

RE Lab 0x03 - Static Analysis

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based on previous work by R. Caragea and C. Rusu

Setup

Get your lab files from moodle: [03-lab-files.zip](#). The password for the zip is `infected` .

As you will recall from the previous session, we learned what the process of going from code to assembly looks like. In this lab we will learn how to go the other way around: from the compiled binary file, back to source code. We'll take a look at the structure of ELF binary files and the tools we can use in order to analyse them.

You can use both Windows and Linux interchangeably for this lab, but Linux will probably be more convenient.

Tools

You won't need all of the tools that I'll mention here, but it's nice to have options right?

Super basic, cli tools

- readelf
- objdump (you have already played with this one in the previous labs)

You can always learn more about them by reading the corresponding man pages.

Advanced

- [IDA \(Pro?\)](#) ^[1]
 - x86/x86_64 only / cloud decompilation only (free)
 - full RE framework with tons of features (very expensive)
- [Binary Ninja](#) ^[2]
 - modern and really promising project (75\$ for students, among others)
- [Ghidra](#) ^[3]
 - NSA-developed, FOSS IDA Pro competitor with huge community support and the ugliest UI on the planet (free)
- [Cutter/Rizin](#) ^[4]
 - minimalist RE tool, built on top of Rizin, a very powerful CLI RE suite of tools. (free)

ELF structure

As described in the course, Linux executables follow the ELF structure. If we want to reverse engineer Linux executables, we must learn some basics about this structure.

Firstly, ELF files can have dual use: as an object file by compilers or as runnable files by the OS interpreter. This is the reason why ELF programs have both segments, also known as program headers and sections. Segments

To investigate program headers on Linux we can use the `readelf` utility in the `binutils` package. Most of the time, the target binary is compiled by gcc or clang. In these cases, the structure (internal organization) of output binaries is basically the same. Let us now look at the structure of a classic “Hello, World” program.

If you are interested in learning more about the ELF file format, or file formats in general, you can check out [corkami](#) and [\[5\]](#) [\[6\]](#) [\[7\]](#).

Program Headers

```
$ readelf --program-headers hello-world
```

Elf file type is DYN (Shared object file)

Entry point 0x580

There are 9 program headers, starting at offset 64

Program Headers:

Type	Offset FileSiz	VirtAddr MemSiz	PhysAddr Flags	Align
PHDR	0x0000000000000040 0x00000000000001f8	0x0000000000000040 0x00000000000001f8	0x0000000000000040 R E	0x8
INTERP	0x0000000000000238 0x000000000000001c	0x0000000000000238 0x000000000000001c	0x0000000000000238 R	0x1
[Requesting program interpreter: /lib64/ld-linux-x86-64.so.2]				
LOAD	0x0000000000000000 0x000000000000089c	0x0000000000000000 0x000000000000089c	0x0000000000000000 R E	0x200000
LOAD	0x0000000000000dd8 0x0000000000000258	0x0000000000200dd8 0x0000000000000260	0x00000000000200dd8 RW	0x200000
DYNAMIC	0x0000000000000df0 0x00000000000001e0	0x0000000000200df0 0x00000000000001e0	0x00000000000200df0 RW	0x8
NOTE	0x0000000000000254 0x0000000000000044	0x0000000000000254 0x0000000000000044	0x0000000000000254 R	0x4
GNU_EH_FRAME	0x0000000000000754 0x000000000000003c	0x0000000000000754 0x000000000000003c	0x0000000000000754 R	0x4
GNU_STACK	0x0000000000000000 0x0000000000000000	0x0000000000000000 0x0000000000000000	0x0000000000000000 RW	0x10
GNU_RELRO	0x0000000000000dd8 0x0000000000000228	0x0000000000200dd8 0x0000000000000228	0x00000000000200dd8 R	0x1

Section to Segment mapping:

Segment Sections...

```

00
01      .interp
02      .interp .note.ABI-tag .note.gnu.build-id .gnu.hash .dynsym .dynstr .gnu.version .gnu.ve
03      .init_array .fini_array .jcr .dynamic .got .got.plt .data .bss
04      .dynamic
05      .note.ABI-tag .note.gnu.build-id
06      .eh_frame_hdr
07
08      .init_array .fini_array .jcr .dynamic .got

```

Sections

```
$ readelf --section-headers hello-world
```

There are 31 section headers, starting at offset 0x1a00:

Section Headers:

[Nr]	Name	Type	Address	Offset
	Size	EntSize	Flags Link Info	Align
[0]		NULL	0000000000000000	00000000
	0000000000000000	0000000000000000	0 0	0
[1]	.interp	PROGBITS	0000000000000238	00000238
	000000000000001c	0000000000000000	A 0 0	1
[2]	.note.ABI-tag	NOTE	0000000000000254	00000254
	0000000000000020	0000000000000000	A 0 0	4
[3]	.note.gnu.build-i	NOTE	0000000000000274	00000274
	0000000000000024	0000000000000000	A 0 0	4
[4]	.gnu.hash	GNU_HASH	0000000000000298	00000298
	000000000000001c	0000000000000000	A 5 0	8
[5]	.dynsym	DYNSYM	00000000000002b8	000002b8
	00000000000000c0	0000000000000018	A 6 1	8
[6]	.dynstr	STRTAB	0000000000000378	00000378

	00000000000000096	00000000000000000	A	0	0	1
[7]	.gnu.version	VERSYM	0000000000000040e	00000040e		
	00000000000000010	00000000000000002	A	5	0	2
[8]	.gnu.version_r	VERNEED	00000000000000420	000000420		
	00000000000000020	00000000000000000	A	6	1	8
[9]	.rela.dyn	RELA	00000000000000440	000000440		
	000000000000000d8	00000000000000018	A	5	0	8
[10]	.rela.plt	RELA	00000000000000518	000000518		
	00000000000000018	00000000000000018	AI	5	24	8
[11]	.init	PROGBITS	00000000000000530	000000530		
	00000000000000017	00000000000000000	AX	0	0	4
[12]	.plt	PROGBITS	00000000000000550	000000550		
	00000000000000020	00000000000000010	AX	0	0	16
[13]	.plt.got	PROGBITS	00000000000000570	000000570		
	00000000000000008	00000000000000000	AX	0	0	8
[14]	.text	PROGBITS	00000000000000580	000000580		
	000000000000001b1	00000000000000000	AX	0	0	16
[15]	.fini	PROGBITS	00000000000000734	000000734		
	00000000000000009	00000000000000000	AX	0	0	4
[16]	.rodata	PROGBITS	00000000000000740	000000740		
	00000000000000012	00000000000000000	A	0	0	4
[17]	.eh_frame_hdr	PROGBITS	00000000000000754	000000754		
	0000000000000003c	00000000000000000	A	0	0	4
[18]	.eh_frame	PROGBITS	00000000000000790	000000790		
	0000000000000010c	00000000000000000	A	0	0	8
[19]	.init_array	INIT_ARRAY	0000000000200dd8	00000dd8		
	00000000000000008	00000000000000008	WA	0	0	8
[20]	.fini_array	FINI_ARRAY	0000000000200de0	00000de0		
	00000000000000008	00000000000000008	WA	0	0	8
[21]	.jcr	PROGBITS	0000000000200de8	00000de8		
	00000000000000008	00000000000000000	WA	0	0	8
[22]	.dynamic	DYNAMIC	0000000000200df0	00000df0		
	000000000000001e0	00000000000000010	WA	6	0	8
[23]	.got	PROGBITS	0000000000200fd0	00000fd0		
	00000000000000030	00000000000000008	WA	0	0	8
[24]	.got.plt	PROGBITS	0000000000201000	00001000		

```

000000000000000020 000000000000000008 WA      0      0      8
[25] .data          PROGBITS          000000000000201020 00001020
      000000000000000010 000000000000000000 WA      0      0      8
[26] .bss           NOBITS            000000000000201030 00001030
      000000000000000008 000000000000000000 WA      0      0      1
[27] .comment        PROGBITS          000000000000000000 00001030
      00000000000000002d 000000000000000001 MS      0      0      1
[28] .symtab          SYMTAB            000000000000000000 00001060
      0000000000000000660 000000000000000018          29      48      8
[29] .strtab          STRTAB            000000000000000000 000016c0
      000000000000000022f 000000000000000000          0      0      1
[30] .shstrtab        STRTAB            000000000000000000 000018ef
      000000000000000010c 000000000000000000          0      0      1

```

Key to Flags:

W (write), A (alloc), X (execute), M (merge), S (strings), I (info),
 L (link order), O (extra OS processing required), G (group), T (TLS),
 C (compressed), x (unknown), o (OS specific), E (exclude),
 l (large), p (processor specific)

Interesting sections

For a reverse engineer the following sections are (usually) of interest:

- `.text` - containing the majority of code, both user-written code and boilerplate generated by the compiler
- `.init` - containing code usually generated by the compiler that is supposed to run before `main()` has been called
- `.fini` - containing code usually generated by the compiler that is supposed to run after `main()` has been called
- `.plt` - containing code generated by the compiler in order to call functions from libraries
- `.rodata` - containing Read Only Data used by the program (strings, constants, etc)
- `.data` - containing Read/Write Data, used for initialized variables, mutable strings etc

- `.bss` - containing Read/Write Data, used for uninitialized global variables

For an exploit developer the following additional sections are of interest for information leakage/control flow hijacking purposes:

- `.got` / `.got.plt` - (Global Offset Table) containing pointers used in library call resolution
- `.init_array` / `.fini_array` - containing pointers used in the code from the `.init/.fini` sections

Check the man page of `readelf` and see what other ways you can use it to learn more about your target binary.

Usage of Advanced Tools

Recall the list of advanced tools from the beginning. All of them have the following functionality (*to various degrees*):

- Code Navigation
 - identifying functions and navigating between them
 - various code views (Listing (Assembly), Decompilation, Graph-View, HexView, etc.) and the option of switching between them, while maintaining context.
 - find usages of the current function/variable/item (cross-references/xrefs)
- Renaming/Redeclaring
 - changing the signature of a function
 - changing the type of a variable/item
 - changing the name of a function/variable/item
- Reorganizing the stack variables
 - change a stack variable into something else (smaller, bigger, structure, array); **be mindful of how much space there is available for your desired actions.**

Note that when turning into an array it's ideal to first change the variable into the array unit (e.g. if you want to change a stack space into `int v[30]` , and `v` is currently of type `char` , first turn `v` into an `int`) and then change its type into the desired array. Your tool of choice might give you suggestions regarding the ideal/maximum array size.

Experiment!

Write a small program in C (or any other compiled language that you like), and open the resulting binary in any of the tools suggested here (or some other that your know of). Learn to navigate your tool, make the output look more pretty, try to work backwards from the binary to the source code.

See how various compiler flags affect the output of your tool (`-g` , `-s` , `0x` , etc).

When in doubt Right Click!, hover, or check the cheatsheet/docs in the [resources section](#)

Lab Tasks

For today's tasks you can use any of the *advanced* tools mentioned in the [setup](#) section above, **but** in the hints along the way we will assume you are using Ghidra.

Task 1: Reverse engineering with spoilers (5p)

Usually, when reverse engineering, all we have is a binary. Starting from it, we need to reconstruct (mainly through guessing/inferences) what the function names could be, what the variables are used for, what the program does as a whole.

For this task, you have a binary, `task1` , and also its corresponding source code, `task1.c` . Using the stripped binary, you will simulate normal reverse engineering by using the source code (instead of guessing).

Your task is to create a near-original replica of the original source in the IDA interface:

1. Rename and retype the 4 functions in the source code (aside from `main()`) (2p)
2. Rename and retype the stack variables in `setup()` and `main()` (1p)
3. Rename and retype the stack variables (including the arrays) in `chance()` and `gen_rand_string()` (2p)

Task 2: Statically linked crackme - graybox analysis (dynamic + static) (4p)

In this task, you will learn to navigate through functions in a statically linked and stripped `crackme`. (*What is a stripped binary?*)

Since the binary has a whopping 783 functions detected, you do not have the time or motivation to go through all of them. As such, you need to approach the problem in a clever and elegant way:

1. Run the program once, note any strings. Go to the `.rodata` segment and find any/all of the strings. Using the xref (cross-reference) functionality, determine where the `main()` function is. (1p)
2. Rename all the functions in `main()` and determine the password checking function. (1p)
3. In the password checking function, observe how the correct password is generated; we want to make this function more legible. Go to the location of any `DAT_XXX` variable in the Ghidra decompiler view. Notice that each of them is a letter of the alphabet. Find the location of the start of the alphabet and redeclare that address as a wide C string (Right Click->Data->TerminatedUnicode). Again, in the password checking function, observe how the decompiler view looks now. Redecclare the alphabet as constant data (Right Click->Data->Mutability). After changing the data type of the corresponding local variable accordingly, the assignment sequence should collapse and reveal the correct password. Finally check that the password is accepted (2p)

Task 3: Data Structures (6p)

In this task, you will learn to use the *structures* functionality of Ghidra.

Only the simplest programs are written without any sort of data structures in mind. Even basic OOP features are implemented using structures; classes themselves are also compiled as structures. However, after compilation, structure and type information is lost (if we don't have debugging symbols) but we can still observe repeated access patterns and infer what various structures might have looked like (*recall how structure fields are accessed in ASM from the previous lab*).

Look at the code in main and the password checking function, analyze the access patterns and verify that it matches the linked list structure below.

offset	size	mnemonic	data type	field name
0x0	0x4	ddw	dword	field_0_idx
0x4	0x1	db	byte	field_4
0x5	0x1	db	byte	
0x6	0x1	db	byte	
0x7	0x1	db	byte	
0x8	0x8	dq	astruct*	*field_8_next

You can either:

- go to DataTypeManager->task3->Right Click->New->Structure
- go to your variable, Right Click->Auto Create Structure, then Edit Data Type

1. Use the Structure Editor and create the list structure above, and make sure to also declare the last field as a `astruct*` pointer. (2p)

2. In main, cast the buffer returned from `malloc` and the head of the list to this struct type and propagate in the password checking function, renaming and retyping where necessary. (2p)
3. Describe what the code does and figure out the correct password (2p)

Task 4: Bonus (20p)

Check out `./bonus/bonus` . Verify that it work on the remote server **TODO**. No partial points for this one.

You can present solutions for this task until the end of the semester.

Resources

1. <https://docs.hex-rays.com/9.0/getting-started/basic-usage> ↩
2. <https://docs.binary.ninja/> ↩
3. <https://ghidra-sre.org/CheatSheet.html> ↩
4. <https://book.rizin.re/> ↩
5. <https://raw.githubusercontent.com/corkami/pics/refs/heads/master/binary/elf101/elf101.svg> ↩
6. https://media.ccc.de/v/31c3_-5930-en-saal_6-201412291400-funky_file_formats-_ange_albertini ↩
7. <https://media.ccc.de/v/38c3-fearsome-file-formats> ↩