

# 1 The fibration zoo

If the reader thought that the theory of Kan fibrations was vibrant, they are in for a treat.

Let  $f: K \rightarrow S$  be a morphism of simplicial sets.

- We call the morphism  $f$  a Kan fibration if it has the right lifting property with respect to all horn inclusion  $\Lambda_i^n \rightarrow \Delta^n$  with  $0 \leq i \leq n$ .
- We call the morphism  $f$  a left fibration if it has the right lifting property with respect to all horn inclusion  $\Lambda_i^n \rightarrow \Delta^n$  with  $0 \leq i < n$ .
- We call the morphism  $f$  a right fibration if it has the right lifting property with respect to all horn inclusion  $\Lambda_i^n \rightarrow \Delta^n$  with  $0 < i \leq n$ .
- We call the morphism  $f$  an inner fibration if it has the right lifting property with respect to all horn inclusion  $\Lambda_i^n \rightarrow \Delta^n$  with  $0 < i < n$ .

We will use the following terminology.

- The morphism  $f$  is anodyne if it has the left-lifting property with respect to all Kan fibrations.
- The morphism  $f$  is left anodyne if it has the left-lifting property with respect to all left fibrations.
- The morphism  $f$  is right anodyne if it has the left-lifting property with respect to all right fibrations.
- The morphism  $f$  is inner anodyne if it has the left-lifting property with respect to all inner fibrations.

Unable to leave well enough alone, we now define the following additional notions.

**Definition 1** (Cartesian fibration). Let  $f: K \rightarrow S$  be a morphism of simplicial sets. The morphism  $f$  is a Cartesian fibration if the following conditions are satisfied.

- The morphism  $f$  is an inner fibration
- For every  $h: x \rightarrow y$  in  $S$  and  $\tilde{y} \in K_0$  there is a  $f$ -Cartesian morphism  $\tilde{h}: \tilde{x} \rightarrow \tilde{y}$  such that  $f(\tilde{h}) = h$ .
- A  $f$ -Cartesian edge in  $K$  is an edge  $f: x \rightarrow y$  such that the map

$$K_{/f} \rightarrow K_{/y} \times_{S_{/f(y)}} S_{/p(f)}$$

is a trivial Kan fibration.

**Theorem 2.** Let  $f$  be a trivial fibration. Then  $f$  is a Kan fibration.

## 1 The fibration zoo

*Proof.* Let  $f$  be a trivial fibration, and consider the following lifting problem.

$$\begin{array}{ccc} \Delta_i^n & \longrightarrow & K \\ \downarrow & & \downarrow \\ \Delta^n & \longrightarrow & S \end{array}$$

Restricting to the face opposite  $i$ , we find the lifting problem

$$\begin{array}{ccc} \partial\Delta^{\{1,\dots,n\}} & \longrightarrow & K \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^{\{1,\dots,n\}} & \longrightarrow & S \end{array}$$

which has a dashed solution. This allows us to turn our map  $\Delta_i^n \rightarrow K$  to a map  $\partial\Delta^n \rightarrow K$  making the following diagram commute.

$$\begin{array}{ccc} \partial\Delta^n & \longrightarrow & K \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^n & \longrightarrow & S \end{array}$$

which has a dashed lift. □

*Note 3.* A map  $f: X \rightarrow Y$  is a left fibration if and only if  $f^{\text{op}}: X^{\text{op}} \rightarrow Y^{\text{op}}$  is a right fibration. Thus, it suffices to study left fibrations; the theory of right fibrations is dual.

**Lemma 4.** Let  $p: \mathcal{C} \rightarrow \mathcal{D}$  be a left fibration of infinity categories, and let  $\tilde{f}: \bar{x} \rightarrow p(y)$  be an equivalence in  $\mathcal{D}$ , and  $y \in \mathcal{C}_0$ . Then there exists  $f: x \rightarrow y$  so that  $p(f) = \tilde{f}$ .

**Lemma 5.** Consider morphisms of simplicial sets as follows

$$\bar{K} \hookrightarrow K \xrightarrow{p} X \xrightarrow{q} S$$

where  $q$  is an inner fibration. Let  $r = p \circ q$ ,  $\bar{p} = p|_{\bar{K}}$ ,  $\bar{r} = r|_{\bar{K}}$ . Then the induced map

$$S_{p/} \rightarrow X_{\bar{p}/} \times_{S_{\bar{p}/}} S_{r/}$$

is an inner fibration. If  $q$  is a left fibration, then the map is a left fibration.

**Corollary 6.** Let  $\mathcal{C}$  be an infinity category, and let  $p: K \rightarrow \mathcal{C}$  be a diagram in  $\mathcal{C}$ . Then the map  $\mathcal{C}_{p/} \rightarrow \mathcal{C}$  is a left fibration, and  $\mathcal{C}_{p/}$  is an infinity category.

*Proof.* Take  $X = \mathcal{C}$ ,  $S = \Delta^0$ , and  $\bar{K} = \emptyset$ . This tells us that the map

$$\mathcal{C}_{p/} \rightarrow \mathcal{C}$$

is a left fibration. Thus, the composition

$$\mathcal{C}_{p/} \rightarrow \mathcal{C} \rightarrow *$$

is the composition of a left fibration and an inner fibration, so is itself an inner fibration. □

**Corollary 7.** Let  $f: x \rightarrow y$  be a morphism in an  $\infty$ -category  $\mathcal{C}$ . Then  $f$  is an equivalence if and only if for every  $n \geq 2$  and left horn  $h: \Lambda_0^n \rightarrow \mathcal{C}$ , with  $h|_{\{0,1\}} = f$ ,  $h$  admits a filler.

**Theorem 8.** The following are equivalent for an  $\infty$ -category  $\mathcal{C}$ .

- The homotopy category  $h\mathcal{C}$  is a groupoid.
- Every left horn in  $\mathcal{C}$  has a filler.
- Every right horn has a filler.
- $\mathcal{C}$  is a Kan complex.

**Corollary 9.** All fibers of a left fibration are Kan complexes.

## 1.1 The classical Grothendieck construction

**Definition 10** (Grothendieck construction). Let  $\mathcal{D}$  be a small category, and let  $\chi: \mathcal{D} \rightarrow \mathbf{Grpd}$ . The Grothendieck construction consists of the following data the following data.

- A category  $\mathcal{C}_\chi$
- A functor  $\pi: \mathcal{C}_\chi \rightarrow \mathcal{D}$

The category  $\mathcal{C}_\chi$  is as follows.

- **Objects:** Pairs  $(d, x)$ , where  $d \in \mathcal{D}$  and  $x \in \chi(d)$ .
- **Morphisms:** A morphism  $(d, x) \rightarrow (d', x')$  consists of
  - A morphism  $f: d \rightarrow d'$  in  $\mathcal{D}$
  - An isomorphism  $\alpha: \chi(f)(x) \simeq x'$

The functor  $\pi$  is the forgetful functor to  $\mathcal{D}$ .

The Grothendieck construction can be understood in the following way. The functor  $\chi$  gives us a way of associating a groupoid to each object  $d \in \mathcal{D}$ . Each of these groupoids is itself a category, and has its own objects. The idea of the Grothendieck construction is to throw all of these objects together into one big category. To keep track of where each object  $x$  comes from, we think of  $x$  as ordered pairs  $(d, x)$ , where the  $x \in \chi(d)$ . The functor  $\pi$  simply remembers where each  $x$  came from.

For two objects  $x$  and  $\tilde{x}$  from the same  $\chi(d)$ , the morphisms  $(d, x) \rightarrow (d, \tilde{x})$  are simply the morphisms between  $x$  and  $\tilde{x}$ . For an object  $x'$  from a different groupoid  $\chi(d')$ , the morphisms  $(d, x) \rightarrow (d', x')$  come from translating, using  $\chi$  and a morphism  $d \rightarrow d'$ , morphisms in  $\chi(d)$  into morphisms in  $\chi(d')$ .

**Example 11.** Let  $G$  and  $H$  be a groups, and  $\mathbf{B}G$  and  $\mathbf{B}H$  their delooping groupoids. Let  $\phi: G \rightarrow H$  be a group homomorphism, or equivalently a functor  $\mathbf{B}G \rightarrow \mathbf{B}H$ . This can also be taken to be a functor  $\mathbf{B}G \rightarrow \mathbf{Grpd}$ .

The Grothendieck construction consists of a category with one object, whose morphisms consist of an element  $g \in G$

There is also a contravariant version of the Grothendieck construction, consisting of the following.

**Definition 12** (contravariant Grothendieck construction). For any functor  $\mathcal{D}^{\text{op}} \rightarrow \mathbf{Grpd}$ , the contravariant Grothendieck construction consists of the following.

1. A category  $\mathcal{C}^{\chi}$
2. A functor  $\pi: \mathcal{C}^{\chi} \rightarrow \mathcal{D}$ .

The category  $\mathcal{C}^{\chi}$  is as follows.

- The objects of  $\mathcal{C}^{\chi}$  are pairs  $(d, x)$ , with  $d \in \mathcal{D}$  and  $x \in \chi(d)$ .
- The morphisms  $(d', x') \rightarrow (d, x)$  are pairs  $(f, \alpha)$ , where  $f: d' \rightarrow d$  and  $\alpha: x' \rightarrow F(f)(x)$ .

There is the following relationship between the covariant and contravariant Grothendieck constructions.  $\mathcal{C}^{\chi} = \mathcal{C}_{\chi}^{\text{op}}$ .

**Example 13.** Let  $X: \mathcal{D}^{\text{op}} \rightarrow \mathbf{Set}$  be a diagram of sets, which we interpret as discrete groupoids. Then there is an equivalence of categories between the contravariant construction and the slice category  $\mathcal{D}/X$ .

The objects of  $\mathcal{D}/X$  are pairs  $(d, \alpha)$ , where  $d \in \mathcal{D}$  and  $\alpha: \text{Hom}_{\mathcal{D}}(-, d) \Rightarrow X$ . The morphisms  $(d, \alpha) \rightarrow (d', \beta)$  are morphisms  $f: d \rightarrow d'$  such that the diagram

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(-, d) & \xrightarrow{f \circ -} & \text{Hom}_{\mathcal{D}}(-, d') \\ & \searrow \alpha & \swarrow \beta \\ & X & \end{array} \quad (1.1)$$

commutes.

Consider any such morphism  $f$ . The above diagram commutes if and only if it commutes component-wise, i.e. if for each  $c$ , the diagram below commutes.

$$\begin{array}{ccc} \text{Hom}_{\mathcal{D}}(c, d) & \xrightarrow{f \circ -} & \text{Hom}_{\mathcal{D}}(c, d') \\ & \searrow \alpha_c & \swarrow \beta_c \\ & X(c) & \end{array} \quad (1.2)$$

In particular, it commutes when  $c = d$ . Following the identity  $\text{id}_d$  around, we find that morphisms  $f$  making the above diagram commute satisfy the following equality.

$$\begin{array}{ccc} \text{id}_d & \xrightarrow{\quad} & f \\ & \searrow & \swarrow \\ & \alpha_d(\text{id}_d) = \beta_d(f) & \end{array}$$

By the naturality of the Yoneda embedding, we have that

$$\beta_d(f) = X(f)(\beta_{d'}(\text{id}_{d'})),$$

so

$$\alpha_d(\text{id}_d) = X(f)(\beta_{d'}(\text{id}_{d'})). \quad (1.3)$$

We can specify a functor  $\mathcal{D}/X \rightarrow \mathcal{C}^X$  on objects by sending

$$(d, \alpha) \mapsto (d, \alpha_b(\text{id}_b)),$$

and on morphisms  $f: (d, \alpha) \rightarrow (d', \beta)$  commute by sending

$$f \mapsto (f, \text{id}_{X(f)(\beta_{d'}(\text{id}_{d'}))}).$$

We are allowed to use the above identity map precisely because it is an isomorphism

$$\alpha_d(\text{id}_d) \rightarrow X(f)(\beta_{d'}(\text{id}_{d'}))$$

as shown in [Equation 1.3](#).

The fully faithfulness of the Yoneda embedding implies that the map between the objects of  $\mathcal{D}/X$  and the objects of  $\mathcal{C}^X$  is a bijection. We will now construct a functor  $\mathcal{C}^X \rightarrow \mathcal{D}/X$ , defined on objects to be the inverse of the one defined above, i.e. sending a pair  $(d, x)$  to the pair  $(d, \alpha^x)$ , where  $\alpha^x$  is the natural transformation uniquely specified by  $\alpha_d^x(\text{id}_d) \mapsto x \in X(d)$ .

Our inverse functor acts on any morphism  $(f, \phi)$  by forgetting  $\phi$ . It only remains to check that this provides a well-defined map of morphisms, as  $f$  has to make [Diagram 1.1](#) commute. We can check that this is so by following a general morphism  $g: c \rightarrow d$  around [Diagram 1.2](#), finding the following condition.

$$\begin{array}{ccc} g & \xrightarrow{\quad} & f \circ g \\ & \searrow & \swarrow \\ & \alpha_c(g) \stackrel{!}{=} \beta_c(f \circ g) & \end{array}$$

By naturality of the Yoneda embedding, we have

$$\alpha_c(g) = (Xg)(\alpha_d(\text{id}_d)), \quad \beta_c(f \circ g) = X(f \circ g)(\beta_{d'}(\text{id}_{d'})) = (Xg \circ Xf)(\beta_{d'}(\text{id}_{d'})),$$

so our condition equivalently reads

$$Xg(\alpha_d(\text{id}_d)) = Xg((Xf)(\beta_{d'}(\text{id}_{d'}))).$$

This is true because by definition of  $f$ ,

$$\alpha_d(\text{id}_d) = (Xf)(\beta_{d'}(\text{id}_{d'})).$$

The two functors defined in this way are manifestly each other's inverses, hence provide an isomorphism of categories, which is stronger than what we are looking for.

### 1.1.1 When can we invert the classical Grothendieck construction?

Suppose we are given a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  exhibiting  $\mathcal{C}$  as cofibered in groupoids over  $\mathcal{D}$ . By [Lemma 27](#), each of the fibers  $\mathcal{C}_d$  is a groupoid. We may hope that we could reverse the Grothendieck construction by defining a functor  $\mathcal{D} \rightarrow \mathbf{Grpd}$  defined on objects by  $d \mapsto \mathcal{C}_d$ .

We almost can. For any morphism  $f: d \rightarrow d'$  in  $\mathcal{D}$ , we can build functor  $f_!: \mathcal{C}_d \rightarrow \mathcal{C}_{d'}$  as follows.

Let  $f: d \rightarrow d'$ , and let  $c \in \mathcal{C}_d$ . By property 2 in [Definition 17](#), we can find an object  $c' \in \mathcal{C}_{d'}$  and a morphism  $\bar{f}: c \rightarrow c'$  in  $\mathcal{C}$  such that  $F(\bar{f}) = f$ . Define  $f_!(c) = c'$ , and  $f_!()$

The above discussion means that the Grothendieck construction defines the following correspondence.

$$\left\{ \begin{array}{l} F: \mathcal{C} \rightarrow \mathcal{D} \text{ exhibits } \mathcal{C} \\ \text{as cofibered in groupoids} \end{array} \right\} \xrightarrow{\cong} \left\{ \begin{array}{l} \text{pseudofunctors} \\ \mathcal{D} \rightarrow \mathbf{Grpd} \end{array} \right\}$$

### 1.1.2 Fibrance and cofibrance in groupoids

**Example 14.** The Grothendieck construction

$$\pi: \mathcal{C}_\chi \rightarrow \mathcal{D}$$

exhibits  $\mathcal{D}$  as cofibered over groupoids. We simply transcribe the definition, making the necessary changes as we go along.

In order for  $\pi$  to exhibit  $\mathcal{C}_\chi$  as cofibered in groupoids over  $\mathcal{D}$ , we have to check the following.

1. For every  $(d, x) \in \mathcal{C}_\chi$  and every morphism

$$\eta: d \rightarrow d' \quad \text{in } \mathcal{D}$$

there exists a lift

$$\tilde{\eta} = (f, \alpha): (d, x) \rightarrow (d', x') \quad \text{in } \mathcal{C}_\chi$$

such that

$$\pi(\tilde{\eta}) = f: d \rightarrow d' = \eta.$$

2. For every morphism  $\eta: (x, d) \rightarrow (x', d')$  in  $\mathcal{C}$  and every object  $(x'', d'') \in \mathcal{C}$ , the map

$$\mathcal{C}_\chi((d', x') \rightarrow (d'', x'')) \rightarrow \mathcal{C}_\chi((d, x), (d'', x'')) \times_{\mathcal{D}(d, d'')} \mathcal{D}(d', d'')$$

is a bijection.

What do these mean?

1. We need to construct from our map  $\eta: d \rightarrow d'$  an isomorphism

$$\alpha: \chi(f)(x) \simeq x'$$

for some  $x'$ . We can pick  $x' = \chi(f)(x)$ . Then  $\tilde{\eta} = (f, \alpha)$  does the job. Indeed,  $\pi(f, \alpha) = f$  as required.

2. Draw a picture. Given the data of

- A morphism

**Example 15.** The contravariant Grothendieck construction exhibits  $\mathcal{C}^\chi \rightarrow \mathcal{D}$  as fibered in groupoids over  $\mathcal{D}$ .

To see this, we need to make the following checks.

- **Check:** For every  $(d, x) \in \mathcal{C}^\chi$ , and every morphism  $\eta: d' \rightarrow d$  in  $\mathcal{D}$ , we should be able to find some object  $x' \in \chi(d')$  and a morphism  $(f, \alpha): (d', x') \rightarrow (d, x)$ .

**Solution:** We can pick  $f = \eta$ ,  $x' = \chi(f)(x)$ , and  $\alpha = \text{id}_{x'}$ .

- **Check:** Fixing any diagram

$$\begin{array}{ccc} & (d', x') & \\ & \searrow (g, \alpha) & \\ (d'', x'') & \xrightarrow{(\eta, \beta)} & (d, x) \end{array}$$

in  $\mathcal{C}^\chi$ , for any morphism  $f: d'' \rightarrow d'$  making the below left diagram commute, we should be able to find a morphism  $(f, \gamma): (d'', x'') \rightarrow (d', x')$  making the below right diagram commute.

$$\begin{array}{ccc} & d' & \\ f \nearrow & & \searrow g \\ d'' & \xrightarrow{\eta} & d \end{array} \rightsquigarrow \begin{array}{ccc} & (d', x') & \\ (f, \gamma) \nearrow & & \searrow (g, \alpha) \\ (d'', x'') & \xrightarrow{(\eta, \beta)} & (d, x) \end{array}$$

**Solution:** We need a morphism  $\gamma: x'' \rightarrow \chi(f)(x')$  making the above diagram commute.

Using the data available to us, we can create the following diagram in  $\chi(d'')$ .

$$\begin{array}{ccc} & \chi(f)(x') & \\ & \searrow \chi(f)(\alpha) & \\ x'' & \xrightarrow{\beta} & \chi(\eta)(x) \end{array}$$

Note that we are entitled to draw  $\chi(f)(\alpha)$  terminating at  $\chi(\eta)(x)$  because

$$\chi(f)(\chi(g)(x)) = \chi(g \circ f)(x) = \chi(\eta)(x).$$

Because  $\chi(d'')$  is a groupoid, we get the morphism  $\gamma$  via

$$\gamma = (\chi(f)(\alpha))^{-1} \circ \beta.$$

## 1.2 Left fibrations

**Lemma 16.** Consider simplicial sets and morphisms as follows.

$$\begin{array}{ccccc}
 & & r_0 & & \\
 & \curvearrowright & & \curvearrowright & \\
 K_0 & \hookrightarrow & K & \xrightarrow{p} & X & \xrightarrow{q} & S \\
 & \curvearrowleft & & \curvearrowleft & \\
 & & p_0 & & 
 \end{array}$$

We can translate a lifting problem of the first type into a lifting problem of the second.

$$\begin{array}{ccc}
 A & \xrightarrow{\quad} & X_{p/} \\
 \downarrow & \nearrow \text{dashed} & \downarrow \\
 B & \xrightarrow{\quad} & X_{p_0/} \times_{S_{r_0/}} S_{r/}
 \end{array}
 \qquad
 \begin{array}{ccc}
 K \star A \amalg_{K_0 \star A} K_0 \star B & \xrightarrow{\quad} & X \\
 \downarrow & \nearrow \text{dashed} & \downarrow \\
 K \star B & \xrightarrow{\quad} & S
 \end{array}$$

## 1.3 Left fibrations and cofibrancy in groupoids

Since left fibrations are of the form

$$\{\Lambda_i^n \hookrightarrow \Delta^n \mid n \geq 1, 0 \leq i < n\}_\perp,$$

Proposition ?? implies that they are closed under pullback.

Let  $F: X \rightarrow Y$  be a left fibration of simplicial sets. For  $x \in X$ , the fiber

### 1.3.1 Cofibrancy in groupoids

**Definition 17** (cofibrant in groupoids). Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor. We say that  $f$  exhibits  $\mathcal{C}$  as cofibrant in groupoids over  $\mathcal{D}$  if the following conditions are satisfied.

1. For every  $c \in \mathcal{C}$  and every morphism

$$\eta: F(c) \rightarrow d \quad \text{in } \mathcal{D}$$

there exists a morphism

$$\tilde{\eta}: c \rightarrow \tilde{d} \quad \text{in } \mathcal{C}$$

such that  $F(\tilde{\eta}) = \eta$ .

2. For every morphism  $\eta: c \rightarrow c'$  in  $\mathcal{C}$  and every object  $c'' \in \mathcal{C}$ , the map

$$\mathcal{C}(c', c'') \rightarrow \mathcal{C}(c, c'') \times_{\mathcal{D}(F(c), F(c''))} \mathcal{D}(F(c'), F(c'')).$$

is a bijection.



Let us unravel what the second condition means. The second set is the following pull-back.

$$\begin{array}{ccc} * & \longrightarrow & \mathcal{D}(F(c'), F(c'')) \\ \downarrow & & \downarrow \eta \\ \mathcal{C}(c, c'') & \xrightarrow{F} & \mathcal{D}(F(c), F(c'')) \end{array}$$

Its elements are pairs

$$(f, \bar{g}), \quad \text{where } f: c \rightarrow c'' \quad \text{and} \quad \bar{g}: F(c') \rightarrow F(c'')$$

such that

$$F(f) = \bar{g} \circ F(\eta).$$

The bijection above is defined as follows. Starting on the LHS with a morphism  $h: c' \rightarrow c''$ , build the following commutative diagram.

$$\begin{array}{ccc} & c' & \\ \eta \nearrow & & \searrow h \\ c & \xrightarrow{h \circ \eta} & c'' \end{array}$$

We can read off a map  $f = h \circ \eta: c \rightarrow c''$  and a map  $\bar{g} = F(h): F(c') \rightarrow F(c'')$ , and indeed  $F(f) = \bar{g} \circ F(\eta)$ , giving us an element of the RHS.

The fact that this is a bijection means that the above operation is invertible. Fixing a morphism  $\eta: c \rightarrow c'$  and an object  $c''$ , whenever we have the following data:

- A morphism  $g: c \rightarrow c''$  as follows

$$\begin{array}{ccc} & c' & \\ \eta \nearrow & & \\ c & \xrightarrow{\bar{g}} & c'' \end{array}$$

- A map  $\bar{h}: F(c') \rightarrow F(c'')$  making the following diagram commute

$$\begin{array}{ccc} & F(c') & \\ F(\eta) \nearrow & & \searrow \bar{h} \\ F(c) & \xrightarrow{F(g)} & F(c'') \end{array}$$

we can lift  $\bar{h}$  to a morphism  $h$  in  $\mathcal{C}$  making the following diagram commute.

$$\begin{array}{ccc} & c' & \\ \eta \nearrow & & \searrow h \\ c & \xrightarrow{g} & c'' \end{array}$$

## 1 The fibration zoo

**Lemma 18.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor exhibiting  $\mathcal{C}$  as cofibered in groupoids over  $\mathcal{D}$ . Let  $f$  be a morphism in  $\mathcal{C}$  such that  $F(f) = \text{id}_d$ . Then  $f$  is an isomorphism.

*Proof.* Let  $f: c \rightarrow c'$  in  $\mathcal{C}_d$  be a morphism such that  $F(f) = \text{id}_d$ . Then we can form the following commutative diagram.

$$\begin{array}{ccc} & c' & \\ f \nearrow & & \\ c & \xrightarrow{\text{id}} & c \end{array}$$

When mapped to  $\mathcal{D}$ , we find the following very simple diagram, with a dashed filler.

$$\begin{array}{ccc} & d & \\ \text{id} \nearrow & \text{--- id} \searrow & \\ d & \xrightarrow{\text{id}} & d \end{array}$$

Lifting this gives us a left inverse  $\tilde{f}$  to  $f$ . Pulling the same trick with  $\tilde{f}$  gives us a left inverse  $\hat{f}$  to  $\tilde{f}$ . Then

$$\begin{aligned} f \circ \tilde{f} &= \text{id} \\ \hat{f} \circ f \circ \tilde{f} &= \hat{f} \\ \hat{f} &= f, \end{aligned}$$

so  $\tilde{f}$  is both a left and a right inverse for  $f$ , i.e.  $f$  is an isomorphism.  $\square$

**Corollary 19.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  exhibit  $\mathcal{C}$  as cofibered in groupoids over  $\mathcal{D}$ . Then the fibers of  $F$  are groupoids.

*Proof.* The morphisms in the fiber  $\mathcal{C}_d$  are precisely those which are mapped to the identity  $\text{id}_d$ .  $\square$

**Lemma 20.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor exhibiting  $\mathcal{C}$  as cofibered in groupoids over  $\mathcal{D}$ . Let  $f$  be a morphism in  $\mathcal{C}$  such that  $F(f)$  is an equivalence. Then  $f$  is an equivalence.

**Definition 21** (coCartesian morphism). Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor of small categories. A morphism  $f: c \rightarrow c'$  is said to be F-coCartesian if, for every  $c'' \in \mathcal{C}$ , the map

$$\mathcal{C}(c', c'') \rightarrow \mathcal{C}(c, c'') \times_{\mathcal{D}(F(c), F(c''))} \mathcal{D}(F(c'), F(c''))$$

is a bijection.

**Definition 22** (coCartesian fibration). A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is said to be a coCartesian fibration if, for each  $x \in \mathcal{C}$  and every morphism  $\tilde{f}: F(x) \rightarrow y$  in  $\mathcal{D}$ , there exists a coCartesian morphism  $f: x \rightarrow \tilde{y}$  such that  $F(f) = \tilde{f}$ .

The reader has probably guessed the next sentence: “There is a dual notion.”

**Definition 23** (fibrant in groupoids). Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor. We say that  $f$  exhibits  $\mathcal{C}$  as fibrant in groupoids over  $\mathcal{D}$  if the following conditions are satisfied.

1. For every  $c' \in \mathcal{C}$  and every morphism

$$\eta: d \rightarrow F(c') \quad \text{in } \mathcal{D}$$

there exists a morphism

$$\tilde{\eta}: \tilde{d} \rightarrow c' \quad \text{in } \mathcal{C}$$

such that  $F(\tilde{\eta}) = \eta$ .

2. For every morphism  $\eta: c \rightarrow c'$  in  $\mathcal{C}$  and every object  $c'' \in \mathcal{C}$ , the map

$$\mathcal{C}(c', c'') \rightarrow \mathcal{C}(c, c'') \times_{\mathcal{D}(F(c), F(c''))} \mathcal{D}(F(c'), F(c'')).$$

is a bijection.

## 1.4 The starred smash

**Definition 24** (starred smash). Let  $f: A \rightarrow A', g: B \rightarrow B'$  be two morphisms of simplicial sets. Define a morphism  $f \overset{\star}{\wedge} g$  using the universal property for pushouts as follows.

$$\begin{array}{ccc}
 A \star B & \xrightarrow{f \star \text{id}} & A' \star B \\
 \text{id} \star g \downarrow & & \downarrow \\
 A \star B' & \hookrightarrow & A' \star B \amalg_{A \star B} A \star B' \\
 & & \searrow f \overset{\star}{\wedge} g \\
 & & A' \star B'
 \end{array}$$

$f \star \text{id}$  (curved arrow from  $A \star B'$  to  $A' \star B'$ ) and  $\text{id} \star g$  (curved arrow from  $A' \star B$  to  $A' \star B'$ )

**Proposition 25.** We have the following facts.

1. If  $f$  is right anodyne, then  $f \overset{\star}{\wedge} g$  and  $f \wedge g$  are inner anodyne.
2. If  $f$  is left anodyne, then  $f \overset{\star}{\wedge} g$  and  $f \wedge g$  are left anodyne.
3. If  $f$  is anodyne, then  $f \overset{\star}{\wedge} g$  is a Kan fibration.

*Proof.* Consider the set

$$\mathcal{S} = \{f \mid f \overset{\star}{\wedge} g \text{ is inner anodyne}\}.$$

## 1 The fibration zoo

That is,  $f$  consists of precisely those maps such that the lifting

$$\begin{array}{ccc} A_0 \star B \coprod_{A_0 \star B_0} A \star B_0 & \longrightarrow & X \\ f \wedge q \downarrow & \nearrow & \downarrow \text{inner fibration} \\ A \star B & \longrightarrow & Y \end{array}$$

has a solution

By a result analogous to Odious Lemma ??, we have immediately that such a lifting problem is equivalent to the following.

□

We have the following ‘perpendicular statements’ (whatever that means).

**Fact 26.** Consider morphisms of simplicial sets as follows

$$\bar{K} \hookrightarrow K \xrightarrow{p} X \xrightarrow{q} S$$

where  $q$  is an inner fibration. Let  $r = p \circ q$ ,  $\bar{p} = p|_{\bar{K}}$ ,  $\bar{r} = r|_{\bar{K}}$ . Denote the

$$S/p \rightarrow X/\bar{p} \times_{S/\bar{p}} S/r.$$

### 1.4.1 Cofibration in groupoids

We would like to find the correct  $\infty$ -categorical notion of being cofibered in groupoids.

**Lemma 27.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor exhibiting  $\mathcal{C}$  as cofibered in groupoids over  $\mathcal{D}$ . Then for each  $d \in \mathcal{D}$ , the fiber  $\mathcal{C}_d$  defined by the pullback below is a groupoid.

$$\begin{array}{ccc} \mathcal{C}_d & \longrightarrow & \mathcal{C} \\ \downarrow & & \downarrow F \\ \{d\} & \hookrightarrow & \mathcal{D} \end{array}$$

*Proof.* Every morphism in  $\mathcal{C}_d$  is by definition mapped to  $\text{id}_d$ , so by

□

**Theorem 28.** Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  be a functor. Then  $F$  exhibits  $\mathcal{C}$  as cofibered in groupoids over  $\mathcal{D}$  if and only if  $N(F): N(\mathcal{C}) \rightarrow N(\mathcal{D})$  is a left fibration.

*Proof.* Suppose  $F$  is a left fibration. We examine some horn filling conditions in small degrees. First, consider a horn filling

$$\begin{array}{ccc} \Lambda_0^1 & \longrightarrow & N(\mathcal{C}) \\ \downarrow & & \downarrow \\ \Delta^1 & \longrightarrow & N(\mathcal{D}) \end{array}$$

The commutativity conditions tell us that we have the following solid diagrams in  $\mathcal{C}$  and  $\mathcal{D}$  respectively.

$$x \xrightarrow{F} F(x) \xrightarrow{\bar{f}} \bar{y}$$

A dashed lift

$$\begin{array}{ccc} \Lambda_0^1 & \longrightarrow & N(\mathcal{C}) \\ \downarrow & \nearrow \text{dashed} & \downarrow \\ \Delta^1 & \longrightarrow & N(\mathcal{D}) \end{array}$$

gives us the following dashed data.

$$x \dashrightarrow^f y \xrightarrow{F} F(x) \xrightarrow{\bar{f}} \bar{y}$$

Thus, the first condition of [Definition 17](#) is satisfied.

Similarly, a horn filling problem  $\Lambda_0^2 \rightarrow \Delta^2$  gives us the following solid data, while a lift gives us the dashed data.

$$\begin{array}{ccc} & y & \\ f \nearrow & & \searrow g \\ x & \xrightarrow{h} & z \end{array} \quad \begin{array}{ccc} & F(y) & \\ F(f) \nearrow & & \searrow \bar{g} \\ F(x) & \xrightarrow{F(h)} & F(z) \end{array}$$

This says that there exists a lift as required, but not that it is unique.

Now consider two possible lifts  $\bar{g}$  and  $\bar{g}'$ , arranged into a  $(3, 0)$ -horn as follows.

$$\begin{array}{ccc} x & \xrightarrow{f} & y \\ f \downarrow & \searrow h & \downarrow g \\ y & \xrightarrow{g'} & z \end{array} \quad \begin{array}{ccc} x & \xrightarrow{f} & y \\ f \downarrow & \searrow \text{id}_y & \downarrow g \\ y & \xrightarrow{g'} & z \end{array}$$

These correspond to a 3-simplex in  $\mathcal{D}$  as follows.

$$\begin{array}{ccc} F(x) & \xrightarrow{F(f)} & F(y) \\ F(f) \downarrow & \searrow F(h) & \downarrow \bar{g}' \\ F(y) & \xrightarrow{\bar{g}} & F(z) \end{array} \quad \begin{array}{ccc} F(x) & \xrightarrow{F(f)} & F(y) \\ F(f) \downarrow & \searrow \text{id}_{F(y)} & \downarrow \bar{g}' \\ F(y) & \xrightarrow{\bar{g}} & F(z) \end{array}$$

The lifting properties say that we can fill the above horn, telling us that  $g = g'$ , i.e. that the lift is unique.  $\square$

## 1.5 The infinity-categorical Grothendieck construction

**Definition 29** (fiber over a point). Let  $f: X \rightarrow S$  be a map of simplicial sets, and let  $s \in S$ . The fiber over  $s$ , denoted  $X_s$  is the following pullback.

$$\begin{array}{ccc} X_s & \longrightarrow & X \\ \downarrow & & \downarrow f \\ \{s\} & \longrightarrow & S \end{array}$$

**Theorem 30.** Let  $p: X \rightarrow S$  be a left fibration of simplicial sets. Then the assignment

$$s \mapsto X_s$$

extends to a functor  $\mathbf{h}S \rightarrow \mathbf{h}\mathcal{S}$ , where  $\mathcal{S}$  is the category of spaces (Definition ??).

*Proof.* By [Corollary 9](#), each fiber  $X_s$  is a Kan complex.

Now, let  $f: s \rightarrow s'$  be any morphism in  $S$ . Consider the following lifting problem.

$$\begin{array}{ccc} \{0\} \times X_s & \xrightarrow{\quad} & X \\ \downarrow & \nearrow Lf & \downarrow p \\ \Delta^1 \times X_s & \xrightarrow{\quad} \Delta^1 \xrightarrow{f} & S \end{array}$$

The left-hand morphism is left anodyne because it is the smash product

$$(\{0\} \hookrightarrow \Delta^1) \wedge^* (\emptyset \hookrightarrow X_s),$$

so we can find a dashed lift. This gives us a map  $Lf: \Delta^1 \times X_s \rightarrow X$  such that  $Lf_{\{0\}} = \text{id}_{X_s}$ . Restricting to  $\{1\} \times X_s$ , we get a map  $Lf_{\{1\} \times X_s}: \{1\} \times X_s \rightarrow X_{s'}$ , giving us a map

$$\tilde{f}: X_s \rightarrow X_{s'}.$$

The functor  $\mathbf{h}S \rightarrow \mathbf{h}\mathcal{S}$  sends  $[f] \mapsto [\tilde{f}]$ . It remains to check that this is well-defined, i.e. respects  $\square$

## 1.6 More model structures

Let  $S$  be a simplicial set.

- Denote by  $(\mathbf{Set}_\Delta)_S$  the category of simplicial sets over  $S$ .
- Denote by  $\mathbf{Set}_\Delta^{\mathcal{C}(S)}$  the category of functors of simplicial categories from  $\mathcal{C}(S) \rightarrow \mathbf{Set}_\Delta$ .

**Theorem 31.** There is a Quillen equivalence

$$(\mathbf{Set}_\Delta)_S \leftrightarrow \mathbf{Set}_\Delta^{\mathcal{C}(S)}$$

where  $(\mathbf{Set}_\Delta)_S$  is taken to have the covariant model structure, and  $\mathbf{Set}_\Delta^{\mathcal{C}(S)}$  is taken to have the projective model structure.

## 1.7 Some correspondences

$$\begin{array}{ccccc}
 \left\{ \begin{array}{c} \text{functors } \mathcal{C}^{\mathcal{X}} \rightarrow \mathcal{D} \\ \text{fibered in groupoids} \end{array} \right\} & \longleftrightarrow & \left\{ \begin{array}{c} \text{pseudofunctors} \\ \mathcal{D} \rightarrow \mathbf{Grpd} \end{array} \right\} & \longleftrightarrow & \left\{ \begin{array}{c} \text{functors } \mathcal{C} \rightarrow \mathcal{D} \\ \text{cofibered in groupoids} \end{array} \right\} \\
 \left\{ \begin{array}{c} \text{coCartesian fibrations} \\ \mathcal{C}^{\mathcal{X}} \rightarrow \mathcal{D} \end{array} \right\} & \longleftrightarrow & \left\{ \begin{array}{c} \text{pseudofunctors} \\ \mathcal{D} \rightarrow \mathbf{Grpd} \end{array} \right\} & \longleftrightarrow & \left\{ \begin{array}{c} \text{coCartesian fibrations} \\ \mathcal{C}^{\mathcal{X}} \rightarrow \mathcal{D} \end{array} \right\}
 \end{array}$$