COMP20003 Workshop Week 7

- 1 Sorting, Insertion Sort and Selection Sort revisited
- 2 have some fun ...
- 3 Quicksort + Q 6.1
- 4 Review for the Test / sample test
- 5 Implementing Hash Table (P6.1)

Sorting: algorithms, stable sort, in-place sort...

	Selection	Insertion	Quick	Merge (top- down)
Basic Idea	Identify the smallest and swap it with the first	From 2 nd element, insert it to the left sub-array so that the extended left sub-array remains sorted.	Choose a pivot, partition array into a <i>lesser</i> and a <i>greater</i> (than pivot) halves	Split to equalsize halves, sort them then merge them.
Best case				
Worst case				
Average				
In-place?				
Stable?				

Sorting by distribution counting – very special!

Condition: $m \le A[i] \le n$ for all i, and r=n-m+1 is not too large.

Special aspect: unlike all other sorting algorithms, we do not compare keys.

How? big-O = ? Example for n=100, m=0, n=10.

The best case of Selection sort is:

- a. $O(n \log n)$.
- b. O(n).
- c. $O(n^2)$.
- d. O(log n).

When?

The best case of Insertion Sort is:

- a. $O(n \log n)$.
- b. O(n).
- c. $O(n^2)$.
- *d. O*(*log n*).

When?

The average case of Insertion Sort is:

- a. $O(n \log n)$.
- b. O(n).
- c. $O(n^2)$.
- d. O(log n).

The average case of Quick Sort is:

- a. $O(n \log n)$.
- b. O(n).
- c. $O(n^2)$.
- d. O(log n).

The worst case of Quick Sort is:

- a. $O(n \log n)$.
- b. O(n).
- c. $O(n^2)$.
- d. O(log n).

When?

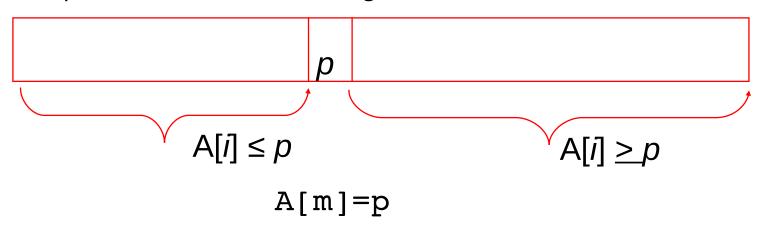
Suppose that A[] is a sorted array of n elements. Which of the following are correct:

- a. We can write a O(n) search algorithm.
- b. We can write a O(log n) search algorithm.
- c. We can write a O(1) search algorithm.
- d. All search algorithms on A[] can be used even if A[] is an unsorted array.

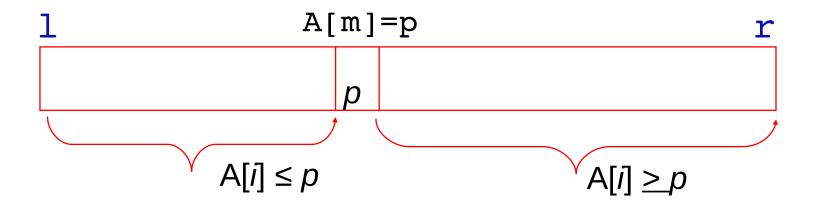
Partitioning (using the rightmost element as pivot)

The task:

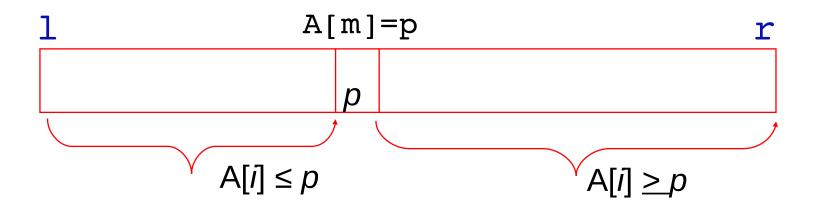
- Input: Given an unsorted slice A[l..r], and supposing that p = a[r].
- Output: A value m and re-arrangement of A so that:



Partitioning (using rightmost element as pivot)



Partitioning variations



- Using element at a random position as pivot
- Using median-of-three as pivot
- Using the leftmost element as pivot

• ...

Quick Sort (concept + algorithm)

```
void quicksort(item A[], int l, int r)
      int i;
      if (r <= 1) return;</pre>
      i = partition(A, 1, r);
      quicksort(A, 1, i-1);
     quicksort(A,i+1,r);
              A[i] \leq p
                                      A[i] \geq p
                                                 Υ
```

Q 6.1

You are asked to show the operation of quicksort on the following keys. For simplicity, use the rightmost element as the partition element:

2 3 97 23 15 21 4 23 29 37 5 23

Comment on the stability of quicksort and its behavior on almost sorted inputs.

week 1-6 in a nutshell (*might be incomplete!*)

Lectures

Introduction & Efficiency

	Fibonacci: recursive & iterative	,1 0
2	 Big-O & Big-O : concepts & definition Grow order of popular functions Best-worst-average case analysis 	 Complexity Arrays: static and dynamic allocation Multi-file C projects and Makefile L: pointer parameters, dynamic arrays
3	 ADT Arrays: selection sort, search Linked lists: arrays vs linked lists 	 C useful tools Tree traversal, BST insert, gdb/valgrind AVL and rotations L: impl BST Insert
4	Tree & BST, operations on BSTAVL & Rotations	 Stacks & Queue: linked lists & array impl L: Assignment 1 and/or free mem used by a stack
5	 Distribution Counting (alg + complexity) Hashing: collisions, chaining, linear probing, double hashing 	 Hashing, dealing with collisions L: Assigment 1 and/or impl simple hash table
6	Sorting, insertion sort, qsort	 sorting algorithms, qsort, distribution counting L: impl simple hash table COMP20003. Workshop. Anh Vo 15
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Followed Workshops (Week W+1)

C review, prog environment, Makefile

Sorting

Concepts: sorting, stable sort, in-place sort Some sorting algorithm (on an array of n elements):

	Selection	Insertion	Quick	Merge
Basic Idea	Identify the smallest and swap it with the first	From 2 nd element, insert it to the left sub-array so that the extended left sub-array remains sorted.	Choose a pivot, partition array into a <i>lesser</i> and a <i>greater</i> (than pivot) halves	Split to equalsize halves, sort them then merge them.
Best case	$O(n^2)$	0(n)	$O(n \log n)$	$O(n \log n)$
Worst case	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(n \log n)$
Average	$O(n^2)$	$O(n^2)$	$O(n \log n)$	$O(n \log n)$
In-place?	✓	✓	✓	*
Stable?	*	✓	*	✓

P6.1 Implementing a simple hash table, OR group work:

Choices:

- Implementing a simple hash table
- Group work with sample MST questions and/or programming tasks, for example:
 - **Big-O** questions
 - Hashing examples
 - AVL rotations
 - Programming: dynamic arrays and strings
 - Programming: linked lists