# 'MATH+ECON+CODE' MASTERCLASS ON MATCHING MODELS, OPTIMAL TRANSPORT AND APPLICATIONS

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Thursday: "Multinomial choice" Block 10. Nonparametric multinomial choice

#### LEARNING OBJECTIVES: BLOCK 10

- ► Emax operator and generalized entropy of choice
- ► The Daly-Zachary-Williams theorem
- ► The GEV class

#### REFERENCES FOR BLOCK 10

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- ▶ Berry, Levinsohn, and Pakes (1995). "Automobile Prices in Market Equilibrium," *Econometrica*.
- ► Train. (2009). *Discrete Choice Methods with Simulation*. 2nd Edition. Cambridge University Press.
- ► G and Salanié (2017). "Cupid's invisible hands". Preprint.
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### Section 1

## EMAX OPERATORS AND DEMAND MAPS

#### **DISCRETE CHOICE MODELS**

- Assume a consumer is facing a number of options  $y \in \mathcal{Y}_0 = \mathcal{Y} \cup \{0\}$ , where y = 0 is a default option. The consumer is drawing a utility shock which is a vector  $\varepsilon = (\varepsilon_0, \dots, \varepsilon_{|\mathcal{Y}|}) \sim \mathbf{P}$  such that the utility of option y is  $U_v + \varepsilon_v$ , while the outside option yields utility  $\varepsilon_0$ .
- ▶ *U* is called vector of *systematic utilities*;  $\varepsilon$  is called vector of *utility shocks*.
- ► We assume thoughout that **P** has a density with respect to the Lebesgue measure, and has full support.
- ▶ The preferred option is the one which attains the maximum in

$$\max_{y\in\mathcal{Y}}\left\{U_{y}+\varepsilon_{y},\varepsilon_{0}\right\}.$$

#### **DEMAND MAP: DEFINITION**

▶ Let  $s_y = \sigma_y(U)$  be the probability of choosing option y, where  $\sigma$  is given by

$$\sigma_{y}\left(U\right) = \Pr(U_{y} + \varepsilon_{y} \geq U_{z} + \varepsilon_{z} \text{ for all } z \in \mathcal{Y}_{0}).$$

The map  $\sigma$  is called *demand map*, and the vector s is called vector of market shares, or vector of choice probabilities.

- ▶ Note that if  $s = \sigma(U)$ , then  $s_y > 0$  for all  $y \in \mathcal{Y}_0$  and  $\sum_{y \in \mathcal{Y}_0} s_y = 1$ .
- Note that because the distribution  ${\bf P}$  of  $\varepsilon$  is continuous, the probability of being indifferent between two options is zero, and hence we could have indifferently replaced weak preference  $\geq$  by strict preference >. Without this, choice probabilities may not have been well defined.

#### **DEMAND MAP: PROPERTIES**

- $ightharpoonup \sigma_{y}(U)$  is increasing in  $U_{y}$ .
- $ightharpoonup \sigma_{V}\left(U\right)$  is weakly decreasing in  $U_{V'}$  for  $y'\neq y$ .
- ▶ If one replaces  $(U_y)$  by  $(U_y + c)$ , for a constant c, one has  $\sigma(U + c) = \sigma(U)$ .

#### **NORMALIZATION**

▶ Because of the last property, we can normalize the utility of one of the alternatives. We will normalize the utility of the utility associated to y = 0, and hence take

$$U_0 = 0.$$

▶ Thus in the sequel,  $\sigma$  will be seen as a mapping from  $\mathbb{R}^{\mathcal{Y}}$  to the set of  $(s_y)_{y \in \mathcal{Y}}$  such that  $s_y > 0$  and  $\sum_{y \in \mathcal{Y}} s_y < 1$ , and the choice probability of alternative y = 0 is recovered by

$$s_0=1-\sum_{y\in\mathcal{Y}}s_y.$$

#### THE DALY-WILLIAMS-ZACHARY THEOREM

Define the expected indirect utility of consumers by

$$G(U) = \mathbb{E}\left[\max_{y \in \mathcal{Y}}(U_y + \varepsilon_y, \varepsilon_0)\right]$$

This is called *Emax operator*, a.k.a. *McFadden's surplus function*.

As the expectation of the maximum of terms which are linear in U, G is convex function in U (strictly convex in fact), and

$$\frac{\partial G}{\partial U_y}(U) = \Pr(U_y + \varepsilon_y \ge U_z + \varepsilon_z \text{ for all } z \in \mathcal{Y}_0).$$

But the right-hand side is simply the probability  $s_y$  of chosing option y; therefore, we get:

**Theorem (Daly-Zachary-Williams)**. The demand map  $\sigma$  is the gradient of the Emax operator G, that is

$$\sigma(U) = \nabla G(U). \tag{1}$$

## Section 2

## **EXAMPLES**

#### **EXAMPE 1: LOGIT**

► Assume that **P** is the distribution of i.i.d. *centered type I extreme value* a.k.a. *centered Gumbel* terms, which has c.d.f.

$$F(z) = \exp(-\exp(-x + \gamma))$$

where  $\gamma = 0.5772...$  (Euler's constant). The mean of this distribution is zero.

- ▶ Basic fact from extreme value theory: if  $\varepsilon_1,...,\varepsilon_n$  are i.i.d. Gumbel distributions, then max  $\{U_y + \varepsilon_y\}$  has the same distribution as  $\log\left(\sum_{y=1}^n \exp U_y\right) + \epsilon$ , where  $\epsilon$  is also a Gumbel. (Proof of this fact later).
- ► Notes:
  - ► This distribution is sometimes called the "Gumbel max" distribution, to contrast it with the distribution of its opposite, which is then called "Gumbel min".
  - ▶ The literature usually calls "standard Gumbel" the distribution with c.d.f.  $\exp(-\exp(-x))$ ; but that distribution has mean  $\gamma$ , which is why we slightly depart from the convention.

#### **EXAMPLE 1: LOGIT, EMAX FUNCTION AND DEMAND MAP**

► The Emax operator associated with the logit model can be given in closed form as

$$G(U) = \log \left(1 + \sum_{y \in \mathcal{Y}} \exp(U_y)\right)$$

where  $s_0 = 1 - \sum_{y \in \mathcal{Y}} s_y$ . This is called a *log-partition function*.

▶ As a result, the choice probability of alternative *y* is proportional to the exponential of the systematic utility associated with *U*, that is

$$\sigma_{y}(U) = \frac{\exp U_{y}}{1 + \sum_{y' \in \mathcal{Y}} \exp(U_{y'})}$$

which is called a Gibbs distribution.

ightharpoonup Assume that the random utility shock is scaled by a factor T. Then

$$\sigma_{y}(U) = \frac{\exp(U_{y}/T)}{1 + \sum_{v' \in \mathcal{V}} \exp(U_{v'}/T)}$$

which is sometimes called the *soft-max operator*, and converges as  $T \rightarrow 0$  toward

$$\max_{y \in \mathcal{V}} \{U_y, 0\}.$$

#### EXAMPLE 2: THE GENERALIZED EXTREME VALUE (GEV) CLASS

Let  $\mathbf{F}$  be a cumulative distribution such that function g defined by

$$g(x_1, ..., x_n) = -\log \mathbf{F}(-\log x_1, ..., -\log x_n)$$
 (2)

is positive homogeneous of degree 1. (This inverts into  $\mathbf{F}(u_1,...,u_n)=\exp\left(-g\left(e^{-u_1},...,e^{-u_n}\right)\right)$ ). We have by a theorem of McFadden (1978):

#### **THEOREM**

Let  $(\varepsilon_y)_{1 \le y \le n}$  be a random vector with c.d.f. **F**, and define

$$Z = \max_{y=1,\ldots,n} \{U_y + \varepsilon_y\}.$$

Then Z has the same distribution as  $\log g\left(e^{U_1},...,e^{U_n}\right) + \gamma + \epsilon$ , where  $\epsilon$  is a standard Gumbel. In particular,

$$\mathbb{E}\left[\max_{y=1,...,n}\left\{U_{y}+\varepsilon_{y}\right\}\right]=\log g\left(e^{U_{1}},...,e^{U_{n}}\right)+\gamma$$

where  $\gamma$  is the Euler constant  $\gamma \simeq 0.5772$ .

#### **EXAMPLE 2: GEV (CONTINUED)**

#### PROOF.

Let  $F_Z$  be the c.d.f. of  $Z = \max_{v=1,...,n} \{U_v + \varepsilon_v\}$ . One has

$$\begin{split} F_{Z}\left(z\right) &= \Pr\left(\max_{y=1,...,n}\left\{U_{y} + \varepsilon_{y}\right\} \leq z\right) = \Pr\left(\forall y: \ \varepsilon_{y} \leq z - U_{y}\right) \\ &= \mathbf{F}\left(z - U_{1},...,z - U_{n}\right) = \exp\left(-g\left(e^{U_{1} - z},...,e^{U_{n} - z}\right)\right) \\ &= \exp\left(-e^{-z}g\left(e^{U_{1}},...,e^{U_{n}}\right)\right) = \varphi\left(z - \log g\left(e^{U_{1}},...,e^{U_{n}}\right) - \gamma\right) \end{split}$$

where  $\varphi(z) := \exp\left(-e^{-(z-\gamma)}\right)$  is the cdf of the standard Gumbel distribution. Hence Z has the distribution of  $\log g\left(e^{U_1},...,e^{U_n}\right) + \gamma + \epsilon$ , where  $\epsilon$  is a standard Gumbel.

#### EXAMPLE 2: GEV, DEMAND MAP

ightharpoonup As a result, the choice probability of alternative y is

$$\sigma_{y}\left(U\right) = \frac{\frac{\partial g}{\partial x_{y}}\left(e^{U_{1}}, ..., e^{U_{n}}\right)}{g\left(e^{U_{1}}, ..., e^{U_{n}}\right)}e^{U_{y}}.$$

- ► The GEV framework has several commonly used examples: logit, nested logit, mixture of logit....
- ► We just saw the logit model, in which  $g(x_1,...,x_n) = e^{-\gamma} \sum_{y=1}^n x_y$ . In this case, the distribution of

$$Z = \max_{y=1,\dots,n} \{U_y + \varepsilon_y\}$$

is  $\log \sum_{y=1}^{n} e^{U_y} + \epsilon$ , where  $\epsilon$  is a standard Gumbel.

#### **EXAMPLE 3: NESTED LOGIT MODEL**

- ► The nested logit model is an instance of GEV model where alternatives can be grouped in nests. Eg, people choose their means of transportation (nest), and within this nest, a particular operator.
- ▶ Let  $\mathcal{X}$  be the set of nests and assume that for each nest x, there is a set  $\mathcal{Y}_x$  alternatives. Let  $U_{xy}$  be utility from alternative y in nest x, and  $\lambda_x \in [0,1]$  and

$$g(U_{xy}) = e^{-\gamma} \sum_{x \in \mathcal{X}} \left( \sum_{y \in \mathcal{Y}_x} U_{xy}^{1/\lambda_x} \right)^{\lambda_x}.$$

► In this case,

$$G(U) = \mathbb{E}\left[\max_{x \in \mathcal{X}} \max_{y \in \mathcal{Y}_{x}} \left\{U_{xy} + \varepsilon_{xy}\right\}\right] = \log \sum_{x \in \mathcal{X}} \left(\sum_{y \in \mathcal{Y}_{x}} e^{U_{xy}/\lambda_{x}}\right)^{\lambda_{x}}$$

$$\sigma_{xy}(U) = \frac{\left(\sum_{y \in \mathcal{Y}_{x}} e^{U_{xy}/\lambda_{x}}\right)^{\lambda_{x}}}{\sum_{x \in \mathcal{X}} \left(\sum_{y \in \mathcal{Y}_{x}} e^{U_{xy}/\lambda_{x}}\right)^{\lambda_{x}}} \frac{e^{U_{xy}/\lambda_{x}}}{\left(\sum_{y \in \mathcal{Y}_{x}} e^{U_{xy}/\lambda_{x}}\right)}$$

so the demand map has an interesting interpretation as "choice of nest then choice of alternative".

#### **EXAMPLE 3: NESTED LOGIT MODEL (CTD)**

Assume that  $(\varepsilon_1, \varepsilon_2)$  have a nested logit distribution with two nests, that is, their cdf is given by

$$\mathbf{F}\left(u_{1},u_{2}
ight)=\exp\left(-e^{-\gamma}\left(e^{-u_{1}/\lambda}+e^{-u_{2}/\lambda}
ight)^{\lambda}
ight).$$

- ▶ Particular cases:
  - ▶ When  $\lambda=1$ ,  $\varepsilon_1$  and  $\varepsilon_2$  are independent and one recovers the logit model.
  - $$\begin{split} & \quad \text{When } \lambda \rightarrow 0, \\ & \quad \mathbf{F}\left(u_1,u_2\right) = \exp\left(-e^{-\gamma}e^{\max\{-u_1,-u_2\}}\right) = \min\left\{\mathbf{F}\left(u_1\right),\mathbf{F}\left(u_2\right)\right\} \text{ and } \\ & \quad \text{therefore } \epsilon_1 \text{ and } \epsilon_2 \text{ are perfectly correlated.} \end{split}$$
- ► In general one can show that

$$\lambda = \sqrt{1 - \textit{cor}\left(\epsilon_1, \epsilon_2\right)}$$

This formula, due to Tiago de Oliviera, is not straightforward to prove and is the object of an optional problemset.

#### OTHER POPULAR EXAMPLES

- ► Probit model (later)
- Berry-Pakes' pure characteristics model (later)
- ► Berry-Levinsohn-Pakes' mixed logit coefficient model (later)

## Section 3

## **DEMAND INVERSION**

#### WHAT IS DEMAND INVERSION?

▶ In many settings, the econometrician observes the market shares  $s_y$  and wants to deduce the corresponding vector of systematic utilities. That is, we would like to solve:

**Problem**. Given a vector s with positive entries satisfying  $\sum_{y \in \mathcal{Y}} s_y < 1$ , characterize and compute the set

$$\sigma^{-1}\left(s\right)=\left\{ U\in\mathbb{R}^{\mathcal{Y}}:\sigma\left(U\right)=s\right\} .$$

► This problem is called "demand inversion," or "conditional choice probability inversion," or "identification problem." It is a central issue in econometrics/industrial organization and will be a key building block for matching models.

#### **DEMAND INVERSION VIA CONVEX ANALYSIS**

▶ We saw in Lecture 3 how to invert gradient of convex functions: if G is strictly convex and  $C^1$ , then

$$\sigma^{-1}\left(s\right) = \nabla G^{-1}(s) = \nabla G^{*}\left(s\right)$$
.

 $ightharpoonup G^*$  is the Legendre-Fenchel transform of G; we call it the *entropy of choice*, defined by

$$G^*(s) = \max_{U} \left\{ \sum_{y \in \mathcal{Y}} s_y U_y - G(U) \right\}.$$
 (3)

► Hence,  $\sigma^{-1}(s)$  is the vector U such that

$$U \in \arg\max_{U} \left\{ \sum_{y \in \mathcal{Y}} s_y U_y - G(U) \right\}.$$

#### ENTROPY OF CHOICE

▶ Convex duality implies that if s and U are related by  $s \in \partial G(U)$ , then

$$G(U) = \sum_{y \in \mathcal{Y}} s_y U_y - G_x^*(s). \tag{4}$$

▶ But letting  $Y = \arg\max_{y} \{U_{y} + \varepsilon_{y}\}, G(U) = \mathbb{E}[U_{Y} + \varepsilon_{Y}]$  implies

$$G(U) = \sum_{y \in \mathcal{Y}} s_y U_y + \mathbb{E}\left[\varepsilon_{Y}\right],$$

thus one has

$$G^*(s) = -\mathbb{E}\left[\varepsilon_{Y}\right]. \tag{5}$$

Hence, the entropy of choice  $G^*(s)$  is interpreted as minus the expected amount of heterogeneity needed to rationalize the choice probabilities s.

► Then

$$G^{*}\left(s\right) = s_{0}\log(s_{0}) + \sum_{y \in \mathcal{Y}} s_{y}\log s_{y}$$

where  $s_0 = 1 - \sum_{y \in \mathcal{Y}} s_y$ . Hence,  $G^*$  is a bona fide entropy function when **P** is Gumbel–hence the name of *entropy of choice* in general.

► As a result,

$$\sigma_{y}^{-1}\left(s\right) = \log \frac{s_{y}}{s_{0}}$$

which is the celebrated log-odds ratio formula: the log of the odds of alternatives y and 0 identify the difference between the systematic utilities of these alternatives.

# EXAMPLE: ENTROPY OF CHOICE AND IDENTIFICATION, NESTED LOGIT MODEL

▶ The entropy of choice  $G^*$  in the nested logit model is given by

$$G^{*}(s) = \sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}_{x}} \lambda_{x} s_{xy} \ln s_{xy} + \sum_{x \in \mathcal{X}} (1 - \lambda_{x}) \left( \sum_{z \in \mathcal{Y}_{x}} s_{xz} \right) \ln \left( \sum_{z \in \mathcal{Y}_{x}} s_{xz} \right)$$
(6)

if  $s_{xy} \geq 0$  and  $\sum_{x \in \mathcal{X}} \sum_{y \in \mathcal{Y}_x} s_{xy} = 1$ ,  $G^*(s) = +\infty$  otherwise.

▶ Identification in the nested logit model: with normalization

$$\begin{split} & \sum_{x \in \mathcal{X}} \left( \sum_{y \in \mathcal{Y}_x} e^{U_{xy}/\lambda_x} \right)^{\lambda_x} = 1, \text{ one has} \\ & s_{xy} = \left( \sum_{y \in \mathcal{Y}_x} e^{U_{xy}/\lambda_x} \right)^{\lambda_x - 1} e^{U_{xy}/\lambda_x}, \text{ thus} \\ & \sum_{y \in \mathcal{Y}_x} e^{U_{xy}/\lambda_x} = \left( \sum_{y \in \mathcal{Y}_x} s_{xy} \right)^{1/\lambda_x}, \text{ therefore} \\ & U_{xy} = \lambda_x \log s_{xy} - (\lambda_x - 1) \log \sum s_{xy}. \end{split}$$

► And now. let's code!