

1 Horn filling via left Kan extensions

Kan extensions of ∞ -categories are usually defined pointwise via the colimit formula [2] [1]. In this section, we will show that such pointwise ∞ -Kan extensions enjoy a horn filling property in $\mathbb{C}at_\infty$ which generalizes the universal property for Kan extensions of functors between 1-categories.

Definition 1.0.1. We will call a 2-simplex $\tau: \Delta_b^2 \rightarrow \mathbb{C}at_\infty$ **left Kan** if it is of the form

$$\begin{array}{ccc} X & \xrightarrow{F} & \mathcal{C} \\ & \searrow f \quad \eta \Downarrow \quad \nearrow f_!F & \\ & Y & \end{array},$$

where \mathcal{C} is cocomplete, X is small, $f_!F$ is a left Kan extension of F along f , and $\eta: F \Rightarrow f^* f_!F$ is the unit map.

The goal of this section is to prove the following.

Theorem 1.0.2. Let $\tau: \Delta^2 \rightarrow \mathbb{C}at_\infty$ be a left Kan 2-simplex. Then for each $n \geq 2$, every solid lifting problem of the form

$$\begin{array}{ccc} \Delta_b^{\{0,1,n+1\}} & \xrightarrow{\tau} & \mathbb{C}at_\infty \\ \downarrow & \searrow & \uparrow \\ (\Lambda_0^{n+1})_b & \xrightarrow{\quad} & \mathbb{C}at_\infty \\ \downarrow & \nearrow \text{dashed} & \\ \Delta_b^{n+1} & & \end{array} \quad (1)$$

admits a dashed solution.

Let us examine this in the case $n = 2$. In this case we have the data of categories and functors

$$\begin{array}{ccccc} & & Y & & \\ & f \nearrow & & f_!F \searrow & \\ X & & & & \mathcal{C} \\ & \xrightarrow{\quad} & & \xrightarrow{F} & \\ & \searrow h & & \nearrow G & \\ & & Z & & \end{array}$$

together with natural transformations $\eta: F \Rightarrow f_!F \circ f$, $\beta: h \Rightarrow g \circ f$, and $\delta: F \Rightarrow G \circ h$. We need to produce a natural transformation $\alpha: f_!F \Rightarrow G \circ g$ and a filling of the full 3-simplex,

which is the data of a homomotopy-commutative diagram

$$\begin{array}{ccc} F & \xrightarrow{\eta} & f_! F \circ f \\ \delta \downarrow & & \downarrow \alpha f \\ G \circ h & \xrightarrow{G\beta} & G \circ g \circ f \end{array}$$

in $\text{Fun}(X, \mathcal{C})$.¹

We can rephrase this as follows. We are given the data

$$\begin{array}{ccc} \overbrace{\begin{array}{ccc} F & \xrightarrow{\eta} & f_! F \circ f \\ \delta \downarrow & & \\ G \circ h & \xrightarrow{G\beta} & G \circ g \circ f \end{array}}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\begin{array}{c} f_! F \\ \\ G \circ g \end{array}}^{\text{in Fun}(Y, \mathcal{C})} \end{array}$$

where the left-hand diagram is a map $LC^2 \rightarrow \text{Fun}(X, \mathcal{C})$, where LC^2 is the simplicial subset of the boundary of $\Delta^1 \times \Delta^1$ except the right-hand face, and the right-hand diagram is a map $\partial\Delta^1 \rightarrow \text{Fun}(Y, \mathcal{C})$. We need to construct a filler $\partial\Delta^1 \hookrightarrow \Delta^1$ on the right, and from it a filler $LC^2 \hookrightarrow \Delta^1 \times \Delta^1$ on the left. We do this in the following sequence of steps.

1. We first can fill the lower-left half of the diagram on the left simply by taking the composition $G\beta \circ \delta$. Doing this, we are left with the filling problem

$$\begin{array}{ccc} \overbrace{\begin{array}{ccc} F & \xrightarrow{\eta} & f_! F \circ f \\ & \searrow & \\ & & G \circ g \circ f \end{array}}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\begin{array}{c} f_! F \\ \\ G \circ g \end{array}}^{\text{in Fun}(Y, \mathcal{C})} \end{array} .$$

2. We note that we can extend the diagram on the right to a diagram which is adjunct to the diagram on the left under the adjunction $f_! \dashv f^*$:

$$\begin{array}{ccc} \overbrace{\begin{array}{ccc} F & \xrightarrow{\eta} & f_! F \circ f \\ & \searrow & \\ & & G \circ g \circ f \end{array}}^{\text{in Fun}(X, \mathcal{C})} & & \overbrace{\begin{array}{ccc} f_! F & \xrightarrow{\text{id}} & f_! F \\ & \searrow & \\ & & G \circ g \end{array}}^{\text{in Fun}(Y, \mathcal{C})} \end{array}$$

¹In particular, if we take $h = f$, $g = \text{id}_Y$, and β to be the identity $f \Rightarrow f$, we recover the classical universal property satisfied by left Kan extension.

3. The diagram on the right has an obvious filler, which is adjunct to a filler on the left.

$$\begin{array}{ccc}
 \text{in Fun}(X, \mathcal{C}) & & \text{in Fun}(Y, \mathcal{C}) \\
 \begin{array}{ccc}
 F & \xrightarrow{\eta} & f_! F \circ f \\
 & \searrow & \downarrow \alpha f \\
 & & G \circ g \circ f
 \end{array} & & \begin{array}{ccc}
 f_! F & \xrightarrow{\text{id}} & f_! F \\
 & \searrow & \downarrow \alpha \\
 & & G \circ g
 \end{array}
 \end{array}$$

The lifting problems which we have to solve in Equation 1 for $n > 2$ amount to replacing $\Delta^1 \times \Delta^1$ by $(\Delta^1)^n$, etc; the basic process remains unchanged, but the combinatorics involved in filling the necessary cubes becomes more involved. In Subsection 1.1 we explain the combinatorics of filling cubes relative to their boundaries. In Subsection 1.2, we give a formalization the concept of adjunct data, and provide a means of for filling partial data to total adjunct data. In Subsection 1.3, we show that we can always solve the lifting problems of Equation 1.

1.1 Filling cubes relative to their boundaries

In this section, we give the combinatorics of filling the n -cube $(\Delta^1)^n$ relative to its boundary. Our main result is Proposition 1.1.21, which writes the inclusion the boundary of the n -cube missing a certain face into the full cube as a composition of an inner anodyne map and a marked anodyne map.

Definition 1.1.1. The n -cube is the simplicial set $C^n := (\Delta^1)^n$.

Note 1.1.2. We will consider the the factors Δ^1 of C^n to be ordered, so that we can speak about the first factor, the second factor, etc.

Our first step is to understand the nondegenerate simplices of the n -cube.

Definition 1.1.3. Denote by $a^i: \Delta^n \rightarrow \Delta^1$ the map which sends $\Delta^{\{0, \dots, i-1\}}$ to $\{0\}$, and $\Delta^{\{i, \dots, n\}}$ to $\{1\}$.

Note 1.1.4. The superscript of a^i counts how many vertices of Δ^n are sent to $\Delta^{\{0\}} \subset \Delta^1$.

Every nondegenerate simplex $\Delta^n \rightarrow C^n$ can be specified by giving a walk along the edges of C^n starting at $(0, \dots, 0)$ and ending at $(1, \dots, 1)$, and any simplex specified in this way is nondegenerate. We now use this to define a bijection between S_n , the symmetric group on the set $\{1, \dots, n\}$, and the nondegenerate simplices $\Delta^n \rightarrow C^n$ as follows.

Definition 1.1.5. For any permutation $\tau: \{1, \dots, n\} \rightarrow \{1, \dots, n\}$, we define an n -simplex

$$\phi(\tau): \Delta^n \rightarrow C^n; \quad \phi(\tau)_i = a^{\tau(i)}.$$

That is,

$$\phi(\tau) = (a^{\tau(1)}, \dots, a^{\tau(n)}).$$

Note 1.1.6. It is easy to check that the above definition really results in a bijection between S_n and the nondegenerate simplices of C^n .

Notation 1.1.7. We will denote any permutatation $\tau: \{1, \dots, n\} \rightarrow \{1, \dots, n\}$ by the corresponding n -tuple $(\tau(1), \dots, \tau(n))$. Thus, the identity permutation is $(1, \dots, n)$, and the permutation $\gamma_{1,2}$ which swaps 1 and 2 is $(2, 1, 3, \dots, n)$.

Given a permutation $\tau \in S_n$ corresponding to an n -simplex

$$\phi(\tau) = (a^{\tau(1)}, \dots, a^{\tau(n)}): \Delta^n \rightarrow C^n,$$

the i th face of $\phi(\tau)$ is a nondegenerate $(n-1)$ -simplex in C^n , i.e. a map $\Delta^{n-1} \rightarrow C^n$. We calculate this as follows: for any $a^i: \Delta^n \rightarrow \Delta^1$ and any face map $\partial_j: \Delta^{n-1} \rightarrow \Delta^n$, we note that

$$\partial_j^* a^i = \begin{cases} a^{i-1}, & j < i \\ a^i, & j \geq i \end{cases},$$

where by minor abuse of notation we denote the map $\Delta^{n-1} \rightarrow \Delta^1$ sending $\Delta^{\{0, \dots, i-1\}}$ to $\{0\}$ and the rest to $\{1\}$ also by a^i . We then have that

$$d_i \tau = d_i(a^{\tau(1)}, \dots, a^{\tau(n)}) := (\partial_i^* a^{\tau(1)}, \dots, \partial_i^* a^{\tau(n)}).$$

Example 1.1.8. Consider the 3-simplex (a^1, a^2, a^3) in C^3 . This has spine

$$(0, 0, 0) \rightarrow (1, 0, 0) \rightarrow (1, 1, 0) \rightarrow (1, 1, 1),$$

as is easy to see quickly:

- The function $\Delta^3 \rightarrow \Delta^1$ corresponding to the first coordinate is a^1 , so the first coordinate of the first vertex is 0, and the first coordinate of the rest of the vertices are 1.
- The function $\Delta^3 \rightarrow \Delta^1$ corresponding to the second coordinate is a^2 , so the second coordinates of the first two points are equal to zero, and the second coordinate of the remaining vertices is 1.
- The function $\Delta^3 \rightarrow \Delta^1$ corresponding to the third coordinate is a^3 , so the third coordinates of the first three points is equal to zero, and the second coordinate of the final vertex is equal to 1.

Using [Equation 1.1](#), we then calculate that

$$d_2(a^1, a^2, a^3) = (a^1, a^2, a^2).$$

This corresponds to the 2-simplex in C^n with spine

$$(0, 0, 0) \rightarrow (1, 0, 0) \rightarrow (1, 1, 1).$$

Our next task is to understand the boundary of the n -cube. We can view the boundary of the cube C^n as the union of its faces.

Definition 1.1.9. The *boundary of the n -cube* is the simplicial subset

$$\partial C^n := \bigcup_{i=1}^n \bigcup_{j=0}^1 \Delta^{(1)} \times \cdots \times \{j\} \times \cdots \times \Delta^{(n)} \subset C^n. \quad (2)$$

The following is easy to see.

Proposition 1.1.10. An $(n-1)$ -simplex

$$\gamma = (\gamma_1, \dots, \gamma_n): \Delta^{n-1} \rightarrow C^n$$

lies entirely within the face

$$\Delta^{(1)} \times \cdots \times \{0\} \times \cdots \times \Delta^{(n)}$$

if and only if $\gamma_i = a^n = \text{const}_0$, and to the face

$$\Delta^{(1)} \times \cdots \times \{1\} \times \cdots \times \Delta^{(n)}$$

if and only if $\gamma_i = a^0 = \text{const}_1$.

Definition 1.1.11. The *left box of the n -cube*, denoted LC^n , is the simplicial subset of C^n given by the union of all of the faces in [Equation 2](#) except for

$$\Delta^1 \times \cdots \times \Delta^1 \times \Delta^{\{1\}}.$$

Note that drawing the box with the final coordinate going from left to right, the right face is open here; the terminology is chosen to match with ‘left horn’.

Example 1.1.12. The simplicial subset $LC^2 \subset C^2$ can be drawn as follows.

$$\begin{array}{ccc} (0, 0) & \longrightarrow & (0, 1) \\ \downarrow & & \\ (1, 0) & \longrightarrow & (1, 1) \end{array}$$

Note that in LC^2 , the morphisms $(0, 0) \rightarrow (1, 0) \rightarrow (1, 1)$ form an inner horn, which we can fill by pushing out along an inner horn inclusion. Our main goal in this section is to show that this is generically true: given a left cube LC^n , we can fill much of C^n using pushouts along inner horn inclusions. To this end, it will be helpful to know how each nondegenerate simplex in C^n intersects LC^n .

Lemma 1.1.13. Let $\phi: \Delta^n \rightarrow C^n$ be a nondegenerate n -simplex corresponding to the permutation $\tau \in S_n$. We have the following.

1. The zeroth face $d_0\phi$ belongs to LC^n if and only if $\tau(n) \neq 1$.
2. For $0 < i < n$, the face $d_i\phi$ never belongs to LC^n .
3. The n th face $d_n\phi$ always belongs to LC^n .

Proof. 1. We have

$$d_0\phi = (a^{\tau(0)-1}, \dots, a^{\tau(n)-1})$$

since $\tau(i) > 0$ for all i . For $j = \tau^{-1}(1)$, we have that $\tau(j) = 1$, so the j th entry of $d_0\phi$ is a^0 . Thus, [Proposition 1.1.10](#) guarantees that $d_0\phi$ is contained in the face

$$\Delta^{(1)} \times \dots \times \Delta^{(j)} \times \dots \times \Delta^{(n)}.$$

This face belongs to LC^n except when $j = n$.

2. In this case, the superscript of each entry of $d_i\phi$ is between 1 and $n - 1$, hence not equal to n or 0.
3. In this case,

$$d_n(a^{\tau(1)}, \dots, a^{\tau(n)}) = (a^{\tau(1)}, \dots, a^{\tau(n)})$$

since $n \geq \tau(i)$ for all $1 \leq i \leq n$. Thus, the $\tau^{-1}(n) = j$ th entry is a^n , so $d_n\phi$ belongs to the face

$$\Delta^{(1)} \times \dots \times \Delta^{(j)} \times \dots \times \Delta^{(n)}.$$

□

This is promising: it means that for many simplices in C^n which we want to fill relative to LC^n , we already have the data of the first and last faces. The following lemma will allow us to take advantage of this.

Notation 1.1.14. For any subset $T \subseteq [n]$, write

$$\Lambda_T^n := \bigcup_{t \in T} d_t \Delta^n \subset \Delta^n.$$

Lemma 1.1.15. For any proper subset $T \subset [n]$ containing 0 and n , the inclusion $\Lambda_T^n \hookrightarrow \Delta^n$ is inner anodyne.

Proof. Induction. For $n = 2$, T must be equal to $\{0, 2\}$, so $\Lambda_T^2 = \Lambda_1^2$.

Assume the result holds for $n - 1$, and let $T \subset [n]$ be a proper subset containing 0 and n . Using the inductive step, we can fill all but one of the faces $d_1 \Delta^n, \dots, d_{n-1} \Delta^n$. The result follows \square

We can use [Lemma 1.1.15](#) to show that we can fill much of C^n as a sequence of inner anodyne pushouts, but to do this we need to pick an order in which to fill our simplices. We do this as follows.

Definition 1.1.16. We define a total order on S_n as follows. For any two permutations $\tau, \tau' \in S_n$, we say that $\tau < \tau'$ if there exists $k \in [n]$ such that the following conditions are satisfied.

- For all $i < k$, we have that $\tau^{-1}(i) = \tau'^{-1}(i)$.
- We have that $\tau^{-1}(k) < \tau'^{-1}(k)$.

We then define a total order on the nondegenerate simplices of C^n by saying that $\phi(\tau) < \phi(\tau')$ if and only if $\tau < \tau'$.

Note that this is *not* the lexicographic order on S_n ; instead, we have that $\tau < \tau'$ if and only if τ^{-1} is less than τ'^{-1} under the lexicographic order.

Example 1.1.17. The elements of the permutation group S_3 have the order

$$(1, 2, 3) < (1, 3, 2) < (2, 1, 3) < (3, 1, 2) < (2, 3, 1) < (3, 2, 1).$$

We use this ordering to define our filtration.

Notation 1.1.18. For $\tau \in S_n$, we mean by

$$\bigcup_{\tau' \in S_n}^{\tau}$$

the union over all $\tau' \in S_n$ for which $\tau' < \tau$ with respect to the ordering defined above. Note the strict inequality.

We would like to show that each step of the filtration

$$\begin{aligned} LC^n &\hookrightarrow LC^n \cup \phi(1, \dots, n) \\ &\hookrightarrow \dots \\ &\hookrightarrow LC^n \cup \bigcup_{\tau \in S_n}^{(2, \dots, n, 1)} \phi(\tau) \end{aligned} \tag{3}$$

is inner anodyne. To do this, it suffices to show that for each $\tau < (2, \dots, n, 1)$, the intersection

$$B_\tau = \left(LC^n \cup \bigcup_{\tau' \in S_n}^{\tau} \phi(\tau') \right) \cap \phi(\tau) \tag{4}$$

is of the form Λ_T^n for some proper $T \subseteq [n]$ containing 0 and n .

Proposition 1.1.19. Each inclusion in Equation 3 is inner anodyne.

Proof. The first inclusion

$$LC^n \hookrightarrow LC^n \cup \phi(0, \dots, n)$$

is inner anodyne because by Lemma 1.1.13 the intersection $LC^n \cap \phi(0, \dots, n)$ is of the form $d_0\Delta^n \cup d_n\Delta^n$.

Each subsequent intersection contains the faces $d_0\Delta^n$ and $d_n\Delta^n$ (again by Lemma 1.1.13), so by Lemma 1.1.15, it suffices to show that for each τ under consideration, there is at least one face not shared with any previous τ' .

To this end, fix $0 < i < n$, and consider $\tau \in S_n$ such that $\tau \neq (0, \dots, n)$. We consider the face $d_i\phi(\tau)$. Let $j = \tau^{-1}(i)$ and $j' = \tau^{-1}(i+1)$. Then $d_i a^{\tau(j)} = d_i a^{\tau(j')}$, so $d_i\phi(\tau) = d_i\phi(\tau \circ \gamma_{jj'})$, where $\gamma_{jj'} \in S_n$ is the permutation swapping j and j' . Thus, the face $d_i\sigma(\tau)$ is equal to the face $d_i\sigma(\tau \circ \gamma_{jj'})$, which is already contained in the union under consideration if and only if $\tau \circ \gamma_{jj'} < \tau$. This in turn is true if and only if $j < j'$.

Assume that every face of $\phi(\tau)$ is contained in the union. Then

$$\tau^{-1}(1) < \tau^{-1}(2) < \dots < \tau^{-1}(n),$$

which implies that $\tau = (1, \dots, n)$. This is a contradiction. Thus, each boundary inclusion under consideration is of the form $\Lambda_T^n \hookrightarrow \Delta^n$, for $T \subset [n]$ a proper subset containing 0 and n , and is thus inner anodyne. \square

Each nondegenerate simplex of C^n which we have yet to fill is of the form $\phi(\tau)$, where $\tau(n) = 1$. Thus, each of their spines begins

$$(0, \dots, 0, 0) \rightarrow (0, \dots, 0, 1) \rightarrow \dots,$$

and then remains confined to the face $\Delta^1 \times \dots \times \Delta^1 \times \{1\}$. This implies that the part of LC^n yet to be filled is of the form $\Delta^0 * C^{n-1}$, and the boundary of this with respect to which we must do the filling is of the form $\Delta^0 * \partial C^{n-1}$.

Our next task is to show that we can also perform this filling, under the assumption that $(0, \dots, 0, 0) \rightarrow (0, \dots, 0, 1)$ is marked. We do not give the details of this proof, at is quite similar to the proof of [Proposition 1.1.19](#). One continues to fill simplices in the order prescribed in [Definition 1.1.16](#), and shows that each of these inclusions is marked anodyne. The main tool is the following lemma, proved using the same inductive argument as [Lemma 1.1.15](#).

Lemma 1.1.20. For any subset $T \subset [n]$ containing 1 and n , and *not* containing 0, the inclusion

$$(\Lambda_T^n, \mathcal{E}) \hookrightarrow (\Delta^n, \mathcal{F})$$

is (cocartesian-)marked anodyne, where by \mathcal{F} we mean the set of degenerate 1-simplices together with $\Delta^{\{0,1\}}$, and by \mathcal{E} we mean the restriction of this marking to Λ_T^n .

We collect our major results from this section in the following proposition.

Proposition 1.1.21. Denote by $J^n \subseteq C^n$ the simplicial subset spanned by those nondegenerate n -simplices whose spines do not begin $(0, \dots, 0, 0) \rightarrow (0, \dots, 0, 1)$. We can write the filling $LC^n \hookrightarrow C^n$ as a composition

$$LC^n \xhookrightarrow{i} J^n \xhookrightarrow{j} C^n,$$

where i is inner anodyne, and j fits into a pushout square

$$\begin{array}{ccc} \Delta^0 * \partial C^{n-1} & \hookrightarrow & \Delta^0 * C^{n-1} \\ \downarrow & & \downarrow \\ J^n & \xhookrightarrow{\quad} & C^n \end{array},$$

where the top inclusion underlies a marked anodyne morphism

$$(\Delta^0 * \partial C^{n-1}, \mathcal{G}) \hookrightarrow (\Delta^0 * C^{n-1}, \mathcal{G}),$$

where the marking \mathcal{G} contains all degenerate morphisms together with the morphism $(0, \dots, 0) \rightarrow (0, \dots, 1)$.

1.2 Adjunct data

In this section, we give a formalization of the the notion of adjunct data. Our main result is [Proposition 1.2.6](#), which shows that under certain conditions, we can use data on one side of an adjunction to solve lifting problems on the other side.

Definition 1.2.1. Let $s: K \rightarrow \Delta^1$ be a map of simplicial sets, and let $p: \mathcal{M} \rightarrow \Delta^1$ be a bicartesian fibration associated to the adjunction

$$f: \mathcal{C} \longleftrightarrow \mathcal{D} : g$$

via equivalences $h_0: \mathcal{C} \rightarrow \mathcal{M}_0$ and $h_1: \mathcal{D} \rightarrow \mathcal{M}_1$. We say that a map $\alpha: K \rightarrow \mathcal{C}$ is **adjunct** to a map $\tilde{\alpha}: K \rightarrow \mathcal{D}$ relative to s , and equivalently that $\tilde{\alpha}$ is adjunct to α relative to s , if there exists a map $A: K \times \Delta^1 \rightarrow \mathcal{M}$ such that the diagram

$$\begin{array}{ccc} K \times \Delta^1 & \xrightarrow{A} & \mathcal{M} \\ & \searrow \text{pr}_{\Delta^1} & \downarrow p \\ & & \Delta^1 \end{array}$$

commutes, and such that the following conditions are satisfied.

1. The restriction $A|_{K \times \{0\}} = h_0 \circ \alpha$.
2. The restriction $A|_{K \times \{1\}} = h_1 \circ \tilde{\alpha}$.
3. For each vertex $k \in K$ in the fiber $s^{-1}\{0\}$, the image of the edge $\{k\} \times \Delta^1 \rightarrow \mathcal{M}$ is p -cocartesian.
4. For each vertex $k' \in K$ in the fiber $s^{-1}\{1\}$, the image of the edge $\{k'\} \times \Delta^1 \rightarrow \mathcal{M}$ is p -cartesian.

We often leave the equivalences h_0 and h_1 implicit to lighten notation.

In the next examples, fix an adjunction

$$f: \mathcal{C} \longleftrightarrow \mathcal{D} : g$$

corresponding to a bicartesian fibration $p: \mathcal{M} \rightarrow \Delta^1$.

Example 1.2.2. Any morphism in \mathcal{D} of the form $fC \rightarrow D$ is adjunct to some morphism in \mathcal{C}

of the form $C \rightarrow gD$, witnessed by any square

$$\begin{array}{ccc} C & \xrightarrow{a} & fC \\ \downarrow & & \downarrow \\ gD & \xrightarrow{b} & D \end{array}$$

in \mathcal{M} where a is p -cocartesian and b is p -cartesian. This corresponds to the map $s = \text{id}: \Delta^1 \rightarrow \Delta^1$.

Example 1.2.3. Pick some object $D \in \mathcal{D}$, and consider the identity morphism $\text{id}: gD \rightarrow gD$ in \mathcal{C} . This morphism is adjoint to the component of the unit map $\eta_D: D \rightarrow fgD$ relative to $s = \text{id}: \Delta^1 \rightarrow \Delta^1$.

The next example will follow from [Proposition 1.2.6](#).

Example 1.2.4. Any diagram $K * K' \rightarrow \mathcal{D}$ such that K is in the image of f is adjoint to some diagram $K * K' \rightarrow \mathcal{C}$ such that K' is in the image of g .

Lemma 1.2.5. The inclusion of marked simplicial sets

$$\left(\Delta^n \times \Delta^{\{0\}} \coprod_{\Delta^n \times \Delta^{\{0\}}} \partial \Delta^n \times \Delta^{\{0\}}, \mathcal{E} \right) \hookrightarrow (\Delta^n \times \Delta^1, \mathcal{F}) \quad (5)$$

where the marking \mathcal{F} is the flat marking together with the edge $\Delta^{\{0\}} \times \Delta^1$, and \mathcal{E} is the restriction of this marking, is marked (cocartesian) anodyne.

The next proposition shows the power of adjunctions in lifting problems: given data on either side of an adjunction, we can fill in adjunct data on the other side.

Proposition 1.2.6. Let $\mathcal{M} \rightarrow \Delta^1$, \mathcal{C} , \mathcal{D} , f , g , h_0 and h_1 be as in [Definition 1.2.1](#), and let

$$\begin{array}{ccc} K' & \xhookrightarrow{i} & K \\ \searrow \text{so } i & & \swarrow s \\ & \Delta^1 & \end{array},$$

be a commuting triangle of simplicial sets, where i is a cofibration such that $i|_{\{1\}}$ is an isomorphism. Let $\tilde{\alpha}': K' \rightarrow \mathcal{D}$ and $\alpha: K \rightarrow \mathcal{C}$ be maps, and denote $\alpha' = \alpha \circ i$. Suppose that $\tilde{\alpha}'$ is

adjunct to α' relative to $s \circ i$. Then there exists a dashed extension

$$\begin{array}{ccc} K' & \xrightarrow{\alpha'} & \mathcal{C} \\ i \downarrow & \nearrow \alpha & \\ K & & \end{array} \rightsquigarrow \begin{array}{ccc} K' & \xrightarrow{\tilde{\alpha}'} & \mathcal{D} \\ i \downarrow & \nearrow \tilde{\alpha} & \\ K & & \end{array} \quad (6)$$

such that $\tilde{\alpha}$ is adjunct to α relative to s .

Proof. Pick some commutative diagram

$$\begin{array}{ccc} K' \times \Delta^1 & \xrightarrow{A} & \mathcal{M} \\ & \searrow \text{pr}_{\Delta^1} & \downarrow \\ & & \Delta^1 \end{array}$$

displaying $\tilde{\alpha}'$ and α' as adjunct. We further consider the map $h_0 \circ \alpha: K \rightarrow \mathcal{M}_0$. From these two pieces of data we can construct the solid commutative square of cocartesian-marked simplicial sets

$$\begin{array}{ccc} (K \times \Delta^{\{0\}} \amalg_{K' \times \Delta^{\{0\}}} K' \times \Delta^1)^{\heartsuit} & \xrightarrow{\quad} & \mathcal{M}^{\natural} \\ \downarrow & \nearrow \ell & \downarrow \\ (K \times \Delta^1)^{\heartsuit} & \xrightarrow{\quad} & (\Delta^1)^{\#} \end{array},$$

where $(K \times \Delta^1)^{\heartsuit}$ is the marked simplicial set where the only nondenerate morphisms marked are those of the form $\{k\} \times \Delta^1$, for $k \in K|_{s^{-1}\{0\}}$, and the \heartsuit -marked pushout-product denotes the restriction of this marking to the pushout-product. Starting from K' and adjoining one nondegenerate simplex of K at a time, we can write the left-hand map in the above square as a transfinite composition of pushouts of morphisms of the form given in Equation 5 (by assumption each simplex σ we adjoin is not completely contained in K' , so the initial vertex of σ must lie in the fiber of s over 0). Each of these is marked anodyne, so the left-hand map is as well. Thus, a dashed lift ℓ exists, giving us a map $h_1^{-1} \circ \ell|_{K \times \Delta^{\{1\}}}: K \rightarrow \mathcal{D}$. This, together with the composition

$$K' \times \Delta^1 \xrightarrow{\tilde{\alpha}' \times \text{id}} \mathcal{D} \times \Delta^1 \xrightarrow{c} \mathcal{D}$$

where $c: h_1^{-1} \circ h_1 \rightarrow \text{id}_{\mathcal{D}}$ is a natural isomorphism, gives a solid diagram

$$\begin{array}{ccc} K \times \Delta^{\{0\}} \amalg_{K' \times \Delta^{\{0\}}} K' \times (\Delta^1)^{\#} & \xrightarrow{\quad} & \mathcal{D}^{\natural} \\ \downarrow & \nearrow m & \\ K \times (\Delta^1)^{\#} & & \end{array},$$

where \mathcal{D}^{\natural} is the marked simplicial set whose underlying simplicial set is \mathcal{D} and where all

equivalences are marked. This admits a dashed lift, and restricting $m|_{K \times \Delta^{\{1\}}}$ gives us a map $\tilde{\alpha}: K \rightarrow \mathcal{D}$ as in [Equation 6](#). However, we still need to display this as adjunct to α : the lift ℓ has $\ell|_{K \times \Delta^{\{0\}}} = h_0 \circ \alpha$, but only $\ell|_{K \times \Delta^{\{1\}}} \simeq h_1 \circ \tilde{\alpha}$, not equality as we require.

Our last task is to remedy this. Denote the spine of the 3-simplex by

$$I_3 := \Delta^{\{0,1\}} \amalg_{\Delta^{\{1\}}} \Delta^{\{1,2\}} \amalg_{\Delta^{\{2\}}} \Delta^{\{2,3\}} \hookrightarrow \Delta^3.$$

We build a map $I_3 \times K \xrightarrow{b} \mathcal{M}$ which:

- On $\Delta^{\{0,1\}} \times K$ is given by the lift ℓ .
- On $\Delta^{\{1,2\}} \times K$ is given by the composition

$$K \times \Delta^1 \xrightarrow{\ell|_{K \times \Delta^{\{1\}}} \times \text{id}} \mathcal{M}_1 \times \Delta^1 \xrightarrow{H} \mathcal{M}_1 \hookrightarrow \mathcal{M},$$

where H is a natural isomorphism $\text{id}_{\mathcal{M}_1} \rightarrow h_1 \circ h_1^{-1}$.

- On $\Delta^{\{2,3\}} \times K$ is given by the natural isomorphism

$$K \times \Delta^1 \xrightarrow{m} \mathcal{D} \xrightarrow{h_1} \mathcal{M}_1 \hookrightarrow \mathcal{M}.$$

The map $I_3 \hookrightarrow \Delta^3$ is inner anodyne, so there exists a map $\Delta^3 \times K \rightarrow \mathcal{M}$ whose restriction to $I_3 \times K$ is equal to b . The restriction of this filling to $\Delta^{\{0,3\}} \times K$ satisfies the conditions of [Definition 1.2.1](#) (since the composition of a (co)cartesian morphism with an equivalence is again (co)cartesian), exhibiting $\tilde{\alpha}$ as adjunct to α . \square

Note 1.2.7. The filler provided by [Proposition 1.2.6](#) is not guaranteed to be unique (even up to equivalence). However, we have some say in how our filler should look. We construct our adjoint filler $\tilde{\alpha}$ by starting with the solid data of [Equation 6](#) and building the dashed extension $\tilde{\alpha}: K \rightarrow \mathcal{D}$ by filling a simplex at a time relative to its boundary, such that each new simplex making up the map $\tilde{\alpha}$ is adjunct to the corresponding one in α . We are free to choose each new adjunct simplex however we like.

Example 1.2.8. Let $f: \mathcal{C} \leftrightarrow \mathcal{D} : g$ be an adjunction of 1-categories, giving an adjunction between quasicategories upon taking nerves. Consider the diagram

$$\begin{array}{ccc} K' = \Delta^{\{0,1,3\}} \cup \Delta^{\{0,2,3\}} & \xhookrightarrow{i} & \Delta^3 = K \\ & \searrow & \swarrow \\ & \Delta^1 & \end{array},$$

where both downwards-facing arrows take $0, 1 \mapsto 0$ and $2, 3 \mapsto 1$, and fix a map $\tilde{\alpha}: K \rightarrow N(\mathcal{C})$

and a map $\alpha': K' \rightarrow N(\mathcal{D})$ such that $\tilde{\alpha} \circ i$ is adjoint to α' . This corresponds to the diagrams

$$\begin{array}{ccc} \text{in } \mathcal{C} & & \text{in } \mathcal{D} \\ \begin{array}{ccc} C & \longrightarrow & gD \\ \downarrow & \nearrow & \downarrow \\ C' & \longrightarrow & gD' \end{array} & & \begin{array}{ccc} fC & \longrightarrow & D \\ \downarrow & & \downarrow \\ fC' & \longrightarrow & D' \end{array} \end{array}$$

where the solid diagrams in each category are adjoint to one another. [Proposition 1.2.6](#) implies that the solution to the lifting problem on the left yields one on the right. Similarly, its dual implies the converse. Thus, [Proposition 1.2.6](#) implies in particular that such lifting problems are equivalent.

1.3 Left Kan implies globally left Kan

In this subsection, we will prove [Theorem 1.0.2](#). In order to do that, we must show that we can solve lifting problems of the form

$$\begin{array}{ccc} \Delta_b^{\{0,1,n+1\}} & & \\ \downarrow & \searrow \tau & \\ (\Lambda_0^{n+1})_b & \longrightarrow & \mathbb{C}at_\infty \\ \downarrow & \nearrow & \\ \Delta_b^{n+1} & & \end{array}$$

for $n \geq 2$, where τ is left Kan. Forgetting for a moment that we have singled out some 2-simplex τ , and using the adjunction $\mathbb{C}^{sc} \dashv N_{sc}$, we will do this by solving equivalent lifting problems of the form²

$$\begin{array}{ccc} \mathbb{C}[\Lambda_0^{n+1}] & \longrightarrow & \mathbb{Q}Cat \\ \downarrow & \nearrow \mathcal{F} & \\ \mathbb{C}[\Delta^{n+1}] & & \end{array} \quad (7)$$

We expect that the reader is familiar with the basics of (scaled) rigidification, so we give only a rough description of this lifting problem here. The objects of the simplicial category $\mathbb{C}[\Delta^{n+1}]$ are given by the set $\{0, \dots, n+1\}$, and the mapping spaces are defined by

$$\mathbb{C}[\Delta^{n+1}](i, j) = \begin{cases} N(P_{ij}), & i \leq j \\ \emptyset, & i > j \end{cases}.$$

²Note that we are implicitly taking $\mathbb{C}[\Lambda_0^{n+1}]$ and $\mathbb{C}[\Delta^{n+1}]$ to carry the flat markings on their mapping spaces.

where P_{ij} is the poset of subsets of the linearly ordered set $\{i, \dots, j\}$ containing i and j , ordered by inclusion. The simplicial category $\mathfrak{C}[\Lambda_0^{n+1}]$ has the same objects as $\mathfrak{C}[\Delta^{n+1}]$, and each morphism space $\mathfrak{C}[\Lambda_0^{n+1}](i, j)$ is a simplicial subset of $\mathfrak{C}[\Delta^{n+1}](i, j)$, as we shall soon describe.

The map $\mathfrak{C}[\Lambda_0^{n+1}](i, j) \hookrightarrow \mathfrak{C}[\Delta^{n+1}](i, j)$ is an isomorphism except for $(i, j) = (1, n+1)$ and $(i, j) = (0, n+1)$. The missing data corresponds in the case of $(1, n+1)$ to the missing face $d_0\Delta^{n+1}$ of Λ_0^{n+1} , and in the case of $(0, n+1)$ to the missing interior. Thus, to find a filling as in Equation 7, we need to solve the lifting problems

$$\begin{array}{ccc} \mathfrak{C}[\Lambda_0^{n+1}](1, n+1) & \longrightarrow & \text{Fun}(\mathcal{F}(1), \mathcal{F}(n+1)) \\ \downarrow & \nearrow \ell' & \\ \mathfrak{C}[\Delta^{n+1}](1, n+1) & & \end{array} \quad (8)$$

and

$$\begin{array}{ccc} \mathfrak{C}[\Lambda_0^{n+1}](0, n+1) & \longrightarrow & \text{Fun}(\mathcal{F}(0), \mathcal{F}(n+1)) \\ \downarrow & \nearrow \ell & \\ \mathfrak{C}[\Delta^{n+1}](0, n+1) & & \end{array} . \quad (9)$$

However, these problems are not independent; the filling ℓ of the full simplex needs to agree with the filling ℓ' we found for the missing face of the horn, corresponding to the condition that the square

$$\begin{array}{ccc} \mathfrak{C}[\Delta^{n+1}](1, n+1) & \xrightarrow{\ell'} & \text{Fun}(\mathcal{F}(1), \mathcal{F}(n+1)) \\ \{0,1\}^* \downarrow & & \downarrow \mathcal{F}(\{0,1\})^* \\ \mathfrak{C}[\Delta^{n+1}](0, n+1) & \xrightarrow{\ell} & \text{Fun}(\mathcal{F}(0), \mathcal{F}(n+1)) \end{array} \quad (10)$$

commute.

Notation 1.3.1. Recall our desired filling $\mathcal{F}: \mathfrak{C}[\Delta^{n+1}] \rightarrow \mathbf{QCat}$ of Equation 7. We will denote $\mathcal{F}(i)$ by X_i , and for each subset $S = \{i_1, \dots, i_k\} \subseteq [n]$, we will denote $\mathcal{F}(S) \in \text{Map}(X_{i_1}, X_{i_k})$ by $f_{i_k \dots i_1}$. For any inclusion $S' \subseteq S \subseteq [n]$ preserving minimum and maximum elements, we will denote the corresponding morphism by $\alpha_S^{S'}$.

Theorem 1.3.2. For any $n \geq 2$ and any globally left Kan 2-simplex τ , the solid extension problem

$$\begin{array}{ccc} \Delta_b^{\{0,1,n+1\}} & & \\ \downarrow & \searrow \tau & \\ (\Lambda_0^{n+1})_b & \longrightarrow & \mathbf{Cat}_\infty \\ \downarrow & \nearrow & \\ \Delta_b^{n+1} & & \end{array}$$

admits a dashed filler.

Proof. We need to solve the lifting problems of Equation 8 and Equation 9, and check that our solutions satisfy the condition of Equation 10. We note the following useful pieces of information.

- The information that the simplex $\tau: \Delta_{\mathfrak{b}}^{\{0,1,n+1\}} \rightarrow \mathbb{C}at_{\infty}$ is left Kan tells us that $f_{n,1}$ is the left Kan extension of $f_{n,0}$ along $f_{1,0}$, and that $\alpha_{\{0,n+1\}}^{\{0,1,n+1\}}: f_{n+1,0} \rightarrow f_{n+1,1} \circ f_{1,0}$ is the unit map.
- We have an adjunction

$$(f_{1,0})! : \text{Fun}(X_0, X_{n+1}) \longleftrightarrow \text{Fun}(X_1, X_{n+1}) : f_{1,0}^*.$$

It is easy to see that we have the following succinct descriptions of the inclusions of Equation 8 and Equation 9 along which we have to extend.

- There is an isomorphism of simplicial sets

$$\mathfrak{C}[\Delta^{n+1}](0, n+1) \xrightarrow{\cong} C^n$$

specified completely by sending a subset $S \subseteq [n+1]$ to the point

$$(z_1, \dots, z_n) \in C^n, \quad z_i = \begin{cases} 1, & i \in S \\ 0, & \text{otherwise.} \end{cases}$$

The inclusion $\mathfrak{C}[\Lambda_0^{n+1}](0, n+1) \hookrightarrow \mathfrak{C}[\Delta^{n+1}](0, n+1)$ corresponds to the simplicial subset $LC^n \hookrightarrow C^n$.

- There is a similar isomorphism of simplicial sets

$$\mathfrak{C}[\Delta^{n+1}](1, n+1) \xrightarrow{\cong} C^{n-1}$$

specified completely by sending $S \subseteq \{1, \dots, n+1\}$ to the point

$$(z_1, \dots, z_{n-1}) \in C^{n-1}, \quad z_i = \begin{cases} 1, & i+1 \in S \\ 0, & \text{otherwise.} \end{cases}$$

The inclusion $\mathfrak{C}[\Lambda_0^{n+1}](1, n+1) \hookrightarrow \mathfrak{C}[\Delta^{n+1}](1, n+1)$ corresponds to the inclusion $\partial C^{n-1} \hookrightarrow C^{n-1}$.

Using these descriptions, we can write our lifting problems in the more inviting form

$$(*) = \begin{array}{ccc} LC^n & \longrightarrow & \text{Fun}(X_0, X_{n+1}) \\ \downarrow & \nearrow & \\ C^n & & \end{array}, \quad (**) = \begin{array}{ccc} \partial C^{n-1} & \xrightarrow{\tilde{\alpha}'} & \text{Fun}(X_1, X_{n+1}) \\ \downarrow & \nearrow & \\ C^{n-1} & & \end{array},$$

and our condition becomes that the square

$$(\star) = \begin{array}{ccc} C^{n-1} & \longrightarrow & \text{Fun}(X_1, X_{n+1}) \\ \downarrow & & \downarrow f_{1,0}^* \\ C^n & \longrightarrow & \text{Fun}(X_0, X_{n+1}) \end{array}$$

commutes, where the left-hand vertical morphism is the inclusion of the right face.

Using [Proposition 1.1.21](#), we can partially solve the lifting problem $(*)$, reducing it to the lifting problem

$$(*)' = \begin{array}{ccc} \Delta^0 * \partial C^{n-1} & \xrightarrow{\alpha} & \text{Fun}(X_0, X_{n+1}) \\ \downarrow & \nearrow & \\ \Delta^0 * C^{n-1} & & \end{array}.$$

The image of the 1-simplex $(0, \dots, 0) \rightarrow (0, \dots, 1)$ of $\Delta^0 * \partial C^{n-1} \subseteq C^n$ under α is the unit map $\alpha_{\{0, n+1\}}^{\{0, 1, n+1\}} : f_{n+1, 0} \rightarrow f_{n+1, 1} \circ f_{1, 0}$. This is not in general an equivalence, so we cannot use [Proposition 1.1.21](#) to solve the lifting problem $(*)'$ directly. However, the unit map is adjunct to the identity map $\text{id}_{f_{n+1, 1}}$ in $\text{Fun}(X_1, X_{n+1})$ relative to $s = \text{id}_{\Delta^1}$. Furthermore, the restriction of α to ∂C^{n-1} is in the image of $f_{1, 0}^*$, so by [Proposition 1.2.6](#) we can augment the map $\tilde{\alpha}'$ to a map $\tilde{\alpha} : \Delta^0 * \partial C^{n-1} \rightarrow \text{Fun}(X_1, X_{n+1})$ which is adjunct to α relative to the map

$$s : \Delta^0 * \partial C^{n-1} \rightarrow \Delta^1$$

sending Δ^0 to $\Delta^{\{0\}}$ and ∂C^{n-1} to $\Delta^{\{1\}}$; by [Note 1.2.7](#), we can choose $\tilde{\alpha}$ such that the image of the morphism $(0, \dots, 0) \rightarrow (0, \dots, 1)$ is an equivalence. This allows us to replace the lifting problem $(**)$ by the superficially more complicated lifting problem

$$(**') = \begin{array}{ccc} \Delta^0 * \partial C^{n-1} & \xrightarrow{\tilde{\alpha}} & \text{Fun}(X_1, X_{n+1}) \\ \downarrow & \nearrow \tilde{\beta} & \\ \Delta^0 * C^{n-1} & & \end{array}$$

However, since image of $(0, \dots, 0) \rightarrow (0, \dots, 1)$ is an equivalence, [Proposition 1.1.21](#) implies that we can solve the lifting problem $(**')$. Again using (the dual to) [Proposition 1.2.6](#), we can transport this filling to a solution to the lifting problem $(*)'$. The condition (\star) amounts

to demanding that the restriction $\beta|_{C^{n-1}}$ be the image of $\tilde{\beta}|_{C^{n-1}}$ under $f_{1,0}^*$, which is true by construction. \square