

The Neo4j Operations Manual v4.0

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This is the operations manual for Neo4j version 4.0, authored by the Neo4j Team.

#### This manual covers the following areas:

- Introduction Introduction of Neo4j Community and Enterprise Editions.
- Installation Instructions on how to install Neo4j in different deployment contexts.
- Docker Instructions on how to use Neo4j on Docker.
- Configuration Instructions on how to configure certain parts of Neo4i.
- Manage databases Instructions on how to manage multiple active databases with Neo4j.
- Clustering Comprehensive descriptions of Neo4j Causal Clustering.
- Fabric Instructions on how to configure and use Neo4j Fabric.
- Upgrade Instructions on upgrading Neo4j.
- Backup Instructions on setting up Neo4j backups.
- Authentication and authorization Instructions on user management and role-based access control.
- Security Instructions on server security.
- Monitoring Instructions on setting up Neo4j monitoring.
- Performance Instructions on how to go about performance tuning for Neo4j.
- Tools Description of Neo4j tools.
- Reference Listings of all Neo4j configuration parameters.
- Tutorials Step-by-step instructions on various scenarios for setting up Neo4j.
- Advanced Causal Clustering Advanced concepts and actions for Neo4j Causal Clustering.
- Deprecated security procedures Deprecated security procedures.

#### Who should read this?

#### This manual is written for:

- the engineer performing the Neo4j production deployment.
- the operations engineer supporting and maintaining the Neo4j production database.
- the enterprise architect investigating database options.
- the infrastructure architect planning the Neo4j production deployment.

# Chapter 1. Introduction

This chapter introduces Neo4j.

Neo4j is the world's leading graph database. Its architecture is designed for optimal management, storage and traversal of nodes and relationships. The database takes a property graph approach which is beneficial for both traversal performance and operations runtime. Neo4j offers dedicated memory management as well as memory efficient operations.

It is scalable and can be deployed as a standalone server or scaled out across multiple machines in a fault-tolerant cluster for production environments. Other features for production applications include hot backups and extensive monitoring.

Cypher is a declarative query language for graphs. Learn the details in the Cypher Manual. The recommended way of interacting with the database programmatically is either through the official drivers, or using the Java API.

### 1.1. Fditions

There are two editions of Neo4j to choose from: Community Edition and Enterprise Edition:

# 1.1.1. Community Edition

The Community Edition is a fully functional edition of Neo4j, suitable for single instance deployments. It has full support for key Neo4j features, such as ACID compliance, Cypher, and programming APIs. It is ideal for learning Neo4j, for do-it-yourself projects, and for applications in small workgroups.

### 1.1.2. Enterprise Edition

The Enterprise Edition extends the functionality of Community Edition to include key features for performance and scalability, such as a clustering architecture and online backup functionality. Additional security features include role-based access control and LDAP support; for example, Active Directory. It is the choice for production systems with requirements for scale and availability, such as commercial solutions and critical internal solutions.

#### 1.1.3. Feature details

Table 1. Features

Edition	Community	Enterprise
Property graph model	0	
Native graph processing & storage	0	
ACID transactions	0	
Cypher graph query language	D	
Neo4j Browser with syntax highlighting		
Bolt binary protocol		
Language drivers for C#, Java, JavaScript & Python [1]		
High-performance native API	0	
High-performance caching	0	

Edition	Community	Enterprise
Cost-based query optimizer	0	0
Graph algorithms library to support Al initiatives $\ ^{\scriptsize{\scriptsize{[1]}}}$		П
Fast writes via native label indexes		
Composite indexes		
Full-text node & relationship indexes	0	
Slotted and Pipelined Cypher runtimes	-	
Property-existence constraints	-	
Node Key constraints	-	
Listing and terminating running queries	-	0
Auto-reuse of space	-	
Role-based access control	-	
Subgraph access control	-	0
LDAP and Active Directory integration	-	0
Kerberos security option	-	0

 $<sup>{}^{{\</sup>scriptscriptstyle [1]}}\!\mathsf{Note}$  that these need to be downloaded and installed separately.

Table 2. Performance & Scalability

Edition	Community	Enterprise
Causal Clustering for global scale applications	-	
Enterprise lock manager accesses all cores on server	-	0
Intra-cluster encryption	-	
Offline backups		
Online backups	-	
Encrypted backups	-	
Rolling upgrades	-	
Automatic cache warming	-	
Routing and load balancing with Neo4j Drivers	-	П
Advanced monitoring	-	
Graph size limitations	34B nodes, 34B relationships, 68B properties	No limit
Bulk import tool		
Bulk import tool, resumable	-	О

# Chapter 2. Installation

This chapter describes installation of Neo4j in different deployment contexts, such as Linux, macOS, and Windows.

#### The topics described are:

- System requirements The system requirements for a production deployment of Neo4j.
- Neo4j Desktop About Neo4j Desktop
- Linux Installation instructions for Linux.
- macOS Installation instructions for macOS.
- Windows Installation instructions for Windows.

As an alternative to installation, you can also run Neo4j in a Docker container. For information on running Neo4j on Docker, see Docker.

# 2.1. System requirements

This section provides an overview of the system requirements for running Neo4j in a production environment.

Neo4j can be installed in many environments and for different scopes, therefore system requirements largely depends on the use of the software. This section distinguishes between a personal/development installation, and a server-based installation.

#### This section contains the following:

- Supported platforms
- Hardware requirements
- Software requirements
- Filesystem
- Java

### 2.1.1. Supported platforms

Neo4j is supported on systems with x86\_64 architectures, whether they are a physical, virtual, or containerized environments.

### 2.1.2. Hardware requirements

In terms of hardware requirements, follow these guidelines:

#### Table 3. Hardware requirement guidelines.

CPU	Performance is generally memory or I/O bound for large graphs, and compute bound for graphs that fit in memory.
-----	---

Memory	More memory allows for larger graphs, but it needs to be configured properly to avoid disruptive garbage collection operations.
Storage	Aside from capacity, the performance characteristics of the disk are the most important when selecting storage:  • Neo4j workloads tend significantly toward random reads.  • Select media with low average seek time: SSD over spinning disks.  • Consult Disks, RAM and other tips for more details.

### For personal use and software development:

Table 4. Hardware requirement guidelines for personal use and software development.

CPU	Intel Core i3 minimum, Intel Core i7 recommended.
Memory	2GB minimum, 16GB or more recommended.
Storage	10GB SATA Minimum, SSD with SATA Express or NVMe recommended.

#### For cloud environments:

Table 5. Hardware requirement guidelines for cloud environments.

CPU	2vCPU minimum, 16+ recommended, possibly Xeon processors.
Memory	2GB minimum, size depends on workloads: in some cases, it is recommended to use instances with memory that fits the size of the graph in use.
Storage	10GB minimum block storage, attached NVMe SSD recommended. Storage size depends on the size of the databases.

### For server-based, on-premise environments:

Table 6. Hardware requirement guidelines for server-based, on-premise environments.

CPU	Intel Xeon processors.
Memory	8GB minimum, size depends on workloads; in some cases, it is recommended to use instances with memory that fits the size of the graph in use.
Storage	SATA i7.2K RPM 6Gbps Hard Drive minimum, NVMe SSD recommended.  Storage size depends on the size of the databases.

### 2.1.3. Software requirements

For personal use and software development:

Table 7. Software requirements for personal use and software development.

Operating System	Supported JDK
MacOS 10.14+	ZulujDK 11 or OpenJDK 11
Ubuntu Desktop 16.04+	ZuluJDK 11, OpenJDK 11 or OracleJDK 11
Fedora 29+	OpenJDK 11 or ZuluJDK 11
Debian 9+	OpenJDK 11 or OracleJDK 11
Windows 10	ZuluJDK 11 or OracleJDK 11

For cloud environments, and server-based, on-premise environments:

Table 8. Software requirements for cloud environments, and server-based, on-premise environments.

Operating System	Supported JDK
Ubuntu Server 16.04+	OpenJDK 11 or OracleJDK 11
Red Hat Enterprise Linux Server 7.5+	OpenJDK 11 or Red Hat OpenJDK 11
CentOS Server 7.5+	OpenJDK 11
SUSE Linux Enterprise Server 15+	OracleJDK 11
Amazon Linux AMI 2018.03+	Amazon Corretto 11, OpenJDK 11 or OracleJDK 11
Windows Server 2012+	OracleJDK 11

# 2.1.4. Filesystem

For proper ACID behavior, the filesystem must support flush (*fsync*, *fdatasync*). See Linux file system tuning for a discussion on how to configure the filesystem in Linux for optimal performance.

# 2.1.5. Java

Neo4j Desktop is available for developers and personal users.

Neo4j Desktop includes a Java Virtual Machine, *JVM*, for convenience. All other versions of Neo4j require Java to be pre-installed.

# 2.2. Neo4j Desktop

This section introduces Neo4j Desktop.

Neo4j Desktop is a convenient way for developers to work with local Neo4j databases.

To install Neo4j Desktop, go to Neo4j Download Center and follow the instructions.



While most functionality is the same, the instructions in this manual are not written for Neo4j Desktop. For example, file locations for a database installed via Neo4j Desktop will be different from those described here.

Neo4j Desktop is not suited for production environments.

# 2.3. Linux installation

This section describes how to install Neo4j on Linux using Debian or RPM packages, or from a Tar archive.

#### This section describes the following:

• Install Neo4j on Debian and Debian-based distributions
□ Installation
□ Upgrade
☐ File locations
□ Operation
Deploy Neo4j using the Neo4j RPM package
☐ Install on Red Hat, CentOS, Fedora or Amazon Linux
☐ Standard installation
<ul> <li>Non-interactive installation of Neo4j Enterprise Edition</li> </ul>
□ Offline installation
□ Install on SUSE
Install Neo4j on Linux from a tarball
<ul><li>Unix console application</li></ul>
□ Linux service
<ul><li>Setting the number of open files</li></ul>
Install Neo4j as a system service
□ Configuration
□ Controlling the service
□ Log

# 2.3.1. Debian

This section describes how to install Neo4j on Debian, and Debian-based distributions like Ubuntu, using the Neo4j Debian package.

#### This section describes the following:

- Installation
  - ☐ Prerequisites (Oracle Java, Debian 9 and Ubuntu 16.04 only)
  - Add the repository
  - □ Install Neo4j
- Upgrade
- File locations
- Operation

#### Installation

To install Neo4j on Debian you need to make sure of the following:

- An OpenJDK Java 11 runtime is installed or available through your package manager.
- The repository containing the Neo4j Debian package is known to the package manager.

Java Prerequisites (Oracle Java, Debian 9+ and Ubuntu 16.04+ only)

Neo4j 4.0 requires the Java 11 runtime. Java 11 is not included in Ubuntu 16.04 LTS or Debian 9 (stretch) and will have to be set up manually prior to installing or upgrading to Neo4j 4.0, as described below. Debian 9 users can find OpenJDK 11 in backports. Debian 10 and Ubuntu 18.04 onwards already have the Openjdk Java 11 package available through apt.

Oracle Java and Debian

Neo4j is compatible with Oracle Java on Debian/Ubuntu Linux, but should be installed via tarball. The Debian installer may still be used, but it will install OpenJDK Java 11 in addition to any existing Java installations.

This is due to changes in Oracle's Debian package manifest between Java versions 8 and 11.

Java 11 on Debian 9

Add the line deb <a href="http://httpredir.debian.org/debian">http://httpredir.debian.org/debian</a> stretch-backports main to a file with the ".list" extension in /etc/apt/sources.list.d/. Then run apt-get update:

```
echo "deb http://httpredir.debian.org/debian stretch-backports main" | sudo tee -a /etc/apt/sources.list.d/stretch-backports.list sudo apt-get update
```

You are now ready to install Neo4j, which will install Java 11 automatically if it is not already installed. See Dealing with multiple installed Java versions to make sure you can start Neo4j after install.

lava 11 on Ubuntu 16.04

Add the official OpenJDK package repository to apt:

```
sudo add-apt-repository -y ppa:openjdk-r/ppa
sudo apt-get update
```

You are now ready to install Neo4j, which will install Java 11 automatically if it is not already installed. See Dealing with multiple installed Java versions to make sure you can start Neo4j after install.

Dealing with multiple installed Java versions

It is important that you configure your default Java version to point to Java 11, or Neo4j 4.0.0 will be unable to start. Do so with the update-java-alternatives command.

• First list all your installed version of Java with update-java-alternatives --list

Your results may vary, but this is an example of the output:

```
java-1.11.0-openjdk-amd64 1071 /usr/lib/jvm/java-1.11.0-openjdk-amd64
java-1.8.0-openjdk-amd64 1069 /usr/lib/jvm/java-1.8.0-openjdk-amd64
```

• Identify your Java 11 version, in this case it is java-1.11.0-openjdk-amd64. Then set it as the default with (replacing java11name with the appropriate name from above)

```
sudo update-java-alternatives --jre --set <java11name>
```

#### Add the repository

The Debian package is available from https://debian.neo4j.com.

• To use the repository for generally available versions of Neo4j, run:

```
wget -0 - https://debian.neo4j.com/neotechnology.gpg.key | sudo apt-key add -
echo 'deb https://debian.neo4j.com stable latest' | sudo tee -a /etc/apt/sources.list.d/neo4j.list
sudo apt-get update
```

To avoid the risk of the apt package manager accidentally forcing a database upgrade, different major and minor releases of Neo4j are also available separately inside the repository. To install Neo4j this way, specify the major and minor version required, in place of latest. We recommend the following method for production or business critical installations:

```
wget -0 - https://debian.neo4j.com/neotechnology.gpg.key | sudo apt-key add -
echo 'deb https://debian.neo4j.com stable {neo4j-version}' | sudo tee -a
/etc/apt/sources.list.d/neo4j.list
sudo apt-get update
```

• To use the repository for milestone releases, alpha, beta, release candidate and preview versions of Neo4j, run:

```
wget -0 - https://debian.neo4j.com/neotechnology.gpg.key | sudo apt-key add -
echo 'deb https://debian.neo4j.com testing latest' | sudo tee -a /etc/apt/sources.list.d/neo4j.list
sudo apt-get update
```

 Once the repository has been added into apt, you can verify which Neo4j versions are available by running:

```
apt list -a neo4j
```

#### Install Neo4j

To install Neo4j Community Edition:

```
sudo apt-get install neo4j=1:4.0.0
```

To install Neo4j Enterprise Edition:

```
sudo apt-get install neo4j-enterprise=1:4.0.0
```

Note that the version includes an epoch version component (1:), in accordance with the Debian policy on versioning.

Versions of Neo4j that are not yet generally available may differ in naming.



The naming structure of packages are normally composed as neo4j-enterprise=1:
enterprise=1:</pr>
crelease</pr>
Release 3 would be: neo4j-enterprise=1:4.0.0~beta03mr03

Refer to the download page for more information regarding the name of packages.

When installing Neo4j Enterprise Edition, you will be prompted to accept the license agreement. Once the license agreement is accepted installation begins. Your answer to the license agreement prompt will be remembered for future installations on the same system.

To forget the stored answer, and trigger the license agreement prompt on subsequent installation, use debconf-communicate to purge the stored answer:

```
echo purge | sudo debconf-communicate neo4j-enterprise
```

Non-interactive installation of Neo4j Enterprise Edition

For Neo4j Enterprise Edition, the license agreement is presented in an interactive prompt. If you require non-interactive installation of Neo4j Enterprise Edition, you can indicate that you have read and accepted the license agreement using <a href="debconf-set-selections">debconf-set-selections</a>:

```
echo "neo4j-enterprise neo4j/question select I ACCEPT" | sudo debconf-set-selections echo "neo4j-enterprise neo4j/license note" | sudo debconf-set-selections
```

#### Upgrade

For upgrade of any 3.x version of Neo4j to 4.0.0, follow instructions in Upgrade.

#### File locations

File locations for all Neo4j packages are documented *here*.

#### Operation

Most Neo4j configuration goes into neo4j.conf.

For operating systems using systemd, some package-specific options are set in *neo4j.service* and can be edited using systemctl\_edit\_neo4j.service.

For operating systems that are not using systemd, some package-specific options are set in /etc/default/neo4j.

Environment variable	Default value	Details
NEO4J_SHUTDOWN_TIMEOUT	120	Timeout in seconds when waiting for Neo4j to stop. If it takes longer than this then the shutdown is considered to have failed. This may need to be increased if the system serves long-running transactions.
NEO4J_ULIMIT_NOFILE	60000	Maximum number of file handles that can be opened by the Neo4j process.

# 2.3.2. Deploy Neo4j using the Neo4j RPM package

This section describes how to deploy Neo4j using the Neo4j RPM package on Red Hat, CentOS, Fedora, or Amazon Linux distributions.

#### This section describes the following:

- Install on Red Hat, CentOS, Fedora or Amazon Linux
  - ☐ Standard installation
  - ☐ Non-interactive installation of Neo4j Enterprise Edition
  - Offline installation
- Install on SUSE

Install on Red Hat, CentOS, Fedora or Amazon Linux

#### Standard installation

1. Add the repository.

Use the following as root to add the repository:

```
rpm --import https://debian.neo4j.com/neotechnology.gpg.key
cat <<EOF> /etc/yum.repos.d/neo4j.repo
[neo4j]
name=Neo4j RPM Repository
baseurl=https://yum.neo4j.com/stable
enabled=1
gpgcheck=1
EOF
```



For milestone, alpha, beta, and release candidate releases, replace the baseurl
line with baseurl=https://yum.neo4j.org/testing

2. Ensure the correct Java version.

Neo4j 4.0 requires the Java 11 runtime. Most of our supported RPM Linux distributions have Java 11 available by default. There is some minor setup required for Amazon Linux, and for compatibility with Oracle Java 11.

☐ Java 11 on Amazon Linux.

To enable OpenJDK 11 on Amazon Linux run the shell command:

```
amazon-linux-extras enable java-openjdk11
```

You are now ready to install Neo4j 4.0.0, which will install Java 11 automatically if it is not already installed.

☐ Oracle Java 11.

Oracle and OpenJDK provide incompatible RPM packages for Java 11. We provide an adapter for Oracle Java 11 which must be installed before Neo4j. The adapter contains no code, but will stop the package manger from installing OpenJDK 11 as a dependency despite an existing Java 11 installation.

This step assumes that you have performed the previous step to set up the yum repository.

- a. Download and install the Oracle java 11 JDK from the Oracle website.
- b. Install the adapter:

```
sudo yum install https://dist.neo4j.org/neo4j-java11-adapter.noarch.rpm
```

The SHA-256 of the adapter package can be verified against https://dist.neo4j.org/neo4j-java11-adapter.noarch.rpm.sha256.

You are now ready to install Neo4j 4.0.0.

- 3. Install Neo4j.
  - ☐ To install Neo4j Community Edition as root:

```
yum install neo4j-4.0.0
```

☐ To install Neo4j Enterprise Edition as root:

```
yum install neo4j-enterprise-4.0.0
```



If you are installing a milestone, alpha, beta, or release candidate release, the name of the package is neo4j-enterprise-<version>-0.<release>.1. For example, Neo4j Enterprise Edition Milestone Release 3 would be: neo4j-enterprise-4.0.0-0.beta03mr03.1

4. Run the following to return the version and edition of Neo4j that has been installed:

```
rpm -qa | grep neo
```

Non-interactive installation of Neo4j Enterprise Edition

When installing Neo4j Enterprise Edition, you will be required to accept the license agreement before installation is allowed to complete. This is an interactive prompt. If you require non-interactive installation of Neo4j Enterprise Edition, you can indicate that you have read and accepted the license agreement by setting the environment variable NEO4J\_ACCEPT\_LICENSE\_AGREEMENT to yes:

```
{\tt NEO4J\_ACCEPT\_LICENSE\_AGREEMENT=yes~yum~install~neo4j-enterprise-4.0.0}
```

#### Offline installation

If you cannot reach <a href="https://yum.neo4j.org/stable">https://yum.neo4j.org/stable</a>, perhaps due to a firewall, you will need to obtain Neo4j via an alternative machine which has the relevant access, and then move the RPM package manually.



It is important to note that using this method will mean that the offline machine will not receive the dependencies that are that are normally downloaded and installed automatically when using RPM for installing Neo4j; Cypher Shell and Java (if not installed already):

- The Cypher Shell RPM package can be downloaded from Neo4j Download Center.
- For information on supported versions of Java, see System requirements.
- 1. Run the following to obtain the required RPM package:
  - ☐ Neo4j Enterprise Edition:

```
curl -O http://yum.neo4j.org/stable/neo4j-enterprise-4.0.0-1.noarch.rpm
```

☐ Neo4j Community Edition:

```
curl -O http://yum.neo4j.org/stable/neo4j-4.0.0-1.noarch.rpm
```

- 2. Manually move the downloaded RPM package to the offline machine.
- 3. Run the following on the offline machine to install Neo4j:

```
rpm -i <rpm file name>
```

### Install on SUSE



SUSE is not certified for production use. These instructions are provided for convenience for those wishing to use Neo4j in non-production environments.

For SUSE-based distributions the steps are as follows:

1. Use the following as root to add the repository:

```
zypper addrepo --refresh https://yum.neo4j.org/stable neo4j-repository
```

- 2. Install Neo4j.
  - ☐ To install Neo4j Community Edition as root:

```
zypper install neo4j-4.0.0
```

☐ To install Neo4j Enterprise Edition as root:

```
zypper install neo4j-enterprise-4.0.0
```

### 2.3.3. Linux tarball installation

This section describes how to install Neo4j on Linux from a tarball, and run it as a console application or service.

#### Unix console application

1. Download the latest release from Neo4j Download Center.

Select the appropriate tar.gz distribution for your platform.

- 2. Check that the SHA hash of the downloaded file is correct:
  - a. To find the correct SHA hash, go to Neo4j Download Center and click on SHA-256 which will be located below your downloaded file.
  - b. Using the appropriate commands for your platform, display the SHA-256 hash for the file that you downloaded.
  - c. Ensure that the two are identical.
- 3. Extract the contents of the archive, using tar -xf <filename>

Refer to the top-level extracted directory as: NEO4J\_HOME

4. Change directory to: \$NEO4J\_HOME

```
Run ./bin/neo4j console
```

5. Stop the server by typing Ctrl-C in the console.

#### Linux service

If you want to run Neo4j as a system service, you can install either the Debian or RPM package.

For more information on configuring and operating the Neo4j system service, see Neo4j system service.

# Setting the number of open files

Linux platforms impose an upper limit on the number of concurrent files a user may have open. This number is reported for the current user and session with the <u>ulimit -n</u> command:

```
user@localhost:~$ ulimit -n
1024
```

It is possible that the default value of 1024 may not be not enough. This is especially true when many indexes are used or a server installation sees too many connections. Network sockets count against the limit as well. Users are therefore encouraged to increase the limit to a healthy value of 40000 or more, depending on usage patterns.

It is possible to set the limit with the <u>ulimit</u> command, but only for the root user, and it only affects the current session. To set the value system wide, follow the instructions for your platform.

The following steps will set the open file descriptor limit to 60000 for user *neo4j* under Ubuntu 16.04 LTS, Debian 8, Centos 7 or later versions of those operating systems:

- 1. If you run Neo4j as a service, you must complete the following:
  - a. Run the following command, which will enable you to edit the *neo4j.service* file:

```
user@localhost:~$ sudo systemctl edit neo4j.service
```

b. Append the [Service] section in the neo4j.service file:

```
[Service]
LimitNOFILE=60000
```

- 2. If you run Neo4j as an interactive user (which can be the case for testing purposes), you must complete the following:
  - a. Run the following command, which will enable you to edit the *user.conf* file:

```
user@localhost:~$ sudo vi /etc/systemd/user.conf
```

b. Uncomment and define the value for DefaultLimitNOFILE, found in the [MANAGER] section:

```
[Manager]
...
DefaultLimitNOFILE=60000
```

c. Run the following command which will append the /etc/security/limits.conf file:

```
user@localhost:~$ sudo vi /etc/security/limits.conf
neo4j soft nofile 60000
neo4j hard nofile 60000
```

3. Reload systemd settings with the command:

```
user@localhost:~$ sudo systemctl daemon-reload
```

4. Reboot your machine.

After completing the above procedure, the Neo4j user will have a limit of 60000 simultaneous open files. If you continue experiencing exceptions on Too many open files or Could not stat() directory, you may have to raise the limit further.

# 2.3.4. Neo4j system service

This article covers configuring and operating the Neo4j system service. It assumes that your system has systemd, which is the case for most Linux distributions.



Setting the number of open files.

For instructions on how to set the number of concurrent files that a user can have open, see Setting the number of open files.

# Configuration

Configuration is stored in /etc/neo4j/neo4j.conf. See File locations for a complete catalog of where files are found for the various packages.

Starting the service automatically on system start

If you installed the RPM package and want Neo4j to start automatically on system boot then you need to enable the service. On Debian-based distributions this is done for you at installation time.

```
systemctl enable neo4j
```

### Controlling the service

System services are controlled with the systemctl command. It accepts a number of commands:

```
systemctl {start|stop|restart} neo4j
```

Service customizations can be placed in a service override file. To edit your specific options, do the following command which will open up an editor of the appropriate file:

```
systemctl edit neo4j
```

Then place any customizations under a [Service] section. The following example lists default values which may be interesting to change for some users:

```
[Service]
# The user and group which the service runs as.
User=neo4j
Group=neo4j
# If it takes longer than this then the shutdown is considered to have failed.
# This may need to be increased if the system serves long-running transactions.
TimeoutSec=120
```

You can print the effective service, including possible overrides, with:

```
systemctl cat neo4j
```

Remember to restart neo4j if you change any settings.

```
systemctl restart neo4j
```

#### Log

The neo4j log is written to journald which can be viewed using the journalctl command:

```
journalctl -e -u neo4j
```

journald automatically rotates the log after a certain time and by default it commonly does not persist across reboots. Please see man journald.conf for further details.

# 2.4. macOS installation

This section describes how to install Neo4j on macOS.

# 2.4.1. Unix console application

1. Download the latest release from Neo4j Download Center.

Select the appropriate tar.gz distribution for your platform.

- 2. Check that the SHA hash of the downloaded file is correct:
  - a. To find the correct SHA hash, go to Neo4j Download Center and click on SHA-256 which will be located below your downloaded file.
  - b. Using the appropriate commands for your platform, display the SHA-256 hash for the file that you downloaded.
  - c. Ensure that the two are identical.
- 3. Extract the contents of the archive, using tar -xf <filename>

Refer to the top-level extracted directory as: NEO4J\_HOME

4. Change directory to: \$NEO4J\_HOME

Run ./bin/neo4j console

5. Stop the server by typing Ctrl-C in the console.

When Neo4j runs in console mode, logs are printed to the terminal.

#### 2.4.2. macOS service

Use the standard macOS system tools to create a service based on the neo4j command.

# 2.4.3. macOS file descriptor limits

The limit of open file descriptors may have to be increased if a database has many indexes or if there are many connections to the database. The currently configured open file descriptor limitation on your macOS system can be inspected with the launchctl limit maxfiles command. The method for changing the limit may differ depending on the version of macOS. Consult the documentation for your operating system in order to find out the appropriate command.

If you raise the limit above 10240, then you must also add the following setting to your *neo4j.conf* file:

```
dbms.jvm.additional=-XX:-MaxFDLimit
```

Without this setting, the file descriptor limit for the JVM will not be increased beyond 10240. Note, however, that this only applies to macOS. On all other operating systems, you should always leave the MaxFDLimit JVM setting enabled.

# 2.5. Windows installation

This section describes how to install Neo4j on Windows.

# 2.5.1. Windows console application

1. Download the latest release from Neo4j Download Center.

Select the appropriate ZIP distribution.

- 2. Check that the SHA hash of the downloaded file is correct:
  - a. To find the correct SHA hash, go to Neo4j Download Center and click on SHA-256 which will be located below your downloaded file.

- b. Using the appropriate commands for your platform, display the SHA-256 hash for the file that you downloaded.
- c. Ensure that the two are identical.
- 3. Right-click the downloaded file, click Extract All.
- 4. Change directory to the top-level extracted directory.

Run bin\neo4j console

5. Stop the server by typing Ctrl-C in the console.

#### 2.5.2. Windows service

Neo4j can also be run as a Windows service. Install the service with bin\neo4j install-service, and start it with bin\neo4j start.

The available commands for bin\neo4j are: help, start, stop, restart, status, install-service, uninstall-service, and update-service.

#### Java options

When Neo4j is installed as a service, Java options are stored in the service configuration. Changes to these options after the service is installed will not take effect until the service configuration is updated. For example, changing the setting <a href="mailto:dbms.memory.heap.max\_size">dbms.memory.heap.max\_size</a> in <a href="mailto:neo4j.conf">neo4j.conf</a> will not take effect until the service is updated and restarted. To update the service, run <a href="mailto:bin\neo4j">bin\neo4j</a> update-service. Then restart the service to run it with the new configuration.

The same applies to the path to where Java is installed on the system. If the path changes, for example when upgrading to a new version of Java, it is necessary to run the update-service command and restart the service. Then the new Java location will be used by the service.

#### Example 1. Update service example

١.	Install service
	bin\neo4j install-service
2.	Change memory configuration
	<pre>echo dbms.memory.heap.initial_size=8g &gt;&gt; conf\neo4j.conf echo dbms.memory.heap.max_size=16g &gt;&gt; conf\neo4j.conf</pre>
3.	Update service
	bin\neo4j update-service
4.	Restart service
	bin\neo4j restart

### 2.5.3. Windows PowerShell module

The Neo4j PowerShell module allows administrators to:

- Install, start and stop Neo4j Windows® Services.
- Start tools, such as Neo4j Admin and Cypher Shell.

The PowerShell module is installed as part of the ZIP file distributions of Neo4j.

#### System requirements

- Requires PowerShell v2.0 or above.
- Supported on either 32 or 64 bit operating systems.

#### Managing Neo4j on Windows

On Windows, it is sometimes necessary to *Unblock* a downloaded ZIP file before you can import its contents as a module. If you right-click on the ZIP file and choose "Properties" you will get a dialog which includes an "Unblock" button, which will enable you to import the module.

Running scripts has to be enabled on the system. This can, for example, be achieved by executing the following from an elevated PowerShell prompt:

```
Set-ExecutionPolicy -ExecutionPolicy RemoteSigned
```

For more information, see About execution policies.

The PowerShell module will display a warning if it detects that you do not have administrative rights.

### How do I import the module?

The module file is located in the *bin* directory of your Neo4j installation, i.e. where you unzipped the downloaded file. For example, if Neo4j was installed in *C:\Neo4j* then the module would be imported like this:

```
Import-Module C:\Neo4j\bin\Neo4j-Management.psd1
```

This will add the module to the current session.

Once the module has been imported you can start an interactive console version of a Neo4j Server like this:

```
Invoke-Neo4j console
```

To stop the server, issue Ctrl-C in the console window that was created by the command.

How do I get help about the module?

Once the module is imported you can query the available commands like this:

```
Get-Command -Module Neo4j-Management
```

The output should be similar to the following:

ommandType	Name	Version	Source
Function	Invoke-Neo4j	4.0.0	Neo4j-Management
Function	Invoke-Neo4jAdmin	4.0.0	Neo4j-Management
Function	Invoke-Neo4jBackup	4.0.0	Neo4j-Management
Function	Invoke-Neo4jImport	4.0.0	Neo4j-Management
Function	Invoke-Neo4jShell	4.0.0	Neo4j-Management

The module also supports the standard PowerShell help commands.

```
Get-Help Invoke-Neo4j
```

Run the following to see examples of help commands:

```
Get-Help Invoke-Neo4j -examples
```

# Example usage

• List of available commands:

```
Invoke-Neo4j
```

• Current status of the Neo4j service:

```
Invoke-Neo4j status
```

• Install the service with verbose output:

```
Invoke-Neo4j install-service -Verbose
```

• Available commands for administrative tasks:

```
Invoke-Neo4jAdmin
```

# Common PowerShell parameters

The module commands support the common PowerShell parameter of Verbose.

# Chapter 3. Docker

This chapter describes how run Neo4j in a Docker container.

This chapter describes the following:

- Introduction Introduction to running Neo4j in a Docker container.
- Configuration How to configure Neo4j to run in a Docker container.
- Clustering How to set up Causal Clustering when using Docker.
- Docker specific operations Descriptions of various operations that are specific to using Docker.
- Security Information about using encryption with the Docker image.



Docker does not run natively on macOS or Windows. For running Docker on macOS and Windows, please consult the documentation provided by Docker.

# 3.1. Introduction

An introduction to how Neo4j runs in a Docker container.

The Neo4j Docker image, and instructions on how to start using it, can be found here: https://hub.docker.com/\_/neo4j/.

#### 3.1.1. Ports

By default, the Docker image exposes three ports for remote access:

- 7474 for HTTP.
- 7473 for HTTPS.
- 7687 for Bolt.

Note that when Docker is used, Neo4j is configured automatically to allow remote access to the HTTP, HTTPS, and Bolt services. This is different than the default Neo4j configuration, where the HTTP, HTTPS, and Bolt services do not allow remote connections. For more information on configuring connections, see Configure connectors.

#### 3.1.2. Volumes

The Docker image also exposes the following volumes. Directories on the host can be mounted using the --volume option. See File locations for details about the directories used by default in different Neo4j distributions.

- /conf
- /data
- /import
- /logs
- /metrics
- /plugins
- /ssl

It is often desirable to keep database and logs outside of the container. The following command will start a container with ports for Bolt and HTTP published, and with the /data and /logs volumes mapped to directories on the host.

```
docker run \
    --publish=7474:7474 --publish=7687:7687 \
    --volume=$HOME/neo4j/data:/data \
    --volume=$HOME/neo4j/logs:/logs \
    neo4j:4.0
```

Point your browser at http://localhost:7474 on Linux or http://\$(docker-machine ip default):7474 on macOS.

All the volumes in this documentation are stored under \$HOME in order to work on macOS where \$HOME is automatically mounted into the machine VM. On Linux the volumes can be stored anywhere.



By default Neo4j requires authentication and requires you to login with neo4j/neo4j at the first connection and set a new password. You can set the password for the Docker container directly by specifying --env NEO4J\_AUTH=neo4j/<password> in your run directive. Alternatively, you can disable authentication by specifying --env NEO4J\_AUTH=none instead.

### 3.1.3. Running Neo4j as a non-root user

For security reasons Neo4j will run as the neo4j user inside the container. You can specify which user to run as by invoking docker with the --user argument. For example, the following would run Neo4j as your current user:

```
docker run \
    --publish=7474:7474 --publish=7687:7687 \
    --volume=$HOME/neo4j/data:/data \
    --volume=$HOME/neo4j/logs:/logs \
    --user="$(id -u):$(id -g)" \
    neo4j:4.0
```

# 3.1.4. Neo4j editions

Tags are available for both Community Edition and Enterprise Edition. Version-specific Enterprise Edition tags have an -enterprise suffix, for example: neo4j:4.0.0-enterprise. Community Edition tags have no suffix, for example neo4j:4.0.0. The latest Neo4j Enterprise Edition release is available as neo4j:enterprise.

# Neo4j Enterprise Edition license

In order to use Neo4j Enterprise Edition you must accept the license agreement.

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Email inquiries can be directed to: licensing@neo4j.com

More information is also available at: https://neo4j.com/licensing/

To accept the license agreement set the environment variable NEO4J\_ACCEPT\_LICENSE\_AGREEMENT=yes.

To do this you can use the following docker argument:

```
--env NEO4J_ACCEPT_LICENSE_AGREEMENT=yes
```

# 3.2. Configuration

This chapter describes how configure Neo4j to run in a Docker container.

The default configuration provided by this image is intended for learning about Neo4j, but must be modified to make it suitable for production use. In particular, the default memory assignments to Neo4j are very limited (NEO4J\_dbms\_memory\_pagecache\_size=512M and NEO4J\_dbms\_memory\_heap\_max\_\_size=512M), to allow multiple containers to be run on the same server. You can read more about configuring Neo4j in the Configuration settings.

There are three ways to modify the configuration:

- Set environment variables.
- Mount a /conf volume.
- · Build a new image.

Which one to choose depends on how much you need to customize the image.

#### 3.2.1. Environment variables

Pass environment variables to the container when you run it.

```
docker run \
    --detach \
    --publish=7474:7474 --publish=7687:7687 \
    --volume=$HOME/neo4j/data:/data \
    --volume=$HOME/neo4j/logs:/logs \
    --env NEO4J_dbms_memory_pagecache_size=4G \
    neo4j:4.0
```

Any configuration value (see Configuration settings) can be passed using the following naming scheme:

- Prefix with NEO4J\_.
- Underscores must be written twice: \_ is written as \_\_.
- Periods are converted to underscores: . is written as \_.

As an example, <a href="mailto:dbms.tx\_log.rotation.size">dbms.tx\_log.rotation.size</a> could be set by specifying the following argument to Docker:

```
--env NEO4J_dbms_tx__log_rotation_size
```

Variables which can take multiple options, such as <a href="mailto:dbms\_jvm\_additional">dbms\_jvm\_additional</a>, must be defined just once, and include a concatenation of the multiple values. For example:

```
--env NEO4J_dbms_jvm_additional="-Dcom.sun.management.jmxremote.authenticate=true -Dcom.sun.management.jmxremote.ssl=false -Dcom.sun.management.jmxremote.password.file= $HOME/conf/jmx.password -Dcom.sun.management.jmxremote.access.file=$HOME/conf/jmx.access -Dcom.sun.management.jmxremote.port=3637"
```

#### Neo4j Enterprise Edition

The following environment variables are specific to Causal Clustering, and are available in the Neo4j Enterprise Edition:

- NEO4J\_dbms\_mode: the database mode, defaults to SINGLE, set to CORE or READ\_REPLICA for Causal Clustering.
- NEO4J\_causal\_\_clustering\_expected\_\_core\_\_cluster\_\_size: the initial cluster size (number of Core instances) at startup.
- NEO4J\_causal\_\_clustering\_initial\_\_discovery\_\_members: the network addresses of an initial set of Core cluster members.
- NEO4J\_causal\_\_clustering\_discovery\_\_advertised\_\_address: hostname/ip address and port to advertise for member discovery management communication.
- NEO4J\_causal\_\_clustering\_transaction\_\_advertised\_\_address: hostname/ip address and port to advertise for transaction handling.
- NEO4J\_causal\_\_clustering\_raft\_\_advertised\_\_address: hostname/ip address and port to advertise for cluster communication.

See Clustering for examples of how to configure Causal Clustering on Docker.

# 3.2.2. /conf volume

To make arbitrary modifications to the Neo4j configuration, provide the container with a /conf volume.

```
docker run \
    --detach \
    --publish=7474:7474 --publish=7687:7687 \
    --volume=$HOME/neo4j/data:/data \
    --volume=$HOME/neo4j/logs:/logs \
    --volume=$HOME/neo4j/conf:/conf \
    neo4j:4.0
```

Any configuration files in the /conf volume will override files provided by the image. So if you want to change one value in a file you must ensure that the rest of the file is complete and correct. Environment variables passed to the container by Docker will still override the values in configuration files in /conf volume.



If you use a configuration volume you must make sure to listen on all network interfaces. This can be done by setting dbms.default\_listen\_address=0.0.0.0.

To dump an initial set of configuration files, run the image with the dump-config command.

```
docker run --rm \
    --volume=$HOME/neo4j/conf:/conf \
    neo4j:4.0 dump-config
```

# 3.2.3. Building a new image

For more complex customization of the image you can create a new image based on this one.

```
FROM neo4j:4.0
```

If you need to make your own configuration changes, we provide a hook so you can do that in a script:

```
COPY extra_conf.sh /extra_conf.sh
```

Then you can pass in the EXTENSION\_SCRIPT environment variable at runtime to source the script:

```
docker run -e "EXTENSION_SCRIPT=/extra_conf.sh" cafe12345678
```

When the extension script is sourced, the current working directory will be the root of the Neo4j installation.

# 3.3. Clustering

This chapter describes how to set up Causal Clustering when running Neo4j in a Docker container.

# 3.3.1. Setting up a Causal Cluster

In order to run Neo4j in CC mode under Docker you need to wire up the containers in the cluster so that they can talk to each other. Each container must have a network route to each of the others and the NEO4J\_causal\_\_clustering\_expected\_core\_\_cluster\_\_size and NEO4J\_causal\_\_clustering\_initial\_\_discovery\_\_members environment variables must be set for cores. Read Replicas only need to define NEO4J\_causal\_\_clustering\_initial\_\_discovery\_\_members.

Within a single Docker host, this can be achieved as follows. Note that the default ports for HTTP, HTTPS and Bolt are used. For each container, these ports are mapped to a different set of ports on the Docker host.

```
docker network create --driver=bridge cluster
docker run --name=core1 --detach --network=cluster \
    --publish=7474:7474 --publish=7473:7473 --publish=7687:7687 \
    --env NEO4J_dbms_mode=CORE \
    --env NEO4J_causal__clustering_expected__core__cluster__size=3 \
    --env NEO4J_causal__clustering_initial__discovery__members=core1:5000,core2:5000,core3:5000 \
--env NEO4J_ACCEPT_LICENSE_AGREEMENT=yes \
    --env NEO4J_dbms_connector_bolt_advertised__address=localhost:7687 \
    --env NEO4J_dbms_connector_http_advertised__address=localhost:7474 \
    neo4j:4.0-enterprise
docker run --name=core2 --detach --network=cluster \
    --publish=8474:7474 --publish=8473:7473 --publish=8687:7687 \
    --env NEO4J_dbms_mode=CORE \
    --env NEO4J_causal__clustering_expected__core__cluster__size=3 \
    --env NEO4J_causal__clustering_initial__discovery__members=core1:5000,core2:5000,core3:5000 \
    --env NEO4J_ACCEPT_LICENSE_AGREEMENT=yes \
    --env NEO4J_dbms_connector_bolt_advertised__address=localhost:8687 \
    --env NEO4J_dbms_connector_http_advertised__address=localhost:8474 \
    neo4j:4.0-enterprise
docker run --name=core3 --detach --network=cluster \
    --publish=9474:7474 --publish=9473:7473 --publish=9687:7687 \
    --env NEO4J_dbms_mode=CORE
    --env NEO4J_causal__clustering_expected__core__cluster__size=3 \
--env NEO4J_causal__clustering_initial__discovery__members=core1:5000,core2:5000,core3:5000 \
    --env NEO4J_ACCEPT_LICENSE_AGREEMENT=yes \
    --env NEO4J_dbms_connector_bolt_advertised__address=localhost:9687 \
    --env NEO4J_dbms_connector_http_advertised__address=localhost:9474 \
    neo4j:4.0-enterprise
```

Additional instances can be added to the cluster in an ad-hoc fashion. A Read Replica can for example be added with:

When each container is running on its own physical machine and Docker network is not used, it is necessary to define the advertised addresses to enable communication between the physical machines. Each container should also bind to the host machine's network.

Each instance would then be invoked similar to:

```
docker run --name=neo4j-core --detach \
         --network=host
         --publish=7474:7474 --publish=7687:7687 \
         --publish=5000:5000 --publish=6000:6000 --publish=7000:7000 \
         --env NEO4J_dbms_mode=CORE \
         --env NEO4J_causal__clustering_expected__core__cluster__size=3 \
         --env NEO4J_causal__clustering_initial__discovery__members=core1-public-address:5000,core2-
public-address:5000,core3-public-address:5000 \
         --env NEO4J_causal__clustering_discovery__advertised__address=public-address:5000
         --env NEO4J_causal__clustering_transaction__advertised__address=public-address:6000 \
         --env NEO4J_causal__clustering_raft__advertised__address=public-address:7000 \
         --env NEO4J_dbms_connectors_default__advertised__address=public-address \
         --env NEO4J_ACCEPT_LICENSE_AGREEMENT=yes \
         --env NEO4J_dbms_connector_bolt_advertised__address=public-address:7687 \
         --env NEO4J_dbms_connector_http_advertised__address=public-address:7474 \
         neo4j:4.0-enterprise
```

Where public-address is the public hostname or ip-address of the machine.

See Deploy a cluster for more details of Neo4j Causal Clustering.

# 3.4. Docker specific operations

This chapter describes various operations that are specific to running Neo4j in a Docker container.

### 3.4.1. Using /plugins

To install user-defined procedures, provide a /plugins volume containing the jars.

```
docker run --publish=7474:7474 --publish=7687:7687 --volume=$HOME/neo4j/plugins:/plugins neo4j:4.0
```

See Java Reference | Procedures for more details on procedures.

# 3.4.2. Using Cypher Shell

The Neo4j Shell can be run locally within a container using a command like this:

```
docker exec --interactive --tty <container> bin/cypher-shell
```

To determine the *<container>*, you can run docker ps to list the currently running Docker containers. From the results, use the CONTAINER ID in place of *<container>*.

To pass a local file into Cypher Shell, you can run the following, again using the CONTAINER ID identified above to replace <container>:

```
cat <local-file> | sudo docker exec --interactive <container> bin/cypher-shell -u <username> -p <password>
```

When you run the above, the contents of the *<local-file>* will be passed into the Docker container.

It will then run bin/cypher-shell, and authenticate with the provided <username> and <password>.

For more information on using Cypher Shell, see Cypher Shell.

## 3.4.3. Upgrading Neo4j on Docker

To enable upgrades, set NEO4J\_dbms\_allow\_upgrade to true. For more details on upgrading, see:

- Single-instance upgrade
- Upgrade a Neo4j Causal Cluster

# 3.5. Security

This chapter describes security in Neo4j when running in a Docker container.

# 3.5.1. Encryption

The Docker image can expose Neo4j's native TLS support. To use your own key and certificate, provide an /ssl volume with the key and certificate inside. The files must be called *neo4j.key* and *neo4j.cert*. You must also publish port 7473 to access the HTTPS endpoint.

```
docker run --publish=7473:7473 --publish=7687:7687 --volume=$HOME/neo4j/ssl:/ssl neo4j:4.0
```

# Chapter 4. Configuration

This chapter describes the configuration of Neo4j components.

### The topics described are:

- The neo4j.conf file An introduction to the primary configuration file in Neo4j.
- File locations An overview of where files are stored in the different Neo4j distributions and the necessary file permissions for running Neo4j.
- Ports An overview of the ports relevant to a Neo4j installation.
- Set initial password How to set an initial password.
- Password and user recovery How to recover after a lost admin password.
- Configure Neo4j connectors How to configure Neo4j connectors.
- Configure dynamic settings How to configure certain Neo4j parameters while Neo4j is running.
- Transaction logs How to configure transaction logs.

For a complete reference of Neo4j configuration settings, see Configuration settings.

# 4.1. The neo4j.conf file

This section introduces the neo4j.conf file, and its syntax.

### This section contains the following:

- Introduction
- Syntax
- JVM-specific configuration settings
- List currently active settings

For a complete reference of Neo4j configuration settings, see Configuration settings.

## 4.1.1. Introduction

The *neo4j.conf* file is the main source of configuration settings in Neo4j, and includes the mappings of configuration setting keys to values. The location of the *neo4j.conf* file in the different configurations of Neo4j is listed in The locations of important files.

Most of the configuration settings in the *neo4j.conf* file apply directly to Neo4j itself, but there are also some settings which apply to the Java Runtime (the JVM) on which Neo4j runs. For more information, see the JVM specific configuration settings below. Many of the configuration settings are also used by the neo4j launcher scripts.

## 4.1.2. Syntax

- The equals sign (=) maps configuration setting keys to configuration values.
- Lines that start with the number sign (#) are handled as comments.
- · Empty lines are ignored.
- Configuring a setting in neo4j.conf will overwrite any default values. In case a setting can define a

list of values, and you wish to amend the default values with custom values, you will have to explicitly list the default values along with the new values.

• There is no order for configuration settings, and each setting in the *neo4j.conf* file must be uniquely specified. If you have multiple configuration settings with the same key, but different values, this can lead to unpredictable behavior.

The only exception to this is dbms.jvm.additional. If you set more than one value for dbms.jvm.additional, then each setting value will add another custom JVM argument to the java launcher.

## 4.1.3. JVM-specific configuration settings

```
dbms.memory.heap.initial_size
```

- dbms.memory.heap.max\_size
- dbms.jvm.additional

# 4.1.4. List currently active settings

You can use the procedure <a href="mailto:dbms.listConfig">dbms.listConfig</a>() to list the currently active configuration settings and their values.

Example 2. List currently active configuration settings

See also Dynamic settings for information about dynamic settings.

## 4.2. File locations

This section provides an overview of where files are stored in the different Neo4j distributions, and the necessary file permissions for running Neo4j.

### This section describes the following:

- Where to find important files
- Log files
- File permissions

# 4.2.1. Where to find important files

The table below lists the location of important files, and whether the location is customizable:

Table 9. The locations of important files

Package	Configurati on [1]	Data <sup>[2]</sup>	Logs	Metrics	Import	Bin	Lib	Plugins
Linux or macOS tarball	<neo4j- home&gt;/conf /neo4j.conf</neo4j- 	<neo4j- home&gt;/data</neo4j- 	<neo4j- home&gt;/logs</neo4j- 	<neo4j- home&gt;/metr ics</neo4j- 	<neo4j- home&gt;/imp ort</neo4j- 	<neo4j- home&gt;/bin</neo4j- 	<neo4j- home&gt;/lib</neo4j- 	<neo4j- home&gt;/plugi ns</neo4j- 
Windows zip	<neo4j- home&gt;\conf \neo4j.conf</neo4j- 	<neo4j- home&gt;\data</neo4j- 	<neo4j- home&gt;\logs</neo4j- 	<neo4j- home&gt;\metr ics</neo4j- 	<neo4j- home&gt;\imp ort</neo4j- 	<neo4j- home&gt;\bin</neo4j- 	<neo4j- home&gt;\lib</neo4j- 	<neo4j- home&gt;\plugi ns</neo4j- 
Debian	/etc/neo4j/n eo4j.conf	/var/lib/neo 4j/data	/var/log/neo 4j/ <sup>[3]</sup>	/var/lib/neo 4j/metrics	/var/lib/neo 4j/import	/usr/bin	/usr/share/n eo4j/lib	/var/lib/neo 4j/plugins
RPM	/etc/neo4j/n eo4j.conf	/var/lib/neo 4j/data	/var/log/neo 4j/ <sup>[3]</sup>	/var/lib/neo 4j/metrics	/var/lib/neo 4j/import	/usr/bin	/usr/share/n eo4j/lib	/var/lib/neo 4j/plugins
Neo4j Desktop <sup>[4]</sup>		Use the Ope	<i>n Folder</i> butto	n in the appli	cation to loca	te the relevar	nt directories.	
Customizab le by option		dbms.direct ories.data	dbms.direct ories.logs	dbms.direct ories.metri cs	dbms.direct ories.impor t	Not applicable	dbms.direct ories.lib	dbms.direct ories.plugi ns

<sup>[1]</sup> For details about *neo4j.conf*, see: The neo4j.conf file.

The locations of <neo4j-home> and conf can be configured using environment variables, as described below:

Table 10. Configuration of <neo4j-home> and conf

Location	Default	Environment variable	Notes
<neo4j-home></neo4j-home>	Root directory of the installation	NEO4J_HOME	Must be an absolute path.
conf	<neo4j-home>/conf</neo4j-home>	NEO4J_CONF	Must be set explicitly if it is not a subdirectory of <neo4j-home>.</neo4j-home>

# 4.2.2. Log files

Filename	Description
neo4j.log	The standard log, where general information about Neo4j is written. Not written for Debian and RPM packages. See relevant sections.
debug.log	Information useful when debugging problems with Neo4j.
http.log	Request log for the HTTP API.
gc.log	Garbage Collection logging provided by the JVM.

<sup>&</sup>lt;sup>[2]</sup> Please note that the data directory is internal to Neo4j and its structure is subject to change between versions without notice.

<sup>[3]</sup> To view the *neo4j.log* for Debian and RPM, use journalctl --unit=neo4j.

<sup>[4]</sup> Applicable to all operating systems where Neo4j Desktop is supported.

Filename	Description
query.log	Log of executed queries (Enterprise Edition only).
security.log	Log of security events (Enterprise Edition only).
service-error.log	Log of errors encountered when installing or running the Windows service. (Windows only.)

# 4.2.3. File permissions

The user that Neo4j runs as must have the following permissions:

## Read only

- conf
- import
- bin
- lib
- plugins

### Read and write

- data
- logs
- metrics

### Execute

• all files in bin

## 4.3. Ports

This section lists ports relevant to a Neo4j installation.

This section provides an overview for determining which Neo4j-specific ports should be opened up in your firewalls. Note that these ports are in addition to those necessary for ordinary network operation. Specific recommendations on port openings cannot be made, as the firewall configuration must be performed taking your particular conditions into consideration.

Name	Default port number	Related settings	Comments
Backups	6362-6372	dbms.backup.enabled dbms.backup.listen_address	Backups are enabled by default. In production environments, external access to the backup port(s) should be blocked by a firewall. See also Backup.
НТТР	7474	See Configure connectors.	It is recommended to not open up this port for external access in production environments, since traffic is unencrypted. Used by Neo4j Browser.
HTTPS	7473	See Configure connectors.	
Bolt	7687	See Configure connectors.	Used by Cypher Shell and by Neo4j Browser.

Name	Default port number	Related settings	Comments
Causal Cluster	5000, 6000, 7000	causal_clustering.discover y_listen_address causal_clustering.transact ion_listen_address causal_clustering.raft_lis ten_address	The listed ports are the default ports in neo4j.conf. The ports are likely be different in a production installation; therefore the potential opening of ports must be modified accordingly. See also Settings reference.
Graphite monitoring	2003	metrics.graphite.server	This is an outbound connection in order for the Neo4j database to communicate with the Graphite server. See also Metrics.
Prometheus monitoring	2004	metrics.prometheus.enabled and metrics.prometheus.endpoin t	See also Metrics.
JMX monitoring	3637	dbms.jvm.additional=- Dcom.sun.management.jmxrem ote.port=3637	This setting is for exposing the JMX. This is not the recommended way of inspecting a Neo4j database. It is not enabled by default.

# 4.4. Set an initial password

This section describes how to set an initial password for Neo4j.

Use the set-initial-password command of neo4j-admin to define the password for the native user neo4j. This must be performed before starting up the database for the first time.

### Syntax:

neo4j-admin set-initial-password <password> [--require-password-change]

Example 3. Use the set-initial-password command of neo4j-admin

Set the password for the native neo4j user to 'h6u4%kr' before starting the database for the first time.

\$neo4j-home> bin/neo4j-admin set-initial-password h6u4%kr

Example 4. Use the set-initial-password command of neo4j-admin with the optional --require-password -change flag

Set the password for the native neo4j user to 'secret' before starting the database for the first time. You will be prompted to change this password to one of your own choice at first login.

\$neo4j-home> bin/neo4j-admin set-initial-password secret --require-password-change

If the password is not set explicitly using this method, it will be set to the default password neo4j. In that case, you will be prompted to change the default password at first login.

# 4.5. Password and user recovery

This section describes how to recover from a lost password, specifically for an admin user.

It is recommended to block network connections during the recovery phase, so users should connect to Neo4j only via localhost. This can be achieved by temporarily commenting/setting this parameter in neo4j.conf:



#dbms.connectors.default\_listen\_address=your\_configuration

or

dbms.connectors.default\_listen\_address=127.0.0.1

Use the following steps to set a new password (assuming your admin user is named neo4j):

1. Stop Neo4j

\$ bin/neo4j stop

2. Modify neo4j.conf:

dbms.security.auth\_enabled=false

3. Start Neo4j

\$ bin/neo4j start

4. Connect via Cypher Shell to modify the admin user password:

```
$ bin/cypher-shell -d system
neo4j@system> ALTER USER neo4j SET PASSWORD 'mynewpass';
neo4j@system> :exit
```

5. Stop Neo4j

\$ bin/neo4j stop

6. Modify neo4j.conf by commenting out dbms.security.auth\_enabled like this:

# dbms.security.auth\_enabled=false

or setting dbms.security.auth\_enabled to true like this:

dbms.security.auth\_enabled=true

7. Start Neo4j

\$ bin/neo4j start

# 4.6. Configure connectors

This section describes how to configure connectors for Neo4j.

This section describes the following:

- Available connectors
- Configuration options
- · Options for Bolt thread pooling
- Defaults for addresses

## 4.6.1. Available connectors

The table below lists the available Neo4j connectors:

Table 11. Default connectors and their ports

Connector name	Protocol	Default port number
dbms.connector.bolt	Bolt	7687
dbms.connector.http	НТТР	7474
dbms.connector.https	HTTPS	7473

When configuring the HTTPS or Bolt connector, see also SSL framework for details on how to work with SSL certificates.

# 4.6.2. Configuration options

The connectors are configured by settings on the format dbms.connector.<connector-name>.<setting-suffix>>. The available suffixes are described in the table below:

Table 12. Configuration option suffixes for connectors

Option name	Default	Setting(s)	Description
enabled	true <sup>[1]</sup>	dbms.connector.bolt.enabled, dbms.connector.http.enabled, dbms.connector.https.enabled [2]	This setting allows the client connector to be enabled or disabled. When disabled, Neo4j does not listen for incoming connections on the relevant port.
			<sup>[1]</sup> When Neo4j is used in embedded mode, the default value is false.
			<sup>[2]</sup> The default value for dbms.connector.https.enabled is false.
listen_ad dress	127.0.0.1: <conne ctor-default- port&gt;</conne 	<pre>dbms.connector.bolt.listen_address, dbms.connector.https.listen_address, dbms.connector.http.listen_address</pre>	This setting specifies how Neo4j listens for incoming connections. It consists of two parts; an IP address (e.g. 127.0.0.1 or 0.0.0.0) and a port number (e.g. 7687), and is expressed in the format <ip-address>:<port-number>. See below for an example of usage.</port-number></ip-address>

Option name	Default	Setting(s)	Description
advertise d_address	<pre>localhost:<conne ctor-default-="" port=""></conne></pre>	<pre>dbms.connector.bolt.advertised_address, dbms.connector.https.advertised_address, dbms.connector.http.advertised_address</pre>	This setting specifies the address that clients should use for this connector. This is useful in a Causal Cluster as it allows each server to correctly advertise addresses of the other servers in the cluster. The advertised address consists of two parts; an address (fully qualified domain name, hostname, or IP address) and a port number (e.g. 7687), and is expressed in the format <address>:<port-number>. See below for an example of usage.</port-number></address>
tls_level	DISABLED	dbms.connector.bolt.tls_level	This setting is only applicable to the Bolt connector. It allows the connector to accept encrypted and/or unencrypted connections. The default value is DISABLED, where only unencrypted client connections are to be accepted by this connector, and all encrypted connections will be rejected.  Other values are REQUIRED and OPTIONAL. Use REQUIRED when only encrypted client connections are to be accepted by this connector, and all unencrypted connections will be rejected. Use OPTIONAL where either encrypted or unencrypted client connections are accepted by this connector.

### Example 5. Specify listen\_address for the Bolt connector

To listen for Bolt connections on all network interfaces (0.0.0.0) and on port 7000, set the listen\_address for the Bolt connector:

 ${\tt dbms.connector.bolt.listen\_address=0.0.0.0:7000}$ 

### Example 6. Specify advertised\_address for the Bolt connector

If routing traffic via a proxy, or if port mappings are in use, it is possible to specify <a href="mailto:advertised\_address">advertised\_address</a> for each connector individually. For example, if port 7687 on the Neo4j Server is mapped from port 9000 on the external network, specify the <a href="mailto:advertised\_address">advertised\_address</a> for the Bolt connector:

dbms.connector.bolt.advertised\_address=<server-name>:9000

# 4.6.3. Options for Bolt thread pooling

See Bolt thread pool configuration to learn more about Bolt thread pooling and how to configure it on the connector level.

## 4.6.4. Defaults for addresses

It is possible to specify defaults for the configuration options with listen\_address and advertised\_address suffixes, as described below. Setting a default value will apply to all the connectors, unless specifically configured for a certain connector.

dbms.default\_listen\_address

This configuration option defines a default IP address of the settings with the <u>listen\_address</u> suffix for all connectors. If the IP address part of the <u>listen\_address</u> is not specified, it is inherited from the shared setting <u>dbms.default\_listen\_address</u>.

### Example 7. Specify listen\_address for the Bolt connector

To listen for Bolt connections on all network interfaces (0.0.0.0) and on port 7000, set the listen\_address for the Bolt connector:

```
dbms.connector.bolt.listen_address=0.0.0.0:7000
```

This is equivalent to specifying the IP address by using the <a href="mailto:default\_listen\_address">default\_listen\_address</a> setting, and then specifying the port number for the Bolt connector.

```
dbms.default_listen_address=0.0.0.0
dbms.connector.bolt.listen_address=:7000
```

### dbms.default\_advertised\_address

This configuration option defines a default address of the settings with the <a href="advertised\_address">advertised\_address</a> is not specified, it is inherited from the shared setting <a href="advertised\_address">advertised\_address</a>.

### Example 8. Specify advertised\_address for the Bolt connector

Specify the address that clients should use for the Bolt connector:

```
dbms.connector.bolt.advertised_address=server1:9000
```

This is equivalent to specifying the address by using the <a href="mailto:default\_advertised\_address">default\_advertised\_address</a> setting, and then specifying the port number for the Bolt connector.

```
dbms.default_advertised_address=server1
dbms.connector.bolt.advertised_address=:9000
```



The default address settings can only accept the hostname or IP address portion of the full socket address. Port numbers are protocol-specific, and can only be added by the protocol-specific connector configuration.

For example, if you configure the default address value to be example.com: 9999, Neo4j will fail to start and you will get an error in neo4j.log.

# 4.7. Dynamic settings

This section describes how to change your Neo4j configuration while Neo4j is running, and which settings can be changed.

This section contains the following:

- Introduction
- Discover dynamic settings
- · Update dynamic settings
- Dynamic settings reference

## 4.7.1. Introduction

Neo4j Enterprise Edition supports changing some configuration settings at runtime, without restarting the service.



Changes to the configuration at runtime are not persisted. To avoid losing changes when restarting Neo4j make sure to update *neo4j.conf* as well.

## 4.7.2. Discover dynamic settings

Use the procedure <a href="mailto:discover">dbms.listConfig()</a> to discover which configuration values can be dynamically updated, or consult <a href="mailto:Dynamic settings">Dynamic settings</a> reference.

### Example 9. Discover dynamic settings

## 4.7.3. Update dynamic settings

An administrator is able to change some configuration settings at runtime, without restarting the service.

#### Syntax:

CALL dbms.setConfigValue(setting, value)

#### **Returns:**

Nothing on success.

### **Exceptions:**

Unknown or invalid setting name.

The setting is not dynamic and can not be changed at runtime.

```
Invalid setting value.
```

The following example shows how to dynamically enable query logging.

### Example 10. Set a config value

```
CALL dbms.setConfigValue('dbms.logs.query.enabled', 'info')
```

If an invalid value is passed, the procedure will show a message to that effect.

### Example 11. Try to set invalid config value

```
CALL dbms.setConfigValue('dbms.logs.query.enabled', 'yes')

Failed to invoke procedure `dbms.setConfigValue`: Caused by: org.neo4j.graphdb.config.InvalidSettingException: Bad value 'yes' for setting 'dbms.logs.query.enabled': 'yes' not one of [OFF, INFO, VERBOSE]
```

To reset a config value to its default, pass an empty string as the value argument.

## Example 12. Reset a config value to default

```
CALL dbms.setConfigValue('dbms.logs.query.enabled', '')
```

## 4.7.4. Dynamic settings reference

## Dynamic settings reference

- cypher.query\_max\_allocations: The maximum amount of heap memory allocations to for cypher to perform on a single query, in bytes (or kilobytes with the 'k' suffix, megabytes with 'm' and gigabytes with 'g').
- dbms.checkpoint.iops.limit: Limit the number of IOs the background checkpoint process will consume per second.
- dbms.logs.debug.level: Debug log level threshold.
- dbms.logs.query.allocation\_logging\_enabled: Log allocated bytes for the executed queries being logged.
- dbms.logs.query.enabled: Log executed queries.
- dbms.logs.query.page\_logging\_enabled: Log page hits and page faults for the executed queries being logged.
- dbms.logs.query.parameter\_logging\_enabled: Log parameters for the executed queries being logged.
- dbms.logs.query.rotation.keep\_number: Maximum number of history files for the query log.
- dbms.logs.query.rotation.size: The file size in bytes at which the query log will auto-rotate.
- dbms.logs.query.runtime\_logging\_enabled: Logs which runtime that was used to run the query.
- dbms.logs.query.threshold: If the execution of query takes more time than this threshold, the query is logged once completed provided query logging is set to INFO.

- dbms.logs.query.time\_logging\_enabled: Log detailed time information for the executed queries being logged.
- dbms.track\_query\_allocation: Enables or disables tracking of how many bytes are allocated by the execution of a query.
- dbms.track\_query\_cpu\_time: Enables or disables tracking of how much time a query spends actively executing on the CPU.
- dbms.transaction.concurrent.maximum: The maximum number of concurrently running transactions.
- dbms.transaction.sampling.percentage: Transaction sampling percentage.
- dbms.transaction.timeout: The maximum time interval of a transaction within which it should be completed.
- dbms.transaction.tracing.level: Transaction creation tracing level.
- dbms.tx\_log.preallocate: Specify if Neo4j should try to preallocate logical log file in advance.
- dbms.tx\_log.rotation.retention\_policy: Make Neo4j keep the logical transaction logs for being able to backup the database.
- dbms.tx\_log.rotation.size: Specifies at which file size the logical log will auto-rotate.

## 4.8. Transaction logs

This section explains the retention and rotation policies for the Neo4j transaction logs, and how to configure them.

The transaction logs record all write operations in the database. This includes additions or modifications to data, as well as the addition or modification of any indexes or constraints. The transaction logs are the "source of truth" in scenarios where the database needs to be recovered. They are used to provide for incremental backups, as well as for cluster operations. For any given configuration, at least the latest non-empty transaction log will be kept.

#### Log location

By default, transaction logs for a database are located at <neo4j-home>/data/transactions/<database-name>. Each database keeps its own directory with transaction logs. The root directory where those folders are located is configured by dbms.directories.transaction.logs.root. For maximum performance, it is recommended to configure transaction logs to be stored on a dedicated device.

### Log rotation

Log rotation is configured using the parameter dbms.tx\_log.rotation.size. By default, log switches happen when log sizes surpass 250 MB.

#### Log retention

There are several different means of controlling the amount of transaction logs that are kept, using the parameter  $dbms.tx_{log.rotation.retention_policy}$ . This parameter can be configured in two different ways:

dbms.tx\_log.rotation.retention\_policy=<true/false>

If this parameter is set to true, transaction logs will be kept indefinitely. This option is not recommended due to the effectively unbounded storage usage. Old transaction logs cannot be safely archived or removed by external jobs, since safe log pruning requires knowledge about the most recent successful checkpoint.

If this parameter is set to false, only the most recent non-empty log will be kept. This option is

not recommended in production Enterprise Edition environments, as incremental backups rely on the presence of the transaction logs since the last backup.

• dbms.tx\_log.rotation.retention\_policy=<amount> <type>

### Log pruning

Transaction log pruning refers to the safe and automatic removal of old, unnecessary transaction log files. The transaction log can be pruned when one or more files fall outside of the configured retention policy. Two things are necessary for a file to be removed:

- The file must have been rotated.
- At least one checkpoint must have happened in a more recent log file.

Observing that you have more transaction log files than you expected is likely due to checkpoints either not happening frequently enough, or taking too long. This is a temporary condition and the gap between expected and observed number of log files will be closed on the next successful checkpoint. The interval between checkpoints can be configured using <a href="mailto:dbms.checkpoint.interval.time">dbms.checkpoint.interval.time</a> and <a href="mailto:dbms.checkpoint.interval.time">dbms.checkpoint.interval.time</a> and also help to speed up the checkpoint process itself. The configuration parameter <a href="mailto:dbms.checkpoint.iops.limit">dbms.checkpoint.iops.limit</a> controls the number of IOs per second the checkpoint process is allowed to use. Setting the value of this parameter to -1 allows unlimited IOPS, which can speed up checkpointing. Note that disabling the IOPS limit can cause transaction processing to slow down a bit.

Table 13. Types that can be used to control log retention

٠,	•	
Туре	Description	Example
files	Number of most recent logical log files to keep	"10 files"
size	Max disk size to allow log files to occupy	"300M size" or "1G size"
txs	Number of transactions to keep	"250k txs" or "5M txs"
hours	Keep logs which contains any transaction committed within N hours from current time	"10 hours"
days	Keep logs which contains any transaction committed within N days from current time	"50 days"

This example shows some different ways to configure the log retention policy.

• Keep transaction logs indefinitely:

```
dbms.tx_log.rotation.retention_policy=true
```

• Keep only the most recent non-empty log:

```
dbms.tx_log.rotation.retention_policy=false
```

• Keep logical logs which contain any transaction committed within 30 days:

```
dbms.tx_log.rotation.retention_policy=30 days
```

• Keep logical logs which contain any of the most recent 500 000 transactions:

```
dbms.tx_log.rotation.retention_policy=500k txs
```

# Chapter 5. Manage databases

This chapter describes how to create and manage multiple active databases.

## This chapter describes the following:

- Introduction
- · Administration and configuration
- Queries
- Databases in a Causal Cluster

## 5.1. Introduction

This section gives an introduction to managing multiple active databases with Neo4j.

### This section describes the following:

- Concepts
- The system database
- The default database

## 5.1.1. Concepts

With Neo4j 4.0 you can create and use more than one active database at the same time.

#### DBMS

Neo4j is a Database Management System, or *DBMS*, capable of managing multiple databases. The DBMS can manage a standalone server, or a group of servers in a Causal Cluster.

### Instance

A Neo4j instance is a Java process that is running the Neo4j server code.

#### Transaction domain

A transaction domain is a collection of graphs that can be updated within the context of a single transaction.

#### **Execution context**

An execution context is a runtime environment for the execution of a request. In practical terms, a request may be a query, a transaction, or an internal function or procedure.

#### Database

A database is an administrative partition of a DBMS. In practical terms, it is a physical structure of files organized within a directory or folder, that has the same name of the database. In logical terms, a database is a container for one or more graphs.

A database defines a *transaction domain* and an *execution context*. This means that a transaction cannot span across multiple databases. Similarly, a procedure is called within a database, although its logic may access data that is stored in other databases.

A default installation of Neo4j 4.0 contains two databases:

- system the system database, containing metadata on the DBMS and security configuration.
- neo4j the default database, a single database for user data. This has a default name of neo4j. A different name can be configured before starting Neo4j for the first time.

The following image illustrates a default installation, including the system database and a single database named neo4j for user data:

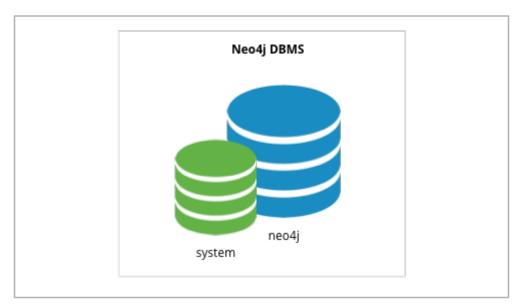


Figure 1. A default Neo4j installation.

#### **Editions**

The edition of Neo4j determines the number of possible databases:



- Installations of Community Edition can have exactly **one** user database.
- Installations of Enterprise Edition can have any number of user databases.

All installations include the system database.

## 5.1.2. The system database

All installations include a built-in database named system, which contains meta-data and security configuration.

The system database behaves differently than all other databases. In particular, when connected to this database you can only perform a specific set of administrative functions, as described in detail in Cypher Manual [] Administration.

Most of the available administrative commands are restricted to users with specific administrative privileges. An example of configuring security privileges is described in Fine-grained access control. Security administration is described in detail in Cypher Manual 

Security of administration.

The following image illustrates an installation of Neo4j with multiple active databases, named marketing, sales, and hr:

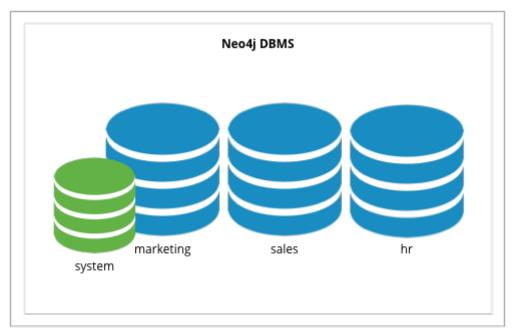


Figure 2. A multiple database Neo4j installation.

## 5.1.3. The default database

Each Neo4j instance has a default database. If a user connects to Neo4j without specifying a database, it will connect to the default database.

The default database is configurable. See configuration parameters for details.

The following image illustrates an installation of Neo4j containing the three databases for user data, named marketing, sales and hr, and the system database. The default database is sales:

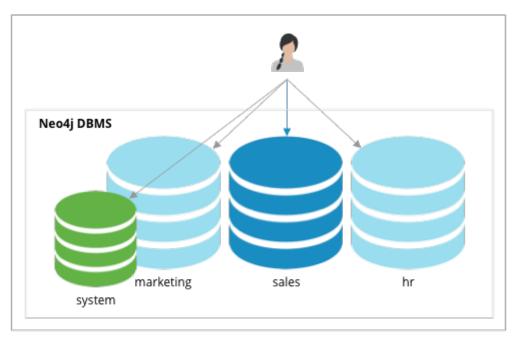


Figure 3. A multiple database Neo4j installation, with a default database.

# 5.2. Administration and configuration

This section describes how to manage multiple active databases.

This section describes the following:

- Administrative commands
- Configuration parameters

## 5.2.1. Administrative commands

The following Cypher commands are used on the system database to manage multiple databases:

Command	Description
CREATE DATABASE name	Create and start a new database.
DROP DATABASE name	Drop (remove) an existing database.
START DATABASE name	Start a database that has been stopped.
STOP DATABASE name	Shut down a database.
SHOW DATABASE name	Show the status of a specific database.
SHOW DATABASES	Show the name and status of all the databases.
SHOW DEFAULT DATABASE	Show the name and status of the default database.

Naming rules for databases are as follows:

- Length must be between 3 and 63 characters.
- The first character of a name must be an ASCII alphabetic character.
- Subsequent characters must be ASCII alphabetic or numeric characters, dots or dashes; [a..z][0..9].-
- Names are case-insensitive, and normalized to lowercase.
- Names that begin with an underscore, and with the prefix system are reserved for internal use.
- All of the above commands are executed as Cypher commands, and the database name is subject to the standard Cypher restrictions on valid identifiers. In particular, the (dash) character is not legal in Cypher variables, and therefore names with dashes need to be escaped by enclosing with back-ticks. For example, CREATE DATABASE `main-db`.

For detailed information on Cypher administrative commands, see Cypher Manual 

Administration.

For examples of using the Cypher administrative commands to manage multiple active databases, see Queries.

## 5.2.2. Configuration parameters

Configuration parameters are defined in the neo4j.conf file.

The following configuration parameters are applicable for managing databases:

Parameter name	Description	Default value
dbms.default_database	Name of the default database for the Neo4j instance. The database is created if it does not exist when the instance starts.	neo4j
dbms.max_databases	Maximum number of databases that can be used in a Neo4j single instance or Causal Cluster. The number includes all the online and offline databases. The value is an integer with a minimum value of 2.  Note that once the limit has been reached, it is not possible to create any additional databases. Similarly, if the limit is changed to a number lower than the total number of existing databases, no additional databases can be created.	100

# 5.3. Queries

This section provides examples of queries and Cypher commands that can be used to create and manage multiple active databases.

## This section describes the following:

- Show the status of a specific database
- Show the status of all databases
- Show the status of the default database
- Create a database
- Stop a database
- Start a database
- Drop, or remove a database

For detailed information on Cypher administrative commands, see Cypher Manual 

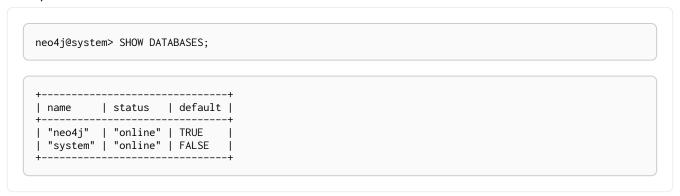
Administration.

# 5.3.1. Show the status of a specific database

### Example 14. SHOW DATABASE

## 5.3.2. Show the status of all databases

### Example 15. SHOW DATABASES



Switching between online and offline states is achieved using the START DATABASE and STOP DATABASE commands.

## 5.3.3. Show the status of the default database

The config setting dbms.default\_database defines which database is created and started by default when Neo4j starts. The default value of this setting is neo4j.

### Example 16. SHOW DEFAULT DATABASE



You can change the default database by using dbms.default\_database, and restarting the server.



In Community Edition, the default database is the only database available, other than the system database.

## 5.3.4. Create a database

## Example 17. CREATE DATABASE

# 5.3.5. Stop a database

## Example 18. STOP DATABASE

## 5.3.6. Start a database

### Example 19. START DATABASE

## 5.3.7. Drop, or remove a database

### Example 20. DROP DATABASE

# 5.4. Databases in a Causal Cluster

This section describes how to manage multiple active databases in a Causal Cluster.

Multiple databases in a Causal Cluster are managed the same way as a single instance. Administrators can use the same Cypher commands described in Administrative commands to manage databases.

This is based on two main principles:

- All databases are available on all members of a cluster this applies to Core servers and Read Replicas.
- Administrative commands must be executed on the system database, on the Leader member of the cluster.

# 5.4.1. Running Cypher administrative commands from Cypher Shell on a Causal Cluster.

For the following examples, consider a Causal Cluster environment formed by 5 members, 3 Core servers, and 2 Read Replicas:

Example 21. View the members of a Causal Cluster

Administrators can connect and execute Cypher commands in the following ways:

```
$ bin/cypher-shell -a bolt://localhost:7681 -d system -u neo4j -p neo4j1
Connected to Neo4j 4.0.0 at bolt://localhost:7681 as user neo4j.
Type :help for a list of available commands or :exit to exit the shell.
Note that Cypher queries must end with a semicolon.
neo4j@system> SHOW DATABASES;
| name | status | default |
| "neo4j" | "online" | TRUE
| "system" | "online" | FALSE
2 rows available after 34 ms, consumed after another 0 ms
neo4j@system> CREATE DATABASE data001;
0 rows available after 378 ms, consumed after another 12 ms
Added 1 nodes, Set 4 properties, Added 1 labels neo4j@system> SHOW DATABASES;
name
           | status | default |
| "neo4j" | "online" | TRUE
| "system" | "online" | FALSE
| "data001" | "online" | FALSE
3 rows available after 2 ms, consumed after another 1 ms
```

```
$ bin/cypher-shell -a neo4j://localhost:7683 -d system -u neo4j -p neo4j1
Connected to Neo4j 4.0.0 at neo4j://localhost:7683 as user neo4j.
Type :help for a list of available commands or :exit to exit the shell.
Note that Cypher queries must end with a semicolon.
neo4j@system> SHOW DATABASES;
| "neo4j" | "online" | TRUE
| "system" | "online" | FALSE
| "data001" | "online" | FALSE
3 rows available after 0 ms, consumed after another 0 ms
neo4j@system> CREATE DATABASE data002;
0 rows available after 8 ms, consumed after another 1 ms
Added 1 nodes, Set 4 properties, Added 1 labels
neo4j@system> SHOW DATABASES;
         | status | default |
name
| "neo4j" | "online" | TRUE
 "system"
           | "online" | FALSE
 "data001" | "online" | FALSE
| "data002" | "online" | FALSE
4 rows available after 33 ms, consumed after another 0 ms
```



The neo4j:// scheme is the equivalent to the bolt+routing: scheme available in earlier versions of Neo4j, but it can be used seamlessly with a standalone and clustered DBMS.

# Chapter 6. Clustering

This chapter describes the configuration and operation of a Neo4j Causal Cluster.

### This chapter describes the following:

- Introduction An overview of the Causal Clustering architecture.
- Deploy a cluster The basics of configuring and deploying a new cluster.
- Seed a cluster How to deploy a Causal Cluster with pre-existing data.
- Discovery How members of a cluster discover each other.
- Intra-cluster encryption How to secure the cluster communication.
- Internals A few internals regarding the operation of the cluster.
- Settings reference A summary of the most important Causal Cluster settings.

#### Further information:

- For instructions on setting up Causal Clustering when running Neo4j in a Docker container, see Causal Clustering on Docker.
- For an example of managing multiple databases in a Causal Cluster, see Multiple databases in a Causal Cluster.
- For instructions on how you can upgrade your Neo4j Causal Cluster, see Upgrading a Causal Cluster.
- For a summary of the facilities that are available for monitoring a Neo4j Causal Cluster, see Monitoring (and specifically, Monitoring a Causal Cluster).
- For a tutorial on setting up a test cluster locally on a single machine, see Set up a local Causal Cluster.
- For advanced concepts, including the implementation of the Raft Protocol, see Advanced Causal Clustering

## 6.1. Introduction

This section gives an introduction to the Neo4j Causal Clustering architecture.

### This section includes:

- Overview
- Operational view
  - Core Servers
  - ☐ Read Replicas
- Causal consistency
- Summary

## 6.1.1. Overview

Neo4j's Causal Clustering provides three main features:

1. Safety: Core Servers provide a fault tolerant platform for transaction processing which will remain

available while a simple majority of those Core Servers are functioning.

- 2. **Scale:** Read Replicas provide a massively scalable platform for graph queries that enables very large graph workloads to be executed in a widely distributed topology.
- 3. **Causal consistency:** when invoked, a client application is guaranteed to read at least its own writes.

Together, this allows the end-user system to be fully functional and both read and write to the database in the event of multiple hardware and network failures and makes reasoning about database interactions straightforward.

In the remainder of this section we will provide an overview of how causal clustering works in production, including both operational and application aspects.

## 6.1.2. Operational view

From an operational point of view, it is useful to view the cluster as being composed of servers with two different roles: Cores and Read Replicas.

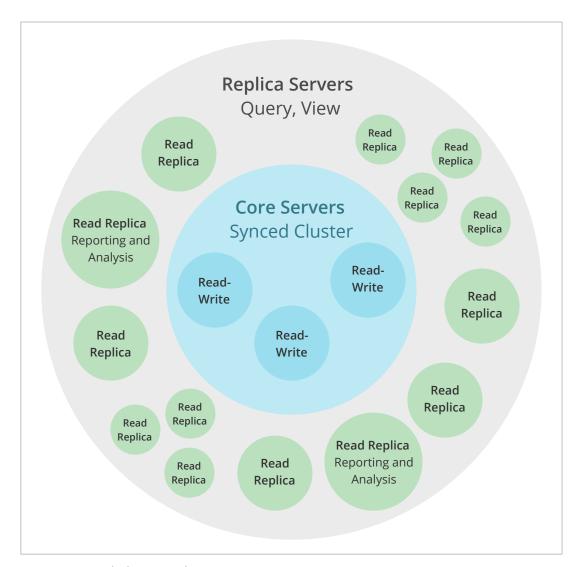


Figure 4. Causal Cluster Architecture

The two roles are foundational in any production deployment but are managed at different scales from one another and undertake different roles in managing the fault tolerance and scalability of the overall cluster.

## Core Servers

The main responsibility of Core Servers is to safeguard data. Core Servers do so by replicating all transactions using the Raft protocol. Raft ensures that the data is safely durable before confirming transaction commit to the end user application. In practice this means once a majority of Core Servers in a cluster (N/2+1) have accepted the transaction, it is safe to acknowledge the commit to the end user application.

The safety requirement has an impact on write latency. Implicitly writes will be acknowledged by the fastest majority, but as the number of Core Servers in the cluster grows so do the size of the majority needed to acknowledge a write.

In practice this means that there are relatively few machines in a typical Core Server cluster, enough to provide sufficient fault tolerance for the specific deployment. This is calculated with the formula M = 2F + 1 where M is the number of Core Servers required to tolerate F faults. For example:

- In order to tolerate two failed Core Servers we would need to deploy a cluster of five Cores.
- The smallest fault tolerant cluster, a cluster that can tolerate one fault, must have three Cores.
- It is also possible to create a Causal Cluster consisting of only two Cores. However, that cluster will not be fault-tolerant. If one of the two servers fails, the remaining server will become read-only.



Should the cluster suffer enough Core failures then it can no longer process writes and it will become read-only to preserve safety.

## Read Replicas

The main responsibility of Read Replicas is to scale out graph workloads. Read Replicas act like caches for the graph data that the Core Servers safeguard and are fully capable of executing arbitrary (read-only) queries and procedures.

Read Replicas are asynchronously replicated from Core Servers via transaction log shipping. They will periodically poll an upstream server for new transactions and have these shipped over. Many Read Replicas can be fed data from a relatively small number of Core Servers, allowing for a large fan out of the guery workload for scale.

Read Replicas should typically be run in relatively large numbers and treated as disposable. Losing a Read Replica does not impact the cluster's availability, aside from the loss of its fraction of graph query throughput. It does not affect the fault tolerance capabilities of the cluster.

## 6.1.3. Causal consistency

While the operational mechanics of the cluster are interesting from an application point of view, it is also helpful to think about how applications will use the database to get their work done. In an application we typically want to read from the graph and write to the graph. Depending on the nature of the workload we usually want reads from the graph to take into account previous writes to ensure causal consistency.



Causal consistency is one of numerous consistency models used in distributed computing. It ensures that causally related operations are seen by every instance in the system in the same order. Consequently, client applications are guaranteed to read their own writes, regardless of which instance they communicate with. This simplifies interaction with large clusters, allowing clients to treat them as a single (logical) server.

Causal consistency makes it possible to write to Core Servers (where data is safe) and read those writes from a Read Replica (where graph operations are scaled out). For example, causal consistency

guarantees that the write which created a user account will be present when that same user subsequently attempts to log in.

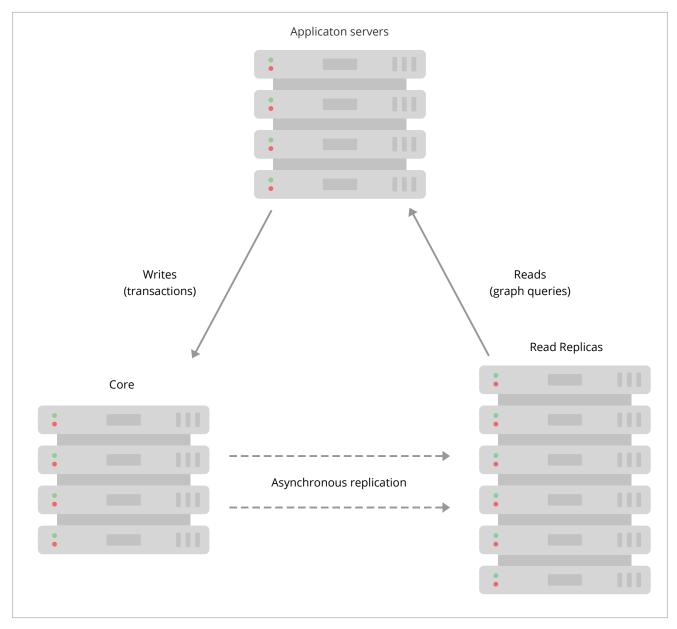


Figure 5. Causal Cluster setup with causal consistency via Neo4j drivers

On executing a transaction, the client can ask for a bookmark which it then presents as a parameter to subsequent transactions. Using that bookmark the cluster can ensure that only servers which have processed the client's bookmarked transaction will run its next transaction. This provides a *causal chain* which ensures correct read-after-write semantics from the client's point of view.

Aside from the bookmark everything else is handled by the cluster. The database drivers work with the cluster topology manager to choose the most appropriate Core Servers and Read Replicas to provide high quality of service.

## 6.1.4. Summary

In this section we have examined Causal Clustering at a high level from an operational and an application development point of view. We now understand that the Core Servers are responsible for the long-term safekeeping of data while the more numerous Read Replicas are responsible for scaling out graph query workloads. Reasoning about this powerful architecture is greatly simplified by the Neo4j drivers which abstract the cluster topology to easily provide read levels like causal consistency.

# 6.2. Deploy a cluster

This section describes how to deploy a new Neo4j Causal Cluster.

### This section includes:

- Introduction
- · Configure a Core-only cluster
- · Add a Core Server to an existing cluster
- · Add a Read Replica to an existing cluster

## 6.2.1. Introduction

In this section we describe how to set up a new Causal Cluster consisting of three Core instances. We then proceed to show how more Core Servers as well as Read Replicas can be added to a running cluster.

Three Cores is the minimum number of servers needed in order to form a fault-tolerant Causal Cluster. See Core Servers for a discussion on the number of servers required in various scenarios.

Refer to Set up a local Causal Cluster for a tutorial on how to set up a Causal Cluster on a local machine.

# 6.2.2. Configure a Core-only cluster

The following configuration settings are important to consider when deploying a new Causal Cluster. See also Settings reference for more detailed descriptions and examples.

Table 14. Important settings for a new Causal Cluster

Option name	Description
dbms.default_listen_address	The address or network interface this machine uses to listen for incoming messages. Setting this value to 0.0.0.0 makes Neo4j bind to all available network interfaces.
dbms.default_advertised_address	The address that other machines are told to connect to. In the typical case, this should be set to the fully qualified domain name or the IP address of this server.
dbms.mode	The operating mode of a single server instance. For Causal Clustering, there are two possible modes: CORE or READ_REPLICA.
<pre>causal_clustering.minimum_core_cluster_size_at_formati on</pre>	The minimum number of Core machines in the cluster at formation. A cluster will not form without the number of Cores defined by this setting, and this should in general be configured to the full and fixed amount.
causal_clustering.minimum_core_cluster_size_at_runtime	The minimum number of Core instances which will exist in the consensus group.
causal_clustering.initial_discovery_members	The network addresses of an initial set of Core cluster members that are available to bootstrap this Core or Read Replica instance. In the default case, the initial discovery members are given as a comma-separated list of address/port pairs, and the default port for the discovery service is :5000. It is good practice to set this parameter to the same value on all Core Servers.  The behavior of this setting can be modified by configuring the setting causal_clustering.discovery_type. This is described in detail in Discovery.



### Listen configuration

Listening on 0.0.0.0 makes the ports publicly available. Make sure you understand the security implications and strongly consider setting up encryption.

The following example shows how to set up a simple cluster with three Core servers:

### Example 24. Configure a Core-only cluster

In this example, we will configure three Core Servers named <code>core01.example.com</code>, <code>core02.example.com</code> and <code>core03.example.com</code>. We have already installed Neo4j Enterprise Edition on all three servers. We configure them by preparing <code>neo4j.conf</code> on each server. Note that they are all identical, except for the configuration of <code>dbms.default\_advertised\_address</code>:

*neo4j.conf on core01.example.com:* 

```
dbms.default_listen_address=0.0.0.0
dbms.default_advertised_address=core01.example.com
dbms.mode=CORE
causal_clustering.initial_discovery_members=core01.example.com:5000,core02.example.com:5000
ample.com:5000
```

### neo4j.conf on core02.example.com:

```
dbms.default_listen_address=0.0.0.0
dbms.default_advertised_address=core02.example.com
dbms.mode=CORE
causal_clustering.initial_discovery_members=core01.example.com:5000,core02.example.com:5000
ample.com:5000
```

### neo4j.conf on core03.example.com:

```
dbms.default_listen_address=0.0.0.0
dbms.default_advertised_address=core03.example.com
dbms.mode=CORE
causal_clustering.initial_discovery_members=core01.example.com:5000,core02.example.com:5000,core03.ex
ample.com:5000
```

Now we are ready to start the Neo4j servers. The startup order does not matter.

After the cluster has started, we can connect to any of the instances and run CALL dbms.cluster.overview() to check the status of the cluster. This will show information about each member of the cluster.

We now have a Neo4j Causal Cluster of three instances running.



#### Startup time

The instance may appear unavailable while it is joining the cluster. If you want to follow along with the startup, you can follow the messages in *neo4j.log*.

## 6.2.3. Add a Core Server to an existing cluster

Core Servers are added to an existing cluster by starting a new Neo4j instance with the appropriate configuration. The new server will join the existing cluster and become available once it has copied the data from its peers. It may take some time for the new instance to perform the copy if the existing cluster contains large amounts of data.

The setting causal\_clustering.initial\_discovery\_members shall be updated on all the servers in the cluster to include the new server.

In this example, we will add a Core Server, core04.example.com, to the cluster that we created in Configure a Core-only cluster.

We configure the following entries in *neo4j.conf*:

neo4j.conf on core04.example.com:

```
dbms.default_listen_address=0.0.0.0
dbms.default_advertised_address=core04.example.com
dbms.mode=CORE
causal_clustering.minimum_core_cluster_size_at_formation=3
causal_clustering.minimum_core_cluster_size_at_runtime=3
causal_clustering.discovery_members=core01.example.com:5000,core02.example.com:5000,core03.example.co
m:5000,core04.example.com:5000
```

Note that the configuration is very similar to that of the previous servers. In this example, the new server is not intended to be a permanent member of the cluster, thus it is not included in causal\_clustering.discovery\_members.

Now we can start the new Core Server and let it add itself to the existing cluster.

## 6.2.4. Add a Read Replica to an existing cluster

Initial Read Replica configuration is provided similarly to Core Servers via *neo4j.conf*. Since Read Replicas do not participate in cluster quorum decisions, their configuration is shorter; they only need to know the addresses of some of the Core Servers which they can bind to in order to discover the cluster. They can then choose an appropriate Core Server from which to copy data.

Example 26. Add a Read Replica to an existing cluster

In this example, we will add a Read Replica, replica01.example.com, to the cluster that we created in Configure a Core-only cluster.

We configure the following entries in *neo4j.conf*:

neo4j.conf on replica01.example.com:

```
dbms.mode=READ_REPLICA
causal_clustering.discovery_members=core01.example.com:5000,core02.example.com:5000,core03.example.co
m:5000
```

Now we can start the new Read Replica and let it add itself to the existing cluster.

## 6.3. Seed a cluster

This section describes how to seed a new Neo4j Causal Cluster with existing data.

#### This section includes:

- Introduction
- Seed from backups
- · Seed using the import tool

## 6.3.1. Introduction

In Deploy a cluster we learned how to create a cluster with empty databases. However, regardless of whether we are just playing around with Neo4j or setting up a production environment, it is likely that we have some existing data that we wish to transfer into our cluster.

This section outlines how to create a Causal Cluster containing data either seeded from an existing online or offline Neo4j database, or imported from some other data source using the import tool. The general steps to seed a cluster will follow the same pattern, regardless of which format our data is in:

- 1. Create a new Neo4j Core-only cluster.
- 2. Seed the cluster.
- 3. Start the cluster.



The databases which you are using to seed the cluster must be of the same version of Neo4j as the cluster itself.

## 6.3.2. Seed from backups

For this section, it is assumed that we already have healthy backups of an existing Neo4j deployment. This could be online or offline backups from a standalone Neo4j instance or a Neo4j Causal Cluster. For details on performing online backups, please refer to Backup.



Moving files and directories manually in or out of a Neo4j installation is not recommended and considered unsupported usage. If you have an existing Neo4j database which you wish to use for a new cluster, then use neo4j-admin dump to create an offline backup.

The process described here can also be used to seed a new Causal Cluster from an existing Read Replica. This can be useful, for example, in disaster recovery where some servers have retained their data during a catastrophic event.

1. Create a new Neo4j Core-only cluster.

Follow the instructions in Configure a Core-only cluster to create a new Neo4j Core-only cluster.



You could start the cluster now in order to test that everything is correctly configured, but this will create default databases as part of cluster formation. Since you are trying to seed with a set of databases you have to subsequently stop every instance, unbind them from the cluster using neo4j-admin unbind and remove those databases so that the correct seeds can be used instead.

### 2. Seed the cluster.

Use neo4j-admin restore or neo4j-admin load to seed all the Core instances in the cluster.

The examples assume that we are restoring one user database with the default name of neo4j in addition to the system database which contains replicated configuration state. Modify the command line arguments to match your exact setup.

### Example 27. Seed using neo4j-admin restore.

```
neo4j-01$ ./bin/neo4j-admin restore --from=/path/to/system-backup-dir --database=system neo4j-01$ ./bin/neo4j-admin restore --from=/path/to/neo4j-backup-dir --database=neo4j

neo4j-02$ ./bin/neo4j-admin restore --from=/path/to/system-backup-dir --database=system neo4j-02$ ./bin/neo4j-admin restore --from=/path/to/neo4j-backup-dir --database=neo4j

neo4j-03$ ./bin/neo4j-admin restore --from=/path/to/system-backup-dir --database=system neo4j-03$ ./bin/neo4j-admin restore --from=/path/to/neo4j-backup-dir --database=neo4j
```

### Example 28. Seed using neo4j-admin load.

```
neo4j-01$ ./bin/neo4j-admin load --from=/path/to/system.dump --database=system neo4j-01$ ./bin/neo4j-admin load --from=/path/to/neo4j.dump --database=neo4j

neo4j-02$ ./bin/neo4j-admin load --from=/path/to/system.dump --database=system neo4j-02$ ./bin/neo4j-admin load --from=/path/to/neo4j.dump --database=neo4j

neo4j-03$ ./bin/neo4j-admin load --from=/path/to/system.dump --database=system neo4j-03$ ./bin/neo4j-admin load --from=/path/to/neo4j.dump --database=system neo4j-03$ ./bin/neo4j-admin load --from=/path/to/neo4j.dump --database=neo4j
```

### 3. Start the cluster.

At this point, all of the instances in the Core cluster have been seeded. Between them, the Core Servers have everything necessary to form a cluster. We are ready to start all instances. The cluster will form and the replicated Neo4j DBMS deployment will come online.

### Example 29. Start each of the Core instances.

```
neo4j-01$ ./bin/neo4j start

neo4j-02$ ./bin/neo4j start

neo4j-03$ ./bin/neo4j start
```

## 6.3.3. Seed using the import tool

In order to create a cluster based on imported data, it is recommended to first import the data into a standalone Neo4j DBMS and then use an offline backup to seed the cluster.

- 1. Import the data.
  - a. Deploy a standalone Neo4j DBMS.
  - b. Import the data using the import tool.
- 2. Use neo4j-admin dump to create an offline backup of the neo4j database.

3. Seed a new cluster using the instructions in Seed from backups.

Skip the system database in this scenario since it is not needed.

## 6.4. Discovery

This section explains how members of a cluster discover each other.

#### This section includes:

- Overview
- Discovery using a list of server addresses
- · Discovery using DNS with multiple records
- · Discovery in Kubernetes

### 6.4.1. Overview

In order to form or connect to a running cluster, a Core Server or a Read Replica needs to know the addresses of some of the Core Servers. This information is used to bind to the Core Servers in order to run the discovery protocol and get the full information about the cluster. The way in which this is best done depends on the configuration in each specific case.

If the addresses of the other cluster members are known upfront, they can be listed explicitly. This is convenient, but has limitations:

- If Core members are replaced and the new members have different addresses, the list will become outdated. An outdated list can be avoided by ensuring that the new members can be reached via the same address as the old members, but this is not always practical.
- Under some circumstances the addresses are unknown when configuring the cluster. This can be the case, for example, when using container orchestration to deploy a Causal Cluster.

Additional mechanisms for using DNS are provided for the cases where it is not practical or possible to explicitly list the addresses of cluster members to discover.

The discovery configuration is just used for initial discovery and a running cluster will continuously exchange information about changes to the topology. The behavior of the initial discovery is determined by the parameters <code>causal\_clustering.discovery\_type</code> and <code>causal\_clustering.initial\_discovery\_members</code>, and is described in the following sections.

## Discovery using a list of server addresses

If the addresses of the other cluster members are known upfront, they can be listed explicitly. We use the default causal\_clustering.discovery\_type=LIST and hard code the addresses in the configuration of each machine. This alternative is illustrated by Configure a Core-only cluster.

## Discovery using DNS with multiple records

When using initial discovery with DNS, a DNS record lookup is performed when an instance starts up. Once an instance has joined a cluster, further membership changes are communicated amongst Core members as part of the discovery service.

The following DNS-based mechanisms can be used to get the addresses of Core Cluster members for discovery:

#### causal\_clustering.discovery\_type=DNS

With this configuration, the initial discovery members will be resolved from *DNS A* records to find the IP addresses to contact. The value of causal\_clustering.initial\_discovery\_members should be set to a single domain name and the port of the discovery service. For example: causal\_clustering.initial\_discovery\_members=cluster01.example.com:5000. The domain name should return an A record for every Core member when a DNS lookup is performed. Each A record returned by DNS should contain the IP address of the Core Server. The configured Core Server will use all the IP addresses from the A records to join or form a cluster.

The discovery port must be the same on all Cores when using this configuration. If this is not possible, consider using the discovery type SRV instead.

#### causal\_clustering.discovery\_type=SRV

With this configuration, the initial discovery members will be resolved from *DNS SRV* records to find the IP addresses/hostnames and discovery service ports to contact. The value of causal\_clustering.initial\_discovery\_members should be set to a single domain name and the port set to 0. For example: causal\_clustering.initial\_discovery\_members=cluster01.example.com:0. The domain name should return a single SRV record when a DNS lookup is performed. The SRV record returned by DNS should contain the IP address or hostname, and the discovery port, for the Core Servers to be discovered. The configured Core Server will use all the addresses from the SRV record to join or form a cluster.

### Discovery in Kubernetes

A special case is when a Causal Cluster is running in Kubernetes and each Core Server is running as a Kubernetes service. Then the addresses of Core Cluster members can be obtained using the List Service API.

The following settings are used to configure for this scenario:

- Set causal\_clustering.discovery\_type=K8S.
- Set causal\_clustering.kubernetes.label\_selector to a label selector for the Causal Cluster services.
- Set causal\_clustering.kubernetes.service\_port\_name to the name of the service port used in the Kubernetes service definition for the Core's discovery port.

With this configuration, causal\_clustering.initial\_discovery\_members is not used and any value assigned to it will be ignored.



- The pod running Neo4j must use a service account which has permission to list services. For further information, see the Kubernetes documentation on RBAC authorization or ABAC authorization.
- The configured causal\_clustering.discovery\_advertised\_address must exactly match the Kubernetes-internal DNS name, which will be of the form <service-name>.<namespace>.svc.cluster.local.

As with DNS-based methods, the Kubernetes record lookup is only performed at start up.

## 6.5. Intra-cluster encryption

This chapter describes how to secure the cluster communication between server instances.



Securing client to server communication is not covered in this chapter (e.g. Bolt, HTTPS, Backup).

#### This section includes:

- Introduction
- Example deployment
  - ☐ Generate and install cryptographic objects
  - ☐ Configure the cluster SSL policy
  - ☐ Validate the secure operation of the cluster

### 6.5.1. Introduction

The security solution for cluster communication is based on standard SSL/TLS technology (referred to jointly as SSL). Encryption is in fact just one aspect of security, with the other cornerstones being authentication and integrity. A secure solution will be based on a key infrastructure which is deployed together with a requirement of authentication.

The SSL support in the platform is documented in detail in SSL framework. This section will cover the specifics as they relate to securing a cluster.

Under SSL, an endpoint can authenticate itself using certificates managed by a Public Key Infrastructure (*PKI*).

It should be noted that the deployment of a secure key management infrastructure is beyond the scope of this manual, and should be entrusted to experienced security professionals. The example deployment illustrated below is for reference purposes only.

## 6.5.2. Example deployment

The following steps will create an example deployment, and each step is expanded in further detail below.

- Generate and install cryptographic objects
- Configure Causal Clustering with the SSL policy
- Validate the secure operation of the cluster

## Generate and install cryptographic objects

The generation of cryptographic objects is for the most part outside the scope of this manual. It will generally require having a PKI with a Certificate Authority (CA) within the organization and they should be able to advise here. Please note that the information in this manual relating to the PKI is mainly for illustrative purposes.

When the certificates and private keys have been obtained they can be installed on each of the servers. Each server will have a certificate of its own, signed by a CA, and the corresponding private key. The certificate of the CA is installed into the trusted directory, and any certificate signed by the CA will thus be trusted. This means that the server now has the capability of establishing trust with other servers.



Please be sure to exercise caution when using CA certificates in the trusted directory, as any certificates signed by that CA will then be trusted to join the cluster. For this reason, never use a public CA to sign certificates for your cluster. Instead, use an intermediate certificate or a CA certificate which originates from and is controlled by your organization.

In this example we will deploy a mutual authentication setup, which means that both ends of a channel have to authenticate. To enable mutual authentication the SSL policy must have client\_auth

set to REQUIRE (which is the default). Servers are by default required to authenticate themselves, so there is no corresponding server setting.

If the certificate for a particular server is compromised it is possible to revoke it by installing a Certificate Revocation List (CRL) in the revoked directory. It is also possible to redeploy using a new CA. For contingency purposes, it is advised that you have a separate intermediate CA specifically for the cluster which can be substituted in its entirety should it ever become necessary. This approach would be much easier than having to handle revocations and ensuring their propagation.

#### Example 30. Generate and install cryptographic objects

In this example we assume that the private key and certificate file are named *private.key* and *public.crt*, respectively. If you want to use different names you may override the policy configuration for the key and certificate names/locations. We want to use the default configuration for this server so we create the appropriate directory structure and install the certificate:

```
$neo4j-home> mkdir certificates/cluster
$neo4j-home> mkdir certificates/cluster/trusted
$neo4j-home> mkdir certificates/cluster/revoked

$neo4j-home> cp $some-dir/private.key certificates/cluster
$neo4j-home> cp $some-dir/public.crt certificates/cluster
```

### Configure the cluster SSL policy

By default, cluster communication is unencrypted. To configure a Causal Cluster to encrypt its intracluster communication, set <a href="mailto:dbms.ssl.policy.cluster.enabled">dbms.ssl.policy.cluster.enabled</a> to true.

An SSL policy utilizes the installed cryptographic objects and additionally allows parameters to be configured. We will use the following parameters in our configuration:

Table 15. Example settings

Setting suffix	Value	Comment
client_auth	REQUIRE	Setting this to REQUIRE effectively enables mutual authentication for servers.
ciphers	TLS_ECDHE_RSA_WITH_AES_256_CBC_SHA38 4	We can enforce a particular single strong cipher and remove any doubt about which cipher gets negotiated and chosen. The cipher chosen above offers Perfect Forward Secrecy (PFS) which is generally desirable. It also uses Advanced Encryption Standard (AES) for symmetric encryption which has great support for acceleration in hardware and thus allows performance to generally be negligibly affected.
tls_versions	TLSv1.2	Since we control the entire cluster we can enforce the latest TLS standard without any concern for backwards compatibility. It has no known security vulnerabilities and uses the most modern algorithms for key exchanges, etc.

In the following example we will create and configure an SSL policy that we will use in our cluster.

In this example we assume that the directory structure has been created, and certificate files have been installed, as per the previous example.

We add the following content to our *neo4j.conf* file:

```
dbms.ssl.policy.cluster.enabled=true
dbms.ssl.policy.cluster.tls_versions=TLSv1.2
dbms.ssl.policy.cluster.ciphers=TLS_ECDHE_RSA_WITH_AES_256_CBC_SHA384
dbms.ssl.policy.cluster.client_auth=REQUIRE
```

Any user data communicated between instances will now be secured. Please note that an instance which is not correctly setup would not be able to communicate with the others.

Note that the policy must be configured on every server with the same settings. The actual cryptographic objects installed will be mostly different since they do not share the same private keys and corresponding certificates. The trusted CA certificate will be shared however.

### Validate the secure operation of the cluster

To make sure that everything is secured as intended it makes sense to validate using external tooling such as, for example, the open source assessment tools nmap or OpenSSL.

#### Example 32. Validate the secure operation of the cluster

In this example we will use the nmap tool to validate the secure operation of our cluster. A simple test to perform is a cipher enumeration using the following command:

```
nmap --script ssl-enum-ciphers -p <port> <hostname>
```

The hostname and port have to be adjusted according to our configuration. This can prove that TLS is in fact enabled and that the only the intended cipher suites are enabled. All servers and all applicable ports should be tested.

For testing purposes we could also attempt to utilize a separate testing instance of Neo4j which, for example, has an untrusted certificate in place. The expected result of this test is that the test server is not able to participate in replication of user data. The debug logs will generally indicate an issue by printing an SSL or certificate-related exception.

## 6.6. Internals of clustering

This section details a few selected internals of a Neo4j Causal Cluster. Understanding the internals is not vital but can be helpful in diagnosing and resolving operational issues.

#### This section includes:

- Elections and leadership
- Multi-database and the reconciler
- Store copy
- · On-disk state

## 6.6.1. Elections and leadership

The Core Servers in a Causal Cluster use the Raft protocol to ensure consistency and safety. An implementation detail of Raft is that it uses a *Leader* role to impose an ordering on an underlying log with other instances acting as *Followers* which replicate the leader's state. In Neo4j terms this means writes to the database are ordered by the Core instance currently playing the leader role.

If a follower has not heard from the leader for a while, then it can initiate an election and attempt to become the new leader. The follower makes itself a *Candidate* and asks other Cores to vote for it. If it can get a majority of the votes, then it assumes the leader role. Cores will not vote for a candidate which is less up-to-date than itself. There can only be one leader at any time, and that leader is guaranteed to have the most up-to-date log.

It is expected for elections to occur during the normal running of a cluster and they do not pose an issue in and of itself. If you are experiencing frequent re-elections and they are disturbing the operation of the cluster then you should try to figure out what is causing them. Some common causes are environmental issues (e.g. a flaky networking) and work overload conditions (e.g. more concurrent queries and transactions than the hardware can handle).

### 6.6.2. Multi-database and the reconciler

Databases operate as independent entities in a Neo4j DBMS, both in standalone and in a cluster. Since a cluster consists of multiple independent server instances, the effects of administrative operations like creating a new database happen asynchronously and independently for each server. However, the immediate effect of an administrative operation is to safely commit the desired state in the system database.

The desired state committed in the system database gets replicated and is picked up by an internal component called the reconciler. It runs on every instance and takes the appropriate actions required locally on that instance for reaching the desired state; creating, starting, stopping and dropping databases.

Every database runs in an independent Raft group and since there are two databases in a fresh cluster, system and neo4j, this means that it also has two Raft groups. Every Raft group also has an independent leader and thus a particular Core server could be the leader for one database and a follower for another.

## 6.6.3. Store copy

Store copies are initiated when an instance does not have an up-to-date copy of the database. For example, this will be the case when a new instance is joining a cluster (without a seed). It can also happen as a consequence of falling behind the rest of the cluster, for reasons such as connectivity issues or having been shutdown. Upon re-establishing connection with the cluster, an instance will recognize that it is too far behind and fetch a new copy from the rest of the cluster.

A store copy is a major operation which may disrupt the availability of instances in the cluster. Store copies should not be a frequent occurrence in a well-functioning cluster, but rather be an exceptional operation that happens due to specific causes, e.g. network outages or planned maintenance outages. If store copies happen during regular operation, then the configuration of the cluster, or the workload directed at it, might have to be reviewed so that all instances can keep up, and that there is enough of a buffer of Raft logs and transaction logs to handle smaller transient issues.

The protocol used for store copies is robust and configurable. The network requests will be directed at an upstream member according to configuration and they will be retried despite transient failures. The maximum amount of time to retry every request can be modified by setting causal\_clustering.store\_copy\_max\_retry\_time\_per\_request. If a request fails and the maximum retry time has elapsed then it will stop retrying and the store copy will fail.

Use causal\_clustering.catch\_up\_client\_inactivity\_timeout to configure the inactivity timeout for any particular request.



This setting is for all requests from the catchup client, including the pulling of transactions.

The default upstream strategy differs for Cores and Read Replicas. Cores will always send the initial request to the leader to get the most up-to-date information about the store. The strategy for the file and index requests for Cores is to vary every other request to a random Read Replica and every other to a random Core member.

Read Replicas use the same strategy for store copies as it uses for pulling transactions. The default is to pull from a random Core member.

If you are running a multi-data center cluster, then upstream strategies for both Cores and Read Replicas can be configured. Remember that for Read Replicas this also affects from where transactions are pulled. See more in Configure for multi-data center operations.

## 6.6.4. On-disk state

The on-disk state of cluster instances is different to that of standalone instances. The biggest difference being the existence of additional cluster state. Most of the files there are relatively small, but the Raft logs can become quite large depending on the configuration and workload.

It is important to understand that once a database has been extracted from a cluster and used in a standalone deployment, it must not be put back into an operational cluster. This is because the cluster and the standalone deployment now have separate databases, with different and irreconcilable writes applied to them.



If you try to reinsert the modified database back into the cluster, then the logs and stores will mismatch. Operators should not try to merge standalone databases into the cluster in the optimistic hope that their data will become replicated. That will not happen and will likely lead to unpredictable cluster behavior.

## 6.7. Settings reference

This section lists the important settings related to running a Neo4j Causal Cluster.

Parameter	Explanation
dbms.mode	This setting configures the operating mode of the database. For Causal Clustering, there are two possible modes: CORE or READ_REPLICA.
	<b>Example:</b> dbms.mode=READ_REPLICA will define this server as a Read Replica.
<pre>causal_clustering.minimum_ core_cluster_size_at_forma</pre>	Minimum number of Core machines required to form a cluster.
tion	<b>Example:</b> causal_clustering.minimum_core_cluster_size_at_formation=3 will specify that the cluster will form when at least three Core members have discovered each other.

### **Parameter Explanation** causal clustering.minimum The minimum size of the dynamically adjusted voting set (which only Core core\_cluster\_size\_at\_runti members may be a part of). Adjustments to the voting set happen automatically as the availability of Core members changes, due to explicit operations such as starting or stopping a member, or unintended issues such as network partitions. Please note that this dynamic scaling of the voting set is generally desirable, as under some circumstances it can increase the number of instance failures which may be tolerated. A majority of the voting set must be available before members are voted in or out. **Example:** causal\_clustering.minimum\_core\_cluster\_size\_at\_runtime=3 will specify that the cluster should not try to dynamically adjust below three Core members in the voting set. causal\_clustering.discover This setting specifies the strategy that the instance will use to determine v\_tvpe the addresses for other instances in the cluster to contact for bootstrapping. Possible values are: LIST, DNS, SRV, and K8S. LIST Treat causal\_clustering.initial\_discovery\_members as a list of addresses of Core Servers to contact for discovery. DNS Treat causal\_clustering.initial\_discovery\_members as a domain name to resolve via DNS. Expect DNS resolution to provide A records with hostnames or IP addresses of Cores to contact for discovery, on the port specified by causal\_clustering.initial\_discovery\_members. SRV Treat causal\_clustering.initial\_discovery\_members as a domain name to resolve via DNS. Expect DNS resolution to provide SRV records with hostnames or IP addresses, and ports, of Cores to contact for discovery. K8S Access the Kubernetes list service API to derive addresses of Cores to contact for discovery. Requires causal\_clustering.kubernetes.label\_selector to be a Kubernetes label selector for Kubernetes services running a Core each and causal\_clustering.kubernetes.service\_port\_name to be a service port name identifying the discovery port of Core services. The value of causal\_clustering.initial\_discovery\_members is ignored for this option. The value of this setting determines how causal\_clustering.initial\_discovery\_members is interpreted. Detailed information about discovery and discovery configuration options is given in Discovery using DNS with multiple records. **Example:** causal\_clustering.discovery\_type=DNS combined with causal\_clustering.initial\_discovery\_members=cluster01.example.com:5000 will fetch all DNS A records for *cluster01.example.com* and attempt to reach Neo4j instances listening on port 5000 for each A record's IP address.

Parameter	Explanation
<pre>causal_clustering.initial_ discovery_members</pre>	The network addresses of an initial set of Core cluster members that are available to bootstrap this Core or Read Replica instance. In the default case, the initial discovery members are given as a comma-separated list of address/port pairs, and the default port for the discovery service is :5000. It is good practice to set this parameter to the same value on all Core Servers.
	It is good practice to set this parameter to the same value on all Core Servers.
	The behavior of this setting can be modified by configuring the setting causal_clustering.discovery_type. This is described in detail in Discovery using DNS with multiple records.
	<b>Example:</b> causal_clustering.discovery_type=LIST combined with core01.example.com:5000,core02.example.com:5000,core03.example.com:5000 will attempt to reach Neo4j instances listening on <i>core01.example.com</i> , core01.example.com and core01.example.com; all on port 5000.
<pre>causal_clustering.raft_adv ertised_address</pre>	The address/port setting that specifies where the Neo4j instance advertises to other members of the cluster that it will listen for Raft messages within the Core cluster.
	<b>Example:</b> causal_clustering.raft_advertised_address=192.168.33.20:7000 will listen for cluster communication in the network interface bound to 192.168.33.20 on port 7000.
<pre>causal_clustering.transact ion_advertised_address</pre>	The address/port setting that specifies where the instance advertises where it will listen for requests for transactions in the transaction-shipping catchup protocol.
	<b>Example:</b> <pre>causal_clustering.transaction_advertised_address=192.168.33.20:6001</pre> will listen for transactions from cluster members on the network interface bound to 192.168.33.20 on port 6001.
<pre>causal_clustering.discover y_listen_address</pre>	The address/port setting for use by the discovery protocol. This is the setting which will be included in the setting causal_clustering.initial_discovery_members which are set in the configuration of the other members of the cluster.
	<b>Example:</b> causal_clustering.discovery_listen_address=0.0.0.0:5001 will listen for cluster membership communication on any network interface at port 5001.
<pre>causal_clustering.raft_lis ten_address</pre>	The address/port setting that specifies which network interface and port the Neo4j instance will bind to for cluster communication. This setting must be set in coordination with the address this instance advertises it will listen at in the setting causal_clustering.raft_advertised_address.
	<b>Example:</b> causal_clustering.raft_listen_address=0.0.0.0:7000 will listen for cluster communication on any network interface at port 7000.

Parameter	Explanation
causal_clustering.transaction_listen_address	The address/port setting that specifies which network interface and port the Neo4j instance will bind to for cluster communication. This setting must be set in coordination with the address this instance advertises it will listen at in the setting causal_clustering.transaction_advertised_address.  Example: causal_clustering.transaction_listen_address=0.0.0.0:6001 will
	listen for cluster communication on any network interface at port 7000.
<pre>causal_clustering.store_co py_max_retry_time_per_requ est</pre>	Condition for when store copy should eventually fail. A request is allowed to retry for any amount of attempts as long as the configured time has not been met. For very large stores or other reason that might make transferring of files slow this could be increased.
	<pre>Example: causal_clustering.store_copy_max_retry_time_per_request=60min</pre>

## 6.7.1. Multi-data center settings

Parameter	Explanation
causal_clustering.multi_dc	Enables multi-data center features. Requires appropriate licensing.
_license	<b>Example:</b> causal_clustering.multi_dc_license=true will enable the multi-data center features.
<pre>causal_clustering.server_g roups</pre>	A list of group names for the server used when configuring load balancing and replication policies.
	<b>Example:</b> causal_clustering.server_groups=us,us-east will add the current instance to the groups us and us-east.
<pre>causal_clustering.upstream _selection_strategy</pre>	An ordered list in descending preference of the strategy which Read Replicas use to choose upstream database server from which to pull transactional updates.
	<b>Example:</b> causal_clustering.upstream_selection_strategy=connect-randomly-within-server-group, typically-connect-to-random-read-replica will configure the behavior so that the Read Replica will first try to connect to any other instance in the group(s) specified in causal_clustering.server_groups. Should we fail to find any live instances in those groups, then we will connect to a random Read Replica. A value of user-defined will enable custom strategy definitions using the setting causal_clustering.user_defined_upstream_strategy.
<pre>causal_clustering.user_def ined_upstream_strategy</pre>	Defines the configuration of upstream dependencies. Can only be used if causal_clustering.upstream_selection_strategy is set to user-defined.
	<b>Example:</b> causal_clustering.user_defined_upstream_strategy=groups(north2); groups(north); halt() will look for servers in the north2. If none are available it will look in the north server group. Finally, if we cannot resolve any servers in any of the previous groups, then rule chain will be stopped via halt().
causal_clustering.load_bal ancing.plugin	The load balancing plugin to use. One pre-defined plugin named server_policies is available by default.
	<b>Example:</b> causal_clustering.load_balancing.plugin=server_policies will enable custom policy definitions.
<pre>causal_clustering.load_bal ancing.config.server_polic ies.<pre>policy-name&gt;</pre></pre>	Defines a custom policy under the name <policy-name>. Note that load balancing policies are cluster-global configurations and should be defined the exact same way on all core machines.</policy-name>
	Example: causal_clustering.load_balancing.config.server_policies.north1_only=groups(north1)  >min(2); halt(); will define a load balancing policy named north1_only. Queries are only sent to servers in the north1 server group, provided there are two of them available. If there are less than two servers in north1 then the chain is halted.

## Chapter 7. Fabric

This chapter describes the configuration and operation of Neo4j Fabric.

#### This chapter describes the following:

- Introduction
- Configuration
- Queries
- Further Considerations

### 7.1. Introduction

This section gives an introduction of Neo4j Fabric.

#### This section describes the following:

- Overview
- Fabric concepts
- · Deployment examples

#### 7.1.1. Overview

Fabric, introduced in Neo4j 4.0, is a way to store and retrieve data in multiple databases, whether they are on the same Neo4j DBMS or in multiple DBMSs, using a single Cypher query. Fabric achieves a number of desirable objectives:

- a unified view of local and distributed data, accessible via a single client connection and user session
- increased scalability for read/write operations, data volume and concurrency
- predictable response time for queries executed during normal operations, a failover or other infrastructure changes
- High Availability and No Single Point of Failure for large data volume.

In practical terms, Fabric provides the infrastructure and tooling for:

- **Data Federation**: the ability to access data available in distributed sources in the form of **disjointed graphs**.
- **Data Sharding**: the ability to access data available in distributed sources in the form of a **common** graph partitioned on multiple databases.

With Fabric, a Cypher query can store and retrieve data in multiple federated and sharded graphs.

## 7.1.2. Fabric concepts

#### The fabric database

A Fabric setup includes a Fabric database, that acts as the entry point to a federated or sharded graph infrastructure. This database is also referred in Fabric as *the virtual database*. Drivers and client applications access and use the Fabric database like any other Neo4j database. The exception is that it

cannot store any data, and queries against it don't return any data. The Fabric database can be configured only on a standalone Neo4j DBMS, i.e. on a Neo4j DBMS where the configuration setting dbms.mode must be set to SINGLE.

### Fabric graphs

In a Fabric database, data is organized in the form of graphs. Graphs are seen by client applications as local logical structures, where physically data is stored in one or more databases. Databases configured as Fabric graphs may be local, i.e. in the same Neo4j DBMS, or they may be located in external Neo4j DBMSes. The databases are accessible by client applications also from regular local connections in their respective Neo4j DBMSs.

## 7.1.3. Deployment examples

Fabric constitutes an extremely versatile environment that provides scalability and availability with no single point of failure in various topologies. Users and developers may use applications that can work on a standalone DBMS as well on a very complex and largely distributed infrastructure without the need to apply any change to the queries accessing the Fabric graphs.

### Development deployment

In its simplest deployment, Fabric can be used on a single instance, where Fabric graphs are associated to local databases. This approach is commonly used by software developers to create applications that will be deployed on multiple Neo4j DBMSs, or by power users who intend to execute Cypher queries against local disjoint graphs.

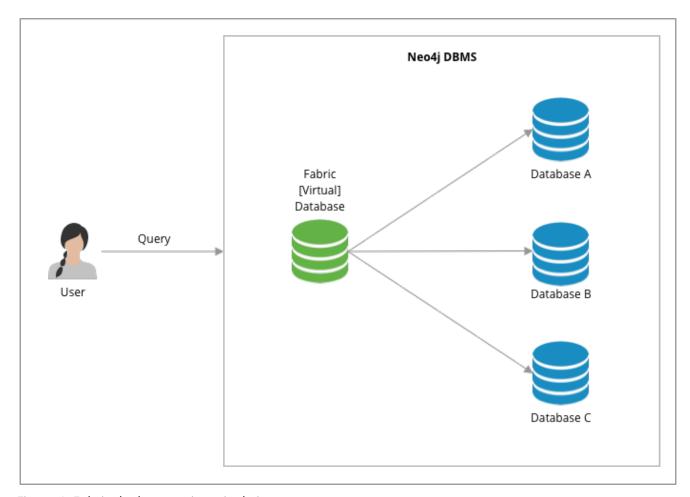


Figure 6. Fabric deployment in a single instance

### Cluster deployment with no single point of failure

In this deployment Fabric guarantees access to disjoint graphs in high availability with no single point of failure. Availability if reached by creating redundant entry points for the Fabric Database (i.e. two standalone Neo4j DBMSs with the same Fabric configuration) and a minimum Causal Cluster of three members for data storage and retrieval. This approach is suitable for production environments and it can be used by power users who intend to execute Cypher queries against disjoint graphs.

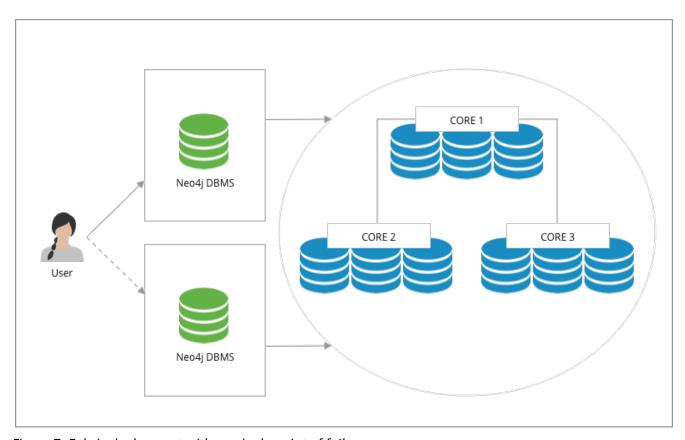


Figure 7. Fabric deployment with no single point of failure

## Multi-cluster deployment

In this deployment Fabric provides high scalability and availability with no single point of failure. Disjoint clusters can be sized according to the expected workload and Databases may be colocated in the same cluster or they can be hosted in their own cluster to provide higher throughput. This approach is suitable for production environments where database can be sharded, federated or a combination of the two.

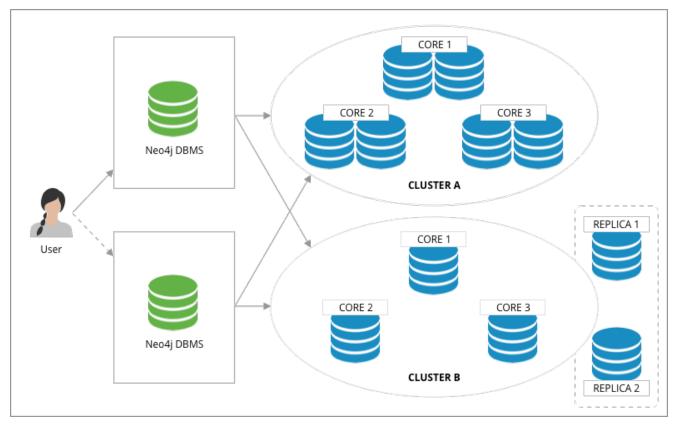


Figure 8. Fabric deployment for scalability with no single point of failure

## 7.2. Configuration

This section describes how to configure Neo4j Fabric.

This section describes the following:

- Fabric database setup
- Authentication and authorization
- Important settings

## 7.2.1. Fabric database setup

Fabric must be set on a standalone Neo4j DBMS: the settings in *neo4j.conf* are identified by the fabric namespace. The minimal requirements to setup Fabric are:

- A **virtual database name**: this is the entry point used by the client applications to access the Fabric environment.
- One or more **Fabric graph URI and database**: this a reference of a URI and a database for each graph set in the Fabric environment.

### Development setup example

Consider a standalone Neo4j DBMS which has two databases, db1 and db2. Note that all databases except for the default and system need to be created using the CREATE DATABASE.

The simplest configuration of Fabric is:

```
fabric.database.name=example
fabric.graph.0.uri=neo4j://localhost:7687
fabric.graph.0.database=db1

fabric.graph.1.uri=neo4j://localhost:7687
fabric.graph.1.database=db2
```

The configuration associates the Fabric database to an existing database named example, and it is accessible using the default URI, i.e. neo4j://localhost:7687. The Fabric graphs are uniquely identified with an ID 0 and 1.

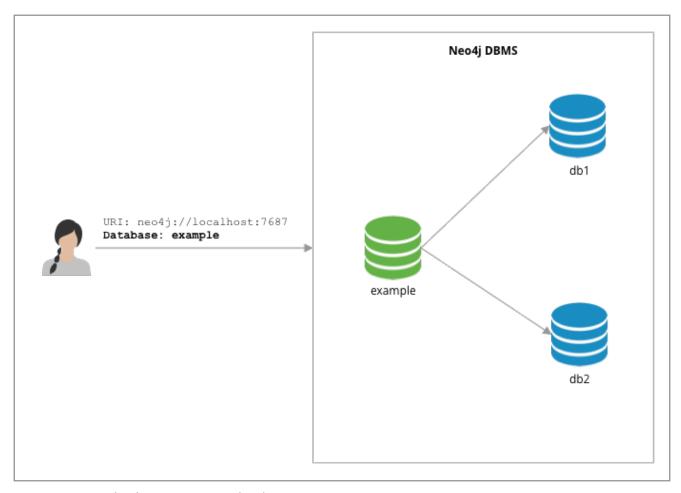


Figure 9. Minimal Fabric setting in a development setup

## Naming graphs

Graphs may be identified by their ID or by a name. A graph can be named by adding an extra configuration setting, fabric.graph.<ID>.name. In the previous example, assuming that the given names are graphA (associated to db1) and graphB (associated to db2), the two additional settings would be:

```
fabric.graph.0.name=graphA
fabric.graph.1.name=graphB
```

## Cluster setup with no single point of failure example

In this example, all components are redundant and data is stored in a Causal Cluster. In addition to the settings described in the previous example, a setting with no single point of failure requires the use of the *routing servers* parameter, which specifies a list of standalone Neo4j DBMSs that expose the same Fabric database and configuration. This parameter is required in order to simulate the same

connectivity that client applications use with Causal Cluster, i.e. in case of fault of one instance the client application may revert to another existing instance.

Assume that in this example, the data is stored in three databases: db1, db2 and db3. The configuration of Fabric would be:

```
dbms.mode=SINGLE

fabric.database.name=example
fabric.routing.servers=server1:7687,server2:7687

fabric.graph.0.name=graphA
fabric.graph.0.uri=neo4j://core1:7687,neo4j://core2:7687,neo4j://core3:7687
fabric.graph.1.name=graphB
fabric.graph.1.uri=neo4j://core1:7687,neo4j://core2:7687,neo4j://core3:7687
fabric.graph.1.database=db2
fabric.graph.1.name=graphC
fabric.graph.1.uri=neo4j://core1:7687,neo4j://core2:7687,neo4j://core3:7687
fabric.graph.1.database=db3
```

The configuration above must be added to the *neo4j.conf* file of the Neo4j DBMSs server1 and server2. The parameter fabric.routing.servers contains the list of available standalone Neo4j DBMSs hosting the Fabric database. The parameter fabric.graph.<ID>.uri can contain a list of URIs, so in case the first server does not respond to the request, the connection can be established to another server that is part of the cluster. The URIs refer to the neo4j:// schema so that Fabric can retrieve a routing table and can use one of the members of the cluster to connect.

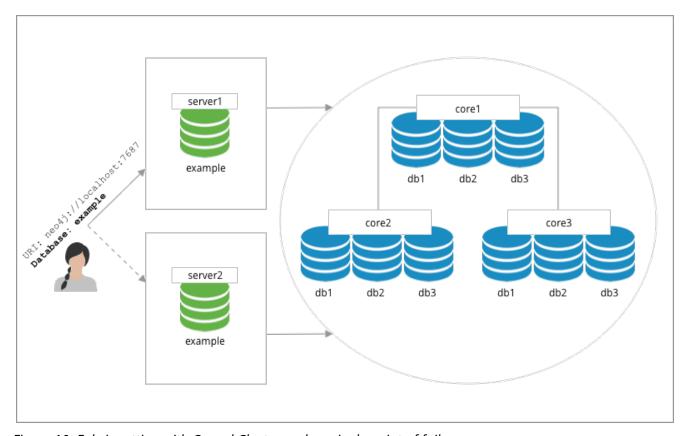


Figure 10. Fabric setting with Causal Cluster and no single point of failure

## Cluster routing context

The URIs in the graph settings may include routing contexts. This can be used to associate a Fabric graph with a filtered subset of Causal Cluster members, by selecting a routing policy.

As an example, assuming we have a server policy called read\_replicas defined in the configuration of the cluster we are targeting, we might set up a Fabric graph that accesses only the read replicas of the cluster.

```
fabric.graph.0.name=graphA
fabric.graph.0.uri=neo4j://core1:7687?policy=read_replicas
fabric.graph.0.database=db1
```

This enables scenarios where queries executed through Fabric are explicitly offloaded to specific instances in clusters.

#### 7.2.2. Authentication and authorization

#### Credentials

Connections between the Fabric database and the Neo4j DBMSs hosting the data are created using the same credentials that are supplied in the client connection to the Fabric database. It is recommended to maintain a set of user credentials on all the Neo4j DBMSs; if required, a subset of credentials may be set for local access on the remote DBMSs.

#### User and role administration

User and role administration actions are not automatically propagated to the Fabric environment, therefore security settings must be executed on any DBMS that is part of Fabric.

### Privileges on the Fabric database

In order to use all Fabric features, users of Fabric databases need ACCESS and READ privileges.

## 7.2.3. Important settings

This section provides general information about Fabric settings and describes the ones important for creating a fabric set-up. Please refer to Configuration settings fort the full list of Fabric configuration options.

Fabric settings are divided in the following categories:

- System Settings: DBMS-level settings.
- **Graph Settings**: definition and configuration of Fabric graphs.
- **Drivers Settings**: configuration of drivers used to access Neo4j DBMSs and databases associated to Fabric graphs.

## System settings

Table 16. Fabric system settings

Parameter	Description
fabric.database.name	Name of the Fabric database. Neo4j Fabric currently supports one Fabric database in a standalone Neo4j DBMS.
fabric.routing.servers	A comma-separated list of Neo4j DBMSs that share the same Fabric configuration. These DBMSs form a routing group. A client application will route transactions through a Neo4j driver or connector to one of the members of the routing group. A Neo4j DBMS is represented by its Bolt connector address. Example: fabric.routing.servers=server1:7687, server2:7687.

### Graph settings



The <ID> in the following settings is the integer associated to each Fabric graph.

Table 17. Fabric graph settings

Parameter	Description
fabric.graph. <id>.uri</id>	URI of the Neo4j DBMS hosting the database associated to the Fabric graph. Example: neo4j://somewhere:7687
fabric.graph. <id>.database</id>	Name of the database associated to the Fabric graph.
fabric.graph. <id>.name</id>	Name assigned to the Fabric graph. The name can be used in Fabric queries.
<pre>fabric.graph.<id>.driver.*</id></pre>	Any specific driver setting, i.e. any setting related to a connection to a specific Neo4j DBMS and database. This setting overrides a global driver setting.



When configuring access to a remote DBMS, please make sure that the remote is configured to advertise its address correctly. This is done through either dbms.default\_advertised\_address Or dbms.connector.bolt.advertised\_address. Fabric reads the routing table from the remote DBMS and then connects back using an appropriate entry in that table.

### Drivers settings

Fabric uses the Neo4j Java driver to connect to and access the data stored in Neo4j databases associated to Fabric graphs. This section presents the most important parameters available to configure the driver.

Drivers settings are configured with parameters with names of the format:

#### fabric.driver.<suffix>

A setting can be global, i.e. be valid for all the drivers used in Fabric, or it can be specific for a given connection to a Neo4j database associated to a graph. The graph-specific setting overrides the global configuration for that graph.

Example 33. Global drivers setting versus graph-specific drivers setting

A drivers setting for Fabric as the following is valid for all the connections established with the Neo4j DBMSs set in Fabric:

```
fabric.driver.api=RX
```

A graph-specific connection for the database with ID=6 will override the fabric.driver.api setting for that database:

fabric.graph.6.driver.api=ASYNC

Table 18. Fabric drivers setting suffixes

Parameter suffix	Explanation
ssl_enabled	SSL for Fabric drivers is configured using the fabric SSL policy. This setting can be used to instruct the driver not to use SSL even though the fabric SSL policy is configured. The driver will use SSL if the fabric SSL policy is configured, and this setting is set to true. This parameter can only be used in fabric.graph. <pre>graph</pre> ID>.driver.ssl_enabled and not fabric.driver.ssl_enabled.
api	Determines which driver API will be used. Supported values are RX and ASYNC. ASYNC must be used when the remote instance is 3.5.



Most driver options described in Driver Manual 
Configuration have an equivalent in Fabric configuration.

## 7.3. Queries

This section provides examples of queries and Cypher commands that can be used with Neo4j Fabric.

This section describes the following:

- Query a single graph
- Query multiple graphs
- Query all graphs
- Query result aggregation
- Correlated subquery
- Updating query
- Mapping functions
- Fabric built-in functions

In this section we will look at a few example queries that show how to perform a range of different tasks.

The examples in this section make use of the two Cypher clauses: USE and CALL {}. The syntax is explained in detail in the Cypher Manual:

- See Cypher Manual [] CALL {} for details about the CALL {} clause.
- See Cypher Manual 

  USE for details about the USE clause.

## 7.3.1. Query a single graph

Example 34. Reading and returning data from a single graph.

USE example.graphA MATCH (movie:Movie) RETURN movie.title AS title The USE clause at the beginning of the query causes the rest of the query to execute against the example.graphA graphA.

## 7.3.2. Query multiple graphs

Example 35. Reading and returning data from two named graphs

```
USE example.graphA
MATCH (movie:Movie)
RETURN movie.title AS title
UNION
USE example.graphB
MATCH (movie:Movie)
RETURN movie.title AS title
```

The first part of the UNION query executes against the example.graphA graph and the second part executes against the example.graphB graph.

## 7.3.3. Query all graphs

Example 36. Reading and returning data from all graphs

```
UNWIND example.graphIds() AS graphId
CALL {
    USE example.graph(graphId)
    MATCH (movie:Movie)
    RETURN movie.title AS title
}
RETURN title
```

We call the built-in function example.graphIds() to get the graph IDs for all remote graphs in our Fabric setup. We UNWIND the result of that function to get one record per graph ID. The CALL {} subquery is executed once per incoming record. We use a USE clause in the subquery with a dynamic graph lookup, causing the subquery to execute once against each remote graph. At the end of the main query we simply RETURN the title variable.

## 7.3.4. Query result aggregation

Example 37. Getting the earliest release year of all movies in all graphs

```
UNWIND example.graphIds() AS graphId
CALL {
    USE example.graph(graphId)
    MATCH (movie:Movie)
    RETURN movie.released AS released
}
RETURN min(released) AS earliest
```

From each remote graph we return the released property of each movie. At the end of the main query we aggregate across the full result to calculate the global minimum.

## 7.3.5. Correlated subquery

Assume that graphA contains American movies and graphB contains European movies. Find all European movies released in the same year as the latest released American movie:

```
CALL {
    USE example.graphA
    MATCH (movie:Movie)
    RETURN max(movie.released) AS usLatest
}
CALL {
    USE example.graphB
    WITH usLatest
    MATCH (movie:Movie)
    WHERE movie.released = usLatest
    RETURN movie
}
RETURN movie
```

We query the example.graphA and return the release year of the latest release. We then query the example.graphB. WITH usLatest is an import clause which lets us refer to the usLatest variable inside the subquery. We find all the movies in this graph that fulfill our condition and return them.

## 7.3.6. Updating query

Example 39. Create a new movie node

```
USE example.graphB
CREATE (m:Movie)
SET m.title = 'Léon: The Professional'
SET m.tagline = 'If you want the job done right, hire a professional.'
SET m.released = 1994
```

## 7.3.7. Mapping functions

Mapping functions are a common Fabric usage pattern. In the previous examples, graphs were identified by providing static graph names in the query. Fabric may be used in scenarios where graphs are identified by a mapping mechanism that can, for example, identify a key of an object contained in a graph. This can be achieved by using user defined functions or other functions that may be already available. These functions ultimately return the ID of a graph in Fabric.

Mapping functions are commonly used in sharding scenarios. In Fabric, shards are associated to graphs, hence mapping functions are used to identify a graph, i.e. a shard.



Refer to Java Reference 

User-defined functions for details on how to create user-defined functions.

Let's assume that Fabric is setup in order to store and retrieve data associated to nodes with the label user. User nodes are partitioned in several graphs (shards) in Fabric. Each user has a numerical userId, which is unique in all Fabric. We decide on a simple scheme where each user is located on a graph determined by taking the userId modulo the number of graphs. We create a user-defined function which implements the following pseudo code:

```
sharding.userIdToGraphId(userId) = userId % NUM_SHARDS
```

Assuming we have supplied a query parameter \$userId with the specific userId that we are interested in, we use our function in this way:

USE example.graph( sharding.userIdToGraphId(\$userId) )
MATCH (u:User) WHERE u.userId = \$userId
RETURN u

### 7.3.8. Fabric built-in functions

Fabric functions are located in a namespace corresponding to a Fabric database in which they are used. The following table provides a description of Fabric built-in functions:

Table 19. Fabric built-in functions

Function	Explanation
<pre><fabric database="" name="">.graphIds()</fabric></pre>	Provides a list of IDs of all remote graph configured for the given Fabric database. This function is supported only in USE clauses
<pre><fabric database="" name="">.graph(graphId)</fabric></pre>	Maps a graph ID to a Graph. It accepts a graph ID as a parameter and returns a graph representation accepted by USE clause. This function is supported only in USE clauses

### 7.4. Further considerations

This section presents considerations about Fabric that developers and administrators must be aware of.

#### DBMS mode

The DBMS hosting the Fabric virtual database cannot be part of a Causal Cluster: it can only be a DBMS with dbms.mode=SINGLE.

#### Database compatibility

Fabric is part of Neo4j DBMS and does not require any special installation or plugin. Fabric graphs can be associated to databases available on Neo4j DBMS version 3.5 or 4.0.

#### Fabric configuration

The Neo4j DBMSs that host the same Fabric virtual database must have the same configuration settings. The configuration must be kept in-sync and applied by the Database Administrator.

#### Security credentials

The Neo4j DBMSs that host the same Fabric virtual database must have the same user credentials. Any change of password on a machine that is part of Fabric, must be kept in-sync and applied to all the Neo4j DBMSs that are part of Fabric.

#### Transactions in Fabric

In Fabric, ACID compliance is guaranteed only within a single graph. This means, that the current version of Fabric does not support transactions that span across multiple graphs. To avoid common mistakes that may lead to data corruption, Fabric does not allow write operations on more than one graph within the same transaction. Transactions with queries that read from multiple graphs, or read from multiple graphs and write to a single graph, are allowed.

#### Administation commands

Fabric does not support issuing Cypher administration commands, on, or through the Fabric virtual database. Any database management commands, index and constraint management commands or user and security management commands must be issued directly to the DBMSs and databases that are part of the Fabric setup.

Neo4j embedded Fabric is not available when Neo4j is used as an embedded database in Java applications. Fabric can be used only in a typical client/server mode, when users connect to a Neo4j DBMS from their client application or tool, via Bolt or HTTP protocol.

# Chapter 8. Upgrade

This chapter describes how to upgrade Neo4i from an earlier version.

It is recommended that your installation of Neo4j is kept up to date. Upgrading your installation to Neo4j 4.0.0 will ensure that you are provided with improvements in performance and security, as well as any latest bug fixes.

#### This chapter describes the following:

- Upgrade planning
  - Supported upgrade paths
  - Limitations
  - ☐ Prepare to upgrade
- Single-instance upgrade
- Upgrade a Causal Cluster

## 8.1. Upgrade planning

This section describes how to plan for an upgrade of Neo4j.

#### This guide describes the following:

- Supported upgrade paths
- Limitations
- Prepare to upgrade

## 8.1.1. Supported upgrade paths

The following upgrade path is supported:

3.5.any [] 4.0.0

The following steps are required if you need to upgrade from a version earlier than 3.5:

- 1. Upgrade to version 3.5. latest by following the instructions in the Neo4j Operations Manual for 3.5.
- 2. Upgrade to version 4.0.0 as per instructions in this manual.

#### 8.1.2. Limitations

- Neo4j does not support downgrades. If the upgrade is not successful, you have to do a full rollback, including restoring a pre-upgrade backup.
- A Neo4j upgrade must be performed as an isolated operation. If you are planning to upgrade from a single-instance installation to a Causal Cluster, this must be performed separately from the migration to 4.0.0.
- In order to further minimize risk, it is recommended that while migrating, you do not switch from Community Edition to Enterprise Edition, change configuration, perform architectural restructuring, or similar tasks.

## 8.1.3. Prepare to upgrade

Refer to the 4.0 Migration Guide for details on how to prepare for this upgrade.

## 8.2. Upgrade a single instance

This section describes how to upgrade a single Neo4j instance.

For instructions on upgrading a Neo4j Causal Cluster, see Upgrade a Causal Cluster.

#### Pre-upgrade steps

- Read Supported upgrade paths and Limitations regarding supported upgrade paths before you start planning your upgrade.
- Read the 4.0 Migration Guide thoroughly and perform all the steps listed there.

#### Shutdown and backup

- 1. If the database is running, shut it down cleanly.
- 2. Perform and verify backups:
  - ☐ Back up *neo4j.conf*.
  - ☐ Back up all the files used for encryption, i.e. private key, public certificate, and the contents of the *trusted* and *revoked* directories. The locations of these are described in SSL framework.
  - ☐ Verify that you have a full backup that is stored in a safe location, either using the online backup tool or offline backups.

#### Upgrade

- 1. Install Neo4j 4.0.0 using one of the following methods, specific to your technology:
  - a. If using a tarball or zipfile for installation:
    - i. Untar or unzip Neo4j 4.0.0.
    - ii. Transfer the new *neo4j.conf* that you prepared in the *Apply configuration changes* step in Upgrade planning.
    - iii. Set dbms.allow\_upgrade=true in *neo4j.conf* of the 4.0.0 installation. Neo4j will fail to start without this configuration.
    - iv. Copy the files used for encryption from the old installation to the new one.
    - v. Copy the *data* directory from the old installation to the new one. This step is not applicable if you have <a href="dbms.directories.data">dbms.directories.data</a> pointing to a directory outside of <a href="https://www.new.directories.data">NEO4J\_HOME</a>.
    - vi. If using custom plugins, place the plugins that are adjusted for the new version in the */plugins* directory.
  - b. If using a Debian or RPM distribution:
    - i. Set dbms.allow\_upgrade=true in neo4j.conf.
    - ii. Install Neo4j 4.0.0.
    - iii. When prompted, review the differences between the *neo4j.conf* files of the previous version and Neo4j 4.0.0. Transfer any custom settings to the 4.0.0 installation, as noted under the *Apply configuration changes* step in 4.0 Migration Guide [] Prepare to upgrade. Make sure to preserve dbms.allow\_upgrade=true as set in the instruction above. Neo4j will fail to start without this configuration.
    - iv. If using custom plugins, place the plugins that are adjusted for the new version in the

/plugins directory.

2. Start up Neo4j 4.0.0. The database upgrade will take place during startup.

The *neo4j.log* file contains valuable information on how many steps the upgrade will involve and how far it has progressed. For large upgrades, it is a good idea to monitor this log continuously.

#### Post-upgrade steps

- 1. When the upgrade has finished, dbms.allow\_upgrade should be set to false or be removed.
- 2. Restart the database.
- 3. It is good practice to make a full backup immediately after the upgrade.

## 8.3. Upgrade a Causal Cluster

This section describes how to upgrade a Neo4j Causal Cluster.

#### Pre-upgrade steps

- 1. Read Supported upgrade paths and Limitations regarding supported upgrade paths before you start planning your upgrade.
- 2. Read the 4.0 Migration Guide thoroughly and perform all the steps listed there.
- 3. Perform and verify backups:
  - ☐ Back up *neo4j.conf*.
  - Back up all the files used for encryption, i.e. private key, public certificate, and the contents of the *trusted* and *revoked* directories. The locations of these are described in SSL framework. You should back up these files on each server in the cluster.
  - ☐ Verify that you have a full backup that is stored in a safe location.
- 4. Prepare a new *neo4j.conf* file for each of the servers in the cluster, following the instructions under the *Apply configuration changes* step in 4.0 Migration Guide 

  Prepare to upgrade.
- 5. If using custom plugins, make sure that you have the plugins that are adjusted for the new version in an accessible location.

#### Limitations

- Neo4j does not support downgrades. If the upgrade is not successful, you have to do a full rollback, including restoring a pre-upgrade backup.
- In order to minimize risk, it is recommended that while upgrading, you do not change configuration, perform architectural restructuring, or similar tasks.

#### **Downtime**

The migration from 3.5 to 4.x will involve downtime. A test upgrade on a production-like equipment provides information on the duration of the downtime.

#### **Upgrade steps**

- 1. Shut down all the servers in the cluster
- 2. On one of the Core Servers:
  - a. Install Neo4j using one of the following methods, specific to your technology:
    - ☐ If using a tarball or zipfile for installation:
      - i. Untar or unzip the version of Neo4j that you want to upgrade to.
      - ii. Transfer the new *neo4j.conf* that you prepared in the pre-upgrade steps.

- iii. Set dbms.mode=SINGLE in neo4j.conf.
- iv. Set dbms.allow\_upgrade=true in *neo4j.conf*. Neo4j will fail to start without this configuration.
- v. Copy the files used for encryption from the old installation to the new one.
- vi. Copy the *data* directory from the old installation to the new one. This step is not applicable if you have dbms.directories.data pointing to a directory outside of *NEO4J HOME*.
- vii. If using custom plugins, place the plugins that are adjusted for the new version in the */plugins* directory.
- ☐ If using a Debian or RPM distribution:
  - i. Set dbms.mode=SINGLE in neo4j.conf.
  - ii. Set dbms.allow\_upgrade=true in neo4j.conf.
  - iii. Install the version of Neo4j that you want to upgrade to.
  - iv. When prompted, review the differences between the *neo4j.conf* files of the previous version and Neo4j 4.0.0. Transfer any custom settings to the new installation, as prepared in the pre-upgrade step. Make sure to preserve dbms.mode=SINGLE and dbms.allow\_upgrade=true as set in the instruction above.
  - v. If using custom plugins, place the plugins that are adjusted for the new version in the */plugins* directory.
- b. Start up Neo4j. The database upgrade will take place during startup.

The *neo4j.log* file contains valuable information on how many steps the upgrade will involve and how far it has progressed. For large upgrades, it is a good idea to monitor this log continuously.

- c. Stop your Neo4j database once again.
- d. Set dbms.allow\_upgrade=false, or remove it.
- e. Set dbms.mode=CORE in neo4i.conf to re-enable Causal Clustering in the configuration.
- f. Use neo4j-admin dump to make a copy of the database.
- g. Do not yet restart the database.
- 3. On each of the other Core Servers:
  - a. Delete the database directory (in a default configuration, this is the directory *databases/neo4j* which is located in the *data* directory).
  - b. Install the version of Neo4j that you want to upgrade to.
  - c. Transfer any custom settings to the new installation, as prepared in the pre-upgrade step.
  - d. If using a tarball or zipfile for installation: Copy the files used for encryption from the old installation to the new one.
  - e. If using custom plugins, place the plugins that are adjusted for the new version in the */plugins* directory.
  - f. Perform neo4j-admin unbind on the instance.
  - g. Using neo4j-admin load, restore the upgraded database onto this server.
- 4. Startup all the Core Servers and see the cluster form.
- 5. On each of the Read Replica servers:
  - a. Stop Neo4j.
  - b. Delete the database directory (in a default configuration, this is the directory *databases/neo4j* which is located in the *data* directory).

- c. Install the version of Neo4j that you want to upgrade to.
- d. Transfer any custom settings to the new installation, as prepared in the pre-upgrade step.
- e. If using a tarball or zipfile for installation: Copy the files used for encryption from the old installation to the new one.
- f. Perform neo4j-admin unbind on the instance.
- g. If using custom plugins, place the plugins that are adjusted for the new version in the */plugins* directory.
- h. Using neo4j-admin dump/load, restore the upgraded database onto this server. Alternatively, you can omit this step and let the Read Replica do a complete store copy.
- i. Start the Read Replica and see it join the cluster.

## Chapter 9. Backup

This chapter covers how to perform and restore backups of a Neo4j DBMS deployed as a Causal Cluster or a single instance.

#### This chapter describes the following:

- Backup planning
  - □ Introduction
  - □ Online and offline backups
  - □ Server configuration
  - □ Databases to backup
  - ☐ Storage considerations
  - □ Cluster considerations
  - Using SSL/TLS for backups
  - Additional files to back up
- Perform a backup
  - Backup command
  - Backup process
  - Memory configuration
- Restore a backup
  - □ Restore commands
  - ☐ Restore a single database
  - □ Restore a cluster

## 9.1. Backup planning

This section explains how to prepare for backing up a Neo4j deployment.

#### This section includes:

- Introduction
- Online and offline backups
- Server configuration
- Databases to backup
- Storage considerations
- Cluster considerations
- Using SSL/TLS for backups
- · Additional files to back up

### 9.1.1. Introduction

Designing an appropriate backup strategy for your Neo4j DBMS is a fundamental part of operations. The backup strategy should take into account elements such as:

- Demands on performance during backup actions.
- Tolerance for data loss in case of failure.
- Tolerance for downtime in case of failure.
- · Data volumes.

The backup strategy will answer question such as:

- What type of backup method used; online or offline backups?
- What physical setup meets our demands?
- What backup media offline or remote storage, cloud storage etc. should we use?
- · How long do we archive backups for?
- · With what frequency should we perform backups;

If using online backups:

- ☐ How often should we perform full backups?
- ☐ How often should we perform incremental backups?
- How do we test recovery routines, and how often?

## 9.1.2. Online and offline backups

Online backups are typically required for production environments, but it is also possible to perform offline backups.

Offline backups are a more limited method for backing up a database. For example:

- Online backups run against a live Neo4j instance, while offline backups require that the database is shut down.
- Online backups can be full or incremental, but there is no support for backing up incrementally with offline backups.

For more details about offline backups, see Dump and load databases.

The remainder of this chapter is dedicated to describing *online* backups.

## 9.1.3. Server configuration

The table below lists the basic server parameters relevant to backups. Note that by default the backup service is enabled but only listens on localhost (127.0.0.1) and this needs to be changed if backups are to be taken from another machine.

Table 20. Server parameters for backups

Parameter name	Default value	Description
dbms.backup.enabled	true	Enable support for running online backups.
dbms.backup.listen_address	127.0.0.1:6362	Listening server for online backups.

## 9.1.4. Databases to backup

Since a Neo4j DBMS can host multiple databases and they are backed up independently of one another, it is important to plan a backup strategy for every database and to not forget any databases. In a new deployment there are two databases by default, neo4j and system. The system database

contains configuration, e.g. operational states of databases, security configuration, etc.

## 9.1.5. Storage considerations

For any backup it is important that you store your data separately from the production system, where there are no common dependencies, and preferably off-site. If you are running Neo4j in the cloud, you could for example use a different availability zone or even a separate cloud provider.

Since backups are kept for a long time, the longevity of archival storage should be considered as part of backup planning.

You may also want to override the settings used for pruning and rotation of transaction log files. The transaction log files are files that keep track of recent changes. Please note that removing transaction logs manually can result in a broken backup.

Recovered servers do not need all of the transaction log files that have already been applied, so it is possible to reduce storage size even further by reducing the size of the files to the bare minimum.

This can be done by setting dbms.tx\_log.rotation.size=1M and dbms.tx\_log.rotation.retention\_policy=3 files. Alternatively you can use the --additional-config override.

### 9.1.6. Cluster considerations

In a cluster it is possible to take a backup from any server, and each server has two configurable ports capable of serving a backup. These ports are configured by <a href="mailto:dbms.backup.listen.address">dbms.backup.listen.address</a> and <a href="mailto:causal\_clustering.transaction\_listen\_address">causal\_clustering.transaction\_listen\_address</a> respectively. Functionally they are equivalent for backups, but separating them can allow some operational flexibility, while using just a single port can simplify the configuration.

It is generally recommended to select Read Replicas to act as backup servers, since they are more numerous than Core Servers in typical cluster deployments. Furthermore, the possibility of performance issues on a Read Replica, caused by a large backup, will not affect the performance or redundancy of the Core Cluster. If a Read Replica is not available, then a Core can be picked based on factors such as its physical proximity, bandwidth, performance, and liveness.

Note that both Read Replicas and Cores can fall behind the leader and be out-of-date. We can look at transaction IDs in order to avoid taking a backup from a server that has lagged too far behind. The latest transaction ID can be found by exposing Neo4j metrics or via Neo4j Browser. To view the latest processed transaction ID (and other metrics) in Neo4j Browser, type : sysinfo at the prompt.

## 9.1.7. Using SSL/TLS for backups

The backup server can be configured to require SSL/TLS. If that is the case then the backup client must also be configured to use it with a compatible policy. Refer to SSL framework to learn how SSL is configured in general. See below table for more details about how configured SSL policies map to the configured ports.

Table 21. Mapping backup configuration to SSL policies

Backup target address on database server	SSL policy setting on database server	SSL policy setting on backup client	Default port
dbms.backup.listen_address	dbms.ssl.policy.backup	dbms.ssl.policy.backup	6362
<pre>causal_clustering.transact ion_listen_address</pre>	dbms.ssl.policy.cluster	dbms.ssl.policy.backup	6000

## 9.1.8. Additional files to back up

The files listed below are not included in online nor offline backups. Make sure to back them up separately.

- If you have a cluster, it may be relevant to back up *neo4j.conf* on each server.
- Back up all the files used for SSL/TLS, i.e. private keys, public certificates, and the contents of the *trusted* and *revoked* directories. The locations of these are described in SSL framework. If you have a cluster, you should back up these files on each server in the cluster.

## 9.2. Perform a backup

This section describes how to perform an online backup of a Neo4j database.



Remember to back up all of your created databases, including the system database.

#### This section includes:

- Backup command
- Backup process
- Memory configuration

## 9.2.1. Backup command

A Neo4j database can be backed up in online mode using the backup command of neo4j-admin. The machine that runs the backup command must have Neo4j installed, but does not need to run a Neo4j server.

#### **Syntax**

#### **Options**

Option	Default	Description
backup-dir		Directory to place backup in.
verbose	false	Enable verbose output.
from	localhost:6362	Host and port of Neo4j.
database	neo4j	Name of the database to back up. If a backup of the specified database exists in the target directory, then an incremental backup will be attempted.

Option	Default	Description
fallback-to-full	true	If an incremental backup fails backup will move the old backup to <name>.err.<n> and fallback to a full backup instead.</n></name>
pagecache	8M	The size of the page cache to use for the backup process.
check-consistency	true	If a consistency check should be made.
check-graph	true	Perform checks between nodes, relationships, properties, types and tokens.
check-indexes	true	Perform checks on indexes.
check-label-scan-store	true	Perform checks on the label scan store.
check-property-owners	false	Perform additional checks on property ownership. This check is <b>very</b> expensive in time and memory.
report-dir		Directory where consistency report will be written.
additional-config		Configuration file to supply additional configuration in.

#### Exit codes

neo4j-admin backup will exit with different codes depending on success or error. In the case of error, this includes details of what error was encountered.

Table 22. Neo4j Admin backup exit codes

Code	Description
0	Success.
1	Backup failed.
2	Backup succeeded but consistency check failed.
3	Backup succeeded but consistency check found inconsistencies.

## 9.2.2. Backup process

The backup client can operate in two slightly different modes referred to as performing a *full backup* or an *incremental backup*. A full backup is always required initially for the very first backup into a target location. Subsequent backups will attempt to use the incremental mode where just the delta of the transcation logs since the last backup are transferred and applied onto the target location. If the required transaction logs aren't available on the backup server then the backup client will fallback to performing a full backup instead, unless --fallback-to-full is disabled.

After the backup has been successfully performed the *consistency checker* will be invoked by default. Checking the consistency of the backup is a major operation which can consume significant computational resources, e.g. memory, CPU, I/O.



It is strongly discouraged to run the backup client on a live Neo4j server, especially together with a consistency check. Doing so can adversely affect the server.

To avoid adversely affecting a running server with the resource demands of the backup client it is recommended to take the backup and perform the consistency check on a dedicated machine which has sufficient free resources to perform the consistency check. Another alternative is to decouple the backup operation from the consistency checking and schedule that part of the workflow to happen at

a later point in time on a dedicated machine. The value of consistency checking a backup should not be underestimated as it is vital for safe guarding and ensuring the quality of your data.

The transaction log files in the backup are rotated and pruned based on the provided configuration. For example, setting dbms.tx\_log.rotation.retention\_policy=3 files will keep 3 transaction log files in the backup. You can use the --additional-config parameter to override this configuration.

#### Example 40. Full backup

In this example, set environment variables in order to control memory usage.

The page cache is defined by using the command line option --pagecache. Further, the HEAP\_SIZE environment variable will specify the maximum heap size allocated to the backup process.

Now you can perform a full backup:

```
$neo4j-home> export HEAP_SIZE=2G
$neo4j-home> mkdir /mnt/backups
$neo4j-home> bin/neo4j-admin backup --from=192.168.1.34 --backup-dir=/mnt/backups/neo4j --database
=neo4j --pagecache=4G
Doing full backup...
2017-02-01 14:09:09.510+0000 INFO [o.n.c.s.StoreCopyClient] Copying neostore.nodestore.db.labels
2017-02-01 14:09:09.537+0000 INFO [o.n.c.s.StoreCopyClient] Copied neostore.nodestore.db.labels 8.00
kB
2017-02-01 14:09:09.538+0000 INFO [o.n.c.s.StoreCopyClient] Copying neostore.nodestore.db
2017-02-01 14:09:09.538+0000 INFO [o.n.c.s.StoreCopyClient] Copying neostore.nodestore.db
2017-02-01 14:09:09.540+0000 INFO [o.n.c.s.StoreCopyClient] Copied neostore.nodestore.db
16.00 kB
...
...
```

If you do a directory listing of /mnt/backups you will now see that you have a backup in a directory called neo4j.

#### Example 41. Incremental backup

This example assumes that you have performed a full backup as per the previous example. In the same way as before, make sure to control the memory usage.

To perform an incremental backup you need to specify the location of your previous backup:

```
$neo4j-home> export HEAP_SIZE=2G
$neo4j-home> bin/neo4j-admin backup --from=192.168.1.34 --backup-dir=/mnt/backups/neo4j --database
=neo4j --pagecache=4G
Destination is not empty, doing incremental backup...
Backup complete.
```

## 9.2.3. Memory configuration

The following options are available for configuring the memory allocated to the backup client:

#### Configure heap size for the backup

This is done by setting the environment variable HEAP\_SIZE before starting the backup program. If not specified by HEAP\_SIZE, the Java Virtual Machine will choose a value based on server resources. HEAP\_SIZE configures the maximum heap size allocated for the backup process.

#### Configure page cache for the backup

The page cache size can be determined for the backup program by using the --pagecache option to the neo4j-admin backup command. If not explicitly defined, the page cache will default to 8MB.

## 9.3. Restore a backup

This section describes how to restore from backups of a Neo4j deployment.

#### This section includes:

- · Restore command
- Restore a standalone server
- Restore a cluster

#### 9.3.1. Restore command

A Neo4j database can be restored using the restore command of neo4j-admin.

Note that restoring a database does not configure it as being created in the Neo4j DBMS. That configuration lives in the *system* database and if a backup of the system database is restored then any configuration for other databases in it will also be restored. For the case where the restored database is not created yet, use the CREATE DATABASE command after restoring to create it.

#### **Syntax**

```
neo4j-admin restore --from=<path> [--verbose] [--database=<database>] [--force]
```

#### **Options**

Option	Default	Description
from		Path to backup to restore from.
database	neo4j	Name of database.
force		If an existing database should be replaced.

#### 9.3.2. Restore a standalone server

To restore from backups, follow these steps:

- 1. If the server is running, shut it down.
- 2. Run neo4j-admin restore for every database.
- 3. Start up the server.

#### Example 42. Restore a standalone server

Restore the databases system and neo4j from the backups located in /mnt/backups. Note that the server must be shut down.

```
neo4j-home> bin/neo4j stop
neo4j-home> bin/neo4j-admin restore --from=/mnt/backups/system --database=system --force
neo4j-home> bin/neo4j-admin restore --from=/mnt/backups/neo4j --database=neo4j --force
neo4j-home> bin/neo4j start
```

## 9.3.3. Restore a cluster

To restore a Causal Cluster from backups, follow these steps:

- 1. Shut down all server instances in the cluster.
- 2. Run the neo4j-admin unbind command on each of the Core Servers.
- 3. Restore the backups on each instance, using neo4j-admin restore.
- 4. If you are restoring onto new hardware, please review the Causal Clustering settings in neo4j.conf.

In particular, check the settings causal\_clustering.initial\_discovery\_members, causal\_clustering.minimum\_core\_cluster\_size\_at\_formation, and causal\_clustering.minimum\_core\_cluster\_size\_at\_runtime, and ensure that they correctly reflect the new setup.

5. Start the server instances.

Refer to Seed a cluster for more detailed information.

# Chapter 10. Authentication and authorization

This chapter describes authentication and authorization in Neo4j.

Ensure that your Neo4j deployment adheres to your company's information security guidelines by setting up the appropriate authentication and authorization rules.

Ths section describes the following:

- Introduction
- Built-in roles
- Fine-grained access control
- Integration with LDAP
- Manage procedure permissions
- Terminology



The functionality described in this section is applicable to Enterprise Edition. A limited set of user management functions are also available in Community Edition. Native roles overview gives a quick overview of these.

## 10.1. Introduction

This section provides an overview of authentication and authorization in Neo4j.

Authentication is the process of ensuring that a user is who the user claims to be, while authorization pertains to checking whether the authenticated user is allowed to perform a certain action. Authorization is managed using role-based access control (*RBAC*). Permissions that define access control are assigned to roles, which are in turn assigned to users.

Neo4j has the following auth providers, that can perform user authentication and authorization:

#### Native auth provider

Neo4j provides a native auth provider that stores user and role information in the system database. This option is controlled by the parameter dbms.security.auth\_enabled, which is set to true by default. The Cypher commands to manage users, roles and permissions are described in detail in Cypher Manual 

Administration. Various scenarios that describe the use of the native auth provider are available in Fine-grained access control.

### LDAP auth provider

Another way of controlling authentication and authorization is through external security software such as Active Directory or OpenLDAP, which is accessed via the built-in LDAP connector. A description of the LDAP plugin using Active Directory is available in Integration with LDAP.

#### Custom-built plugin auth providers

For clients with specific requirements not satisfied with either native or LDAP, Neo4j provides a plugin option for building custom integrations. It is recommended that this option is used as part of a custom delivery as negotiated with Neo4j Professional Services. The plugin is described in Java Reference 

Authentication and authorization plugins.

#### Kerberos authentication and single sign-on

In addition to LDAP, Native and custom providers, Neo4j supports Kerberos for authentication and single sign-on. Kerberos support is provided via the Neo4j Kerberos Add-On.

## 10.2. Built-in roles

This section describes the roles that come pre-defined with Neo4j.

Neo4j provides the following native roles:

#### reader

Read-only access to the data graph (all nodes, relationships, properties).

#### editor

- Read/write access to the data graph.
- Write access limited to creating and changing existing properties key, node labels, and relationship types of the graph.

#### publisher

• Read/write access to the data graph.

#### architect

- Read/write access to the data graph.
- Set/delete access to indexes along with any other future schema constructs.

#### admin

- Read/write access to the data graph.
- Set/delete access to indexes along with any other future schema constructs.
- · View/terminate queries.

A user who has no assigned roles will not have any rights or capabilities regarding the data, not even read privileges. A user may have more than one assigned role, and the union of these determine what action(s) on the data may be undertaken by the user.

When an administrator suspends or deletes another user, the following rules apply:

- Administrators can suspend or delete any other user (including other administrators), but not themselves.
- Deleting a user terminates all of the user's running queries and sessions.
- All queries currently running for the deleted user are rolled back.
- The user will no longer be able to log back in (until re-activated by an administrator if suspended).
- There is no need to remove assigned roles from a user prior to deleting the user.

The set of actions on the data and database prescribed by each role are described below. The subset of the functionality which is available with Community Edition is also included:

Table 23. Native roles overview

Action	reader	editor	publisher	architect	admin	(no role)	Available in Community Edition
Change own password	X	X	X	X	X	X	Х
View own details	X	X	X	X	X	X	Х
Read data	X	X	X	X	X		X

Action	reader	editor	publisher	architect	admin	(no role)	Available in Community Edition
View own queries	Х	Х	Х	Х	Х		
Terminate own queries	Х	X	X	Х	Х		
Write/update /delete data		X	X	X	X		Х
Create new types of properties key			X	Х	Х		X
Create new types of nodes labels			X	X	X		X
Create new types of relationship types			X	X	X		X
Create/drop index/constra int				X	X		X
Create/delete user					X		X
Change another user's password					X		
Assign/remov e role to/from user					Х		
Suspend/acti vate user					X		
View all users					X		Х
View all roles					X		
View all roles for a user					X		
View all users for a role					X		
View all queries					X		
Terminate all queries					X		
Dynamically change configuration (see Dynamic settings)					X		

# 10.3. Fine-grained access control

This section contains a worked example that illustrates various aspects of security and finegrained access control.

The topics described in this section are:

- The data model
- Security
- Access control using built-in roles
- Sub-graph access control using privileges

## 10.3.1. The data model

Consider a *healthcare* database, as could be relevant in a medical clinic or hospital. A simple version of this might contain only three labels, representing three entity types:

#### (:Patient)

Nodes of this type represent patients that visit the clinic because they have some symptoms. Information specific to the patient can be captured in properties:

- Name
- SSN
- Address
- · Date of birth

#### (:Symptom)

A medical database contains a catalog of known illnesses and associated symptoms, which can be described using properties:

- Name
- Description

#### (:Disease)

A medical database contains a catalog of known illnesses and associated symptoms, which can be described using properties:

- Name
- Description

These entities will be modelled as nodes, and connected using relationships of the following types:

#### (:Patient)-[:HAS]→(:Symptom)

When a patient reports to the clinic, they will describe their symptoms to the nurse or the doctor. The nurse or doctor will then enter this information into the database in the form of connections between the patient node and a graph of known symptoms. Possible properties of interest on this relationship could be:

Date - date when symptom was reported

#### (:Symptom)-[:OF]→(:Disease)

The graph of known symptoms is part of a graph of diseases and their symptoms. The relationship between a symptom and a disease can include a probability factor for how likely or common it is for people with that disease to express that symptom. This will make it easier for the doctor to make a diagnosis using statistical queries.

Probability - probability of symptom matching disease

#### (:Patient)-[:DIAGNOSIS]→(:Disease)

The doctor can use the graph of diseases and their symptoms to perform an initial investigation into the most likely diseases to match the patient. Based on this, and their own assessment of the

patient, they may make a diagnosis which they would persist to the graph through the addition of this relationship with appropriate properties:

- By: doctor's name
- · Date: date of diagnosis
- · Description: additional doctors' notes

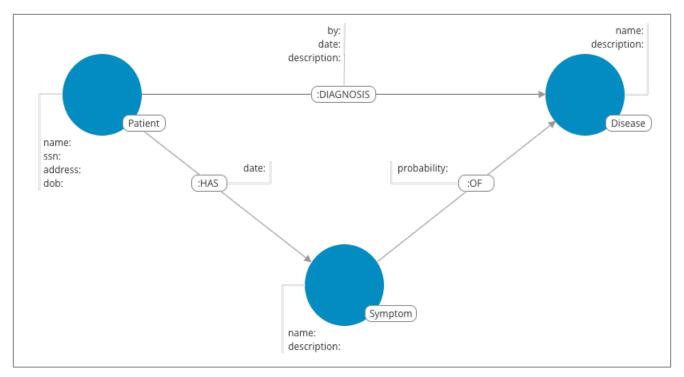


Figure 11. Healthcare use case

The database would be used by a number of different user types, with different needs for access.

- **Doctors** who need to perform diagnosis on patients.
- Nurses who need to treat patients.
- Receptionists who need to identify and record patient information.
- Researchers who need to perform statistical analysis of medical data.
- IT administrators who need to administer the database, creating and assigning users.

# 10.3.2. Security

When building an application for a specific domain, it usual to model the different users within the application itself. However, when working with a database that provides rich user management with roles and privileges, it is possible to model these entirely within the database security model. This results in separation of concerns for the access control to the data and the data itself. We will show two approaches to using Neo4j security features to support the *healthcare* database application. First, a simple approach using built-in roles, and then a more advanced approach using fine-grained privileges for sub-graph access control.

Our *healthcare* example will involve five users of the database:

- · Alice the doctor
- Daniel the nurse
- Bob the receptionist
- Charlie the researcher

· Tina the IT administrator

These users can be created using the CREATE USER command:

#### Example 43. Creating users

```
CREATE USER charlie SET PASSWORD $secret1 CHANGE NOT REQUIRED;
CREATE USER alice SET PASSWORD $secret2 CHANGE NOT REQUIRED;
CREATE USER daniel SET PASSWORD $secret3 CHANGE NOT REQUIRED;
CREATE USER bob SET PASSWORD $secret4 CHANGE NOT REQUIRED;
CREATE USER tina SET PASSWORD $secret5 CHANGE NOT REQUIRED;
```

At this point the users have no ability to interact with the database, so we need to grant those capabilities using roles. There are two different ways of doing this, either by using the built-in roles or a more fine-grained access control using privileges and custom roles.

# 10.3.3. Access control using built-in roles

Neo4j 4.0 comes with a number of built-in roles that cover a number of common needs:

- reader Can only read data from the database.
- editor Can read and update the database, but not expand the schema with new labels, relationship types or property names.
- publisher Can read and edit, as well as add new labels, relationship types and property names.
- architect Has all the capabilities of the publisher as well as the ability to manage indexes and constraints.
- admin Can perform architect actions as well as manage database, users, roles and privileges.

Charlie is a researcher and will not need write access to the database, and so he is assigned the reader role. Alice the doctor, Daniel the nurse and Bob the receptionist all need to update the database with new patient information, but do not need to expand the schema with new labels, relationship types, property names or index. We assign them all the editor role. Tina is the IT administrator that installs and manages the database. In order to create all other users, Tina is assigned the admin role.

#### Example 44. Granting roles

```
GRANT ROLE reader TO charlie;
GRANT ROLE editor TO alice;
GRANT ROLE editor TO daniel;
GRANT ROLE editor TO bob;
GRANT ROLE admin TO tina;
```

A limitation of this approach is that it does allow all users to see all data in the database, and in many real-world scenarios we would prefer to restrict the users' access. For example, we would want to restrict the researcher from being able to read any personal information about the patients, and the receptionist should only be able to see the patient records and no more.

These, and more restrictions, could be coded into the application layer. However, it is possible and more secure to enforce these kinds of fine-grained restrictions directly within the Neo4j security model, by creating custom roles and assigning specific privileges to those roles. 2 Since we will be creating new custom roles, the first thing to do is revoke the current roles from the users:

#### Example 45. Revoking roles

```
REVOKE ROLE reader FROM charlie;
REVOKE ROLE editor FROM alice;
REVOKE ROLE editor FROM daniel;
REVOKE ROLE editor FROM bob;
REVOKE ROLE admin FROM tina;
```

Now the users are unable to do anything, and so we can start building the set of new privileges based on a complete understanding of what we want each user to be able to do.

# 10.3.4. Sub-graph access control using privileges

With privileges, we can take much more control over what each user is capable of doing. We start by identifying each type of user:

#### Doctor

Should be able to read and write most of the graph. We would, however, like to prevent the doctor from reading the patient's address.

#### Receptionist

Should be able to read and write all patient data, but not be able to see the symptoms, diseases or diagnoses.

#### Researcher

Should be able to perform statistical analysis on all data, except patients' personal information, and as such should not be able to read most patient properties.

#### Nurse

The nurse should be able to perform all tasks that both the doctor and the receptionist can do. For this reason, we do not need to create a dedicated role, but can assign nurses to both doctor and receptionist roles.

#### IT administrator

This role is very similar to the built-in admin role, except that we wish to restrict access to the patients SSN. To achieve this, we can create this role by copying the built-in admin role and modifying the privileges of the copy.

#### Example 46. Creating custom roles

```
CREATE ROLE doctor;
CREATE ROLE receptionist;
CREATE ROLE researcher;
CREATE ROLE itadmin AS COPY OF admin;
```

Before we assign the new roles to Alice, Bob, Daniel, Charlie and Tina, we should define the privileges of each role.

## Privileges of itadmin

This role was created as a copy of the built-in admin role, and so all we need to do is restrict access to the patient's SSN:

```
DENY READ {ssn} ON GRAPH healthcare NODES Patient TO itadmin;
```

The complete set of privileges available to users assigned the itadmin role can be viewed using the following command:

```
SHOW ROLE itadmin PRIVILEGES
| access
          | action
                         | resource
                                            | graph
                                                            | segment
                                                                                | role
                      | "all_properties" |
                                                              "NODE(*)"
                                                                                  "itadmin"
  "GRANTED" | "read"
            | "write"
                                                              "NODE(*)"
  "GRANTED"
                                              "*"
                         | "all_properties"
                                                                                  "itadmin"
  "GRANTED"
                                                              "NODE(*)"
             "traverse" |
                           "graph"
                                                                                  "itadmin"
                                              "*"
             "read"
                           "all_properties" |
                                                              "RELATIONSHIP(*)"
                                                                                  "itadmin"
  "GRANTED"
                                                              "RELATIONSHIP(*)"
  "GRANTED"
             "write"
                           "all_properties"
                                              "*"
                                                                                  "itadmin"
  "GRANTED"
                                              "*"
                                                             "RELATIONSHIP(*)"
              "traverse" |
                           "graph"
                                                                                  "itadmin"
  "GRANTED"
             "access"
                                              "*"
                                                                                  "itadmin"
                           "database"
                                                             "database"
                                              "*"
  "GRANTED"
              "admin"
                           "database"
                                                              "database"
                                                                                  "itadmin"
  "GRANTED"
              "schema"
                           "database"
                                                              "database"
                                                                                  "itadmin"
  "GRANTED" |
                           "database"
                                              11 🛧 11
                                                              "database"
             "token"
                                                                                  "itadmin"
  "DENIED" | "read"
                         | "property(ssn)" | "healthcare" | "NODE(Patient)"
                                                                                | "itadmin"
```

In order for the IT-Admin tina to be provided these privileges, she must be assigned the new role itadmin.

```
GRANT ROLE itadmin TO tina;
```

To demonstrate that Tina is not able to see the patients SSN, we can login to healthcare as tina and run the query:

```
MATCH (n:Patient)
WHERE n.dateOfBirth < date('1972-06-12')
RETURN n.name, n.ssn, n.address, n.dateOfBirth;
```

```
"n.name" "n.ssn" "n.address" "n.dateOfBirth"

"Mark Smith" null "1 secret way, downtown" "1970-01-21"

"Mary Smith" null "1 secret way, downtown" "1970-12-30"

"Sally Stone" null "1 secret way, downtown" "1971-06-17"

"Jane Anderson" null "1 secret way, downtown" "1971-10-22"
```

# Privileges of researcher

Charlie the researcher was previously our only read-only user. We could do something similar to what we did with the <code>itadmin</code> role, by copying and modifying the <code>reader</code> role. However, we would like to explicitly illustrate how to build a role from scratch. There are various possibilities for building this role using the concepts of <code>whitelisting</code> and <code>blacklisting</code>:

#### · Blacklisting:

We could grant the role the ability to find all nodes and read all properties (much like the reader role) and then deny read access to the Patient properties we want to restrict the researcher from seeing, such as name, SSN and address. This approach is simple but suffers from one problem. If Patient nodes are assigned additional properties, *after* we have restricted access, these new properties will automatically be visible to the researcher, which may not be desirable.

#### Example 47. Blacklisting

```
GRANT ACCESS ON DATABASE healthcare TO researcherB;

// First grant access to everything
GRANT MATCH {*}

ON GRAPH healthcare

TO researcherB;

// Then deny read on specific node properties
DENY READ {name, address, ssn}

ON GRAPH healthcare

NODES Patient

TO researcherB;

// And deny traversal of the doctors diagnosis
DENY TRAVERSE

ON GRAPH healthcare

RELATIONSHIPS DIAGNOSIS
TO researcherB;
```

#### · Whitelisting:

An alternative is to only provide specific access to the properties we wish the researcher to see. Then, the addition of new properties will not automatically make them visible to the researcher. In this case, adding new properties to a Patient will not mean that the researcher can see them by default. If we wish to have them visible, we need to explicitly grant read access.

#### Example 48. Whitelisting

```
GRANT ACCESS ON DATABASE healthcare TO researcherW;
// We allow the researcher to find all nodes
GRANT TRAVERSE
   ON GRAPH healthcare
    NODES *
    TO researcherW:
// Now only allow the researcher to traverse specific relationships
GRANT TRAVERSE
    ON GRAPH healthcare
    RELATIONSHIPS HAS, OF
    TO researcherW;
// Allow reading of all properties of medical metadata
GRANT READ {*}
    ON GRAPH healthcare
    NODES Symptom, Disease
    TO researcherW;
// Only allow reading dateOfBirth for research purposes
GRANT READ {dateofbirth}
   ON GRAPH healthcare
    NODES Patient
    TO researcherW;
```

In order to test that Charlie now has the privileges we have specified, we assign him to the researcherB role with blacklisting:

```
GRANT ROLE researcherB TO charlie;
```

We can use a version of the SHOW PRIVILEGES command to see Charlies access rights:

```
neo4j@system> SHOW USER charlie PRIVILEGES;
| access | action | resource
                                                       | graph
                                                                         | segment
                                                                                                             | role
| "GRANTED" | "read"
                              | "all_properties"
                                                      | "healthcare" | "NODE(*)"
"researcherB" | "charlie"
| "GRANTED" | "traverse" | "graph"
                                                         | "healthcare" | "NODE(*)"
"researcherB" | "charlie" |
| "DENIED" | "read" |
"researcherB" | "charlie" |
                                "property(address)" | "healthcare" | "NODE(Patient)"
| "DENIED" | "read" | "researcherB" | "charlie" |
                             | "property(name)"
                                                       | "healthcare" | "NODE(Patient)"
| "DENIED" | "read"
                            | "property(ssn)"
                                                       | "healthcare" | "NODE(Patient)"
"researcherB" | "charlie" |

"GRANTED" | "read" | "all_properties" | "healthcare" | "RELATIONSHIP(*)"

"researcherB" | "charlie" |
| "GRANTED" | "traverse" | "graph"
"researcherB" | "charlie" |
                                                         | "healthcare" | "RELATIONSHIP(*)"
| "DENIED" | "traverse" | "graph"
                                                         | "healthcare" | "RELATIONSHIP(DIAGNOSIS)" |
"researcherB" | "charlie" |
| "GRANTED" | "access" | "database"
"researcherB" | "charlie" |
                                                         | "healthcare" | "database"
```

Now when Charlie logs into the healthcare database and tries to run a command similar to the one used by the itadmin above, we will see different results:

```
MATCH (n:Patient)
WHERE n.dateOfBirth < date('1972-06-12')
RETURN n.name, n.ssn, n.address, n.dateOfBirth;
```

```
"n.name" "n.ssn" "n.address" "n.dateOfBirth"
null
         null
                 null
                              "1970-01-21"
                 null
                              "1970-12-30"
nul1
         nul1
null
         null
                 null
                              "1971-06-17"
null
         null
                 null
                              "1971-10-22"
```

Only the date of birth is available, so Charlie the researcher may perform statistical analysis, for example. Another query Charlie could try is to find the ten diseases a patient younger than 25 is most likely to be diagnosed with, listed by probability:

```
WITH datetime() - duration({years:25}) AS timeLimit
MATCH (n:Patient)
WHERE n.dateOfBirth > date(timeLimit)
MATCH (n)-[h:HAS]->(s:Symptom)-[o:OF]->(d:Disease)
WITH d.name AS disease, o.probability AS prob
RETURN disease, sum(prob) AS score ORDER BY score DESC LIMIT 10;
```

```
"disease"
                      "score"
"Chronic Whatitis"
                      111.30050876448621
"Chronic Someitis"
                      110.56964390091147
"Acute Yellowitis"
                      98.82266316401365
"Chronic Otheritis"
                      80.41346486864003
"Acute Otheritis"
                      79.82679831362869
"Acute Placeboitis"
                      78.86090865510758
"Chronic Yellowitis" 77.53519713418886
"Chronic Argitis"
                      70.04150610048167
"Acute Someitis"
                      69.45011166554933
"Chronic Placeboitis" 64.36353437805441
```

## Privileges of doctor

Doctors should be given the ability to read and write almost everything. We would, however, like to remove the ability to read the patients' address property. This role can be built from scratch by assigning full read and write access, and then specifically denying access to the address property:

```
GRANT ACCESS ON DATABASE healthcare TO doctor;
GRANT TRAVERSE ON GRAPH healthcare TO doctor;
GRANT READ {*} ON GRAPH healthcare TO doctor;
GRANT WRITE ON GRAPH healthcare TO doctor;
DENY READ {address} ON GRAPH healthcare NODES Patient TO doctor;
```

To allow Alice to have these privileges, we grant her this new role:

```
GRANT ROLE doctor TO alice;
```

To demonstrate that Alice is not able to see patient addresses, we can run the query:

```
MATCH (n:Patient)
WHERE n.dateOfBirth < date('1972-06-12')
RETURN n.name, n.ssn, n.address, n.dateOfBirth;
```

```
"n.name" "n.ssn" "n.address" "n.dateOfBirth"

"Mark Smith" 1234610 null "1970-01-21"

"Mary Smith" 1234640 null "1970-12-30"

"Sally Stone" 1234641 null "1971-06-17"

"Jane Anderson" 1234652 null "1971-10-22"
```

As we can see, the doctor has the expected privileges, including being able to see the SSN, but not the address of each patient.

## Privileges of receptionist

Receptionists should only be able to manage patient information. They are not allowed to find or read any other parts of the graph:

```
GRANT ACCESS ON DATABASE healthcare TO receptionist;
GRANT MATCH {*} ON GRAPH healthcare NODES Patient TO receptionist;
GRANT WRITE {*} ON GRAPH healthcare TO receptionist;
```



It is currently not possible to be specific on WRITE access, and therefore a user that is granted write access is able to write to all nodes and relationships. For example, the receptionist could create a new Symptom node, even if they are then not able to find that in the database due to the restricted read access.

```
GRANT ROLE receptionist TO bob;
```

With these privileges, if Bob tries to read the entire database, he will still only see the patients:

```
MATCH (n) WITH labels(n) AS labels
RETURN labels, count(*);

"labels" "count(*)"

["Patient"] 101
```

However, Bob is able to see all fields of the Patient records:

```
MATCH (n:Patient)
WHERE n.dateOfBirth < date('1972-06-12')
RETURN n.name, n.ssn, n.address, n.dateOfBirth;
```

```
"n.name" "n.ssn" "n.address" "n.dateOfBirth"

"Mark Smith" 1234610 "1 secret way, downtown" "1970-01-21"

"Mary Smith" 1234640 "1 secret way, downtown" "1970-12-30"

"Sally Stone" 1234641 "1 secret way, downtown" "1971-06-17"

"Jane Anderson" 1234652 "1 secret way, downtown" "1971-10-22"
```

Let say that Bob the receptionist wants to remove a patient from the database:

```
MATCH (n:Patient)
WHERE n.SSN = 1234610
DETACH DELETE n;
```

```
org.neo4j.graphdb.ConstraintViolationException: Cannot delete node<42>, because it still has relationships. To delete this node, you must first delete its relationships.
```

The reason this fails is that Bob can find the (:Patient) node, but does not have sufficient traverse rights to find the outgoing relationships from it. Either he needs to ask Tina the itadmin for help for this task, or we can add more privileges to the receptionist role:

```
GRANT TRAVERSE ON GRAPH healhcare NODES Symptom, Disease TO receptionist;
GRANT TRAVERSE ON GRAPH healthcare RELATIONSHIPS HAS, DIAGNOSIS TO receptionist;
```

# Privileges of nurses

The nurse is not defined as a separate role because, as we have established, nurses have the capabilities of both doctors *and* receptionists. Therefore we can assign both those roles to Daniel the nurse and achieve desired behaviour for a nurse.

```
GRANT ROLE doctor, receptionist TO daniel;
```

Now we can see that the user 'Daniel' has a combined set of privileges:

```
neo4j@system> SHOW USER daniel PRIVILEGES;
access
         | action
                      resource
                                            | graph
                                                           | segment
user |
| "GRANTED" | "read"
                        | "all_properties"
                                           | "healthcare" | "NODE(*)"
                                                                               | "doctor"
"daniel" |
                       | "all_properties" | "healthcare" | "NODE(*)"
| "GRANTED" | "write"
                                                                              | "doctor"
"daniel"
 "GRANTED" | "traverse" | "graph"
                                             | "healthcare" | "NODE(*)"
                                                                               | "doctor"
"daniel" |
| "DENIED"
                        | "property(address)" | "healthcare" | "NODE(Patient)" | "doctor"
          | "read"
"daniel" |
| "GRANTED" | "read"
                        | "all_properties"
                                            | "healthcare" | "RELATIONSHIP(*)" | "doctor"
"daniel"
                        | "all_properties" | "healthcare" | "RELATIONSHIP(*)" | "doctor"
| "GRANTED" | "write"
"daniel" |
 "GRANTED" | "traverse" | "graph"
                                             | "healthcare" | "RELATIONSHIP(*)" | "doctor"
"daniel" |
| "GRANTED" | "access"
                                             | "healthcare" | "database"
                        | "database"
                                                                              | "doctor"
"daniel" |
                        | "all_properties"
                                            | "healthcare" | "NODE(*)" | "receptionist" |
| "GRANTED" | "write"
"daniel"
 "GRANTED" | "read"
                        | "all_properties" | "healthcare" | "NODE(Patient)" | "receptionist" |
"daniel" |
| "GRANTED" | "traverse" | "graph"
                                             | "healthcare" | "NODE(Patient)" | "receptionist" |
"daniel" |
| "GRANTED" | "write"
                        | "all_properties"
                                             | "healthcare" | "RELATIONSHIP(*)" | "receptionist" |
"daniel" |
| "GRANTED" | "access"
                                             | "healthcare" | "database"
                                                                               | "receptionist" |
                      | "database"
"daniel" |
```

# 10.4. Integration with LDAP

This section describes Neo4j support for integrating with LDAP systems.

#### This section describes the following:

- Introduction
- Configure the LDAP auth provider
  - □ Configuration for Active Directory
  - □ Configuration for openLDAP
- Use 'Idapsearch' to verify the configuration
- The auth cache
- Available methods of encryption
  - ☐ Use LDAP with encryption via StartTLS
  - ☐ Use LDAP with encrypted LDAPS
- Use a self-signed certificate in a test environment

## 10.4.1. Introduction

Neo4j supports LDAP which allows for integration with Active Directory, OpenLDAP or other LDAP-compatible authentication services. We will show example configurations where management of federated users is deferred to the LDAP service, using that service's facilities for administration. This means that we completely turn off native Neo4j user and role administration and map LDAP groups to the Neo4j roles.

# 10.4.2. Configure the LDAP auth provider

All settings need to be defined at server startup time in the default configuration file neo4j.conf.

First, configure Neo4j to use LDAP as authentication and authorization provider:

```
# Turn on security:
dbms.security.auth_enabled=true

# Choose LDAP connector as security provider for both authentication and authorization:
dbms.security.authentication_providers=ldap
dbms.security.authorization_providers=ldap
```

## Configuration for Active Directory

See below for an example configuration for Active Directory:

```
# Configure LDAP to point to the AD server:
dbms.security.ldap.host=ldap://myactivedirectory.example.com
# Provide details on user structure within the LDAP system:
dbms.security.ldap.authentication.user_dn_template=cn={0},cn=Users,dc=example,dc=com
dbms.security.ldap.authorization.user_search_base=cn=Users,dc=example,dc=com
dbms.security.ldap.authorization.user_search_filter=(&(objectClass=*)(cn={0}))
dbms.security.ldap.authorization.group_membership_attributes=memberOf
# Configure the actual mapping between groups in the LDAP system and roles in Neo4j:
dbms.security.ldap.authorization.group_to_role_mapping=\
  "cn=Neo4j Read Only,cn=Users,dc=neo4j,dc=com"
                                                     = reader
  "cn=Neo4j Read-Write,cn=Users,dc=neo4j,dc=com"
                                                     = publisher
  "cn=Neo4j Schema Manager,cn=Users,dc=neo4j,dc=com" = architect
  "cn=Neo4j Administrator,cn=Users,dc=neo4j,dc=com"
                                                    = admin
  "cn=Neo4j Procedures,cn=Users,dc=neo4j,dc=com"
                                                     = allowed_role
# In case defined users are not allowed to search for themselves, we can specify credentials for a user
with read access to all users and groups.
# Note that this account only needs read-only access to the relevant parts of the LDAP directory and does
not need to have access rights to Neo4j or any other systems.
# dbms.security.ldap.authorization.use_system_account=true
# dbms.security.ldap.authorization.system_username=cn=search-account,cn=Users,dc=example,dc=com
# dbms.security.ldap.authorization.system_password=secret
```

Below is an alternative configuration for Active Directory that allows for logging in with sAMAccountName:

```
# Configure LDAP to point to the AD server:
dbms.security.ldap.host=ldap://myactivedirectory.example.com
# Provide details on user structure within the LDAP system:
dbms.security.ldap.authorization.user_search_base=cn=Users,dc=example,dc=com
dbms.security.ldap.authorization.user_search_filter=(&(objectClass=*)(samaccountname={0}))
dbms.security.ldap.authorization.group_membership_attributes=memberOf
# Configure the actual mapping between groups in the LDAP system and roles in Neo4j:
dbms.security.ldap.authorization.group_to_role_mapping=\
  "cn=Neo4j Read Only,cn=Users,dc=neo4j,dc=com"
                                                      = reader
  "cn=Neo4j Read-Write,cn=Users,dc=neo4j,dc=com"
                                                      = publisher
  "cn=Neo4j Schema Manager,cn=Users,dc=neo4j,dc=com" = admin
"cn=Neo4j Administrator,cn=Users,dc=neo4j,dc=com" = allowed_role
  "cn=Neo4j Procedures,cn=Users,dc=neo4j,dc=com"
# In case defined users are not allowed to search for themselves, we can specify credentials for a user
with read access to all users and groups.
# Note that this account only needs read-only access to the relevant parts of the LDAP directory and does
not need to have access rights to Neo4j or any other systems.
dbms.security.ldap.authorization.use_system_account=true
dbms.security.ldap.authorization.system_username=cn=search-account,cn=Users,dc=example,dc=com
dbms.security.ldap.authorization.system_password=secret
# Perform authentication with sAMAccountName instead of DN.
# Using this setting requires dbms.security.ldap.authorization.system_username and
dbms.security.ldap.authorization.system_password to be used, since there is no way to log in through LDAP
directly with the sAMAccountName.
# Instead, the login name will be resolved to a DN that will be used to log in with.
dbms.security.ldap.authentication.use_samaccountname=true
```

Below is an alternative configuration for Active Directory that allows for authenticating users from different organizational units by using the Active Directory attribute samacountName:

```
# Configure LDAP to point to the AD server:
dbms.security.ldap.host=ldap://myactivedirectory.example.com

dbms.security.ldap.authentication.user_dn_template={0}@example.com
dbms.security.ldap.authorization.user_search_base=dc=example,dc=com
dbms.security.ldap.authorization.user_search_filter=(&(objectClass=user)(sAMAccountName={0}))
dbms.security.ldap.authorization.group_membership_attributes=memberOf

# Configure the actual mapping between groups in the LDAP system and roles in Neo4j:
dbms.security.ldap.authorization.group_to_role_mapping=\
"cn=Neo4j Read Only,cn=Users,dc=example,dc=com" = reader ;\
"cn=Neo4j Read-Write,cn=Users,dc=example,dc=com" = architect ;\
"cn=Neo4j Schema Manager,cn=Users,dc=example,dc=com" = architect ;\
"cn=Neo4j Administrator,cn=Users,dc=example,dc=com" = admin ;\
"cn=Neo4j Procedures,cn=Users,dc=example,dc=com" = allowed_role
```

Specifying the <code>{0}@example.com</code> pattern in the user\_dn\_template enables the authentication to start at the root domain. The whole tree is checked to find the user, regardless of where it is located within the tree.

Note that the setting dbms.security.ldap.authentication.use\_samaccountname is not configured in this example.

# Configuration for openLDAP

See below for an example configuration for openLDAP:

```
# Configure LDAP to point to the OpenLDAP server:
dbms.security.ldap.host=myopenldap.example.com
# Provide details on user structure within the LDAP system:
dbms.security.ldap.authentication.user_dn_template=cn={0},ou=users,dc=example,dc=com
dbms.security.ldap.authorization.user_search_base=ou=users,dc=example,dc=com
dbms.security.ldap.authorization.user_search_filter=(&(objectClass=*)(uid={0}))
{\tt dbms.security.ldap.authorization.group\_membership\_attributes=} {\tt gidnumber}
# Configure the actual mapping between groups in the OpenLDAP system and roles in Neo4j:
dbms.security.ldap.authorization.group_to_role_mapping=\
  101 = reader
  102 = publisher
  103 = architect
  104 = admin
  105 = allowed_role
# In case defined users are not allowed to search for themselves, we can specify credentials for a user
with read access to all users and groups.
# Note that this account only needs read-only access to the relevant parts of the LDAP directory and does
not need to have access rights to Neo4j or any other systems.
# dbms.security.ldap.authorization.use_system_account=true
# dbms.security.ldap.authorization.system_username=cn=search-account,ou=users,dc=example,dc=com
# dbms.security.ldap.authorization.system_password=search-account-password
```

We would like to draw attention to some details in the configuration examples. A comprehensive overview of LDAP configuration options is available in Configuration settings.

Parameter name	Default value	Description
dbms.security.ldap.authentication.use r_dn_template	<pre>uid={0},ou=users,dc=example,dc=com</pre>	Converts usernames into LDAP- specific fully qualified names required for logging in.
dbms.security.ldap.authorization.user _search_base	ou=users,dc=example,dc=com	Sets the base object or named context to search for user objects.
dbms.security.ldap.authorization.user _search_filter	(&(objectClass=*)(uid={0}))	Sets up an LDAP search filter to search for a user principal.
dbms.security.ldap.authorization.group_membership_attributes	[memberOf]	Lists attribute names on a user object that contains groups to be used for mapping to roles.
dbms.security.ldap.authorization.group_to_role_mapping		Lists an authorization mapping from groups to the pre-defined built-in roles admin, architect, publisher and reader, or to any other custom-defined roles.

## 10.4.3. Use **Idapsearch** to verify the configuration

We can use the LDAP command-line tool <a href="Ldapsearch">1</a>dapsearch to verify that the configuration is correct, and that the LDAP server is actually responding. We do this by issuing a search command that includes LDAP configuration setting values.

These example searches verify both the authentication (using the simple mechanism) and authorization of user 'john'. See the ldapsearch documentation for more advanced usage and how to use SASL authentication mechanisms.

With dbms.security.ldap.authorization.use\_system\_account=false (default):

With dbms.security.ldap.authorization.use\_system\_account=true:

```
#ldapsearch -v -H ldap://<dbms.security.ldap.host> -x -D 
<dbms.security.ldap.authorization.system_username> -w <dbms.security.ldap.authorization.system_password> -b <dbms.security.ldap.authorization.user_search_base> 
"<dbms.security.ldap.authorization.user_search_filter>" 
<dbms.security.ldap.authorization.group_membership_attributes> 
ldapsearch -v -H ldap://myactivedirectory.example.com:389 -x -D cn=search- 
account,cn=Users,dc=example,dc=com -w secret -b cn=Users,dc=example,dc=com "(&(objectClass=*)(cn=john))" 
memberOf
```

Then verify that we get a successful response, and that the value of the returned membership attribute is a group that is mapped to a role in

dbms.security.ldap.authorization.group\_to\_role\_mapping.

```
# extended LDIF
#
# LDAPv3
# base <cn=Users,dc=example,dc=com> with scope subtree
# filter: (cn=john)
# requesting: memberOf
#

# john, Users, example.com
dn: CN=john,CN=Users,DC=example,DC=com
memberOf: CN=Neo4j Read Only,CN=Users,DC=example,DC=com
# search result
search: 2
result: 0 Success
# numResponses: 2
# numEntries: 1
```

## 10.4.4. The auth cache

The *auth cache* is the mechanism by which Neo4j caches the result of authentication via the LDAP server in order to aid performance. It is configured with the

dbms.security.ldap.authentication.cache\_enabled and dbms.security.auth\_cache\_ttl parameters.

```
# Turn on authentication caching to ensure performance
dbms.security.ldap.authentication.cache_enabled=true
dbms.security.auth_cache_ttl=10m
```

Parameter name	Default value	Description
dbms.security.ldap.authentication.cac he_enabled	true	Determines whether or not to cache the result of authentication via the LDAP server. Whether authentication caching should be enabled or not must be considered in view of your company's security guidelines.
dbms.security.auth_cache_ttl	600 seconds	Is the time to live (TTL) for cached authentication and authorization info. Setting the TTL to 0 will disable all auth caching. A short TTL will require more frequent re-authentication and reauthorization, which can impact performance. A very long TTL will also mean that changes to the users settings on an LDAP server may not be reflected in the Neo4j authorization behaviour in a timely manner. Valid units are ms, s, m; default unit is s.

An administrator can clear the auth cache to force the re-querying of authentication and authorization

information from the federated auth provider system.

#### Example 49. Clear the auth cache

Use Neo4j Browser or Neo4j Cypher Shell to execute this statement.

```
CALL dbms.security.clearAuthCache()
```

# 10.4.5. Available methods of encryption

All the following ways of specifying the dbms.security.ldap.host parameter are valid. Doing so will configure using LDAP without encryption. Not specifying the protocol or port will result in ldap being used over the default port 389.

```
dbms.security.ldap.host=myactivedirectory.example.com
dbms.security.ldap.host=myactivedirectory.example.com:389
dbms.security.ldap.host=ldap://myactivedirectory.example.com
dbms.security.ldap.host=ldap://myactivedirectory.example.com:389
```

# Use LDAP with encryption via StartTLS

To configure Active Directory with encryption via StartTLS, set the following parameters:

```
dbms.security.ldap.use_starttls=true
dbms.security.ldap.host=ldap://myactivedirectory.example.com
```

## Use LDAP with encrypted LDAPS

To configure Active Directory with encrypted LDAPS, set dbms.security.ldap.host to one of the following. Not specifying the port will result in ldaps being used over the default port 636.

```
dbms.security.ldap.host=ldaps://myactivedirectory.example.com
dbms.security.ldap.host=ldaps://myactivedirectory.example.com:636
```

This method of securing Active Directory is being deprecated and is therefore not recommended. Instead, use Active Directory with encryption via StartTLS.

# 10.4.6. Use a self-signed certificate in a test environment

Production environments should always use an SSL certificate issued by a Certificate Authority for secure access to the LDAP server. However, there are scenarios, for example in test environments, where you may wish to use a self-signed certificate on the LDAP server. In these scenarios you will have to tell Neo4j about the local certificate. This is done by entering the details of the certificate using dbms.jvm.additional in neo4j.conf.

This example shows how to specify details for a self-signed certificate on an LDAP server. The path to the certificate file MyCert. jks is an absolute path on the Neo4j server.

```
dbms.jvm.additional=-Djavax.net.ssl.keyStore=/path/to/MyCert.jks
dbms.jvm.additional=-Djavax.net.ssl.keyStorePassword=secret
dbms.jvm.additional=-Djavax.net.ssl.trustStore=/path/to/MyCert.jks
dbms.jvm.additional=-Djavax.net.ssl.trustStorePassword=secret
```

# 10.5. Manage procedure permissions

This section describes how to configure subgraph access control in Neo4j.

This section describes the following:

- Introduction
- Configuration steps
  - ☐ Create a custom role
  - Manage procedure permissions

## 10.5.1. Introduction

It is possible to restrict a user's access to, and subsequent actions on, specified portions of the graph. For example, a user can be allowed to read, but not write, nodes with specific labels and relationships with certain types.

To implement subgraph access control, you must complete the following steps:

- 1. Put a procedure or function in place that performs the reads from, and/or writes to, the portion of the graph that you wish to control. This can either be developed in-house, or be made available as a third-party library. Please refer to Java Reference 

  Extending Neo4j for a description on creating and using user-defined procedures and functions.
- 2. Create one, or several, custom roles, with which to run the procedure described above. These roles can subsequently be assigned the relevant privileges.
- 3. Configure the procedure so that it can be executed by users with the custom roles.

The steps below assume that the procedure or function is already developed and installed.

## 10.5.2. Configuration steps

#### Create a custom role

Create a custom role and manage it either through native user management or through federated user management with LDAP.

#### Native users scenario

In the native users scenario, a custom role is created and assigned to the relevant user(s).

#### Example 51. Native users scenario

In this example, we will use Cypher to create a custom accounting role and assign it to a preexisting user, billsmith.

```
CALL dbms.security.createRole('accounting')

CALL dbms.security.addRoleToUser('accounting', 'billsmith')
```

#### Federated users scenario (LDAP)

In the LDAP scenario, the LDAP user group is mapped to a custom role in Neo4j.

Example 52. Federated users scenario (LDAP)

```
In this example, we will use Cypher to create a custom accounting role.

CALL dbms.security.createRole('accounting')

We will then map the accounting role to the LDAP group with groupID 101.

dbms.security.realms.ldap.authorization.group_to_role_mapping=101=accounting
```

## Manage procedure permissions

In standard use, procedures and functions are executed according to the same security rules as regular Cypher statements. For example, users assigned any one of the native roles publisher, architect and admin will be able to execute a procedure with mode=WRITE, whereas a user assigned only the reader role will not be allowed to execute the procedure.

For the purpose of subgraph access control, we allow specific roles to execute procedures that they would otherwise be prevented from accessing through their assigned native roles. The user is given the privilege that comes with the mode of the procedure, during the execution of the procedure only. The following two parameters are used to configure the desired behavior:

#### dbms.security.procedures.default\_allowed

The setting dbms.security.procedures.default\_allowed defines a single role that is allowed to execute any procedure or function that is not matched by the dbms.security.procedures.roles configuration.

Assume that we have the following configuration:

dbms.security.procedures.default\_allowed=superAdmin

This will have the following effects:

- If the setting dbms.security.procedures.roles is left unconfigured, the role superAdmin will be able to execute all custom procedures and functions.
- If the setting dbms.security.procedures.roles has some roles and functions defined, the role superAdmin will be able to execute all custom procedures and functions that are *not* configured by dbms.security.procedures.roles.

#### dbms.security.procedures.roles

The dbms.security.procedures.roles setting provides fine-grained control over procedures.

Example 54. Configure roles for the execution of specific procedures

Assume that we have the following configuration:

 $\label{lower} {\bf dbms.security.procedures.roles} = {\tt apoc.convert.*:Converter;apoc.load.json.*:Converter,DataSource;apoc.trigger.add:TriggerHappy}$ 

This will have the following effects:

- All users with the role Converter will be able to execute all procedures in the apoc.convert namespace.
- All users with the roles Converter and DataSource will be able to execute procedures in the apoc.load.json namespace.
- All users with the role TriggerHappy will be able to execute the specific procedure apoc.trigger.add.



A procedure will fail if it attempts to execute database operations that violates its mode. For example, a procedure assigned the mode of READ will fail if it is programmed to do write actions. This will happen regardless of user or role configuration.

# 10.6. Terminology

This section lists the relevant terminology related to authentication and authorization in Neo4i.

The following terms are relevant to role-based access control within Neo4j:

active user

A user who is active within the system and can perform actions prescribed by any assigned roles on the data. This is in contrast to a suspended user.

#### administrator

This is a user who has been assigned the admin role.

#### current user

This is the currently logged-in user invoking the commands described in this chapter.

#### password policy

The password policy is a set of rules of what makes up a valid password. For Neo4j, the following rules apply:

- The password cannot be the empty string.
- When changing passwords, the new password cannot be the same as the previous password.

#### role

This is a collection of actions — such as read and write — permitted on the data.

#### suspended user

A user who has been suspended is not able to access the database in any capacity, regardless of any assigned roles.

#### user

- A user is composed of a username and credentials, where the latter is a unit of information, such as a password, verifying the identity of a user.
- A user may represent a human, an application etc.

# Chapter 11. Security

This chapter covers important security aspects in Neo4j.

Ensure your physical data security by following industry best practices with regard to server and network security.

This chapter includes the following:

- Securing extensions
- SSL framework
- Credentials handling in Neo4j Browser
- Security checklist

Additionally, logs can be useful for continuous analysis, or for specific investigations. Facilities are available for producing security event logs as well as query logs as described in Monitoring.



Refer to Authentication and authorization for information on how to manage users and their authentication and authorization.

# 11.1. Securing extensions

This section describes how to ensure the security of custom-written additions in Neo4j.

Neo4j can be extended by writing custom code which can be invoked directly from Cypher, as described in Java Reference 
Procedures and functions. This section describes how to ensure the security of these additions.

# 11.1.1. Sandboxing

Neo4j provides sandboxing to ensure that procedures do not inadvertently use insecure APIs. For example, when writing custom code it is possible to access Neo4j APIs that are not publicly supported, and these internal APIs are subject to change, without notice. Additionally, their use comes with the risk of performing insecure actions. The sandboxing functionality limits the use of extensions to publicly supported APIs, which exclusively contain safe operations, or contain security checks.

The configuration setting dbms.security.procedures.unrestricted provides the possibility to circumvent the sandboxing functionality by defining a comma-separated list of procedures and/or user-defined functions that are allowed full access to the database. The list may contain fully-qualified procedure names, and partial names with the wildcard \*.

Any attempt to load an extension which contains an unsupported API, which has not been marked as allowed in <a href="mailto:documents-unrestricted">documents-unrestricted</a>, will result in a warning in the security log. The warning will point out that the extension does not have access to the components it is trying to load. Additionally, a mocked procedure will be loaded with the procedure's name. Calling the mocked procedure will result in an error, saying that the procedure failed to load due to needing more permissions.

#### Example 55. Sandboxing

In this example we assume that the procedure my.extensions.example, as well as some procedures in the my.procedures library, make use of unsupported APIs. In order to allow the running of these procedures we configure the setting as shown below.

```
# Example sandboxing dbms.security.procedures.unrestricted=my.extensions.example,my.procedures.*
```

# 11.1.2. White listing

White listing can be used to allow loading only a few extensions from a larger library.

The configuration setting dbms.security.procedures.whitelist is used to name certain procedures that should be available from a library. It defines a comma-separated list of procedures that are to be loaded. The list may contain both fully-qualified procedure names, and partial names with the wildcard \*.

#### Example 56. White listing

In this example we wish to allow the use of the method apoc.load.json as well as all the methods under apoc.coll. We do not want to make available any additional extensions from the apoc library, other than the ones matching these criteria.

```
# Example white listing
dbms.security.procedures.whitelist=apoc.coll.*,apoc.load.*
```

There are a few things that should be noted about <a href="mailto:dbms.security.procedures.whitelist">dbms.security.procedures.whitelist</a>:

- If using this setting, no extensions other than those listed will be loaded. In particular, if it is set to the empty string, no extensions will be loaded.
- The default of the setting is \*. This means that if you do not explicitly give it a value (or no value), all libraries in the *plugins* directory will be loaded.
- If the extensions pointed out by this parameter are programmed to access internal APIs, they also have to be explicitly allowed, as described in Sandboxing.

# 11.2 SSI framework

This section describes SSL/TLS integration for securing communication channels in Neo4j.

Neo4j supports the securing of communication channels using standard SSL/TLS technology.

This section describes the following:

- Introduction
- Certificates
- Configuration
  - Settings
- Choosing an SSL provider

Terminology

# 11.2.1. Introduction

The SSL support is enabled per communication channel and requires SSL certificates encoded in the PEM format. The process is described in the following sections.

## 11.2.2. Certificates

The instructions in this section assume that you have already acquired the required certificates.

All certificates must be in the PEM format, and they can be combined into one file. The private key is also required to be in the PEM format. Multi-host and wildcard certificates are supported.

# 11.2.3. Configuration

The SSL policies are configured by assigning values to parameters of the following format:

```
dbms.ssl.policy.<scope>.<setting-suffix>
```

The scope is the name of the communication channel, and must be one of bolt, https, cluster, backup or fabric.

Each policy needs to be explicitly enabled by setting:

dbms.ssl.policy.<scope>.enabled=true

## Settings

The valid values for setting-suffix are described below.

Setting suffix	Description	Default value		
	Basic			
enabled	Setting this to true will enable this policy.	false		
base_directory	The base directory under which cryptographic objects are searched for by default.	certificates/ <scope></scope>		
private_key	The private key used for authenticating and securing this instance.	private.key		
private_key_password	The passphrase to decode the private key. Only applicable for encrypted private keys.			
<pre>public_certificate</pre>	A public certificate matching the private key signed by a Certificate Authority (CA).	public.crt		
trusted_dir	A directory populated with certificates of trusted parties.	trusted/		
revoked_dir	A directory populated with certificate revocation lists (CRLs).	revoked/		
Advanced				

Setting suffix	Description	Default value
verify_hostname	Enabling this setting will turn on client- side hostname verification. After the client has received the servers public certificate, it will compare the address it used against the certificate Common Name (CN) and Subject Alternative Names (SAN) fields. If the address used doesn't match those fields, the client will disconnect.	false
ciphers	A comma-separated list of ciphers suits that will be allowed during cipher negotiation. Valid values depend on the current JRE and SSL provider, see note below for examples.	Java platform default allowed cipher suites
tls_versions	A comma-separated list of allowed TLS versions.	TLSv1.2
client_auth	Whether or not clients must be authenticated. Setting this to REQUIRE effectively enables mutual authentication for servers. Available values given to this setting are NONE, OPTIONAL, or REQUIRE.	OPTIONAL for bolt and https and REQUIRE for cluster and backup.
trust_all	Setting this to true will result in all clients and servers being trusted. The content of the trusted_dir directory will be ignored. Use of this is discouraged, since it will not offer security. It is provided as a mean of debugging.	false



Ciphers supported by the Oracle JRE can be found here: https://docs.oracle.com/en/java/javase/11/docs/specs/security/standard-names.html#jsse-cipher-suite-names

For security reasons, Neo4j will not attempt to automatically create any of these directories. The creation of an SSL policy therefore requires the appropriate file system structure to be set up manually. Note that the existence of the directories is mandatory, as well as the presence of the certificate file and the private key. Ensure correct permissions are set on the private key, such that only the Neo4j user can read it.

In this example we will configure Bolt to use SSL/TLS. As the simplest configuration possible, we will just enable it in *neo4j.conf* and rely on the default values:

```
dbms.ssl.policy.bolt.enabled=true
```

Then create the mandatory directories:

```
$neo4j-home> mkdir certificates/bolt
$neo4j-home> mkdir certificates/bolt/trusted
$neo4j-home> mkdir certificates/bolt/revoked
```

Finally, place the files *private.key* and *public.crt* into the base directory:

```
$neo4j-home> cp /path/to/certs/private.key certificates/bolt
$neo4j-home> cp /path/to/certs/public.crt certificates/bolt
```

The base directory should now show the following listings:

```
$neo4j-home> ls certificates/bolt
-r----- ... private.key
-rw-r--r- ... public.crt
drwxr-xr-x ... revoked
drwxr-xr-x ... trusted
```

# 11.2.4. Choosing an SSL provider

The secure networking in Neo4j is provided through the Netty library, which supports both the native JDK SSL provider as well as Netty-supported OpenSSL derivatives.

Follow these steps to utilize OpenSSL:

- 1. Install a suitable dependency into the plugins/ folder of Neo4j. Dependencies can be downloaded from https://netty.io/wiki/forked-tomcat-native.html.
- 2. Set dbms.netty.ssl.provider=OPENSSL.



Using OpenSSL can significantly improve performance, especially for AES-GCM-cryptos, e.g. TLS\_ECDHE\_RSA\_WITH\_AES\_128\_GCM\_SHA256.

# 11.2.5. Terminology

The following terms are relevant to SSL support within Neo4j:

#### Certificate Authority (CA)

A trusted entity that issues electronic documents that can verify the identity of a digital entity. The term commonly refers to globally recognized CAs, but can also include internal CAs that are trusted inside of an organization. The electronic documents are digital certificates. They are an essential part of secure communication, and play an important part in the Public Key Infrastructure.

#### Certificate Revocation List (CRL)

In the event of a certificate being compromised, that certificate can be revoked. This is done by means of a list (located in one or several files) spelling out which certificates are revoked. The CRL

is always issued by the CA which issues the corresponding certificates.

#### cipher

An algorithm for performing encryption or decryption. In the most general implementation of encryption of Neo4j communications, we make implicit use of ciphers that are included as part of the Java platform. The configuration of the SSL framework also allows for the explicit declaration of allowed ciphers.

#### communication channel

A means for communicating with the Neo4j database. Available channels are:

- Bolt client traffic
- · HTTPS client traffic
- intra-cluster communication
- backup traffic

#### cryptographic objects

A term denoting the artifacts private keys, certificates and CRLs.

#### configuration parameters

These are the parameters defined for a certain ssl policy in neo4j.conf.

#### certificate

SSL certificates are issued by a trusted certificate authority (*CA*). The public key can be obtained and used by anyone to encrypt messages intended for a particular recipient. The certificate is commonly stored in a file named *<file name>.crt*. This is also referred to as the public key.

#### SAN

SAN is an acronym for *Subject Alternative Names*. It is an extension to certificates that one can include optionally. When presented with a certificate that includes SAN entries, it is recommended that the address of the host is checked against this field. Verifying that the hostname matches the certificate SAN helps prevent attacks where a rogue machine has access to a valid key pair.

#### SSL

SSL is an acronym for *Secure Sockets Layer*, and is the predecessor of TLS. It is common to refer to SSL/TLS as just SSL. However, the modern and secure version is TLS, and this is also the default in Neo4j.

#### SSL policy

An SSL policy in Neo4j consists of a a digital certificate and a set of configuration parameters defined in *neo4j.conf*.

## private key

The private key ensures that encrypted messages can be deciphered only by the intended recipient. The private key is commonly stored in a file named *file name*. It is important to protect the private key to ensure the integrity of encrypted communication.

### Public Key Infrastructure (PKI)

A set of roles, policies, and procedures needed to create, manage, distribute, use, store, and revoke digital certificates and manage public-key encryption.

#### public key

The public key can be obtained and used by anyone to encrypt messages intended for a particular recipient. This is also referred to as the certificate.

#### TLS protocol

The cryptographic protocol that provides communications security over a computer network. The Transport Layer Security (TLS) protocol and its predecessor, the Secure Sockets Layer (SSL) protocol are both frequently referred to as "SSL".

#### TLS version

A version of the TLS protocol.

# 11.3. Browser credentials handling

This section explains how to control how credentials are handled in Neo4j Browser.

Neo4j Browser has two mechanisms for avoiding users having to repeatedly enter their Neo4j credentials.

First, while the Browser is open in a web browser tab, it ensures that the existing database session is kept alive. This is subject to a timeout. The timeout is configured in the setting browser.credential\_timeout. The timeout is reset whenever there is user interaction with the Browser.

Second, the Browser can also cache the user's Neo4j credentials locally. When credentials are cached, they are stored unencrypted in the web browser's local storage. If the web browser tab is closed and then re-opened, the session is automatically re-established using the cached credentials. This local storage is also subject to the timeout configured in the setting <a href="mailto:browser.credential\_timeout">browser.credential\_timeout</a>. In addition, caching credentials in browser local storage can be disabled altogether. To disable credentials caching, set <a href="mailto:browser.retain\_connection\_credentials=false">browser.retain\_connection\_credentials=false</a> in the server configuration.

If the user issues a :server disconnect command then any existing session is terminated and the credentials are cleared from local storage.

# 11.4. Security checklist

This section provides a summary of recommendations regarding security in Neo4j.

Below is a simple checklist highlighting the specific areas within Neo4j that may need some extra attention in order to ensure the appropriate level of security for your application.

- 1. Deploy Neo4j on safe servers in safe networks:
  - a. Use subnets and firewalls.
  - b. Only open up for the necessary ports. For a list of relevant ports see Ports.

In particular, ensure that there is no external access to the port specified by the setting dbms.backup.listen\_address. Failing to protect this port may leave a security hole open by which an unauthorized user can make a copy of the database onto a different machine.

#### 2. Protect data-at-rest:

- a. Use volume encryption (e.g. Bitlocker).
- b. Manage access to database dumps (refer to Dump and load databases) and backups (refer to Perform a backup).
- c. Manage access to data files and transaction logs by ensuring the correct file permissions on the Neo4j files. Refer to File permissions for instructions on permission levels.

#### 3. Protect data-in-transit:

a. For remote access to the Neo4j database, only open up for encrypted Bolt or HTTPS.

- b. Use SSL certificates issued from a trusted Certificate Authority.
  - i. For configuring your Neo4j installation to use encrypted communication, refer to SSL framework.
  - ii. If using Causal Clustering, configure and use encryption for intra-cluster communication. For details, see Intra-cluster encryption.
  - iii. If using Causal Clustering, configure and use encryption for backups. This ensures that only servers with the specified SSL policy and SSL certificates will be able to access the server and perform the backup.
  - iv. For configuring your Bolt and/or HTTPS connectors, refer to Configure connectors.
  - v. If using LDAP, configure your LDAP system with encryption via StartTLS; see Use LDAP with encryption via StartTLS.
- 4. Be on top of the security for custom extensions:
  - a. Validate any custom code that you deploy (procedures and unmanaged extensions) and ensure that they do not expose any parts of the product or data unintentionally.
  - b. Survey the settings dbms.security.procedures.unrestricted and dbms.security.procedures.whitelist to ensure that they exclusively contain intentionally exposed extensions.
- 5. Ensure the correct file permissions on the Neo4j files.
- 6. Protect against the execution of unauthorized extensions by restricting access to the *bin*, *lib*, and *plugins* directories. Only the operating system user that Neo4j runs as should have permissions to those files. Refer to File permissions for instructions on permission levels.
- 7. If LOAD CSV is enabled, ensure that it does not allow unauthorized users to import data. How to configure LOAD CSV is described in Cypher Manual \( \Bigcap \) LOAD CSV.
- 8. Use Neo4j authentication. The setting dbms.security.auth\_enabled controls native authentication. The default value is true, which enables the native auth provider.
- Survey your neo4j.conf file for ports relating to deprecated functions such as remote JMX (controlled by the parameter setting dbms.jvm.additional=-Dcom.sun.management.jmxremote.port=3637).
- Review Browser credentials handling to determine whether the default credentials handling in Neo4j Browser complies with your security regulations. Follow the instructions to configure it if necessary.
- 11. Use the latest patch version of Neo4j.

# Chapter 12. Monitoring

This chapter describes the tools that are available for monitoring Neo4j.

Neo4j provides mechanisms for continuous analysis through the output of metrics as well as the inspection and management of currently-executing queries.

Logs can be harvested for continuous analysis, or for specific investigations. Facilities are available for producing security event logs as well as query logs. The query management functionality is provided for specific investigations into query performance. Monitoring features are also provided for ad-hoc analysis of a Causal Cluster.

This chapter describes the following:

- Metrics
  - Expose metrics
  - Metrics reference
- Logging
  - □ Security events logging
  - Query logging
- Query management
  - ☐ List all running queries
  - ☐ List all active locks for a query
  - ☐ Terminate multiple queries
  - ☐ Terminate a single query
- Transaction management
  - □ Configure transaction timeout
  - ☐ Configure lock acquisition timeout
  - List all running transactions
- Connection management
  - ☐ List all network connections
  - ☐ Terminate multiple network connections
  - ☐ Terminate a single network connection
- Monitoring a Causal Cluster
  - Procedures for monitoring a Causal Cluster
  - ☐ Endpoints for status information
- Monitoring the state of individual databases

## 12.1. Metrics

This section describes the Neo4j metrics output facilities.

## This section describes the following:

• Types of metrics

- Global metrics
- Database metrics
- Expose metrics
  - □ Enable metrics logging
  - □ Graphite
  - Prometheus
  - □ CSV files
  - □ JMX Beans
- · Metrics reference
  - ☐ General-purpose metrics
  - Metrics specific to Causal Clustering
  - ☐ Java Virtual Machine metrics

## 12.1.1. Types of metrics

This section describes the types of metrics available in Neo4j.

Neo4j provides a built-in metrics subsystem. Reported metrics can be queried via JMX, retrieved from CSV files, or consumed by third-party monitoring tools.

Neo4j has the following types of metrics:

- Global
- Per database covers individual database

#### Global metrics

Global metrics cover the whole database management system, and represents the status of the system as a whole.

Global metrics have following name format: <user-configured-prefix>.metric.name

Metrics of this type are reported as soon as the database management system is available. For example, all JVM related metrics are global. In particular, the neo4j.vm.thread.count metric has a default user-configured-prefix neo4j and the metric name is vm.thread.count.

By default, global metrics include:

- · Page cache metrics
- · GC metrics
- · Thread metrics
- · Memory pool metrics
- · Memory buffers metrics
- File descriptor metrics
- Bolt metrics
- · Web Server metrics

#### Database metrics

Each database metric is reported for a particular database only. Database metrics are only available during the lifetime of the database. When a database becomes unavailable, all of its metrics become unavailable also.

Database metrics have following name format: <user-configured-prefix>.<database-name>.metric.name

For example, any transaction metric is a database metric. In particular, the neo4j.mydb.transaction.started metric has a default user-configured-prefix neo4j and it is a metric for the mydb database.

By default, database metrics include:

- Transaction metrics
- Checkpoint metrics
- · Log rotation metrics
- · Database data metrics
- Cypher metrics
- Causal clustering metrics

## 12.1.2. Expose metrics

This section describes how to log and display various metrics by using the Neo4j metrics output facilities.

This section describes the following:

- · Enable metrics logging
- Graphite
- Prometheus
- CSV files
- IMX Beans

## Enable metrics logging

By default, metrics logging into CSV files and exposure via JMX beans are enabled. All metrics are enabled once metrics.enabled=true is set and you can use the individual settings to disable specific metrics.

```
# Setting for enabling all supported metrics.
metrics.enabled=true

# Setting for exposing metrics about transactions; number of transactions started, committed, etc.
metrics.neo4j.tx.enabled=false

# Setting for exposing metrics about the Neo4j page cache; page faults, evictions, flushes and exceptions, etc.
metrics.neo4j.pagecache.enabled=false

# Setting for exposing metrics about approximately entities are in the database; nodes, relationships, properties, etc.
metrics.neo4j.counts.enabled=false
```

## Graphite

Send metrics to Graphite or any monitoring tool based on the Graphite protocol.

Add the following settings to *neo4j.conf* in order to enable integration with Graphite:

```
# Enable the Graphite integration. Default is 'false'.
metrics.graphite.enabled=true
# The IP and port of the Graphite server on the format <hostname or IP address>:<port number>.
# The default port number for Graphite is 2003.
metrics.graphite.server=localhost:2003
# How often to send data. Default is 3 seconds.
metrics.graphite.interval=3s
# Prefix for Neo4j metrics on Graphite server.
metrics.prefix=Neo4j_1
```

Start Neo4j and connect to Graphite via a web browser in order to monitor your Neo4j metrics.



If you configure the Graphite server to be a hostname or DNS entry you should be aware that the JVM resolves hostnames to IP addresses and by default caches the result indefinitely for security reasons. This is controlled by the value of <a href="networkaddress.cache.ttl">networkaddress.cache.ttl</a> in the JVM Security properties. See <a href="https://docs.oracle.com/javase/8/docs/technotes/guides/net/properties.html">https://docs.oracle.com/javase/8/docs/technotes/guides/net/properties.html</a> for more information.

### Prometheus

Publish metrics for polling as Prometheus endpoint.

Add the following settings to neo4j.conf in order to enable the Prometheus endpoint.

```
# Enable the Prometheus endpoint. Default is 'false'.
metrics.prometheus.enabled=true
# The IP and port the endpoint will bind to in the format <hostname or IP address>:<port number>.
# The default is localhost:2004.
metrics.prometheus.endpoint=localhost:2004
```

When Neo4j is fully started a Prometheus endpoint will be available at the configured address.

#### CSV files

Export metrics to CSV files.

Add the following settings to neo4j.conf in order to enable export of metrics into local .CSV files:

```
# Enable the CSV exporter. Default is 'true'.
metrics.csv.enabled=true
# Directory path for output files.
# Default is a "metrics" directory under NEO4J_HOME.
#dbms.directories.metrics='/local/file/system/path'
# How often to store data. Default is 3 seconds.
metrics.csv.interval=3s
# The maximum number of CSV files that will be saved. Default is 7.
metrics.csv.rotation.keep_number
# The file size at which the csv files will auto-rotate. Default is 10M.
metrics.csv.rotation.size=10M
```

## **IMX** Beans

Expose metrics over JMX beans.

In order to enable metrics exposure via JMX, add the following setting to neo4j.conf:

```
# Enable settings export via JMX. Default is 'true'.
metrics.jmx.enabled=true
```

# 12.1.3. Metrics reference

This section provides a listing of available metrics.

## This section describes the following:

- General-purpose metrics
- Metrics specific to Causal Clustering
- Java Virtual Machine metrics

## General-purpose metrics

## Table 24. Database checkpointing metrics

Name	Description
<pre><pre><pre>fix&gt;.check_point. events</pre></pre></pre>	The total number of check point events executed so far.
<pre><pre><pre>fix&gt;.check_point. total_time</pre></pre></pre>	The total time spent in check pointing so far.
<pre><pre><pre><pre>fix&gt;.check_point. duration</pre></pre></pre></pre>	The duration of the last check point event.

### Table 25. Database data metrics

Name	Description
<pre><pre><pre>fix&gt;.ids_in_use.r elationship_type</pre></pre></pre>	The total number of different relationship types stored in the database.
<pre><pre>fix&gt;.ids_in_use.p roperty</pre></pre>	The total number of different property names used in the database.
<pre><pre><pre>fix&gt;.ids_in_use.r elationship</pre></pre></pre>	The total number of relationships stored in the database.
<pre><prefix>.ids_in_use.n ode</prefix></pre>	The total number of nodes stored in the database.

## Table 26. Database page cache metrics

Name	Description
<pre><pre><pre>cache.e viction_exceptions</pre></pre></pre>	The total number of exceptions seen during the eviction process in the page cache.
<pre><prefix>.page_cache.f lushes</prefix></pre>	The total number of flushes executed by the page cache.
<pre><pre><pre>cache.u npins</pre></pre></pre>	The total number of page unpins executed by the page cache.
<pre><prefix>.page_cache.p ins</prefix></pre>	The total number of page pins executed by the page cache.
<pre><pre><pre>cache.e victions</pre></pre></pre>	The total number of page evictions executed by the page cache.
<pre><pre><pre>fix&gt;.page_cache.p age_faults</pre></pre></pre>	The total number of page faults happened in the page cache.

Name	Description
<pre><prefix>.page_cache.h its</prefix></pre>	The total number of page hits happened in the page cache.
<pre><prefix>.page_cache.h it_ratio</prefix></pre>	The ratio of hits to the total number of lookups in the page cache.
<pre><prefix>.page_cache.u sage_ratio</prefix></pre>	The ratio of number of used pages to total number of available pages.

### Table 27. Database transaction metrics

Name	Description
<pre><prefix>.transaction. started</prefix></pre>	The total number of started transactions.
<pre><pre><pre>fix&gt;.transaction. peak_concurrent</pre></pre></pre>	The highest peak of concurrent transactions.
<pre><pre><pre>fix&gt;.transaction. active</pre></pre></pre>	The number of currently active transactions.
<pre><pre><pre><pre>fix&gt;.transaction. active_read</pre></pre></pre></pre>	The number of currently active read transactions.
<pre><pre><pre><pre>fix&gt;.transaction. active_write</pre></pre></pre></pre>	The number of currently active write transactions.
<pre><pre><pre><pre>fix&gt;.transaction. committed</pre></pre></pre></pre>	The total number of committed transactions.
<pre><pre><pre>fix&gt;.transaction. committed_read</pre></pre></pre>	The total number of committed read transactions.
<pre><pre><pre><pre>fix&gt;.transaction. committed_write</pre></pre></pre></pre>	The total number of committed write transactions.
<pre><pre><pre>fix&gt;.transaction. rollbacks</pre></pre></pre>	The total number of rolled back transactions.
<pre><pre><pre>fix&gt;.transaction. rollbacks_read</pre></pre></pre>	The total number of rolled back read transactions.
<pre><pre><pre>fix&gt;.transaction. rollbacks_write</pre></pre></pre>	The total number of rolled back write transactions.
<pre><pre><pre>fix&gt;.transaction. terminated</pre></pre></pre>	The total number of terminated transactions.
<pre><pre><pre>fix&gt;.transaction. terminated_read</pre></pre></pre>	The total number of terminated read transactions.
<pre><pre><pre>fix&gt;.transaction. terminated_write</pre></pre></pre>	The total number of terminated write transactions.
<pre><pre><pre>fix&gt;.transaction. last_committed_tx_id</pre></pre></pre>	The ID of the last committed transaction.
<pre><pre><pre>fix&gt;.transaction. last_closed_tx_id</pre></pre></pre>	The ID of the last closed transaction.

# Table 28. Cypher metrics

Name	Description
<pre><pre><pre>fix&gt;.cypher.repla n_events</pre></pre></pre>	The total number of times Cypher has decided to re-plan a query.
<pre><pre><pre>fix&gt;.cypher.repla n_wait_time</pre></pre></pre>	The total number of seconds waited between query replans.

### Table 29. Database transaction log metrics

Name	Description
<pre><pre><pre><pre>prefix&gt;.log.rotation _events</pre></pre></pre></pre>	The total number of transaction log rotations executed so far.
<pre><pre><pre>fix&gt;.log.rotation _total_time</pre></pre></pre>	The total time spent in rotating transaction logs so far.
<pre><pre><pre><pre>fix&gt;.log.rotation _duration</pre></pre></pre></pre>	The duration of the last log rotation event.
<pre><pre><pre>fix&gt;.log.appended _bytes</pre></pre></pre>	The total number of bytes appended to transaction log.

### Table 30. Bolt metrics

Name	Description
<pre><pre><pre><pre>s_started</pre></pre></pre></pre>	The total number of Bolt sessions started since this instance started. This includes both succeeded and failed sessions (deprecated, use connections_opened instead).
<pre><pre><pre>fix&gt;.bolt.connect ions_opened</pre></pre></pre>	The total number of Bolt connections opened since this instance started. This includes both succeeded and failed connections.
<pre><pre><pre>fix&gt;.bolt.connect ions_closed</pre></pre></pre>	The total number of Bolt connections closed since this instance started. This includes both properly and abnormally ended connections.
<pre><pre><pre>fix&gt;.bolt.connect ions_running</pre></pre></pre>	The total number of Bolt connections currently being executed.
<pre><pre><pre>fix&gt;.bolt.connect ions_idle</pre></pre></pre>	The total number of Bolt connections sitting idle.
<pre><pre><pre><pre>s_received</pre></pre></pre></pre>	The total number of messages received via Bolt since this instance started.
<pre><pre><pre><pre><pre><pre>s_started</pre></pre></pre></pre></pre></pre>	The total number of messages that began processing since this instance started. This is different from messages received in that this counter tracks how many of the received messages have been taken on by a worker thread.
<pre><pre><pre>fix&gt;.bolt.message s_done</pre></pre></pre>	The total number of messages that completed processing since this instance started. This includes successful, failed and ignored Bolt messages.
<pre><pre><pre><pre>s_failed</pre></pre></pre></pre>	The total number of messages that failed processing since this instance started.
<pre><pre><pre>fix&gt;.bolt.accumul ated_queue_time</pre></pre></pre>	The accumulated time messages have spent waiting for a worker thread.
<pre><pre><pre><pre>fix&gt;.bolt.accumul ated_processing_time</pre></pre></pre></pre>	The accumulated time worker threads have spent processing messages.

### Table 31. Server metrics

Name	Description
<pre><pre><pre><pre><pre><pre><pre>ds.jetty.idle</pre></pre></pre></pre></pre></pre></pre>	The total number of idle threads in the jetty pool.
<pre><prefix>.server.threa ds.jetty.all</prefix></pre>	The total number of threads (both idle and busy) in the jetty pool.

# Metrics specific to Causal Clustering

### Table 32. Core metrics

Name	Description
<pre><pre><pre><pre>core.append_ind ex</pre></pre></pre></pre>	Append index of the RAFT log.
<pre><pre><pre><pre>core.commit_ind ex</pre></pre></pre></pre>	Commit index of the RAFT log.

Name	Description
<pre><pre><pre>cprefix&gt;.causal_clust ering.core.term</pre></pre></pre>	RAFT Term of this server.
<pre><pre><pre>cprefix&gt;.causal_clust ering.core.tx_retries</pre></pre></pre>	Transaction retries.
<pre><pre><pre>cprefix&gt;.causal_clust ering.core.is_leader</pre></pre></pre>	Is this server the leader?
<pre><pre><pre><pre>core.in_flight_ cache.total_bytes</pre></pre></pre></pre>	In-flight cache total bytes.
<pre><pre><pre><pre><pre><pre><pre><pre></pre></pre></pre></pre></pre></pre></pre></pre>	In-flight cache max bytes.
<pre><pre><pre>cprefix&gt;.causal_clust ering.core.in_flight_ cache.element_count</pre></pre></pre>	In-flight cache element count.
<pre><pre><pre><pre>core.in_flight_ cache.max_elements</pre></pre></pre></pre>	In-flight cache maximum elements.
<pre><pre><pre>cprefix&gt;.causal_clust ering.core.in_flight_ cache.hits</pre></pre></pre>	In-flight cache hits.
<pre><pre><pre><pre>causal_clust ering.core.in_flight_ cache.misses</pre></pre></pre></pre>	In-flight cache misses.
<pre><pre><pre><pre>core.message_pr ocessing_delay</pre></pre></pre></pre>	Delay between RAFT message receive and process.
<pre><pre><pre><pre>core.message_pr ocessing_timer</pre></pre></pre></pre>	Timer for RAFT message processing.
<pre><pre><pre><pre>core.discovery. replicated_data</pre></pre></pre></pre>	Size of replicated data structures.
<pre><pre><pre><pre>causal_clust ering.core.replicatio n_new</pre></pre></pre></pre>	Raft replication new request count.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.replicatio n_attempt</pre></pre></pre></pre>	Raft replication attempt count.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.replicatio n_fail</pre></pre></pre></pre>	Raft Replication fail count.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.replicatio n_maybe</pre></pre></pre></pre>	Raft Replication maybe count.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.replicatio n_success</pre></pre></pre></pre>	Raft Replication success count.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.discovery. cluster.members</pre></pre></pre></pre>	Discovery cluster member size.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.discovery. cluster.unreachable</pre></pre></pre></pre>	Discovery cluster unreachable size.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.core.discovery. cluster.converged</pre></pre></pre></pre>	Discovery cluster convergence.

Table 33. Read Replica Metrics

Name	Description
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.read_replica.pu ll_updates</pre></pre></pre></pre>	The total number of pull requests made by this instance.
<pre><pre>cprefix&gt;.causal_clust ering.read_replica.pu ll_update_highest_tx_ id_requested</pre></pre>	The highest transaction id requested in a pull update by this instance.
<pre><pre><pre><pre>cprefix&gt;.causal_clust ering.read_replica.pu ll_update_highest_tx_ id_received</pre></pre></pre></pre>	The highest transaction id that has been pulled in the last pull updates by this instance.

### Java Virtual Machine Metrics

These metrics are environment dependent and they may vary on different hardware and with JVM configurations. Typically these metrics will show information about garbage collections (for example the number of events and time spent collecting), memory pools and buffers, and finally the number of active threads running.

### Table 34. GC metrics.

Name	Description
<pre><prefix>.vm.gc.time.% s</prefix></pre>	Accumulated garbage collection time in milliseconds.
<pre><prefix>.vm.gc.count. %s</prefix></pre>	Total number of garbage collections.

### Table 35. JVM memory buffers metrics.

Name	Description
<pre><pre><pre><pre><pre><pre><pre><pre></pre></pre></pre></pre></pre></pre></pre></pre>	Estimated number of buffers in the pool.
<pre><pre><pre><pre><pre><pre><pre><pre></pre></pre></pre></pre></pre></pre></pre></pre>	Estimated amount of memory used by the pool.
<pre><pre><pre><pre><pre><pre><pre><pre></pre></pre></pre></pre></pre></pre></pre></pre>	Estimated total capacity of buffers in the pool.

### Table 36. JVM memory pools metrics.

Name	Description
<pre><prefix>.vm.memory.po ol.%s</prefix></pre>	Estimated number of buffers in the pool.

### Table 37. JVM threads metrics.

Name	Description
<pre><prefix>.vm.thread.co unt</prefix></pre>	Estimated number of active threads in the current thread group.
<pre><prefix>.vm.thread.to tal</prefix></pre>	The total number of live threads including daemon and non-daemon threads.

# 12.2. Logging

This section describes security and query logging in Neo4j.

Neo4j provides two types of logs for inspection of queries that are run in the database, and of security

events that have occurred.

The section describes the following:

- Query logging
- · Security events logging

# 12.2.1. Query logging

This section describes Neo4j support for query logging.

Neo4j can be configured to log queries executed in the database.

Query logging is enabled by default and is controlled by the setting dbms.logs.query.enabled.

### Configuration options are:

Option	Default	Description
off		Will completely disable logging.
info		Will log at the end of queries that have either succeeded or failed. The dbms.logs.query.threshold parameter is used to determine the threshold for logging a query. If the execution of a query takes a longer time than this threshold, it will be logged. Setting the threshold to 0 will result in all queries being logged.
verbose	Yes	Will log all queries at both start and finish, regardless of dbms.logs.query.threshold.

# Log configuration

The name of the log file is query. log and it resides in the logs directory (see File locations).

Rotation of the query log can be configured in the *neo4j.conf* configuration file. The following parameters are available:

Parameter name	Default value	Description
<pre>dbms.logs.query.allocation_logging_e nabled</pre>	false	Log allocated bytes for the executed queries being logged.
dbms.logs.query.enabled	verbose	Log executed queries.
dbms.logs.query.page_logging_enabled	false	Log page hits and page faults for the executed queries being logged.
<pre>dbms.logs.query.parameter_logging_en abled</pre>	true	Log parameters for executed queries that take longer than the configured threshold.
dbms.logs.query.rotation.keep_number	7	Sets number of historical log files kept.
dbms.logs.query.rotation.size	20M	Sets the file size at which the query log will auto-rotate.
dbms.logs.query.threshold	0	If the execution of query takes a longer time than this threshold, the query is logged (provided query logging is set to info).

Parameter name	Default value	Description
dbms.logs.query.time_logging_enabled	false	Log detailed time information for the executed queries being logged.

### Example 58. Configure for simple query logging

In this example we set query logging to info, but leave all other query log parameters at their defaults.

```
dbms.logs.query.enabled=info
```

Below is an example of the query log with this basic configuration:

```
2017-11-22 14:31 ... INFO 9 ms: bolt-session
                                                        johndoe neo4j-javascript/1.4.1
                                               bolt
client/127.0.0.1:59167
2017-11-22 14:31 ... INFO 0 ms: bolt-session
                                               bolt
                                                        johndoe neo4j-javascript/1.4.1
client/127.0.0.1:59167
2017-11-22 14:32 ... INFO 3 ms: server-session http
                                                                   /db/data/cypher neo4j - CALL
dbms.procedures() - {}
2017-11-22 14:32 ... INFO 1 ms: server-session http
                                                        127.0.0.1 /db/data/cypher neo4j - CALL
dbms.showCurrentUs...
2017-11-22 14:32 ... INFO 0 ms: bolt-session
                                                        johndoe neo4j-javascript/1.4.1
                                               bolt
client/127.0.0.1:59167
2017-11-22 14:32 ... INFO 0 ms: bolt-session
                                                bolt
                                                        johndoe neo4j-javascript/1.4.1
client/127.0.0.1:59167
2017-11-22 14:32 ... INFO 2 ms: bolt-session
                                                bolt
                                                        johndoe neo4j-javascript/1.4.1
client/127.0.0.1:59261 ...
```

### Example 59. Configure for query logging with more details

In this example we turn query logging on, and also enable some additional logging.

```
dbms.logs.query.parameter_logging_enabled=true
dbms.logs.query.time_logging_enabled=true
dbms.logs.query.allocation_logging_enabled=true
dbms.logs.query.page_logging_enabled=true
```

Below is an example of the query log with these configuration parameters enabled:

```
2017-11-22 12:38 ... INFO 3 ms: bolt-session bolt johndoe neo4j-javascript/1.4.1 ... 2017-11-22 22:38 ... INFO 61 ms: (planning: 0, cpu: 58, waiting: 0) - 6164496 B - 0 page hits, 1 page faults ... 2017-11-22 12:38 ... INFO 78 ms: (planning: 40, cpu: 74, waiting: 0) - 6347592 B - 0 page hits, 0 page faults ... 2017-11-22 12:38 ... INFO 44 ms: (planning: 9, cpu: 25, waiting: 0) - 1311384 B - 0 page hits, 0 page faults ... 2017-11-22 12:38 ... INFO 6 ms: (planning: 2, cpu: 6, waiting: 0) - 420872 B - 0 page hits, 0 page faults - ...
```

### Attach metadata to a query

It is possible to attach metadata to a query and have it printed in the query log, using the built-in procedure tx.setMetaData. This is typically done programmatically, but can be illustrated as follows, using cypher-shell.

Start a transaction and call tx.setMetaData with a list of meta data.

```
neo4j> :begin
neo4j# CALL tx.setMetaData({ User: 'jsmith', AppServer: 'app03.dc01.company.com'});
neo4j# CALL dbms.procedures() YIELD name RETURN COUNT(name);
COUNT(name)
39
neo4j# :commit
```

Below are the corresponding results in the query log:

```
... CALL tx.setMetaData({ User: 'jsmith', AppServer: 'app03.dc01.company.com'}); - {} - {} ... CALL dbms.procedures() YIELD name RETURN COUNT(name); - {} - {User: 'jsmith', AppServer: 'app03.dc01.company.com'}
```

## 12.2.2. Security events logging

This section describes Neo4j support for security events logging.

Neo4j provides security event logging that records all security events.

For native user management, the following actions are recorded:

- Login attempts per default both successful and unsuccessful logins are recorded.
- All administration commands run towards the system database.
- All security procedures run towards the system database.

### Log configuration

The name of the log file is security.log and it resides in the logs directory (see File locations).

Rotation of the security events log can be configured in the *neo4j.conf* configuration file. The following parameters are available:

Parameter name	Default value	Description
dbms.logs.security.rotation.size	20M	Sets the file size at which the security event log will auto-rotate.
dbms.logs.security.rotation.delay	300s	Sets the minimum time interval after the last log rotation occurred, before the log may be rotated again.
<pre>dbms.logs.security.rotation.keep_num ber</pre>	7	Sets number of historical log files kept.

If using LDAP as the authentication method, some cases of LDAP misconfiguration will also be logged, as well as LDAP server communication events and failures.

If many programmatic interactions are expected, it is advised to disable the logging of successful logins. Logging of successful logins is disabled by setting the

 ${\tt dbms.security.log\_successful\_authentication}\ parameter\ in\ the\ {\tt neo4j.conf}\ file:$ 

dbms.security.log\_successful\_authentication=false

### Below is an example of the security log:

```
2019-12-09 13:45:00.796+0000 INFO [AsyncLog @ 2019-12-09 ...] [johnsmith]: logged in 2019-12-09 13:47:53.443+0000 ERROR [AsyncLog @ 2019-12-09 ...] [johndoe]: failed to log in: invalid principal or credentials 2019-12-09 13:48:28.566+0000 INFO [AsyncLog @ 2019-12-09 ...] [johnsmith]: CREATE USER janedoe SET PASSWORD '***** CHANGE REQUIRED 2019-12-09 13:48:32.753+0000 INFO [AsyncLog @ 2019-12-09 ...] [johnsmith]: CREATE ROLE custom 2019-12-09 13:49:11.880+0000 INFO [AsyncLog @ 2019-12-09 ...] [johnsmith]: GRANT ROLE custom TO janedoe 2019-12-09 13:49:34.979+0000 INFO [AsyncLog @ 2019-12-09 ...] [johnsmith]: GRANT TRAVERSE ON GRAPH * NODES A, B (*) TO custom 2019-12-09 13:49:37.053+0000 INFO [AsyncLog @ 2019-12-09 ...] [johnsmith]: DROP USER janedoe
```

# 12.3. Query management

This section describes facilities for query management.

### This section describes the following:

- List all running queries
- List all active locks for a query
- Terminate multiple queries
- Terminate a single query

### 12.3.1. List all running queries

An administrator is able to view all queries that are currently executing within the instance. Alternatively, the current user may view all of their own currently-executing queries.

### Syntax:

CALL dbms.listQueries()

Name	Туре	Description
queryId	String	This is the ID of the query.
username	String	This is the username of the user who is executing the query.
metaData	Мар	This is any metadata associated with the transaction.
query	String	This is the query itself.
parameters	Мар	This is a map containing all the parameters used by the query.
planner	String	Planner used by the query
runtime	String	Runtime used by the query
indexes	List	Indexes used by the query
startTime	String	This is the time at which the query was started.
elapsedTime	String	Deprecated: Use <a href="mailto:elapsed">elapsed</a> This is the time that has elapsed since the query was started.

Name	Туре	Description
connectionDetails	String	Deprecated: Use requestScheme, clientAddress,requestUri These are the connection details pertaining to the query.
protocol	String	The protocol used by connection issuing the query.
connectionId	String	The ID of the connection issuing the query. This field will be null if the query was issued using embedded API.
clientAddress	String	The client address of the connection issuing the query.
requestUri	String	The request URI used by the client connection issuing the query.
status	String	Status of the executing query.
resourceInformation	Мар	Status of the executing query.
activeLockCount	Integer	Count of active locks held by transaction executing the query.
elapsedTimeMillis	Integer	This is the time in milliseconds that has elapsed since the query was started.
cpuTimeMillis	Integer	CPU time in milliseconds that has been actively spent executing the query. This field will be null unless the config parameter dbms.track_query_cpu_time is set to true.
waitTimeMillis	Integer	Wait time in milliseconds that has been spent waiting to acquire locks.
idleTimeMillis	Integer	Idle time in milliseconds. This field will be null unless the config parameter dbms.track_query_cpu_time is set to true.
allocatedBytes	Integer	Bytes allocated for the executing query. This number is cumulative over the duration of the query. For memory-intense or long-running queries the value may be larger than the current memory allocation. This field will be null unless the config parameter dbms.track_query_allocation is set to true.
pageHits	Integer	Page hits occurred during the execution.
pageFaults	Integer	Page faults occurred during the execution.
database	String	This is the name of the database the query is executing against.

The following example shows that the user alwood is currently running dbms.listQueries() yielding specific variables, namely queryId, username, query, elapsedTimeMillis, requestUri, and status. It is executed against the database myDb, which is indicated by the prefix in the queryId.

```
CALL dbms.listQueries() YIELD queryId, username, query, elapsedTimeMillis, requestUri, status
```

# 12.3.2. List all active locks for a query

An administrator is able to view all active locks held by the transaction executing the query with the queryId.

### Syntax:

CALL dbms.listActiveLocks(queryId)

Name	Туре	Description
mode	String	Lock mode corresponding to the transaction.
resourceType	String	Resource type of the locked resource
resourceId	Integer	Resource id of the locked resource .

The following example shows the active locks held by transaction executing query with id myDb-query-614

```
CALL dbms.listActiveLocks( "myDb-query-614" )
```

The following example shows the active locks for all currently executing queries by yielding the queryId from dbms.listQueries procedure

```
CALL dbms.listQueries() YIELD queryId, query
CALL dbms.listActiveLocks( queryId ) YIELD resourceType, resourceId, mode
RETURN queryId, query, resourceType, resourceId, mode
```

### 12.3.3. Terminate multiple queries

An administrator is able to terminate within the instance all transactions executing a query with any of the given query IDs. Alternatively, the current user may terminate all of their own transactions executing a query with any of the given query IDs.

#### Syntax:

CALL dbms.killQueries(queryIds)

#### **Arguments:**

Name	Туре	Description
ids	List <string></string>	This is a list of the IDs of all the queries to be terminated.

Name	Туре	Description
queryId	String	This is the ID of the terminated query.
username	String	This is the username of the user who was executing the (now terminated) query.

### Example 63. Terminating multiple queries

The following example shows that the administrator has terminated the queries with IDs joeDb-query-378 and anneDb-query-765, started by the users joesmith and annebrown, respectively.

This command can target queries from multiple databases at the same time, as noted by the prefixes joeDb- and anneDb-:

```
CALL dbms.killQueries(['joeDb-query-378','anneDb-query-765'])
```

# 12.3.4. Terminate a single query

An administrator is able to terminate within the instance any transaction executing the query whose ID is provided. Alternatively, the current user may terminate their own transaction executing the query whose ID is provided.

### Syntax:

CALL dbms.killQuery(queryId)

### **Arguments:**

Name	Туре	Description
id	String	This is the ID of the query to be terminated.

Name	Туре	Description
queryId	String	This is the ID of the terminated query.
username	String	This is the username of the user who was executing the (now terminated) query.
message	String	A message stating whether the query was successfully found.

The following example shows that the user joesmith has terminated his query with the ID joeDb-query-502.

```
CALL dbms.killQuery('joeDb-query-502')
```

The following example shows the output when trying to kill a query with an ID that does not exist.

```
CALL dbms.killQuery('myDb-query-502')
```

# 12.4. Transaction management

This section describes facilities for transaction management.

This section describes the following:

- Configure transaction timeout
- Configure lock acquisition timeout
- List all running transactions

# 12.4.1. Configure transaction timeout

It is possible to configure Neo4j to terminate transactions whose execution time has exceeded the configured timeout. To enable this feature, set <a href="mailto:dbms.transaction.timeout">dbms.transaction.timeout</a> to some positive time interval value denoting the default transaction timeout. Setting <a href="mailto:dbms.transaction.timeout">dbms.transaction.timeout</a> to <a href="mailto:dbms.transaction.timeout">— which is the default value — disables the feature.

Example 65. Configure transaction timeout

```
Set the timeout to ten seconds.
```

```
dbms.transaction.timeout=10s
```

Configuring transaction timeout will have no effect on transactions executed with custom timeouts (via the Java API), as a custom timeout will override the value set for dbms.transaction.timeout.

The transaction timeout feature is also known as the transaction guard.

### 12.4.2. Configure lock acquisition timeout

An executing transaction may get stuck while waiting for some lock to be released by another transaction. A transaction in such state is not desirable, and in some cases it is better for the transaction to instead give up and fail.

To enable this feature, set <a href="mailto:dock.acquisition.timeout">dock.acquisition.timeout</a> to some positive time interval value denoting the maximum time interval within which any particular lock should be acquired, before failing the transaction. Setting <a href="mailto:dock.acquisition.timeout">dock.acquisition.timeout</a> to 0 — which is the default value — disables the lock acquisition timeout.

### Example 66. Configure lock acquisition timeout

Set the timeout to ten seconds.

dbms.lock.acquisition.timeout=10s

# 12.4.3. List all running transactions

An administrator is able to view all transactions that are currently executing within the instance. Alternatively, the current user may view all of their own currently-executing transactions.

### Syntax:

CALL dbms.listTransactions()

Name	Туре	Description
transactionId	String	This is the ID of the transaction.
username	String	This is the username of the user who is executing the transaction.
metaData	Мар	This is any metadata associated with the transaction.
startTime	String	This is the time at which the transaction was started.
protocol	String	The protocol used by connection issuing the transaction.
connectionId	String	The ID of the connection issuing the transaction. This field will be null if the transaction was issued using embedded API.
clientAddress	String	The client address of the connection issuing the transaction.
requestUri	String	The request URI used by the client connection issuing the transaction.
currentQueryId	String	This is the ID of the current query executed by transaction.
currentQuery	String	This is the current query executed by transaction.

Name	Туре	Description
activeLockCount	Integer	Count of active locks held by transaction.
status	String	Status of the executing transaction.
resourceInformation	Мар	Information about what transaction is waiting for when it is blocked.
elapsedTimeMillis	Integer	This is the time in milliseconds that has elapsed since the transaction was started.
cpuTimeMillis	Integer	CPU time in milliseconds that has been actively spent executing the transaction.
waitTimeMillis	Integer	Wait time in milliseconds that has been spent waiting to acquire locks.
idleTimeMillis	Integer	Idle time in milliseconds.
allocatedBytes	Integer	Bytes allocated for the executing transaction. This number is cumulative over the duration of the transaction. For memory-intense or long-running transactions the value may be larger than the current memory allocation.
allocatedDirectBytes	Integer	Direct bytes used by the executing transaction.
pageHits	Integer	Page hits occurred during the execution.
pageFaults	Integer	Page faults occurred during the execution.
database	String	This is the name of the database the transaction is executing against.

### Example 67. Viewing transactions that are currently executing

The following example shows that the user 'alwood' is currently running dbms.listTransactions(). The procedure call yields specific information about the running transaction, namely transactionId, username, currentQuery, elapsedTimeMillis, requestUri, and status.

 ${\tt CALL\ dbms.listTransactions()\ YIELD\ transactionId,\ username,\ currentQuery,\ elapsedTimeMillis,\ requestUri,\ status}$ 

```
"transactionId" "username" "currentQuery" "elapsedTimeMillis"

"myDb-transaction-22" "alwood" "CALL dbms.listTransactions() YIELD "1"

"127.0.0.1:7687" "Running" transactionId, username, currentQuery

elapsedTime, requestUri, status"
```

# 12.5. Connection management

This section describes facilities for connection management.

### This section describes the following:

- List all network connections
- Terminate multiple network connections
- Terminate a single network connection

### 12.5.1. List all network connections

An administrator is able to view all network connections within the database instance. Alternatively, the current user may view all of their own network connections.

The procedure dbms.listConnections lists all accepted network connections for all configured connectors, including Bolt, HTTP, and HTTPS. Some listed connections might never perform authentication. For example, HTTP GET requests to the Neo4j Browser endpoint fetches static resources and does not need to authenticate. However, connections made using Neo4j Browser require the user to provide credentials and perform authentication.

### Syntax:

CALL dbms.listConnections()

Name	Туре	Description
connectionId	String	This is the ID of the network connection.
connectTime	String	This is the time at which the connection was started.
connector	String	Name of the connector that accepted the connection.
username	String	This is the username of the user who initiated the connection. This field will be null if the transaction was issued using embedded API. It can also be null if connection did not perform authentication.
userAgent	String	Name of the software that is connected. This information is extracted from the User-Agent request header for HTTP and HTTPS connections. It is available natively for Bolt connections which supply the agent string in an initialization message.
serverAddress	String	The server address this connection is connected to.
clientAddress	String	The client address of the connection.

The following example shows that the user 'alwood' is connected using Java driver and a Firefox web browser. The procedure call yields specific information about the connection, namely connectionId, connectTime, connector, username, userAgent, and clientAddress.

```
{\tt CALL \ dbms.listConnections() \ YIELD \ connectionId, \ connectTime, \ connector, \ username, \ userAgent, \ clientAddress}
```

### 12.5.2. Terminate multiple network connections

An administrator is able to terminate within the instance all network connections with any of the given IDs. Alternatively, the current user may terminate all of their own network connections with any of the given IDs.

### Syntax:

CALL dbms.killConnections(connectionIds)

### **Arguments:**

Name	Туре	Description
ids	List <string></string>	This is a list of the IDs of all the connections to be terminated.

#### **Returns:**

Name	Туре	Description
connectionId	String	This is the ID of the terminated connection.
username	String	This is the username of the user who initiated the (now terminated) connection.
message	String	A message stating whether the connection was successfully found.

#### **Considerations:**

Bolt connections are stateful. Termination of a Bolt connection results in termination of the ongoing query/transaction.

Termination of an HTTP/HTTPS connection can terminate the ongoing HTTP/HTTPS request.

The following example shows that the administrator has terminated the connections with IDs 'bolt-37' and 'https-11', started by the users 'joesmith' and 'annebrown', respectively. The administrator also attempted to terminate the connection with ID 'http-42' which did not exist.

```
CALL dbms.killConnections(['bolt-37', 'https-11', 'http-42'])

"connectionId" "username" "message"

"bolt-37"   "joesmith" "Connection found"

"https-11"   "annebrown" "Connection found"

"http-42"   "n/a"   "No connection found with this id"

3 rows
```

# 12.5.3. Terminate a single network connection

An administrator is able to terminate within the instance any network connection with the given ID. Alternatively, the current user may terminate their own network connection with the given ID.

### Syntax:

CALL dbms.killConnection(connectionId)

### **Arguments:**

Name	Туре	Description
id		This is the ID of the connection to be terminated.

#### **Returns:**

Name	Туре	Description
connectionId	String	This is the ID of the terminated connection.
username	String	This is the username of the user who initiated the (now terminated) connection.
message	String	A message stating whether the connection was successfully found.

#### **Considerations:**

Bolt connections are stateful. Termination of a Bolt connection results in termination of the ongoing query/transaction. Termination of an HTTP/HTTPS connection can terminate the ongoing HTTP/HTTPS request.

The following example shows that the user 'joesmith' has terminated his connection with the ID 'bolt-4321'.

```
"bolt-4321'.

CALL dbms.killConnection('bolt-4321')

"connectionId" "username" "message"
   "bolt-4321"   "joesmith" "Connection found"
1 row

The following example shows the output when trying to kill a connection with an ID that does not exist.

CALL dbms.killConnection('bolt-987')

"connectionId" "username" "message"
   "bolt-987"   "n/a"   "No connection found with this id"
1 row
```

# 12.6. Monitoring a Causal Cluster

This section covers additional facilities available for monitoring a Neo4j Causal Cluster.

In addition to specific metrics as described in previous sections, Neo4j Causal Clusters provide an infrastructure that operators will wish to monitor. The procedures can be used to inspect the cluster state and to understand its current condition and topology. Additionally, there are HTTP endpoints for checking health and status.

This section describes the following:

- Procedures for monitoring a Causal Cluster
  - ☐ Find out the role of a cluster member
  - Gain an overview over the instances in the cluster
  - Get routing recommendations
- Endpoints for status information
  - Adjusting security settings for Causal Clustering endpoints
  - Unified endpoints

### 12.6.1. Procedures for monitoring a Causal Cluster

This section covers procedures for monitoring a Neo4j Causal Cluster.

A number of procedures are available that provide information about a cluster. This section describes

### the following:

- Find out the role of a cluster member
- Gain an overview over the instances in the cluster
- Get routing recommendations

### Find out the role of a cluster member

The procedure dbms.cluster.role(databaseName) can be called on every instance in a Causal Cluster to return the role of the instance. Each instance holds multiple databases and participates in multiple independent Raft groups. The role returned by the procedure is for the database denoted by the databaseName parameter.

### Syntax:

CALL dbms.cluster.role(databaseName)

### **Arguments:**

Name	Туре	Description
databaseName	String	The name of the database to get the cluster role for.

### **Returns:**

Name	Туре	Description
role	String	This is the role of the current instance, which can be LEADER, FOLLOWER, or READ_REPLICA.

### **Considerations:**

• While this procedure is useful in and of itself, it serves as basis for more powerful monitoring procedures.

### Example 71. Check the role of this instance

The following example shows how to find out the role of the current instance for database neo4j, which in this case is FOLLOWER.

CALL dbms.cluster.role("neo4j")

role

FOLLOWER

### Gain an overview over the instances in the cluster

The procedure <a href="mailto:dbms.cluster.overview">dbms.cluster.overview</a>() provides an overview of cluster topology by returning details on all the instances in the cluster.

### Syntax:

CALL dbms.cluster.overview()

### **Returns:**

Name	Туре	Description
id	String	This is id of the instance.
addresses	List	This is a list of all the addresses for the instance.
groups	List	This is a list of all the server groups which an instance is part of.
databases	Мар	This is a map of all databases with corresponding roles which the instance is hosting. The keys in the map are database names. The values are roles of this instance in the corresponding Raft groups, which can be LEADER, FOLLOWER, or READ_REPLICA.

### Example 72. Get an overview of the cluster

The following example shows how to explore the cluster topology.

CALL dbms.cluster.overview()

id	addresses	groups	databases
08eb9305-53b9-4394- 9237-0f0d63bb05d5	[bolt://neo20:7687, http://neo20:7474, https://neo20:7473]	D	{system: LEADER, neo4j: FOLLOWER}
cb0c729d-233c-452f-8f06- f2553e08f149	[bolt://neo21:7687, http://neo21:7474, https://neo21:7473]	D	{system: FOLLOWER, neo4j: FOLLOWER}
ded9eed2-dd3a-4574- bc08-6a569f91ec5c	[bolt://neo22:7687, http://neo22:7474, https://neo22:7473]		{system: FOLLOWER, neo4j: LEADER}
00000000-0000-0000-0000-0000-0000	[bolt://neo34:7687, http://neo34:7474, https://neo34:7473]	О	{system: READ_REPLICA, neo4j: READ_REPLICA}
00000000-0000-0000-0000- 00000000000	[bolt://neo28:7687, http://neo28:7474, https://neo28:7473]	П	{system: READ_REPLICA, neo4j: READ_REPLICA}
00000000-0000-0000-0000-0000-0000	[bolt://neo31:7687, http://neo31:7474, https://neo31:7473]	П	{system: READ_REPLICA, neo4j: READ_REPLICA}

### Get routing recommendations

From the application point of view it is not interesting to know about the role a member plays in the cluster. Instead, the application needs to know which instance can provide the wanted service. The procedure dbms.routing.getRoutingTable(routingContext, databaseName) provides this information.

### Syntax:

CALL dbms.routing.getRoutingTable(routingContext, databaseName)

### **Arguments:**

Name	Туре	Description
routingContext	Мар	The routing context used for multi- data center deployments. It should be used in combination with multi-data center load balancing.
databaseName	String	The name of the database to get the routing table for.

### Example 73. Get routing recommendations

The following example shows how discover which instances in the cluster can provide which services for database neo4j.

```
CALL dbms.routing.getRoutingTable({}, "neo4j")
```

The procedure returns a map between a particular service, READ, WRITE and ROUTE, and the addresses of instances that provide this service. It also returns a Time To Live (TTL) in seconds as a suggestion on how long the client could cache the response.

The result is not primarily intended for human consumption. Expanding it this is what it looks like.

### 12.6.2. Endpoints for status information

This section describes HTTP endpoints for monitoring the health of a Neo4j Causal Cluster.

A Causal Cluster exposes some HTTP endpoints which can be used to monitor the health of the cluster. In this section we will describe these endpoints and explain their semantics.

#### The section includes:

- Adjusting security settings for Causal Clustering endpoints
- Unified endpoints

### Adjusting security settings for Causal Clustering endpoints

If authentication and authorization is enabled in Neo4j, the Causal Clustering status endpoints will also require authentication credentials. The setting <a href="mailto:dbms.security.auth\_enabled">dbms.security.auth\_enabled</a> controls whether the native auth provider is enabled. For some load balancers and proxy servers, providing authentication

credentials with the request is not an option. For those situations, consider disabling authentication of the Causal Clustering status endpoints by setting

dbms.security.causal\_clustering\_status\_auth\_enabled=false in neo4j.conf.

### Unified endpoints

A unified set of endpoints exist, both on Core Servers and on Read Replicas, with the following behavior:

- /db/<databasename>/cluster/writable Used to direct write traffic to specific instances.
- /db/<databasename>/cluster/read-only Used to direct read traffic to specific instances.
- /db/<databasename>/cluster/available Available for the general case of directing arbitrary request types to instances that are available for processing read transactions.
- /db/<databasename>/cluster/status Gives a detailed description of this instance's view of its own status within the cluster. See Status endpoint for further details.

Every endpoint targets a specific database with its own Raft-group. The databaseName path parameter represents the name of the database. By default, a fresh Neo4j installation will have endpoints for two databases system and neo4j:

```
http://localhost:7474/db/system/cluster/writable
http://localhost:7474/db/system/cluster/read-only
http://localhost:7474/db/system/cluster/available
http://localhost:7474/db/system/cluster/status

http://localhost:7474/db/neo4j/cluster/writable
http://localhost:7474/db/neo4j/cluster/read-only
http://localhost:7474/db/neo4j/cluster/available
http://localhost:7474/db/neo4j/cluster/status
```

#### Table 38. Unified HTTP endpoint responses

Endpoint	Instance state	Returned code	Body text
	Leader	200 OK	true
/db/ <databasename>/cluster/writable</databasename>	Follower	404 Not Found	false
	Read Replica	404 Not Found	false
	Leader	404 Not Found	false
/db/ <databasename>/cluster/read-only</databasename>	Follower	200 OK	true
	Read Replica	200 OK	true

Endpoint	Instance state	Returned code	Body text
	Leader	200 OK	true
/db/ <databasename>/cluster/availab le</databasename>	Follower	200 OK	true
	Read Replica	200 OK	true
	Leader	200 OK	JSON - See Status endpoint for details.
/db/ <databasename>/cluster/status</databasename>	Follower	200 OK	JSON - See Status endpoint for details.
	Read Replica	200 OK	JSON - See Status endpoint for details.

Example 74. Use a Causal Clustering monitoring endpoint

From the command line, a common way to ask those endpoints is to use curl. With no arguments, curl will do an HTTP GET on the URI provided and will output the body text, if any. If you also want to get the response code, just add the -v flag for verbose output. Here are some examples:

• Requesting writable endpoint on a Core Server that is currently elected leader with verbose output:

```
#> curl -v localhost:7474/db/neo4j/cluster/writable
* About to connect() to localhost port 7474 (#0)
* Trying ::1...
* connected
* Connected to localhost (::1) port 7474 (#0)
> GET /db/neo4j/clustering/writable HTTP/1.1
> User-Agent: curl/7.24.0 (x86_64-apple-darwin12.0) libcurl/7.24.0 OpenSSL/0.9.8r zlib/1.2.5
> Host: localhost:7474
> Accept: */*
< HTTP/1.1 200 OK
< Content-Type: text/plain
< Access-Control-Allow-Origin: *
< Transfer-Encoding: chunked
< Server: Jetty(9.4.17)
* Connection #0 to host localhost left intact
true* Closing connection #0
```

### Status endpoint

Typically, you will want to have some guarantee that a Core is safe to shutdown before removing it from a cluster. The status endpoint provides the following information in order to help resolve such issues:

### Example 75. Example status response

```
{
    "lastAppliedRaftIndex":0,
    "votingMembers":["30edc1c4-519c-4030-8348-7cb7af44f591","80a7fb7b-c966-4ee7-88a9-35db8b4d68fe",
,"f9301218-1fd4-4938-b9bb-a03453e1f779"],
    "memberId":"80a7fb7b-c966-4ee7-88a9-35db8b4d68fe",
    "leader":"30edc1c4-519c-4030-8348-7cb7af44f591",
    "millisSinceLastLeaderMessage":84545,
    "participatingInRaftGroup":true,
    "core":true,
    "healthy":true,
    "raftCommandsPerSecond":124
}
```

Table 39. Status endpoint descriptions

Field	Туре	Optional	Example	Description
core	boolean	no	true	Used to distinguish between Core Servers and Read Replicas.
lastAppliedRaftIndex	number	no	4321	Every transaction in a cluster is associated with a raft index.  Gives an indication of what the latest applied raft log index is.
participatingInRaftGr oup	boolean	no	false	A participating member is able to vote. A Core is considered participating when it is part of the voter membership and has kept track of the leader.
votingMembers	string[]	no	[]	A member is considered a voting member when the leader has been receiving communication with it.  List of member's memberId that are considered
				part of the voting set by this Core.
healthy	boolean	no	true	Reflects that the local database of this member has not encountered a critical error preventing it from writing locally.
memberId	string	no	30edc1c4-519c-4030- 8348-7cb7af44f591	Every member in a cluster has it's own unique member id to identify it. Use memberId to distinguish between Core and Read Replica.
leader	string	yes	80a7fb7b-c966-4ee7- 88a9-35db8b4d68fe	Follows the same format as <pre>memberId</pre> , but if it is null or missing, then the leader is unknown.
millisSinceLastLeader Message	number	yes	1234	The number of milliseconds since the last heartbeat-like leader message. Not relevant to Read Replicas, and hence is not included.
raftCommandsPerSecond	number	yes	124	An estimate of the average Raft state machine throughput over a sampling windown configurable via causal_clustering.status_throughput_window setting.

In general, you will want to follow the pattern of first adding a new, updated instance, and then removing an old instance. After an instance has been switched on, you can access the status endpoint in order to make sure all the guarantees listed in the table below are met. This process can then be repeated until all old Cores have been removed.

Table 40. Measured values, accessed via the status endpoint

Name of check	Method of calculation	Description
allServersAreHealthy	Every Core's status endpoint indicates dbHealth==true.	We want to make sure the data across the entire cluster is healthy. Whenever any Cores are false that indicates a larger problem.
allVotingSetsAreEqual	For any 2 Cores (A and B), status endpoint A's votingMembers== status endpoint B's votingMembers.	When the voting begins, all the Cores are equal to each other, and you know all members agree on membership.
allVotingSetsContainA tLeastTargetCluster	For all Cores (S), excluding Core Z (to be switched off), every member in S contains S in their voting set. Membership is determined by using the memberId and votingMembers from the status endpoint.	Sometimes network conditions will not be perfect and it may make sense to switch off a different Core to the one we originally wanted to switch off. If you run this check for all Cores, the ones that match this condition can be switched off (providing other conditions are also met).
hasOneLeader	For any 2 Cores (A and B), A.leader == B.leader && leader!=null.	If the leader is different then there may be a partition (alternatively, this could also occur due to bad timing). If the leader is unknown, that means the leader messages have actually timed out.
noMembersLagging	For Core A with lastAppliedRaftIndex = min, and Core B with lastAppliedRaftIndex = max, B.lastAppliedRaftIndex- A.lastAppliedRaftIndex <raftindexlagthreshol d.</raftindexlagthreshol 	If there is a large difference in the applied indexes between Cores, then it could be dangerous to switch off a Core.

# 12.7. Monitoring individual database states

This section covers the use of SHOW DATABASES, and other related Cypher commands.

In addition to the system-wide metrics and logs described in previous sections, operators may wish to monitor the state of individual databases being hosted within a Neo4j instance. The SHOW DATABASES command may be used for this purpose.

# 12.7.1. Listing Databases

First ensure that you are executing queries against the system database, either by running the command :use system (if using the Cypher shell or Neo4j Browser) or by creating a session against the system database using a Neo4j driver. Subsequently, run the SHOW DATABASES command.

### Syntax:

SHOW DATABASES

Name	Туре	Description
name	String	The human-readable name of the database.
address	String	The bolt address of the Neo4j instance hosting the database.
role	String	The cluster role which the Neo4j instance fulfils for this database.

Name	Туре	Description
requestedStatus	String	The state that an operator has requested the database to be in.
currentStatus	String	The state the database is actually in on this Neo4j instance.
error	String	Error encountered by the Neo4j instance when transitioning the database to requestedStatus, if any.
default	String	Whether this database is the default for this DBMS.

When executing SHOW DATABASES against a standalone instance of Neo4j, you should see output like the following:

name	address	role	requestedStatu s	currentStatus	error	default
"neo4j"	"localhost:768 7"	"standalone"	"online"	"online"	1111	true
"system"	"localhost:768 7"	"standalone"	"online"	"online"	""	false

Note that the role and address columns are primarily intended to distinguish between the states of a given database, across multiple Neo4j instances deployed in a Causal Cluster. In a standalone deployment where you have a single Neo4j instance, your address field should be the same for every database, and your role field should always be "standalone".

If an error occurred whilst creating (or stopping, dropping etc.) a database, you should see output like the following:

name	address	role	requestedStatu s	currentStatus	error	default
"neo4j"	"localhost:768 7"	"standalone"	"online"	"online"	ш	true
"system"	"localhost:768 7"	"standalone"	"online"	"online"	ш	false
"foo"	"localhost:768 7"	"standalone"	"online"	"Reached maximum number of active databases"	1111	false

Note that for failed databases, the <u>currentStatus</u> and <u>requestedStatus</u> are different. This can imply an error. For example:

- a database may take a while to transition from "offline" to "online", due to performing recovery.
- during normal operation, the currentStatus of a database may be transiently different from its requestedStatus, due to a necessary automatic process, such as one Neo4j instance copying store files from another.

The possible statuses are "initial", "online", "offline", "store copying", and "unknown".

### Example 77. Listing databases in a Neo4j Causal Cluster

When running SHOW DATABASES against a Neo4j Causal Cluster you might see output like the following:

name	address	role	requestedStatu s	currentStatus	error	default
"neo4j"	"localhost:200 31"	"follower"	"online"	"online"	""	true
"neo4j"	"localhost:200 10"	"follower"	"online"	"online"	IIII	true
"neo4j"	"localhost:200 05"	"leader"	"online"	"online"	···	true

name	address	role	requestedStatu s	currentStatus	error	default
"neo4j"	"localhost:200 34"	"read_replica"	"online"	"online"	ш	true
"system"	"localhost:200 31"	"follower"	"online"	"online"	111	false
"system"	"localhost:200 10"	"follower"	"online"	"online"	ш	false
"system"	"localhost:200 05"	"leader"	"online"	"online"	1111	false
"system"	"localhost:200 34"	"read_replica"	"online"	"online"	1111	false
"foo"	"localhost:200 31"	"leader"	"online"	"online"	1111	false
"foo"	"localhost:200 10"	"follower"	"online"	"online"	1111	false
"foo"	"localhost:200 05"	"follower"	"online"	"online"	1111	false
"foo"	"localhost:200 34"	"read_replica"	"online"	"online"	1111	false

Note that SHOW DATABASES does **not** return 1 row per database. Instead, it returns 1 row per database, per Neo4j instance in the cluster. Therefore, if you have a 4-instance cluster, hosting 3 databases, you will have 12 rows.

If an error occurred whilst creating (or stopping, dropping etc.) a database, you should see output like the following:

name	address	role	requestedStatu s	currentStatus	error	default
"neo4j"	"localhost:200 31"	"follower"	"online"	"online"	III	true
"neo4j"	"localhost:200 10"	"follower"	"online"	"online"	nn	true
"neo4j"	"localhost:200 05"	"leader"	"online"	"online"	nn	true
"neo4j"	"localhost:200 34"	"read_replica"	"online"	"online"	nn	true
"system"	"localhost:200 31"	"follower"	"online"	"online"	III	false
"system"	"localhost:200 10"	"follower"	"online"	"online"	1111	false
"system"	"localhost:200 05"	"leader"	"online"	"online"	1111	false
"system"	"localhost:200 34"	"read_replica"	"online"	"online"	1111	false
"foo"	"localhost:200 31"	"unknown"	"online"	"initial"	1111	false
"foo"	"localhost:200 10"	"leader"	"online"	"online"	III	false
"foo"	"localhost:200 05"	"follower"	"online"	"online"	""	false
"foo"	"localhost:200 34"	"unknown"	"online"	"initial"	1111	false

Note that different instances may have different roles for each database.

If a database is offline on a particular Neo4j instance, either because it was stopped by an operator or an error has occurred, its cluster role is "unknown". This is because the cluster role of a given instance/database combination cannot be assumed in advance. This differs from standalone Neo4j instances, where the role of that instance for each database can always be assumed to be "standalone".

The possible roles are "standalone", "leader", "follower", "read\_replica", and "unknown".

# 12.7.2. Listing a single database

The number of rows returned by SHOW DATABASES can be quite large, especially when run in a cluster. You can filter the rows returned by database name (e.g. "foo") by using the command SHOW DATABASE foo.

### Syntax:

SHOW DATABASE databaseName

### **Arguments:**

Name	Туре	Description
databaseName	$\mathbf{c}$	The name of the database whose status to report

Name	Туре	Description
name	String	The human-readable name of the database.
address	String	The bolt address of the Neo4j instance hosting the database.
role	String	The cluster role which the Neo4j instance fulfils for this database.
requestedStatus	String	The state that an operator has requested the database to be in.
currentStatus	String	The state the database is actually in on this Neo4j instance.
error	String	Error encountered by Neo4j instance when transitioning the database to requestedStatus, if any.

### Example 78. Listing statuses for database foo

# When running $SHOW\ DATABASE\ foo\ in\ a\ Neo4j\ Causal\ Cluster,\ you\ should\ see\ output\ like\ the\ following:$

name	address	role	requestedStatu s	currentStatus	error	default
"foo"	"localhost:200 31"	"unknown"	"online"	"initial"	ш	false
"foo"	"localhost:200 10"	"leader"	"online"	"online"	ш	false
"foo"	"localhost:200 05"	"follower"	"online"	"online"	ш	false
"foo"	"localhost:200 34"	"unknown"	"online"	"initial"	ш	false

# Chapter 13. Performance

This chapter describes factors that affect operational performance, and how to tune Neo4j for optimal throughput.

The topics described in this chapter are:

- Memory configuration How to configure memory settings for efficient operations.
- Index configuration How to configure indexes.
- Garbage collector How to configure the Java Virtual Machine's garbage collector.
- Bolt thread pool configuration How to configure the Bolt thread pool.
- Linux file system tuning How to configure the Linux file system.
- Disks, RAM and other tips Disks, RAM and other tips.
- Statistics and execution plans How schema statistics and execution plans affect Cypher query performance.

# 13.1. Memory configuration

This section describes the different aspects of Neo4j memory configuration and use.

### 13.1.1. Overview

Consider the image below. As you can see, the RAM of the Neo4j server has a number of usage areas, with some sub-areas:

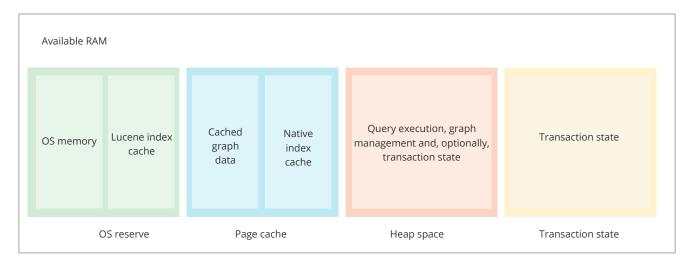


Figure 12. Neo4j memory management

### **OS** memory

Some memory must be reserved for running the processes of the operating system itself. It is not possible to explicitly configure the amount of RAM that should be reserved for the operating system, as this is what RAM remains available after configuring page cache and heap space. However, if we do not leave enough space for the OS it will start swapping to disk, which will heavily affect performance.

1GB is a good starting point for a server that is dedicated to running Neo4j. However, there are cases where the amount reserved for the OS is significantly larger than 1GB, such as servers with

exceptionally large RAM.

#### Lucene index cache

Neo4j uses Apache Lucene for some of its indexing functionality. Index lookup performance is optimized by ensuring that as much as possible of the indexes are cached into memory. Similar to the OS memory, the Lucene index cache can not be explicitly configured. Instead we estimate the memory needed and make sure that there is enough headroom left for Lucene indexes after the page cache and heap cache have been assigned.

### Page cache

The page cache is used to cache the Neo4j data and native indexes. The caching of graph data and indexes into memory will help avoid costly disk access and result in optimal performance.

The parameter for specifying how much memory Neo4j is allowed to use for the page cache is: dbms.memory.pagecache.size.

### Heap size

The heap space is used for query execution, transaction state, management of the graph etc. The size needed for the heap is very dependent on the nature of the usage of Neo4j. For example, long-running queries, or very complicated queries, are likely to require a larger heap than simpler queries.

Generally speaking, in order to aid performance, we want to configure a large enough heap to sustain concurrent operations.

In case of performance issues we may have to tune our queries, and monitor their memory usage, in order to determine whether the heap needs to be increased.

The heap memory size is determined by the parameters dbms.memory.heap.initial\_size and dbms.memory.heap.max\_size. It is recommended to set these two parameters to the same value. This will help avoid unwanted full garbage collection pauses.

#### **Transaction state**

Transaction state is the memory that is needed to hold data and intermediate results in transactions that update records in the database. Queries that only read data do not require transaction state memory allocation. By default, transaction state is allocated off-heap. When the transaction state is allocated off-heap, the maximum size of the transaction state can be defined using the parameter <a href="max.tx\_state.max\_off\_heap\_memory">dbms.tx\_state.max\_off\_heap\_memory</a>. Note that the transaction state memory is not pre-allocated; it will grow and shrink as needed by the activity in the database. Keeping transaction state off-heap is particularly beneficial to applications characterized by large, write-intensive transactions.

Transaction state can also be configured to be allocated on-heap, by using the parameter dbms.tx\_state.memory\_allocation. Note that when the transaction state is configured on-heap, its maximum size cannot be specified.

### 13.1.2. Considerations

### Always use explicit configuration

In order to have good control of a system's behavior, it is recommended that you always define the page cache and heap size parameters explicitly in *neo4j.conf*. If these parameters are not explicitly defined, some heuristic values will be computed at startup based on available system resources.

### *Initial memory recommendation*

Use the neo4j-admin memrec command to get an initial recommendation for how to distribute a certain amount of memory. The values may need to be adjusted to cater for each specific use case.

Inspect the memory settings of a database

The neo4j-admin memrec --database command is useful for inspecting the current distribution of data and indexes in an existing database.

### Example 79. Use neo4j-admin memrec to inspect the memory settings of a database

We wish to estimate the total size of the database files.

```
$neo4j-home> bin/neo4j-admin memrec --database=neo4j
...
...
# Lucene indexes: 6690m
# Data volume and native indexes: 17050m
```

We can see that the Lucene indexes take up approximately 6.7GB of data, and that the data volume and native indexes combined take up approximately 17GB.

Using this information, we can do a sanity check of our memory configuration:

- Compare the value for data volume and native indexes to the value of dbms.memory.pagecache.size.
- For cases when *off-heap* transaction state is used, estimate transactional workload and how much memory is left to the value of dbms.tx\_state.max\_off\_heap\_memory.
- Compare the value for Lucene indexes to how much memory is left after assigning dbms.memory.pagecache.size and dbms.memory.heap.initial\_size.

Note that even though we strive for caching as much of our data and indexes as possible, in some production systems the access to memory is limited and must be negotiated between different areas. Then there will will be a certain amount of testing and tuning to figure out the optimal division of the available memory.

# 13.1.3. Capacity planning

In many use cases, it is advantageous to try to cache as much of the data and indexes as possible. The following examples illustrate methods for estimating the page cache size, depending on whether we are already running in production or planning for a future deployment:

First estimate the total size of data and indexes, and then multiply with some factor, for example 20%, to allow for growth.

```
$neo4j-home> bin/neo4j-admin memrec --database=neo4j
...
...
# Lucene indexes: 6690m
# Data volume and native indexes: 35050m
```

We can see that the data volume and native indexes combined take up approximately 35GB. In our specific use case we estimate that 20% will provide sufficient head room for growth.

```
dbms.memory.pagecache.size = 1.2 * (35GB) = 42GB
```

We configure the page cache by adding the following to *neo4j.conf*:

```
dbms.memory.pagecache.size=42GB
```

### Example 81. Estimate page cache for a new Neo4j database

When planning for a future database, it is useful to run an import with a fraction of the data, and then multiply the resulting store size by that fraction plus some percentage for growth. For example, import 1/100th of the data and measure its data volume and native indexes. Then multiply that number by 120 to size up the result, and allow for 20% growth.

Assume that we have imported 1/100th of the data into a test database.

```
$neo4j-home> bin/neo4j-admin memrec --database=neo4j
...
...
# Lucene indexes: 425.0
# Data volume and native indexes: 251100k
```

We can see that the data volume and native indexes combined take up approximately 250MB. We size up the result and additionally reserve 20% for growth:

```
dbms.memory.pagecache.size = 120 * (250MB) = 30GB
```

We configure the page cache by adding the following to neo4j.conf:

```
dbms.memory.pagecache.size=30G
```

### 13.1.4. Configure query heap usage

When running a Cypher query, Neo4j will utilize the heap internally for storing results. It may be difficult to predict how much memory a query needs and if a query ends up using too much memory it could severely hamper the overall performance of the database.

There are two settings that can be enabled in *neo4j.conf* that can help improve the heap utilization of Neo4j:

```
dbms.track_query_allocation=true
cypher.query_max_allocations.size=1G
```

If the former setting is enabled, Neo4j will be tracking the heap utilization of all Cypher queries. You can view the utilization of running queries by calling:

```
CALL dbms.listQueries()
```

Or alternatively, you can enable dbms.logs.query.allocation\_logging\_enabled and the memory usage of each query will be logged in the query log.

By setting the latter configuration, cypher.query\_max\_allocations.size, you can limit the amount of memory each query can use. Whenever that limit is reached, the query will be gracefully terminated without affecting the overall health of the database.



The heap-usage of query is only an estimate and the actual heap utilization might be slightly bigger or slightly smaller than the estimated value.

# 13.2. Index configuration

This section describes how to configure indexes to enhance performance in search, and to enable full-text search.

This section contains the following:

- Introduction
- B-tree indexes
  - Limitations
    - ☐ Limitations for queries using CONTAINS and ENDS WITH
    - ☐ Limitations on key size
    - Workarounds to address limitations
  - □ Index migration
  - Procedures to create index and index backed constraint
- Full-text indexes
  - □ Configuration

### 13.2.1. Introduction

In Neo4j there are two different index types: b-tree and full-text.

B-tree indexes can be created and dropped using Cypher. Users typically do not have to know about the index in order to use it, since Cypher's query planner decides which index to use in which situation. B-tree indexes are good at exact look-ups on all types of values, and range scans, full scans, and prefix searches.

For details on the configuration aspects of b-tree indexes, see B-tree indexes.

Full-text indexes differ from b-tree indexes, in that they are optimised for indexing and searching text. They are used for queries that demand an understanding of language, and they only index string data. They must also be queried explicitly via procedures, as Cypher will not make plans that rely on

them.

An example of a use case for full-text indexes is parsing a book for a certain term, and taking advantage of the knowledge that the book is written in a certain language. The use of an *analyzer* for that language will, among other things, enable you to exclude words that are not relevant for the search (for example "if" and "and"), and include conjugations of words that are.

Another use case example is indexing the various address fields and text data in a corpus of emails. Indexing this data using the email analyzer would enable someone to find all emails that are sent from, or to, or mentions, an email account.

In contrast to b-tree indexes, full-text indexes are created, queried, and dropped using built-in procedures. The use of full-text indexes do require a familiarity with how the indexes operate.

For details on the configuration aspects of full-text indexes, see Full-text indexes.

For details on creating, querying and dropping full-text indexes, see Cypher Manual 

Indexes to support full-text search.

The type of an index can be identified according to the table below:

Index type	Procedure	Core API
B-tree index	db.indexes#BTREE	org.neo4j.graphdb.schema.IndexType#B TREE
Full-text index	db.indexes#FULLTEXT	org.neo4j.graphdb.schema.IndexType#F ULLTEXT

## 13.2.2. B-tree indexes

B-tree indexes can be backed by two different index providers, <a href="native-btree-1.0">native-btree-1.0</a> and the deprecated <a href="lucene+native-3.0">lucene+native-3.0</a>. If not explicitly set, <a href="native-btree-1.0">native-btree-1.0</a> will be used.

For more information on the different index types, refer to Cypher Manual Indexes.

Deprecated index providers

Index provider <u>lucene+native-3.0</u> has been deprecated, and will be removed in a future release.



The recommended index provider to use is native-btree-1.0.

The only reason to use a deprecated provider should be due to the limitations, as described in Limitations for queries using CONTAINS and ENDS WITH or Limitations on key size.

## Limitations

In this section a few limitations for b-tree indexes are described, together with suggested workarounds.

Limitations for queries using CONTAINS and ENDS WITH

The index provider native-btree-1.0 has limited support for ENDS WITH and CONTAINS queries. These queries will not be able to do an optimized search as per queries that use STARTS WITH, =, and <>. Instead, the index result will be a stream of an index scan with filtering.

In the future, ENDS WITH and CONTAINS queries will be supported with full-text indexes, but for now the

deprecated index provider <u>lucene+native-3.0</u> can be used instead. Please note that <u>lucene+native-3.0</u> only has support for <u>ENDS WITH and CONTAINS</u> for single property strings.

- For details about execution plans, refer to Cypher Manual 

  Execution plans.
- For details about string operators, refer to Cypher Manual 

  Operators.

## Limitations on key size

The index provider native-btree-1.0 has a key size limit of around 8kB.

If a transaction reaches the key size limit for one or more of its changes, that transaction will fail before committing any changes. If the limit is reached during index population, the resulting index will be in a failed state, and as such will not be usable for any queries.

If this is an issue, you can use the deprecated index provider <a href="lucene+native-3.0">lucene+native-3.0</a> instead. This provider has a key size limit for single property strings of around 32kB.

Please note that <a href="lucene+native-3.0">lucene+native-3.0</a> has been deprecated and will be removed in the future, whereby 8kB will be the key size limit for b-tree indexes. The recommended option for such large values is to use full-text index.

#### Workarounds to address limitations

To workaround problems with key size, or performance issues related to ENDS WITH or CONTAINS, you can use the deprecated index provider <a href="lucene+native-3.0">lucene+native-3.0</a>. This only works for single-property string indexes.

This can be done using either of the following methods:

## Option 1. Use a built-in procedure (recommended)

There are built-in procedures that can be used to specify index provider on index creation, unique property constraint creation, and node key creation (for details on constraints, see Cypher manual Constraints.

For more information, see Built-in procedures.

## Option 2. Change the config

- 1. Configure the setting dbms.index.default\_schema\_provider to the one required.
- 2. Restart Neo4j.
- 3. Drop and recreate the relevant index.
- 4. Change dbms.index.default\_schema\_provider back to the original value.
- 5. Restart Neo4j.



Please note that the index setting <a href="mailto:default\_schema\_provider">default\_schema\_provider</a> has been deprecated, and will be removed in a future release. It will be a fully internal concern which index provider an index is using.

The recommended way to set index provider for an index is to use the built in procedures for index creation, unique property constraint creation, and node key creation.

For more information, see Built-in procedures

## Index migration

When upgrading a 3.5 store to 4.0.0, all indexes will be upgraded to the latest index version, and rebuilt automatically, with the exception for the indexes that were previously using Lucene for single-property strings. They will be upgraded to a fallback version which still use Lucene for those properties. Please note that they will still need to be rebuilt.

The table below shows the migration mapping:

Index provider in 3.5	Index provider in 4.0
native-btree-1.0	native-btree-1.0
lucene+native-2.0	native-btree-1.0
lucene+native-1.0	lucene+native-3.0
lucene-1.0	lucene+native-3.0

The caching of indexes takes place in different memory areas for different index providers. See Memory configuration. It can be useful to run neo4j-admin memrec --database before and after the rebuilding of indexes, and adjust memory settings in accordance with the findings.

## Procedures to create index and index backed constraint.

Indexes and constraints are best created through Cypher, but when these indexes or constraints need to be more specifically configured than what is possible through Cypher, then you can use the procedures described in the example below.

## Example 82. Example of procedures to create index and index backed constraint

The following procedures provide the option to specify both index provider and index settings (optional). Note that settings keys need to be escaped with back-ticks if they contain dots.

Use db. createIndex procedure to create an index:

```
CALL db.createIndex("MyIndex", ["Person"], ["name"], "native-btree-1.0", {+`spatial.cartesian.max+`: [100.0,100.0], +`spatial.cartesian.min+`: [-100.0,-100.0]})
```

If a settings map is not provided, the settings will be picked up from the Neo4j config file, the same way as when creating an index or constraint through Cypher.

```
CALL db.createIndex("MyIndex", ["Person"], ["name"], "native-btree-1.0")
```

Use db.createUniquePropertyConstraint to create a node property uniqueness constraint (the example is without settings map, left out for abbreviation):

```
CALL db.createUniquePropertyConstraint("MyIndex", ["Person"], ["name"], "native-btree-1.0", {+`spatial.cartesian.max+`: [100.0,100.0], +`spatial.cartesian.min+`: [-100.0,-100.0]})
```

Use <a href="db.createNodeKey">db.createNodeKey</a> to create node key constraint (the example is without settings map, left out for abbreviation):

```
CALL db.createNodeKey("MyIndex", ["Person"], ["name"], "native-btree-1.0", {+`spatial.cartesian.max+`: [100.0,100.0], +`spatial.cartesian.min+`: [-100.0,-100.0]})
```

## 13.2.3. Full-text indexes

Full-text indexes are powered by the Apache Lucene indexing and search library. A full-text index enables you to write queries that matches within the *contents* of indexed string properties. A full description on how to create and use full-text indexes is provided in the Cypher Manual [] Indexes to support full-text search.

## Configuration

The following options are available for configuring full-text indexes:

#### dbms.index.fulltext.default\_analyzer

The name of the analyzer that the full-text indexes should use by default. This setting only has effect when a full-text index is created, and will be remembered as an index-specific setting from then on.

The list of possible analyzers is available through the db.index.fulltext.listAvailableAnalyzers() Cypher procedure.

Unless otherwise specified, the default analyzer is standard-no-stop-words, which is the same as the StandardAnalyzer from Lucene, except no stop-words are filtered out.

#### dbms.index.fulltext.eventually\_consistent

Used to declare whether full-text indexes should be eventually consistent, or not. This setting only has effect when a full-text index is created, and will be remembered as an index-specific setting from then on.

Indexes are normally fully consistent, and the committing of a transaction does not return until both the store and the indexes have been updated. Eventually consistent full-text indexes, on the other hand, are not updated as part of commit, but instead have their updates queued up and applied in a background thread. This means that there can be a short delay between committing a change, and that change becoming visible via any eventually consistent full-text indexes. This delay is just an artifact of the queueing, and will usually be quite small since eventually consistent indexes are updated "as soon as possible".

By default, this is turned off, and full-text indexes are fully consistent.

#### dbms.index.fulltext.eventually\_consistent\_index\_update\_queue\_max\_length

Eventually consistent full-text indexes have their updates queued up and applied in a background thread, and this setting determines the maximum size of that update queue. If the maximum queue size is reached, then committing transactions will block and wait until there is more room in the queue, before adding more updates to it.

This setting applies to all eventually consistent full-text indexes, and they all use the same queue. The maximum queue length must be at least 1 index update, and must be no more than 50 million due to heap space usage considerations.

The default maximum queue length is 10.000 index updates.

## 13.3. Tuning of the garbage collector

This section discusses the effect of the Java Virtual Machine's garbage collector with regards to Neo4j performance.

The heap is separated into an *old generation* and a *young generation*. New objects are allocated in the young generation, and then later moved to the old generation, if they stay live (in use) for long

enough. When a generation fills up, the garbage collector performs a collection, during which all other threads in the process are paused. The young generation is quick to collect since the pause time correlates with the *live set* of objects. In the old generation, pause times roughly correlates with the size of the heap. For this reason, the heap should ideally be sized and tuned such that transaction and query state never makes it to the old generation.

The heap size is configured with the <code>dbms.memory.heap.max\_size</code> (in MBs) setting in the <code>neo4j.conf</code> file. The initial size of the heap is specified by the <code>dbms.memory.heap.initial\_size</code> setting, or with the <code>-Xms???m</code> flag, or chosen heuristically by the JVM itself if left unspecified. The JVM will automatically grow the heap as needed, up to the maximum size. The growing of the heap requires a full garbage collection cycle. It is recommended to set the initial heap size and the maximum heap size to the same value. This way the pause that happens when the garbage collector grows the heap can be avoided.

If the new generation is too small, short-lived objects may be moved to the old generation too soon. This is called premature promotion and will slow the database down by increasing the frequency of old generation garbage collection cycles. If the new generation is too big, the garbage collector may decide that the old generation does not have enough space to fit all the objects it expects to promote from the new to the old generation. This turns new generation garbage collection cycles into old generation garbage collection cycles, again slowing the database down. Running more concurrent threads means that more allocations can take place in a given span of time, in turn increasing the pressure on the new generation in particular.



The *Compressed OOPs* feature in the JVM allows object references to be compressed to use only 32 bits. The feature saves a lot of memory but is only available for heaps up to 32 GB. The maximum applicable size varies from platform and JVM version. The -XX:+UseCompressedOops option can be used to verify whether the system can use the *Compressed OOPs* feature. If it cannot, this will will be logged in the default process output stream.

How to tune the specific garbage collection algorithm depends on both the JVM version and the workload. It is recommended to test the garbage collection settings under realistic load for days or weeks. Problems like heap fragmentation can take a long time to surface.

To gain good performance, these are the things to look into first:

- Make sure the JVM is not spending too much time performing garbage collection. The goal is to have a large enough heap to make sure that heavy/peak load will not result in so called GC-trashing. Performance can drop as much as two orders of magnitude when GC-trashing happens. Having too large heap may also hurt performance so you may have to try some different heap sizes.
- Neo4j needs enough heap memory for the transaction state and query processing, plus some head-room for the garbage collector. As heap memory requirements are so workload-dependent, it is common to see heap memory configurations from 1 GB, up to 32 GB.

Edit the following properties:

Table 41. neo4j.conf JVM tuning properties

Property Name	Meaning
dbms.memory.heap.initial_size	initial heap size (in MB)
dbms.memory.heap.max_size	maximum heap size (in MB)
dbms.jvm.additional	additional literal JVM parameter

# 13.4. Bolt thread pool configuration

This section discusses the thread pool infrastructure built into Bolt connectors, and how it can be configured.

The Bolt connector is backed by a thread pool on the server side. The thread pool is constructed as part of the server startup process.

# 13.4.1. How thread pooling works

The Bolt thread pool has a minimum and a maximum capacity. It starts with a minimum number of threads available, and grows up to the maximum count depending on the workload. Threads that sit idle for longer than a specified time period are stopped and removed from the pool in order to free up resources. However, the size of the pool will never go below the minimum.

Each connection being established is assigned to the connector's thread pool. Idle connections do not consume any resources on the server side, and they are monitored against messages arriving from the client. Each message arriving on a connection triggers the scheduling of a connection on an available thread in the thread pool. If all the available threads are busy, and there is still space to grow, a new thread is created and the connection is handed over to it for processing. If the pool capacity is filled up, and no threads are available to process, the job submission is rejected and a failure message is generated to notify the client of the problem.

The default values assigned to the Bolt thread pool will fit most workloads, so it is generally not necessary to configure the connection pool explicitly. If the maximum pool size is set too low, an exception will be thrown with an error message indicating that there are no available threads to serve. The message will also be written to neo4j.log.



Any connection with an active explicit, or implicit, transaction will stick to the thread that starts the transaction, and will not return that thread to the pool until the transaction is closed. Therefore, in applications that are making use of explicit transactions, it is important to close the transactions appropriately. To learn more about transactions, refer to the Neo4j Driver Manual.

## 13.4.2. Configuration options

The following configuration options are available for configuring the Bolt connector:

Table 42. Thread pool options

Option name	Default	Description
<pre>dbms.connector.bolt.thread_pool_min_ size</pre>	5	The minimum number of threads that will always be up even if they are idle.
<pre>dbms.connector.bolt.thread_pool_max_ size</pre>	400	The maximum number of threads that will be created by the thread pool.
<pre>dbms.connector.bolt.thread_pool_keep _alive</pre>	5m	The duration that the thread pool will wait before killing an idle thread from the pool. However, the number of threads will never go below dbms.connector.bolt.thread_pool_min_size.

## 13.4.3. How to size your Bolt thread pool

Select values for thread pool sizing based on your workload. Since each active transaction will borrow a thread from the pool until the transaction is closed, it is basically the minimum and maximum active

transaction at any given time that determine the values for pool configuration options. You can use the monitoring capabilities (see Monitoring) of the database to discover more about your workload.

Configure dbms.connector.bolt.thread\_pool\_min\_size based on your minimum or average workload. Since there will always be this many amount of threads in the thread pool, sticking with lower values may be more resource-friendly than having too many idle threads waiting for job submissions.

Configure dbms.connector.bolt.thread\_pool\_max\_size based on your maximum workload. This should basically be set after the maximum number of active transactions that is expected on the server. You should also account for non-transaction operations that will take place on the thread pool, such as connection and disconnection of clients.

Example 83. Configure the thread pool for a Bolt connector

In this example we configure the Bolt thread pool to be of minimum size 5, maximum size 100, and have a keep-alive time of 10 minutes.

```
dbms.connector.bolt.thread_pool_min_size=10
dbms.connector.bolt.thread_pool_max_size=100
dbms.connector.bolt.thread_pool_keep_alive=10m
```

# 13.5. Linux file system tuning

This section covers Neo4j I/O behavior, and how to optimize for operations on disk.

Databases often produce many small and random reads when querying data, and few sequential writes when committing changes. For maximum performance, it is recommended to store database and transaction logs on separate physical devices.

Often, recommended practice is to disable file and directory access time updates. This way, the file system won't have to issue writes that update this meta-data, thus improving write performance.

Since databases can put a high and consistent load on a storage system for a long time, it is recommended to use a file system that has good aging characteristics. The EXT4 and XFS file systems are recommended as a first choice.

A high read and write I/O load can also degrade SSD performance over time. The first line of defense against SSD wear is to ensure that the working dataset fits in RAM. A database with a high write workload will, however, still cause wear on SSDs. The simplest way to combat this is to over-provision; use SSDs that are at least 20% larger than you strictly need them to be.

To be able to achieve optimum performance, we do not recommend the use of NFS or NAS as database storage.

For more detailed and precise configuration for your operating system and hardware, please consult corresponding manuals.

# 13.6. Disks, RAM and other tips

This section provides an overview of performance considerations for disk and RAM when running Neo4j.

As with any persistence solution, performance depends a lot on the persistence media used. In general, the faster storage you have, and the more of your data you can fit in RAM, the better

performance you will get.

## 13.6.1. Storage

If you have multiple disks or persistence media available it may be a good idea to divide the store files and transaction logs across those disks. Keeping the store files on disks with low seek time can do wonders for read operations.

To achieve maximum performance, it is recommended to provide Neo4j with as much RAM as possible in order to avoid hitting the disk.

Use tools like dstat or vmstat to gather information when your application is running. If the swap or paging numbers are high, that is a sign that the database don't quite fit in memory. In this case, database access can have high latencies.

# 13.6.2. Page cache

When Neo4j starts up, its page cache is empty and needs to warm up. The pages, and their graph data contents, are loaded into memory on demand as queries need them. This can take a while, especially for large stores. It is not uncommon to see a long period with many blocks being read from the drive, and high IO wait times. This will show up in the page cache metrics as an initial spike in page faults. The page fault spike is then followed by a gradual decline of page fault activity, as the probability of queries needing a page that is not yet in memory drops.

## Active page cache warmup

The Neo4j Enterprise Edition has a feature called *active page cache warmup*, which shortens the page fault spike and makes the page cache warm up faster. This is done by periodically recording *cache profiles* of the store files, as the database is running. These profiles contain information about what data is and is not in memory, and are stored in a "profiles" sub-directory of the store directory. When Neo4j is later restarted, it will look for these cache profiles and eagerly load in the same data that was in memory when the profile was created. The profiles are also copied as part of online backup and cluster store-copy, and helps warm up new database instances that join a cluster.

It is also possible to configure *page cache warmup* to prefetch database data. To enable it, you can use the <a href="mailto:dbms.memory.pagecache.warmup.preload">dbms.memory.pagecache.warmup.preload</a> setting. Data prefetching will disregard the cache profile files, and will prefetch data into a page cache as part of the startup. This can be quite useful in scenarios where the database size is smaller than the page cache, and you want to make sure that data is loaded before it will be required.

## Checkpoint IOPS limit

Neo4j flushes its page cache in the background as part of its checkpoint process. This will show up as a period of elevated write IO activity. If the database is serving a write-heavy workload, the checkpoint can slow the database down by reducing the IO bandwidth that is available to query processing. Running the database on a fast SSD, which can service a lot of random IOs, significantly reduces this problem. If a fast SSD is not available in your environment, or if it is insufficient, then an artificial IOPS limit can be placed on the checkpoint process. The <a href="mailto:dbms.checkpoint.iops.limit">dbms.checkpoint.iops.limit</a> restricts the IO bandwidth that the checkpoint process is allowed to use. Each IO is, in the case of the checkpoint process, an 8 KiB write. An IOPS limit of 300, for instance, would thus only allow the checkpoint process to write at a rate of roughly 2.5 MiB per second. This will, on the other hand, make checkpoints take longer to complete. A longer time between checkpoints can cause more transaction log data to accumulate, and can lengthen recovery times. See the transaction logs section for more details on the relationship between checkpoints and log pruning. The IOPS limit can be changed at runtime, making it possible to tune it until you have the right balance between IO usage and checkpoint time.

# 13.7. Statistics and execution plans

This section describes the configuration options that affect the gathering of statistics, and the replanning of query plans in the Cypher query engine.

When a Cypher query is issued, it gets compiled to an execution plan that can run and answer the query. The Cypher query engine uses available information about the database, such as schema information about which indexes and constraints exist in the database. Neo4j also uses statistical information about the database to optimize the execution plan.

For further details, please see Cypher Manual 

Query tuning and Cypher Manual 

Execution plans.

The frequency of statistics gathering and the replanning of execution plans are described in the sections below.

## 13.7.1. Statistics

The statistical information that Neo4j keeps is:

- 1. The number of nodes having a certain label.
- 2. The number of relationships by type.
- 3. The number of relationships by type, ending or starting from a node with a specific label.
- 4. Selectivity per index.

Neo4j keeps the statistics up to date in two different ways. For label and relationship counts, the number is updated whenever you set or remove a label from a node. For indexes, however, Neo4j needs to scan the full index to produce the selectivity number. Since this is potentially a very time-consuming operation, these numbers are collected in the background when enough data on the index has been changed.

The following settings allow you to control whether statistics are collected automatically, and at which rate:

Parameter name	Default value	Description
<pre>dbms.index_sampling.background_enabl ed</pre>	true	Controls whether indexes will automatically be re-sampled when they have been updated enough. The Cypher query planner depends on accurate statistics to create efficient plans, so it is important it is kept up to date as the database evolves.
<pre>dbms.index_sampling.update_percentag e</pre>	5	Controls the percentage of the index that has to have been updated before a new sampling run is triggered.

It is possible to manually trigger index sampling, using the following built-in procedures:

### db.resampleIndex()

Trigger resampling of an index.

### db.resampleOutdatedIndexes()

Trigger resampling of all outdated indexes.

The following example illustrates how to trigger a resampling of the index on the label Person and property name, by calling db.resampleIndex():

```
CALL db.resampleIndex(":Person(name)");
```

The following example illustrates how to call <a href="db.resampleOutdatedIndexes">db.resampleOutdatedIndexes</a>() in order to trigger a resampling of all outdated indexes:

```
CALL db.resampleOutdatedIndexes();
```

## 13.7.2. Execution plans

Execution plans are cached and will not be replanned until the statistical information used to produce the plan has changed. The following setting enables you to control how sensitive replanning should be to updates of the database:

Parameter name	Default value	Description
<pre>cypher.statistics_divergence_thresho ld</pre>	0.75	The threshold when a plan is considered stale. If any of the underlying statistics used to create the plan have changed more than this value, the plan will be considered stale and will be replanned. Change is calculated as abs(a-b)/max(a,b). This means that a value of 0.75 requires the database to approximately quadruple in size before replanning occurs. A value of 0 means replan as soon as possible, with the soonest being defined by the parameter cypher.min_replan_interval, which defaults to 10s. After this interval the divergence threshold will slowly start to decline, reaching 10% after about 7h. This will ensure that long-running databases will still get query replanning on even modest changes, while not replanning frequently unless the changes are very large.

It is possible to manually force the database to replan execution plans that are already in the cache, using the following built-in procedures:

## db.clearQueryCaches()

Clears out all query caches, but does not change the database statistics.

```
CALL db.clearQueryCaches();
```

### db.prepareForReplanning()

Completely recalculates all database statistics, so that they can be used for any subsequent query planning.

It triggers an index resampling and waits for it to complete, and after that it also clears all query caches. After this procedure has finished queries will be planned on empty caches using the latest database statistics.

CALL db.prepareForReplanning();

For more information, see Built-in procedures.

# Chapter 14. Tools

This chapter describes the Neo4j tools Neo4j Admin and Cypher Shell, and explains how to use commands to import data into a new database, dump and load an existing database, and how to check consistency.

## This chapter comprises the following topics:

- A description of the Neo4j Admin tool
- How to check the consistency of a Neo4j database using Neo4j Admin
- How to collect the most common information needed for remote assessments
- · How to display information about a database store
- How to get an initial recommendation for Neo4j memory settings
- · How to import data into Neo4j using the import tool
  - □ Syntax
  - CSV header format
  - Options
- How to dump and load Neo4j databases using Neo4j Admin
- How to remove cluster state data from a Neo4j server.
- How to copy data from an existing database to a new database.
- How to use the Cypher Shell

# 14.1. Neo4j Admin

This section describes the Neo4j Admin tool.

## This section describes the following:

- Introduction
- Syntax and commands
- Environment variables
- Exit codes

## 14.1.1. Introduction

Neo4j Admin is the primary tool for managing your Neo4j instance. It is a command-line tool that is installed as part of the product and can be executed with a number of commands. Some of the commands are described in more detail in separate sections.

## 14.1.2. Syntax and commands

#### **Syntax**

Neo4j Admin is located in the *bin* directory and is invoked as:

neo4j-admin <command>

## **Commands**

Functionality area	Command	Description
General	help help <command/>	Display help text.
		Display help text for <command/>
	check-consistency	Check the consistency of a database. For details, see Consistency checker.
	report	Collect the most common information needed for remote assessments. For details, see Report tool.
	store-info	Print information about a Neo4j database store. For details, see Display store information.
	memrec	Print Neo4j heap and pagecache memory settings recommendations. For details, see Memory recommendations.
	import	Import from a collection of CSV files or a pre-3.0 database. For details, see Import.
Authentication	set-default-admin	Sets the default admin user when no roles are present.
	set-initial-password	Sets the initial password of the initial admin user (neo4j). For details, see Set an initial password.
Offline backup  For details see Dump and load databases.	dump	Dump a database into a single-file archive.
	load	Load a database from an archive created with the dump command.
Online backup	backup	Perform an online backup from a running Neo4j server.
For details see Backup.	restore	Restore a backed-up database.
Clustering	unbind	Removes cluster state data from a stopped Neo4j server.
		For details, see Unbind a Core Server.

## Limitations

neo4j-admin must be invoked as the neo4j user in order to ensure the appropriate file permissions.

## 14.1.3. Environment variables

Neo4j Admin can utilize the following environment variables:

NEO4J\_DEBUG

Set to anything to enable debug output.

NEO4J\_HOME

Neo4j home directory.

NEO4J\_CONF

Path to directory which contains neo4j.conf.

HEAP\_SIZE

Set JVM maximum heap size during command execution. Takes a number and a unit, for example 512m.

## JAVA\_OPTS

Additional JVM arguments.

## 14.1.4. Exit codes

If neo4j-admin finishes as expected it will return an exit code of 0. A non-zero exit code means something undesired happened during command execution. The non-zero exit code can contain further information about the error, such as the backup command's exit codes.

# 14.2. Consistency checker

This section describes the Neo4j consistency checker.

The consistency of a database or a backup can be checked using the <a href="heck-consistency">check-consistency</a> argument to the <a href="heck-consistency">neo4j-admin</a> tool.

## 14.2.1. Check consistency of a database or a backup

The neo4j-admin tool is located in the bin directory. Run it with the check-consistency argument in order to check the consistency of a database.

## **Syntax**

neo4j-admin check-consistency ([--database=<database>] | [--backup=<path>]) [--verbose] [-additional-config=<path>] [--check-graph=<true/false>] [--check-indexes=<true/false>] [--check-index-structure=<true/false>] [--check-label-scan-store=<true/false>] [--check-propertyowners=<true/false>] [--report-dir=<path>]

## **Options**

Option	Default	Description
database	neo4j	Name of database.
backup		Path to backup to check consistency of. Cannot be used together withdatabase.
additional-config		Configuration file to supply additional configuration in.
verbose	false	Enable verbose output.
report-dir		Directory to write report file in.
check-graph	true	Perform checks between nodes, relationships, properties, types and tokens.
check-indexes	true	Perform checks on indexes by comparing content with the store.
check-index-structure	true	Perform physical structure check on indexes. No comparison with the store takes place.
check-label-scan-store	true	Perform checks on the label scan store.

Option	Default	Description
check-property-owners		Perform additional checks on property ownership. This check is <b>very</b> expensive in time and memory.

#### Limitations

The consistency checker cannot be used with a database which is currently in use. If used with a running database, it will stop and print an error.

### Output

If the consistency checker does not find errors, it will exit cleanly and not produce a report. If the consistency checker finds errors, it will exit with an exit code of 1 and write a report file with a name on the format inconsistencies-YYYY-MM-DD.HH24.MI.SS.report. The location of the report file is the current working directory, or as specified by the parameter report-dir.

### Example 85. Run the consistency checker

Run with the --database option to check the consistency of a database. Note that the database must be stopped first.

```
$neo4j-home> bin/neo4j stop
$neo4j-home> bin/neo4j-admin check-consistency --database=neo4j
2019-11-13 12:42:14.479+0000 INFO [o.n.k.i.s.f.RecordFormatSelector] Selected
RecordFormat:StandardV4_0[SF4.0.b] record format from store /data/databases/neo4j
2019-11-13 12:42:14.481+0000 INFO [o.n.k.i.s.f.RecordFormatSelector] Format not configured for store
/data/databases/neo4j. Selected format from the store files: RecordFormat:StandardV4_0[SF4.0.b]
Index structure consistency check
20%
40%
                 50%
Full Consistency Check
..... 20%
60%
                 70%
                 90%
.Checking node and relationship counts
20%
                 30%
                 50%
                 80%
                 90%
```

Run with the --backup option to check the consistency of a backup.

```
bin/neo4j-admin check-consistency --backup backup/neo4j-backup
```

# 14.3. Report tool

This chapter describes the report command of Neo4j Admin.

Use the report command of neo4j-admin to gather information about a Neo4j installation and save it to an archive.

```
neo4j-admin report [--force] [--list] [--verbose] [--pid=<pid>] [--to=<path>] [<classifier>...]
```

The intended usage of the report tool is to simplify the support process by collecting the relevant information in a standard way. This tool does not send any information automatically. To share this information with the Neo4j Support organization, you will have to send it manually.

## **Options**

Option	Default	Description
to	reports/	Specifies to target directory where the report should be written to.
list		Will list available classifiers.
verbose		Will instruct the tool to print more verbose output.
force		Will disable the available disk space check.
pid		Specify process id of running Neo4j instance. Only applicable when used together with classifiers indicated as <i>Online</i> in the table below.

By default, the tool tries to estimate the final size of the report and use that to assert that there is enough disk space available for it. If there is not enough available space the tool will abort. However, this estimation is pessimistic and does not take the compression into consideration, so if you are confident that you do have enough disk space, you can disable this check with the option --force.

#### Classifiers

Classifier	Online	Description
all		Include all of the available classifiers listed below.
activetxs	X	<pre>Include the output of dbms.listTransactions().</pre>
ccstate		Include the current cluster state.
config		Include the neo4j.conf file.
env	X	Include a list of all environment variables.
heap	Х	Include a heap dump.
logs		<pre>Include log files, e.g debug.log, neo4j.log etc.</pre>
metrics		Include the collected metrics.
plugins		Include a text view of the plugin directory (no files are collected).
ps		include a list of running processes.
raft		Include the raft log.

Classifier	Online	Description
sysprop	X	Include a list of Java system properties.
threads	X	Include a thread dump of the running instance.
tree		Include a text view of the folder structure of the data directory (no files are collected).
tx		Include transaction logs.

The classifiers indicated as *Online* only work when you have a running Neo4j instance that the tool can find.

If no classifiers are specified, the following classifiers will be used: logs, config, plugins, tree, metrics, threads, env, sysprop and ps.

The report tool does not read any data from your database. However, the heap, the raft logs, and the transaction logs may contain data. Additionally while the standard neo4j.conf file does not contain password information, for certain configurations it may have this type of information. Therefore, be aware of your organization's data security rules before using the classifiers heap, tx, raft and config.



This tool uses the Java Attach API to gather data from a running Neo4j instance. Therefore, it requires the Java JDK in order to execute properly.

# 14.4. Display store information

This chapter describes the store-info command of Neo4j Admin.

The version of a Neo4j database can be displayed using the store-info command of neo4j-admin:

neo4j-admin store-info <path-to-store>

The store argument takes a path to a Neo4j database or backup. We will know if the store will need to go through a format migration as part of restoring a backup, by the presence of the message Store format superseded in: in the output.



The Neo4j Admin store-info command can not be used on a running database. The store is locked for protection while the database is running, and using the store-info command on it will fail with an error message indicating this.

You have an old backup of Neo4j a folder called /backups/graph-db.2016-11-11/ which you would like to restore. It is unclear what the version of Neo4j was at the time of backup. To find out, you would run:

```
$neo4j-home> bin/neo4j-admin store-info /backups/graph-db.2016-11-11
Store format version: vE.H.0
Introduced in version: 3.0.0
Store format superseded in: 3.0.6
```

The output reveals that the database was configured to use the high-limit format, which is Enterprise Edition only. This means the backup can only be used with Neo4j Enterprise Edition.

The output also explains that the format was introduced in Neo4j 3.0.0, but a newer format was shipped in Neo4j 3.0.6. This means that a format migration must be performed if the backup is restored (see Upgrade for details).

## Example 87. Example usage with newer backup

As with the previous example, let us assume that another backup of Neo4j is present in /backups/graph-db.2016-11-11/. In this case the output is:

The output tells us that the backup is made with the standard store format, which all editions of Neo4j support. It is also the newest version of the format known to Neo4j. This means that no format migration is necessary for the backup to be restored.

# 14.5. Memory recommendations

This chapter describes the memrec command of Neo4j Admin.

Use the memrec command of neo4j-admin to get an initial recommendation on how to configure memory parameters for Neo4j:

```
neo4j-admin memrec [--memory=<memory dedicated to Neo4j>]
```

The recommendations will be given in a format such that it can be copied and pasted straight into neo4j.conf.

## **Options**

Option	Default	Description
memory	The memory capacity of the machine	The amount of memory to allocate to Neo4j. Valid units are: k, K, m, M, g, G.

#### Considerations

The neo4j-admin memrec command calculates a valid starting point for Neo4j memory settings, based on the provided memory. The specific conditions for your use case may warrant adjustment of these values. See Memory configuration for a description of the memory settings in Neo4j.

The following example illustrates how neo4j-admin memrec provides a recommendation on how to use 16g of memory:

```
$neo4j-home> bin/neo4j-admin memrec --memory=16g
...
...
# Based on the above, the following memory settings are recommended:
dbms.memory.heap.initial_size=5g
dbms.memory.heap.max_size=5g
dbms.memory.pagecache.size=7g
```

# 14.6. Import

This section describes how to perform batch imports of data into Neo4j.

You can do batch imports of large amounts of data into a Neo4j database from CSV files, using the import command of neo4j-admin. This tool can only be used to load data into a previously unused database. By default, this database is set to neo4j but can be configured to other names and locations. If you wish to import small to medium-sized CSV files into an existing database, use LOAD CSV (see Cypher Manual [] LOAD CSV).

These are some things you will need to keep in mind when creating your input files:

- Fields are comma-separated by default but a different delimiter can be specified.
- · All files must use the same delimiter.
- Multiple data sources can be used for both nodes and relationships.
- A data source can optionally be provided using multiple files.
- A separate file with a header that provides information on the data fields, must be the first specified file of each data source.
- Fields without corresponding information in the header will not be read.
- UTF-8 encoding is used.
- By default, the importer will trim extra whitespace at the beginning and end of strings. Quote your data to preserve leading and trailing whitespaces.



### *Indexes and constraints*

Indexes and constraints are not created during the import. Instead, you will need to add these afterwards (see Cypher Manual 🛘 Indexes).

This chapter explains how to format the input data and use the import tool. If you wish to see indepth examples of using the import tool, refer to the Import tutorial.

The following sections describe how to use the import tool:

- Syntax The syntax of the neo4j-admin import command.
- CSV header format How to construct the header row of each CSV file.
- Options The options available for use with neo4j-admin import.

## 14.6.1. Syntax

The syntax for importing a set of CSV files is:

```
neo4j-admin import [--verbose]
                   [--high-io[=<true/false>]]
                   [--cache-on-heap]
                   [--ignore-empty-strings[=<true/false>]]
                   [--ignore-extra-columns[=<true/false>]]
                   [--legacy-style-quoting[=<true/false>]]
                   [--multiline-fields[=<true/false>]]
                   [--normalize-types[=<true/false>]]
                   [--skip-bad-entries-logging[=<true/false>]]
                   [--skip-bad-relationships[=<true/false>]]
                   [--skip-duplicate-nodes[=<true/false>]]
                   [--trim-strings[=<true/false>]]
                   [--additional-config=<path>]
                   [--array-delimiter=<char>]
                   [--bad-tolerance=<num>]
                   [--database=<database>]
                   [--delimiter=<char>]
                   [--id-type=<idType>]
                   [--input-encoding=<character-set>]
                   [--max-memory=<size>]
                   [--processors=<num>]
                   [--quote=<char>]
                   [--read-buffer-size=<size>]
                   [--report-file=<path>]
                   --nodes=[<label>[:<label>]...=]<files>...
                   [--nodes=[<label>[:<label>]...=]<files>...]...
                   [--relationships=[<type>=]<files>...]...
```

## Example 89. Import data from CSV files

Assume that we have formatted our data as per CSV file header format so that we have it in six different files: movies\_header.csv, movies.csv, actors\_header.csv, actors\_header.csv, roles\_header.csv, and roles.csv. The following command will import the three datasets:

```
neo4j_home$ bin/neo4j-admin import --nodes import/movies_header.csv,import/movies.csv \
--nodes import/actors_header.csv,import/actors.csv \
--relationships import/roles_header.csv,import/roles.csv
```

## 14.6.2. CSV file header format

This section explains the header format of CSV files when using the Neo4j import tool.

## This section describes the following:

- Header files
- Properties
- Node files
- Relationship files
- Using ID spaces
- Skipping columns
- Compressed files

## Header files

The header file of each data source specifies how the data fields should be interpreted. You must use the same delimiter for the header file and for the data files.

The header contains information for each field, with the format <name>:<field\_type>. The <name> is used for properties and node IDs. In all other cases, the <name> part of the field is ignored.

## **Properties**

For properties, the <name> part of the field designates the property key, while the <field\_type> part assigns a data type (see below). You can have properties in both node data files and relationship data files.

## Data types

Use one of int, long, float, double, boolean, byte, short, char, string, point, date, localtime, time, localdatetime, datetime, and duration to designate the data type for properties. If no data type is given, this defaults to string. To define an array type, append [] to the type. By default, array values are separated by ;. A different delimiter can be specified with --array-delimiter. Boolean values are *true* if they match exactly the text true. All other values are *false*. Values that contain the delimiter character need to be escaped by enclosing in double quotation marks, or by using a different delimiter character with the --delimiter option.

## Example 90. Header format with data types

This example illustrates several different data types specified in the CSV header.

```
:ID,name,joined:date,active:boolean,points:int
user01,Joe Soap,2017-05-05,true,10
user02,Jane Doe,2017-08-21,true,15
user03,Moe Know,2018-02-17,false,7
```

#### Special considerations for the point data type

A point is specified using the Cypher syntax for maps. The map allows the same keys as the input to the Cypher Manual Depoint function. The point data type in the header can be amended with a map of default values used for all values of that column, e.g. point{crs: 'WGS-84'}. Specifying the header this way allows you to have an incomplete map in the value position in the data file. Optionally, a value in a data file may override default values from the header.

This example illustrates various ways of using the **point** data type in the import header and the data files.

We are going to import the name and location coordinates for cities. First, we define the header as:

```
:ID,name,location:point{crs:WGS-84}
```

We then define cities in the data file.

- The first city's location is defined using latitude and longitude, as expected when using the coordinate system defined in the header.
- The second city uses x and y instead. This would normally lead to a point using the coordinate reference system cartesian. Since the header defines crs: WGS-84, that coordinate reference system will be used.
- The third city overrides the coordinate reference system defined in the header, and sets it explicitly to WGS-84-3D.

```
:ID,name,location:point{crs:WGS-84}
city01,"Malmö","{latitude:55.6121514, longitude:12.9950357}"
city02,"London","{y:51.507222, x:-0.1275}"
city03,"San Mateo","{latitude:37.554167, longitude:-122.313056, height: 100, crs:'WGS-84-3D'}"
```

Note that all point maps are within double quotation marks " in order to prevent the enclosed , character from being interpreted as a column separator. An alternative approach would be to use --delimiter='\t' and reformat the file with tab separators, in which case the " characters are not required.

```
:ID name location:point{crs:WGS-84}
city01 Malmö {latitude:55.6121514, longitude:12.9950357}
city02 London {y:51.507222, x:-0.1275}
city03 San Mateo {latitude:37.554167, longitude:-122.313056, height: 100, crs:'WGS-84-3D'}
```

### Special considerations for temporal data types

The format for all temporal data types must be defined as described in Cypher Manual [] Temporal instants syntax and Cypher Manual [] Durations syntax. Two of the temporal types, *Time* and *DateTime*, take a time zone parameter which might be common between all or many of the values in the data file. It is therefor possible to specify a default time zone for *Time* and *DateTime* values in the header, for example: time{timezone:+02:00} and: datetime{timezone:Europe/Stockholm}. If no default time zone is specified, the default timezone is determined by the db.temporal.timezone configuration setting. The default time zone can be explicitly overridden in the values in the data file.

This example illustrates various ways of using the data type in the import header and the data files.

First, we define the header with two *DateTime* columns. The first one defines a time zone, but the second one does not:

```
:ID,date1:datetime{timezone:Europe/Stockholm},date2:datetime
```

We then define dates in the data file.

- The first row has two values that do not specify an explicit timezone. The value for date1 will use the Europe/Stockholm time zone that was specified for that field in the header. The value for date2 will use the configured default time zone of the database.
- In the second row, both date1 and date2 set the time zone explicitly to be Europe/Berlin. This overrides the header definition for date1, as well as the configured default time zone of the database.

```
1,2018-05-10T10:30,2018-05-10T12:30
2,2018-05-10T10:30[Europe/Berlin],2018-05-10T12:30[Europe/Berlin]
```

## Node files

Files containing node data can have an ID field, a LABEL field as well as properties.

ID

Each node must have a unique ID if it is to be connected by any relationships created in the import. The IDs are used to find the correct nodes when creating relationships. Note that the ID has to be unique across all nodes in the import; even for nodes with different labels. The unique ID can be persisted in a property whose name is defined by the <name> part of the field definition <name>: ID. If no such property name is defined, the unique ID will be used for the purpose of the import but not be available for reference later. If no ID is specified, the node will be imported but it will not be able to be connected by any relationships during the import.

#### LABEL

Read one or more labels from this field. Like array values, multiple labels are separated by ;, or by the character specified with --array-delimiter.

We define the headers for movies in the *movies\_header.csv* file. Movies have the properties movieId, year and title. We also specify a field for labels.

```
movieId:ID,title,year:int,:LABEL
```

We define three movies in the *movies.csv* file. They contain all the properties defined in the header file. All the movies are given the label Movie. Two of them are also given the label Sequel.

```
tt0133093,"The Matrix",1999,Movie
tt0234215,"The Matrix Reloaded",2003,Movie;Sequel
tt0242653,"The Matrix Revolutions",2003,Movie;Sequel
```

Similarly, we also define three actors in the *actors\_header.csv* and *actors.csv* files. They all have the properties personId and name, and the label Actor.

```
personId:ID, name,:LABEL
```

```
keanu, "Keanu Reeves", Actor
laurence, "Laurence Fishburne", Actor
carrieanne, "Carrie-Anne Moss", Actor
```

## Relationship files

Files containing relationship data have three mandatory fields and can also have properties. The mandatory fields are:

#### **TYPF**

The relationship type to use for this relationship.

## START\_ID

The ID of the start node for this relationship.

## END\_ID

The ID of the end node for this relationship.

The START\_ID and END\_ID refer to the unique node ID defined in one of the node data sources, as explained in the previous section. None of these takes a name, e.g. if <name>: START\_ID or <name>: END\_ID is defined, the <name> part will be ignored.

In this example we assume that the two nodes files from the previous example are used together with the following relationships file.

We define relationships between actors and movies in the files *roles\_header.csv* and *roles.csv*. Each row connects a start node and an end node with a relationship of relationship type ACTED\_IN. Notice how we use the unique identifiers personId and movieId from the nodes files above. The name of character that the actor is playing in this movie is stored as a *role* property on the relationship.

```
:START_ID,role,:END_ID,:TYPE
```

```
keanu, "Neo", tt0133093, ACTED_IN
keanu, "Neo", tt0234215, ACTED_IN
keanu, "Neo", tt0242653, ACTED_IN
laurence, "Morpheus", tt0133093, ACTED_IN
laurence, "Morpheus", tt0234215, ACTED_IN
laurence, "Morpheus", tt0242653, ACTED_IN
carrieanne, "Trinity", tt0133093, ACTED_IN
carrieanne, "Trinity", tt0234215, ACTED_IN
carrieanne, "Trinity", tt0234215, ACTED_IN
carrieanne, "Trinity", tt0242653, ACTED_IN
```

## Using ID spaces

By default, the import tool assumes that node identifiers are unique across node files. In many cases the ID is only unique across each entity file, for example when our CSV files contain data extracted from a relational database and the ID field is pulled from the primary key column in the corresponding table. To handle this situation we define *ID spaces*. ID spaces are defined in the ID field of node files using the syntax ID(<ID space identifier>). To reference an ID of an ID space in a relationship file, we use the syntax START\_ID(<ID space identifier>) and END\_ID(<ID space identifier>).

```
Define a Movie-ID ID space in the movies_header.csv file.
```

movieId:ID(Movie-ID),title,year:int,:LABEL

personId:ID(Actor-ID),name,:LABEL

```
1,"The Matrix",1999,Movie
2,"The Matrix Reloaded",2003,Movie;Sequel
3,"The Matrix Revolutions",2003,Movie;Sequel
```

Define an Actor-ID ID space in the header of the actors\_header.csv file.

```
1, "Keanu Reeves", Actor
2, "Laurence Fishburne", Actor
3, "Carrie-Anne Moss", Actor
```

Now use the previously defined ID spaces when connecting the actors to movies.

```
:START_ID(Actor-ID),role,:END_ID(Movie-ID),:TYPE
```

```
1,"Neo",1,ACTED_IN
1,"Neo",2,ACTED_IN
1,"Neo",3,ACTED_IN
2,"Morpheus",1,ACTED_IN
2,"Morpheus",2,ACTED_IN
2,"Morpheus",3,ACTED_IN
2,"Morpheus",3,ACTED_IN
3,"Trinity",1,ACTED_IN
3,"Trinity",2,ACTED_IN
3,"Trinity",3,ACTED_IN
```

## Skipping columns

#### **IGNORE**

If there are fields in the data that we wish to ignore completely, this can be done using the IGNORE keyword in the header file. IGNORE must be prepended with a :.

## Example 96. Skip a column

In this example, we are not interested in the data in the third column of the nodes file and wish to skip over it. Note that the IGNORE keyword is prepended by a :.

```
personId:ID,name,:IGNORE,:LABEL

keanu, "Keanu Reeves", "male", Actor
laurence, "Laurence Fishburne", "male", Actor
carrieanne, "Carrie-Anne Moss", "female", Actor
```

If all your superfluous data is placed in columns located to the right of all the columns that you wish to import, you can instead use the command line option --ignore-extra-columns.

## Import compressed files

The import tool can handle files compressed with zip or gzip. Each compressed file must contain a single file.

## Example 97. Perform an import using compressed files

neo4j\_home\$ ls import
actors-header.csv actors.csv.zip movies-header.csv movies.csv.gz roles-header.csv roles.csv.gz
neo4j\_home\$ bin/neo4j-admin import --nodes import/movies-header.csv,import/movies.csv.gz --nodes
import/actors-header.csv,import/actors.csv.zip --relationships import/rolesheader.csv,import/roles.csv.gz

## 14.6.3. Options

This section describes in details the options available when using the Neo4j import tool to import data from CSV files.



Some of the options below are marked as **Advanced**. These options should not be used for experimentation.

For more information, please contact Neo4j Professional Services.

#### --verbose

Enable verbose output.

### --database=<name>

Name of database.

Default: neo4j

#### --additional-config=<config-file-path>

Configuration file with which to supply additional configuration.

## --report-file=<filename>

File in which to store the report of the csv-import.

Default: import.report

### --id-type=<STRING|INTEGER|ACTUAL>

Each node must provide a unique ID in order to be used for creating relationships during the import.

## Possible values are:

- STRING: arbitrary strings for identifying nodes,
- INTEGER: arbitrary integer values for identifying nodes,
- ACTUAL: (Advanced) actual node IDs.

Default: STRING

## --input-encoding=<character-set>

Character set that input data is encoded in.

#### Default: UTF-8

### --ignore-extra-columns[=<true/false>]

If unspecified columns should be ignored during the import.

Default: false

#### --multiline-fields[=<true/false>]

Determines whether or not fields from input source can span multiple lines, i.e. contain newline characters.

Setting --multiline-fields=true can severely degrade performance of the importer. Therefore, use it with care, especially with large imports.

Default: false

## --ignore-empty-strings[=<true/false>]

Determines whether or not empty string fields, such as "", from input source are ignored (treated as null).

Default: false

#### --trim-strings[=<true/false>]

Determines whether or not strings should be trimmed for whitespaces.

Default: false

## --legacy-style-quoting[=<true/false>]

Determines whether or not backslash-escaped quote e.g. \" is interpreted as inner quote.

Default: false

### --delimiter=<char>

Delimiter character between values in CSV data.

Unicode character encoding can be used if prepended by \. For example, \44 is equivalent to ,.

Default:,

#### --array-delimiter=<char>

Delimiter character between array elements within a value in CSV data.

Unicode character encoding can be used if prepended by \. For example, \59 is equivalent to ;.

Default: ;

### --quote=<char>

Character to treat as quotation character for values in CSV data.

Quotes can be escaped as per RFC 4180 by doubling them, for example "" would be interpreted as a literal ".

You cannot escape using \.

## --read-buffer-size=<size>

Size of each buffer for reading input data.

It has to at least be large enough to hold the biggest single value in the input data. Value can be a plain number or byte units string, e.g. 128k, 1m.

Default: 4m

### --max-memory=<size>

Maximum memory that neo4j-admin can use for various data structures and caching to improve performance.

Values can be plain numbers such as 10000000, or 20G for 20 gigabyte. It can also be specified as a percentage of the available memory, for example 70%.

Default: 90%

### --high-io[=<true/false>]

Ignore environment-based heuristics, and specify whether the target storage subsystem can support parallel IO with high throughput.

Typically this is true for SSDs, large raid arrays and network-attached storage.

#### --cache-on-heap[=<true/false>] Advanced

Determines whether or not to allow allocating memory for the cache on heap.

If false, then caches will still be allocated off-heap, but the additional free memory inside the JVM will not be allocated for the caches.

Use this to have better control over the heap memory.

Default: false

## --processors=<num> Advanced

Max number of processors used by the importer.

Defaults to the number of available processors reported by the JVM. There is a certain amount of minimum threads needed, so for that reason there is no lower bound for this value.

For optimal performance, this value shouldn't be greater than the number of available processors.

#### --bad-tolerance=<num>

Number of bad entries before the import is considered failed.

This tolerance threshold is about relationships referring to missing nodes. Format errors in input data are still treated as errors.

Default: 1000

### --skip-bad-entries-logging[=<true/false>]

Determines whether or not to skip logging bad entries detected during import.

Default: false

#### --skip-bad-relationships[=<true/false>]

Determines whether or not to skip importing relationships that refer to missing node IDs, i.e. either start or end node ID/group referring to node that wasn't specified by the node input data.

Skipped nodes will be logged, containing at most the number of entities specified by bad-tolerance, unless otherwise specified by the skip-bad-entries-logging option.

Default: false

#### --skip-duplicate-nodes[=<true/false>]

Determines whether or not to skip importing nodes that have the same ID/group.

In the event of multiple nodes within the same group having the same ID, the first encountered will be imported, whereas consecutive such nodes will be skipped.

Skipped nodes will be logged, containing at most the number of entities specified by bad-tolerance, unless otherwise specified by the skip-bad-entries-logging option.

Default: false

#### --normalize-types[=<true/false>]

Determines whether or not to normalize property types to Cypher types, e.g. int becomes long and float becomes double.

Default: true

### --nodes=[<label>[:<label>]...=]<files>...

Node CSV header and data.

- Multiple files will be logically seen as one big file from the perspective of the importer.
- The first line must contain the header.
- Multiple data sources like these can be specified in one import, where each data source has its own header.
- Files can also be specified using regular expressions.

For an example, see Using regular expressions for specifying multiple input files.

#### --relationships=[<type>=]<files>...

Relationship CSV header and data.

- Multiple files will be logically seen as one big file from the perspective of the importer.
- The first line must contain the header.
- Multiple data sources like these can be specified in one import, where each data source has its own header.
- Files can also be specified using regular expressions.

For an example, see Using regular expressions for specifying multiple input files.

## @<arguments-file-path>

File containing all arguments, used as an alternative to supplying all arguments on the command line directly.

Each argument can be on a separate line, or multiple arguments per line and separated by space.

Arguments containing spaces must be quoted.

Heap size for the import



You want to set the maximum heap size to a relevant value for the import. This is done by defining the HEAP\_SIZE environment parameter before starting the import. For example, 2G is an appropriate value for smaller imports.

If doing imports in the order of magnitude of 100 billion entities, 20G will be an appropriate value.

## Record format



If your import data will result in a graph that is larger than 34 billion nodes, 34 billion relationships, or 68 billion properties you will need to configure the importer to use the high limit record format. This is achieved by setting the parameter dbms.record\_format=high\_limit in a configuration file, and supplying that file to the importer with --additional-config.

The high\_limit format is available for Enterprise Edition only.

## Output

The location of the import log file can be controlled using the --report-file option. If you run large imports of CSV files that have low data quality, the import log file can grow very large. For example, CSV files that contain duplicate node IDs, or that attempt to create relationships between non-existent nodes, could be classed as having low data quality. In these cases, you may wish to direct the output to a location that can handle the large log file.

If you are running on a UNIX-like system and you are not interested in the output, you can get rid of it altogether by directing the report file to /dev/null.

If you need to debug the import, it might be useful to collect the stack trace. This is done by using --verbose option.

# 14.7. Dump and load databases

This section describes the dump and load commands of Neo4j Admin.

A Neo4j database can be dumped and loaded using the following commands:

```
neo4j-admin dump --database=<database> --to=<destination-path>
neo4j-admin load --from=<archive-path> --database=<database> [--force]
```

#### Limitations

- The database should be shutdown before running the dump and load commands.
- neo4j-admin must be invoked as the neo4j user in order to ensure the appropriate file permissions.

## Usage

Using the dump and load commands are the recommended, and safe, way of transferring databases between environments. They understand which files need to be exported and imported and which should not.

+ By contrast, file system copy-and-paste of databases is not supported.

#### **Examples**

Dump the database called <a href="neo4j">neo4j</a> into a file called <a href="neo4j">/backups/neo4j</a>/ — must exist before calling the command.

```
$neo4j-home> bin/neo4j-admin dump --database=neo4j --to=/backups/neo4j/2016-10-02.dump
$neo4j-home> ls /backups/neo4j
$neo4j-home> 2016-10-02.dump
```

## Example 99. Use the load command of neo4j-admin

Load the backed-up database contained in the file /backups/neo4j/2016-10-02.dump into database neo4j. If you have a database running, it needs to be shutdown first. When we use the --force option, any existing database gets overwritten.

\$neo4j-home> bin/neo4j-admin load --from=/backups/neo4j/2016-10-02.dump --database=neo4j --force



If using the load command to seed a Causal Cluster, you must first perform neo4j-admin unbind on each of the cluster instances. The procedure is described in Seed from backups.

## 14.8. Unbind a Core Server

This section describes how to remove cluster state data for a Neo4j server.

The cluster state of a Causal Cluster member can be removed by using the following command:

neo4j-admin unbind

### Limitations

The Neo4j server process should be shutdown before running the unbind command.

#### **Examples of usage**

• When transforming a Causal Cluster member to a standalone server:

The <u>unbind</u> command can be used to turn a Causal Cluster server into a standalone server. To start the database in single (standalone) mode after unbinding it from the cluster, first set dbms.mode=SINGLE in neo4j.conf.

• When seeding a Causal Cluster with existing store files:

If you wish to seed a new Causal Cluster using the store files of a previous cluster, you must first run neo4j-admin unbind on each server. For more information about seeding Causal Clusters, see Seed a cluster.

When recovering a Causal Cluster:

In the event of serious failures you may need to recover an entire Causal Cluster from backups.

Before restoring those backups, you must first run neo4j-admin unbind on each server. For more information about recovering databases from online backups, see Restore a backup.



Unlike versions of Neo4j prior to v4.0.0, you **must** run the unbind command on both Read Replicas and Core members.

# 14.9. Copy a database

This chapter describes the copy command of Neo4j Admin.

The copy command of neo4j-admin is used to copy data from an existing database to a new database.

The syntax follows:

The existing database must be stopped before copying from it, and the destination database must not yet exist.

The copy command can process an optional set of filters. This can be used to remove data that is unwanted in the destination database.

The schema definitions, i.e. index and constraint, are not automatically transferred. However, they will be extracted and presented as Cypher statements so you can recreate the ones you want.

#### **Options**

Option	Description
from-database	The database name to copy from. Will assume the database is in the configured location.
from-path	The path to the database to copy from. It can be used to target databases outside of the installation, e.g. backups.
from-path-tx	The path to the transaction log files. You only need to use this if the command is unable to determine where they are located.
to-database	The destination database name.

Option	Description
to-format	The store format of the destination database. Valid values are same, standard, high_limit. The default value for this option is the format of the source database.
delete-nodes-with-labels	A list of labels. Any node matching any of the labels will be ignored during copy.
skip-properties	A list of property keys to ignore during the copy.
skip-labels	A list of labels to ignore during the copy.
skip-relationships	A list of relationship types to ignore during the copy.
verbose	Will instruct the tool to print more verbose output.
force	Will force the command to proceed even if the integrity of the database can not be verified.



Due to the way filters are processed, the  ${\tt id}$  of the node might change. This is true even if no filters are specified.

## Examples

Example 100. Use the copy command to take a copy of the database neo4j.

To begin, you must stop the database named neo4j; this can be done by issuing the following Cypher statement:

```
STOP DATABASE neo4j
```

You can now copy the data from neo4j, to a new database called copy:

```
$neo4j-home> bin/neo4j-admin copy --from-database=neo4j --to-database=copy
```

A new database with the name copy now exists on the server, but it is not automatically picked up by Neo4j. To start the new database you have to insert it into Neo4j with the following Cypher query:

```
CREATE DATABASE copy
```

Neo4j will then detect the copied database and begin to use that.

Remember to start the database you copied from, if you still want it, using:

```
START DATABASE neo4j
```

The console output is saved to *logs/neo4j-admin-copy-<date>.log*.

Example 101. Use the copy command with filters.

The command can perform some basic forms of processing. You can remove nodes, labels, properties, and/or relationships.

The difference with --skip-labels and --delete-nodes-with-labels is that --skip-labels will just remove the labels, potentially leaving nodes without any labels.

```
\label{local-phome} $$ \sin/neo4j-admin copy --from-database=neo4j --to-database=copy --delete-nodes-with-labels="Cat,Dog" $$
```

After this command, you will have a copy of the database neo4j, without nodes with the labels :Cat and :Dogs.



Labels are processed independently, i.e. the filter described above will delete any node with either :Cat or :Dogs, and not only nodes that have both of the labels.

# 14.10. Cypher Shell

This section describes Neo4j Cypher Shell.

This section describes the following:

- About Cypher Shell
- Syntax
- Query parameters
- Transactions
- Procedures

## 14.10.1. About Cypher Shell

Cypher Shell is a command-line tool that comes with the Neo4j installation. It can also be downloaded from Neo4j Download Center and installed separately.

Cypher Shell is used to run queries and perform administrative tasks. It communicates via the encrypted binary protocol Bolt.

## 14.10.2. Syntax

Cypher Shell is located in the bin directory if installed as part of the product. The syntax is:

```
cypher-shell [-h] [-a ADDRESS] [-u USERNAME] [-p PASSWORD] [--encryption {true,false}] [{-d
DATABASE,--database DATABASE}] [--fail-fast | --fail-at-end] [--format {verbose,plain}] [{--P
PARAM,--param PARAM}] [--debug] [--non-interactive] [--sample-rows SAMPLE-ROWS] [--wrap
{true,false}] [-v] [--driver-version] [{-f FILE,--file FILE}] [cypher]
```

### **Arguments**

Positional arguments:	
cypher	An optional string of Cypher to execute and then exit.
Optional arguments:	
-h,help	Show help message and exit.
fail-fast	Exit and report failure on first error when reading from file (this is the default behavior).
fail-at-end	Exit and report failures at end of input when reading from file.
format {auto,verbose,plain}	Desired output format. auto displays results in tabular format if you use the shell interactively and with minimal formatting if you use it for scripting. verbose displays results in tabular format and prints statistics and plain displays data with minimal formatting (default: auto).
P PARAM,param PARAM	Add a parameter to this session. Example: -P "number => 3". This argument can be specified multiple times.
debug	Print additional debug information (default: false).
non-interactive	Force non-interactive mode; only useful if auto-detection fails (default: false).
sample-rows SAMPLE-ROWS	Number of rows sampled to compute table widths (only for format=VERBOSE) (default: 1000).
wrap {true,false}	Wrap table column values if column is too narrow (only for format=VERBOSE) (default: true).
-v,version	Print version of cypher-shell and exit (default: false).
driver-version	Print version of the Neo4j Driver used and exit (default: false).
-f FILE,file FILE	Pass a file with Cypher statements to be executed. After the statements have been executed cypher-shell will be shut down.

Connection arguments:	
-a ADDRESS,address ADDRESS	Address and port to connect to (default: bolt://localhost:7687).
-u USERNAME,username USERNAME	Username to connect as. Can also be specified using environment variable NEO4J_USERNAME (default: ).
-p PASSWORD,password PASSWORD	Password to connect with. Can also be specified using environment variable NEO4J_PASSWORD (default: ).
encryption {true,false}	Whether the connection to Neo4j should be encrypted; must be consistent with Neo4j's configuration (default: true).
-d DATABASE,database DATABASE	Database to connect to. Can also be specified using environment variable NEO4J_DATABASE (default: ).

#### Example 102. Invoke Cypher Shell with username and password

```
$neo4j-home> bin/cypher-shell -u johndoe -p secret

Connected to Neo4j at bolt://localhost:7687 as user neo4j.
Type :help for a list of available commands or :exit to exit the shell.
Note that Cypher queries must end with a semicolon.
neo4j>
```

#### Example 103. Invoke help from within Cypher Shell

```
neo4j> :help
Available commands:
 :begin
          Open a transaction
 :commit
          Commit the currently open transaction
         Exit the logger
 :exit
 :help
         Show this help message
 :history Print a list of the last commands executed
          Set the value of a query parameter
 :param
 :rollback Rollback the currently open transaction
 :source Interactively executes Cypher statements from a file
          Set the active database
 :use
For help on a specific command type:
   :help command
```

#### Example 104. Execute a query from within Cypher Shell

Below is the contents of a file called examples.cypher:

```
MATCH (n) RETURN n;

MATCH (batman:Person {name: 'Bruce Wayne'}) RETURN batman;
```

Invoke the 'examples.cypher' script from the command-line. All the examples in the remainder of this section will use the --format plain flag for a simple output.

```
$neo4j-home> cat examples.cypher | bin/cypher-shell -u neo4j -p secret --format plain

n
  (:Person {name: "Bruce Wayne", alias: "Batman"})
  (:Person {name: "Selina Kyle", alias: ["Catwoman", "The Cat"]})
batman
  (:Person {name: "Bruce Wayne", alias: "Batman"})
```

# 14.10.3. Query parameters

Cypher Shell supports querying based on parameters. This is often used while scripting.

Example 106. Use parameters within Cypher Shell

Set the parameter 'thisAlias' to 'Robin' using the ':param' keyword. Check the parameter using the ':params' keyword.

```
neo4j> :param thisAlias => 'Robin'
neo4j> :params
:param thisAlias => 'Robin'
```

Now use the parameter 'this Alias' in a Cypher query. Verify the result.

#### 14.10.4. Transactions

Cypher Shell supports explicit transactions. Transaction states are controlled using the keywords :begin, :commit, and :rollback:

Start a transaction in your first Cypher Shell session:

If you now open up a second Cypher Shell session, you will notice no changes from the latest CREATE statement:

Go back to the first session and commit the transaction:

#### 14.10.5. Procedures

Cypher Shell supports running any procedures for which the current user is authorized. Here, we are using the natively built-in procedure <a href="mailto:dbms.showCurrentUser">dbms.showCurrentUser</a>().

# Example 108. Call a procedure from within Cypher Shell

# Appendix A: Reference

This appendix contains the Neo4j configuration settings reference, the list of built-in procedures bundled with Neo4j, and a description of user management for Community Edition.

#### This appendix contains the following:

- Configuration settings
- Built-in procedures

# A.1. Configuration settings

This section contains a complete reference of Neo4j configuration settings. They can be set in neo4j.conf. Refer to The neo4j.conf file for details on how to use configuration settings.

#### All settings

- browser.allow\_outgoing\_connections: Configure the policy for outgoing Neo4j Browser connections.
- browser.credential\_timeout: Configure the Neo4j Browser to time out logged in users after this idle period.
- browser.post\_connect\_cmd: Commands to be run when Neo4j Browser successfully connects to this server.
- browser.remote\_content\_hostname\_whitelist: Whitelist of hosts for the Neo4j Browser to be allowed to fetch content from.
- browser.retain\_connection\_credentials: Configure the Neo4j Browser to store or not store user credentials.
- causal\_clustering.catch\_up\_client\_inactivity\_timeout: The catch up protocol times out if the given duration elapses with no network activity.
- causal\_clustering.catchup\_batch\_size: The maximum batch size when catching up (in unit of entries).
- causal\_clustering.cluster\_allow\_reads\_on\_followers: Configure if the dbms.routing.getRoutingTable() procedure should include followers as read endpoints or return only read replicas.
- causal\_clustering.cluster\_topology\_refresh: Time between scanning the cluster to refresh current server's view of topology.
- causal\_clustering.connect-randomly-to-server-group: Comma separated list of groups to be used by the connect-randomly-to-server-group selection strategy.
- causal\_clustering.discovery\_advertised\_address: Advertised cluster member discovery management communication.
- causal\_clustering.discovery\_listen\_address: Host and port to bind the cluster member discovery management communication.
- causal\_clustering.discovery\_type: Configure the discovery type used for cluster name resolution.
- causal\_clustering.enable\_pre\_voting: Enable pre-voting extension to the Raft protocol (this is breaking and must match between the core cluster members).
- causal\_clustering.global\_session\_tracker\_state\_size: The maximum file size before the global session tracker state file is rotated (in unit of entries).

- causal\_clustering.handshake\_timeout: Time out for protocol negotiation handshake.
- causal\_clustering.in\_flight\_cache.max\_bytes: The maximum number of bytes in the in-flight cache.
- causal\_clustering.in\_flight\_cache.max\_entries: The maximum number of entries in the in-flight cache.
- causal\_clustering.in\_flight\_cache.type: Type of in-flight cache.
- causal\_clustering.initial\_discovery\_members: A comma-separated list of other members of the cluster to join.
- causal\_clustering.join\_catch\_up\_timeout: Time out for a new member to catch up.
- causal\_clustering.kubernetes.address: Address for Kubernetes API.
- causal\_clustering.kubernetes.ca\_crt; File location of CA certificate for Kubernetes API.
- causal\_clustering.kubernetes.label\_selector: LabelSelector for Kubernetes API.
- causal clustering.kubernetes.namespace: File location of namespace for Kubernetes API.
- causal\_clustering.kubernetes.service\_port\_name: Service port name for discovery for Kubernetes API.
- causal\_clustering.kubernetes.token: File location of token for Kubernetes API.
- causal\_clustering.last\_applied\_state\_size: The maximum file size before the storage file is rotated (in unit of entries).
- causal\_clustering.leader\_election\_timeout: The time limit within which a new leader election will occur if no messages are received.
- causal\_clustering.load\_balancing.plugin: The load balancing plugin to use.
- causal\_clustering.load\_balancing.shuffle: Enables shuffling of the returned load balancing result.
- causal\_clustering.log\_shipping\_max\_lag: The maximum lag allowed before log shipping pauses (in unit of entries).
- causal\_clustering.middleware.logging.level: The level of middleware logging.
- causal\_clustering.minimum\_core\_cluster\_size\_at\_formation: Minimum number of Core machines initially required to form a cluster.
- causal\_clustering.minimum\_core\_cluster\_size\_at\_runtime: The minimum size of the dynamically adjusted voting set (which only core members may be a part of).
- causal\_clustering.multi\_dc\_license: Enable multi-data center features.
- causal\_clustering.protocol\_implementations.catchup: Catchup protocol implementation versions that this instance will allow in negotiation as a comma-separated list.
- causal\_clustering.protocol\_implementations.compression: Network compression algorithms that this instance will allow in negotiation as a comma-separated list.
- causal\_clustering.protocol\_implementations.raft: Raft protocol implementation versions that this instance will allow in negotiation as a comma-separated list.
- causal clustering, pull interval: Interval of pulling updates from cores.
- causal\_clustering.raft\_advertised\_address: Advertised hostname/IP address and port for the RAFT server.
- causal clustering.raft in queue max batch bytes: Largest batch processed by RAFT in bytes.
- causal\_clustering.raft\_in\_queue\_max\_bytes: Maximum number of bytes in the RAFT in-queue.
- causal\_clustering.raft\_listen\_address: Network interface and port for the RAFT server to listen on.
- causal\_clustering.raft\_log\_implementation: RAFT log implementation.
- causal\_clustering.raft\_log\_prune\_strategy: RAFT log pruning strategy.
- causal\_clustering.raft\_log\_pruning\_frequency: RAFT log pruning frequency.

- causal clustering.raft log reader pool size: RAFT log reader pool size.
- causal\_clustering.raft\_log\_rotation\_size: RAFT log rotation size.
- causal\_clustering.raft\_membership\_state\_size: The maximum file size before the membership state file is rotated (in unit of entries).
- causal\_clustering.raft\_term\_state\_size: The maximum file size before the term state file is rotated (in unit of entries).
- causal\_clustering.raft\_vote\_state\_size: The maximum file size before the vote state file is rotated (in unit of entries).
- causal\_clustering.refuse\_to\_be\_leader: Prevents the current instance from volunteering to become Raft leader.
- causal\_clustering.replicated\_lease\_state\_size: The maximum file size before the replicated lease state file is rotated (in unit of entries).
- causal\_clustering.replication\_leader\_await\_timeout: The duration for which the replicator will await a new leader.
- causal\_clustering.replication\_retry\_timeout\_base: The initial timeout until replication is retried.
- causal\_clustering.replication\_retry\_timeout\_limit: The upper limit for the exponentially incremented retry timeout.
- causal\_clustering.server\_groups: A list of group names for the server used when configuring load balancing and replication policies.
- causal\_clustering.state\_machine\_apply\_max\_batch\_size: The maximum number of operations to be batched during applications of operations in the state machines.
- causal\_clustering.state\_machine\_flush\_window\_size: The number of operations to be processed before the state machines flush to disk.
- causal\_clustering.status\_throughput\_window: Sampling window for throughput estimate reported in the status endpoint.
- causal clustering.store copy chunk size: Store copy chunk size.
- causal\_clustering.store\_copy\_max\_retry\_time\_per\_request: Maximum retry time per request during store copy.
- causal\_clustering.transaction\_advertised\_address: Advertised hostname/IP address and port for the transaction shipping server.
- causal\_clustering.transaction\_listen\_address: Network interface and port for the transaction shipping server to listen on.
- causal\_clustering.unknown\_address\_logging\_throttle: Throttle limit for logging unknown cluster member address.
- causal\_clustering.upstream\_selection\_strategy: An ordered list in descending preference of the strategy which read replicas use to choose the upstream server from which to pull transactional updates.
- causal\_clustering.user\_defined\_upstream\_strategy: Configuration of a user-defined upstream selection strategy.
- cypher.default\_language\_version: Set this to specify the default parser (language version).
- cypher.forbid\_exhaustive\_shortestpath: This setting is associated with performance optimization.
- cypher.forbid\_shortestpath\_common\_nodes: This setting is associated with performance optimization.
- cypher.hints\_error: Set this to specify the behavior when Cypher planner or runtime hints cannot be fulfilled.
- cypher.lenient\_create\_relationship: Set this to change the behavior for Cypher create relationship when the start or end node is missing.

- cypher.min replan interval: The minimum time between possible cypher query replanning events.
- cypher.planner: Set this to specify the default planner for the default language version.
- cypher.query\_max\_allocations: The maximum amount of heap memory allocations to for cypher to perform on a single query, in bytes (or kilobytes with the 'k' suffix, megabytes with 'm' and gigabytes with 'g').
- cypher.statistics\_divergence\_threshold: The threshold when a plan is considered stale.
- db.temporal.timezone: Database timezone for temporal functions.
- dbms.allow\_upgrade: Whether to allow an upgrade in case the current version of the database starts against an older version.
- dbms.backup.enabled: Enable support for running online backups.
- dbms.backup.listen\_address: Network interface and port for the backup server to listen on.
- dbms.checkpoint: Configures the general policy for when check-points should occur.
- dbms.checkpoint.interval.time: Configures the time interval between check-points.
- dbms.checkpoint.interval.tx: Configures the transaction interval between check-points.
- dbms.checkpoint.iops.limit: Limit the number of IOs the background checkpoint process will consume per second.
- dbms.config.strict\_validation: A strict configuration validation will prevent the database from starting up if unknown configuration options are specified in the neo4j settings namespace (such as dbms., cypher., etc).
- dbms.connector.bolt.advertised\_address: Advertised address for this connector.
- dbms.connector.bolt.enabled: Enable the bolt connector.
- dbms.connector.bolt.listen address: Address the connector should bind to.
- dbms.connector.bolt.thread\_pool\_keep\_alive: The maximum time an idle thread in the thread pool bound to this connector will wait for new tasks.
- dbms.connector.bolt.thread\_pool\_max\_size: The maximum number of threads allowed in the thread pool bound to this connector.
- dbms.connector.bolt.thread\_pool\_min\_size: The number of threads to keep in the thread pool bound to this connector, even if they are idle.
- dbms.connector.bolt.tls\_level: Encryption level to require this connector to use.
- dbms.connector.bolt.unsupported\_thread\_pool\_shutdown\_wait\_time: The maximum time to wait for the thread pool to finish processing its pending jobs and shutdown.
- dbms.connector.http.advertised\_address: Advertised address for this connector.
- dbms.connector.http.enabled: Enable the http connector.
- dbms.connector.http.listen\_address: Address the connector should bind to.
- dbms.connector.https.advertised address: Advertised address for this connector.
- dbms.connector.https.enabled: Enable the https connector.
- dbms.connector.https.listen\_address: Address the connector should bind to.
- dbms.db.timezone: Database timezone.
- dbms.default\_advertised\_address: Default hostname or IP address the server uses to advertise itself
- dbms.default\_database: Name of the default database.
- dbms.default\_listen\_address: Default network interface to listen for incoming connections.
- dbms.directories.data: Path of the data directory.
- dbms.directories.import: Sets the root directory for file URLs used with the Cypher LOAD CSV

#### clause.

- · dbms.directories.lib: Path of the lib directory.
- dbms.directories.logs: Path of the logs directory.
- dbms.directories.metrics: The target location of the CSV files: a path to a directory wherein a CSV file per reported field will be written.
- dbms.directories.neo4j\_home: Root relative to which directory settings are resolved.
- dbms.directories.plugins: Location of the database plugin directory.
- dbms.directories.run: Path of the run directory.
- dbms.directories.transaction.logs.root: Root location where Neo4j will store transaction logs for configured databases.
- dbms.dynamic.setting.whitelist: A list of setting name patterns (comma separated) that are allowed to be dynamically changed.
- dbms.filewatcher.enabled: Allows the enabling or disabling of the file watcher service.
- dbms.import.csv.buffer\_size: The size of the internal buffer in bytes used by LOAD CSV.
- dbms.import.csv.legacy\_quote\_escaping: Selects whether to conform to the standard https://tools.ietf.org/html/rfc4180 for interpreting escaped quotation characters in CSV files loaded using LOAD CSV.
- dbms.index.default\_schema\_provider: Index provider to use for newly created schema indexes.
- dbms.index.fulltext.default\_analyzer: The name of the analyzer that the fulltext indexes should use by default.
- dbms.index.fulltext.eventually\_consistent: Whether or not fulltext indexes should be eventually consistent by default or not.
- dbms.index.fulltext.eventually\_consistent\_index\_update\_queue\_max\_length: The
  eventually\_consistent mode of the fulltext indexes works by queueing up index updates to be
  applied later in a background thread.
- dbms.index\_sampling.background\_enabled: Enable or disable background index sampling.
- dbms.index\_sampling.sample\_size\_limit: Index sampling chunk size limit.
- dbms.index\_sampling.update\_percentage: Percentage of index updates of total index size required before sampling of a given index is triggered.
- dbms.index\_searcher\_cache\_size: The maximum number of open Lucene index searchers.
- dbms.jvm.additional: Additional JVM arguments.
- dbms.lock.acquisition.timeout: The maximum time interval within which lock should be acquired.
- dbms.logs.debug.level: Debug log level threshold.
- dbms.logs.debug.path: Path to the debug log file.
- dbms.logs.debug.rotation.delay: Minimum time interval after last rotation of the debug log before it may be rotated again.
- dbms.logs.debug.rotation.keep\_number: Maximum number of history files for the debug log.
- dbms.logs.debug.rotation.size: Threshold for rotation of the debug log.
- dbms.logs.gc.enabled: Enable GC Logging.
- dbms.logs.gc.options: GC Logging Options.
- dbms.logs.gc.rotation.keep\_number: Number of GC logs to keep.
- dbms.logs.gc.rotation.size: Size of each GC log that is kept.
- dbms.logs.http.enabled: Enable HTTP request logging.
- dbms.logs.http.path: Path to HTTP request log.

- dbms.logs.http.rotation.keep\_number: Number of HTTP logs to keep.
- dbms.logs.http.rotation.size: Size of each HTTP log that is kept.
- dbms.logs.query.allocation\_logging\_enabled: Log allocated bytes for the executed queries being logged.
- dbms.logs.query.enabled: Log executed queries.
- dbms.logs.query.page\_logging\_enabled: Log page hits and page faults for the executed queries being logged.
- dbms.logs.query.parameter\_logging\_enabled: Log parameters for the executed queries being logged.
- dbms.logs.query.path: Path to the query log file.
- dbms.logs.query.rotation.keep\_number: Maximum number of history files for the query log.
- dbms.logs.query.rotation.size: The file size in bytes at which the query log will auto-rotate.
- dbms.logs.query.runtime\_logging\_enabled: Logs which runtime that was used to run the query.
- dbms.logs.query.threshold: If the execution of query takes more time than this threshold, the query is logged once completed provided query logging is set to INFO.
- dbms.logs.query.time\_logging\_enabled: Log detailed time information for the executed queries being logged.
- dbms.logs.security.level: Security log level threshold.
- dbms.logs.security.path: Path to the security log file.
- dbms.logs.security.rotation.delay: Minimum time interval after last rotation of the security log before it may be rotated again.
- dbms.logs.security.rotation.keep\_number: Maximum number of history files for the security log.
- dbms.logs.security.rotation.size: Threshold for rotation of the security log.
- dbms.logs.user.path: Path to the user log file.
- dbms.logs.user.rotation.delay: Minimum time interval after last rotation of the user log before it may be rotated again.
- dbms.logs.user.rotation.keep\_number: Maximum number of history files for the user log.
- dbms.logs.user.rotation.size: Threshold for rotation of the user log.
- dbms.logs.user.stdout enabled: Send user logs to the process stdout.
- dbms.max databases: The maximum number of databases.
- dbms.memory.heap.initial\_size: Initial heap size.
- dbms.memory.heap.max\_size: Maximum heap size.
- dbms.memory.pagecache.size: The amount of memory to use for mapping the store files, in bytes (or kilobytes with the 'k' suffix, megabytes with 'm' and gigabytes with 'g').
- dbms.memory.pagecache.swapper: Specify which page swapper to use for doing paged IO.
- dbms.memory.pagecache.warmup.enable: Page cache can be configured to perform usage sampling of loaded pages that can be used to construct active load profile.
- dbms.memory.pagecache.warmup.preload: Page cache warmup can be configured to prefetch files, preferably when cache size is bigger than store size.
- dbms.memory.pagecache.warmup.preload.whitelist: Page cache warmup prefetch file whitelist regex.
- dbms.memory.pagecache.warmup.profile.interval: The profiling frequency for the page cache.
- dbms.mode: Configure the operating mode of the database 'SINGLE' for stand-alone operation, 'CORE' for operating as a core member of a Causal Cluster, or 'READ\_REPLICA' for operating as a

read replica member of a Causal Cluster.

- dbms.netty.ssl.provider: Netty SSL provider.
- dbms.query\_cache\_size: The number of Cypher query execution plans that are cached.
- dbms.read\_only: Only allow read operations from this Neo4j instance.
- dbms.reconciler.max\_backoff: Defines the maximum amount of time to wait before retrying after the dbms fails to reconcile a database to its desired state.
- dbms.reconciler.may\_retry: Defines whether the dbms may retry reconciling a database to its desired state.
- dbms.reconciler.min\_backoff: Defines the minimum amount of time to wait before retrying after the dbms fails to reconcile a database to its desired state.
- dbms.record\_format: Database record format.
- dbms.recovery.fail on missing files: If true, Neo4j will abort recovery if logical log files are missing.
- dbms.relationship\_grouping\_threshold: Relationship count threshold for considering a node to be dense.
- dbms.rest.transaction.idle\_timeout: Timeout for idle transactions in the REST endpoint.
- dbms.routing\_ttl: How long callers should cache the response of the routing procedure dbms.routing.getRoutingTable().
- dbms.security.allow\_csv\_import\_from\_file\_urls: Determines if Cypher will allow using file URLs when loading data using LOAD CSV.
- dbms.security.auth\_cache\_max\_capacity: The maximum capacity for authentication and authorization caches (respectively).
- dbms.security.auth\_cache\_ttl: The time to live (TTL) for cached authentication and authorization info when using external auth providers (LDAP or plugin).
- dbms.security.auth\_cache\_use\_ttl: Enable time-based eviction of the authentication and authorization info cache for external auth providers (LDAP or plugin).
- dbms.security.auth\_enabled: Enable auth requirement to access Neo4j.
- dbms.security.auth\_lock\_time: The amount of time user account should be locked after a configured number of unsuccessful authentication attempts.
- dbms.security.auth\_max\_failed\_attempts: The maximum number of unsuccessful authentication attempts before imposing a user lock for the configured amount of time. The locked out user will not be able to log in until the lock period expires, even if correct credentials are provided.
- dbms.security.authentication\_providers: A list of security authentication providers containing the users and roles.
- dbms.security.authorization\_providers: A list of security authorization providers containing the users and roles.
- dbms.security.causal\_clustering\_status\_auth\_enabled: Require authorization for access to the Causal Clustering status endpoints.
- dbms.security.http\_access\_control\_allow\_origin: Value of the Access-Control-Allow-Origin header sent over any HTTP or HTTPS connector.
- dbms.security.http\_auth\_whitelist: Defines a whitelist of http paths where Neo4j authentication is not required.
- dbms.security.http\_strict\_transport\_security: Value of the HTTP Strict-Transport-Security (HSTS) response header.
- dbms.security.ldap.authentication.cache\_enabled: Determines if the result of authentication via the LDAP server should be cached or not.
- dbms.security.ldap.authentication.mechanism: LDAP authentication mechanism.

- dbms.security.ldap.authentication.use\_samaccountname: Perform authentication with sAMAccountName instead of DN. Using this setting requires dbms.security.ldap.authorization.system\_username and dbms.security.ldap.authorization.system\_password to be used since there is no way to log in through ldap directly with the sAMAccountName, instead the login name will be resolved to a DN that will be used to log in with.
- dbms.security.ldap.authentication.user\_dn\_template: LDAP user DN template.
- dbms.security.ldap.authorization.group\_membership\_attributes: A list of attribute names on a
  user object that contains groups to be used for mapping to roles when LDAP authorization is
  enabled.
- dbms.security.ldap.authorization.group\_to\_role\_mapping: An authorization mapping from LDAP group names to Neo4j role names.
- dbms.security.ldap.authorization.system\_password: An LDAP system account password to use for authorization searches when dbms.security.ldap.authorization.use\_system\_account is true.
- dbms.security.ldap.authorization.system\_username: An LDAP system account username to use for authorization searches when dbms.security.ldap.authorization.use\_system\_account is true.
- dbms.security.ldap.authorization.use\_system\_account: Perform LDAP search for authorization
  info using a system account instead of the user's own account. If this is set to false (default), the
  search for group membership will be performed directly after authentication using the LDAP
  context bound with the user's own account.
- dbms.security.ldap.authorization.user\_search\_base: The name of the base object or named context to search for user objects when LDAP authorization is enabled.
- dbms.security.ldap.authorization.user\_search\_filter: The LDAP search filter to search for a user principal when LDAP authorization is enabled.
- dbms.security.ldap.connection\_timeout: The timeout for establishing an LDAP connection.
- dbms.security.ldap.host: URL of LDAP server to use for authentication and authorization.
- dbms.security.ldap.read\_timeout: The timeout for an LDAP read request (i.e.
- dbms.security.ldap.referral: The LDAP referral behavior when creating a connection.
- dbms.security.ldap.use\_starttls: Use secure communication with the LDAP server using opportunistic TLS.
- dbms.security.log\_successful\_authentication: Set to log successful authentication events to the security log.
- dbms.security.procedures.default\_allowed: The default role that can execute all procedures and user-defined functions that are not covered by the dbms.security.procedures.roles setting.
- dbms.security.procedures.roles: This provides a finer level of control over which roles can execute procedures than the dbms.security.procedures.default\_allowed setting.
- dbms.security.procedures.unrestricted: A list of procedures and user defined functions (comma separated) that are allowed full access to the database.
- dbms.security.procedures.whitelist: A list of procedures (comma separated) that are to be loaded.
- dbms.security.property level.enabled: This has been replaced by privilege management on roles.
- dbms.shutdown\_transaction\_end\_timeout: The maximum amount of time to wait for running transactions to complete before allowing initiated database shutdown to continue.
- dbms.threads.worker\_count: Number of Neo4j worker threads.
- dbms.track\_query\_allocation: Enables or disables tracking of how many bytes are allocated by the execution of a query.
- dbms.track\_query\_cpu\_time: Enables or disables tracking of how much time a query spends actively executing on the CPU.
- dbms.transaction.bookmark\_ready\_timeout: The maximum amount of time to wait for the

database state represented by the bookmark.

- dbms.transaction.concurrent.maximum: The maximum number of concurrently running transactions.
- dbms.transaction.monitor.check.interval: Configures the time interval between transaction monitor checks.
- dbms.transaction.sampling.percentage: Transaction sampling percentage.
- dbms.transaction.timeout: The maximum time interval of a transaction within which it should be completed.
- dbms.transaction.tracing.level: Transaction creation tracing level.
- dbms.tx\_log.preallocate: Specify if Neo4j should try to preallocate logical log file in advance.
- dbms.tx\_log.rotation.retention\_policy: Make Neo4j keep the logical transaction logs for being able to backup the database.
- dbms.tx\_log.rotation.size: Specifies at which file size the logical log will auto-rotate.
- dbms.tx\_state.max\_off\_heap\_memory: The maximum amount of off-heap memory that can be used to store transaction state data; it's a total amount of memory shared across all active transactions.
- dbms.tx\_state.memory\_allocation: Defines whether memory for transaction state should be allocated on- or off-heap.
- dbms.tx\_state.off\_heap.block\_cache\_size: Defines the size of the off-heap memory blocks cache.
- dbms.tx\_state.off\_heap.max\_cacheable\_block\_size: Defines the maximum size of an off-heap memory block that can be cached to speed up allocations for transaction state data.
- dbms.unmanaged\_extension\_classes: Comma-separated list of <classname>=<mount point> for unmanaged extensions.
- dbms.windows service name: Name of the Windows Service.
- fabric.database.name: Name of the Fabric database.
- fabric.driver.api: Determines which driver API will be used.
- fabric.driver.connection.connect\_timeout: Socket connection timeout. A timeout of zero is treated as an infinite timeout and will be bound by the timeout configured on the operating system level.
- fabric.driver.connection.max\_lifetime: Pooled connections older than this threshold will be closed and removed from the pool. Setting this option to a low value will cause a high connection churn and might result in a performance hit. It is recommended to set maximum lifetime to a slightly smaller value than the one configured in network equipment (load balancer, proxy, firewall, etc.
- fabric.driver.connection.pool.acquisition\_timeout: Maximum amount of time spent attempting to acquire a connection from the connection pool. This timeout only kicks in when all existing connections are being used and no new connections can be created because maximum connection pool size has been reached. Error is raised when connection can't be acquired within configured time. Negative values are allowed and result in unlimited acquisition timeout.
- fabric.driver.connection.pool.idle\_test: Pooled connections that have been idle in the pool for longer than this timeout will be tested before they are used again, to ensure they are still alive. If this option is set too low, an additional network call will be incurred when acquiring a connection, which causes a performance hit. If this is set high, no longer live connections might be used which might lead to errors. Hence, this parameter tunes a balance between the likelihood of experiencing connection problems and performance Normally, this parameter should not need tuning. Value 0 means connections will always be tested for validity.
- fabric.driver.connection.pool.max\_size: Maximum total number of connections to be managed by a connection pool. The limit is enforced for a combination of a host and user.
- fabric.driver.logging.level: Sets level for driver internal logging.
- fabric.routing.servers: A comma-separated list of Fabric instances that form a routing group.

- fabric.routing.ttl: The time to live (TTL) of a routing table for fabric routing group.
- fabric.stream.buffer.low\_watermark: Number of records in prefetching buffer that will trigger prefetching again.
- fabric.stream.buffer.size: Maximal size of a buffer used for pre-fetching result records of remote queries. To compensate for latency to remote databases, the Fabric execution engine pre-fetches records needed for local executions. This limit is enforced per fabric query.
- fabric.stream.concurrency: Maximal concurrency within Fabric queries. Limits the number of iterations of each subquery that are executed concurrently.
- metrics.bolt.messages.enabled: Enable reporting metrics about Bolt Protocol message processing.
- metrics.csv.enabled: Set to true to enable exporting metrics to CSV files.
- metrics.csv.interval: The reporting interval for the CSV files.
- metrics.csv.rotation.keep\_number: Maximum number of history files for the csv files.
- metrics.csv.rotation.size: The file size in bytes at which the csv files will auto-rotate.
- metrics.cypher.replanning.enabled: Enable reporting metrics about number of occurred replanning events.
- metrics.enabled: Enable metrics.
- metrics.graphite.enabled: Set to true to enable exporting metrics to Graphite.
- metrics.graphite.interval: The reporting interval for Graphite.
- metrics.graphite.server: The hostname or IP address of the Graphite server.
- metrics.jmx.enabled: Set to true to enable the JMX metrics endpoint.
- metrics.jvm.buffers.enabled: Enable reporting metrics about the buffer pools.
- metrics.jvm.file.descriptors.enabled: Enable reporting metrics about the number of open file descriptors.
- metrics.jvm.gc.enabled: Enable reporting metrics about the duration of garbage collections.
- metrics.jvm.heap.enabled: Enable reporting metrics about the heap memory usage.
- metrics.jvm.memory.enabled: Enable reporting metrics about the memory usage.
- metrics.jvm.threads.enabled: Enable reporting metrics about the current number of threads running.
- metrics.neo4j.causal\_clustering.enabled: Enable reporting metrics about Causal Clustering mode.
- metrics.neo4j.checkpointing.enabled: Enable reporting metrics about Neo4j check pointing.
- metrics.neo4j.counts.enabled: Enable reporting metrics about approximately how many entities are in the database.
- metrics.neo4j.data.counts.enabled: Enable reporting metrics about number of entities in the database.
- metrics.neo4j.logs.enabled: Enable reporting metrics about the Neo4j transaction logs.
- metrics.neo4j.pagecache.enabled: Enable reporting metrics about the Neo4j page cache.
- metrics.neo4j.server.enabled: Enable reporting metrics about Server threading info.
- metrics.neo4j.size.enabled: Enable reporting metrics about the store size of each database.
- metrics.neo4j.tx.enabled: Enable reporting metrics about transactions.
- metrics.prefix: A common prefix for the reported metrics field names.
- metrics.prometheus.enabled: Set to true to enable the Prometheus endpoint.
- metrics.prometheus.endpoint: The hostname and port to use as Prometheus endpoint.

#### Table 43. browser.allow\_outgoing\_connections

Description	Configure the policy for outgoing Neo4j Browser connections.
Valid values	browser.allow_outgoing_connections, a boolean
Default value	true

# Table 44. browser.credential\_timeout

Description	Configure the Neo4j Browser to time out logged in users after this idle period. Setting this to 0 indicates no limit.
Valid values	browser.credential_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	0ns

#### Table 45. browser.post\_connect\_cmd

Description	Commands to be run when Neo4j Browser successfully connects to this server. Separate multiple commands with semi-colon.
Valid values	browser.post_connect_cmd, a string
Default value	

# Table 46. browser.remote\_content\_hostname\_whitelist

Description	Whitelist of hosts for the Neo4j Browser to be allowed to fetch content from.
Valid values	browser.remote_content_hostname_whitelist, a string
Default value	guides.neo4j.com,localhost

#### Table 47. browser.retain\_connection\_credentials

Description	Configure the Neo4j Browser to store or not store user credentials.
Valid values	browser.retain_connection_credentials, a boolean
Default value	true

# Table 48. causal\_clustering.catch\_up\_client\_inactivity\_timeout

Description	The catch up protocol times out if the given duration elapses with no network activity. Every message received by the client from the server extends the time out duration.
Valid values	causal_clustering.catch_up_client_inactivity_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10m

#### Table 49. causal\_clustering.catchup\_batch\_size

Description	The maximum batch size when catching up (in unit of entries)
Valid values	causal_clustering.catchup_batch_size, an integer
Default value	64

#### Table 50. causal\_clustering.cluster\_allow\_reads\_on\_followers

Description	Configure if the `dbms.routing.getRoutingTable()` procedure should include followers as read endpoints or return only read replicas. Note: if there are no read replicas in the cluster, followers are returned as read end points regardless the value of this setting. Defaults to true so that followers are available for read-only queries in a typical heterogeneous setup.
Valid values	causal_clustering.cluster_allow_reads_on_followers, a boolean
Default value	true

# Table 51. causal\_clustering.cluster\_topology\_refresh

Description	Time between scanning the cluster to refresh current server's view of topology.
Valid values	causal_clustering.cluster_topology_refresh, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is minimum PT1S
Default value	5s

#### Table 52. causal\_clustering.connect-randomly-to-server-group

Description	+Comma separated list of groups to be used by the connect-randomly-to-server-group selection strategy. The connect-randomly-to-server-group strategy is used if the list of strategies (causal_clustering.upstream_selection_strategy) includes the value connect-randomly-to-server-group.
Valid values	causal_clustering.connect-randomly-to-server-group, a ',' separated list with elements of type 'a string'.
Default value	

# Table 53. causal\_clustering.discovery\_advertised\_address

Description	Advertised cluster member discovery management communication.
Valid values	causal_clustering.discovery_advertised_address, a socket address. If missing port or hostname it is acquired from dbms.default_advertised_address
Default value	:5000

#### Table 54. causal\_clustering.discovery\_listen\_address

Description	Host and port to bind the cluster member discovery management communication.
Valid values	causal_clustering.discovery_listen_address, a socket address. If missing port or hostname it is acquired from dbms.default_listen_address
Default value	:5000

# Table 55. causal\_clustering.discovery\_type

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Description	Configure the discovery type used for cluster name resolution.
Valid values	causal_clustering.discovery_type, one of [DNS, LIST, SRV, K8S]
Default value	LIST

#### Table 56. causal\_clustering.enable\_pre\_voting

Description	Enable pre-voting extension to the Raft protocol (this is breaking and must match between the core cluster members)
Valid values	causal_clustering.enable_pre_voting, a boolean
Default value	true

# Table 57. causal\_clustering.global\_session\_tracker\_state\_size

Description	The maximum file size before the global session tracker state file is rotated (in unit of entries)
Valid values	causal_clustering.global_session_tracker_state_size, an integer
Default value	1000

# Table 58. causal\_clustering.handshake\_timeout

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Description	Time out for protocol negotiation handshake.
Valid values	causal_clustering.handshake_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	20s

#### Table 59. causal\_clustering.in\_flight\_cache.max\_bytes

Description	The maximum number of bytes in the in-flight cache.
Valid values	causal_clustering.in_flight_cache.max_bytes, a byte size (valid multipliers are $k$ , $m$ , $g$ , $K$ , $M$ , $G$ )
Default value	2147483648

#### Table 60. causal\_clustering.in\_flight\_cache.max\_entries

Description	The maximum number of entries in the in-flight cache.
Valid values	causal_clustering.in_flight_cache.max_entries, an integer
Default value	1024

#### Table 61. causal\_clustering.in\_flight\_cache.type

Description	Type of in-flight cache.
Valid values	causal_clustering.in_flight_cache.type, one of [NONE, CONSECUTIVE, UNBOUNDED]
Default value	CONSECUTIVE

#### Table 62. causal\_clustering.initial\_discovery\_members

Description	A comma-separated list of other members of the cluster to join.
Valid values	causal_clustering.initial_discovery_members, a ',' separated list with elements of type 'a socket address'.

#### Table 63. causal\_clustering.join\_catch\_up\_timeout

Description	Time out for a new member to catch up.
Valid values	causal_clustering.join_catch_up_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10m

#### Table 64. causal\_clustering.kubernetes.address

Description	Address for Kubernetes API.
Valid values	causal_clustering.kubernetes.address, a socket address
Default value	kubernetes.default.svc:443

# Table 65. causal\_clustering.kubernetes.ca\_crt

Description	File location of CA certificate for Kubernetes API.
Valid values	causal_clustering.kubernetes.ca_crt, a path
Default value	/var/run/secrets/kubernetes.io/serviceaccount/ca.crt

#### Table 66. causal\_clustering.kubernetes.label\_selector

Description	LabelSelector for Kubernetes API.

Valid values	causal_clustering.kubernetes.label_selector, a string

#### Table 67. causal\_clustering.kubernetes.namespace

Description	File location of namespace for Kubernetes API.
Valid values	causal_clustering.kubernetes.namespace, a path
Default value	/var/run/secrets/kubernetes.io/serviceaccount/namespace

#### Table 68. causal\_clustering.kubernetes.service\_port\_name

Description	Service port name for discovery for Kubernetes API.
Valid values	causal_clustering.kubernetes.service_port_name, a string

#### Table 69. causal\_clustering.kubernetes.token

Description	File location of token for Kubernetes API.
Valid values	causal_clustering.kubernetes.token, a path
Default value	/var/run/secrets/kubernetes.io/serviceaccount/token

#### Table 70. causal\_clustering.last\_applied\_state\_size

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Description	The maximum file size before the storage file is rotated (in unit of entries)
Valid values	causal_clustering.last_applied_state_size, an integer
Default value	1000

#### Table 71. causal\_clustering.leader\_election\_timeout

Description	The time limit within which a new leader election will occur if no messages are received.
Valid values	causal_clustering.leader_election_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	7s

# Table 72. causal\_clustering.load\_balancing.plugin

Description	The load balancing plugin to use.
Valid values	causal_clustering.load_balancing.plugin, a string
Default value	server_policies

#### Table 73. causal\_clustering.load\_balancing.shuffle

Description	Enables shuffling of the returned load balancing result.
Valid values	causal_clustering.load_balancing.shuffle, a boolean
Default value	true

# Table 74. causal\_clustering.log\_shipping\_max\_lag

Description	The maximum lag allowed before log shipping pauses (in unit of entries)
Valid values	causal_clustering.log_shipping_max_lag, an integer
Default value	256

# Table 75. causal\_clustering.middleware.logging.level

Description	The level of middleware logging.
Valid values	causal_clustering.middleware.logging.level, one of [DEBUG, INFO, WARN, ERROR, NONE]
Default value	WARN

#### Table 76. causal\_clustering.minimum\_core\_cluster\_size\_at\_formation

Description	Minimum number of Core machines initially required to form a cluster. The cluster will form when at least this many Core members have discovered each other.
Valid values	causal_clustering.minimum_core_cluster_size_at_formation, an integer which is minimum 2
Default value	3

# Table 77. causal\_clustering.minimum\_core\_cluster\_size\_at\_runtime

Description	The minimum size of the dynamically adjusted voting set (which only core members may be a part of). Adjustments to the voting set happen automatically as the availability of core members changes, due to explicit operations such as starting or stopping a member, or unintended issues such as network partitions. Note that this dynamic scaling of the voting set is generally desirable as under some circumstances it can increase the number of instance failures which may be tolerated. A majority of the voting set must be available before voting in or out members.
Valid values	causal_clustering.minimum_core_cluster_size_at_runtime, an integer which is minimum 2
Default value	3

# Table 78. causal\_clustering.multi\_dc\_license

Description	Enable multi-data center features. Requires appropriate licensing.
Valid values	causal_clustering.multi_dc_license, a boolean
Default value	false

# Table 79. causal\_clustering.protocol\_implementations.catchup

Description	Catchup protocol implementation versions that this instance will allow in negotiation as a comma-separated list. Order is not relevant: the greatest value will be preferred. An empty list will allow all supported versions. Example value: "1.1, 1.2, 2.1, 2.2"
Valid values	causal_clustering.protocol_implementations.catchup, a ',' separated list with elements of type 'an application protocol version'.
Default value	

# $Table~80.~causal\_clustering.protocol\_implementations.compression$

Description	Network compression algorithms that this instance will allow in negotiation as a comma-separated list. Listed in descending order of preference for incoming connections. An empty list implies no compression. For outgoing connections this merely specifies the allowed set of algorithms and the preference of the remote peer will be used for making the decision. Allowable values: [Gzip, Snappy, Snappy_validating, LZ4, LZ4_high_compression, LZ_validating, LZ4_high_compression_validating]
Valid values	causal_clustering.protocol_implementations.compression, a ',' separated list with elements of type 'a string'.
Default value	

# Table 81. causal\_clustering.protocol\_implementations.raft

Description	Raft protocol implementation versions that this instance will allow in negotiation as a comma-separated list. Order is not relevant: the greatest value will be preferred. An empty list will allow all supported versions. Example value: "1.0, 1.3, 2.0, 2.1"
Valid values	causal_clustering.protocol_implementations.raft, a ',' separated list with elements of type 'an application protocol version'.
Default value	

# Table 82. causal\_clustering.pull\_interval

Description	Interval of pulling updates from cores.
Valid values	causal_clustering.pull_interval, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1s

#### Table 83. causal\_clustering.raft\_advertised\_address

Description	Advertised hostname/IP address and port for the RAFT server.
Valid values	causal_clustering.raft_advertised_address, a socket address. If missing port or hostname it is acquired from dbms.default_advertised_address
Default value	:7000

#### Table 84. causal\_clustering.raft\_in\_queue\_max\_batch\_bytes

Description	Largest batch processed by RAFT in bytes.
Valid values	causal_clustering.raft_in_queue_max_batch_bytes, a byte size (valid multipliers are k, m, g, K, M, G)
Default value	8388608

# Table 85. causal\_clustering.raft\_in\_queue\_max\_bytes

Description	Maximum number of bytes in the RAFT in-queue.
Valid values	causal_clustering.raft_in_queue_max_bytes, a byte size (valid multipliers are $k$ , $m$ , g, $K$ , M, G)
Default value	2147483648

#### Table 86. causal\_clustering.raft\_listen\_address

Description	Network interface and port for the RAFT server to listen on.
Valid values	causal_clustering.raft_listen_address, a socket address. If missing port or hostname it is acquired from dbms.default_listen_address
Default value	:7000

#### Table 87. causal\_clustering.raft\_log\_implementation

Description	RAFT log implementation.
Valid values	causal_clustering.raft_log_implementation, a string
Default value	SEGMENTED

#### Table 88. causal\_clustering.raft\_log\_prune\_strategy

Description	RAFT log pruning strategy.
Valid values	causal_clustering.raft_log_prune_strategy, a string
Default value	1g size

#### Table 89. causal\_clustering.raft\_log\_pruning\_frequency

Description	RAFT log pruning frequency.
Valid values	causal_clustering.raft_log_pruning_frequency, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10m

# Table 90. causal\_clustering.raft\_log\_reader\_pool\_size

Description	RAFT log reader pool size.
Valid values	causal_clustering.raft_log_reader_pool_size, an integer
Default value	8

#### Table 91. causal\_clustering.raft\_log\_rotation\_size

Description	RAFT log rotation size.
Valid values	causal_clustering.raft_log_rotation_size, a byte size (valid multipliers are k, m, g, K, M, G) which is minimum 1024
Default value	262144000

# Table 92. causal\_clustering.raft\_membership\_state\_size

Description	The maximum file size before the membership state file is rotated (in unit of entries)
Valid values	causal_clustering.raft_membership_state_size, an integer
Default value	1000

#### Table 93. causal\_clustering.raft\_term\_state\_size

Description	The maximum file size before the term state file is rotated (in unit of entries)
Valid values	causal_clustering.raft_term_state_size, an integer
Default value	1000

#### Table 94. causal\_clustering.raft\_vote\_state\_size

Description	The maximum file size before the vote state file is rotated (in unit of entries)
Valid values	causal_clustering.raft_vote_state_size, an integer
Default value	1000

#### Table 95. causal\_clustering.refuse\_to\_be\_leader

Description	Prevents the current instance from volunteering to become Raft leader. Defaults to false, and should only be used in exceptional circumstances by expert users. Using this can result in reduced availability for the cluster.
Valid values	causal_clustering.refuse_to_be_leader, a boolean
Default value	false

# Table 96. causal\_clustering.replicated\_lease\_state\_size

Description	The maximum file size before the replicated lease state file is rotated (in unit of entries)
Valid values	causal_clustering.replicated_lease_state_size, an integer
Default value	1000

# Table 97. causal\_clustering.replication\_leader\_await\_timeout

Description	The duration for which the replicator will await a new leader.
Valid values	causal_clustering.replication_leader_await_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10s

#### Table 98. causal\_clustering.replication\_retry\_timeout\_base

Description	The initial timeout until replication is retried. The timeout will increase exponentially.
Valid values	causal_clustering.replication_retry_timeout_base, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10s

#### Table 99. causal\_clustering.replication\_retry\_timeout\_limit

Description	The upper limit for the exponentially incremented retry timeout.
Valid values	causal_clustering.replication_retry_timeout_limit, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1m

# Table 100. causal\_clustering.server\_groups

Description	A list of group names for the server used when configuring load balancing and replication policies.
Valid values	causal_clustering.server_groups, a ',' separated list with elements of type 'a string'.

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#### Table 101. causal\_clustering.state\_machine\_apply\_max\_batch\_size

Description	The maximum number of operations to be batched during applications of operations in the state machines.
Valid values	causal_clustering.state_machine_apply_max_batch_size, an integer
Default value	16

# Table 102. causal\_clustering.state\_machine\_flush\_window\_size

Description	The number of operations to be processed before the state machines flush to disk.
Valid values	causal_clustering.state_machine_flush_window_size, an integer
Default value	4096

#### Table 103. causal\_clustering.status\_throughput\_window

Description	Sampling window for throughput estimate reported in the status endpoint.
Valid values	causal_clustering.status_throughput_window, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is in the range PT1S to PT5M
Default value	5s

#### Table 104. causal\_clustering.store\_copy\_chunk\_size

Description	Store copy chunk size.
Valid values	causal_clustering.store_copy_chunk_size, an integer which is in the range 4096 to 1048576
Default value	32768

#### Table 105. causal\_clustering.store\_copy\_max\_retry\_time\_per\_request

Description	+Maximum retry time per request during store copy. Regular store files and indexes are downloaded in separate requests during store copy. This configures the maximum time failed requests are allowed to resend.
Valid values	causal_clustering.store_copy_max_retry_time_per_request, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	20m

#### Table 106. causal\_clustering.transaction\_advertised\_address

Description	Advertised hostname/IP address and port for the transaction shipping server.

Valid values	causal_clustering.transaction_advertised_address, a socket address. If missing port or hostname it is acquired from dbms.default_advertised_address
Default value	:6000

#### Table 107. causal\_clustering.transaction\_listen\_address

Description	Network interface and port for the transaction shipping server to listen on. Please note that it is also possible to run the backup client against this port so always limit access to it via the firewall and configure an ssl policy.
Valid values	causal_clustering.transaction_listen_address, a socket address. If missing port or hostname it is acquired from dbms.default_listen_address
Default value	:6000

#### Table 108. causal\_clustering.unknown\_address\_logging\_throttle

Description	Throttle limit for logging unknown cluster member address.
Valid values	causal_clustering.unknown_address_logging_throttle, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10s

#### Table 109. causal\_clustering.upstream\_selection\_strategy

Description	An ordered list in descending preference of the strategy which read replicas use to choose the upstream server from which to pull transactional updates.
Valid values	causal_clustering.upstream_selection_strategy, a ',' separated list with elements of type 'a string'.
Default value	default

# Table 110. causal\_clustering.user\_defined\_upstream\_strategy

Description	+Configuration of a user-defined upstream selection strategy. The user-defined strategy is used if the list of strategies (causal_clustering.upstream_selection_strategy) includes the value user_defined.
Valid values	causal_clustering.user_defined_upstream_strategy, a string
Default value	

# Table 111. cypher.default\_language\_version

Description	Set this to specify the default parser (language version).
Valid values	cypher.default_language_version, one of [default, 3.5, 4.0]
Default value	default

Table 112. cypher.forbid\_exhaustive\_shortestpath

Description	This setting is associated with performance optimization. Set this to `true` in situations where it is preferable to have any queries using the 'shortestPath' function terminate as soon as possible with no answer, rather than potentially running for a long time attempting to find an answer (even if there is no path to be found). For most queries, the 'shortestPath' algorithm will return the correct answer very quickly. However there are some cases where it is possible that the fast bidirectional breadth-first search algorithm will find no results even if they exist. This can happen when the predicates in the `WHERE` clause applied to 'shortestPath' cannot be applied to each step of the traversal, and can only be applied to the entire path. When the query planner detects these special cases, it will plan to perform an exhaustive depth-first search if the fast algorithm finds no paths. However, the exhaustive search may be orders of magnitude slower than the fast algorithm. If it is critical that queries terminate as soon as possible, it is recommended that this option be set to `true`, which means that Neo4j will never consider using the exhaustive search for shortestPath queries. However, please note that if no paths are found, an error will be thrown at run time, which will need to be handled by the application.
Valid values	cypher.forbid_exhaustive_shortestpath, a boolean
Default value	false

Table 113. cypher.forbid\_shortestpath\_common\_nodes

Description	This setting is associated with performance optimization. The shortest path algorithm does not work when the start and end nodes are the same. With this setting set to `false` no path will be returned when that happens. The default value of `true` will instead throw an exception. This can happen if you perform a shortestPath search after a cartesian product that might have the same start and end nodes for some of the rows passed to shortestPath. If it is preferable to not experience this exception, and acceptable for results to be missing for those rows, then set this to `false`. If you cannot accept missing results, and really want the shortestPath between two common nodes, then re-write the query using a standard Cypher variable length pattern expression followed by ordering by path length and limiting to one result.
Valid values	cypher.forbid_shortestpath_common_nodes, a boolean
Default value	true

# Table 114. cypher.hints\_error

Description	Set this to specify the behavior when Cypher planner or runtime hints cannot be fulfilled. If true, then non-conformance will result in an error, otherwise only a warning is generated.
Valid values	cypher.hints_error, a boolean
Default value	false

Table 115. cypher.lenient\_create\_relationship

Description	Set this to change the behavior for Cypher create relationship when the start or end node is missing. By default this fails the query and stops execution, but by setting this flag the create operation is simply not performed and execution continues.
Valid values	cypher.lenient_create_relationship, a boolean
Default value	false

# Table 116. cypher.min\_replan\_interval

Description	The minimum time between possible cypher query replanning events. After this time, the graph statistics will be evaluated, and if they have changed by more than the value set by < <config_cypher.statistics_divergence_threshold,cypher.statistics_divergence_th reshold="">&gt;, the query will be replanned. If the statistics have not changed sufficiently, the same interval will need to pass before the statistics will be evaluated again. Each time they are evaluated, the divergence threshold will be reduced slightly until it reaches 10% after 7h, so that even moderately changing databases will see query replanning after a sufficiently long time interval.</config_cypher.statistics_divergence_threshold,cypher.statistics_divergence_th>
Valid values	cypher.min_replan_interval, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10s

# Table 117. cypher.planner

Description	Set this to specify the default planner for the default language version.
Valid values	cypher.planner, one of [DEFAULT, COST]
Default value	DEFAULT

# Table 118. cypher.query\_max\_allocations

Description	The maximum amount of heap memory allocations to for cypher to perform on a single query, in bytes (or kilobytes with the 'k' suffix, megabytes with 'm' and gigabytes with 'g'). Zero means 'unlimited'. If a query exceeds this limit, it will be terminated. Determining the heap memory allocations done by a query is a rough estimate and not an exact measurement. If no memory limit is configured, queries will be allowed to allocate as much heap memory as needed. This could potentially lead to queries consuming more heap memory than available, which will kill the Neo4j server.
Valid values	cypher.query_max_allocations, a byte size (valid multipliers are k, m, g, K, M, G) which is minimum 0
Dynamic	true
Default value	0

Table 119. cypher.statistics\_divergence\_threshold

Description	The threshold when a plan is considered stale. If any of the underlying statistics used to create the plan have changed more than this value, the plan will be considered stale and will be replanned. Change is calculated as abs(a-b)/max(a,b). This means that a value of 0.75 requires the database to approximately quadruple in size. A value of 0 means replan as soon as possible, with the soonest being defined by the <config_cypher.min_replan_interval,cypher.min_replan_interval>&gt; which defaults to 10s. After this interval the divergence threshold will slowly start to decline, reaching 10% after about 7h. This will ensure that long running databases will still get query replanning on even modest changes, while not replanning frequently unless the changes are very large.</config_cypher.min_replan_interval,cypher.min_replan_interval>
Valid values	cypher.statistics_divergence_threshold, a double which is in the range 0.0 to 1.0
Default value	0.75

#### Table 120. db.temporal.timezone

Description	Database timezone for temporal functions. All Time and DateTime values that are created without an explicit timezone will use this configured default timezone.
Valid values	db.temporal.timezone, a string describing a timezone, either described by offset (e.g. '+02:00') or by name (e.g. 'Europe/Stockholm')
Default value	Z

#### Table 121. dbms.allow\_upgrade

Description	Whether to allow an upgrade in case the current version of the database starts against an older version.
Valid values	dbms.allow_upgrade, a boolean
Default value	false

# Table 122. dbms.backup.enabled

Description	Enable support for running online backups.
Valid values	dbms.backup.enabled, a boolean
Default value	true

#### Table 123. dbms.backup.listen\_address

Description	Network interface and port for the backup server to listen on.
Valid values	dbms.backup.listen_address, a socket address
Default value	127.0.0.1:6362

# Table 124. dbms.checkpoint

Description	Configures the general policy for when check-points should occur. The default policy is the 'periodic' check-point policy, as specified by the '< <config_dbms.checkpoint.interval.tx,dbms.checkpoint.interval.tx>&gt;' and '&lt;<config_dbms.checkpoint.interval.time,dbms.checkpoint.interval.time>&gt;' settings. The Neo4j Enterprise Edition provides two alternative policies: The first is the 'continuous' check-point policy, which will ignore those settings and run the check-point process all the time. The second is the 'volumetric' check-point policy, which makes a best-effort at check-pointing often enough so that the database doesn't get too far behind on deleting old transaction logs in accordance with the '&lt;<config_dbms.tx_log.rotation.retention_policy,dbms.tx_log.rotation.retention_policy>&gt;' setting.</config_dbms.tx_log.rotation.retention_policy,dbms.tx_log.rotation.retention_policy></config_dbms.checkpoint.interval.time,dbms.checkpoint.interval.time></config_dbms.checkpoint.interval.tx,dbms.checkpoint.interval.tx>
Valid values	dbms.checkpoint, one of [PERIODIC, CONTINUOUS, VOLUMETRIC]
Default value	PERIODIC

# Table 125. dbms.checkpoint.interval.time

Description	Configures the time interval between check-points. The database will not check-point more often than this (unless check pointing is triggered by a different event), but might check-point less often than this interval, if performing a check-point takes longer time than the configured interval. A check-point is a point in the transaction logs, from which recovery would start from. Longer check-point intervals typically means that recovery will take longer to complete in case of a crash. On the other hand, a longer check-point interval can also reduce the I/O load that the database places on the system, as each check-point implies a flushing and forcing of all the store files.
Valid values	dbms.checkpoint.interval.time, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	15m

#### Table 126. dbms.checkpoint.interval.tx

Description	Configures the transaction interval between check-points. The database will not check-point more often than this (unless check pointing is triggered by a different event), but might check-point less often than this interval, if performing a check-point takes longer time than the configured interval. A check-point is a point in the transaction logs, from which recovery would start from. Longer check-point intervals typically means that recovery will take longer to complete in case of a crash. On the other hand, a longer check-point interval can also reduce the I/O load that the database places on the system, as each check-point implies a flushing and forcing of all the store files. The default is '100000' for a check-point every 100000 transactions.
Valid values	dbms.checkpoint.interval.tx, an integer which is minimum 1
Default value	100000

Table 127. dbms.checkpoint.iops.limit

Description	Limit the number of IOs the background checkpoint process will consume per second. This setting is advisory, is ignored in Neo4j Community Edition, and is followed to best effort in Enterprise Edition. An IO is in this case a 8 KiB (mostly sequential) write. Limiting the write IO in this way will leave more bandwidth in the IO subsystem to service random-read IOs, which is important for the response time of queries when the database cannot fit entirely in memory. The only drawback of this setting is that longer checkpoint times may lead to slightly longer recovery times in case of a database or system crash. A lower number means lower IO pressure, and consequently longer checkpoint times. Set this to -1 to disable the IOPS limit and remove the limitation entirely; this will let the checkpointer flush data as fast as the hardware will go. Removing the setting, or commenting it out, will set the default value of 300.
Valid values	dbms.checkpoint.iops.limit, an integer
Dynamic	true
Default value	300

# Table 128. dbms.config.strict\_validation

Description	A strict configuration validation will prevent the database from starting up if unknown configuration options are specified in the neo4j settings namespace (such as dbms., cypher., etc).
Valid values	dbms.config.strict_validation, a boolean
Default value	false

#### Table 129. dbms.connector.bolt.advertised\_address

Description	Advertised address for this connector.
Valid values	dbms.connector.bolt.advertised_address, a socket address. If missing port or hostname it is acquired from dbms.default_advertised_address
Default value	:7687

#### Table 130. dbms.connector.bolt.enabled

Description	Enable the bolt connector.
Valid values	dbms.connector.bolt.enabled, a boolean
Default value	true

# Table 131. dbms.connector.bolt.listen\_address

Description	Address the connector should bind to.
Valid values	dbms.connector.bolt.listen_address, a socket address. If missing port or hostname it is acquired from dbms.default_listen_address

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#### Table 132. dbms.connector.bolt.thread\_pool\_keep\_alive

Description	The maximum time an idle thread in the thread pool bound to this connector will wait for new tasks.
Valid values	dbms.connector.bolt.thread_pool_keep_alive, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	5m

#### Table 133. dbms.connector.bolt.thread\_pool\_max\_size

Description	The maximum number of threads allowed in the thread pool bound to this connector.
Valid values	dbms.connector.bolt.thread_pool_max_size, an integer
Default value	400

#### Table 134. dbms.connector.bolt.thread\_pool\_min\_size

Description	The number of threads to keep in the thread pool bound to this connector, even if they are idle.
Valid values	dbms.connector.bolt.thread_pool_min_size, an integer
Default value	5

#### Table 135. dbms.connector.bolt.tls\_level

Description	Encryption level to require this connector to use.
Valid values	dbms.connector.bolt.tls_level, one of [REQUIRED, OPTIONAL, DISABLED]
Default value	DISABLED

#### Table 136. dbms.connector.bolt.unsupported\_thread\_pool\_shutdown\_wait\_time

Description	The maximum time to wait for the thread pool to finish processing its pending jobs and shutdown.
Valid values	dbms.connector.bolt.unsupported_thread_pool_shutdown_wait_time, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	5s

#### Table 137. dbms.connector.http.advertised\_address

Description	Advertised address for this connector.

Valid values	dbms.connector.http.advertised_address, a socket address. If missing port or hostname it is acquired from dbms.default_advertised_address
Default value	:7474

#### Table 138. dbms.connector.http.enabled

Description	Enable the http connector.
Valid values	dbms.connector.http.enabled, a boolean
Default value	true

#### Table 139. dbms.connector.http.listen\_address

Description	Address the connector should bind to.
Valid values	dbms.connector.http.listen_address, a socket address. If missing port or hostname it is acquired from dbms.default_listen_address
Default value	:7474

# Table 140. dbms.connector.https.advertised\_address

Description	Advertised address for this connector.
Valid values	dbms.connector.https.advertised_address, a socket address. If missing port or hostname it is acquired from dbms.default_advertised_address
Default value	:7473

#### Table 141. dbms.connector.https.enabled

Description	Enable the https connector.
Valid values	dbms.connector.https.enabled, a boolean
Default value	false

# Table 142. dbms.connector.https.listen\_address

Description	Address the connector should bind to.
Valid values	dbms.connector.https.listen_address, a socket address. If missing port or hostname it is acquired from dbms.default_listen_address
Default value	:7473

# Table 143. dbms.db.timezone

Description	Database timezone. Among other things, this setting influences which timezone the logs and monitoring procedures use.
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Valid values	dbms.db.timezone, one of [UTC, SYSTEM]
Default value	<b>UTC</b>

# Table 144. dbms.default\_advertised\_address

Description	Default hostname or IP address the server uses to advertise itself.
Valid values	dbms.default_advertised_address, a socket address which has no specified port
Default value	localhost

# Table 145. dbms.default\_database

Description	Name of the default database.
Valid values	dbms.default_database, A valid database name. Containing only alphabetic characters, numbers, dots and dashes, with a length between 3 and 63 characters. It should be starting with an alphabetic character but not with the name 'system'.
Default value	neo4j

# Table 146. dbms.default\_listen\_address

Description	+Default network interface to listen for incoming connections. To listen for connections on all interfaces, use "0.0.0.0".
Valid values	dbms.default_listen_address, a socket address which has no specified port
Default value	localhost

#### Table 147. dbms.directories.data

Description	Path of the data directory. You must not configure more than one Neo4j installation to use the same data directory.
Valid values	dbms.directories.data, a path. If relative it is resolved from dbms.directories.neo4j_home
Default value	data

# Table 148. dbms.directories.import

Description	Sets the root directory for file URLs used with the Cypher `LOAD CSV` clause. This should be set to a directory relative to the Neo4j installation path, restricting access to only those files within that directory and its subdirectories. For example the value "import" will only enable access to files within the 'import' folder. Removing this setting will disable the security feature, allowing all files in the local system to be imported. Setting this to an empty field will allow access to all files within the Neo4j installation folder.
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dbms.directories.import, a path. If relative it is resolved from dbms.directories.neo4j_home

#### Table 149. dbms.directories.lib

Description	Path of the lib directory.
Valid values	dbms.directories.lib, a path. If relative it is resolved from dbms.directories.neo4j_home
Default value	lib

# Table 150. dbms.directories.logs

Description	Path of the logs directory.
Valid values	dbms.directories.logs, a path. If relative it is resolved from dbms.directories.neo4j_home
Default value	logs

#### Table 151. dbms.directories.metrics

Description	The target location of the CSV files: a path to a directory wherein a CSV file per reported field will be written.
Valid values	dbms.directories.metrics, a path. If relative it is resolved from dbms.directories.neo4j_home
Default value	metrics

# Table 152. dbms.directories.neo4j\_home

Description	Root relative to which directory settings are resolved.
Valid values	dbms.directories.neo4j_home, a path which is absolute
Default value	Defaults to current working directory

# Table 153. dbms.directories.plugins

Description	Location of the database plugin directory. Compiled Java JAR files that contain database procedures will be loaded if they are placed in this directory.
Valid values	dbms.directories.plugins, a path. If relative it is resolved from dbms.directories.neo4j_home
Default value	plugins

#### Table 154. dbms.directories.run

Description	Path of the run directory. This directory holds Neo4j's runtime state, such as a pidfile when it is running in the background. The pidfile is created when starting neo4j and removed when stopping it. It may be placed on an in-memory filesystem such as tmpfs.
Valid values	dbms.directories.run, a path. If relative it is resolved from dbms.directories.neo4j_home
Default value	run

#### Table 155. dbms.directories.transaction.logs.root

Description	Root location where Neo4j will store transaction logs for configured databases.
Valid values	dbms.directories.transaction.logs.root, a path. If relative it is resolved from dbms.directories.data
Default value	transactions

# Table 156. dbms.dynamic.setting.whitelist

Description	A list of setting name patterns (comma separated) that are allowed to be dynamically changed. The list may contain both full setting names, and partial names with the wildcard '*'. If this setting is left empty all dynamic settings updates will be blocked.
Valid values	dbms.dynamic.setting.whitelist, a ',' separated list with elements of type 'a string'.
Default value	*

# Table 157. dbms.filewatcher.enabled

Description	Allows the enabling or disabling of the file watcher service. This is an auxiliary service but should be left enabled in almost all cases.
Valid values	dbms.filewatcher.enabled, a boolean
Default value	true

# Table 158. dbms.import.csv.buffer\_size

Description	The size of the internal buffer in bytes used by `LOAD CSV`. If the csv file contains huge fields this value may have to be increased.
Valid values	dbms.import.csv.buffer_size, a long which is minimum 1
Default value	2097152

# Table 159. dbms.import.csv.legacy\_quote\_escaping

Description	Selects whether to conform to the standard https://tools.ietf.org/html/rfc4180 for interpreting escaped quotation characters in CSV files loaded using `LOAD CSV`. Setting this to `false` will use the standard, interpreting repeated quotes """ as a single in-lined quote, while `true` will use the legacy convention originally supported in Neo4j 3.0 and 3.1, allowing a backslash to include quotes in-lined in fields.
Valid values	dbms.import.csv.legacy_quote_escaping, a boolean
Default value	true

### Table 160. dbms.index.default\_schema\_provider

Description	Index provider to use for newly created schema indexes. An index provider may store different value types in separate physical indexes. native-btree-1.0: All value types and arrays of all value types, even composite keys, are stored in one native index. lucene+native-3.0: Like native-btree-1.0 but single property strings are stored in Lucene. A native index has faster updates, less heap and CPU usage compared to a Lucene index. A native index has some limitations around key size and slower execution of CONTAINS and ENDS WITH string index queries, compared to a Lucene index. Deprecated: Which index provider to use will be a fully internal concern.
Valid values	dbms.index.default_schema_provider, a string
Default value	native-btree-1.0
Deprecated	The dbms.index.default_schema_provider configuration setting has been deprecated.

## Table 161. dbms.index.fulltext.default\_analyzer

Description	The name of the analyzer that the fulltext indexes should use by default.
Valid values	dbms.index.fulltext.default_analyzer, a string
Default value	standard-no-stop-words

### Table 162. dbms.index.fulltext.eventually\_consistent

Description	Whether or not fulltext indexes should be eventually consistent by default or not.
Valid values	dbms.index.fulltext.eventually_consistent, a boolean
Default value	false

## $\textit{Table 163. dbms.} index. \textit{fulltext.eventually\_consistent\_index\_update\_queue\_max\_length}$

Description	The eventually_consistent mode of the fulltext indexes works by queueing up index updates to be applied later in a background thread. This newBuilder sets an upper bound on how many index updates are allowed to be in this queue at any one point in time. When it is reached, the commit process will slow down and wait for the index update applier thread to make some more room in the queue.
Valid values	dbms.index.fulltext.eventually_consistent_index_update_queue_max_length, an integer which is in the range 1 to 50000000
Default value	10000

### Table 164. dbms.index\_sampling.background\_enabled

Description	Enable or disable background index sampling.
Valid values	dbms.index_sampling.background_enabled, a boolean
Default value	true

### Table 165. dbms.index\_sampling.sample\_size\_limit

Description	Index sampling chunk size limit.
Valid values	dbms.index_sampling.sample_size_limit, an integer which is in the range 1048576 to 2147483647
Default value	8388608

## Table 166. dbms.index\_sampling.update\_percentage

Description	Percentage of index updates of total index size required before sampling of a given index is triggered.
Valid values	dbms.index_sampling.update_percentage, an integer which is minimum 0
Default value	5

### Table 167. dbms.index\_searcher\_cache\_size

Description	The maximum number of open Lucene index searchers.
Valid values	dbms.index_searcher_cache_size, an integer which is minimum 1
Default value	2147483647
Deprecated	The <a href="mailto:dbms.index_searcher_cache_size">dbms.index_searcher_cache_size</a> configuration setting has been deprecated.

### Table 168. dbms.jvm.additional

Description	Additional JVM arguments. Argument order can be significant. To use a Java commercial feature, the argument to unlock commercial features must precede the argument to enable the specific feature in the config value string. For example, to use Flight Recorder, `-XX:+UnlockCommercialFeatures` must come before `-XX:+FlightRecorder`.
Valid values	dbms.jvm.additional, a string
Default value	

### Table 169. dbms.lock.acquisition.timeout

Description	The maximum time interval within which lock should be acquired.
Valid values	dbms.lock.acquisition.timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	0ns

### Table 170. dbms.logs.debug.level

Description	Debug log level threshold.
Valid values	dbms.logs.debug.level, one of [DEBUG, INFO, WARN, ERROR, NONE]
Dynamic	true
Default value	INFO

### Table 171. dbms.logs.debug.path

Description	Path to the debug log file.
Valid values	dbms.logs.debug.path, a path. If relative it is resolved from dbms.directories.logs
Default value	debug.log

## Table 172. dbms.logs.debug.rotation.delay

Description	Minimum time interval after last rotation of the debug log before it may be rotated again.
Valid values	dbms.logs.debug.rotation.delay, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	5m

### Table 173. dbms.logs.debug.rotation.keep\_number

Description	Maximum number of history files for the debug log.

Valid values	dbms.logs.debug.rotation.keep_number, an integer which is minimum 1
Default value	7

### Table 174. dbms.logs.debug.rotation.size

Description	Threshold for rotation of the debug log.
Valid values	dbms.logs.debug.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is in the range 0 to 9223372036854775807
Default value	20971520

### Table 175. dbms.logs.gc.enabled

Description	Enable GC Logging.
Valid values	dbms.logs.gc.enabled, a boolean
Default value	false

### Table 176. dbms.logs.gc.options

Description	GC Logging Options.
Valid values	dbms.logs.gc.options, a string

# Table 177. dbms.logs.gc.rotation.keep\_number

Description	Number of GC logs to keep.
Valid values	dbms.logs.gc.rotation.keep_number, an integer
Default value	0

### Table 178. dbms.logs.gc.rotation.size

Description	Size of each GC log that is kept.
Valid values	dbms.logs.gc.rotation.size, a long

### Table 179. dbms.logs.http.enabled

Description	Enable HTTP request logging.
Valid values	dbms.logs.http.enabled, a boolean
Default value	false

### Table 180. dbms.logs.http.path

Description	Path to HTTP request log.
Valid values	dbms.logs.http.path, a path. If relative it is resolved from dbms.directories.logs
Default value	http.log

## Table 181. dbms.logs.http.rotation.keep\_number

Description	Number of HTTP logs to keep.
Valid values	dbms.logs.http.rotation.keep_number, an integer
Default value	5

# Table 182. dbms.logs.http.rotation.size

Description	Size of each HTTP log that is kept.
Valid values	dbms.logs.http.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is in the range 0 to $9223372036854775807$
Default value	20971520

## Table 183. dbms.logs.query.allocation\_logging\_enabled

Description	Log allocated bytes for the executed queries being logged. The logged number is cumulative over the duration of the query, i.e. for memory intense or long-running queries the value may be larger than the current memory allocation. Requires `< <config_dbms.track_query_allocation,dbms.track_query_allocation>&gt;=true`</config_dbms.track_query_allocation,dbms.track_query_allocation>
Valid values	dbms.logs.query.allocation_logging_enabled, a boolean
Dynamic	true
Default value	false

## Table 184. dbms.logs.query.enabled

Description	Log executed queries. Valid values are 'OFF', 'INFO' & 'VERBOSE'. OFF: no logging. INFO: log queries at the end of execution, that take longer than the configured threshold, <pre>&lt;<config_dbms.logs.query.threshold,dbms.logs.query.threshold>&gt;. VERBOSE: log queries at the start and end of execution, regardless of &lt;&gt;config_dbms.logs.query.threshold,dbms.logs.query.threshold&gt;&gt;. Log entries are by default written to the filequery.log located in the Logs directory. For location of the Logs directory, see &lt;<file-locations>&gt;. This feature is available in the Neo4j Enterprise Edition.</file-locations></config_dbms.logs.query.threshold,dbms.logs.query.threshold></pre>
Valid values	dbms.logs.query.enabled, one of [OFF, INFO, VERBOSE]

Dynamic	true
Default value	VERBOSE

## Table 185. dbms.logs.query.page\_logging\_enabled

Description	Log page hits and page faults for the executed queries being logged.
Valid values	dbms.logs.query.page_logging_enabled, a boolean
Dynamic	true
Default value	false

## Table 186. dbms.logs.query.parameter\_logging\_enabled

Description	Log parameters for the executed queries being logged.
Valid values	dbms.logs.query.parameter_logging_enabled, a boolean
Dynamic	true
Default value	true

## Table 187. dbms.logs.query.path

Description	Path to the query log file.
Valid values	dbms.logs.query.path, a path. If relative it is resolved from dbms.directories.logs
Default value	query.log

### Table 188. dbms.logs.query.rotation.keep\_number

Description	Maximum number of history files for the query log.
Valid values	dbms.logs.query.rotation.keep_number, an integer which is minimum 1
Dynamic	true
Default value	7

## Table 189. dbms.logs.query.rotation.size

Description	The file size in bytes at which the query log will auto-rotate. If set to zero then no rotation will occur. Accepts a binary suffix `k`, `m` or `g`.
Valid values	dbms.logs.query.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is in the range 0 to $9223372036854775807$

Dynamic	true
Default value	20971520

## Table 190. dbms.logs.query.runtime\_logging\_enabled

Description	Logs which runtime that was used to run the query.
Valid values	dbms.logs.query.runtime_logging_enabled, a boolean
Dynamic	true
Default value	false

## Table 191. dbms.logs.query.threshold

Description	If the execution of query takes more time than this threshold, the query is logged once completed - provided query logging is set to INFO. Defaults to 0 seconds, that is all queries are logged.
Valid values	dbms.logs.query.threshold, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Dynamic	true
Default value	0ns

### Table 192. dbms.logs.query.time\_logging\_enabled

Description	Log detailed time information for the executed queries being logged. Requires `< <config_dbms.track_query_cpu_time,dbms.track_query_cpu_time>&gt;=true`</config_dbms.track_query_cpu_time,dbms.track_query_cpu_time>
Valid values	dbms.logs.query.time_logging_enabled, a boolean
Dynamic	true
Default value	false

## Table 193. dbms.logs.security.level

Description	Security log level threshold.
Valid values	dbms.logs.security.level, one of [DEBUG, INFO, WARN, ERROR, NONE]
Default value	INFO

### Table 194. dbms.logs.security.path

Description	Path to the security log file.

Valid values	dbms.logs.security.path, a path. If relative it is resolved from dbms.directories.logs
Default value	security.log

### Table 195. dbms.logs.security.rotation.delay

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Description	Minimum time interval after last rotation of the security log before it may be rotated again.
Valid values	dbms.logs.security.rotation.delay, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	5m

### Table 196. dbms.logs.security.rotation.keep\_number

Description	Maximum number of history files for the security log.
Valid values	dbms.logs.security.rotation.keep_number, an integer which is minimum 1
Default value	7

### Table 197. dbms.logs.security.rotation.size

Description	Threshold for rotation of the security log.
Valid values	dbms.logs.security.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is in the range 0 to $9223372036854775807$
Default value	20971520

### Table 198. dbms.logs.user.path

Description	Path to the user log file.
Valid values	dbms.logs.user.path, a path. If relative it is resolved from dbms.directories.logs
Default value	neo4j.log

#### Table 199. dbms.logs.user.rotation.delay

Description	Minimum time interval after last rotation of the user log before it may be rotated again.
Valid values	dbms.logs.user.rotation.delay, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	5m

### Table 200. dbms.logs.user.rotation.keep\_number

Description	Maximum number of history files for the user log.
Valid values	dbms.logs.user.rotation.keep_number, an integer which is minimum 1
Default value	7

### Table 201. dbms.logs.user.rotation.size

Description	Threshold for rotation of the user log. If set to 0 log rotation is disabled.
Valid values	dbms.logs.user.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is in the range 0 to 9223372036854775807
Default value	0

### Table 202. dbms.logs.user.stdout\_enabled

Description	Send user logs to the process stdout. If this is disabled then logs will instead be sent to the fileneo4j.log located in the logs directory. For location of the Logs directory, see < <file-locations>&gt;.</file-locations>
Valid values	dbms.logs.user.stdout_enabled, a boolean
Default value	true

#### Table 203. dbms.max\_databases

Description	The maximum number of databases.
Valid values	dbms.max_databases, a long which is minimum 2
Default value	100

### Table 204. dbms.memory.heap.initial\_size

Description	Initial heap size. By default it is calculated based on available system resources. (valid units are $k^, K^, M^, M^, G^, G^.$ ).
Valid values	dbms.memory.heap.initial_size, a string
Default value	

### Table 205. dbms.memory.heap.max\_size

Description	Maximum heap size. By default it is calculated based on available system resources. (valid units are `k`, `K`, `m`, `M`, `g`, `G`).
Valid values	dbms.memory.heap.max_size, a string
Default value	

### Table 206. dbms.memory.pagecache.size

Description	The amount of memory to use for mapping the store files, in bytes (or kilobytes with the 'k' suffix, megabytes with 'm' and gigabytes with 'g'). If Neo4j is running on a dedicated server, then it is generally recommended to leave about 2-4 gigabytes for the operating system, give the JVM enough heap to hold all your transaction state and query context, and then leave the rest for the page cache. If no page cache memory is configured, then a heuristic setting is computed based on available system resources.
Valid values	dbms.memory.pagecache.size, a string

# Table 207. dbms.memory.pagecache.swapper

Description	Specify which page swapper to use for doing paged IO. This is only used when integrating with proprietary storage technology.
Valid values	dbms.memory.pagecache.swapper, a string

### Table 208. dbms.memory.pagecache.warmup.enable

Description	Page cache can be configured to perform usage sampling of loaded pages that can be used to construct active load profile. According to that profile pages can be reloaded on the restart, replication, etc. This setting allows disabling that behavior. This feature available in Neo4j Enterprise Edition.
Valid values	dbms.memory.pagecache.warmup.enable, a boolean
Default value	true

### Table 209. dbms.memory.pagecache.warmup.preload

Description	+Page cache warmup can be configured to prefetch files, preferably when cache size is bigger than store size. Files to be prefetched can be filtered by 'dbms.memory.pagecache.warmup.preload.whitelist'. Enabling this disables warmup by profile
Valid values	dbms.memory.pagecache.warmup.preload, a boolean
Default value	false

### Table 210. dbms.memory.pagecache.warmup.preload.whitelist

Description	Page cache warmup prefetch file whitelist regex. By default matches all files.
Valid values	dbms.memory.pagecache.warmup.preload.whitelist, a string
Default value	.*

### Table 211. dbms.memory.pagecache.warmup.profile.interval

Description	The profiling frequency for the page cache. Accurate profiles allow the page cache to do active warmup after a restart, reducing the mean time to performance. This feature available in Neo4j Enterprise Edition.
Valid values	dbms.memory.pagecache.warmup.profile.interval, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1m

### Table 212. dbms.mode

Description	Configure the operating mode of the database 'SINGLE' for stand-alone operation, 'CORE' for operating as a core member of a Causal Cluster, or 'READ_REPLICA' for operating as a read replica member of a Causal Cluster.
Valid values	dbms.mode, one of [SINGLE, CORE, READ_REPLICA]
Default value	SINGLE

## Table 213. dbms.netty.ssl.provider

Description	Netty SSL provider.
Valid values	dbms.netty.ssl.provider, one of [JDK, OPENSSL, OPENSSL_REFCNT]
Default value	JDK

### Table 214. dbms.query\_cache\_size

Description	The number of Cypher query execution plans that are cached.
Valid values	dbms.query_cache_size, an integer which is minimum 0
Default value	1000

### Table 215. dbms.read\_only

Description	Only allow read operations from this Neo4j instance. This mode still requires write access to the directory for lock purposes.
Valid values	dbms.read_only, a boolean
Default value	false

## Table 216. dbms.reconciler.max\_backoff

Description	Defines the maximum amount of time to wait before retrying after the dbms fails to reconcile a database to its desired state.
Valid values	dbms.reconciler.max_backoff, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is minimum PT1M

Default value	1h

### Table 217. dbms.reconciler.may\_retry

Description	Defines whether the dbms may retry reconciling a database to its desired state.
Valid values	dbms.reconciler.may_retry, a boolean
Default value	false

### Table 218. dbms.reconciler.min\_backoff

Description	Defines the minimum amount of time to wait before retrying after the dbms fails to reconcile a database to its desired state.
Valid values	dbms.reconciler.min_backoff, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is minimum PT1S
Default value	2s

### Table 219. dbms.record\_format

Description	Database record format. Valid values: `standard`, `high_limit`. The `high_limit` format is available for Enterprise Edition only. It is required if you have a graph that is larger than 34 billion nodes, 34 billion relationships, or 68 billion properties. A change of the record format is irreversible. Certain operations may suffer from a performance penalty of up to 10%, which is why this format is not switched on by default.
Valid values	dbms.record_format, a string
Default value	

### Table 220. dbms.recovery.fail\_on\_missing\_files

Description	If `true`, Neo4j will abort recovery if logical log files are missing. Setting this to `false` will allow Neo4j to create new empty missing files for already existing database, but, the integrity of the database might be compromised.
Valid values	dbms.recovery.fail_on_missing_files, a boolean
Default value	true

### Table 221. dbms.relationship\_grouping\_threshold

Description	Relationship count threshold for considering a node to be dense.
Valid values	dbms.relationship_grouping_threshold, an integer which is minimum 1
Default value	50

### Table 222. dbms.rest.transaction.idle\_timeout

Description	Timeout for idle transactions in the REST endpoint.
Valid values	dbms.rest.transaction.idle_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1m

### Table 223. dbms.routing\_ttl

Description	How long callers should cache the response of the routing procedure `dbms.routing.getRoutingTable()`
Valid values	dbms.routing_ttl, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is minimum PT1S
Default value	5m

### Table 224. dbms.security.allow\_csv\_import\_from\_file\_urls

Description	Determines if Cypher will allow using file URLs when loading data using `LOAD CSV`. Setting this value to `false` will cause Neo4j to fail `LOAD CSV` clauses that load data from the file system.
Valid values	dbms.security.allow_csv_import_from_file_urls, a boolean
Default value	true

#### Table 225. dbms.security.auth\_cache\_max\_capacity

Description	The maximum capacity for authentication and authorization caches (respectively).
Valid values	dbms.security.auth_cache_max_capacity, an integer
Default value	10000

### Table 226. dbms.security.auth\_cache\_ttl

Description	The time to live (TTL) for cached authentication and authorization info when using external auth providers (LDAP or plugin). Setting the TTL to 0 will disable auth caching. Disabling caching while using the LDAP auth provider requires the use of an LDAP system account for resolving authorization information.
Valid values	dbms.security.auth_cache_ttl, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10m

### Table 227. dbms.security.auth\_cache\_use\_ttl

Description	Enable time-based eviction of the authentication and authorization info cache for external auth providers (LDAP or plugin). Disabling this setting will make the cache live forever and only be evicted when `< <config_dbms.security.auth_cache_max_capacity,dbms.security.auth_cache_max_capacity>&gt;` is exceeded.</config_dbms.security.auth_cache_max_capacity,dbms.security.auth_cache_max_capacity>
Valid values	dbms.security.auth_cache_use_ttl, a boolean
Default value	true

### Table 228. dbms.security.auth\_enabled

Description	Enable auth requirement to access Neo4j.
Valid values	dbms.security.auth_enabled, a boolean
Default value	true

## Table 229. dbms.security.auth\_lock\_time

Description	The amount of time user account should be locked after a configured number of unsuccessful authentication attempts. The locked out user will not be able to log in until the lock period expires, even if correct credentials are provided. Setting this configuration option to a low value is not recommended because it might make it easier for an attacker to brute force the password.
Valid values	dbms.security.auth_lock_time, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is minimum PT0S
Default value	5s

## Table 230. dbms.security.auth\_max\_failed\_attempts

Description	The maximum number of unsuccessful authentication attempts before imposing a user lock for the configured amount of time. The locked out user will not be able to log in until the lock period expires, even if correct credentials are provided. Setting this configuration option to values less than 3 is not recommended because it might make it easier for an attacker to brute force the password.
Valid values	dbms.security.auth_max_failed_attempts, an integer which is minimum ₀
Default value	3

### *Table 231. dbms.security.authentication\_providers*

Description	A list of security authentication providers containing the users and roles. This can be any of the built-in `native` or `ldap` providers, or it can be an externally provided plugin, with a custom name prefixed by `plugin-`, i.e. `plugin- <auth_provider_name>`. They will be queried in the given order when login is attempted.</auth_provider_name>
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Valid values	dbms.security.authentication_providers, a ',' separated list with elements of type 'a string'.
Default value	native

### Table 232. dbms.security.authorization\_providers

Description	A list of security authorization providers containing the users and roles. This can be any of the built-in `native` or `ldap` providers, or it can be an externally provided plugin, with a custom name prefixed by `plugin-`, i.e. `plugin- <auth_provider_name>`. They will be queried in the given order when login is attempted.</auth_provider_name>
Valid values	dbms.security.authorization_providers, a ',' separated list with elements of type 'a string'.
Default value	native

### Table 233. dbms.security.causal\_clustering\_status\_auth\_enabled

Description	Require authorization for access to the Causal Clustering status endpoints.
Valid values	dbms.security.causal_clustering_status_auth_enabled, a boolean
Default value	true

### Table 234. dbms.security.http\_access\_control\_allow\_origin

Description	Value of the Access-Control-Allow-Origin header sent over any HTTP or HTTPS connector. This defaults to '*', which allows broadest compatibility. Note that any URI provided here limits HTTP/HTTPS access to that URI only.
Valid values	dbms.security.http_access_control_allow_origin, a string
Default value	*

#### Table 235. dbms.security.http\_auth\_whitelist

Description	Defines a whitelist of http paths where Neo4j authentication is not required.
Valid values	dbms.security.http_auth_whitelist, a ',' separated list with elements of type 'a string'.
Default value	/,/browser.*

### Table 236. dbms.security.http\_strict\_transport\_security

Description	Value of the HTTP Strict-Transport-Security (HSTS) response header. This header tells browsers that a webpage should only be accessed using HTTPS instead of HTTP. It is attached to every HTTPS response. Setting is not set by default so 'Strict-Transport-Security' header is not sent. Value is expected to contain directives like 'max-age', 'includeSubDomains' and 'preload'.

Valid values	dbms.security.http_strict_transport_security, a string

# Table 237. dbms.security.ldap.authentication.cache\_enabled

Description	Determines if the result of authentication via the LDAP server should be cached or not. Caching is used to limit the number of LDAP requests that have to be made over the network for users that have already been authenticated successfully. A user can be authenticated against an existing cache entry (instead of via an LDAP server) as long as it is alive (see `< <config_dbms.security.auth_cache_ttl,dbms.security.auth_cache_ttl>&gt;`). An important consequence of setting this to `true` is that Neo4j then needs to cache a hashed version of the credentials in order to perform credentials matching. This hashing is done using a cryptographic hash function together with a random salt. Preferably a conscious decision should be made if this method is considered acceptable by the security standards of the organization in which this Neo4j instance is deployed.</config_dbms.security.auth_cache_ttl,dbms.security.auth_cache_ttl>
Valid values	dbms.security.ldap.authentication.cache_enabled, a boolean
Default value	true

### Table 238. dbms.security.ldap.authentication.mechanism

Description	LDAP authentication mechanism. This is one of `simple` or a SASL mechanism supported by JNDI, for example `DIGEST-MD5`. `simple` is basic username and password authentication and SASL is used for more advanced mechanisms. See RFC 2251 LDAPv3 documentation for more details.
Valid values	dbms.security.ldap.authentication.mechanism, a string
Default value	simple

### Table 239. dbms.security.ldap.authentication.use\_samaccountname

Description	Perform authentication with sAMAccountName instead of DN. Using this setting requires `< <config_dbms.security.ldap.authorization.system_username,dbms.security.ld ap.authorization.system_username="">&gt;` and &lt;<config_dbms.security.ldap.authorization.system_password,dbms.security.ldap .authorization.system_password="">&gt; to be used since there is no way to log in through ldap directly with the sAMAccountName, instead the login name will be resolved to a DN that will be used to log in with.</config_dbms.security.ldap.authorization.system_password,dbms.security.ldap></config_dbms.security.ldap.authorization.system_username,dbms.security.ld>
Valid values	dbms.security.ldap.authentication.use_samaccountname, a boolean
Default value	false

Table 240. dbms.security.ldap.authentication.user\_dn\_template

Description	LDAP user DN template. An LDAP object is referenced by its distinguished name (DN), and a user DN is an LDAP fully-qualified unique user identifier. This setting is used to generate an LDAP DN that conforms with the LDAP directory's schema from the user principal that is submitted with the authentication token when logging in. The special token {0} is a placeholder where the user principal will be substituted into the DN string.
Valid values	dbms.security.ldap.authentication.user_dn_template, a string
Default value	uid={0},ou=users,dc=example,dc=com

### Table 241. dbms.security.ldap.authorization.group\_membership\_attributes

Description	A list of attribute names on a user object that contains groups to be used for mapping to roles when LDAP authorization is enabled.
Valid values	dbms.security.ldap.authorization.group_membership_attributes, a ',' separated list with elements of type 'a string'.
Default value	memberOf

### Table 242. dbms.security.ldap.authorization.group\_to\_role\_mapping

Description	+An authorization mapping from LDAP group names to Neo4j role names. The map should be formatted as a semicolon separated list of key-value pairs, where the key is the LDAP group name and the value is a comma separated list of corresponding role names. For example: group1=role1;group2=role2;group3=role3,role4,role5 You could also use whitespaces and quotes around group names to make this mapping more readable, for example:
	<pre>dbms.security.ldap.authorization.group_to_role_mapping=\     "cn=Neo4j Read Only,cn=users,dc=example,dc=com" = reader;     "cn=Neo4j Read-Write,cn=users,dc=example,dc=com" = publisher; \     "cn=Neo4j Schema Manager,cn=users,dc=example,dc=com" = architect; \     "cn=Neo4j Administrator,cn=users,dc=example,dc=com" = admin</pre>
	Deprecated: This will be replaced by dynamic configuration in the system graph in 4.0, including a migration step for the existing setting value.
Valid values	dbms.security.ldap.authorization.group_to_role_mapping, a string
Deprecated	The dbms.security.ldap.authorization.group_to_role_mapping configuration setting has been deprecated.

### Table 243. dbms.security.ldap.authorization.system\_password

Description	An LDAP system account password to use for authorization searches when `< <config_dbms.security.ldap.authorization.use_system_account,dbms.security.ldap.authorization.use_system_account>&gt;` is `true`.</config_dbms.security.ldap.authorization.use_system_account,dbms.security.ldap.authorization.use_system_account>
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Valid values	dbms.security.ldap.authorization.system_password, a secure string

### Table 244. dbms.security.ldap.authorization.system\_username

Description	An LDAP system account username to use for authorization searches when `< <config_dbms.security.ldap.authorization.use_system_account,dbms.security.ldap.authorization.use_system_account>&gt;` is `true`. Note that the `&lt;<config_dbms.security.ldap.authentication.user_dn_template,dbms.security.ld ap.authentication.user_dn_template="">&gt;` will not be applied to this username, so you may have to specify a full DN.</config_dbms.security.ldap.authentication.user_dn_template,dbms.security.ld></config_dbms.security.ldap.authorization.use_system_account,dbms.security.ldap.authorization.use_system_account>
Valid values	dbms.security.ldap.authorization.system_username, a string

### Table 245. dbms.security.ldap.authorization.use\_system\_account

Description	Perform LDAP search for authorization info using a system account instead of the user's own account. If this is set to `false` (default), the search for group membership will be performed directly after authentication using the LDAP context bound with the user's own account. The mapped roles will be cached for the duration of ` <config_dbms.security.auth_cache_ttl,dbms.security.auth_cache_ttl>&gt;`, and then expire, requiring re-authentication. To avoid frequently having to reauthenticate sessions you may want to set a relatively long auth cache expiration time together with this option. NOTE: This option will only work if the users are permitted to search for their own group membership attributes in the directory. If this is set to `true`, the search will be performed using a special system account user with read access to all the users in the directory. You need to specify the username and password using the settings `<config_dbms.security.ldap.authorization.system_username,dbms.security.ldap.authorization.system_username,dbms.security.ldap.authorization.system_password,dbms.security.ldap.authorization.system_password,dbms.security.ldap.authorization.system_password>&gt;` with this option. Note that this account only needs read access to the relevant parts of the LDAP directory and does not need to have access rights to Neo4j, or any other systems.</config_dbms.security.ldap.authorization.system_username,dbms.security.ldap.authorization.system_username,dbms.security.ldap.authorization.system_password,dbms.security.ldap.authorization.system_password,dbms.security.ldap.authorization.system_password></config_dbms.security.auth_cache_ttl,dbms.security.auth_cache_ttl>
Valid values	dbms.security.ldap.authorization.use_system_account, a boolean
Default value	false

### Table 246. dbms.security.ldap.authorization.user\_search\_base

Description	The name of the base object or named context to search for user objects when LDAP authorization is enabled. A common case is that this matches the last part of `< <config_dbms.security.ldap.authentication.user_dn_template,dbms.security.ld ap.authentication.user_dn_template="">&gt;`.</config_dbms.security.ldap.authentication.user_dn_template,dbms.security.ld>
Valid values	dbms.security.ldap.authorization.user_search_base, a string
Default value	ou=users,dc=example,dc=com

Table 247. dbms.security.ldap.authorization.user\_search\_filter

Description	The LDAP search filter to search for a user principal when LDAP authorization is enabled. The filter should contain the placeholder token {0} which will be substituted for the user principal.
Valid values	dbms.security.ldap.authorization.user_search_filter, a string
Default value	(&(objectClass=*)(uid={0}))

## Table 248. dbms.security.ldap.connection\_timeout

Description	The timeout for establishing an LDAP connection. If a connection with the LDAP server cannot be established within the given time the attempt is aborted. A value of 0 means to use the network protocol's (i.e., TCP's) timeout value.
Valid values	dbms.security.ldap.connection_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	30s

### Table 249. dbms.security.ldap.host

Description	URL of LDAP server to use for authentication and authorization. The format of the setting is ` <pre>cprotocol&gt;://<hostname>:<port>`, where hostname is the only required field. The supported values for protocol are `ldap` (default) and `ldaps`. The default port for `ldap` is 389 and for `ldaps` 636. For example: `ldaps://ldap.example.com:10389`. You may want to consider using STARTTLS (`<config_dbms.security.ldap.use_starttls,dbms.security.ldap.use_starttls>&gt;`) instead of LDAPS for secure connections, in which case the correct protocol is `ldap`.</config_dbms.security.ldap.use_starttls,dbms.security.ldap.use_starttls></port></hostname></pre>
Valid values	dbms.security.ldap.host, a string
Default value	localhost

### Table 250. dbms.security.ldap.read\_timeout

Description	The timeout for an LDAP read request (i.e. search). If the LDAP server does not respond within the given time the request will be aborted. A value of 0 means wait for a response indefinitely.
Valid values	dbms.security.ldap.read_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	30s

## Table 251. dbms.security.ldap.referral

Description	The LDAP referral behavior when creating a connection. This is one of `follow`, `ignore` or `throw`. * `follow` automatically follows any referrals * `ignore` ignores any referrals * `throw` throws an exception, which will lead to authentication failure.
Valid values	dbms.security.ldap.referral, a string

Default value	follow

### Table 252. dbms.security.ldap.use\_starttls

Description	Use secure communication with the LDAP server using opportunistic TLS. First an initial insecure connection will be made with the LDAP server, and a STARTTLS command will be issued to negotiate an upgrade of the connection to TLS before initiating authentication.
Valid values	dbms.security.ldap.use_starttls, a boolean
Default value	false

### Table 253. dbms.security.log\_successful\_authentication

Description	Set to log successful authentication events to the security log. If this is set to `false` only failed authentication events will be logged, which could be useful if you find that the successful events spam the logs too much, and you do not require full auditing capability.
Valid values	dbms.security.log_successful_authentication, a boolean
Default value	true

### Table 254. dbms.security.procedures.default\_allowed

Description	The default role that can execute all procedures and user-defined functions that are not covered by the `< <config_dbms.security.procedures.roles,dbms.security.procedures.roles>&gt;`setting. If the ``dbms.security.procedures.default_allowed`` setting is the empty string (default), procedures will be executed according to the same security rules as normal Cypher statements. Deprecated: This will be replaced by dynamic configuration in the system graph in 4.0, including a migration step for the existing setting value.</config_dbms.security.procedures.roles,dbms.security.procedures.roles>
Valid values	dbms.security.procedures.default_allowed, a string
Default value	
Deprecated	The <pre>dbms.security.procedures.default_allowed</pre> configuration setting has been deprecated.

Table 255. dbms.security.procedures.roles

Description	This provides a finer level of control over which roles can execute procedures than the `< <config_dbms.security.procedures.default_allowed,dbms.security.procedures.default_allowed>&gt;` setting. For example: `+dbms.security.procedures.roles=apoc.convert.*:reader;apoc.load.json*:writer;apoc.trigger.add:TriggerHappy` will allow the role reader to execute all procedures in the apoc.convert namespace, the role writer to execute all procedures in the apoc.load namespace that starts with json and the role TriggerHappy to execute the specific procedure apoc.trigger.add. Procedures not matching any of these patterns will be subject to the dbms.security.procedures.default_allowed setting. Deprecated: This will be replaced by dynamic configuration in the system graph in 4.0, including a migration step for the existing setting value.+</config_dbms.security.procedures.default_allowed,dbms.security.procedures.default_allowed>
Valid values	dbms.security.procedures.roles, a string
Default value	
Deprecated	The dbms.security.procedures.roles configuration setting has been deprecated.

## Table 256. dbms.security.procedures.unrestricted

Description	A list of procedures and user defined functions (comma separated) that are allowed full access to the database. The list may contain both fully-qualified procedure names, and partial names with the wildcard '*'. Note that this enables these procedures to bypass security. Use with caution.
Valid values	dbms.security.procedures.unrestricted, a ',' separated list with elements of type 'a string'.
Default value	

## Table 257. dbms.security.procedures.whitelist

Description	A list of procedures (comma separated) that are to be loaded. The list may contain both fully-qualified procedure names, and partial names with the wildcard '*'. If this setting is left empty no procedures will be loaded.
Valid values	dbms.security.procedures.whitelist, a ',' separated list with elements of type 'a string'.
Default value	*

## Table 258. dbms.security.property\_level.enabled

Description	This has been replaced by privilege management on roles. Setting it to true will prevent the server from starting.
Valid values	dbms.security.property_level.enabled, a boolean
Default value	false

Deprecated	The dbms.security.property_level.enabled configuration setting has been deprecated.

## $\textit{Table 259. dbms.shutdown\_transaction\_end\_timeout}$

Description	The maximum amount of time to wait for running transactions to complete before allowing initiated database shutdown to continue.
Valid values	dbms.shutdown_transaction_end_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	10s

## Table 260. dbms.threads.worker\_count

Description	Number of Neo4j worker threads. This setting is only valid for REST, and does not influence bolt-server. It sets the amount of worker threads for the Jetty server used by neo4j-server. This option can be tuned when you plan to execute multiple, concurrent REST requests, with the aim of getting more throughput from the database. Your OS might enforce a lower limit than the maximum value specified here.
Valid values	dbms.threads.worker_count, an integer which is in the range 1 to 44738
Default value	Number of available processors, or 500 for machines which have more than 500 processors.

### Table 261. dbms.track\_query\_allocation

Description	Enables or disables tracking of how many bytes are allocated by the execution of a query. If enabled, calling `dbms.listQueries` will display the allocated bytes. If enabled, the maximum allocated bytes of a query can be limited using `< <config_cypher.query_max_allocations,cypher.query_max_allocations>&gt;`. This can also be logged in the query log by using `&lt;<config_dbms.logs.query.allocation_logging_enabled,dbms.logs.query.allocation_logging_enabled>&gt;`.</config_dbms.logs.query.allocation_logging_enabled,dbms.logs.query.allocation_logging_enabled></config_cypher.query_max_allocations,cypher.query_max_allocations>
Valid values	dbms.track_query_allocation, a boolean
Dynamic	true
Default value	false

### Table 262. dbms.track\_query\_cpu\_time

Description	Enables or disables tracking of how much time a query spends actively executing on the CPU. Calling `dbms.listQueries` will display the time. This can also be logged in the query log by using `log_queries_detailed_time_logging_enabled`.
Valid values	dbms.track_query_cpu_time, a boolean

Dynamic	true
Default value	false

## $\textit{Table 263. dbms.} transaction. bookmark\_ready\_timeout$

Description	The maximum amount of time to wait for the database state represented by the bookmark.
Valid values	dbms.transaction.bookmark_ready_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's') which is minimum PT1S
Default value	30s

#### Table 264. dbms.transaction.concurrent.maximum

Description	The maximum number of concurrently running transactions. If set to 0, limit is disabled.
Valid values	dbms.transaction.concurrent.maximum, an integer
Dynamic	true
Default value	1000

#### Table 265. dbms.transaction.monitor.check.interval

Description	Configures the time interval between transaction monitor checks. Determines how often monitor thread will check transaction for timeout.
Valid values	dbms.transaction.monitor.check.interval, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	2s

### Table 266. dbms.transaction.sampling.percentage

Description	Transaction sampling percentage.
Valid values	dbms.transaction.sampling.percentage, an integer which is in the range 1 to 100
Dynamic	true
Default value	5

#### Table 267. dbms.transaction.timeout

Description	The maximum time interval of a transaction within which it should be completed.

Valid values	dbms.transaction.timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Dynamic	true
Default value	0ns

## Table 268. dbms.transaction.tracing.level

Description	Transaction creation tracing level.
Valid values	dbms.transaction.tracing.level, one of [DISABLED, SAMPLE, ALL]
Dynamic	true
Default value	DISABLED

### Table 269. dbms.tx\_log.preallocate

Description	Specify if Neo4j should try to preallocate logical log file in advance.
Valid values	dbms.tx_log.preallocate, a boolean
Dynamic	true
Default value	true

## Table 270. dbms.tx\_log.rotation.retention\_policy

Description	Make Neo4j keep the logical transaction logs for being able to backup the database. Can be used for specifying the threshold to prune logical logs after. For example "10 days" will prune logical logs that only contains transactions older than 10 days from the current time, or "100k txs" will keep the 100k latest transactions and prune any older transactions.
Valid values	dbms.tx_log.rotation.retention_policy, a string which matches the pattern ^(true keep_all false keep_none (\d+[KkMmGg]?( (files size txs entries hours days))))\$ (must be true, false or of format <number><optional unit=""> <type>. Valid units are k, M and G. Valid types are files, size, txs, entries, hours and days. For example, 100M size will limiting logical log space on disk to 100Mb, or 200k txs will limiting the number of transactions to keep to 200 000)</type></optional></number>
Dynamic	true
Default value	7 days

## Table 271. dbms.tx\_log.rotation.size

Description	+Specifies at which file size the logical log will auto-rotate. Minimum accepted value is 128 KiB.
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Valid values	dbms.tx_log.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is minimum 131072
Dynamic	true
Default value	262144000

## Table 272. dbms.tx\_state.max\_off\_heap\_memory

Description	The maximum amount of off-heap memory that can be used to store transaction state data; it's a total amount of memory shared across all active transactions. Zero means 'unlimited'. Used when < <config_dbms.tx_state.memory_allocation,dbms.tx_state.memory_allocation>&gt; is set to 'OFF_HEAP'.</config_dbms.tx_state.memory_allocation,dbms.tx_state.memory_allocation>
Valid values	dbms.tx_state.max_off_heap_memory, a byte size (valid multipliers are k, m, g, K, M, G) which is minimum 0
Default value	2147483648

### Table 273. dbms.tx\_state.memory\_allocation

Description	Defines whether memory for transaction state should be allocated on- or off-heap.
Valid values	dbms.tx_state.memory_allocation, one of [ON_HEAP, OFF_HEAP]
Default value	OFF_HEAP

### Table 274. dbms.tx\_state.off\_heap.block\_cache\_size

Description	Defines the size of the off-heap memory blocks cache. The cache will contain this number of blocks for each block size that is power of two. Thus, maximum amount of memory used by blocks cache can be calculated as 2 * < <config_dbms.tx_state.off_heap.max_cacheable_block_size,dbms.tx_state.off_heap.max_cacheable_block_size>&gt; * `dbms.tx_state.off_heap.block_cache_size`</config_dbms.tx_state.off_heap.max_cacheable_block_size,dbms.tx_state.off_heap.max_cacheable_block_size>
Valid values	dbms.tx_state.off_heap.block_cache_size, an integer which is minimum 16
Default value	128

### Table 275. dbms.tx\_state.off\_heap.max\_cacheable\_block\_size

Description	Defines the maximum size of an off-heap memory block that can be cached to speed up allocations for transaction state data. The value must be a power of 2.
Valid values	dbms.tx_state.off_heap.max_cacheable_block_size, a byte size (valid multipliers are k, m, g, K, M, G) which is minimum 4096 and is power of 2
Default value	524288

### Table 276. dbms.unmanaged\_extension\_classes

Description	Comma-separated list of <classname>=<mount point=""> for unmanaged extensions.</mount></classname>
Valid values	dbms.unmanaged_extension_classes, a ',' separated list with elements of type ' <classname>=<mount point=""> string'.</mount></classname>
Default value	

### Table 277. dbms.windows\_service\_name

Description	Name of the Windows Service.
Valid values	dbms.windows_service_name, a string
Default value	neo4j

### Table 278. fabric.database.name

Description	Name of the Fabric database. Only one Fabric database is currently supported per Neo4j instance.
Valid values	fabric.database.name, A valid database name. Containing only alphabetic characters, numbers, dots and dashes, with a length between 3 and 63 characters. It should be starting with an alphabetic character but not with the name 'system'.

## Table 279. fabric.driver.api

Description	Determines which driver API will be used. ASYNC must be used when the remote instance is 3.5.
Valid values	fabric.driver.api, one of [RX, ASYNC]
Default value	RX

## Table 280. fabric.driver.connection.connect\_timeout

Description	Socket connection timeout. A timeout of zero is treated as an infinite timeout and will be bound by the timeout configured on the operating system level.
Valid values	fabric.driver.connection.connect_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	5s

### Table 281. fabric.driver.connection.max\_lifetime

Description	Pooled connections older than this threshold will be closed and removed from the pool. Setting this option to a low value will cause a high connection churn and might result in a performance hit. It is recommended to set maximum lifetime to a slightly smaller value than the one configured in network equipment (load balancer, proxy, firewall, etc. can also limit maximum connection lifetime). Zero and negative values result in lifetime not being checked.
Valid values	fabric.driver.connection.max_lifetime, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1h

## $Table~282.~fabric.driver.connection.pool.acquisition\_timeout$

Description	Maximum amount of time spent attempting to acquire a connection from the connection pool. This timeout only kicks in when all existing connections are being used and no new connections can be created because maximum connection pool size has been reached. Error is raised when connection can't be acquired within configured time. Negative values are allowed and result in unlimited acquisition timeout. Value of 0 is allowed and results in no timeout and immediate failure when connection is unavailable.
Valid values	fabric.driver.connection.pool.acquisition_timeout, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1m

### Table 283. fabric.driver.connection.pool.idle\_test

Description	Pooled connections that have been idle in the pool for longer than this timeout will be tested before they are used again, to ensure they are still alive. If this option is set too low, an additional network call will be incurred when acquiring a connection, which causes a performance hit. If this is set high, no longer live connections might be used which might lead to errors. Hence, this parameter tunes a balance between the likelihood of experiencing connection problems and performance Normally, this parameter should not need tuning. Value 0 means connections will always be tested for validity.
Valid values	fabric.driver.connection.pool.idle_test, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	No connection liveliness check is done by default.

### Table 284. fabric.driver.connection.pool.max\_size

Description	Maximum total number of connections to be managed by a connection pool. The limit is enforced for a combination of a host and user. Negative values are allowed and result in unlimited pool. Value of 0is not allowed.
Valid values	fabric.driver.connection.pool.max_size, an integer
Default value	Unlimited

### Table 285. fabric.driver.logging.level

Description	Sets level for driver internal logging.
Valid values	fabric.driver.logging.level, one of [DEBUG, INFO, WARN, ERROR, NONE]
Default value	Value of dbms.logs.debug.level

### Table 286. fabric.routing.servers

Description	A comma-separated list of Fabric instances that form a routing group. A driver will route transactions to available routing group members. A Fabric instance is represented by its Bolt connector address.
Valid values	fabric.routing.servers, a ',' separated list with elements of type 'a socket address'.
Default value	localhost:7687

## Table 287. fabric.routing.ttl

Description	The time to live (TTL) of a routing table for fabric routing group.
Valid values	fabric.routing.ttl, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	1m

### Table 288. fabric.stream.buffer.low\_watermark

Description	Number of records in prefetching buffer that will trigger prefetching again. This is strongly related to < <config_fabric.stream.buffer.size,fabric.stream.buffer.size>&gt;</config_fabric.stream.buffer.size,fabric.stream.buffer.size>
Valid values	fabric.stream.buffer.low_watermark, an integer
Default value	300

### Table 289. fabric.stream.buffer.size

Description	Maximal size of a buffer used for pre-fetching result records of remote queries. To compensate for latency to remote databases, the Fabric execution engine pre-fetches records needed for local executions. This limit is enforced per fabric query. If a fabric query uses multiple remote stream at the same time, this setting represents the maximal number of pre-fetched records counted together for all such remote streams.
Valid values	fabric.stream.buffer.size, an integer
Default value	1000

### Table 290. fabric.stream.concurrency

Description	Maximal concurrency within Fabric queries. Limits the number of iterations of each subquery that are executed concurrently. Higher concurrency may consume more memory and network resources simultaneously, while lower concurrency may force sequential execution, requiring more time.
Valid values	fabric.stream.concurrency, an integer
Default value	The number of remote graphs

### Table 291. metrics.bolt.messages.enabled

Description	Enable reporting metrics about Bolt Protocol message processing.
Valid values	metrics.bolt.messages.enabled, a boolean
Default value	true

### Table 292. metrics.csv.enabled

Description	Set to `true` to enable exporting metrics to CSV files.
Valid values	metrics.csv.enabled, a boolean
Default value	true

### Table 293. metrics.csv.interval

Description	The reporting interval for the CSV files. That is, how often new rows with numbers are appended to the CSV files.
Valid values	metrics.csv.interval, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	3s

### Table 294. metrics.csv.rotation.keep\_number

Description	Maximum number of history files for the csv files.
Valid values	metrics.csv.rotation.keep_number, an integer which is minimum 1
Default value	7

## Table 295. metrics.csv.rotation.size

Description	The file size in bytes at which the csv files will auto-rotate. If set to zero then no rotation will occur. Accepts a binary suffix `k`, `m` or `g`.
Valid values	metrics.csv.rotation.size, a byte size (valid multipliers are k, m, g, K, M, G) which is in the range 0 to $9223372036854775807$
Default value	10485760

### Table 296. metrics.cypher.replanning.enabled

Description	Enable reporting metrics about number of occurred replanning events.
Valid values	metrics.cypher.replanning.enabled, a boolean
Default value	true

#### Table 297. metrics.enabled

Description	Enable metrics. Setting this to `false` will to turn off all metrics.
Valid values	metrics.enabled, a boolean
Default value	true

### Table 298. metrics.graphite.enabled

Description	Set to `true` to enable exporting metrics to Graphite.
Valid values	metrics.graphite.enabled, a boolean
Default value	false

### Table 299. metrics.graphite.interval

Description	The reporting interval for Graphite. That is, how often to send updated metrics to Graphite.
Valid values	metrics.graphite.interval, a duration (Valid units are: 'ns', 'ms', 's', 'm' and 'h'; default unit is 's')
Default value	3s

### Table 300. metrics.graphite.server

Description	The hostname or IP address of the Graphite server.
Valid values	metrics.graphite.server, a hostname and port
Default value	:2003

### Table 301. metrics.jmx.enabled

Description	Set to `true` to enable the JMX metrics endpoint.
Valid values	metrics.jmx.enabled, a boolean
Default value	true

# Table 302. metrics.jvm.buffers.enabled

Description	Enable reporting metrics about the buffer pools.
Valid values	metrics.jvm.buffers.enabled, a boolean
Default value	true

## Table 303. metrics.jvm.file.descriptors.enabled

Description	Enable reporting metrics about the number of open file descriptors.
Valid values	metrics.jvm.file.descriptors.enabled, a boolean
Default value	true

### Table 304. metrics.jvm.gc.enabled

Description	Enable reporting metrics about the duration of garbage collections.
Valid values	metrics.jvm.gc.enabled, a boolean
Default value	true

### Table 305. metrics.jvm.heap.enabled

Description	Enable reporting metrics about the heap memory usage.
Valid values	metrics.jvm.heap.enabled, a boolean
Default value	true

### Table 306. metrics.jvm.memory.enabled

Description	Enable reporting metrics about the memory usage.
Valid values	metrics.jvm.memory.enabled, a boolean
Default value	true

### Table 307. metrics.jvm.threads.enabled

Description	Enable reporting metrics about the current number of threads running.
Valid values	metrics.jvm.threads.enabled, a boolean
Default value	true

## Table 308. metrics.neo4j.causal\_clustering.enabled

Description	Enable reporting metrics about Causal Clustering mode.
Valid values	metrics.neo4j.causal_clustering.enabled, a boolean

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### Table 309. metrics.neo4j.checkpointing.enabled

Description	Enable reporting metrics about Neo4j check pointing; when it occurs and how much time it takes to complete.
Valid values	metrics.neo4j.checkpointing.enabled, a boolean
Default value	true

### Table 310. metrics.neo4j.counts.enabled

Description	Enable reporting metrics about approximately how many entities are in the database; nodes, relationships, properties, etc.
Valid values	metrics.neo4j.counts.enabled, a boolean
Default value	true

### Table 311. metrics.neo4j.data.counts.enabled

Description	Enable reporting metrics about number of entities in the database.
Valid values	metrics.neo4j.data.counts.enabled, a boolean
Default value	true

### Table 312. metrics.neo4j.logs.enabled

Description	Enable reporting metrics about the Neo4j transaction logs.
Valid values	metrics.neo4j.logs.enabled, a boolean
Default value	true

### Table 313. metrics.neo4j.pagecache.enabled

Description	Enable reporting metrics about the Neo4j page cache; page faults, evictions, flushes, exceptions, etc.
Valid values	metrics.neo4j.pagecache.enabled, a boolean
Default value	true

#### Table 314. metrics.neo4j.server.enabled

Description	Enable reporting metrics about Server threading info.			
Valid values	metrics.neo4j.server.enabled, a boolean			
Default value	true			

#### Table 315. metrics.neo4j.size.enabled

Description	Enable reporting metrics about the store size of each database.
Valid values	metrics.neo4j.size.enabled, a boolean
Default value	true

#### Table 316. metrics.neo4j.tx.enabled

Description	Enable reporting metrics about transactions; number of transactions started, committed, etc.
Valid values	metrics.neo4j.tx.enabled, a boolean
Default value	true

#### Table 317. metrics.prefix

Description	A common prefix for the reported metrics field names.			
Valid values	metrics.prefix, a string			
Default value	neo4j			

#### Table 318. metrics.prometheus.enabled

Description	Set to `true` to enable the Prometheus endpoint.			
Valid values	metrics.prometheus.enabled, a boolean			
Default value	false			

#### Table 319. metrics.prometheus.endpoint

Description	The hostname and port to use as Prometheus endpoint.
Valid values	metrics.prometheus.endpoint, a hostname and port
Default value	localhost:2004

# A.2. Built-in procedures

This section contains a reference of Neo4j built-in procedures.

#### This section includes:

- Procedures, editions and modes
- Procedures available in standalone Neo4j Enterprise Edition
- Procedures available in Neo4j Community Edition

# A.2.1. Procedures, editions and modes

The procedures available depends on the type of installation. Enterprise Edition provides a fuller set of procedures than Community Edition. Cluster members have procedures that are not available in standalone mode.

The cluster-specific procedures are not included in this reference, instead see Procedures for monitoring a Causal Cluster. To check which procedures are available in your Neo4j instance, use the dbms.procedures() procedure.

#### Example 109. List available procedures

To list the procedures available on your particular installation, run the following command in Neo4j Browser or in Cypher Shell:

CALL dbms.procedures()

# A.2.2. Enterprise Edition procedures

#### Table 320. Enterprise Edition procedures

Name	Description	Signature	Mode	Roles
db.awaitIndex()	Wait for an index to come online (for example: CALL db.awaitIndex("My Index", 300)).	<pre>db.awaitIndex(indexNa me :: STRING?, timeOutSeconds = 300 :: INTEGER?) :: VOID</pre>	READ	N/A
db.awaitIndexes()	Wait for all indexes to come online (for example: CALL db.awaitIndexes(3 00)).	<pre>db.awaitIndexes(time0 utSeconds = 300 :: INTEGER?) :: VOID</pre>	READ	N/A

Name	Description	Signature	Mode	Roles
db.checkpoint()	Initiate and wait for a new check point, or wait any already on-going check point to complete. Note that this temporarily disables the dbms.checkpoint.i ops.limit setting in order to make the check point complete faster. This might cause transaction throughput to degrade slightly, due to increased IO load.	<pre>db.checkpoint() ::   (success :: BOOLEAN?,   message :: STRING?)</pre>	DBMS	N/A
db.clearQueryCach es()	Clears all query caches.	<pre>db.clearQueryCaches() :: (value :: STRING?)</pre>	DBMS	N/A
db.constraints()	List all constraints in the database.	<pre>db.constraints() ::   (name :: STRING?,   description ::   STRING?)</pre>	READ	N/A
db.createIndex()	Create a named schema index with specified index provider and configuration (optional). Yield: name, labels, properties, providerName, status	<pre>db.createIndex(indexN ame :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, config = {} :: MAP?) :: (name :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, status :: STRING?)</pre>	SCHEMA	N/A
db.createLabel()	Create a label	<pre>db.createLabel(newLab el :: STRING?) :: VOID</pre>	WRITE	N/A
db.createNodeKey(	Create a named node key constraint. Backing index will use specified index provider and configuration (optional). Yield: name, labels, properties, providerName, status	<pre>db.createNodeKey(cons traintName :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, config = {} :: MAP?) :: (name :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, properties :: LIST? STRING?, status :: STRING?, status :: STRING?)</pre>	SCHEMA	N/A

Name	Description	Signature	Mode	Roles
db.createProperty( )	Create a Property	<pre>db.createProperty(new Property :: STRING?) :: VOID</pre>	WRITE	N/A
db.createRelations hipType()	Create a RelationshipType	<pre>db.createRelationship Type(newRelationshipT ype :: STRING?) :: VOID</pre>	WRITE	N/A
db.createUniquePr opertyConstraint()	Create a named unique property constraint. Backing index will use specified index provider and configuration (optional). Yield: name, labels, properties, providerName, status	<pre>db.createUniqueProper tyConstraint(constrai ntName :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, config = {} :: MAP?) :: (name :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, status :: STRING?)</pre>	SCHEMA	N/A
db.index.fulltext.a waitEventuallyCon sistentIndexRefres h()	Wait for the updates from recently committed transactions to be applied to any eventually-consistent full-text indexes.	<pre>db.index.fulltext.awa itEventuallyConsisten tIndexRefresh() :: VOID</pre>	READ	N/A

Name	Description	Signature	Mode	Roles
db.index.fulltext.cr eateNodeIndex()	Create a node full-text index for the given labels and properties. The optional 'config' map parameter can be used to supply settings to the index. Supported settings are 'analyzer', for specifying what analyzer to use when indexing and querying. Use the db.index.fulltext.listAvailableAnalyzers procedure to see what options are available. And 'eventually_consist ent' which can be set to 'true' to make this index eventually consistent, such that updates from committing transactions are applied in a background thread.	<pre>db.index.fulltext.cre ateNodeIndex(indexNam e :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, config = {} :: MAP?) :: VOID</pre>	SCHEMA	N/A

Name	Description	Signature	Mode	Roles
db.index.fulltext.cr eateRelationshipIn dex()	Create a relationship full-text index for the given relationship types and properties. The optional 'config' map parameter can be used to supply settings to the index. Supported settings are 'analyzer', for specifying what analyzer to use when indexing and querying. Use the db.index.fulltext.listAvailableAnalyzers procedure to see what options are available. And 'eventually_consist ent' which can be set to 'true' to make this index eventually consistent, such that updates from committing transactions are applied in a background thread.	<pre>db.index.fulltext.cre   ateRelationshipIndex(   indexName :: STRING?,   relationshipTypes ::   LIST? OF STRING?,   properties :: LIST?   OF STRING?, config =   {} :: MAP?) :: VOID</pre>	SCHEMA	N/A
db.index.fulltext.dr op()	Drop the specified index.	<pre>db.index.fulltext.dro p(indexName :: STRING?) :: VOID</pre>	SCHEMA	N/A
db.index.fulltext.lis tAvailableAnalyzer s()	List the available analyzers that the full-text indexes can be configured with.	<pre>db.index.fulltext.lis tAvailableAnalyzers() :: (analyzer :: STRING?, description :: STRING?, stopwords :: LIST? OF STRING?)</pre>	READ	N/A
db.index.fulltext.q ueryNodes()	Query the given full-text index. Returns the matching nodes and their Lucene query score, ordered by score.	<pre>db.index.fulltext.que ryNodes(indexName :: STRING?, queryString :: STRING?) :: (node :: NODE?, score :: FLOAT?)</pre>	READ	N/A

Name	Description	Signature	Mode	Roles
db.index.fulltext.q ueryRelationships()	Query the given full-text index. Returns the matching relationships and their Lucene query score, ordered by score.	<pre>db.index.fulltext.que ryRelationships(index Name :: STRING?, queryString :: STRING?) :: (relationship :: RELATIONSHIP?, score :: FLOAT?)</pre>	READ	N/A
db.indexDetails()	Detailed description of specific index.	<pre>db.indexDetails(index Name :: STRING?) ::   (id :: INTEGER?, name</pre>	READ	N/A
db.indexes()	List all indexes in the database.	<pre>db.indexes() :: (id :: INTEGER?, name :: STRING?, state :: STRING?, populationPercent :: FLOAT?, uniqueness :: STRING?, type :: STRING?, entityType :: STRING?, labelsOrTypes :: LIST? OF STRING?, properties :: LIST? OF STRING?, provider :: STRING?)</pre>	READ	N/A
db.labels()	List all available labels in the database.	<pre>db.labels() :: (label :: STRING?)</pre>	READ	N/A
db.prepareForRepl anning()	Triggers an index resample and waits for it to complete, and after that clears query caches. After this procedure has finished queries will be planned using the latest database statistics.	<pre>db.prepareForReplanni ng(timeOutSeconds = 300 :: INTEGER?) :: VOID</pre>	READ	N/A
db.propertyKeys()	List all property keys in the database.	<pre>db.propertyKeys() ::   (propertyKey ::   STRING?)</pre>	READ	N/A

Name	Description	Signature	Mode	Roles
db.relationshipTyp es()	List all available relationship types in the database.	<pre>db.relationshipTypes( ) ::   (relationshipType ::   STRING?)</pre>	READ	N/A
db.resampleIndex(	Schedule resampling of an index (for example: CALL db.resampleIndex( "MyIndex")).	<pre>db.resampleIndex(inde xName :: STRING?) :: VOID</pre>	READ	N/A
db.resampleOutda tedIndexes()	Schedule resampling of all outdated indexes.	<pre>db.resampleOutdatedIn dexes() :: VOID</pre>	READ	N/A
db.schema.nodeTy peProperties()	Show the derived property schema of the nodes in tabular form.	<pre>db.schema.nodeTypePro perties() ::   (nodeType :: STRING?,   nodeLabels :: LIST?   OF STRING?,   propertyName ::   STRING?,   propertyTypes ::   LIST? OF STRING?,   mandatory ::   BOOLEAN?)</pre>	READ	N/A
db.schema.relType Properties()	Show the derived property schema of the relationships in tabular form.	<pre>db.schema.relTypeProp erties() :: (relType :: STRING?, propertyName :: STRING?, propertyTypes :: LIST? OF STRING?, mandatory :: BOOLEAN?)</pre>	READ	N/A
db.schema.visualiz ation()	Visualize the schema of the data.	db.schema.visualizati on() :: (nodes :: LIST? OF NODE?, relationships :: LIST? OF RELATIONSHIP?)	READ	N/A
db.schemaStatem ents()	List all statements for creating and dropping existing indexes and constraints.	<pre>db.schemaStatements() :: (name :: STRING?, type :: STRING?, createStatement :: STRING?, dropStatement :: STRING?)</pre>	READ	N/A
db.stats.clear()	Clear collected data of a given data section. Valid sections are 'QUERIES'	<pre>db.stats.clear(sectio n :: STRING?) :: (section :: STRING?, success :: BOOLEAN?, message :: STRING?)</pre>	READ	N/A

Name	Description	Signature	Mode	Roles
db.stats.collect()	Start data collection of a given data section. Valid sections are 'QUERIES'	<pre>db.stats.collect(sect ion :: STRING?, config = {} :: MAP?) :: (section :: STRING?, success :: BOOLEAN?, message :: STRING?)</pre>	READ	N/A
db.stats.retrieve()	Retrieve statistical data about the current database. Valid sections are 'GRAPH COUNTS', 'TOKENS', 'QUERIES', 'META'	<pre>db.stats.retrieve(sec tion :: STRING?, config = {} :: MAP?) :: (section :: STRING?, data :: MAP?)</pre>	READ	N/A
db.stats.retrieveAll Anonymized()	Retrieve all available statistical data about the current database, in an anonymized form.	<pre>db.stats.retrieveAllA nonymized(graphToken :: STRING?, config = {} :: MAP?) :: (section :: STRING?, data :: MAP?)</pre>	READ	N/A
db.stats.status()	Retrieve the status of all available collector daemons, for this database.	<pre>db.stats.status() ::   (section :: STRING?,   status :: STRING?,   data :: MAP?)</pre>	READ	N/A
db.stats.stop()	Stop data collection of a given data section. Valid sections are 'QUERIES'	<pre>db.stats.stop(section :: STRING?) :: (section :: STRING?, success :: BOOLEAN?, message :: STRING?)</pre>	READ	N/A
dbms.cluster.routi ng.getRoutingTabl e()	Returns endpoints of this instance.	<pre>dbms.cluster.routing. getRoutingTable(conte xt :: MAP?, database = null :: STRING?) :: (ttl :: INTEGER?, servers :: LIST? OF MAP?)</pre>	DBMS	N/A
dbms.components ()	List DBMS components and their versions.	<pre>dbms.components() :: (name :: STRING?, versions :: LIST? OF STRING?, edition :: STRING?)</pre>	DBMS	N/A
dbms.database.sta te()	The actual status of the database with the provided name on this neo4j instance.	<pre>dbms.database.state(d atabaseName :: STRING?) :: (role :: STRING?, address :: STRING?, status :: STRING?, error :: STRING?)</pre>	DBMS	N/A

Name	Description	Signature	Mode	Roles
dbms.functions()	List all functions in the DBMS.	<pre>dbms.functions() ::   (name :: STRING?,   signature :: STRING?,   description ::   STRING?, aggregating   :: BOOLEAN?,   defaultBuiltInRoles   :: LIST? OF STRING?)</pre>	DBMS	N/A
dbms.killConnection()	Kill network connection with the given connection id.	<pre>dbms.killConnection(i d :: STRING?) :: (connectionId :: STRING?, username :: STRING?, message :: STRING?)</pre>	DBMS	N/A
dbms.killConnections()	Kill all network connections with the given connection ids.	<pre>dbms.killConnections( ids :: LIST? OF STRING?) :: (connectionId :: STRING?, username :: STRING?, message :: STRING?)</pre>	DBMS	N/A
dbms.killQueries()	Kill all transactions executing a query with any of the given query ids.	<pre>dbms.killQueries(ids :: LIST? OF STRING?) :: (queryId :: STRING?, username :: STRING?, message :: STRING?)</pre>	DBMS	N/A
dbms.killQuery()	Kill all transactions executing the query with the given query id.	<pre>dbms.killQuery(id :: STRING?) :: (queryId :: STRING?, username :: STRING?, message :: STRING?)</pre>	DBMS	N/A
dbms.killTransacti on()	Kill transaction with provided id.	<pre>dbms.killTransaction( id :: STRING?) ::   (transactionId ::   STRING?, username ::   STRING?, message ::   STRING?)</pre>	DBMS	N/A
dbms.killTransacti ons()	Kill transactions with provided ids.	<pre>dbms.killTransactions (ids :: LIST? OF STRING?) :: (transactionId :: STRING?, username :: STRING?, message :: STRING?)</pre>	DBMS	N/A
dbms.listActiveLoc ks()	List the active lock requests granted for the transaction executing the query with the given query id.	<pre>dbms.listActiveLocks( queryId :: STRING?) :: (mode :: STRING?, resourceType :: STRING?, resourceId :: INTEGER?)</pre>	DBMS	N/A
dbms.listConfig()	List the currently active config of Neo4j.	<pre>dbms.listConfig(searc hString = :: STRING?) :: (name :: STRING?, description :: STRING?, value :: STRING?, dynamic :: BOOLEAN?)</pre>	DBMS	N/A

Name	Description	Signature	Mode	Roles
dbms.listConnecti ons()	List all accepted network connections at this instance that are visible to the user.	<pre>dbms.listConnections( ) :: (connectionId :: STRING?, connectTime :: STRING?, connector :: STRING?, username :: STRING?, userAgent :: STRING?, serverAddress :: STRING?, clientAddress :: STRING?)</pre>	DBMS	N/A
dbms.listQueries()	List all queries currently executing at this instance that are visible to the user.	<pre>dbms.listQueries() ::   (queryId :: STRING?,     username :: STRING?,   metaData :: MAP?,   query :: STRING?,   parameters :: MAP?,   planner :: STRING?,   runtime :: STRING?,   runtime :: STRING?,   indexes :: LIST? OF   MAP?, startTime ::   STRING?, protocol ::   STRING?, protocol ::   STRING?, requestUri   :: STRING?, status ::   STRING?, requestUri   :: STRING?, reacuestUri   :: STRING?, reduestUri   :: STRING?, status ::   STRING?,   resourceInformation   :: MAP?,   activeLockCount ::   INTEGER?,   elapsedTimeMillis ::   INTEGER?,   cultimeMillis ::   INTEGER?,   idleTimeMillis ::   INTEGER?,   allocatedBytes ::   INTEGER?,   allocatedBytes ::   INTEGER?, pageFaults   :: INTEGER?,   connectionId ::   STRING?, database ::   STRING?)</pre>	DBMS	N/A

Name	Description	Signature	Mode	Roles
dbms.listTransacti ons()	List all transactions currently executing at this instance that are visible to the user.	dbms.listTransactions () :: (transactionId :: STRING?, username :: STRING?, metaData :: MAP?, startTime :: STRING?, protocol :: STRING?, protocol :: STRING?, clientAddress :: STRING?, requestUri :: STRING?, currentQuery :: STRING?, currentQuery :: STRING?, currentQuery :: STRING?, requestUri :: STRING?, currentQuery :: STRING?, resourceInformation :: MAP?, elapsedTimeMillis :: INTEGER?, cpuTimeMillis :: INTEGER?, idleTimeMillis :: INTEGER?, allocatedBytes :: INTEGER?, allocatedDirectBytes :: INTEGER?, pageFaults :: INTEGER?, pageFaults :: INTEGER?, pageFaults :: INTEGER?, connectionId :: STRING?, initializationStackTr ace :: STRING?, database :: STRING?)	DBMS	N/A
dbms.procedures()	List all procedures in the DBMS.	<pre>dbms.procedures() :: (name :: STRING?, signature :: STRING?, description :: STRING?, mode :: STRING?, defaultBuiltInRoles :: LIST? OF STRING?, worksOnSystem :: BOOLEAN?)</pre>	DBMS	N/A
dbms.queryJmx()	Query JMX management data by domain and name. For instance, "org.neo4j:*"	<pre>dbms.queryJmx(query :: STRING?) :: (name :: STRING?, description :: STRING?, attributes :: MAP?)</pre>	DBMS	N/A
dbms.routing.getR outingTable()	Returns endpoints of this instance.	<pre>dbms.routing.getRouti ngTable(context :: MAP?, database = null :: STRING?) :: (ttl :: INTEGER?, servers :: LIST? OF MAP?)</pre>	DBMS	N/A
dbms.scheduler.gr oups()	List the job groups that are active in the database internal job scheduler.	<pre>dbms.scheduler.groups () :: (group :: STRING?, threads :: INTEGER?)</pre>	DBMS	N/A

Name	Description	Signature	Mode	Roles
dbms.scheduler.pr ofile()	Begin profiling all threads within the given job group, for the specified duration. Note that profiling incurs overhead to a system, and will slow it down.	<pre>dbms.scheduler.profil e(method :: STRING?, group :: STRING?, duration :: STRING?) :: (profile :: STRING?)</pre>	DBMS	N/A
dbms.security.acti vateUser()	Activate a suspended user.	<pre>dbms.security.activat eUser(username :: STRING?, requirePasswordChange = true :: BOOLEAN?) :: VOID</pre>	DBMS	N/A
dbms.security.add RoleToUser()	Assign a role to the user.	<pre>dbms.security.addRole ToUser(roleName :: STRING?, username :: STRING?) :: VOID</pre>	DBMS	N/A
dbms.security.cha ngePassword()	Change the current user's password.	<pre>dbms.security.changeP assword(password :: STRING?, requirePasswordChange = false :: BOOLEAN?) :: VOID</pre>	DBMS	N/A
dbms.security.cha ngeUserPassword( )	Change the given user's password.	<pre>dbms.security.changeU serPassword(username :: STRING?, newPassword :: STRING?, requirePasswordChange = true :: BOOLEAN?) :: VOID</pre>	DBMS	N/A
dbms.security.clea rAuthCache()	Clears authentication and authorization cache.	<pre>dbms.security.clearAu thCache() :: VOID</pre>	DBMS	N/A
dbms.security.crea teRole()	Create a new role.	<pre>dbms.security.createR ole(roleName :: STRING?) :: VOID</pre>	DBMS	N/A
dbms.security.crea teUser()	Create a new user.	<pre>dbms.security.createU ser(username :: STRING?, password :: STRING?, requirePasswordChange = true :: BOOLEAN?) :: VOID</pre>	DBMS	N/A
dbms.security.dele teRole()	Delete the specified role. Any role assignments will be removed.	<pre>dbms.security.deleteR ole(roleName :: STRING?) :: VOID</pre>	DBMS	N/A
dbms.security.dele teUser()	Delete the specified user.	<pre>dbms.security.deleteU ser(username :: STRING?) :: VOID</pre>	DBMS	N/A

Name	Description	Signature	Mode	Roles
dbms.security.listR oles()	List all available roles.	<pre>dbms.security.listRol es() :: (role :: STRING?, users :: LIST? OF STRING?)</pre>	DBMS	N/A
dbms.security.listR olesForUser()	List all roles assigned to the specified user.	<pre>dbms.security.listRol esForUser(username :: STRING?) :: (value :: STRING?)</pre>	DBMS	N/A
dbms.security.listU sers()	List all native users.	<pre>dbms.security.listUse rs() :: (username :: STRING?, roles :: LIST? OF STRING?, flags :: LIST? OF STRING?)</pre>	DBMS	N/A
dbms.security.listU sersForRole()	List all users currently assigned the specified role.	<pre>dbms.security.listUse rsForRole(roleName :: STRING?) :: (value :: STRING?)</pre>	DBMS	N/A
dbms.security.rem oveRoleFromUser()		<pre>dbms.security.removeR oleFromUser(roleName :: STRING?, username :: STRING?) :: VOID</pre>	DBMS	N/A
dbms.security.sus pendUser()	Suspend the specified user.	<pre>dbms.security.suspend User(username :: STRING?) :: VOID</pre>	DBMS	N/A
dbms.setConfigVal ue()	Updates a given setting value. Passing an empty value will result in removing the configured value and falling back to the default value. Changes will not persist and will be lost if the server is restarted.	<pre>dbms.setConfigValue(s etting :: STRING?, value :: STRING?) :: VOID</pre>	DBMS	N/A
dbms.showCurren tUser()	Show the current user.	<pre>dbms.showCurrentUser( ) :: (username :: STRING?, roles :: LIST? OF STRING?, flags :: LIST? OF STRING?)</pre>	DBMS	N/A
tx.getMetaData()	Provides attached transaction metadata.	<pre>tx.getMetaData() :: (metadata :: MAP?)</pre>	DBMS	N/A

Name	Description	Signature	Mode	Roles
tx.setMetaData()	Attaches a map of data to the transaction. The data will be printed when listing queries, and inserted into the query log.	tx.setMetaData(data :: MAP?) :: VOID	DBMS	N/A

# A.2.3. Community Edition procedures

Table 321. Community Edition procedures

Name	Description	Signature	Mode
db.awaitIndex()	Wait for an index to come online (for example: CALL db.awaitIndex("MyIndex ", 300)).	<pre>db.awaitIndex(indexName :: STRING?, timeOutSeconds = 300 :: INTEGER?) :: VOID</pre>	READ
db.awaitIndexes()	Wait for all indexes to come online (for example: CALL db.awaitIndexes(300)).	<pre>db.awaitIndexes(timeOutSec onds = 300 :: INTEGER?) :: VOID</pre>	READ
db.clearQueryCaches()	Clears all query caches.	<pre>db.clearQueryCaches() ::   (value :: STRING?)</pre>	DBMS
db.constraints()	List all constraints in the database.	<pre>db.constraints() :: (name :: STRING?, description :: STRING?)</pre>	READ
db.createIndex()	Create a named schema index with specified index provider and configuration (optional). Yield: name, labels, properties, providerName, status	<pre>db.createIndex(indexName :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, config = {} :: MAP?) :: (name :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, status :: STRING?)</pre>	SCHEMA
db.createLabel()	Create a label	<pre>db.createLabel(newLabel :: STRING?) :: VOID</pre>	WRITE
db.createProperty()	Create a Property	<pre>db.createProperty(newPrope rty :: STRING?) :: VOID</pre>	WRITE
db.createRelationshipTy pe()	Create a RelationshipType	<pre>db.createRelationshipType( newRelationshipType :: STRING?) :: VOID</pre>	WRITE

Name	Description	Signature	Mode
db.createUniqueProper tyConstraint()	Create a named unique property constraint. Backing index will use specified index provider and configuration (optional). Yield: name, labels, properties, providerName, status	<pre>db.createUniquePropertyCon straint(constraintName :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, config = {} :: MAP?) :: (name :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, providerName :: STRING?, status :: STRING?)</pre>	SCHEMA
db.index.fulltext.awaitE ventuallyConsistentInde xRefresh()	Wait for the updates from recently committed transactions to be applied to any eventually-consistent full-text indexes.	<pre>db.index.fulltext.awaitEve ntuallyConsistentIndexRefr esh() :: VOID</pre>	READ
db.index.fulltext.create NodeIndex()	Create a node full-text index for the given labels and properties. The optional 'config' map parameter can be used to supply settings to the index. Supported settings are 'analyzer', for specifying what analyzer to use when indexing and querying. Use the db.index.fulltext.list AvailableAnalyzers procedure to see what options are available. And 'eventually_consistent' which can be set to 'true' to make this index eventually consistent, such that updates from committing transactions are applied in a background thread.	<pre>db.index.fulltext.createNo deIndex(indexName :: STRING?, labels :: LIST? OF STRING?, properties :: LIST? OF STRING?, config = {} :: MAP?) :: VOID</pre>	SCHEMA

Name	Description	Signature	Mode
db.index.fulltext.create RelationshipIndex()	Create a relationship full-text index for the given relationship types and properties. The optional 'config' map parameter can be used to supply settings to the index. Supported settings are 'analyzer', for specifying what analyzer to use when indexing and querying. Use the db.index.fulltext.list AvailableAnalyzers procedure to see what options are available. And 'eventually_consistent' which can be set to 'true' to make this index eventually consistent, such that updates from committing transactions are applied in a background thread.	<pre>db.index.fulltext.createRe lationshipIndex(indexName :: STRING?, relationshipTypes :: LIST? OF STRING?, properties :: LIST? OF STRING?, config = {} :: MAP?) :: VOID</pre>	SCHEMA
db.index.fulltext.drop()	Drop the specified index.	<pre>db.index.fulltext.drop(ind exName :: STRING?) :: VOID</pre>	SCHEMA
db.index.fulltext.listAvai lableAnalyzers()	List the available analyzers that the full- text indexes can be configured with.	<pre>db.index.fulltext.listAvai lableAnalyzers() :: (analyzer :: STRING?, description :: STRING?, stopwords :: LIST? OF STRING?)</pre>	READ
db.index.fulltext.queryN odes()	Query the given full-text index. Returns the matching nodes and their Lucene query score, ordered by score.	<pre>db.index.fulltext.queryNod es(indexName :: STRING?, queryString :: STRING?) :: (node :: NODE?, score :: FLOAT?)</pre>	READ
db.index.fulltext.queryR elationships()	Query the given full-text index. Returns the matching relationships and their Lucene query score, ordered by score.	<pre>db.index.fulltext.queryRel ationships(indexName :: STRING?, queryString :: STRING?) :: (relationship :: RELATIONSHIP?, score :: FLOAT?)</pre>	READ

Name	Description	Signature	Mode
db.indexDetails()	Detailed description of specific index.	db.indexDetails(indexName :: STRING?) :: (id :: INTEGER?, name :: STRING?, state :: STRING?, populationPercent :: FLOAT?, uniqueness :: STRING?, type :: STRING?, entityType :: STRING?, labelsOrTypes :: LIST? OF STRING?, properties :: LIST? OF STRING?, provider :: STRING?, indexConfig :: MAP?, failureMessage :: STRING?)	READ
db.indexes()	List all indexes in the database.	<pre>db.indexes() :: (id :: INTEGER?, name :: STRING?, state :: STRING?, populationPercent :: FLOAT?, uniqueness :: STRING?, type :: STRING?, entityType :: STRING?, labelsOrTypes :: LIST? OF STRING?, properties :: LIST? OF STRING?, provider :: STRING?)</pre>	READ
db.labels()	List all available labels in the database.	<pre>db.labels() :: (label :: STRING?)</pre>	READ
db.prepareForReplanni ng()	Triggers an index resample and waits for it to complete, and after that clears query caches. After this procedure has finished queries will be planned using the latest database statistics.	<pre>db.prepareForReplanning(ti meOutSeconds = 300 :: INTEGER?) :: VOID</pre>	READ
db.propertyKeys()	List all property keys in the database.	<pre>db.propertyKeys() :: (propertyKey :: STRING?)</pre>	READ
db.relationshipTypes()	List all available relationship types in the database.	<pre>db.relationshipTypes() ::   (relationshipType ::    STRING?)</pre>	READ
db.resampleIndex()	Schedule resampling of an index (for example: CALL db.resampleIndex("MyI ndex")).	<pre>db.resampleIndex(indexName :: STRING?) :: VOID</pre>	READ
db.resampleOutdatedIn dexes()	Schedule resampling of all outdated indexes.	<pre>db.resampleOutdatedIndexes () :: VOID</pre>	READ

Name	Description	Signature	Mode
db.schema.nodeTypePr operties()	Show the derived property schema of the nodes in tabular form.	<pre>db.schema.nodeTypeProperti es() :: (nodeType :: STRING?, nodeLabels :: LIST? OF STRING?, propertyName :: STRING?, propertyTypes :: LIST? OF STRING?, mandatory :: BOOLEAN?)</pre>	READ
db.schema.relTypeProp erties()	Show the derived property schema of the relationships in tabular form.	<pre>db.schema.relTypePropertie s() :: (relType :: STRING?, propertyName :: STRING?, propertyTypes :: LIST? OF STRING?, mandatory :: BOOLEAN?)</pre>	READ
db.schema.visualization ()	Visualize the schema of the data.	<pre>db.schema.visualization() :: (nodes :: LIST? OF NODE?, relationships :: LIST? OF RELATIONSHIP?)</pre>	READ
db.schemaStatements()	List all statements for creating and dropping existing indexes and constraints.	<pre>db.schemaStatements() ::   (name :: STRING?, type ::   STRING?, createStatement   :: STRING?, dropStatement   :: STRING?)</pre>	READ
db.stats.clear()	Clear collected data of a given data section. Valid sections are 'QUERIES'	<pre>db.stats.clear(section :: STRING?) :: (section :: STRING?, success :: BOOLEAN?, message :: STRING?)</pre>	READ
db.stats.collect()	Start data collection of a given data section. Valid sections are 'QUERIES'	<pre>db.stats.collect(section :: STRING?, config = {} :: MAP?) :: (section :: STRING?, success :: BOOLEAN?, message :: STRING?)</pre>	READ
db.stats.retrieve()	Retrieve statistical data about the current database. Valid sections are 'GRAPH COUNTS', 'TOKENS', 'QUERIES', 'META'	<pre>db.stats.retrieve(section :: STRING?, config = {} :: MAP?) :: (section :: STRING?, data :: MAP?)</pre>	READ
db.stats.retrieveAllAnon ymized()	Retrieve all available statistical data about the current database, in an anonymized form.	<pre>db.stats.retrieveAllAnonym ized(graphToken :: STRING?, config = {} :: MAP?) :: (section :: STRING?, data :: MAP?)</pre>	READ
db.stats.status()	Retrieve the status of all available collector daemons, for this database.	<pre>db.stats.status() ::   (section :: STRING?,   status :: STRING?, data :: MAP?)</pre>	READ
db.stats.stop()	Stop data collection of a given data section. Valid sections are 'QUERIES'	<pre>db.stats.stop(section :: STRING?) :: (section :: STRING?, success :: BOOLEAN?, message :: STRING?)</pre>	READ

Name	Description	Signature	Mode
dbms.cluster.routing.ge tRoutingTable()	Returns endpoints of this instance.	<pre>dbms.cluster.routing.getRo utingTable(context :: MAP?, database = null :: STRING?) :: (ttl :: INTEGER?, servers :: LIST? OF MAP?)</pre>	DBMS
dbms.components()	List DBMS components and their versions.	<pre>dbms.components() :: (name :: STRING?, versions :: LIST? OF STRING?, edition :: STRING?)</pre>	DBMS
dbms.database.state()	The actual status of the database with the provided name on this neo4j instance.	<pre>dbms.database.state(databa seName :: STRING?) :: (role :: STRING?, address :: STRING?, status :: STRING?, error :: STRING?)</pre>	DBMS
dbms.functions()	List all functions in the DBMS.	<pre>dbms.functions() :: (name :: STRING?, signature :: STRING?, description :: STRING?, aggregating :: BOOLEAN?, defaultBuiltInRoles :: LIST? OF STRING?)</pre>	DBMS
dbms.listConfig()	List the currently active config of Neo4j.	<pre>dbms.listConfig(searchStri ng = :: STRING?) :: (name :: STRING?, description :: STRING?, value :: STRING?, dynamic :: BOOLEAN?)</pre>	DBMS
dbms.procedures()	List all procedures in the DBMS.	<pre>dbms.procedures() :: (name :: STRING?, signature :: STRING?, description :: STRING?, mode :: STRING?, defaultBuiltInRoles :: LIST? OF STRING?, worksOnSystem :: BOOLEAN?)</pre>	DBMS
dbms.queryJmx()	Query JMX management data by domain and name. For instance, "org.neo4j:*"	<pre>dbms.queryJmx(query :: STRING?) :: (name :: STRING?, description :: STRING?, attributes :: MAP?)</pre>	DBMS
dbms.routing.getRoutin gTable()	Returns endpoints of this instance.	<pre>dbms.routing.getRoutingTab le(context :: MAP?, database = null :: STRING?) :: (ttl :: INTEGER?, servers :: LIST? OF MAP?)</pre>	DBMS
dbms.security.changeP assword()	Change the current user's password.	<pre>dbms.security.changePasswo rd(password :: STRING?) :: VOID</pre>	DBMS
dbms.security.createUs er()	Create a new user.	<pre>dbms.security.createUser(u sername :: STRING?, password :: STRING?, requirePasswordChange = true :: BOOLEAN?) :: VOID</pre>	DBMS
dbms.security.deleteUs er()	Delete the specified user.	<pre>dbms.security.deleteUser(u sername :: STRING?) :: VOID</pre>	DBMS
dbms.security.listUsers( )	List all native users.	<pre>dbms.security.listUsers() :: (username :: STRING?, roles :: LIST? OF STRING?, flags :: LIST? OF STRING?)</pre>	DBMS

Name	Description	Signature	Mode
dbms.showCurrentUser ()	Show the current user.	<pre>dbms.showCurrentUser() ::   (username :: STRING?,   roles :: LIST? OF STRING?,   flags :: LIST? OF STRING?)</pre>	DBMS
tx.getMetaData()	Provides attached transaction metadata.	<pre>tx.getMetaData() ::   (metadata :: MAP?)</pre>	DBMS
tx.setMetaData()	Attaches a map of data to the transaction. The data will be printed when listing queries, and inserted into the query log.	tx.setMetaData(data :: MAP?) :: VOID	DBMS

# Appendix B: Tutorials

This appendix contains examples and tutorials that further describe usages of Neo4j.

The following step-by-step tutorials cover common operational tasks or otherwise exemplify working with Neo4j.

- Set up a local Causal Cluster
- Use the Import tool to import data into Neo4j

## B.1. Set up a local Causal Cluster

This tutorial walks through the basics of setting up a Neo4j Causal Cluster. The result is a local cluster of six instances: three Cores and three Read Replicas.

#### This tutorial describes the following:

- Introduction
- Download Neo4j and configure the Core instances
  - □ Configure the first Core instance
  - ☐ Configure the second Core instance
  - ☐ Configure the third Core instance
- Start the Core Servers
- Check the status of the cluster
- Configure the Read Replicas
  - ☐ Configure the first Read Replica
  - ☐ Configure the second Read Replica
  - ☐ Configure the third Read Replica
- Start the Read Replicas
- Test the cluster with Read Replicas

### Introduction

In this tutorial we will learn how to deploy a Causal Cluster locally, on a single machine. This is a useful learning exercise, and will get you started quickly with developing an application against a Neo4j Causal Cluster. Please note that in practice, a cluster on a single machine has no fault tolerance and is therefore not suitable for production use.

We will begin by configuring and starting a cluster of three Core instances. This is the minimal deployment for a fault-tolerant Causal Cluster (for a discussion on the number of servers required for a Causal Cluster, see Core Servers). The Core instances are responsible for keeping the data safe.

After the Core of the cluster is operational we will add three Read Replicas. The Read Replicas are responsible for scaling the capacity of the cluster.

The Core of the Causal Cluster is intended to remain stable over time. The roles within the Core will change as needed, but the Core itself is long-lived and stable. Read Replicas live at the edge of the cluster and can be brought up and taken down without affecting the Core. They can be added as

needed to increase the operational capacity of the cluster as a whole.

In this tutorial we will be running all instances in the cluster on a single machine. Many of the default configuration settings work well out of the box in a production deployment, with multiple machines. Some of these we have to change when deploying multiple instances on a single machine, so that instances do not try to use the same network ports. We call out the settings that are specific to this scenario as we go along.

## Download Neo4j and configure the Core instances

- 1. Create a local working directory.
- 2. Download a copy of Neo4j Enterprise Edition from the Neo4j download site.
- 3. Unpack Neo4j in the working directory.
- 4. Make a copy of the *neo4j-enterprise-4.0.0* directory, and name it *core-01*. We will need to keep the original directory to use when setting up the Read Replicas later. The *core-01* directory will contain the first Core instance.

### Configure the first Core instance

All configuration that we will do takes place in the Neo4j configuration file, *conf/neo4j.conf*. If you used a different package than in the download instructions above, see File locations to locate the configuration file.

The first settings we will change represent the minimum configuration for a Core instance:

- 1. Locate and uncomment the setting dbms.mode=CORE.
- 2. Locate and uncomment the setting causal\_clustering.minimum\_core\_cluster\_size\_at\_formation=3.
- 3. Locate and uncomment the setting causal\_clustering.minimum\_core\_cluster\_size\_at\_runtime=3.
- 4. Locate and uncomment the setting causal\_clustering.initial\_discovery\_members=localhost:5000,localhost:5001,localhost:5002.

Since we are setting up the Causal Cluster to run on a single machine, we must do some additional configuration. Please note that these steps would not be necessary if the instances are running on different servers.

- 1. Locate and uncomment the setting causal\_clustering.discovery\_listen\_address=:5000.
- 2. Locate and uncomment the setting causal\_clustering.transaction\_listen\_address=:6000.
- Locate and uncomment the setting causal\_clustering.raft\_listen\_address=:7000.
- 4. Locate and uncomment the setting dbms.connector.bolt.listen\_address=:7687.
- 5. Locate and uncomment the setting dbms.connector.http.listen\_address=:7474.
- 6. Locate the dbms.connector.https.listen\_address setting and change the value to :6474.
- 7. Locate the dbms.backup.listen\_address setting and change the value to :6362.

### Configure the second Core instance

We can now create the second Core instance.

As mentioned above, we also need to amend some of the additional values in the *neo4j.conf* file so that our cluster can run on a single machine:

1. Make a copy of the *core-01* directory and rename it *core-02*.

- 2. Open the *neo4j.conf* file in the new *core-02* directory.
  - a. Locate the causal\_clustering.discovery\_listen\_address setting and change the value to :5001.
  - b. Locate the causal\_clustering.transaction\_listen\_address setting and change the value to :6001.
  - C. Locate the causal\_clustering.raft\_listen\_address setting and change the value to :7001.
  - d. Locate the dbms.connector.bolt.listen\_address setting and change the value to :7688.
  - e. Locate the dbms.connector.http.listen\_address setting and change the value to :7475.
  - f. Locate the dbms.connector.https.listen\_address setting and change the value to :6475.
  - g. Locate the dbms.backup.listen\_address setting and change the value to :6363.

### Configure the third Core instance

We can now create the third, and final Core instance.

Again, we also need to amend some of the additional values in the *neo4j.conf* file so that our cluster can run on a single machine:

- 1. Make a copy of the *core-02* directory and rename it *core-03*.
- 2. Open the *neo4j.conf* file in the new *core-03* directory.
  - a. Locate the causal\_clustering.discovery\_listen\_address setting and change the value to :5002.
  - b. Locate the causal\_clustering.transaction\_listen\_address setting and change the value to :6002.
  - C. Locate the causal\_clustering.raft\_listen\_address setting and change the value to :7002.
  - d. Locate the dbms.connector.bolt.listen\_address setting and change the value to :7689.
  - e. Locate the dbms.connector.http.listen\_address setting and change the value to :7476.
  - f. Locate the dbms.connector.https.listen\_address setting and change the value to :6476.
  - g. Locate the dbms.backup.listen\_address setting and change the value to :6364.

### Start the Core servers

In any order, we can now start each of the Neo4j instances:

```
core-01$ ./bin/neo4j start

core-02$ ./bin/neo4j start

core-03$ ./bin/neo4j start
```

#### Startup Time

If you want to follow along with the startup of a server you can follow the messages in *logs/neo4j.log*:



- On a Unix system issue the command tail -f logs/neo4j.log.
- On Windows Server run Get-Content .\logs\neo4j.log -Tail 10 -Wait.

While an instance is joining the cluster, the server may appear unavailable. In the case where an instance is joining a cluster with lots of data, it may take a number of minutes for the new instance to download the data from the cluster and become available.

### Check the status of the cluster

Now the minimal cluster of three Core Servers is operational and is ready to serve requests.

- 1. Connect to any of the three Core instances to check the cluster status. For example, for *core-01* point your web browser to http://localhost:7474.
- 2. Authenticate with the default neo4j/neo4j credentials, and set a new password when prompted. These credentials are not shared between cluster members. A new password must be set on each instance when connecting for the first time.
- 3. Check the status of the cluster by running the following in Neo4j Browser:

:sysinfo

Example 110. Example of a Cluster overview, achieved by running :sysinfo

The **Causal Cluster Members** table shows the status of instances in the cluster. The table below is an example of a test cluster:

Roles	Addresses	Actions
LEADER	bolt://localhost:7687, http://localhost:7474, https://localhost:6474	0pen
FOLLOWER	bolt://localhost:7688, http://localhost:7475, https://localhost:6475	0pen
FOLLOWER	bolt://localhost:7689, http://localhost:7476, https://localhost:6476	0pen

The three Core instances in this cluster are now operational.

Now you can run queries to create nodes and relationships, and see that the data gets replicated in the cluster.

- 4. Click the Open action on the instance that has the *LEADER* role. This will open a new Neo4j Browser session against the Leader of the cluster.
- 5. Authenticate and set a new password, as before.
- 6. Run the following query to create nodes and relationships:

```
UNWIND range(0, 100) AS value

MERGE (person1:Person {id: value})

MERGE (person2:Person {id: toInteger(100.0 * rand())})

MERGE (person1)-[:FRIENDS]->(person2)
```

- 7. When the query has executed, choose an instance with the *FOLLOWER* role from the sysinfo view. Click the Open action to connect.
- 8. Run the following query to see that the data has been replicated:

```
MATCH path = (person:Person)-[:FRIENDS]-(friend)
RETURN path
LIMIT 10
```

## Configure the Read Replicas

Read Replicas instances do not participate in quorum decisions, so their configuration is simpler than the configuration of Core Servers as there are fewer settings to amend.

All that a Read Replica needs to know is the addresses of Core Servers which they can bind to in order to discover the cluster. See Discovery protocol for the details of how this works. Once it has completed the initial discovery, the Read Replica becomes aware of the currently available Core Servers and can choose an appropriate one from which to catch up. See Catchup protocol for the details of how this works.

### Configure the first Read Replica

- 1. In your working directory, make a copy of the *neo4j-enterprise-4.0.0* directory and name it *replica-01*.
- 2. Open the *neo4j.conf* file in the new *replicα-01* directory. The first settings we will change represent the minimum configuration for a Read Replica:
  - a. Locate and uncomment the dbms.mode setting and change the value to READ\_REPLICA.
  - b. Locate and uncomment the setting causal\_clustering.initial\_discovery\_members=localhost:5000,localhost:5001,localhost:5002
- 3. Since we are setting up the Causal Cluster to run on a single machine, we must do some additional configuration. Please note that the following steps would not be necessary if the instances are running on different servers:
  - a. Locate and uncomment the causal\_clustering.transaction\_listen\_address setting and change the value to :6003.
  - b. Locate and uncomment the dbms.connector.bolt.listen\_address setting and change the value to :7690.
  - c. Locate and uncomment the dbms.connector.http.listen\_address setting and change the value to:7477.
  - d. Locate and uncomment the dbms.connector.https.listen\_address setting and change the value to :6477.
  - e. Locate and uncomment the dbms.backup.listen\_address setting and change the value to :6365.

### Configure the second Read Replica

We can now create the second Read Replica.

As mentioned above, we also need to amend some of the additional values in the *neo4j.conf* file so that our cluster can run on a single machine:

- 1. Make a copy of the *replica-01* directory and rename it *replica-02*.
- 2. Open the *neo4j.conf* file in the new *replica-02* directory.
  - a. Locate the causal\_clustering.transaction\_listen\_address setting and change the value to :6004.
  - b. Locate the dbms.connector.bolt.listen\_address setting and change the value to :7691.
  - C. Locate the dbms.connector.http.listen\_address setting and change the value to :7478.
  - d. Locate the dbms.connector.https.listen\_address setting and change the value to :6478.
  - e. Locate the dbms.backup.listen\_address setting and change the value to :6366.

### Configure the third Read Replica

We can now create the third, and final Read Replica.

Again, we also need to amend some of the additional values in the *neo4j.conf* file so that our cluster can run on a single machine:

- 1. Make a copy of the *replica-02* directory and rename it *replica-03*.
- 2. Open the *neo4j.conf* file in the new *replica-03* directory.
  - a. Locate the causal\_clustering.transaction\_listen\_address setting and change the value to :6005.
  - b. Locate the dbms.connector.bolt.listen\_address setting and change the value to :7692.
  - C. Locate the dbms.connector.http.listen\_address setting and change the value to :7479.
  - d. Locate the dbms.connector.https.listen\_address setting and change the value to :6479.
  - e. Locate the dbms.backup.listen\_address setting and change the value to :6367.

## Start the Read Replicas

In any order, we can now start the Read Replica instances:

```
replica-01$ ./bin/neo4j start

replica-02$ ./bin/neo4j start

replica-03$ ./bin/neo4j start
```

## Test the cluster with Read Replicas

To test the status of the cluster now that the Read Replicas are running, we will repeat the steps from earlier, but via a Read Replica:

- 1. Connect to any of the three Read Replica instances. For example, for *replica-01* point your web browser to http://localhost:7477.
- 2. Authenticate with the default neo4j/neo4j credentials. Once again, you will need to set a new password when prompted.

3. Check the status of the cluster by running the following in Neo4j Browser:

```
:sysinfo
```

Example 111. Example of a cluster with both Core instances and Read Replicas, achieved by running :sysinfo

The following table shows the status of a test cluster which now includes Read Replicas:

Roles	Addresses	Actions	
LEADER	bolt://localhost:7687, http://localhost:7474, https://localhost:6474	Open	
FOLLOWER	bolt://localhost:7688, http://localhost:7475, https://localhost:6475	Open	
FOLLOWER	bolt://localhost:7689, http://localhost:7476, https://localhost:6476	Open	
READ_REPLICA	bolt://localhost:7690, http://localhost:7477, https://localhost:6477	Open	
READ_REPLICA	bolt://localhost:7691, http://localhost:7478, https://localhost:6478	Open	
READ_REPLICA	bolt://localhost:7692, http://localhost:7479, https://localhost:6479	Open	

- 4. Click the Open action to connect to any of the Read Replicas.
- 5. Run the same query as before:

```
MATCH path = (person:Person)-[:FRIENDS]-(friend)
RETURN path
LIMIT 10
```

## B.2. Use the Import tool

This tutorial provides detailed examples of using the Neo4j import tool.

This tutorial walks us through a series of examples to illustrate the capabilities of the Import tool.

When using CSV files for loading a database, each node must have a unique ID to be able to be referenced when creating relationships between nodes in the same import. In cases where the node ID is only unique within files, using *ID spaces* is a way to ensure uniqueness across all nodes files. Please see below and Using ID spaces for more information.

Relationships are created by connecting the node IDs. In the examples below, the node IDs are stored as properties on the nodes. Node IDs may be of interest later for cross-reference to other systems, traceability etc., but they are not mandatory. If you do not want the IDs to persist after a completed import, then do not specify a property name in the :ID field.

It is possible to import only nodes using the import tool by omitting a relationships file when calling

neo4j-admin import. Any relationships between the imported nodes will have to be created later by another method, since the import tool works for initial graph population only.

For this tutorial we will use a data set containing movies, actors and roles. If not stated otherwise, the examples assume that the name of the database is neo4j (the default name) and that all CSV files are located in the *import* directory. Note that if you wish to run one example after another you have to remove the database in between.



#### Header files

For the basic examples we will use the first row of the data file for the header. This is fine when experimenting, but when working with anything but the smallest datasets we recommend keeping the header in a separate file.

### B.2.1. Basic example

First we will look at the movies. Each movie has an ID, which is used for referring to it from other data sources. Moreover, each movie also has a title and a year. Along with these properties we also add the node labels Movie and Sequel:

#### movies.csv

```
movieId:ID,title,year:int,:LABEL
tt0133093,"The Matrix",1999,Movie
tt0234215,"The Matrix Reloaded",2003,Movie;Sequel
tt0242653,"The Matrix Revolutions",2003,Movie;Sequel
```

Next up are the actors. They have an ID - in this case a shorthand of their name - and a name. All the actors have the node label Actor:

#### actors.csv

```
personId:ID,name,:LABEL
keanu,"Keanu Reeves",Actor
laurence,"Laurence Fishburne",Actor
carrieanne,"Carrie-Anne Moss",Actor
```



The node ID is needed to create relationships, but the node labels are optional.

Finally we have the roles that an actor plays in a movie, which will be represented by relationships in the database. There are three mandatory fields for relationships; START\_ID, END\_ID and :TYPE.

In order to create a relationship between nodes we use the IDs defined in <a href="actors.csv">actors.csv</a> and <a href="movies.csv">movies.csv</a> for the <a href="movies.csv">START\_ID</a> and <a href="movies.csv">END\_ID</a> fields. We also need to provide a relationship type (in this case <a href="movies.csv">ACTED\_IN</a>) for the <a href="movies.csv">TYPE</a> field:

#### roles.csv

```
:START_ID,role,:END_ID,:TYPE
keanu,"Neo",tt0133093,ACTED_IN
keanu,"Neo",tt0234215,ACTED_IN
keanu,"Neo",tt0242653,ACTED_IN
laurence,"Morpheus",tt0133093,ACTED_IN
laurence,"Morpheus",tt0234215,ACTED_IN
laurence,"Morpheus",tt0242653,ACTED_IN
carrieanne,"Trinity",tt0133093,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
```

The call to neo4j-admin import would look like this:

neo4j\_home\$ bin/neo4j-admin import --nodes=import/movies.csv --nodes=import/actors.csv
--relationships=import/roles.csv



In Community Edition, the database neo4j must be empty before the import, or else the import will fail. If neo4j is not empty, the entire directory may be deleted as it is created on import if non-existent. However, in Enterprise Edition, if there is an unfinished import in the neo4j folder, the tool will attempt to complete that import but fail if it cannot.

Now start up a database from the target directory:

```
neo4j_home$ ./bin/neo4j start
```

To see your imported data in the graph, try a simple command:

```
MATCH (n)
RETURN count(n) as count
```

## B.2.2. Customizing configuration options

We can customize the configuration options that the import tool uses (see Options) if our data does not fit the default format. The following CSV files are delimited by ;, use | as the array delimiter and use ' for quotes:

#### movies2.csv

```
movieId:ID;title;year:int;:LABEL
tt0133093;'The Matrix';1999;Movie
tt0234215;'The Matrix Reloaded';2003;Movie|Sequel
tt0242653;'The Matrix Revolutions';2003;Movie|Sequel
```

#### actors2.csv

```
personId:ID;name;:LABEL
keanu;'Keanu Reeves';Actor
laurence;'Laurence Fishburne';Actor
carrieanne;'Carrie-Anne Moss';Actor
```

#### roles2.csv

```
:START_ID;role;:END_ID;:TYPE
keanu;'Neo';tt0133093;ACTED_IN
keanu;'Neo';tt0234215;ACTED_IN
keanu;'Neo';tt0242653;ACTED_IN
laurence;'Morpheus';tt0133093;ACTED_IN
laurence;'Morpheus';tt0234215;ACTED_IN
laurence;'Morpheus';tt0242653;ACTED_IN
carrieanne;'Trinity';tt0133093;ACTED_IN
carrieanne;'Trinity';tt0234215;ACTED_IN
carrieanne;'Trinity';tt0242653;ACTED_IN
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies2.csv --nodes=import/actors2.csv
--relationships=import/roles2.csv --delimiter=";" --array-delimiter="|" --quote="'"
```

## B.2.3. Using separate header files

When dealing with very large CSV files it is more convenient to have the header in a separate file. This makes it easier to edit the header as you avoid having to open a huge data file just to change it. The header file must be specified before the rest of the files in each file group.

The import tool can also process single file compressed archives, for example: . --nodes nodes.csv.gz . --relationships rels.zip

We will use the same data as in the previous example but put the headers in separate files:

#### movies3-header.csv

```
movieId:ID,title,year:int,:LABEL
```

#### movies3.csv

```
tt0133093,"The Matrix",1999,Movie
tt0234215,"The Matrix Reloaded",2003,Movie;Sequel
tt0242653,"The Matrix Revolutions",2003,Movie;Sequel
```

#### actors3-header.csv

```
personId:ID, name,:LABEL
```

#### actors3.csv

```
keanu, "Keanu Reeves", Actor
laurence, "Laurence Fishburne", Actor
carrieanne, "Carrie-Anne Moss", Actor
```

#### roles3-header.csv

```
:START_ID,role,:END_ID,:TYPE
```

#### roles3.csv

```
keanu, "Neo", tt0133093, ACTED_IN
keanu, "Neo", tt0234215, ACTED_IN
keanu, "Neo", tt0242653, ACTED_IN
laurence, "Morpheus", tt0133093, ACTED_IN
laurence, "Morpheus", tt0234215, ACTED_IN
laurence, "Morpheus", tt0234215, ACTED_IN
carrieanne, "Trinity", tt0133093, ACTED_IN
carrieanne, "Trinity", tt0133093, ACTED_IN
carrieanne, "Trinity", tt0234215, ACTED_IN
carrieanne, "Trinity", tt0242653, ACTED_IN
```

The call to neo4j-admin import would look as follows, note how the file groups are enclosed in quotation marks in the command:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies3-header.csv,import/movies3.csv
--nodes=import/actors3-header.csv,import/actors3.csv --relationships=import/roles3
-header.csv,import/roles3.csv
```



The header line for a file group, whether it is the first line of a file in the group or a dedicated header file, must be the *first* line in the file group.

## B.2.4. Multiple input files

In addition to using a separate header file you can also provide multiple nodes or relationships files. This may be useful for example for processing the output from a Hadoop pipeline. Files within such an input group can be specified with multiple match strings, delimited by ,, where each match string can be either the exact file name or a regular expression matching one or more files. Multiple matching files will be sorted according to their characters and their natural number sort order for file names containing numbers:

#### movies4-header.csv

```
movieId:ID,title,year:int,:LABEL
```

### movies4-part1.csv

```
tt0133093,"The Matrix",1999,Movie
tt0234215,"The Matrix Reloaded",2003,Movie;Sequel
```

#### movies4-part2.csv

```
tt0242653, "The Matrix Revolutions", 2003, Movie; Sequel
```

#### actors4-header.csv

```
personId:ID,name,:LABEL
```

#### actors4-part1.csv

```
keanu, "Keanu Reeves", Actor
laurence, "Laurence Fishburne", Actor
```

#### actors4-part2.csv

```
carrieanne, "Carrie-Anne Moss", Actor
```

#### roles4-header.csv

```
:START_ID,role,:END_ID,:TYPE
```

#### roles4-part1.csv

```
keanu, "Neo", tt0133093, ACTED_IN
keanu, "Neo", tt0234215, ACTED_IN
keanu, "Neo", tt0242653, ACTED_IN
laurence, "Morpheus", tt0133093, ACTED_IN
laurence, "Morpheus", tt0234215, ACTED_IN
```

#### roles4-part2.csv

```
laurence, "Morpheus", tt0242653, ACTED_IN carrieanne, "Trinity", tt0133093, ACTED_IN carrieanne, "Trinity", tt0234215, ACTED_IN carrieanne, "Trinity", tt0242653, ACTED_IN carrieanne, "Trinity", tt0242653, ACTED_IN
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies4-header.csv,import/movies4
-part1.csv,import/movies4-part2.csv --nodes=import/actors4-header.csv,import/actors4
-part1.csv,import/actors4-part2.csv --relationships=import/roles4-header.csv,import/roles4
-part1.csv,import/roles4-part2.csv
```

See also Using regular expressions for specifying multiple input files.

Using regular expressions for specifying multiple input files

File names can be specified using regular expressions in order to simplify using the command line when there are many data source files. Each file name that matches the regular expression will be included.

As mentioned in a previous section, for the import to work correctly, the header file must be first in the file group. When using regular expressions to specify the input files, the list of files will be sorted according to the names of the files that match the expression. The matching is aware of numbers inside the file names and will sort them accordingly, without the need for padding with zeros.

For example, let's assume that we have the following files:

- movies4-header.csv
- movies4-data1.csv
- movies4-data2.csv
- movies4-data12.csv

If we use the regular expression movies4.\*, the sorting will place the header file last and the import will fail. A better alternative would be to name the header file explicitly and use a regular expression that only matches the names of the data files. For example: --nodes "import/movies4-header.csv,movies-data.\*" will accomplish this.

Using the same data files as in the previous example, the call to neo4j-admin import can be simplified to:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies4-header.csv,import/movies4-part.*
--nodes=import/actors4-header.csv,import/actors4-part.* --relationships=import/roles4
-header.csv,import/roles4-part.*
```

The use of regular expressions should not be confused with file globbing.



The expression .\* means: "zero or more occurrences of any character except line break". Therefore, the regular expression movies4.\* will list all files starting with movies4. Conversely, with file globbing, 1s movies4.\* will list all files starting with movies4..

Another important difference to pay attention to is the sorting order. The result of a regular expression matching will place the file movies4-part2.csv before the file movies4-part12.csv. If doing 1s movies4-part\* in a directory containing the above listed files, the file movies4-part12.csv will be listed before the file movies4-part2.csv.

## B.2.5. Types and labels

Using the same label for every node

If you want to use the same node label(s) for every node in your nodes file you can do this by specifying the appropriate value as an option to neo4j-admin import. There is then no need to specify

the :LABEL field in the node file if you pass it as a command line option. If you do then both the label provided in the file and the one provided on the command line will be added to the node.

In this example we put the label Movie on every node specified in movies5a.csv, and we put the labels Movie and Sequel on the nodes specified in sequels5a.csv:

#### movies5a.csv

```
movieId:ID,title,year:int
tt0133093,"The Matrix",1999
```

#### sequels5a.csv

```
movieId:ID,title,year:int
tt0234215,"The Matrix Reloaded",2003
tt0242653,"The Matrix Revolutions",2003
```

#### actors5a.csv

```
personId:ID,name
keanu,"Keanu Reeves"
laurence,"Laurence Fishburne"
carrieanne,"Carrie-Anne Moss"
```

#### roles5a.csv

```
:START_ID,role,:END_ID,:TYPE
keanu,"Neo",tt0133093,ACTED_IN
keanu,"Neo",tt0234215,ACTED_IN
keanu,"Neo",tt0242653,ACTED_IN
laurence,"Morpheus",tt0133093,ACTED_IN
laurence,"Morpheus",tt0234215,ACTED_IN
laurence,"Morpheus",tt0234215,ACTED_IN
carrieanne,"Trinity",tt0133093,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=Movie=import/movies5a.csv
--nodes=Movie:Sequel=import/sequels5a.csv --nodes=Actor=import/actors5a.csv
--relationships=import/roles5a.csv
```

### Using the same relationship type for every relationship

If you want to use the same relationship type for every relationship in your relationships file this can be done by specifying the appropriate value as an option to neo4j-admin import. If you provide a relationship type both on the command line and in the relationships file, the one in the file will be applied. In this example we put the relationship type ACTED\_IN on every relationship specified in roles5b.csv:

#### movies5b.csv

```
movieId:ID,title,year:int,:LABEL
tt0133093,"The Matrix",1999,Movie
tt0234215,"The Matrix Reloaded",2003,Movie;Sequel
tt0242653,"The Matrix Revolutions",2003,Movie;Sequel
```

#### actors5b.csv

```
personId:ID,name,:LABEL
keanu,"Keanu Reeves",Actor
laurence,"Laurence Fishburne",Actor
carrieanne,"Carrie-Anne Moss",Actor
```

#### roles5b.csv

```
:START_ID,role,:END_ID
keanu,"Neo",tt0133093
keanu,"Neo",tt0234215
keanu,"Neo",tt0242653
laurence,"Morpheus",tt0133093
laurence,"Morpheus",tt0234215
laurence,"Morpheus",tt0242653
carrieanne,"Trinity",tt0133093
carrieanne,"Trinity",tt0234215
carrieanne,"Trinity",tt0234215
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies5b.csv --nodes=import/actors5b.csv
--relationships=ACTED_IN=import/roles5b.csv
```

## B.2.6. Property types

The type for properties specified in nodes and relationships files is defined in the header row. (see CSV file header format)

The following example creates a small graph containing one actor and one movie connected by an ACTED\_IN relationship. There is a roles property on the relationship which contains an array of the characters played by the actor in a movie:

#### movies6.csv

```
movieId:ID,title,year:int,:LABEL tt0099892,"Joe Versus the Volcano",1990,Movie
```

#### actors6.csv

```
personId:ID,name,:LABEL
meg,"Meg Ryan",Actor
```

#### roles6.csv

```
:START_ID,roles:string[],:END_ID,:TYPE
meg,"DeDe;Angelica Graynamore;Patricia Graynamore",tt0099892,ACTED_IN
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies6.csv --nodes=import/actors6.csv
--relationships=import/roles6.csv
```

## B.2.7. ID handling

A node ID is used to find the correct nodes when creating relationships. Each node processed by neo4j-admin import must provide an ID if it is to be connected in any relationships.

### Working with sequential or auto incrementing identifiers

The import tool makes the assumption that identifiers are unique across node files. This may not be the case for data sets which use sequential, auto incremented or otherwise colliding identifiers. Those data sets can define ID spaces where identifiers are unique within their respective ID space.

For example, if movies and people both use sequential identifiers then we would define Movie and Actor ID spaces:

#### movies7.csv

```
movieId:ID(Movie-ID),title,year:int,:LABEL
1,"The Matrix",1999,Movie
2,"The Matrix Reloaded",2003,Movie;Sequel
3,"The Matrix Revolutions",2003,Movie;Sequel
```

#### actors7.csv

```
personId:ID(Actor-ID),name,:LABEL
1,"Keanu Reeves",Actor
2,"Laurence Fishburne",Actor
3,"Carrie-Anne Moss",Actor
```

We also need to reference the appropriate ID space in our relationships file so it knows which nodes to connect together:

#### roles7.csv

```
:START_ID(Actor-ID),role,:END_ID(Movie-ID)

1,"Neo",1

1,"Neo",2

1,"Neo",3

2,"Morpheus",1

2,"Morpheus",2

2,"Morpheus",3

3,"Trinity",1

3,"Trinity",2

3,"Trinity",3
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies7.csv --nodes=import/actors7.csv
--relationships=ACTED_IN=import/roles7.csv
```

### B.2.8. Bad input data

The import tool has no tolerance for bad entities (relationships or nodes) and will fail the import on the first bad entity. You can specify explicitly that you want it to ignore rows that contain bad entities.

There are two different types of bad input: bad relationships and bad nodes. We will have a closer look at these in the following examples.

### Relationships referring to missing nodes

Relationships that refer to missing node IDs, either for :START\_ID or :END\_ID are considered bad relationships. Whether or not such relationships are skipped is controlled with --skip-bad -relationships flag which can have the values true or false or no value, which means true. The default is false, which means that any bad relationship is considered an error and will fail the import. For more information, see the --skip-bad-relationships option.

In the following example there is a missing emil node referenced in the roles file:

#### movies8a.csv

```
movieId:ID,title,year:int,:LABEL
tt0133093,"The Matrix",1999,Movie
tt0234215,"The Matrix Reloaded",2003,Movie;Sequel
tt0242653,"The Matrix Revolutions",2003,Movie;Sequel
```

#### actors8a.csv

```
personId:ID,name,:LABEL
keanu,"Keanu Reeves",Actor
laurence,"Laurence Fishburne",Actor
carrieanne,"Carrie-Anne Moss",Actor
```

#### roles8a.csv

```
:START_ID,role,:END_ID,:TYPE
keanu,"Neo",tt0133093,ACTED_IN
keanu,"Neo",tt0234215,ACTED_IN
keanu,"Neo",tt0242653,ACTED_IN
laurence,"Morpheus",tt0133093,ACTED_IN
laurence,"Morpheus",tt0234215,ACTED_IN
laurence,"Morpheus",tt0242653,ACTED_IN
carrieanne,"Trinity",tt0133093,ACTED_IN
carrieanne,"Trinity",tt0234215,ACTED_IN
carrieanne,"Trinity",tt0242653,ACTED_IN
carrieanne,"Trinity",tt0242653,ACTED_IN
emil,"Emil",tt0133093,ACTED_IN
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies8a.csv --nodes=import/actors8a.csv
--relationships=import/roles8a.csv
```

Since there was a bad relationship in the input data, the import process will fail completely.

Let's see what happens if we append the --skip-bad-relationships flag:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/movies8a.csv --nodes=import/actors8a.csv --relationships=import/roles8a.csv --skip-bad-relationships
```

The data files are successfully imported and the bad relationship is ignored. An entry is written to the *import.report* file.

#### ignore bad relationships

```
InputRelationship:
    source: roles8a.csv:11
    properties: [role, Emil]
    startNode: emil (global id space)
    endNode: tt0133093 (global id space)
    type: ACTED_IN
    referring to missing node emil
```

Multiple nodes with same ID within same ID space

Nodes that specify: ID which has already been specified within the ID space are considered bad nodes. Whether or not such nodes are skipped is controlled with --skip-duplicate-nodes flag which can have the values true or false or no value, which means true. The default is false, which means that any duplicate node is considered an error and will fail the import. For more information, see the

--skip-duplicate-nodes option.

In the following example there is a node ID that is specified twice within the same ID space:

#### actors8b.csv

```
personId:ID,name,:LABEL
keanu,"Keanu Reeves",Actor
laurence,"Laurence Fishburne",Actor
carrieanne,"Carrie-Anne Moss",Actor
laurence,"Laurence Harvey",Actor
```

The call to neo4j-admin import would look like this:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/actors8b.csv
```

Since there was a bad node in the input data, the import process will fail completely.

Let's see what happens if we append the --skip-duplicate-nodes flag:

```
neo4j_home$ bin/neo4j-admin import --nodes=import/actors8b.csv --skip-duplicate-nodes
```

The data files are successfully imported and the bad node is ignored. An entry is written to the *import.report* file.

### ignore bad nodes

ID 'laurence' is defined more than once in global ID space, at least at actors8b.csv:3 and actors8b.csv:5

# Appendix C: Advanced Causal Clustering

This appendix describes advanced features of Neo4j Causal Clustering.

This section includes information about advanced deployments and configuration options for multidata center operations.

- Causal Clustering lifecycle A walk-through of the lifecycle of a cluster.
- Multi-data center Overview of the multi-data center section.
  - ☐ Licensing for multi-data center operations Information about licensing for multi-data center operations.
  - ☐ Multi-data center design Patterns for multi-data center deployments.
  - ☐ Multi-data center operations Configuration options for multi-data center deployments.
  - ☐ Multi-data center load balancing Configuration options for making client applications aware of multi-data center topologies.
  - ☐ Data center disaster recovery How to recover a cluster to full working capability after data center loss.
- Advanced settings reference Advanced settings related to running a Neo4j Causal Cluster.

## C.1. Causal Clustering lifecycle

This section describes the lifecycle of a Neo4j Causal Cluster.

#### This section includes:

- Introduction
- Discovery protocol
- Core membership
- Read Replica membership
- Transacting via the Raft protocol
- Catchup protocol
- Read Replica shutdown
- Core shutdown

### C.1.1. Introduction

In this section we will develop some deeper knowledge of how the cluster operates. By developing our understanding of how the cluster works we will be better equipped to design, deploy, and troubleshoot our production systems.

Our in-depth tour will follow the lifecycle of a cluster. We will boot a Core cluster and pick up key architectural foundations as the cluster forms and transacts. We will then add in Read Replicas and show how they bootstrap join the cluster and then catchup and remain caught up with the Core Servers. We will then see how backup is used in live cluster environments before shutting down Read Replicas and Core Servers.

## C.1.2. Discovery protocol

The discovery protocol is the first step in forming a Causal Cluster. It takes in some information about existing *Core* cluster servers, and uses this to initiate a network join protocol.

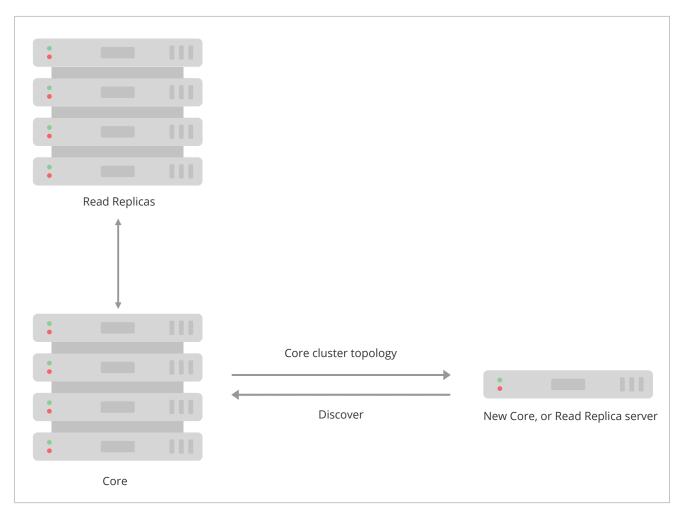


Figure 13. Causal Cluster discovery protocol: Core-to-Core or Read Replica-to-Core only.

Using this information, the server will either join an existing cluster or form one of its own.



The discovery protocol targets Core Servers only regardless of whether it is a Core Server or Read Replica performing discovery. It is because we expect Read Replicas to be both numerous and, relatively speaking, transient whereas Core Servers will likely be fewer in number and relatively stable over time.

The discovery protocol takes information from <code>causal\_clustering.initial\_discovery\_members</code> in <code>neo4j.conf</code>, which lists which IP addresses and ports that form the cluster on startup. Detailed information about discovery and discovery configuration options is given in the <code>Initial discovery</code> of <code>cluster members</code> section. When consuming this information, the server will try to handshake with the other listed servers. On successful handshake with another server (or servers), the current server will discover the whole current topology.

The discovery protocol continues to run throughout the lifetime of the Causal Cluster and is used to maintain the current state of available servers and to help clients route queries to an appropriate server via the client-side drivers.

### C.1.3. Core membership

If it is a Core Server that is performing discovery, once it has made a connection to the one of the

existing Core Servers, it then joins the Raft protocol.



Raft is a distributed algorithm for maintaining a consistent log across multiple shared-nothing servers designed by Diego Ongaro for his 2014 Ph.D. thesis. See the Raft thesis for details.

Raft handles cluster membership by making it a normal part of keeping a distributed log in sync. Joining a cluster involves the insertion of a cluster membership entry into the Raft log which is then reliably replicated around the existing cluster. Once that entry is applied to enough members of the Raft consensus group (those machines running the specific instance of the algorithm), they update their view of the cluster to include the new server. Thus membership changes benefit from the same safety properties as other data transacted via Raft (see Transacting via the Raft protocol for more information).

The new Core Server must also catch up its own Raft log with respect to the other Core Servers as it initializes its internal Raft instance. This is the normal case when a cluster is first booted and has performed few operations. There will be a delay before the new Core Server becomes available if it also needs to catch up (as per Catchup protocol) graph data from other servers. This is the normal case for a long lived cluster where the servers holds a great deal of graph data.

When an instance establishes a connection to any other instance, it determines the current state of the cluster and ensures that it is eligible to join. To be eligible the Neo4j instance must host the same database store as other members of the cluster (although it is allowed to be in an older, outdated, state), or be a new deployment without a database store.

### C.1.4. Read Replica membership

When a Read Replica performs discovery, once it has made a connection to any of the available Core clusters it proceeds to add itself into a shared whiteboard.

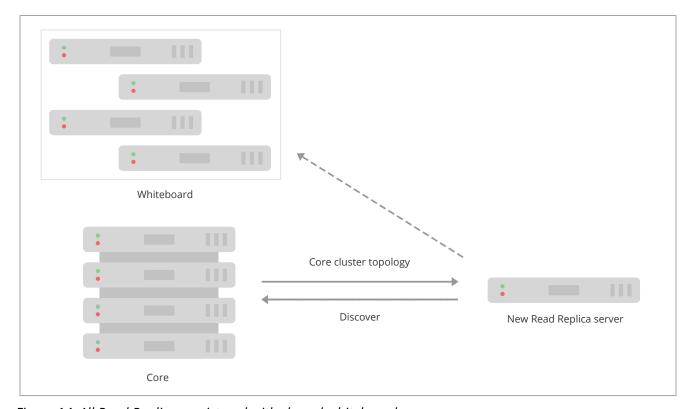


Figure 14. All Read Replicas registered with shared whiteboard.

This whiteboard provides a view of all live Read Replicas and is used both for routing requests from database drivers that support end-user applications and for monitoring the state of the cluster.

The Read Replicas are not involved in the Raft protocol, nor are they able to influence cluster topology. Hence a shared whiteboard outside of Raft comfortably scales to very large numbers of Read Replicas.

The whiteboard is kept up to date as Read Replicas join and leave the cluster, even if they fail abruptly rather than leaving gracefully.

### C.1.5. Transacting via the Raft protocol

Once bootstrapped, each Core Server spends its time processing database transactions. Updates are reliably replicated around Core Servers via the Raft protocol. Updates appear in the form of a (committed) Raft log entry containing transaction commands which is subsequently applied to update the database.



One of Raft's primary design goals is to be easily understandable so that there are fewer places for tricky bugs to hide in implementations. As a side-effect, it is also possible for database operators to reason about their Core Servers in their Causal Clusters.

The Raft Leader for the current term (a logical clock) appends the transaction (an 'entry' in Raft terminology) to the head of its local log and asks the other instances to do the same. When the Leader can see that a majority instances have appended the entry, it can be considered committed into the Raft log. The client application can now be informed that the transaction has safely committed since there is sufficient redundancy in the system to tolerate any (non-pathological) faults.



The Raft protocol describes three roles that an instance can be playing: *Leader*, *Follower*, and *Candidate*. These are transient roles and any Core Server can expect to play them throughout the lifetime of a cluster. While it is interesting from a computing science point of view to understand those states, operators should not be overly concerned: they are an implementation detail.

For safety, within any Raft protocol instance there is only one Leader able to make forward progress in any given term. The Leader bears the responsibility for imposing order on Raft log entries and driving the log forward with respect to the *Followers*.

Followers maintain their logs with respect to the current Leader's log. Should any participant in the cluster suspect that the Leader has failed, then they can instigate a leadership election by entering the *Candidate* state. In Neo4j Core Servers this is happens at ms timescale, around 500ms by default.

Whichever instance is in the best state (including the existing Leader, if it remains available) can emerge from the election as Leader. The "best state" for a Leader is decided by highest term, then by longest log, then by highest committed entry.

The ability to fail over roles without losing data allows forward progress even in the event of faults. Even where Raft instances fail, the protocol can rapidly piece together which of the remaining instances is best placed to take over from the failed instance (or instances) without data loss. This is the essence of a *non-blocking* consensus protocol which allows Neo4j Causal Clustering to provide continuous availability to applications.

# C.1.6. Catchup protocol

Read Replicas spend their time concurrently processing graph queries and applying a stream of transactions from the Core Servers to update their local graph store.

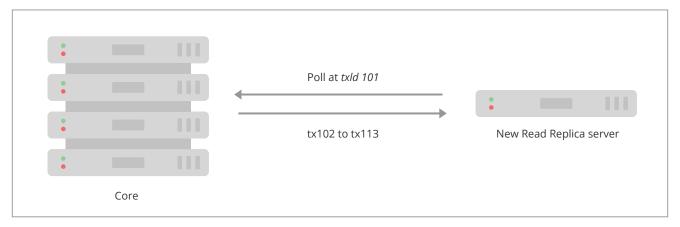


Figure 15. Transactions shipped from Core to Read Replica.

Updates from Core Servers to Read Replicas are propagated by transaction shipping. Transaction shipping is instigated by Read Replicas frequently *polling* any of the Core Servers specifying the ID of the last transaction they received and processed. The frequency of polling is an operational choice.



Neo4j transaction IDs are strictly monotonic integer values (they always increase). This makes it possible to determine whether or not a transaction has been applied to a Read Replica by comparing its last processed transaction ID with that of a Core Server.

If there is a large difference between an Read Replica's transaction history and that of a Core Server, polling may not result in any transactions being shipped. This is quite expected, for example when a new Read Replica is introduced to a long-running cluster or where a Read Replica has been down for some significant period of time. In such cases the catchup protocol will realize the gap between the Core Servers and Read Replica is too large to fill via transaction shipping and will fall back to copying the database store directly from Core Server to Read Replica. Since we are working with a live system, at the end of the database store copy the Core Server's database is likely to have changed. The Read Replica completes the catchup by asking for any transactions missed during the copy operation before becoming available.



A very slow database store copy could conceivably leave the Read Replica too far behind to catch up via transaction log shipping as the Core Server has substantially moved on. In such cases the Read Replica server repeats the catchup protocol. In pathological cases the operator can intervene to snapshot, restore, or file copy recent store files from a fast backup.

# C.1.7. Read Replica shutdown

On clean shutdown, a Read Replica will invoke the discovery protocol to remove itself from the shared whiteboard overview of the cluster. It will also ensure that the database is cleanly shutdown and consistent, immediately ready for future use.

On an unclean shutdown such as a power outage, the Core Servers maintaining the overview of the cluster will notice that the Read Replica's connection has been abruptly been cut. The discovery machinery will initially hide the Read Replica's whiteboard entry, and if the Read Replica does not reappear quickly its modest memory use in the shared whiteboard will be reclaimed.

On unclean shutdown it is possible the Read Replica will not have entirely consistent store files or transaction logs. On subsequent reboot the Read Replica will rollback any partially applied transactions such that the database is in a consistent state.

### C.1.8. Core shutdown

A clean Core Server shutdown, like Core Server booting, is handled via the Raft protocol. When a Core Server is shut down, it appends a membership entry to the Raft log which is then replicated around the Core Servers. Once a majority of Core Servers have committed that membership entry the leaver has logically left the cluster and can safely shut down. All remaining instances accept that the cluster has grown smaller, and is therefore less fault tolerant. If the leaver happened to be playing the Leader role at the point of leaving, it will be transitioned to another Core Server after a brief election.

An unclean shutdown does not directly inform the cluster that a Core Server has left. Instead the Core cluster size remains the same for purposes of computing majorities for commits. Thus an unclean shutdown in a cluster of 5 Core Servers now requires 3/4 members to agree to commit which is a tighter margin than 3/5 before the unclean shutdown.



Of course when Core Servers fail, operators or monitoring scripts can be alerted so that they can intervene in the cluster if necessary.

If the leaver was playing the Leader role, there will be a brief election to produce a new Leader. Once the new Leader is established, the Core cluster continues albeit with less redundancy. However even with this failure, a Core cluster of 5 servers reduced to 4 can still tolerate one more fault before becoming read-only.

### C.2. Multi-data center

This section introduces the multi-data center functionality in Neo4j.

Some use cases present high needs for availability, redundancy, locality of client applications, or simply scale. In these cases it is important that the cluster is aware of its physical topology so that it can optimize for workload. This makes configuring a single cluster to span multiple data centers a necessary proposition.

The following sections are dedicated to describing the different aspects of multi-data center operations of a Causal Cluster.

- Licensing for multi-data center operations
- Multi-data center design
  - □ Introduction
  - ☐ Core Server deployment scenarios
  - ☐ Allowing Read Replicas to catch up from other Read Replicas
- Multi-data center operations
  - ☐ Enable multi-data center operations
  - Server groups
  - Strategy plugins
- Multi-data center load balancing
  - □ Introduction
  - ☐ Prerequisite configuration
  - ☐ The load balancing framework
  - □ Load balancing examples
- Data center disaster recovery

□ Data center loss scenario
☐ Procedure for recovering from data center loss
C.2.1. Licensing for multi-data center operations
Multi-data center functionality is intended for very demanding users of Neo4j who typically operate under a commercial database license. As a result, multi-data center functionality is licensed separately from the single-data center Causal Clustering features.
In order to confirm that you are operating under a suitable license, you must explicitly set the following in <i>neo4j.conf</i> :
causal_clustering.multi_dc_license=true
Without this configuration, all of the multi-data center features will remain disabled.
C.2.2. Multi-data center design
This section describes common patterns for multi-data center deployments that can act as building blocks for your own multi-data center production environments.
This section describes the following:
• Introduction
Core Server deployment scenarios
<ul> <li>Allowing Read Replicas to catch up from other Read Replicas</li> </ul>
☐ Hierarchical Read Replica deployment
☐ Catch up (mostly) from peer Read Replicas
☐ Maintaining causal consistency in scale-out topologies
Introduction
This section is based on a series of examples to illustrate the different considerations we should take into account when designing our Causal Cluster for a multi-data center environment. We'll come to understand the weaknesses and benefits of common multi-data center deployment scenarios. Each scenario is presented at a high architectural level for clarity. In subsequent sections we will go into more detail on how such deployments are configured.
Core Server deployment scenarios
We will start with the conceptually simplest multi-data center scenario where we deploy the same number and kind of instances into each DC. This is a <i>homogeneous</i> deployment because each data center is identical to the other.
Example 112. Homogeneous three data center deployment

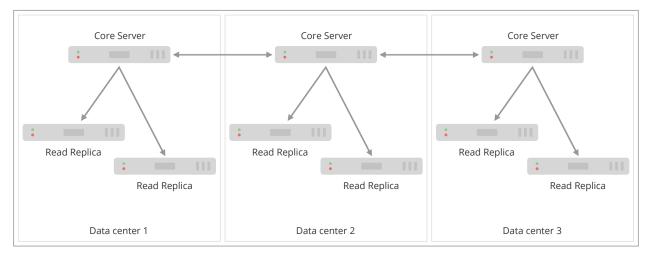


Figure 16. Homogeneous deployment across three data centers with one Core instance in each

In diagram above we have three data centers, each identically equipped with a single Core Server and a small number of Read Replicas.

Since Raft only requires a majority of the instances to acknowledge a write before it is safely committed, the latency of the commit path for this pattern involves only the two fastest data centers. As such the cost of committing to this setup is two WAN messages: one to send the transaction and one ACK message. In a non-failure case the other data center will not be far behind and will apply the transaction as well.

Within each of the data centers we can increase machine-level redundancy by adding more Core instances. For example we could add two more machines in each data center so that we can tolerate the spontaneous loss of up to four machines anywhere in the cluster or a single data center as a whole.

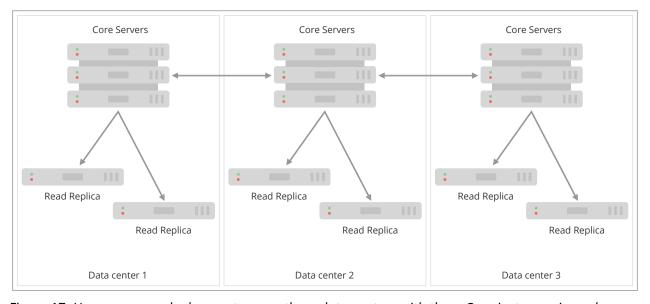


Figure 17. Homogeneous deployment across three data centers with three Core instances in each

To recap the strengths and weaknesses of this deployment pattern:

- We can lose an entire data center without losing availability and, depending on the number of machines in each data center, we may still be able to tolerate the loss of individual servers regardless of which data center they are in.
- The commit path for transactions is short, just two WAN messages exchanged.
- While the loss of majority data centers will need to be recovered, the operational procedure is identical irrespective of which of the data centers are lost.

As will be shown in the section on multi-data center configuration the Read Replicas can be biased to catchup from their data center-local Core Servers to minimize catchup latency. Data center-local client applications would also likely be routed to those same Read Replicas both for topological locality and scaling. More details are available in the section on multi-data center load balancing.

In the two data center case, our first instinct is to balance the available servers for operational consistency. An example of a homogeneous deployment across two data centers with two Core instances in each is illustrated in the diagram below:

Example 113. Homogeneous two data center deployment

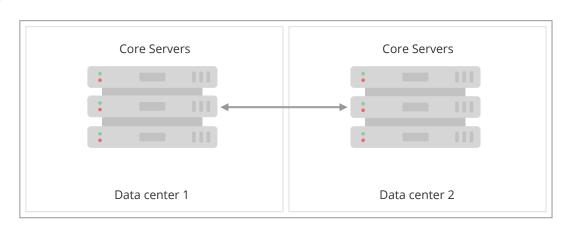


Figure 18. Homogeneous deployment across two data centers

The problem with this configuration is that while architecturally simple, it does not play to the strengths of the Raft protocol which is based on majority consensus. In the non-failure case, we incur two WAN messages to commit any transaction because a majority commit implies at least one response from the non-local data center instances. Worse, if we lose either data center the cluster will become read-only because it is impossible to achieve a majority.

As seen in the example above, the homogeneous deployment over two data centers does not take full advantage of the strengths of Causal Clustering. However it guarantees that the full Raft log will be present in either data center in the case of total data center loss.

The opposite of spreading Core Servers around our data centers, is to have them all hosted in a single one. This may be for technical or governance reasons, but either way has the advantage of LAN commit latencies for writes.

While our Core Servers are colocated, we spread out our Read Replicas close to the client applications to enable fan-out scaling.

The diagram below shows an example of a heterogeneous deployment directing writes to one data center, and reads to all. This pattern provides high survivability for data because of georeplication. It also provides locality for client applications. However, if the Core Server data center is lost, we must immediately instigate recovery and turn one of the remaining Read Replica data centers into a new Core cluster.

It is possible that none of the Read Replicas have received all of the confirmed transactions prior to losing Data Center 1. While this is a convenient pattern for geo-replication, its semantics are best-effort. Cluster designers must take this aspect under consideration when deciding on recovery strategy.

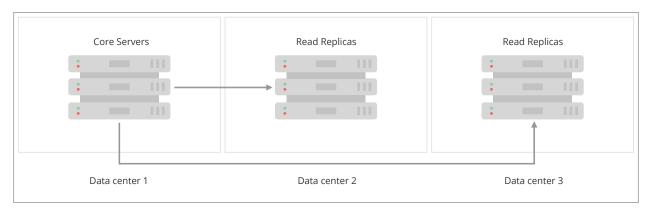


Figure 19. Heterogeneous deployment separating Read Replicas from the Core cluster

An operational tweak to this approach would be to host a Core Server in Data Center 2 and 3 as the starting point for recovery. During normal operations, these extra Core Servers should be configured with causal\_clustering.refuse\_to\_be\_leader=true. Should we lose Data Center 1, then we can use one of these Core Servers to quickly bootstrap a new Core cluster and return to full service rapidly.

To recap the strengths of this deployment pattern:

- Core Servers commit at LAN latencies if using the setup with Core Servers exclusively in one data center.
- Read Replicas provide scale and locality for client applications.
- Geo-replication provides high survivability for data.

### Allowing Read Replicas to catch up from other Read Replicas

With an understanding of the basic multi-data center patterns at our disposal, we can refine our deployment models to embrace local catchup within data centers. This means that any server, including Read Replicas, can act as a source of transactions for Read Replica server. When catching up from data center-local instances we aim to amortize the cost of WAN traffic catchup across many local replications.

Allowing Read Replicas to choose a data center-local Core Server or even another Read Replica gives us a great deal of design freedom, and importantly allows us to scale to truly huge numbers of Read Replicas. Using this feature we might choose to fan-out Read Replicas so that the catchup load on the Core Servers grows (approximately) logarithmically rather than linearly.

### Hierarchical Read Replica deployment

The primary motivation for Read Replicas catching up from other Read Replicas is to allow for fan-out scale. To achieve a fan-out we arrange the Read Replicas in a hierarchy, with each layer of the

hierarchy being broader than the one above.

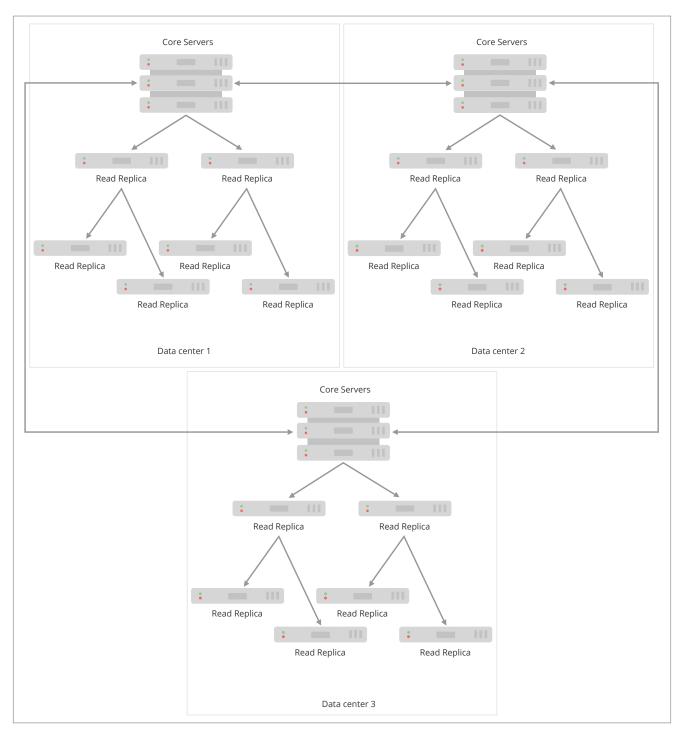


Figure 20. Fan out from Core Servers for scale at log cost

An illustrative hierarchy is presented in the diagram above. The Core Servers supply transactions to a relatively small number of Read Replicas at the first tier. This results in a relatively modest load on the Core Servers, freeing up resources to focus on the commit path. Those Read Replicas in the first tier in turn feed a larger number of Read Replicas in the second tier. This pattern can be reasonably extended to several tiers to provide enormous fan-out.

At each tier we expand the scalability of the Read Replicas, but we add another level of catchup latency. By careful measurement we can ascertain the appropriate depth and breadth of the hierarchy to match the application requirements.

We should also take care that each tier in the hierarchy has sufficient redundancy so that failures do not compromise transmission of data from the Core Servers. A strategy for keeping Read Replicas current in the presence of failures is to occasionally have them subvert the hierarchy. That is, if a

given Read Replica occasionally goes to its grandparents or even directly to the Core Servers then we can avoid pathologically high replication latencies under fault conditions.

#### Catch up (mostly) from peer Read Replicas

Another strategy for Read Replica catchup is to treat them all as peers and have peer-to-peer catchup. This avoids the need to manage tiers of replicas to maintain availability since the Read Replicas catch up from one another in a mesh.

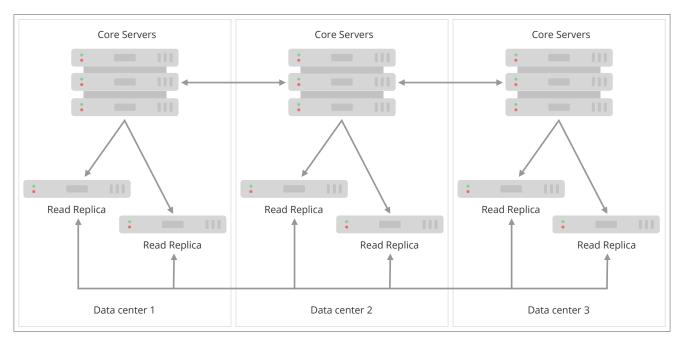


Figure 21. Peer-to-peer Read Replica catchup

Having a reduced load on the Core Servers allows us to scale out. For example if only one in ten catchup requests goes to the Core Servers, we could expand the number of Read Replicas by approximately a factor of 10.

To avoid groups of orphans in the mesh, Read Replicas will occasionally catch up directly from Core Servers. Having Read Replicas catch up with Core Servers ensures that no Read Replica is left behind indefinitely, placing an upper bound on replication latency. While this places some load on the Core Servers, it is far less than if all catch up attempts from Read Replicas were directed to a Core Server.

The upper bound on replication latency for this mode of operation is the number of catchup attempts served by Read Replicas before trying core. The average replication latency will be half the number of attempts to replicate. This is because on average half the Read Replicas will be ahead and half behind any given Read Replica.



Connecting to a random Core Server on failure to retrieve updates from other sources is the default behavior of Read Replicas.

#### Maintaining causal consistency in scale-out topologies

Causal consistency is always maintained, even in extreme situations with chains of Read Replicas catching up from other upstream Read Replicas. The key trade-off to understand, as so often in distributed systems, is that of latency for scale.

In Fan out from Core Servers for scale at log cost we see that number of hops required for a transaction to propagate to the lowest tier is 2: the highest latency in this topology. Equally we see how the bottommost tier has far more members than any other tier giving it scale advantages.

Correspondingly, in the middle tier we have better latency (one hop) but less scale. At the top most

tier (Core Servers) we have very little latency (just the Raft commit path) but the fewest available servers. This means we should target queries at the most appropriate tier based on latency, scale, and locality.

#### Summary on latency versus scalability:

- Issuing read queries to a Core Server generally has the lowest latency in principle but may have the highest contention.
- Issuing read queries to a Read Replica topologically closest to Core Servers typically has higher latency but also higher scalability.
- Issuing read queries to a Read Replica topologically further from Core Servers typically has the highest latency but also the highest scalability.

In large systems like the scale-out hierarchy above, we are conventionally used to having relaxed or *eventual* consistency semantics. With Neo4j multi-data center setups, that is also possible. Where we don't care about causality we can read from any Read Replica and accept that we might see older values. However the causal consistency semantics are maintained.

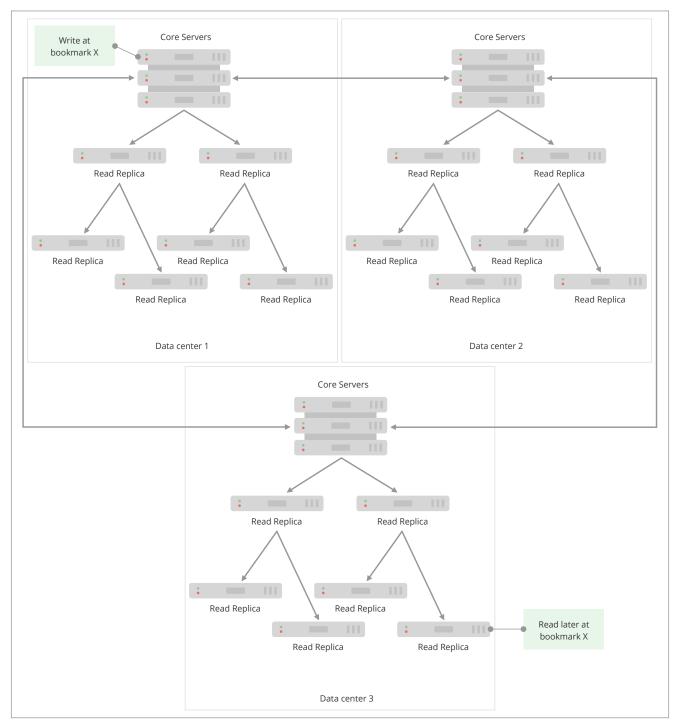


Figure 22. Each tier in the Read Replicas is further behind the source of truth, but offers greater scale-out

As we can see in diagram above, even if the client binds to a Read Replica that is multiple hops/data centers away from the source of truth, causal consistency is maintained. While the query may be suspended while the necessary transaction propagates to the Read Replica, the benefit is that there will be more Read Replicas available and so overall client throughput is higher than with a single-tier configuration.

# C.2.3. Multi-data center operations

This section shows how to configure Neo4j servers so that they are topology/data center-aware. It describes the precise configuration needed to achieve a scalable multi-data center deployment.

This section describes the following:

- Enable multi-data center operations
- Server groups
- Strategy plugins
  - ☐ Configuring upstream selection strategy using pre-defined strategies
  - □ Configuring user-defined strategies
  - ☐ Building upstream strategy plugins using Java
  - □ Favoring data centers

### Enable multi-data center operations

Before doing anything else, we must enable the multi-data center functionality. This is described in Licensing for multi-data center operations.



Licensing for multi-data center

The multi-data center functionality is separately licensed and must be specifically enabled.

### Server groups

In order to optimize the use of our Causal Cluster servers according to our specific requirements, we sort them into *Server Groups*. Server Group membership can map to data centers, availability zones, or any other significant topological elements from the operator's domain. Server Groups can also overlap.

Server Groups are defined as a key that maps onto a set of servers in a Causal Cluster. Server Group membership is defined on each server using the causal\_clustering.server\_groups parameter in neo4j.conf. Each server in a Causal Cluster can belong to zero or more server groups.

#### Example 115. Definition of Server Group membership

The membership of a server group or groups can be set in *neo4j.conf* as in the following examples:

```
# Add the current instance to the groups `us` and `us-east`
causal_clustering.server_groups=us,us-east

# Add the current instance into the group `london`
causal_clustering.server_groups=london

# Add the current instance into the group `eu`
causal_clustering.server_groups=eu
```

We must be aware that membership of each server group is explicit. For example, a server in the gb-london group is not automatically part of some gb or eu group unless that server is explicitly added to those groups. That is, any (implied) relationship between groups is reified only when those groups are used as the basis for requesting data from upstream systems.

Server Groups are not mandatory, but unless they are present, we cannot set up specific upstream transaction dependencies for servers. In the absence of any specified server groups, the cluster

defaults to its most pessimistic fall-back behavior: each Read Replica will catch up from a random Core Server.

### Strategy plugins

Strategy plugins are sets of rules that define how Read Replicas contact servers in the cluster in order to synchronize transaction logs. Neo4j comes with a set of pre-defined strategies, and also provides a Design Specific Language, DSL, to flexibly create user-defined strategies. Finally, Neo4j supports an API which advanced users may use to enhance upstream recommendations.

Once a strategy plugin resolves a satisfactory upstream server, it is used for pulling transactions to update the local Read Replica for a single synchronization. For subsequent updates, the procedure is repeated so that the most preferred available upstream server is always resolved.

Configuring upstream selection strategy using pre-defined strategies

Neo4j ships with the following pre-defined strategy plugins. These provide coarse-grained algorithms for choosing an upstream instance:

Plugin name	Resulting behavior
connect-to-random-core-server	Connect to any <b>Core Server</b> selecting at random from those currently available.
typically-connect-to-random-read-replica	Connect to any available <b>Read Replica</b> , but around 10% of the time connect to any random Core Server.
connect-randomly-to-server-group	Connect at random to any available <b>Read Replica</b> in any of the server groups specified in the comma-separated list causal_clustering.connect-randomly-to-server-group.
leader-only	Connect only to the current Raft leader of the Core Servers.
connect-randomly-within-server-group	Connect at random to any available <b>Read Replica</b> in any of the server groups to which this server belongs. Deprecated, please use connect-randomly-to-servergroup.

Pre-defined strategies are used by configuring the <code>causal\_clustering.upstream\_selection\_strategy</code> option. Doing so allows us to specify an ordered preference of strategies to resolve an upstream provider of transaction data. We provide a comma-separated list of strategy plugin names with preferred strategies earlier in that list. The upstream strategy is chosen by asking each of the strategies in list-order whether they can provide an upstream server from which transactions can be pulled.

Consider the following configuration example:

 ${\bf causal\_clustering.upstream\_selection\_strategy=} connect-randomly-to-server-group, typically-connect-to-random-read-replica$ 

With this configuration the instance will first try to connect to any other instance in the group(s) specified in config\_causal\_clustering.connect-randomly-to-server-group. Should we fail to find any live instances in those groups, then we will connect to a random Read Replica.

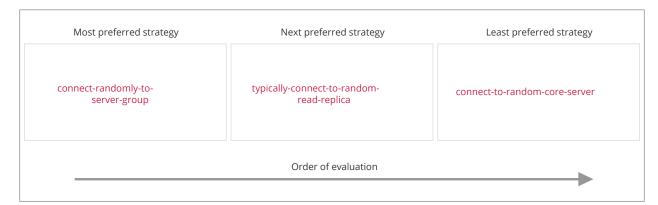


Figure 23. The first satisfactory response from a strategy will be used.

To ensure that downstream servers can still access live data in the event of upstream failures, the last resort of any instance is always to contact a random Core Server. This is equivalent to ending the causal\_clustering.upstream\_selection\_strategy configuration with connect-to-random-core-server.

#### Configuring user-defined strategies

Neo4j Causal Clusters support a small DSL for the configuration of client-cluster load balancing. This is described in detail in Policy definitions and Filters. The same DSL is used to describe preferences for how an instance binds to another instance to request transaction updates.

The DSL is made available by selecting the user-defined strategy as follows:

```
causal_clustering.upstream_selection_strategy=user-defined
```

Once the user-defined strategy has been specified, we can add configuration to the causal\_clustering.user\_defined\_upstream\_strategy setting based on the server groups that have been set for the cluster.

We will describe this functionality with two examples:

For illustrative purposes we propose four regions: north, south, east, and west and within each region we have a number of data centers such as north1 or west2. We configure our server groups so that each data center maps to its own server group. Additionally we will assume that each data center fails independently from the others and that a region can act as a supergroup of its constituent data centers. So an instance in the north region might have configuration like causal\_clustering.server\_groups=north2, north which puts it in two groups that match to our physical topology as shown in the diagram below.



Figure 24. Mapping regions and data centers onto server groups

Once we have our server groups, our next task is to define some upstream selection rules based on them. For our design purposes, let's say that any instance in one of the north region data centers prefers to catchup within the data center if it can, but will resort to any northern instance otherwise. To configure that behavior we add:

```
causal_clustering.user_defined_upstream_strategy=groups(north2); groups(north); halt()
```

The configuration is in precedence order from left to right. The <code>groups()</code> operator yields a server group from which to catch up. In this case only if there are no servers in the <code>north2</code> server group will we proceed to the <code>groups(north)</code> rule which yields any server in the <code>north</code> server group. Finally, if we cannot resolve any servers in any of the previous groups, then we will stop the rule chain via <code>halt()</code>.

Note that the use of halt() will end the rule chain explicitly. If we don't use halt() at the end of the rule chain, then the all() rule is implicitly added. all() is expansive: it offers up all servers and so increases the likelihood of finding an available upstream server. However all() is indiscriminate and the servers it offers are not guaranteed to be topologically or geographically local, potentially increasing the latency of synchronization.

The example above shows a simple hierarchy of preferences. But we can be more sophisticated if we so choose. For example we can place conditions on the server groups from which we catch up.

In this example we wish to roughly qualify cluster health before choosing from where to catch up. For this we use the min() filter as follows:

```
causal_clustering.user_defined_upstream_strategy=groups(north2)->min(3), groups(north)->min(3);
all();
```

groups(north2)->min(3) states that we want to catch up from the north2 server group if it has three available machines, which we here take as an indicator of good health. If north2 can't meet that requirement (is not healthy enough) then we try to catch up from any server across the north region provided there are at least three of them available as per groups(north)->min(3). Finally, if we cannot catch up from a sufficiently healthy north region, then we'll (explicitly) fall back to the whole cluster with all().

The min() filter is a simple but reasonable indicator of server group health.

#### Building upstream strategy plugins using Java

Neo4j supports an API which advanced users may use to enhance upstream recommendations in arbitrary ways: load, subnet, machine size, or anything else accessible from the JVM. In such cases we are invited to build our own implementations of

org.neo4j.causalclustering.readreplica.UpstreamDatabaseSelectionStrategy to suit our own needs, and register them with the strategy selection pipeline just like the pre-packaged plugins.

#### We have to override the

org.neo4j.causalclustering.readreplica.UpstreamDatabaseSelectionStrategy#upstreamDatabase() method in our code. Overriding that class gives us access to the following items:

Resource	Description
org.neo4j.causalclustering.discovery.TopologyService	This is a directory service which provides access to the addresses of all servers and server groups in the cluster.
org.neo4j.kernel.configuration.Config	This provides the configuration from <i>neo4j.conf</i> for the local instance. Configuration for our own plugin can reside here.
org.neo4j.causalclustering.identity.MemberId	This provides the unique cluster MemberId of the current instance.

Once our code is written and tested, we have to prepare it for deployment.

UpstreamDatabaseSelectionStrategy plugins are loaded via the Java Service Loader. This means when we package our code into a jar file, we'll have to create a file META-

*INF.services/org.neo4j.causalclustering.readreplica.UpstreamDatabaseSelectionStrategy* in which we write the fully qualified class name(s) of the plugins, e.g. org.example.myplugins.PreferServersWithHighIOPS.

To deploy this jar into the Neo4j server we copy it into the *plugins* directory and restart the instance.

#### Favoring data centers

In a multi-DC scenario, while it remains a rare occurrence, it is possible to bias which data centers are used to host Raft leaders (and thus where writes are directed). To do so, we can apply causal\_clustering.refuse\_to\_be\_leader=true on the leaders in the data centers where we do not want leaders to materialize. In doing so we implicitly prefer the instances where we have **not** applied that setting.

This may be useful when planning for highly distributed multi-data center deployments. However this must be very carefully considered because in failure scenarios it limits the availability of the cluster. It

is advisable to engage Neo4j Professional Services to help design a suitably resilient topology.

### C.2.4. Multi-data center load balancing

This section describes the topology-aware load balancing options available for client applications in a multi-data center Neo4j deployment. It describes how to configure the load balancing for the cluster so that client applications can direct its workload at the most appropriate cluster members, such as those nearby.

#### This section describes the following:

- Introduction
- Prerequisite configuration
  - ☐ Enable multi-data center operations
  - □ Server groups
  - Cores for reading
- The load balancing framework
  - Policy definitions
  - Policy names
  - ☐ Filters
- Load balancing examples

### Enabling load balancing



The load balancing functionality is part of the separately licensed multi-data center package and must be specifically enabled. See Licensing for multi-data center operations for details.

#### Introduction

When deploying a multi-data center cluster we often wish to take advantage of locality to reduce latency and improve performance. For example, we would like our graph-intensive workloads to be executed in the local data center at LAN latencies rather than in a faraway data center at WAN latencies. Neo4j's enhanced load balancing for multi-data center scenarios facilitates precisely this and can also be used to define fall-back behaviors. This means that failures can be planned for upfront and persistent overload conditions be avoided.

The load balancing system is a cooperative system where the driver asks the cluster on a recurring basis where it should direct the different classes of its workload (e.g. writes and reads). This allows the driver to work independently for long stretches of time, yet check back from time to time to adapt to changes like for example a new server having been added for increased capacity. There are also failure situations where the driver will ask again immediately, for example when it cannot use any of its allocated servers.

This is mostly transparent from the perspective of a client. On the server side we configure the load balancing behaviors and expose them under a named *load balancing policy* which the driver can bind to. All server-side configuration is performed on the Core Servers.

Use load balancing from Neo4j drivers



This chapter describes how to configure a Causal Cluster to use custom load balancing policies Once enabled and configured, the custom load balancing feature is used by drivers to route traffic as intended. See the Driver Manual for instructions on how to configure drivers to use custom load balancing.

# Prerequisite configuration

#### Enable multi-data center operations

In order to configure a cluster for load balancing we must enable the multi-data center functionality. This is described in Licensing for multi-data center operations.

#### Server groups

In common with server-to-server catchup, load balancing across multiple data centers is predicated on the *server group* concept. Servers can belong to one or more potentially overlapping server groups, and decisions about where to route requests from client to cluster member are parameterized based on that configuration. For details on server group configuration, refer to Server groups.

#### Cores for reading

Depending on the deployment and the available number of servers in the cluster different strategies make sense for whether or not the reading workload should be routed to the Core Servers. The following configuration will allow the routing of read workload to Core Servers. Valid values are true and false.

causal\_clustering.cluster\_allow\_reads\_on\_followers=true

### The load balancing framework

The load balancing system is based on a plugin architecture for future extensibility and for allowing user customizations. The current version ships with exactly one such canned plugin called the *server policies* plugin.

The server policies plugin is selected by setting the following property:

causal\_clustering.load\_balancing.plugin=server\_policies

Under the server policies plugin, a number of load balancing policies can be configured server-side and be exposed to drivers under unique names. The drivers, in turn, must on instantiation select an appropriate policy by specifying its name. Common patterns for naming policies are after geographical regions or intended application groups.

It is of crucial importance to define the exact same policies on all core machines since this is to be regarded as cluster-wide configuration and failure to do so will lead to surprising behavior. Similarly, policies which are in active use should not be removed or renamed since it will break applications trying to use these policies. It is perfectly acceptable and expected however that policies be modified under the same name.

If a driver asks for a policy name which is not available, then it will not be able to use the cluster. A driver which does not specify any name at all will get the behavior of the default policy as configured. The default policy, if left unchanged, distributes the load across all servers. It is possible to change the default policy to any behavior that a named policy can have.

A misconfigured driver or load balancing policy will result in suboptimal routing choices or even prevent successful interactions with the cluster entirely.



The details of how to write a custom plugin is not documented here. Please get in contact with Neo4j Professional Services if you think that you need a custom plugin.

#### Policy definitions

The configuration of load balancing policies is transparent to client applications and expressed via a simple DSL. The syntax consists of a set of rules which are considered in order. The first rule to produce a non-empty result will be the final result.

```
rule1; rule2; rule3
```

Each rule in turn consists of a set of filters which limit the considered servers, starting with the complete set. Note that the evaluation of each rule starts fresh with the complete set of available servers.

There is a fixed set of filters which compose a rule and they are chained together using arrows

```
filter1 -> filter2 -> filter3
```

If there are any servers still left after the last filter then the rule evaluation has produced a result and this will be returned to the driver. However, if there are no servers left then the next rule will be considered. If no rule is able to produce a usable result then the driver will be signalled a failure.

#### Policy names

The policies are configured under the namespace of the *server policies* plugin and named as desired. Policy names can contain alphanumeric characters and underscores, and they are case sensitive. Below is the property key for a policy with the name mypolicy.

```
causal_clustering.load_balancing.config.server_policies.mypolicy=
```

The actual policy is defined in the value part using the DSL.

The default policy name is reserved for the default policy. It is possible to configure this policy like any other and it will be used by driver clients which do not specify a policy.

Additionally, any number of policies can be created using unique policy names. The policy name can suggest a particular region or an application for which it is intended to be used.

#### Filters

There are four filters available for specifying rules, detailed below. The syntax is similar to a method call with parameters.

- groups(name1, name2, …)
  - $\ \square$  Only servers which are part of any of the specified groups will pass the filter.
  - ☐ The defined names must match those of the *server groups*.
- min(count)
  - ☐ Only the minimum amount of servers will be allowed to pass (or none).
  - ☐ Allows overload conditions to be managed.

- all()
  - ☐ No need to specify since it is implicit at the beginning of each rule.
  - ☐ Implicitly the last rule (override this behavior using halt).
- halt()
  - ☐ Only makes sense as the last filter in the last rule.
  - ☐ Will stop the processing of any more rules.

The groups filter is essentially an OR-filter, e.g. groups(A,B) which will pass any server in either A, B or both (the union of the server groups). An AND-filter can also be created by chaining two filters as in groups(A) -> groups(B), which will only pass servers in both groups (the intersect of the server groups).

### Load balancing examples

In our discussion on multi-data center clusters we introduced a four region, multi-data center setup. We used the cardinal compass points for regions and numbered data centers within those regions. We'll use the same hypothetical setup here too.

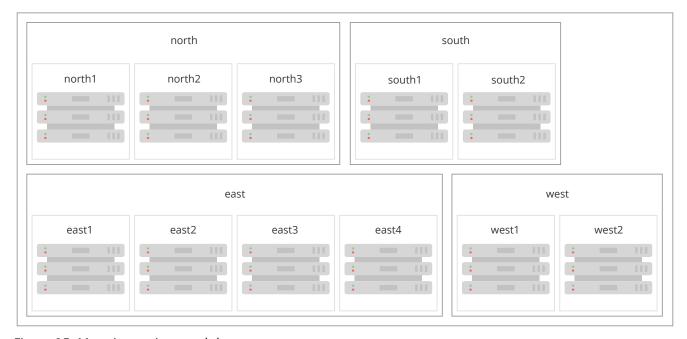


Figure 25. Mapping regions and data centers onto server groups

We configure the behavior of the load balancer in the property causal\_clustering.load\_balancing.config.server\_policies.<policy-name>. The rules we specify will allow us to fine tune how the cluster routes requests under load.

In the examples we will make use of the line continuation character \ for better readability. It is valid syntax in *neo4j.conf* as well and it is recommended to break up complicated rule definitions using this and a new rule on every line.

The most restrictive strategy would be to insist on a particular data center to the exclusion of all others:

```
causal_clustering.load_balancing.config.server_policies.north1_only=\
groups(north1)->min(2); halt();
```

In this case we're stating that we are only interested in sending queries to servers in the north1 server group, which maps onto a specific physical data center, provided there are two of them available. If we cannot provide at least two servers in north1 then we should halt(), i.e. not try any other data center.

While the previous example demonstrates the basic form of our load balancing rules, we can be a little more expansive:

Example 120. Specific data center preferably

```
causal_clustering.load_balancing.config.server_policies.north1=\
groups(north1)->min(2);
```

In this case if at least two servers are available in the <a href="north1">north1</a> data center then we will load balance across them. Otherwise we will use any server in the whole cluster, falling back to the implicit, final <a href="all()">all()</a> rule.

The previous example considered only a single data center before resorting to the whole cluster. If we have a hierarchy or region concept exposed through our server groups we can make the fall back more graceful:

Example 121. Gracefully falling back to neighbors

```
causal_clustering.load_balancing.config.server_policies.north_app1=\
groups(north1,north2)->min(2);\
groups(north);\
all();
```

In this case we're saying that the cluster should load balance across the north1 and north2 data centers provided there are at least two machines available across them. Failing that, we'll resort to any instance in the north region, and if the whole of the north is offline we'll resort to any instances in the cluster.

### C.2.5. Data center disaster recovery

This section describes how to recover your Neo4j Causal Cluster following a data center failure. Specifically it covers safely turning a small number of surviving instances from a read-only state back into a fully operational cluster of read/write instances.

This section describes the following:

- Data center loss scenario
- Procedure for recovering from data center loss

#### Data center loss scenario

This section describes how to recover a multi-data center deployment which owing to external circumstances has reduced the cluster below half of its members. It is most easily typified by a 2x2 deployment with 2 data centers each containing two instances. This deployment topology can either arise because of other data center failures, or be a deliberate choice to ensure the geographic survival of data for catastrophe planning. However, by distributing an instance over three data centers instead, you could avoid having the cluster lose quorum through a single data center failure. For example, in a 1x1x1 deployment.

Under normal operation this provides a stable majority quorum where the fastest three out of four machines will execute users' transactions, as we see highlighted in Two Data Center Deployment with Four Core Instances.

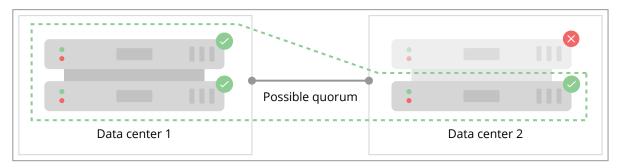


Figure 26. Two Data Center Deployment with Four Core Instances

However if an entire data center becomes offline because of some disaster, then a *majority quorum* cannot be formed in this case.



Neo4j Core clusters are based on the Raft consensus protocol for processing transactions. The Raft protocol requires a majority of cluster members to agree in order to ensure the safety of the cluster and data. As such, the loss of a majority quorum results in a read-only situation for the remaining cluster members.

When data center is lost abruptly in a disaster rather than having the instances cleanly shut down, the surviving members still believe that they are part of a larger cluster. This is different from even the case of rapid failures of individual instances in a live data center which can often be detected by the underlying cluster middleware, allowing the cluster to automatically reconfigure.

Conversely if we lose a data center, there is no opportunity for the cluster to automatically reconfigure. The loss appears instantaneous to other cluster members. However, because each remaining machine has only a partial view of the state of the cluster (its own), it is not safe to allow any individual machine to make an arbitrary decision to reform the cluster.

In this case we are left with two surviving machines which cannot form a quorum and thus make progress.

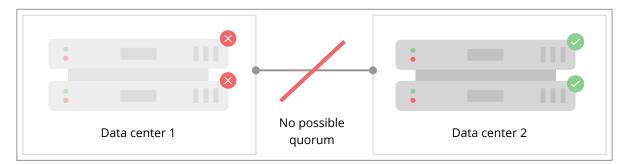


Figure 27. Data Center Loss Requires Guided Recovery

But, from a birds's eye view, it's clear we have surviving machines which are sufficient to allow a non-fault tolerant cluster to form under operator supervision.



Groups of individual cluster members (e.g. those in a single data center) may become isolated from the cluster during network partition for example. If they arbitrarily reformed a new, smaller cluster there is a risk of *split-brain*. That is from the clients' point of view there may be *two* or more smaller clusters that are available for reads and writes depending on the nature of the partition. Such situations lead to divergence that is tricky and laborious to reconcile and so best avoided.

To be safe, an operator or other out-of-band agent (e.g. scripts triggered by well-understood, trustworthy alerts) that has a trusted view on the whole of the system estate must make that decision. In the surviving data center, the cluster can be rebooted into a smaller configuration whilst retaining all data committed to that point. While end users may experience unavailability during the switch over, no committed data will be lost.

### Procedure for recovering from data center loss

The following procedure for performing recovery of a data center should not be done lightly. It assumes that we are completely confident that a disaster has occurred and our previously data center-spanning cluster has been reduced to a read-only cluster in a single data center, where there is no possible way to repair a connection to the lost instances. Further it assumes that the remaining cluster members are fit to provide a seed from which a new cluster can be created from a data quality point of view.

Having acknowledged the above, the procedure for returning the cluster to full availability following catastrophic loss of all but one data centers can be done using one of the following options, depending on your infrastructure.

Please note that the main difference between the options is that Option 2 will allow read-availability during recovery.

#### Option 1.

If you are unable to add instances to the current data-center, and can only use the current read-only cluster, the following steps are recommended:

- 1. Verify that a catastrophe has occurred, and that access to the surviving members of the cluster in the surviving data center is possible. Then for each instance:
  - a. Stop the instance with bin/neo4j stop or shut down the service.
  - b. Change the configuration in *neo4j.conf* such that the causal\_clustering.initial\_discovery\_members property contains the DNS names or IP addresses of the other surviving instances.
  - c. Optional: you may need to update causal\_clustering.minimum\_core\_cluster\_size\_at\_formation, depending on the current size of the cluster (in the current example, two cores).
  - d. Unbind the instance using neo4j-admin unbind.
  - e. Start the instance with bin/neo4j start or start the neo4j service.

#### Option 2.

If it is possible to create a new cluster while the previous read-only cluster is still running, then the following steps will enable you to keep read-availability during recovery:

- 1. Verify that a catastrophe has occurred, and that access to the surviving members of the cluster in the surviving data center is possible.
- 2. Perform an online backup of the currently running, read-only, cluster.

- 3. Seed a new cluster (in the current example, two new cores) using the backup from the read-only cluster, as described in Seed a cluster.
- 4. When the new cluster is up, load balance your workload over to the new cluster.
- 5. Shutdown the old, read-only, cluster.

Once your chosen recovery procedure is completed for each instance, they will form a cluster that is available for reads and writes. It recommended at this point that other cluster members are incorporated into the cluster to improve its load handling and fault tolerance. See Deploy a cluster for details of how to configure instances to join the cluster from scratch.

# C.3. Advanced settings reference

This section lists the advanced settings related to running a Neo4j Causal Cluster.

Parameter	Explanation
causal_clustering.multi_dc	Enables multi-data center features. Requires appropriate licensing.
	<b>Example:</b> causal_clustering.multi_dc_license=true will enable the multi-data center features.
<pre>causal_clustering.server_g roups</pre>	A list of group names for the server used when configuring load balancing and replication policies.
	<b>Example:</b> causal_clustering.server_groups=us,us-east will add the current instance to the groups us and us-east.
causal_clustering.upstream _selection_strategy	An ordered list in descending preference of the strategy which Read Replicas use to choose upstream database server from which to pull transactional updates.
	<b>Example:</b> causal_clustering.upstream_selection_strategy=connect-randomly-within-server-group, typically-connect-to-random-read-replica will configure the behavior so that the Read Replica will first try to connect to any other instance in the group(s) specified in causal_clustering.server_groups. Should we fail to find any live instances in those groups, then we will connect to a random Read Replica. A value of user_defined will enable custom strategy definitions using the setting causal_clustering.user_defined_upstream_strategy.
<pre>causal_clustering.user_def ined_upstream_strategy</pre>	Defines the configuration of upstream dependencies. Can only be used if causal_clustering.upstream_selection_strategy is set to user_defined.
	<b>Example:</b> causal_clustering.user_defined_upstream_strategy=groups(north2); groups(north); halt() will look for servers in the north2. If none are available it will look in the north server group. Finally, if we cannot resolve any servers in any of the previous groups, then rule chain will be stopped via halt().
causal_clustering.load_bal ancing.plugin	The load balancing plugin to use. One pre-defined plugin named server_policies is available by default.
	<b>Example:</b> causal_clustering.load_balancing.plugin=server_policies will enable custom policy definitions.
<pre>causal_clustering.load_bal ancing.config.server_polic ies.<policy-name></policy-name></pre>	Defines a custom policy under the name <policy-name>. Note that load balancing policies are cluster-global configurations and should be defined the exact same way on all core machines.</policy-name>
	Example: causal_clustering.load_balancing.config.server_policies.north1_only=groups(north1) →min(2); halt(); will define a load balancing policy named north1_only. Queries are only sent to servers in the north1 server group, provided there are two of them available. If there are less than two servers in north1 then the chain is halted.

# Appendix D: Deprecated security procedures

This appendix describes deprecated procedures for security management.

This appendix describes deprecated procedures for security management:

- Enterprise Edition
- Community Edition

The procedures described in this appendix have been deprecated and will be removed in a future release.



It is strongly recommended to migrate to the security features as described in Cypher Manual [] Administration [] Security

See also a worked example in Fine-grained access control.

# D.1. Enterprise Edition

This section describes deprecated procedures for native user and role management for Neo4j Enterprise Edition.

A subset of this functionality is also available in Community Edition. The table below includes an indication of which functions this is valid for. Refer to Community Edition for a complete description.

In Neo4j, native user and role management are managed by using built-in procedures through Cypher. This section gives a list of all the security procedures for user management along with some simple examples. Use Neo4j Browser or Neo4j Cypher Shell to run the examples provided.

The following table lists the available procedures:

Procedure name	Description	Executable by role(s)	Available in Community Edition
dbms.security.activateUser	Activate a suspended user	[admin]	
<pre>dbms.security.addRoleToUse r</pre>	Assign a role to the user	[admin]	
dbms.security.changePasswo rd	Change the current user's password	[reader,editor,publisher,arc hitect,admin]	0
dbms.security.changeUserPa ssword	Change the given user's password	[admin]	
dbms.security.createRole	Create a new role	[admin]	
dbms.security.createUser	Create a new user	[admin]	
dbms.security.deleteRole	Delete the specified role. Any role assignments will be removed	[admin]	
dbms.security.deleteUser	Delete the specified user	[admin]	
dbms.security.listRoles	List all available roles	[admin]	
dbms.security.listRolesFor User	List all roles assigned to the specified user	[admin]	
dbms.security.listUsers	List all local users	[admin]	

Procedure name	Description	Executable by role(s)	Available in Community Edition
dbms.security.listUsersFor Role	List all users currently assigned the specified role	[admin]	
<pre>dbms.security.removeRoleFr omUser</pre>	Unassign a role from the user	[admin]	
dbms.security.suspendUser	Suspend the specified user	[admin]	
dbms.showCurrentUser	Show the current user	[reader,editor,publisher,arc hitect,admin]	П
dbms.procedures	List roles per procedure	[reader,editor,publisher,arc hitect,admin]	Not applicable

# D.1.1. Activate a suspended user

An administrator is able to activate a suspended user so that the user is once again able to access the data in their original capacity.

#### Syntax:

CALL dbms.security.activateUser(username, requirePasswordChange)

#### **Arguments:**

Name	Туре	Description
username	String	This is the username of the user to be activated.
requirePasswordChange	Boolean	This is optional, with a default of true. If this is true, (i) the user will be forced to change their password when they next log in, and (ii) until the user has changed their password, they will be forbidden from performing any other operation.

#### **Exceptions:**

The current user is not an administrator.

The username does not exist in the system.

The username matches that of the current user (i.e. activating the current user is not permitted).

#### **Considerations:**

This is an idempotent procedure.

#### Example 122. Activate a suspended user

The following example activates a user with the username 'jackgreen'. When the user 'jackgreen' next logs in, he will be required to change his password.

CALL dbms.security.activateUser('jackgreen')

# D.1.2. Assign a role to the user

An administrator is able to assign a role to any user in the system, thus allowing the user to perform a series of actions upon the data.

#### Syntax:

CALL dbms.security.addRoleToUser(roleName, username)

#### **Arguments:**

Name	Туре	Description
roleName	String	This is the name of the role to be assigned to the user.
username	String	This is the username of the user who is to be assigned the role.

#### **Exceptions:**

The current user is not an administrator.

The username does not exist in the system.

The username contains characters other than alphanumeric characters and the '\_' character.

The role name does not exist in the system.

The role name contains characters other than alphanumeric characters and the '\_' character.

#### **Considerations:**

This is an idempotent procedure.

#### Example 123. Assign a role to the user

The following example assigns the role publisher to the user with username 'johnsmith'.

CALL dbms.security.addRoleToUser('publisher', 'johnsmith')

### D.1.3. Change the current user's password



The procedure dbms.security.changePassword(newPassword, requirePasswordChange) has been entirely removed since the corresponding Cypher administration command also requires the old password, and thus is more secure. Please use ALTER CURRENT USER SET PASSWORD FROM 'oldPassword' TO 'newPassword', documented in the Cypher Manual, instead.

### D.1.4. Change the given user's password

An administrator is able to change the password of any user within the system. Alternatively, the current user may change their own password.

#### Syntax:

CALL dbms.security.changeUserPassword(username, newPassword, requirePasswordChange)

#### **Arguments:**

Name	Туре	Description
username	String	This is the username of the user whose password is to be changed.
newPassword	String	This is the new password for the user.
requirePasswordChange	Boolean	This is optional, with a default of true. If this is true, (i) the user will be forced to change their password when they next log in, and (ii) until the user has changed their password, they will be forbidden from performing any other operation.

#### **Exceptions:**

The current user is not an administrator and the username does not match that of the current user.

The username does not exist in the system.

The password is the empty string.

The password is the same as the user's previous password.

#### **Considerations:**

This procedure may be invoked by the current user to change their own password, irrespective of whether or not the current user is an administrator.

This procedure may be invoked by an administrator to change another user's password.

In addition to changing the user's password, this will terminate with immediate effect all of the user's sessions and roll back any running transactions.

#### Example 124. Change a given user's password

The following example changes the password of the user with the username 'joebloggs' to 'h6u4%kr'. When the user 'joebloggs' next logs in, he will be required to change his password.

CALL dbms.security.changeUserPassword('joebloggs', 'h6u4%kr')

### D.1.5. Create a new role

An administrator is able to create custom roles in the system.

#### Syntax:

CALL dbms.security.createRole(roleName)

#### **Arguments:**

Name	Туре	Description
roleName	String	This is the name of the role to be created.

#### **Exceptions:**

The current user is not an administrator.

The role name already exists in the system.

The role name is empty.

The role name contains characters other than alphanumeric characters and the '\_' character.

The role name matches one of the native roles: reader, publisher, architect, and admin.

#### Example 125. Create a new role

The following example creates a new custom role.

CALL dbms.security.createRole('operator')

### D.1.6. Create a new user

An administrator is able to create a new user. This action ought to be followed by assigning a role to the user, which is described here.

#### Syntax:

CALL dbms.security.createUser(username, password, requirePasswordChange)

#### **Arguments:**

Name	Туре	Description
username	String	This is the user's username.
password	String	This is the user's password.
requirePasswordChange	Boolean	This is optional, with a default of true. If this is true, (i) the user will be forced to change their password when they log in for the first time, and (ii) until the user has changed their password, they will be forbidden from performing any other operation.

#### **Exceptions:**

The current user is not an administrator.

The username either contains characters other than the ASCII characters between ! and  $\sim$ , or contains : and  $\downarrow$ .

The username is already in use within the system.

The password is the empty string.

#### Example 126. Create a new user

The following example creates a user with the username 'johnsmith' and password 'h6u4%kr'. When the user 'johnsmith' logs in for the first time, he will be required to change his password.

CALL dbms.security.createUser('johnsmith', 'h6u4%kr')

# D.1.7. Delete the specified role

An administrator is able to delete custom roles from the system. The native roles reader, publisher, architect, and admin cannot be deleted.

#### Syntax:

CALL dbms.security.deleteRole(roleName)

#### **Arguments:**

Name	Туре	Description
roleName	String	This is the name of the role to be deleted.

#### **Exceptions:**

The current user is not an administrator.

The role name does not exist in the system.

The role name matches one of the native roles: reader, publisher, architect, and admin.

#### **Considerations:**

Any role assignments will be removed.

#### Example 127. Delete the specified role

The following example deletes the custom role 'operator' from the system.

CALL dbms.security.deleteRole('operator')

# D.1.8. Delete the specified user

An administrator is able to delete permanently a user from the system. It is not possible to undo this action, so, if in any doubt, consider suspending the user instead.

#### Syntax:

CALL dbms.security.deleteUser(username)

#### **Arguments:**

Name	Туре	Description
username	String	This is the username of the user to be deleted.

#### **Exceptions:**

The current user is not an administrator.

The username does not exist in the system.

The username matches that of the current user (i.e. deleting the current user is not permitted).

#### **Considerations:**

It is not necessary to remove any assigned roles from the user prior to deleting the user.

Deleting a user will terminate with immediate effect all of the user's sessions and roll back any running transactions.

As it is not possible for the current user to delete themselves, there will always be at least one administrator in the system.

#### Example 128. Delete the specified user

The following example deletes a user with the username 'janebrown'.

```
CALL dbms.security.deleteUser('janebrown')
```

### D.1.9. List all available roles

An administrator is able to view all assigned users for each role in the system.

#### Syntax:

CALL dbms.security.listRoles()

#### **Returns:**

Name	Туре	Description
role	String	This is the name of the role.
users	List <string></string>	This is a list of the usernames of all users who have been assigned the role.

#### **Exceptions:**

The current user is not an administrator.

#### Example 129. List all available roles

The following example shows, for each role in the system, the name of the role and the usernames of all assigned users.

```
CALL dbms.security.listRoles()
```

# D.1.10. List all roles assigned to the specified user

Any active user is able to view all of their assigned roles. An administrator is able to view all assigned roles for any user in the system.

#### Syntax:

CALL dbms.security.listRolesForUser(username)

#### **Arguments:**

Name	Туре	Description
username	String	This is the username of the user.

#### **Returns:**

Name	Туре	Description
value		This returns all roles assigned to the requested user.

#### **Exceptions:**

The current user is not an administrator and the username does not match that of the current user.

The username does not exist in the system.

#### **Considerations:**

This procedure may be invoked by the current user to view their roles, irrespective of whether or not the current user is an administrator.

This procedure may be invoked by an administrator to view the roles for another user.

Example 130. List all roles assigned to the specified user

The following example lists all the roles for the user with username 'johnsmith', who has the roles reader and publisher.

### D.1.11. List all local users

An administrator is able to view the details of every user in the system.

### Syntax:

CALL dbms.security.listUsers()

#### **Returns:**

Name	Туре	Description
username	String	This is the user's username.
roles	List <string></string>	This is a list of roles assigned to the user.
flags	List <string></string>	This is a series of flags indicating whether the user is suspended or needs to change their password.

#### **Exceptions:**

The current user is not an administrator.

#### Example 131. List all local users

The following example shows, for each user in the system, the username, the roles assigned to the user, and whether the user is suspended or needs to change their password.

# D.1.12. List all users currently assigned the specified role

An administrator is able to view all assigned users for a role.

#### Syntax:

CALL dbms.security.listUsersForRole(roleName)

#### **Arguments:**

Name	Туре	Description
roleName	String	This is the name of the role.

#### **Returns:**

Name	Туре	Description
value	String	This returns all assigned users for the requested role.

#### **Exceptions:**

The current user is not an administrator.

The role name does not exist in the system.

The following example lists all the assigned users - 'bill' and 'anne' - for the role publisher.

CALL dbms.security.listUsersForRole('publisher')

+-----+
| value |

# D.1.13. Unassign a role from the user

An administrator is able to remove a role from any user in the system, thus preventing the user from performing upon the data any actions prescribed by the role.

#### Syntax:

CALL dbms.security.removeRoleFromUser(roleName, username)

#### **Arguments:**

"bill" "anne"

2 rows

Name	Туре	Description
roleName	String	This is the name of the role which is to be removed from the user.
username	String	This is the username of the user from which the role is to be removed.

#### **Exceptions:**

The current user is not an administrator.

The username does not exist in the system.

The role name does not exist in the system.

The username is that of the current user and the role is admin.

#### **Considerations:**

If the username is that of the current user and the role name provided is admin, an error will be thrown; i.e. the current user may not be demoted from being an administrator.

As it is not possible for the current user to remove the admin role from themselves, there will always be at least one administrator in the system.

This is an idempotent procedure.

#### Example 133. Unassign a role from the user

The following example removes the role publisher from the user with username 'johnsmith'.

CALL dbms.security.removeRoleFromUser('publisher', 'johnsmith')

# D.1.14. Suspend the specified user

An administrator is able to suspend a user from the system. The suspended user may be activated at a later stage.

#### Syntax:

CALL dbms.security.suspendUser(username)

#### **Arguments:**

Name	Туре	Description
username	String	This is the username of the user to be suspended.

#### **Exceptions:**

The current user is not an administrator.

The username does not exist in the system.

The username matches that of the current user (i.e. suspending the current user is not permitted).

#### **Considerations:**

Suspending a user will terminate with immediate effect all of the user's sessions and roll back any running transactions.

All of the suspended user's attributes — assigned roles and password — will remain intact.

A suspended user will not be able to log on to the system.

As it is not possible for the current user to suspend themselves, there will always be at least one active administrator in the system.

This is an idempotent procedure.

#### Example 134. Suspend the specified user

The following example suspends a user with the username 'billjones'.

CALL dbms.security.suspendUser('billjones')

# D.1.15. Show the current user

The current user is able to view whether or not they need to change their password.

#### Syntax:

CALL dbms.showCurrentUser()

#### **Returns:**

Name	Туре	Description
username	String	This is the user's username.
flags	List <string></string>	This is a flag indicating whether the user needs change their password.

#### Example 135. Show the current user

The following example shows that the current user — with the username '**johnsmith**' — does not need to change his password.

# D.1.16. List roles per procedure

Any active user is able to view all procedures in the system, including which role(s) have the privilege to execute them. Note that the defaultBuiltInRoles column does not include roles created by administration commands, unless they are included in the procedure security settings.

#### Syntax:

CALL dbms.procedures()

#### **Returns:**

Name	Туре	Description
name	String	This is the name of the procedure.
signature	String	This is the signature of the procedure.
description	String	This is a description of the procedure.
defaultBuiltInRoles	List <string></string>	This is a list of roles having the privilege to execute the procedure.

The following example shows, for four of the security procedures, the procedure name, the description, and which roles have the privilege to execute the procedure.

```
CALL dbms.procedures()
YIELD name, signature, description, roles
WITH name, description, roles
WHERE name contains 'security'
RETURN name, description, roles
ORDER BY name
LIMIT 4
```

```
| toles | roles | role
```

# D.2. Community Edition

This section describes deprecated procedures for user and password management for Neo4j Community Edition.

User and password management for Community Edition is a subset of the functionality available in Enterprise Edition. The following is true for user management in Community Edition:

- It is possible to create multiple users.
- All users assume the privileges of an admin for the available functionality.

Users are managed by using built-in procedures through Cypher. This section gives a list of all the security procedures for user management along with some simple examples. Use Neo4j Browser or Neo4j Cypher Shell to run the examples provided. Unless stated otherwise, all arguments to the procedures described in this section must be supplied.

Name	Description
dbms.security.changePassword	Change the current user's password
dbms.security.createUser	Add a user
dbms.security.deleteUser	Delete a user
dbms.security.listUsers	List all users
dbms.showCurrentUser	Show details for the current user

# D.2.1. Change the current user's password



The procedure dbms.security.changePassword(newPassword, requirePasswordChange) has been entirely removed since the corresponding Cypher administration command also requires the old password, and thus is more secure. Please use ALTER CURRENT USER SET PASSWORD FROM 'oldPassword' TO 'newPassword', documented in the Cypher Manual, instead.

### D.2.2. Add a user

The current user is able to add a user to the system.

#### Syntax:

CALL dbms.security.createUser(username, password, requirePasswordChange)

#### **Arguments:**

Name	Туре	Description
username	String	This is the user's username.
password	String	This is the user's password.
requirePasswordChange	Boolean	This is optional, with a default of true. If this is true, (i) the user will be forced to change their password when they log in for the first time, and (ii) until the user has changed their password, they will be forbidden from performing any other operation.

#### **Exceptions:**

The username either contains characters other than the ASCII characters between! and ~, or contains: and ,.

The username is already in use within the system.

The password is the empty string.

#### Example 137. Add a user

The following example creates a user with the username 'johnsmith' and password 'h6u4%kr'. When the user 'johnsmith' logs in for the first time, he will be required to change his password.

CALL dbms.security.createUser('johnsmith', 'h6u4%kr', true)

### D.2.3. Delete a user

The current user is able to delete permanently a user from the system.

#### Syntax:

CALL dbms.security.deleteUser(username)

#### **Arguments:**

Name	Туре	Description
username	String	This is the username of the user to be deleted.

#### **Exceptions:**

The username does not exist in the system.

The username matches that of the current user (i.e. deleting the current user is not permitted).

#### **Considerations:**

Deleting a user will terminate with immediate effect all of the user's sessions and roll back any running transactions.

As it is not possible for the current user to delete themselves, there will always be at least one user in the system.

#### Example 138. Delete a user

The following example deletes a user with the username 'janebrown'.

```
CALL dbms.security.deleteUser('janebrown')
```

### D.2.4. List all native users

The current user is able to view the details of every user in the system.

#### Syntax:

CALL dbms.security.listUsers()

#### **Returns:**

Name	Туре	Description
username	String	This is the user's username.
flags	List <string></string>	This is a flag indicating whether the user needs to change their password.

#### Example 139. List all users

The following example shows the username for each user in the system, and whether the user needs to change their password.

```
CALL dbms.security.listUsers()
```

### D.2.5. Show details for the current user

The current user is able to view whether or not they need to change their password.

#### Syntax:

CALL dbms.showCurrentUser()

#### **Returns:**

Name	Туре	Description
username	String	This is the user's username.
flags	List <string></string>	This is a flag indicating whether the user needs change their password.

#### Example 140. Show details for the current user

The following example shows that the current user — with the username 'johnsmith' — does not need to change his password.

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