

COMP20003 Workshop Week 7

- 1** Sorting, Insertion Sort and Selection Sort revisited
- 2** Quicksort + Q 6.1
- 3** Review for the Test / sample test
- 4** Implementing Hash Table (P6.1)

MST:
Week 8

Sorting Algorithms

Any problem with:

- Selection Sort?
- Insertion Sort?
- Quick Sort?

Properties of sorting algorithms:

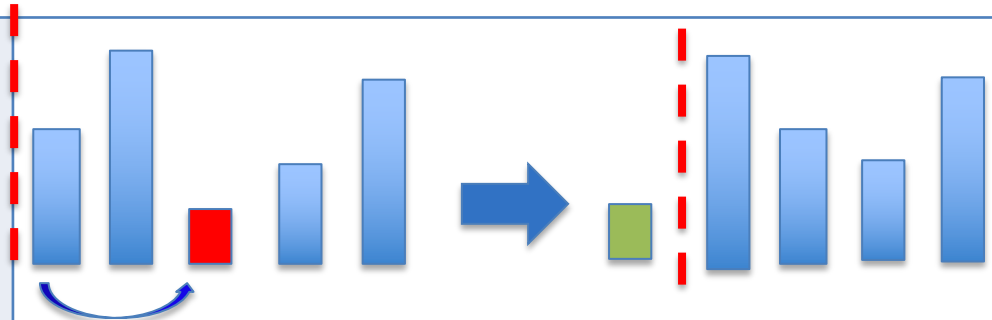
- *in-place*
- *stable*

Note: *In this workshop, suppose we need to sort an array in non-decreasing order.*

Selection Sort: $n=5$, (by selecting the smallest)

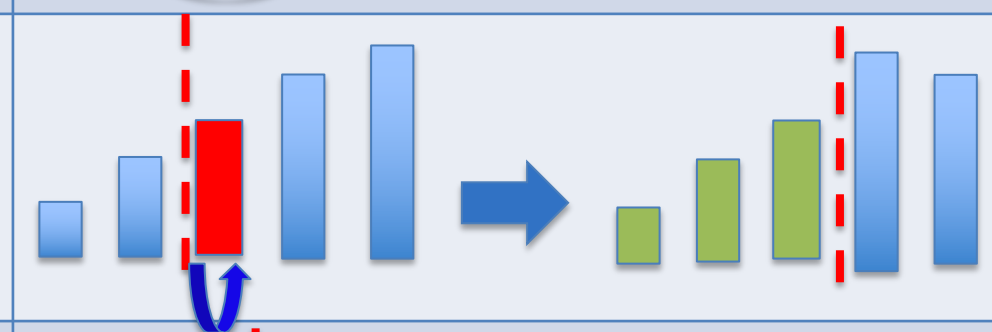
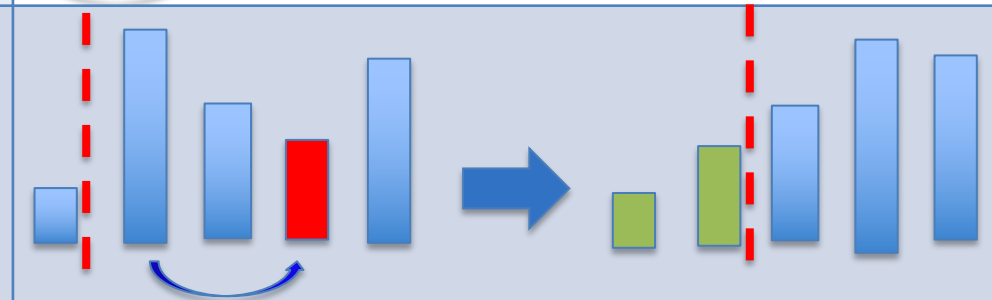
Round 1: consider $A[0..4]$

- *determine position of the smallest*
- *swap with the first, ie. $A[0]$*



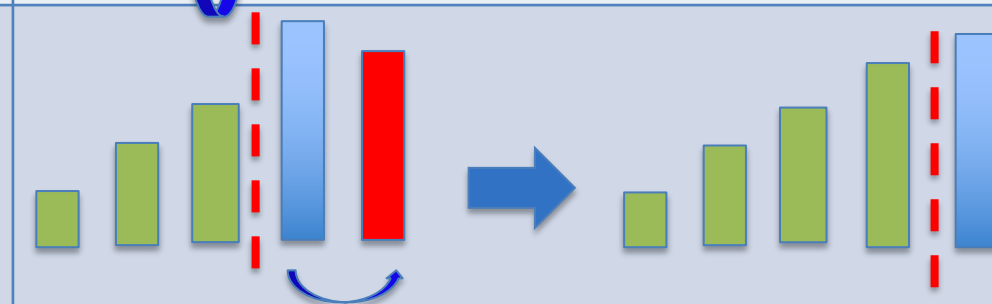
Round 2: consider $A[1..4]$

- *determine position of the smallest*
- *swap with the first, ie. $A[1]$*



Round 4: consider $A[3..4]$

- *determine position of the smallest*
- *swap with the first, ie. $A[3]$*



Selection Sort

-	A N A L Y S I S
1	
2	
3	
4	
5	
6	
7	

- i. Run the algorithm on the input array: [A N A L Y S I S]
- ii. What is the time complexity of the algorithm?
- iii. Is the sorting algorithm stable?
- iv. Is the algorithm in-place?

Selection Sort

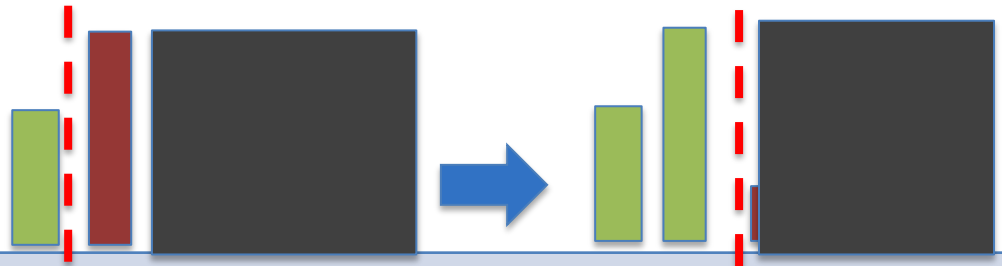
-	A	N	A	L	Y	S	I	S
1	A	N	A	L	Y	S	I	S
2	A	A	N	L	Y	S	I	S
3	A	A	I	L	Y	S	N	S
4	A	A	I	L	Y	S	N	S
5	A	A	I	L	N	S	Y	S
6	A	A	I	L	N	S	Y	S
7	A	A	I	L	N	S	S	Y

- Run the algorithm on the input array: **[A N A L Y S I S]**
- What is the time complexity of the algorithm?
- Is the sorting algorithm stable?
- Is the algorithm in-place?

Insertion Sort: understanding (n=5)

Round 1: consider $A[1]$

- $A[0..0]$ is sorted
- insert $A[1]$ to the left so that $A[0..1]$ is sorted



Round 2: consider $A[2]$

- $A[0..1]$ is sorted
- insert $A[2]$ to the left so that $A[0..2]$ is sorted



Round i: consider $A[4]$

- $A[0..4]$ is sorted
- insert $A[i]$ to the left so that $A[0..4]$ is sorted



Insertion Sort

	A N A L Y S I S
1	
2	
3	
4	
5	
6	
7	

- i. Run the algorithm on the input array: [A N A L Y S I S]
- ii. What is the time complexity of the algorithm?
- iii. Is the sorting algorithm stable?
- iv. Does the algorithm sort in-place?

Insertion Sort

	A		N	A	L	Y	S	I	S
1	A	N		A	L	Y	S	I	S
2	A	A	N		L	Y	S	I	S
3	A	A	L	N		Y	S	I	S
4	A	A	L	N	Y		S	I	S
5	A	A	L	N	S	Y		I	S
6	A	A	I	L	N	S	Y		S
7	A	A	I	L	N	S	S	Y	

- i. Run the algorithm on the input array: [A N A L Y S I S]
- ii. What is the time complexity of the algorithm?
- iii. Is the sorting algorithm stable?
- iv. Does the algorithm sort in-place?

Quicksort (usage: `Quicksort(A[0..n-1])`)

```
function QUICKSORT( $A[l..r]$ )  
  if  $l < r$  then  
     $s \leftarrow \text{PARTITION}(A[l..r])$   
    QUICKSORT( $A[l..s-1]$ )  
    QUICKSORT( $A[s+1..r]$ )
```

`Partition($A[1..r]$)`

$\leq A[s]$ $\geq A[s]$

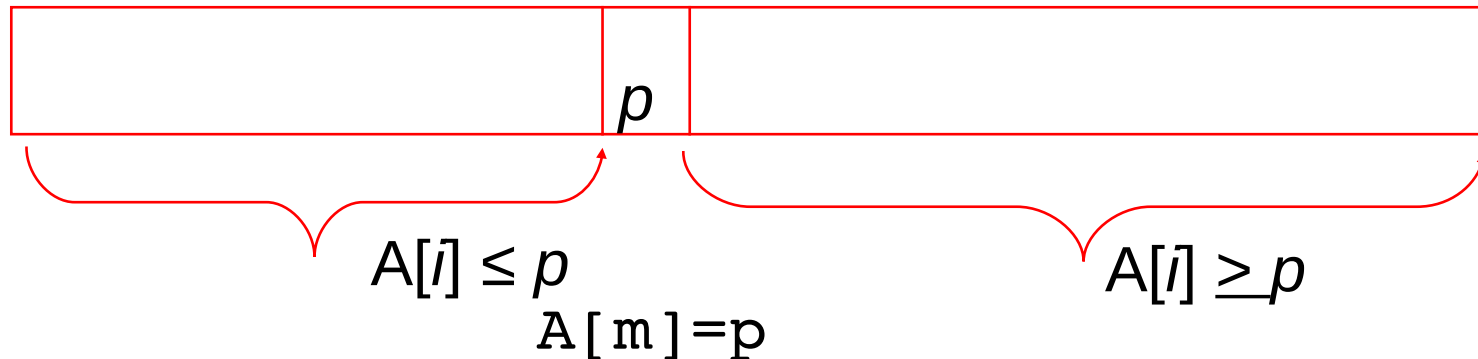
 $A[1..s-1]$ s $s+1..r$]

return s

Partitioning (using the rightmost element as pivot)

Input: Given an unsorted slice $A[1..r]$,

Output: A value m and re-arrangement of A so that:





Algorithm:


1. set $p = A[r]$ ($= \text{rightmost}$), keep $A[r]$ unchanged
2. Start with $i = 1, j = r - 1$
3. Moving i forward, stop when $A[i] \geq p$
4. Moving j backward, stop when $A[j] \leq p$
5. If $i < j$: swap $A[i], A[j]$ then go back to step 3
6. Swap $A[r]$ and $A[i]$, return $m = i$

Partitioning (using the rightmost element as pivot)

1. set $p = A[r]$ (=rightmost), keep $A[r]$ un`changed
2. Start with $i = l, j = r - 1$
3. Moving i forward, stop when $A[i] \geq p$
4. Moving j backward, stop when $A[j] \leq p$
5. If $i < j$: swap $A[i], A[j]$ then go back to step 3
6. Swap $A[r]$ and $A[i]$, return $m = i$

Example

start 1 2 3 4 5 6 7
  

moved 1 2 3 4 5 6 7
  

$i > j$: no swap

Step 6: 1 2 3 4 5 6 7 (swap 7 with 7)

 [1 2 3 4 5 6] 7 []

1. set $p = A[r]$ ($= \text{rightmost}$), keep $A[r]$ unchanged
2. Start with $i = l, j = r - 1$
3. Moving i forward, stop when $A[i] \geq p$
4. Moving j backward, stop when $A[j] \leq p$
5. If $i < j$: swap $A[i], A[j]$ then go back to step 3
6. Swap $A[r]$ and $A[i]$, return $m = i$

Example

start

A N A L Y S N

moved

A N A L Y S N

$i < j$: swap

A L A N Y S N

moved:

A L A N Y S N

$i \geq j$: stops looping

Step 6: swap N with N, return $m = i = 3$

A L A N Y S N

answer=3, $[\text{left}]m[\text{right}] = [A \ L \ A] \ N \ [Y \ S \ N]$

Q 6.1

You are asked to show the operation of quicksort on the following keys. For simplicity, use the rightmost element as the partition element:

2 3 97 23 15 21 4 23 29 37 5 23

Comment on the stability of quicksort and its behavior on almost sorted inputs.

Quiz 1

The best case of Selection sort is:

a. $O(n \log n)$.

b. $O(n)$.

c. $O(n^2)$.

d. $O(\log n)$.

When?

Quiz 2

The best case of Insertion Sort is:

a. $O(n \log n)$.

b. $O(n)$.

c. $O(n^2)$.

d. $O(\log n)$.

When?

Quiz 3

The average case of Insertion Sort is:

a. $O(n \log n)$.

b. $O(n)$.

c. $O(n^2)$.

d. $O(\log n)$.

Quiz 4

The average case of Quick Sort is:

a. $O(n \log n)$.

b. $O(n)$.

c. $O(n^2)$.

d. $O(\log n)$.

Quiz 5

The big-O complexity of Quick Sort is:

a. $O(n \log n)$.

b. $O(n)$.

c. $O(n^2)$.

d. $O(\log n)$.

Sorting algorithms: complexity; being stable, in-place

	Selection	Insertion	Quick	Merge (top-down)
Basic Idea	For $A[0..n-1]$ Identify the smallest and swap it with $A[0]$. Repeat for $A[1..n-1]$ and so on.	From 2 nd element, insert it to the left sub-array so that the extended left sub-array remains sorted.	Choose a pivot, partition array into a <i>lesser</i> and a <i>greater</i> (than pivot) halves. Do recursively with each half.	Split to equal-size halves, sort them then merge them.
Complexity				
Best case				
Worst case				
Average				
In-place?				
Stable?				

Sorting algorithms: complexity; being stable, in-place

	Selection	Insertion	Quick	Merge (top-down)
Basic Idea	For A[0..n-1] Identify the smallest and swap it with A[0]. Repeat for A[1..n-1] and so on.	From 2 nd element, insert it to the left sub-array so that the extended left sub-array remains sorted.	Choose a pivot, partition array into a <i>lesser</i> and a <i>greater</i> (than pivot) halves. Do recursively with each half.	Split to equal-size halves, sort them then merge them.
Complexity	$\theta(n^2)$	$O(n^2)$	$O(n^2)$	
Best case	$O(n^2)$	$O(n)$	$O(n \log n)$	
Worst case	$O(n^2)$	$O(n^2)$	$O(n^2)$	
Average	$O(n^2)$	$O(n^2)$	$O(n \log n)$	
In-place?				
Stable?				

Sorting

Concepts: sorting, stable sort, in-place sort

Some sorting algorithm (on an array of n elements):

	Selection	Insertion	Quick	Merge
Basic Idea	Identify the smallest and swap it with the first...	From 2 nd element, insert it to the left sub-array so that the extended left sub-array remains sorted.	Choose a pivot, partition array into a <i>lesser</i> and a <i>greater</i> (than pivot) halves...	Split to equal-size halves, sort them then merge them.
Best case	$O(n^2)$	$O(n)$	$O(n \log n)$	$O(n \log n)$
Worst case	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(n \log n)$
Average	$O(n^2)$	$O(n^2)$	$O(n \log n)$	$O(n \log n)$
In-place?				
Stable?				

Assignment 2: Report

A report should

- be concise : <2 pages including table/graphs
- be clear as instructed in specs
- have:
 - proper introduction (purpose of the report)
 - expts description (data, queries, outcome)
 - expts outcome (graphs/table)
 - discussion: compare with theory, compare stage 1 and 2
 - short conclusion

Experiments

What to show? How?

What can affect the performance?

How to (automatically) organize experiments? Create data files? Accumulate the results? Build table/graphs?

?

Experiments

n= 1000 2000 4000 8000 16000 ?

n= 3000 6000 9000 12000 15000 18000 ?

Data: sorted, random [no need median]

query: run 100 queries each times and compute average?

Stage 1 vs stage 2?

stage	data_type	n	comparison/query
1	sortx	1000	???
1	sortx	2000	???
...			
1	rand	1000	???

Creating data file

1. Create a sorted data file with no header line

```
tail -n +2 CLUEdata2018_sortx.csv > sortx_all.csv
```

2. Create 2 data files of 1000 lines:

```
head -n 1000 sortx_all.csv > sortx_1000.csv
```

```
shuf sortx_1000.csv > rand_1000.csv
```

Repeat step 2 for all the sizes you want. You can also run the following single-line command:

```
for size in 1000 10000 100000;  
do head -n $size sortx_all.csv > sortx_$size.csv;  
shuf sortx_$size.csv > rand_$size.csv; done
```

Note: The above 3 parts are in a single line, if you want to break into sub-lines, add a single slash \ to the end of each sub-line.

The data files don't have the header line. It's OK for the experiment purpose.

Creating query files

Create 2 query files of 100 lines each (q1.txt and q2.txt for Stage 1 and Stage 2) from sortx_all.csv :

```
cat sortx_all.csv  
| awk -F ',' '{if (NF==11) print $9,$10}'  
| shuf | head -n 100  
| awk '{printf("%lf %lf\n", $1, $2)}'  
  
> q1.txt
```

For Stage 2 queries, change the last 2 components to

```
| awk '{printf("%lf %lf 0.0005\n", $1, $2)}'  
  
> q2.txt
```

Note: each command is in a single line, if you want to break into sub-lines, add a single slash \ to the end of each sub-line.

Summing up the output of you program

Suppose that the output into stdout of Stage 1 are *only* of format:

```
144.959522 -37.800095 --> 4000
0 0 --> 300
```

And we run, say:

```
./dict1 sortx-1000 o.txt > Slout_sortx_1000.txt
```

Then (**\$4** for “column 4”, which is the number 4000 and 300 in the top 2 lines):

```
cat Slout_sortx_1000.txt | awk 'BEGIN{sum=0; n=0}{sum = sum+$4;
n++}END{print "1000 " sum/n}' >> Stage1_sortx.txt
```

will *append* line

```
1000 2150
```

to file `Stage1_sortx.txt`. This line represents:

```
n      average_cmp_per_search
```

Note:

- Make sure that your output to stdout is in correct format: `0<SPACE>0<SPACE>--><SPACE>300`
- For Stage 2 output, we need to relace **\$4** by **\$5** because the output has format where num_cmp is in column 5:

```
144.959522 -37.800095 0.0005 --> 4000
```

SUMMARY: Process for expts with sizes 100 200 300

Creating a sorted all-record data file:

```
tail -n +2 CLUEdata2018_sortx.csv > sortx_all.csv
```

Creating data files for experiments with n= **100 200 300** (you should change these to your designed values such as : **3000 6000 9000 12000 15000 18000**)

```
for size in 100 200 300; do head -n $size sortx_all.csv >  
sortx_$size.csv; shuf sortx_$size.csv > rand_$size.csv; done
```

Run experiments (warnings: it might take long, depending on how fast is your code):

```
for map in 1 2; do for size in 100 200 300; do for type in  
sortx rand; do ./map$map $type-$size.csv out-$map-$type-  
$size.txt < q$map.txt > map$map-$type-$size.txt; done ; done;  
done
```

Summary all expt outcome into file **expts.txt**

```
for map in 1 2; do for type in sortx rand; do for size in 100  
200 300; do cat map$map-$type-$size.txt | awk -v size=$size -  
v map=$map -v type=$type 'BEGIN{sum=0; n=0}{sum = sum+$4;  
n++}END{print "map"map, type, size, sum/n}' >> expts.txt ;  
done ; done; done
```


On Canvas

2015 Mid-sem

Warning: Way too hard!

2017 Mid-sem

2019 Mid-sem

P6.1 Implementing a simple hash table, OR group work:

Choices:

- Implementing a simple hash table
- Individual working on ass2
- Group work with sample MST questions and/or programming tasks, for example:
 - Big-O questions
 - Hashing examples
 - AVL rotations
 - Programming: dynamic arrays and strings
 - Programming: linked lists