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CATEGORY-THEORETIC PRELIMINARIES

The theory of infinity categories rests on the shoulders of the theory of ordinary category theory. In this appendix, we recall some category-theoretic results that we will use. In general, we provide no proofs for these results; they are standard, and proofs can be found in any introductory text on category theory.

1 *Limits and colimits*

1.1 *Formulae for limits*

One can calculate limits using products and equalizers, and colimits using coproducts and coequalizers.

Fact 1.1. Let \mathcal{C} be a small category, and let \mathcal{D} be locally small. Let $F: \mathcal{C} \rightarrow \mathcal{D}$ be a functor.

- If \mathcal{D} has small products and equalizers, then $\lim F$ is given by the following equalizer,

$$\lim F \hookrightarrow \prod_{c \in \mathcal{C}} F(c) \begin{array}{c} \xrightarrow{R} \\ \xleftarrow{S} \end{array} \prod_{\alpha \in \text{Mor}(\mathcal{C})} F(\text{codom}(\alpha))$$

where the morphisms R and S are defined by their components

$$R_\alpha = \pi_{F(\text{codom}(\alpha))} : \prod_{c \in \mathcal{C}} F(c) \rightarrow F(\text{codom}(\alpha))$$

and

$$S_\alpha = F(\alpha) \circ \pi_{\text{dom}(\alpha)} : \prod_{c \in \mathcal{C}} F(c) \rightarrow F(\text{codom}(\alpha))$$

- If \mathcal{D} has small coproducts and coequalizers, then $\text{colim } F$ is given by the following coequalizer,

$$\prod_{\alpha \in \text{Mor}(\mathcal{C})} F(\text{dom}(\alpha)) \begin{array}{c} \xrightarrow{A} \\ \xleftarrow{B} \end{array} \prod_{c \in \mathcal{C}} F(c) \twoheadrightarrow \text{colim } F$$

where the morphisms A and B are defined by their components

$$A_\alpha = \iota_{F(\text{dom}(\alpha))} : F(\text{dom}(\alpha)) \rightarrow \coprod_{c \in \mathcal{C}} F(c)$$

and

$$B_\alpha = \iota_{\text{codom}(\alpha)} \circ F(\alpha) : F(\text{dom}(\alpha)) \rightarrow \coprod_{c \in \mathcal{C}} F(c)$$

1.2 Limits in functor categories

In this section, we review facts about limits and colimits in functor categories. Much of this text is devoted to one particular presheaf category \mathbf{Set}_Δ , the category of functors $\Delta^{\text{op}} \rightarrow \mathbf{Set}$,¹ and it will pay to understand categories of this form.

¹ Where Δ is the *simplex category*, Definition ??.

Fact 1.2. Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be a functor.

- If \mathcal{C} has an initial object i , then $\lim F = F(i)$.
- If \mathcal{C} has a terminal object i , then $\text{colim } F = F(i)$.

Theorem 1.3. Let \mathcal{C} and \mathcal{D} be categories, with \mathcal{D} small and \mathcal{C} locally small.

- If \mathcal{D} is complete, then so is $[\mathcal{C}, \mathcal{D}]$.
- If \mathcal{D} is cocomplete, then so is $[\mathcal{C}, \mathcal{D}]$.

Corollary 1.4. Limits and colimits in functor categories are computed pointwise. That is, for any functor $F : J \rightarrow [\mathcal{C}, \mathcal{D}]$ and any $j \in J$,

$$(\lim F)(j) = \lim(Fj),$$

where the first limit is taken in $[\mathcal{C}, \mathcal{D}]$, and the second is taken in \mathcal{D} .

Corollary 1.5. A morphism $\eta : F \rightarrow G$ in a functor category is a monomorphism (resp. epimorphism) if and only if each component η_x is a monomorphism (resp. epimorphism).

Proof. A morphism $\eta : F \rightarrow G$ is a monomorphism if and only if the square

$$\begin{array}{ccc} F & \xrightarrow{\text{id}} & F \\ \text{id} \downarrow & & \downarrow \eta \\ F & \xrightarrow{\eta} & G \end{array}$$

is a pullback square. But by [Corollary 1.4](#), the above square is a pullback square if and only if its components are pullback squares, which in turn is true if and only if each component is a monomorphism.

The epimorphism case is dual. □

2 Kan extensions

2.1 Pointwise formulae for Kan extensions

Let $F: \mathcal{C} \rightarrow \mathcal{D}$ be a functor, and let $d \in \mathcal{D}$ be an object.

- We will denote by $\mathcal{C}_{/d}$ the category whose objects are pairs (c, f) , where $c \in \mathcal{C}$ and $f: F(c) \rightarrow d$, and whose morphisms $(c, f) \rightarrow (c', f')$ are morphisms $h: c \rightarrow c'$ such that the triangle

$$\begin{array}{ccc} F(c) & \xrightarrow{F(h)} & F(c') \\ & \searrow f & \swarrow f' \\ & d & \end{array}$$

commutes. Note that there is an obvious forgetful functor $\mathcal{C}_{/d} \rightarrow \mathcal{C}$.

- We will denote by $\mathcal{C}_{d/}$ the category whose objects are pairs (c, f) , where $c \in \mathcal{C}$ and $f: d \rightarrow F(c)$, and whose morphisms $(c, f) \rightarrow (c', f')$ are morphisms $h: c \rightarrow c'$ such that the triangle

$$\begin{array}{ccc} & d & \\ f \swarrow & & \searrow f' \\ F(c) & \xrightarrow{F(h)} & F(c') \end{array}$$

commutes. Note that there is an obvious forgetful functor $\mathcal{C}_{d/} \rightarrow \mathcal{C}$.

Fact 1.6. Consider categories and functors as follows, where \mathcal{C} is small.

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{D} & \mathcal{D} \\ & \searrow F & \\ & \mathcal{C}' & \end{array}$$

- If \mathcal{D} has all small colimits, then the left Kan extension $F_!D$ exists and has values

$$(F_!D)(c') = \operatorname{colim} \left[\mathcal{C}_{c'/} \longrightarrow \mathcal{C} \xrightarrow{D} \mathcal{D} \right].$$

- If \mathcal{D} has all small limits, then the right Kan extension F_*D exists and has values

$$(F_*D)(c') = \lim \left[\mathcal{C}_{/c'} \longrightarrow \mathcal{C} \xrightarrow{D} \mathcal{D} \right].$$

2.2 Properties of Kan extensions

Theorem 1.7. Let \mathcal{C} be a small category, and \mathcal{D} locally small. Let $F: \mathcal{C} \rightarrow \mathcal{C}'$ be fully faithful. If \mathcal{D} has small colimits, then the unit

$$\epsilon: \operatorname{id}_{\operatorname{Fun}(\mathcal{C}, \mathcal{D})} \Rightarrow F^* \circ F_!$$

is a natural isomorphism.

The categories $\mathcal{C}_{/d}$ and $\mathcal{C}_{d/}$ are often referred to as (co)slice categories, comma categories, and under- and overcategories. There seems to be little agreement in the literature as to which of these terms refers to precisely which concept. We will call $\mathcal{C}_{d/}$ an *overcategory* and $\mathcal{C}_{/d}$ an *undercategory*.

The notations used for these categories are as diverse as the names. What we are calling $\mathcal{C}_{/d}$ is often called $(F \downarrow d)$, for example.

Corollary 1.8. Let \mathcal{C} be a small category, and \mathcal{D} locally small. Let $F: \mathcal{C} \rightarrow \mathcal{C}'$ be fully faithful.

- If \mathcal{D} has small limits, then F_* is fully faithful.
- If \mathcal{D} has small colimits, then $F_!$ is fully faithful.

2.3 Yoneda extensions

For any locally small category \mathcal{C} , the *Yoneda embedding* is the functor

$$\mathcal{Y}: \mathcal{C} \rightarrow \mathbf{Fun}(\mathcal{C}^{\mathrm{op}}, \mathbf{Set}) := \mathbf{Set}_{\mathcal{C}}; \quad c \mapsto \mathrm{Hom}_{\mathcal{C}}(-, c).$$

The category $\mathbf{Set}_{\mathcal{C}}$ is the category of (**Set**-valued) presheaves on \mathcal{C} . We will frequently be interested in Kan extensions along the Yoneda embedding; these are called *Yoneda extensions*. That is, for any functor $F: \mathcal{C} \rightarrow \mathcal{D}$, we will often want to compute the Kan extension $\mathcal{Y}_! F$.

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{D} \\ & \searrow \mathcal{Y} & \nearrow \mathcal{Y}_! F \\ & \mathbf{Set}_{\mathcal{C}} & \end{array}$$

In this section, we prove some basic results about Yoneda extensions.

Lemma 1.9. Let \mathcal{C} be a locally small category. For any functor $F \in \mathbf{Set}_{\mathcal{C}}$ we have an isomorphism

$$F \cong \mathrm{colim} \left[\mathcal{C}_{/F} \rightarrow \mathcal{C} \xrightarrow{\mathcal{Y}} \mathbf{Set}_{\mathcal{C}} \right].$$

Proof. By [Theorem 1.7](#), the fully faithfulness of the Yoneda embedding, and the definition of the restriction functor, we have that

$$F \cong (\mathcal{Y}^* \circ \mathcal{Y}_!)(F) \cong (\mathcal{Y}_! \mathcal{Y})(F).$$

The functor $\mathcal{Y}_! \mathcal{Y}$ comes from the following Kan extension.

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{\mathcal{Y}} & \mathbf{Set}_{\mathcal{C}} \\ & \searrow \mathcal{Y} & \nearrow \mathcal{Y}_! \mathcal{Y} \\ & \mathbf{Set}_{\mathcal{C}} & \end{array}$$

The colimit formula for left Kan extensions ([Fact 1.6](#)) then tells us that

$$F \cong \mathrm{colim} \left[\mathcal{C}_{/F} \rightarrow \mathcal{C} \xrightarrow{\mathcal{Y}} \mathbf{Set}_{\mathcal{C}} \right]$$

as required. □

Lemma 1.10. Let \mathcal{J} be a small category, and let \mathcal{C} be locally small and cocomplete. Let F be a functor as below. Then the Yoneda extension $\mathcal{Y}_! F$ is left adjoint to the functor

$$G: \mathcal{C} \rightarrow \mathbf{Set}_{\mathcal{J}}; \quad c \mapsto \mathcal{C}(F(-), c).$$

$$\begin{array}{ccc}
 \mathcal{J} & \xrightarrow{F} & \mathcal{C} \\
 \searrow \mathcal{Y} & & \nearrow \mathcal{Y}_! F \\
 & \mathbf{Set}_{\mathcal{J}} &
 \end{array}
 \quad \mathcal{Y}_! F : \mathbf{Set}_{\mathcal{J}} \leftrightarrow \mathcal{C} : G$$

Proof. Since \mathcal{C} is cocomplete, we are entitled to use the colimit formula for left Kan extensions. That is, we can write

$$(\mathcal{Y}_! F)(H) = \operatorname{colim} \left[(\mathcal{Y}/H) \xrightarrow{\mathcal{U}} \mathcal{J} \xrightarrow{F} \mathcal{C} \right] = \operatorname{colim} [F \circ \mathcal{U}] .$$

Thus, we have the following string of natural isomorphisms.

$$\begin{aligned}
 \mathcal{C}((\mathcal{Y}_! F)(H), c) &\cong \mathcal{C}(\operatorname{colim}[F \circ \mathcal{U}], c) \\
 &\cong \lim \mathcal{C}(F(-), c) \circ \mathcal{U} \\
 &\cong \lim G(c) \circ \mathcal{U} \\
 &\cong \lim \mathbf{Set}_{\mathcal{J}}(\mathcal{Y}(-), G(c)) \circ \mathcal{U} . \\
 &\cong \mathbf{Set}_{\mathcal{J}}(\operatorname{colim} \mathcal{Y} \circ \mathcal{U}, G(c)) \\
 &\cong \mathbf{Set}_{\mathcal{J}}(H, G(c)) .
 \end{aligned}$$

□

Corollary 1.11. Let \mathcal{J} be a small category, and let \mathcal{C} be locally small and cocomplete. Let $F : \mathcal{J} \rightarrow \mathcal{C}$. Then the Yoneda extension $\mathcal{Y}_! F$ preserves colimits.

Proof. By [Lemma 1.10](#), $\mathcal{Y}_! F$ is a left adjoint.

□

Theorem 1.12. Let \mathcal{C} be a small category, and let \mathcal{D} be cocomplete. Let $F : \mathbf{Set}_{\mathcal{C}} \rightarrow \mathcal{D}$ be any functor. Then F preserves colimits if and only if it is the left Yoneda extension of its restriction $\mathcal{Y}^* F : \mathcal{C} \rightarrow \mathcal{D}$, i.e. if $F = (\mathcal{Y}_! \mathcal{Y}^*)(F)$.

Proof. Assume that $F = (\mathcal{Y}_! \mathcal{Y}^*)(F)$. By [Lemma 1.10](#), $(\mathcal{Y}_! \mathcal{Y}^*)(F)$ is a left adjoint, hence preserves colimits.

Now assume that F preserves colimits. Pick some functor $G : \mathcal{C}^{\operatorname{op}} \rightarrow \mathbf{Set}$. By [Lemma 1.9](#), we can write

$$\begin{aligned}
 F(G) &\cong F \left(\lim_{\rightarrow} \left[(\mathcal{Y}/G) \rightarrow \mathcal{C} \xrightarrow{\mathcal{Y}} \mathbf{Set}_{\mathcal{C}} \right] \right) \\
 &\cong \lim_{\rightarrow} \left[(\mathcal{Y}/G) \rightarrow \mathcal{C} \xrightarrow{\mathcal{Y}} \mathbf{Set}_{\mathcal{C}} \xrightarrow{F} \mathcal{D} \right] \\
 &\cong \lim_{\rightarrow} \left[(\mathcal{Y}/G) \rightarrow \mathcal{C} \xrightarrow{\mathcal{Y}^* F} \mathcal{D} \right] \\
 &\cong \mathcal{Y}_! (\mathcal{Y}^* F)(G) .
 \end{aligned}$$

□

Corollary 1.13. Let \mathcal{C} be a small category, and \mathcal{D} locally small. There is an equivalence of categories

$$[\mathbf{Set}_{\mathcal{C}}, \mathcal{D}]_{\text{cocontinuous}} \cong [\mathcal{C}, \mathcal{D}].$$

That is, colimit preserving functors $\mathbf{Set}_{\mathcal{C}} \rightarrow \mathcal{D}$ are equivalent to functors $\mathcal{C} \rightarrow \mathcal{D}$ via restriction along the Yoneda embedding.

Proof. We have an adjunction

$$\mathcal{Y}_! : [\mathcal{C}, \mathcal{D}] \leftrightarrow [\mathbf{Set}_{\mathcal{C}}, \mathcal{D}] : \mathcal{Y}^*.$$

By [Corollary 1.11](#), the image of $\mathcal{Y}_!$ is completely contained in $[\mathbf{Set}_{\mathcal{C}}, \mathcal{D}]_{\text{cocontinuous}}$, so the above adjunction extends to an adjunction

$$\mathcal{Y}_! : [\mathcal{C}, \mathcal{D}] \leftrightarrow [\mathbf{Set}_{\mathcal{C}}, \mathcal{D}]_{\text{cocontinuous}} : \mathcal{Y}^*.$$

We know that the Yoneda embedding is fully faithful. Thus, by [Theorem 1.7](#), we have that the unit is an isomorphism. Also, by [Theorem 1.12](#), the counit is an isomorphism, so the above adjunction is in fact an equivalence of categories. \square

2.4 Day convolution

Let (\mathcal{C}, \otimes) be a monoidal category. The category $\mathbf{Set}_{\mathcal{C}}$ of functors from \mathcal{C}^{op} to \mathbf{Set} inherits a monoidal structure as follows.

Given functors $F, G : \mathcal{C} \rightarrow \mathbf{Set}$, we can define a functor

$$F \bar{\times} G : \mathcal{C}^{\text{op}} \times \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}; \quad (c, d) \mapsto F(c) \times G(d).$$

Definition 1.14 (Day convolution). The Day convolution of F and G , denoted $F \otimes G$, is given by the left Kan extension of $F \bar{\times} G$ along \otimes^{op} .

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} \times \mathcal{C}^{\text{op}} & \xrightarrow{F \bar{\times} G} & \mathbf{Set} \\ \searrow \otimes^{\text{op}} & & \nearrow F \otimes G \\ & \mathcal{C}^{\text{op}} & \end{array}$$

The colimit formula for left Kan extensions tells us that

$$(F \otimes G)(c) = \text{colim} \left[(c \downarrow \otimes^{\text{op}}) \rightarrow \mathcal{C}^{\text{op}} \times \mathcal{C}^{\text{op}} \xrightarrow{F \bar{\times} G} \mathbf{Set} \right]$$

3 Enriched category theory

Definition 1.15 (enriched category). Let $(\mathcal{M}, \otimes, 1_{\mathcal{M}}, \alpha, \lambda, \rho)$ be a monoidal category. A \mathcal{M} -enriched category \mathcal{C} consists of the following data.

- A set $\text{ob}(\mathcal{C})$ of objects.

- For each pair x, y of objects, an object $\mathcal{C}(x, y) \in \mathcal{M}$ of morphisms.
- For each object $x \in \mathcal{C}$, a morphism

$$1_{\mathcal{M}} \rightarrow \mathcal{C}(x, x)$$

called the *unit*

- For objects x, y , and z , a morphism

$$\mathcal{C}(x, y) \otimes \mathcal{C}(y, z) \rightarrow \mathcal{C}(x, z)$$

called the *composition*.

This composition must be unital and associative, i.e. the diagrams

$$\begin{array}{ccc} \mathcal{C}(x, y) & \xrightarrow{\quad} & \mathcal{C}(y, y) \otimes \mathcal{C}(x, y) \\ \downarrow & \searrow \text{id} & \downarrow \\ \mathcal{C}(x, y) \otimes \mathcal{C}(x, x) & \xrightarrow{\quad} & \mathcal{C}(x, y) \end{array}$$

and

$$\begin{array}{ccc} \mathcal{C}(x, y) \otimes \mathcal{C}(y, z) \otimes \mathcal{C}(z, w) & \xrightarrow{\quad} & \mathcal{C}(x, y) \otimes \mathcal{C}(y, w) \\ \downarrow & & \downarrow \\ \mathcal{C}(x, z) \otimes \mathcal{C}(z, w) & \xrightarrow{\quad} & \mathcal{C}(x, w) \end{array}$$

must commute (where unitors and associators are notationally suppressed).

Lemma 1.16. Let \mathcal{C} be an \mathcal{M} -enriched category, and let $F: \mathcal{M} \rightarrow \mathcal{N}$ be a monoidal functor with data

$$\Phi_{x,y}: F(x) \otimes_{\mathcal{N}} F(y) \rightarrow F(x \otimes_{\mathcal{M}} y), \quad \phi: 1_{\mathcal{N}} \rightarrow F(1_{\mathcal{M}}).$$

This data allows us to build from \mathcal{C} a category enriched over \mathcal{N} .

Proof.

- We apply F to each hom-object in \mathcal{M} to get a new hom-object in \mathcal{N} ;

$$\mathcal{C}(x, y)_{\mathcal{N}} = F(\mathcal{C}(x, y)_{\mathcal{M}}).$$

- The unit is the composition

$$1_{\mathcal{N}} \longrightarrow F(1_{\mathcal{M}}) \longrightarrow F(\mathcal{C}(x, y)_{\mathcal{M}}) = \mathcal{C}(x, y)_{\mathcal{N}}$$

- The composition is given by the obvious.

□