

## 1 Horn filling via left Kan extensions

In this section, we show that a simplex  $\Delta_b^2 \rightarrow \mathbb{Cat}_\infty$  which is left Kan is also globally left Kan. This will amount to showing that a simplex  $\tau: \Delta_b^2 \rightarrow \mathbb{Cat}_\infty$  of the form

$$\begin{array}{ccc} X & \xrightarrow{F} & \mathcal{C} \\ & \searrow f \quad \Downarrow \eta \quad \nearrow f_! F & \\ & Y & \end{array},$$

where  $\mathcal{C}$  is cocomplete,  $X$  is small,  $f_!F$  is a left Kan extension of  $F$  along  $f$ , and  $\eta: F \Rightarrow f^*f_!F$  is the unit map, has the property that for any  $n \geq 2$ , any solid filling problem

$$\begin{array}{ccc}
\Delta_b^{\{0,1,n+1\}} & & \\
\downarrow & \searrow \tau & \\
(\Lambda_0^{n+1})_b & \longrightarrow & \mathbb{C}at_\infty \\
\downarrow & \nearrow \text{dashed} & \\
\Delta_b^{n+1} & & 
\end{array} \tag{1}$$

has a dashed solution.

Let us examine this in the case  $n = 2$ . In this case we have the data of categories and functors

A commutative diagram with objects  $X$ ,  $Y$ ,  $Z$ , and  $\mathcal{C}$ . The diagram consists of the following morphisms:
 

- $f: X \rightarrow Y$  (top-left arrow)
- $g: Y \rightarrow Z$  (middle-left arrow)
- $h: X \rightarrow Z$  (bottom-left arrow)
- $F: X \rightarrow \mathcal{C}$  (top-right arrow)
- $f_!F: Y \rightarrow \mathcal{C}$  (middle-right arrow)
- $G: Z \rightarrow \mathcal{C}$  (bottom-right arrow)

 The diagram is commutative, meaning that the composition of morphisms along any path between two objects is the same. For example,  $h = g \circ f$  and  $G \circ f_!F = F$ .

together with natural transformations  $\eta: F \Rightarrow f_!F \circ f$ ,  $\beta: h \Rightarrow g \circ f$ , and  $\delta: F \Rightarrow G \circ h$ . We need to produce a natural transformation  $\alpha: f_!F \Rightarrow G \circ g$  and a filling of the full 3-simplex, which is the data of a homomotopy-commutative diagram

$$\begin{array}{ccc} F & \xrightarrow{\eta} & f_! F \circ f \\ \delta \downarrow & & \downarrow \alpha f \\ G \circ h & \xrightarrow{G\beta} & G \circ g \circ f \end{array}$$

in  $\text{Fun}(X, \mathcal{C})$ .

We can rephrase this as follows. We are given the data

$$\begin{array}{ccc}
 \overbrace{\quad}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\quad}^{\text{in Fun}(Y, \mathcal{C})} \\
 F \xrightarrow{\eta} f_! F \circ f & & f_! F \\
 \delta \downarrow & & \\
 G \circ h \xrightarrow{G\beta} G \circ g \circ f & & G \circ g
 \end{array}$$

where the left-hand diagram is a map  $LC^2 \rightarrow \text{Fun}(X, C)$ , where  $LC^2$  is the simplicial subset of the boundary of  $\Delta^1 \times \Delta^1$  except the right-hand face, and the right-hand diagram is a map  $\partial\Delta^1 \rightarrow \text{Fun}(Y, C)$ . We need to construct a filler  $\partial\Delta^1 \hookrightarrow \Delta^1$  on the right, and from it a filler  $LC^2 \hookrightarrow \Delta^1 \times \Delta^1$  on the left. We do this in the following sequence of steps.

1. We first can fill the lower-left half of the diagram on the left simply by taking the composition  $G\beta \circ \delta$ . Doing this, we are left with the filling problem

$$\begin{array}{ccc}
 \overbrace{\quad}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\quad}^{\text{in Fun}(Y, \mathcal{C})} \\
 F \xrightarrow{\eta} f_! F \circ f & & f_! F \\
 \searrow & & \\
 G \circ g \circ f & & G \circ g
 \end{array}$$

2. We note that we can extend the diagram on the right to a diagram which is adjunct to the diagram on the left under the adjunction  $f_! \dashv f^*$ :

$$\begin{array}{ccc}
 \overbrace{\quad}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\quad}^{\text{in Fun}(Y, \mathcal{C})} \\
 F \xrightarrow{\eta} f_! F \circ f & & f_! F \xrightarrow{\text{id}} f_! F \\
 \searrow & & \searrow \\
 G \circ g \circ f & & G \circ g
 \end{array}$$

3. The diagram on the right has an obvious filler, which is adjunct to a filler on the left.

$$\begin{array}{ccc}
 \overbrace{\quad}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\quad}^{\text{in Fun}(Y, \mathcal{C})} \\
 F \xrightarrow{\eta} f_! F \circ f & & f_! F \xrightarrow{\text{id}} f_! F \\
 \searrow \quad \downarrow \alpha f & & \searrow \quad \downarrow \alpha \\
 G \circ g \circ f & & G \circ g
 \end{array}$$

The lifting problems which we have to solve in [Equation 1](#) for  $n > 2$  amount to replacing

$\Delta^1 \times \Delta^1$  by  $(\Delta^1)^n$ , etc; the basic process remains unchanged, but the combinatorics involved in filling the necessary cubes becomes more involved. In [Subsection 1.1](#) we explain the combinatorics of filling cubes relative to their boundaries. In [Subsection 1.2](#), we give a formalization the concept of adjunct data, and provide a means of for filling partial data to total adjunct data. In [Subsection 1.3](#), we show that we can always solve the lifting problems of [Equation 1](#).

## 1.1 Filling cubes relative to their boundary

In this section, we give the combinatorics of filling the  $n$ -cube  $(\Delta^1)^n$  relative to its boundary. Our first step is to understand the nondegenerate simplices of the  $n$ -cube. Our main result is [Proposition 1.1.19](#), which writes the inclusion  $LC^n \hookrightarrow C^n$  of part of the boundary of the  $n$ -cube into the full cube as a composition of an inner anodyne map and a marked anodyne map.

**Definition 1.1.1.** The  $n$ -*cube* is the simplicial set  $C^n := (\Delta^1)^n$ .

*Note 1.1.2.* We will consider the the factors  $\Delta^1$  of  $C^n$  to be ordered, so that we can speak about the first factor, the second factor, etc.

**Definition 1.1.3.** Denote by  $a^i: \Delta^n \rightarrow \Delta^1$  the map which sends  $\Delta^{\{0, \dots, i-1\}}$  to  $\{0\}$ , and  $\Delta^{\{i, \dots, n\}}$  to  $\{1\}$ .

The superscript of  $a^i$  counts how many vertices of  $\Delta^n$  are sent to  $\Delta^{\{0\}} \subset \Delta^1$ .

We would now like to understand which  $n$ -simplices in  $C^n$  are nondegenerate. Every nondegenerate simplex  $\Delta^n \rightarrow C^n$  can be specified by giving a walk along the edges of  $C^n$  starting at  $(0, \dots, 0)$  and ending at  $(1, \dots, 1)$ , and any simplex specified in this way is nondegenerate. We now use this to define a bijection between  $S_n$ , the symmetric group on the set  $\{1, \dots, n\}$ , and the nondegenerate simplices  $\Delta^n \rightarrow C^n$  as follows.

**Definition 1.1.4.** For any permutation  $\tau: \{1, \dots, n\} \rightarrow \{1, \dots, n\}$ , we define an  $n$ -simplex

$$\phi(\tau): \Delta^n \rightarrow C^n; \quad \phi(\tau)_i = a^{\tau(i)}.$$

That is,

$$\phi(\tau) = (a^{\tau(1)}, \dots, a^{\tau(n)}).$$

*Notation 1.1.5.* We will denote any permutatation  $\tau: \{1, \dots, n\} \rightarrow \{1, \dots, n\}$  by the corresponding  $n$ -tuple  $(\tau(1), \dots, \tau(n))$ . Thus, the identity permutation is  $(1, \dots, n)$ , and the permutation  $\gamma_{1,2}$  which swaps 1 and 2 is  $(2, 1, 3, \dots, n)$ .

Given a permutation  $\tau \in S_n$  corresponding to an  $n$ -simplex

$$\phi(\tau) = (a^{\tau(1)}, \dots, a^{\tau(n)}): \Delta^n \rightarrow C^n,$$

the  $i$ th face of  $\phi(\tau)$  is a nondegenerate  $(n-1)$ -simplex in  $C^n$ , i.e. a map  $\Delta^{n-1} \rightarrow C^n$ . We calculate this as follows: for any  $a^i: \Delta^n \rightarrow \Delta^1$  and any face map  $\partial_j: \Delta^{n-1} \rightarrow \Delta^n$ , we note that

$$\partial_j^* a^i = \begin{cases} a^{i-1}, & j < i \\ a^i, & j \geq i \end{cases},$$

where by minor abuse of notation we denote the map  $\Delta^{n-1} \rightarrow \Delta^1$  sending  $\Delta^{\{0, \dots, i-1\}}$  to  $\{0\}$  and the rest to  $\{1\}$  also by  $a^i$ . We then have that

$$d_i \tau = d_i(a^{\tau(1)}, \dots, a^{\tau(n)}) := (\partial_i^* a^{\tau(1)}, \dots, \partial_i^* a^{\tau(n)}).$$

**Example 1.1.6.** Consider the 3-simplex  $(a^1, a^2, a^3)$  in  $C^3$ . This has spine

$$(0, 0, 0) \rightarrow (1, 0, 0) \rightarrow (1, 1, 0) \rightarrow (1, 1, 1),$$

as is easy to see quickly:

- The function  $\Delta^3 \rightarrow \Delta^1$  corresponding to the first coordinate is  $a^1$ , so the first coordinate of the first vertex is 0, and the first coordinate of the rest of the vertices are 1.
- The function  $\Delta^3 \rightarrow \Delta^1$  corresponding to the second coordinate is  $a^2$ , so the second coordinates of the first two points are equal to zero, and the second coordinate of the remaining vertices is 1.
- The function  $\Delta^3 \rightarrow \Delta^1$  corresponding to the third coordinate is  $a^3$ , so the third coordinates of the first three points is equal to zero, and the second coordinate of the final vertex is equal to 1.

Using [Equation 1.1](#), we then calculate that

$$d_2(a^1, a^2, a^3) = (a^1, a^2, a^2).$$

This corresponds to the 2-simplex in  $C^n$  with spine

$$(0, 0, 0) \rightarrow (1, 0, 0) \rightarrow (1, 1, 1).$$

Our next task is to understand the boundary of the  $n$ -cube. We can view the boundary of the cube  $C^n$  as the union of its faces.

**Definition 1.1.7.** The *boundary of the  $n$ -cube* is the simplicial subset

$$\partial C^n := \bigcup_{i=1}^n \bigcup_{j=0}^1 \Delta^{(1)} \times \cdots \times \{j\} \times \cdots \times \Delta^{(n)} \subset C^n. \quad (2)$$

The following is easy to see.

**Proposition 1.1.8.** An  $(n-1)$ -simplex

$$\gamma = (\gamma_1, \dots, \gamma_n) : \Delta^{n-1} \rightarrow C^n$$

lies entirely within the face

$$\Delta^{(1)} \times \cdots \times \{0\} \times \cdots \times \Delta^{(n)}$$

if and only if  $\gamma_i = a^n = \text{const}_0$ , and to the face

$$\Delta^{(1)} \times \cdots \times \{1\} \times \cdots \times \Delta^{(n)}$$

if and only if  $\gamma_i = a^0 = \text{const}_1$ .

**Definition 1.1.9.** The *left box of the  $n$ -cube*, denoted  $LC^n$ , is the simplicial subset of  $C^n$  given by the union of all of the faces in Equation 2 except for

$$\Delta^1 \times \cdots \times \Delta^1 \times \Delta^{\{1\}}.$$

Note that drawing the box with the final coordinate going from left to right, the right face is open here; the terminology is chosen to match with ‘left horn’.

**Example 1.1.10.** The simplicial subset  $LC^2 \subset C^2$  can be drawn as follows.

$$\begin{array}{ccc} (0, 0) & \longrightarrow & (0, 1) \\ \downarrow & & \\ (1, 0) & \longrightarrow & (1, 1) \end{array}$$

Note that in  $LC^2$ , the morphisms  $(0, 0) \rightarrow (1, 0) \rightarrow (1, 1)$  form an inner horn, which we can fill by pushing out along an inner horn inclusion. Our main goal in this section is to show that this is generically true: given a left cube  $LC^n$ , we can fill much of  $C^n$  using pushouts along inner horn inclusions. To this end, it will be helpful to know how each nondegenerate simplex in  $C^n$  intersects  $LC^n$ .

**Lemma 1.1.11.** Let  $\phi: \Delta^n \rightarrow C^n$  be a nondegenerate  $n$ -simplex corresponding to the permutation  $\tau \in S_n$ . We have the following.

1. The zeroth face  $d_0\phi$  belongs to  $LC^n$  if and only if  $\tau(n) \neq 1$ .
2. For  $0 < i < n$ , the face  $d_i\phi$  never belongs to  $LC^n$ .
3. The  $n$ th face  $d_n\phi$  always belongs to  $LC^n$ .

*Proof.* 1. We have

$$d_0\phi = (a^{\tau(0)-1}, \dots, a^{\tau(n)-1})$$

since  $\tau(i) > 0$  for all  $i$ . For  $j = \tau^{-1}(1)$ , we have that  $\tau(j) = 1$ , so the  $j$ th entry of  $d_0\phi$  is  $a^0$ . Thus, [Proposition 1.1.8](#) guarantees that  $d_0\phi$  is contained in the face

$$\Delta^{(1)} \times \dots \times \{1\}^{(j)} \times \dots \times \Delta^{(n)}.$$

This face belongs to  $LC^n$  except when  $j = n$ .

2. In this case, the superscript of each entry of  $d_i\phi$  is between 1 and  $n-1$ , hence not equal to  $n$  or 0.
3. In this case,

$$d_n(a^{\tau(1)}, \dots, a^{\tau(n)}) = (a^{\tau(1)}, \dots, a^{\tau(n)})$$

since  $n \geq \tau(i)$  for all  $1 \leq i \leq n$ . Thus, the  $\tau^{-1}(n) = j$ th entry is  $a^n$ , so  $d_n\phi$  belongs to the face

$$\Delta^{(1)} \times \dots \times \{0\}^{(j)} \times \dots \times \Delta^{(n)}.$$

□

This is promising: it means that for many simplices in  $C^n$  which we want to fill relative to  $LC^n$ , we already have the data of the first and last faces. The following lemma will allow us to take advantage of this.

*Notation 1.1.12.* For any subset  $T \subseteq [n]$ , write

$$\Lambda_T^n := \bigcup_{t \in T} d_t \Delta^n \subset \Delta^n.$$

**Lemma 1.1.13.** For any proper subset  $T \subset [n]$  containing 0 and  $n$ , the inclusion  $\Lambda_T^n \hookrightarrow \Delta^n$  is inner anodyne.

*Proof.* Induction. For  $n = 2$ ,  $T$  must be equal to  $\{0, 2\}$ , so  $\Lambda_T^2 = \Lambda_1^2$ .

Assume the result holds for  $n-1$ , and let  $T \subset [n]$  be a proper subset containing 0 and  $n$ . Using the inductive step, we can fill all but one of the faces  $d_1\Delta^n, \dots, d_{n-1}\Delta^n$ . The result follows  $\square$

We can use [Lemma 1.1.13](#) to show that we can fill much of  $C^n$  as a sequence of inner anodyne pushouts, but to do this we need to pick an order in which to fill our simplices. We do this as follows.

**Definition 1.1.14.** We define a total order on  $S_n$  as follows. For any two permutations  $\tau, \tau' \in S_n$ , we say that  $\tau < \tau'$  if there exists  $k \in [n]$  such that the following conditions are satisfied.

- For all  $i < k$ , we have that  $\tau^{-1}(i) = \tau'^{-1}(i)$ .
- We have that  $\tau^{-1}(k) < \tau'^{-1}(k)$ .

We then define a total order on the nondegenerate simplices of  $C^n$  by saying that  $\phi(\tau) < \phi(\tau')$  if and only if  $\tau < \tau'$ .

Note that this is *not* the lexicographic order on  $S_n$ ; instead, we have that  $\tau < \tau'$  if and only if  $\tau^{-1}$  is less than  $\tau'^{-1}$  under the lexicographic order.

**Example 1.1.15.** The elements of the permutation group  $S_3$  have the order

$$(1, 2, 3) < (1, 3, 2) < (2, 1, 3) < (3, 1, 2) < (2, 3, 1) < (3, 2, 1).$$

We use this ordering to define our filtration.

*Notation 1.1.16.* For  $\tau \in S_n$ , we mean by

$$\bigcup_{\tau' \in S_n}^{\tau}$$

the union over all  $\tau' \in S_n$  for which  $\tau' < \tau$  with respect to the ordering defined above. Note the strict inequality.

We would like to show that each step of the filtration

$$\begin{aligned} LC^n &\hookrightarrow LC^n \cup \phi(0, \dots, n) \\ &\hookrightarrow \dots \\ &\hookrightarrow LC^n \cup \bigcup_{\tau \in S_n}^{(2, \dots, n, 1)} \phi(\tau) \end{aligned} \tag{3}$$

is inner anodyne. To do this, it suffices to show that for each  $\tau < (2, \dots, n, 1)$ , the intersection

$$B_\tau = \left( LC^n \cup \bigcup_{\tau' \in S_n} \phi(\tau') \right) \cap \phi(\tau) \quad (4)$$

is of the form  $\Lambda_T^n$  for some proper  $T \subseteq [n]$  containing 0 and  $n$ .

**Proposition 1.1.17.** Each inclusion in Equation 3 is inner anodyne.

*Proof.* The first inclusion

$$LC^n \hookrightarrow LC^n \cup \phi(0, \dots, n)$$

is inner anodyne because by Lemma 1.1.11 the intersection  $LC^n \cap \phi(0, \dots, n)$  is of the form  $d_0\Delta^n \cup d_n\Delta^n$ .

Each subsequent intersection contains the faces  $d_0\Delta^n$  and  $d_n\Delta^n$  (again by Lemma 1.1.11), so by Lemma 1.1.13, it suffices to show that for each  $\tau$  under consideration, there is at least one face not shared with any previous  $\tau'$ .

To this end, fix  $0 < i < n$ , and consider  $\tau \in S_n$  such that  $\tau \neq (0, \dots, n)$ . We consider the face  $d_i\phi(\tau)$ . Let  $j = \tau^{-1}(i)$  and  $j' = \tau^{-1}(i+1)$ . Then  $d_i a^{\tau(j)} = d_i a^{\tau(j')}$ , so  $d_i\phi(\tau) = d_i\phi(\tau \circ \gamma_{jj'})$ , where  $\gamma_{jj'} \in S_n$  is the permutation swapping  $j$  and  $j'$ . Thus, the face  $d_i\sigma(\tau)$  is equal to the face  $d_i\sigma(\tau \circ \gamma_{jj'})$ , which is already contained in the union under consideration if and only if  $\tau \circ \gamma_{jj'} < \tau$ . This in turn is true if and only if  $j < j'$ .

Assume that every face of  $\phi(\tau)$  is contained in the union. Then

$$\tau^{-1}(1) < \tau^{-1}(2) < \dots < \tau^{-1}(n),$$

which implies that  $\tau = (1, \dots, n)$ . This is a contradiction. Thus, each boundary inclusion under consideration is of the form  $\Lambda_T^n \hookrightarrow \Delta^n$ , for  $T \subset [n]$  a proper subset containing 0 and  $n$ , and is thus inner anodyne.  $\square$

Each nondegenerate simplex of  $C^n$  which we have yet to fill is of the form  $\phi(\tau)$ , where  $\tau(n) = 1$ . Thus, each of their spines begins

$$(0, \dots, 0, 0) \rightarrow (0, \dots, 0, 1) \rightarrow \dots,$$

and then remains confined to the face  $\Delta^1 \times \dots \times \Delta^1 \times \{1\}$ . This implies that the part of  $LC^n$  yet to be filled is of the form  $\Delta^0 * C^{n-1}$ , and the boundary of this with respect to which we must do the filling is of the form  $\Delta^0 * \partial C^{n-1}$ .

Our next task is to show that we can also perform this filling, under the assumption that  $(0, \dots, 0, 0) \rightarrow (0, \dots, 0, 1)$  is marked. We do not give the details of this proof, as it is quite sim-



ilar to the proof of [Proposition 1.1.17](#). One continues to fill simplices in the order prescribed in [Definition 1.1.14](#), and shows that each of these inclusions is marked anodyne. The main tool is the following lemma, proved using the same inductive argument as [Lemma 1.1.13](#).

**Lemma 1.1.18.** For any subset  $T \subset [n]$  containing 1 and  $n$ , and *not* containing 0, the inclusion

$$(\Lambda_T^n, \mathcal{E}) \hookrightarrow (\Delta^n, \mathcal{F})$$

is (cocartesian-)marked anodyne, where by  $\mathcal{F}$  we mean the set of degenerate 1-simplices together with  $\Delta^{\{0,1\}}$ , and by  $\mathcal{E}$  we mean the restriction of this marking to  $\Lambda_T^n$ .

We collect our major results from this section in the following proposition.

**Proposition 1.1.19.** Denote by  $J^n \subseteq C^n$  the simplicial subset spanned by those nondegenerate  $n$ -simplices whose spines do not begin  $(0, \dots, 0, 0) \rightarrow (0, \dots, 0, 1)$ . We can write the filling  $LC^n \hookrightarrow C^n$  as a composition

$$LC^n \xrightarrow{i} J^n \xrightarrow{j} C^n,$$

where  $i$  is inner anodyne, and  $j$  fits into a pushout square

$$\begin{array}{ccc} \Delta^0 * \partial C^{n-1} & \hookrightarrow & \Delta^0 * C^{n-1} \\ \downarrow & & \downarrow \\ J^n & \xrightarrow{\quad} & C^n \end{array},$$

where the top inclusion underlies a marked anodyne morphism

$$(\Delta^0 * \partial C^{n-1}, \mathcal{G}) \hookrightarrow (\Delta^0 * C^{n-1}, \mathcal{G}),$$

where the marking  $\mathcal{G}$  contains all degenerate morphisms together with the morphism  $(0, \dots, 0) \rightarrow (0, \dots, 1)$ .

## 1.2 Adjunct data

In this section, we give a formalization of the the notion of adjunct data. Our main result is [Proposition 1.2.6](#), which shows that under certain conditions, we can use data on one side of an adjunction to solve lifting problems on the other side.

**Definition 1.2.1.** Let  $s: K \rightarrow \Delta^1$  be a map of simplicial sets, and let  $p: \mathcal{M} \rightarrow \Delta^1$  be a bicartesian fibration associated to the adjunction

$$f: \mathcal{C} \longleftrightarrow \mathcal{D} : g$$

via equivalences  $h_0: \mathcal{C} \rightarrow \mathcal{M}_0$  and  $h_1: \mathcal{D} \rightarrow \mathcal{M}_1$ . We say that a map  $\alpha: K \rightarrow \mathcal{C}$  is **adjunct** to a map  $\tilde{\alpha}: K \rightarrow \mathcal{D}$  relative to  $s$ , and equivalently that  $\tilde{\alpha}$  is adjunct to  $\alpha$  relative to  $s$ , if there exists a map  $A: K \times \Delta^1 \rightarrow \mathcal{M}$  such that the diagram

$$\begin{array}{ccc} K \times \Delta^1 & \xrightarrow{A} & \mathcal{M} \\ & \searrow \text{pr}_{\Delta^1} & \downarrow p \\ & & \Delta^1 \end{array}$$

commutes, and such that the following conditions are satisfied.

1. The restriction  $A|_{K \times \{0\}} = h_0 \circ \alpha$ .
2. The restriction  $A|_{K \times \{1\}} = h_1 \circ \tilde{\alpha}$ .
3. For each vertex  $k \in K$  in the fiber  $s^{-1}\{0\}$ , the image of the edge  $\{k\} \times \Delta^1 \rightarrow \mathcal{M}$  is  $p$ -cocartesian.
4. For each vertex  $k' \in K$  in the fiber  $s^{-1}\{1\}$ , the image of the edge  $\{k'\} \times \Delta^1 \rightarrow \mathcal{M}$  is  $p$ -cartesian.

We often leave the equivalences  $h_0$  and  $h_1$  implicit to lighten notation.

In the next examples, fix an adjunction

$$f: \mathcal{C} \longleftrightarrow \mathcal{D} : g$$

corresponding to a bicartesian fibration  $p: \mathcal{M} \rightarrow \Delta^1$ .

**Example 1.2.2.** Any morphism in  $\mathcal{D}$  of the form  $fC \rightarrow D$  is adjunct to some morphism in  $\mathcal{C}$  of the form  $C \rightarrow gD$ , witnessed by any square

$$\begin{array}{ccc} C & \xrightarrow{a} & fC \\ \downarrow & & \downarrow \\ gD & \xrightarrow{b} & D \end{array}$$

in  $\mathcal{M}$  where  $a$  is  $p$ -cocartesian and  $b$  is  $p$ -cartesian. This corresponds to the map  $s = \text{id}: \Delta^1 \rightarrow \Delta^1$ .

**Example 1.2.3.** Pick some object  $D \in \mathcal{D}$ , and consider the identity morphism  $\text{id}: gD \rightarrow gD$  in  $\mathcal{C}$ . This morphism is adjunct to the component of the unit map  $\eta_D: D \rightarrow fgD$  relative to  $s = \text{id}: \Delta^1 \rightarrow \Delta^1$ .

The next example will follow from [Proposition 1.2.6](#).

**Example 1.2.4.** Any diagram  $K * K' \rightarrow \mathcal{D}$  such that  $K$  is in the image of  $f$  is adjunct to some diagram  $K * K' \rightarrow \mathcal{C}$  such that  $K'$  is in the image of  $g$ .

**Lemma 1.2.5.** The inclusion of marked simplicial sets

$$\left( \Delta^n \times \Delta^{\{0\}} \coprod_{\Delta^n \times \Delta^{\{0\}}} \partial \Delta^n \times \Delta^{\{0\}}, \mathcal{E} \right) \hookrightarrow (\Delta^n \times \Delta^1, \mathcal{F}) \quad (5)$$

where the marking  $\mathcal{F}$  is the flat marking together with the edge  $\Delta^{\{0\}} \times \Delta^1$ , and  $\mathcal{E}$  is the restriction of this marking, is marked (cocartesian) anodyne.

The next proposition shows the power of adjunctions in lifting problems: given data on either side of an adjunction, we can fill in adjunct data on the other side.

**Proposition 1.2.6.** Let  $\mathcal{M} \rightarrow \Delta^1$ ,  $\mathcal{C}$ ,  $\mathcal{D}$ ,  $f$ ,  $g$ ,  $h_0$  and  $h_1$  be as in [Definition 1.2.1](#), and let

$$\begin{array}{ccc} K' & \xrightarrow{i} & K \\ & \searrow_{s \circ i} & \swarrow_s \\ & \Delta^1 & \end{array},$$

be a commuting triangle of simplicial sets, where  $i$  is a cofibration such that  $i|_{\{1\}}$  is an isomorphism. Let  $\tilde{\alpha}': K' \rightarrow \mathcal{D}$  and  $\alpha: K \rightarrow \mathcal{C}$  be maps, and denote  $\alpha' = \alpha \circ i$  relative to  $s \circ i$ . Suppose that  $\tilde{\alpha}'$  is adjunct to  $\alpha'$  relative to  $s \circ i$ . Then there exists a dashed extension

$$\begin{array}{ccc} K' & \xrightarrow{\alpha'} & \mathcal{C} \\ i \downarrow & \nearrow \alpha & \\ K & & \end{array} \rightsquigarrow \begin{array}{ccc} K' & \xrightarrow{\tilde{\alpha}'} & \mathcal{D} \\ i \downarrow & \nearrow \tilde{\alpha} & \\ K & & \end{array} \quad (6)$$

such that  $\tilde{\alpha}$  is adjunct to  $\alpha$  relative to  $s$ .

*Proof.* Pick some commutative diagram

$$\begin{array}{ccc} K' \times \Delta^1 & \xrightarrow{A} & \mathcal{M} \\ & \searrow_{\text{pr}_{\Delta^1}} & \downarrow \\ & & \Delta^1 \end{array}$$

displaying  $\tilde{\alpha}'$  and  $\alpha'$  as adjunct. We further consider the map  $h_0 \circ \alpha: K \rightarrow \mathcal{M}_0$ . From these two pieces of data we can construct the solid commutative square of cocartesian-marked

simplicial sets

$$\begin{array}{ccc} \left( K \times \Delta^{\{0\}} \amalg_{K' \times \Delta^{\{0\}}} K' \times \Delta^1 \right)^{\heartsuit} & \xrightarrow{\quad} & \mathcal{M}^{\sharp} \\ \downarrow & \nearrow \ell & \downarrow \\ (K \times \Delta^1)^{\heartsuit} & \xrightarrow{\quad} & (\Delta^1)^{\sharp} \end{array},$$

where  $(K \times \Delta^1)^{\heartsuit}$  is the marked simplicial set where the only nondenerate morphisms marked are those of the form  $\{k\} \times \Delta^1$ , for  $k \in K|_{s^{-1}\{0\}}$ , and the  $\heartsuit$ -marked pushout-product denotes the restriction of this marking to the pushout-product. Starting from  $K'$  and adjoining one simplex of  $K$  at a time, we can write the left-hand map in the above square as a transfinite composition of pushouts of morphisms of the form given in Equation 5 (by assumption each simplex  $\sigma$  we adjoin is not completely contained in  $K'$ , so the initial vertex of  $\sigma$  must lie in the fiber of  $s$  over 0). Each of these is marked anodyne, so the left-hand map is as well. Thus, a dashed lift  $\ell$  exists, giving us a map  $h_1^{-1} \circ \ell|_{K \times \Delta^{\{1\}}} : K \rightarrow \mathcal{D}$ . This, together with the composition

$$K' \times \Delta^1 \xrightarrow{\tilde{\alpha}' \times \text{id}} \mathcal{D} \times \Delta^1 \xrightarrow{c} \mathcal{D}$$

where  $c : h_1^{-1} \circ h_1 \rightarrow \text{id}_{\mathcal{D}}$  is a natural isomorphism, gives a solid diagram

$$\begin{array}{ccc} K \times \Delta^{\{0\}} \amalg_{K' \times \Delta^{\{0\}}} K' \times (\Delta^1)^{\sharp} & \xrightarrow{\quad} & \mathcal{D}^{\sharp} \\ \downarrow & \nearrow m & \\ K \times (\Delta^1)^{\sharp} & & \end{array},$$

where  $\mathcal{D}^{\sharp}$  is the marked simplicial set whose underlying simplicial set is  $\mathcal{D}$  and where all equivalences are marked. This admits a dashed lift, and restricting  $m|_{K \times \Delta^{\{1\}}}$  gives us a map  $\tilde{\alpha} : K \rightarrow \mathcal{D}$  as in Equation 6. However, we still need to display this as adjunct to  $\alpha$ : the lift  $\ell$  has  $\ell|_{K \times \Delta^{\{0\}}} = h_0 \circ \alpha$ , but only  $\ell|_{K \times \Delta^{\{1\}}} \simeq h_1 \circ \tilde{\alpha}$ , not equality as we require.

Our last task is to remedy this. Denote the spine of the 3-simplex by

$$I_3 := \Delta^{\{0,1\}} \amalg_{\Delta^{\{1\}}} \Delta^{\{1,2\}} \amalg_{\Delta^{\{2\}}} \Delta^{\{2,3\}} \hookrightarrow \Delta^3.$$

We build a map  $I_3 \times K \xrightarrow{b} \mathcal{M}$  which:

- On  $\Delta^{\{0,1\}} \times K$  is given by the lift  $\ell$ .
- On  $\Delta^{\{1,2\}} \times K$  is given by the composition

$$K \times \Delta^1 \xrightarrow{\ell|_{K \times \Delta^{\{1\}}} \times \text{id}} \mathcal{M}_1 \times \Delta^1 \xrightarrow{H} \mathcal{M}_1 \hookrightarrow \mathcal{M},$$

where  $H$  is a natural isomorphism  $\text{id}_{\mathcal{M}_1} \rightarrow h_1 \circ h_1^{-1}$ .

- On  $\Delta^{\{2,3\}} \times K$  is given by the natural isomorphism

$$K \times \Delta^1 \xrightarrow{m} \mathcal{D} \xrightarrow{h_1} \mathcal{M}_1 \hookrightarrow \mathcal{M}.$$

The map  $I_3 \hookrightarrow \Delta^3$  is inner anodyne, so there exists a map  $\Delta^3 \times K \rightarrow \mathcal{M}$  whose restriction to  $I_3 \times K$  is equal to  $b$ . The restriction of this filling to  $\Delta^{\{0,3\}} \times K$  satisfies the conditions of [Definition 1.2.1](#) (since the composition of a (co)cartesian morphism with an equivalence is again (co)cartesian), exhibiting  $\tilde{\alpha}$  as adjunct to  $\alpha$ .  $\square$

*Note 1.2.7.* The filler provided by [Proposition 1.2.6](#) is not guaranteed to be unique (even up to equivalence). However, we have some say in how our filler should look. We construct our adjoint filler  $\tilde{\alpha}$  by starting with the solid data of [Equation 6](#) and building the dashed extension  $\tilde{\alpha}: K \rightarrow \mathcal{D}$  a simplex at a time, such that each new simplex making up the map  $\tilde{\alpha}$  is adjunct to the corresponding one in  $\alpha$ . We are free to choose each new adjunct simplex however we like.

**Example 1.2.8.** Let  $f: \mathcal{C} \leftrightarrow \mathcal{D}: g$  be an adjunction of 1-categories, giving an adjunction between quasicategories upon taking nerves. Consider the diagram

$$\begin{array}{ccc} K' = \Delta^{\{0,1,3\}} \cup \Delta^{\{0,2,3\}} & \xhookrightarrow{i} & \Delta^3 = K \\ & \searrow & \swarrow \\ & \Delta^1 & \end{array},$$

where both downwards-facing arrows take  $0, 1 \mapsto 0$  and  $2, 3 \mapsto 1$ , and fix a map  $\tilde{\alpha}: K \rightarrow N(\mathcal{C})$  and a map  $\alpha': K' \rightarrow N(\mathcal{D})$  such that  $\tilde{\alpha} \circ i$  is adjunct to  $\alpha'$ . This corresponds to the diagrams

$$\begin{array}{ccc} \text{in } \mathcal{C} & & \text{in } \mathcal{D} \\ \begin{array}{ccc} C & \xrightarrow{\quad} & gD \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ C' & \xrightarrow{\quad} & gD' \end{array} & & \begin{array}{ccc} fC & \xrightarrow{\quad} & D \\ \downarrow & & \downarrow \\ fC' & \xrightarrow{\quad} & D' \end{array} \end{array}$$

where the solid diagrams in each category are adjunct to one another. [Proposition 1.2.6](#) implies that the solution to the lifting problem on the left yields one on the right. Similarly, its dual implies the converse. Thus, [Proposition 1.2.6](#) implies in particular that such lifting problems are equivalent.

### 1.3 Left Kan implies globally left Kan

In this subsection, we will show that we can solve lifting problems of the form

$$\begin{array}{ccc}
 \Delta_b^{\{0,1,n+1\}} & & \\
 \downarrow & \searrow \tau & \\
 (\Lambda_0^{n+1})_b & \longrightarrow & \mathbb{C}at_\infty \\
 \downarrow & \nearrow & \\
 \Delta_b^{n+1} & & 
 \end{array}$$

for  $n \geq 2$ , where  $\tau$  is left Kan. Forgetting for a moment that have singled out some 2-simplex  $\tau$ , and using the adjunction  $\mathbb{C}^{sc} \dashv N_{sc}$ , we will do this by solving equivalent lifting problems of the form<sup>1</sup>

$$\begin{array}{ccc}
 \mathbb{C}[\Lambda_0^{n+1}] & \longrightarrow & \mathbb{Q}Cat \\
 \downarrow & \nearrow \mathcal{F} & \\
 \mathbb{C}[\Delta^{n+1}] & & 
 \end{array} . \tag{7}$$

We expect that the reader is familiar with the basics of (scaled) rigidification, so we give only a rough description of this lifting problem here. The objects of the simplicial category  $\mathbb{C}[\Delta^{n+1}]$  are given by the set  $\{0, \dots, n+1\}$ , and the mapping spaces are defined by

$$\mathbb{C}[\Delta^{n+1}](i, j) = \begin{cases} N(P_{ij}), & i \leq j \\ \emptyset, & i > j \end{cases} .$$

where  $P_{ij}$  is the poset of subsets of the linearly ordered set  $\{i, \dots, j\}$  containing  $i$  and  $j$ , ordered by inclusion. The simplicial category  $\mathbb{C}[\Lambda_0^{n+1}]$  has the same objects as  $\mathbb{C}[\Delta^{n+1}]$ , and each morphism space  $\mathbb{C}[\Lambda_0^{n+1}](i, j)$  is a simplicial subset of  $\mathbb{C}[\Delta^{n+1}](i, j)$ , as we shall now describe.

The map  $\mathbb{C}[\Lambda_0^{n+1}](i, j) \hookrightarrow \mathbb{C}[\Delta^{n+1}](i, j)$  is an isomorphism except for  $(i, j) = (1, n+1)$  and  $(i, j) = (0, n+1)$ . The missing data corresponds in the case of  $(1, n+1)$  to the missing face  $d_0\Delta^{n+1}$  of  $\Lambda_0^{n+1}$ , and in the case of  $(0, n+1)$  to the missing interior of  $\Delta^{n+1}$ . Thus, to find a

<sup>1</sup>Note that we are implicitly taking  $\mathbb{C}[\Lambda_0^{n+1}]$  and  $\mathbb{C}[\Delta^{n+1}]$  to carry the flat markings on their mapping spaces.

filling as in Equation 7, we need to solve the lifting problems

$$\begin{array}{ccc}
\mathfrak{C}[\Lambda_0^{n+1}](1, n+1) & \longrightarrow & \text{Fun}(\mathcal{F}(1), \mathcal{F}(n+1)) \\
\downarrow & \nearrow \ell' & \\
\mathfrak{C}[\Delta^{n+1}](1, n+1) & & 
\end{array} \quad (8)$$

and

$$\begin{array}{ccc}
\mathfrak{C}[\Lambda_0^{n+1}](0, n+1) & \longrightarrow & \text{Fun}(\mathcal{F}(0), \mathcal{F}(n+1)) \\
\downarrow & \nearrow \ell & \\
\mathfrak{C}[\Delta^{n+1}](0, n+1) & & 
\end{array} . \quad (9)$$

However, these problems are not independent; the filling  $\ell$  of the full simplex needs to agree with the filling  $\ell'$  we found for the missing face of the horn, corresponding to the condition that the square

$$\begin{array}{ccc}
\mathfrak{C}[\Delta^{n+1}](1, n+1) & \xrightarrow{\ell'} & \text{Fun}(\mathcal{F}(1), \mathcal{F}(n+1)) \\
\{0,1\}^* \downarrow & & \downarrow \mathcal{F}(\{0,1\})^* \\
\mathfrak{C}[\Delta^{n+1}](0, n+1) & \xrightarrow{\ell} & \text{Fun}(\mathcal{F}(0), \mathcal{F}(n+1))
\end{array} \quad (10)$$

commute.

*Notation 1.3.1.* Recall our desired filling  $\mathcal{F}: \mathfrak{C}[\Delta^{n+1}] \rightarrow \mathbf{QCat}$  of Equation 7. We will denote  $\mathcal{F}(i)$  by  $X_i$ , and for each subset  $S = \{i_1, \dots, i_k\} \subseteq [n]$ , we will denote  $\mathcal{F}(S) \in \text{Map}(X_{i_1}, X_{i_k})$  by  $f_{i_k \dots i_1}$ . For any inclusion  $S' \subseteq S \subseteq [n]$  preserving minimum and maximum elements, we will denote the corresponding morphism by  $\alpha_{S'}^{S'}$ .

**Theorem 1.3.2.** For any  $n \geq 2$  and any globally left Kan 2-simplex  $\tau$ , the solid extension problem

$$\begin{array}{ccc}
\Delta_b^{\{0,1,n+1\}} & & \\
\downarrow & \searrow \tau & \\
(\Lambda_0^{n+1})_b & \longrightarrow & \text{Cat}_\infty \\
\downarrow & \nearrow & \\
\Delta_b^{n+1} & & 
\end{array}$$

admits a dashed filler.

*Proof.* We need to solve the lifting problems Equation 8 and Equation 9, and check that our solutions satisfy the condition of Equation 10. We note the following useful pieces of information.

- The information that the simplex  $\Delta_b^{\{0,1,n+1\}}$  is left Kan tells us that  $f_{n,1}$  is the left Kan extension of  $f_{n,0}$  along  $f_{1,0}$ , and that  $\alpha_{\{0,n+1\}}^{\{0,1,n+1\}} : f_{n+1,0} \rightarrow f_{n+1,1} \circ f_{1,0}$  is the unit map.
- We have an adjunction

$$(f_{1,0})_! : \text{Fun}(X_0, X_{n+1}) \longleftrightarrow \text{Fun}(X_1, X_{n+1}) : f_{\{0,1\}}^*.$$

It is easy to see that we have the following succinct descriptions of the inclusions of [Equation 8](#) and [Equation 9](#) along which we have to extend.

- There is an isomorphism of simplicial sets

$$\mathfrak{C}[\Delta^{n+1}](0, n+1) \xrightarrow{\cong} C^n$$

specified completely by sending a subset  $S \subseteq [n+1]$  to the point

$$(z_1, \dots, z_n) \in C^n, \quad z_i = \begin{cases} 1, & i \in S \\ 0, & \text{otherwise.} \end{cases}$$

The inclusion  $\mathfrak{C}[\Lambda_0^{n+1}](0, n+1) \hookrightarrow \mathfrak{C}[\Delta^{n+1}](0, n+1)$  corresponds to the simplicial subset  $LC^n \hookrightarrow C^n$ .

- There is a similar isomorphism of simplicial sets

$$\mathfrak{C}[\Delta^{n+1}](1, n+1) \xrightarrow{\cong} C^{n-1}$$

specified completely by sending  $S \subseteq \{1, \dots, n+1\}$  to the point

$$(z_1, \dots, z_{n-1}) \in C^{n-1}, \quad z_i = \begin{cases} 1, & i+1 \in S \\ 0, & \text{otherwise.} \end{cases}$$

The inclusion  $\mathfrak{C}[\Lambda_0^{n+1}](1, n+1) \hookrightarrow \mathfrak{C}[\Delta^{n+1}](1, n+1)$  corresponds to the inclusion  $\partial C^{n-1} \hookrightarrow C^{n-1}$ .

Using these descriptions, we can write our lifting problems in the more inviting form

$$(*) = \begin{array}{ccc} LC^n & \longrightarrow & \text{Fun}(X_0, X_{n+1}) \\ \downarrow & \nearrow & \\ C^n & & \end{array}, \quad (**) = \begin{array}{ccc} \partial C^{n-1} & \xrightarrow{\tilde{\alpha}'} & \text{Fun}(X_1, X_{n+1}) \\ \downarrow & \nearrow & \\ C^{n-1} & & \end{array},$$



and our condition becomes that the square

$$(\star) = \begin{array}{ccc} C^{n-1} & \longrightarrow & \text{Fun}(X_1, X_{n+1}) \\ \downarrow & & \downarrow f_{1,0}^* \\ C^n & \longrightarrow & \text{Fun}(X_0, X_{n+1}) \end{array}$$

commutes, where the left-hand vertical morphism is the inclusion of the right face.

Using [Proposition 1.1.19](#), we can partially solve the lifting problem  $(\star)$ , reducing it to the lifting problem

$$(*)' = \begin{array}{ccc} \Delta^0 * \partial C^{n-1} & \xrightarrow{\alpha} & \text{Fun}(X_0, X_{n+1}) \\ \downarrow & \nearrow & \\ \Delta^0 * C^{n-1} & & \end{array}.$$

The image of the 1-simplex  $(0, \dots, 0) \rightarrow (0, \dots, 1)$  of  $\Delta^0 * \partial C^{n-1} \subseteq C^n$  under  $\alpha$  is the unit map  $\alpha_{\{0, n+1\}}^{\{0, 1, n+1\}} : f_{n+1,0} \rightarrow f_{n+1,1} \circ f_{1,0}$ . This is not an equivalence, so we cannot use [Proposition 1.1.19](#) to solve the lifting problem  $(*)'$  directly. However, the unit map is adjunct to the identity map  $\text{id}_{f_{n+1,1}}$  in  $\text{Fun}(X_1, X_{n+1})$  relative to  $s = \text{id}_{\Delta^1}$ . Furthermore, the restriction of  $\alpha$  to  $\partial C^{n-1}$  is in the image of  $f_{1,0}^*$ , so by [Proposition 1.2.6](#) we can augment the map  $\tilde{\alpha}'$  to a map  $\tilde{\alpha} : \Delta^0 * \partial C^{n-1} \rightarrow \text{Fun}(X_1, X_{n+1})$  which is adjunct to  $\alpha$  relative to the map

$$s : \Delta^0 * \partial C^{n-1} \rightarrow \Delta^1$$

sending  $\Delta^0$  to  $\Delta^{\{0\}}$  and  $\partial C^{n-1}$  to  $\Delta^{\{1\}}$ ; by [Note 1.2.7](#), we can choose  $\tilde{\alpha}$  such that the image of the morphism  $(0, \dots, 0) \rightarrow (0, \dots, 1)$  is an equivalence. This allows us to replace the lifting problem  $(**)$  by the superficially more complicated lifting problem

$$(**)' = \begin{array}{ccc} \Delta^0 * \partial C^{n-1} & \xrightarrow{\tilde{\alpha}} & \text{Fun}(X_1, X_{n+1}) \\ \downarrow & \nearrow \tilde{\beta} & \\ \Delta^0 * C^{n-1} & & \end{array}$$

However, since image of  $(0, \dots, 0) \rightarrow (0, \dots, 1)$  is an equivalence, [Proposition 1.1.19](#) implies that we can solve the lifting problem  $(**')$ . Again using (the dual to) [Proposition 1.2.6](#), we can transport this filling to a solution to the lifting problem  $(*)'$ . The condition  $(\star)$  amounts to demanding that the restriction  $\beta|_{C^{n-1}}$  be the image of  $\tilde{\beta}|_{C^{n-1}}$  under  $f_{1,0}^*$ , which is true by construction.  $\square$