

### Natural Language Processing

Angel Xuan Chang angelxuanchang.github.io/nlp-class adapted from lecture slides from Anoop Sarkar

Simon Fraser University

January 14th, 2020

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Part 1: Probability and Language

#### Probability and Language

Quick guide to probability theory

Entropy and Information Theory

### Probability and Language

#### Assign a probability to an input sequence

Given a URL: choosespain.com. What is this website about?

Input	Scoring function
choose spain	-8.35
chooses pain	-9.88
:	:

#### The Goal

Find a good **scoring function** for input sequences.

# Scoring Hypotheses in Speech Recognition

#### From acoustic signal to candidate transcriptions

Hypothesis	Score
the station signs are in deep in english	-14732
the stations signs are in deep in english	-14735
the station signs are in deep into english	-14739
the station 's signs are in deep in english	-14740
the station signs are in deep in the english	-14741
the station signs are indeed in english	-14757
the station 's signs are indeed in english	-14760
the station signs are indians in english	-14790
the station signs are indian in english	-14799
the stations signs are indians in english	-14807
the stations signs are indians and english	-14815

# Scoring Hypotheses in Machine Translation

#### From source language to target language candidates

Hypothesis	Score
we must also discuss a vision .	-29.63
we must also discuss on a vision .	-31.58
it is also discuss a vision .	-31.96
we must discuss on greater vision .	-36.09
÷	:

# Scoring Hypotheses in Spelling Correction

#### Substitute spelling variants to generate hypotheses

Hypothesis	Score
stellar and versatile acress whose combination	-18920
of sass and glamour has defined her	
stellar and versatile acres whose combination	-10209
of sass and glamour has defined her	
stellar and versatile actress whose combination	-9801
of sass and glamour has defined her	

# Scoring Hypotheses in Decryption

### Character substitutions on ciphertext to plaintext candidates

Hypothesis	Score
Heopaj, zk ukq swjp pk gjks w oaynap?	-93
Urbcnw, mx hxd fjwc cx twxf j bnlanc?	-92
Wtdepy, oz jzf hlye ez vyzh I dpncpe?	-91
Mjtufo, ep zpv xbou up lopx b tfdsfu?	-89
Nkuvgp, fq aqw ycpv vq mpqy c ugetgv?	-87
Gdnozi, yj tjp rvio oj fijr v nzxmzo?	-86
Czjkve, uf pfl nrek kf befn r jvtivk?	-85
Yvfgra, qb lbh jnag gb xabj n frperg?	-84
Zwghsb, rc mci kobh hc ybck o gsqfsh?	-83
Byijud, te oek mqdj je adem q iushuj?	-77
Jgqrcl, bm wms uylr rm ilmu y qcapcr?	-76
Listen, do you want to know a secret?	-25

#### The Goal

- ▶ Write down a **model** over sequences of words or letters.
- ▶ **Learn** the parameters of the model from data.
- Use the model to predict the probability of new sequences.

### Probabilistic Language Modeling

- Compute the probability of a sentence of sequence of words  $P(W) = P(w_1, w_2, w_3, w_4, \dots, w_n)$
- Build the sequence up over time and use chain rule
  - $P(w_1, w_2, w_3, w_3, w_4, \ldots, w_n) = P(w_1)P(w_2, \ldots, w_n|w_1)$
  - $P(w_1, w_2, w_3, w_3, w_4, \dots, w_n) = P(w_1)P(w_2|w_1)P(w_3|w_1, w_2)\dots P(w_n|w_1, \dots, w_{n-1})$
  - $P(w_1, w_2, w_3, w_3, w_4, \dots, w_n) = \prod_i P(w_i | w_1, \dots w_{i-1})$
  - ► P(it is very cold today) = P(it)P(is|it)P(very|it is)P(cold|it is very)P(today|it is very cold)
- ▶ Need to compute:  $P(w_i|w_1,...w_{i-1})$ 
  - Let's count!
  - $P(\text{cold}|\text{it is very}) = \frac{\text{Count}(\text{it is very cold})}{\text{Count}(\text{it is very})}$

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Part 2: Quick guide to probability theory

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Entropy and Information Theory

# Probability: The Basics

- ► Sample space
- Event space
- ► Random variable

### Probability distributions

- ightharpoonup P(X): probability of random variable X having a certain value.
  - P(X = killer) = 1.05e-05
  - ightharpoonup P(X = app) = 1.19e-05

### Joint probability

- P(X,Y): probability that X and Y each have a certain value.
  - Let Y stand for choice of a word
  - Let X stand for the choice of a word that occurs before Y
  - ▶ P(X = killer, Y = app) = 1.24e-10

### Joint Probability: P(X=value AND Y=value)

- Since X=value AND Y=value, the order does not matter
- $ightharpoonup P(X = killer, Y = app) \Leftrightarrow P(Y = app, X = killer)$
- ▶ In both cases it is P(X,Y) = P(Y,X) = P('killer app')
- ▶ In NLP, we often use numerical indices to express this:  $P(W_{i-1} = \text{killer}, W_i = \text{app})$

### Joint probability

#### Joint probability table

$W_{i-1}$	$W_i = app$	$P(W_{i-1}, W_i)$
$\langle S \rangle$	арр	1.16e-19
an	арр	1.76e-08
killer	арр	1.24e-10
the	арр	2.68e-07
this	арр	3.74e-08
your	арр	2.39e-08

There will be a similar table for each choice of  $W_i$ .

Get 
$$P(W_i)$$
 from  $P(W_{i-1}, W_i)$ 

$$P(W_i = \mathsf{app}) = \sum_{x} P(W_{i-1} = x, W_i = \mathsf{app}) = 1.19e - 05$$

# Conditional probability

- ▶  $P(W_i | W_{i-1})$ : probability that  $W_i$  has a certain value after fixing value of  $W_{i-1}$ .
- $ightharpoonup P(W_i = app \mid W_{i-1} = killer)$
- $P(W_i = \mathsf{app} \mid W_{i-1} = \mathsf{the})$

#### Conditional probability from Joint probability

$$P(W_i \mid W_{i-1}) = \frac{P(W_{i-1}, W_i)}{P(W_{i-1})}$$

- ▶ P(killer) = 1.05e-5
- ▶ P(killer, app) = 1.24e-10
- ▶ P(app | killer) = 1.18e-5

#### **Basic Terms**

- ▶ P(e) a priori probability or just prior
- ▶  $P(f \mid e)$  conditional probability. The chance of f given e
- ▶ P(e, f) *joint* probability. The chance of e and f both happening.
- If e and f are independent then we can write  $P(e, f) = P(e) \times P(f)$
- ▶ If e and f are not independent then we can write  $P(e, f) = P(e) \times P(f \mid e)$  $P(e, f) = P(f) \times ?$

#### **Basic Terms**

► Addition of integers:

$$\sum_{i=1}^{n} i = 1 + 2 + 3 + \ldots + n$$

Product of integers:

$$\prod_{i=1}^{n} i = 1 \times 2 \times 3 \times \ldots \times n$$

Factoring:

$$\sum_{i=1}^{n} i \times k = k + 2k + 3k + \ldots + nk = k \sum_{i=1}^{n} i$$

Product with constant:

$$\prod_{i=1}^{n} i \times k = 1k \times 2k \dots \times nk = k^{n} \times \prod_{i=1}^{n} i$$

### Probability: Axioms

- P measures total probability of a set of events
- $ightharpoonup P(\emptyset) = 0$
- ightharpoonup P(all events) = 1
- ▶  $P(X) \le P(Y)$  for any  $X \subseteq Y$
- ▶  $P(X) + P(Y) = P(X \cup Y)$  provided that  $X \cap Y = \emptyset$

### **Probability Axioms**

▶ All events sum to 1:

$$\sum_{e} P(e) = 1$$

 $\blacktriangleright$  Marginal probability P(f):

$$P(f) = \sum_{e} P(e, f)$$

Conditional probability:

$$\sum_{e} P(e \mid f) = \sum_{e} \frac{P(e, f)}{P(f)} = \frac{1}{P(f)} \sum_{e} P(e, f) = 1$$

ightharpoonup Computing P(f) from axioms:

$$P(f) = \sum_{e} P(e) \times P(f \mid e)$$

# Probability: The Chain Rule

- $\triangleright$   $P(a, b, c, d \mid e)$
- ▶ We can simplify this using the Chain Rule:
- $P(a,b,c,d \mid e) = P(d \mid e) \cdot P(c \mid d,e) \cdot P(b \mid c,d,e) \cdot P(a \mid b,c,d,e)$
- ▶ To see why this is possible, recall that  $P(X \mid Y) = \frac{p(X,Y)}{p(Y)}$
- We can approximate the probability by removing some random variables from the context. For example, we can keep at most two variables to get:

$$P(a, b, c, d \mid e) \approx P(d \mid e) \cdot P(c \mid d, e) \cdot P(b \mid c, e) \cdot P(a \mid b, e)$$

### Probability: The Chain Rule

 $P(e_1, e_2, \ldots, e_n) = P(e_1) \times P(e_2 \mid e_1) \times P(e_3 \mid e_1, e_2) \ldots$ 

$$P(e_1, e_2, \ldots, e_n) = \prod_{i=1}^n P(e_i \mid e_{i-1}, e_{i-2}, \ldots, e_1)$$

- ► In NLP, we call:
  - $\triangleright$   $P(e_i)$ : unigram probability
  - $\triangleright$   $P(e_i \mid e_{i-1})$ : bigram probability
  - $ightharpoonup P(e_i \mid e_{i-1}, e_{i-2})$ : trigram probability
  - ►  $P(e_i | e_{i-1}, e_{i-2}, ..., e_{i-(n-1)})$ : n-gram probability

# Probability: Random Variables and Events

- $\blacktriangleright$  What is y in P(y)?
- Shorthand for value assigned to a random variable Y, e.g. Y = y
- ightharpoonup y is an element of some implicit **event space**:  $\mathcal{E}$

### Probability: Random Variables and Events

► The marginal probability P(y) can be computed from P(x, y) as follows:

$$P(y) = \sum_{x \in \mathcal{E}} P(x, y)$$

Finding the value that maximizes the probability value:

$$\hat{x} = \arg\max_{x \in \mathcal{E}} P(x)$$

# Log Probability Arithmetic

- Practical problem with tiny P(e) numbers: underflow
- One solution is to use log probabilities:

$$\log(P(e)) = \log(p_1 \times p_2 \times \ldots \times p_n)$$
  
= 
$$\log(p_1) + \log(p_2) + \ldots + \log(p_n)$$

Note that:

$$x = \exp(\log(x))$$

Also more efficient: addition instead of multiplication

# Log Probability Arithmetic

р	$\log(p)$
0.0	$-\infty$
0.1	-3.32
0.2	-2.32
0.3	-1.74
0.4	-1.32
0.5	-1.00
0.6	-0.74
0.7	-0.51
0.8	-0.32
0.9	-0.15
1.0	0.00

# Log Probability Arithmetic

- So:  $(0.5 \times 0.5 \times \dots 0.5) = (0.5)^n$  might get too small but (-1-1-1-1) = -n is manageable
- Another useful fact when writing code (log<sub>2</sub> is log to the base 2):

$$\log_2(x) = \frac{\log_{10}(x)}{\log_{10}(2)}$$

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Part 3: Entropy and Information Theory

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# Information Theory

- ► Information theory is the use of probability theory to quantify and measure "information".
- Consider the task of efficiently sending a message. Sender Alice wants to send several messages to Receiver Bob. Alice wants to do this as efficiently as possible.
- ▶ Let's say that Alice is sending a message where the entire message is just one character a, e.g. aaaa.... In this case we can save space by simply sending the length of the message and the single character.

# Information Theory

- Now let's say that Alice is sending a completely random signal to Bob. If it is random then we cannot exploit anything in the message to compress it any further.
- The expected number of bits it takes to transmit some infinite set of messages is what is called entropy.
- ▶ This formulation of entropy by Claude Shannon was adapted from thermodynamics, converting information into a quantity that can be measured.
- Information theory is built around this notion of message compression as a way to evaluate the amount of information.

### Expectation

- For a probability distribution p
- **Expectation** with respect to *p* is a weighted average:

$$E_{p}[x] = \frac{x_{1} \cdot p(x_{1}) + x_{2} \cdot p(x_{2}) + \dots + x_{n}p(x_{n})}{p(x_{1}) + p(x_{2}) + \dots + p(x_{n})}$$

$$= x_{1} \cdot p(x_{1}) + x_{2} \cdot p(x_{2}) + \dots + x_{n}p(x_{n})$$

$$= \sum_{x \in \mathcal{E}} x \cdot p(x)$$

Example: for a six-sided die the expectation is:

$$E_p[x] = 1 \cdot \frac{1}{6} + 2 \cdot \frac{1}{6} + \ldots + 6 \cdot \frac{1}{6} = 3.5$$

### Entropy

- For a probability distribution p
- **Entropy** of *p* is:

$$H(p) = -\sum_{x \in \mathcal{E}} p(x) \cdot \log_2 p(x)$$

- Any base can be used for the log, but base 2 means that entropy is measured in bits.
- ▶ What is the *expected* number of bits with respect to *p*:

$$-E_p[\log_2 p(x)] = H(p)$$

▶ Entropy answers the question: What is the expected number of bits needed to transmit messages from event space  $\mathcal{E}$ , where p(x) defines the probability of observing x?

# Perplexity

- ▶ The value  $2^{H(p)}$  is called the **perplexity** of a distribution p
- Perplexity is the weighted average number of choices a random variable has to make.
- ▶ Choosing between 8 equally likely options (H=3) is  $2^3 = 8$ .

### Relative Entropy

- We often wish to determine the divergence of a distribution q from another distribution p
- Let's say q is the estimate and p is the true probability
- We define the *divergence* from q to p as the **relative** entropy: written as D(p||q)

$$D(p||q) = -\sum_{x \in \mathcal{E}} p(x) \log_2 \frac{q(x)}{p(x)}$$

Note that

$$D(p||q) = -E_{p(x)} \left[ \log_2 \frac{q(x)}{p(x)} \right]$$

► The relative entropy is also called the Kullback-Leibler divergence.

## Cross Entropy and Relative Entropy

The relative entropy can be written as the sum of two terms:

$$D(p||q) = -\sum_{x \in \mathcal{E}} p(x) \log_2 \frac{q(x)}{p(x)}$$
$$= -\sum_{x} p(x) \log_2 q(x) + \sum_{x} p(x) \log_2 p(x)$$

- We know that  $H(p) = -\sum_{x} p(x) \log_2 p(x)$
- ► Similarly define  $H(p,q) = -\sum_{x} p(x) \log_2 q(x)$

$$D(p||q) = H(p,q) - H(p)$$
  
relative entropy $(p,q)$ =cross entropy $(p,q)$ -entropy $(p)$ 

▶ The term H(p,q) is called the **cross entropy**.

# Cross Entropy and Relative Entropy

- $\vdash$   $H(p,q) \ge H(p)$  always.
- ▶  $D(p||q) \ge 0$  always, and D(p||q) = 0 iff p = q
- ▶ D(p||q) is not a true distance:
  - ▶ It is asymmetric:  $D(p||q) \neq D(q||p)$ ,
  - ► It does not obey the triangle inequality:  $D(p||q) \nleq D(p||r) + D(r||q)$

# Conditional Entropy and Mutual Information

Entropy of a random variable X:

$$H(X) = -\sum_{x \in \mathcal{E}} p(x) \log_2 p(x)$$

► Conditional Entropy between two random variables X and Y:

$$H(X \mid Y) = -\sum_{x,y \in \mathcal{E}} p(x,y) \log_2 p(x \mid y)$$

Mutual Information between two random variables X and Y:

$$I(X; Y) = D(p(x, y) || p(x)p(y)) = \sum_{x} \sum_{y} p(x, y) \log_2 \frac{p(x, y)}{p(x)p(y)}$$

#### Acknowledgements

Many slides borrowed or inspired from lecture notes by Anoop Sarkar, Dan Jurafsky, Michael Collins, Chris Dyer, Kevin Knight, Chris Manning, Philipp Koehn, Adam Lopez, Graham Neubig, Richard Socher and Luke Zettlemoyer from their NLP course materials.

All mistakes are my own.

A big thank you to all the students who read through these notes and helped me improve them.