COMP20007 Workshop Week 2

Welcome to a F2F Workshop!

1   About our workshops and comp20007

2   Arrays and Linked Lists   (Tutorial Q1-Q2)
        ADT: Stacks, Queues   (Q3-Q4)

3   5-min break for networking

4   LAB:
- Programming Environment
- C revision with some exercises

a friend of yours in comp20007

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if email:

- you should send from your uni's email account,

- subject better to start with "COMP20007" or just "C207"

suggested ToDos for Workshops (Learning-By-Doing)

in your own time	in workshop

in your own time

W=1

↓

Attend & Revise lectures of week W

↓

W++

↓

- Use Ed's Workshop W to prepare
- Also note the questions you want to ask

- Finish the outstanding tasks
- Check with the solutions provided

in workshop

in the first part (tutorial):
- be active:
  - do the problems
  - ask, answer questions, vote
- be cooperative.

have 5 min for stretch exercises and social networking

in the second part (lab):
- be cooperate & active
- finish the lab problems or be confident that you can do so

directly (or via emails) tell Anh on what else you want in the workshops

comp20007

Focus: designing algorithms.

Use efficiency performance to evaluate designs:
- mainly, time complexity, and
- sometimes, space complexity.

```
return p
p := p.next
return null
```

Problem → **Algorithm/Pseudocode**   **Complexity**

Searching for a specified element (amongst a series of elements)

```
j := 0
while j < last
  if A[j] == x
    return j
  j := j+1
return null
```

O(?)

```
p := head
while p != null
  if p.val == x
    return p
  p := p.next
return null
```

O(?)

```
function find(A,x,lo,hi)
  if lo > hi
    return null
  else if A[lo] == x
    return lo
  else
    return find(A,x,lo+1,hi)
```

O(?)

Problem → **Pseudocode** → C code ?

Searching for a specified element (amongst a number of elements)

```
p := head
while p != null
   if p.val == x
      return p
   p := p.next
return null
```

```
j := 0
while j < last
   if A[j] == x
      return j
   j := j+1
return null
```

```
function find(A,x,lo,hi)
   if lo > hi
      return null
   else if A[lo] == x
      return lo
   else
      return find(A,x,lo+1,hi)
```

???

???

???

6

Problem → **Pseudocode** → C code

Searching for a specified element (amongst a number of elements)

```
p := head
while p != null
  if p.val == x
    return p
  p := p.next
return null
```

```
node_t *search(int x, list_t *l) {
  node_t *p= l->head;
  while (p) {
    if (p->val==x) return p;
    p= p->next;
  }
  return NULL;
}
```

```
j := 0
while j < last
  if A[j] == x
    return j
  j := j+1
return null
```

```
???
```

```
function find(A,x,lo,hi)
  if lo > hi
    return null
  else if A[lo] == x
    return lo
  else
    return find(A,x,lo+1,hi)
```

```
???
```

Our focus is building algorithms and analysing efficiency.

Pseudocode and big-O are the main tools.
But we also do some C implementation.

Q2.1: Arrays & Linked Lists

Describe how you could perform the following operations on *sorted* and *unsorted arrays*, and decide if they are *O(1), O(log n),* or *O(n)*, where *n* is the number of elements initially in the array. Assume that there is no need to change the size of the array to complete each operation.

Operation	Unsorted Arrays	Sorted Arrays
Searching for a specified element		
Inserting a new element		
Deleting the final element		
Deleting a specified element		

What's your best way to work and take notes in workshops?

- pens and papers: excellent, you will need to both write/type and draw
- e-notes: if you can easily draw; for e-notes, you can also:
  - copy and paste the exercises from ED, then add your solution
  - use Anh's slides and directly add notes

Q2.1: Arrays & Linked Lists

Describe how you could perform the following operations on *sorted* and *unsorted arrays*, and decide if they are *O(1)*, *O(log n)*, or *O(n)*, where *n* is the number of elements initially in the array. Assume that there is no need to change the size of the array to complete each operation.

Operation	Unsorted Arrays	Sorted Arrays
Searching for a specified element	O(n) • do a linear search	O( ) •
Inserting a new element	O( ) •	O( ) •
Deleting the final element	O( ) •	O( ) •
Deleting a specified element	O( ) •	O( ) •

Q2.2: Linked Lists

Describe how you could perform the following operations on singly-linked and doubly-linked lists, and decide if they are *O(1)*, *O(log n)*, or *O(n)*, where *n* is the number of elements initially in the linked list. Assume that the lists need to keep track of their final element.

Operation	Singly	Doubly
Inserting a node at the start	**O( )**	
Inserting a node at the end		
Deleting the first node (at the start)		
Deleting last node (at the end)		

Q2.2: Check your answers

Describe how you could perform the following operations on singly-linked and doubly-linked lists, and decide if they are *O(1)*, *O(log n)*, or *O(n)*, where *n* is the number of elements initially in the linked list. Assume that the lists need to keep track of their final element. **Note: also assume that the links to the first and last element of a list is called `head` and `foot` respectively.**
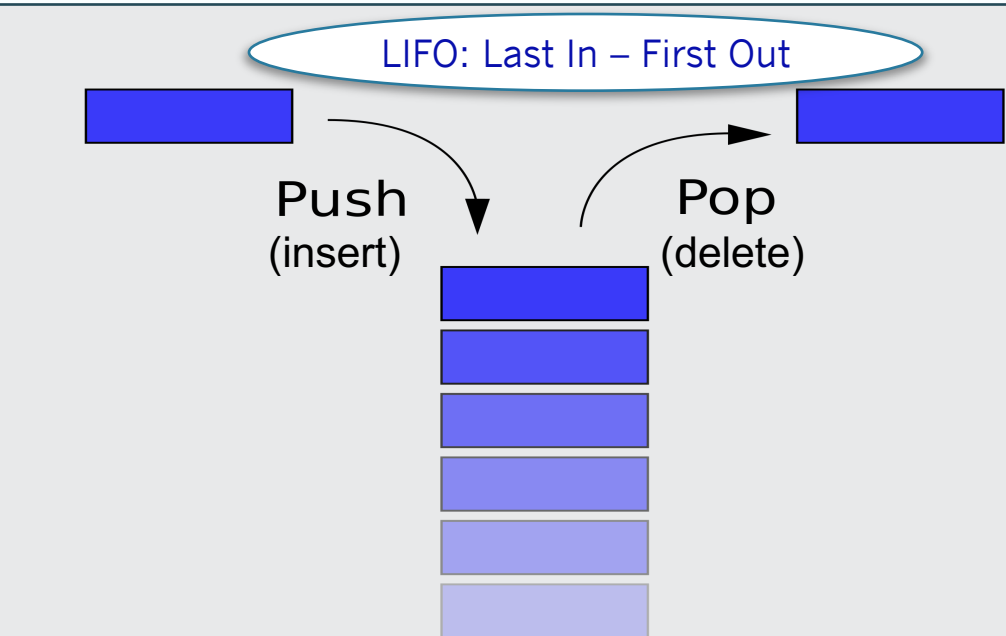
| Operation	Singly	Doubly
Inserting a node at the start	**O(1)** • ensure that the link in the inserted/removed node, the link in the next or previous node are updated • ensure that the links `head` and `foot` of the list are updated accordingly; note that when inserting to an empty list, and when deleting from a list that has a single element, both `head` and `foot` need to be updated	**O(1)** • as in the first row of LHS, but note that each node has 2 links, and both links of the inserted/removed node need to be updated. note that all operations, including deleting the last node, is O(1).
Inserting a node at the end		
Deleting the first node (at the start)		
Deleting the last node (at the end)	**O(n)** • just like the above, but we need to identify the second last element in order to disconnect the last element and to update the link `foot`. To do that, we need to the follow the list all the way from the start to the second last element, and that causes O(n).	

An Abstract Data Type (ADT): Stack (LIFO)

LIFO: Last In – First Out

Push
(insert)

Pop
(delete)

http://www.123rf.com/stock-photo/tyre.html

adapted from https://simple.wikipedia.org/wiki/Stack_(data_structure)

Stack Operations

push(x): insert element **x** to (the top of) stack

pop()   : remove and return an element from (the top of) stack

isEmpty(): check if stack is empty

create() : create a new, empty stack

delete() : delete (free all associated memory)

Another ADT: Queue (FIFO)
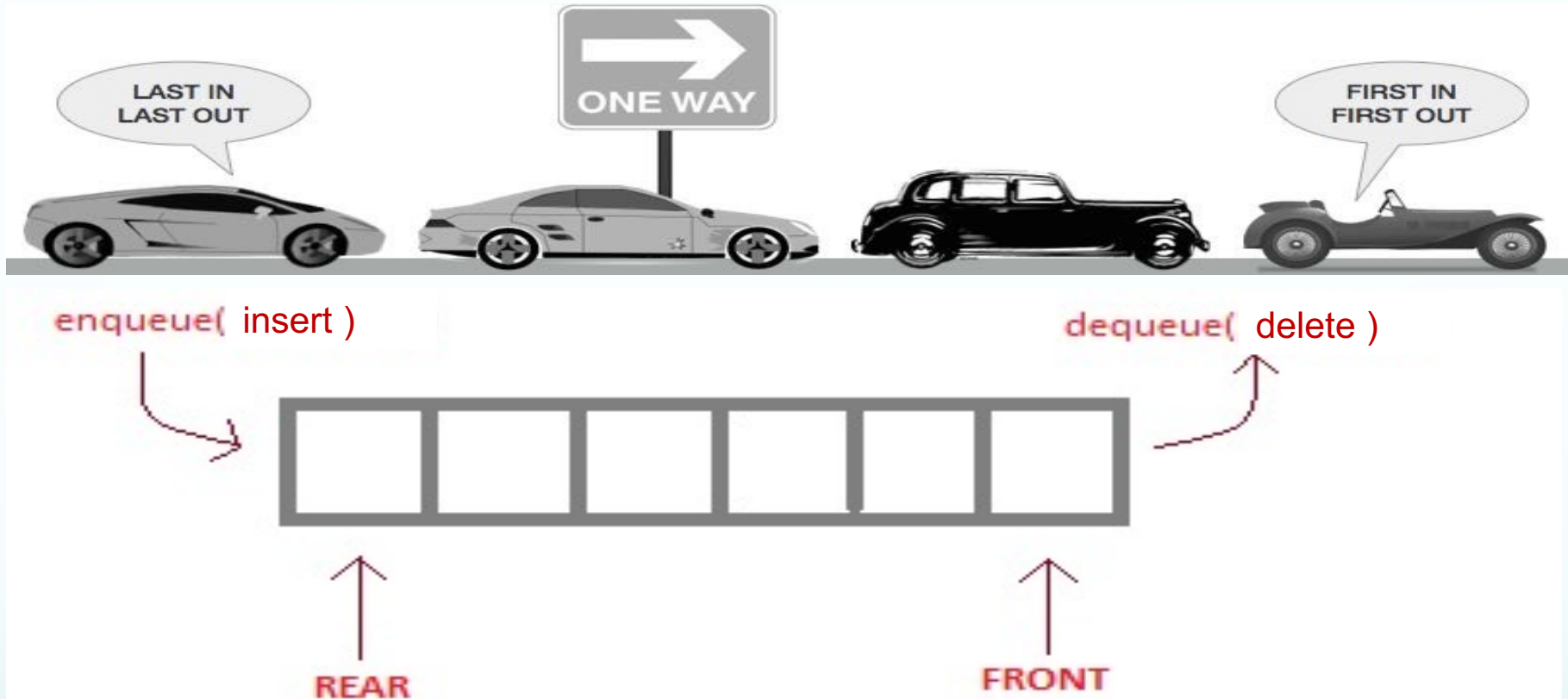
Queue Operations	**enqueue(x):** add **x** to (the rear of) the queue
	dequeue(): remove and return the element from (the front of) the queue
	create(): create a new, empty queue
	isEmpty(): check if queue is empty, or
	delete(): delete a queue (free all associated memory)

Q2.3: Stacks

- Describe how to implement `push` and `pop` using an unsorted array, and using a singly-linked list.

Using an (unsorted) array	Using a (singly-)linked list

Q2.4: Queues

- Describe how to implement **enqueue** and **dequeue** using an unsorted array, and using a singly-linked list. Is it possible to perform each operation in constant time?

Using an array	Using a linked list

Q2.5 [optional]: Stacks & Queues

If you have access only to stacks and stack operations, can you faithfully implement a queue? How about the other way around?

~~You may assume that your stacks and queues also come with a size operation, which returns the number of elements currently stored.~~

Your answer: using stacks to implement a queue

enqueue	dequeue

Your answer: using queues to implement a stack

push	pop

5-minute break

- stretch exercises

- networking

Just for fun (perhaps at home)

google "algorithm for making friends"
and watch "The Friendship Algorithm"
(a 2.5-minute videos).

Lab Time: Use Ed for exercises and assignments

1. *Start* with `helloworld.c`

2. (Together) Implement functions in `functions.c`, which reviews *function and function parameters*

3. *dynamically resizing arrays* with `malloc`/`calloc` and `free`. Forgot `malloc`? Try command "`man malloc`" in your terminal.

Why `Ed`?
- Strong: powerful editor, shell, compilers, `valgrind`, `gdb`, …
- Safe   :   codes and files will never be lost
- Sound :   codes/files can be accessed from any devices
- Sane   :   your assignments will be tested on `Ed`

4. *Optional*: download Alistair's `listops.c` (google it!), then add a *least-effort* implementation of stack with:

- data type `mystack_t` ,

- functions `createStack`, `freeStack`, `push`, `pop`.

Do, or ask Anh how to, **implement stacks in just 5 minutes**!

Wrap Up

- array and linked list as concrete data types.

- stack and queue:
  - as and Abstract Data Type (ADT),
  - operations,
  - implementation using array and linked list.

- C revision, especially:
  - functions, pointers, `malloc`/`realloc`, `free;`
  - dynamically resizing arrays (or just *dynamic arrays*).

- Technical stuffs:
  - Use Ed for programming exercises & assignments!
  - Self-Learn to use `Ed`'s `gcc` and debugging tools `gdb`, `valgrind` at home.

Have Fun with comp20007 and avo

Additional Pages

Data Types & ADT

- A *concrete data type*, such as array or linked list, specifies a representation of data, and programmers can rely on that to implement operations (such as `insert`, `delete`).

- An *abstract data type* specifies possible operations, but not representation. Examples: stacks, queues, dictionaries.

  - When implementing an ADT, programmers use a concrete data type. For example, we might attempt to employ array to implement stack.

  - When using an ADT, programmers just use its facilities and ignore the actual representation and the underlined concrete data type.

Two concrete data types: Arrays & Linked Lists

Access A[k] in O(1) time!
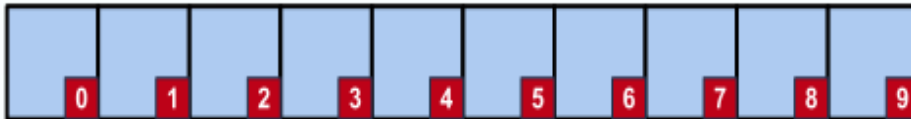
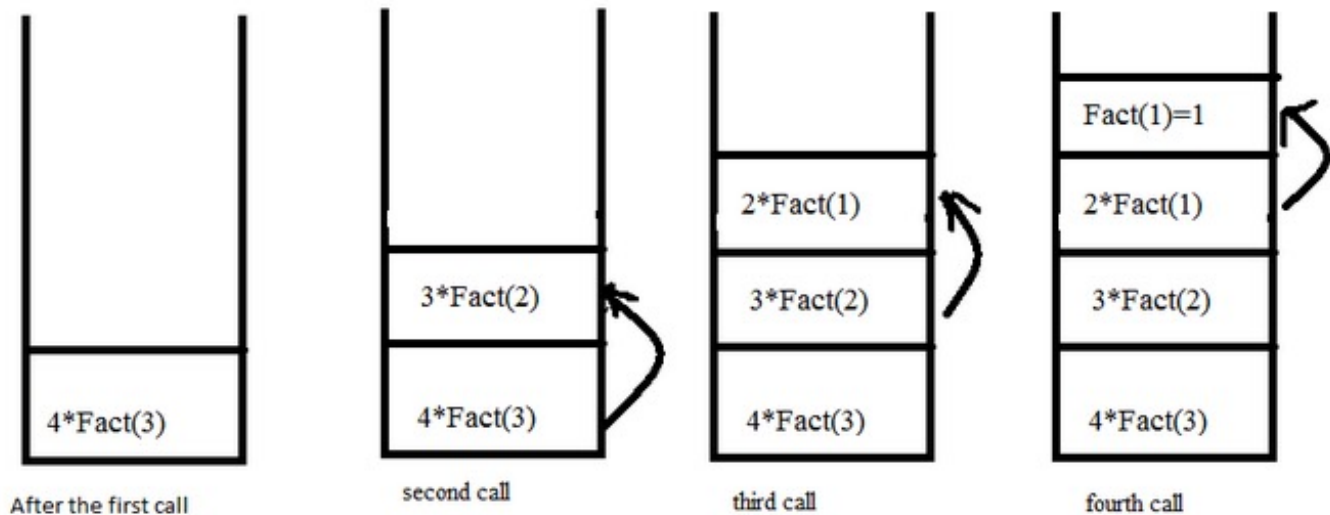Access L[k] in O(n) time!
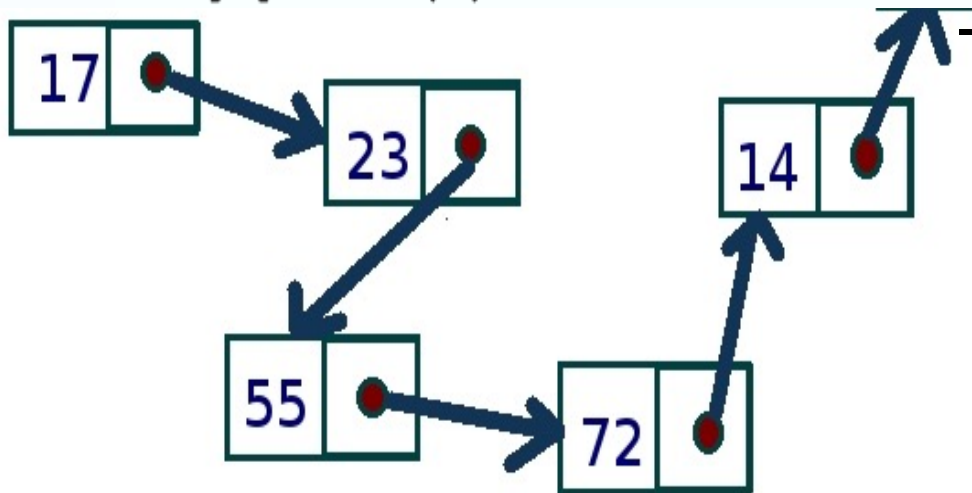
In C:

- How to specify an array? How to traverse it?

- How to specify a linked list? How to traverse it?

Example of using Stacks ?

Stack is widely used in implementation of programming systems. For example, compilers employ stacks for keeping track of function calls and execution.

Stack for :

fact(4)

```
int fact( int n ) {
  if ( n<=1 )
      return 1;
  return n*fact(n-1);
```


When function call happens previous variables gets stored in stack

			Fact(1)=1
		2*Fact(1)	2*Fact(1)
	3*Fact(2)	3*Fact(2)	3*Fact(2)
4*Fact(3)	4*Fact(3)	4*Fact(3)	4*Fact(3)
After the first call	second call	third call	fourth call

Returning values from base case to caller function

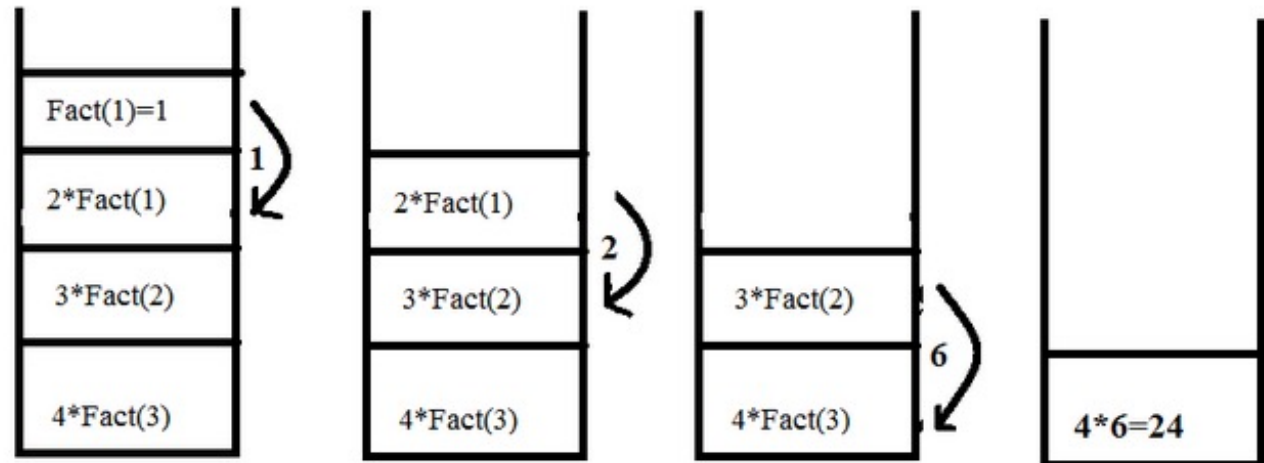| Fact(1)=1 | | | |
2*Fact(1)	2*Fact(1)		
3*Fact(2)	3*Fact(2)	3*Fact(2)	
4*Fact(3)	4*Fact(3)	4*Fact(3)	4*6=24

Image source: http://stackoverflow.com/questions/19865503/can-recursion-be-named-as-a-simple-function-call

ADT: Queue (FIFO)