Tarski Fixed Point Computation and the Arrival Problem

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Overview

The integer lattice, monotone functions, Tarski's theorem

Definition (Bounded Integer Lattice)

The *bounded d-dimensional integer lattice* $[N]^d = \{1, ..., N\}^d$ is equipped with a lattice ordering where for $(x_1, ..., x_d), (x'_1, ..., x'_d) \in [N]^d$, $(x_1, ..., x_d) \leq (x'_1, ..., x'_d)$ if $x_i \leq x'_i$ for each $i \in [d]$.

Definition (Monotone function)

A function $f : [N]^d \to [N]^d$ is monotone if whenever $x, x' \in [N]^d$ with $x \leq x'$, $f(x) \leq f(x')$.

Theorem (Tarski)

If $f : [N]^d \to [N]^d$ *is monotone, then there is a point* $x \in [N]^d$ *with* $f(x) = x$.

The problem, Lower Bounds

Definition (TARSKI)

The problem $\text{TARSKI}(N, d)$ is, given oracle access to a monotone function $f : [N]^d \to [N]^d$, find a point $x \in [N]^d$ such that $f(x) = x$.

Theorem (Etessami, Papadimitriou, Rubinstein, Yannakis)

The query complexity of $\text{TARSKI}(N, 2)$ is $\Theta(\log^2 N)$.

Corollary

The query complexity of $\text{TARSKI}(N, d)$ is $\Omega(\log^2 N)$ for $d \geq 2$.
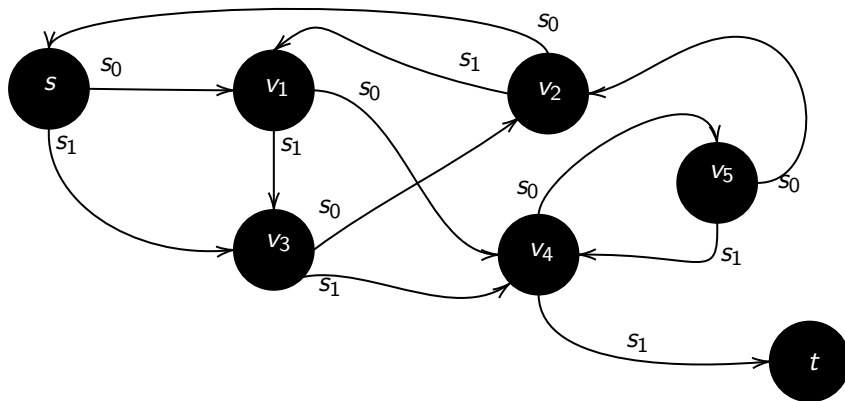
Known Upper Bounds

Theorem (Fearnley, Pálvölgyi, Savani)

The query complexity of $\textsc{Tarski}$ $(N, 3)$ *is* $\Theta(\log^2 N)$.

Theorem (Chen, Li)

The query complexity of $\textsc{Tarski}$ (N, d) *is* $O(\log^{\lceil (k+1)/2 \rceil} N)$.

Reduction to TARSKI

Theorem (Gärtner, Haslebacher, Hoang)

ARRIVAL *is polynomial-time reducible to* TARSKI.

Progress and Next Steps

- Progress
  - Explored solving the 4-dimensional TARSKI problem,
  - Implemented Fearnley, Pálvölgyi, Savani algorithm,
  - Applied the algorithm to randomly generated instances of ARRIVAL.
- Next Steps
  - Explore solving the 4-dimensional TARSKI problem in the special case of monotone functions corresponding to arrival instances.

Experimental Data

Average queries for random arrival instance (n=20)

Average time for random arrival instance (n=20)