

Optimization Without Derivatives

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The Plan

Motivation

Black-box Optimization

Direct Search Methods Model-Based Methods Some Global Optimization

Simulation-Based Optimization and Structure

Nonlinear Least Squares Least Squares and Partials Constraints Nonsmoothness



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Simulation-Based Optimization

$$\min_{x\in\mathbb{R}^n}\left\{f(x)=F[S(x)]:c(S(x))\leq 0\right\}$$

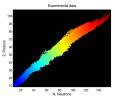
- S (numerical) simulation output, often "noisy" (even when deterministic)
- \diamond Derivatives $\nabla_x S$ often unavailable or prohibitively expensive to obtain/approximate directly
- \diamond S can contribute to objective and/or constraints
- Single evaluation of S could take seconds/minutes/hours/days Evaluation is a bottleneck for optimization

Functions of complex numerical simulations arise everywhere









Computing is Responsible for Pervasiveness of Simulations in Sci&Eng



Argonne's Blue Gene/P (2008: 163,840 cores)
Currently 67th fastest in the world



Argonne's Blue Gene/Q (2012: 786,432 cores) Currently 5th fastest in the world

- ◆ Parallel/multi-core environments increasingly common
 - Small clusters/multi-core desktops/multi-core laptops pervasive
 - Leadership class machines increasingly parallel
- Simulations (the "forward problem") become faster/more realistic/more complex

Improvements from Algorithms Can Trump Those From Hardware

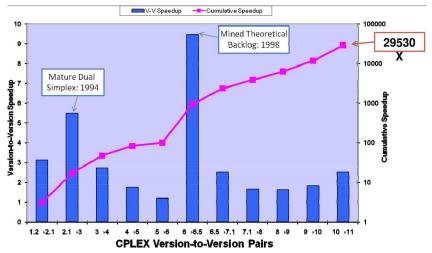
Martin Grötschel's production planning benchmark problem (a MIP):

1988 solve time using current computers and LP algorithms: 82 years

2003 solve time using current computers and LP algorithms: 1 minute

 $^{\diamond}$ Speed up of 43,000,000X $10^3 \mathrm{X}$ from processor improvements $10^4 \mathrm{X}$ additional from algorithmic improvements

Improvements from Algorithms Can Trump Those From Hardware



1991 (v1.2) to 2007 (v11.0): Moore's Law transistor speedup: $\approx 256 \text{X}$ [Slide from Bixby (CPLEX/GUROBI)]: Solves 1,852 MIPs

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Derivative-Free Optimization

"Some derivatives are unavailable for optimization purposes"

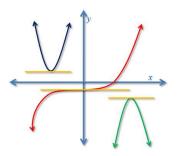
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The Challenge: Optimization is tightly coupled with derivatives

Typical optimality (no noise, smooth functions)

$$\nabla_x f(x_*) + \lambda^T \nabla_x c_E(x_*) = 0, c_E(x_*) = 0$$



(sub)gradients $\nabla_x f$, $\nabla_x c$ enable:

- Faster feasibility
- Faster convergence
 - Guaranteed descent
 - Approximation of nonlinearities
- Better termination
 - Measure of criticality $\|\nabla_x f\|$ or $\|\mathcal{P}_{\Omega}(\nabla_x f)\|$
- Sensitivity analysis
 - Correlations, standard errors, UQ, . . .

Ways to Get Derivatives

(assuming they exist)

Handcoding (HC)

- "Army of students/programmers"
 - ? Prone to errors/conditioning
 - ? Intractable as number of ops increases

Algorithmic/Automatic Differentiation (AD)

- "Exact* derivatives!"
 - ? No black boxes allowed
 - ? Not always automatic/cheap/well-conditioned

Finite Differences (FD)

- "Nonintrusive"
 - ? Sensitive to stepsize choice/noise
 - ? Expense grows with n







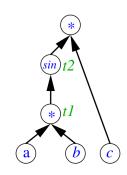
... then apply derivative-based method (that handles inexact derivatives)

Algorithmic Differentiation

Computational Graph

- $\diamond y = sin(a * b) * c$
- Forward and reverse modes
- AD tool provides code for your derivatives

You should write codes and formulate problems with AD in mind!



Many tools (see www.autodiff.org):

F/C Tapenade, Rapsodia

C/C++ ADOL-C, ADIC

Matlab ADiMat, INTLAB

Also done in AMPL and GAMS!

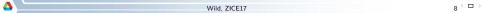


Caveats

You will pay a price for not having derivatives

We will focus primarily on:

- Minimization (sorry, economists, you'll need to stand on your head)
- Continuous domains
- Unconstrained problems
 - Simulation-based constraints bring additional challenges
- Underlying smooth behavior (whatever that means)
- Deterministic objective (or function of psuedorandom number generators)
- Algorithms with convergence guarantees to local minima
 - GAs and other heuristics often require far too many evaluations
- → Can address these in detail during lunch, dinner, TA/office hours



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Please: Interrupt, Ask Questions



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Global Optimization, $\min_{x \in \Omega} f(x)$

Careful:

- \diamond Global convergence: Convergence (to a local solution/stationary point) from anywhere in Ω
- ♦ Convergence to a global minimizer: Obtain x_* with $f(x_*) \le f(x) \, \forall x \in \Omega$

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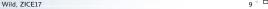
Anyone selling you global solutions when derivatives are unavailable:

either assumes more about your problem (e.g., convex f)

or expects you to wait forever

Törn and Žilinskas: An algorithm converges to the global minimum for any continuous f if and only if the sequence of points visited by the algorithm is dense in Ω .

or cannot be trusted



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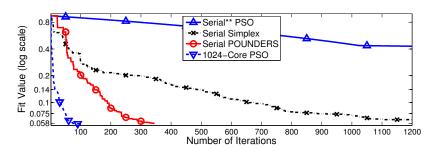
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Instead:

- Rapidly find good local solutions and/or be robust to poor solutions
- Consider multistart approaches and/or structure of multimodality

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(One Reason) Why We Won't Be Talking About Heuristics



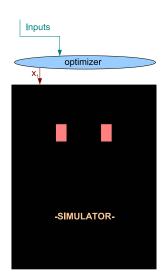
- Heuristics often "embarrassingly/naturally parallel";
 PS0= particle swarm method
 - Typically through stochastic sampling/evolution
 - 1024 function evaluations per iteration
- Simplex is Nelder-Mead; POUNDERS is model-based trust-region algorithm
 - one function evaluation per iteration
- → Is this an effective use of resources?
- → How many cores would have sufficed?



Black-box Optimization Problems

NB- My "black box" is different from Ken's "black box"

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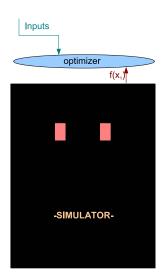
Only knowledge about f is obtained by sampling

- $\diamond f = S$ a black box (running some executable-only code or performing an experiment in the lab)
- \diamond Only give a single output (no derivatives $abla_x S(x)$)

Good solutions guaranteed in the limit, but:

- Usually have <u>computational budget</u> (due to scheduling, finances, deadlines)
- Limited number of evaluations

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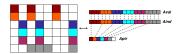
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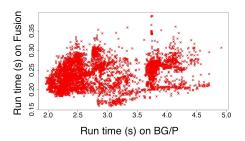
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A Black Box: Automating Empirical Performance Tuning

Given semantically equivalent codes C_1, C_2, \ldots , minimize run time subject to energy consumption





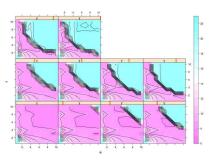
$\min \{ f(x) : (x_{\mathcal{C}}, x_{\mathcal{I}}, x_{\mathcal{B}}) \in \Omega_{\mathcal{C}} \times \Omega_{\mathcal{I}} \times \Omega_{\mathcal{B}} \}$

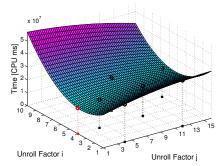
- x multidimensional parameterization (internal tolerances, unroll/tiling factors, compiler flags, compiler type, ...)
- Ω search domain (feasible transformation, no errors)
- f quantifiable performance objective (requires a run)

Optimization for Automatic Tuning of HPC Codes

Evaluation of f requires:

transforming source, compilation, (repeated?) execution, checking for correctness





Challenges:

- Evaluating $f(\Omega)$ prohibitively expensive (e.g., 10^{19} discrete decisions)
- f noisy

- Discrete x unrelaxable
- $\nabla_x f$ unavailable/nonexistent
- Many distinct/local solutions

Black-box Algorithms

Solve general problems $\min\{f(x): x \in \mathbb{R}^n\}$:

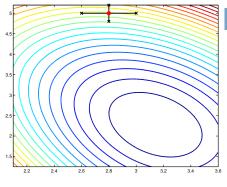
- Only require function values (no $\nabla f(x)$)
- \diamond Don't rely on finite-difference approximations to $\nabla f(x)$
- Seek greedy and rapid decrease of function value
- Have asymptotic convergence guarantees
- Assume parallel resources are used within function evaluation

Main styles of DFO algorithms

- Randomized methods (this afternoon!)
- ◇ Direct search methods (pattern search, Nelder-Mead, ...)
- Model-based methods (quadratics, radial basis functions, ...)

Choose a set of directions (pattern or mesh) \mathcal{D}^k

Ex.- \pm coordinate directions (2n directions)



Basic iteration $(k \ge 0)$

- \diamond Evaluate $f(x^k + \Delta_k d^j)$, $j = 1, \dots, |\mathcal{D}^k|$
- $\label{eq:force_eq} \begin{tabular}{l} $ & $\inf\left[f(x^k+\Delta_k d^j) < f(x^k)\right],$ \\ & $ & $ & move to \ x^{k+1} = x^k + \Delta_k d^j $ \\ \end{tabular}$

Otherwise shrink Δ_k

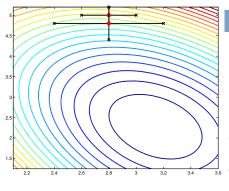
 \diamond Update \mathcal{D}^k

This is an indicator function, does not say anything about the magnitude of f values, just the ordering



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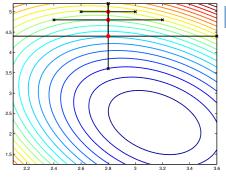
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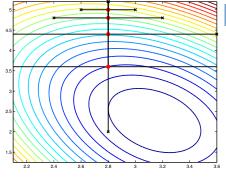
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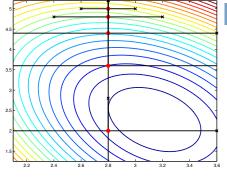
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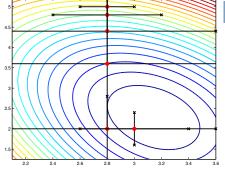
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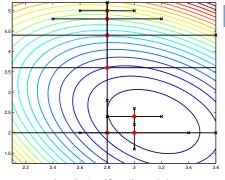
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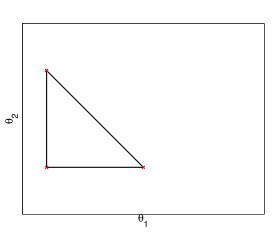
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- \rightarrow Lends itself well to doing concurrent function evaluations
- → See also mesh-adaptive direct search methods (e.g., NOMAD)
- ightarrow Can establish convergence for nonsmooth f

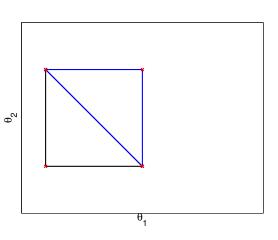
Basic iteration $(k \ge 0)$

- \diamond Evaluate f on the n+1 vertices of the simplex $x^k + \Delta_k \mathcal{S}^{(k)}$
- Reflect worst vertex about the best face
- $^{\diamond}$ Shrink, contract, or expand $\Delta_k \mathcal{S}^{(k)}$



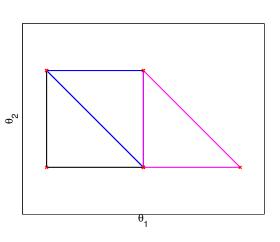
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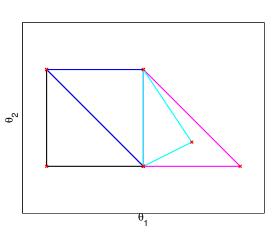
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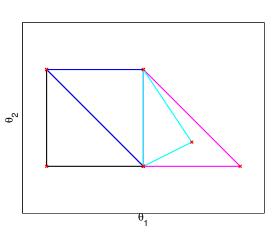
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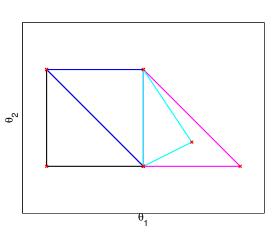
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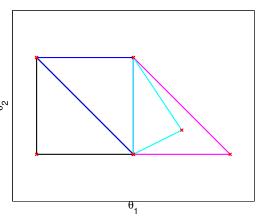
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The Nelder-Mead Method [1965]

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Only the order of the function values matter:

$$f(\hat{x}) = 1, \ \overline{f(\tilde{x})} = 1.0001$$
 is the same as $f(\hat{x}) = 1, \ f(\tilde{x}) = 10000$.

→ A very popular (due to "Numerical Recipes"), robust first choice

What Are We Missing?

These methods will (eventually) find a local solution

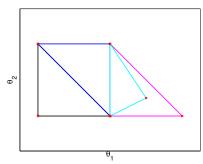
Overview: → [Kolda, Lewis, Torczon, SIREV 2003]



Each evaluation of f is expensive (valuable

N-M:

- 1. Only remembers the last n+1 evaluations
- 2. Neglects the <u>magnitudes</u> of the function values (order only)
- 3. Doesn't take into account the special (LS) problem structure

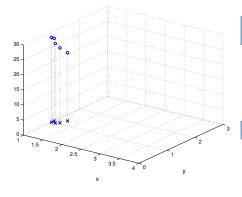


ightarrow This is the reason many direct search methods use a <u>search</u> phase on top of the usual poll phase



Making the Most of Little Information on f

- \diamond f is expensive \Rightarrow can afford to make better use of points
- Overhead of the optimization routine is negligible relative to the cost of evaluating the simulation.



Bank of data, $\{x^i, f(x^i)\}_{i=1}^k$:

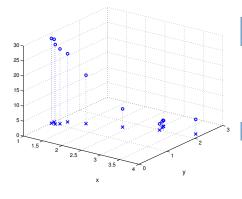
- Points (& function values) evaluated so far
- = Everything known about f

Goal

 Make use of growing Bank as optimization progresses

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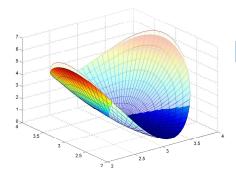
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Trust-Region Methods Use Models Instead of f

To reduce the number of expensive f evaluations

ightarrow Replace difficult optimization problem $\min f(x)$ with a much simpler one $\min \{m(x) : x \in \mathcal{B}\}$

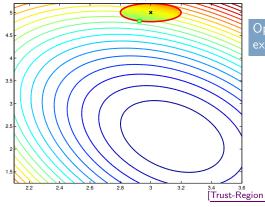


Classic NLP Technique:

- f Original function: computationally expensive, no derivatives
- *m* Surrogate model: computationally attractive, analytic derivatives





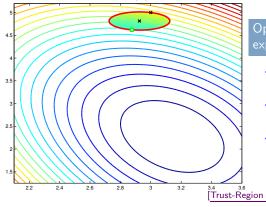


Optimize over m to avoid expense of f:

- ⋄ Trust m to approximate f within $\mathcal{B} = \{x \in \mathbb{R}^n : ||x x^k|| \le \Delta_k\},$
- ♦ Obtain next point from $\min \{m(x) : x \in \mathcal{B}\},\$
- \diamond Evaluate function and update (x^k, Δ_k) based on how good the model's prediction was.

[Trust-Region Methods; Conn, Gould, Toint; SIAM, 2000]





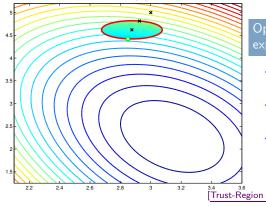
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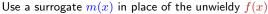


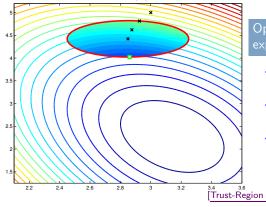


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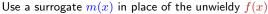


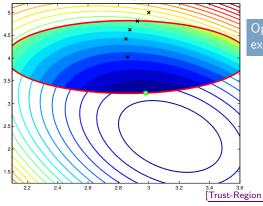


Optimize over m to avoid expense of f:

- ⋄ Trust m to approximate f within $\mathcal{B} = \{x \in \mathbb{R}^n : ||x x^k|| \le \Delta_k\},$
- ♦ Obtain next point from $\min \{m(x) : x \in \mathcal{B}\},\$
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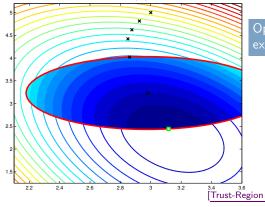


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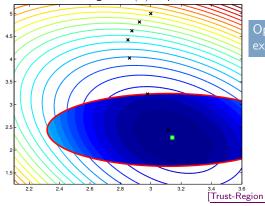


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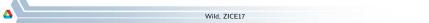
Trust-Region Methods; Conn, Gould, Toint; SIAM, 2000

Where Does the Model Come From?

When derivatives are available

Taylor-based model
$$m(x^k+s) = c + (g^k)^T s + \frac{1}{2} s^T H^k s$$

 $\diamond g^k = \nabla_x f(x^k)$
 $\diamond H^k \approx \nabla^2_{x,x} f(x^k)$



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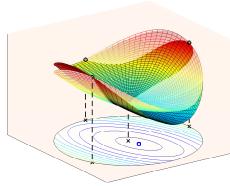
 $\diamond g^k = \nabla_x f(x^k)$
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Without derivatives

- Interpolation-based models
- Regression-based models
- Stochastic/randomized models



Interpolation-Based Quadratic Models



An interpolating quadratic in \mathbb{R}^2

$$m(x^k + s) = c + g^T s + \frac{1}{2} s^T H s$$
:

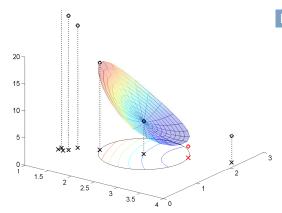
Get the model parameters $c,g,H=H^T$ by demanding interpolation:

$$m(x^k + y^i) = f(x^k + y^i)$$

for all $y^i \in \mathcal{Y} = \text{interpolation set}$

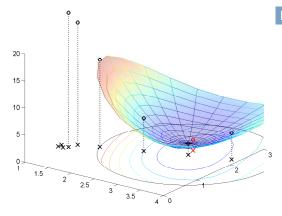
Main difficulty is \mathcal{Y} :

- Use prior function evaluations,
- m well-defined and approximates f locally.



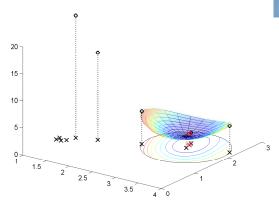
Iteration $\it k$

- $^{\diamond}$ Build a model m_k interpolating f on ${\mathcal Y}$
- \diamond Trust m_k within region \mathcal{B}_k
- \diamond Minimize m_k within \mathcal{B}_k to obtain next point for evaluation
- Do expensive evaluation
- \diamond Update m_k and \mathcal{B}_k based on how good model prediction was



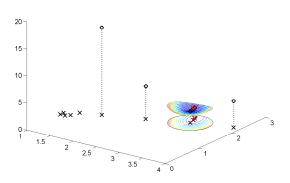
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Quick Diversion: Polynomial Bases

- \diamond Let ϕ denote a basis for some space of polynomials of n variables
 - Linear:

$$\phi(x) = [1, x_1, \cdots, x_n]$$



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• Full quadratics:

$$\phi(x) = \begin{bmatrix} 1, x_1, \cdots, x_n x_1^2, \cdots, x_n^2 x_1 x_2, \cdots, x_{n-1} x_n \end{bmatrix}$$

Quick Diversion: Polynomial Bases

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 \diamond Given a collection of $p = |\mathcal{Y}|$ points $\mathcal{Y} = \{y^1, \cdots, y^p\}$:

$$\Phi(\mathcal{Y}) = \begin{bmatrix} 1 & y_1^1 & \cdots & y_n^1 & (y_1^1)^2 & \cdots & (y_n^1)^2 & y_1^1 y_2^1 & \cdots & y_{n-1}^1 y_n^1 \\ \vdots & & & & & \vdots \\ 1 & y_1^p & \cdots & y_n^p & (y_1^p)^2 & \cdots & (y_n^p)^2 & y_1^p y_2^p & \cdots & y_{n-1}^p y_n^p \end{bmatrix}$$

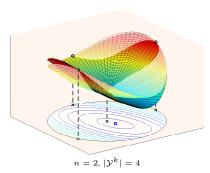
This is a matrix of size $p imes rac{(n+1)(n+2)}{2}$



Building Models Without Derivatives

Given $(\mathcal{Y}^k, f(\mathcal{Y}^k))$, "solve"

$$\Phi(\mathcal{Y}^k)z = \left[egin{array}{cc} \Phi_c & \Phi_g & \Phi_H \end{array}
ight] \left[egin{array}{c} z_c \ z_g \ z_H \end{array}
ight] = \underline{\mathsf{f}} = f \left(\mathcal{Y}^k
ight)$$



Full quadratics,
$$|\mathcal{Y}^k| = rac{(n+1)(n+2)}{2}$$

 \diamond Geometric conditions on points in \mathcal{Y}^k

Undetermined interpolation, $|\mathcal{Y}^k| < \frac{(n+1)(n+2)}{2}$

 \diamond Use (Powell) Hessian updates $\min_{g^k,H^k} \ \|H^k-H^{k-1}\|_F^2$ s.t. $q_k=\underline{\mathbf{f}}$ on \mathcal{Y}^k

Regression,
$$|\mathcal{Y}^k| > \frac{(n+1)(n+2)}{2}$$

 \bullet Solve $\min_z \|\Phi z - \underline{\mathbf{f}}\|$

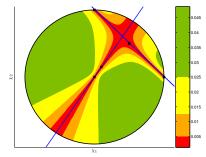
Multivariate (Scattered Data) Interpolation is a Different Kind of Animal

$$m(x^k + y^i) = f(x^k + y^i) \qquad \forall y^i \in \mathcal{Y}$$

- n=1 Given p distinct points, can find a unique degree p-1 polynomial m
- n > 1 Not true! (see Mairhuber-Curtis Theorem)

For quadratic models in \mathbb{R}^n :

- $\diamond \frac{(n+1)(n+2)}{2}$ coefficients
- \diamond Unique interpolant may not exist, even when $|\mathcal{Y}| = \frac{(n+1)(n+2)}{2}$
- Locations of the points in Y must satisfy additional geometric conditions (has nothing to do with f values)



 \rightarrow [Scattered Data Approximation; Wendland; Cambridge University Press, 2010]



Notions of Nonlinear Model Quality

"Taylor-like" Error Bounds

- 1. Assuming underlying f is sufficiently smooth
 - = derivatives of f exist but are unavailable
- 2. A model m_k is locally fully linear if:

For all
$$x \in \mathcal{B}_k = \{x \in \Omega : ||x - x^k|| \le \Delta_k\}$$

- $|m_k(x) f(x)| \le \kappa_1 \Delta_k^2$
- $\|\nabla m_k(x) \nabla f(x)\| \le \kappa_2 \Delta_k$

for constants κ_i independent of x and Δ_k .

→[Intro. to DFO; Conn, Scheinberg, Vicente; SIAM 2009]

Notions of Nonlinear Model Quality

"Taylor-like" Error Bounds

- 1. Assuming underlying f is sufficiently smooth
- 2. A model m_k is locally fully quadratic if:

For all
$$x \in \mathcal{B}_k = \{x \in \Omega : ||x - x^k|| \le \Delta_k\}$$

- $|m_k(x) f(x)| \le \kappa_1 \Delta_k^3$
- $\|\nabla m_k(x) \nabla f(x)\| \le \kappa_2 \Delta_k^2$
- $\|\nabla^2 m_k(x) \nabla^2 f(x)\| \le \kappa_3 \Delta_k$

for constants κ_i independent of x and Δ_k .

→[Intro. to DFO; Conn, Scheinberg, Vicente; SIAM 2009]



Ingredients for Convergence to Stationary Points

Assuming underlying f is sufficiently smooth and regular (e.g., has bounded level sets)

$$\lim_{k\to\infty} \nabla f(x^k) = 0$$
 provided:

- 1. Control \mathcal{B}_k based on model quality
- 2. (Occasional) approximation within \mathcal{B}_k Our quadratics satisfy
 - $|q_k(x) f(x)| \le \kappa_1(\gamma_f + ||H^k||)\Delta_k^2$, $x \in \mathcal{B}_k$
 - $\|g^k + H^k(x x^k) \nabla f(x)\| \le \kappa_2 (\gamma_f + \|H^k\|) \Delta_k, \quad x \in \mathcal{B}_k$
- 3. Sufficient decrease

At least model gradient should be good

Radial Basis Function (RBF) models also fit in this framework

→ [W. & Shoemaker, SIREV 2013]

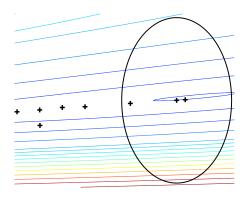
General introduction → [Intro. to DFO; Conn, Scheinberg, Vicente; SIAM 2009]



Greed. Alone. Can. Hurt.

Model-improvement may be needed when:

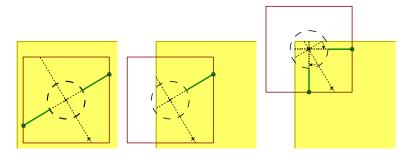
- Nearby points line up
- May not have enough points to ensure model quality in all directions



 \rightarrow May need n additional evaluations

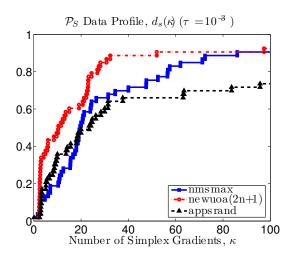
Constraints and Model Quality

Constraints complicate matters ...if one does not allow evaluation of infeasible points



→ May need directions normal to nearby constraints

Performance Comparisons on Test Functions



phase (nmsmax, appsrand) → [Moré & W., SIOPT 2009]

search methods without a search

 When evaluations are sequential,

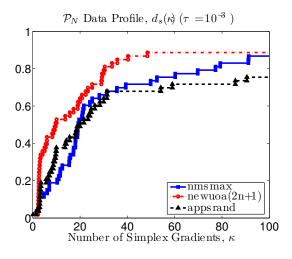
model-based methods

(NEWUOA) regularly outperform direct

Smooth problems



Performance Comparisons on Test Functions



without a search phase (nmsmax, appsrand) → [Moré & W., SIOPT 2009]

search methods

When evaluations are sequential.

model-based methods

(NEWUOA) regularly outperform direct

Noisy problems

Many Practical Details In Implementations

- \diamond Choice of interpolation points \mathcal{Y}^k
- \diamond Updating of trust region \mathcal{B}_k
- Improvement of models



Many Practical Details In Implementations

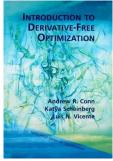
- \diamond Choice of interpolation points \mathcal{Y}^k
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```
BOBYQA [Powell], DFO [Scheinberg], POUNDer [W.] Initialization p = |\mathcal{Y}^k| structured evaluations Based on input, \approx 2n+1 Based on input, no more than n+1 Interpolation Set p = |\mathcal{Y}^k|, \forall k Bootstrap to |\mathcal{Y}^k| = \frac{(n+1)(n+2)}{2}, then fixed Varies in \{n+1,\cdots,\frac{(n+1)(n+2)}{2}\} based on available points Linear Algebra If p = \mathcal{O}(n), model formation costs only \mathcal{O}(n^2) Expensive Expensive
```

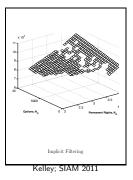


Growing Body of Tools and Resources for Local DFO

? What to use on problems with characteristics X, Y, and Z?



Conn, Scheinberg, Vicente: SIAM 2009



Many solvers

Sample considered by Rios & Sahinidis, 2010:

ASA,	BOBYQA,
CMA-ES,	DFO,
DAKOTA/*,	TOMLAB/*,
FMINSEARCH,	GLOBAL,
HOPSPACK,	IMFIL,
MCS,	NEWUOA,
NOMAD,	PSWARM,
STD-PSM.	SNOBFIT



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Toward Global Optimization

A quick sketch of a multistart methods and some practical details

- useful in derivative-based and derivative-free cases
- obtain a list of distinct minimizers (for post-processing, etc.)
- simple to get started

! simple to abuse/misuse ("I found all minimizers")



Why Multistart?

Multiple local minima are often of interest in practice:

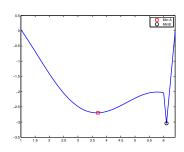
Design: Multiple objectives (or even constraints) might later be of

interest

Simulation Errors: Could have spurious local minima from anomalies in the

simulator

Uncertainty: Some minima are more sensitive to perturbations than others (gentle valleys versus steep cliffs)



d, ZICE17 36

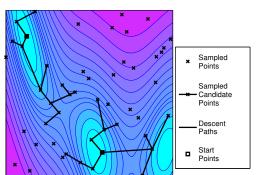
Global Optimization Multistart Methods

Two phase iterative method

Global Exploration: Sample N points in \mathcal{D} . \leftarrow Guarantees convergence Local Refinement: Start a local minimization algorithm \mathcal{A} from some promising subset of the sample points.

Want to find many (good) local minima while avoiding repeatedly finding the same local minima.

Where to start \mathcal{A} in kth iteration [Rinnooy Kan & Timmer, Math. Programming 1987]



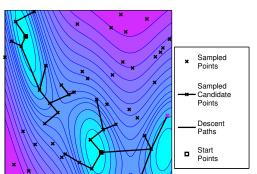
Ex.: It. 1 Exploration

Start \mathcal{A} at each sample point x^i provided:

- $^{\diamond}$ ${\cal A}$ has not been started from x^i , and
- $^{\diamond}$ no other sample point x^{j} with $f(x^{j}) < f(x^{i})$ is within a distance

$$r_k = \frac{1}{\sqrt{\pi}} \sqrt[n]{\operatorname{vol}(\mathcal{D})} \frac{5\Gamma\left(1 + \frac{n}{2}\right) \log(kN)}{kN}$$

Where to start \mathcal{A} in kth iteration [Rinnooy Kan & Timmer, Math. Programming 1987]



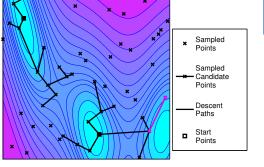
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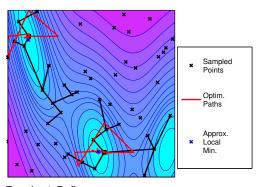
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Ex.: It. 1 Exploration
Thm [RK-T]- Will start finitely many local runs with probability 1.

Where to start \mathcal{A} in kth iteration [Rinnooy Kan & Timmer, Math. Programming 1987]



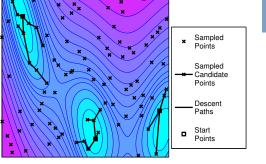
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$$r_k = \frac{1}{\sqrt{\pi}} \sqrt[n]{\operatorname{vol}(\mathcal{D})} \frac{5\Gamma\left(1 + \frac{n}{2}\right) \log(kN)}{kN},$$

Ex.: It. 1 Refinement Thm [RK-T]- Will start finitely many local runs with probability 1.

Where to start \mathcal{A} in kth iteration [Rinnooy Kan & Timmer, Math. Programming 1987]



Start \mathcal{A} at each sample point x^i provided:

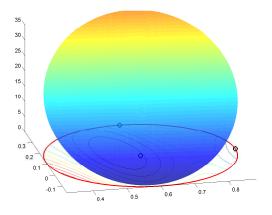
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$$r_k = \frac{1}{\sqrt{\pi}} \sqrt[n]{\operatorname{vol}(\mathcal{D})} \frac{5\Gamma\left(1 + \frac{n}{2}\right) \log(kN)}{kN}$$

Ex.: It. 2 Exploration
Thm [RK-T]- Will start finitely many local runs with probability 1.

Using External Points To Form Better Local Models

Initial model interpolating n+1=3 points



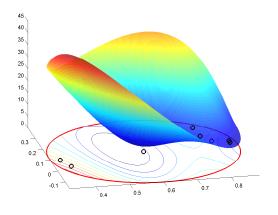
Take advantage of global history:

- from current minimization,
- from previous minimizations, and
- from the global sampling.

More information means rapider progress.

Using External Points To Form Better Local Models

Initial model interpolating 8 sample points



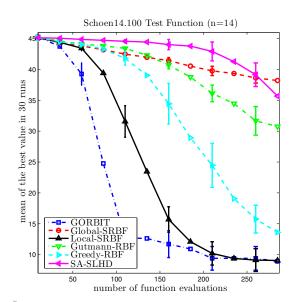
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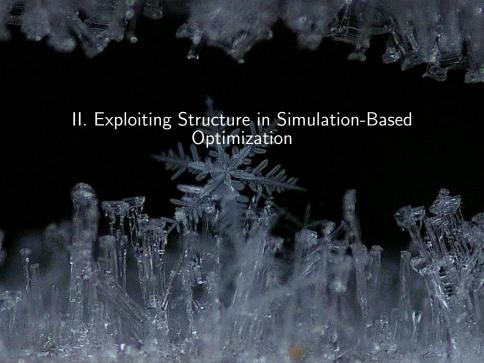
More information means rapider progress.

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Performance Comparisons on Test Functions



- GORBIT is multistart with RBF model-based method
- SA-SLHD is a heuristic (simulated annealing with a symmetric Latin hypercube design as initialization)
- \rightarrow [W., Cornell University, 2009]



Beyond the Black Box

$$\min f(x) = F[S(x)]$$

So far, $f={\cal S}$

△ Wild, ZICE17

Beyond the Black Box

$$\min f(x) = F[S(x)]$$

So far, f = S

Your problems are not black-box problems

Beyond the Black Box

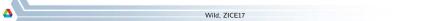
$$\min f(x) = F[S(x)]$$

So far, f = S

Your problems are not black-box problems

You formulated the problem

 \Rightarrow You know more than nothing



Structure in Simulation-Based Optimization, $\min f(x) = F[S(x)]$

f is often not a black box S

Nonlinear least squares

$$f(x) = \frac{1}{2} \sum_{i} (S_i(x) - d_i)^2$$

Not all variables enter simulation

$$f(x) = g(x_I, x_J) + h(S(x_J))$$

Only some constraints depend on simulation

$$\min\{f(x): c_1(x) = 0, c_{\mathbf{S}}(x) = 0\}$$

Slack variables

$$\Omega_S = \{(x_I, x_J) : S(x_J) + x_I = 0, x_I \ge 0\}$$

Model-based methods can be a great way to exploit this structure



Ex. 1- Least Squares
$$f(x) = \frac{1}{2} \sum_{i} F_i(x)^2$$

Obtain a vector of output $F_1(x), \ldots, F_p(x)$

 \diamond Model each F_i

$$F_i(x) \approx q_k^{(i)}(x) = F_i(x^k) + (x - x^k)^{\top} g^{(i,k)} + \frac{1}{2} (x - x^k)^{\top} H^{(i,k)}(x - x^k)$$

Approximate:

$$\nabla f(x) = \sum_{i} \nabla \mathbf{F_{i}}(\mathbf{x}) F_{i}(x) \rightarrow \sum_{i} \nabla q_{k}^{(i)}(x) F_{i}(x)$$

$$\nabla^{2} f(x) = \sum_{i} \nabla \mathbf{F_{i}}(\mathbf{x}) \nabla \mathbf{F_{i}}(\mathbf{x})^{T} + \sum_{i} F_{i}(x) \nabla^{2} \mathbf{F_{i}}(\mathbf{x})$$

$$\rightarrow \sum_{i} \nabla q_{k}^{(i)}(x) \nabla q_{k}^{(i)}(x)^{T} + \sum_{i} F_{i}(x) \nabla^{2} q_{k}^{(i)}(x)$$

Model f via Gauss-Newton or similar

 $\rightarrow [\mathsf{DFLS}; \, \mathsf{Zhang}, \, \mathsf{Conn}, \, \mathsf{Scheinberg}] \, \, \mathbf{regularized} \, \, \mathbf{Hessians} \\ \rightarrow [\mathsf{POUNDERS}; \, \mathsf{W.}, \, \mathsf{Mor\'e}] \, \, \mathbf{uses} \, \, \mathbf{full} \, \, \mathbf{Newton}$

Ex. 1- Consequences for $f(x) = \frac{1}{2} \sum_{i} F_i(x)^2$

- \diamond Save linear algebra using interpolation set \mathcal{Y}^k common to all models
 - Single system solve, multiple right hand sides

$$\Phi(\mathcal{Y}^k) \begin{bmatrix} z^{(1)} & \cdots & z^{(p)} \end{bmatrix} = \begin{bmatrix} \underline{\mathsf{F}}_1 & \cdots & \underline{\mathsf{F}}_p \end{bmatrix}$$

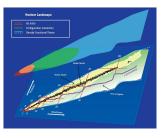
- fully linear $q^{(0)} \Rightarrow$ all $q^{(i)}$ fully linear
- \diamond (nearly) exact gradients for F_i (nearly) linear
- No longer interpolate function at data points

$$\begin{split} m(x^k+\delta) &= & f(x^k) + \delta^T \sum_i g^{(i,k)} F_i(x^k) \\ &+ \frac{1}{2} \delta^T \sum_i \left(g^{(i,k)} (g^{(i,k)})^T + F_i(x^k) H^{(i,k)} \right) \delta \\ &+ & \text{missing h.o. terms} \end{split}$$

Ex. 1- Calibrating Energy Density Functionals

$$\min_{x} \left\{ f(x) = \sum_{i=1}^{p} \left(\frac{s_i(x) - d_i}{w_i} \right)^2 \right\}$$

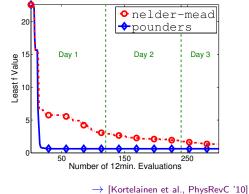
- $s_i(x)$ Simulated (DFT) nucleus property
 - d_i Nucleus i experimental data
 - w_i Weight for data type i
 - p Parallel simulations
 - \diamond Engineering starting point x_0
 - Bound constraints scaled to unit cube
 - ◆ 12 CPU minutes per evaluation

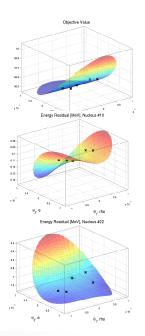


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Ex. 1- POUNDERS in Practice

- \diamond Models for each residual $\frac{d_i s_{t,i}(x)}{\sigma_i}$. $\left\{m^{i}(x)\right\}_{i=1}^{90}$
- Further reduces # of expensive evaluations





Ex. 2- Some Known Partials

$$x=(x_I,x_J)$$
; have $\frac{\partial f}{\partial x_I}$ but not $\frac{\partial f}{\partial x_J}$

"Solve"

$$\Phi z = \underline{\mathsf{f}}$$

with known $z_{g,I}, z_{H,I}$

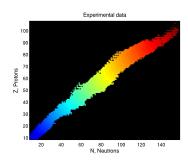
$$\left[\begin{array}{cc} \Phi_c & \Phi_{g,J} & \Phi_{H,J} \end{array} \right] \left[\begin{array}{c} z_c \\ z_{g,J} \\ z_{H,J} \end{array} \right] = \underline{\mathbf{f}} - \Phi_{g,I} z_{g,I} - \Phi_{H,I} z_{H,I}$$

- \diamond Effectively lowers dimension to |J|=n-|I| for
 - approximation
 - model-improving evaluations
 - linear algebra
- Still have interpolation where required



Ex. 2- Multi-level Optimization Structure

$$\min_{x} \left\{ f(x) = \sum_{i=1}^{p} \left(s_i(x) - d_i \right)^2 \right\}$$



[Bertolli, Papenbrock, W., PhysRevC '12]

$s_i(x)$ solution to lower level problem

$$s_{i}(x) = \tilde{g}_{i}(x) + \min_{y} \{h_{i}(x_{J}; y) : y \in \mathcal{D}_{i}\}$$

$$= \tilde{g}_{i}(x) + h_{i}(x_{J}; y_{i,*}[x_{J}])$$

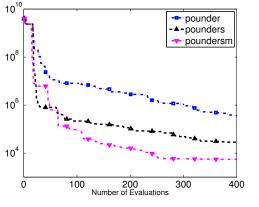
$$\nabla_{\mathbf{x}_{J}} \mathbf{s}_{i}(\mathbf{x}) \approx \nabla_{x_{J}} \tilde{g}_{i}(x) + \nabla_{x_{J}} q^{(i)}(x_{J})$$

For $x = (x_I, x_J)$

- $\diamond \nabla_{x_I} s_i(x_I, x_J)$ available
- \diamond $s_i(x)$ continuous and smooth in x_I
- $\tilde{g}_i(x)$ cheap to compute!
- No noise/errors introduced in $\tilde{g}_i(x)$

♦ Wild, ZICE17 49 ⁴ F

Ex. 2- Numerical Results With Some Partials



s exploit least squares m use ∇_{x_I} derivatives

- n = 16, |I| = 3
- 5-10 seconds/evaluation on 8 cores

"Same algorithmic framework", performance advantages from exploiting structure

Ex. 2- Convergence With Known Partials $abla_{x_I}$

$$\lim_{k\to\infty} \nabla f(x^k) = 0$$
 as before

Approximation bounds

$$|q_k(x) - f(x)| \le \kappa_1(\gamma_f + ||H^k||)\Delta_k^2, \quad x \in \mathcal{B}_k$$

with most constants now a function of $n-\vert I\vert.$

Guaranteed

- strict descent
- optimality

in some directions.

Ex. 3- General Constraints

$$\min\{f(x): c_1(x) = 0, c_S(x) = 0\}$$

Approximate Lagrangian:

$$\nabla L = \nabla f + \lambda_1^T \nabla c_1 + \lambda_2^T \nabla \mathbf{c_S}$$

$$\rightarrow \nabla f + \lambda_1^T \nabla c_1 + \lambda_2^T \nabla m$$

- Use favorite method: filters, augmented Lagrangian, . . .
- Slack variables
 - Do not increase effective dimension
 - Subproblems can treat separately
 - Know derivatives

→[Diniz-Ehrhardt, Martínez, Pedroso, C&A Math. 2011]: modified AL methods



Ex. 3- What Constraint Derivatives Buy You

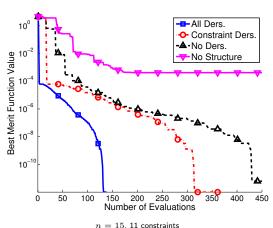
Augmented Lagrangian methods, $L_A(x,\lambda;\mu) = f(x) - \lambda^T c(x) + \frac{1}{\mu} \|c(x)\|^2$

$\min_{x} \left\{ f(x) : c(x) = 0 \right\}$

Four choices:

- 1. Penalize constraints
- 2. Treat c and f both as (separate) black boxes
- 3. Work with f and $\nabla_x c$
- 4. Have both $\nabla_x f$ and $\nabla_x c$

 \rightarrow With Slava Kungurtsev (UCSD)



n=15, 11 constraints

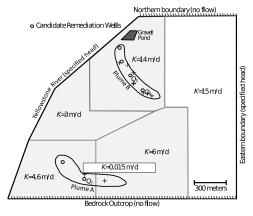
Ex. 4- Remediation of Chlorinated Solvents

Determine extraction rates for 6 installed wells

Structure

- Minimize operating cost (linear)
- Plume flux constraints (expensive simulation)

$$c_P(x) = \sum_{i=1}^p |S_{P,i}(x)|$$



Lockwood Solvent Ground Water Plume Site (LSGPS)

Wild, ZICE17 54 To The Control of th

Ex. 4- Exploiting Nonsmoothness

What if derivatives of f(x) do not always exist?

$$f(x) = g(x, S(x))$$

- ♦ g Nonsmooth, known
- ♦ S Smooth, black-box
 - Examples:
 - $f(x) = \sum_{i=1}^{p} |F_i[S(x)]|$ $f(x) = \max(S_1(x), S_2(x))$

→ With Aswin Kannan



Ex. 4- Targeting Nonsmoothness in $f(x) = \sum_{i=1}^{p} |F_i(x)|$

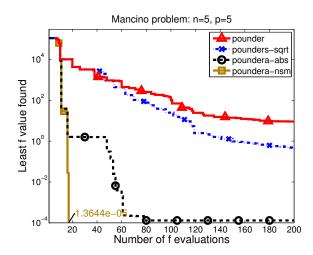
Model-based Approaches:

pounder Ignore structure, model
$$f$$
 as usual pounders-sqrt $f = \sum\limits_{i=1}^p \sqrt{|F_i|}^2$, model $\sqrt{|F_i|}$ by Q_i subproblem $\min \sum_{i=1}^p \tilde{Q}_i(x)^2$ poundera-abs $f = \sum\limits_{i=1}^p |F_i|$, model $|F_i|$ by Q_i subproblem $\min \sum_{i=1}^p Q_i(x)$ poundera-nsm $f = \sum\limits_{i=1}^p |F_i|$, model F_i by Q_i subproblem $\min \sum_{i=1}^p |Q_i(x)|$

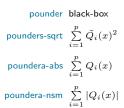
Ex. 4- Preliminary Results, $\min \sum_{i=1}^{p} |F_i(x)|$

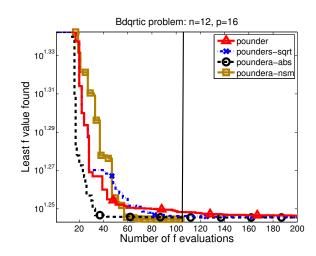
 $\begin{array}{ll} \text{pounder} & \text{black-box} \\ \\ \text{pounders-sqrt} & \sum\limits_{i=1}^p \tilde{Q_i}(x)^2 \\ \\ \text{poundera-abs} & \sum\limits_{i=1}^p Q_i(x) \end{array}$

poundera-nsm $\sum\limits_{i=1}^{p}|Q_{i}(x)|$



Ex. 4- Preliminary Results, $\min \sum_{i=1}^{p} |F_i(x)|$







So You Want To Solve A Hard Optimization Problem?

Mathematically unwrap problems to expose the deepest black boxes!

- It is easy to get started with derivative-free methods
- You should strive to obtain derivatives & apply methods from Todd's talk
- Model-based methods can make use of expensive function values
- Structure is everywhere, even in "black-box" / legacy code-driven optimization problems
- By exploiting structure, optimization can solve grand-challenge problems in (insert your field here):
 - Model residuals $\{r_i(x)\}_i$, not ||r(x)||
 - Model constraints $\{c_i(x)\}_i$, not a penalty P(c(x))
 - Explicitly handle nonsmoothness (and noise)