



Design and Analysis
of Algorithms I

Introduction

Merge Sort

(Pseudocode)

Merge Sort: Pseudocode

- recursively sort 1st half of input array
 - " " 2nd " " " "
 - merge two sorted sublists into one
- [ignores base cases]

Pseudocode for Merge:

C = output array [length = n]

A = 1st sorted array [$n/2$]

B = 2nd " " [$n/2$]



$i \uparrow$



$j \uparrow$

$i = 1$
 $j = 1$

for $k = 1$ to n

if $A(i) < B(j)$

$C(k) = A(i)$

$i++$

else [$B(j) < A(i)$]

$C(k) = B(j)$

$j++$

end [ignore end cases]

Pseudocode for Merge:

C = output [length = n]

A = 1st sorted array [n/2]

B = 2nd sorted array [n/2]

i = 1

j = 1

for k = 1 to n

if A(i) < B(j)

C(k) = A(i)

i++

else [B(j) < A(i)]

C(k) = B(j)

j++

end

(ignores end cases)

Merge Sort Running Time?

Key Question: running time of MergeSort on
array of n numbers?

Running time \approx # of lines of
code executed

Pseudocode for Merge:

C = output [length = n]

A = 1st sorted array [n/2]

B = 2nd sorted array [n/2]

i = 1

j = 1

2 operations

for k = 1 to n ✓

if A(i) < B(j) ✓

C(k) = A(i) —

i++ —

else [B(j) < A(i)]

C(k) = B(j) —

j++ —

end

(ignores end cases)

Running Time of Merge

Upshot: running time of Merge on a array
of m numbers is $\leq 4m + 2$
 $\leq 6m$ (since $m \geq 1$)

Running Time of Merge Sort

Claim: Merge Sort requires

$$\leq 6n \log_2 n + 6n \text{ operations}$$

to sort n numbers.

Recall: $\log_2 n =$
of times you divide
by 2 until you get
down to 1.

