



Microeconomics III: Problem Set 5^a

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^aSlides created for exercise class 3 and 4, with reservation for possible errors.

Kahoot!

PS5, Ex. 1 (A): Dynamic game (backwards induction)

PS5, Ex. 2 (A): Dynamic game (strategy sets)

PS5, Ex. 3 (A): Stackelberg Duopoly (empty threats)

PS5, Ex. 4: The Mutated Seabass (backwards induction)

PS5, Ex. 5: Three player game (backwards induction)

PS5, Ex. 6: Stackelberg assignment (backwards induction)

PS5, Ex. 7: Dynamic game (proper subgames)

PS5, Ex. 8:

PS5, Ex. 9:

Code examples

Kahoot!

Form a group for each table:

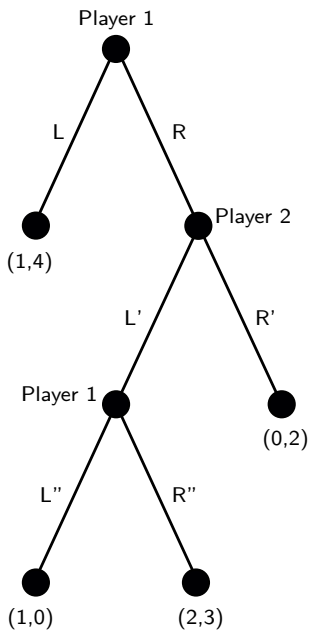
- Get prepared to answer the three A exercises as a team (5 min).



**PS5, Ex. 1 (A): Dynamic game
(backwards induction)**

PS5, Ex. 1 (A): Dynamic game (backwards induction)

- Consider the dynamic game shown in extensive form. Solve it by backwards induction.

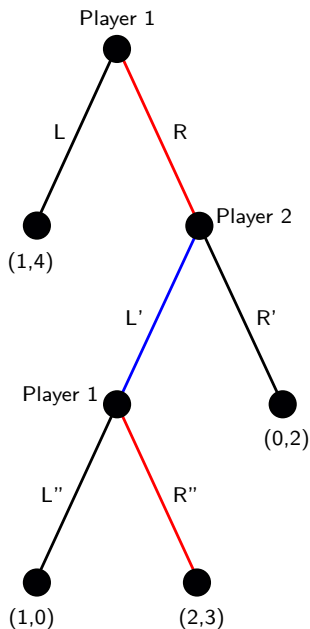


PS5, Ex. 1 (A): Dynamic game (backwards induction)

- Consider the dynamic game shown in extensive form. Solve it by backwards induction.

The backwards induction solution is the full strategy profile given by the subgame perfect NE:

$$BI = (s_1^*, s_2^*) = (R \ R'', \ L')$$

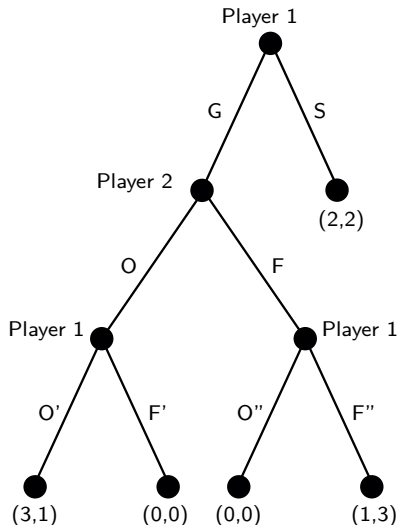


**PS5, Ex. 2 (A): Dynamic game
(strategy sets)**

PS5, Ex. 2 (A): Extended Battle of the Sexes Game (strategy sets)

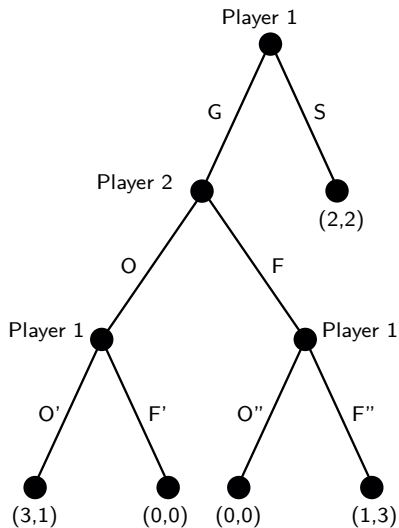
Consider the game in the figure.

- (a) Write up the strategy sets of the players.
- (b) Write up the normal form (bi-matrix).
- (c) Find the Nash Equilibria.
- (d) Find the backwards induction outcome.



PS5, Ex. 2.a (A): Extended Battle of the Sexes Game (strategy sets)

- (a) Write up the strategy sets of the players.



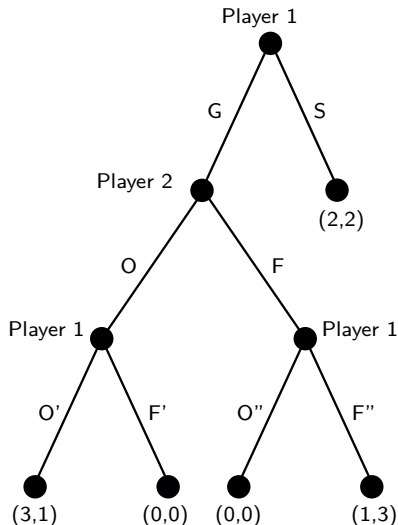
PS5, Ex. 2.a (A): Extended Battle of the Sexes Game (strategy sets)

- (a) Write up the strategy sets of the players.

The two strategy sets are:

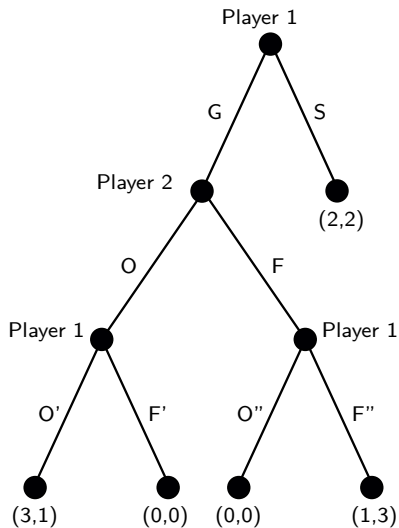
$$S_1 = \{ (G, O', O''); (G, O', F''); \\ (G, F', O''); (G, F', F''); \\ (S, O', O''); (S, O', F''); \\ (S, F', O''); (S, F', F'') \}$$

$$S_2 = \{ O; F \}$$



PS5, Ex. 2.b (A): Extended Battle of the Sexes Game (strategy sets)

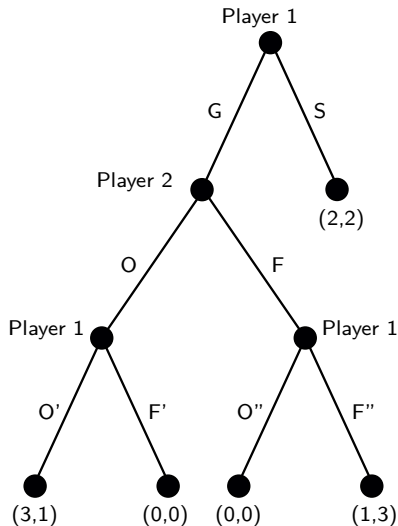
- (b) Write up the normal form (bi-matrix).



PS5, Ex. 2.b (A): Extended Battle of the Sexes Game (strategy sets)

- (b) Write up the normal form (bi-matrix).

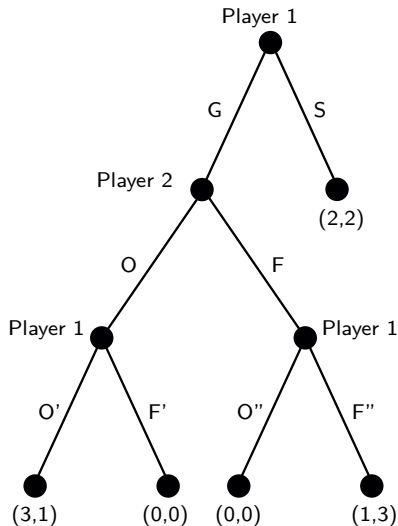
		Player 2	
		O	F
Player 1	(G O' O'')	3, 1	0, 0
	(G O' F'')	3, 1	1, 3
	(G F' O'')	0, 0	0, 0
	(G F' F'')	0, 0	1, 3
	(S O' O'')	2, 2	2, 2
	(S O' F'')	2, 2	2, 2
	(S F' O'')	2, 2	2, 2
	(S F' F'')	2, 2	2, 2



PS5, Ex. 2.c (A): Extended Battle of the Sexes Game (strategy sets)

(c) Find the Nash Equilibria.

		Player 2	
		O	F
Player 1	(G O' O'')	3, 1	0, 0
	(G O' F'')	3, 1	1, 3
	(G F' O'')	0, 0	0, 0
	(G F' F'')	0, 0	1, 3
	(S O' O'')	2, 2	2, 2
	(S O' F'')	2, 2	2, 2
	(S F' O'')	2, 2	2, 2
	(S F' F'')	2, 2	2, 2



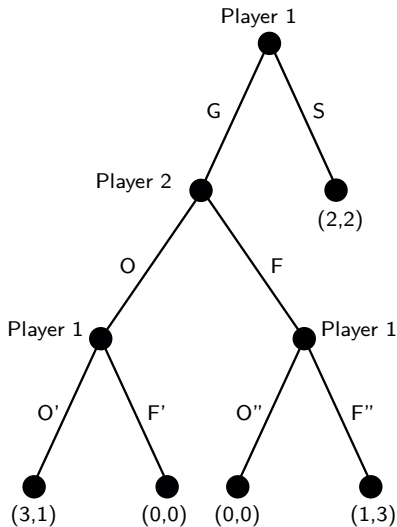
PS5, Ex. 2.c (A): Extended Battle of the Sexes Game (strategy sets)

(c) Find the Nash Equilibria.

The five NE are:

$$NE = \{ (G \ O' \ O'', O); (S \ O' \ O'', F); \\ (S \ O' \ F'', F); (S \ F' \ O'', F); \\ (S \ F' \ F'', F) \}$$

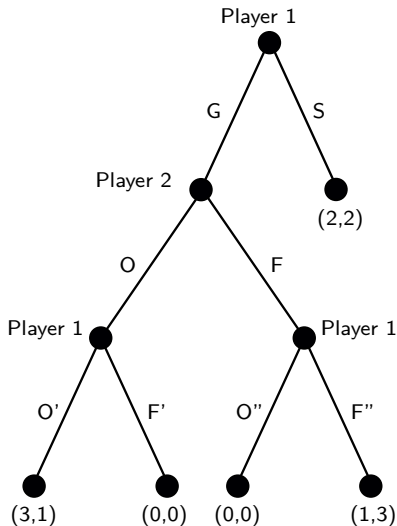
		Player 2	
		O	F
Player 1	(G O' O'')	3, 1	0, 0
	(G O' F'')	3, 1	1, 3
	(G F' O'')	0, 0	0, 0
	(G F' F'')	0, 0	1, 3
	(S O' O'')	2, 2	2, 2
	(S O' F'')	2, 2	2, 2
	(S F' O'')	2, 2	2, 2
	(S F' F'')	2, 2	2, 2



PS5, Ex. 2.d (A): Extended Battle of the Sexes Game (strategy sets)

- (d) Find the backwards induction outcome.

		Player 2	
		O	F
Player 1	(G O' O'')	3, 1	0, 0
	(G O' F'')	3, 1	1, 3
	(G F' O'')	0, 0	0, 0
	(G F' F'')	0, 0	1, 3
	(S O' O'')	2, 2	2, 2
	(S O' F'')	2, 2	2, 2
	(S F' O'')	2, 2	2, 2
	(S F' F'')	2, 2	2, 2



PS5, Ex. 2.d (A): Extended Battle of the Sexes Game (strategy sets)

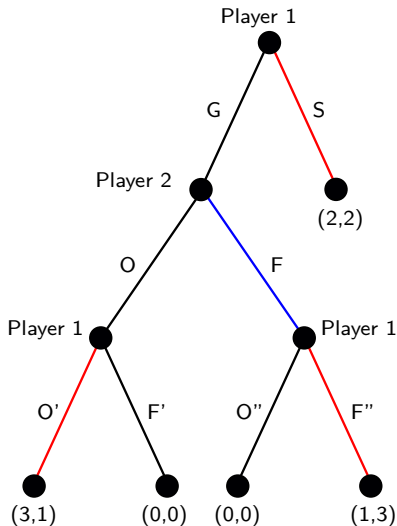
(d) Find the backwards induction outcome.

BI gives the unique SPNE:

$$SPNE = (s_1^*, s_2^*) = (S \ O' \ F'', F)$$

The NE $(G \ O' \ O'', O)$ is not subgame perfect as player 1's strategy is weakly dominated by $(G \ O' \ F'')$. A SPNE needs to be rational on and off the equilibrium path, thus, O'' is an empty threat.

		Player 2	
		O	F
Player 1	(G O' O'')	3, 1	0, 0
	(G O' F'')	3, 1	1, 3
	(G F' O'')	0, 0	0, 0
	(G F' F'')	0, 0	1, 3
	(S O' O'')	2, 2	2, 2
	(S O' F'')	2, 2	2, 2
	(S F' O'')	2, 2	2, 2
	(S F' F'')	2, 2	2, 2



**PS5, Ex. 3 (A): Stackelberg Duopoly
(empty threats)**

Consider the Stackelberg game we saw in the lecture (Stackelberg Duopoly presented in the end of Lecture 3). We solved for the backwards induction outcome. But if we were to look for Nash Equilibria, there are many. In fact there is an infinite amount. But they rely on 'empty threats'. To see why, consider the following equilibrium. The *follower* says to the *leader*: I want you to produce \hat{q}_1 (where $\hat{q}_1 < a$) and then I will produce $\hat{q}_2 = BR_2(\hat{q}_1)$. If you produce $q_1 \neq \hat{q}_1$ then I will set $q_2 = a - q_1$ such that you make zero profit.

- (a) Write up the equilibria described above formally, using for the strategies the notation of the 'Simple Dynamic Game'.
- (b) Explain why this kind of equilibrium does not survive backwards induction unless $\hat{q}_1 = BR_1(BR_2(q_1))$, which is the Stackelberg outcome we derived in the lecture.

PS5, Ex. 3.a (A): Stackelberg Duopoly (empty threats)

(a) Write up the equilibria described above formally, using for the strategies the notation of the 'Simple Dynamic Game'.

In the equilibrium, player 1 plays any quantity \hat{q}_1 he is dictated, player 2 then play his BR.

To finalize the equilibrium, we need players 2's response off the equilibrium path, i.e. if player one plays $q_1 \neq \hat{q}_1$

The NE consists of the strategies:

$$q_1^* = \hat{q}_1, \quad \text{where } q_1^* < a$$

$$q_2^* = \begin{cases} BR_2(\hat{q}_1) & \text{if } q_1 = \hat{q}_1 \\ a - q_1 & \text{if } q_1 \neq \hat{q}_1 \end{cases}$$

(b) Explain why this kind of equilibrium does not survive backwards induction unless $\hat{q}_1 = BR_1(BR_2(q_1))$, which is the Stackelberg outcome we derived in the lecture.

PS5, Ex. 3.b (A): Stackelberg Duopoly (empty threats)

- (a) Write up the equilibria described above formally, using for the strategies the notation of the 'Simple Dynamic Game'.

In the equilibrium, player 1 plays any quantity \hat{q}_1 he is dictated, player 2 then play his BR.

To finalize the equilibrium, we need players 2's response off the equilibrium path, i.e. if player one plays $q_1 \neq \hat{q}_1$

The NE consists of the strategies:

$$q_1^* = \hat{q}_1, \quad \text{where } q_1^* < a$$
$$q_2^* = \begin{cases} BR_2(\hat{q}_1) & \text{if } q_1 = \hat{q}_1 \\ a - q_1 & \text{if } q_1 \neq \hat{q}_1 \end{cases}$$

- (b) Explain why this kind of equilibrium does not survive backwards induction unless $\hat{q}_1 = BR_1(BR_2(q_1))$, which is the Stackelberg outcome we derived in the lecture.

For the case where player 2 selects \hat{q}_1 to be the quantity selected in the Stackelberg game, the outcome will be exactly as in the Stackelberg game.

For the case where player 2 selects another outcome, player 1 will know that no matter his choice of q_1 , player 2 will always maximize his profits, selecting his q_2 using his best response function. Thus, the threat is empty.

PS5, Ex. 4: The Mutated Seabass (backwards induction)

PS5, Ex. 4: The Mutated Seabass (backwards induction)

Consider a game where two evil organizations, rather prosaically named A and B, are battling for world domination. The battle takes the form of a three-stage game. Organization A is on the verge of acquiring a new powerful weapon, the *mutated seabass*. In stage 1 of the game, they decide whether to acquire the weapon or not. Their choice is observed by organization B. In stage 2, organization B decides whether to attack organization A. If an attack occurs, the game stops. If no attack occurs, it moves to stage 3, where organization A decides whether or not to attack organization B. The payoffs are as follows. If no-one attacks the other, the payoffs to both organizations are 0. If B attacks A, then the payoffs to both organizations are .1. The same if A attacks B, without having acquired the seabass weapon. If, on the other hand, A acquires the weapon, the payoffs from A attacking B are 2 to A and -2 to B.

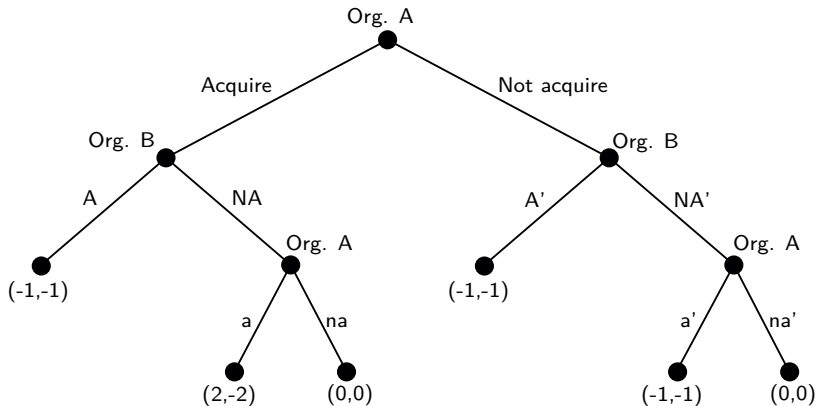
- (a) Draw the game tree that corresponds to the game. What are the strategies of the players?
- (b) What is the backwards induction outcome?
- (c) What is the intuition for the outcome? What role do you think it plays that B observes if A acquires the weapon or not?

PS5, Ex. 4.a: The Mutated Seabass (backwards induction)

(a) Draw the game tree that corresponds to the game. What are the strategies of the players?

$S_A = \{(Acquire, a, a'), (Acquire, a, na'), (Acquire, na, a'), (Acquire, na, na'),$
 $(Not\ acquire, a, a'), (Not\ acquire, a, na'), (Not\ acquire, na, a'), (Not\ acquire, na, na')\}$

$S_B = \{(A, A'), (A, NA'), (NA, A'), (NA, NA')\}$



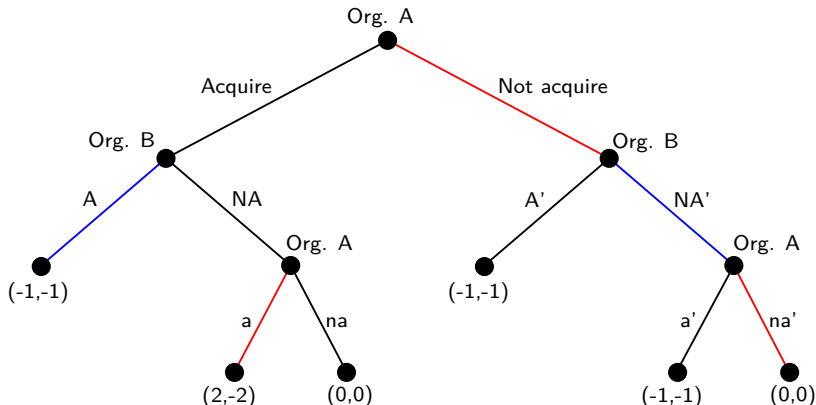
PS5, Ex. 4.b: The Mutated Seabass (backwards induction)

(b) What is the backwards induction outcome?

3rd stage: Org. A will choose to attack if having acquired the weapon and not attack if not having acquired the weapon.

2nd stage: Org. B will choose to attack if Org. A has acquired the weapon and not attack if they have not acquired the weapon.

1st stage: Org. A will choose to not acquire the weapon in order to signal peaceful intentions to Org. B, i.e. giving the payoffs (0, 0).



$$SPNE = \{S_A, S_B\} = \{(Not\ acquire, a, na'), (A, NA')\}$$

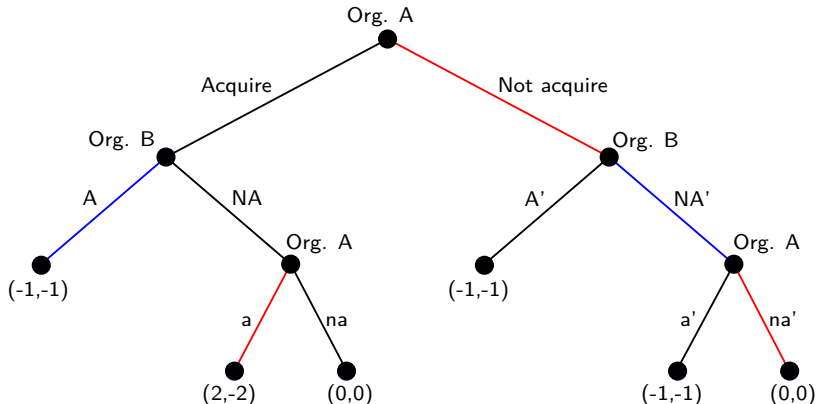
PS5, Ex. 4.c: The Mutated Seabass (backwards induction)

(c) What is the intuition for the outcome?

3rd stage: Org. A does only benefit from attacking if having acquired the weapon.

2nd stage: Org. B will only choose to attack if Org. A has acquired the weapon.

1st stage: Not acquiring the weapon is a credible signal that Org. A will not attack.



What role do you think it plays that B observes if A acquires the weapon or not?

I.e. what is the outcome if Organization A cannot send a signal in the 1st stage?

PS5, Ex. 4.c: The Mutated Seabass (backwards induction)

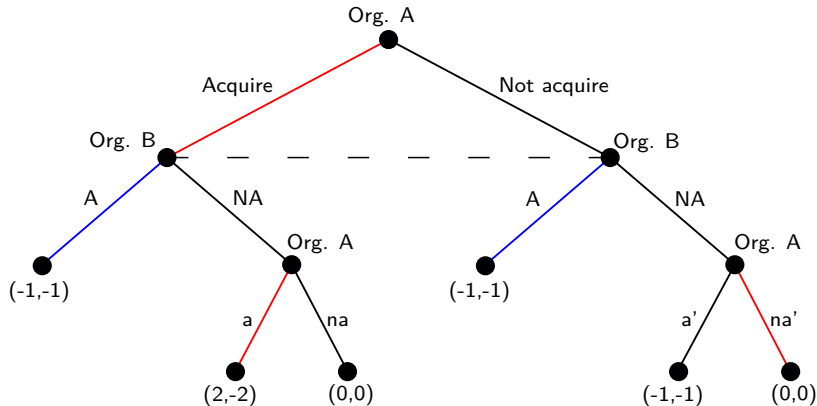
(c) What role do you think it plays that B observes if A acquires the weapon or not?

I.e. what is the outcome if Organization A cannot send a signal in the 1st stage?

3rd stage: [unchanged] Org. A will choose to attack if having acquired the weapon and not attack if not having acquired the weapon.

2nd stage: Knowing that Org. A will attack if having acquired the weapon, Org. B chooses to attack first, giving the payoffs $(-1, -1)$ regardless of stage one.

1st stage: Org. A cannot affect the outcome, but acquires it in case Org. B deviates.



$$SPNE = \{S_A, S_B\} = \{(Acquire, a, na'), A\}$$

**PS5, Ex. 5: Three player game
(backwards induction)**

PS5, Ex. 5: Three player game (backwards induction)

Consider the game below where player 1 chooses the matrix (A or B), player 2 chooses the row (C or D), and player 3 chooses the column (E or F). In each cell, the first number gives the payoff of Player 1, the second number the payoff of Player 2, and the third number the payoff of Player 3.

	E	F
C	5, 2, 2	2, 1, 1
D	0, 1, 1	1, 0, 0

A

	E	F
C	6, 0, 1	3, 1, 2
D	1, 1, 0	2, 2, 1

B

- (a) Suppose first that the game is static, such that all three players move simultaneously. Find all the pure-strategy Nash Equilibria.
- (b) Now suppose the game is dynamic: Player 1 moves first, and then, after having observed his move, Player 2 moves, and, finally, after having observed the first two moves, Player 3 moves. Draw the game tree and solve by backwards induction.
- (c) Discuss the differences in the results you find

PS5, Ex. 5.a: Three player game (backwards induction)

- (a) Suppose first that the game is static, such that all three players move simultaneously. Find all the pure-strategy Nash Equilibria.

	E	F		E	F
C	5, 2, 2	2, 1, 1	C	6, 0, 1	3, 1, 2
D	0, 1, 1	1, 0, 0	D	1, 1, 0	2, 2, 1
	A			B	

Iterated Elimination of Strictly Dominated Strategies (IESDS): For player 1, A is strictly dominated by B. In the reduced form game, C and E are strictly dominated for Player 2 and 3 respectively.

⇒ Pure Strategy Nash Equilibrium: $\{B, D, F\}$ with outcome (2,2,1).

PS5, Ex. 5.b: Three player game (backwards induction)

Consider the game below where player 1 chooses the matrix (A or B), player 2 chooses the row (C or D), and player 3 chooses the column (E or F). In each cell, the first number gives the payoff of Player 1, the second number the payoff of Player 2, and the third number the payoff of Player 3.

	E	F
C	5, 2, 2	2, 1, 1
D	0, 1, 1	1, 0, 0

A

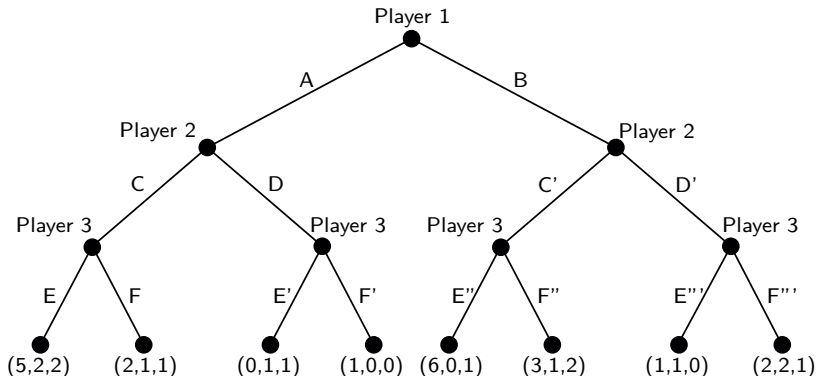
	E	F
C	6, 0, 1	3, 1, 2
D	1, 1, 0	2, 2, 1

B

- (b) Now suppose the game is dynamic: Player 1 moves first, and then, after having observed his move, Player 2 moves, and, finally, after having observed the first two moves, Player 3 moves. **Draw the game tree** and solve by backwards induction.

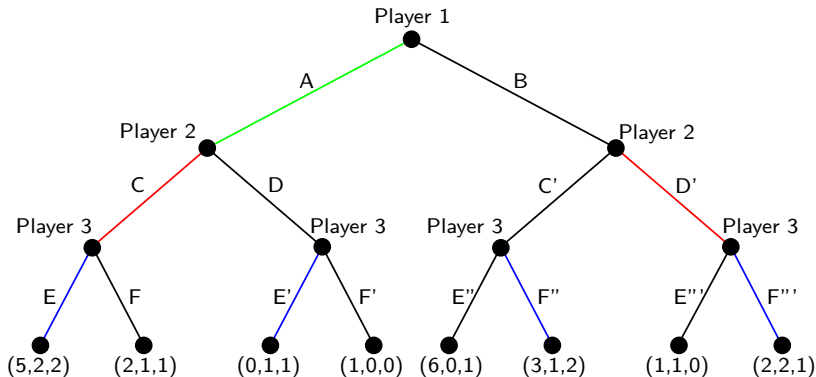
PS5, Ex. 5.b: Three player game (backwards induction)

- (b) Now suppose the game is dynamic: Player 1 moves first, and then, after having observed his move, Player 2 moves, and, finally, after having observed the first two moves, Player 3 moves. Draw the game tree and ***solve by backwards induction***.



PS5, Ex. 5.b: Three player game (backwards induction)

- (b) Now suppose the game is dynamic: Player 1 moves first, and then, after having observed his move, Player 2 moves, and, finally, after having observed the first two moves, Player 3 moves. Draw the game tree and solve by backwards induction.

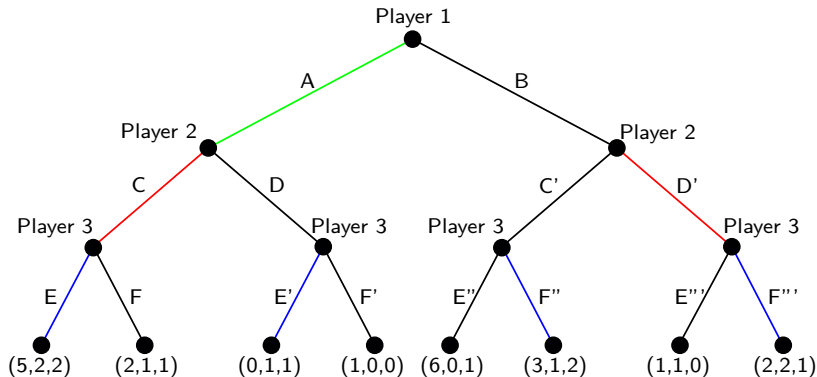


\Rightarrow Subgame Perfect NE: $\{A, (C, D'), (E, E', F'', F''')\}$ with outcome (5,2,2).

PS5, Ex. 5.c: Three player game (backwards induction)

(a) $PSNE = \{B, D, F\}$ with outcome $(2,2,1)$.

(b) $SPNE = \{A, (C, D'), (E, E', F'', F''')\}$ with outcome $(5,2,2)$.

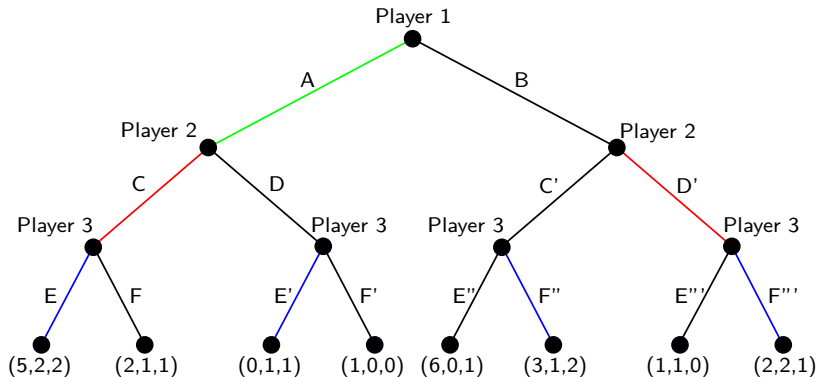


(c) *Discuss the differences in the results you find.*

PS5, Ex. 5.c: Three player game (backwards induction)

(a) $PSNE = \{B, D, F\}$ with outcome $(2,2,1)$.

(b) $SPNE = \{A, (C, D'), (E, E', F'', F''')\}$ with outcome $(5,2,2)$.



(c) Discuss the differences in the results you find.

In the static game: (A, C, E) with outcome $(5,2,2)$ cannot be a solution. Player 2 and 3 will not play C and E as they expect player 1 to play B instead and get $(6,0,1)$.

In the dynamic game: Player 1 can expect at higher payoff on the left side of the tree than on the right side, thus, commits to A , allowing Player 2 and 3 to play C and E .

**PS5, Ex. 6: Stackelberg assignment
(backwards induction)**

PS5, Ex. 6: Stackelberg assignment (backwards induction)

Two students are working together on the next assignment. Student i , $i = 1, 2$, exerts an effort $y_i \geq 0$. The resulting quality of the assignment is

$$q(y_1, y_2) = y_1 y_2.$$

Exerting effort is costly, but the costs differ, since one student likes game theory more than the other. More precisely, the cost functions are

$$C_1(y_1) = \frac{1}{3}(y_1)^3,$$

$$C_2(y_2) = (y_2)^2.$$

The payoff for student i , U_i , is equal to the quality of the assignment less his cost of effort.

$$U_1(y_1, y_2) = q(y_1, y_2) - C_1(y_1),$$

$$U_2(y_1, y_2) = q(y_1, y_2) - C_2(y_2).$$

(a) Consider the game where both of them choose their effort levels simultaneously and independently. Derive the best response functions. Find the (pure strategy) Nash equilibrium (y_1^{NE}, y_2^{NE}) with $y_1^{NE}, y_2^{NE} > 0$.

(b) Suppose now that Student 1 chooses his effort first, then sends the assignment on to Student 2. Student 2 observes how much effort Student 1 has exerted, makes his own choice of effort, and then submits. Solve by backwards induction.

(c) Compare the outcomes in (a) and (b) with respect to the payoffs of the students. Which game does each of the two students prefer? Give an intuitive explanation of your answer.

(d) Find the socially optimal levels of effort (y_1^{SO}, y_2^{SO}) , i.e., the levels that maximize the sum of the two students' payoffs. Calculate the payoff that the two students get in the social optimum.

- (a) Consider the game where both of them choose their effort levels simultaneously and independently. Derive the best response functions. Find the (pure strategy) Nash equilibrium (y_1^{NE}, y_2^{NE}) with $y_1^{NE}, y_2^{NE} > 0$.

Information so far:

Quality: $q(y_1, y_2) = y_1 y_2$.

Costs: $C_1(y_1) = \frac{1}{3}(y_1)^3$, $C_2(y_2) = (y_2)^2$.

Payoffs: $U_i(y_i, y_j) = q(y_i, y_j) - C_i(y_i)$.

- (b) Suppose now that Student 1 chooses his effort first, then sends the assignment on to Student 2. Student 2 observes how much effort Student 1 has exerted, makes his own choice of effort, and then submits. Solve by backwards induction.

Information so far:

Quality: $q(y_1, y_2) = y_1 y_2$.

Costs: $C_1(y_1) = \frac{1}{3}(y_1)^3$, $C_2(y_2) = (y_2)^2$.

Payoffs: $U_i(y_i, y_j) = q(y_i, y_j) - C_i(y_i)$.

- (c) Compare the outcomes in (a) and (b) with respect to the payoffs of the students. Which game does each of the two students prefer? Give an intuitive explanation of your answer.

Information so far:

Quality: $q(y_1, y_2) = y_1 y_2$.

Costs: $C_1(y_1) = \frac{1}{3}(y_1)^3$, $C_2(y_2) = (y_2)^2$.

Payoffs: $U_i(y_i, y_j) = q(y_i, y_j) - C_i(y_i)$.

- (d) Find the socially optimal levels of effort (y_1^{SO}, y_2^{SO}) , i.e., the levels that maximize the sum of the two students' payoffs. Calculate the payoff that the two students get in the social optimum.

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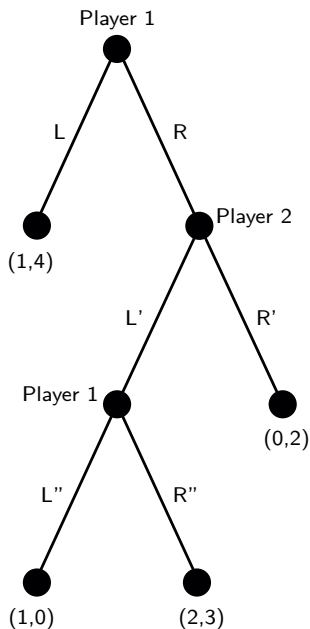
Payoffs: $U_i(y_i, y_j) = q(y_i, y_j) - C_i(y_i)$.

PS5, Ex. 7: Dynamic game (proper subgames)

PS5, Ex. 7: Dynamic game (proper subgames)

Consider the game from exercise 1.

- (a) Write down the strategies of the two players. How many proper subgames are there (so not including the entire game itself)?
- (b) Write down the Subgame-Perfect Nash Equilibrium (SPNE). Compare the SPNE to the backwards-induction outcome which you found in question 1.
- (c) Write down the normal form of the game and find all (pure strategy) Nash equilibria. Compare to the set of SPNE and comment.



PS5, Ex. 7.a: Dynamic game (proper subgames)

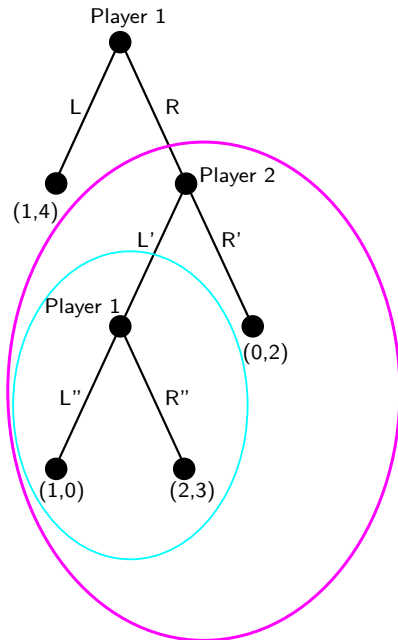
- (a) Write down the strategies of the two players. How many proper subgames are there (so not including the entire game itself)?

The two strategy sets are:

$$S_1 = \{ (L, L''); (L, R''); (R, L''); (R, R'') \}$$

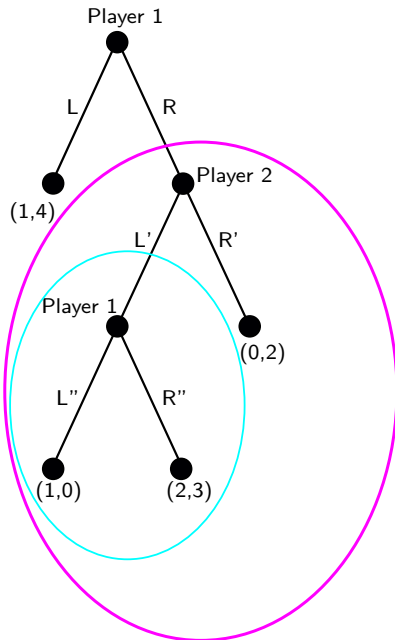
$$S_2 = \{ L' ; R' \}$$

There are two proper subgames.



PS5, Ex. 7.b: Dynamic game (proper subgames)

- (b) Write down the Subgame-Perfect Nash Equilibrium (SPNE). Compare the SPNE to the backwards-induction outcome which you found in question 1.

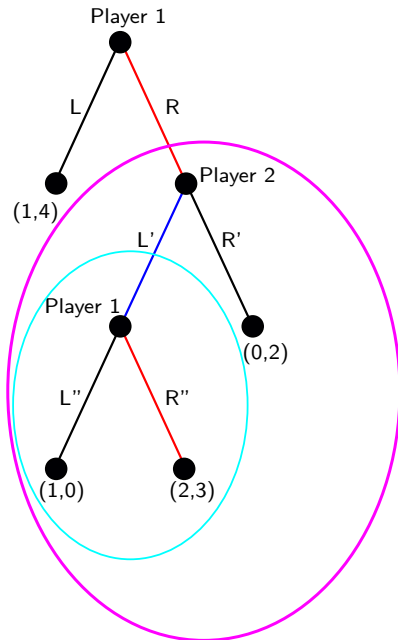


PS5, Ex. 7.b: Dynamic game (proper subgames)

- (b) Write down the Subgame-Perfect Nash Equilibrium (SPNE). Compare the SPNE to the backwards-induction outcome which you found in question 1.

$$SPNE = (s_1^*, s_2^*) = (R, R'', L')$$

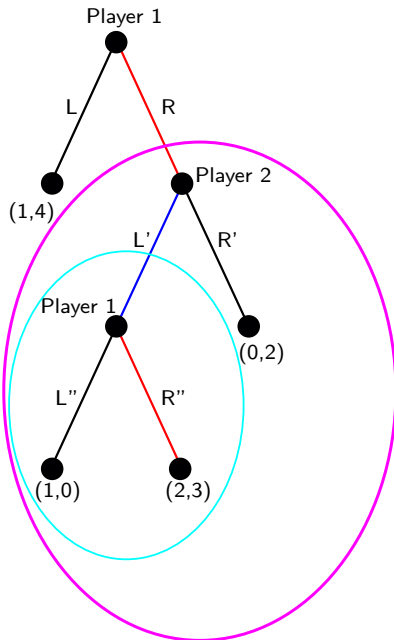
Which coincides with the backwards induction solution from Exercise 1!



PS5, Ex. 7.c: Dynamic game (proper subgames)

Consider the game from exercise 1.

- (c) *Write down the normal form of the game* and find all (pure strategy) Nash equilibria. Compare to the set of SPNE and comment.

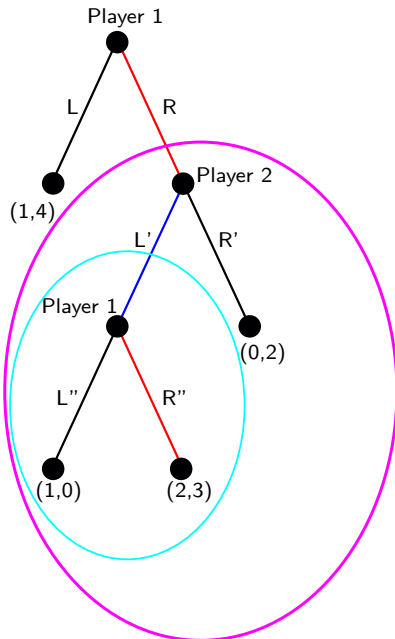


PS5, Ex. 7.c: Dynamic game (proper subgames)

Consider the game from exercise 1.

- (c) Write down the normal form of the game and **find all (pure strategy) Nash equilibria**. Compare to the set of SPNE and comment.

		Player 2	
		L'	R'
Player 1	(L, L'')	1, 4	1, 4
	(L, R'')	1, 4	1, 4
	(R, L'')	1, 0	0, 2
	(R, R'')	2, 3	0, 2



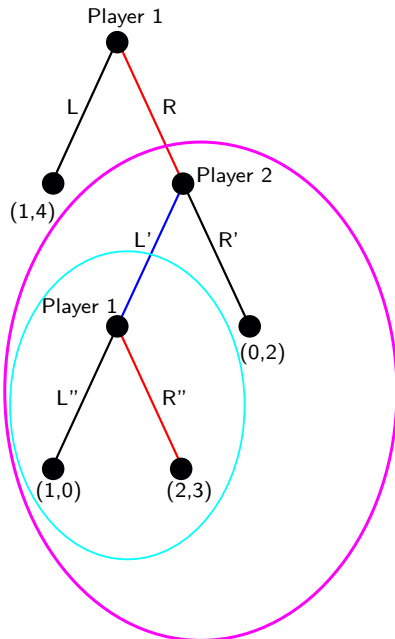
PS5, Ex. 7.c: Dynamic game (proper subgames)

Consider the game from exercise 1.

- (c) Write down the normal form of the game and **find all (pure strategy) Nash equilibria**. Compare to the set of SPNE and comment.

		Player 2	
		L'	R'
Player 1	(L, L'')	1, 4	1, 4
	(L, R'')	1, 4	1, 4
	(R, L'')	1, 0	0, 2
	(R, R'')	2, 3	0, 2

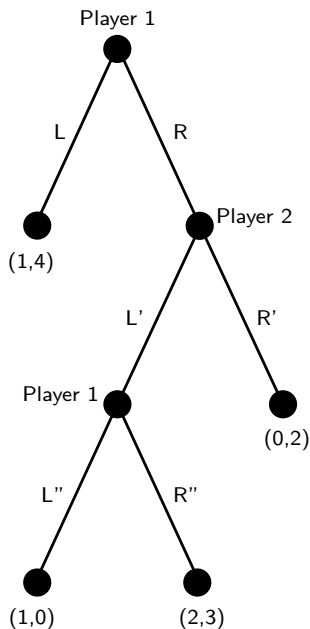
PSNE =



PS5, Ex. 8:

PS5, Ex. 9:

Code examples



Matrix, no player names:

	L (q)	R (1-q)
T (p)		
B (1-p)		

Matrix, no colors:

		Player 2	
		L (q)	R (1-q)
Player 1	T (p)		
	B (1-p)		

Matrix, with colors:

		Player 2	
		L (q)	R (1-q)
Player 1	T (p)	,	
	B (1-p)		