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1. Syntax and Semantics.

- (a) Yes
- (b) Yes
- (c) No
- (d) No,because an expression can have multiple expression rules applied to it at one time. For example:
 - i. $(3 + 7) + (1 + 2) \Rightarrow 10 + (1 + 2)$
 - ii. $(3 + 7) + (1 + 2) \Rightarrow (3 + 7) + 3$
- (e) $e := \dots | \text{ if0 } e_1 \text{ then } e_2 \text{ else } e_3$
- (f) $\text{if0 } 0 \text{ then } e_2 \text{ else } e_3 \rightarrow e_2$
- (g) $\text{if0 } i \text{ then } e_2 \text{ else } e_3 \rightarrow e_3 (i \neq 0)$
- (h) $K := \dots | \text{ if0 } K \text{ then } e_1 \text{ else } e_2$
- (i)
$$\frac{e_1 \Downarrow 0 \ e_2 \Downarrow i_2}{\text{if0 } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow i_2} \quad \frac{e_1 \Downarrow i_1 (i_1 \neq 0) \ e_3 \Downarrow i_3}{\text{if0 } e_1 \text{ then } e_2 \text{ else } e_3 \Downarrow i_3}$$
- (j) $\text{if0 } y + x + x \text{ then } 42 \text{ else } 0$

let $x := -3$ in let $y := 6$ in
 $\text{if } 0 \neq x + x \text{ then } 42 \text{ else } 0$

$\Rightarrow (\text{let } y := 6 \text{ in } \text{if } 0 \neq y + x + x \text{ then } 42$
 $\text{else } 0) [-3/x]$

$\Rightarrow (\text{let } y := 6 \text{ if } 0 \neq y + (-3) + (-3)$
then $42 \text{ else } 0$

$\Rightarrow (\text{if } 0 \neq y + (-3) + (-3)$
then $42 \text{ else } 0) [6/y]$

$\Rightarrow \text{if } 0 \neq 6 + (-3) + (-3)$
then $42 \text{ else } 0$

$\Rightarrow \text{if } 0 \neq 3 + (-3)$
then $42 \text{ else } 0$

$\Rightarrow \text{if } 0 \neq 0 \text{ then } 42 \text{ else } 0$

$\Rightarrow 42$

(k)

$$\begin{array}{c}
 \begin{array}{ccccccc}
 & 6 & \textcircled{\text{U}} & 6 & & (-3) & \textcircled{\text{U}} & (-3) \\
 \hline
 -3 & \textcircled{\text{U}} & -3 & & 6 & + & (-3) & \textcircled{\text{U}} & 3 \\
 \hline
 & 6 & + & (-3) & \textcircled{\text{U}} & (-3) & \textcircled{\text{U}} & 0
 \end{array} \\
 \hline
 \text{let } y := 6 \text{ in } y + \overbrace{(-3)}^{\text{---}} + \overbrace{(-3)}^{\text{---}} \textcircled{\text{U}} 0 \\
 \hline
 \text{let } x := -3 \text{ in let } y := 6 \text{ in } y + \cancel{x} + x \textcircled{\text{U}} 0
 \end{array}$$

(l)

2. Untyped λ -calculus

- (a) $\text{hd} := \lambda x . \text{fst} (\text{snd} x)$
- (b) $\text{tl} := \lambda x . \text{snd} (\text{snd} x)$
- (c) $\text{isnil} := \lambda x . \text{fst} x$
- (d) $\begin{aligned} \text{hd} (\text{cons} x l) &\equiv \text{fst} (\text{snd} (\text{pair} \text{ false} (\text{pair} x l))) \\ &\equiv \text{fst} (\text{pair} x l) \\ &\equiv x \end{aligned}$
- (e) $\begin{aligned} \text{tl} (\text{cons} x l) &\equiv \text{snd} (\text{snd} (\text{pair} \text{ false} (\text{pair} x l))) \\ &\equiv \text{snd} (\text{pair} x l) \\ &\equiv l \end{aligned}$
- (f) $\begin{aligned} \text{isnil nil} &\equiv \text{fst} (\lambda x . x) \\ &\equiv (\lambda x . x) (\text{true}) \\ &\equiv \text{true} \end{aligned}$
- (g) $\begin{aligned} \text{isnil} (\text{cons} x l) &\equiv \text{fst} (\text{pair} \text{ false} (\text{pair} x l)) \\ &\equiv \text{false} \end{aligned}$
- (h) $\begin{aligned} \text{hd nil} &\equiv \text{fst} (\text{snd} (\lambda x . x)) \\ &\equiv \text{fst} ((\lambda p . p \text{ false}) (\lambda x . x)) \\ &\equiv (\lambda p . p \text{ true}) (\text{false}) \\ &\equiv \text{false true} \end{aligned}$