

1 Application to presheaves

1.1 Introduction

Consider the ∞ -category \mathcal{S} of spaces. By [3, Lemma 6.1.1.1] and [3, Lemma 2.4.7.12], the functor

$$\mathrm{ev}_1 : \mathrm{Map}(\Delta^1, \mathcal{S}) \rightarrow \mathcal{S} \quad (1)$$

is a bicartesian fibration. For each edge $f : \Delta^1 \rightarrow \mathcal{S}$, corresponding to a morphism $f : X \rightarrow Y$ of spaces, the fiber of ev_1 over f corresponds to a pair of adjoint functors

$$f_! : \mathcal{S}/X \longleftrightarrow \mathcal{S}/Y : f^*.$$

Interpreting \mathcal{S}/X and \mathcal{S}/Y as models for the $\mathcal{P}(X)$ and $\mathcal{P}(Y)$, the ∞ -categories of presheaves on X and Y , the functor f^* corresponds to restriction, taking a presheaf $Y \rightarrow \mathcal{S}$ to the pullback $X \rightarrow Y \rightarrow \mathcal{S}$, and the functor $f_!$ takes a presheaf $X \rightarrow \mathcal{S}$ to its left Kan extension along $X \rightarrow Y$.

In this section, we apply the construction of the previous chapter to build from these data a functor

$$\hat{r} : \mathrm{Span}(\mathcal{S}) \rightarrow \mathrm{Cat}_\infty. \quad (2)$$

This functor will take a space X to the category $\mathcal{P}(X)$ of presheaves on X , and a span of spaces

$$\begin{array}{ccc} & Y & \\ g \swarrow & & \searrow f \\ X & & X' \end{array}$$

to the map $\mathcal{P}(X) \rightarrow \mathcal{P}(Y)$ given by the composition

$$\begin{array}{ccc} & \mathcal{P}(Y) & \\ g_! \nearrow & & \searrow f^* \\ \mathcal{P}(X) & \xrightarrow{f^* \circ g_!} & \mathcal{P}(X'). \end{array}.$$

Note that the category \mathcal{S} has products. This induces a symmetric monoidal structure on \mathcal{S} , the Cartesian monoidal structure, which (as we will see) induces a symmetric monoidal structure on $\mathrm{Span}(\mathcal{S})$. When we take Cat_∞ also to carry the Cartesian monoidal structure, we will see that our functor \hat{r} (Equation 2) is even lax monoidal.

The idea behind the construction of the functor \hat{r} , in broad strokes, is as follows. We dress up the functor ev_1 , making it into a symmetric monoidal functor as follows (the monoidal structures on $\mathrm{Map}(\Delta^1, \mathcal{S})^\otimes$ and \mathcal{S}^\otimes are not exactly the standard cocartesian monoidal functors, as we will later discuss; they are

represented by *cartesian*, rather than *cocartesian* fibrations).

$$\begin{array}{ccc} \mathrm{Map}(\Delta^1, \mathcal{S})^\otimes & \xrightarrow{r} & \mathcal{S}^\otimes \\ & \searrow q \quad \swarrow p & \\ & \mathcal{F}\mathrm{in}_*^{\mathrm{op}} & \end{array}$$

We then take spans in each of these categories, restricting the backwards- and forwards-facing legs judiciously so that we get a diagram

$$\begin{array}{ccc} \mathrm{Span}(\mathrm{Map}(\Delta^1, \mathcal{S}))^\otimes & \xrightarrow{\rho} & \mathrm{Span}(\mathcal{S})^\otimes \\ & \searrow \varpi \quad \swarrow \pi & \\ & \mathcal{F}\mathrm{in}_* & \end{array},$$

where the ϖ and π are symmetric monoidal structures and ρ is a cocartesian fibration. In the language of [2], ρ exhibits $\mathrm{Span}(\mathrm{Map}(\Delta^1, \mathcal{S}))$ as a $\mathrm{Span}(\mathcal{S})$ -monoidal category. Restricting to the component of ρ over $\langle 1 \rangle \in \mathcal{F}\mathrm{in}_*$, one finds a cocartesian fibration, which classifies a functor $\mathrm{Span}(\mathcal{S}) \rightarrow \mathcal{C}\mathrm{at}_\infty$. By [2], 2.4.2.4–2.4.2.6, this functor is lax monoidal, giving us our functor \hat{r} .

Much of the technology in this section is present in [1]. However, the manner in which we present it allows for generalization in a different direction than that carried out there. This will be the subject of future work.

1.2 Defining the maps p , q , and r

Recall that for any ∞ -category \mathcal{C} with finite products, one can construct a symmetric monoidal structure $\mathcal{C}^\times \rightarrow \mathcal{F}\mathrm{in}_*$ whose corresponding tensor product $\times: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ is the cartesian product; for more details, see [2, Sec. 2.4.1]. Similarly, as described in [2, Sec. 2.4.2], for any ∞ -category with coproducts there is a symmetric monoidal structure $\mathcal{C}^\amalg \rightarrow \mathcal{F}\mathrm{in}_*$, called the *cocartesian monoidal structure*, whose tensor product $\amalg: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ is the coproduct. Because cocartesian fibrations give better control of the morphisms out of an object than morphisms into an object, the cocartesian monoidal structure is much more tractable than the cartesian monoidal structure.

Since taking coproducts in $\mathcal{C}^{\mathrm{op}}$ is the same as taking products in \mathcal{C} , we can express the cartesian monoidal structure on \mathcal{C} as a cartesian fibration rather than a cocartesian fibration by taking the opposite of the cocartesian monoidal structure on $\mathcal{C}^{\mathrm{op}}$. In the case $\mathcal{C} = \mathcal{S}$, this gives us the monoidal structure we are interested in.

Definition 1.2.1. We define the map $p: \mathcal{S}_\times \rightarrow \mathcal{F}\mathrm{in}_*^{\mathrm{op}}$ to be the functor

$$((\mathcal{S}^{\mathrm{op}})^\amalg \rightarrow \mathcal{F}\mathrm{in}_*)^{\mathrm{op}},$$

where $(\mathcal{S}^{\mathrm{op}})^\amalg \rightarrow \mathcal{F}\mathrm{in}_*$ is the cocartesian monoidal structure, as described in [2,

Sec. 2.4.2].

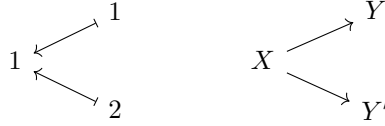
In low degrees, the category \mathcal{S}_\times admits the following description.

- The objects are pairs $(\langle n \rangle, \vec{X})$, where $\langle n \rangle$ is an object of $\mathcal{F}\text{in}_*$, and $\vec{X} = [X_1, \dots, X_n]$ is an ordered n -tuple of objects of \mathcal{S} .
- The morphisms $(\langle m \rangle, \vec{X}) \rightarrow (\langle n \rangle, \vec{Y})$ are pairs (ϕ, Ψ) , where $\phi: \langle n \rangle \rightarrow \langle m \rangle$ is a morphism in $\mathcal{F}\text{in}_*$, and Ψ consists, for each $i \in \langle m \rangle^\circ$, of a collection of morphisms $X_i \rightarrow Y_j$ for each $j \in \langle n \rangle^\circ$ with $\phi(j) = i$.

The functor $\mathcal{S}_\times \rightarrow \mathcal{F}\text{in}_*$ is the obvious forgetful functor.

Example 1.2.2. An object $(\langle 1 \rangle, \vec{X})$ of \mathcal{S}_\times in the fiber over $\langle 1 \rangle$ is the same thing as an object $X \in \langle S \rangle$. An object $(\langle 2 \rangle, \vec{Y})$ in the fiber over $\langle 2 \rangle$ is the same as a pair of objects $[Y, Y']$ in \mathcal{S} .

A morphism $(\phi, \Psi): (\langle 1 \rangle, \vec{X}) \rightarrow (\langle 2 \rangle, \vec{Y})$ where $\phi: \langle 2 \rangle \rightarrow \langle 1 \rangle$ is the morphism in $\mathcal{F}\text{in}_*$ which sends $1 \mapsto 1, 2 \mapsto 1$, consists of a pair of maps $X \rightarrow Y, X \rightarrow Y'$.



Lemma 1.2.3. The map p is a cartesian fibration, and a morphism $(\phi, \Psi): (\langle m \rangle, \vec{X}) \rightarrow (\langle n \rangle, \vec{Y})$ in \mathcal{S}_\times is p -cartesian if and only if for each $i \in \langle m \rangle^\circ$, the maps $X_i \rightarrow Y_j$ exhibit X_i as the product of the Y_j .

Proof. This follows easily from the definition of the cocartesian monoidal structure. \square

Example 1.2.4. The morphism described in [Example 1.2.2](#) is cartesian if and only if $X \simeq Y \times Y'$, and the maps $X \rightarrow Y$ and $X \rightarrow Y'$ are the canonical projections.

We can now define the maps q and r in one fell swoop.

Definition 1.2.5. We define the maps q and r via the following diagram, where the left-hand square is pullback. In particular, the space $\text{Map}(\Delta^1, \mathcal{S})_\times$ is defined to be the (strict) pullback below.

$$\begin{array}{ccccc}
 & & r & & \\
 & \text{Map}(\Delta^1, \mathcal{S})_\times & \hookrightarrow & \text{Map}(\Delta^1, \mathcal{S}_\times) & \xrightarrow{\text{ev}_1} \mathcal{S}_\times \\
 q \downarrow & & & \downarrow & \downarrow p \\
 \mathcal{F}\text{in}_*^{\text{op}} & \hookrightarrow & \text{Map}(\Delta^1, \mathcal{F}\text{in}_*^{\text{op}}) & \xrightarrow{\text{ev}_1} & \mathcal{F}\text{in}_*^{\text{op}} \\
 & & \text{id} & &
 \end{array}$$

Lemma 1.2.6. The map q is a Cartesian fibration, and the cartesian morphisms are those which are level-wise cartesian morphisms in \mathcal{S}_\times .

Proof. The map q is given by the pullback

$$\begin{array}{ccc} \mathrm{Map}(\Delta^1, \mathcal{S})_\times & \hookrightarrow & \mathrm{Map}(\Delta^1, \mathcal{S}_\times) \\ q \downarrow & & \downarrow q' \\ \mathcal{F}\mathrm{in}_*^{\mathrm{op}} & \hookrightarrow & \mathrm{Map}(\Delta^1, \mathcal{F}\mathrm{in}_*^{\mathrm{op}}) \end{array} .$$

By [3, Prop. 3.1.2.1], the map q' is a cartesian fibration whose cartesian morphisms are pointwise cartesian morphisms in \mathcal{S}_\times . The map q is also a cartesian fibration, where a morphism in $\mathrm{Map}(\Delta^1, \mathcal{S})_\times$ is q -cartesian if and only if its image in $\mathrm{Map}(\Delta^1, \mathcal{S}_\times)$ is q' -cartesian. That every q cartesian morphism is q' cartesian follows from [3, Prop 2.4.1.3]; to see other direction, consider a q -cartesian morphism in $\mathrm{Map}(\Delta^1, \mathcal{S})_\times$. \square

Lemma 1.2.7. The map r is a bicartesian fibration with r -cartesian morphisms pullback squares in \mathcal{S}_\times and r -cocartesian morphisms squares of the form

$$\begin{array}{ccc} X & \xrightarrow{\sim} & Y \\ \downarrow & & \downarrow \\ X' & \longrightarrow & Y' \end{array} .$$

Proof. The map r factors into

$$\mathrm{Map}(\Delta^1, \mathcal{S})_\times \xhookrightarrow{j} \mathrm{Map}(\Delta^1, \mathcal{S}_\times) \xrightarrow{\mathrm{ev}_1} \mathcal{S}_\times$$

The map ev_1 is a bicartesian fibration with the cartesian and cocartesian morphisms described above; the map j is a full inclusion (in the sense of Definition ??), and therefore has the right-lifting property with respect to *all* horn inclusions $\Lambda_i^n \hookrightarrow \Delta^n$, $n \geq 2$ (because Λ_i^n contains all vertices for $n \geq 2$). Every edge in $\mathrm{Map}(\Delta^1, \mathcal{S})_\times$ is therefore both j -cartesian and j -cocartesian. [3, Prop 2.4.1.3.3] implies that an edge in $\mathrm{Map}(\Delta^1, \mathcal{S})_\times$ is r -cartesian (resp. cocartesian) if and only if its image in $\mathrm{Map}(\Delta^1, \mathcal{S}_\times)$ is ev_1 -cartesian (resp. cocartesian). \square

Let us take stock. We have now defined the data of the commutative diagram below. Here, p and q are cartesian fibrations, and r is a bicartesian fibration.

$$\begin{array}{ccc} \mathrm{Map}(\Delta^1, \mathcal{S})_\times & \xrightarrow{r} & \mathcal{S}_\times \\ & \searrow q \quad \swarrow p & \\ & \mathcal{F}\mathrm{in}_*^{\mathrm{op}} & \end{array}$$

Our next step is to take spans with judiciously-chosen legs in each of these quascategories. To do this, we must define the subcategories to which the legs

of our spans are allowed to belong. As in Section ??, we do this by defining triples of categories.

1.3 Restricting the legs of the spans

In this section, we define triple structures on the quasicategories $\text{Map}(\Delta^1, \mathcal{S})_\times$, \mathcal{S}_\times , and Fin_*^{op} .

We start by defining a triple structure on Fin_*^{op} .

Definition 1.3.1. We define a triple $(\mathcal{F}, \mathcal{F}_\dagger, \mathcal{F}^\dagger)$ as follows.

- $\mathcal{F} = \text{Fin}_*^{\text{op}}$
- $\mathcal{F}_\dagger = (\text{Fin}_*^{\text{op}})^\simeq$
- $\mathcal{F}^\dagger = \text{Fin}_*^{\text{op}}$

Recall that in Example ??, we considered a triple structure on any quasicategory \mathcal{C} which produced spans in \mathcal{C} whose backwards-pointing legs were equivalences; this yielded a quasicategory $\text{Span}^\simeq(\mathcal{C})$ which was equivalent to \mathcal{C} via a weak categorical equivalence $\mathcal{C} \rightarrow \text{Span}^\simeq(\mathcal{C})$, which we constructed explicitly. The triple $(\mathcal{F}, \mathcal{F}_\dagger, \mathcal{F}^\dagger)$ is similar to the triple structure used to define $\text{Span}^\simeq(\mathcal{C})$, except that the forward-facing legs are equivalences rather than the backward-facing legs. A dual construction thus produces the following.

Lemma 1.3.2. There is a weak Joyal equivalence

$$\text{Fin}_* \rightarrow \text{Span}(\mathcal{F}, \mathcal{F}_\dagger, \mathcal{F}^\dagger),$$

which sends an n -simplex $\langle m \rangle \rightarrow \langle n \rangle \rightarrow \cdots \rightarrow \langle p \rangle$ to a diagram

$$\begin{array}{ccccc} & & \langle p \rangle & & \\ & \nearrow \cdots & \uparrow & \nwarrow \cdots & \\ \langle m \rangle & \nearrow & \langle n \rangle & \nwarrow & \langle p \rangle \\ & \searrow \cdots & \downarrow & \swarrow \cdots & \\ & & \langle n \rangle & & \langle p \rangle \end{array} \quad ,$$

where we have drawn the morphisms in the span as belonging to Fin_* rather than Fin_*^{op} .

The reasoning in Example ?? also shows the following.

Lemma 1.3.3. The triple $(\mathcal{F}, \mathcal{F}_\dagger, \mathcal{F}^\dagger)$ is adequate.

Next, we define a triple structure on \mathcal{S}_\times .

Definition 1.3.4. We define a triple $(\mathcal{P}, \mathcal{P}_\dagger, \mathcal{P}^\dagger)$ as follows.

- $\mathcal{P} = \mathcal{S}_\times$
- $\mathcal{P}_\dagger = \mathcal{S}_\times \times_{\mathcal{F}\text{in}_*^{\text{op}}} (\mathcal{F}\text{in}_*^{\text{op}})^{\simeq}$
- $\mathcal{P}^\dagger = \mathcal{S}_\times$

Here, we do not restrict the legs of our spans at all, except to ensure that p restricts to a map $p_\dagger: \mathcal{P}_\dagger \rightarrow \mathcal{F}_\dagger$.

Next, we need to show that the triple $(\mathcal{P}, \mathcal{P}_\dagger, \mathcal{P}^\dagger)$ is adequate. For this we need a few lemmas.

Lemma 1.3.5. The cartesian fibration $p: \mathcal{S}_\times \rightarrow \mathcal{F}\text{in}_*^{\text{op}}$ admits relative pullbacks.

Proof. By [3, Cor. 4.3.1.1], it suffices to show that each fiber $(\mathcal{S}_\times)_{\langle n \rangle}$ admits pullbacks, and that for each map $\phi: \langle m \rangle \leftarrow \langle n \rangle$ in $\mathcal{F}\text{in}_*^{\text{op}}$, the map $\phi^*: (\mathcal{S}_\times)_{\langle m \rangle} \rightarrow (\mathcal{S}_\times)_{\langle n \rangle}$ preserves pullbacks. \square

Lemma 1.3.6. The triple $(\mathcal{P}, \mathcal{P}_\dagger, \mathcal{P}^\dagger)$ is adequate.

Proof. It will suffice to show that \square

Lastly, we define a triple structure on $\text{Map}(\Delta^1, \mathcal{S})_\times$.

Definition 1.3.7. We define a triple $(\mathcal{Q}, \mathcal{Q}_\dagger, \mathcal{Q}^\dagger)$ as follows.

- $\mathcal{Q} = \text{Map}(\Delta^1, \mathcal{S})_\times$.
- \mathcal{Q}_\dagger is given by the pullback

$$\begin{array}{ccc} \mathcal{Q}_\dagger & \longrightarrow & \mathcal{Q} \\ \downarrow & & \downarrow \\ (\mathcal{F}\text{in}_*^{\text{op}})^{\simeq} & \longrightarrow & \mathcal{F}\text{in}_*^{\text{op}} \end{array} .$$

- \mathcal{Q}^\dagger is given by the subcategory of \mathcal{Q} on morphisms

$$\begin{array}{ccc} S & \xrightarrow{f} & T \\ \downarrow & & \downarrow \\ X & \longrightarrow & Y \end{array}$$

such that f is p -cartesian.

Lemma 1.3.8. The triple $(\mathcal{Q}, \mathcal{Q}_\dagger, \mathcal{Q}^\dagger)$ is adequate.

The maps p , q , and r are in particular maps $\mathcal{M} \rightarrow \mathcal{P} \rightarrow \mathcal{F}$ as follows:

$$\begin{array}{ccc} \text{Map}(\Delta^1, \mathcal{S})_{\times} & \xrightarrow{r} & \mathcal{S}_{\times} \\ & \searrow q \quad \swarrow p & \\ & \mathcal{F}\text{in}_{*}^{\text{op}} & \end{array}$$

By definition, p , q , and r restrict to maps

$$\begin{array}{ccc} \mathcal{M}_{\dagger} & \xrightarrow{r_{\dagger}} \mathcal{P}_{\dagger} & \xrightarrow{q_{\dagger}} \mathcal{F}_{\dagger} \\ & \searrow q_{\dagger} & \swarrow q_{\dagger} \end{array} \quad \text{and} \quad \begin{array}{ccc} \mathcal{M}^{\dagger} & \xrightarrow{r^{\dagger}} \mathcal{P}^{\dagger} & \xrightarrow{q^{\dagger}} \mathcal{F}^{\dagger} \\ & \searrow q^{\dagger} & \swarrow q^{\dagger} \end{array} .$$

1.4 Building categories of spans

We now have the following diagram of adequate triples.

$$\begin{array}{ccc} (\mathcal{Q}, \mathcal{Q}_{\dagger}, \mathcal{Q}^{\dagger}) & \xrightarrow{\quad} & (\mathcal{P}, \mathcal{P}_{\dagger}, \mathcal{P}^{\dagger}) \\ & \searrow \quad \swarrow & \\ & (\mathcal{F}, \mathcal{F}_{\dagger}, \mathcal{F}^{\dagger}) & \end{array}$$

Lemma 1.4.1. The morphism p satisfies the conditions of Theorem ??.

Proof. We check the conditions.

1. We have to show that for any isomorphism in $\mathcal{F}\text{in}_{*}^{\text{op}}$, there exists a p -cocartesian lift. But since every morphism in $\mathcal{F}_{\dagger} = (\mathcal{F}\text{in}_{*}^{\text{op}})^{\simeq}$ is an equivalence, it suffices to find a p -cartesian lift, which is possible because p is a Cartesian fibration.
2. We need to show that for any commutative square

$$\sigma = \begin{array}{ccc} S' & \xrightarrow{f} & T' \\ \downarrow & & \downarrow \\ S & \xrightarrow{\simeq} & T \end{array}$$

lying over a square

$$\begin{array}{ccc} \langle n \rangle & \xleftarrow{\cong} & \langle n \rangle \\ \uparrow & & \uparrow \\ \langle m \rangle & \xleftarrow{\cong} & \langle m \rangle \end{array}$$

in $\mathcal{F}\text{in}_{*}^{\text{op}}$ (where we have drawn the arrows as they are in $\mathcal{F}\text{in}_{*}$), σ is a pullback square if and only if f is p -cocartesian. But f is lying over an equivalence, so it is p -cocartesian if and only if it is an equivalence, which in turn is true if and only if the square in question is pullback.

□

Lemma 1.4.2. The morphism q satisfies the conditions of the Omnibus theorem.

Proof. We have to check the conditions of the Omnibus Theorem:

1. We have to show that for any isomorphism in $\mathcal{F}\text{in}_*^{\text{op}}$, there exists a q -cocartesian lift. But it suffices to find a q -cartesian lift, which is possible because q is a Cartesian fibration.
2. We need to show that for any commutative square in $\text{Map}(\Delta^1, \mathcal{S})_\times$ corresponding to a cube in \mathcal{S}_\times of the form

$$\begin{array}{ccccc}
 X' & \xrightarrow{F} & Y' & & \\
 \downarrow & \searrow \phi & \downarrow & \searrow \psi & \\
 & X & \xrightarrow{\cong} & Y & \\
 \downarrow & \downarrow & \downarrow & \downarrow & \\
 S' & \xrightarrow{f} & T' & & \\
 \downarrow & \searrow & \downarrow & \searrow & \\
 & S & \xrightarrow{\cong} & T &
 \end{array}$$

(where ϕ is p -cartesian) and lying over a square in $\mathcal{F}\text{in}_*^{\text{op}}$ of the form

$$\begin{array}{ccc}
 \langle n \rangle & \xleftarrow{\cong} & \langle n \rangle \\
 \uparrow & & \uparrow \\
 \langle m \rangle & \xleftarrow{\cong} & \langle m \rangle
 \end{array} ,$$

the morphism (F, f) is p -cocartesian if and only if the top and bottom squares are pullback and ψ is p -cocartesian. But this again follows because equivalences are p -cocartesian, and cocartesian morphisms are closed under composition.

□

Lemma 1.4.3. The map r satisfies the conditions of the Omnibus Theorem.

Proof.

1. It suffices to show that for each morphism $T \rightarrow S'$ in \mathcal{S}_\times lying over $\langle m \rangle \xleftarrow{\cong} \langle m \rangle$ and each $Y \rightarrow T$ in the fiber over $\langle m \rangle$, the square

$$\begin{array}{ccc}
 Y & \xlongequal{\quad} & Y \\
 \downarrow & & \downarrow \\
 T & \longrightarrow & S'
 \end{array}$$

is both p -cocartesian and p_{\dagger} -cocartesian, both of which follow from [Lemma 1.2.7](#).

2. We need to show that for any commutative square in $\text{Map}(\Delta^1, \mathcal{S})_{\times}$ corresponding to a cube in \mathcal{S}_{\times} of the form

$$\begin{array}{ccccc}
 X' & \xrightarrow{F} & Y' & & \\
 \downarrow & \searrow \simeq & \downarrow & \searrow \psi & \\
 & X & \xrightarrow{\simeq} & Y & \\
 \downarrow & \downarrow & \downarrow & \downarrow & \\
 S' & \xrightarrow{f} & T' & & \\
 \downarrow & \searrow & \downarrow & \searrow & \\
 & S & \xrightarrow{\quad} & T &
 \end{array}$$

where the bottom square is pullback, the morphism (F, f) is r -cocartesian if and only if the top square is pullback and ψ is an equivalence. This is clear.

□

This gives us the following commutative diagram.

$$\begin{array}{ccc}
 \text{Span}(\mathcal{M}, \mathcal{M}_{\dagger}, \mathcal{M}^{\dagger}) & \xrightarrow{\rho} & \text{Span}(\mathcal{P}, \mathcal{P}_{\dagger}, \mathcal{P}^{\dagger}) \\
 \searrow \varpi & & \swarrow \pi \\
 & \text{Span}(\mathcal{F}, \mathcal{F}_{\dagger}, \mathcal{F}^{\dagger}) &
 \end{array}$$

Recall that there is an equivalence $\text{Fin}_{*} \rightarrow \text{Span}((\mathcal{F}, \mathcal{F}_{\dagger}, \mathcal{F}^{\dagger}))$. Pulling back along this map gives