

## Tutorial: Getting Started with Optimization: Computational Noise, Noisy Derivatives, Stochastic Methods

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#### Do this now:

- Obtain matlab files (from Simon's cluster or at www.mcs.anl.gov/~wild/codes/zice16.zip)
- Open matlab (ideally on your own machine so that you can view graphics, otherwise on the cluster)
- $\circ$  [Optional:] Have your function ready (a matlab function that receives x and outputs f(x))



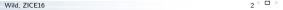
#### I. Computational Noise



- What is computational noise?
- How can noise be estimated efficiently?
- How does noise affect numerical differentiation?
- How accurate are near-optimal finite-difference estimates?

#### Questions To Ask Yourself

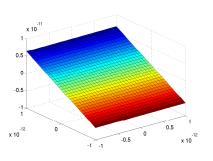
- 1. Do you know how "noisy" your function is?
- 2. Do you know how accurate your derivatives are?
- 3. Is the noise/accuracy stationary (independent of x)?
- 4. What do you do with this information?

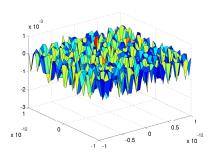




## Noise May Hurt You, Or It May Not

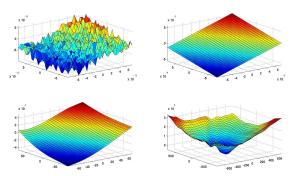
#### These are the same problem:





## Noise May Hurt You, Or It May Not

#### So are these:





#### Computational Noise is not a Newcomer

#### From Hamming's 1971 Introduction to Numerical Analysis:

Where does this noise come from? ...infinite processes in mathematics which of necessity must be approximated by finite processes.

Truncation vs. roundoff Finite number length leads to roundoff. Finite processes lead to truncation.



Competing errors Smaller steps usually reduce truncation error and may increase roundoff error.

Deterministic In practice, the same input, barring machine failures, gives the same result.



#### Computational Noise is not a Newcomer

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Competing errors Smaller steps usually reduce truncation error and may increase roundoff error.

Deterministic In practice, the same input, barring machine failures, gives the same result. 

— changing!



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## Living In A Finite-Precision World

Roundoff Error

$$f_{\infty}(x) - f(x)$$

#### Floating Point Arithmetic

Commutative:

$$A + B = B + A$$
 and  $A * B = B * A$ 

Non-associative:

$$A + (B + C) \neq (A + B) + C$$

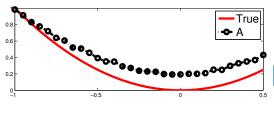
 This is likely to affect the reproducibility of your calculations in the future (for performance reasons)

Many details  $\rightarrow$  [What Every Computer Scientist Should Know About Floating-Point Arithmetic, Goldberg, 1991]

Wild, ZICE16 5 4 □ >



## Truncation/Approximation Error

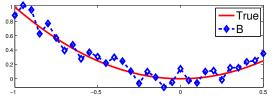




$$R_{m+1}(x) = f_a(x) - \sum_{i=0}^{m} P_i(x)$$

#### Which do you prefer?

- A less noise, more error
- B less error, more noise



## Computational Noise in Deterministic Simulations

Finite precision + finite processes

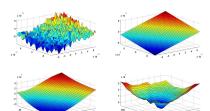
- Iteratively solving systems of PDEs or estimating eigenvalues
- Adaptively computing integrals
- Discretizations/meshes

destroy underlying smoothness

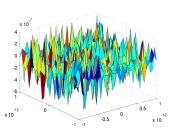
# $\underline{\mathsf{Goal}}$ : estimate the "variation" in $f(\mathbf{x})$

- ⋄ a few f evaluations
- deterministic and stochastic noise

## Difference $|f(x) - f(x + Z\omega)|$ ,







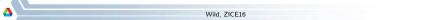
Sparse linear large-scale system

#### Matlab Time

#### Basic tips

(Examples in runexamples.m)

- $\diamond$  Moving from n-d to 1-d
- Deterministic function (probnum=1)
- Stochastic function (probnum=2)
- Scaling (probnum=3)
- Constraint cautions



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## Estimating Computational Noise: The Noise Level $\epsilon_f$

#### Simple model for the noise

$$f(t) = f_s(t) + \varepsilon(t), \quad t \in \mathcal{I}$$

- the computed function
- a smooth, deterministic function
- $\varepsilon$  is the noise with  $\{\varepsilon(t):t\in\mathcal{I}\}$  iid

← only assumption

The noise level of f is  $\varepsilon_f = (\operatorname{Var} \{ \varepsilon(t) \})^{1/2}$ 

(independent of t)



## The k-th Order Difference $\Delta^k f(t)$

$$\Delta^{k+1} f(t) = \Delta^k f(t+h) - \Delta^k f(t), \qquad \Delta^0 f(t) = f(t)$$

$$\Delta^k f(t) = \Delta^k f_s(t) + \Delta^k \varepsilon(t)$$

- 1. Differences of smooth  $f_s$  tend to zero rapidly
- 2. Differences of noise are bounded away from zero
  - ullet If h is sufficiently small,

$$\Delta^k f(t) \approx \Delta^k \varepsilon(t)$$

• If  $f_s$  is k-times differentiable,

$$\Delta^k f(t) = f_s^{(k)}(\xi_k) h^k + \Delta^k \varepsilon(t), \qquad \xi_k \in (t, t + kh)$$

Goal: make h small enough to remove smooth component

△ Wild, ZICE16

## Theory Underlying the ECNoise Algorithm

## For $\{\varepsilon(t+ih): i=0,\ldots,m\}$ iid and $k\leq m$ :

- 1.  $\mathrm{E}\left\{\Delta^k\varepsilon(t)\right\}=0$
- 2.  $\gamma_k \mathbf{E}\left\{ \left[ \Delta^k \varepsilon(t) \right]^2 \right\} = \varepsilon_f^2 \qquad \gamma_k = \frac{(k!)^2}{(2k)!}$
- 3. If  $f_s$  is continuous at t, then

$$\lim_{h \to 0} \gamma_k \mathbf{E} \left\{ \left[ \Delta^k f(t) \right]^2 \right\} = \varepsilon_f^2$$

4. If  $f_s$  is k-times continuously differentiable at t, then

$$\lim_{h \to 0} \frac{\gamma_k \mathbf{E} \left\{ [\boldsymbol{\Delta}^k f(t)]^2 \right\} - \boldsymbol{\varepsilon}_f^2}{h^{2k}} = \gamma_k \left[ f_s^{(k)}(t) \right]^2$$

$$\Rightarrow \varepsilon_f^2 \approx \gamma_k \mathbf{E} \left\{ [\Delta^k f(t)]^2 \right\},\,$$

when the sampling distance h is sufficiently small

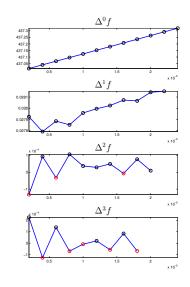


## The ECNoise Algorithm

Uses 
$$\sigma_k = \left(\frac{\gamma_k}{m+1-k}\sum_{i=0}^{m-k}[\Delta^k f(t+ih)]^2\right)^{1/2}$$

- 1. Chooses k
- 2. Verifies h is small enough
- $\diamond$  Works for deterministic f

[Estimating Computational Noise. Moré & W., SISC 2011]





ECNoise Estimator 
$$\sigma_k = \left(\frac{\gamma_k}{m+1-k}\sum_{i=0}^{m-k}[\Delta^k f(t_i)]^2\right)^{1/2}$$

For 
$$f(t) = \cos(t) + \sin(t) + 10^{-3} U_{[0,2\sqrt{3}]} \ \left( m = 6, t_i = \frac{i}{100} \right)$$

$f(t_i)$	$\Delta f(t_i)$	$\Delta^2 f(t_i)$	$\Delta^3 f(t_i)$	$\Delta^4 f(t_i)$	$\Delta^5 f(t_i)$	$\Delta^6 f(t_i)$
1.003	7.54e-3	2.15e-3	1.87e-4	-5.87e-3	1.46e-2	-2.49e-2
1.011	9.69e-3	2.33e-3	-5.68e-3	8.73e-3	-1.03e-2	
1.021	1.20e-2	-3.35e-3	3.05e-3	-1.61e-3		
1.033	8.67e-3	-2.96e-4	1.44e-3			
1.041	8.38e-3	1.14e-3				
1.050	9.52e-3					
1.059						
$\sigma_k$	6.78e-3	8.96e-4	9.02e-4	9.93e-4	1.10e-3	1.14e-3



## Extension to Multivariate $g: \mathbb{R}^n \mapsto \mathbb{R}$

Given base point  $x_b \in \mathbb{R}^n$ , unit direction  $p \in \mathbb{R}^n$ , consider

$$f_p(t) = g(x_b + tp), \quad t \ge 0$$

Apply univariate theory

- Directional differences, directional derivatives
- $\diamond$   $\varepsilon_f$  may now depend on a direction  $p \in \mathbb{R}^n$
- $\diamond$  ECnoise uses  $T_{i,0} = f(x_b + ihp)$  with random unit direction  $p \in \mathbb{R}^n$



## Computational Experience with Stochastic Noise

Validate ECnoise and empirical properties of

$$\sigma_k^2 = \frac{\gamma_k}{m+1-k} \sum_{i=0}^{m-k} T_{i,k}^2$$

under known conditions:

- ♦ Known noise level  $\varepsilon_f$
- Theory directly applies

Target: every estimate within a factor  $\eta=4$  of the mean  $% \left( 1\right) =1$ 



Noisy Quadratic, 
$$f(x) = (x^T x)(1 + R), \quad x \in \mathbb{R}^{10}$$

#### Estimate relative noise

$$\frac{\sigma_k}{f(x_h)} \approx \sqrt{\operatorname{Var}\{R\}} = 10^{-3}$$

- $x_b$  random base point
  - p 10000 random unit directions
- m evaluations

Noisy Quadratic, 
$$f(x) = (x^T x)(1 + R), \quad x \in \mathbb{R}^{10}$$

$$R \sim \! \mathsf{Uniform} \big[ -\sqrt{3} \cdot 10^{-3}, \sqrt{3} \cdot 10^{-3} \big]$$

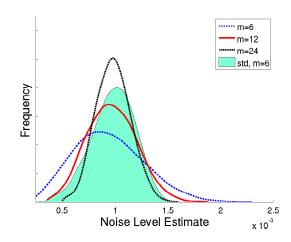
Estimate relative noise 
$$\frac{\sigma_k}{f(x_k)} \approx \sqrt{\operatorname{Var}\left\{R\right\}} = 10^{-3}$$

 $x_b$  random base point

 $p \ 10000$  random unit directions

m evaluations

99.2% within a factor  $\eta=4$  for m=6





Noisy Quadratic, 
$$f(x) = (x^T x)(1 + R), \quad x \in \mathbb{R}^{10}$$

$$R \sim \mathsf{Normal}(0, 10^{-6})$$

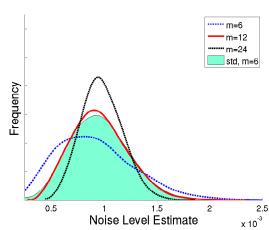
Estimate relative noise  $\frac{\sigma_k}{f(x_k)} \approx \sqrt{\operatorname{Var}\left\{R\right\}} = 10^{-3}$ 

 $x_b$  random base point

p 10000 random unit directions

m evaluations

98.9% within a factor  $\eta=4$  for m=6





## MC Finance Example with Higher Order Derivatives

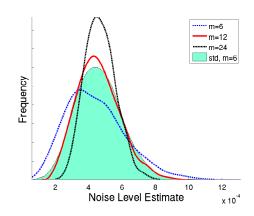
Today's value of a \$1 payment n years from now rates [Caflisch]:

$$f(x) = \frac{1}{(2\pi)^{n/2}} \int_{\mathbb{R}^n} \prod_{i=0}^n \frac{e^{-\frac{\|u\|^2}{2}}}{1 + r_i(u, x)} du, \quad r_i(u, x) = \begin{cases} \frac{1}{10} & i = 0\\ r_{i-1}(u, x) e^{x_i u_i - x_i^2/2} & i \ge 1 \end{cases}$$

10000 MC integrations (directions p) with

- n=3 years,  $x_b = [.1, .1, .1]$
- tol = 5000 standard normal random variables
- no variance reduction

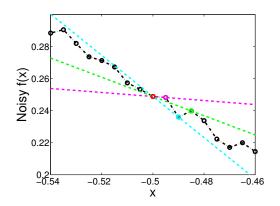
99.6% within a factor 4 for m=6





#### Finite Differences Sensitive to Choice of h

$$\frac{f(t_0+h)-f(t_0)}{h}\approx f_s'(t_0)$$





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## Noisy Forward Differences

$$\mathbb{E}\left\{\mathcal{E}(h)\right\} = \mathbb{E}\left\{\left(\frac{f(t_0+h)-f(t_0)}{h} - f_s'(t_0)\right)^2\right\}$$

#### Our h will depend on

- Loose estimate of noise
- $\diamond$  Loose estimate of |f''|
- Stochastic theory:
  - 1.  $f(t) = f_s(t) + \epsilon$  on  $I = \{t_0 + h : 0 \le h \le h_0\}$
  - 2.  $f_s$  twice differentiable
  - 3.  $\mu_L < |f_s''| < \mu_M$  on I

[Estimating Noisy Derivatives. Moré & W., TOMS 2012]]



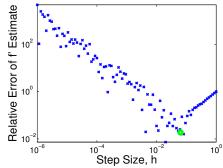
## Optimal Forward Difference Parameter h

$$\frac{1}{4}\mu_{\scriptscriptstyle L}^2h^2 + 2\frac{\varepsilon_f^2}{h^2} \leq \operatorname{E}\left\{\mathcal{E}(h)\right\} \leq \frac{1}{4}\mu_{\scriptscriptstyle M}^2h^2 + 2\frac{\varepsilon_f^2}{h^2}$$

- $h \downarrow Variance (noise) dominates$
- $h \uparrow \text{ Bias } (f'') \text{ dominates}$

For  $h_0$  sufficiently large

- 1. Upper bound minimized by  $h_M = 8^{1/4} \left( \frac{\varepsilon_f}{u_M} \right)^{1/2}$
- 2. When  $\mu_L > 0$ ,  $h_M$  is near-optimal:



$$\mathrm{E}\left\{\mathcal{E}(h_{M})\right\} = \sqrt{2}\mu_{M}\varepsilon_{f} \leq \left(\frac{\mu_{M}}{\mu_{L}}\right) \min_{0 \leq h \leq h_{0}} \mathrm{E}\left\{\mathcal{E}(h)\right\}.$$

Given uniform bound on roundoff error,

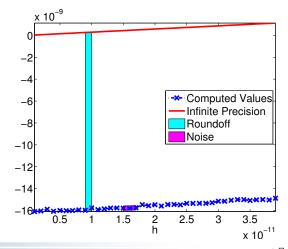
$$|f(t) - f_{\infty}(t)| \le \varepsilon_A \qquad t \in I,$$

Minimizer of (upper bound on)  $l_1$  error is

$$h_A = 2 \left( \frac{arepsilon_A}{\mu_M} \right)^{1/2}$$

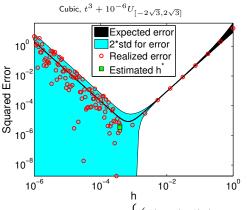
Assumes:

- $h_A \leq h_0$
- $\diamond$  Estimate of  $\varepsilon_A$  available



#### Stochastic Examples

Estimate 
$$f_s'(t) = E\{f(t)\}'$$
 at  $t=1$   $(\varepsilon_f = 10^{-6})$ 



Log-log realizations of 
$$\mathcal{E}(h) = \mathrm{E}\left\{\left(\frac{f(t_0+h)-f(t_0)}{h} - f_s'(t_0)\right)^2\right\}$$

Expected error and uncertainty regions predicted by the theory

#### Extension: Central Differences

# First derivatives, $\frac{f(t_0+h)-f(t_0-h)}{2h}$

$$|h_M| = \gamma_5 \left(\frac{\varepsilon_f}{\mu_M}\right)^{1/3}, \qquad \gamma_5 = 3^{1/3} \approx 1.44$$

$$\diamond \ \mathrm{E}\left\{\mathcal{E}_{c}(h_{M})\right\} \leq \left(\frac{\mu_{M}}{\mu_{L}}\right)^{2/3} \min_{|h| \leq h_{0}} \mathrm{E}\left\{\mathcal{E}_{c}(h)\right\}$$

## Second derivatives, $\frac{f(t_0+h)-2f(t_0)+f(t_0-h)}{h^2}$

$$|h_M| = \gamma_7 \left(\frac{\varepsilon_f}{\mu_M}\right)^{1/4}, \qquad \gamma_7 = 2^{5/8} \, 3^{1/8} \approx 2.33$$

$$\diamond \operatorname{E} \left\{ \mathcal{E}_{2}(h_{M}) \right\} \leq \left( \frac{\mu_{M}}{\mu_{L}} \right) \min_{|h| \leq h_{0}} \operatorname{E} \left\{ \mathcal{E}_{2}(h) \right\}$$

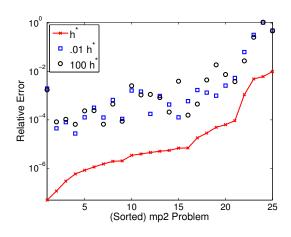
ullet use to obtain rough estimate of  $|f_s''|$  for forward-difference h



## Ex.- Highly Nonlinear MINPACK-2 Problems

## 25 problems, $n \le 64 \cdot 10^4$

♦ Accurate estimates obtained even when f" not constant



Compared with hand-coded derivative

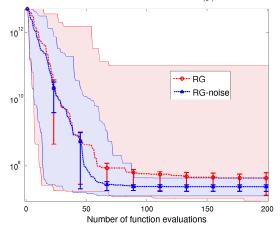


## Using the Noise in Nesterov's Random Gradient Method

#### General RG iteration

- 1. Generate direction  $d_k$
- 2. Evaluate gradient-free oracle  $g(x_k; h_k) = \frac{f(x_k + h_k d_k) f(x_k)}{h} d_k$
- 3. Compute  $x_{k+1} = x_k \delta_k g(x_k; h_k)$ , evaluate  $f(x_{k+1})$

## bicgstab quadratic: tol= $10^{-2}$ , $\frac{\varepsilon_f}{|f|} \approx$ 5e-3





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#### Matlab Time II

- Start playing around with stochastic algorithms
- Notice cost of getting fd parameter wrong
- Notice cost of not using derivatives



#### Summary: How Loud Are Your Functions?

- Computational noise complicates analysis of real-world functions, worst-case bounds overly pessimistic
- With a few (6-8) additional evaluations, ECNoise reliably estimates the noise
- Stochastic theory for near-optimal difference parameters
- Coarse estimates of |f"| (2-4 evaluations) yield more accurate directional derivatives
- Both work on deterministic functions in practice



[Estimating Computation Noise, SISC 2011] [Estimating Derivatives of Noisy Simulations, TOMS 2012] [Do You Trust Derivatives or Differences?, JCP 2014] [Obtaining Quadratic Models of Noisy Functions, Preprint, 2014]

Computing http://mcs.anl.gov/~wild/cnoise







Part II?

## Stochastic Methods for Two Types of Problems

- A. Stochastic optimization
  - Modeling and algorithms for optimization under uncertainty
  - Stochasticity from problem and/or algorithm
- B. Deterministic optimization
  - Objectives and constraints deterministic
  - Methods are "randomized"



## Stochastic Methods for Two Types of Problems

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- B. Deterministic optimization
  - Objectives and constraints deterministic
  - Methods are "randomized"
- → Methods and analysis are related



# A. Stochastic Optimization Problems and Methods

### Stochastic Optimization

#### General problem

$$\min\left\{f(x) = \mathbb{E}_{\xi}\left[F(x,\xi)\right] : x \in X\right\} \tag{1}$$

- $x \in \mathbb{R}^n$  decision variables
- $\diamond$   $\xi$  vector of random variables
  - ullet independent of x
  - $P(\xi)$  distribution function for  $\xi$
  - $\xi$  has support  $\Xi$
- $\diamond$   $F(x,\cdot)$  functional form of uncertainty for decision x
- $\diamond \ X \subseteq \mathbb{R}^n$  set defined by deterministic constraints
  - Also: stochastic/probabilistic constraints

(not addressed here)



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# Approach of Sampling Methods for $f(x) = \mathbb{E}_{\xi}\left[F(x,\xi)\right]$

- $\diamond$  Let  $\xi^1, \xi^2, \cdots, \xi^N \sim P$
- $\diamond$  For  $x \in X$ , define:

$$f_N(x) = \frac{1}{N} \sum_{i=1}^{N} F(x, \xi^i)$$

- $f_N$  is a random variable (really, a stochastic process) (depends on  $(\xi^1, \xi^2, \cdots, \xi^N)$ )
- Motivated by  $\mathbb{E}_{\xi}\left[f_{N}(x)\right]=f(x)$

**A** 

# Bias of Sampling Methods

$$\diamond$$
 Let  $f^* = f(x^*)$  for  $x^* \in X^* \subseteq X$ 

# Bias of Sampling Methods

- $\diamond$  Let  $f^* = f(x^*)$  for  $x^* \in X^* \subseteq X$
- $\diamond$  For any  $N \geq 1$ :

$$\mathbb{E}_{\xi} \left[ f_N^* \right] \le f^* = \mathbb{E}_{\xi} \left[ F(x^*, \xi) \right]$$

because

$$\mathbb{E}_{\xi}\left[f_1^*\right] = \mathbb{E}_{\xi}\left[\min\left\{F(x,\xi): x \in X\right\}\right] \leq \min\left\{\mathbb{E}_{\xi}\left[F(x,\xi)\right]: x \in X\right\} = f^*$$



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# Bias of Sampling Methods

- $\diamond$  Let  $f^* = f(x^*)$  for  $x^* \in X^* \subseteq X$
- For any  $N \geq 1$ :

$$\mathbb{E}_{\xi}\left[f_{N}^{*}\right] \leq f^{*} = \mathbb{E}_{\xi}\left[F(x^{*}, \xi)\right]$$

because

$$\mathbb{E}_{\xi} \left[ f_1^* \right] = \mathbb{E}_{\xi} \left[ \min \left\{ F(x, \xi) : x \in X \right\} \right] \le \min \left\{ \mathbb{E}_{\xi} \left[ F(x, \xi) \right] : x \in X \right\} = f^*$$

- $\diamond$  Sampling problems result in optimal values below  $f^*$
- $\diamond$   $f_N^*$  is biased estimator of  $f^*$



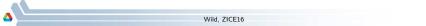
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# Sample Average Approximation

- $\diamond$  Draw realizations  $\hat{\xi}^1, \hat{\xi}^2, \cdots, \hat{\xi}^N \sim P$  of  $(\xi^1, \xi^2, \cdots, \xi^N)$
- Replace (1) with

$$\min\left\{\frac{1}{N}\sum_{i=1}^{N}F(x,\hat{\xi}^{i}):\ x\in X\right\} \tag{2}$$

- $\hat{f}_N(x) = \frac{1}{N} \sum_{i=1}^N F(x, \hat{\xi}^i)$  deterministic
- ullet Follows mean of the N sample paths defined by the (fixed)  $\hat{\xi}^i$



### SAA Algorithm

Input N, (maybe  $x^0 \in X$ )

- 1. Generate  $\hat{\xi}^1, \hat{\xi}^2, \cdots, \hat{\xi}^N \sim P$
- 2. Solve the deterministic problem

$$\min\left\{\frac{1}{N}\sum_{i=1}^N F(x,\hat{\xi}^i):\ x\in X\right\}$$

Output  $x_N^*$  (or  $X_N^*$ ).



# Convergence with ${\cal N}$

- A sufficient condition:
  - For any  $\epsilon > 0$  there exists  $N_{\epsilon}$  so that

$$\left| \hat{f}_N(x) - f(x) \right| < \epsilon \quad \forall N \ge N_{\epsilon} \quad \forall x \in X$$

with probability 1 (wp1).

- Then  $\hat{f}_N^* \to f^*$  wp1.
- $\diamond$  (With additional assumptions on f and  $X^* \subset X$ ):

$$\mathsf{dist}(x_N^*,X^*)\to 0$$

 $\diamond$  (+ uniqueness,  $X^* = x^*$ ):

$$x_N^* \to x^*$$

# Stochastic Approximation Method

### Basically just:

### Input $x^0$

1. 
$$x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k - \alpha_k s^k \right\}$$
,

$$k = 0, 1, \dots$$

- $\diamond$   $\alpha_k$  a step size
- $\diamond$   $s^k$  a random direction

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,

$$k = 0, 1, \dots$$

- $\diamond$   $\alpha_k$  a step size
- $\diamond$   $s^k$  a random direction

#### Generally assume:

$$\alpha_k$$
:  $\sum_{k=0}^{\infty} \alpha_k = \infty$ ,  $\sum_{k=0}^{\infty} \alpha_k^2 < \infty$ 

$$s^k$$
:  $\mathbb{E}\left\{\nabla f(x^k)^T s^k\right\} > 0$ 

 $\boldsymbol{s}^k$  is an ascent direction (in expectation) at  $\boldsymbol{x}^k$ 

# Stochastic Approximation Method

#### Basically just:

#### Input $x^0$

1. 
$$x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k - \alpha_k s^k \right\}$$
,

$$k=0,1,\ldots$$

- $\diamond \alpha_k$  a step size
- $\diamond$   $s^k$  a random direction

### Generally assume:

$$\alpha_k$$
:  $\sum_{k=0}^{\infty} \alpha_k = \infty$ ,  $\sum_{k=0}^{\infty} \alpha_k^2 < \infty$ 

$$s^k$$
:  $\mathbb{E}\left\{\nabla f(x^k)^T s^k\right\} > 0$ 

 $s^k$  is an ascent direction (in expectation) at  $x^k$ 

 $\diamond$  "Exact" Stochastic Gradient Descent:  $s^k = \nabla f(x^k)$ 



### Classic SA Algorithms

- "Original" method is Robbins-Monro (1951)
- Without derivatives: Kiefer-Wolfowitz (1952) replaces gradient with finite-difference approximation, e.g.,

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$$x^{k+1} \leftarrow x^k - \alpha_k s^k$$
,  $k = 0, 1, \dots$ 

where

$$s^k = \frac{F(x^k + h_k I_n; \hat{\xi}^k) - F(x^k - h_k I_n; \hat{\xi}^{k+1/2})}{2h_k}$$



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- Requires 2n evaluations every iteration
- Can appeal to variance reduction techniques (e.g., common RNs)
- Convergence  $x^k \to x^*$  if f strongly convex (near  $x^*$ ), usual conditions on  $\alpha_k$ ,  $h_k \to 0$ ,  $\sum_k \frac{\alpha_k^2}{h_i^2} < \infty$
- K-W recommend:  $\alpha_k = \frac{1}{k}$ ,  $h_k = \frac{1}{k^{1/3}}$



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- K-W recommend:  $\alpha_k = \frac{1}{k}$ ,  $h_k = \frac{1}{k^{1/3}}$
- Extensions such as SPSA (Spall) reduce number of evaluations (see randomized methods slides...)



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### Derivative-Based Stochastic Gradient Descent

- 1. Draw realization  $\hat{\xi}^k \sim P$  of  $\xi^k$
- 2. Compute  $s^k = \nabla_x F(x^k; \hat{\xi}^k)$
- 3. Update  $x^{k+1} \leftarrow \mathcal{P}_X \left\{ x^k \alpha_k s^k \right\}$



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- $\diamond \ \nabla_x F(x^k; \hat{\xi}^k)$  is an unbiased estimator for  $\nabla f(x^k)$



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  abla_x F(x^k; \hat{\xi}^k)$  is an unbiased estimator for  $abla f(x^k)$
- Can incorporate curvature if desired
  - e.g.,  $B^k s^k$  an unbiased estimator for  $\left(\nabla^2 f(x^k)\right)^{-1} \nabla f(x^k)$
- Can work with subgradients
- $\diamond$  Can even output  $x^N = \frac{1}{N} \sum_{k=1}^N x^k$



I, ZICE16 39

#### Modern Stochastic Gradient Descent Codes

#### Stochastic gradient descent seems inherently sequential

Better in special cases, e.g.,

$$f(x) = \sum_{e \in \mathcal{E}} f_e(x_e), \qquad e \subset \{1, \dots, n\}$$

 $|\mathcal{E}|$  and n large



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 $|\mathcal{E}|$  and n large

- HOGWILD! (Niu, Recht, Ré, Wright)
  - parallel, asynchronous implementation
  - http://i.stanford.edu/hazy/victor/Hogwild/



B. Randomized Algorithms for Deterministic Problems

# Randomized Algorithms for Deterministic Problems

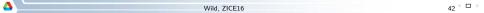
$$\min \left\{ f(x) : x \in X \subseteq \mathbb{R}^n \right\}$$

- f deterministic
- Random variables are now generated by the method, not from the problem
- Often assume properties of f
  - e.g.,  $\nabla f$  is L'-Lipschitz:

$$\|\nabla f(x) - \nabla f(y)\| \le L' \|x - y\| \qquad \forall x, y \in X$$

e.g., f is strongly convex (with parameter  $\tau$ ):

$$f(x) \ge f(y) + (x - y)^T \nabla f(y) + \frac{\tau}{2} ||x - y||^2 \forall x, y \in X$$



# Basic Algorithms

### Matyas (e.g., 1965):

- Input  $x^0$ ; repeat:
  - 1. Generate Gaussian  $u^k$  (centered about 0) 2. Evaluate  $f(x^k+u^k)$

3. 
$$x^{k+1} = \begin{cases} x^k + u^k & \text{if } f(x^k + u^k) < f(x^k) \\ x^k & \text{otherwise.} \end{cases}$$



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#### Poljak (e.g., 1987)

- $\diamond$  Input  $x^0$ ,  $\{h_k, \mu_k\}_k$ ; repeat:
  - 1. Generate a random  $u^k \in \mathbb{R}^n$

2. 
$$x^{k+1} = x^k - h_k \frac{f(x^k + \mu_k u^k) - f(x^k)}{\mu_k} u^k$$

- $h_k > 0$  is the step size
- $\mu_k > 0$  is called the smoothing parameter



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Componentwise Lipschitz parameter M > 0:

$$|\nabla_i f(x + he_i) - \nabla_i f(x)| \le M|h|, \quad \forall h \in \mathbb{R}, \quad i = 1, \dots, n$$

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- 1. Choose  $i_k = \arg \max_{i=1,...,n} |\nabla_i f(x^k)|$
- 2. Update  $x^{k+1} = x^k \frac{1}{M} \nabla_{i_k} f(x^k) e_{i_k}$

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- $\diamond$  Generates  $f(x^k) f^* \leq \frac{2nMR^2}{k+4}$ , where  $R \geq \|x^0 x^*\|$



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- $\diamond$  Good: only updates  $x_{i_k}$
- $\diamond$  Bad: requires entire gradient  $\nabla f(x^k)$



Component-wise Lipschitz parameter M>0:

$$|\nabla_i f(x + he_i) - \nabla_i f(x)| \le L_i |h|, \quad \forall h \in \mathbb{R}, \quad i = 1, \dots, n$$

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- Generates  $E\{f(x^k)\} f^* \leq \frac{2nR_1^2}{k+4}$ , where  $R_1 = \max\{\|x - x^*\|_1 : f(x) < f(x_0)\}$
- $\diamond$  Good: only updates  $x_{i_k}$
- $\diamond$  Better: requires only component  $i_k$  of gradient  $\nabla f(x^k)$
- Can also:
  - generate  $i_k$  proportional to coordinate Lipschitz parameters  $\{L_i\}_i$
  - perform block-coordinate (and other subspace) operations



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# Gaussian Smoothing

- $\diamond$  Let  $f: \mathbb{R}^n \to \mathbb{R}$  be deterministic
- $\circ \ u \in \mathbb{R}^n$  from a Gaussian distribution,  $\mathbb{E}_u\left[u\right] = 0$ 
  - Here: Covariance matrix  $I_n$ , general C OK
- $\diamond$  For scalar  $\mu > 0$ , Gaussian-smoothed version of f:

$$f_{\mu}(x) = \mathbb{E}_{u} \left[ f(x + \mu u) \right]$$

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- ♦ For scalar  $\mu > 0$ , Gaussian-smoothed version of f:

$$f_{\mu}(x) = \mathbb{E}_{u} \left[ f(x + \mu u) \right]$$

- If f is convex, then  $f_{\mu}(x) \geq f(x)$
- If f is convex and  $\nabla f$  is L'-Lipschitz, then

$$|f_{\mu}(x) - f(x)| \le \frac{\mu^2}{2} L' n$$



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# Gaussian Smoothing and Directional Derivatives

$$f_{\mu}(x) = \mathbb{E}_{u} \left[ f(x + \mu u) \right]$$

 $\diamond$  Derivative of f in the direction u:  $f'(x;u) = \lim_{h \downarrow 0} \frac{f(x+hu) - f(x)}{h}$ 



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- $q_0(x) = f'_u(x)u$ 
  - If f is convex, then  $\mathbb{E}_u\left[g_0(x)\right]$  is a subgradient of f
  - If f is differentiable at x, then

$$\mathbb{E}_{u}\left[\|g_{0}(x)\|^{2}\right] \leq (n+4)\|\nabla f(x)\|^{2}$$



47 □ ▶ Wild, ZICE16

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- - If f is differentiable at x, then  $\mathbb{E}_u\left[g_\mu(x)\right] = \nabla f_\mu(x)$
  - ullet If f is differentiable at x and  $\nabla f$  is L'-Lipschitz, then

$$\mathbb{E}_{u}\left[\|g_{\mu}(x)\|^{2}\right] \leq 2(n+4)\|\nabla f(x)\|^{2} + \frac{\mu^{2}}{2}L'^{2}(n+6)^{3}$$

Wild, ZICE16

#### Random Gradient Method

Input  $x^0 \in X$ ,  $\{h_k\}_k$ ; repeat:

- 1. Generate Gaussian  $u^k \in \mathbb{R}^n$  and compute  $g_0(x^k) = f_{u^k}'(x^k)u^k$
- 2.  $x^{k+1} = \mathcal{P}_X \left\{ x^k h_k g_0(x^k) \right\}$

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- ♦ Key result (Nesterov) for convex (but possibly nonsmooth) f: For fixed  $h_k = \frac{R}{\sqrt{n+4}\sqrt{N+1}L}$  and any  $\epsilon > 0$ ,

$$\mathbb{E}_u\left[f(\hat{x}^N)\right] - f^* \le \epsilon, \quad \text{where} \quad \hat{x}^N = \arg\min_{i=1,\dots,N} f(x^i)$$

- in  $\mathcal{O}\left(\frac{n}{\epsilon^2}\right)$  iterations
- Also works for convex stochastic optimization and convex smooth f (with improved bounds and rates)



#### Random Gradient-Free Method

Input  $x^0 \in X$ ,  $\mu > 0$ ,  $\{h_k\}_k$ ; repeat:

- 1. Generate Gaussian  $u^k \in R^n$  and compute  $g_\mu(x^k) = \frac{f(x^k + u^k) f(x^k)}{\mu} u^k$
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- $^{\diamond}$  Key result (Nesterov) for convex (but possibly nonsmooth) f: For fixed  $h_k = \frac{R}{(n+4)\sqrt{N+1}L}$ ,  $\mu = \frac{\epsilon}{2L\sqrt{n}}$ , and any  $\epsilon > 0$ ,

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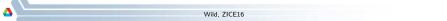
#### Accelerated Random Gradient-Free Method

f strongly convex (with convexity parameter au)

Input  $v^0=x^0$ ,  $\mu>0$ ,  $\gamma_0\geq \tau$ ,  $\{h_k\}_k$ ; repeat:

- 1. Obtain  $\alpha_k > 0$  satisfying  $16(n+1)^2 L' \alpha_k^2 = (1-\alpha_k)\gamma_k + \tau \alpha_k$
- 2. Set  $\gamma_{k+1}=(1-\alpha_k)\gamma_k+\tau\alpha_k$ ,  $\lambda_k=\frac{\alpha_k\tau}{\gamma_{k+1}}$ ,  $\beta_k=\frac{\alpha_k\gamma_k}{\gamma_k+\alpha_k\tau}$
- 3. Set  $y^k = (1 \beta_k)x^k + \beta_k v^k$
- 4. Generate Gaussian  $u^k \in R^n$  and compute  $g_\mu(y^k) = \frac{f(y^k + u^k) f(y^k)}{\mu} u^k$
- 5. Update

$$\begin{array}{ll} x^{k+1} = & y^k - \frac{1}{4(n+4)L'} g_{\mu}(y^k) \\ v^{k+1} = & (1 - \lambda_k) v^k + \lambda_k y^k - \frac{1}{16(n+1)^2 L' \alpha_k} g_{\mu}(y^k) \end{array}$$



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 $\diamond$  Key result (Nesterov): for  $\tau=0$  functions  $\exists \mu>0$  so that

$$\mathbb{E}_u\left[f(\hat{x}^N)\right] - f^* \leq \epsilon, \qquad \text{where} \quad \hat{x}^N = \arg\min_{i=1,\dots,N} f(x^i)$$

in  $\mathcal{O}\left(\frac{n}{\epsilon^{1/2}}\right)$  iterations

Wild, ZICE16 50

# Applying SA-Like Ideas to Special Cases

$$\min \left\{ f(x) = \frac{1}{m} \sum_{i=1}^{m} F_i(x) : x \in X \right\}$$

$$m \text{ huge}$$



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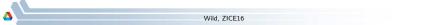
Ex.- Nonlinear Least Squares

 $F_i(x) = \|\phi(x; \theta^i) - d^i\|^2$ 

Warning: likely nonconvex!

Evaluating  $\phi(\cdot,\cdot)$  requires solving a large PDE

Ex.- Sample Average Approximation  $F_i(x) = R(x; \xi^i)$  $\hat{\xi}^i \in \Omega$  a scenario/RV realization (and R depends nontrivially on  $\hat{\xi}^i$ )



# Applying SA-Like Ideas to Special Cases

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 $\hat{\xi}^i \in \Omega$  a scenario/RV realization  
(and  $R$  depends nontrivially on  $\hat{\xi}^i$ )

The good:

$$\diamond \nabla f(x) = \sum_{i=1}^{m} \nabla F_i(x)$$

The bad:

 $\diamond \ m$  still huge

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$$\min \left\{ f(x) = \frac{1}{m} \sum_{i=1}^{m} F_i(x) : x \in X \right\}$$

" $F_i(x)$  is a member of a population of size m"

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- " $F_i(x)$  is a member of a population of size m"
  - $\diamond$  Randomly sample  $\mathcal{S}$ , a subset of size  $|\mathcal{S}|$ , from  $\{1,\cdots,m\}$

Δ

$$\min \left\{ f(x) = \frac{1}{m} \sum_{i=1}^{m} F_i(x) : x \in X \right\}$$

" $F_i(x)$  is a member of a population of size m"

- $\diamond$  Randomly sample S, a subset of size |S|, from  $\{1, \dots, m\}$
- Under minimal assumptions:

$$\mathrm{E}\left\{\frac{1}{|\mathcal{S}|}\sum_{i\in\mathcal{S}}F_i(x)\right\} = f(x) \qquad \text{and} \qquad \mathrm{E}\left\{\frac{1}{|\mathcal{S}|}\sum_{i\in\mathcal{S}}\nabla F_i(x)\right\} = \nabla f(x)$$



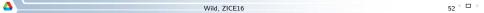
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 $\diamond$  Use  $-\nabla f_{\mathcal{S}} = -\frac{1}{|\mathcal{S}|} \sum_{i \in \mathcal{S}} \nabla F_i(x)$  as direction  $s^k$ 



$$\min \left\{ f(x) = \frac{1}{m} \sum_{i=1}^{m} F_i(x) : x \in X \right\}$$

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- $\diamond$  Randomly sample S, a subset of size |S|, from  $\{1, \dots, m\}$
- Under minimal assumptions:

$$\mathrm{E}\left\{\frac{1}{|\mathcal{S}|}\sum_{i\in\mathcal{S}}F_i(x)\right\} = f(x) \qquad \text{and} \qquad \mathrm{E}\left\{\frac{1}{|\mathcal{S}|}\sum_{i\in\mathcal{S}}\nabla F_i(x)\right\} = \nabla f(x)$$

- Use  $-\nabla f_{\mathcal{S}} = -\frac{1}{|\mathcal{S}|} \sum_{i \in \mathcal{S}} \nabla F_i(x)$  as direction  $s^k$
- ♦ How to choose S?

$$\mathbb{E}\left\{\left\|\nabla f_{\mathcal{S}_n} - \nabla f\right\|^2\right\} = \left(1 - \frac{|\mathcal{S}|}{m}\right) \mathbb{E}\left\{\left\|\nabla f_{\mathcal{S}_r} - \nabla f\right\|^2\right\}$$

 $\Rightarrow$  sampling without replacement  $(S_n)$  gives lower variance than does sampling with replacement  $(S_r)$ 

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#### Summary

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- + Incredibly simple to code basic implementation
- + Well-studied complexity bounds, especially for convex cases; can show that asymptotic rates are optimal
- + Even useful when gradient/subgradient unavailable



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#### Summary

- Methods for stochastic optimization and randomized methods for deterministic optimization closely related
- + Incredibly simple to code basic implementation
- + Well-studied complexity bounds, especially for convex cases; can show that asymptotic rates are optimal
- + Even useful when gradient/subgradient unavailable
- Bounds and parameters depend on characteristics of function (e.g., Lipschitz parameters, level set diameters, strong convexity)
- (Some) Practitioners remain nervous about performance deviations from the mean (active research area)

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