

1 Categories of spans

This is a rephrasing of [3].

1.1 Motivation

Given an ∞ -category \mathcal{C} , we would like to create an ∞ -category whose objects are spans in \mathcal{C} . That is, we would like to create a category whose objects are the objects of \mathcal{C} , whose morphisms are spans in \mathcal{C} , and whose higher simplices parametrize compositions of spans.

$$\{x\} \begin{matrix} \xleftarrow{\quad} \\ \xrightarrow{\quad} \end{matrix} \left\{ \begin{array}{ccc} & y & \\ x \swarrow & & \searrow \\ & x' & \end{array} \right\} \begin{matrix} \xleftarrow{\quad} \\ \xrightarrow{\quad} \end{matrix} \left\{ \begin{array}{ccccc} & & z & & \\ & \swarrow & \vee & \searrow & \\ x & \swarrow y & & \searrow y' & \\ & \searrow & x' & \swarrow & \\ & & & & x'' \end{array} \right\} \dots$$

Our first job is to formulate this in rigorous, combinatorial language.

Definition 1. Define a functor $T: \Delta \rightarrow \Delta$ sending

$$[n] \mapsto [n] \oplus [n]^{\text{op}} = \{0, \dots, n, \bar{n}, \dots, \bar{0}\}.$$

We can pull any simplicial set $K: \Delta^{\text{op}} \rightarrow \mathbf{Set}$ back by T , giving a simplicial set T^*K with n -simplices

$$(T^*K)_n = K([n] \oplus [n]^{\text{op}}).$$

Example 2. With $K = \Delta^1$, we have

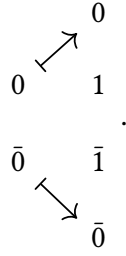
$$(T^*\Delta^1)_0 = \{\phi: \{0, \bar{0}\} \mapsto \{0, 1\} \mid \phi \text{ weakly monotone}\} = \{00, 01, 11\}$$

and

$$(T^*\Delta^1)_1 = \{0000, 0001, 0011, 0111, 1111\}.$$

1 Categories of spans

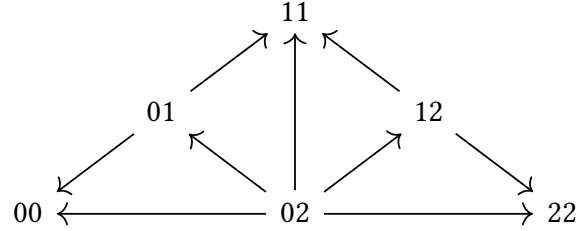
The map d_0 acts by pulling back by $T\partial_0$, which is given by the diagram



That is, the source of an edge $abcd$ is ad , and the target is bc . We can draw $T^*\Delta^1$ as follows.

$$00 \xleftarrow{0001} 01 \xrightarrow{0111} 11$$

Similarly, we can draw $T^*\Delta^2$ as follows.



We would like the span $00 \leftarrow 02 \rightarrow 22$ to correspond to the composition of the spans $00 \leftarrow 01 \rightarrow 11$ and $11 \leftarrow 12 \rightarrow 22$, but it is not immediately clear how to impose this condition in a general way. We turn instead to a different model.

Definition 3. Denote by \mathbb{Z}^n the poset with objects (i, j) , with $0 \leq i \leq j \leq n$, with

$$(i, j) \leq (i', j') \iff i \leq i' \text{ and } j' \leq j.$$

Since any map of totally ordered sets $\phi: [m] \rightarrow [n]$ induces a functor

$$\mathbb{Z}^m \rightarrow \mathbb{Z}^n, (i, j) \mapsto (\phi(i), \phi(j)),$$

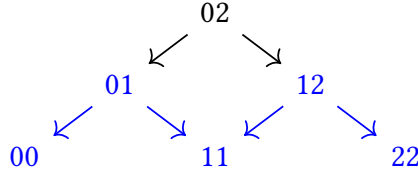
the categories \mathbb{Z}^n form a simplicial object $\mathbb{Z}^\bullet: \Delta \rightarrow \mathbf{Cat}$.

Similarly, let \mathbb{A}^n be the full subcategory of \mathbb{Z}^n on objects of the form

$$(i, j): j - i \leq 1.$$

There is an obvious inclusion $i: \mathbb{A}^\bullet \hookrightarrow \mathbb{Z}^\bullet$.

Example 4. We can draw \mathbb{Z}^2 as follows. The subcategory $\mathbb{A}^2 \hookrightarrow \mathbb{Z}^2$ is in blue.



As simplicial sets, $N(\mathbb{Z}^n) \cong T^*\Delta^n$. The k -simplices of \mathbb{Z}^n are maps $[k] \mapsto \mathbb{Z}^n$, i.e. strings $(i_0, j_0) \leq \dots \leq (i_k, j_k)$. This is a collection of integers such that

$$0 \leq i_k \leq \dots \leq i_0 \leq j_0 \leq \dots \leq j_k \leq n.$$

Providing such a string is equivalent to providing a map $[k]^{\text{op}} \oplus [k] \rightarrow [n]$.

Definition 5 (Cartesian). Let \mathcal{C} be an ∞ -category with finite limits. A functor $f: \mathbb{Z}^n \rightarrow \mathcal{C}$ is **Cartesian** if it is a right Kan extension of its restriction to \mathbb{A}^n .

For any ∞ -category \mathcal{C} , we might try to procede by defining a simplicial set

$$\text{Hom}^{\text{Cart}}(\mathbb{Z}^\bullet, \mathcal{C}): \Delta^{\text{op}} \rightarrow \mathbf{Set},$$

whose n -simplices are Cartesian diagrams $\mathbb{Z}^n \rightarrow \mathcal{C}$. However, it is not at all clear how to show that this is an infinity category. It turns out to be profitable to change models; we will be able to show that the functor

$$\mathbf{Fun}^{\text{Cart}}(\mathbb{Z}^\bullet, \mathcal{C})^\simeq; \quad \Delta^{\text{op}} \rightarrow \mathcal{S}$$

is a complete Segal space.

For any ∞ -category \mathcal{C} , we define a simplicial object

$$\mathbf{Fun}(\mathbb{Z}^\bullet, \mathbf{Cat}_\infty): \Delta^{\text{op}} \rightarrow \mathbf{Cat}_\infty. \quad (1.1)$$

Applying the relative nerve to [Equation 1.1](#) gives a coCartesian fibration

$$\pi: \overline{\mathbf{SPAN}}^+ \rightarrow \Delta^{\text{op}}.$$

The ∞ -category $\overline{\mathbf{SPAN}}^+$ has

- 0-simplices given by functors $\alpha: \mathbb{Z}^m \rightarrow \mathcal{C}$.
- 1-simplices given by pairs (ϕ, η) , where $\phi: [n] \rightarrow [m]$ and η is a 1-simplex

$$\phi^* \alpha \xrightarrow{\eta} \beta$$

in $\mathbf{Fun}(\mathbb{Z}^n, \mathcal{C})$.

• ...

Definition 6. We define SPAN^+ to be the full subcategory of $\overline{\text{SPAN}}^+$ on Cartesian functors.

Lemma 7. Let $\alpha: \mathbb{Z}^m \rightarrow \mathcal{C}$ be a Cartesian functor. Then for any $\phi: [n] \rightarrow [m]$, the functor $\phi^* \alpha$ given by the composition

$$\mathbb{Z}^n \xrightarrow{\phi} \mathbb{Z}^m \xrightarrow{\alpha} \mathcal{C}$$

is Cartesian.

Corollary 8. The restriction

$$\pi': \text{SPAN}^+ \rightarrow \Delta^{\text{op}}$$

is a coCartesian fibration.

Proof. We need to show that we can have a sufficient supply of π' -coCartesian morphisms. A morphism f in $\overline{\text{SPAN}}^+$ is π -coCartesian if and only if every horn Λ_i^n with $\Delta^{\{0,1\}} = f$ has a filler. Because $\text{SPAN}^+ \hookrightarrow \overline{\text{SPAN}}^+$ is a full subcategory inclusion, any \square

Definition 9 (category of spans). Let \mathcal{C} be an ∞ -category with finite limits. We define a category

$$\text{Span}(\mathcal{C})$$

by

1.2 Monoidal structures

1.2.1 1-categories of spans

Let \mathcal{C}^\otimes be a symmetric monoidal infinity category, which we model by a commutative monoid in Cat_∞ , i.e. a functor

$$\mathcal{C}^\otimes: N(\text{Fin}_*) \rightarrow \text{Cat}_\infty$$

such that

$$\mathcal{C}_{\langle n \rangle}^\otimes \xrightarrow{\Pi_i \rho_i} (\mathcal{C}_{\langle 1 \rangle}^\otimes)^n$$

is an equivalence for all n . Further suppose that

1. Each ∞ -category $\mathcal{C}_{\langle n \rangle}^\otimes$ admits finite limits; and
2. For each morphism $\rho: \langle m \rangle \rightarrow \langle n \rangle$ in Fin_* , the functor $\rho_*: \mathcal{C}_{\langle m \rangle}^\otimes \rightarrow \mathcal{C}_{\langle n \rangle}^\otimes$ preserves finite limits.

Definition 10. We will call such a symmetric monoidal ∞ -category *pullback-compatible*.

Note 11. By presenting a symmetric monoidal ∞ -category as a Cartesian fibration admitting relative finite limits, these conditions are satisfied automatically by [4, Prop. 4.3.1.10]. For example, for any ∞ -category \mathcal{C} with finite limits, the construction \mathcal{C}_\times in [2] yields such a Cartesian fibration.

Definition 12. Define a functor $F: \Delta^{\text{op}} \times \mathcal{F}\text{in}_* \rightarrow \mathbf{Kan}$ (where we are viewing $\Delta^{\text{op}} \times \mathcal{F}\text{in}_*$ as a discrete simplicially enriched category) by

$$F([m], \langle n \rangle) = \mathbf{Span}(\mathcal{C}_{\langle n \rangle}^\otimes)_m = \mathbf{Fun}^{\text{Cart}}(\mathbb{Z}^m, \mathcal{C}_{\langle n \rangle}^\otimes)^\simeq.$$

Taking the homotopy coherent nerve gives a functor $N(\Delta^{\text{op}} \times \mathcal{F}\text{in}_*) \rightarrow \mathcal{S}$. This is the adjunct to a functor

$$G: N(\mathcal{F}\text{in}_*) \rightarrow \mathbf{Fun}(N(\Delta^{\text{op}}), \mathcal{S}); \quad \langle n \rangle \mapsto \mathbf{Span}(\mathcal{C}_{\langle n \rangle}^\otimes).$$

One might worry that this construction is not well-defined. At the moment, it's seeming like one is smarter than the author.

Lemma 13. The functor F is well-defined.

Proof. The functor F is well-defined in the first slot by [3, Prop. 5.9], and well-defined in the second because if $f: \mathbb{Z}^n \rightarrow \mathcal{C}_{\langle m \rangle}^\otimes$ is Cartesian then so is $\rho_* \circ f$. To see this, we need to show that

$$(\rho_* \circ f)(i, j) \simeq \lim \left[\mathbb{A}_{/(i, j)}^m \rightarrow \mathbb{A}^m \xrightarrow{f} \mathcal{C}_{\langle m \rangle}^\otimes \xrightarrow{\rho_*} \mathcal{C}_{\langle n \rangle}^\otimes \right].$$

But ρ_* preserves limits, so this is equivalently

$$\lim \left[\mathbb{A}_{/(i, j)}^m \rightarrow \mathbb{A}^m \xrightarrow{f} \mathcal{C}_{\langle m \rangle}^\otimes \right] \xrightarrow{\rho_*} \mathcal{C}_{\langle n \rangle}^\otimes$$

which is the same as $(\rho_* \circ f)(i, j)$ because f is Cartesian. \square

Proposition 14. The functor G takes its values in \mathbf{CSS} , the ∞ -category of complete Segal spaces.

Proof. This is [3, Prop. 8.1] \square

Note 15. I'm not sure what the best way of writing this next proposition is. It sort of feels like I'm missing the point here.

Proposition 16. The functor G is a commutative monoid in complete Segal spaces, i.e. for each $n \geq 0$, the maps ρ_i give an equivalence

$$G(\langle n \rangle) \simeq G(\langle 1 \rangle)^n.$$

Proof. Limits in functor categories are computed pointwise. Therefore, it suffices to check this claim for fixed $[m]$. We have equivalences

$$\begin{aligned} G(\langle n \rangle)_m &\simeq \mathbf{Fun}^{\mathrm{Cart}}(\mathbb{Z}^m, \mathcal{C}_{\langle n \rangle}^{\otimes})^{\simeq} \\ &\simeq \mathbf{Fun}^{\mathrm{Cart}}(\mathbb{Z}^m, (\mathcal{C}_{\langle 1 \rangle}^{\otimes})^n)^{\simeq} \\ &\simeq \left(\mathbf{Fun}^{\mathrm{Cart}}(\mathbb{Z}^m, \mathcal{C}_{\langle 1 \rangle}^{\otimes}) \right)^n \\ &\simeq \left(\mathbf{Fun}^{\mathrm{Cart}}(\mathbb{Z}^m, \mathcal{C}_{\langle 1 \rangle}^{\otimes}) \right)^n \\ &\simeq G(\langle 1 \rangle)_m^n \end{aligned}$$

The second equivalence holds because \mathcal{C}^{\otimes} is a monoid, the third holds because the hom functor commutes with limits, and the fourth holds because $(-)^{\simeq}$ is a right adjoint and hence preserves limits. \square

I suspect that it will be more useful to cast everything in fibration-y language than commutative monoid in CSS language. However, we stop at this point to remark that our construction almost reproduces Toby's construction.

Using the equivalence $\mathrm{CSS} \rightarrow \mathrm{Cat}_{\infty}$ coming from forgetting higher simplices, we get a functor

$$\mathcal{F}\mathrm{in}_* \rightarrow \mathrm{CSS} \rightarrow \mathrm{Cat}_{\infty}; \quad \langle n \rangle \mapsto \mathrm{Hom}^{\mathrm{Cart}}(\mathbb{Z}^{\bullet}, \mathcal{C}_{\langle n \rangle}^{\otimes}),$$

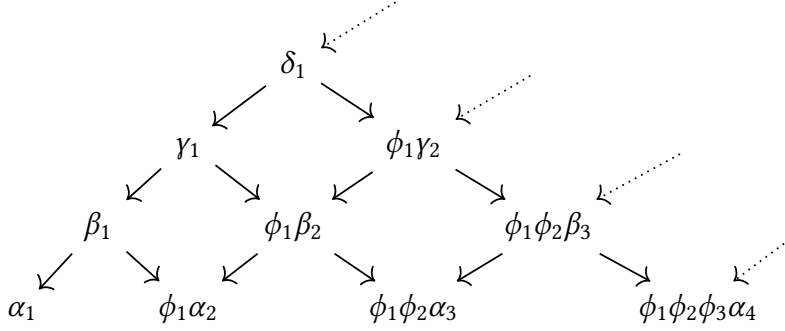
where $\mathrm{Hom}^{\mathrm{Cart}}$ denotes the set of Cartesian functors. Using the relative nerve we can unstraighten this to a Cartesian fibration

$$\chi \rightarrow \mathcal{F}\mathrm{in}_*^{\mathrm{op}}.$$

The k -simplices of this are pairs of an n -simplex in $N(\mathcal{F}\mathrm{in}_*^{\mathrm{op}})$ and an n -simplex F in $\mathrm{Span}(\mathcal{C}_{\langle n_1 \rangle}^{\otimes})$

$$\left(\langle n_0 \rangle \xleftarrow{\phi_1} \langle n_1 \rangle \xleftarrow{\phi_2} \cdots \xleftarrow{\phi_k} \langle n_k \rangle, \quad F: \mathbb{Z}^n \rightarrow \mathcal{C}_{\langle n_0 \rangle}^{\otimes} \right),$$

where F is a Cartesian functor of the form



and α_i, β_i , etc. are in $\mathcal{C}_{\langle n_i \rangle}^\otimes$. This is tantalizingly close to Toby's construction, the difference being that in my construction the tensor products are taken before we take spans.

1.2.2 Higher categories of spans

Let $\mathcal{C}^\otimes: \mathcal{F}\text{in}_* \rightarrow \mathbf{Cat}_\infty$ be a pullback-compatible symmetric monoidal ∞ -category. We can again define a functor

$$(\Delta^{\text{op}})^k \times \mathcal{F}\text{in}_* \rightarrow \mathcal{S}; \quad ([m_1], \dots, [m_k], \langle n \rangle) \mapsto \mathbf{Fun}^{\text{Cart}}(\Sigma^{m_1, \dots, m_k}, \mathcal{C}_{\langle n \rangle}^\otimes)$$

This is well-defined in the first slot by the reasoning in [3, Cor. 5.12], and in the second by exactly the same reasoning as above.

1.3 Push/pull and the Grothendieck construction

Let X be a space. The category of presheaves over X is the ∞ -category $\mathbf{Fun}(X^{\text{op}}, \mathcal{S})$. However, by the Grothendieck construction, the ∞ -category of functors $X^{\text{op}} \rightarrow \mathcal{S}$ is equivalent to the ∞ -category of spaces over X . We therefore make the following definition.

Definition 17 (category of presheaves). Let X be a space. The *category of presheaves* on X is the category $\mathcal{P}(X)$, given by the pullback

$$\begin{array}{ccc} \mathcal{P}(X) & \longrightarrow & \mathbf{Fun}(\Delta^1, \mathcal{S}) \\ \downarrow & & \downarrow \text{ev}_1 \\ \Delta^0 & \xrightarrow{X} & \mathcal{S} \end{array}$$

Note: fix ops.

Fix once and for all $f: \Delta^1 \rightarrow \mathcal{S}$ be a map of spaces $f: X \rightarrow Y$. This gives us a functor

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of ∞ -categories

$$f^*: \mathcal{P}(X) \rightarrow \mathcal{P}(Y)$$

between categories of presheaves via pullback. We can view this as a map $(\Delta^1)^{\text{op}} \rightarrow \mathbf{Cat}_\infty$. The Grothendieck construction again tells us that this can be thought of as a Cartesian fibration $p: \mathcal{M} \rightarrow \Delta^1$.

Definition 18. Define a morphism of simplicial sets $p: \mathcal{M} \rightarrow \Delta^1$ by the following pullback.

$$\begin{array}{ccc} \mathcal{M} & \longrightarrow & \mathbf{Fun}(\Delta^1, \mathcal{S}) \\ p \downarrow & & \downarrow \text{ev}_1 \\ \Delta^1 & \xrightarrow{f} & \mathcal{S} \end{array} .$$

Proposition 19. We have:

1. The map p is a Cartesian fibration, and the p -Cartesian morphisms in \mathcal{M} are pullback squares in \mathcal{S} .
2. The map p is a coCartesian fibration, and the p -coCartesian morphisms in \mathcal{M} are squares

$$\begin{array}{ccc} K & \xrightarrow{b} & S \\ \downarrow & & \downarrow \\ A & \xrightarrow{a} & A' \end{array} ,$$

in \mathcal{S} , where $a = \text{id}_X$, f , or id_Y , and b is an equivalence.

Proof. This is wrong!!!!

1. This follows from [4, Lem. 6.1.1.1], which shows that

$$\mathbf{Fun}(\Delta^1, \mathcal{S}) \rightarrow \mathcal{S}$$

is a Cartesian fibration whose Cartesian morphisms are pullback squares

$$\begin{array}{ccc} X' & \xrightarrow{f'} & Y' \\ \downarrow & & \downarrow g \\ X & \xrightarrow{f} & Y \end{array}$$

in \mathcal{S} . Here, we expand [4, Lem. 6.1.1.1], because the proof is funky.

Denote $\mathbf{Fun}(\Delta^1, \mathcal{S}) = \mathcal{O}_{\mathcal{S}}$. We will call an injective map $[k] \rightarrow \mathcal{P}$, for \mathcal{P} a poset, a *string of length k* in \mathcal{P} .

For every simplicial set K , let K^+ be the full simplicial subset of $(K \star \{x\} \star \{y\}) \times \Delta^1$ spanned by all of the vertices except $(x, 0)$. Two examples, which we will use later on, $(\Delta^n)^+$ and $(\partial \Delta^n)^+$, are as follows.

- The k -simplices of $(\Delta^n)^+$ are in bijective correspondence with those maps of posets

$$[k] \rightarrow ([n] \oplus \{x\} \oplus \{y\}) \times [1]$$

which never hit $(x, 0)$. The nondegenerate k -simplices, as usual, correspond to the *injective* maps of posets which never hit $(x, 0)$, i.e. strings of length k not containing $(x, 0)$.

- For $n \geq 1$, the nondegenerate k -simplices of $(\partial\Delta^n)^+$ are in bijective correspondence with strings

$$[k] \rightarrow ([n] \oplus \{x\} \oplus \{y\}) \times [1]$$

which never hit $(x, 0)$, and contain neither

$$(0, 0) \rightarrow (1, 0) \rightarrow \cdots \rightarrow (n, 0) \quad \text{nor} \quad (0, 1) \rightarrow \cdots \rightarrow (n, 1) \quad (1.2)$$

as contiguous substrings.

For $n = 0$, $(\partial\Delta^n)^+ \cong \emptyset^+ \cong \Lambda_2^2$.

Note that $(\partial\Delta^n)_k \cong (\Delta^n)_k$ for $k < n$, because no string of length less than n contain either of the strings in [Equation 1.2](#) as substrings.

Define a simplicial set \mathcal{C} by setting

$$\mathbf{Fun}(K, \mathcal{C}) = \{m: K^+ \rightarrow \mathcal{C} \mid m|_{(\{x\} \star \{y\}) \times \{1\}} = f, m|_{\{y\} \times \Delta^1} = g\}.$$

The square of inclusions

$$\begin{array}{ccc} (\Delta^n \star \{y\}) \times \{1\} & \hookrightarrow & (\Delta^n \star \{y\}) \times \Delta^1 \\ \downarrow & & \downarrow \\ (\Delta^n \star \{x\} \star \{y\}) \times \{1\} & \hookrightarrow & (\Delta^n)^+ \end{array}$$

gives a commuting square

$$\begin{array}{ccc} \mathcal{C} & \longrightarrow & (\mathcal{O}_{\mathcal{S}})_{/g} \\ \downarrow & & \downarrow \\ \mathcal{S}_{/f} & \longrightarrow & \mathcal{S}_{/Y} \end{array}$$

This in turn gives us a map

$$q: \mathcal{C} \rightarrow (\mathcal{O}_{\mathcal{S}})_{/g} \times_{\mathcal{S}_{/Y}} \mathcal{S}_{/f}.$$

We claim that q is a trivial fibration. To check this, we show that for any $n \geq 0$

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we can always find a lift ℓ as below.

$$\begin{array}{ccc} \partial\Delta^n & \xrightarrow{a} & \mathcal{C} \\ \downarrow & \nearrow \ell & \downarrow q \\ \Delta^n & \xrightarrow{b} & (\mathcal{O}_{\mathcal{S}})_{/g} \times_{\mathcal{S}_{/Y}} \mathcal{S}_{/f} \end{array}$$

The case $n = 0$ is somewhat special since $\partial\Delta^n = \emptyset$, but as none of the simplices of $(\Delta^0)^+$ have dimension greater than three one can find a solution simply by drawing pictures. From now on, we fix $n \geq 1$.

We can read the solid diagram as equivalently specifying a map

$$\tilde{a}: (\partial\Delta^n)^+ \rightarrow \mathcal{S}$$

and a map

$$\tilde{b}: (\Delta^n \star \{x\} \star \{y\}) \times \{1\} \coprod_{(\Delta^n \star \{y\}) \times \{1\}} (\Delta^n \star \{y\}) \times \Delta^1 \equiv K_n \rightarrow \mathcal{S}$$

which are compatible on their overlap

$$(\partial\Delta^n)^+ \cap K_n \subset (\Delta^n)^+.$$

and which map $(\{x\} \star \{y\}) \times \{1\}$ to f and $\{y\} \times \Delta^1$ to g .

We would like to produce from this data a filler

$$\tilde{\ell}: (\Delta^n)^+ \rightarrow \mathcal{S}.$$

Let us first see which nondegenerate simplices we are missing.

- The simplicial set $(\Delta^n)^+$ does not contain any nondegenerate simplices of dimension $n + 4$ or higher.
- There are two nondegenerate $(n+3)$ -simplices in $(\Delta^n)^+$ missing from $(\partial\Delta^n)^+$:

$$\sigma = \begin{array}{c} (0, 0) \\ \downarrow \\ (0, 1) \rightarrow (1, 1) \rightarrow \cdots \rightarrow (n, 1) \rightarrow (x, 1) \rightarrow (y, 1) \end{array}$$

and

$$\tau = \begin{array}{c} (0, 0) \rightarrow (1, 0) \rightarrow \cdots \rightarrow (n, 0) \\ \downarrow \\ (n, 1) \rightarrow (x, 1) \rightarrow (y, 1) \end{array}.$$

One finds that σ is contained in K_n , but τ is not.

- Apart from the faces of σ and τ , there is only one nondegenerate $(n + 2)$ -simplex in $(\Delta^n)^+$ not contained in $(\partial\Delta^n)^+$:

$$\chi = \begin{array}{ccccccc} (0, 0) & \rightarrow & (1, 0) & \rightarrow & \cdots & \rightarrow & (n, 0) & \xrightarrow{\quad} & (y, 0) \\ & & & & & & & & \downarrow \\ & & & & & & & & (y, 1) \end{array}$$

However, χ is contained in K_n .

- There are no nondegenerate simplices of $(\Delta^n)^+$ of dimension $n + 1$ or lower missing from both $(\partial\Delta^n)^+$ and K_n .

Therefore, the only data we have to fill in is the simplex $\tau \rightarrow \mathcal{S}$.

- The faces $d_0\tau, d_1\tau, \dots, d_n\tau$ are all contained in $(\partial\Delta^n)^+$.
- The face $d_{n+1}\tau$ is missing from both $(\partial\Delta^n)^+$ and K_n .
- The face $d_{n+2}\tau$ is contained in K_n .
- The face $d_{n+3}\tau$ is contained in K_n .

Thus, in order to find our lift ℓ we need only fill $\Lambda_{n+1}^{n+3} \hookrightarrow \Delta^{n+3}$, which we can do because \mathcal{S} is an ∞ -category.

2. This follows from the dual to [4, Cor. 2.4.7.12] with $f = \text{id}_{\mathcal{S}}$.

□

1.3.1 Barwick is Kafka for simplicial sets

Lemma 20. Let \mathcal{C} be an ∞ -category, and consider a solid diagram

$$\begin{array}{ccc} \Delta^n \star \Delta^1 & \xrightarrow{\quad} & \mathcal{C} \\ \text{id} \star \{1, 2\} \downarrow & \nearrow \text{dashed} & \\ \Delta^n \star \Delta^2 & & \end{array}$$

such that $\Delta^{\{1, 2\}} \hookrightarrow \Delta^2$ is mapped to an equivalence. Then for any $n \geq 0$, we can find a dashed lift.

Proof. Induction. Call the images of the vertices of Δ^2 x , y , and z .

For $n = 0$, we need to fill the following diagram.

PLACEHOLDER

First, fill $0 \rightarrow x \rightarrow y$. This is $\Lambda_2^2 \hookrightarrow \Delta^2$ with $\Delta^{\{1, 2\}}$ an equivalence. Then fill $0 \rightarrow x \rightarrow$

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$y \rightarrow z$, which is $\Lambda_2^3 \hookrightarrow \Delta^3$.

TBC

□

Lemma 21. Let \mathcal{C} be an infinity category, and let σ be the below square in \mathcal{C} with f an equivalence.

$$\begin{array}{ccc} x & \xrightarrow{f'} & y \\ \downarrow & & \downarrow \\ y' & \xrightarrow{f} & z \end{array}$$

Then σ is a pullback square if and only if f' is an equivalence.

Proof. First, suppose that f' is an equivalence. We need to show that $\mathcal{C}_{/\sigma} \hookrightarrow \mathcal{C}_{\sigma|\Lambda_2^2}$ is a trivial fibration.

Denote by K the simplicial subset of $\Delta^1 \times \Delta^1$ given by the nerve of the poset

$$\begin{array}{ccc} 0 & \longrightarrow & 1 \\ & \searrow & \downarrow \\ 1' & \longrightarrow & 2 \end{array}.$$

We first show that

$$\mathcal{C}_{/\sigma|K} \rightarrow \mathcal{C}_{/\sigma|\Lambda_2^2}$$

is a trivial fibration. It is sufficient to be able to solve lifting problems

$$\partial$$

First, suppose that σ is pullback. Consider the following Λ_2^2 horn associated to σ .

$$\begin{array}{ccc} & y' & \\ & \downarrow g & \\ x & \xrightarrow{f} & x' \end{array}$$

Let f^{-1} be any homotopy inverse to f , and pick a composite $f^{-1} \circ g$. By our previous work, the square

$$\tau = \begin{array}{ccc} y & \xrightarrow{\text{id}} & y \\ f^{-1} \circ g \downarrow & & \downarrow g \\ x' & \xrightarrow{f} & x \end{array}$$

is a pullback square in \mathcal{C} . Any other pullback square

$$\begin{array}{ccc} y \times_x x' & \xrightarrow{f'} & y \\ \downarrow & & \downarrow \\ x' & \longrightarrow & x \end{array}$$

factors through τ via an equivalence g , giving us in particular a 2-simplex

$$\begin{array}{ccc} & x' \times_x y & \\ f' \swarrow & & \searrow g \\ y & \xrightarrow{\text{id}} & y \end{array} .$$

But by 2/3, this means that f' is an equivalence.

□

Lemma 22. Denote $\mathbf{Fun}(\Delta^1, \mathcal{S})$ by $\mathcal{O}_{\mathcal{S}}$. The bicartesian fibration $p = \text{ev}_1: \mathcal{O}_{\mathcal{S}} \rightarrow \mathcal{S}$ satisfies the following conditions, which are conditions 12.2.1 and 12.2.2 of Barwick, with all morphisms ingressive and egressive.

1. For any morphism $g: X \rightarrow Y$ in \mathcal{S} and any object $K \rightarrow X$ in $\mathcal{O}_{\mathcal{S}}$, there exists a morphism

$$\begin{array}{ccc} K & \longrightarrow & S \\ \downarrow & & \downarrow \\ X & \longrightarrow & Y \end{array}$$

in $\mathcal{O}_{\mathcal{S}}$ covering f which is p -coCartesian.

2. Suppose σ is a commutative square in $\mathcal{O}_{\mathcal{S}}$, corresponding to a commutative cube Σ in \mathcal{S}

$$\begin{array}{ccccc} K' & \xrightarrow{F'} & S' & & \\ \downarrow & \searrow & \downarrow & \searrow & \\ & K & \xrightarrow{F} & S & \\ \downarrow & \downarrow & \downarrow & \downarrow & \\ X' & \xrightarrow{f'} & Y' & & \\ \downarrow & \searrow & \downarrow & \searrow & \\ & X & \xrightarrow{f} & Y & \end{array}$$

such that the bottom face of Σ is a pullback and F is an equivalence. Then the back face of Σ is p -coCartesian if and only if σ is a pullback in $\mathcal{O}_{\mathcal{S}}$.

Proof. 1. The map p is a bicartesian fibration.

2. The square σ is a pullback square in $\mathcal{O}_{\mathcal{S}}$ if and only if both the top and bottom

faces of Σ are pullbacks. The bottom face of Σ is pullback by assumption, and by [Lemma 21](#) the top face of Σ is a pullback if and only if F' is an equivalence. But F' is an equivalence if and only if the back face of Σ is p -coCartesian.

□

1.3.2 Lifting things we need to lift

Everything in this section is in [1].

Notation 23. We will distinguish integers expressed in binary notation with a subscript ‘2’. For example, $5 = 101_2$. We will often refer to the i th digit of a number N expressed in binary as d_i ; in this notation,

$$N = (d_1 \dots d_n)_2.$$

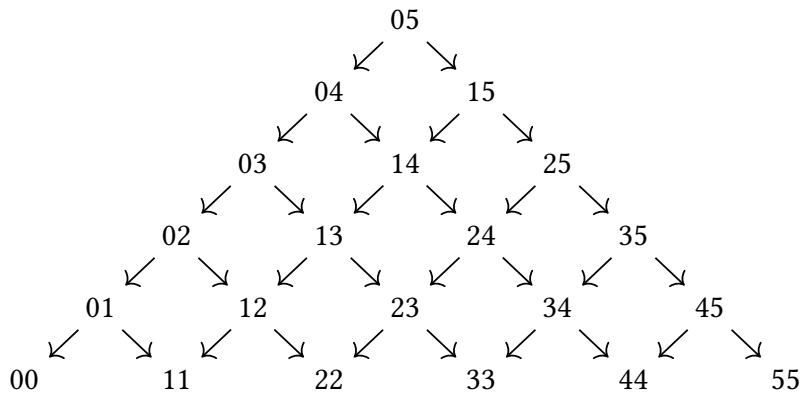
We are trying to find lifts of the following form. The maps f , g , and ℓ should be Segal.

$$\begin{array}{ccc} \mathrm{asd}(\Lambda_k^n) & \xrightarrow{f} & \mathcal{O}_S \\ \downarrow & \nearrow \ell & \downarrow \\ \mathrm{asd}(\Delta^n) & \xrightarrow{g} & \mathcal{S} \end{array}$$

For the entirety of this section, fix $n \geq 2$.

As we have seen in [Definition 3](#), we can think of $\mathrm{asd}(\Delta^n)$ as the nerve of the opposite of the poset of subintervals of $[n]$, which we have denoted $(I_n)^{\mathrm{op}}$. The poset $(I_n)^{\mathrm{op}}$ has objects ab , with $0 \leq a \leq b \leq n$, and $ab \leq a'b'$ if $a \leq a' \leq b' \leq b$.

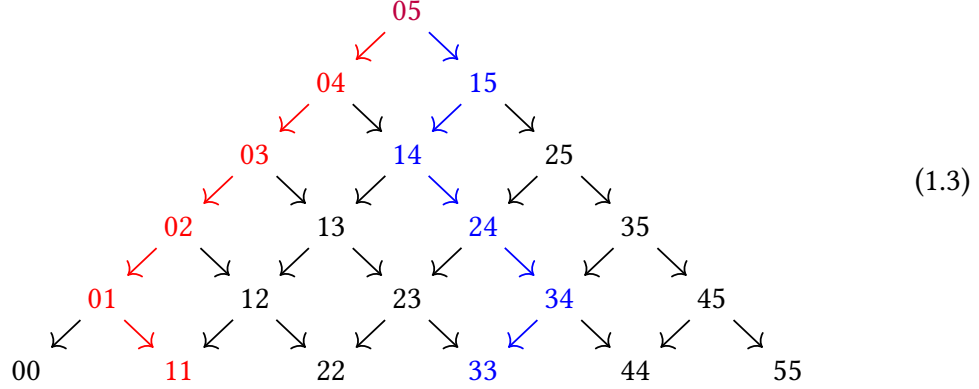
Example 24. The subdivision of the 5-simplex, $\mathrm{asd}(\Delta^5)$, is the nerve of the poset $(I_5)^{\mathrm{op}}$, which can be drawn as follows.



The nondegenerate n -simplices in $\mathrm{asd}(\Delta^n)$ are injective maps $[n] \rightarrow (I_n)^{\mathrm{op}}$, which we can think of as walks of length n starting at $0n$ and ending at ll for some $0 \leq l \leq n$.

These are in one-to-one correspondence with binary strings of length n , where the 0 or 1 in position i (read from left to right) corresponds to whether one goes left or right at the i th row.

Example 25. The nondegenerate 5-simplices of $\text{asd}(\Delta^5)$ corresponding to $N = 00001_2$ and $N = 10110_2$ are below, in red and blue respectively.



Binary strings of length n in turn correspond to numbers N with $0 \leq N < 2^n$. We will make frequent use of the bijection between numbers $0 \leq N < 2^n$ and nondegenerate simplices of $\text{asd}(\Delta^n)$. It will be helpful to have a notation for this bijection.

Notation 26. For any $0 \leq N < 2^n$, denote by $\sigma(N)$ the nondegenerate n -simplex of $\text{asd}(\Delta^n)$ corresponding to N . We will denote the i th binary digit of N , read left-to-right, by d_i ; that is, $N = (d_1 d_2 \dots d_n)_2$.

Example 27. The red and blue simplices in [Diagram 1.3](#) are $\sigma(00001_2)$ and $\sigma(10110_2)$ respectively.

To sum up, we may specify a nondegenerate n -simplex in $\text{asd}(\Delta^n)$ by providing any one of the following equivalent data.

- A walk of length n on the poset $(I_n)^{\text{op}}$ of the form

$$0n \rightarrow \dots \rightarrow a_i b_i \rightarrow \dots \rightarrow ll, \quad 0 \leq l \leq n.$$

- A binary string $(d_1 \dots d_n)_2$ of length n .
- An integer N with $0 \leq N < 2^n$.

Armed with this information, our plan of attack is as follows: starting with $\text{asd}(\Lambda_k^n)$, we will construct a filtration

$$\text{asd}(\Lambda_k^n) \subset \text{asd}(\Lambda_k^n) \cup \sigma(0) \subset \dots \subset \text{asd}(\Lambda_k^n) \cup \bigcup_{0 \leq N < 2^n} \sigma(N) = \text{asd}(\Delta^n). \quad (1.4)$$

It will be helpful to have a name for the N th step of this filtration.

Definition 28. Let $n \geq 1$, and $0 \leq k < n$. Define

$$P_0(k) = \text{asd}(\Lambda_k^n), \quad P_N(k) = \sigma(N) \cup P_{N-1}(k), \quad 0 \leq N < 2^n.$$

That is,

$$P_N(k) = \text{asd}(\Lambda_k^n) \cup \bigcup_{0 \leq N' < N} \sigma(N'). \quad (1.5)$$

We first need to figure out what each of these inclusions in Equation 1.4 looks like. More precisely, we must compute the intersections $\sigma(N) \cap P_N(k)$, $0 \leq N \leq 2^k - 1$.

From Equation 1.5 and the fact that $\Lambda_k^n = \bigcup_{j \neq k} d_j \Delta^n$, we have

$$\sigma(N) \cap P_N(k) = \overbrace{\left(\bigcup_{0 \leq K < N} \sigma(N) \cap \sigma(K) \right)}^{(a)} \cup \overbrace{\left(\bigcup_{j \neq k} \sigma(N) \cap \text{asd}(d_j \Delta^n) \right)}^{(b)}. \quad (1.6)$$

We can understand (a) and (b) individually. First, let's tackle (a). We need some terminology.

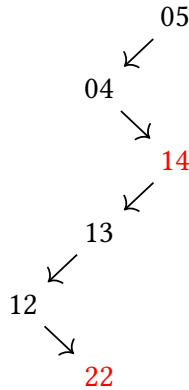
Definition 29. Fix an integer N with $0 \leq N < 2^n$, with N expressed in binary notation as $N = (d_1 \dots d_n)_2$. A **jut** of $\sigma(N)$ is an integer $1 \leq j \leq n$ such that one of the following conditions is satisfied:

- For $1 \leq j < n$, j is a jut of $\sigma(N)$ if $d_j = 1$ and $d_{j+1} = 0$.
- For $j = n$, n is a jut of $\sigma(N)$ if $d_n = 1$.

The set of all juts of $\sigma(N)$ is denoted $Z(N)$. By definition, $Z(N) \subset \{1, \dots, n\}$.

Note 30. The juts of $\sigma(N)$ are the places where the walk corresponding to N is going right, and then either heads left or ends. In terms of the binary representation $N = (d_1 \dots d_n)_2$, an integer $j < n$ is a jut of $\sigma(N)$ every time the string '10' appears at the j th position, i.e. if $d_j d_{j+1} = 10$. There is a jut at n every time the last digit of N is a 1.

Example 31. The juts of the simplex $\sigma(01011_2) \subset \text{asd}(\Delta^5)$ are 2 and 5, marked in red below.



Thus, $Z(01011_2) = \{2, 5\}$.

Knowledge of the juts of $\sigma(N)$ allows us to calculate (a).

Lemma 32. The expression (a) in Equation 1.6 can be written

$$\bigcup_{0 \leq N' < N} \sigma(N) \cap \sigma(N') = \bigcup_{j \in Z(N)} d_j \sigma(N),$$

where $Z(N) \subset \{1, \dots, n\}$ is the set of juts of $\sigma(N)$.

Proof. For any jut $j \in Z(N)$, define a number $N^j = (d'_1 \dots d'_n)_2$ as follows.

- If $j < n$, define d'_i to be

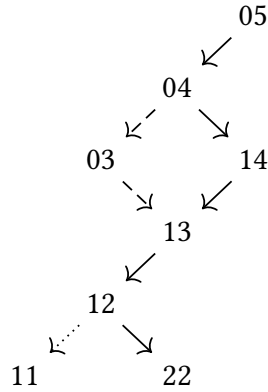
$$d'_i = \begin{cases} d_i, & i \neq j, j+1 \\ 0, & i = j \\ 1, & i = j+1 \end{cases}.$$

- If $j = n$, define d'_i to be

$$d'_i = \begin{cases} d_i, & i \neq n \\ 0, & i = n \end{cases}.$$

The walk corresponding to $\sigma(N^j)$ agrees with the walk corresponding to $\sigma(N)$, except at the j th position, where it goes left instead of right.

For $N = 01011_2$, $\sigma(N)$ is the solid walk below, $\sigma(N^2)$ is the dashed walk, and $\sigma(N^5)$ is the dotted walk.



Consider the poset of simplicial subsets of $\sigma(N)$ of the form

$$\sigma(N) \cap \sigma(K), \quad 0 \leq K < N.$$

The maximal elements of this poset are intersections $\sigma(N) \cap \sigma(N^j)$, where $j \in Z(N)$. Because $\sigma(N)$ and $\sigma(N^j)$ agree except at the j th vertex, the intersection $\sigma(N) \cap \sigma(N^j)$ is given by $d_j \sigma(N)$.

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Thus, we find that

$$\bigcup_{0 \leq N' < N} \sigma(N) \cap \sigma(N') = \bigcup_{j \in Z(N)} d_j \sigma(N).$$

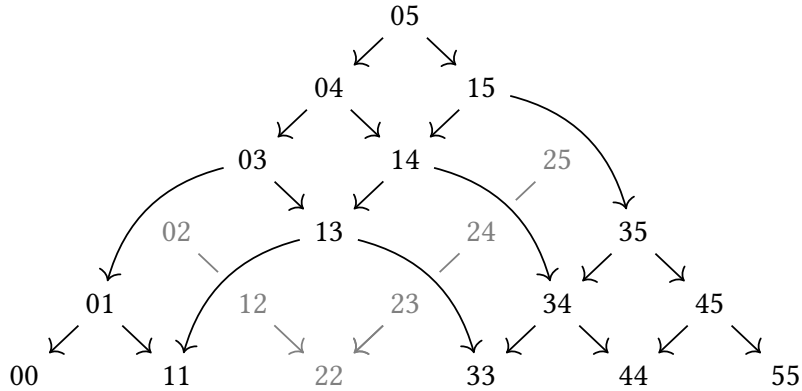
□

Now we turn our attention to (b). First, we should understand $\text{asd}(\Lambda_k^n) = \bigcup_{j \neq k} \text{asd}(d_j \Delta^n)$. To do this we must understand $\text{asd}(d_j \Delta^n)$. This is the nerve of the full subcategory of $(I_n)^{\text{op}}$ on objects ab where $a \neq j$ and $b \neq j$.

We will also pay attention to the full subcategory of $(I_n)^{\text{op}}$ on objects containing js .

Definition 33. Denote the full subcategory of $(I_n)^{\text{op}}$ on objects ab , where either $a = j$ or $b = j$, by V_j^n .

Example 34. The subdivision $\text{asd}(d_2 \Delta^5)$ is the nerve of the black diagram below. The category V_2^5 is pictured gray.



The i th face of any simplex $\sigma(N)$ corresponds to the walk on $(I_n)^{\text{op}}$ which skips the i th position of the walk corresponding to $\sigma(N)$. That is, if $\sigma(N)$ corresponds to the walk

$$a_0 b_0 \rightarrow \cdots \rightarrow a_n b_n,$$

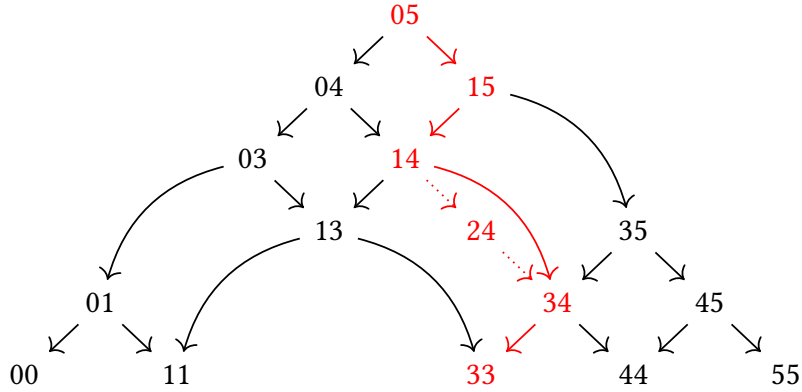
then $d_i \sigma(N)$ corresponds to the walk

$$a_0 b_0 \rightarrow \cdots \rightarrow a_{i-1} b_{i-1} \rightarrow a_{i+1} b_{i+1} \rightarrow \cdots \rightarrow a_n b_n.$$

Thus, the i th face of $\sigma(N)$ is contained in $\text{asd}(d_j \Delta^n)$ if and only if

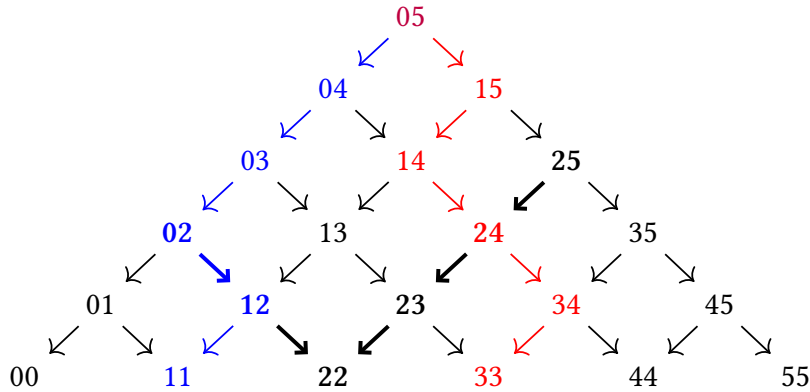
- $a_i = j$ or $b_i = j$ (or both); and
- $a_{i'} \neq j$ and $b_{i'} \neq j$ for $i' \neq i$.

Example 35. The simplex $d_3\sigma(10110_2)$ is contained in $\text{asd}(d_2\Delta^n)$.



The faces of $\sigma(N)$ contained in $\text{asd}(d_j\Delta^n)$ correspond to *crossings of $\sigma(N)$ at j* : places at which $\sigma(N)$ intersects V_j^n exactly once.

Example 36. The simplex $\sigma(10110_2)$, in red, has a crossing at 2. The simplex $\sigma(00010_2)$, in blue, does not. The subcategory $V_2^5 \subset (I_5)^{\text{op}}$ is in bold.



Definition 37. Let $0 \leq j < n$, and fix some $0 \leq N < 2^n$. Let the binary representation of N be $(d_1 \dots d_n)_2$, and the walk corresponding to $\sigma(N)$ be

$$a_0b_0 \rightarrow a_1b_1 \rightarrow \dots \rightarrow a_nb_n.$$

An integer $0 \leq i \leq n-1$ is said to be a **crossing of $\sigma(N)$ at j** if one of the following holds:

- For $j = 0$, i is a crossing at j if $d_1 = 1$.
- For $0 < j < n$, i is a crossing at j if $d_i = d_{i+1} = 0$ and $b_i = j$, or if $d_i = d_{i+1} = 1$ and $a_i = j$.

We have already essentially shown the following.

Lemma 38. Let $0 \leq j < n$, and let $0 \leq N < 2^n$. The simplex $\sigma(N)$ has a face in common with $\text{asd}(d_j\Delta^n)$ if and only if $\sigma(N)$ has a crossing at j .

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We can now understand the faces of $\sigma(N)$ contained in $\bigcup_{j \neq k} \sigma(N) \cap \text{asd}(d_j \Delta^n)$ as corresponding to *crossings of $\sigma(N)$ away from k* ; that is, crossings at some $j \neq k$. This amounts to the following.

Definition 39. Let $0 \leq k \leq n$, and let $0 \leq N < 2^n$, with binary representation $(d_1 \dots d_n)_2$ and walk

$$a_0 b_0 \rightarrow \dots \rightarrow a_n b_n.$$

An integer $0 < i < n$ is said to be a **crossing of $\sigma(N)$ away from k** if any of the following conditions are satisfied.

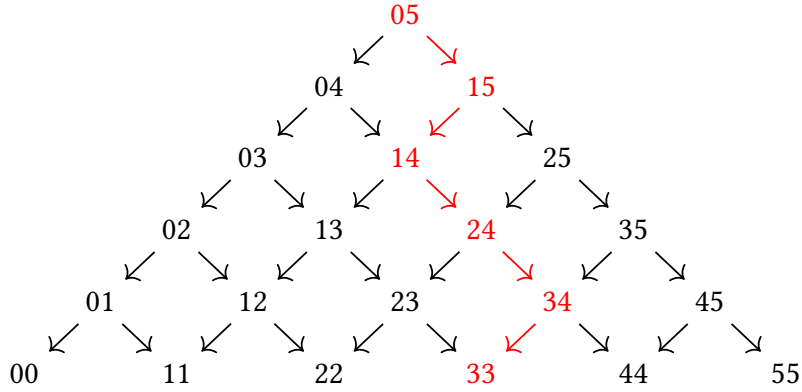
- Either $d_i = d_{i+1} = 1$ and $a_i \neq k$; or
- $d_i = d_{i+1} = 0$ and $b_i \neq k$.

The integer $i = 0$ is said to be a **crossing of $\sigma(N)$ away from k** if either of the following conditions are satisfied.

- $d_1 = 0$; or
- $d_1 = 1$ and $k \neq 0$.

We will denote the set of crossings of $\sigma(N)$ away from k by $X(N, k)$. By definition, $X(N, k) \subset \{0, \dots, n-1\}$.

Example 40. Let $N = 10110_2$, and consider the simplex $\sigma(N)$.



The 0th vertex of the simplex $\sigma(N)$ is a crossing at 0, and the 3rd vertex is a crossing at 2. None of the other vertices are crossings. Therefore:

- $X(10110_2, 0) = \{3\}$,
- $X(10110_2, 1) = \{0, 3\}$,
- $X(10110_2, 2) = \{0\}$,
- $X(10110_2, 3) = \{0, 3\}$, and
- $X(10110_2, 4) = \{0, 3\}$.

We are now ready to understand the intersection $\sigma(N) \cap P_N(k)$.

Proposition 41. Let $0 \leq k < N$, and let $0 \leq N < 2^n$. The intersection $\sigma(N) \cap P_N(k)$ is given by the union of faces

$$\bigcup_{i \in Z(N) \cup X(N, k)} d_i \sigma(N).$$

Proof. Denote the walk corresponding to N by

$$a_0 b_0 \rightarrow \cdots \rightarrow a_n b_n.$$

Recall from Equation 1.6 the decomposition $\sigma(N) \cap P_N(k) = (a) \cup (b)$. We have seen in Lemma 32 that (a) is given by $\bigcup_{i \in Z(n)} d_i \sigma(N)$, and in Lemma 38 that the faces of $\sigma(N)$ contained in (b) are of the form $d_i \sigma(N)$, where $i \in X(N, k)$. Therefore, we will be done if we can show that (b) contains only faces, apart from lower-dimensional simplices which are also contained in (a) . More explicitly, we need to show that any subsimplex of $\sigma(N)$ of codimension 2 or greater which

- is contained in (b) , and
- is not contained in a higher-dimensional subsimplex of $\sigma(N)$ which is contained in (b) ,

must be contained in (a) .

For any subsimplex τ of $\sigma(N)$ meeting the conditions above, there must exist a $j \neq k$ such that $\sigma(N)$ intersects V_j^n more than once; the subsimplex τ then has as its vertices the vertices s of $\sigma(N)$ such that $a_s \neq j$ and $b_s \neq j$. If τ has dimension m , then $\sigma(N)$ and V_j^n must intersect exactly m times. If a walk on $(I_n)^{\text{op}}$ leaves V_j^n it cannot come back, so for any simplex which intersects V_j^n in $m < 1$ places, the intersections are contiguous; that is, there exists some i such that for each i' with $i \leq i' < i + m$, either $a_{i'} = j$ or $b_{i'} = j$. It is easy to convince oneself that in this situation either i or $i + m$ must be a jut of $\sigma(N)$.¹ Without loss of generality, suppose i is a jut of $\sigma(N)$. Then $\tau \subset \sigma(N^i) \cap \sigma(N)$, which is contained in (a) . \square

Notation 42. Let T be a finite totally ordered set, and let $S \subset T$ be a proper subset. Define

$$\Lambda_S^T = \bigcup_{i \notin S} d_i \Delta^n.$$

Take $T = [n]$. We will write Λ_S^n instead of $\Lambda_S^{[n]}$. Interpreting $S \subset T$ as a set of vertices of Δ^n , Λ_S^n is the union of all faces of Δ^n which contain every vertex in S . This is because $d_i \Delta^n$ contains S if and only if $S \subseteq [n] \setminus \{i\}$, which in turn is true if and only if $i \notin S$.

Example 43. For $T = [m]$ and $S = \{i\}$ a singleton, $\Lambda_S^T = \Lambda_i^m$, the i th horn of Δ^n .

¹Apart from the trivial case of $\sigma(0)$'s intersection with V_0^n .

Definition 44. Fix some $0 \leq k < n$, and let $0 \leq N < 2^n$. The **exceptional vertices** of $\sigma(N)$ are the subset

$$E(N, k) = [m] \setminus (Z(N) \cup X(N, k)).$$

To put it another way, a vertex of $\sigma(N)$ is exceptional if it is neither a jut nor a crossing away from k . In our new notation, [Proposition 41](#) tells us that $\sigma(N) \cap P_N(k) = \Lambda_{E(N, k)}^n$.

We are computing the intersections $\sigma(N) \cap P_N(k)$ in the hope that we can show that $P_N(k) \hookrightarrow P_{N+1}(k)$ is inner anodyne. We now have some idea when this might be so. We have a pushout square

$$\begin{array}{ccc} \Lambda_{E(N, k)}^n & \hookrightarrow & \Delta^n \\ \downarrow & & \downarrow \\ P_N(k) & \hookrightarrow & P_{N+1}(k) \end{array} .$$

If $\Lambda_{E(N, k)}^n \hookrightarrow \Delta^n$ is inner anodyne, then $P_N(k) \hookrightarrow P_{N+1}(k)$ will be as well. Our next goal will be to establish a sufficient condition for $\Lambda_{E(N, k)}^n \hookrightarrow \Delta^n$ to be inner anodyne.

Lemma 45. Let T be a finite, totally ordered set, and $S \subset T$ a nonempty proper subset. A sufficient condition for $\Lambda_S^T \hookrightarrow \Delta^T$ to be inner anodyne as follows:

- (*) There exist elements $a < s < b$ in T with $s \in S$ and $a, b \notin S$.

It will be useful to denote the elements of S by \star , and the elements of $T \setminus S$ by \circ . We can then draw the inclusion $S \subset T$ as a word on the letters \star and \circ . Then [Lemma 45](#) tells us that $\Lambda_S^T \hookrightarrow \Delta^T$ is inner anodyne if our drawing of $S \subset T$ is of the form

$$\dots \circ \dots \star \dots \circ \dots .$$

Example 46. With $T = [4]$ and $S = \{0, 2\}$, we can draw $S \subset T$ as

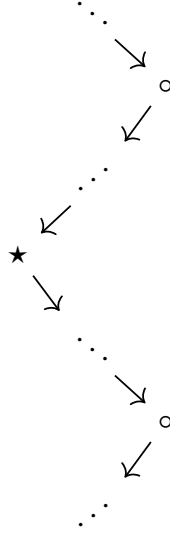
$$\star \circ \star \circ \circ .$$

[Lemma 45](#) guarantees us that the inclusion $\Lambda_{\{0, 2\}}^4 \hookrightarrow \Delta^4$ is inner anodyne.

Definition 47. We will call a simplex $\sigma(N)$ **exceptional** if $E(N, k) \subset [n]$ does not satisfy [Condition \(*\)](#).

In Barwick's words, exceptional simplices are indeed exceptional. Any exceptional simplex cannot have more than one jut, since between any two juts there is a vertex which

is neither a jut nor a crossing away from k .

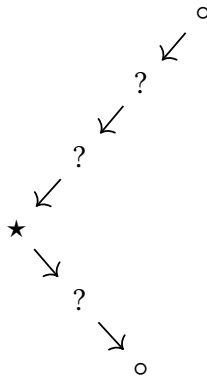


As follows easily from [Note 30](#), the only N such that $\sigma(N)$ contains at most one jut are of the form

- $N_r = \overbrace{1 \cdots 1}^r \overbrace{0 \cdots 0}^{n-r}$, for $0 \leq r \leq n$; or
- $N'_r = \overbrace{0 \cdots 0}^r \overbrace{1 \cdots 1}^{n-r}$, for $0 < r < n$.

In fact, we can easily see that all of the simplices $\sigma(N'_r)$ are not exceptional, since for any k and r , 0 is a crossing away from k of $\sigma(N'_r)$, n is a jut, and r is an exceptional vertex. This situation is pictured in [Example 48](#).

Example 48. For $n = 5$, $\sigma(N'_3)$ is pictured below. Exceptional vertices are marked with a '★' and juts and crossings away from k are marked with a 'o'. The vertices marked '?' could be either exceptional or crossings away from k depending on k . However, the information we have is already enough to tell us that $\sigma(N'_3)$ is not exceptional.



Lemma 49. For $k = 0$, each simplex $\sigma(N_r)$ is exceptional, and we have

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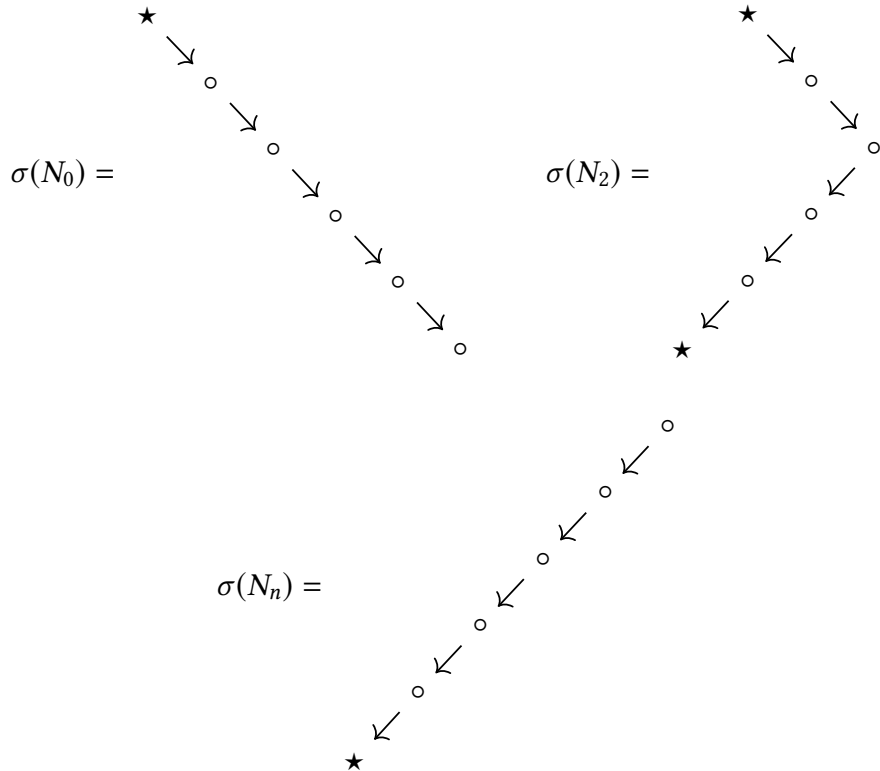
- $E(0, N_0) = \{n\}$;
- $E(0, N_r) = \{0, n\}$, for $0 \leq r \leq n$; and
- $E(0, N_n) = \{0\}$.

For $0 < k < n$, the simplices N_{k-1} and N_k are exceptional, and we have

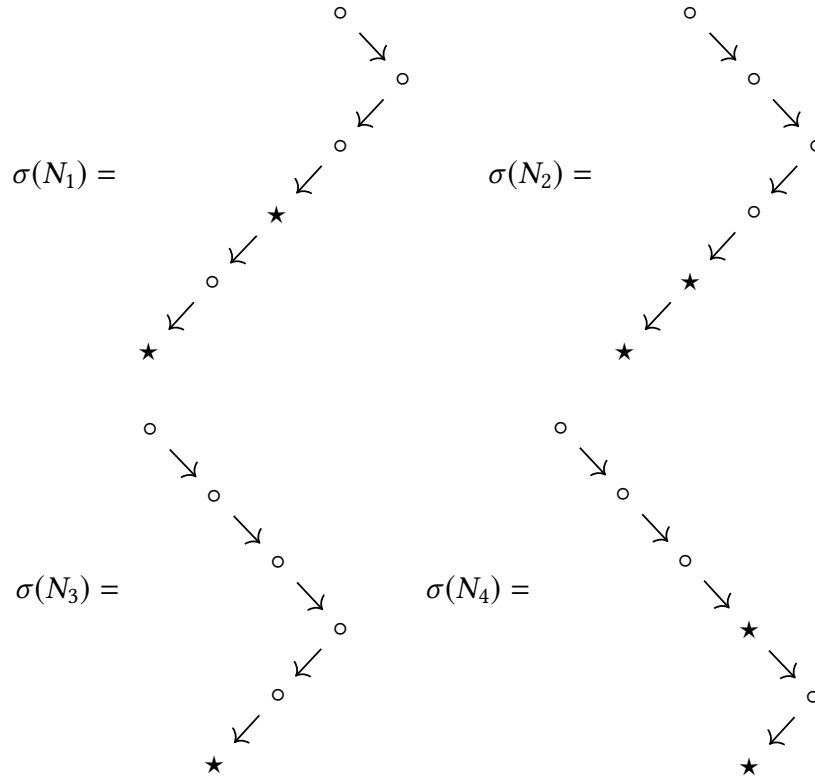
- $E(k, N_{k-1}) = \{n-1, n\}$; and
- $E(k, N_k) = \{n\}$.

We do not prove this statement explicitly, but the example below should make clear why it is true.

Example 50. Let $n = 5$ and $k = 0$. The simplices $\sigma(N_0)$, $\sigma(N_2)$, and $\sigma(N_n)$ are pictured below.



For $k = 3$, $\sigma(N_1)$, $\sigma(N_2)$, $\sigma(N_3)$, and $\sigma(N_4)$ are pictured below.



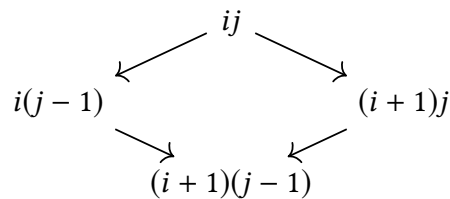
Recall our goal: we are trying to find Segal lifts ℓ as below, where f and g are Segal.

$$\begin{array}{ccc}
 \mathrm{asd}(\Lambda_k^n) & \xrightarrow{f} & \mathcal{O}_S \\
 \downarrow & \nearrow \ell & \downarrow \\
 \mathrm{asd}(\Delta^n) & \xrightarrow{g} & \mathcal{S}
 \end{array} \tag{1.7}$$

It turns out that we needn't worry ourselves that a lift might not be Segal; in most situations, any lift we can construct will be Segal automatically. Morally, this is because we are taking a limit over a very simple category, and demanding that $\mathrm{asd}(\Delta^n) \rightarrow \mathcal{C}$ be Segal is equivalent to demanding that each square

Lemma 51. Let $n \geq 3$, $k \neq 0$, and let ℓ be any lift as in [Diagram 1.7](#). Then ℓ is automatically Segal.

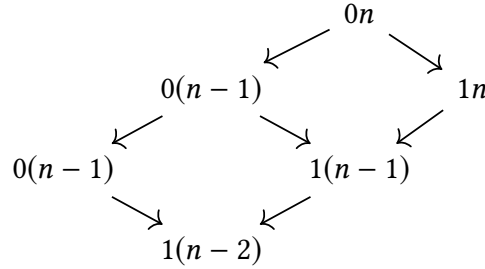
Proof. A lift $\ell : \mathrm{asd}(\Delta^n) \rightarrow \mathcal{O}_S$ is Segal if and only if the image of each of the squares



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is pullback, for $0 \leq i < n$, $0 < j \leq n$.

The image of each square of the above form except the top square ($i = 0, j = n$) is contained in at least one of $\text{asd}(d_0\Delta^n)$ and $\text{asd}(d_n\Delta^n)$, and is hence pullback by assumption. Thus, we need only verify that either the top square is already in some $\text{asd}(d_j\Delta^n)$, and hence is pullback by assumption, or that any way of filling the top square is guaranteed to be pullback. We have fixed some k , so either $k \neq 1$ or $k \neq n - 1$. If $k \neq n - 1$, then the outer square of



is contained in $\text{asd}(d_{n-1}\Delta^n) \subset \text{asd}(\Lambda_k^n)$, and is hence pullback, and the lower-left square is contained in $\text{asd}(d_n(\Delta^n)) \subset \text{asd}(\Lambda_k^n)$, and hence is also pullback. Thus, the pasting lemma implies that any way of filling in the top square will be pullback. Similar reasoning works in the case $k \neq 1$. \square

1.3.3 Lifting things we need to lift the right way

Warm up:

Some classical knowledge: There is a Quillen equivalence

$$p_1^* : \mathbf{Set}_\Delta^{\text{Joyal}} \longleftrightarrow \mathbf{Set}_{\Delta^2}^{\text{CSS}} : i_1^*.$$

This means that the counit

$$p_1^* i_1^* \mathcal{C} \rightarrow \mathcal{C}$$

is a weak equivalence when evaluated on a fibrant object.

Definition 52. We say that a map $\mathcal{C} \rightarrow \mathcal{D}$ of complete Segal spaces is a **inner fibration** if the following square is homotopy pullback in the Kan model structure.

$$\begin{array}{ccc} \mathcal{C}_2 & \longrightarrow & \mathcal{C}_1 \times_{\mathcal{C}_0} \mathcal{C}_1 \\ \downarrow & & \downarrow \\ \mathcal{D}_2 & \longrightarrow & \mathcal{D}_1 \times_{\mathcal{D}_0} \mathcal{D}_1 \end{array}$$

Proposition 53. Let \mathcal{C} and \mathcal{D} be quasicategories with pullbacks, and let $\pi : \mathcal{C} \rightarrow \mathcal{D}$ be an inner fibration admitting relative pullbacks. Then $\mathbf{Span}(\mathcal{C}) \rightarrow \mathbf{Span}(\mathcal{D})$ is an inner fibration of complete Segal spaces.

Proof. The square □

Definition 54. Let $\pi: \mathcal{C} \rightarrow \mathcal{D}$ be a map of complete Segal spaces. We say that a morphism $f: x \rightarrow y$ in \mathcal{C} is *p -coCartesian* if the following square is homotopy pullback.

$$\begin{array}{ccc} \mathcal{C}_2 \times_{\mathcal{C}_{\{0,1\}}} \{f\} & \longrightarrow & \mathcal{C}_1 \times_{\mathcal{C}_{\{0\}}} \{x\} \\ \downarrow & & \downarrow \\ \mathcal{D}_2 \times_{\mathcal{D}_{\{0,1\}}} \{\pi f\} & \longrightarrow & \mathcal{D}_1 \times_{\mathcal{D}_{\{0\}}} \{\pi x\} \end{array}$$

Proposition 55. Let $\pi: \mathcal{C} \rightarrow \mathcal{D}$ be a bicartesian fibration admitting relative pullbacks with the following property:

For any square

$$\sigma = \begin{array}{ccc} x & \xrightarrow{f} & y \\ \downarrow & & \downarrow \\ x' & \xrightarrow{f'} & y' \end{array}$$

in \mathcal{C} , such that f' is p -Cartesian and $\pi(\sigma)$ is pullback in \mathcal{D} , the following are equivalent.

- The morphism f is π -Cartesian.
- The square σ is pullback.

Definition 56. Let $\pi: \mathcal{C} \rightarrow \mathcal{D}$ be a Reedy fibration of complete segal spaces.

Proposition 57. Let $\mathcal{C} \rightarrow \mathcal{D}$ be an inner fibration admitting relative pullbacks. Then $\text{Span}(\mathcal{C}) \rightarrow \text{Span}(\mathcal{D})$ is an inner fibration.