#1 Reader Writer (Writers priority implementation)

```
int writers; // # writer threads that want to enter the critical section (some or
all of these may be blocked)
int writing; // Number of threads that are actually writing inside the C.S. (can
only be zero or one - can you see why?)
int reading; // Number of threads that are reading inside the C.S.
int readers; // Number of threads that are or want to read
// if writing !=0 then reading must be zero (and vice versa)
```

```
reader() {
                                          writer() {
    lock (&m)
                                              lock (&m)
    readers ++
                                              writers++
    while (
                                              while (
        cond wait(&r cv, &m)
                                                  cond wait(&w cv, &m)
    Do we need to wait for
    both 'writers' and 'writing'?
    reading++
                                              writing++
    unlock (&m)
                                              unlock (&m)
  // perform reading here
                                              // perform writing here
    lock (&m)
                                              lock (&m)
    reading--
                                              writing--
    readers--
                                              writers--
    wake up who here? (and how many)
                                              wake up who here? (and how many)
    unlock (&m)
                                              unlock (&m)
    return result
```

DEADLOCK

#2 Deadlock Definition:

#3 Coffman Conditions

```
Necessary? Y/N
Sufficient? Y/N
```

2

1

3

#4 Resource Allocation Graphs



Figure 1. Deadlock do not confuse it with dreadlocks. Assume processes acquire locks in the order specified and release resources only when finished. Create a resource allocation graph to determine if and when there is deadlock.

When a process waits for a resource it will acquire an exclusive lock on resource as soon as no other process has an exclusive lock. Assume locks are fair (earliest waiting process obtains the lock).

| Q1 Process 1 ("P1") requests (and obtains) Resource A and then Resource B Process 2 requests C and then B. Deadlock for P1? P2? | |
|--|--|
| Q2 P1 requests (and obtains?) A P2 requests (and obtains?) B P3 requests (and obtains?) C P2 requests (and obtains?) C P3 requests (and obtains?) A P1 requests (and obtains?) C | |
| Q3 P1 requests A then B P2 requests C then B P3 requests B P4 requests C then B Deadlock for P1? P2? P3? P4? | |
| Q4 P1 requests A then B P2 requests C, D then B P4 requests D P3 requests B P1 requests C Deadlock for P1? P2? P3? P4? | |
| Q5 P1 requests A and B P2 requests C and D then B P4 requests D P3 requests B P1 releases B (thus P2 acquires B) P1 requests C Deadlock for P1? P2? P3? P4? | |

- #5 What is the Banker's Algorithm?
- #6 Deadlock Avoidance
- #7 Linux/Windows strategy for deadlock avoidance?
- #8 Acquiring resources in same rank