# **Numerical Optimization**

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# Part I

# Principles of Numerical Computing



# Types of Computing

- Symbolic computing (e.g. mathematica)
  - ▶ Represent and manipulate mathematical expressions
  - ► Apply rules of algebra to simplify expressions
  - ▶ Irrational numbers are represented exactly

# Types of Computing

- ► Symbolic computing (e.g. mathematica)
  - Represent and manipulate mathematical expressions
  - Apply rules of algebra to simplify expressions
  - Irrational numbers are represented exactly
- Numerical computing (e.g. matlab)
  - Finite precision arithmetic
  - Apply basic algebra operations
  - Irrational numbers are approximated
  - Numerical analysis of algorithms required

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  - ▶ Shift for common exponent  $(9.123 \times 10^{-4} + 0.9123 \times 10^{-4})$
  - Add  $(10.0353 \times 10^{-4})$
  - ▶ Round result  $(10.04 \times 10^{-4})$
  - Normalize  $(1.004 \times 10^{-3})$
  - ► Accumulate roundoff error  $(-4.700 \times 10^{-7})$ 
    - ▶ Importance of  $\epsilon$ :  $1.0 + \epsilon = 1.0$
    - Adding many numbers:  $x^{k+1} = x^k + \epsilon$ ,  $x^0 = 1.0$

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- Subtraction results in loss of precision
- Multiplication and division similar

# Dealing with Roundoff Errors

- ▶ Structure computation to minimize roundoff error
  - ▶ Adding positive numbers from least magnitude to largest
  - ► Requires keeping sorted list of numbers
  - Minimum is very difficult in general
  - ► Reduce the error when possible
- ▶ Perform computations tolerant to roundoff error
  - Analyze algorithms
  - Prove statements about the error
  - ► Field of numerical analysis
- Numerical linear algebra widely studied

# **Notation**

- ▶ Scalar:  $a \in \Re$
- ▶ Vector:  $b \in \Re^n$ ,  $b_i$  is the *i*th element
- ▶ Matrix:  $A \in \Re^{m \times n}$ 
  - $\triangleright$   $A_{i,.}$  is the *i*th row
  - $\triangleright$   $A_{\cdot,j}$  is the *j*th column
  - ▶  $A_{i,j}$  is the element in row i and column j

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  - Symmetric:  $A = A^T$
  - Skew symmetric:  $A = -A^T$
- ▶ Storage of matrices
  - ▶ Dense matrices store *mn* entries
  - Sparse matrices store
    - ▶ Nonzero structure: (i,j) such that  $A_{i,j} \neq 0$
    - Nonzero entries:  $A_{i,j} \neq 0$
  - Sparse linear algebra used for performance

# Linear Algebra

► Matrix-vector products

$$y = Ax + b$$

► Matrix-matrix products

$$C = AB$$

► Linear systems of equations

$$Ax = b$$

► Condition number estimation



# Direct Methods for Medium Problems

$$Ax = b$$

- ▶ Use backslash  $x = A \setminus b$  rather than inverse x = inv(A)b
- ► LL' Cholesky factorization (sparse, symmetric, positive definite)
- ► LDL' factorization (sparse, symmetric, indefinite)
- ► LU factorization (sparse, generic, full rank)
- QR factorization (dense, generic)
- ► USV' singular value decomposition (dense, generic)

# Iterative Methods for Large Problems

$$Ax = b$$

- ► Conjugate gradient method (sparse, symmetric, positive definite)
- Generalized minimum residual (sparse, generic)
- Quasi-minimum residual (sparse, generic)
- ► Transpose-free quasi-minimum residual (sparse, generic)



# **Functions**

► Objective function

$$f: \Re^n \to \Re$$

- ▶ Linear:  $f(x) = a^T x$  for vector  $a \in \Re^n$
- Quadratic:  $f(x) = \frac{1}{2}x^TQx + a^Tx$  for symmetric matrix  $Q \in \Re^{n \times n}$
- ► Nonlinear:
  - Continuous versus discontinuous
  - ► Smooth versus nonsmooth

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  - ▶ Smooth versus nonsmooth
- Constraints

$$c: \Re^n \to \Re^m$$

- ▶ Linear: c(x) = Ax + b for matrix  $A \in \Re^{m \times n}$  and vector  $b \in \Re^m$
- Quadratic:  $c_i(x) = \frac{1}{2}x^TQ_ix + a_i^Tx + b_i$  for symmetric  $Q_i \in \Re^{n \times n}$
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- Nonlinear
  - ► Continuous versus discontinuous
  - Smooth versus nonsmooth
- Code functions using intrinsic operations
  - ► Basic algebra operations
  - Transcendental functions
  - Reduce roundoff errors

## **Derivatives**

▶ Gradient of f (vector in  $\Re^n$ )

$$[\nabla f(x)]_j = \frac{\partial f}{\partial x_j}$$

▶ Hessian of f (symmetric matrix in  $\Re^{n \times n}$ )

$$[\nabla^2 f(x)]_{j,k} = \frac{\partial^2 f}{\partial x_j \partial x_k}$$

▶ Jacobian of c (matrix in  $\Re^{m \times n}$ )

$$[\nabla c(x)]_{i,j} = \frac{\partial c_i}{\partial x_j}$$

▶ Second derivatives of c (tensor in  $\Re^{m \times n \times n}$ ,  $\nabla^2 c_i(x)$  symmetric)

$$[\nabla^2 c(x)]_{i,j,k} = \frac{\partial^2 c_i}{\partial x_j \partial x_k}$$



Numerical derivatives via finite differences

$$\frac{\partial f(x)}{\partial x_i} \approx \frac{f(x + he_i) - f(x)}{h}$$

- ▶ Can requires n+1 function evaluations
- Quality depends on function and stepsize
- Subtraction results in loss of precision
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  - ▶ Use calculus rules for the intrinsic functions
  - Apply the chain rule to assemble partials
    - ► Forward mode easy to implement, but can be slow
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  - Precision approximately equal to function precision
  - Provable bounds on cost that does not grow in n

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- Assume derivatives available for now

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  - Approximation methods for nonsmooth functions

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- Derivative-free methods are a gateway or last resort

Locally Convergent Newton Method

- Most good algorithms are variants of Newton's method
- ▶ Compute solutions F(x) = 0 where  $F: \Re^n \to \Re^n$ 
  - Form Taylor series approximation around  $x^k$

$$F(x) \approx \nabla F(x^k)(x - x^k) + F(x^k)$$

Solve for x and iterate

$$x^{k+1} = x^k - \nabla F(x^k) \backslash F(x^k)$$



Locally Convergent Newton Method

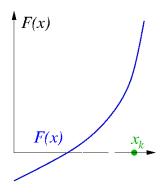
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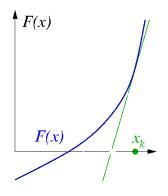
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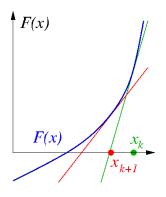
- Several possible outcomes
  - ► Convergence:  $\lim_{k\to\infty} x^k = x^*$
  - Iterate divergence:  $\lim_{k\to\infty} ||x^k|| \to \infty$
  - Sequence cycles:
    - Multiple convergent subsequences (limit points)
    - Limit points are not solutions













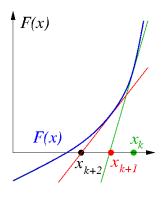




Illustration of Iterate Divergence

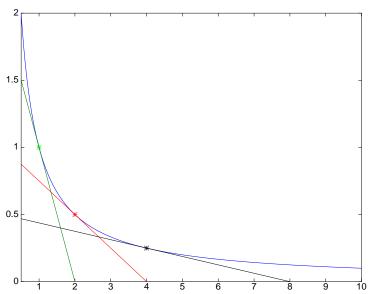
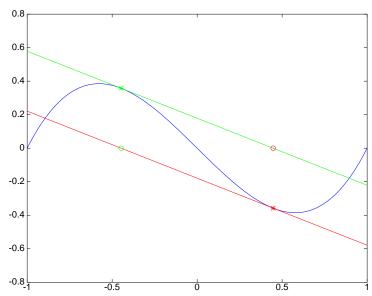


Illustration of Cycling



#### Globally Convergent Newton Method

Use Newton method to compute a step

$$s^k = -\nabla F(x^k) \backslash F(x^k)$$

▶ Check direction for descent using merit function  $\Psi(x) = ||F(x)||_2^2$ 

$$\nabla \Psi(x^k)' s^k < 0$$

▶ Minimize merit function along step

$$t_k \in \arg\min_{t \in (0,1]} \Psi(x^k + ts^k)$$

▶ Update iterate

$$x^{k+1} = x^k + t_k s^k$$

and repeat until convergence

- ► Several possible outcomes
  - ▶ Good convergence:  $\lim_{k\to\infty} x^k \to x^*$  with  $F(x^*) = 0$
  - ▶ Bad convergence:  $\lim_{k\to\infty} x^k \to x^*$  with  $\Psi(x^*) > 0$  and  $\nabla \Psi(x^*) = 0$
  - ▶ Iterate divergence:  $\lim_{k\to\infty} ||x^k|| \to \infty$

#### Implementation Details

- Numerical linear algebra
  - Direct methods compute a (sparse) factorization

$$\nabla F(x^k) = LU$$

where L is lower triangular and U is upper triangular

- Iterative methods compute an approximation
  - Krylov subspace method such as generalize minimum residual
  - ▶ Large part of the computation is matrix-vector products
  - Accelerate convergence using a preconditioner
- ▶ Recourse when  $\nabla F(x^k)$  is not invertible



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- ▶ Recourse when  $\nabla F(x^k)$  is not invertible
- Minimizing merit function along the step
  - Golden section and Fibonacci search
  - ▶ Armijo backtracking line search: find smallest integer *i* with

$$\Psi(x^k + \beta^i s^k) \le \Psi(x^k) + \sigma \beta^i \nabla \Psi(x^k)' s^k$$

Moré-Thuente: cubic interpolation and refinement

Convergence Results

- ▶ Global convergence under suitable conditions
- ▶ Local fast convergence under suitable conditions
  - ▶ If  $x^k$  is near a solution, method converges to a solution  $x^*$
  - ► The distance to the solution decreases quickly; ideally,

$$||x^{k+1} - x^*|| \le c||x^k - x^*||^2$$

#### Newton Method with Proximal Perturbation

- ▶ Recourse when  $\nabla F(x^k)$  is not invertible or bad direction
- ► Solve linear system of equations

$$(\nabla F(x^k) + \lambda_k I)s^k = -F(x^k)$$

▶ Check direction for descent using merit function  $\Psi(x) = ||F(x)||_2^2$ 

$$\nabla \Psi(x^k)' s^k < -p_1 ||s^k||^{p_2}$$

otherwise use steepest descent direction  $s^k = -\nabla \Psi(x^k)$ 

▶ Minimize merit function along step

$$t_k \in \arg\min_{t \in (0,1]} \Psi(x^k + ts^k)$$

- Update perturbation
- ► Update iterate

$$x^{k+1} = x^k + t_{\nu} s^k$$

and repeat until convergence

### Nonsquare Nonlinear Systems of Equations

▶ Given  $F: \Re^n \to \Re^m$ , compute x such that

$$F(x) = 0$$

- ▶ System is underdetermined if m < n
  - More variables than constraints
  - Solution typically not unique
  - Need to select one solution

$$\min_{x} ||x||_2$$
 subject to  $F(x) = 0$ 

- ▶ System is overdetermined if m > n
  - More constraints than variables
  - Solution typically does not exist
  - ▶ Need to select approximate solution

$$\min_{x} \|F(x)\|_{2}$$

- ▶ System is square if m = n
  - Jacobian has full rank then solution is unique
  - If Jacobian is rank deficient then
    - Underdetermined when compatible
    - Overdetermined when incompatible

### Part II

# Foundations of Optimization



## Foundations of Optimization

Generic nonlinear optimization problem

$$\min_{\substack{x \\ \text{subject to}}} f(x)$$

- $x \in \Re^n$  are the optimization variables
- $f: \Re^n \to \Re$  represents the objective function
- ▶  $c: \Re^n \to \Re^m$  represents the constraints
- Least squares problems
- ▶ Further classification
  - Types of functions
  - Types of constraints
- Model transformations

## Ordinary Least Squares

Minimize two norm:  $||y||_2^2 = \sum_i y_i^2$  $\min_x ||Ax - b||_2^2$ 

▶ Data representation: 
$$A \in \Re^{m \times n}$$

- ▶ Vector of observations:  $b \in \Re^n$ 
  - m is the number of measurements
  - n is the number of variables
  - Both can be very large

### Ordinary Least Squares

Basic Algorithm

► Assuming A has full column rank

$$x = (A^T A) \setminus (A^T b)$$

where  $A^T A \in \Re^{n \times n}$  is symmetric

- ▶ Standard errors use  $diag((A^TA)^{-1})$
- ► Two cases
  - Number of variables is small
  - Number of variables is large

## **Ordinary Least Squares**

Small Number of Variables

Form the matrix  $B = A^T A$ 

$$B_{i,j} = \sum_{k} A_{k,i} A_{k,j}$$

- Requires O(mn²) operations
- Control roundoff errors for large m
- Dense linear algebra
- ▶ Solve  $x = B \setminus (A^T b)$ 
  - Requires  $O(n^3 + mn)$  operations
  - Dense linear algebra
- ▶ Compute standard errors using diag( $B^{-1}$ )
  - ▶ Find diagonal entries of  $B \setminus I$
  - ► Compute *inv*(*B*)
  - Dense linear algebra



### Generalized Least Squares

#### Basic Algorithm

▶ Minimize matrix norm:  $||y||_M^2 = y^T M y$ , M positive definite

$$\min_{x} \|Ax - b\|_{M}^{2}$$

- ▶ Data representation:  $A \in \Re^{m \times n}$
- ▶ Matrix representation:  $M \in \Re^{m \times m}$
- ▶ Vector of observations:  $b \in \Re^n$ 
  - m is the number of measurements
  - n is the number of variables
- Assuming A has full column rank

$$x = (A^T MA) \backslash A^T Mb$$

where  $A^TMA \in \Re^{n \times n}$  is symmetric

- ▶ Typical weight is inverse correlation matrix  $\Sigma^{-1}$ 
  - ► Large dense matrix
  - Must have a special representation
- ▶ Forming  $A^T(\Sigma \backslash A)$  is a bad idea
- ▶ Utilize matrix-free iterative methods

#### Convex Functions

▶  $f: \Re^n \to \Re$  is convex if

$$\forall x, y : f(tx + (1-t)y) \le tf(x) + (1-t)f(y)$$

for all  $t \in [0,1]$ 

▶ Differentiable  $f: \Re^n \to \Re$  is convex if

$$\forall x, y : f(y) \ge f(x) + \nabla f(x)^T (y - x)$$

▶ Twice differentiable  $f: \Re^n \to \Re$  is convex if

$$\forall x, y : y^T \nabla^2 f(x) y \ge 0$$

▶ A matrix  $Q \in \Re^{n \times n}$  is positive semidefinite if

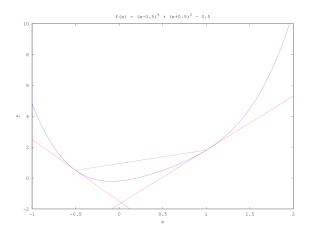
$$\forall y : y^T Q y \geq 0$$

- Many equivalent characterizations
- $\triangleright$  For symmetric Q, eigenvalues are nonnegative
- ▶ A set  $X \subseteq \Re^n$  is convex if

$$\forall x \in X, y \in X : tx + (1-t)y \in X$$

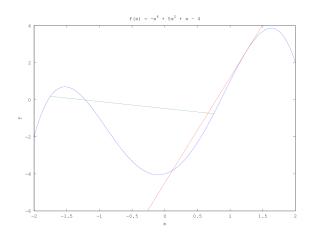
for all  $t \in [0,1]$ 

## Illustration of Convexity



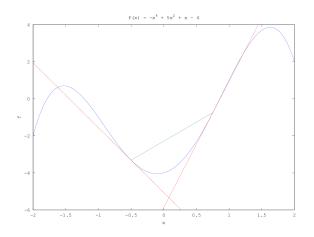


## Illustration of Nonconvexity





## Illustration of Local Convexity





### Convexity

Proving and Disproving

► Generally a hard problem

### Convexity

#### Proving and Disproving

- Generally a hard problem
- Proving convexity analyzes functional form
  - Representation as an abstract syntax tree
  - Apply convexity rules
  - ▶ E.g. If g(x) and h(x) are convex, then g(x) + h(x) is convex
  - Cannot disprove convexity

### Convexity

#### Proving and Disproving

- Generally a hard problem
- Proving convexity analyzes functional form
  - Representation as an abstract syntax tree
  - Apply convexity rules
  - ▶ E.g. If g(x) and h(x) are convex, then g(x) + h(x) is convex
  - Cannot disprove convexity
- Disproving convexity uses the definitions
  - Choose finite points  $x^k$  and  $y^k$
  - ▶ Choose finite steps  $\{t_\ell\}$   $\subset$  (0,1)
  - ▶ Not convex if there exists k and  $\ell$  such that

$$f(t_{\ell}x^{k} + (1-t_{\ell})y^{k}) > t_{\ell}f(x^{k}) + (1-t_{\ell})f(y^{k})$$

Cannot prove convexity



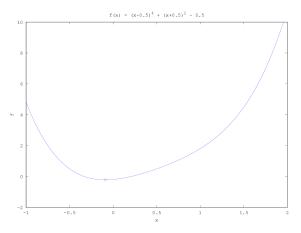
#### Unconstrained Optimization

$$\min_{x} f(x)$$

- ▶ Linear:  $f(x) = c^T x$  (unbounded solution)
- Quadratic:  $f(x) = \frac{1}{2}x^TQx + c^Tx$ 
  - Convex if Q is positive semidefinite
  - ightharpoonup Concave if -Q is positive semidefinite (unbounded solution)
  - Otherwise it is neither convex nor concave
  - Conjugate gradient iterative methods
- Nonlinear function
  - Convex under suitable conditions
  - ▶ Concave if -f(x) is convex (unbounded solution)
  - Otherwise it is neither convex nor concave
- Implications of convexity
  - Convex implies local solution is a global solution
  - Strictly convex implies at most one solution
  - Strongly convex implies one solution
  - Convex implies convex solution set
- Provable global solutions for nonconvex problems is difficult
- Algorithms discussed in next session

## Convex Unconstrained Optimization

#### Critical Points



• Global minimizer:  $\nabla f(x) = 4(x - 0.5)^3 + 2(x + 0.5) = 0$ 



#### Nonconvex Unconstrained Optimization

$$\min_{x} f(x)$$

Global solution x\* if

$$f(x^*) \leq f(x) \quad \forall \ x$$

- Global solutions need not be unique
- All global solutions have the same objective value
- ▶ Local solution *x*\* if

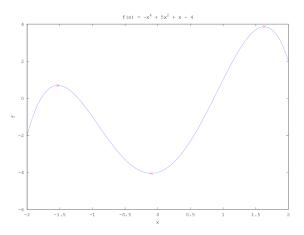
$$f(x^*) \le f(x) \quad \forall \ x \in \mathcal{N}(x^*)$$

where  $\mathcal{N}(x^*)$  is a neighborhood of  $x^*$ 

- ► Local solutions need not be unique
- Local solutions need not have the same objective value

## Nonconvex Unconstrained Optimization

#### Critical Points



- ► Stationary:  $\nabla f(x) = -4x^3 + 10x + 1 = 0$
- ▶ Local maximizer:  $\nabla^2 f(x) = -12x^2 + 10 < 0$
- ▶ Local minimizer:  $\nabla^2 f(x) = -12x^2 + 10 > 0$

#### **Bound Constrained Optimization**

- ▶ Linear:  $f(x) = c^T x$
- Quadratic:  $f(x) = \frac{1}{2}x^TQx + c^Tx$ 
  - Convex if Q is positive semidefinite (special methods)
  - ightharpoonup Concave if -Q is positive semidefinite
  - Otherwise it is neither convex nor concave
  - Conjugate gradient iterative methods
- Nonlinear function
  - Convex under suitable conditions
  - Concave if -f(x) is convex
  - ▶ Otherwise it is neither convex nor concave
- Implications of convexity
  - Convex implies local solution is a global solution
  - Strictly convex implies at most one solution
  - Strongly convex implies one solution
  - Convex implies convex solution set
- ▶ Provable global solutions for nonconvex problems is difficult

#### Linear Program

- Convex optimization problem
- Specialized solution methods
  - Simplex method and variants (restarts well on related problems)
  - ► Interior-point methods (may not restart well)
  - Available software
    - ► CLP open source from COIN-OR
    - CPLEX and GUROBI commercial

#### Linear Program

Primal linear program

$$\min_{\substack{x \\ \text{subject to}}} c^{T}x \\
Ax + b \le 0 \\
x \ge 0$$

Dual linear program

$$\begin{array}{ll} \max_{\mu} & b^{T}\mu \\ \text{subject to} & A^{T}\mu + c \geq 0 \\ & \mu \geq 0 \end{array}$$

- Strong duality theory applies
  - Both feasible and optimal solutions are equal
  - ▶ One is infeasible and the other is unbounded
  - ▶ Both are infeasible
- Representation depends variables and constraints
  - ▶ Primal has *n* variables and *m* constraints
  - Dual has m variables and n constraints

#### Quadratic Program

$$\min_{\substack{x \text{subject to} \\ x \text{subject}}} \frac{\frac{1}{2}x^TQx + c^Tx}{Ax + b \le 0}$$

- Q is positive semidefinite implies
  - Convex optimization problem
  - Local solutions are global solutions
  - Specialized solution methods
    - Active-set methods (restarts well for related problems)
    - Interior-point methods (may not restart well)
  - Available software
    - OOQP source available
    - CPLEX and GUROBI commercial
- Otherwise problem is nonconvex
  - Attempt to compute local solution
  - May only find a stationary point
  - Provable global solutions for nonconvex problems is difficult

#### Nonlinear Program

$$\min_{x} f(x)$$
 subject to  $c(x) \le 0$ 

- Convex functions f and c<sub>i</sub>
  - Convex optimization problem
  - ▶ Feasible set  $X = \{x \mid c(x) \le 0\}$  is convex
  - Local solutions are global solutions
- Solutions methods
  - Attempt to compute local solution
  - Approximation methods
  - Interior-point methods
  - Provable global solutions for nonconvex problems is difficult
  - ► Algorithms discussed in next session

### **Taxonomy**

#### Nonlinear Program

$$\min_{x \in X} f(x)$$

with 
$$X = \{x \mid c(x) \le 0\}$$

▶ Global solution *x*\*:

$$f(x^*) \le f(x) \ \forall \ x \in X$$

- Global solutions need not be unique
- ▶ All global solutions have the same objective value
- ▶ Local solution *x*\*:

$$f(x^*) \le f(x) \ \forall \ x \in \mathcal{N}(x^*) \cap X$$

- ▶ Local solutions need not be unique
- ▶ Local solutions need not have the same objective value

### **Taxonomy**

#### Other Types of Nonlinear Programs (Maybe Discussed)

- Dynamic programming see other lectures
- ► Optimal control problem

Bilevel optimization problem

$$\begin{aligned} & \min_{x,y} & & f(x,y) \\ & \text{subject to} & & y \in \arg \left\{ \begin{array}{ll} \min_{\bar{y}} & & g(x,\bar{y}) \\ & \text{subject to} & c(x,\bar{y}) \leq 0 \end{array} \right. \end{aligned}$$

Multi-objective optimization (Pareto surfaces)



### **Taxonomy**

#### Other Types of Nonlinear Programs (Not Discussed)

- Semidefinite programs
- Second-order cone constraints
- ► Two-stage stochastic programs

$$\min_{x,y} f(x) + \sum_{s \in S} p_s g_s(x, y_s)$$
 subject to  $c_s(x, y_s) \leq 0 \quad \forall \ s \in S$ 

Semi-infinite optimization problem

Discrete optimization and categorical variables



#### Transformations of Nonsmooth Problems

▶ One norm:  $||x||_1 = \sum_i |x_i|$ 

$$\min_{x} f(x) + \|x\|_{1}$$

Smooth reformulation

$$\min_{\substack{x,y\\ \text{subject to}}} f(x) + e^T y$$

where  $y \in \Re^n$  is a vector

- ▶ Add *n* variables and 2*n* constraints
- Minimization is important
- ▶ Not applicable to

$$\max_{x} f(x) + \|x\|_{1}$$



#### Transformations of Nonsmooth Problems

▶ Infinity norm:  $||x||_{\infty} = \max_i |x_i|$ 

$$\min_{x} f(x) + ||x||_{\infty}$$

Smooth reformulation

$$\min_{\substack{x,y\\ \text{subject to}}} f(x) + y$$

where  $y \in \Re$  is a scalar

- ▶ Add one variable and 2*n* constraints
- Minimization is important
- ▶ Not applicable to

$$\max_{x} f(x) + \|x\|_{1}$$



#### Transformations of Nonsmooth Problems

▶ Pointwise maximum

$$\min_{x} f(x) + \max_{i} \{c_{i}(x)\}$$

Smooth reformulation

$$\min_{\substack{x,y\\ \text{subject to}}} f(x) + y \\
c_i(x) \le y \quad \forall i$$

where  $y \in \Re$  is a scalar

- Add one variable and several constraints
- Minimization is important
- Not applicable to

$$\max_{x} f(x) + \max_{i} \{c_{i}(x)\}$$



#### Transformations of Constraints

▶ One norm constraint:  $||x||_1 = \sum_i |x_i|$ 

$$\min_{x} f(x) 
\text{subject to} ||x||_{1} \le \Delta$$

Smooth reformulation

$$\begin{aligned} & \min_{x,y} & & f(x) \\ & \text{subject to} & & -y \leq x \leq y \\ & & & e^T y \leq \Delta \end{aligned}$$

where  $y \in \Re^n$  is a vector

- ▶ Add *n* variables and 2*n* constraints
- Not applicable to

$$\min_{\substack{x \\ \text{subject to}}} f(x)$$

#### Transformations of Constraints

▶ Infinity norm constraint:  $||x||_{\infty} = \max_i |x_i|$ 

$$\min_{x} f(x) 
\text{subject to} \|x\|_{\infty} \le \Delta$$

Smooth reformulation

$$\min_{\substack{x \\ \text{subject to}}} f(x) \\
-\Delta \le x \le \Delta$$

Not applicable to

$$\min_{\substack{x \\ \text{subject to}}} f(x)$$



#### Smooth Approximations of Nonsmooth Functions

Some examples

$$|x| \approx \sqrt{x^2 + \epsilon^2} - \epsilon$$

$$\sqrt{x^2 + y^2} \approx \sqrt{x^2 + y^2 + \epsilon^2} - \epsilon$$

$$(tx^p + (1-t)y^p)^{\frac{1}{p}} \approx (t(x+\epsilon)^p + (1-t)(y+\epsilon)^p)^{\frac{1}{p}} - \epsilon$$

▶ Derivatives can become bad as  $\epsilon \to 0$ 



### **Bad Transformations**

#### Do Not Make Them!

Square transformation

$$\min_{x} f(x)$$
subject to  $x \ge 0$ 

is equivalent to

$$\min_{y} f(y^2)$$

Exponential transformation

$$\min_{x} f(x)$$
subject to  $x > 0$ 

is equivalent to

$$\min_{y} f(a^{y})$$

for a > 0



# Student Discussion



# Part III

# Continuous Optimization Algorithms



# Overview of Continuous Optimization

- One-dimensional unconstrained optimization
  - Characterization of critical points
  - Basic algorithms
- Multi-dimensional unconstrained optimization
  - Critical points and their types
  - Computation of local minimizers
- Multi-dimensional constrained optimization
  - Critical points and Lagrange multipliers
  - Second-order sufficiency conditions
  - Globally-convergent algorithms
- Complementarity constraints
  - Stationarity concepts
  - Constraint qualifications
  - Numerical methods

#### Model Formulation

- Classify m people into two groups using v variables
  - $c \in \{0,1\}^m$  is the known classification
  - $d \in \Re^{m \times v}$  are the observations
  - $\beta \in \Re^{v+1}$  defines the separator
  - ▶ logit distribution function
- Maximum likelihood problem

$$\max_{\beta} \quad \sum_{i=1}^{m} c_i \log(f(\beta, d_{i, \cdot})) + (1 - c_i) \log(1 - f(\beta, d_{i, \cdot}))$$

where

$$f(\beta, x) = \frac{\exp\left(\beta_0 + \sum_{j=1}^{\nu} \beta_j x_j\right)}{1 + \exp\left(\beta_0 + \sum_{j=1}^{\nu} \beta_j x_j\right)}$$



#### Model Formulation

- Classify m people into two groups using v variables
  - $c \in \{0,1\}^m$  is the known classification
  - $d \in \Re^{m \times v}$  are the observations
  - ▶  $\beta \in \Re^{v+1}$  defines the separator
  - ▶ logit distribution function
- Maximum likelihood problem

$$\min_{eta} \quad -\left(\sum_{i=1}^m c_i \log(f(eta, d_{i,\cdot})) + (1-c_i) \log(1-f(eta, d_{i,\cdot}))\right)$$

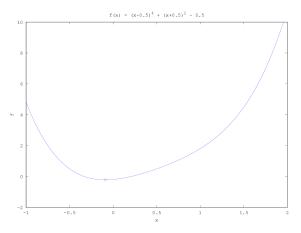
where

$$f(\beta, x) = \frac{\exp\left(\beta_0 + \sum_{j=1}^{\nu} \beta_j x_j\right)}{1 + \exp\left(\beta_0 + \sum_{j=1}^{\nu} \beta_j x_j\right)}$$



# Convex Unconstrained Optimization

#### Critical Points

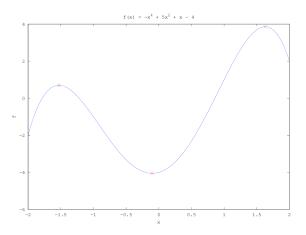


• Global minimizer:  $\nabla f(x) = 4(x - 0.5)^3 + 2(x + 0.5) = 0$ 



# Nonconvex Unconstrained Optimization

#### Critical Points



- ► Stationary:  $\nabla f(x) = -4x^3 + 10x + 1 = 0$
- ▶ Local maximizer:  $\nabla^2 f(x) = -12x^2 + 10 < 0$
- ▶ Local minimizer:  $\nabla^2 f(x) = -12x^2 + 10 > 0$

# Locally Convergent Newton Method for Optimization

- Most good algorithms are variants of Newton's method
- ▶ Attempt to compute local minimizers for  $f: \Re^n \to \Re$
- ▶ Settle for stationary points  $\nabla f(x) = 0$ 
  - Form Taylor series approximation around x<sup>k</sup>

$$f(x) \approx f(x^k) + \nabla f(x^k)^T (x - x^k) + \frac{1}{2} (x - x^k)^T \nabla^2 f(x^k) (x - x^k)$$

Solve quadratic optimization problem for s<sup>k</sup>

$$\min_{s} f(x^{k}) + \nabla f(x^{k})^{\mathsf{T}} s + \frac{1}{2} s^{\mathsf{T}} \nabla^{2} f(x^{k}) s$$

Convex case solutions satisfy

$$\nabla^2 f(x^k) s^k = -\nabla f(x^k)$$

- Nonconvex case can be unbounded
- ► Update iterate

$$x^{k+1} = x^k + s^k$$

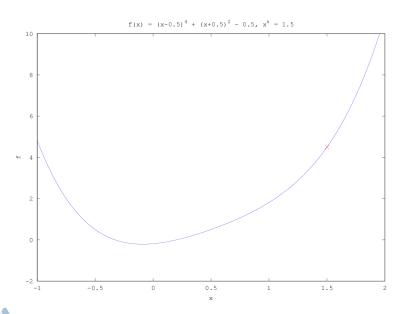
and repeat until convergence



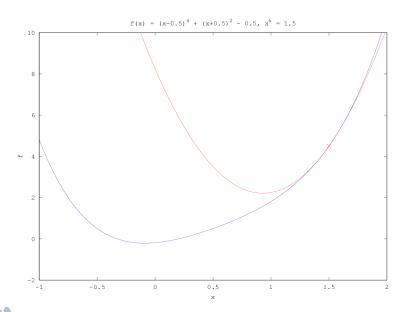
# Locally Convergent Newton Method for Optimization

Some Possible Outcomes

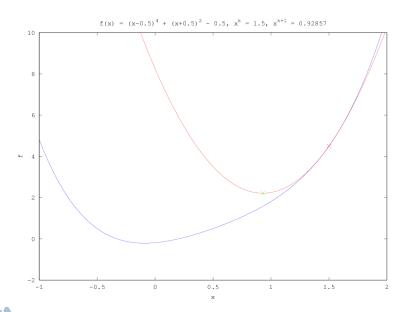
- ► Convergence:  $\lim_{k\to\infty} x^k = x^*$
- ▶ Divergence:  $\lim_{k\to\infty} ||x^k|| \to \infty$
- ▶ Unbounded below:  $\lim_{k\to\infty} f(x^k) \to -\infty$
- Sequence cycles
  - Multiple convergent subsequences (limit points)
  - Limit points are not solutions



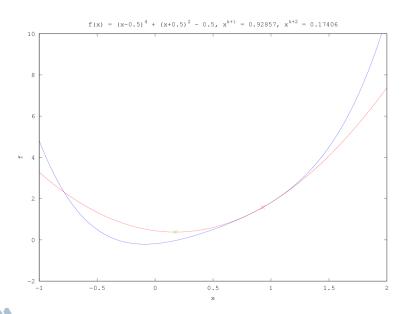


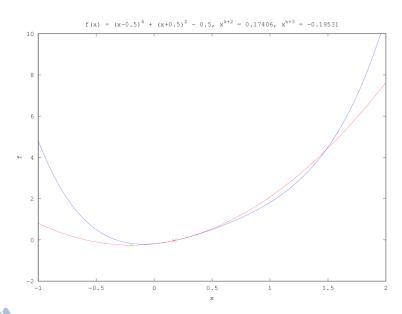




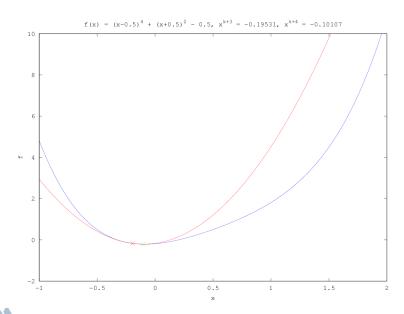




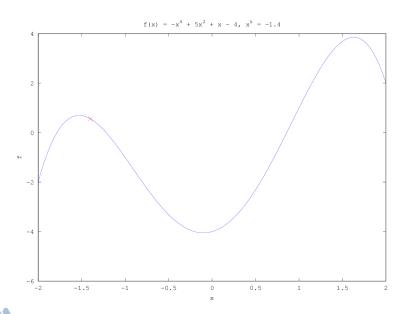






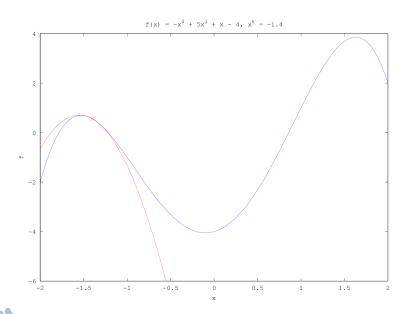


# Illustration on Nonconvex Problem





# Illustration on Nonconvex Problem



# Basic Theory

$$\min_{x} f(x)$$

- ► Convex functions local minimizers are global minimizers
- Nonconvex functions
  - ▶ Stationarity:  $\nabla f(x) = 0$
  - ▶ Local minimizer:  $\nabla^2 f(x)$  is positive definite (min eig positive)
  - ▶ Local maximizer:  $\nabla^2 f(x)$  is negative definite (max eig negative)



# Solving Unconstrained Optimization Problems

$$\min_{x} f(x)$$

Main ingredients of solution approaches:

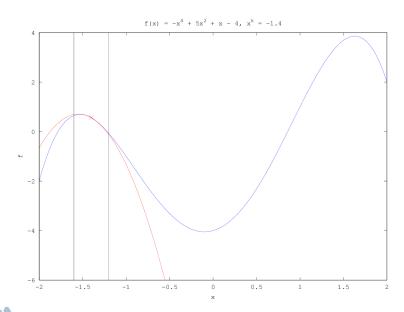
- ▶ Local method: given  $x^k$  (solution guess) compute a step  $s^k$ 
  - Gradient Descent
  - Quasi-Newton Approximation
  - Sequential Quadratic Programming
- ▶ Globalization strategy: converge from any starting point
  - Trust region
  - Line search

# Trust-Region Method

$$\min_{\substack{s \\ \text{subject to}}} f(x^k) + s^T \nabla f(x^k) + \frac{1}{2} s^T H(x^k) s$$
 where  $H(x^k)$  approximates  $\nabla^2 f(x^k)$ 

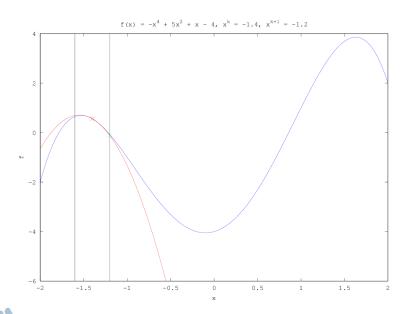


# Illustration on Nonconvex Problem



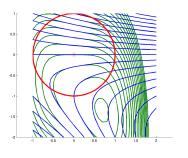


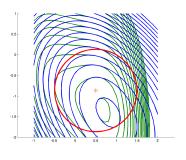
# Illustration on Nonconvex Problem





# Two-Dimensional Example







# Trust-Region Method

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve trust-region subproblem

$$\begin{aligned} & \min_{s} & & f(x^k) + s^T \nabla f(x^k) + \frac{1}{2} s^T H(x^k) s \\ & \text{subject to} & & \|s\|_2 \leq \Delta_k \end{aligned}$$



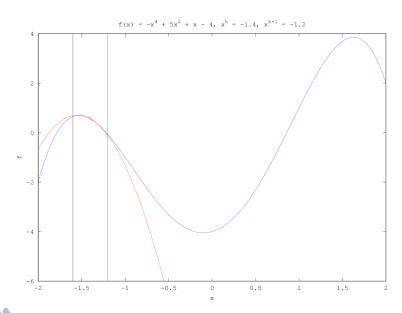
# Trust-Region Method

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve trust-region subproblem

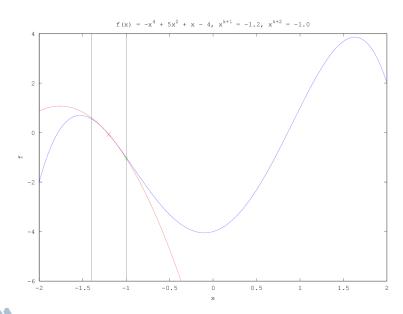
min<sub>s</sub> 
$$f(x^k) + s^T \nabla f(x^k) + \frac{1}{2} s^T H(x^k) s$$
  
subject to  $||s||_2 \le \Delta_k$ 

- 2.2 Accept or reject iterate
- 2.3 Update trust-region radius
  - Increase if actual reduction more than predicted
  - Decrease if actual reduction less than predicted
- 3. Check convergence

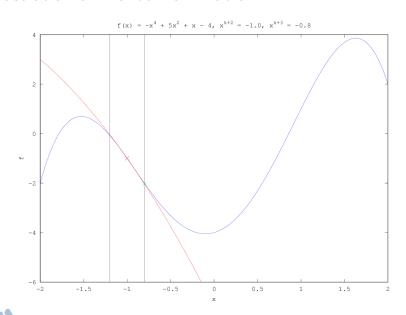
# Illustration on Nonconvex Problem



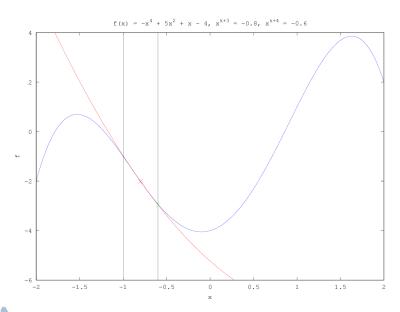




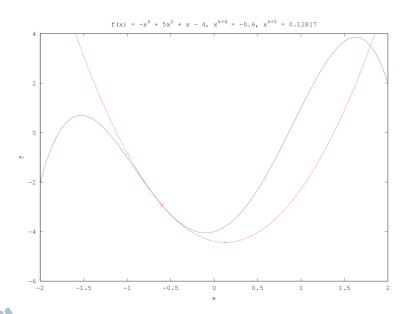


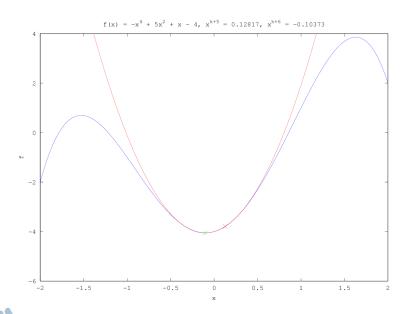












## Solving a Convex Quadratic Program

Assume the quadratic program is strictly convex

$$\min_{s} \quad \frac{1}{2}s^{T}Hs + c^{T}s$$

- ▶ *H* is symmetric and positive definite
- $\rightarrow$   $H^{-1}$  exists
- Stationary points are necessary and sufficient

$$Hs = -c$$

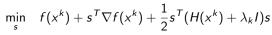
- Cholesky factorization
  - Compute (sparse) lower triangular matrix with  $H = LL^T$
  - Solve  $s = L^{-T}(L^{-1}c)$  exploiting lower triangular property
- Conjugate gradient method
  - ▶ Iteratively compute a set of *H* conjugate directions
  - Analytically minimize quadratic along the directions
  - Objective function decreases monotonically
  - Guaranteed convergence in n steps

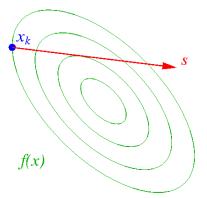
## Solving a Nonconvex Quadratic Program

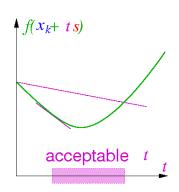
No assumptions on the quadratic program

$$\min_{s} \frac{1}{2} s^{T} H s + c^{T} s$$
subject to  $||s||_{2} \leq \Delta_{k}$ 

- Trust region bounds objective function
- No unbounded solutions
- ► Can detect inertia with a LDL<sup>T</sup> factorization and use direct method
- ► Global solutions computed with Moré-Sorensen method
  - Requires repeated factorization of a matrix
  - Can be expensive to calculate
  - ▶ Little benefit
- Conjugate gradient method with a trust region
  - ▶ Iteratively compute a set of *H* conjugate directions
  - ► Analytically minimize quadratic along the directions
  - Stop when trust region boundary is encountered
  - Objective function decreases monotonically









- 1. Initialize perturbation to zero
- 2. Solve perturbed quadratic model

$$\min_{s} \quad f(x^k) + s^T \nabla f(x^k) + \frac{1}{2} s^T (H(x^k) + \lambda_k I) s$$



- 1. Initialize perturbation to zero
- 2. Solve perturbed quadratic model

$$\min_{s} f(x^{k}) + s^{T} \nabla f(x^{k}) + \frac{1}{2} s^{T} (H(x^{k}) + \lambda_{k} I) s$$

- 3. Find new iterate
  - 3.1 Search along Newton direction
  - 3.2 Search along gradient-based direction

- 1. Initialize perturbation to zero
- 2. Solve perturbed quadratic model

$$\min_{s} f(x^{k}) + s^{T} \nabla f(x^{k}) + \frac{1}{2} s^{T} (H(x^{k}) + \lambda_{k} I) s$$

- 3. Find new iterate
  - 3.1 Search along Newton direction
  - 3.2 Search along gradient-based direction
- 4. Update perturbation
  - Decrease perturbation if the following hold
    - Iterative method succeeds
    - Search along Newton direction succeeds
  - Otherwise increase perturbation
- 5. Check convergence

## Solving the Subproblem

- ▶ Use LDL<sup>T</sup> to determine inertia and update perturbation
- ▶ Apply conjugate gradient method and stop on unbounded directions

## Solving the Subproblem

- ▶ Use LDL<sup>T</sup> to determine inertia and update perturbation
- Apply conjugate gradient method and stop on unbounded directions
- Conjugate gradient method with trust region
  - Initialize radius
  - Update radius

## Performing the Line Search

- Backtracking Armijo line search
  - Find t to satisfy sufficient decrease condition

$$f(x^k + ts) \le f(x^k) + \sigma t \nabla f(x^k)^T s$$

- ▶ Try  $t = 1, \beta, \beta^2, ...$  for  $0 < \beta < 1$
- More-Thuente line search
  - Find t to satisfy strong Wolfe conditions

$$\begin{array}{rcl} f(x^k + ts) & \leq & f(x^k) + \sigma t \nabla f(x^k)^T s \\ |\nabla f(x^k + ts)^T s| & \leq & \delta |\nabla f(x^k)^T s| \end{array}$$

- ► Construct cubic interpolant
- ▶ Compute *t* to minimize interpolant
- ► Refine interpolant



## Updating the Perturbation

1. If increasing and  $\lambda_k = 0$ 

$$\lambda_{k+1} = \mathsf{Proj}_{[\ell_0, u_0]} \left( \alpha_0 \| \nabla f(x^k) \| \right)$$

2. If increasing and  $\lambda_k > 0$ 

$$\lambda_{k+1} = \mathsf{Proj}_{[\ell_i, u_i]} \left( \mathsf{max} \left( \alpha_i \| \nabla f(x^k) \|, \beta_i \lambda_k \right) \right)$$

3. If decreasing

$$\lambda_{k+1} = \min \left( \alpha_d \| \nabla f(x^k) \|, \beta_d \lambda_k \right)$$

4. If  $\lambda_{k+1} < \ell_d$ , then  $\lambda_{k+1} = 0$ 



## Trust-Region Line-Search Method

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve trust-region subproblem

min<sub>s</sub> 
$$f(x^k) + s^T \nabla f(x^k) + \frac{1}{2} s^T H(x^k) s$$
  
subject to  $||s|| \le \Delta_k$ 

- 2.2 Search along direction
- 2.3 Update trust-region radius
- 3. Check convergence

### Iterative Methods

- ► Conjugate gradient method
  - Stop if negative curvature encountered
  - ► Stop if residual norm is small

#### Iterative Methods

- Conjugate gradient method
  - Stop if negative curvature encountered
  - ► Stop if residual norm is small
- ► Conjugate gradient method with trust region
  - Nash
    - Follow direction to boundary if first iteration
    - Stop at base of direction otherwise
  - Steihaug-Toint
    - Follow direction to boundary
  - Generalized Lanczos
    - Compute tridiagonal approximation
    - Find global solution to approximate problem on boundary
    - Initialize perturbation with approximate minimum eigenvalue

### Preconditioners to Improve Performance

Modify system of equations solved

$$MHs = -Mc$$

- M is symmetric positive definite
- ▶ M<sup>-1</sup> can be easily applied to vector
- ▶ MH is well conditioned or has clustered eigenvalues
- Corresponds to changing to an elliptic trust region

$$\begin{aligned} \min_{s} & & \frac{1}{2}s^{T}Hs + c^{T}s \\ \text{subject to} & & \|s\|_{M} \leq \Delta_{k} \end{aligned}$$

where 
$$||x||_M = \sqrt{x^T M x}$$
 and  $M$  defines the ellipse

- ▶ Preconditioners are problem specific
- Many possibly preconditioners
  - No preconditioner − M = I
  - ▶ Diagonal of Hessian  $M = |\text{diag}(H(x^k))|$
  - ▶ Diagonal of perturbed Hessian  $M = |\text{diag}(H(x^k) + \lambda_k I)|$
  - Quasi-newton approximation to Hessian matrix
  - Incomplete Cholesky factorization of Hessian
  - Block Jacobi with Cholesky factorization of blocks

#### **Termination**

- ► Typical convergence criteria
  - ▶ Absolute residual  $\|\nabla f(x^k)\| < \tau_a$
  - ▶ Relative residual  $\frac{\|\nabla f(x^k)\|}{\|\nabla f(x_0)\|} < \tau_r$
  - ▶ Unbounded objective  $f(x^k) < \kappa$
  - ▶ Slow progress  $|f(x^k) f(x_{k-1})| < \epsilon$
  - ▶ Iteration limit
  - ► Time limit
- Check the solver status

### Convergence Issues

- Quadratic convergence best outcome
- Linear convergence
  - ▶ Far from a solution  $-\|\nabla f(x^k)\|$  is large
  - ► Hessian is incorrect disrupts quadratic convergence
  - ▶ Hessian is rank deficient  $-\|\nabla f(x^k)\|$  is small
  - Limits of finite precision arithmetic
    - 1.  $\|\nabla f(x^k)\|$  converges quadratically to small number
    - 2.  $\|\nabla f(x^k)\|$  hovers around that number with no progress
- ▶ Domain violations such as  $\frac{1}{x}$  when x = 0
  - Make implicit constraints explicit
- Nonglobal solution
  - Apply a multistart heuristic
  - Use global optimization solver

#### Some Available Software

- ► TRON Newton method with trust-region
- ▶ LBFGS Limited-memory quasi-Newton method with line search
- ► TAO Toolkit for Advanced Optimization
  - NLS Newton line-search method
  - ▶ NTR Newton trust-region method
  - NTL Newton line-search/trust-region method
  - LMVM Limited-memory quasi-Newton method
  - CG Nonlinear conjugate gradient methods

## Social Planning Model

- ▶ Economy with *n* agents and *m* commodities
  - $ightharpoonup e \in \Re^{n \times m}$  are the endowments
  - $\bullet$   $\alpha \in \Re^{n \times m}$  and  $\beta \in \Re^{n \times m}$  are the utility parameters
  - lacksquare  $\lambda \in \Re^n$  are the social weights
- Social planning problem

$$\max_{\substack{k \geq 0 \\ x \geq 0}} \sum_{i=1}^{n} \lambda_{i} \left( \sum_{k=1}^{m} \frac{\alpha_{i,k} (1 + x_{i,k})^{1 - \beta_{i,k}}}{1 - \beta_{i,k}} \right)$$
subject to 
$$\sum_{i=1}^{n} x_{i,k} \leq \sum_{i=1}^{n} e_{i,k} \qquad \forall k = 1, \dots, m$$



## Life-Cycle Saving Model

- Maximize discounted utility
  - $u(\cdot)$  is the utility function
  - R is the retirement age
  - ► T is the terminal age
  - w is the wage
  - $\triangleright$   $\beta$  is the discount factor
  - r is the interest rate
- Optimization problem

$$\max_{s,c} \qquad \sum_{t=0}^{r} \beta^{t} u(c_{t})$$
subject to 
$$s_{t+1} = (1+r)s_{t} + w - c_{t} \quad t = 0, \dots, R-1$$

$$s_{t+1} = (1+r)s_{t} - c_{t} \quad t = R, \dots, T$$

$$s_{0} = s_{T+1} = 0$$



## Theory Revisited

▶ Strict descent direction s

$$\nabla f(x)^T s < 0$$

- ► Stationarity conditions (first-order conditions)
  - ► No feasible, strict descent directions
  - For all feasible directions s

$$\nabla f(x)^T s \geq 0$$

▶ Unconstrained case,  $s \in \Re^n$  and

$$\nabla f(x) = 0$$

- Constrained cases
  - Characterize feasible directions
  - ▶ Requires constraint qualification

## Basic Theory

$$\min_{x} f(x) \\
\text{subject to} c(x) \le 0$$

- Feasible and no strict descent directions
  - Constraint qualification LICQ, MFCQ
  - Linearized active constraints characterize directions
  - Objective gradient is a linear combination of constraint gradients
    - Under a constraint qualification, s = 0 solves the linear program

min 
$$\nabla f(x^*)^T s + f(x^*)$$
 subject to  $\nabla c(x^*) s + c(x^*) \le 0$ 

Note:  $x^*$  feasible, implies the linear program is feasible

▶ Dual linear program

$$\begin{array}{ll} \max_{\lambda} & c(x^*)^T \lambda \\ \text{subject to} & \nabla f(x^*) + \nabla c(x^*)^T \lambda = 0 \\ & \lambda \geq 0 \end{array}$$

- Linear programming theory
- ▶ Lagrangian:  $\mathcal{L}(x,\lambda) = f(x) + \lambda^T c(x)$

## Convergence Criteria

Satisfies Constraint Qualification (Slater)

has solution x = -1

Primal linear program

$$\min_{s} s$$
subject to  $-2s \le 0$ 

- ▶ Optimal solution is s = 0
- ▶ Dual program produces  $\lambda = \frac{1}{2}$

## Convergence Criteria

Lacks Constraint Qualification (Slater/LICQ)

has solution x = 0

► Primal linear program

$$\min_{s} s$$
subject to  $0 \le 0$ 

- ▶ Optimal solution is  $s = -\infty$
- ► Dual program is infeasible

### Convergence Criteria

Lacks Constraint Qualification (Slater/LICQ)

has solution x = 0

Primal linear program

$$\min_{s}$$
subject to  $0 \le 0$ 

- ▶ Optimal solution is  $s \in (-\infty, \infty)$
- ▶ Dual program produces  $\lambda \in [0, \infty)$
- Constraint qualification is sufficient but not required

## **Optimality Conditions**

▶ If  $x^*$  is a local minimizer and a constraint qualification holds, then there exist multipliers  $\lambda^* \ge 0$  such that

$$\nabla f(x^*) + \nabla c_{\mathcal{A}}(x^*)^T \lambda_{\mathcal{A}}^* = 0$$

- ▶ Lagrangian function  $\mathcal{L}(x,\lambda) = f(x) + \lambda^T c(x)$
- Optimality conditions can be written as

$$\nabla f(x) + \nabla c(x)^{T} \lambda = 0$$
  
 
$$0 \le \lambda \quad \perp \quad -c(x) \ge 0$$

Complementarity problem

## Solving Constrained Optimization Problems

#### Main ingredients of solution approaches:

- ▶ Local method: given  $x^k$  (solution guess) find a step  $s^k$ 
  - Sequential Quadratic Programming (SQP)
  - Sequential Linear/Quadratic Programming (SLQP)
  - ► Interior-Point Method (IPM)
- Globalization strategy: converge from any starting point
  - Trust region
  - ▶ Line search
- Acceptance criteria: filter or penalty function

# Sequential Linear Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate

# Sequential Linear Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve linear program

$$\begin{aligned} \min_{s} & f(x^k) + \nabla f(x^k)^T s \\ \text{subject to} & c(x^k) + \nabla c(x^k) s \leq 0 \\ & \|s\| \leq \Delta_k \end{aligned}$$



## Sequential Linear Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve linear program

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- 2.2 Accept or reject iterate
- 2.3 Update trust-region radius
- 3. Check convergence

# Sequential Quadratic Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate



# Sequential Quadratic Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve quadratic program

$$\min_{s} f(x^{k}) + \nabla f(x^{k})^{T} s + \frac{1}{2} s^{T} W(x^{k}, \lambda^{k}) s$$
subject to 
$$c(x^{k}) + \nabla c(x^{k}) s \leq 0$$

$$||s|| \leq \Delta_{k}$$

where  $W(x^k, \lambda^k)$  approximates  $\nabla^2_{x,x} \mathcal{L}(x^k, \lambda^k)$ 



# Sequential Quadratic Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve quadratic program

$$\min_{s} f(x^{k}) + \nabla f(x^{k})^{T} s + \frac{1}{2} s^{T} W(x^{k}, \lambda^{k}) s$$
subject to 
$$c(x^{k}) + \nabla c(x^{k}) s \leq 0$$

$$||s|| \leq \Delta_{k}$$

where 
$$W(x^k, \lambda^k)$$
 approximates  $\nabla^2_{x,x} \mathcal{L}(x^k, \lambda^k)$ 

- 2.2 Accept or reject iterate
- 2.3 Update trust-region radius
- 3. Check convergence



# Sequential Linear Quadratic Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate

# Sequential Linear Quadratic Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve linear program to predict active set

$$\begin{aligned} \min_{d} & f(x^{k}) + \nabla f(x^{k})^{T} d \\ \text{subject to} & c(x^{k}) + \nabla c(x^{k}) d \leq 0 \\ & \|d\| \leq \Delta_{k} \end{aligned}$$

where 
$$A_k = \{i \mid c(x^k) + \nabla c_i(x^k)^T d^k = 0\}$$



# Sequential Linear Quadratic Programming

- 1. Initialize trust-region radius
- 2. Compute a new iterate
  - 2.1 Solve linear program to predict active set

$$\begin{aligned} \min_{d} & f(x^{k}) + \nabla f(x^{k})^{T} d \\ \text{subject to} & c(x^{k}) + \nabla c(x^{k}) d \leq 0 \\ & \|d\| \leq \Delta_{k} \end{aligned}$$

where 
$$A_k = \{i \mid c(x^k) + \nabla c_i(x^k)^T d^k = 0\}$$

2.2 Solve equality constrained quadratic program

$$\min_{s} f(x^{k}) + \nabla f(x^{k})^{T} s + \frac{1}{2} s^{T} W(x^{k}, \lambda^{k}) s$$
subject to 
$$c_{\mathcal{A}_{k}}(x^{k}) + \nabla c_{\mathcal{A}_{k}}(x^{k}) s = 0$$

- 2.3 Accept or reject iterate
- 2.4 Update trust-region radius
- 3. Check convergence

# Acceptance Criteria

- ▶ Constraint violation:  $h(x) = \| \max(c(x), 0) \|$
- ▶ Decrease objective function value:  $f(x^k + s^k) \le f(x^k)$
- ▶ Decrease constraint violation:  $h(x^k + s^k) \le h(x^k)$

### Acceptance Criteria

- ► Constraint violation:  $h(x) = \| \max(c(x), 0) \|$
- ▶ Decrease objective function value:  $f(x^k + s^k) \le f(x^k)$
- ▶ Decrease constraint violation:  $h(x^k + s^k) \le h(x^k)$
- Four possibilities

-our possibilities	
1. step can decrease both $f(x)$ and $h(x)$	GOOD
2. step can decrease $f(x)$ and increase $h(x)$	???
3. step can increase $f(x)$ and decrease $h(x)$	???
4. step can increase both $f(x)$ and $h(x)$	BAD



### Acceptance Criteria

- ► Constraint violation:  $h(x) = \| \max(c(x), 0) \|$
- ▶ Decrease objective function value:  $f(x^k + s^k) \le f(x^k)$
- ▶ Decrease constraint violation:  $h(x^k + s^k) \le h(x^k)$
- ► Four possibilities
  - 1. step can decrease both f(x) and h(x)
  - 2. step can decrease f(x) and increase h(x)
  - 3. step can increase f(x) and decrease h(x)
  - 4. step can increase both f(x) and h(x)
- ▶ Filter uses concept from multi-objective optimization

$$(h_{k+1},f_{k+1})$$
 dominates  $(h_\ell,f_\ell)$  iff  $h_{k+1}\leq h_\ell$  and  $f_{k+1}\leq f_\ell$ 

GOOD

???

???

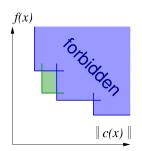
BAD

SAD

### Filter Framework

Filter  $\mathcal{F}$ : list of non-dominated pairs  $(h_{\ell}, f_{\ell})$ 

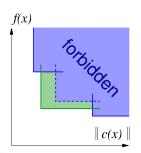
- ▶ new  $x^{k+1} = x^k + s^k$  is acceptable to filter  $\mathcal F$  iff for all  $\ell \in \mathcal F$ 
  - 1.  $h_{k+1} \le h_{\ell}$  or
  - $2. \ f_{k+1} \leq f_{\ell}$



### Filter Framework

Filter  $\mathcal{F}$ : list of non-dominated pairs  $(h_{\ell}, f_{\ell})$ 

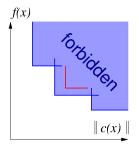
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  - 1.  $h_{k+1} \le h_{\ell}$  or
  - $2. \ f_{k+1} \leq f_{\ell}$
- ► remove redundant filter entries



### Filter Framework

Filter  $\mathcal{F}$ : list of non-dominated pairs  $(h_{\ell}, f_{\ell})$ 

- ▶ new  $x^{k+1} = x^k + s^k$  is acceptable to filter  $\mathcal F$  iff for all  $\ell \in \mathcal F$ 
  - 1.  $h_{k+1} \le h_{\ell}$  or
  - $2. \ f_{k+1} \leq f_{\ell}$
- remove redundant filter entries
- ▶ new  $x^{k+1}$  is rejected if for some  $\ell \in \mathcal{F}$ 
  - 1.  $h_{k+1} > h_\ell$  and
  - 2.  $f_{k+1} > f_{\ell}$



### **Termination**

- ▶ Feasible and complementary  $\|\min(-c(x^k), \lambda^k)\| \le \tau_f$
- ▶ Optimal  $\|\nabla_{\mathbf{x}}\mathcal{L}(\mathbf{x}^k,\lambda^k)\| \leq \tau_o$
- ▶ Other possible conditions
  - Slow progress
  - ► Iteration limit
  - ► Time limit
- Multipliers and reduced costs



### Convergence Issues

- Quadratic convergence best outcome
- Globally infeasible linear constraints infeasible
- ▶ Locally infeasible nonlinear constraints locally infeasible
- Unbounded objective hard to detect
- Unbounded multipliers constraint qualification not satisfied
- Linear convergence rate
  - Far from a solution
  - Hessian is incorrect disrupts quadratic convergence
  - Hessian is rank deficient
  - Limits of finite precision arithmetic
- ▶ Domain violations such as  $\frac{1}{x}$  when x = 0
  - Make implicit constraints explicit
- Nonglobal solutions
  - ► Apply a multistart heuristic
  - Use global optimization solver

### Some Available Software

- filterSQP
  - trust-region SQP; robust QP solver
  - filter to promote global convergence
- ▶ SNOPT
  - line-search SQP; null-space CG option
  - $\ell_1$  exact penalty function
- ► SLIQUE part of KNITRO
  - SLP-EQP
  - trust-region with  $\ell_1$  penalty
  - ▶ use with knitro\_options = "algorithm=3";

► Original optimization problem

Original optimization problem

Reformulate by adding slacks

$$\begin{aligned} \min_{x,u} & f(x) \\ \text{subject to} & c(x) + u = 0 \\ & x \geq 0, \ u \geq 0 \end{aligned}$$



Equality constrained optimization problem

Construct perturbed optimality conditions

$$F_{\tau}(x,\lambda,\nu) = \left[ egin{array}{c} 
abla f(x) + 
abla c(x)^T \lambda - \mu \\ 
-c(x) \\ 
x \odot \mu - \tau e 
\end{array} 
ight]$$

- ▶ Central path  $\{x(\tau), \lambda(\tau), \mu(\tau) \mid \tau > 0\}$
- ▶ Apply Newton's method for sequence  $\tau \searrow 0$



- 1. Compute a new iterate
  - 1.1 Solve linear system of equations

$$\begin{bmatrix} W^k & \nabla c(x^k)^T & -I \\ -\nabla c(x^k) & 0 & 0 \\ \operatorname{diag}(\mu^k) & 0 & \operatorname{diag}(x^k) \end{bmatrix} \begin{pmatrix} s_x \\ s_\lambda \\ s_\mu \end{pmatrix} = -F_\tau(x^k, \lambda^k, \mu^k)$$

- 1.2 Accept or reject iterate
- 1.3 Update parameters
- 2. Check convergence

### Convergence Issues

- Quadratic convergence best outcome
- ▶ Globally infeasible linear constraints infeasible
- ► Locally infeasible nonlinear constraints locally infeasible
- ▶ Dual infeasible dual problem is locally infeasible
- Unbounded objective hard to detect
- Unbounded multipliers constraint qualification not satisfied
- Duality gap
- ▶ Domain violations such as  $\frac{1}{x}$  when x = 0
  - Make implicit constraints explicit
- Nonglobal solutions
  - Apply a multistart heuristic
  - Use global optimization solver

### **Termination**

- ▶ Feasible and complementary  $\|\min(-c(x^k), \lambda^k)\| \le \tau_f$
- ▶ Optimal  $\|\nabla_{\mathbf{x}}\mathcal{L}(\mathbf{x}^k,\lambda^k)\| \leq \tau_o$
- ▶ Other possible conditions
  - Slow progress
  - ► Iteration limit
  - ► Time limit
- Multipliers and reduced costs



### Some Available Software

- ▶ IPOPT open source in COIN-OR
  - ▶ line-search filter algorithm
- ► KNITRO
  - trust-region Newton to solve barrier problem
  - $\ell_1$  penalty barrier function
  - Newton system: direct solves or null-space CG
- ► LOQO
  - ▶ line-search method
  - Newton system: modified Cholesky factorization

### Part IV

# Equilibrium Problems



### Nash Games

- ▶ Non-cooperative game played by *n* individuals
  - ► Each player selects a strategy to optimize their objective
  - Strategies for the other players are fixed
- ▶ Equilibrium reached when no improvement is possible

### Nash Games

- ▶ Non-cooperative game played by *n* individuals
  - ▶ Each player selects a strategy to optimize their objective
  - Strategies for the other players are fixed
- Equilibrium reached when no improvement is possible
- ▶ Characterization of two player equilibrium  $(x^*, y^*)$

$$x^* \in \left\{ egin{array}{ll} rg \min_{x \geq 0} & f_1(x,y^*) \ 
subject to & c_1(x) \leq 0 \ 
subset 
s$$

### Nash Games

- ▶ Non-cooperative game played by *n* individuals
  - ► Each player selects a strategy to optimize their objective
  - Strategies for the other players are fixed
- Equilibrium reached when no improvement is possible
- ▶ Characterization of two player equilibrium  $(x^*, y^*)$

$$x^* \in \left\{ egin{array}{ll} \mathop{\mathrm{arg\,min}}_{x \geq 0} & f_1(x,y^*) \\ \mathop{\mathsf{subject\ to}} & c_1(x) \leq 0 \\ \mathop{\mathsf{arg\,min}}_{y \geq 0} & f_2(x^*,y) \\ \mathop{\mathsf{subject\ to}} & c_2(y) \leq 0 \end{array} 
ight.$$

- Many applications in economics
  - Bimatrix games
  - Cournot duopoly models
  - General equilibrium models
  - Arrow-Debreau models

### Complementarity Formulation

- Assume each optimization problem is convex
  - $f_1(\cdot, y)$  is convex for each y
  - $f_2(x,\cdot)$  is convex for each x
  - $ightharpoonup c_1(\cdot)$  and  $c_2(\cdot)$  are convex and satisfy constraint qualification
- ▶ Then first-order conditions are necessary and sufficient

$$\begin{array}{lll} \min\limits_{\substack{x \geq 0 \\ \text{subject to}}} & f_1(x,y^*) \\ & c_1(x) \leq 0 \end{array} \Leftrightarrow \quad \begin{array}{lll} 0 \leq x & \bot & \nabla_x f_1(x,y^*) + \lambda_1^T \nabla_x c_1(x) \geq 0 \\ & 0 \leq \lambda_1 & \bot & -c_1(x) \geq 0 \end{array}$$

### Complementarity Formulation

- Assume each optimization problem is convex
  - $f_1(\cdot, y)$  is convex for each y
  - $f_2(x, \cdot)$  is convex for each x
  - $c_1(\cdot)$  and  $c_2(\cdot)$  are convex and satisfy constraint qualification
- ▶ Then first-order conditions are necessary and sufficient

$$\begin{array}{lll} \min\limits_{\substack{y \geq 0 \\ \text{subject to}}} & f_2(x^*,y) \\ & c_2(y) \leq 0 \end{array} \quad \Leftrightarrow \quad \begin{array}{lll} 0 \leq y & \bot & \nabla_y f_2(x^*,y) + \lambda_2^T \nabla_y c_2(y) \geq 0 \\ & 0 \leq \lambda_2 & \bot & -c_2(y) \geq 0 \end{array}$$

### Complementarity Formulation

- Assume each optimization problem is convex
  - $f_1(\cdot, y)$  is convex for each y
  - $f_2(x,\cdot)$  is convex for each x
  - $c_1(\cdot)$  and  $c_2(\cdot)$  are convex and satisfy constraint qualification
- ▶ Then first-order conditions are necessary and sufficient

$$0 \le x \qquad \bot \qquad \nabla_{x} f_{1}(x, y) + \lambda_{1}^{T} \nabla_{x} c_{1}(x) \ge 0$$

$$0 \le y \qquad \bot \qquad \nabla_{y} f_{2}(x, y) + \lambda_{2}^{T} \nabla_{y} c_{2}(y) \ge 0$$

$$0 \le \lambda_{1} \qquad \bot \qquad -c_{1}(y) \ge 0$$

$$0 \le \lambda_{2} \qquad \bot \qquad -c_{2}(y) \ge 0$$

▶ Nonlinear complementarity problem

$$0 \le x \perp F(x) \ge 0$$

- Square system number of variables and constraints the same
- ► Each solution is an equilibrium for the Nash game

#### Model Formulation

- ▶ Economy with *n* agents and *m* commodities
  - $e \in \Re^{n \times m}$  are the endowments
  - $\bullet$   $\alpha \in \Re^{n \times m}$  and  $\beta \in \Re^{n \times m}$  are the utility parameters
  - ▶  $p \in \Re^m$  are the commodity prices
- ▶ Agent *i* maximizes utility with budget constraint

$$\max_{\substack{x_{i,*} \geq 0}} \qquad \sum_{\substack{k=1 \\ m}}^m \frac{\alpha_{i,k} (1+x_{i,k})^{1-\beta_{i,k}}}{1-\beta_{i,k}}$$
 subject to 
$$\sum_{k=1}^m p_k \left(x_{i,k} - e_{i,k}\right) \leq 0$$

▶ Market *k* sets price for the commodity

$$0 \leq p_k \quad \perp \quad \sum_{i=1}^n \left(e_{i,k} - x_{i,k}\right) \geq 0$$



# Methods for Complementarity Problems

- Sequential linearization methods (PATH)
  - 1. Solve the linear complementarity problem

$$0 \le x \quad \perp \quad F(x_k) + \nabla F(x_k)(x - x_k) \ge 0$$

- 2. Perform a line search along merit function
- 3. Repeat until convergence

# Methods for Complementarity Problems

- Sequential linearization methods (PATH)
  - 1. Solve the linear complementarity problem

$$0 \le x \perp F(x_k) + \nabla F(x_k)(x - x_k) \ge 0$$

- 2. Perform a line search along merit function
- 3. Repeat until convergence
- Semismooth reformulation methods (SEMI)
  - Solve linear system of equations to obtain direction
  - Globalize with a trust region or line search
  - Less robust in general
- Interior-point methods

#### Semismooth Reformulation

▶ Define Fischer-Burmeister function

$$\phi(a, b) := a + b - \sqrt{a^2 + b^2}$$

- $\phi(a, b) = 0$  iff  $a \ge 0$ ,  $b \ge 0$ , and ab = 0
- ▶ Define the system

$$[\Phi(x)]_i = \phi(x_i, F_i(x))$$

- $x^*$  solves complementarity problem iff  $\Phi(x^*) = 0$
- Nonsmooth system of equations
  - Jacobian does not exist on set of measure zero
  - ▶ Degeneracy:  $x_i = F_i(x) = 0$
- ▶ Merit function  $\Psi(x) = \|\Phi(x)\|_2^2$  is differentiable!

# Semismooth Algorithm

▶ Calculate  $H^k \in \partial_B \Phi(x^k)$  and solve

$$H^k s^k = -\Phi(x^k)$$

▶ If this system has no solution or

$$\nabla \Psi(x^k)^T s^k > -p_1 \|s^k\|^{p_2}$$

then let 
$$s^k = -\nabla \Psi(x^k)$$

## Semismooth Algorithm

▶ Calculate  $H^k \in \partial_B \Phi(x^k)$  and solve

$$H^k s^k = -\Phi(x^k)$$

▶ If this system has no solution or

$$\nabla \Psi(x^k)^T s^k > -p_1 \|s^k\|^{p_2}$$

then let  $s^k = -\nabla \Psi(x^k)$ 

ightharpoonup Compute smallest nonnegative integer  $i_k$  such that

$$\Psi(x^k + \beta^{i_k} s^k) \le \Psi(x^k) + \sigma \beta^{i_k} \nabla \Psi(x^k) s^k$$

▶ Update the iterate

$$x^{k+1} = x^k + \beta^{i_k} s^k$$

and repeat until convergence



### Convergence Issues

- Quadratic convergence best outcome
- ► Linear convergence
  - ▶ Far from a solution  $-\Psi(x^k)$  is large
  - ▶ Jacobian is incorrect disrupts quadratic convergence
  - ▶ Jacobian is rank deficient  $-\|\nabla \Psi(x^k)\|$  is small
  - Converge to local minimizer guarantees rank deficiency
  - Limits of finite precision arithmetic
    - 1.  $\Psi(x^k)$  converges quadratically to small number
    - 2.  $\Psi(x^k)$  hovers around that number with no progress
- ▶ Domain violations such as  $\frac{1}{x}$  when x = 0

### Some Available Software

- ▶ PATH sequential linearization method
- MILES sequential linearization method
- ▶ SEMI semismooth linesearch method
- ► TAO Toolkit for Advanced Optimization
  - SSLS full-space semismooth linesearch methods
  - ► ASLS active-set semismooth linesearch methods
  - RSCS reduced-space method

### Definition

- Leader-follower game
  - Dominant player (leader) selects a strategy y\*
  - ▶ Then followers respond by playing a Nash game

$$x_i^* \in \left\{ egin{array}{ll} \mathop{\mathrm{arg\,min}}_{x_i \geq 0} & f_i(x,y) \\ \mathop{\mathrm{subject\ to}} & c_i(x_i) \leq 0 \end{array} \right.$$

Leader solves optimization problem with equilibrium constraints

$$\begin{array}{ll} \min\limits_{y\geq 0,x,\lambda} & g(x,y) \\ \text{subject to} & h(y)\leq 0 \\ & 0\leq x_i \perp \nabla_{x_i}f_i(x,y) + \lambda_i^T\nabla_{x_i}c_i(x_i)\geq 0 \\ & 0\leq \lambda_i \perp -c_i(x_i)\geq 0 \end{array}$$

- Many applications in economics
  - Optimal taxation
  - Tolling problems

#### Model Formulation

- ▶ Economy with *n* agents and *m* commodities
  - $e \in \Re^{n \times m}$  are the endowments
  - $\bullet$   $\alpha \in \Re^{n \times m}$  and  $\beta \in \Re^{n \times m}$  are the utility parameters
  - ▶  $p \in \Re^m$  are the commodity prices
- ▶ Agent *i* maximizes utility with budget constraint

$$\max_{\substack{x_{i,*} \geq 0}} \qquad \sum_{\substack{k=1 \\ m}}^m \frac{\alpha_{i,k} (1+x_{i,k})^{1-\beta_{i,k}}}{1-\beta_{i,k}}$$
 subject to 
$$\sum_{k=1}^m p_k \left(x_{i,k} - e_{i,k}\right) \leq 0$$

▶ Market *k* sets price for the commodity

$$0 \leq p_k \quad \perp \quad \sum_{i=1}^n \left(e_{i,k} - x_{i,k}\right) \geq 0$$



## Nonlinear Programming Formulation

$$\begin{aligned} \min_{\substack{x,y,\lambda,s,t\geq 0\\ \text{subject to}}} & g(x,y)\\ \text{subject to} & h(y) \leq 0\\ & s_i = \nabla_{x_i} f_i(x,y) + \lambda_i^T \nabla_{x_i} c_i(x_i)\\ & t_i = -c_i(x_i)\\ & \sum_i \left(s_i^T x_i + \lambda_i t_i\right) \leq 0 \end{aligned}$$

- Constraint qualification fails
  - Lagrange multiplier set unbounded
  - Constraint gradients linearly dependent
  - ► Central path does not exist
- ▶ Able to prove convergence results for some methods
- ▶ Reformulation very successful and versatile in practice

## Penalization Approach

$$\begin{aligned} \min_{\substack{x,y,\lambda,s,t\geq 0 \\ \text{subject to}}} & g(x,y) + \pi \sum_{i} \left( s_i^T x_i + \lambda_i t_i \right) \\ \text{subject to} & h(y) \leq 0 \\ & s_i = \nabla_{x_i} f_i(x,y) + \lambda_i^T \nabla_{x_i} c_i(x_i) \\ & t_i = -c_i(x_i) \end{aligned}$$

- ▶ Optimization problem satisfies constraint qualification
- ▶ Need to increase  $\pi$

## Relaxation Approach

$$\begin{aligned} \min_{\substack{x,y,\lambda,s,t\geq 0\\ \text{subject to}}} & g(x,y)\\ & h(y) \leq 0\\ & s_i = \nabla_{x_i} f_i(x,y) + \lambda_i^T \nabla_{x_i} c_i(x_i)\\ & t_i = -c_i(x_i)\\ & \sum_i \left(s_i^T x_i + \lambda_i t_i\right) \leq \tau \end{aligned}$$

ightharpoonup Need to decrease au



#### Limitations

- Multipliers may not exist
- Solvers can have a hard time computing solutions
  - ► Try different algorithms
  - Compute feasible starting point
- Stationary points may have descent directions
  - Checking for descent is an exponential problem
  - Strong stationary points found in certain cases
- Many stationary points global optimization

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  - Checking for descent is an exponential problem
  - Strong stationary points found in certain cases
- Many stationary points global optimization
- Formulation of follower problem
  - Multiple solutions to Nash game
  - Nonconvex objective or constraints
  - Existence of multipliers

### Part V

## Optimization Without Derivatives



## Part VI

# Stochastic Approximation Methods



## Stochastic Methods for Two Types of Problems

- A. Stochastic optimization
  - ▶ Modeling and algorithms for optimization under uncertainty
  - Stochasticity from problem and/or algorithm
- B. Deterministic optimization
  - Objectives and constraints deterministic
  - Methods are "randomized"

## Stochastic Methods for Two Types of Problems

- A. Stochastic optimization
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- B. Deterministic optimization
  - Objectives and constraints deterministic
  - Methods are "randomized"
- $\rightarrow$  Methods and analysis are related

# A. Stochastic Optimization Problems and Methods



## Stochastic Optimization

#### General problem

$$\min \left\{ f(x) = \mathbb{E}_{\xi} \left[ F(x, \xi) \right] : x \in X \right\} \tag{1}$$

- $\mathbf{x} \in \mathbb{R}^n$  decision variables
- $\triangleright$   $\xi$  vector of random variables
  - independent of x
  - $P(\xi)$  distribution function for  $\xi$
  - $\xi$  has support  $\Xi$
- $ightharpoonup F(x,\cdot)$  functional form of uncertainty for decision x
- ▶  $X \subseteq \mathbb{R}^n$  set defined by deterministic constraints
  - Also: stochastic/probabilistic constraints (not addressed here)

## Approach of Sampling Methods for $f(x) = \mathbb{E}_{\xi} [F(x, \xi)]$

- ▶ Let  $\xi^1, \xi^2, \cdots, \xi^N \sim P$
- ightharpoonup For  $x \in X$ , define:

$$f_N(x) = \frac{1}{N} \sum_{i=1}^N F(x, \xi^i)$$

- $f_N$  is a random variable (really, a stochastic process)
  - (depends on  $\left(\xi^1,\xi^2,\cdots,\xi^{\it N}\right)$ )

▶ Motivated by  $\mathbb{E}_{\xi} [f_N(x)] = f(x)$ 



## Bias of Sampling Methods

▶ Let 
$$f^* = f(x^*)$$
 for  $x^* \in X^* \subseteq X$ 

## Bias of Sampling Methods

- ▶ Let  $f^* = f(x^*)$  for  $x^* \in X^* \subseteq X$
- ▶ For any  $N \ge 1$ :

$$\mathbb{E}_{\xi}\left[f_{N}^{*}\right] \leq f^{*} = \mathbb{E}_{\xi}\left[F(x^{*}, \xi)\right]$$

because

$$\mathbb{E}_{\xi} [f_1^*] = \mathbb{E}_{\xi} [\min \{ F(x, \xi) : x \in X \}] \le \min \{ \mathbb{E}_{\xi} [F(x, \xi)] : x \in X \} = f^*$$



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- ightharpoonup Sampling problems result in optimal values below  $f^*$
- ▶  $f_N^*$  is biased estimator of  $f^*$

## Sample Average Approximation

- ▶ Draw realizations  $\hat{\xi}^1, \hat{\xi}^2, \cdots, \hat{\xi}^N \sim P$  of  $(\xi^1, \xi^2, \cdots, \xi^N)$
- ▶ Replace (1) with

$$\min \left\{ \frac{1}{N} \sum_{i=1}^{N} F(x, \hat{\xi}^i) : x \in X \right\}$$
 (2)

- $\hat{f}_N(x) = \frac{1}{N} \sum_{i=1}^N F(x, \hat{\xi}^i)$  deterministic
- lacktriangle Follows mean of the N sample paths defined by the (fixed)  $\hat{\xi}^i$



## SAA Algorithm

Input N, (maybe  $x^0 \in X$ )

- 1. Generate  $\hat{\xi}^1, \hat{\xi}^2, \cdots, \hat{\xi}^N \sim P$
- 2. Solve the deterministic problem

$$\min\left\{\frac{1}{N}\sum_{i=1}^{N}F(x,\hat{\xi}^{i}):\ x\in X\right\}$$

Output  $x_N^*$  (or  $X_N^*$ ).



## Convergence with N

- ► A sufficient condition:
  - ▶ For any  $\epsilon > 0$  there exists  $N_{\epsilon}$  so that

$$\left|\hat{f}_N(x) - f(x)\right| < \epsilon \qquad \forall N \ge N_\epsilon \quad \forall x \in X$$

with probability 1 (wp1).

- ▶ Then  $\hat{f}_N^* \to f^*$  wp1.
- ▶ (With additional assumptions on f and  $X^* \subset X$ ):

$$\operatorname{dist}(x_N^*, X^*) \to 0$$

 $(+ \text{ uniqueness, } X^* = x^*):$ 

$$x_N^* \rightarrow x^*$$

## Stochastic Approximation Method

#### Basically just:

#### Input $x^0$

1. 
$$x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k - \alpha_k s^k \right\}$$
,

$$k=0,1,\ldots$$

- $ightharpoonup \alpha_k$  a step size
- $\triangleright$   $s^k$  a random direction

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#### Generally assume:

$$\alpha_k$$
:  $\sum_{k=0}^{\infty} \alpha_k = \infty$ ,  $\sum_{k=0}^{\infty} \alpha_k^2 < \infty$   
 $s^k$ :  $\mathbb{E}\left[\nabla f(x^k)^T s^k\right] > 0$   
 $s^k$  is an ascent direction (in expectation) at  $x^k$ 

## Stochastic Approximation Method

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# Input $x^0$ 1. $x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k - \alpha_k s^k \right\}$ , $k = 0, 1, \dots$

- $ightharpoonup \alpha_k$  a step size
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#### Generally assume:

$$\begin{array}{l} \alpha_k\colon \sum_{k=0}^\infty \alpha_k = \infty, \ \sum_{k=0}^\infty \alpha_k^2 < \infty \\ s^k\colon \mathbb{E}\left[\nabla f(x^k)^T s^k\right] > 0 \\ s^k \text{ is an ascent direction (in expectation) at } x^k \end{array}$$

• "Exact" Stochastic Gradient Descent:  $s^k = \nabla f(x^k)$ 

## Classic SA Algorithms

- "Original" method is Robbins-Monro (1951)
- Without derivatives: Kiefer-Wolfowitz (1952)
   replaces gradient with finite-difference approximation, e.g.,

$$1. x^{k+1} \leftarrow x^k - \alpha_k s^k, \qquad k = 0, 1, \dots$$

where

$$s^{k} = \frac{F(x^{k} + h_{k}I_{n}; \hat{\xi}^{k}) - F(x^{k} - h_{k}I_{n}; \hat{\xi}^{k+1/2})}{2h_{k}}$$

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- ▶ Requires 2*n* evaluations every iteration
- Can appeal to variance reduction techniques (e.g., common RNs)
- ► Convergence  $x^k \to x^*$  if f strongly convex (near  $x^*$ ), usual conditions on  $\alpha_k$ ,  $h_k \to 0$ ,  $\sum_k \frac{\alpha_k^2}{h_k^k} < \infty$
- K-W recommend:  $\alpha_k = \frac{1}{k}$ ,  $h_k = \frac{1}{k^{1/3}}$

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- K-W recommend:  $\alpha_k = \frac{1}{k}$ ,  $h_k = \frac{1}{k^{1/3}}$
- Extensions such as SPSA (Spall) reduce number of evaluations (see randomized methods slides...)

#### Derivative-Based Stochastic Gradient Descent

- 1. Draw realization  $\hat{\xi}^k \sim P$  of  $\xi^k$
- 2. Compute  $s^k = \nabla_x F(x^k; \hat{\xi}^k)$
- 3. Update  $x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k \alpha_k s^k \right\}$

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- 3. Update  $x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k \alpha_k s^k \right\}$
- ▶  $\nabla_x F(x^k; \hat{\xi}^k)$  is an unbiased estimator for  $\nabla f(x^k)$
- ► Can incorporate curvature if desired e.g.,  $B^k s^k$  an unbiased estimator for  $(\nabla^2 f(x^k))^{-1} \nabla f(x^k)$
- Can work with subgradients
- ► Can even output  $x^N = \frac{1}{N} \sum_{k=1}^{N} x^k$



#### Modern Stochastic Gradient Descent Codes

#### Stochastic gradient descent seems inherently sequential

▶ Better in special cases, e.g.,

$$f(x) = \sum_{e \in \mathcal{E}} f_e(x_e), \qquad e \subset \{1, \cdots, n\}$$

 $|\mathcal{E}|$  and n large

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 $|\mathcal{E}|$  and n large

- ► HOGWILD! (Niu, Recht, Ré, Wright)
  - parallel, asynchronous implementation
  - http://i.stanford.edu/hazy/victor/Hogwild/
  - ightharpoonup Basic idea: Each processor samples an e uniformly from  $\mathcal E$  and updates the coordinates  $x_e$ , ... ties broken arbitrarily

# B. Randomized Algorithms for Deterministic Problems

## Randomized Algorithms for Deterministic Problems

$$\min \{f(x) : x \in X \subseteq \mathbb{R}^n\}$$

- ▶ f deterministic
- ► Random variables are now generated by the method, *not from the problem*
- ▶ Often assume properties of *f* 
  - e.g.,  $\nabla f$  is L'-Lipschitz:

$$\|\nabla f(x) - \nabla f(y)\| \le L' \|x - y\| \qquad \forall x, y \in X$$

e.g., f is strongly convex (with parameter  $\tau$ ):

$$f(x) \ge f(y) + (x - y)^T \nabla f(y) + \frac{\tau}{2} ||x - y||^2 \quad \forall x, y \in X$$



## Basic Algorithms

#### Matyas (e.g., 1965):

- ▶ Input  $x^0$ ; repeat:
  - 1. Generate Gaussian  $u^k$  (centered about 0)
  - 2. Evaluate  $f(x^k + u^k)$

3. 
$$x^{k+1} = \begin{cases} x^k + u^k & \text{if } f(x^k + u^k) < f(x^k) \\ x^k & \text{otherwise.} \end{cases}$$

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#### Poljak (e.g., 1987)

- ▶ Input  $x^0$ ,  $\{h_k, \mu_k\}_k$ ; repeat:
  - 1. Generate a random  $u^k \in R^n$
  - 2.  $x^{k+1} = x^k h_k \frac{f(x^k + \mu_k u^k) f(x^k)}{\mu_k} u^k$
  - $h_k > 0$  is the step size
  - $\mu_k > 0$  is called the smoothing parameter

#### Basic Coordinate Descent Method

Componentwise Lipschitz parameter M > 0:

$$|\nabla_i f(x + he_i) - \nabla_i f(x)| \le M|h|, \quad \forall h \in \mathbb{R}, \quad i = 1, \dots, n$$



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- 1. Choose  $i_k = \arg \max_{i=1,...,n} |\nabla_i f(x^k)|$
- 2. Update  $x^{k+1} = x^k \frac{1}{M} \nabla_{i_k} f(x^k) e_{i_k}$

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- ▶ Generates  $f(x^k) f^* \le \frac{2nMR^2}{k+4}$ , where  $R \ge ||x^0 x^*||$



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- ▶ Good: only updates  $x_{i_k}$
- ▶ Bad: requires entire gradient  $\nabla f(x^k)$



Component-wise Lipschitz parameter M > 0:

$$|\nabla_i f(x + he_i) - \nabla_i f(x)| \le L_i |h|, \quad \forall h \in \mathbb{R}, \quad i = 1, \dots, n$$

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- ► Generates  $\mathbb{E}\left[f(x^k)\right] f^* \le \frac{2nR_1^2}{k+4}$ , where  $R_1 = \max\{\|x x^*\|_1 : f(x) \le f(x_0)\}$

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- ▶ Good: only updates  $x_{i_k}$
- ▶ Better: requires only component  $i_k$  of gradient  $\nabla f(x^k)$
- Can also:
  - generate  $i_k$  proportional to coordinate Lipschitz parameters  $\{L_i\}_i$
  - perform block-coordinate (and other subspace) operations

## Gaussian Smoothing

- ▶ Let  $f : \mathbb{R}^n \to \mathbb{R}$  be deterministic
- ▶  $u \in \mathbb{R}^n$  from a Gaussian distribution,  $\mathbb{E}_u[u] = 0$ 
  - ▶ Here: Covariance matrix *I<sub>n</sub>*, general *C* OK
- ▶ For scalar  $\mu > 0$ , Gaussian-smoothed version of f:

$$f_{\mu}(x) = \mathbb{E}_{u}\left[f(x + \mu u)\right]$$

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$$f_{\mu}(x) = \mathbb{E}_{u}\left[f(x + \mu u)\right]$$

- ▶ If f is convex, then  $f_{\mu}(x) \ge f(x)$
- ▶ If f is convex and  $\nabla f$  is L'-Lipschitz, then

$$|f_{\mu}(x)-f(x)|\leq \frac{\mu^2}{2}L'n$$

## Gaussian Smoothing and Directional Derivatives

$$f_{\mu}(x) = \mathbb{E}_{u}\left[f(x + \mu u)\right]$$

▶ Derivative of f in the direction u:  $f'(x; u) = \lim_{h\downarrow 0} \frac{f(x+hu)-f(x)}{h}$ 



## Gaussian Smoothing and Directional Derivatives

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- ▶ Derivative of f in the direction u:  $f'(x; u) = \lim_{h \downarrow 0} \frac{f(x+hu) f(x)}{h}$
- $g_0(x) = f_u'(x)u$ 
  - ▶ If f is convex, then  $\mathbb{E}_u[g_0(x)]$  is a subgradient of f
  - ▶ If f is differentiable at x, then

$$\mathbb{E}_{u}\left[\|g_{0}(x)\|^{2}\right] \leq (n+4)\|\nabla f(x)\|^{2}$$



## Gaussian Smoothing and Directional Derivatives

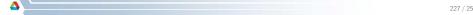
$$f_{\mu}(x) = \mathbb{E}_{u}\left[f(x + \mu u)\right]$$

- ▶ Derivative of f in the direction u:  $f'(x; u) = \lim_{h \downarrow 0} \frac{f(x+hu) f(x)}{h}$
- $g_0(x) = f'_u(x)u$ 
  - ▶ If f is convex, then  $\mathbb{E}_u[g_0(x)]$  is a subgradient of f
  - ▶ If f is differentiable at x, then

$$\mathbb{E}_{u}\left[\|g_{0}(x)\|^{2}\right] \leq (n+4)\|\nabla f(x)\|^{2}$$

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  - ▶ If f is differentiable at x, then  $\mathbb{E}_u[g_\mu(x)] = \nabla f_\mu(x)$
  - ▶ If f is differentiable at x and  $\nabla f$  is L'-Lipschitz, then

$$\mathbb{E}_{u}\left[\|g_{\mu}(x)\|^{2}\right] \leq 2(n+4)\|\nabla f(x)\|^{2} + \frac{\mu^{2}}{2}L'^{2}(n+6)^{3}$$



#### Random Gradient Method

Input  $x^0 \in X$ ,  $\{h_k\}_k$ ; repeat:

- 1. Generate Gaussian  $u^k \in R^n$  and compute  $g_0(x^k) = f'_{u^k}(x^k)u^k$
- 2.  $x^{k+1} = \mathcal{P}_X \{ x^k h_k g_0(x^k) \}$

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- ► Key result (Nesterov) for convex (but possibly nonsmooth) f: For fixed  $h_k = \frac{R}{\sqrt{n+4}\sqrt{N+1}L}$  and any  $\epsilon > 0$ ,

$$\mathbb{E}_{u}\left[f(\hat{x}^{N})\right] - f^{*} \leq \epsilon, \qquad \text{where} \quad \hat{x}^{N} = \arg\min_{i=1,\dots,N} f(x^{i})$$

in  $\mathcal{O}\left(\frac{n}{\epsilon^2}\right)$  iterations

 Also works for convex stochastic optimization and convex smooth f (with improved bounds and rates)



#### Random Gradient-Free Method

Input  $x^0 \in X$ ,  $\mu > 0$ ,  $\{h_k\}_k$ ; repeat:

- 1. Generate Gaussian  $u^k \in R^n$  and compute  $g_\mu(x^k) = \frac{f(x^k + u^k) f(x^k)}{\mu} u^k$
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- ► Key result (Nesterov) for convex (but possibly nonsmooth) f: For fixed  $h_k = \frac{R}{(n+4)\sqrt{N+1}L}$ ,  $\mu = \frac{\epsilon}{2L\sqrt{n}}$ , and any  $\epsilon > 0$ ,

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- in  $\mathcal{O}\left(\frac{n^2}{\epsilon^2}\right)$  iterations
- Also works for convex stochastic optimization and convex smooth f (with improved bounds and rates)

### Accelerated Random Gradient-Free Method

f strongly convex (with convexity parameter  $\tau$ )

Input  $v^0 = x^0$ ,  $\mu > 0$ ,  $\gamma_0 \ge \tau$ ,  $\{h_k\}_k$ ; repeat:

- 1. Obtain  $\alpha_k > 0$  satisfying  $16(n+1)^2 L' \alpha_k^2 = (1-\alpha_k)\gamma_k + \tau \alpha_k$
- 2. Set  $\gamma_{k+1} = (1 \alpha_k)\gamma_k + \tau \alpha_k$ ,  $\lambda_k = \frac{\alpha_k \tau}{\gamma_{k+1}}$ ,  $\beta_k = \frac{\alpha_k \gamma_k}{\gamma_k + \alpha_k \tau}$
- 3. Set  $y^k = (1 \beta_k)x^k + \beta_k v^k$
- 4. Generate Gaussian  $u^k \in R^n$  and compute  $g_{\mu}(y^k) = \frac{f(y^k + u^k) f(y^k)}{\mu} u^k$
- 5. Update

$$x^{k+1} = y^k - \frac{1}{4(n+4)L'} g_{\mu}(y^k) v^{k+1} = (1 - \lambda_k) v^k + \lambda_k y^k - \frac{1}{16(n+1)^2 L' \alpha_k} g_{\mu}(y^k)$$



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▶ Key result (Nesterov): for  $\tau = 0$  functions  $\exists \mu > 0$  so that

$$\mathbb{E}_{u}\left[f(\hat{x}^{N})\right] - f^{*} \leq \epsilon, \qquad \text{where} \quad \hat{x}^{N} = \arg\min_{i=1,\dots,N} f(x^{i})$$

in  $\mathcal{O}\left(\frac{n}{\epsilon^{1/2}}\right)$  iterations

## Applying SA-Like Ideas to Special Cases

$$\min \left\{ f(x) = \frac{1}{m} \sum_{i=1}^{m} F_i(x) : x \in X \right\}$$

$$m \text{ huge}$$



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- Ex.- Nonlinear Least Squares Warning: likely nonconvex!  $F_i(x) = \|\phi(x; \theta^i) d^i\|^2$  Evaluating  $\phi(\cdot, \cdot)$  requires solving a large PDE
- Ex.- Sample Average Approximation  $F_i(x) = R(x; \hat{\xi}^i)$   $\hat{\xi}^i \in \Omega$  a scenario/RV realization (and R depends nontrivially on  $\hat{\xi}^i$ )

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$$\hat{\xi}^i \in \Omega \text{ a scenario/RV realization}$$
 (and  $R$  depends nontrivially on  $\hat{\xi}^i$ )

The good:

The bad:

► *m* still huge

$$\min\left\{f(x) = \frac{1}{m}\sum_{i=1}^{m}F_i(x): x \in X\right\}$$

" $F_i(x)$  is a member of a population of size m"

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- ▶ Use  $-\nabla f_S = -\frac{1}{|S|} \sum_{i \in S} \nabla F_i(x)$  as direction  $s^k$
- ▶ How to choose S?

$$\mathbb{E}\left[\|\nabla f_{\mathcal{S}_n} - \nabla f\|^2\right] = \left(1 - \frac{|\mathcal{S}|}{m}\right) \mathbb{E}\left[\|\nabla f_{\mathcal{S}_r} - \nabla f\|^2\right]$$

 $\Rightarrow$  sampling without replacement  $(S_n)$  gives lower variance than does sampling with replacement  $(S_r)$ 



## Summary

- Methods for stochastic optimization and randomized methods for deterministic optimization closely related
- + Incredibly simple to code basic implementation
- + Well-studied complexity bounds, especially for convex cases; can show that asymptotic rates are optimal
- + Even useful when gradient/subgradient unavailable

## Summary

- Methods for stochastic optimization and randomized methods for deterministic optimization closely related
- + Incredibly simple to code basic implementation
- + Well-studied complexity bounds, especially for convex cases; can show that asymptotic rates are optimal
- + Even useful when gradient/subgradient unavailable
  - Bounds and parameters depend on characteristics of function (e.g., Lipschitz parameters, level set diameters, strong convexity)
  - (Some) Practitioners remain nervous about performance deviations from the mean (active research area)

## Part VII

# **Advanced Topics**



Optimize energy production schedule and transition between old and new reduced-carbon technology to meet carbon targets

- Maximize social welfare
- Constraints
  - Limit total greenhouse gas emissions
  - ▶ Low-carbon technology less costly as it becomes widespread
- ▶ Assumptions on emission rates, economic growth, and energy costs



#### Model Formulation

- Finite time:  $t \in [0, T]$
- ▶ Instantaneous energy output:  $q^{o}(t)$  and  $q^{n}(t)$
- ▶ Cumulative energy output:  $x^{o}(t)$  and  $x^{n}(t)$

$$x^n(t) = \int_0^t q^n(\tau) d\tau$$

Discounted greenhouse gases emissions

$$\int_0^T e^{-at} \left( b_o q^o(t) + b_n q^n(t) \right) dt \le z_T$$

- ▶ Consumer surplus S(Q(t), t) derived from utility
- Production costs
  - c<sub>o</sub> per unit cost of old technology
  - $ightharpoonup c_n(x^n(t))$  per unit cost of new technology (learning by doing)

#### Continuous-Time Model

$$\max_{\{q^o,q^n,x^n,z\}(t)} \qquad \int_0^T e^{-rt} \left[ S(q^o(t) + q^n(t),t) - c_o q^o(t) - c_n(x^n(t)) q^n(t) \right] dt$$
 subject to 
$$\dot{x^n}(t) = q^n(t) \quad x(0) = x_0 = 0$$
 
$$\dot{z}(t) = e^{-st} \left( b_o q^o(t) + b_n q^n(t) \right) \quad z(0) = z_0 = 0$$
 
$$z(T) \le z_T$$
 
$$q^o(t) \ge 0, \quad q^n(t) \ge 0.$$



#### Discretization:

- ▶  $t \in [0, T]$  replaced by N + 1 equally spaced points  $t_i = ih$
- ho h := T/N time integration step-length
- ▶ approximate  $q_i^n \simeq q^n(t_i)$  etc.

#### Replace differential equation

$$\dot{x}(t) = q^n(t)$$

by

$$x_{i+1} = x_i + hq_i^n$$

#### Discretization:

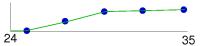
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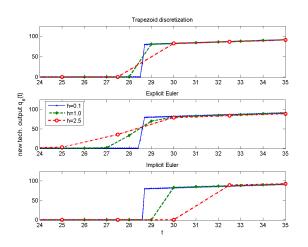
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$$x_{i+1} = x_i + hq_i^n$$



Output of new technology between t=24 and t=35

## Solution with Varying h



Output for different discretization schemes and step-sizes

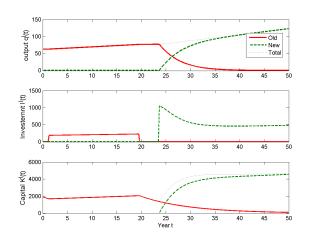
# Add adjustment cost to model building of capacity: Capital and Investment:

- $K^{j}(t)$  amount of capital in technology j at t.
- ▶  $I^{j}(t)$  investment to increase  $K^{j}(t)$ .
- ▶ initial capital level as  $\bar{K}_0^j$ :

#### Notation:

- $Q(t) = q^{o}(t) + q^{n}(t)$
- $C(t) = C^{o}(q^{o}(t), K^{o}(t)) + C^{n}(q^{n}(t), K^{n}(t))$
- $I(t) = I^o(t) + I^n(t)$
- $K(t) = K^{o}(t) + K^{n}(t)$

$$\begin{aligned} & \underset{\{q^j,K^j,l^j,x,z\}(t)}{\text{maximize}} & \left\{ \int_0^T e^{-rt} \left[ \tilde{S}(Q(t),t) - C(t) - K(t) \right] dt + e^{-rT} K(T) \right\} \end{aligned}$$
 subject to 
$$\dot{x}(t) = q^n(t), \quad x(0) = x_0 = 0$$
 
$$\dot{K}^j(t) = -\delta K^j(t) + l^j(t), \quad K^j(0) = \bar{K}^j_0, \quad j \in \{o,n\}$$
 
$$\dot{z}(t) = e^{-at} [b_o q^o(t) + b_n q^n(t)], \quad z(0) = z_0 = 0$$
 
$$z(T) \leq z_T$$
 
$$q^j(t) \geq 0, \ j \in \{o,n\}$$
 
$$l^j(t) \geq 0, \ j \in \{o,n\}$$



Optimal output, investment, and capital for 50% CO2 reduction.



## Pitfalls of Discretizations [Hager, 2000]

#### Optimal Control Problem

minimize 
$$\frac{1}{2}\int_0^1 u^2(t) + 2y^2(t)dt$$

subject to

$$\dot{y}(t) = \frac{1}{2}y(t) + u(t), \ t \in [0, 1],$$
  
 $y(0) = 1.$ 

$$\Rightarrow y^*(t) = \frac{2e^{3t} + e^3}{e^{3t/2}(2 + e^3)},$$
$$u^*(t) = \frac{2(e^{3t} - e^3)}{e^{3t/2}(2 + e^3)}.$$



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#### Discretize with 2nd order RK

minimize 
$$\frac{1}{2} \int_0^1 u^2(t) + 2y^2(t) dt$$
 minimize  $\frac{h}{2} \sum_{k=0}^{K-1} u_{k+1/2}^2 + 2y_{k+1/2}^2$ 

subject to 
$$(k = 0, ..., K)$$
:

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$$y(0) = 1.$$

$$y_{k+1/2} = y_k + \frac{h}{2}(\frac{1}{2}y_k + u_k),$$

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subject to (k = 0, ..., K):

$$y_{k+1/2} = y_k + \frac{h}{2}(\frac{1}{2}y_k + u_k),$$
  
 $y_{k+1} = y_k + h(\frac{1}{2}y_{k+1/2} + u_{k+1/2}),$ 

Discrete solution (k = 0, ..., K):

$$y_k = 1, \quad y_{k+1/2} = 0,$$
  
 $u_k = -\frac{4+h}{2h}, \quad u_{k+1/2} = 0,$ 

DOES NOT CONVERGE!



## Tips to Solve Continuous-Time Problems

- Use discretize-then-optimize with different schemes
- $\blacktriangleright$  Refine discretization: h=1 discretization is nonsense
- Check implied discretization of adjoints

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#### Alternative: Optimize-Then-Discretize

- Consistent adjoint/dual discretization
- Discretized gradients can be wrong!
- Harder for inequality constraints