If $w = x_1 x_2 \dots x_n$ then we can represent $w = w' 1 = x_1 x_2 \dots x_{n-1} 0^z 1$ with $z \geq 0$.

However, $w' = x_1 x_2 \dots x_{n-1} 0^z \in L[(0+1)^*]$, implying $w' 1 \in L[(0+1)^*] 1 \subseteq L[r]$.

$$L[r] \subseteq L[s]$$

$$r = \varepsilon + (0+1)^*1$$

$$s = (0^*1)^*$$

- Let $w \in L[r] = \varepsilon + (0+1)^*1$.
- If $w = \varepsilon$ then $w \in L[s]$ (by definition of $*$).
- If $w \neq \varepsilon$ then w can be represented as $w = w'1$ where $w' \in L[(0+1)^*]$. Assume that w' contains k instances of the letter 1. This means that w' can be written as $w' = x_11x_21\dots x_k1x_{k+1}$ where $x_i \in 0^*$

But then $w = w'1 =$

$$= (x_11)(x_21)\dots(x_{k+1}1) = (0^*1)(0^*1)\dots(0^*1)$$

So $w \in L[(0^*1)^*]$.

Another example

Are r and s equivalent?

$$r = (0+1)^*1+0^*$$

$$s = (1+0)(0^*1)^*$$

Answer: No.

- Consider the word $w = \varepsilon$.
- $w \in L[r] = (0+1)^*1+0^*$, because $w \in 0^*$.
- But $w \notin L[s] = (1+0)(0^*1)^*$, as all words in $L[s]$ have at least one letter.

Equivalence of RE with FA

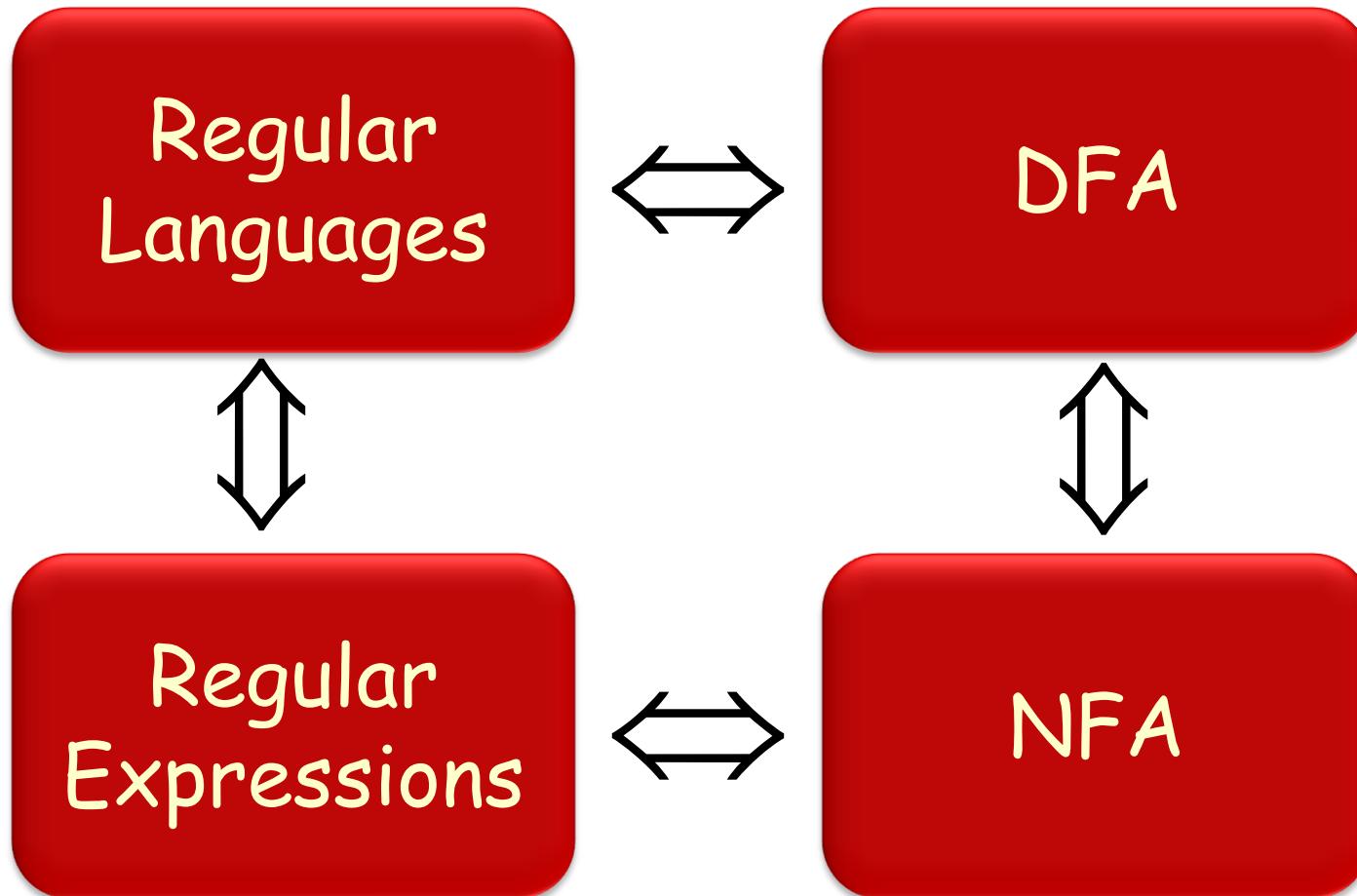
- Regular expressions and finite automata are equivalent in terms of the languages they describe.

Theorem:

A language is regular iff some regular expression describes it.

This theorem has two directions. We prove each direction separately.

Equivalences so far...



Converting RE into FA

- If a language is described by a regular expression, then it is regular.

Proof idea:

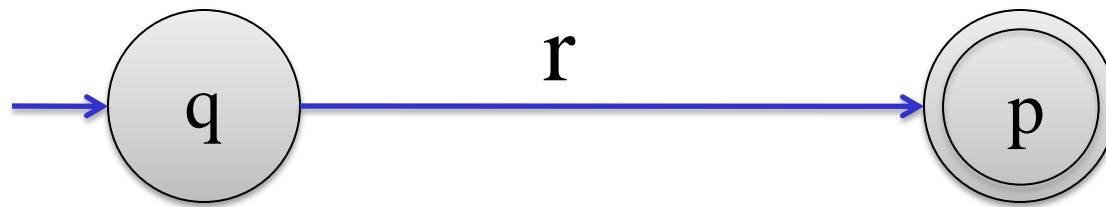
Build a NFA by transforming some regular expression r into a non-deterministic finite automaton N that accepts the language $L(r)$.

Converting RE into RL



RE to FA Algorithm

- Given r we start the algorithm with N having a start state, a single accepting state and an edge labeled r :

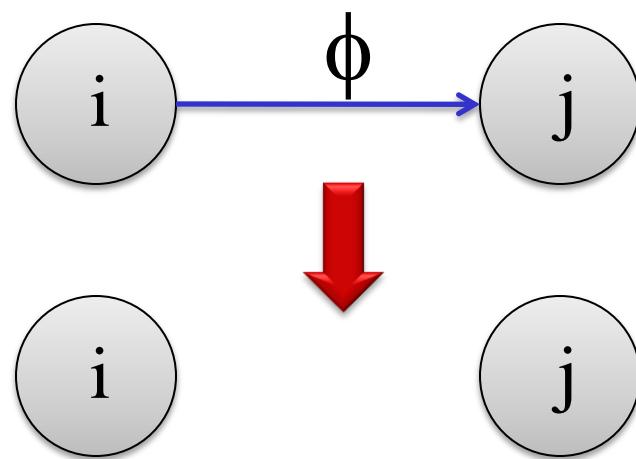


Note: We assume at a moment that transactions can be labeled with RE, not just letters.

RE to FA Algorithm (cont.)

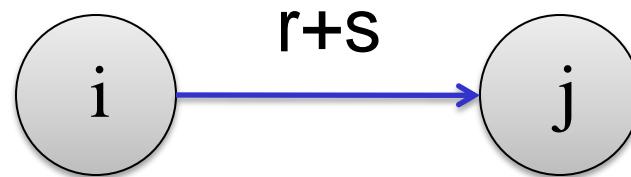
Now transform this machine into an NFA N by applying the following rules until all the edges are labeled with either a letter σ from Σ or ε :

1. If an edge is labeled \emptyset , then delete this edge.

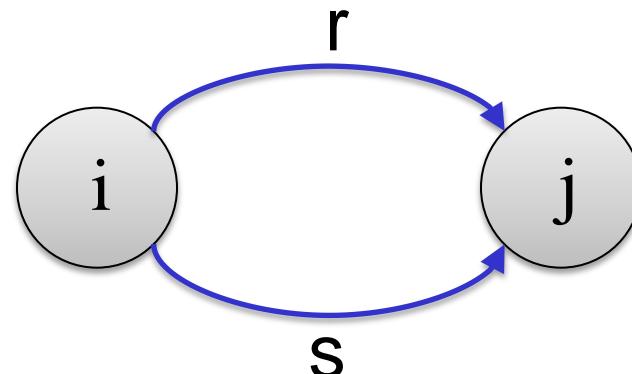


RE to FA Algorithm (cont.)

2. Transform any diagram of the type

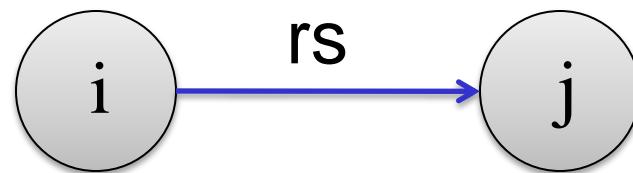


into the diagram

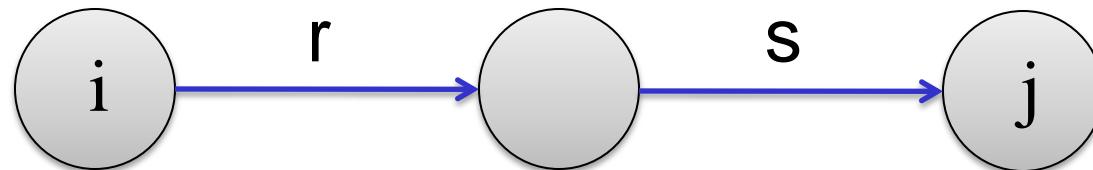


RE to FA Algorithm (cont.)

3. Transform any diagram of the type

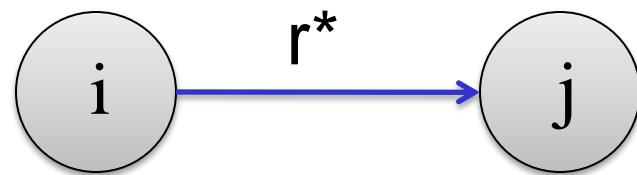


into the diagram

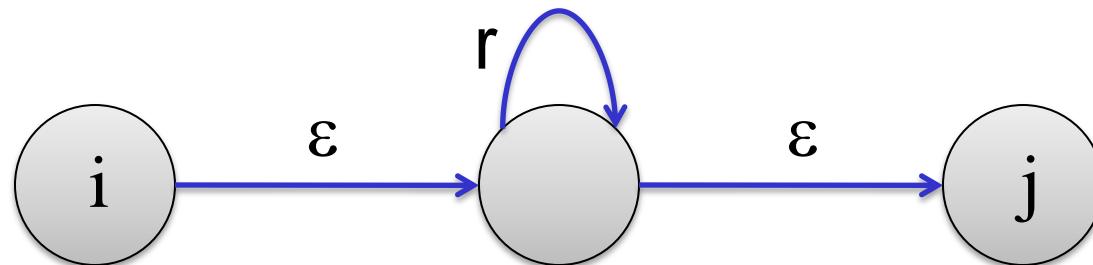


RE to FA Algorithm (cont.)

4. Transform any diagram of the type



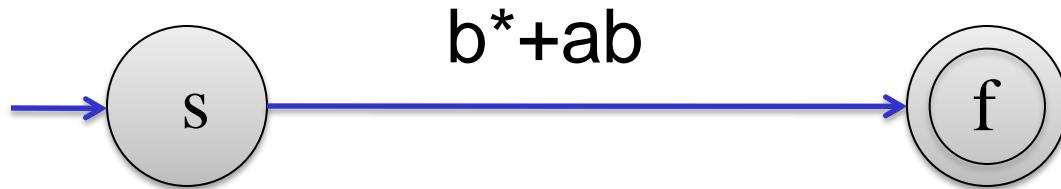
into the diagram



Example

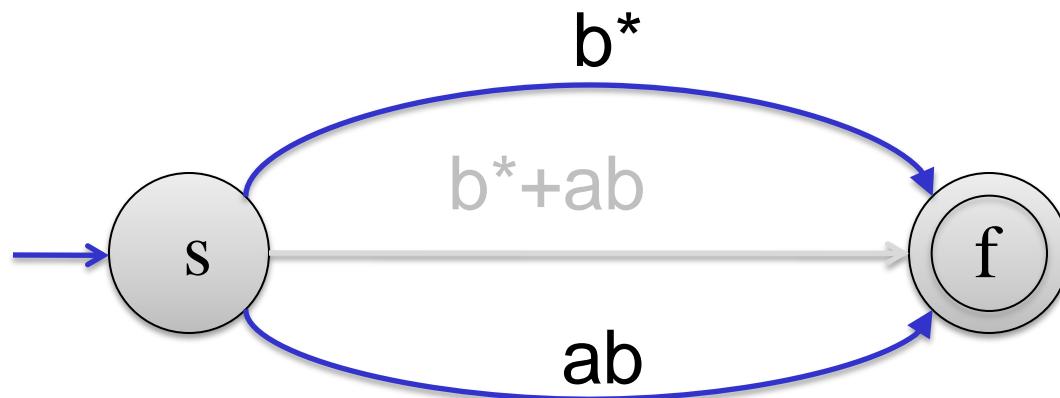
Construct an NFA for the regular expression
 $b^* + ab.$

Solution: We start by drawing the diagram:



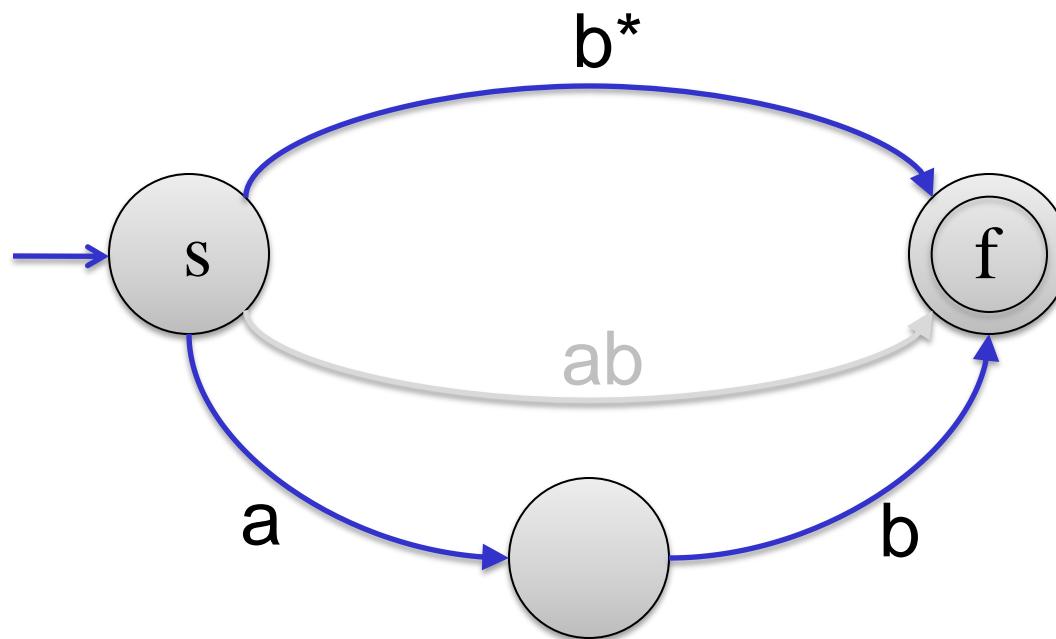
Example (cont.)

Next we apply rule 2 for b^*+ab :



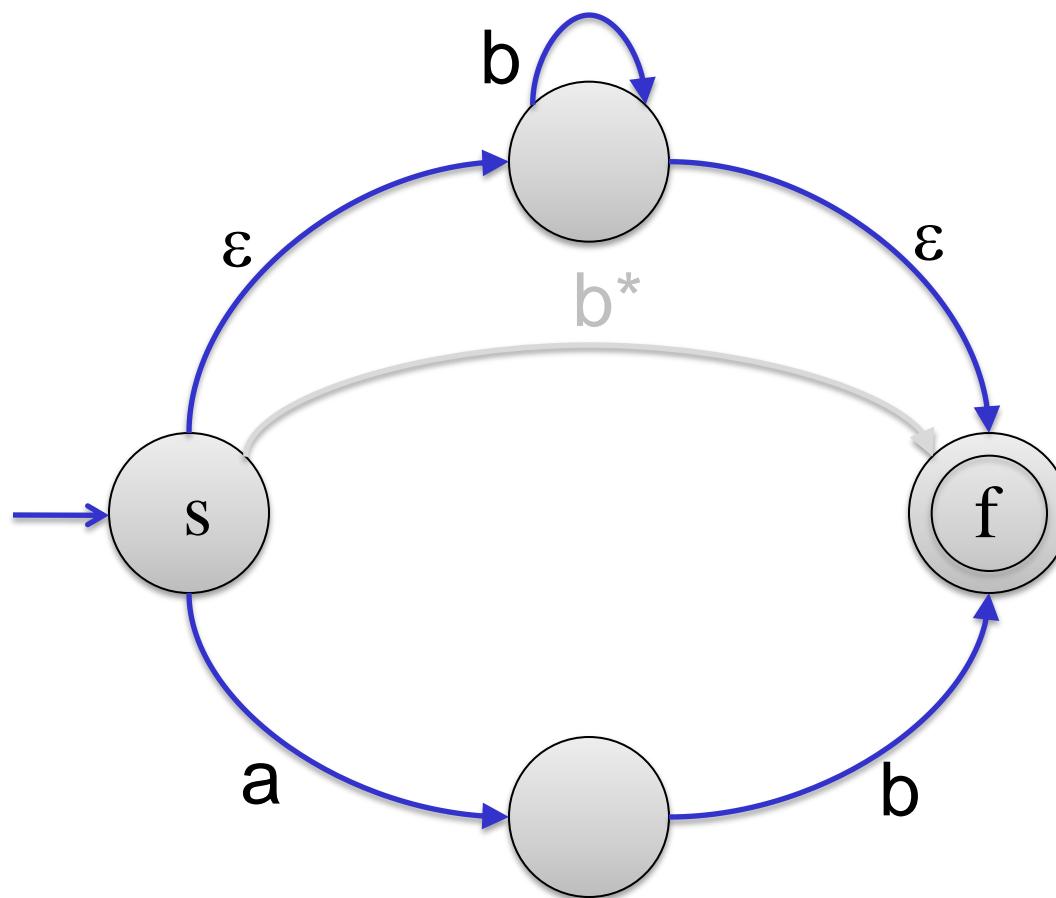
Example (cont.)

Next we apply rule 3 for ab:

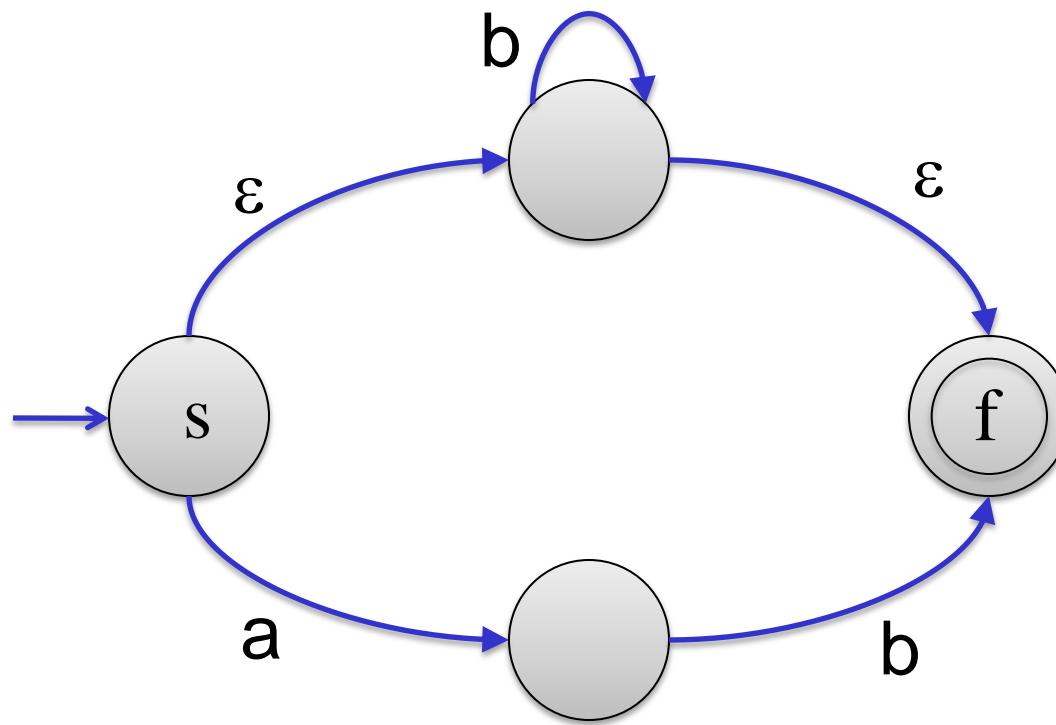


Example (cont.)

Next we apply rule 4 for b^* :



The final NFA



Example 2: Draw an NFA for $(ab+a)^*$

Example 3: Draw an NFA for $\Sigma^*101\Sigma^*$

Converting FA into RE

- If a language is regular, then it can be described by a regular expression.

Proof idea:

Transform some DFA N into a regular expression r that s.t. $L(r)=L(N)$.

Converting RL into RE



Converting FA into RE

- The algorithm will perform a sequence of transformations that contract the DFA into new machines with edges labeled with regular expressions
- It stops when the machine has:
 1. two states: Start, Finish
 2. one edge with a regular expression on it.

Generalized NFA

- Before we start we first convert the DFA into a *Generalized NFA* (GNFA):
 - GNFA might have RE as labels.
 - GNFA has a single accept state.
 - The start state has arrows going out but none coming in.
 - The accept state has arrows coming in but none going out.
 - Except for the start and accept state, one arrow goes from every state to every other state (including itself).

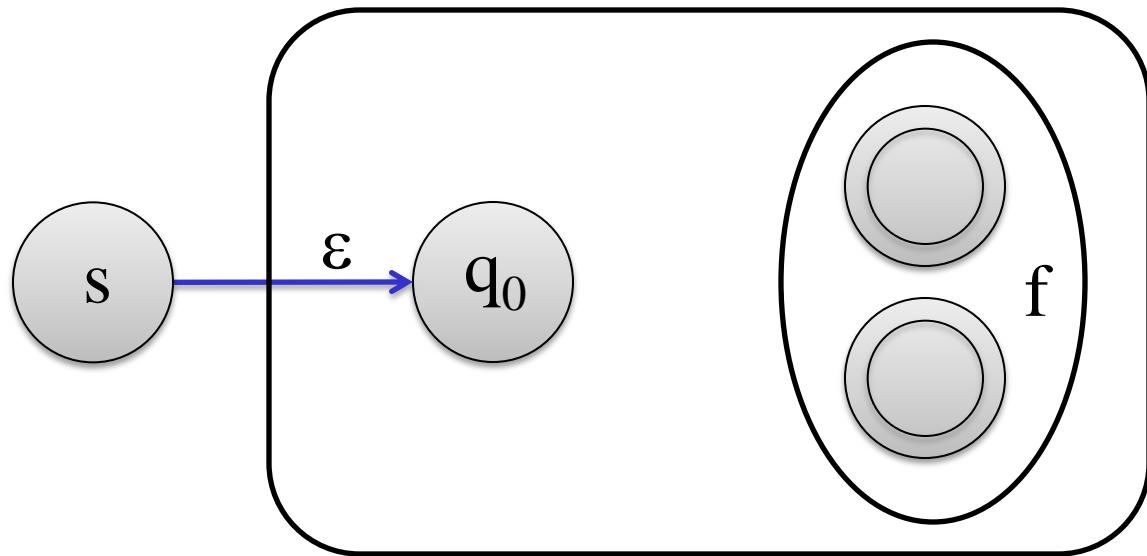


Converting DFA into GNFA

The input is a DFA or an NFA $N=(Q, \Sigma, \delta, q_0, F)$.

Perform the following steps:

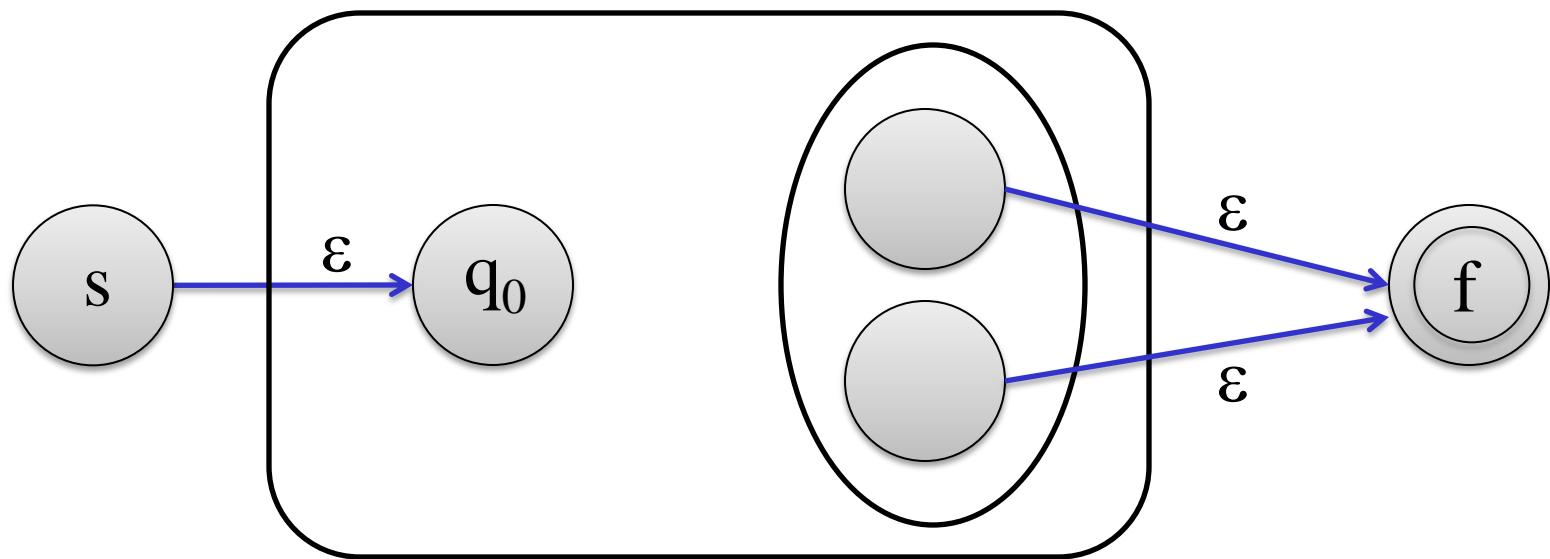
1. Create a new start state s and draw a new edge labeled with ε from s to the q_0 . $\delta(s, \varepsilon) = q_0$.



Converting DFA into GNFA (cont.)

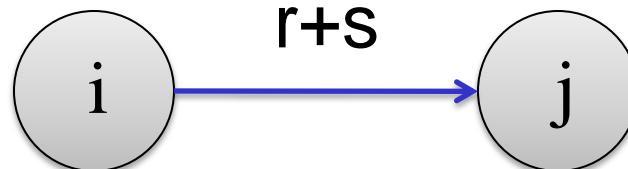
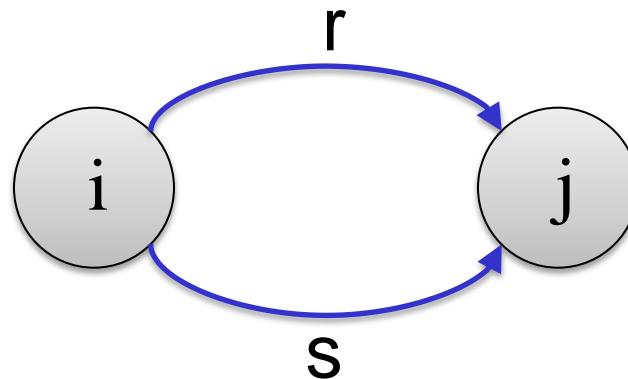
2. Create a new **single** accepting state f and draw new edges labeled ε from all the original accepting states to f .

Formally: For each $q \in F$, $\delta(q, \varepsilon) = f$.

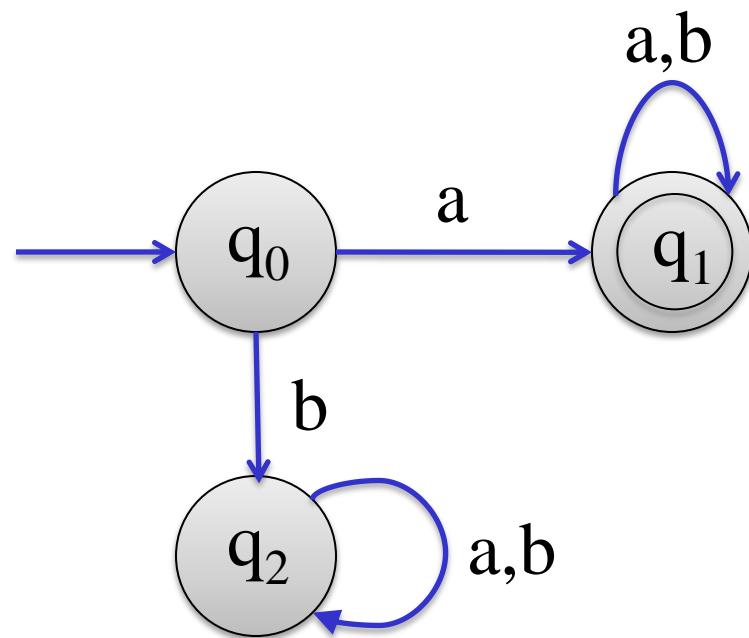


Converting DFA into GNFA (cont.)

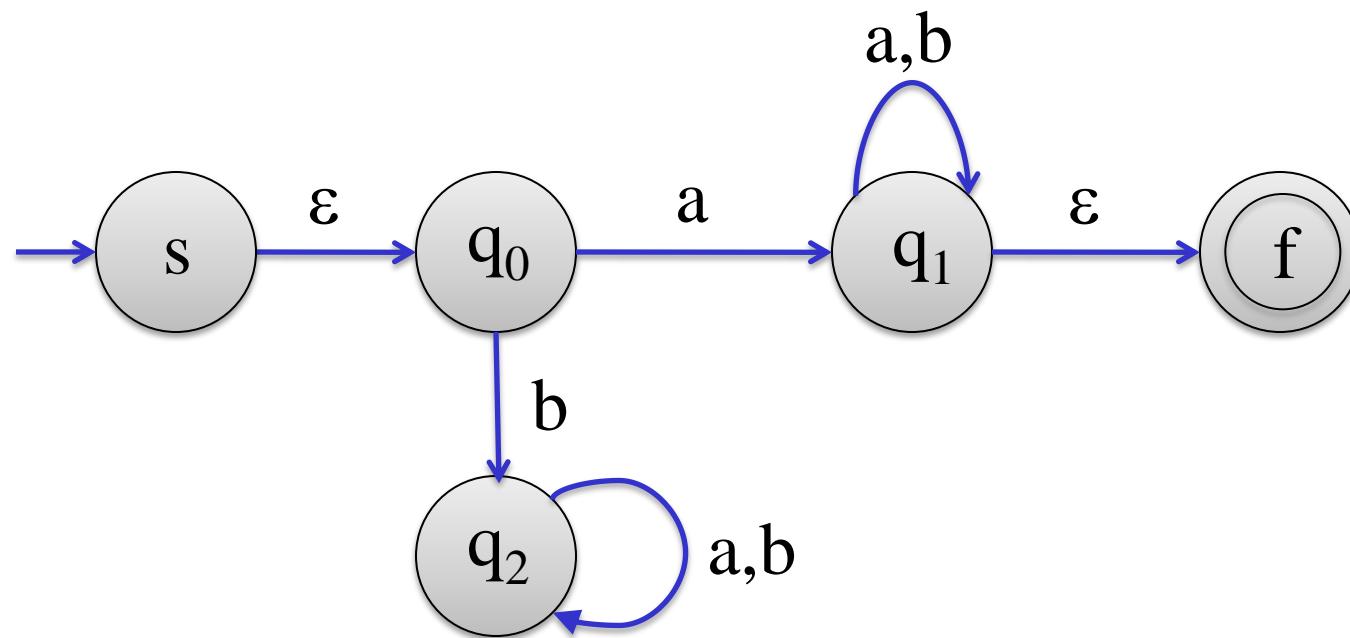
3. For each pair of states i and j that have more than one edge between them (in the same direction), replace all the edges by a *single* edge labeled with the RE formed by the sum of the labels of these edges.
4. If there is no edge $\langle i, j \rangle$ then $\text{label}(i, j) = \emptyset$.



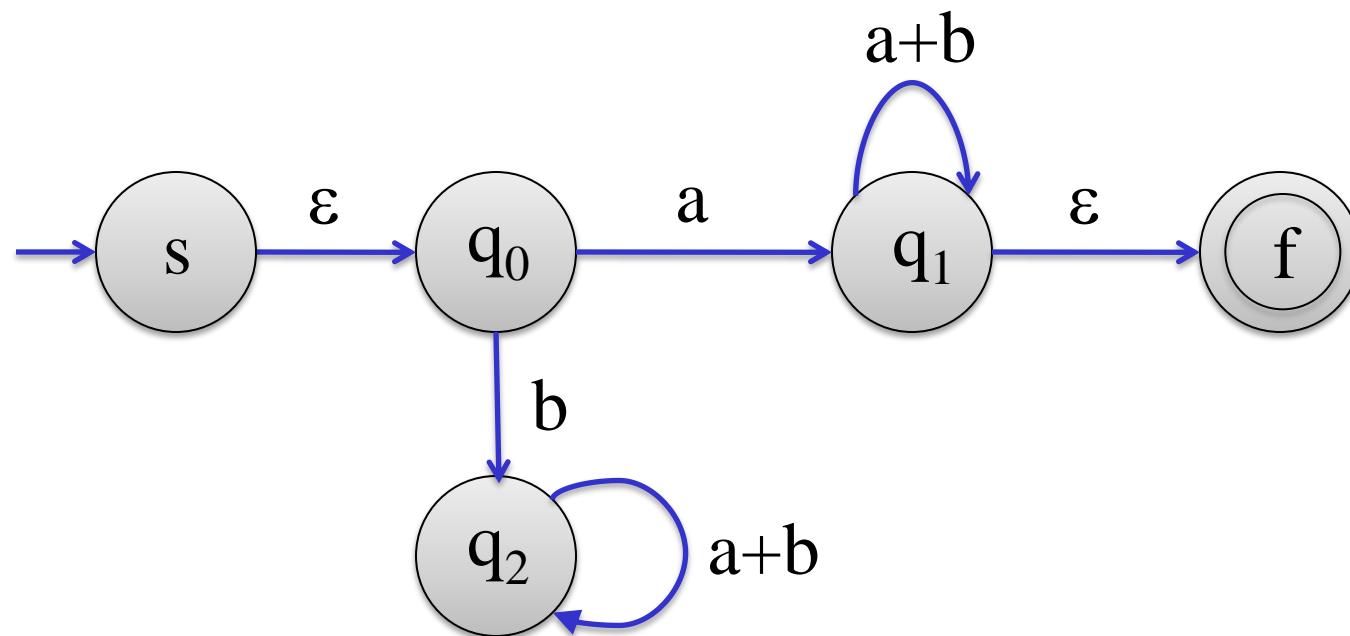
Example: DFA to GNFA



Example: DFA to GNFA (cont.)



Example: DFA to GNFA (cont.)



Converting GNFA into RE

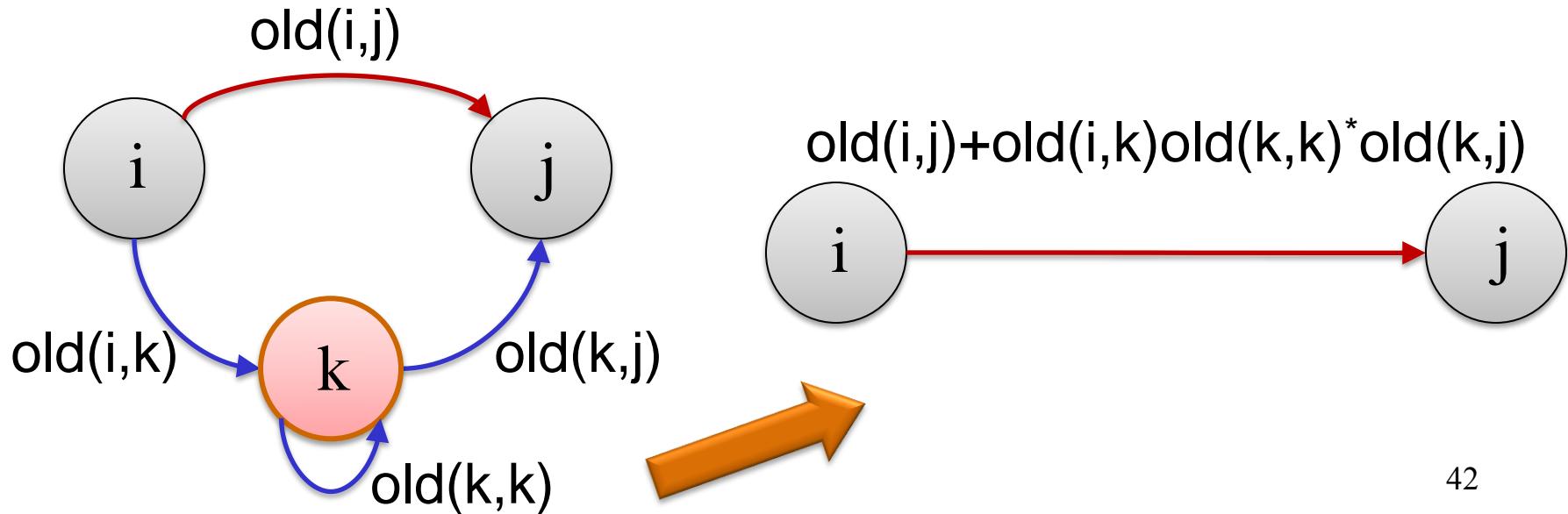
- Let $\text{old}(i,j)$ denote the label on edge $\langle i,j \rangle$ of the current GNFA.
- Construct a sequence of new machines by *eliminating one state at a time* until the only two states remaining are s and f .
- The state elimination order is arbitrary.
- When a state is eliminated, a new (equivalent) machine is constructed.
- The label on $\langle s,t \rangle$ in the final machine is the required *RE*.

Converting GNFA into RE (cont.)

Eliminating state k

- For each pair of states (i,j) where $i,j \neq k$, the label of (i,j) will be updated as follows:

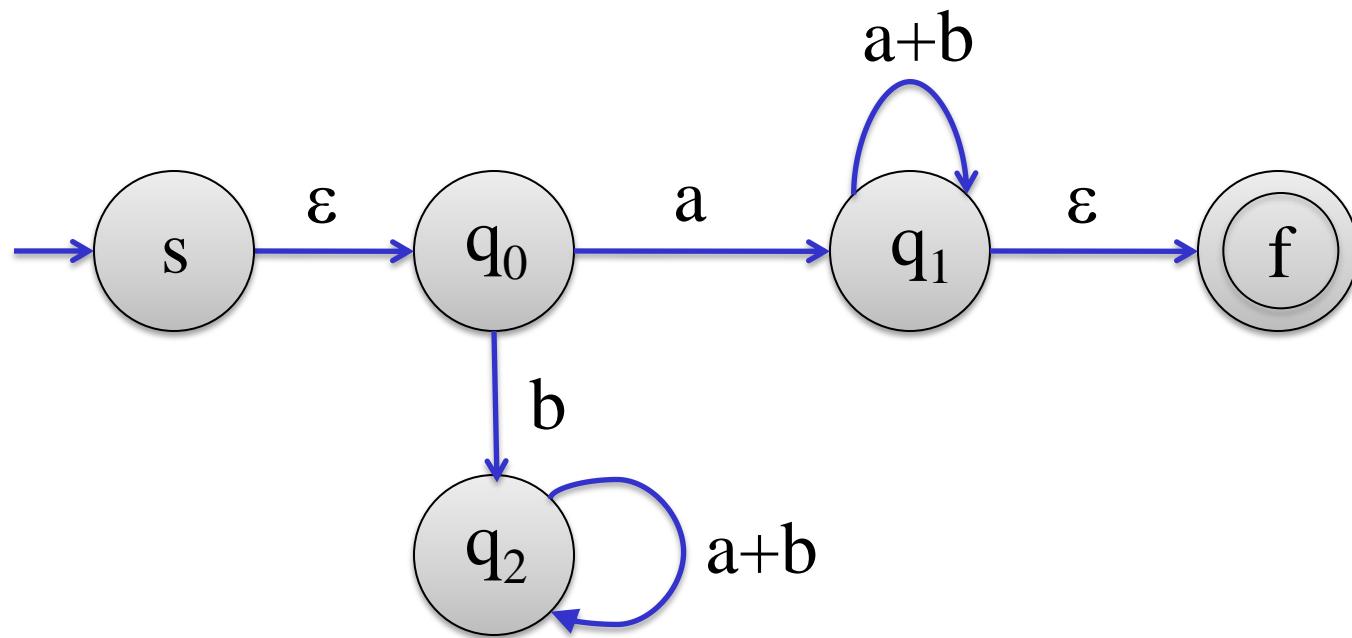
$$\text{new}(i,j) = \text{old}(i,j) + \text{old}(i,k)\text{old}(k,k)^*\text{old}(k,j)$$



Converting GNFA into RE (cont.)

- The **states** of the new machine are those of the current machine with state k eliminated.
- The **edges** of the new machine are those edges (i,j) for which $\text{new}(i,j)$ has been calculated.
- The algorithm terminates when s and f are the only two remaining states. The regular expression $\text{new}(s,f)$ represents the language of the original automaton.

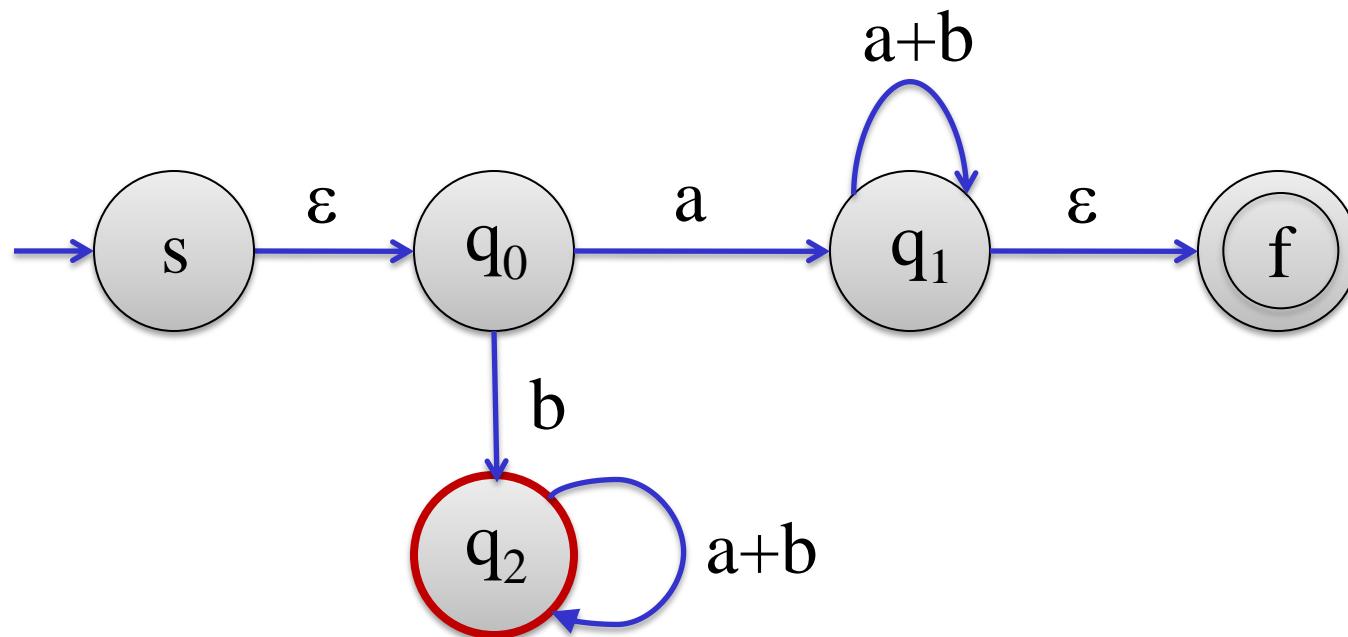
Example: GNFA to RE



Example: GNFA to RE

Eliminate state q_2

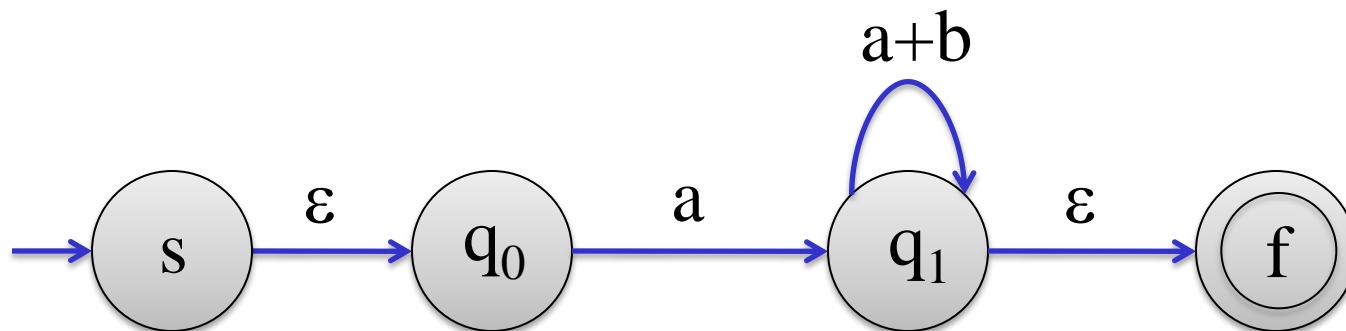
- No paths pass through q_2 . There are no states that connect through 2. So no need to change anything after deletion of state 2.



Example: GNFA to RE

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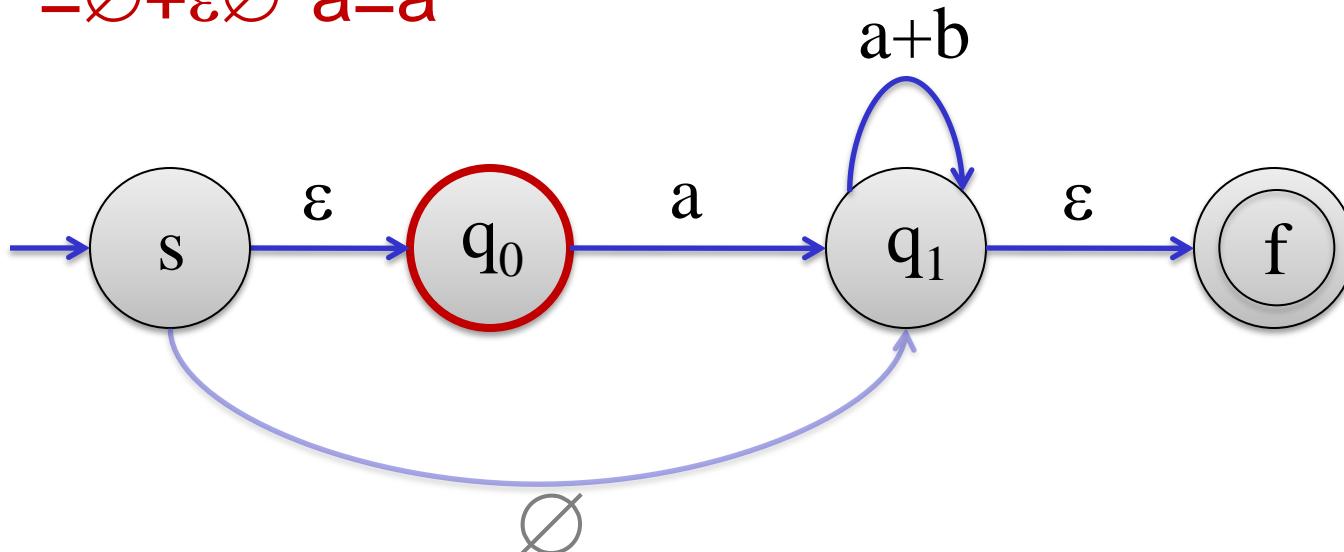


Example: GNFA to RE

Eliminate state q_0

- The only path through it is $s \rightarrow q_1$.
- We add an edge that is labeled by regular expression associated with the deleted edges:

$$\begin{aligned}\text{new}(s, q_1) &= \text{old}(s, q_1) + \text{old}(s, q_0)\text{old}(q_0, q_0)^*\text{old}(q_0, q_1) = \\ &= \emptyset + \varepsilon \emptyset^* a = a\end{aligned}$$

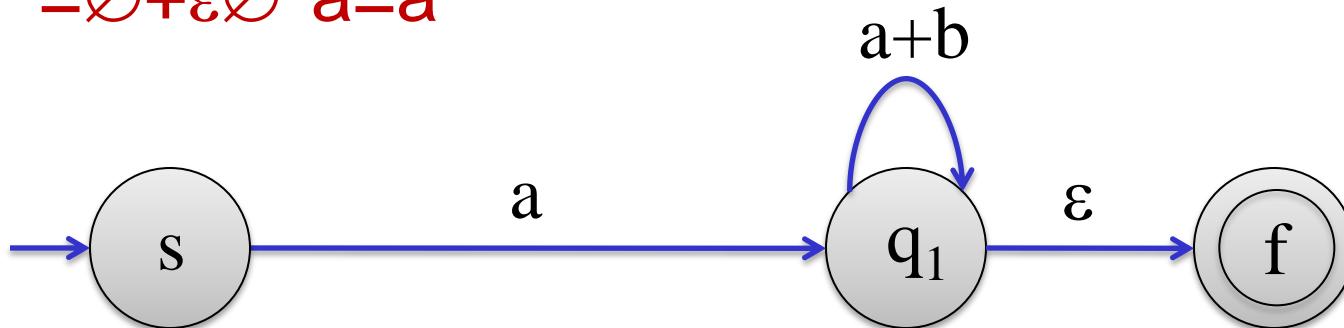


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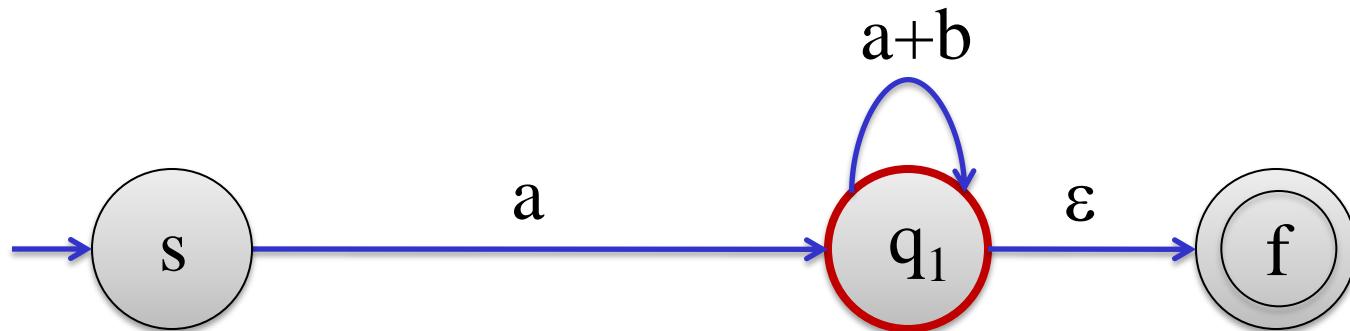


Example: GNFA to RE

Eliminate state q_1

- The only path through it is $s \rightarrow f$.

$$\begin{aligned} \text{new}(s,f) &= \text{old}(s,f) + \text{old}(s,q_1)\text{old}(q_1,q_1)^*\text{old}(q_1,f) = \\ &= \emptyset + a(a+b)^*\varepsilon = a(a+b)^* \end{aligned}$$

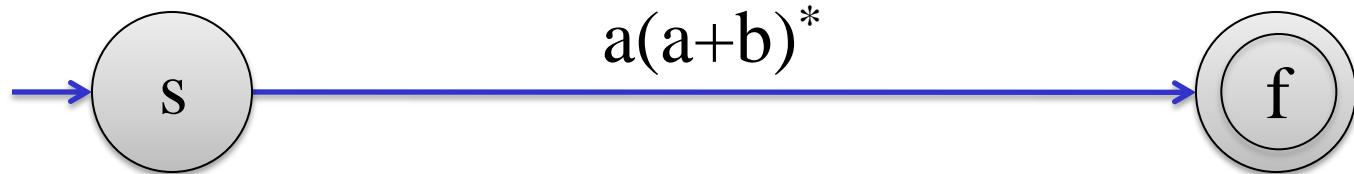


Example: GNFA to RE

Eliminate state q_1

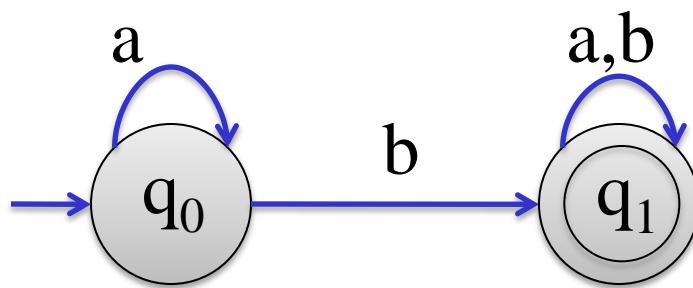
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Example II

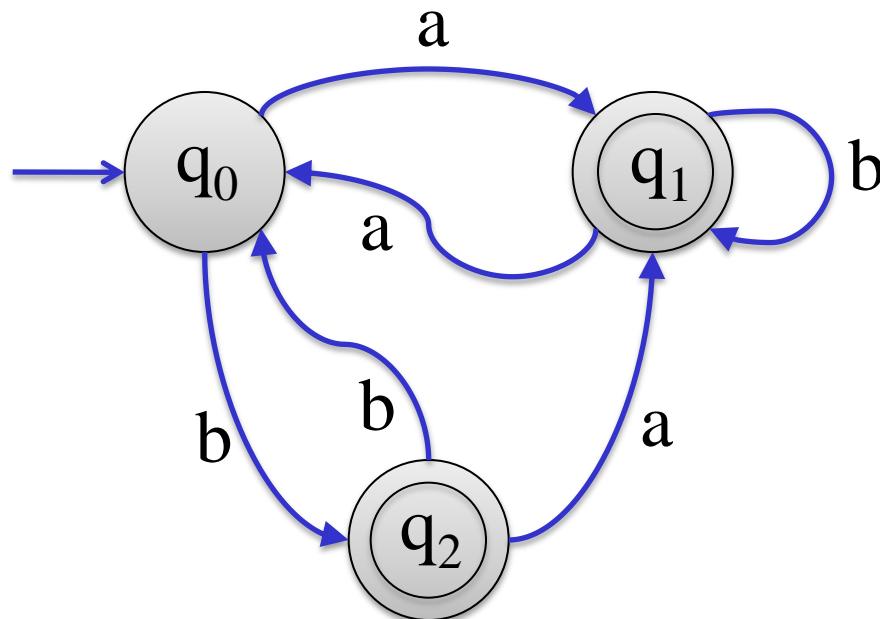
What is the regular expression of $L(A)$?



Solution: In class

Example III

What is the regular expression of $L(A)$?



Solution: In class