

1 Infinity categories

I will edit this chapter heavily. It is in desperate need of a good de-clawing.

1.1 Basic definition

Definition 1 (infinity category). A simplicial set \mathcal{C} is called an ∞ -category if each inner horn

$$\begin{array}{ccc} \Lambda_i^n & \xrightarrow{\quad} & \mathcal{C} \\ \downarrow & \nearrow \text{---} & \\ \Delta^n & & \end{array} \quad 0 < i < n$$

has a filler.

Example 2. For any category \mathcal{C} , we have seen in Proposition ?? the nerve $N(\mathcal{C})$ has *unique* inner horn fillers, thus certainly has inner horn fillers.

Example 3. For any topological space X , the singular set $\text{Sing}(x)$ has *all* horn fillers, hence certainly inner horn fillers.

In the previous chapter, we mainly worked with Kan complexes, so we used the Kan model structure on Set_Δ . There is, however, a different model structure on Set_Δ , which is better suited to infinity-categorical thinking.

1.2 The homotopy category of an infinity category

We would like to be able to view infinity-categories as a generalization of categories. This means that we would like to be able to throw away data from an infinity category to get a regular category. This throwing-away process creates from any infinity category \mathcal{C} the so-called *homotopy category* $h\mathcal{C}$. The homotopy category inherits a lot of information from the infinity category from whence it came.

Definition 4 (homotopy category). Let \mathcal{C} be an ∞ -category. Define a category $h\mathcal{C}$, called the homotopy category of \mathcal{C} , as follows.

- The objects $\text{ob}(\mathcal{C})$ are the vertices \mathcal{C}_0 .
- For objects x and y , the set $\mathcal{C}(x, y)$ consists of all edges f, g with $d_0 f = d_0 g = y$ and $d_1 f = d_1 g = x$, modulo the equivalence relation

$$f \sim g \iff \exists \sigma \in \mathcal{C}_2 \quad \text{with} \quad d_1 \sigma = g, \quad d_2 \sigma = f, \quad \text{and} \quad d_0 \sigma = s_0 y.$$

$$\sigma = \begin{array}{ccc} & y & \\ f \nearrow & & \searrow s_0 y \\ x & \xrightarrow{g} & y \end{array}$$

- We have already seen in Theorem ?? that this is an equivalence relation.
- Composition is induced by inner horn fillings $\Lambda_1^2 \rightarrow \Delta^2$; we define the composition $[f] \circ [g]$ to be $[h]$, where h is some edge given to us by filling the following solid horn.

$$\begin{array}{ccc} & y & \\ f \nearrow & & \searrow g \\ x & \dashrightarrow_h & z \end{array}$$

To see that this is well-defined, i.e. respects equivalence classes, consider two fillings as follows.

$$\sigma = \begin{array}{ccc} & y & \\ f \nearrow & & \searrow g \\ x & \xrightarrow{h} & z \end{array}, \quad \sigma' = \begin{array}{ccc} & y & \\ f \nearrow & & \searrow g \\ x & \xrightarrow{h'} & z \end{array}.$$

We need to show that $h \sim h'$. We get this by filling the following horn.

$$\begin{array}{ccc} y & \xrightarrow{g} & z \\ f \uparrow & \nearrow h & \searrow s_0 z \\ x & \xrightarrow{h'} & z \end{array} \quad \begin{array}{ccc} y & \xrightarrow{g} & z \\ f \uparrow & \nearrow f & \searrow s_0 z \\ x & \xrightarrow{h'} & z \end{array}$$

First suppose that $g_1 \sim g_2$. Then filling the horn

$$\begin{array}{ccc} x & \xrightarrow{f} & y \\ h_2 \downarrow & \nearrow h_1 & \searrow g_1 \\ z & \xrightarrow{s_0 z} & z \end{array} \quad \begin{array}{ccc} x & \xrightarrow{f} & y \\ h_2 \downarrow & \nearrow g_2 & \searrow g_1 \\ z & \xrightarrow{s_0 z} & z \end{array}$$

tells us that $h_1 \sim h_2$.

Now suppose that $f_1 \sim f_2$. In this case we prove something stronger, i.e. that if h is a filler for $g \circ f_1$, then it is also a filler for $g \circ f_2$. We find this from the following commutative diagram.

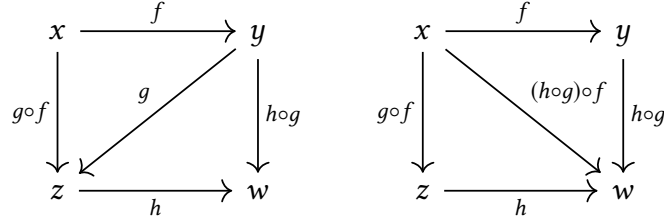
$$\begin{array}{ccc} x & \xrightarrow{f_1} & y \\ f_2 \downarrow & \nearrow h & \searrow g \\ z & \xrightarrow{g} & z \end{array} \quad \begin{array}{ccc} x & \xrightarrow{f} & y \\ h_2 \downarrow & \nearrow g_2 & \searrow g_1 \\ z & \xrightarrow{s_0 z} & z \end{array}$$

This gives us a well-defined composition law. For any $x \in \mathcal{C}_0$, the identity on x is $s_0 x$; it is a left and right identity because the solid diagrams

$$\begin{array}{ccc} & y & \\ f \nearrow & & \searrow s_0 y \\ x & \dashrightarrow_f & y \end{array} \quad \begin{array}{ccc} & x & \\ s_0 x \nearrow & & \searrow g \\ x & \dashrightarrow_g & y \end{array}$$

can be filled by $s_1 f$ and $s_0 g$ respectively.

To see that composition is associative consider the horn



where the bottom triangle on the right is unfilled. Filling it gives a homotopy

$$[(f \circ g) \circ h] = [f \circ (g \circ h)].$$

Lemma 5. The following are properties of $\mathbf{h}\mathcal{C}$.

1. Let \mathcal{C} be a small category. Then $\mathbf{h}N(\mathcal{C}) \simeq \mathcal{C}$.
2. For every topological space X ,

$$\mathbf{h}\mathbf{Sing}(X) \simeq \pi_{\geq 1}(X)$$

3. A simplicial set X is the nerve of a category if and only if it has unique inner horn fillers.
4. For every Kan complex K , the infinity-category $\mathbf{h}K$ is a groupoid.

Proof.

1. The 0-simplices of $N(\mathcal{C})$ are the objects of \mathcal{C} , and these are defined to be the objects of $\mathbf{h}N(\mathcal{C})$.

The 1-simplices of $N(\mathcal{C})$ are precisely the morphisms of \mathcal{C} . In particular, identity morphisms correspond to degenerate simplices. In passing to $\mathbf{h}N(\mathcal{C})$, two morphisms are identified if there is a 2-simplex between them in the sense described above. But this says precisely that the only identifications are of the form $f \circ \text{id} = f$ for all f .

Simplices of degree 2 or higher are discarded.

2. The zero-simplices of both of these are the points of X , and the morphisms of both of these correspond to paths modulo homotopy.
3. Let f be an edge. By filling an outer horn, we find a left inverse g . By filling again, we find a left inverse h to this. But h has to be f , so we are done.

□

Theorem 6. Denote by \mathbf{qCat} the full subcategory of \mathbf{Set}_Δ on ∞ -categories.

1. There is a functor $\mathbf{h} : \mathbf{qCat} \rightarrow \mathbf{Cat}$ which sends $\mathcal{C} \rightarrow \mathbf{h}\mathcal{C}$.
2. There is an adjunction

$$\mathbf{h} : \mathbf{qCat} \leftrightarrow \mathbf{Cat} : N,$$

where N denotes the nerve (Definition ??).

Proof.

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1. Let $f: X \rightarrow Y$ be a map of ∞ -categories. We need to define a map $hf: hX \rightarrow hY$. On objects, this is easy, since $f_0: X_0 \rightarrow Y_0$. On morphisms, we get a map $f_1: X_1 \rightarrow Y_1$ which sends morphisms $x \rightarrow y$ to morphisms $f_0(x) \rightarrow f_0(y)$, as can be checked by substituting d_0 and d_1 into the downward arrows below.

$$\begin{array}{ccc} X_1 & \xrightarrow{f_1} & Y_1 \\ \downarrow & & \downarrow \\ X_0 & \xrightarrow{f_0} & Y_0 \end{array}$$

We still have to check that the map f_1 respects equivalence classes. To this end, suppose $g \sim g': x \rightarrow y$, i.e. there is a simplex as follows.

$$\begin{array}{ccc} & y & \\ g \nearrow & & \searrow s_0 y \\ x & \xrightarrow{g'} & y \end{array}$$

Then we have a simplex as follows

$$\begin{array}{ccc} & f_0(y) & \\ f_1(g) \nearrow & & \searrow s_0(f_0(y))=f_1(s_0(y)) \\ f_0(x) & \xrightarrow{f_1(g')} & f_0(y) \end{array}$$

so $f_1(g) \sim f_1(g')$.

To see functoriality of h , suppose that if $F: X \rightarrow Y$ and $G: Y \rightarrow Z$ are morphisms of simplicial sets. We need to check that $h(G \circ F) = hG \circ hF$. This is clearly true on objects, and on morphisms we have

$$h(G \circ F)([f]) = [(G \circ F)(f)] = [G(F(f))] = hG([F(f)]) = (hG \circ hF)([f]).$$

2. We exhibit a unit-counit adjunction. We first define a natural transformation ϵ with components

$$\epsilon_X: X \rightarrow N(hX).$$

□

1.3 The coherent nerve

1.3.1 Topologically enriched categories

Definition 7 (compactly generated Hausdorff space). A topological space X is known as a compactly generated Hausdorff space if for every subset $U \subset X$, the following are equivalent.

- U is closed.
- For each compact subset $C \subset X$, $C \cap U$ is closed

We denote the full subcategory of **Top** on compactly generated Hausdorff spaces by **CGHaus**.

Our reason for working in the category **CGHaus** rather than **Top** is a technical one: the category **Top** is not Cartesian closed, and therefore is not a good candidate for an enriching category.

Definition 8 (topological category). A category enriched over **CGHaus** (with the Cartesian monoidal structure) is known as a topological category.

1.3.2 Topological and simplicial enrichment

Lemma 9. The functor $\mathbf{Sing}: \mathbf{CGHaus} \rightarrow \mathbf{Set}_\Delta$ is strong monoidal.

Proof. To see this, we build a map

$$\Phi_{X,Y}: \mathbf{Sing}(X \times Y) \rightarrow \mathbf{Sing}(X) \times \mathbf{Sing}(Y)$$

as follows.

$$\begin{aligned} \mathbf{Sing}(X \times Y) &\simeq \mathbf{Top}(\|\Delta^\bullet\|, X \times Y) \\ &\simeq \mathbf{Top}(\|\Delta^\bullet\|, X) \times \mathbf{Top}(\|\Delta^\bullet\|, Y) \\ &\simeq \mathbf{Sing}(X) \times \mathbf{Sing}(Y). \end{aligned}$$

The singular set $\mathbf{Sing}(\|\Delta^0\|) \simeq \Delta^0$, so we don't have any choice in defining ϕ . \square

Definition 10 (simplicial category). A category enriched of simplicial sets (with the Cartesian monoidal structure) is called a simplicial category. Denote the category whose objects are simplicial categories and whose morphisms are enriched functors by \mathbf{Cat}_Δ .

Note 11. This notation is potentially confusing! We are *not* talking about the category of functors $\Delta^{\text{op}} \rightarrow \mathbf{Cat}$.

Definition 12 (singular nerve). Let \mathcal{T} be a topological category. By Lemma ??, we can use $\mathbf{Sing}: \mathbf{Top} \rightarrow \mathbf{Set}_\Delta$ to turn \mathcal{T} into a simplicial category. We will call this the singular nerve of \mathcal{T} and denote it by $\mathbf{Sing}(\mathcal{T})$.

Definition 13 (topological nerve). Let \mathcal{T} be a topological category. The topological nerve of \mathcal{T} , denoted $N_{\text{top}}(\mathcal{T})$, is the infinity category

$$N_\Delta \circ \mathbf{Sing}(\mathcal{T}).$$

Joke 14. How do you get a simplicially enriched category from a topologically enriched category? It's easy—just sing!

1.3.3 Thickening

Definition 15 (thickening). The thickening functor is the functor

$$\mathfrak{C}: \Delta \rightarrow \mathbf{Cat}_\Delta$$

defined as follows.

It sends objects $[n]$ to the simplicially enriched category $\mathfrak{C}[\Delta^n]$ defined as follows.

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- $\text{Obj}(\mathfrak{C}[\Delta^n]) = \{0, \dots, n\}$.
- For $i, j \in \{0, \dots, n\}$

$$\mathfrak{C}[\Delta^n](i, j) = N(P_{i,j}),$$

where $P_{i,j}$ is the poset of subsets of $\{i, \dots, j\}$ containing i and j .

- To define composition, we need a morphism of simplicial sets

$$N(P_{i,j}) \times N(P_{j,k}) \rightarrow N(P_{i,k}).$$

We have a multiplication $P_{i,j} \times P_{j,k} \rightarrow P_{i,k}$ given by $(I, J) \mapsto I \cup J$. This induces the map we need.

A morphism $f: [m] \rightarrow [n]$ induces a functor $\mathfrak{C}[\Delta^m] \rightarrow \mathfrak{C}[\Delta^n]$ on objects by $i \mapsto f(i)$, and on hom-simplicial-sets $N(P_{i,j})$ as follows. On the underlying posets, we send

$$P_{i,j} \supset \{i, k_1, \dots, k_{r-1}, j\} \mapsto \{f(i), f(k_1), \dots, f(k_{r-1}), f(j)\} \subset P_{f(i), f(j)}.$$

This extends to a functor of poset categories, hence a morphism of nerves.

Definition 16 (simplicial nerve). Let \mathcal{C} be a simplicial category. The simplicial nerve of \mathcal{C} , denoted N_Δ , is the functor

$$N_\Delta(\mathcal{C}) = \text{Cat}_\Delta(\mathfrak{C}(-), \mathcal{C}): \Delta^{\text{op}} \rightarrow \text{Set}_\Delta.$$

By Theorem ??, we have an adjunction

$$\mathcal{Y}_! \mathfrak{C} : \text{Set}_\Delta \leftrightarrow \text{Cat}_\Delta : N_\Delta.$$

By abuse of notation, we also denote $\mathcal{Y}_! \mathfrak{C}$ by \mathfrak{C} .

Proposition 17. Let \mathcal{C} be a simplicial category such that for all x, y in \mathcal{C} , $\mathcal{C}(x, y)$ is a Kan complex. Then $N_\Delta(\mathcal{C})$ is an ∞ -category.

Proof. We need to show that we can solve the following lifting problem.

$$\begin{array}{ccc} \Lambda_i^n & \longrightarrow & N_\Delta(\mathcal{C}) \\ \downarrow & \nearrow & \\ \Delta^n & & \end{array}$$

Passing to adjoints, we find the following lifting problem.

$$\begin{array}{ccc} \mathfrak{C}[\Lambda_i^n] & \longrightarrow & \mathcal{C} \\ \downarrow & \nearrow & \\ \mathfrak{C}[\Delta^n] & & \end{array}$$

□

Example 18. Let $\mathbf{Kan} \subset \text{Set}_\Delta$ be the full subcategory on Kan complexes. We may consider \mathbf{Kan} to be simplicially enriched by taking $\mathbf{Kan}(K, S) = \text{Maps}(K, S)$.

Definition 19 (category of spaces). The ∞ -category of spaces is the category

$$\mathcal{S} = N_\Delta(\mathbf{Kan}).$$

1.4 Infinity-groupoids

Definition 20 (equivalence). A morphism $f: x \rightarrow y$ in an infinity category is said to be an equivalence if the corresponding morphism in the homotopy category (Example 4) is an isomorphism.

Definition 21 (infinity-groupoid). An ∞ -category is said to be an ∞ -groupoid if every morphism is an equivalence.

Theorem 22. The following are equivalent.

1. \mathcal{C} is a Kan complex.
2. \mathcal{C} is an infinity-groupoid.

Proof. The direction $1 \implies 2$ is obvious by Lemma 5. We will see why the other direction holds later. \square

Definition 23 (maximal infinity groupoid). Let \mathcal{C} be an infinity category. The maximal infinity groupoid in \mathcal{C} , denoted $\mathcal{C}^\simeq \subset \mathcal{C}$, is the simplicial subset of \mathcal{C} consisting of simplices all of whose edges are equivalences.

1.5 Functors and diagrams

Definition 24 (infinity-functor). An infinity-functor (hereafter simply called *functor*) between infinity-categories \mathcal{C} and \mathcal{D} is simply a map $\mathcal{C} \rightarrow \mathcal{D}$ of simplicial sets.

Here we list some theorems which we will prove later, after we have some more advanced notions.

Proposition 25. Let K be a simplicial set and let \mathcal{C} be an infinity category. Then the simplicial set $\text{Maps}(K, \mathcal{C})$ is an infinity-category.

Proof. Later. \square

Definition 26 (homotopy coherent diagram). Let \mathcal{T} be a topological category and I an ordinary category. Let $\mathcal{C} = N_{\text{top}}(\mathcal{T})$.

1.6 Limits in infinity categories

In this section, we build up the notions necessary to define the limit of a diagram.

In Subsection 1.6.1, we define a monoidal product called the *join* on the category of categories, and show that it reproduces cones and cocones. We then generalize the construction to the category of simplicial sets, and show¹ that the two notions are compatible in the sense that the nerve is a monoidal functor between them; that is we have a natural isomorphism

$$N(\mathcal{C} \star \mathcal{D}) \simeq N(\mathcal{C}) \star N(\mathcal{D}).$$

¹Well, if there's time.

In [Subsection 1.6.2](#), we generalize the notion of the *opposite category* to the context of simplicial sets, and note that it defines an *opposite functor* on the category of simplicial sets, when one extends simplex category Δ in an appropriate way.

In [Subsection 1.6.3](#), we see that joins allow us to define the notion of a cone in an infinity-categorical context. We also define the infinity-categorical generalization of comma categories, and note that there is a duality between the join and the construction of over- and undercategories. We then show that over- and undercategories are ‘topologically’ very simple—they are contractible Kan complexes.

In [Subsection 1.6.4](#) and [Subsection 1.6.5](#), we define the notion of initial (resp. terminal) object in an infinity-category \mathcal{C} . We set up our definitions in such a way that initial objects in an infinity category are initial objects in the homotopy category, and they are unique in the sense that their ‘convex hull’ is a contractible Kan complex.

In [Subsection 1.6.6](#), we define when an object in an infinity category is a limit or colimit of a functor.

1.6.1 Joins

Definition 27 (join of categories). Let \mathcal{C} and \mathcal{C}' be ordinary categories. The join of \mathcal{C} and \mathcal{C}' , denoted $\mathcal{C} \star \mathcal{C}'$, is the category with objects

$$\text{ob}(\mathcal{C} \star \mathcal{C}') = \text{ob}(\mathcal{C}) \amalg \text{ob}(\mathcal{C}')$$

and morphisms

$$(\mathcal{C} \star \mathcal{C}')(x, y) = \begin{cases} \mathcal{C}(x, y), & x, y \in \mathcal{C} \\ \mathcal{C}'(x, y), & x, y \in \mathcal{C}' \\ \{*\}, & x \in \mathcal{C}, y \in \mathcal{C}' \\ \emptyset, & x \in \mathcal{C}', y \in \mathcal{C} \end{cases}.$$

Example 28. We have

$$[m] \star [n] = [m + n + 1].$$

Example 29. Here is an alternative definition for the category of cones over a functor $F: \mathcal{C} \rightarrow \mathcal{D}$, given in terms of the join: it is the category

$$\mathbf{Fun}_F(\Delta^0 \star \mathcal{C}, \mathcal{D}),$$

i.e. the subcategory of $\mathbf{Fun}(\Delta^0 \star \mathcal{C}, \mathcal{D})$ whose objects are functors which agree with F when restricted to \mathcal{C} , and whose morphisms are natural transformations

Definition 30 (join of simplicial sets). Let K and K' be simplicial sets. Define the join of K and K' , denoted $K \star K'$ to be the simplicial set whose n -simplices are

$$(K \star K')_n = \coprod_{[i] \oplus [j] \simeq [n]} K_i \times K'_j,$$

and which acts on a morphism $\phi: [m] \rightarrow [n]$ by sending a simplex

$$\sigma = (\alpha \in K_i, \beta \in K'_j, [i] \oplus [j] \simeq [n]) \in (K \star K')_n$$

to a simplex

$$\phi^* \sigma = (\psi_1^* \alpha, \psi_2^* \beta, \phi^{-1}([i]) \oplus \phi^{-1}([j]) \simeq [m]) \in (K \star K')_m,$$

where

$$\psi_1 = \phi|_{\phi^{-1}([i])}, \quad \text{and} \quad \psi_2 = \phi|_{\phi^{-1}([j])}.$$

This is functorial since

Proposition 31. The join of two ∞ -categories is again an ∞ -category.

Proposition 32. The join is functorial in each argument separately.

Proof. Given a map $f: K \rightarrow K'$ of simplicial sets, we need to specify a map

$$(f, \text{id}_S): K \star S \rightarrow K' \star S.$$

By the universal property of coproducts, we can do this by defining maps

$$K_i \times S_j \rightarrow K'_i \times S_j$$

for each $[i] \oplus [j] \simeq [n]$. We have a map $f_i: K_i \rightarrow K'_i$, which gives a map $f_i \times \text{id}$ as required. We need only check that this is a morphism of simplicial sets, i.e. makes the following square commute for each $\phi: [n] \rightarrow [m]$.

$$\begin{array}{ccc} (K \star S)_m & \longrightarrow & (K \star S)_n \\ \downarrow & & \downarrow \\ (K' \star S)_m & \longrightarrow & (K' \star S)_n \end{array}$$

Let $\sigma = (\alpha, \beta, [i] \oplus [j] \simeq [n]) \in (K \star S)_m$. We follow it around.

$$\begin{array}{ccc} (\alpha, \beta) & \longmapsto & (\psi_i^* \alpha, \psi_j^* \beta) \\ \downarrow & & \downarrow \\ (f_i(\alpha), \beta) & \longmapsto & (f_{\phi^{-1}([i])} \psi_i^* \alpha \stackrel{!}{=} \psi_i^*(f_i \alpha), \beta) \end{array}$$

The condition that $f_{\phi^{-1}([i])} \psi_i^* \alpha = \psi_i^*(f_i \alpha)$ is precisely the condition that the following diagram commute, which it does by naturality of f .

$$\begin{array}{ccc} K_i & \xrightarrow{\psi_i^*} & K_{\phi^{-1}([i])} \\ f_i \downarrow & & \downarrow f_{\phi^{-1}([i])} \\ K'_i & \xrightarrow{\psi_i^*} & K'_{\phi^{-1}([i])} \end{array}$$

Functoriality in the second argument is exactly analogous. □

1.6.2 Opposite categories

Henceforth, we have been considering the domain of simplicial sets to be the category of functors $\Delta^{\text{op}} \rightarrow \mathbf{Set}$. However, we have sometimes relaxed this when it was convenient to do so. For example, the n -simplex Δ^n was defined by

$$\Delta^n = \Delta(-, [n]).$$

We have notated the subsimplex missing the 0th vertex by $\Delta^{\{1, \dots, n\}}$, despite the fact that $\{1, \dots, n\}$ is not an object in Δ , so it is not clear how to make sense of this; in particular, it does not make any sense to write

$$\Delta^{\{1, \dots, n\}} = \Delta(-, \{1, \dots, n\}).$$

What we are really doing is extending Δ to the category of finite, non-empty linearly ordered sets. We could do this because the inclusion functor $\Delta \hookrightarrow \mathbf{FinLinOrd}$ is an equivalence, so passing from Δ to $\mathbf{FinLinOrd}$ is just a shift of perspective.

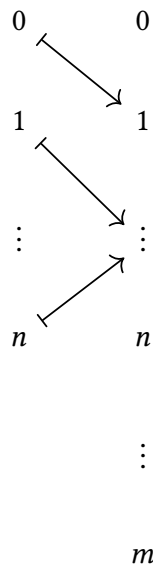
On $\mathbf{FinLinOrd}$, there is a map Swap , which reverses the linear order. That is, it sends the set

$$\{0, \dots, n\} \mapsto \{n, \dots, 0\}.$$

Let us try to understand the behavior of Swap . Any map

$$f: \{0, \dots, n\} \rightarrow \{0, \dots, m\}$$

can be drawn diagrammatically as follows



The condition that the map f be weakly monotonic is equivalent to the demand that none of the arrows in the above diagram cross. The swap operation turns the above diagram on its head, taking uncrossed arrows to uncrossed arrows. Thus, morphisms $f: [n] \rightarrow [m]$ are taken to morphisms $f^{\text{op}}: [\bar{n}] \rightarrow [\bar{m}]$. Furthermore, composition is respected because composition is done by stacking diagrams left to right.

This discussion means that Swap is an endofunctor on $\mathbf{FinLinOrd}$, and further that it is an equivalence of categories, since $\text{Swap}^2 = \text{id}$.²

²We are now solidly in evil territory. I wonder if there is a way out.

Let \mathcal{C} and \mathcal{D} be regular categories. Any endofunctor $E: \mathcal{C} \rightarrow \mathcal{C}$ gives an endofunctor on the functor category $[\mathcal{C}, \mathcal{D}]$ as follows.

- On objects F , it is defined via the pullback.

$$F \mapsto F \circ E$$

- On morphisms η , it is defined via whiskering.

$$\mathcal{C} \xrightarrow{E} \mathcal{C} \begin{array}{c} \xrightarrow{F} \\ \eta \\ \xrightarrow{F'} \end{array} \mathcal{D} \quad \mapsto \quad \mathcal{C} \begin{array}{c} \xrightarrow{F \circ E} \\ E^*(\eta) \\ \xrightarrow{F \circ E'} \end{array} \mathcal{D}$$

Definition 33 (opposite functor). We call the pullback by Swap the opposite functor, and denote it by $(-)^{\text{op}}$.

Theorem 34. For any simplicial sets K and S , we have

$$(K \star S)^{\text{op}} = S^{\text{op}} \star K^{\text{op}}.$$

Proof.

□

1.6.3 Cones and cocones

Definition 35 (left cone, right cone). let K be a simplicial set, and let $x \in K$.

- The left cone of x is the simplicial set

$$x^{\triangleleft} = \Delta^0 \star x.$$

- The right cone of x is the simplicial set

$$x^{\triangleright} = x \star \Delta^0.$$

Definition 36 (overcategory, undercategory). Let K and S be simplicial sets and let $p: K \rightarrow S$ be a morphism.

- The overcategory $S_{p/}$ is the category whose n -simplices are

$$(S_{p/})_n = \{\text{maps } f: K \star \Delta^n \rightarrow \mathcal{C} \text{ satisfying } F|_K = p\}.$$

- The undercategory $S_{/p}$ is the category whose n -simplices are

$$(S_{/p})_n = \{\text{maps } f: \Delta^n \star K \rightarrow \mathcal{C} \text{ satisfying } F|_K = p\}.$$

Theorem 37. There is a pair of adjoint functors

$$K \star - : \text{Set}_\Delta \leftrightarrow (\text{Set}_\Delta)_{/p} : /p$$

More explicitly, given a map $p: X \rightarrow \mathcal{C}$, there is a natural bijection

$$\text{Hom}_p(K \star X, \mathcal{C}) \simeq \text{Hom}(K, \mathcal{C}_{/p}),$$

Corollary 38. A map

For regular categories, it is possible to define colimits as limits in opposite categories; the two notions are equivalent. In infinity categories, this is also the case.

In the Kan model structure, trivial fibrations are Kan fibrations which are also homotopy equivalences. However, we have immediately that $W \cap \text{Fib} = \text{Cof}_\perp$, i.e. trivial Kan fibrations are precisely those morphisms of simplicial sets that have the right lifting property with respect to monomorphisms. However, since Cof is saturated, this is equivalent to having the left lifting property with respect to boundary inclusions.

Lemma 39. Let \mathcal{C} be a Kan complex, and $x \in \mathcal{C}$. Then $\mathcal{C}_{/x}$ and $\mathcal{C}_{x/}$ are contractible Kan complexes.

Proof. We need to show that we can solve lifting problems of the following form.

$$\begin{array}{ccc} \partial\Delta^n & \longrightarrow & \mathcal{C}_{/x} \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^n & \longrightarrow & \{*\} \end{array}$$

Consider

□

1.6.4 Initial objects

Definition 40 (initial object). Let \mathcal{C} be an ∞ -category and $x \in \mathcal{C}$ an object. We say that x is an initial object if the map

$$\mathcal{C}_{x/} \rightarrow \mathcal{C}; \quad \Delta^0 \star \Delta^n \rightarrow \Delta^n$$

is a trivial Kan fibration.

We would like initial/terminal objects to be objects that have an edge to each other object, and we would like this edge to be unique in some sense. Demanding uniqueness on the nose, however, would go against the spirit of the theory of infinity categories.

We have, however, the following.

Proposition 41. Let \mathcal{C} be an ∞ -category, and let x be an initial object in \mathcal{C} . Then for each other object $y \in \mathcal{C}$, there is an edge $x \rightarrow y$.

Furthermore, this edge is unique up to homotopy.

Proof. Consider the following solid lifting problem.

$$\begin{array}{ccc} \emptyset & \longrightarrow & \mathcal{C}_{x/} \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^0 & \xrightarrow{y} & \mathcal{C} \end{array}$$

The left-hand arrow is a boundary inclusion $\partial\Delta^0 \rightarrow \Delta^0$, so we can find a dashed lift, giving us a zero-simplex in $\mathcal{C}_{x/}$. But by definition, a zero simplex in $\mathcal{C}_{x/}$ is a 1-simplex $\sigma: \Delta^0 \star \Delta^0 \rightarrow \mathcal{C}$

such that $\sigma|_{\text{first factor}} = x$. Furthermore, the commutativity of the lower triangle tells us that $\sigma|_{\text{second factor}} = y$, i.e. σ is an edge $x \rightarrow y$.

Now suppose we have two edges $f, g: x \rightarrow y$. These give us the following solid lifting problem.

$$\begin{array}{ccc} \partial\Delta^1 & \xrightarrow{(f,g)} & \mathcal{C}_{x/} \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^1 & \xrightarrow{\text{id}_y} & \mathcal{C} \end{array}$$

Again, this is a boundary inclusion, giving us a dashed lift, i.e. a 2-simplex α in \mathcal{C} whose 0th vertex is x . The commutativity of the upper square tells us that $d_2\alpha = f$ and $d_0\alpha = g$, and the commutativity of the lower square tells us that $d_1\alpha = \text{id}_y$. Thus, we have the following simplex in \mathcal{C} .

$$\alpha = \begin{array}{ccc} & y & \\ f \nearrow & & \searrow \text{id}_y \\ x & \xrightarrow{g} & y \end{array}$$

But this provides precisely a homotopy $f \rightarrow g$. □

Corollary 42. Let \mathcal{C} be an ∞ -category. If x is an initial object in \mathcal{C} , then x is also initial in $\text{h}\mathcal{C}$.

One might expect that these conditions were equivalent. However, that is not the case; an object can be initial in $\text{h}\mathcal{C}$ without being initial in \mathcal{C} .

Example 43. Take any simply connected space X such that there exists a point $x \in X$ with $\pi_2(X, x)$ non-trivial. Then there is a 2-simplex $\alpha: \Delta^2 \rightarrow \mathbf{Sing}(X)$ with boundary $\partial\alpha = x$ which is not homotopy equivalent to x .

Since a 2-simplex with its zeroth vertex x is equivalently a map $\Delta^1 \rightarrow \mathcal{C}_{x/}$, this data organizes into the following lifting problem.

$$\begin{array}{ccc} \partial\Delta^2 & \xrightarrow{(x, \alpha)} & \mathbf{Sing}(X)_{x/} \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^2 & \xrightarrow{x} & \mathbf{Sing}(X) \end{array}$$

The solution to this lifting problem would give us a homotopy from α to x . However, by assumption, there is no such homotopy. Thus, x is not initial in $\mathbf{Sing}(X)$. However, every point $x \in X$ is initial in $\text{h}\mathbf{Sing}(X)$ because $\text{h}\mathbf{Sing}(X) \simeq \pi_{\geq 1}(X)$, and there is exactly one homotopy class of paths between any two points of X because X is simply connected.

Suppose we are given any two initial objects in an infinity category. We have seen that we can find an edge between them, and that given two edge between them, we can find a homotopy. In turn, given homotopies, we can find a higher homotopy between them, etc. One might therefore think of the set of all initial objects as forming the vertices of some polyhedron, and the filling conditions above telling us that the polyhedron in question has no gaps; that is, that it is homotopy equivalent to the point. This intuition turns out to be correct.

Theorem 44. Let \mathcal{C} be a simplicial set with at least one initial object, and let \mathcal{C}' be the simplicial subset consisting of all simplices all of whose vertices are initial objects in \mathcal{C} . Then \mathcal{C}' is a contractible Kan complex.

Proof. The statment that \mathcal{C}' is contractible means that the map $\mathcal{C}' \rightarrow *$ is a weak equivalence. The fact that it is a Kan complex means precisely that the map $\mathcal{C}' \rightarrow *$ is a fibration. Therefore, we need to check that the map $\mathcal{C}' \rightarrow *$ is a trivial fibration, i.e. has the right lifting property with respect to all boundary fillings.

$$\begin{array}{ccc} \partial\Delta^n & \longrightarrow & \mathcal{C}' \\ \downarrow & \nearrow & \\ \Delta^n & & \end{array}$$

For $n = 0$, the lifting problem demands that there exists at least one initial object, which is true by assumption.

For $n \geq 1$, consider an n -boundary $\beta: \partial\Delta^n \rightarrow \mathcal{C}'$. Denote by x the 0th vertex; by assumption, this is an initial object. By forgetting the 0th face, we can consider this as a boundary inclusion

$$\partial\Delta^{\{1,\dots,n\}} \rightarrow \mathcal{C}_{x/}$$

making the following diagram commute.

$$\begin{array}{ccc} \partial\Delta^{\{1,\dots,n\}} & \longrightarrow & \mathcal{C}_{x/} \\ \downarrow & & \downarrow \\ \Delta^{\{1,\dots,n\}} & \xrightarrow{d_0\beta} & \mathcal{C} \end{array}$$

However, we can find a lift, i.e. a map $\Delta^{\{1,\dots,n\}} \rightarrow \mathcal{C}_{x/}$, i.e. an n -simplex in \mathcal{C} which fills β . \square

1.6.5 Terminal objects

Definition 45 (terminal object). Let \mathcal{C} be an ∞ -category and $x \in \mathcal{C}$ an object. We say that x is a terminal object if the map

$$\mathcal{C}_{/x} \rightarrow \mathcal{C}; \quad \Delta^n \star \Delta^0 \mapsto \Delta^n$$

is a trivial Kan fibration.

The theory of terminal objects is completely dual to that of initial objects. For that reason, we list only results here.

Let \mathcal{C} be an infinity category.

- If y is a terminal object in \mathcal{C} , for each object $x \in \mathcal{C}$ there is an edge $x \rightarrow y$, which is unique up to homotopy.
- These homotopies are in turn unique up to homotopy.
- Terminal objects in an ∞ -category are always terminal in the homotopy category; however, the converse is not true.
- The simplicial subset of all simplices whose vertices are terminal is a contractible Kan complex.

1.6.6 limits and colimits

Definition 46 (limit cone). Let \mathcal{C} be an infinity category, K a simplicial set, and $p: K \rightarrow \mathcal{C}$ a diagram.

- A colimit cone for p is an initial object in the category $\mathcal{C}_{p/}$.
- A limit cone for p is a terminal object in the category $\mathcal{C}_{/p}$.

Suppose that I and \mathcal{D} are regular categories, and let $F: I \rightarrow \mathcal{D}$ be a functor. Passing to nerves, we get a functor $N(F)$ between infinity categories. Let us try and understand what a limit cone for $N(F)$ means in this context.

First, let us understand the zero-simplices of $N(\mathcal{D}_{/F})$. By definition, zero-simplices are given by maps

$$\sigma: \Delta^0 \star N(I) \rightarrow N(\mathcal{D}) \quad \text{such that } \sigma|_{N(I)} = N(F).$$

By the monoidality of the nerve with respect to the join, such a map is the nerve of a functor

$$\tilde{\sigma}: \{*\} \star I \rightarrow \mathcal{D} \quad \text{such that } \tilde{\sigma}|_I = F.$$

But this is precisely a natural transformation from the constant functor $\sigma|_{\{*\}} \rightarrow \mathcal{D}$ to F .

Similarly, a 1-simplex τ is a map

$$\tau: \Delta^1 \star N(I) \rightarrow N(\mathcal{D}) \quad \text{such that } \tau|_{N(I)} = N(F).$$

This is the nerve of a functor

$$\tilde{\tau}: \{\bullet \rightarrow *\} \star I \rightarrow \mathcal{D} \quad \text{such that } \tilde{\tau}|_I = F$$

Such functors are precisely morphisms of cones.

Corollary 47. Limit and colimit cones are unique up to contractible choice.

Proof. [Theorem 48](#). □

Theorem 48. Let S be a simplicial set, \mathcal{C} an infinity category, and $p: K \rightarrow \mathcal{C}$ a map of simplicial sets.

1. There is an equivalence of categories

$$(\mathcal{C}_{/p})^{\text{op}} \simeq (\mathcal{C}^{\text{op}})_{p^{\text{op}}/}.$$

2. The limit of p in \mathcal{C} is the colimit of p^{op} in \mathcal{C}^{op} .

Proof.

1. First we show the result on simplices. The n -simplices of $(\mathcal{C}^{\text{op}})_{p^{\text{op}}/}$ are

$$\{f: K^{\text{op}} \star \Delta^n \rightarrow \mathcal{C}^{\text{op}} \mid f|_{K^{\text{op}}} = p^{\text{op}}\} \simeq \{f^{\text{op}}: (\Delta^n)^{\text{op}} \star K \rightarrow \mathcal{C} \mid f^{\text{op}}|_K = p\}$$

But these are precisely the n -simplices of $(\mathcal{C}_{/p})^{\text{op}}$. Thus, the map on n -simplices is given by $f \mapsto f^{\text{op}}$.

1 Infinity categories

Let $f: [m] \rightarrow [n]$. We need to check that the following naturality square commutes.

$$\begin{array}{ccc}
 (\otimes \mathcal{C}_{p^{\text{op}}/}^{\text{op}})_n & \xrightarrow{\text{op}} & (\mathcal{C}_{/p})_n^{\text{op}} \\
 \downarrow & & \downarrow \\
 (\otimes \mathcal{C}_{p^{\text{op}}/}^{\text{op}})_m & \xrightarrow{\text{op}} & (\mathcal{C}_{/p})_m^{\text{op}}
 \end{array}$$

$f \mapsto f^{\text{op}}$
 $f \circ (\text{id}_{K^{\text{op}}} \star \phi) \mapsto f^{\text{op}} \circ (\phi^{\text{op}} \star \text{id}_K)$

But it does.

2. Let x be a colimit of p^{op} in \mathcal{C}^{op} . Then the map

$$\mathcal{C}_{x/}^{\text{op}} \rightarrow \mathcal{C}^{\text{op}}$$

is a trivial Kan fibration. Now consider the map

□