

# NLP - Fall 2018 - Sample Midterm Exam

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- (1) TrueCasing is the process of taking text with missing or unreliable case information and producing the proper case for each word, e.g. if the input looks like:

as previously reported , target letters were issued last month to  
michael milken , drexel 's chief of junk-bond operations ; mr. milken 's  
brother lowell ; cary maultasch , a drexel trader ; james dahl , a  
drexel bond salesman ; and bruce newberg , a former drexel trader .

Then the output of the TrueCasing program should be:

As previously reported , target letters were issued last month to  
Michael Milken , Drexel 's chief of junk-bond operations ; Mr. Milken 's  
brother Lowell ; Cary Maultasch , a Drexel trader ; James Dahl , a  
Drexel bond salesman ; and Bruce Newberg , a former Drexel trader .

Assume **we can only use** the following two probability distributions:

- A *translation probability*  $P(w | W)$  where  $w$  is the lowercase variant of the TrueCase word  $W$  (note that the TrueCase word might still be lowercase). The function `lower` can be used to lowercase a word, e.g. `“HAL9001”.lower() = “hal9001”`
- A *bigram probability*  $P(W | W')$ . A language model  $P(W_1, \dots, W_n)$  is used to provide the probability of a sentence. A bigram language model approximates the probability of a sentence as follows:

$$\Pr(W_1, \dots, W_n) \approx \prod_{i=1}^n P(W_i | W_{i-1})$$

Assume that  $W_{-1} = w_{-1} = \text{none}$  is a dummy word that begins each sentence.

- Assume that  $c(\cdot)$  gives the frequency of unigrams, bigrams, etc.

- a. Complete the following formula to provide a model of the TrueCasing task by using only the translation probability  $P(w | W)$  and the bigram probability  $P(W | W')$ :

$$\begin{aligned} W_1^*, \dots, W_n^* &= \arg \max_{W_1, \dots, W_n} \Pr(W_1, \dots, W_n | w_1, \dots, w_n) \\ &= \text{provide this formula} \end{aligned}$$

*Answer:*

$$\begin{aligned} W_1^*, \dots, W_n^* &= \arg \max_{W_1, \dots, W_n} P(W_1, \dots, W_n | w_1, \dots, w_n) \\ &= \frac{P(W_1, \dots, W_n) \cdot P(w_1, \dots, w_n | W_1, \dots, W_n)}{P(w_1, \dots, w_n)} \\ &\approx P(W_1, \dots, W_n) \cdot P(w_1, \dots, w_n | W_1, \dots, W_n) \\ &= \prod_{i=1}^n \underbrace{P(w_i | W_i)}_{\text{translation probability}} \cdot \underbrace{P(W_i | W_{i-1})}_{\text{bigram language model}} \end{aligned}$$

- b. Using maximum likelihood, provide a formula to estimate the the translation probability parameters  $P(w | W)$  for lowercase words  $w$  and TrueCase words  $W$ . Assume you **only** have access to a sufficient amount of TrueCase text.

*Answer:* For each TrueCase word  $W$  convert it to lowercase  $w$  using  $w = W.lower()$  and count  $f(w, W)$  and  $f(W)$ . Then,

$$P(w | W) = \frac{f(w, W)}{\sum_{w'} f(w', W)}$$

- c. Provide the equation that correctly computes add one smoothing for  $P(w | W)$ .

*Answer:*

$$P(w | W) = \frac{1 + f(w, W)}{|w| + \sum_{w'} f(w', W)}$$

- d. Backoff smoothing for  $P(W_i | W_{i-1})$  is defined as follows:

$$P_{bo}(W_i | W_{i-1}) = \begin{cases} \frac{c^*(W_{i-1}, W_i)}{c(W_{i-1})} & \text{if } c(W_{i-1}, W_i) > 0 \\ \alpha(W_{i-1})P_{bo}(W_i) & \text{otherwise} \end{cases}$$

where  $c^*(W_{i-1}, W_i) = c(W_{i-1}, W_i) - D$  for some  $0 < D < 1$  and  $\alpha(W_{i-1})$  is chosen to make sure that  $P_{bo}(W_i | W_{i-1})$  is a proper probability. Provide the equation to compute  $\alpha(W_{i-1})$ . Assume that  $\sum_{W_i} P_{bo}(W_i) = 1$ .

*Answer:*

$$\alpha(W_{i-1}) = 1 - \sum_{W_i} \frac{c^*(W_{i-1}, W_i)}{c(W_{i-1})}$$

## (2) Language Models

For the CFG  $G$  given below:

$$\begin{aligned} S &\rightarrow A | c \\ A &\rightarrow B a \\ B &\rightarrow b S \end{aligned}$$

- a. What is the language  $L(G)$ ?

*Answer:*  $b^n c a^n : n \geq 0$

- b. Assign probabilities to each rule in the CFG above so that for each string  $w \in L(G)$ :

$$P(w) = \exp\left(\frac{|w| - 1}{2} \times \ln(0.3) + \ln(0.7)\right)$$

where,  $|w|$  is the length of string  $w$ ,  $exp$  is exponentiation, and  $ln$  is  $log$  base  $e$ . Using an example, briefly explain *why* your PCFG provides the desired  $P(w)$  for any  $w$ .

*Answer:*

0.3  $S \rightarrow A$

0.7  $S \rightarrow c$

1.0  $A \rightarrow B a$

1.0  $B \rightarrow b S$

Since  $w \in \{b^n c a^n : n \geq 0\}$  then we always have one  $c$  in  $w$  which is derived using  $S \rightarrow c$  with prob 0.7. Removing  $c$  from  $w$  we get length  $|w| - 1$ . Each  $b \dots a$  pair in  $w$  is derived by first using rule  $S \rightarrow A$  with prob 0.3. The subsequent  $A \rightarrow Ba$  and  $B \rightarrow bS$  for each  $b \dots a$  pair each get prob 1.0 as there are no competing lhs  $A$  and  $B$  rules. There are  $\frac{|w|-1}{2}$  matching  $b \dots a$  pairs in each  $w$ . Hence deriving all of them will take probability  $(0.3)^{\frac{|w|-1}{2}}$ . Since:

$$(0.3)^{\frac{|w|-1}{2}} \times 0.7 = \exp\left(\frac{|w|-1}{2} \times \ln(0.3) + \ln(0.7)\right)$$

we obtain the required definition for  $P(w)$ .