

Sardar Patel Institute of Technology Bhavan's Campus, Munshi Nagar, Andheri (West), Mumbai-400058, India (Autonomous College Affiliated to University of Mumbai)

Duration: 3 Hrs.

Branch: M.C.A.

Semester: IV

Make up Examination

Synoptic-May 2018

Max. Marks: 100

Class: S.Y.

Course Code:MCA43

Name of the Course: Design and Analysis of Algorithms

Instruction:

(1) All questions are compulsory

(2) Draw neat diagrams

(3) Assume suitable data if necessary

Q. No. Q.1	Evaluin the second seco	Max. Marks	CO
a)	Explain the various Asymptotic notations with diagrams.	5	CO1
	O-notation		
	$O(g(n)) = \{f(n): \text{ there exist positive constants c and n0 such that } 0 \le f(n) \le cg(n) \text{ for all } n \ge n0\}.$ Ω -notation		
	$\Omega(g(n)) = \{f(n): \text{ there exist positive constants c and n0 such that } 0 \le cg(n) \le f(n) \text{ for all } n \ge n0\}.$		
	Θ -notation $\Theta(g(n)) = \{f(n) : \text{there exist positive constants c1, c2, and n0 such that } 0 \le c1g(n) \le f(n) \le c2g(n)$		
	for all $n \ge n0$. (1910) (1910) (1910) (1910) (1910)	C .	
	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$		
)	What is the time complexity of following function fun (). Time Complexity of the above function O(n).	7	CO1

c)	Draw the recursion tree for the recurrence $T(n) = 2T(n/2) + n2$. up to first 3 levels.		CO1
	$(n/2)^2$ $(n/2)^2$		
	$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	001 a	
		10154 10154	
	OR Compare P and NP problems with suitable examples.		
Q. 2 a)	Find the LCS of string 1: AGGTAB string2: GXTXAYB	6	CO2
b)	UCS of AGGTAB and GXTXAYB is GTAB Write Kruskal's algorithm for minimum spanning tree.	8	CO2
	Given a chain of four matrices A1, A2, A3, A4 with p0=5, p1=4, p2=6, p3=2, p4=7. Find m[1, 4]. Ans: m[1,4]=158 A1		
c)	Describe the Dynamic 0/1 Knapsack Problem.		
	instance for n=4, m=8, profits are (p1, p2, p3, p4) = (15,10,9,5), weights are (w1,w2,w3,w4)=(1,5,3,4).	6	CO2
	The solution is $x = (1,0,1,1)$ i.e. items 1,3, and 4 are selected. value of		

Discuss the Graph coloring Problem. What technique is used to solve the problem? Write the algorithm to solve above problem. Let C[1j-1] be a partial coloring for the first j-1 nodes. Color(C,j,k,n) if j = n+1 then output C return or quit for i = 1 to k C[j] = i if valid(C,j,n) then Color(C,j+1,k,n) where Valid(C,j,n) for all neighbors v of j with v < j if C[v] = C[j] then eturn false return true OR Explain the branch and bound strategy and write an algorithm for asic branch and bound strategy.		7
Let $C[1j-1]$ be a partial coloring for the first $j-1$ nodes. Color(C,j,k,n) if $j=n+1$ then output C return or quit for $i=1$ to k $C[j]=i$ if $valid(C,j,n)$ then $color(C,j+1,k,n)$ where Valid(C,j,n) for all neighbors v of j with $v < j$ if $C[v] = C[j]$ then eturn false return true		
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eturn false return true		
return true		1
OB	1	
OR OR	1	
Xniain the branch and i		
and bound strategy and write an algorithm for	, "	
asic branch and bound strategy.		
Bound (S, U)		
19. 3 15 d leaving ctotal II -		
(I DOLD ID (G)		
$(LBOUND(S) >= U) \qquad return(Y);$		
(OBOUND(S) < U)U = UBOUND(S); if (LBOUND(S) < U)		
$r(i=0) : \langle k_1, \ldots, k_k;$		1
(BB(S;U) < U)U = DB(S)		
urn(U):		
scuss the Hamiltonian Cycles P. II		
ve the problem? Write the algorithm to solve above problem. gorithm HC (Node, String, edges)	7	CO3
(as, sumg, edges)		
f (edges=N and E(SN,Node))		
nen		
Print (String) //HC found	[1]	
EXIL		
(edges=N)		
or each child X of Node and Notln(X,String)		
HC(X String address to		
Done Done		
	(LBOUND(S) >= U) return(\(\frac{1}{2}\); (UBOUND(S) < U) U = UBOUND(S); if (LBOUND(S) < U) RANCH(S) -> S ₁ ,, S _k ; r (i = 0; i < k; i++) (BB(S _i , U) < U) U = BB(S _i); urn(U); scuss the Hamiltonian Cycles Problem. What technique is used to ve the problem? Write the algorithm to solve above problem. gorithm HC (Node, String, edges) f (edges=N and E(SN,Node)) Then Print (String) //HC found Exit f (edges=N) Then eturn or each child X of Node and NotIn(X,String) String = String + info(X) HC(X,String, edges+1) Done Return	(LBOUND(S) >= U) return(¥); (UBOUND(S) < U) U = UBOUND(S); if (LBOUND(S) < U) RANCH(S) -> S ₁ ,, S _k ; r (i = 0; i < k; i++) (BB(S _i ,U) < U) U = BB(S _i); surn(U); scuss the Hamiltonian Cycles Problem. What technique is used to ve the problem? Write the algorithm to solve above problem. gorithm HC (Node, String, edges) f (edges=N and E(SN,Node)) Then Print (String) //HC found Exit f (edges=N) Then eturn or each child X of Node and NotIn(X,String) String = String + info(X) HC(X,String, edges+1) Done

c)	Write an algorithm for the subset sum problem. Algorithm sumofsubset(s,k,r)	6	CO
	{		
B	//generate the left child. note $s+w(k) \le Bk-1$ is true. $X\{k\}=1$;		
	If $(S+W[k]=m)$ then write $(X[1:k])$; // there is no recursive call here as		
	W 0,1 \- \- .		
	Else if $(S+W[k]+W[k+1] \le m)$ then sum of sub $(S+W[k], k+1,r-1)$		
	\(\(\lambda_{\kappa}\),		
	//generate right child and evaluate Bk.		
	If $((S+r-W[k]>=m)$ and $(S+W[k+1]<=m))$ then		
	$X_{k]=0}$		
	sum of sub $(S, k+1, r- W[k])$;		
	}		
	} the part of and the promise and many of another than		47 .
	Cind the Assert		
	Find the target goal state for the 15-puzzel problem upto 2 levels for the given initial state.		
	1 2 3 4		
	7 8 9 10		
	5 6 11		
	12 13 14 15	*	
Q. 4	Write an algorithm to find the shortest path using Dijkstra's	7	C04
a)	argorumi.	,	C04
	function Dijkstra(Graph, source):		
	create vertex set Q		
	for each vertex v in Graph: $dist[v] \leftarrow INFINITY$		
	for each vertex v in Graph: $dist[v] \leftarrow INFINITY$ $prev[v] \leftarrow UNDEFINED$		
	add v to Q		
	$dist[source] \leftarrow 0$		
	while Q is not empty:		
	$u \leftarrow \text{vertex in } Q \text{ with min } \text{dist}[u]$		
	remove u from Q		
	for each neighbor v of u:		
	alt \leftarrow dist[u] + length(u, v).		
	:f -14 < 1' +5 ?		-1-
	$\begin{array}{ccc} & \text{fi alt } < \text{dist[v]:} & \text{dist[v]} \leftarrow \text{alt} \\ & \text{prev[v]} \leftarrow \text{u} \end{array}$	- 1	
	return dist[], prev[]		
)	Make use of Flyod Warshalls algorithm to find all pair shortest path		
	from every node	6	CO4

	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
c)	Ans: 1 2 3 4 5 1 0 1 -3 2 -4 2 3 0 -4 1 -1 3 7 4 0 5 3 4 2 -1 -5 0 -2 5 8 5 1 6 0
	Write an algorithm for Knuth Morris Pratt string matching with 7 CO4 example. Algorithm kmp_search: let $nP \leftarrow 0$ while $j < length(S)$ do if $W[k] = S[j]$ then
	OR Write an algorithm for Rabin Carp string matching algorithm with example. function RabinKarp(string s[1n], string pattern[1m]) hpattern := hash(pattern[1m]); for i from 1 to n-m+1

	T			
	hs := hash(s[ii+m-1])			
	if hs = hpattern	allar)		
	if $s[ii+m-1] = patter$	n[1m]		
	return i	and maybe		
	return not found	Mark XI		
Q. 5 a)	Derive the Best, Worst an sorting technique.	d Average time complexities of Quick	5	СО
	Worst-case performance	O(n2)		
	Best-case performance	O(n log n) (simple partition) or O(n) (three-equal keys)		
	Average performance	O(n log n)		
		OR		
	Write an algorithm for Ve VERTEX-COVER(G):	rtex cover Problem.APPROXIMATION-		
	C = Ø			
	E'= G.E			
	while E'≠ Ø:			
	let (u, v) be an a	rbitrary edge of E'		
	$C = C \cup \{u, v\}$			
	remove from E' every edge incident on either u			
	or v			
	return C			
)	Discuss Dynamic program	ming and write an algorithm for Octional		
	Discuss Dynamic programming and write an algorithm for Optimal binary search tree.		5	CO2
)	Calculate the time complexity for N-Queens problem.		5	C03
	N- Queens : O(N!)		3	003
)	For following deterministic	finite automaton obtain 5 –tuple DFA	5	C04
	generate 3 strings which are	e accepted by this DFA	5	004
	A B			