

1 Segal spaces of spans

1.1 Basic definitions

We recall the basic definitions of Segal spaces of spans. Note that spans as we will define them form not only a Segal space, but a *complete* Segal space; for the most part, this will not concern us. For more information, we direct the reader to [2]. This section is intended to be a summary of the results found there which we will need.

The objects of our study will be ∞ -categories whose morphisms are spans in some quasicategory \mathcal{C} , i.e. diagrams in \mathcal{C} of the form

$$\begin{array}{ccc} & y & \\ \phi \swarrow & & \searrow \psi \\ x & & x' \end{array} .$$

We will want to be able to place certain conditions on the legs ϕ and ψ of our spans. For example, we may want to restrict our attention to spans such that ϕ is an equivalence. More precisely, we pick out two subcategories of \mathcal{C} to which the respective legs of our spans must belong.

Definition 1.1.1. A *triple* of categories is a triple $(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$, where \mathcal{C} is a quasicategory where \mathcal{C}_\dagger and \mathcal{C}^\dagger are subcategories (in the sense of [3, Sec. 1.2.11]), each of which contain all equivalences.

We will use the following terminology and notation for the morphism in our subcategories.

- We denote the morphisms in \mathcal{C}_\dagger with tails (as in $x \rightarrowtail y$), and call them *ingressive*.
- We denote the morphisms in \mathcal{C}^\dagger with two heads (as in $x \twoheadrightarrow y$), and call them *egressive*.

The egressive morphisms will correspond to the backwards (i.e. leftwards) facing legs of our spans, and the ingressive morphisms will correspond to the forwards (i.e. rightwards) facing legs. We provide the following diagram as a sort of cheat-sheet.

$$\begin{array}{ccc} & Y & \\ \text{egressive} \swarrow & & \searrow \text{ingressive} \\ \mathcal{C}^\dagger & & \mathcal{C}_\dagger \\ X & & X' \end{array}$$

In order for a triple of categories to be able to support a category of spans, it will have to satisfy certain properties. Following Barwick, we will call such triples *adequate*.

Definition 1.1.2. A triple $(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ is said to be *adequate* if it has the following properties:

1. For any ingressive morphism $f \in \mathcal{C}_\dagger$ and any egressive morphism $g \in \mathcal{C}^\dagger$, there exists a pullback square

$$\begin{array}{ccc} y' & \longrightarrow & x' \\ \downarrow & & \downarrow g \\ y & \xrightarrow{f} & x \end{array}.$$

2. For any pullback square

$$\begin{array}{ccc} y' & \xrightarrow{f'} & x' \\ \downarrow & & \downarrow \\ y & \xrightarrow{f} & x \end{array},$$

if the arrow f belongs to \mathcal{C}_\dagger (resp. \mathcal{C}^\dagger), then the arrow f' belongs to \mathcal{C}_\dagger (resp. \mathcal{C}^\dagger).

We will call a square of the form

$$\begin{array}{ccc} y' & \xrightarrow{\quad} & x' \\ \downarrow & & \downarrow \\ y & \xrightarrow{\quad} & x \end{array}$$

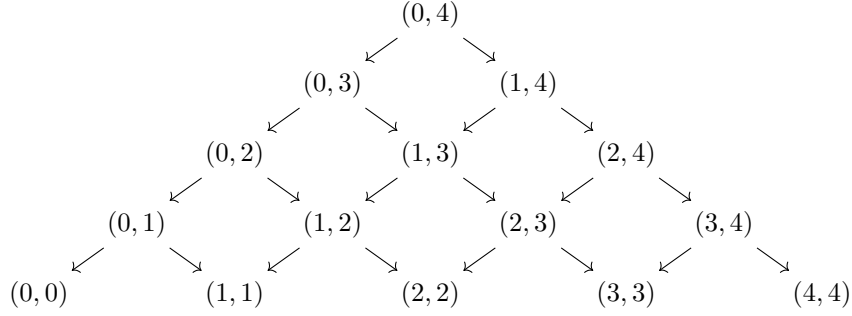
ambigressive. If such an ambigressive square is also a pullback square, we will call it *ambigressive pullback*.

It is now time to set about building our categories of spans. We give only a sketch. For each $n \geq 0$, we define a poset Σ_n as follows:

- The elements of Σ_n are pairs of integers (i, j) , with $0 \leq i \leq j \leq n$.
- For $(i, j), (i', j') \in \Sigma_n$, we define

$$(i, j) \leq (i', j') \iff i \leq i' \leq j' \leq j.$$

Example 1.1.3. We can draw Σ_4 as follows.



These Σ_n assemble into a cosimplicial object

$$\Delta \rightarrow \text{Set}_\Delta; \quad [n] \mapsto N(\Sigma_n),$$

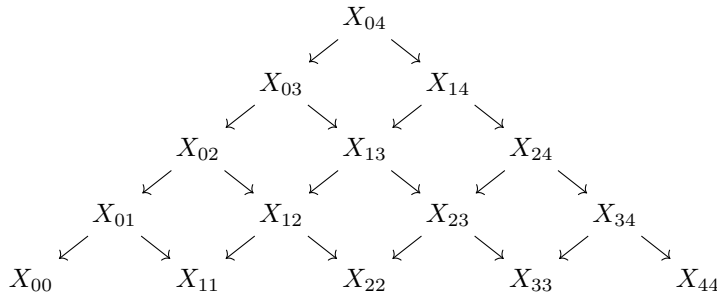
Left Kan extending along the Yoneda embedding $\Delta \hookrightarrow \text{Set}_\Delta$ yields a functor $\text{sd} : \text{Set}_\Delta \rightarrow \text{Set}_\Delta$. General abstract nonsense gives us right adjoint $\text{Span}' : \text{Set}_\Delta \rightarrow \text{Set}_\Delta$ forming an adjunction

$$\text{sd} : \text{Set}_\Delta \longleftrightarrow \text{Set}_\Delta : \text{Span}'.$$

For any simplicial set A , the simplicial set $\text{Span}'(A)$ has n -simplices

$$\text{Span}'(A)_n = \text{Hom}_{\text{Set}_\Delta}(\text{sd}(\Delta^n), A).$$

For example, for $n = 4$, such an n -simplex consists of a diagram of the form



For any n -simplex $\text{sd}(\Delta^n) \rightarrow \mathcal{C}$ corresponding to a diagram as above, we will call the morphisms $X_{ij} \rightarrow X_{i'j'}$ *backward-facing*, and the morphisms $X_{ij} \rightarrow X_{i'j}$ *forward-facing*.

Fix some adequate triple $(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$. We will call a functor $\text{sd}(\Delta^n) \rightarrow \mathcal{C}$ *ambi-*

gressively Cartesian if each square in $\text{sd}(\Delta^n) = N(\Sigma_n)$ of the form

$$\begin{array}{ccc} & (i, j) & \\ \swarrow & & \searrow \\ (i, j - \ell) & & (i + k, j) \\ \searrow & & \swarrow \\ & (i + k, j - \ell) & \end{array}$$

(where we include the possibilities $k = 0$ and $\ell = 0$) is mapped to an ambigressive pullback square,

$$\begin{array}{ccc} & X_{ij} & \\ \swarrow & & \searrow \\ X_{i(j-\ell)} & & X_{(i+k)(j)} \\ \searrow & & \swarrow \\ & X_{(i+k)(j-\ell)} & \end{array}$$

where the backwards-facing morphisms are egressive and the forwards-facing morphisms are ingressive.

Definition 1.1.4. We define a simplicial set $\text{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ level-wise to be the subset

$$\text{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)_n \subseteq \text{Span}'(\mathcal{C})_n$$

on functors $\text{sd}(\Delta^n) \rightarrow \mathcal{C}$ which are ambigressive Cartesian.

While it is not hard to check that the face and degeneracy maps on $\text{Span}'(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ restrict to $\text{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ (and thus that the above construction really produces a simplicial set), solving the lifting problems necessary to prove directly that $\text{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ is a quasicategory turns out to be combinatorially strenuous. It turns out to be easier to switch to a new model of quasicategories, (complete) Segal spaces, which give easier access to homotopical data.

To that end, for an adequate triple $(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ and any simplicial set A , we make the following definition.

Definition 1.1.5. We define a simplicial set $\text{Map}^{\text{aCart}}(\text{sd}(A), \mathcal{C})^\simeq$ to be the full simplicial subset

$$\text{Map}^{\text{aCart}}(\text{sd}(A), \mathcal{C})^\simeq \subseteq \text{Map}(\text{sd}(A), \mathcal{C})^\simeq$$

on functors $\text{sd}(A) \rightarrow \mathcal{C}$ such that for each standard simplex σ of A , the image of $\text{sd}(\sigma)$ in \mathcal{C} is ambigressive Cartesian.

Definition 1.1.6. The (complete) Segal space of spans $\mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ is de-

fined level-wise by

$$\mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)_n = \mathrm{Map}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^n), \mathcal{C})^\simeq.$$

For a proof that this really is a Segal space, the reader is once again referred to [2]. Note that the quasicategory of spans $\mathrm{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ is the bottom row of the complete Segal space of span $\mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$; that is,

$$\mathrm{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) = \mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)/\Delta^0.$$

Proposition ?? thus immediately implies that $\mathrm{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ is a quasicategory.

1.2 Important example: spans and equivalences

Let \mathcal{C} be a quasicategory. One can think about spans in \mathcal{C} as ‘fractions’ of morphisms in \mathcal{C} . Consider a span in \mathcal{C}

$$\begin{array}{ccc} & Y & \\ g \swarrow & & \searrow f \\ X & & X' \end{array}.$$

If the morphism g were invertible, we could fill this to a full 2-simplex

$$\begin{array}{ccc} & Y & \\ g \swarrow & & \searrow f \\ X & \xrightarrow{f \circ g^{-1}} & X' \end{array}.$$

Even if g is not invertible, we can think of our span as representing formally a map ‘ $f \circ g^{-1}: X \rightarrow X'$ ’, which may not exist in \mathcal{C} .

With this point of view in mind, if we build from a quasicategory \mathcal{C} a category of spans whose forwards-facing legs are arbitrary morphisms in \mathcal{C} , and whose backward-facing legs are equivalences, then we have not really done anything; each span, thought of as a formal morphism in \mathcal{C} , is represented by an actual morphism in \mathcal{C} . Thus, the category of spans in \mathcal{C} whose backward-facing legs are equivalences should be equivalent to the category \mathcal{C} . In the following example, we construct this equivalence.

Example 1.2.1. For any quasicategory \mathcal{C} , we can define a triple $(\mathcal{C}, \mathcal{C}_\dagger = \mathcal{C}, \mathcal{C}^\dagger = \mathcal{C}^\simeq)$. These choices of ingressive and egressive morphisms correspond to spans of the form

$$\begin{array}{ccc} & Y & \\ \simeq \swarrow & & \searrow \\ X & & X' \end{array},$$

i.e. spans such that the backwards-facing map $Y \rightarrow X$ is an equivalence in \mathcal{C} .

With this triple, a functor $\text{sd}(\Delta^n) \rightarrow \mathcal{C}$ is ambigressive cartesian if and only if each backward-facing leg is an equivalence. Note that the triple $(\mathcal{C}, \mathcal{C}, \mathcal{C}^\simeq)$ defined above is adequate even if \mathcal{C} does not admit pullbacks: for any solid diagram below, where g is an equivalence, the dashed square completion is a pullback, where g^{-1} is any homotopy inverse for g .

$$\begin{array}{ccc} X & \xrightarrow{g^{-1} \circ f} & X' \\ \text{id} \downarrow & & \downarrow g \\ X & \xrightarrow{f} & Y' \end{array}$$

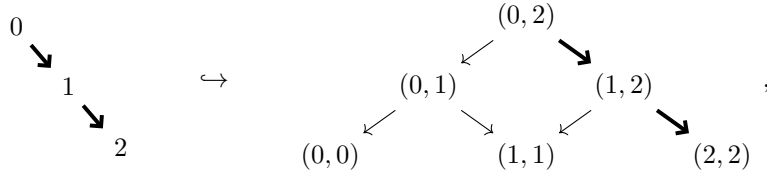
We will denote the resulting Segal space of spans by $\mathbf{Span}^\simeq(\mathcal{C})$, and the quasi-category $\mathbf{Span}^\simeq(\mathcal{C})/\Delta^0$ by $\text{Span}^\simeq(\mathcal{C})$.

The quasicategory $\text{Span}^\simeq(\mathcal{C})$ is categorically equivalent to the category \mathcal{C} from which we started. More specifically, there exists a weak categorical equivalence $\mathcal{C} \rightarrow \text{Span}^\simeq(\mathcal{C})$. We will now construct this map.

For each $n \geq 0$ there is a retraction of posets

$$[n] \hookrightarrow \Sigma_n \twoheadrightarrow [n],$$

where the first map takes $i \mapsto (i, n)$, and the second takes $(i, j) \mapsto i$. In the case $n = 2$, the inclusion can be drawn



and the map $\text{sd}(\Delta^n) \rightarrow \Delta^n$ ‘collapses’ the backward-facing legs of the spans. The general case is hopefully clear.

Applying the nerve, we find a retraction of simplicial sets

$$\Delta^n \hookrightarrow \text{sd}(\Delta^n) \twoheadrightarrow \Delta^n.$$

Pulling back along these maps gives us, for each n , a retraction

$$\text{Map}(\Delta^n, \mathcal{C})^\simeq \xrightarrow{i'_n} \text{Map}(\text{sd}(\Delta^n), \mathcal{C})^\simeq \xrightarrow{r'_n} \text{Map}(\Delta^n, \mathcal{C})^\simeq.$$

Note that the image of the inclusion i'_n consists of spans whose backwards-facing legs are identities (hence certainly equivalences), and hence that we can restrict the total space of our retraction:

$$\text{Map}(\Delta^n, \mathcal{C})^\simeq \xrightarrow{i_n} \text{Map}^{\text{aCart}}(\text{sd}(\Delta^n), \mathcal{C})^\simeq \xrightarrow{r_n} \text{Map}(\Delta^n, \mathcal{C})^\simeq.$$

It is not hard to convince oneself that a map $\mathrm{sd}(\Delta^n) \rightarrow \mathcal{C}$ is ambigressive cartesian if and only if it is a right Kan extension of its restriction along $\Delta^n \hookrightarrow \mathrm{sd}(\Delta^n)$. In the case $n = 2$, for example, it is clear from the limit formula that any right Kan extension of a diagram $\Delta^2 \rightarrow \mathcal{C}$ along the inclusion $\Delta^2 \hookrightarrow \mathrm{sd}(\Delta^2)$ is of the form

$$\begin{array}{ccc} X & & \\ \searrow & & \\ & Y & \\ & \searrow & \\ & & Z \end{array} \hookrightarrow \begin{array}{ccccc} & & X & & \\ & & \swarrow & \searrow & \\ & X' & & Y & \\ \swarrow & \searrow & & \swarrow & \searrow \\ X'' & & Y' & & Z \end{array},$$

(and conversely, any diagram of the above form is a right Kan extension), and diagrams of the above form are precisely those which are ambigressive cartesian. Put differently, $\mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^n), \mathcal{C})$ is the full subcategory of $\mathrm{Fun}(\mathrm{sd}(\Delta^n), \mathcal{C})$ on functors which are right Kan extensions of their restrictions to Δ^n . This, together with [3, Prop. 4.3.2.15] implies that r_n is a trivial fibration for all n . The identity $\mathrm{id}_{\mathrm{Fun}(\Delta^n, \mathcal{C})}$ is certainly a weak Kan equivalence, so by the 2/3 property for weak equivalences, i_n is also a weak equivalence for all n .

The projections $\mathrm{sd}(\Delta^n) \rightarrow \Delta^n$ assemble into a cosimplicial object $\mathrm{sd}(\Delta^\bullet) \rightarrow \Delta^\bullet$. From this it easily follows that the maps i_n assemble into a map of Segal spaces

$$i: \Gamma(\mathcal{C}) = \mathrm{Fun}(\Delta^\bullet, \mathcal{C}) \rightarrow \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^\bullet), \mathcal{C}) = \mathbf{Span}^\simeq(\mathcal{C})$$

which is a level-wise Kan weak equivalence, and hence a weak equivalence in the complete Segal spaces model structure.¹ Here $\Gamma(\mathcal{C})$ is the complete Segal space defined in [1] in the discussion preceding Thm. 4.11. In [1] it is also discussed that the functor $(-)/\Delta^0: \mathrm{Set}_{\Delta^2}^{\mathrm{CSS}} \rightarrow \mathrm{Set}_{\Delta}^{\mathrm{Joyal}}$ is the right Quillen functor in a Quillen equivalence between the Complete Segal Space model structure and the Joyal model structure. Thus, because $\Gamma(\mathcal{C})$ and $\mathbf{Span}^\simeq(\mathcal{C})$ are complete Segal spaces, the restriction of i to first rows i/Δ^0 is a weak categorical equivalence, which is what we wanted to show.

Example 1.2.1 allows us to find a copy of \mathcal{C} in a category of spans. Similarly, if we look at spans whose forward-facing legs are equivalences, we find a copy of $\mathcal{C}^{\mathrm{op}}$.

Example 1.2.2. For any quasicategory \mathcal{C} , we can define an adequate triple $(\mathcal{C}, \mathcal{C}_\dagger = \mathcal{C}^\simeq, \mathcal{C}^\dagger = \mathcal{C})$. These selections of ingressive and egressive morphisms correspond to spans of the form

$$\begin{array}{ccc} & Y & \\ \swarrow & & \searrow \\ X & & X' \end{array},$$

¹For a brief review of the model structures in question, see Appendix ???. For more information, see [1], especially the discussion following Thm. 4.1.

i.e. spans such that the forward-facing map $Y \rightarrow X$ is an equivalence in \mathcal{C} . We will denote the corresponding category of spans by $\text{Span}_{\simeq}(\mathcal{C})$. Exactly analogous reasoning to that in [Example 1.2.1](#) gives us a categorical equivalence

$$\mathcal{C}^{\text{op}} \rightarrow \text{Span}_{\simeq}(\mathcal{C}).$$

1.3 Maps between triples

A functor between adequate triples $(\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \rightarrow (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$ is a functor

$$p: \mathcal{C} \rightarrow \mathcal{D}$$

such that $p(\mathcal{C}_{\dagger}) \subseteq \mathcal{D}_{\dagger}$, $p(\mathcal{C}^{\dagger}) \subseteq \mathcal{D}^{\dagger}$, and p preserves ambigressive pullbacks. We will denote the restriction of p to \mathcal{C}_{\dagger} by

$$p_{\dagger}: \mathcal{C}_{\dagger} \rightarrow \mathcal{D}_{\dagger},$$

and similarly for p^{\dagger} . Clearly, a functor p between adequate triples gives functors $\text{Span}(p)$ and $\mathbf{Span}(p)$ respectively between quasicategories and Segal spaces of spans.

The goal of the remainder of this section is to provide a proof of the following theorem. This is [\[2, Thm. 12.2\]](#).

Theorem 1.3.1. Let $p: (\mathcal{C}, \mathcal{C}_{\dagger}, \mathcal{C}^{\dagger}) \rightarrow (\mathcal{D}, \mathcal{D}_{\dagger}, \mathcal{D}^{\dagger})$ be a functor of adequate triples which is an inner fibration. Further suppose that p satisfies the following conditions:

1. For any ingressive morphism $g: s \rightarrowtail t$ of \mathcal{D} and any object $x \in \mathcal{C}_s$, there exists an ingressive morphism $f: x \rightarrowtail y$ covering g which is both p -cocartesian and p_{\dagger} -cocartesian.
2. Suppose σ is a commutative square

$$\begin{array}{ccc} y' & \xrightarrow{f'} & x' \\ \downarrow & & \downarrow \psi \\ y & \xrightarrow{f} & x \end{array}$$

in \mathcal{C} such that $p(\sigma)$ is ambigressive pullback in \mathcal{D} , with in- and egressive morphisms as marked, and f is p -cocartesian. Then f' is p -cocartesian if and only if σ is an ambigressive pullback square (and in particular ψ is egressive).

Then π is an inner fibration, and if a morphism in $\text{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger)$ is of the form

$$\begin{array}{ccc} & y & \\ \phi \swarrow & & \searrow \psi \\ x & & x' \end{array},$$

where ϕ is egressive and p^\dagger -cartesian, and ψ is ingressive and p -cocartesian, it is π -cocartesian.

This theorem originally appears in [2]. However, the proof there is combinatorial in nature. Our proof uses the material developed in Section ??, leveraging the fact that ∞ -categories of spans have natural incarnations as (complete) Segal spaces.

Before we get to the meat of the proof, however, we have to prove some preliminary results.

1.4 Housekeeping

In this section we recall some basic facts we will need in the course of our proof of the main theorem. It is recommended to skip this section on first reading, and refer back to the results as necessary.

Isofibrations

We will need some results about isofibrations, sometimes called categorical fibrations. These are standard, and proofs can be found on Kerodon.

Proposition 1.4.1. Let $f: \mathcal{C} \rightarrow \mathcal{D}$ be a map between quasicategories. The following are equivalent:

- The map f is an isofibration.
- The map f is an inner fibration, and each equivalence in \mathcal{D} has a lift in \mathcal{C} which is an equivalence.

We will denote by $\text{Map}(-, -)$ the internal hom on Set_Δ .

Proposition 1.4.2. Let $f: \mathcal{C} \rightarrow \mathcal{D}$ be an isofibration between quasicategories, and let $i: A \hookrightarrow B$ be an inclusion of simplicial sets. Then the map

$$\text{Map}(B, \mathcal{C}) \rightarrow \text{Map}(A, \mathcal{C}) \times_{\text{Map}(A, \mathcal{D})} \text{Map}(B, \mathcal{D})$$

is an isofibration.

We will denote by $(-)^{\simeq}$ the *core* functor, i.e. the functor which takes a simplicial set X to the largest Kan complex it contains.

Proposition 1.4.3. Let $f: X \rightarrow Y$ be an isofibration of simplicial sets. Then

$$f^\simeq: X^\simeq \rightarrow Y^\simeq$$

is a Kan fibration.

We will denote the core of the mapping space functor $\text{Map}(-, -)$ by $\text{Fun}(-, -)$. That is, for simplicial sets A and B , we have

$$\text{Fun}(A, B) \cong \text{Map}(A, B)^\simeq.$$

Corollary 1.4.4. Let $f: \mathcal{C} \rightarrow \mathcal{D}$ be an isofibration between quasicategories, and let $A \hookrightarrow B$ be an inclusion of simplicial sets. Then the map

$$\text{Fun}(B, \mathcal{C}) \rightarrow \text{Fun}(A, \mathcal{C}) \times_{\text{Fun}(A, \mathcal{D})} \text{Fun}(B, \mathcal{D})$$

is a Kan fibration.

Connected components

Our proof will rely heavily on the fact that the many properties that morphisms can have (ingressive, egressive, p -cocartesian, etc.) are well-behaved with respect to the homotopical structure of the complete Segal spaces in which they live; for example, if a morphism is p -cocartesian, then every morphism in its path component is p -cocartesian.

Let X' and X be simplicial sets, and let $f: X' \hookrightarrow X$ be an inclusion. We will say that f is an *inclusion of connected components* if for all simplices $\sigma \in X$, if σ has any vertex in common with X' , then σ is wholly contained in X' .

Lemma 1.4.5. Let A be a simplicial set such that between any two vertices x and y of A there exists a finite zig-zag of 1-simplices of A connecting x to y . Let $A_0 \subseteq A$ be a nonempty simplicial subset of A . Let $f: X' \hookrightarrow X$ be a morphism between simplicial sets which is an inclusion of connected components. Then the following dashed lift always exists.

$$\begin{array}{ccc} A_0 & \xrightarrow{f_0} & X' \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ A & \xrightarrow{f} & X \end{array}$$

Proof. We have a map $f: A \rightarrow X$; in order to construct our dashed lift, it suffices to show that under the above assumptions, f takes every simplex of A to a simplex which belongs to X' . To this end, let $\sigma \in A$ be a simplex. There exists a finite zig-zag of 1-simplices connecting some vertex of σ to a vertex α of A_0 ; under f , this is mapped to a zig-zag of 1-simplices connecting $f(\alpha)$ to $f(\sigma)$. The vertex $f(\alpha)$ belongs to X' ; proceeding inductively, we find that the

image of every 1-simplex belonging to the zig-zag also belongs to X' , and that $f(\sigma)$ therefore also belongs to X' . \square

Corollary 1.4.6. Given any commuting square of simplicial sets

$$\begin{array}{ccc} X' & \xhookrightarrow{i} & X \\ f' \downarrow & & \downarrow f \\ Y' & \longrightarrow & Y \end{array}$$

with monomorphisms as marked, where i is an inclusion of connected components between Kan complexes and f is a Kan fibration, the map f' is a Kan fibration.

Furthermore, if f is a trivial fibration and f' is surjective on vertices, then f' is a trivial fibration.

Proof. First, we show that if f is a Kan fibration, then f' is a Kan fibration. We need to show that a dashed lift below exists for each $n \geq 1$ and each $0 \leq i \leq n$.

$$\begin{array}{ccccc} \Lambda_i^n & \longrightarrow & X' & \xhookrightarrow{i} & X \\ \downarrow & \nearrow \text{dashed} & \downarrow f' & & \downarrow f \\ \Delta^n & \longrightarrow & Y' & \longrightarrow & Y \end{array}$$

We can always solve our outer lifting problem, and by [Lemma 1.4.5](#), our lift factors through X' .

Now suppose that f is a trivial Kan fibration. The logic above says that f' has the right lifting property with respect to $\partial\Delta^n \hookrightarrow \Delta^n$ for $n \geq 1$; the right lifting property with respect to $\partial\Delta^0 \hookrightarrow \Delta^0$ is equivalent to the surjectivity of f' on vertices. \square

Contractibility

We will need at certain points to show that various spaces of lifts are contractible. These are mostly common-sense results, but it will be helpful to have them written down somewhere so we can refer to them later.

Lemma 1.4.7. Let $f: \mathcal{C} \rightarrow \mathcal{D}$ be an inner fibration between quasicategories. Then for any commuting square

$$\begin{array}{ccc} \Lambda_1^2 & \xrightarrow{\alpha} & \mathcal{C} \\ \downarrow & & \downarrow f \\ \Delta^2 & \xrightarrow{\beta} & \mathcal{D} \end{array},$$

the fiber F in the pullback square below is contractible.

$$\begin{array}{ccc} F & \longrightarrow & \mathrm{Fun}(\Delta^2, \mathcal{C}) \\ \downarrow & & \downarrow \\ \Delta^0 & \xrightarrow{(\alpha, f\alpha, \beta)} & \mathrm{Fun}(\Lambda_1^2, \mathcal{C}) \times_{\mathrm{Fun}(\Lambda_1^2, \mathcal{D})} \mathrm{Fun}(\Delta^2, \mathcal{D}) \end{array}$$

Proof. The right-hand map is a trivial Kan fibration. \square

One should interpret this as telling us that given a Λ_2^2 -horn α in \mathcal{C} lying over a 2-simplex β in \mathcal{D} , the space of fillings of α lying over β is contractible. We will need a pantheon of similar results.

Lemma 1.4.8. Let $f: \mathcal{C} \rightarrow \mathcal{D}$ be an inner fibration between quasicategories. Let e be any f -cartesian edge. Then for any square

$$\begin{array}{ccc} \Lambda_2^2 & \xrightarrow{\alpha} & \mathcal{C} \\ \downarrow & & \downarrow f \\ \Delta^2 & \xrightarrow{\beta} & \mathcal{D} \end{array}$$

such that $\alpha|_{\{1,2\}} = e$, the fiber F in the pullback diagram

$$\begin{array}{ccc} F & \longrightarrow & \mathrm{Fun}(\Delta^2, \mathcal{C}) \\ \downarrow & & \downarrow \\ \Delta^0 & \xrightarrow{(\alpha, f\alpha, \beta)} & \mathrm{Fun}(\Lambda_2^2, \mathcal{C}) \times_{\mathrm{Fun}(\Lambda_1^2, \mathcal{D})} \mathrm{Fun}(\Delta^2, \mathcal{D}) \end{array} \quad (1)$$

is contractible.

Proof. Denote $f(e) = \bar{e}$. Define a marking \mathcal{E}' on \mathcal{D} containing all degenerate edges and \bar{e} . Define a marking \mathcal{E} on \mathcal{C} containing the f -cocartesian lifts of the edges in \mathcal{E} . By [3, Prop. 3.1.1.6], the map $\mathcal{C}^{\mathcal{E}} \rightarrow \mathcal{D}^{\mathcal{E}'}$ has the right-lifting property with respect to all cocartesian-marked anodyne morphisms. It follows that

$$\begin{array}{c} \mathrm{Map}^\#((\Delta^2)^\mathcal{L}, \mathcal{C}^\mathcal{E}) \\ \downarrow \\ \mathrm{Map}^\#((\Lambda_2^2)^\mathcal{L}, \mathcal{C}^\mathcal{E}) \times_{\mathrm{Map}^\#((\Lambda_1^2)^\mathcal{L}, \mathcal{D}^{\mathcal{E}'})} \mathrm{Map}^\#((\Delta^2)^\mathcal{L}, \mathcal{D}^{\mathcal{E}'}) \end{array}$$

is a trivial Kan fibration; the map on the right-hand side of [Diagram 1](#) is the core of this map, and is thus also a trivial Kan fibration. \square

Lemma 1.4.9. Let $(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow (\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$ be an inner fibration between triples satisfying the conditions of [Theorem 1.3.1](#). Denote by $\mathrm{Fun}'(\Delta^1, \mathcal{C}^\dagger)$ the full simplicial subset of $\mathrm{Fun}(\Delta^1, \mathcal{C}^\dagger)$ on edges $\Delta^1 \rightarrow \mathcal{C}^\dagger$ which are both

p -cocartesian and p^\dagger -cocartesian. Then the fibers of the map

$$\mathrm{Fun}'(\Delta^1, \mathcal{C}^\dagger) \rightarrow \mathrm{Fun}(\Delta^{\{0\}}, \mathcal{C}^\dagger) \times_{\mathrm{Fun}(\Delta^{\{0\}}, \mathcal{D}^\dagger)} \mathrm{Fun}(\Delta^1, \mathcal{D}^\dagger)$$

are contractible.

Proof. Denote by $\mathrm{Fun}''(\Delta^1, \mathcal{C}^\dagger)$ the full simplicial subset of $\mathrm{Fun}(\Delta^1, \mathcal{C}^\dagger)$ on edges which are p^\dagger -cartesian. Then we have a square

$$\begin{array}{ccc} \mathrm{Fun}'(\Delta^1, \mathcal{C}^\dagger) & \xrightarrow{\quad\quad\quad} & \mathrm{Fun}''(\Delta^1, \mathcal{C}^\dagger) \\ \downarrow & & \downarrow \\ \mathrm{Fun}(\Delta^{\{0\}}, \mathcal{C}^\dagger) \times_{\mathrm{Fun}(\Delta^{\{0\}}, \mathcal{D}^\dagger)} \mathrm{Fun}(\Delta^1, \mathcal{D}^\dagger) & \xlongequal{\quad\quad\quad} & \mathrm{Fun}(\Delta^{\{0\}}, \mathcal{C}^\dagger) \times_{\mathrm{Fun}(\Delta^{\{0\}}, \mathcal{D}^\dagger)} \mathrm{Fun}(\Delta^1, \mathcal{D}^\dagger) \end{array}$$

in which the top map is an inclusion of connected components, the right map is a trivial Kan fibration, and the left map is surjective on vertices. Thus, the left map is a trivial Kan fibration by [Corollary 1.4.6](#) \square

1.5 Main theorem

Our goal is now to prove [Theorem 1.3.1](#). As in the statement of the theorem, we fix a map of triples $p: (\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow (\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$ giving us a functor of quasicategories

$$\pi: \mathrm{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow \mathrm{Span}(\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger).$$

We wish to show that morphisms of the form

$$\begin{array}{ccc} & y & \\ \phi \swarrow & & \searrow \psi \\ x & & x' \end{array}, \tag{2}$$

are π -cocartesian, where ϕ is egressive and p_\dagger -cartesian, and ψ is ingressive and p -cocartesian.

The way forward is clear: the functor π is the zeroth row of a functor of Segal spaces

$$\pi': \mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow \mathbf{Span}(\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger).$$

It will suffice to show by [Corollary ??](#) that morphisms of the above form are π' -cocartesian. Morally, this result should not be surprising. We should think of the homotopy pullback condition defining cocartesian morphisms ([Definition ??](#)) as telling us that we can fill relative Λ_2^2 -horns in $\pi': \mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow \mathbf{Span}(\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$, and that such fillings are unique up to contractible choice.

Finding a filling

$$\begin{array}{ccc} \Lambda_2^2 & \longrightarrow & \mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \\ \downarrow & \nearrow & \downarrow \pi' \\ \Delta^2 & \longrightarrow & \mathbf{Span}(\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger) \end{array}$$

is equivalent finding a filling

$$\begin{array}{ccc} \mathrm{sd}(\Lambda_2^2) & \longrightarrow & \mathcal{C} \\ \downarrow & \nearrow & \downarrow p \\ \mathrm{sd}(\Delta^2) & \longrightarrow & \mathcal{D} \end{array}$$

(such that the filling is ambigressive and the necessary square is pullback); the conditions on p guarantee us that we can perform this filling in a series of steps pictured in [Figure 1](#), each of which is unique up to contractible choice.

Reedy fibration

First, however, we have to check that π' is a Reedy fibration. This follows from the following useful fact, implied by the first assumption of [Theorem 1.3.1](#).

Lemma 1.5.1. Let $p: (\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow (\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$ be a functor of triples satisfying the conditions of [Theorem 1.3.1](#). Then the map $p: \mathcal{C} \rightarrow \mathcal{D}$ is an isofibration.

Proof. By assumption, p is an inner fibration and \mathcal{C} and \mathcal{D} are quasicategories, so it suffices by [Proposition 1.4.1](#) to show that p admits lifts of equivalences. Each equivalence in \mathcal{D} is in particular ingressive, and hence admits a p -cocartesian lift by the assumptions of [Theorem 1.3.1](#); these are automatically equivalences. \square

Proposition 1.5.2. For any functor of adequate triples $p: (\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow (\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$, the map

$$\pi': \mathbf{Span}(\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow \mathbf{Span}(\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$$

is a Reedy fibration.

Proof. Consider the following diagram.

$$\begin{array}{ccc} \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^n), \mathcal{C}) & \xhookrightarrow{\quad\quad\quad} & \mathrm{Fun}(\mathrm{sd}(\Delta^n), \mathcal{C}) \\ \downarrow & & \downarrow \\ \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\partial\Delta^n), \mathcal{C}) \times_{\mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\partial\Delta^n), \mathcal{D})} \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^n), \mathcal{D}) & \xhookrightarrow{\quad\quad\quad} & \mathrm{Fun}(\mathrm{sd}(\partial\Delta^n), \mathcal{C}) \times_{\mathrm{Fun}(\mathrm{sd}(\partial\Delta^n), \mathcal{D})} \mathrm{Fun}(\mathrm{sd}(\Delta^n), \mathcal{D}) \end{array}$$

The right-hand map is a Kan fibration by [Corollary 1.4.4](#) (because sd preserves monomorphisms), and the top map is an inclusion of connected components, so [Corollary 1.4.6](#) implies that the left-hand map is a Kan fibration, which is what we needed to show. \square

We are now ready to begin in earnest our proof that morphisms of the form [Equation 2](#) are π' -cocartesian. We first introduce some notation. Let $p: (\mathcal{C}, \mathcal{C}_\dagger, \mathcal{C}^\dagger) \rightarrow (\mathcal{D}, \mathcal{D}_\dagger, \mathcal{D}^\dagger)$ be a functor of triples.

- We will denote p -cartesian morphisms in \mathcal{C} with a circle:

$$x \text{ ---}\circ\text{--}\rightarrow y$$

- We will denote p -cocartesian morphisms in \mathcal{C} with a bullet:

$$x \text{ ---}\bullet\text{--}\rightarrow y$$

- We will denote p^\dagger -cartesian morphisms in \mathcal{C} with a triangle:

$$x \text{ ---}\triangleright\text{--}\rightarrow y$$

- We will denote p_\dagger -cocartesian morphisms in \mathcal{C} with a filled triangle:

$$x \text{ ---}\blacktriangleright\text{--}\rightarrow y$$

Thus, an ingressive morphism which is both p -cocartesian and p_\dagger -cocartesian will be denoted by

$$x \text{ ---}\blacktriangleright\bullet\text{--}\rightarrow y$$

Our proof will rest on the factorization $\text{sd}(\Lambda_0^2) \hookrightarrow \text{sd}(\Delta^2)$ pictured in [Figure 1](#). Denote the underlying factorization of simplicial sets by

$$A_1 \xrightarrow{i_1} A_2 \xrightarrow{i_2} A_3 \xrightarrow{i_3} A_4 \xrightarrow{i_4} A_5 \xrightarrow{i_5} A_6$$

For each of the A_i above, denote by

$$\text{Fun}'(A_i, \mathcal{C}) \subseteq \text{Fun}(A_i, \mathcal{C}) \tag{3}$$

the full simplicial subset on those functors $A_i \rightarrow \mathcal{C}$ which respect each labelling in [Figure 1](#):

- ingressive
- egressive
- p -cartesian
- p -cocartesian
- p^\dagger -cartesian
- p_\dagger -cocartesian

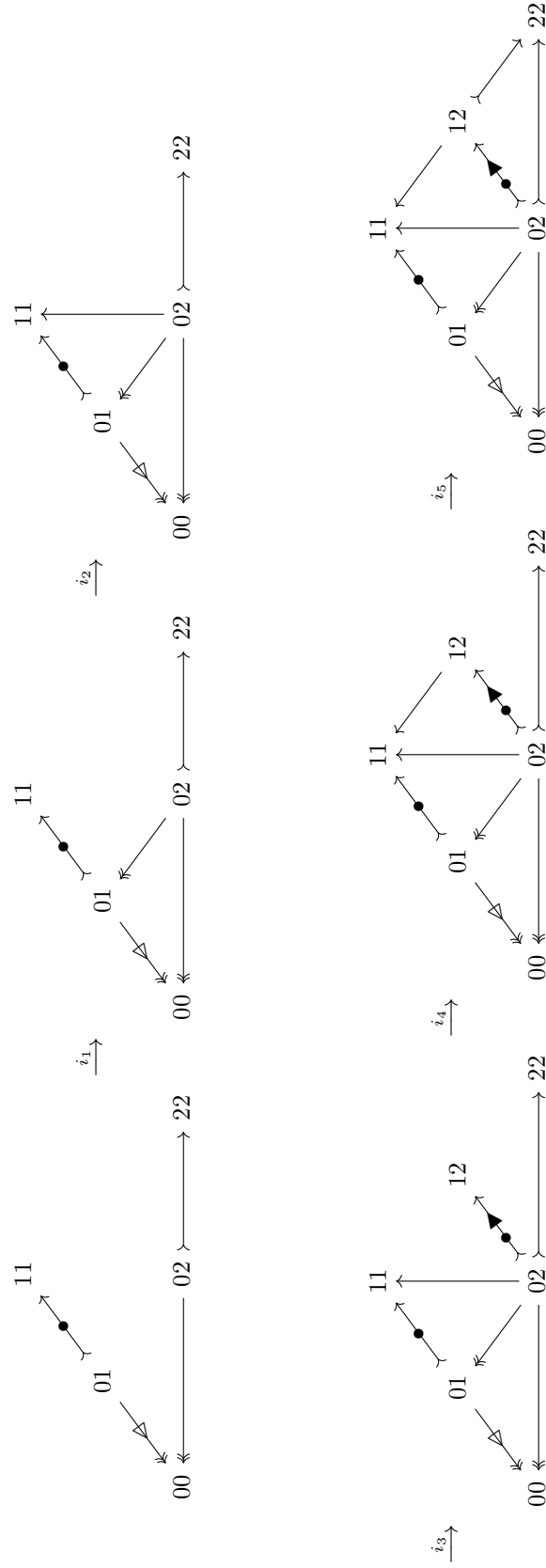


Figure 1: A factorization of the inclusion $\text{sd}(\Lambda_0^2) \hookrightarrow \text{sd}(\Delta^2)$, where certain morphisms have been labelled.

Note that because equivalences in \mathcal{C} belong to each of these classes of morphisms, the inclusion in [Equation 3](#) is an inclusion of connected components.

Similarly, we will denote by $\text{Fun}'(A_i, \mathcal{D})$ the full simplicial subset on functors which respect the ingressive and egressive labellings.

Lemma 1.5.3. The square

$$\begin{array}{ccc} \text{Fun}'(A_6, \mathcal{C}) & \longrightarrow & \text{Fun}'(A_0, \mathcal{C}) \\ \downarrow & & \downarrow \\ \text{Fun}'(A_6, \mathcal{D}) & \longrightarrow & \text{Fun}'(A_0, \mathcal{D}) \end{array}$$

is homotopy pullback.

Proof. The factorization of [Figure 1](#) gives us a factorization of the above square into five squares

$$\begin{array}{ccc} \text{Fun}'(A_{k+1}, \mathcal{C}) & \longrightarrow & \text{Fun}'(A_k, \mathcal{C}) \\ \downarrow & & \downarrow \\ \text{Fun}'(A_{k+1}, \mathcal{D}) & \longrightarrow & \text{Fun}'(A_k, \mathcal{D}) \end{array}, \quad 1 \leq k < 6,$$

each corresponding to one of the inclusions i_k . We will be done if we can show that each of these squares is homotopy pullback. It suffices to show that for each k , the map

$$j_k : \text{Fun}'(A_{k+1}, \mathcal{C}) \rightarrow \text{Fun}'(A_k, \mathcal{C}) \times_{\text{Fun}'(A_{k+1}, \mathcal{D})} \text{Fun}'(A_k, \mathcal{D})$$

is a weak equivalence. First, we show that each of these maps is a Kan fibration. To see this, consider the square

$$\begin{array}{ccc} \text{Fun}'(A_{k+1}, \mathcal{C}) & \xhookrightarrow{\quad} & \text{Fun}(A_{k+1}, \mathcal{C}) \\ j_k \downarrow & & \downarrow \\ \text{Fun}'(A_k, \mathcal{C}) \times_{\text{Fun}'(A_{k+1}, \mathcal{D})} \text{Fun}'(A_k, \mathcal{D}) & \xhookrightarrow{\quad} & \text{Fun}(A_k, \mathcal{C}) \times_{\text{Fun}(A_{k+1}, \mathcal{D})} \text{Fun}(A_k, \mathcal{D}) \end{array}.$$

The right-hand map is a Kan fibration because of [Corollary 1.4.4](#), and we have already seen that the top map is an inclusion of connected components. Hence each j_k is a Kan fibration by [Corollary 1.4.6](#). Thus, in order to show that each j_k is a weak equivalence, it suffices to show that the fibers are contractible.

First consider the case $k = 1$. For any map α below, consider the following pullback square.

$$\begin{array}{ccc} F_\alpha & \longrightarrow & \text{Fun}'(A_1, \mathcal{C}) \\ \downarrow & & \downarrow j_1 \\ \Delta^0 & \xrightarrow{\alpha} & \text{Fun}'(A_0, \mathcal{C}) \times_{\text{Fun}'(A_1, \mathcal{D})} \text{Fun}'(A_1, \mathcal{D}) \end{array}.$$

The fiber F_α over α is the space of ways of completing a diagram of shape A_0 in \mathcal{C} to a diagram of shape A_1 in \mathcal{C} given a diagram of shape A_1 in \mathcal{D} . This is the space of ways of filling $\Lambda_2^2 \hookrightarrow \Delta^2$ in \mathcal{C}^\dagger lying over a 2-simplex in \mathcal{D}^\dagger . This is contractible by [Lemma 1.4.8](#).

The case $k = 2$ is similar, using [Lemma 1.4.7](#).

The case $k = 3$ is similar, using [Lemma 1.4.9](#).

The cases $k = 4$ and $k = 5$ use the dual to [Lemma 1.4.8](#). \square

We are now ready to show that morphisms of the form [Diagram 2](#) are π' -cocartesian.

Proposition 1.5.4. For any morphism e of the form given in [Equation 2](#), the square

$$\begin{array}{ccc} \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^2), \mathcal{C}) \times_{\mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{C})} \{e\} & \longrightarrow & \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Lambda_0^2), \mathcal{C}) \times_{\mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{C})} \{e\} \\ \downarrow & & \downarrow \\ \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^2), \mathcal{D}) \times_{\mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{D})} \{\pi e\} & \longrightarrow & \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Lambda_0^2), \mathcal{D}) \times_{\mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{D})} \{\pi e\} \end{array}$$

is homotopy pullback.

Proof. This square factors into the two squares

$$\begin{array}{ccc} \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^2), \mathcal{C}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{C})} \{e\} & \longrightarrow & \mathrm{Fun}'(A_6, \mathcal{C}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{C})} \{e\} \\ \downarrow & & \downarrow \\ \mathrm{Fun}^{\mathrm{aCart}}(\mathrm{sd}(\Delta^2), \mathcal{D}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{D})} \{\pi e\} & \longrightarrow & \mathrm{Fun}'(A_6, \mathcal{D}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{D})} \{\pi e\} \end{array}$$

and

$$\begin{array}{ccc} \mathrm{Fun}'(A_6, \mathcal{C}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{C})} \{e\} & \longrightarrow & \mathrm{Fun}'(A_0, \mathcal{C}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{C})} \{e\} \\ \downarrow & & \downarrow \\ \mathrm{Fun}'(A_6, \mathcal{D}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{D})} \{\pi e\} & \longrightarrow & \mathrm{Fun}'(A_0, \mathcal{D}) \times_{\mathrm{Fun}(\mathrm{sd}(\Delta^{\{0,1\}}), \mathcal{D})} \{\pi e\} \end{array}.$$

The first is homotopy pullback because the bottom map is an inclusion of connected components, and the second condition of [Theorem 1.3.1](#) guarantees that the fiber over an ambigressive $\mathrm{sd}(\Delta^2) \rightarrow \mathcal{D}$ belonging to $\mathrm{Fun}'(A_6, \mathcal{D})$ whose restriction to $\mathrm{sd}(\Delta^{\{0,1\}})$ is πe is precisely an ambigressive Cartesian functor $\mathrm{sd}(\Delta^2) \rightarrow \mathcal{C}$ whose restriction to $\mathrm{sd}(\Delta^{\{0,1\}})$ is e .

That the second is homotopy pullback follows immediately from [Lemma 1.5.3](#). \square

This proves [Theorem 1.3.1](#).