

Design and Analysis  
of Algorithms I

# Introduction

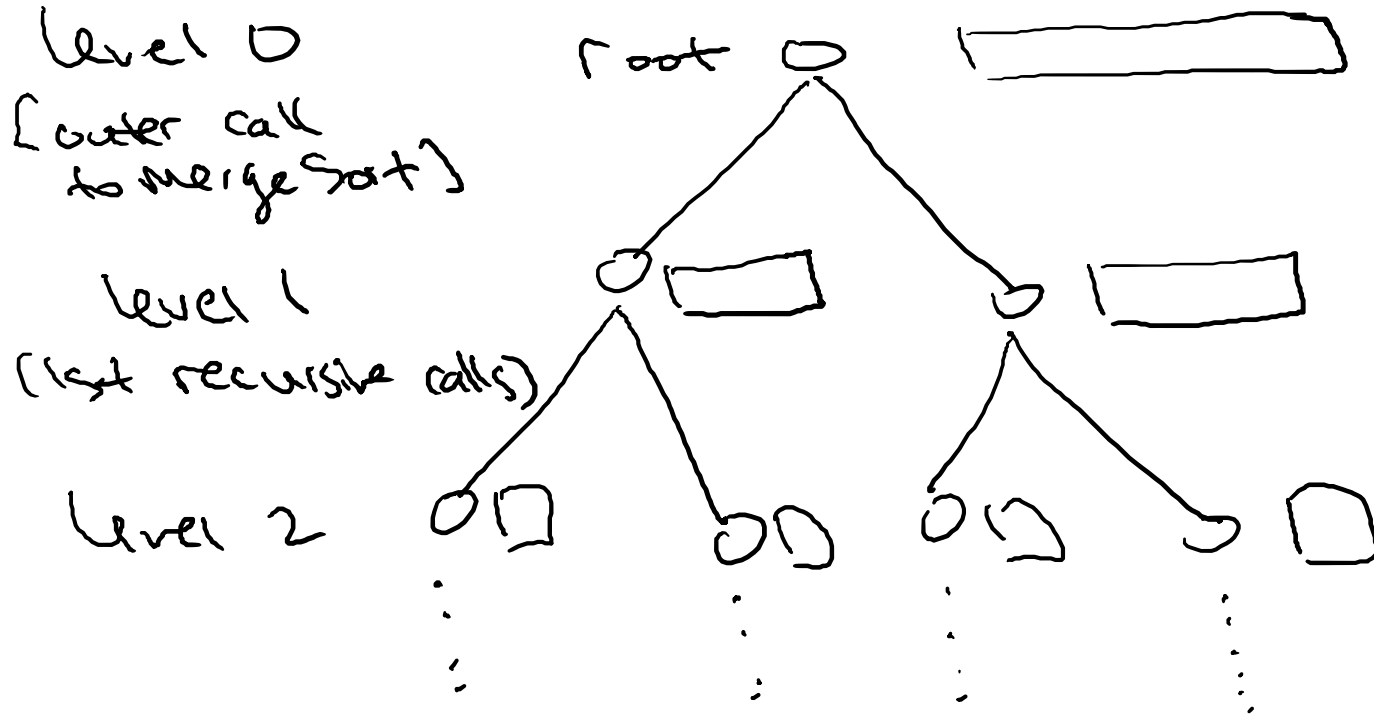
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## Merge Sort (Analysis)

# Running Time of Merge Sort

**Claim:** For every input array of  $n$  numbers, Merge Sort produces a sorted output array and uses at most  $6n \log_2 n + 6n$  operations.

# Proof of claim (assuming $n = \text{power of } 2$ ):



Roughly how many levels does this recursion tree have (as a function of  $n$ , the length of the input array)?

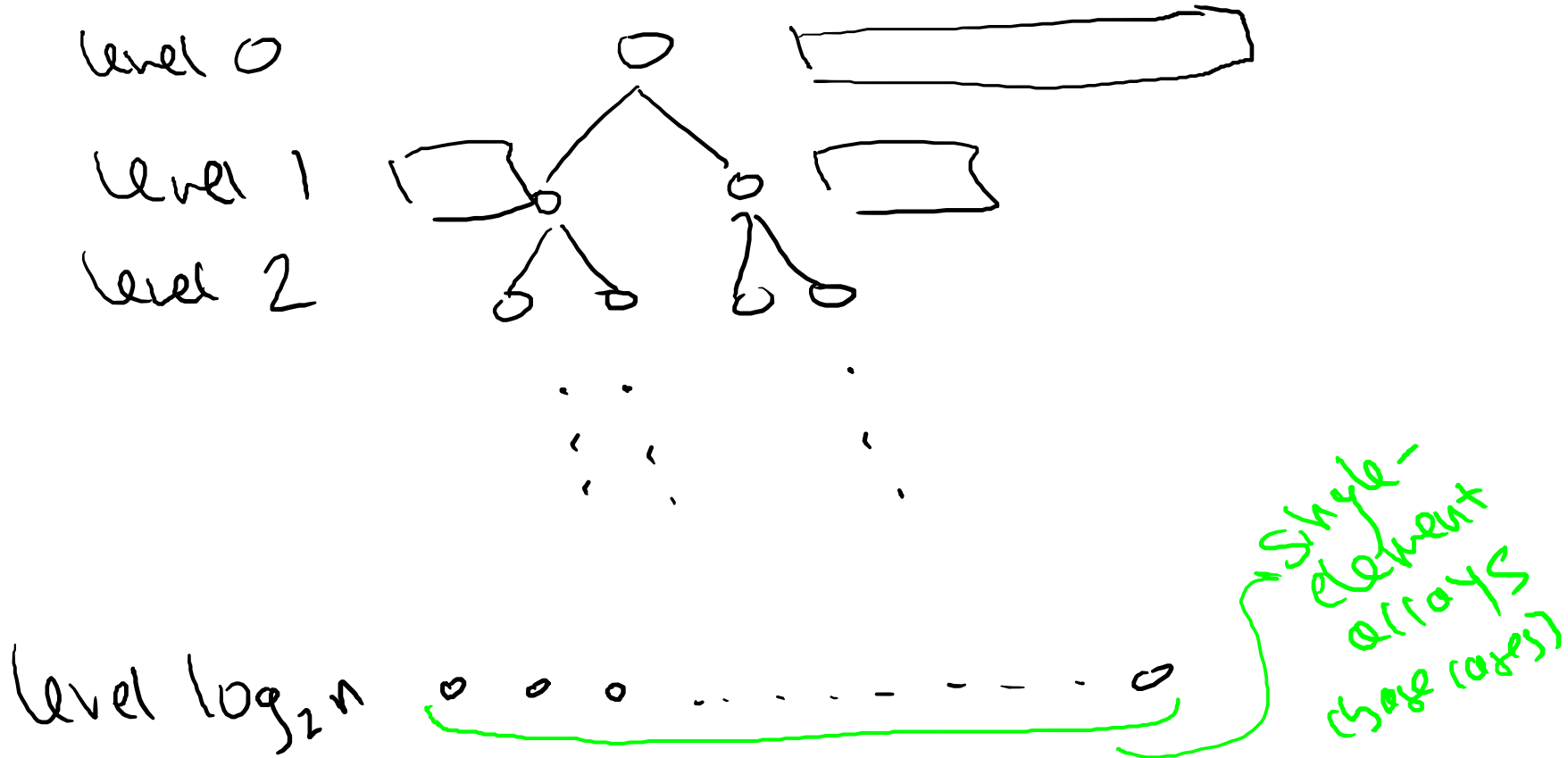
☐ A constant number (independent of  $n$ ).

→ ☒  $\log_2 n$  ( $\log_2 n + 1$ ) to be exact!


☐  $\sqrt{n}$

☐  $n$

# Proof of claim (assuming $n = \text{power of } 2$ ):



What is the pattern? Fill in the blanks in the following statement:  
at each level  $j=0,1,2,\dots, \log_2 n$ , there are *<blank>* subproblems,  
each of size *<blank>*.

- ☐  $2^j$  and  $2^j$ , respectively.
- ☐  $n/2^j$  and  $n/2^j$ , respectively.
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# Proof of claim (assuming $n = \text{power of } 2$ ):

At each level  $j=0,1,2,\dots, \log_2 n$ , there are  $2^j$  subproblems, each of size  $n/2^j$ .

Total # of operations at level  $j=0,1,2,\dots, \log_2 n$ :

$$\leq \cancel{2^j} \cdot 6 \left( \frac{n}{\cancel{2^j}} \right) = 6n$$

# of level- $j$   
subproblems

size of  
level- $j$   
subproblem

work  
per level- $j$  subproblem

Total

$$6n (\log_2 n + 1)$$

work per level

# of levels

# Running Time of Merge Sort

**Claim:** For every input array of  $n$  numbers, Merge Sort produces a sorted output array and uses at most  $6n \log_2 n + 6n$  operations.

QED!