

LAB	<p>Counting & Radix sort: Questions 10.1-10.3</p> <p>Heap: Questions 10.4-10.5</p> <p>Revision : complexity, recurrences, master theorem or: previous weeks' lab</p>
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Given:

```
input array: {0,1,2,0,0,1,2,1,1,0,0,0}
```

How to sort the array efficiently? Is it in-place? Is it stable?

Discuss limitations.

Counting Sort: how

Conditions: keys are integers in a small range (small in comparison with n), for example:
array of positive integers, each ≤ 2 :

input array: {0,1,2,0,0,1,2,1,1,0,0,0}

value	0	1	2
count table	6	4	2
index range in the sorted array	0..5	6..9	10..11

	Option 1	Option 2
cumulative count array	storing end-index of keys: F= {5,9,11} (or F= {6,10,12} as in lecture)	storing start-index of keys: F= {0,6,10}
how to build sorted array B to ensure stability?	scanning input array A <i>right-to-left</i> for i from n-1 to 0 do B[F[A[i]-min]] = A[i] F[A[i]-min]= F[A[i]-min]+1	scanning input array A <i>left-to-right</i> for i from 0 to n-1 do B[F[A[i]-min]] = A[i] F[A[i]-min]= F[A[i]-min]+1

Counting Sort: summary

- Can be applied when $A[i]$ in range min..max, where $k = \text{max} - \text{min} + 1$ is small
- Time complexity: $\theta(n+k)$ (or $\theta(n)$ if k could be considered as a small constant)
- In-place: NO, using additional memory: $\theta(n+k)$, (or $\theta(n)$ if k small)
 - $\theta(k)$ for table of frequencies (and then that becomes the cumulative count table)
 - $\theta(n)$ for building the output sorted array
- Stable: YES, but with careful implementation

Radix Sort= sort digits right-to-left

Applied when all keys can be represented as *same-size strings* over a small-size alphabet σ . Examples:

$\{22, 17, 167, 28, 173, \dots\}$ where $0 \leq x_i < 256$

$= \{ \quad 022, \quad 017, \quad 167, \quad 028, \quad 173, \dots \}$ $\sigma = \{0, 1, \dots, 9\}$ using 3-digit

$= \{ \quad 16, \quad 11, \quad A7, \quad 1C, \quad AD, \dots \}$ $\sigma = \{0, 1, \dots, 9, A, B, \dots, F\}$ using 2-digit

$= \{0001\ 0110, 0001\ 0001, 1010\ 0111, 0001\ 1100, 1010\ 1110, \dots\}$ $\sigma = \{0, 1\}$ using 8-digit

Radix Sort:

From **rightmost to leftmost position** of strings: sort the string by that position by:

- using any *stable* sorting method, and
- counting sort is often the choice

Complexity for n strings of length m : ???

Q10.1 - Counting Sort: Use counting sort to sort the following array of characters:

[a, b, a, a, c, d, a, a, f, c, b]

How much space is required if the array has n characters and our alphabet has k possible letters.

Q10.2 - Radix Sort: Use radix sort to sort the following strings:

abc bab cba ccc bbb aac abb bac bcc cab aba

As a reminder radix sort works on strings of length k by doing k passes of some other (stable) sorting algorithm, each pass sorting by the next most significant element in the string. For example in this case you would first sort by the 3rd character, then the 2nd character and then the 1st character.

Q10.3: Which property is required to use counting sort to sort an array of tuples by only the first element, leaving the original order for tuples with the same first element. For example, the input may be:

(8, campbell), (6, tal), (3, keir), . . . (6, gus), (0, nick), (8, tom)

Discuss how you would ensure that counting sort satisfies this property. Can you achieve this using only arrays? How about using auxiliary linked data structures?

Radix Sort: Use radix sort to sort the following strings:

abc bab cba ccc bbb aac abb bac bcc cab aba

First, sort (using *stable* counting sort) by the last letters:

ab**c** ba**b** cb**a** cc**c** bb**b** aa**c** ab**b** ba**c** bc**c** ca**b** ab**a**

→ cb**a** ab**a** ba**b** bb**b** ab**b** ca**b** ab**c** cc**c** aa**c** ba**c** bc**c**

Next, do the same for the middle letter :

c**b**a a**b**a b**a**b b**b**b a**b**b c**a**b a**b**c c**c**c a**a**c b**a**c b**c**c

→ bab cab aac bac cba aba bbb abb abc ccc bcc

Last, do the same for the first letter:

bab **c**ab **a**ac **b**ac **c**ba **a**ba **b**bb **a**bb **a**bc **c**cc **b**cc

→ aac aba abb abc bab bac bbb bcc cab cba ccc

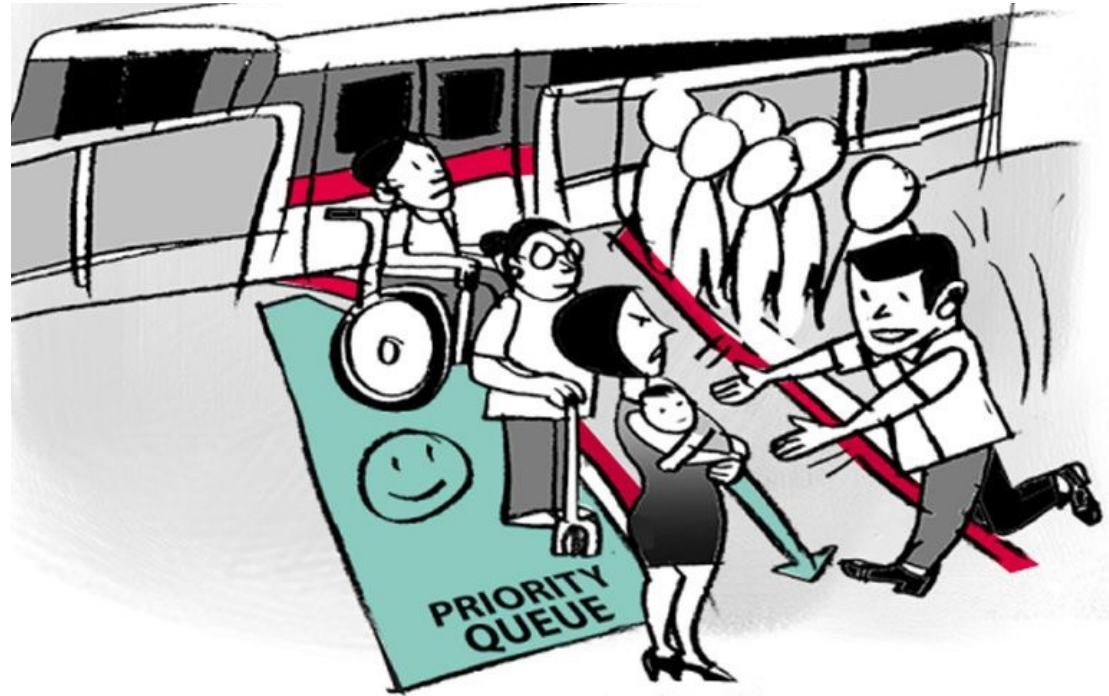
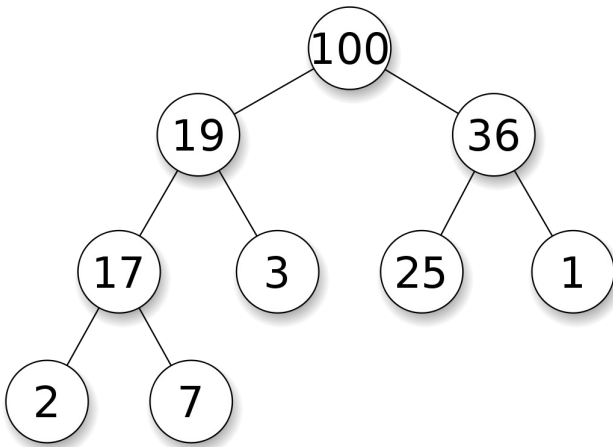
Note: the sorting method is required to be stable, why?

An Implementation for Priority Queue: Binary Heap

Binary Heap as a *concrete data type* (implementation) for PQ.

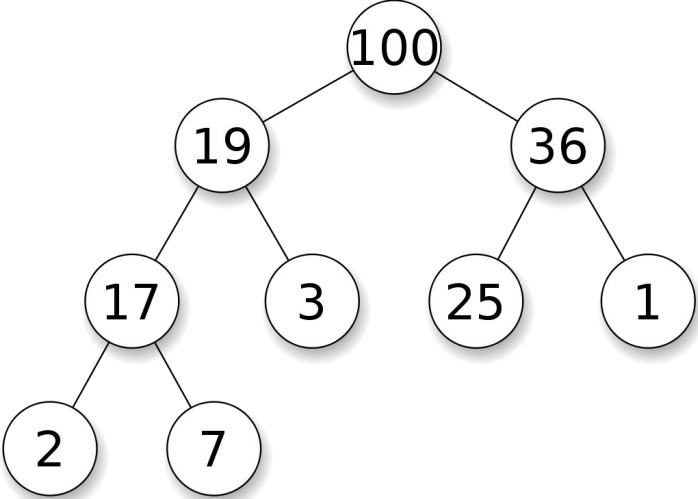
Simple priority: higher, lower → max heap, min heap

What is a, say, max heap?
How is it implemented?



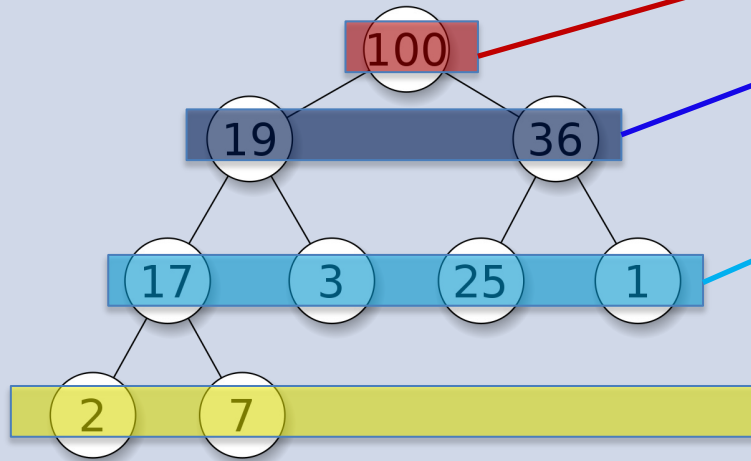
'Can I borrow your baby?...'

Binary Heap: conceptually, is a binary tree satisfying 2 conditions

Example	Conditions
 <pre>graph TD; 100((100)) --- 19((19)); 100 --- 36((36)); 19 --- 17((17)); 19 --- 3((3)); 17 --- 2((2)); 17 --- 7((7)); 36 --- 25((25)); 36 --- 1((1));</pre>	<ol style="list-style-type: none"><li data-bbox="1116 362 2063 668">1. The tree is <i>complete</i>:<ul style="list-style-type: none"><li data-bbox="1212 429 2063 539">• all levels, except for the last, are full<li data-bbox="1212 554 2063 668">• the last level is filled from left to right<li data-bbox="1116 811 2063 1049">2. The <i>heap property</i>: each node has a higher priority (here, is larger) than any of its descendants (or equivalently, just its children).

Binary Heap: is implemented as an array!

Visualisation: as a complete binary tree



Implementation: using arrays

idx	0	1	2	3	4	5	6	7	8	9
val	×	100	19	36	17	3	25	1	2	7

note: $H[1]$ is for the root
 $H[0]$ not used (or
used for a sentinel)

Heap is $H[1..n]$

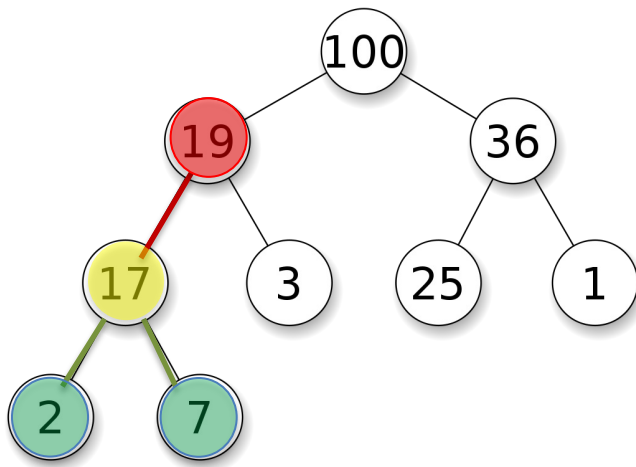
- level i occupies 2^i cells in array $H[1..n]$
(except for the last level)
- if root is level 1, then level i starts from $H[2^{i-1}]$

Binary Heap: parent and children of a node

•

- left child of $H[i]$ is $H[2*i]$
- right child of $H[i]$ is $H[2*i+1]$

parent of $H[i]$ is $H[i/2]$



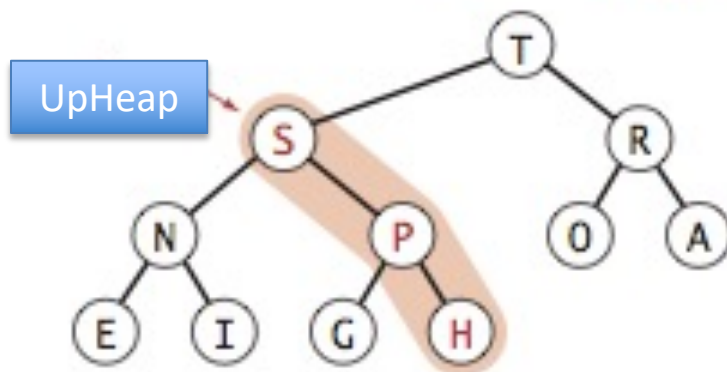
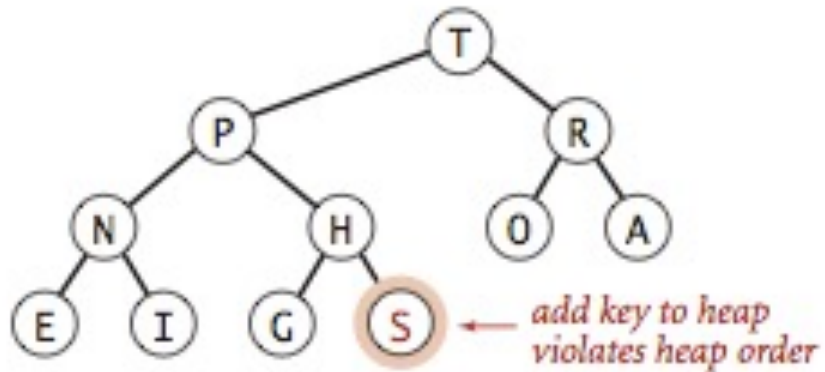
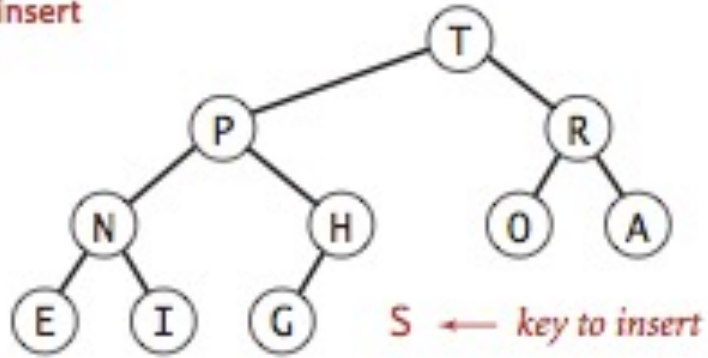
idx	0	1	2	3	4	5	6	7	8	9
val	×	100	19	36	17	3	25	1	2	7

Note:

- The use of indices 1..n is not a rule
- index 0..n-1 can also be used for heaps, but the formula for parent/children are not so nice

inject = enPQ = Insert a new elem into a heap using UpHeap

insert



Notes:

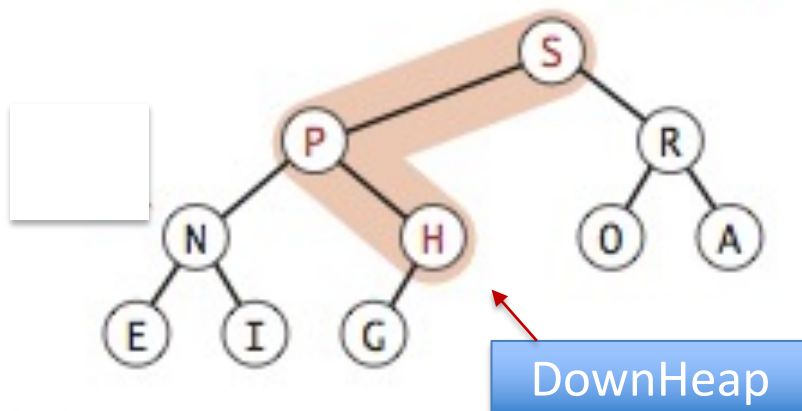
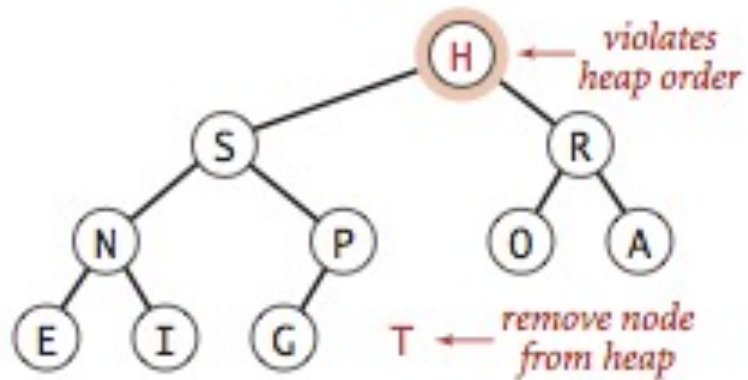
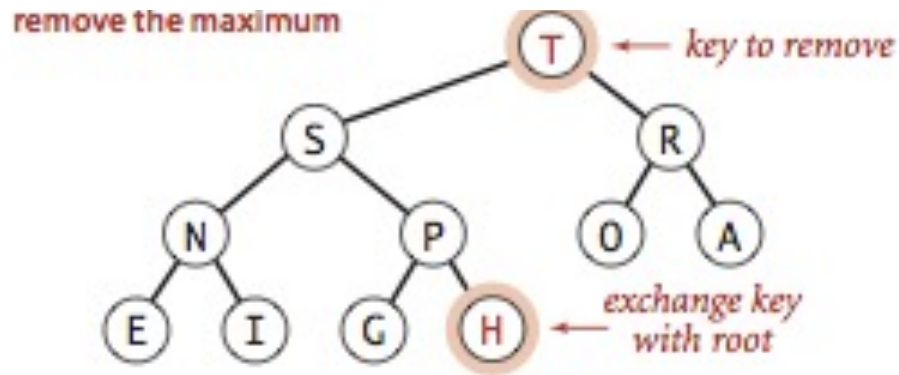
Procedure:

Complexity:

UpHeap (aka. SiftUp)

while (has parent and parent has lower priority): swap up with the parent

ject = delete (and returns the heaviest) and repair using DownHeap



Notes:
Procedure:

Complexity:

DownHeap (aka. SiftDown) :
while (has children and the heavier child
has higher priority): swap down with the
heavier child

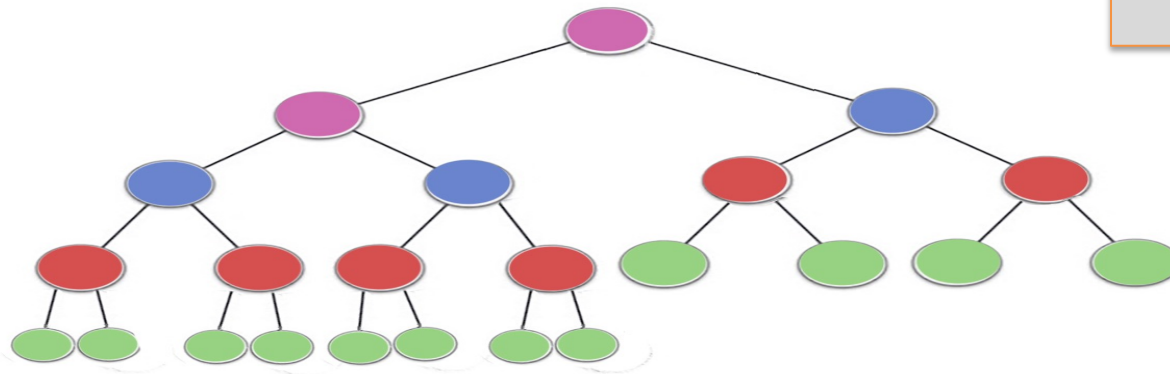
Heapify: Turning an array $H[1..n]$ into a heap

```
function Heapify( $H[1..n]$ )  
  for  $i \leftarrow n/2$  downto 1 do  
    downheap( $H, i$ )
```

= $\Theta(n)$ (see lectures and/or ask Google for a proof)

The operation is aka. **Heapify**/Makeheap/ Bottom-Up Heap Construction

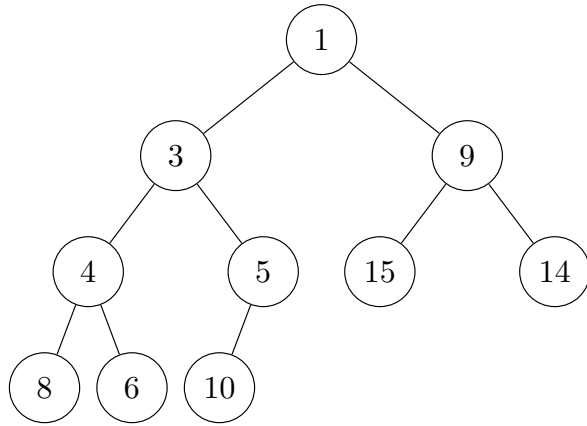
Example: build maxheap for keys E X A M P



Notes:

Complexity:

Q 10.4



a) Show how this heap would be stored in an array as discussed in lectures (root is at index 1; node at index i has children at indices $2i$ and $2i+1$)

b) Run the `RemoveRootFromHeap` (`eject`) algorithm from lectures on this heap by hand

c) Run the `InsertIntoHeap` (`inject`) algorithm and insert the value 2 into the heap

Your answers:

a) array is: [???]

b) Run the `RemoveRootFromHeap`:

c) Run the `InsertIntoHeap(2)`:

NOTES ON HEAPSORT:

Q 10.5 [opt]: k-smallest using min-heap

The *k*-th smallest problem:

- Given an array $A[]$ of n elements, and an integer k
- Find the k -th smallest value (suppose that k is zero-origin, that is, k can be any of $0, 1, 2, \dots, n-1$)

How can we use a min-heap data structure to solve the k th-smallest element problem? What is the time-complexity of this algorithm?

Your answer:

Algorithm	Complexity
<pre>function HeapkthSmallest(A[0..n-1], k) Heap := Heapify(A) for i:= 0 to k-1 do x= DeleteMin(Heap) return x</pre>	

Q:

- How to compare with QuickSelect?
- Can we use max-heap to solve the problem?

$$1 < \log n < n^\epsilon < n^c < n^{\log n} < c^n < n^n \quad \text{where } 0 < \epsilon < 1 < c$$

$$O(f(n) + g(n)) = O(\max\{f(n), g(n)\}) \quad \text{note: these 3 also applied to big-}\theta$$

$$O(cf(n)) = O(f(n))$$

$$O(f(n) \times g(n)) = O(f(n)) \times O(g(n))$$

$$1 + 2 + \dots + n = n(n+1)/2 = \theta(n^2)$$

$$1^2 + 2^2 + \dots + n^2 = n(n+1)(2n+1)/6 = \theta(n^3)$$

$$1 + x + x^2 + \dots + x^n = (x^{n+1} - 1)/(x - 1) \quad (x \neq 1)$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \begin{cases} 0 & f(n) = O(g(n)) \\ c & f(n) = \theta(g(n)) \\ \infty & f(n) = \Omega(g(n)) \end{cases}$$

$$\lim_{n \rightarrow \infty} \frac{t(n)}{g(n)} = \lim_{n \rightarrow \infty} \frac{t'(n)}{g'(n)}$$

Revision: Complexity – Recurrences : Q 10.6 and more

Q 10.6: Solve the following recurrence relations. Give both a **closed form expression** in terms of n and a **Big-Theta bound**.

a) $T(n) = T(n/2) + 1, T(1) = 1$

b) $T(n) = T(n-1) + n/5, T(0) = 0$

The Master Theorem is convenient but:

- constants a, b, d not always available
- not for finding **closed form expression**

Other exercises: review complexity exercises for Workshops Week 3, 4, 7; Levitin C2

Q 10.6+: For each of the following cases, indicate whether $f(n)$ is $O(g(n))$, or $\Omega(g(n))$, or both (that is, $\Theta(g(n))$)

(a) $f(n) = (n^3 + 1)^6$ and $g(n) = (n^6 + 1)^3$,

(b) $f(n) = 3^{3n}$ and $g(n) = 3^{2n}$,

(c) $f(n) = \sqrt{n}$ and $g(n) = 10n^{0.4}$,

(d) $f(n) = 2 \log_2 \{(n + 50)^5\}$ and $g(n) = (\log_e(n))^3$,

(e) $f(n) = (n^2 + 3)!$ and $g(n) = (2n + 3)!$,

(f) $f(n) = \sqrt{n^5}$ and $g(n) = n^3 + 20n^2$.

- Choose the simplest logic:
 - drop free constants & lower members
 - use lim
- Normally, you should show your work-out
- Feel free to add other questions, including the ones in MST

- work other lab questions from earlier weeks, or
- continue with reviewing: complexity, recurrences, divide-and-conquer, DP