

# Econ C103 Problem Set 5

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## Exercise 1

- (a) Each agent  $i \in \{1, 2\}$  has type indicated by the value of the object that they have at the beginning of the game,  $\theta_i \in \Theta = [0, 1]$ , and can send a message from  $M_i = \{0, 1\}$  indicating whether they want to exchange objects. If both agents choose  $m_i = 1$ , then the trade takes place, which we denote with  $x = 1$ . So the allocation rule is

$$x(m) = \begin{cases} 0 & \text{if } m_1 = 0 \text{ or } m_2 = 0 \\ 1 & \text{if } m_1 = m_2 = 1 \end{cases}$$

and each agent's utility is

$$u_i(x, \theta_i) = \begin{cases} \theta_i & \text{if } x = 0 \\ \theta_{-i} & \text{if } x = 1 \end{cases}.$$

- (b) Trading will never be Pareto efficient, so the only Pareto efficient outcome is the original allocation,  $x = 0$ . If  $\theta_1 \neq \theta_2$ , trading would make one agent worse off, since both agents assign the same value to the objects. If  $\theta_1 = \theta_2$ , then trading won't change either agent's utility and thus won't make either agent strictly better off.

Either allocation ( $x = 0$  or  $1$ ) maximizes the sum of the agents' utilities, which is always  $\theta_1 + \theta_2$ .

- (c) The best response of each agent depends on the message the other agent sends and the other agent's type. If  $m_{-i} = 0$ , then trade will never take place, so it doesn't matter what message the agent chooses; we will pick  $m_i = 0$ . Then the best response for agent  $i$  is

$$\sigma_i(m_{-i}, \theta) = \begin{cases} 0 & \text{if } m_{-i} = 0, \text{ or } m_{-i} = 1 \text{ and } \theta_i \geq \theta_{-i} \\ 1 & \text{if } m_{-i} = 1 \text{ and } \theta_i < \theta_{-i} \end{cases}.$$

Thus the agent's best response depends on  $m_{-i}$  and  $\theta_{-i}$ . This means that there is no dominant strategy equilibrium.

- (d) Given that the types  $\theta$  are uniformly distributed on  $[0, 1]$ , each agent initially expects the other agent's type to be  $\mathbb{E}[\theta_{-i}] = 0.5$ . This means it is rational to trade, i.e. send the message  $m_i = 1$ , only if  $\theta_i < 0.5$ . But then the expected payoff from the trade for the *other* agent is  $\mathbb{E}[\theta_{-i} | m_{-i} = 1] = 0.25$ , and if the first agent expects the other agent to play optimally, their expected payoff is now  $\mathbb{E}[\theta_{-i} | m_{-i} = 1] = 0.125$ , and so on. So trading will never be optimal.

Sending the message  $m_i = 0$  is a Bayes-Nash equilibrium, however, because given that every agent expects the other agents to play  $\sigma_i(\theta_i) = 0$ , a trade can never occur, so it is never in anyone's interest to send the message  $m_i = 1$  as it will not change the outcome. Formally,  $\forall m_i \mathbb{E}[u_i(x(\sigma_i, \sigma_{-i}), \theta_i) | \theta_i] = \theta_i = \mathbb{E}[u_i(x(m_i, \sigma_{-i}), \theta_i) | \theta_i]$ . So  $\sigma_i(\theta_i) = 0$  is a Bayes-Nash equilibrium.