

A Appendix

A.1 A brief list of model structures

The category Set_Δ carries two model structures of which we will make frequent use:

- The *Kan* model structure, which has the following description.
 - The fibrations are the Kan fibrations.
 - The cofibrations are the monomorphisms.
 - The weak equivalences are weak homotopy equivalences.
- The *Joyal* model structure. We will not give a complete description, referring the reader to [2, Sec. 2.2.5]. We will make use of the following properties.
 - The cofibrations are monomorphisms.
 - The fibrant objects are quasicategories, and the fibrations between fibrant objects are isofibrations, i.e. inner fibrations with lifts of equivalences (cf. [2, Cor. 2.6.5]).

The category Δ^{op} has a Reedy structure, which gives us, together with the Kan model structure, a model structure on the category $\text{Fun}(\Delta^{\text{op}}, \text{Set}_\Delta) = \text{Set}_{\Delta^2}$ with the following properties.

- The cofibrations are monomorphisms.
- The weak equivalences are level-wise weak homotopy equivalences.
- The fibrations are *Reedy fibrations* (Definition ??).

This model category has two important left Bousfield localizations which we will need.

- The *Segal space model structure* on Set_{Δ^2} is the left Bousfield localization of the Reedy model structure at the set of spine inclusions

$$S = \left\{ \Delta^{\{0,1\}} \amalg_{\Delta^{\{1\}}} \Delta^{\{1,2\}} \amalg_{\Delta^{\{2\}}} \cdots \amalg_{\Delta^{\{n-1\}}} \Delta^{\{n-1,n\}} \hookrightarrow \Delta^n \mid n \geq 2 \right\},$$

which has the following properties.

- The cofibrations are the monomorphisms.
- The fibrant objects are Segal spaces.
- Every Reedy weak equivalence is a weak equivalence in the Segal space model structure.

- The *complete Segal space model structure* is the left Bousfeld localization of the Reedy model structure at the set S of spine inclusions together with the inclusion $\{0\} \hookrightarrow I$, where I is the (nerve of the) walking isomorphism. The complete Segal space model structure has the following properties.
 - The cofibrations are the monomorphisms.
 - The fibrant objects are complete Segal spaces.
 - Every weak equivalence in the Segal space model structure is a weak equivalence in the complete Segal space model structure.

Proposition A.1.1 (Joyal–Tierney). There is a Quillen equivalence

$$p_1^* : \mathbf{Set}_\Delta^{\text{Joyal}} \longleftrightarrow \mathbf{Set}_{\Delta^2}^{\text{CSS}} : i_1^*$$

between the Joyal model structure on simplicial sets and the complete Segal space model structure on bisimplicial spaces. Here, the functors p_1^* and i_1^* are defined as follows.

- The functor p_1^* takes a simplicial set X to the bisimplicial space $(p_1^*X)_{mn} = X_m$. That is, the bisimplicial set p_1^*X is constant in the vertical direction, and each row is equal to the simplicial set X .
- The functor i_1^* takes a bisimplicial set K to the simplicial set $(i_1^*K)_n = K_{n0}$. That is, in the language of [Subsection A.1](#), the simplicial set i_1^*K is the ‘zeroth row’ of K .

Note A.1.2. In the notation of

For more information, the reader is directed to [\[1\]](#).

A.2 Divisibility of bifunctors

In this section, we recall some key results from [\[1\]](#). We refer readers there for more information.

Let $\odot : \mathcal{E}_1 \times \mathcal{E}_2 \rightarrow \mathcal{E}_3$ be a functor of 1-categories. We will say that \odot is *divisible on the left* if for each $A \in \mathcal{E}_1$, the functor $A \odot -$ admits a right adjoint $A \backslash -$. In this case, this construction turns out also to be functorial in A ; that is, we get a functor

$$-\backslash - : \mathcal{E}_1^{\text{op}} \times \mathcal{E}_3 \rightarrow \mathcal{E}_2.$$

Analogously, $\odot : \mathcal{E}_1 \times \mathcal{E}_2 \rightarrow \mathcal{E}_3$ is *divisible on the right* if for each $B \in \mathcal{E}_2$, the functor $- \odot B$ admits a right adjoint $-/B$. In this case we get a of two variables

$$-/ - : \mathcal{E}_3 \times \mathcal{E}_2^{\text{op}} \rightarrow \mathcal{E}_1.$$

Example A.2.1. The reader may find it helpful to keep in mind the cartesian product

$$- \times -: \mathbf{Set}_\Delta \times \mathbf{Set}_\Delta \rightarrow \mathbf{Set}_\Delta.$$

In this case, both $A \backslash X$ and X/A are the mapping space X^A .

If \odot is divisible on both sides, then there is a bijection between maps of the following types:

$$A \odot B \rightarrow X, \quad A \rightarrow X/B, \quad B \rightarrow A \backslash X.$$

In particular, this implies that the functors $X/-$ and $-\backslash X$ are mutually right adjoint.

If both \mathcal{E}_1 and \mathcal{E}_2 are finitely complete and \mathcal{E}_3 is finitely cocomplete, then from a map $u: A \rightarrow A'$ in \mathcal{E}_1 , a map $v: B \rightarrow B'$ in \mathcal{E}_2 , and a map $f: X \rightarrow Y$ in \mathcal{E}_3 , we can build the following maps.

- From the square

$$\begin{array}{ccc} A \odot B & \longrightarrow & A' \odot B \\ \downarrow & & \downarrow \\ A \odot B' & \longrightarrow & A' \odot B' \end{array}$$

we get a map

$$u \odot' v: A \odot B' \amalg_{A \odot B} A' \odot B \rightarrow A' \odot B'.$$

- From the square

$$\begin{array}{ccc} A' \backslash X & \longrightarrow & A \backslash X \\ \downarrow & & \downarrow \\ A' \backslash Y & \longrightarrow & A \backslash Y \end{array}$$

we get a map

$$\langle u \backslash f \rangle: A' \backslash X \rightarrow A \backslash X \times_{A \backslash Y} A' \backslash Y$$

- From the square

$$\begin{array}{ccc} X/B' & \longrightarrow & X/B \\ \downarrow & & \downarrow \\ Y/B' & \longrightarrow & Y/B \end{array}$$

we get a map

$$\langle f/v \rangle: X/B' \rightarrow X/B \times_{Y/B} Y/B'.$$

Proposition A.2.2. With the above notation, the following are equivalent

adjoint lifting problems:

$$\begin{array}{ccc}
 A \odot B' \amalg_{A \odot B} A' \odot B & \longrightarrow & X \\
 \downarrow & \nearrow \text{dashed} & \downarrow \\
 A' \odot B' & \longrightarrow & Y
 \end{array}
 \qquad
 \begin{array}{ccc}
 A & \longrightarrow & X/B' \\
 \downarrow & \nearrow \text{dashed} & \downarrow \\
 A' & \longrightarrow & X/B \times_{Y/B} Y/B'
 \end{array}$$

$$\begin{array}{ccc}
 B & \longrightarrow & A' \backslash X \\
 \downarrow & \nearrow \text{dashed} & \downarrow \\
 B' & \longrightarrow & A' \backslash X \times_{A' \backslash Y} A' \backslash Y
 \end{array}$$