

# COMP20003 Workshop Week 10

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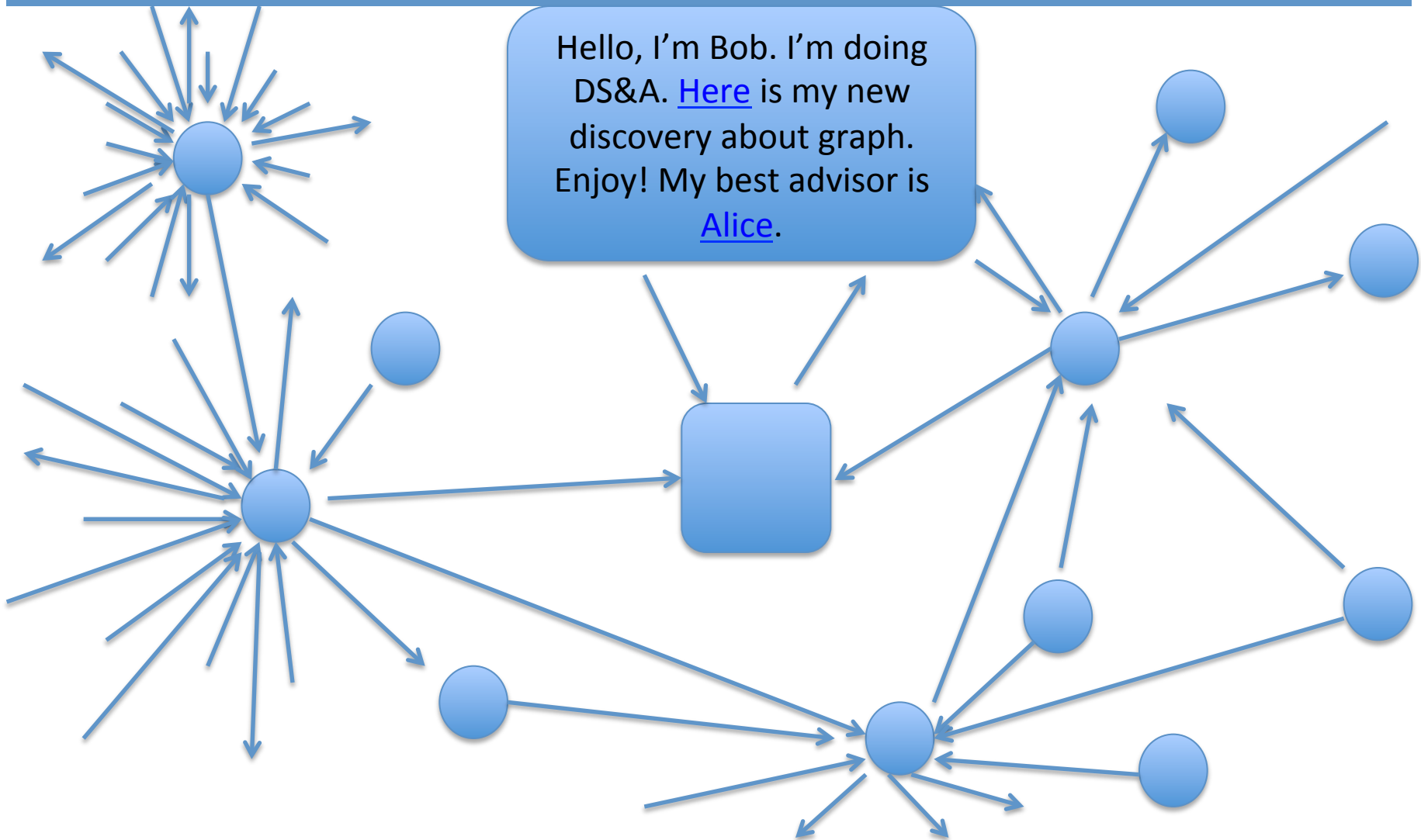
# Graphs

Visualization

Any problem where graph can be applied?

Are link lists and trees graphs? What special about them?

# PageRank: Bob becoming famous!

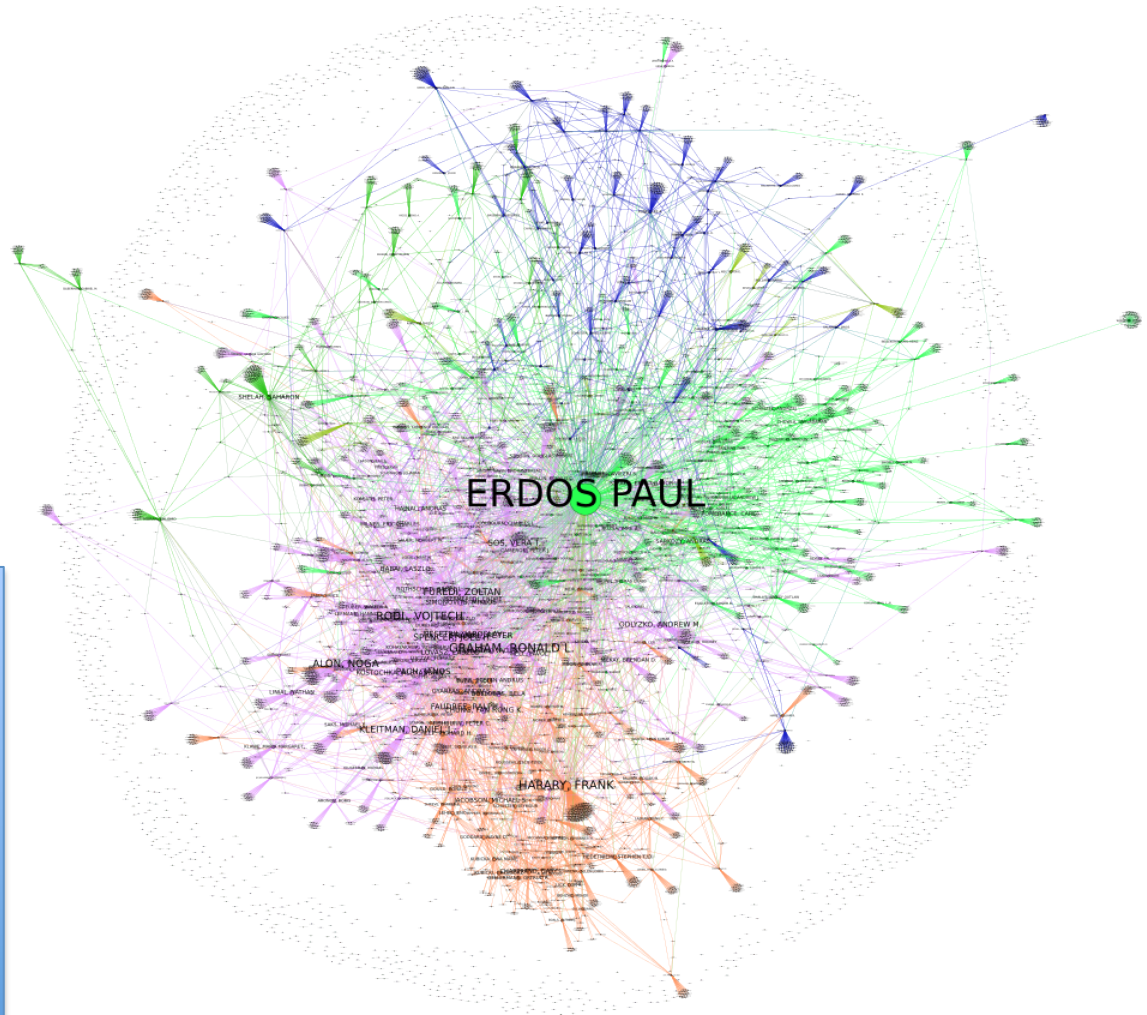


# Erdős' Collaborators Graph (6927 nodes)

(<http://genedan.com/tag/erdos-network-graph/>)



Erdos (1913-1996) is  
a great math-cian.  
He wrote ~1500  
papers, co-authored  
with >500 authors...



# Graphs: Concepts

Formal definition:  $G = (V, E)$  where

$V = \{v_i\}$  : set of *vertices*, or *nodes*

$E = \{(v_i, v_j) \mid v_i \in V, v_j \in V\}$  : set of *edges*, *arcs*, or *links*;

$|V|$  is called the *order* of the graph

$|E|$  is called the *size* of the graph

*dense* and *sparse* graphs

*directed*, *di-graph*, *undirected*, *acyclic*, *DAG*

*connected graph*, *connected component*

*weakly and strongly connected components*

# Graphs: Representation

Representation (general approaches):

1. *Adjacency matrix* (what is that?)
2. Set of *adjacency lists* (what is that?)

What is a suitable representation method for:

- a graph of this class, where nodes represent students, edge  $(a,b)$  means “a knows b”,
- the Erdős’s collaboration graph (check your answer by visiting the website),
- the webgraph?

# Graphs: DFS & BFS



# Breath-first Search (BFS): visit all neighbors first





# BFS from a single vertex

The task:

- Given a weighted graph  $G=(V,E)$ , and  $s \in V$
- Visit all vertices which are reachable from  $s$ .

The algorithm:

```
enqueue(s, Q)
while (Q is not empty) {
    x= dequeue(Q)
    visit x
    for all v that  $(x, v) \in E$ : enqueue(v, Q)
}
```

# Dijkstra's Algorithm

The task:

- Given a weighted graph  $G=(V,E,w(E))$ , and  $s \in V$ , and supposing that *all weights are positive*.
- Find shortest path (path with min total weight) from  $s$  to all other vertices.

# QoCT: do it now

Please do Tutor Quality of Casual Teaching (QoCT) survey (right now!) for COMP20003 by:

Goto Link:

<https://apps.eng.unimelb.edu.au/casmas/index.php?r=qoct/feedback&subjCode=COMP20003>

OR visit the link from googling:

[unimelb tutor survey comp20003](#)

# Dijkstra's Algorithm as a (special) BFS

Basic idea

if  $A \rightarrow B \rightarrow C \rightarrow D$  is a shortest path

then

$A \rightarrow B \rightarrow C$  is a shortest path (from  $A$  to  $C$ )

$A \rightarrow B$  is a shortest path (from  $A$  to  $B$ )

# Dijkstra's Algorithm as a (special) BFS

Basic idea

if  $s \rightarrow A \rightarrow B$  is a shortest path then  $s \rightarrow A$  is a shortest path.

Init:

- start with  $dist[s] = 0$ , and  $dist[*] = \infty$ , set  $unvisited\_set = V$

Round 1:

- choose node with min  $dist[]$ , which is  $s$ ;
- visit all nodes  $u$  adjacent to  $s$  and update  $dist[u]$ ;
- mark  $s$  as visited (remove it from the  $unvisited\_set$ );

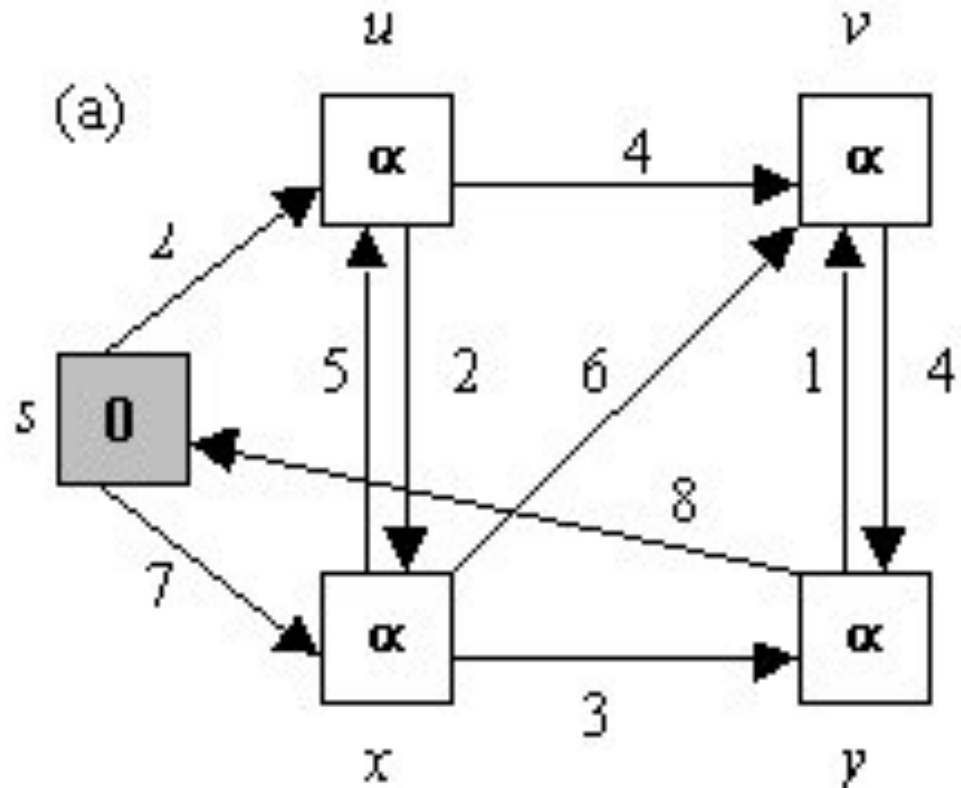
Round 2:

- choose the node with min  $dist[]$  from  $unvisited\_set$
- do the other steps as in Round 1

# Dijkstra algorithm from vertex s: initialization

In each round, one element will be removed from A and added to B.  
Repeat round until A become empty.

B= { }



A= V = { s, u, v, x, y }

dist = { **0**,  $\infty$ ,  $\infty$ ,  $\infty$ ,  $\infty$  }

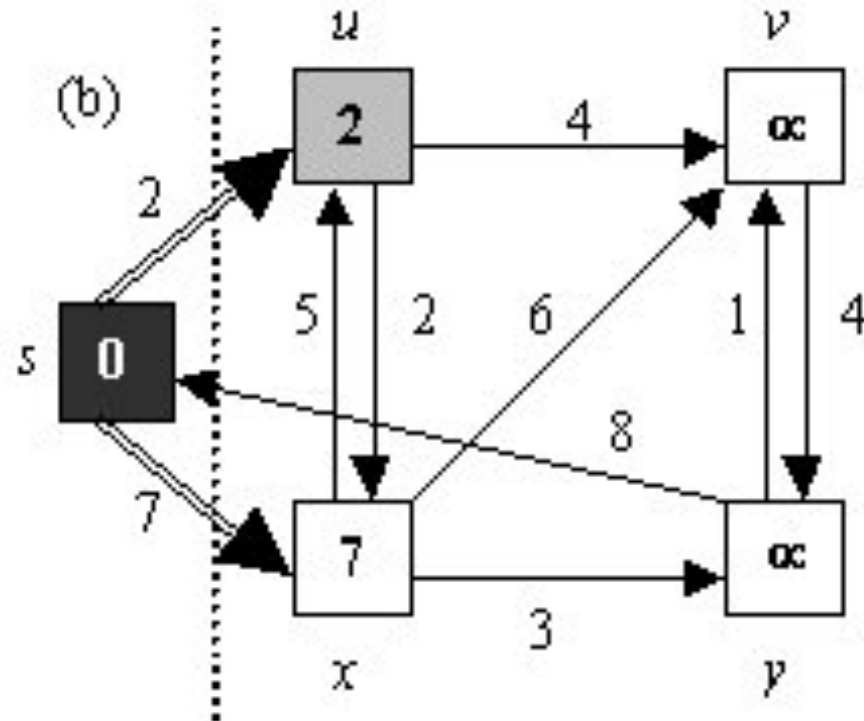
prev = { s, nil, nil, nil, nil }

visited= { 0, 0, 0, 0, 0 }

# Dijkstra algorithm from vertex s

Round 1:

- choose s (node with least dist[] in A)
- update dist[] of nodes adjacent to s *if applicable*



$B = \{s\}$

$\text{dist} = \{0\}$

$A = \{ \quad, u, v, x, y \}$

$\text{dist} = \{ \quad, \mathbf{2}, \infty, 7, \infty \}$

$\text{prev} = \{ \text{nil}, s, \quad, s, \quad \}$



# Dijkstra algorithm from vertex s

Round 1:

	s	u	v	x	y
0	0 s	$\infty$ N	$\infty$ N	$\infty$ N	$\infty$ N
1		2 s	$\infty$ N	7 s	$\infty$ N

visited=1    dist    prev

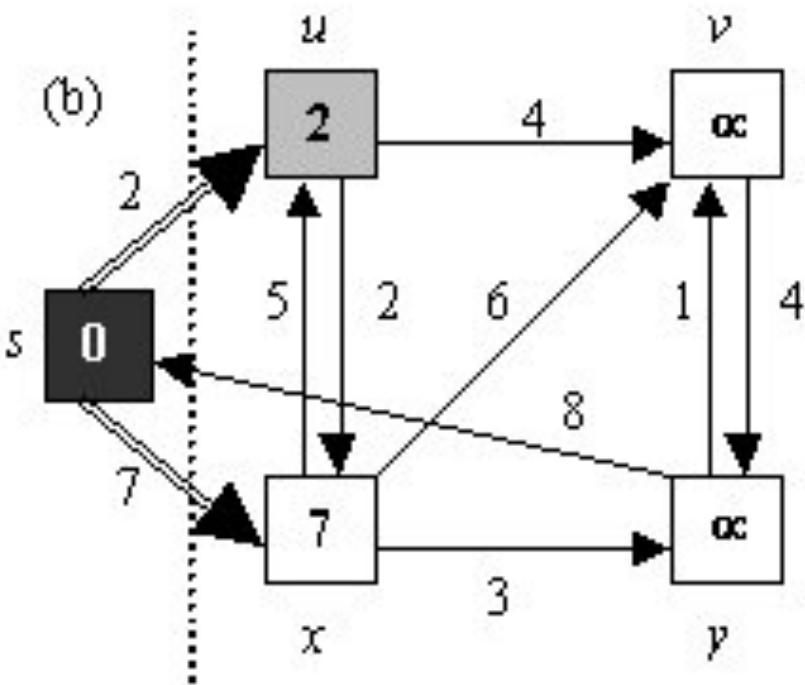
B = { s }

dist= { 0 }

A = {    , u , v , x , y }

dist = {    , **2** ,  $\infty$  , 7 ,  $\infty$  }

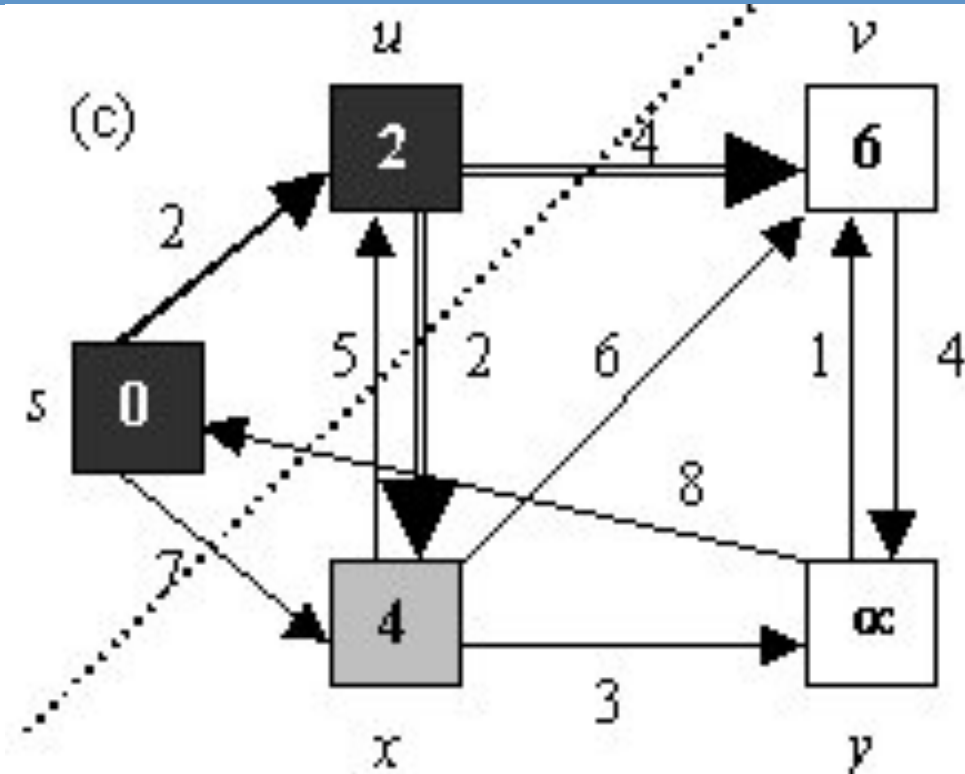
prev = { s , s ,    , s ,    }



# Dijkstra algorithm from vertex s

Round 2:

- choose u (node with least dist[] in A)
- update dist[] of nodes adjacent to s *if applicable*

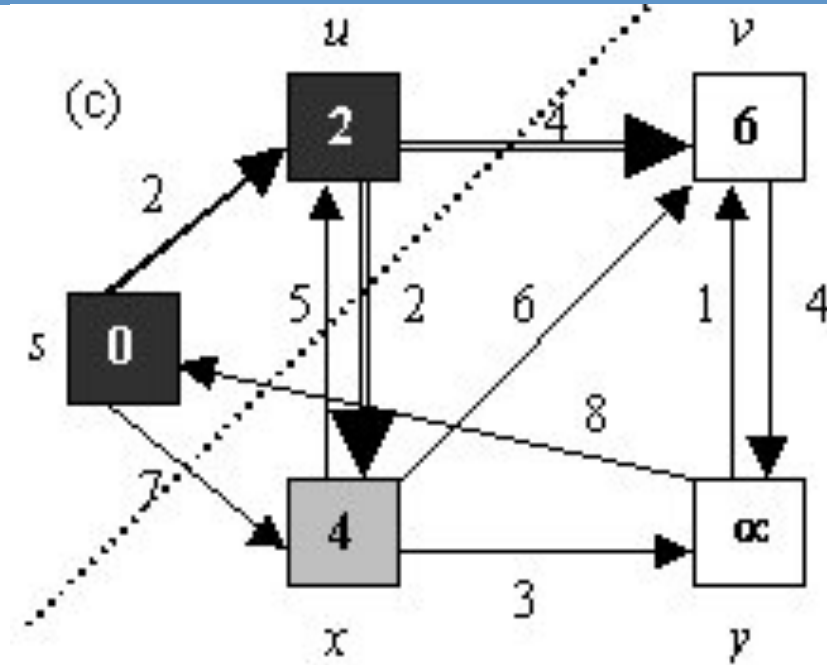


B = { s, u }  
 { 0, 2 }  
 { nil, s }

A = { , , v, x, y }  
 dist = { , , 6, 7 → 4, ∞ }  
 prev = { s, s, u, u, N }

# Dijkstra algorithm from vertex s

	s	u	v	x	y
0	0 s	$\infty$ N	$\infty$ N	$\infty$ N	$\infty$ N
1		2 s	$\infty$ N	7 s	$\infty$ N
2			6 u	<b>4</b> u	$\infty$ N



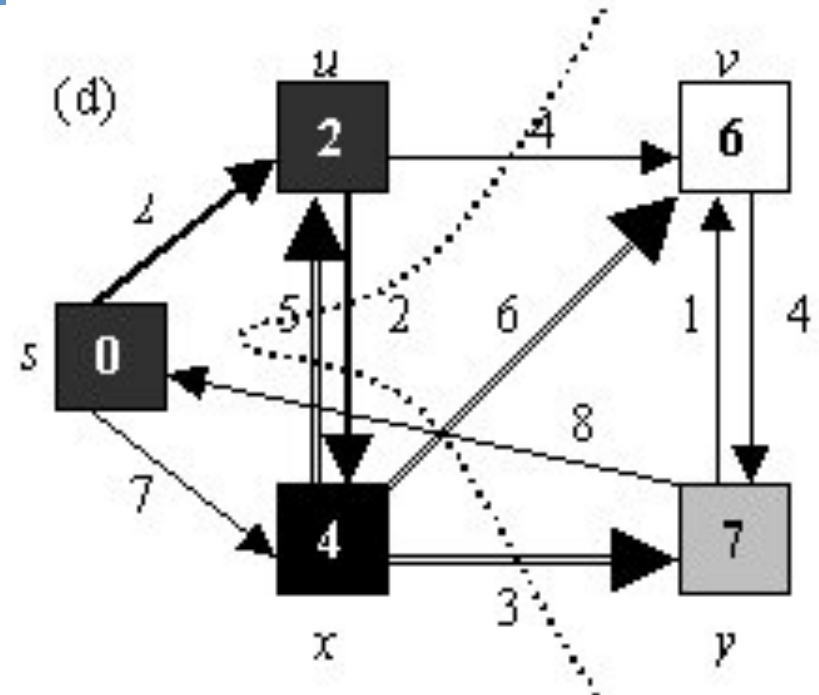
B = { s, u }  
 { 0, 2 }  
 { nil, s }

A = { , , v, x , y }  
 dist = { , , 6, 7  $\rightarrow$  **4** ,  $\infty$  }  
 prev = { s, s , u, u , N }

# Dijkstra algorithm from vertex s

## Round 3

	s	u	v	x	y
0	0 s	$\infty$ N	$\infty$ N	$\infty$ N	$\infty$ N
1		2 s	$\infty$ N	7 s	$\infty$ N
2			6 u	4 u	$\infty$ N
3			6 u		7 x


$$B = \{s, u, x\}$$

$$\{0, 2, 4\}$$

$$\{\text{nil}, s, u\}$$

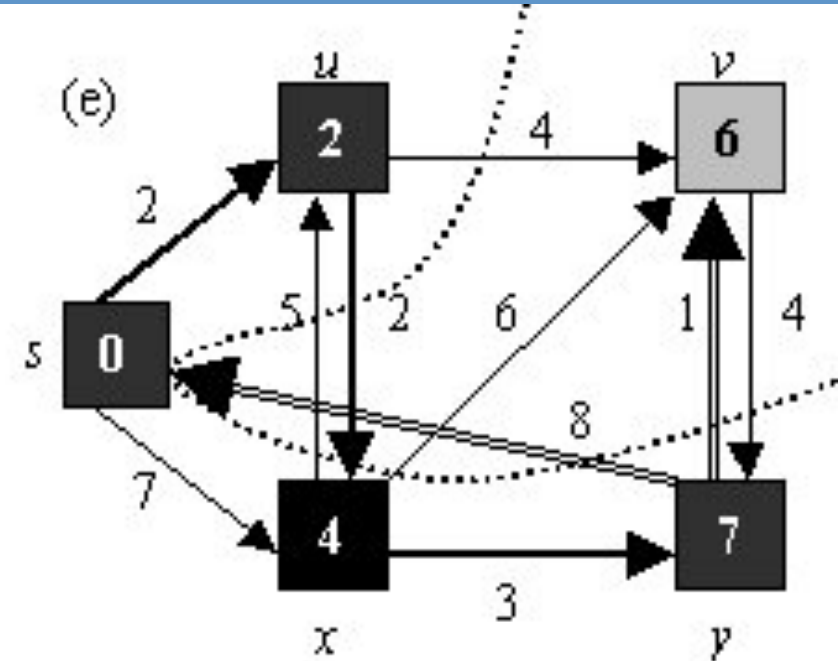
```
A = { , , v , , y }
dist = { , , 6 , , 7 }
prev = { , , u , , x }
```

# Dijkstra algorithm from vertex s

## Round 4

	s	u	v	x	y
0	0 s	$\infty$ N	$\infty$ N	$\infty$ N	$\infty$ N
1		2 s	$\infty$ N	7 s	$\infty$ N
2			6 u	<b>4</b> u	$\infty$ N
3			<b>6</b> u		7 x
4					<b>7</b> x

B = { s, u, x, v }

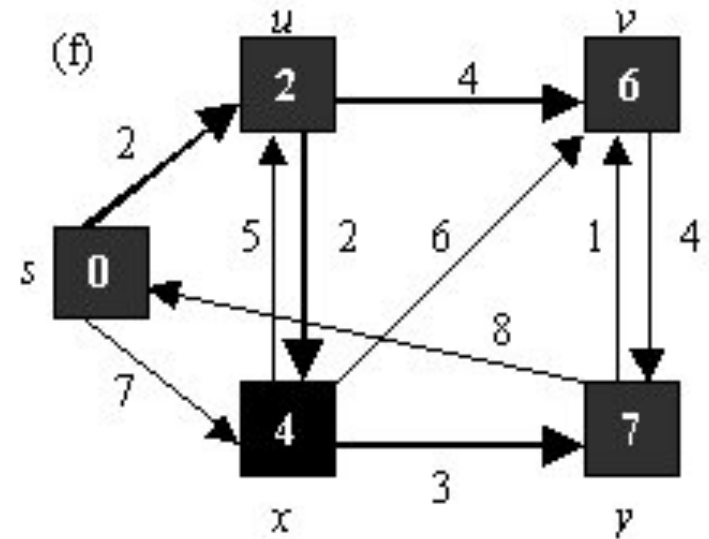


A = { y }

# Dijkstra algorithm from vertex s

# Round 5

	s	u	v	x	y
0	0 s	$\infty$ N	$\infty$ N	$\infty$ N	$\infty$ N
1		2 s	$\infty$ N	7 s	$\infty$ N
2			6 u	4 u	$\infty$ N
3			6 u		7 x
4					7 x
5					


$$B = \{s, u, x, v, y\}$$
$$A = \{ \}$$

Question: what is the shortest path from  $s$  to  $y$ ?

# Dijkstra's algorithm [conceptual only]

Purpose: Find shortest path from vertex  $s$

set  $\text{dist}[u] = \infty$ ,  $\text{pred}[u] = \text{nil}$  for all  $u$ , then  $\text{dist}[s] = 0$ ;

set  $B = \emptyset$ ,  $A = V$ .  $B$  is set of vertices where shortest path from  $s$  found,  $A$  is set of other vertices.

while ( $A$  is not empty):

    select  $u$  from  $A$  such that  $\text{dist}[u]$  is smallest

    remove  $u$  from  $A$  and add it to  $B$

    for all  $(u, v)$  in  $G$ :

        if ( $\text{dist}[v] > \text{dist}[u] + w(u, v)$ ): update  
         $\text{dist}[v]$  and  $\text{pred}[v]$

Q: How to represent  $A$ ?



# Dijkstra's algorithm [conceptual only]

```
set dist[u]=  $\infty$ , pred[u]=nil for all u, then dist[s]= 0;  
set B=  $\emptyset$ , A= makePQ(V)  
while (A not empty)  
    u= deleteMin(A)  
    add u to B    // practice: just mark u as visited  
    for all (u,v) in G:  
        if (dist[v] > dist[u]+w(u,v):  
            update dist[v] and pred[v]
```

**Practical note:** “update dist[v] and pred[v]” means

dist[v]= dist[u]+w(u,v) ; pred[v]= u

decrease weight of v in PQ to dist[v], hence, need  
to locate v in PQ, change weight and upheap

## Q10.1: For a directed graph with the following edges:

### Data

**a b 3**

**a d 7**

**b d 2**

**c e 6**

**d b 2**

**d c 5**

**d e 4**

**e d 2**

**For a directed graph with the edges listed in LHS:**

1. Draw a weighted directed graph that reflects these edges and weights (logical representation).
2. Construct an adjacency matrix for the weighted digraph you have just drawn, including the weights. Be explicit about how you are going to handle matrix cells for which there is no information in the data.
3. Run through Dijkstra's Algorithm starting from the vertex a.

Implement priority queue (see LMS)

- you can use the heap implementation of previous weeks

# Assignment 2:

to fill in code within `puzzle.c`

- Write functions at FILL WITH YOUR CODE
- You might want to add to existed data structures, see comments in `typedef struct node`
- You might want to add some additional functions

Interface:

```
make
```

```
./15puzzle 1.puzzle
```