

Mid Semester Examination

March 2019

Max. Marks: 20

Class: FYMCA

Course Code: MCA23

Name of the Course: Data Structures

Duration: 1 Hr

Semester: II

Branch: MCA

Synopsis

1) Worst case complexity calculation of Insertion sort (2 mks)

Worst case complexity calculation of Bubble sort (2 mks)

Example : 12, 10, 8, 6, 5

Insertion sort : Number of passes and comparison in first pass = 4

Number of passes and comparison in second pass = 8

Number of passes and comparison in third pass = 12

Number of passes and comparison in forth pass = 16

Hence the complexity = $4 + 8 + 12 + 16 = O(n^2)$

Bubble sort : Number of passes and comparison in first pass = 16

Number of passes and comparison in second pass = 12

Number of passes and comparison in third pass = 8

Number of passes and comparison in forth pass = 4

Hence the complexity = $16 + 12 + 8 + 4 = O(n^2)$

OR

Number of passes and Iterations calculation for Insertion sort (2 mks)

Number of passes and Iterations calculation for Selection sort (2 mks)

Example: 10, 2, -20, 30, 12, 34

Insertion sort :

| | 0 | 1 | 2 | 3 | 4 | 5 | No. Of Iterations |
|-----------------------|-----|----|-----|----|----|----|-------------------|
| 0 (Original Array) | 10 | 2 | -20 | 30 | 12 | 34 | |
| 1 | 2 | 10 | | | | | 1 |
| 2 | -20 | 2 | 10 | | | | 2 |
| 3 | -20 | 2 | 10 | 30 | | | 1 |
| 4 | -20 | 2 | 10 | 12 | 30 | | 2 |
| 5 | -20 | 2 | 10 | 12 | 30 | 34 | 1 |

Selection sort:

| | 0 | 1 | 2 | 3 | 4 | 5 | No. Of Iterations |
|-----------------------|-----|---|-----|----|----|----|-------------------|
| 0 (Original Array) | 10 | 2 | -20 | 30 | 12 | 34 | |
| 1 | 10 | 2 | -20 | 30 | 12 | 34 | 5 |
| 2 | 10 | 2 | -20 | 12 | 30 | 34 | 4 |
| 3 | 10 | 2 | -20 | 12 | 30 | 34 | 3 |
| 4 | -20 | 2 | 10 | 12 | 30 | 34 | 2 |
| 5 | -20 | 2 | 10 | 12 | 30 | 34 | 1 |

Comparison Table:

| | No. Of passes | No. Of Iterations |
|----------------|---------------|-------------------|
| Insertion Sort | 5 | 7 |
| Selection Sort | 5 | 15 |

2) Applying double hashing for each key value.

(1 mk)

Formula is,

$$h(k, i) = (h_1(k) + i h_2(k)) \bmod n$$

$$h_1(k) = k \bmod n$$

and

$$h_2(k) = 1 + (k \bmod (n-1))$$

For 12,

$$h(12, 0) = (h_1(12) + 0 h_2(12)) \bmod 10$$

$$h_1(12) = 12 \% 10 = 2$$

Hence 12 is stored at 2nd Position

For 2,

$$h(2, 0) = (h_1(2) + 0 h_2(2)) \bmod 10$$

$$h_1(2) = 2 \% 10 = 2$$

Here, 2 position is already full. So we need to calculate $h(2, 1)$

$$h(2, 1) = (h_1(2) + 1 h_2(2)) \bmod 10$$

$$h_1(2) = 2 \quad \text{and} \quad h_2(2) = 1 + (2 \% 9) = 3$$

$$h(2, 1) = (2 + 3) \% 10 = 5$$

Hence 2 is stored at 5th position

Likewise 18 will be stored at 8th location and 45 will be stored at 6th location.

3) Solution of Construct an algorithm for reversing of singly linked list elements is

(4 mks)

Begin

Initialize q and s pointer with list pointer

Initialize temp pointer will NULL

Initialize r pointer with q->next

Repeat

Set temp to q

Set q to r

Set r to q->next

Set q->next to temp

Until r not equal to NULL

Set list to q

Set s->next to NULL

End

4) Pushing each symbol onto stack.

(0.25 mk each)

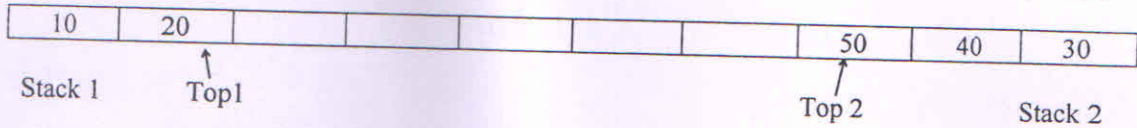
| Symbol | (| A | + | B |) | * | C | / | D | / | E |) |
|--------|---|---|---|---|---|---|---|---|---|---|----|----|
| Index | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |

| Index | Symbol | Stack | Expression |
|-------|--------|-------|------------|
| 0 | (| ((| |
| 1 | A | ((| A |
| 2 | + | ((+ | A |
| 3 | B | ((+ | AB |
| 4 |) | (| AB+ |
| 5 | * | (* | AB+ |
| 6 | C | (* | AB+C |
| 7 | / | (*/ | AB+C |
| 8 | D | (*/ | AB+CD |
| 9 | / | (*/ | AB+CD/ |
| 10 | E | (*/ | AB+CD/E |
| 11 |) | | AB+CD/E/* |

Required Postfix expression is AB+CD/E/*

OR

Solution of construct an algorithm to implement 2 Stacks using only one array is attached here with. (4 mks)



Algorithm :

Step 1 : Start two indexes one at left end and other at right end.
Step 2 : The left index jumps to the right and the right index jumps to the left.

Step 3 : If we want to push an element into the 2nd stack

index.

Step 4 : Similarly, if we want to push an element into the second stack then put at right index.

Step 5 : First stack grows towards right and second stack grows towards left.

5) Solution of construct enqueue and dequeue algorithms to implement queue using 2 stacks. (2 mks each)

Let S1 and S2 be two stacks to be used to implement Queue.

Struct Queue

```

{
    Struct Stack *s1;           // for Enqueue
    Struct Stack *s2;           // for Dequeue
};

```

Enqueue Algorithm :

Just push on to Stack S1

```
void Enqueue(Struct Queue *Q, int data)
```

```

{
    Push (Q -> S1, data);
}

```

Dequeuing Algorithm:

- ✧ If stack S2 is not empty, then pop S2 and return that element.
- ✧ If stack is empty; then transfer all elements from S1 to S2 and pop the top element from S2 and return that top element.
- ✧ If stack S1 is empty then throw an error.

```
int Dequeue(Struct Queue *Q)
{
    if(!IsEmptyStack(Q -> S2))
        return pop(Q -> S2);
    else
    {
        while(!IsEmptyStack(Q -> S1))
            push(Q -> S2, pop(Q -> S1));
        return pop(Q -> S2);
    }
}
```