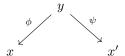
# 1 Segal spaces of spans

### 1.1 Basic definitions

We recall the basic definitions of Segal spaces of spans. Note that spans as we will define them form not only a Segal space, but a *complete* Segal space; this will not concern us. For more information, we direct the reader to [1]. This section is intended to be a summary of the results found there which we will need.

The objects of our study will be  $\infty$ -categories whose morphisms are spans



in some quasicategory  $\mathcal{C}$  (this means that x, x', and y are objects in  $\mathcal{C}$ , and that  $\phi$  and  $\psi$  are morphisms in  $\mathcal{C}$ ). We will want to be able to place certain conditions on the legs  $\phi$  and  $\psi$  of our spans. For example, we may want to restrict our attention to spans such that  $\phi$  is an equivalence, or has contractible fibers. More precisely, we pick out two subcategories of  $\mathcal{C}$  to which the respective legs of our spans must belong.

**Definition 1.1.1.** A *triple* of categories is a triple  $(\mathfrak{C}, \mathfrak{C}_{\dagger}, \mathfrak{C}^{\dagger})$ , where  $\mathfrak{C}$  is a quasicategory where  $\mathfrak{C}_{\dagger}$  and  $\mathfrak{C}^{\dagger}$  are subcategories (in the sense of [2, Sec. 1.2.11]), each of which contain all equivalences.

We will use the following terminology and notation for the morphsism in our subcategories.

- We denote the morphisms in  $C_{\dagger}$  with tails (as in  $x \mapsto y$ ), and call them ingressive.
- We denote the morphisms in  $C^{\dagger}$  with two heads (as in  $x \rightarrow y$ ), and call them *egressive*.

In order for a triple of categories to be able to support a category of spans, it will have to satisfy certain properties. We will call such categories *adequate*.

**Definition 1.1.2.** A triple  $(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger})$  is said to be **adequate** if it has the following properties:

1. For any morphism  $y \mapsto x \in \mathcal{C}_{\dagger}$  and any morphism  $x' \twoheadrightarrow x \in \mathcal{C}^{\dagger}$ , there exists a pullback square



### 2. For any pullback square

$$y' \xrightarrow{f'} x' \\ \downarrow \qquad \downarrow \\ y \xrightarrow{f} x$$

if the arrow f belongs to  $\mathcal{C}_{\dagger}$  (resp.  $\mathcal{C}^{\dagger}$ ), then the arrow f' belongs to  $\mathcal{C}_{\dagger}$  (resp.  $\mathcal{C}^{\dagger}$ ).

We will call a square of the form

$$\begin{array}{ccc}
y' & \longrightarrow & x' \\
\downarrow & & \downarrow \\
y & \longmapsto & x
\end{array}$$

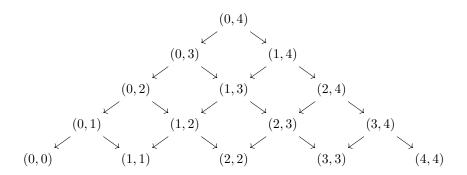
ambigressive. If such an ambigressive square is also a pullback square, we will call it ambigressive pullback.

It is now time to set about building our categories of spans. We give only a sketch. For each  $n \geq 0$ , we define a poset  $\Sigma_n$  as follows:

- The elements of  $\Sigma_n$  are pairs of integers (i,j), with  $0 \le i \le j \le n$ .
- We define a partial order via

$$(i,j) \le (i',j') \iff i \le i' \le j' \le j.$$

**Example 1.1.3.** We can draw  $\Sigma_4$  as follows.



These  $\Sigma_n$  assemble into a cosimplicial object

$$\Delta \to \operatorname{Set}_{\Delta}; \qquad [n] \mapsto N(\Sigma_n),$$

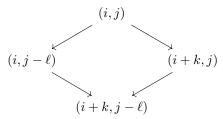
Left Kan extending yields a functor sd:  $\operatorname{Set}_{\Delta} \to \operatorname{Set}_{\Delta}$ . General abstract nonsense gives us right adjoint  $\operatorname{Span}' : \operatorname{Set}_{\Delta} \to \operatorname{Set}_{\Delta}$  forming an adjunction

$$\operatorname{sd}:\operatorname{Set}_\Delta \longleftrightarrow \operatorname{Set}_\Delta:\operatorname{Span}'.$$

For any simplicial set A, the simplicial set  $\operatorname{Span}'(A)$  has n-simplices

$$\operatorname{Span}'(A)_n = \operatorname{Hom}_{\operatorname{Set}_{\Delta}}(\operatorname{sd}(\Delta^n), A).$$

Fix some triple  $(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger})$ . We will call a functor  $\operatorname{sd}(\Delta^n) \to \mathcal{C}$  is called *ambigressive Cartesian* if the image of each square in  $\operatorname{sd}(\Delta^n) = N(\Sigma_n)$  of the form



(where we include the possibilities k=0 and  $\ell=0$ ) is an ambigressive pullback. The quasicategory Span( $\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}$ ) is then defined level-wise to be the subset

$$\operatorname{Span}(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \subseteq \operatorname{Span}'(\mathcal{C})$$

on functors  $\operatorname{sd}(\Delta^n) \to \mathcal{C}$  which are ambigressive Cartesian.

Analogously, for an adequate triple  $(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger})$  and any simplicial set A, define

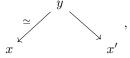
$$\operatorname{Map}^{\operatorname{aCart}}(\operatorname{sd}(A), \mathfrak{C})^{\simeq} \subseteq \operatorname{Map}(\operatorname{sd}(A), \mathfrak{C})^{\simeq}$$

to be the full simplicial subset on functors  $\operatorname{sd}(A)$  such that for each standard simplex  $\sigma$  of A, the image of  $\operatorname{sd}(\sigma)$  in  $\mathcal{C}$  is ambigressive Cartesian. The (complete) Segal space of spans  $\operatorname{\mathbf{Span}}(\mathcal{C},\mathcal{C}_{\dagger},\mathcal{C}^{\dagger})$  is defined level-wise by

$$\mathbf{Span}(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger})_n = \mathrm{Map}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^n), \mathcal{C})^{\simeq}.$$

For a proof that this really is a Segal space, the reader is once again referred to [1].

**Example 1.1.4.** For any quasicategory, we can define a triple  $(\mathfrak{C}, \mathfrak{C}_{\dagger} = \mathfrak{C}, \mathfrak{C}^{\dagger} = \mathfrak{C}^{\simeq})$ . The corresponding spans are of the form



where  $Y \to X$  is an equivalence in  $\mathcal{C}$ .

The resulting quasicategory of spans is categorically equivalent to the category  $\mathcal{C}$  from which we started. To see this, note that for each  $n \geq 0$  there is a retraction of posets

$$[n] \hookrightarrow \Sigma_n \twoheadrightarrow [n],$$

where the first map takes  $i \mapsto (i, n)$ , and the second takes  $(i, j) \mapsto i$ . Under the nerve, this corresponds to a retraction of simplicial sets

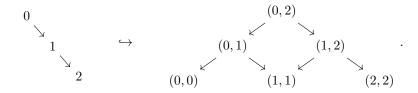
$$\Delta^n \hookrightarrow \operatorname{sd}(\Delta^n) \twoheadrightarrow \Delta^n$$
.

Pulling back along these maps gives us, for each n, a retraction

$$\operatorname{Fun}(\Delta^n, \mathfrak{C}) \stackrel{i_n}{\longrightarrow} \operatorname{Fun}(\operatorname{sd}(\Delta^n), \mathfrak{C}) \stackrel{r_n}{\smile} \operatorname{Fun}(\Delta^n, \mathfrak{C}) \ .$$

Note that the image of the inclusion i

It is not hard to convince oneself that a map  $\operatorname{sd}(\Delta^n) \to \mathcal{C}$  is ambigressive cartesian if and only if it is a right Kan extension of its restriction along  $\Delta^n \hookrightarrow \operatorname{sd}(\Delta^n)$ . In the case n=2, the inclusion  $[2] \hookrightarrow \Sigma_2$  is of the form



The right Kan extension of a diagram  $\Delta^2 \to \mathcal{C}$ 

The inclusions  $\Delta^n \to \operatorname{sd}(\Delta^n)$  does not yield a map of cosimplicial objects  $\Delta^{\bullet} \to \operatorname{sd}(\Delta^{\bullet})$ , but the projections  $\operatorname{sd}(\Delta^n) \to \Delta^n$  do, giving us a map of cosimplicial objects  $\operatorname{sd}(\Delta^{\bullet}) \to \Delta^{\bullet}$ .

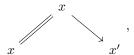
There is an equivalence of Segal spaces

$$\Gamma(\mathcal{C}) \to \mathbf{Span}(\mathcal{C}, \mathcal{C}, \mathcal{C}^{\simeq}),$$

where  $\Gamma(\mathcal{C})$  is the Segal space defined level-wise by

$$\Gamma(\mathcal{C})_n = \operatorname{Fun}(\Delta^n, \mathcal{C}).$$

This sends a morphism  $f : x \to y$  in  $\mathcal{C}$  to the span



A functor between adequate triples  $(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$  is a functor

$$p \colon \mathcal{C} \to \mathcal{D}$$

such that  $p(\mathcal{C}_{\dagger}) \subseteq \mathcal{D}_{\dagger}$ ,  $p(\mathcal{C}^{\dagger}) \subseteq \mathcal{D}^{\dagger}$ , and p preserves ambigressive pullbacks. We

will denote the restriction of p to  $\mathcal{C}_{\dagger}$  by

$$p_{\dagger} \colon \mathcal{C}_{\dagger} \to \mathcal{D}_{\dagger},$$

and similarly for  $p^{\dagger}$ . Clearly, a functor p between adequate triples gives functors  $\mathrm{Span}(p)$  and  $\mathrm{Span}(p)$  respectively between quasicategories and Segal spaces of spans.

The goal of the remainder of this document is to provide a proof of the following theorem. This is [1, Thm. 12.2].

**Theorem 1.1.5.** Let  $p: (\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$  be a functor of adequate triples. Further suppose that p satisfies the following conditions:

- 1. For any ingressive morphism  $g \colon s \to t$  of  $\mathcal{D}$  and any object  $x \in \mathcal{C}_s$ , there exists an ingressive morphism  $f \colon x \to y$  covering g which is both p-cocartesian and  $p_{\dagger}$ -cocartesian.
- 2. Suppose  $\sigma$  is a commutative square

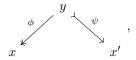
$$y' \xrightarrow{f'} x'$$

$$\downarrow \qquad \qquad \downarrow \psi$$

$$y \xrightarrow{f} x$$

in  $\mathcal{C}$  such that  $p(\sigma)$  is ambigressive pullback in  $\mathcal{D}$ , with in- and egressive morphisms as marked, and f is p-cocartesian. Then f' is p-cocartesian if and only if  $\sigma$  is an ambigressive pullback square (and in particular  $\psi$  is egressive).

Then if a morphism in  $\mathrm{Span}(\mathfrak{C},\mathfrak{C}_\dagger,\mathfrak{C}^\dagger)$  is of the form



where  $\phi$  is egressive and  $p^{\dagger}$ -cartesian, and  $\psi$  is ingressive and p-cocartesian, it is  $\pi$ -cocartesian.<sup>1</sup>

This theorem originally appears in [1] (although, as noted above, the assumptions are slightly different). However, the proof there is combinatorial in nature. Our proof uses the material developed in Section ??, leveraging the fact that  $\infty$ -categories of spans have natural incarnations as (complete) Segal spaces.

Before we get to the meat of the proof, however, we have to prove some preliminary results.

<sup>&</sup>lt;sup>1</sup>This is not precisely the same condition that Barwick gives: he demands that  $\phi$  be p-cartesian rather than  $p^{\dagger}$ -cartesian. As far as I can tell, this is simply an error; Barwick seems to use the condition as stated above in his proof.

# 1.2 Housekeeping

In this section we recall some basic facts we will need in the course of our proof of the main theorem.

#### **Isofibrations**

We will need some results about isofibrations, sometimes called categorical fibrations. These are standard, and proofs can be found on Kerodon.

**Proposition 1.2.1.** Let  $f: \mathcal{C} \to \mathcal{D}$  be a map between quasicategories. The following are equivalent:

- The map f is an isofibration.
- The map f is an inner fibration, and each equivalence in  $\mathcal{D}$  has a lift in  $\mathcal{C}$  which is an equivalence.

We will denote by Map(-, -) the internal hom on  $Set_{\Delta}$ .

**Proposition 1.2.2.** Let  $f: \mathcal{C} \to \mathcal{D}$  be an isofibration between quasicategories, and let  $i: A \hookrightarrow B$  be an inclusion of simplicial sets. Then the map

$$\operatorname{Map}(B, \mathcal{C}) \to \operatorname{Map}(A, \mathcal{C}) \times_{\operatorname{Map}(A, \mathcal{D})} \operatorname{Map}(B, \mathcal{D})$$

is an isofibration.

We will denote by  $(-)^{\sim}$  the *core* functor, i.e. the functor which takes a simplicial set X to the largest Kan complex it contains.

**Proposition 1.2.3.** Let  $f: X \to Y$  be an isofibration of simplicial sets. Then

$$f^{\simeq} \colon X^{\simeq} \to Y^{\simeq}$$

is a Kan fibration.

We will denote the core of the mapping space functor Map(-,-) by Fun(-,-). That is, for simplicial sets A and B, we have

$$\operatorname{Fun}(A,B) \cong \operatorname{Map}(A,B)^{\sim}.$$

**Corollary 1.2.4.** Let  $f: \mathcal{C} \to \mathcal{D}$  be an isofibration between quasicategories, and let  $A \hookrightarrow B$  be an inclusion of simplicial sets. Then the map

$$\operatorname{Fun}(B,{\mathfrak C}) \to \operatorname{Fun}(A,{\mathfrak C}) \times_{\operatorname{Fun}(A,{\mathfrak D})} \operatorname{Fun}(B,{\mathfrak D})$$

is a Kan fibration.

#### Connected components

Our proof will rely heavily on the fact that the many properties that morphisms can have (ingressive, egressive, p-cocartesian, etc.) are well-behaved with respect to the homotopical structure of the complete Segal spaces in which they live; for example, if a morphism is p-cocartesian, then every morphism in its path component is p-cocartesian.

Let X' and X be simplicial sets, and let  $f: X' \hookrightarrow X$  be an inclusion. We will say that f is an *inclusion of connected components* if for all simplices  $\sigma \in X$ , if  $\sigma$  has any vertex in common with X', then  $\sigma$  is wholly contained in X'.

**Lemma 1.2.5.** Let A be a simplicial set such that between any two vertices x and y of A there exists a finite zig-zag of 1-simplices of A connecting x to y. Let  $A_0 \subseteq A$  be a nonempty simplicial subset of A. Let  $f: X' \hookrightarrow X$  be a morphism between simplicial sets which is an inclusion of connected components. Then the following dashed lift always exists.

$$\begin{array}{ccc}
A_0 & \xrightarrow{f_0} & X' \\
\downarrow & & \downarrow \\
A & \xrightarrow{f} & X
\end{array}$$

Proof. We have a map  $f: A \to X$ ; in order to construct our dashed lift, it suffices to show that under the above assumptions, f takes every simplex of A to a simplex which belongs to X'. To this end, let  $\sigma \in A$  be a simplex. There exists a finite zig-zag of 1-simplices connecting some vertex of  $\sigma$  to a vertex  $\alpha$  of  $A_0$ ; under f, this is mapped to a zig-zag of 1-simplices connecting  $f(\alpha)$  to  $f(\sigma)$ . The vertex  $f(\alpha)$  belongs to X'; proceeding inductively, we find that the image of every 1-simplex belonging to the zig-zag also belongs to X', and that  $f(\sigma)$  therefore also belongs to X'.

Corollary 1.2.6. Given any commuting square of simplicial sets

$$X' \stackrel{i}{\longleftarrow} X$$

$$f' \downarrow \qquad \qquad \downarrow f$$

$$Y' \stackrel{f}{\longleftrightarrow} Y$$

with monomorphisms as marked, where i is an inclusion of connected components between Kan complexes and f is a Kan fibration, the map f' is a Kan fibration.

Furthermore, if f is a trivial fibration and f' is surjective on vertices, then f' is a trivial fibration.

*Proof.* First, we show that if f is a Kan fibration, then f' is a Kan fibration. We

need to show that a dashed lift below exists for each  $n \ge 1$  and each  $0 \le i \le n$ .

We can always solve our outer lifting problem, and by Lemma 1.2.5, our lift factors through X'.

Now suppose that f is a trivial Kan fibration. The logic above says that f' has the right lifting property with respect to  $\partial \Delta^n \hookrightarrow \Delta^n$  for  $n \geq 1$ ; the right lifting property with respect to  $\partial \Delta^0 \hookrightarrow \Delta^0$  is equivalent to the surjectivity of f' on vertices.

### Contractibility

We will need at certain points to show that various spaces of lifts are contractible. These are mostly common-sense results, but it will be helpful to have them written down somewhere so we can refer to them later.

**Lemma 1.2.7.** Let  $f: \mathcal{C} \to \mathcal{D}$  be an inner fibration between quasicategories. Then for any commuting square

$$\Lambda_1^2 \xrightarrow{\alpha} \mathcal{C} 
\downarrow_f, 
\Delta^2 \xrightarrow{\beta} \mathcal{D}$$

the fiber F in the pullback square below is contractible.

*Proof.* The right-hand map is a trivial Kan fibration.

One should interpret this as telling us that given a  $\Lambda_2^2$ -horn  $\alpha$  in  $\mathcal{C}$  lying over a 2-simplex  $\beta$  in  $\mathcal{D}$ , the space of fillings of  $\alpha$  lying over  $\beta$  is contractible. We will need a pantheon of similar results.

**Lemma 1.2.8.** Let  $f: \mathcal{C} \to \mathcal{D}$  be an inner fibration between quasicategories. Let e be any f-cartesian edge. Then for any square

$$\begin{array}{ccc} \Lambda_2^2 & \stackrel{\alpha}{\longrightarrow} & \mathcal{C} \\ \downarrow & & \downarrow^f \\ \Delta^2 & \stackrel{\beta}{\longrightarrow} & \mathcal{D} \end{array}$$

such that  $\alpha|_{\{1,2\}} = e$ , the fiber F in the pullback diagram

$$F \xrightarrow{Fun(\Delta^{2}, \mathcal{C})} Fun(\Delta^{2}, \mathcal{C}) \downarrow \qquad \qquad (1)$$

$$\Delta^{0} \xrightarrow{(\alpha, f\alpha, \beta)} Fun(\Lambda^{2}_{2}, \mathcal{C}) \times_{Fun(\Lambda^{2}_{1}, \mathcal{D})} Fun(\Delta^{2}, \mathcal{D})$$

is contractible.

*Proof.* Denote  $f(e) = \bar{e}$ . Define a marking  $\mathcal{E}'$  on  $\mathcal{D}$  containing all degenerate edges and  $\bar{e}$ . Define a marking  $\mathcal{E}$  on  $\mathcal{C}$  containing the f-cocartesian lifts of the edges in  $\mathcal{E}$ . By [2, Prop. 3.1.1.6], the map  $\mathcal{C}^{\mathcal{E}} \to \mathcal{D}^{\mathcal{E}'}$  has the right-lifting property with respect to all cocartesian-marked anodyne morphisms. It follows that

$$\begin{split} \operatorname{Map}^{\sharp}((\Delta^{2})^{\mathcal{L}}, \mathcal{C}^{\mathcal{E}}) \\ \downarrow \\ \operatorname{Map}^{\sharp}((\Lambda_{2}^{2})^{\mathcal{L}}, \mathcal{C}^{\mathcal{E}}) \times_{\operatorname{Map}^{\sharp}((\Lambda_{1}^{2})^{\mathcal{L}}, \mathcal{D}^{\mathcal{E}'})} \operatorname{Map}^{\sharp}((\Delta^{2})^{\mathcal{L}}, \mathcal{D}^{\mathcal{E}'}) \end{split}$$

is a trivial Kan fibration; the map on the right-hand side of Diagram 1 is the core of this map, and is thus also a trivial Kan fibration.  $\Box$ 

**Lemma 1.2.9.** Let  $(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$  be an inner fibration between triples satisfying the conditions of Theorem 1.1.5. Denote by Fun' $(\Delta^1, \mathcal{C}^{\dagger})$  the full simplicial subset of Fun $(\Delta^1, \mathcal{C}^{\dagger})$  on edges  $\Delta^1 \to \mathcal{C}^{\dagger}$  which are both p-cocartesian and  $p^{\dagger}$ -cocartesian. Then the fibers of the map

$$\operatorname{Fun}'(\Delta^1, \mathcal{C}^\dagger) \to \operatorname{Fun}(\Delta^{\{0\}}, \mathcal{C}^\dagger) \times_{\operatorname{Fun}(\Delta^{\{0\}}, \mathcal{D}^\dagger)} \operatorname{Fun}(\Delta^1, \mathcal{D}^\dagger)$$

are contractible.

*Proof.* Denote by Fun"  $(\Delta^1, \mathcal{C}^{\dagger})$  the full simplicial subset of Fun $(\Delta^1, \mathcal{C}^{\dagger})$  on edges which are  $p^{\dagger}$ -cartesian. Then we have a square

$$\begin{split} \operatorname{Fun}'(\Delta^1, \operatorname{\mathcal{C}}^\dagger) & & \longrightarrow \operatorname{Fun}''(\Delta^1, \operatorname{\mathcal{C}}^\dagger) \\ & & \downarrow & & \downarrow \\ \operatorname{Fun}(\Delta^{\{0\}}, \operatorname{\mathcal{C}}^\dagger) \times_{\operatorname{Fun}(\Delta^{\{0\}}, \operatorname{\mathcal{D}}^\dagger)} \operatorname{Fun}(\Delta^1, \operatorname{\mathcal{D}}^\dagger) & = = \operatorname{Fun}(\Delta^{\{0\}}, \operatorname{\mathcal{C}}^\dagger) \times_{\operatorname{Fun}(\Delta^{\{0\}}, \operatorname{\mathcal{D}}^\dagger)} \operatorname{Fun}(\Delta^1, \operatorname{\mathcal{D}}^\dagger) \end{split}$$

in which the top map is an inclusion of connected components, the right map is a trivial Kan fibration, and the left map is surjective on vertices. Thus, the left map is a trivial Kan fibration by Corollary 1.2.6

### 1.3 Main theorem

Our goal is now to prove Theorem 1.1.5. As in the statement of the theorem, we fix a map of triples  $p: (\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$  giving us a functor of quasicategories

$$\pi \colon \mathrm{Span}(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to \mathrm{Span}(\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger}).$$

We wish to show that morphsims of the form

$$\begin{array}{cccc}
 & y & & \\
 & & \downarrow & \\
 & x & & x'
\end{array}$$
(2)

are  $\pi$ -cocartesian, where  $\phi$  is egressive and  $p_{\dagger}$ -cartesian, and  $\psi$  is ingressive and p-cocartesian.

The way forward is clear: the functor  $\pi$  is the zeroth row of a functor of Segal spaces

$$\pi' : \mathbf{Span}(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to \mathbf{Span}(\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger}).$$

It will suffice to show by Corollary ?? that morphisms of the above form are  $\pi'$ -cocartesian. Morally, this result should not be surprising. We should think of the homotopy pullback condition defining cocartesian morphisms (Definition ??) as telling us that we can fill relative  $\Lambda_2^2$ -horns in  $\pi'$ : **Span**( $\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}$ )  $\to$  **Span**( $\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger}$ ), and that such fillings are unique up to contractible choice. Finding a filling

$$\begin{array}{ccc} \Lambda_2^2 & \longrightarrow \mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \\ & & & \downarrow^{\pi'} \\ \Delta^2 & \longrightarrow \mathbf{Span}(\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger) \end{array}$$

is equivalent finding a filling

$$sd(\Lambda_2^2) \longrightarrow \mathcal{C}$$

$$\downarrow p$$

$$sd(\Delta^2) \longrightarrow \mathcal{D}$$

(such that the filling is ambigressive and the necessary square is pullback); the conditions on p guarantee us that we can perform this filling in a series of steps pictured in Figure 1, each of which is unique up to contractible choice.

### Reedy fibration

First, however, we have to check that  $\pi'$  is a Reedy fibration. This follows from the following useful fact, implied by the first assumption of Theorem 1.1.5.

**Lemma 1.3.1.** Let  $p: (\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$  be a functor of triples satisfying

the conditions of Theorem 1.1.5. Then the map  $p: \mathcal{C} \to \mathcal{D}$  is an isofibration.

*Proof.* By assumption, p is an inner fibration and  $\mathbb{C}$  and  $\mathbb{D}$  are quasicategories, so it suffices by Proposition 1.2.1 to show that p admits lifts of equivalences. Each equivalence in  $\mathcal{D}$  is in particular ingressive, and hence admits a p-cocartesian lift by the assumptions of Theorem 1.1.5; these are automatically equivalences.

**Proposition 1.3.2.** For any functor of adequate triples  $p: (\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger}),$ the map

$$\pi' \colon \mathbf{Span}(\mathfrak{C}, \mathfrak{C}_{\dagger}, \mathfrak{C}^{\dagger}) \to \mathbf{Span}(\mathfrak{D}, \mathfrak{D}_{\dagger}, \mathfrak{D}^{\dagger})$$

is a Reedy fibration.

*Proof.* Consider the following diagram.

$$\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^n), \mathfrak{C}) \hookrightarrow \operatorname{Fun}(\operatorname{sd}(\Delta^n), \mathfrak{C}) \hookrightarrow \operatorname{Fun}(\operatorname{sd}(\Delta^n), \mathfrak{C}) \hookrightarrow \operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\partial \Delta^n), \mathfrak{C}) \times_{\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\partial \Delta^n), \mathfrak{D})} \operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^n), \mathfrak{D}) \hookrightarrow \operatorname{Fun}(\operatorname{sd}(\partial \Delta^n), \mathfrak{C}) \times_{\operatorname{Fun}(\operatorname{sd}(\partial \Delta^n), \mathfrak{D})} \operatorname{Fun}(\operatorname{sd}(\Delta^n), \mathfrak{D})$$

The right-hand map is a Kan fibration by Corollary 1.2.4 (because sd preserves monomorphisms), and the top map is an inclusion of connected components, so Corollary 1.2.6 implies that the left-hand map is a Kan fibration, which is what we needed to show. 

We are now ready to begin in earnest our proof that morphisms of the form Equation 2 are  $\pi'$ -cocartesian. We first introduce some notation. Let  $p: (\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \to$  $(\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$  be a functor of triples.

• We will denote p-cartesian morphisms in  $\mathcal{C}$  with a circle:

$$x \longrightarrow y$$

• We will denote p-cocartesian morphisms in  $\mathcal{C}$  with a bullet:

$$x \longrightarrow y$$

• We will denote  $p^{\dagger}$ -cartesian morphisms in  $\mathcal{C}$  with a triangle:

$$x \longrightarrow y$$

• We will denote  $p_{\dagger}$ -cocartesian morphisms in  $\mathcal C$  with a filled triangle:

$$x \longrightarrow y$$

Thus, an ingressive morphism which is both p-cocartesian and  $p_{\dagger}$ -cocartesian will be denoted by

$$x \longrightarrow \longrightarrow y$$

Our proof will rest on the factorization  $sd(\Lambda_0^2) \hookrightarrow sd(\Delta^2)$  pictured in Figure 1. Denote the underlying factorization of simplicial sets by

$$A_1 \stackrel{i_1}{\hookrightarrow} A_2 \stackrel{i_2}{\hookrightarrow} A_3 \stackrel{i_3}{\hookrightarrow} A_4 \stackrel{i_4}{\hookrightarrow} A_5 \stackrel{i_5}{\hookrightarrow} A_6$$

For each of the  $A_i$  above, denote by

$$\operatorname{Fun}'(A_i, \mathfrak{C}) \subseteq \operatorname{Fun}(A_i, \mathfrak{C}) \tag{3}$$

the full simplicial subset on those functors  $A_i \to \mathcal{C}$  which respect each labelling in Figure 1:

- ingressive
- egressive
- $\bullet$  *p*-cartesian
- p-cocartesian
- $p^{\dagger}$ -cartesian
- $p_{\dagger}$ -cocartesian

Note that because equivalences in C belong to each of these classes of morphisms, the inclusion in Equation 3 is an inclusion of connected components.

Similarly, we will denote by  $\operatorname{Fun}'(A_i, \mathcal{D})$  the full simplicial subset on functors which respect the ingressive and egressive labellings.

## Lemma 1.3.3. The square

$$\operatorname{Fun}'(A_6, \mathfrak{C}) \longrightarrow \operatorname{Fun}'(A_0, \mathfrak{C})$$

$$\downarrow \qquad \qquad \downarrow$$

$$\operatorname{Fun}'(A_6, \mathfrak{D}) \longrightarrow \operatorname{Fun}'(A_0, \mathfrak{D})$$

is homotopy pullback.

Proof. The factorization of Figure 1 gives us a factorization of the above square into five squares

$$\operatorname{Fun}'(A_{k+1}, \mathbb{C}) \longrightarrow \operatorname{Fun}'(A_k, \mathbb{C})$$

$$\downarrow \qquad \qquad \downarrow \qquad , \qquad 1 \le k < 6,$$

$$\operatorname{Fun}'(A_{k+1}, \mathbb{D}) \longrightarrow \operatorname{Fun}'(A_k, \mathbb{D})$$

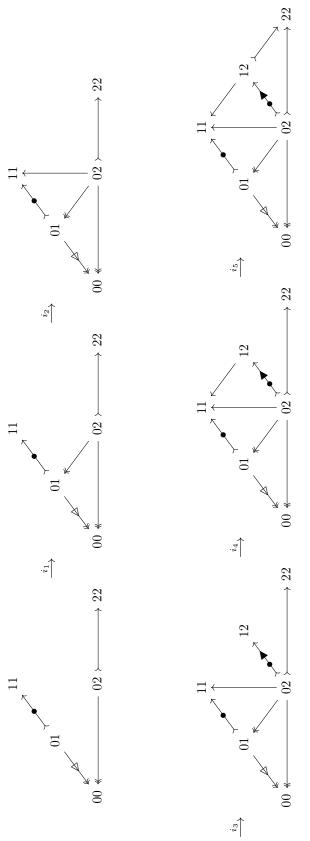


Figure 1: A factorization of the inclusion  $\operatorname{sd}(\Lambda_0^2) \hookrightarrow \operatorname{sd}(\Delta^2)$ , where certain morphisms have been labelled.

each corresponding to one of the inclusions  $i_k$ . We will be done if we can show that each of these squares is homotopy pullback. It suffices to show that for each k, the map

$$j_k \colon \operatorname{Fun}'(A_{k+1}, \mathbb{C}) \to \operatorname{Fun}'(A_k, \mathbb{C}) \times_{\operatorname{Fun}'(A_{k+1}, \mathbb{D})} \operatorname{Fun}'(A_k, \mathbb{D})$$

is a weak equivalence. First, we show that each of these maps is a Kan fibration. To see this, consider the square

$$\operatorname{Fun}'(A_{k+1}, \mathfrak{C}) \hookrightarrow \operatorname{Fun}(A_{k+1}, \mathfrak{C})$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \downarrow$$

$$\operatorname{Fun}'(A_k, \mathfrak{C}) \times_{\operatorname{Fun}'(A_{k+1}, \mathfrak{D})} \operatorname{Fun}'(A_k, \mathfrak{D}) \hookrightarrow \operatorname{Fun}(A_k, \mathfrak{C}) \times_{\operatorname{Fun}(A_{k+1}, \mathfrak{D})} \operatorname{Fun}(A_k, \mathfrak{D})$$

The right-hand map is a Kan fibration because of Corollary 1.2.4, and we have already seen that the top map is an inclusion of connected componenents. Hence each  $j_k$  is a Kan fibration by Corollary 1.2.6. Thus, in order to show that each  $j_k$  is a weak equivalence, it suffices to show that the fibers are contractible.

First consider the case k=1. For any map  $\alpha$  below, consider the following pullback square.

$$F_{\alpha} \longrightarrow \operatorname{Fun}'(A_{1}, \mathfrak{C})$$

$$\downarrow \qquad \qquad \downarrow_{j_{1}} \qquad .$$

$$\Delta^{0} \stackrel{\alpha}{\longrightarrow} \operatorname{Fun}'(A_{0}, \mathfrak{C}) \times_{\operatorname{Fun}'(A_{1}, \mathfrak{D})} \operatorname{Fun}'(A_{1}, \mathfrak{D})$$

The fiber  $F_{\alpha}$  over  $\alpha$  is the space of ways of completing a diagram of shape  $A_0$  in  $\mathcal{C}$  to a diagram of shape  $A_1$  in  $\mathcal{C}$  given a diagram of shape  $A_1$  in  $\mathcal{D}$ . This is the space of ways of filling  $\Lambda_2^2 \hookrightarrow \Delta^2$  in  $\mathcal{C}^{\dagger}$  lying over a 2-simplex in  $\mathcal{D}^{\dagger}$ . This is contractible by Lemma 1.2.8.

The case k = 2 is similar, using Lemma 1.2.7.

The case k = 3 is similar, using Lemma 1.2.9.

The cases 
$$k = 4$$
 and  $k = 5$  use the dual to Lemma 1.2.8.

We are now ready to show that morphisms of the form Diagram 2 are  $\pi'$ -cocartesian.

**Proposition 1.3.4.** For any morphism e of the form given in Equation 2, the square

$$\begin{split} \operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^2), \operatorname{\mathcal{C}}) \times_{\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^{\{0,1\}}), \operatorname{\mathcal{C}})} \{e\} & \longrightarrow \operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Lambda_0^2), \operatorname{\mathcal{C}}) \times_{\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^{\{0,1\}}), \operatorname{\mathcal{C}})} \{e\} \\ & \downarrow & \downarrow \\ \operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^2), \operatorname{\mathcal{D}}) \times_{\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^{\{0,1\}}), \operatorname{\mathcal{D}})} \{\pi e\} & \longrightarrow \operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Lambda_0^2), \operatorname{\mathcal{D}}) \times_{\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^{\{0,1\}}), \operatorname{\mathcal{D}})} \{\pi e\} \end{split}$$

is homotopy pullback.

*Proof.* This square factors into the two squares

$$\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^2), \mathfrak{C}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathfrak{C})} \{e\} \longrightarrow \operatorname{Fun}'(A_6, \mathfrak{C}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathfrak{C})} \{e\}$$

$$\downarrow \qquad \qquad \downarrow$$

$$\operatorname{Fun}^{\operatorname{aCart}}(\operatorname{sd}(\Delta^2), \mathfrak{D}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathfrak{D})} \{\pi e\} \to \operatorname{Fun}'(A_6, \mathfrak{D}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathfrak{C})} \{\pi e\}$$

and

$$\operatorname{Fun}'(A_6, \mathbb{C}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathbb{C})} \{e\} \longrightarrow \operatorname{Fun}'(A_0, \mathbb{C}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathbb{C})} \{e\}$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \cdot$$

$$\operatorname{Fun}'(A_6, \mathbb{D}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathbb{C})} \{\pi e\} \to \operatorname{Fun}'(A_0, \mathbb{D}) \times_{\operatorname{Fun}(\operatorname{sd}(\Delta^{\{0,1\}}), \mathbb{C})} \{\pi e\}$$

The first is homotopy pullback because the bottom map is an inclusion of connected components, and the second condition of Theorem 1.1.5 guarantees that the fiber over an ambigressive  $\operatorname{sd}(\Delta^2) \to \mathcal{D}$  belonging to  $\operatorname{Fun}'(A_6, \mathcal{D})$  whose restriction to  $\operatorname{sd}(\Delta^{\{0,1\}})$  is  $\pi e$  is precisely an ambigressive Cartesian functor  $\operatorname{sd}(\Delta^2) \to \mathcal{C}$  whose restriction to  $\operatorname{sd}(\Delta^{\{0,1\}})$  is e.

That the second is homotopy pullback follows immediately from Lemma 1.3.3.

This proves Theorem 1.1.5.