Escuela de Ingeniería Informática



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Activity 1. [Time measurements for sorting algorithms]

1. Task 1 - Insertion.

Results

N	sorted(10 ⁻⁷ s)	inverse(10 ⁻³ s)	random(10 ⁻³ s)
10,000	58	60	46
20,000	64	217	136
40,000	126	295	150
80,000	210	1,215	572
160,000	435	4,876	2,349
320,000	848	19,041	9,130
640,000	1,669	76,347	36,442
1,280,000	3,854	309,526	158,185
2,560,000	9,808	1,243,104	620,628
5,120,000	19,403	4,946,423	2,479,540
10,240,000		19,802,890	9,900,099
Complexity	0(n)	$O(n^2)$	$O(n^2)$

Conclusion

sorted(t)	inverse(t)	random(t)
$t_1 = 210, n_1 = 80,000 \text{ and } n_2 = 320,000$ $t_3 = 126 n_3 = 40,000 \text{ and } n_4 = 640,000$	$t_1 = 76,347, n_1 = 160,000 \ and \ n_2 = 640,000$ $t_3 = 19,041 \ n_3 = 320,000 \ and \ n_4 = 1,280,000$	$t_1 = 2,349, n_1 = 160,000 \ and \ n_2 = 1,280,000$ $t_3 = 9,130, n_3 = 320,000 \ and \ n_4 = 640,000$
	• $t_2 = \left(\frac{640,000}{160,000}\right)^2 \cdot 76,347 = 1,221,552$	• $t_2 = \left(\frac{1,280,000}{160,000}\right)^2 \cdot 2,349 = 150,336$
• $t_4 = \frac{640,000}{40,000} \cdot 126 = 2,016$	• $t_4 = \left(\frac{1,280,000}{320,000}\right)^2 \cdot 19,041 = 304,656$	• $t_4 = \left(\frac{640,000}{320,000}\right)^2 \cdot 9,130 = 36,520$

Values	sorted(t)	inverse(t)	random(t)
Actual₁	848	1,243,104	158,185
Expected ₁	840	1,221,552	150,336
Actual ₂	1,669	309,526	36,442
Expected ₂	2,016	304,656	36,520

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2. Task 2 - Selection.

• Results

N	sorted(10 ⁻³ s)	inverse(10 ⁻³ s)	random(10 ⁻³ s)
10,000	191	176	205
20,000	527	484	518
40,000	1,897	2,127	2,141
80,000	7,558	8,557	8,466
160,000	29,829	34,115	33,340
320,000	119,689	138,498	135,676
640,000	473,457	553,992	542,704
1,280,000	1,870,023	2,671,256	2,190,816
2,560,000	7,702,096	8,863,872	8,673,784
5,120,000	30,634,978	35,467,670	34,729,056
10,240,000	122,561,536	141,812,972	139,000,344
Complexity	$O(n^2)$	$O(n^2)$	$O(n^2)$

Conclusion

sorted(t)	inverse(t)	random(t)
$t_1 = 1,870,023, n_1 = 1,280,000 $ and $n_2 = 10,240,000$	$t_1 = 34,115, n_1 = 160,000 \text{ and } n_2 = 640,000$	$t_1 = 8,466, n_1 = 80,000 \text{ and } n_2 = 320,000$
$t_3 = 29,829, n_3 = 160,000 \ and \ n_4 = 640,000$	$t_3 = 138,498 \ n_3 = 320,000 \ and \ n_4 = 1,280,000$	$t_1 = 33,340, n_1 = 160,000 \ and \ n_2 = 1,280,000$
• $t_2 = \left(\frac{10,240,000}{1,280,000}\right)^2 \cdot 1,870,023 = 119,681,472$	• $t_2 = \left(\frac{640,000}{160,000}\right)^2 \cdot 34,115 = 545,840$	• $t_2 = \left(\frac{320,000}{80,000}\right)^2 \cdot 8,466 = 135,456$
• $t_4 = \left(\frac{640,000}{160,000}\right)^2 \cdot 29,829 = 477,264$	• $t_4 = \left(\frac{1,280,000}{320,000}\right)^2 \cdot 138,498 = 2,215,968$	• $t_4 = \left(\frac{1,280,000}{160,000}\right)^2 \cdot 33,340 = 2,133,760$

Values	sorted(t)	inverse(t)	random(t)
Actual ₁	122,561,536	553,992	135,676
Expected ₁	119,681,472	545,840	135,456
Actual ₂	273,457	2,671,256	2,190,816
Expected ₂	477,264	2,215,968	2,133,760

3. Task 3 - Bubble.

• Results

N	sorted(10 ⁻³ s)	inverse(10 ⁻³ s)	random(10 ⁻³ s)
10,000	123	152	186
20,000	460	530	788
40,000	2,993	1,984	4,960

N	sorted(10 ⁻³ s)	inverse(10 ⁻³ s)	random(10 ⁻³ s)
80,000	12,057	7,998	19,820
160,000	48,404	32,048	79,341
320,000	194,581	128,929	317,974
640,000	778,324	515,716	1,271,896
1,280,000	3,113,296	2,062,864	5,034,567
2,560,000	12,453,184	8,251,456	20,350,336
5,120,000	49,812,736	33,005,824	81,401,344
10,240,000	199,250,944	132,123,224	325,605,376
Complexity	$O(n^2)$	$O(n^2)$	$O(n^2)$

Conclusion

sorted(t)	inverse(t)	random(t)
$t_1 = 3,113,296, n_1 = 1,280,000 \ and \ n_2 = 10,240,000$	$t_1 = 2,062,864, n_1 = 1,280,000 \ and \ n_2 = 10,240,000$	$t_1 = 79,341, n_1 = 160,000 \ and \ n_2 = 640,000$
$t_3 = 48,404, n_3 = 160,000 \ and \ n_4 = 640,000$	$t_3 = 32,048, n_3 = 160,000 \text{ and } n_4 = 640,000$	$t_3 = 317,974 n_3 = 320,000 and n_4 = 1,280,000$
• $t_2 = \left(\frac{10,240,000}{1,280,000}\right)^2 \cdot 3,113,296 = 199,250,944$	• $t_2 = \left(\frac{10,240,000}{1,280,000}\right)^2 \cdot 2,062,864 = 132,023,296$	• $t_2 = \left(\frac{640,000}{160,000}\right)^2 \cdot 79,341 = 1,269,456$
• $t_4 = \left(\frac{640,000}{160,000}\right)^2 \cdot 48,404 = 774,464$	• $t_4 = \left(\frac{640,000}{160,000}\right)^2 \cdot 32,048 = 512,768$	• $t_4 = \left(\frac{1,280,000}{320,000}\right)^2 \cdot 317,974 = 5,087,584$

Values	sorted(t)	inverse(t)	random(t)
Actual₁	199,250,944	132,123,224	1,271,896
Expected ₁	199,250,944	132,023,296	1,269,456
Actual ₂	778,324	515,716	5,034,567
Expected ₂	774,464	512,768	5,087,584

4. Task 4 – Quicksort with the central element as the pivot.

• Results

N	sorted(10 ⁻⁶ s)	inverse(10 ⁻⁶ s)	random(10 ⁻⁶ s)
10,000	96	95	588
20,000	139	157	1,251
40,000	382	340	2,605
80,000	740	712	5,349
160,000	1,567	1,476	11,148
320,000	2,788	2,867	23,272
640,000	5,694	5,908	49,703
1,280,000	11,970	12,209	99,942

N	sorted(10 ⁻⁶ s)	inverse(10 ⁻⁶ s)	random(10 ⁻⁶ s)
2,560,000	24,872	25,353	204,067
5,120,000	54,417	58,101	469,408
10,240,000	118,601	131,890	1,068,294
Complexity	$O(n \log{(n)})$	$O(n \log{(n)})$	$O(n \log{(n)})$

Conclusion

sorted(t)	inverse(t)	random(t)
$t_1 = 34,115, n_1 = 160,000 \ and \ n_2 = 640,000$	$t_1 = 1,870,023, n_1 = 1,280,000 \ and \ n_2 = 10,240,000$	$t_1 = 1,870,023, n_1 = 1,280,000 \ and \ n_2 = 10,240,000$
$t_3 = 138,498 n_3 = 320,000 $ and $n_4 = 1,280,000$	$t_3 = 29,829, n_3 = 160,000 $ and $n_4 = 640,000$	$t_3 = 29,829, n_3 = 160,000 $ and $n_4 = 640,000$
• $t_2 = \frac{\log k + \log n_2}{\log n_1} \cdot t_1 = 5,708$	• $t_2 = \frac{\log k + \log n_2}{\log n_1} \cdot t_1 = 115,888$	• $t_2 = \frac{\log k + \log n_2}{\log n_1} \cdot t_1 = 1,023,061$
• $t_4 = \frac{\log k + \log n_2}{\log n_1} \cdot t_1 = 11,970$	• $t_4 = \frac{\log k + \log n_2}{\log n_1} \cdot t_1 = 6,086$	• $t_4 = \frac{\log k + \log n_2}{\log n_1} \cdot t_1 = 51,634$

Values	sorted(t)	inverse(t)	random(t)
Actual ₁	5,694	131,890	1,068,294
Expected ₁	5,708	115,888	1,023,061
Actual ₂	11,970	5,908	49,703
Expected ₂	12,663	6,086	51,634

Activity 2. [QuicksortFateful]

In this case it has been chosen the left-most element of the vector — or a partition of it — as a pivot. This will cause the algorithm to slice the array into unbalanced partitions. Notice that this scenario is the worst-case for this kind of algorithm. Imagine we have a set of elements whose size is n, as you may now, the recursive calls of the quicksort algorithm have the following scheme:

- quickSort(left, i-1);
- quickSort(i+1, right);

In the best situation both calls would have the same size, this way, we are dividing by 2 the size with each iteration. But that's the best case, now, let's suppose the worst scenario, the one at which with every and each iteration, the pivot is the smallest element; for example, in the case of choosing as a pivot the left-most element and having the vector already sorted. This way, the size of the problem won't be chopped into 2 halves, what will happen

is that one partition will contain no elements, and the other one with n-1 elements. During its second iteration, it will happen the same thing as before, the next item, the "new" pivot is already the smallest element; this will cause the algorithm to split the array into two portions, one containing 0 elements, but the other n-1 elements (n-2 if we compare it to the starting size). This will occur with every iteration performing this process as many times as needed until the size is 2, when we will divide the array into two pieces one with no elements and the last with 1. This thing will also happen when using the left value as a pivot.