

1 Tor and ext

1.1 The Tor functor

Let R be a ring, and N a right R -module. There is an adjunction

$$- \otimes_R N : \mathbf{Mod}\text{-}R \leftrightarrow \mathbf{Ab} : \text{Hom}(N, -).$$

Thus, the functor $- \otimes_R N$ preserves colimits, hence by Proposition ??, is right exact. Thus, we may form the left derived functor.

Definition 1 (Tor functor). The left derived functor of $- \otimes_R N$, is called the Tor functor and denoted

$$\text{Tor}_i^R(-, N).$$

Example 2. Let R be a ring, and let $r \in R$. Assume that left multiplication by r is injective, and consider the right R -module R/rR . We have the short exact sequence

$$0 \longrightarrow R \xrightarrow{r \cdot} R \xrightarrow{\pi} R/rR \longrightarrow 0$$

exhibiting

$$0 \longrightarrow R \xrightarrow{r \cdot} R \longrightarrow 0$$

as a free resolution of R/rR . Thus we can calculate

$$\text{Tor}_i^R(R/rR, N) \simeq H_i \left(0 \longrightarrow R \otimes_R N \xrightarrow{r \cdot} R \otimes_R N \longrightarrow 0 \right)$$

That is, we have

$$\text{Tor}_i^R(R/rR, N) \simeq \begin{cases} N/rN, & n = 0 \\ rN, & n = 1 \\ 0, & \text{otherwise.} \end{cases}$$

There is a worrying asymmetry to our definition of Tor. The tensor product is, morally speaking, symmetric; that is, it should not matter whether we derive the first factor or the second. Pleasingly, Tor respects this.

Proposition 3. The functor Tor is balanced; that is,

$$L_i \text{Hom}(M, -)(N) \simeq L_i \text{Hom}(-, N)(M).$$

1.2 Ext and extensions

We know that $\text{Hom}(N, -): R\text{-Mod} \rightarrow \mathbf{Ab}$ is left exact, so we may form right derived functors.

Definition 4 (Ext functor). Let R be a ring, and $N \in R\text{-Mod}$ an R -module. The right derived functors

$$R^i \text{Hom}(N, -): R\text{-Mod} \rightarrow \mathbf{Ab}$$

are called the Ext functors, and denoted

$$R^i \text{Hom}(N, -) = \text{Ext}_R^i(N, -).$$

An argument very similar to that given for Tor in [Proposition 3](#) shows that Ext is balanced; that is, when computing $\text{Ext}_R^i(M, N)$, it doesn't matter whether we take an injective resolution of N or a projective resolution of M ; we will get the same result either way. Rather than spelling this out in detail, we will wait for Chapter ??, where a spectral sequences argument gives the result almost immediately.

There turns out to be a deep connection between the Ext groups and extensions.

Definition 5 (extension). Let A and B be R -modules. An extension of A by B is a short exact sequence

$$\xi: \quad 0 \longrightarrow B \hookrightarrow X \twoheadrightarrow A \longrightarrow 0.$$

Two extensions ξ and ξ' are said to be equivalent if there is a morphism of short exact sequences

$$\begin{array}{ccccccc} \xi: & 0 & \longrightarrow & B & \hookrightarrow & X & \twoheadrightarrow A \longrightarrow 0 \\ & & & \downarrow \text{id}_B & & \downarrow \varphi & \downarrow \text{id}_A \\ \xi': & 0 & \longrightarrow & B & \hookrightarrow & X' & \twoheadrightarrow A \longrightarrow 0 \end{array}.$$

The five lemma (Theorem ??) implies that φ is an isomorphism.

Let ξ be an extension of A by B . Applying the Ext functor we find an associated long exact sequence on cohomology.

$$\begin{array}{ccccccc} \text{Ext}^1(A, B) & \longrightarrow & \text{Ext}^1(X, B) & \longrightarrow & \cdots \\ \uparrow & & & & \uparrow \\ \text{Hom}(A, B) & \longrightarrow & \text{Hom}(X, B) & \longrightarrow & \text{Hom}(B, B) \end{array}$$

δ

Definition 6 (obstruction class). The image of id_B under the connecting homomorphism $\delta(\text{id}_B) \in \text{Ext}^1(A, B)$ is known as the obstruction class of ξ . For any extension ξ , we denote the obstruction class of ξ by $\Theta(\xi)$.

Proposition 7. An extension ξ is split if and only if the corresponding obstruction class is zero.

Proof. We are given an extension

$$\xi: \quad 0 \longrightarrow B \xrightarrow{f} X \xrightarrow{g} A \longrightarrow 0 .$$

First suppose that $\Theta(\xi) = 0$. Consider the following portion of the exact sequence on cohomology.

$$\mathrm{Hom}(X, B) \longrightarrow \mathrm{Hom}(B, B) \longrightarrow \mathrm{Ext}^1(A, B)$$

If $\delta: \mathrm{id} \rightarrow 0$, then there exists $a \in \mathrm{Hom}(X, B)$ such that $a \circ f = \mathrm{id}$. Thus, by the splitting lemma (Lemma ??), the sequence splits.

Now suppose that the sequence ξ splits; that is, that we have a morphism $a: A \rightarrow X$ such that $a \circ f = \mathrm{id}_B$.

$$0 \longrightarrow B \xrightarrow{f} X \xrightarrow{g} A \longrightarrow 0 .$$

\xleftarrow{a}

Since $\mathrm{Ext}^1(-, B)$ is additive, the sequence

$$0 \longrightarrow \mathrm{Ext}^1(A, B) \xrightarrow{\mathrm{Ext}^1(g, B)} \mathrm{Ext}^1(X, B) \xrightarrow{\mathrm{Ext}^1(f, B)} \mathrm{Ext}^1(B, B) \longrightarrow 0$$

$\xleftarrow{\mathrm{Ext}^1(a, B)}$

is split exact. Thus, $\mathrm{Ext}^1(g, B)$ is a monomorphism, so it has trivial kernel.

$$\begin{array}{ccc} \mathrm{Ext}^1(A, B) & \xrightarrow{\mathrm{Ext}^1(g, B)} & \dots \\ \uparrow & & \downarrow \delta \\ \mathrm{Hom}(A, B) & \longrightarrow & \mathrm{Hom}(X, B) \longrightarrow \mathrm{Hom}(B, B) \end{array}$$

Exactness then forces $\delta = 0$. □

Proposition 7 explains the name *obstruction class*: $\Theta(\xi)$ is the obstruction to the splitness of ξ . The first Ext group contains some information about when a sequence can split, or fail to split.

In fact, the group $\mathrm{Ext}^1(A, B)$ parametrizes extensions of A by B .

Theorem 8. There is a natural bijection

$$\{\text{Extension classes of } A \text{ by } B\} \xrightarrow{\cong} \mathrm{Ext}^1(A, B); \quad [\xi] \mapsto \Theta(\xi).$$

Proof. We first show well-definedness. Let ξ and ξ' be equivalent extensions. By naturality of δ , the following diagram commutes.

$$\begin{array}{ccccc} \mathrm{Hom}(X, B) & \longrightarrow & \mathrm{Hom}(B, B) & \xrightarrow{\delta} & \mathrm{Ext}^1(A, B) \\ \mathrm{id} \uparrow & & \mathrm{id} \uparrow & & \uparrow \mathrm{id} \\ \mathrm{Hom}(X, B) & \longrightarrow & \mathrm{Hom}(B, B) & \xrightarrow{\delta} & \mathrm{Ext}^1(A, B) \end{array}$$

Thus, $\Theta(\xi) = \Theta(\xi')$.

We now construct an explicit inverse Ψ to Φ . Since $R\text{-}\mathbf{Mod}$ has enough projectives, we can find a projective P which surjects onto A . Fix such a P and take the kernel of the surjection, giving the following short exact sequence.

$$0 \longrightarrow K \xrightarrow{\psi} P \xrightarrow{\gamma} A \longrightarrow 0 \quad (1.1)$$

We construct from such an exact sequence a map which takes an element of $\mathrm{Ext}^1(A, B)$ and gives an extension of A by B . First applying $\mathrm{Ext}^*(-, B)$ and taking the long exact sequence on cohomology we find the following exact fragment.

$$\mathrm{Hom}(K, B) \xrightarrow{\delta} \mathrm{Ext}^1(A, B) \longrightarrow \mathrm{Ext}^1(P, B)$$

Since P is projective, the group $\mathrm{Ext}^1(P, B)$ vanishes, so $\delta: \mathrm{Hom}(K, B) \rightarrow \mathrm{Ext}^1(A, B)$ is surjective. Thus, given any $e \in \mathrm{Ext}^1(A, B)$, we can find a (not unique!) lift to $\phi \in \mathrm{Hom}(K, B)$. Take the pushout in the following diagram.

$$\begin{array}{ccccccc} 0 & \longrightarrow & K & \xrightarrow{\psi} & P & \xrightarrow{\gamma} & A \longrightarrow 0 \\ & & \phi \downarrow & & \downarrow i_P & & \\ & & B & \xrightarrow{i_B} & B \amalg_K^\phi P & & \end{array}$$

By the dual to Lemma ??, the cokernel of i_B is equal to the cokernel of ψ , so we get an induced short exact sequence as follows.

$$\begin{array}{ccccccc} 0 & \longrightarrow & K & \xrightarrow{\psi} & P & \xrightarrow{\gamma} & A \longrightarrow 0 \\ & & \phi \downarrow & & \downarrow i_P & & \parallel \\ 0 & \longrightarrow & B & \xrightarrow{i_B} & B \amalg_K^\phi P & \xrightarrow{\pi} & A \longrightarrow 0 \end{array}$$

It remains to show that Ψ as defined is independent of the choice of ϕ , and that it really is an inverse to Φ . To this end, pick a different $\tilde{\phi} \in \mathrm{Hom}(K, B)$ such that $\delta(\tilde{\phi}) = e$.

This gives us a new morphism of exact sequences as follows

$$\begin{array}{ccccccc}
 0 & \longrightarrow & K & \xrightarrow{\psi} & P & \xrightarrow{\gamma} & A \longrightarrow 0 \\
 & & \downarrow \tilde{\phi} & & \downarrow \tilde{i}_P & & \parallel \\
 0 & \longrightarrow & B & \xrightarrow{\tilde{i}_B} & B \amalg_K^\phi P & \xrightarrow{\tilde{\pi}} & A \longrightarrow 0
 \end{array} \tag{1.2}$$

Because of the way we chose $\tilde{\phi}$, there exists some $\rho \in \text{Hom}(P, B)$ so that $\rho \circ \psi = \tilde{\phi} - \phi$. Consider the following outer square.

$$\begin{array}{ccc}
 K & \xrightarrow{\psi} & P \\
 \downarrow \phi & & \downarrow i_P \\
 B & \xrightarrow{i_B} & B \amalg_N^\phi P
 \end{array}
 \begin{array}{c}
 \searrow \tilde{i}_P - \tilde{i}_B \circ \rho \\
 \downarrow \exists! \alpha \\
 B \amalg_N^\phi P
 \end{array}
 \tag{1.3}$$

Using the commutativity of the left-hand square in [Diagram 1.2](#), we find

$$\begin{aligned}
 (\tilde{i}_P - \tilde{i}_B \circ \rho) \circ \psi &= \tilde{i}_P \circ \psi - \tilde{i}_B \circ \rho \circ \psi \\
 &= \tilde{i}_P \circ \psi - \tilde{i}_B \circ \tilde{\phi} + \tilde{i}_B \circ \phi \\
 &= \tilde{i}_B \circ \phi,
 \end{aligned}$$

so this outer square commutes, giving us the unique morphism α .

The claim that our two extensions are equivalent unwraps to the claim that the diagram

$$\begin{array}{ccccccc}
 0 & \longrightarrow & B & \xrightarrow{i_B} & B \amalg_K^\phi P & \xrightarrow{\pi} & A \longrightarrow 0 \\
 & & \parallel & & \downarrow \alpha & & \parallel \\
 0 & \longrightarrow & B & \xrightarrow{\tilde{i}_B} & B \amalg_K^\phi P & \xrightarrow{\tilde{\pi}} & A \longrightarrow 0
 \end{array}$$

commutes. The left-hand square is the lower commuting triangle in [Diagram 1.3](#).

Reminding ourselves of the proof of Lemma ??, we recall that the morphism π is defined by a pushout: it is the unique map $P \amalg_K^\phi P \rightarrow A$ such that $\pi \circ i_B = 0$ and $\pi \circ i_P = \gamma$. If we can show that $\tilde{\pi} \circ \alpha$ also satisfies these equations, then we can content ourselves that the diagram commutes. Indeed, we have

$$\begin{aligned}
 (\tilde{\pi} \circ \alpha) \circ i_B &= \tilde{\pi} \circ \tilde{i}_B \\
 &= 0
 \end{aligned}$$

and

$$\begin{aligned} (\tilde{\pi} \circ \alpha) \circ i_P &= \tilde{\pi} \circ (\tilde{i}_P - \tilde{i}_B \circ \rho) \\ &= \gamma. \end{aligned}$$

Our last job, which was really the whole point of the endeavor, is to check that the maps Φ and Ψ are inverse to each other. To that end, fix some A and B , and pick $e \in \text{Ext}^1(A, B)$. By construction of Ψ , we get a morphism of exact sequences

$$\begin{array}{ccccccccc} 0 & \longrightarrow & K & \xhookrightarrow{\psi} & P & \xrightarrow{\gamma} \twoheadrightarrow & A & \longrightarrow & 0 \\ & & \downarrow \phi & & \downarrow i_P & & \parallel & & \\ 0 & \longrightarrow & B & \xhookrightarrow{i_B} & B \amalg_K^\phi P & \xrightarrow{\pi} \twoheadrightarrow & A & \longrightarrow & 0 \end{array} \quad (1.4)$$

where ϕ is an element of $\text{Hom}(K, A)$ which is mapped to by e under δ . The next step in computing $\Phi(\Psi(e))$ would now be to stick $\Psi(e)$ into the hom functor $\text{Hom}(-, B)$, and look at the image of id_B under the connecting homomorphism. However, using the fact that Ext is a homological δ -functor, we can stick the whole morphism of exact sequences [Equation 1.4](#) in and get a morphism of long exact sequences out. Picking out the square containing the relevant boundary map, we find the following.

$$\begin{array}{ccc} \text{Hom}(B, B) & \xrightarrow{\delta} & \text{Ext}^1(A, B) \\ \text{Hom}(\phi, B) \downarrow & & \parallel \\ \text{Hom}(K, B) & \xrightarrow{\delta'} & \text{Ext}^1(A, B) \end{array}$$

Chasing the identity around in both directions tells us that

$$\delta(\text{id}_B) = \delta'(\phi) = e.$$

Next, we start with an extension

$$\xi: \quad 0 \longrightarrow B \xhookrightarrow{f} X \xrightarrow{g} \twoheadrightarrow A \longrightarrow 0.$$

Using the connecting homomorphism, we get a representative for the obstruction class $e = \delta(\text{id}_B) \in \text{Ext}^1(A, B)$. Using our short exact sequence from [Equation 1.1](#), we can find the fo □

Example 9. Let's do a sort of capstone computation, which pulls in just about everything we've done so far.

Consider the ring $R = \mathbb{Z}[x]$. We will compute the obstruction class of the extension

$$\xi: \quad 0 \longrightarrow \mathbb{Z} \xrightarrow{\cdot x} \mathbb{Z}[x]/(x^2) \longrightarrow \mathbb{Z} \longrightarrow 0,$$

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where we have written $\mathbb{Z} = \mathbb{Z}[x]/(x)$.

At a very high level of abstraction, we compute the obstruction class of ξ by taking the long exact sequence associated to it under the right derived functor of $\text{Hom}_{\mathbb{Z}[x]}(-, \mathbb{Z})$, namely

$$\begin{array}{c} \text{Ext}_{\mathbb{Z}[x]}^1(\mathbb{Z}, \mathbb{Z}) \longrightarrow \cdots \\ \uparrow \qquad \qquad \qquad \delta \qquad \qquad \qquad \searrow \\ \text{Hom}_{\mathbb{Z}[x]}(\mathbb{Z}, \mathbb{Z}) \longrightarrow \text{Hom}_{\mathbb{Z}[x]}(\mathbb{Z}[x]/(x^2), \mathbb{Z}) \longrightarrow \text{Hom}_{\mathbb{Z}[x]}(\mathbb{Z}, \mathbb{Z}) \end{array}$$

and finding the image of the identity $\text{id}_{\mathbb{Z}}$ under δ .

In order to find this long exact sequence, we need to lift ξ to a short exact sequence of projective resolutions. If we can find a projective resolution of \mathbb{Z} , the horseshoe lemma will do the rest of the work for us.

Our work on syzygies will help us here. We note that \mathbb{Z} has a single generator, namely 1, so we begin our resolution

$$\mathbb{Z}[x] \xrightarrow{\pi} \mathbb{Z}.$$

The kernel of this is the ideal (x) .

$$\begin{array}{c} (x) \\ \searrow \\ \mathbb{Z}[x] \xrightarrow{\pi} \mathbb{Z} \end{array}.$$

This again has a single generator, namely x .

$$\begin{array}{ccccc} & & (x) & & \\ & \nearrow \cdot x & \searrow & & \\ \mathbb{Z}[x] & \xrightarrow{\cdot x} & \mathbb{Z}[x] & \xrightarrow{\pi} & \mathbb{Z} \end{array}.$$

The composition gives us our next differential. Note that multiplication by x is injective, so our projective resolution terminates.

$$\cdots \longrightarrow 0 \longrightarrow \mathbb{Z}[x] \xrightarrow{\cdot x} \mathbb{Z}[x] \longrightarrow \cdots.$$

Now we apply the horseshoe lemma. We first build our horseshoe.

$$\begin{array}{ccccccc}
 & & \mathbb{Z}[x] & & & & \mathbb{Z}[x] \\
 & & \downarrow & & & & \downarrow \\
 & & \mathbb{Z}[x] & & & & \mathbb{Z}[x] \\
 & & \downarrow & & & & \downarrow \\
 0 & \longrightarrow & \mathbb{Z} & \xrightarrow{\cdot x} & \mathbb{Z}[x]/(x^2) & \xrightarrow{\pi} & \mathbb{Z} \longrightarrow 0
 \end{array}$$

The center column is then given by the direct sum of the outer columns, and the horizontal maps are the canonical injection and projection respectively.

$$\begin{array}{ccccccc}
 \mathbb{Z}[x] & \hookrightarrow & \mathbb{Z}[x]^2 & \twoheadrightarrow & \mathbb{Z}[x] \\
 \downarrow & (3) & \downarrow \alpha & (4) & \downarrow \\
 \mathbb{Z}[x] & \hookrightarrow & \mathbb{Z}[x]^2 & \twoheadrightarrow & \mathbb{Z}[x] \\
 \downarrow & (1) & \downarrow \beta & (2) & \downarrow \\
 0 & \longrightarrow & \mathbb{Z} & \xrightarrow{\cdot x} & \mathbb{Z}[x]/(x^2) & \xrightarrow{\pi} & \mathbb{Z} \longrightarrow 0
 \end{array}$$

The horseshoe lemma guarantees the existence of maps α and β above such that all the squares commute and the middle column forms a projective resolution, but being assured of existence is of no help to us; we must find them explicitly. We are helped tremendously that what we have found are free resolutions, so we can write the maps as matrices. We first figure out what we can about the morphisms α and β from the commutativity of the squares labelled (1)-(4).

(1) Taking the low road maps

$$1 \mapsto 1 \mapsto x,$$

while taking the high road maps

$$1 \mapsto (1, 0) \mapsto \beta(1, 0).$$

Thus,

$$\beta: (1, 0) \mapsto x$$

(2) The high road maps

$$(0, 1) \mapsto 1 \mapsto 1,$$

and the low road maps

$$(0, 1) \mapsto \beta(0, 1) \mapsto \beta(0, 1).$$

Thus,

$$\beta: (0, 1) \mapsto 1.$$

Already, this is enough to tell us that

$$\beta(a, b) = ax + b.$$

1.2.1 The Baer sum

By definition, $\text{Ext}^1(A, B)$ has an abelian group structure. We have provided a bijection between equivalence classes of extensions of B by A and $\text{Ext}^1(A, B)$, which allows us to transport our addition law from $\text{Ext}^1(A, B)$ to the set of extension classes. However, it is not at all clear how to interpret the sum of two classes of extensions.

The group operation on the set of extension classes is known as the *Baer sum*. Let

$$\xi: \quad 0 \longrightarrow B \xrightarrow{f} X \xrightarrow{g} \twoheadrightarrow A \longrightarrow 0$$

and

$$\xi': \quad 0 \longrightarrow B \xrightarrow{f'} X' \xrightarrow{g'} \twoheadrightarrow A \longrightarrow 0$$

be extensions. We shall form a new extension from ξ and ξ' called the *Baer sum*, which will turn out to be in the class of $\Psi(\Phi(\xi + \xi'))$.

A word of apology: things are about to get very element-theoretic. Of course, this is not a problem thanks to the Freyd-Mitchell embedding theorem, but it's still a bit sad.

From the above extensions above, take the following pushout, and form $X'' = \ker \alpha$.

$$\begin{array}{ccccc}
 B & \xrightarrow{f'} & X' & & \\
 \downarrow f & \searrow \exists! f'' & \downarrow & \searrow -g' & \\
 & & X'' & & A \\
 & & \downarrow & & \\
 X & \xrightarrow{\quad} & X \amalg_B X' & \xrightarrow{\exists! \alpha} & A \\
 & \searrow g & & & \\
 & & & &
 \end{array}$$

We have immediately the following morphism.

$$0 \quad B \xrightarrow{f''} X'' \quad A \quad 0$$

We would like to fill in morphisms so that we have an exact sequence. The obvious

choice of morphism $X'' \rightarrow A$, namely the composition $\alpha \circ \ker \alpha$, won't work because it's zero, so we have to define something else.

Warning: brutal element-wise computation ahead.

We have an explicit description of X'' as

$$X \amalg_B X' = X \oplus X' / \langle f(b) = f'(b) \rangle.$$

The kernel X'' is therefore the part of this which goes to zero under α , namely $(x, x') \in X \amalg_B X'$ such that

$$\alpha(x, x') = g(x) - g'(x') = 0,$$

i.e.

$$X'' = \{(x, x') \in X \amalg_B X' \mid g(x) = g'(x')\}.$$

The map f'' then simply sends $b \mapsto (f(b), 0)$, which by definition is equivalent to $(0, f'(b))$.

Note that f' is clearly injective.

Next, we define a map $g'': X'' \rightarrow A$ by sending

$$(x, x') \mapsto g(x) = g'(x').$$

This is clearly surjective since g and g' are. The kernel of g'' consists precisely of those (x, x') such that $g(x) = 0$, which is to say, such that $x = f(b)$ by exactness of ξ . Thus, we have constructed a bona fide short exact sequence

$$\xi'': \quad 0 \longrightarrow B \xrightarrow{f''} X'' \xrightarrow{g''} A \longrightarrow 0. \quad (1.5)$$

Definition 10 (Baer sum). The extension ξ'' from [Equation 1.5](#) is called the Baer sum of ξ and ξ' .

Lemma 11. The Baer sum is compatible with equivalences; that is, if $\xi \sim \zeta$, then $\xi + \xi' \sim \zeta + \xi'$.

Proof. Note that equivalence implies in particular \square

This means that the Baer sum descends to an addition law on equivalence classes.

Lemma 12. Let

$$\eta: \quad 0 \longrightarrow B \xrightarrow{i} B \oplus A \xrightarrow{p} A \longrightarrow 0$$

be a split extension. Then for any extension ξ , $\eta + \xi \sim \xi$.

Proof. \square

Lemma 13. Let

$$\xi: \quad 0 \longrightarrow B \xhookrightarrow{f} X \xrightarrow{g} \twoheadrightarrow A \longrightarrow 0 .$$

be an extension. Then

$$-\xi: \quad 0 \longrightarrow B \xhookrightarrow{-f} X \xrightarrow{g} \twoheadrightarrow A \longrightarrow 0$$

is the inverse under the Baer sum, i.e. $\xi + (-\xi) \sim \eta$.