Laboratorium 5: Zastosowanie szeregowania afinicznego do implementacji techniki blokowania pętli

Wariant petli 2

```
for(i=1; i<=n; i++)
    for(j=2; j<=n; j++)
    a[i][j] = a[i][j-2];</pre>
```

Zadanie 1.

Dla wskazanej pętli za pomocą kalkulatora ISCC znaleźć relację zależności R, przestrzeń iteracji LD, oraz zrobić rysunek grafu zależności w przestrzeni 6 x 6.

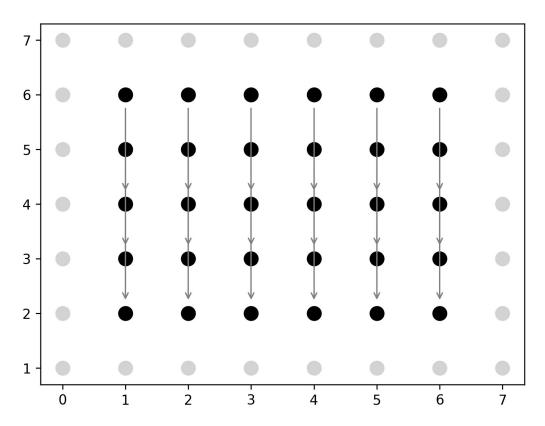
Relacja R:

$$[n] \rightarrow \{ S_0[i, j] \rightarrow S_0[i' = i, j' = 2 + j] : 0 < i < n \text{ and } 2 <= j <= -2 + n \}$$

Przestrzeń iteracji LD (Loop Domain):

$$[n] \rightarrow \{ S_0[i, j] : 0 < i <= n \text{ and } 2 <= j <= n \}$$

Rysunek grafu w przestrzeni 6x6:



Zadanie 2.

Za pomocą operatora kalkulatora ISCC: IslSchedule := schedule LD respecting R minimizing R znaleźć szeregowanie afiniczne w postaci drzewa.

Szeregowanie afiniczne w postaci drzewa:

```
domain: "[n] -> { S_0[i, j] : 0 < i <= n \text{ and } 2 <= j <= n }"
```

child:

```
schedule: "[n] -> [{ S_0[i, j] -> [(i)] }, { S_0[i, j] -> [(j)] }]"
permutable: 1
coincident: [ 1, 1 ]
```

W powyższym drzewie znajdują się dwa następujące szeregowania ISL:

- {[i, j] -> [(i)] }
- {[i, j] -> [(j)] }

Zadanie 3.

Za pomocą operatora kalkulatora ISCC: map przekonwertować szeregowanie afiniczne w postaci drzewa na szeregowanie w postaci relacji.

Szeregowanie w postaci relacji:

```
[n] -> { S_0[i, j] -> [i, j] }
```

Zadanie 4.

Utworzyć szeregowanie, które pozwala na implementację techniki fali frontowej na poziomie iteracji (wave-fronting).

```
"WAVE_FR"
[n] -> { S_0[i, j] -> [i + j, j] : 0 < i <= n and 2 <= j <= n }
```

Zadanie 5.

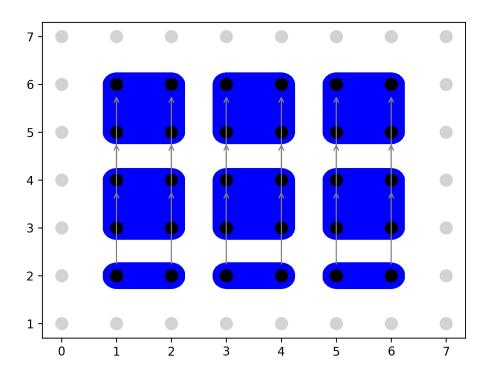
Utworzyć szeregowanie, które pozwala na implementację techniki kafelkowania (tiling) z rozmiarem kafelka 2x2, sekwencyjny sposób wykonywania kafelków.

```
"CODE_TILING"  [n] \rightarrow \{ [i, j] \rightarrow [it, jt, i, j] : 0 < i <= n \ and \ 2 <= j <= n \ and \ -2 + i <= 2it < i \ and \ -2 + j <= 2jt < j \ \}
```

Zadanie 6.

Stosując uzyskane w punkcie 5 szeregowanie za pomocą operatora scan znaleźć wszystkie kafelki w przestrzeni 6x6 (12x12) i zaznaczyć je na rysunku utworzonego w p.1.

```
[n] -> { [i = 6, j = 6] -> [it = 2, jt = 2, 6, 6] : n = 6; [i = 5, j = 6] -> [it = 2, jt = 2, 5, 6] : n = 6; [i = 4, j = 6] -> [it = 1, jt = 2, 4, 6] : n = 6; [i = 3, j = 6] -> [it = 1, jt = 2, 3, 6] : n = 6; [i = 2, j = 6] -> [it = 0, jt = 2, 2, 6] : n = 6; [i = 1, j = 6] -> [it = 0, jt = 2, 1, 6] : n = 6; [i = 6, j = 5] -> [it = 2, jt = 2, 6, 5] : n = 6; [i = 5, j = 5] -> [it = 2, jt = 2, 5, 5] : n = 6; [i = 4, j = 5] -> [it = 1, jt = 2, 4, 5] : n = 6; [i = 3, j = 5] -> [it = 1, jt = 2, 3, 5] : n = 6; [i = 2, j = 5] -> [it = 0, jt = 2, 2, 5] : n = 6; [i = 6, j = 4] -> [it = 2, jt = 1, 6, 4] : n = 6; [i = 5, j = 4] -> [it = 2, jt = 1, 5, 4] : n = 6; [i = 4, j = 4] -> [it = 1, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] -> [it = 1, jt = 1, 3, 4] : n = 6; [i = 2, j = 4] -> [it = 0, jt = 1, 2, 4] : n = 6; [i = 5, j = 3] -> [it = 2, jt = 1, 5, 3] : n = 6; [i = 5, j = 3] -> [it = 2, jt = 1, 5, 3] : n = 6; [i = 3, j = 3] -> [it = 1, jt = 1, 3, 3] : n = 6; [i = 3, j = 3] -> [it = 1, jt = 1, 3, 3] : n = 6; [i = 2, j = 3] -> [it = 0, jt = 1, 2, 3] : n = 6; [i = 1, j = 3] -> [it = 0, jt = 1, 1, 3] : n = 6; [i = 2, j = 3] -> [it = 2, jt = 2] -> [it = 2, jt = 0, 6, 2] : n = 6; [i = 5, j = 2] -> [it = 2, jt = 0, 5, 2] : n = 6; [i = 4, j = 2] -> [it = 2, jt = 0, 5, 2] : n = 6; [i = 4, j = 2] -> [it = 2, jt = 2] -> [it = 2, jt = 0, 5, 2] : n = 6; [i = 4, j = 2] -> [it = 2, jt = 0, 5, 2] : n = 6; [i = 4, j = 2] -> [it = 2, jt = 2] -> [it = 2,
```



Zadanie 7.

Wygenerować pseudokod sekwencyjny i odpowiedni kod kompilowany (sekwencyjny) implementujący kafelkowanie.

```
Wygenerowany pseudokod sekwencyjny:
```

Kod kompilowalny:

```
for (int c0 = 0; c0 < floord(n + 1, 2); c0 += 1)
  #pragma openmp parallel for
  for (int c1 = 0; c1 < (n + 1) / 2; c1 += 1)
    for (int c2 = 2 * c0 + 1; c2 <= min(n, 2 * c0 + 2); c2 += 1)
        for (int c3 = max(2, 2 * c1 + 1); c3 <= min(n, 2 * c1 + 2); c3 += 1)
        a[c2][c3] = a[c2][c3-2];</pre>
```

Zadanie 8.

Zastosować program porównujący wyniki obliczeń (zadanie 7, L2) do sprawdzania poprawności kodu docelowego wygenerowanego w p.6 w przestrzeni 6x6.

```
00 01 00 01 00 01 00 00 00 01 00 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00
```

Zadanie 9.

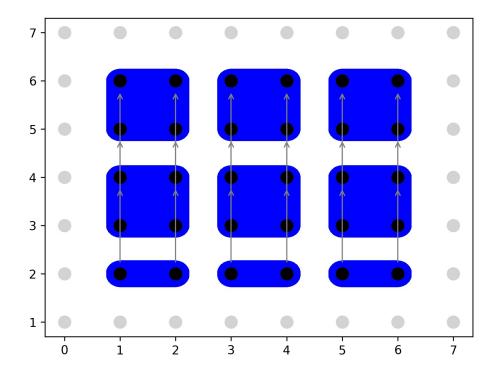
Utworzyć szeregowanie, które pozwala na implementację techniki wave-fronting na poziomie kafli (tiles) z rozmiarem kafelka 2x2, równoległy sposób wykonywania kafelków.

```
"CODE_TILING_PAR"  [n] \rightarrow \{ [i, j] \rightarrow [it, jt, i, j] : 0 < i <= n \ and \ 0 < j <= n \ and \ 2jt >= -2 + j \ and \ -i + 2it < 2jt <= 2 - i + 2it \ and \ 2jt < j \}
```

Zadanie 10.

Stosując uzyskane w punkcie 9 szeregowanie za pomocą operatora scan znaleźć wszystkie partycje czasu na poziomie kafelków w przestrzeni 6x6 (12x12) i zaznaczyć je na rysunku utworzonego w p.1. Czasem wykonania kafelka jest wartość iteratora it.

```
"scan (CODE TILING*[n]->{:n=6})"
[n] \rightarrow \{ [i = 6, j = 6] \rightarrow [it = 4, jt = 2, 6, 6] : n = 6; [i = 5, j = 6] \rightarrow [it = 4, jt = 6] \}
= 2, 5, 6] : n = 6; [i = 4, j = 6] -> [it = 3, jt = 2, 4, 6] : <math>n = 6; [i = 3, j = 6] ->
[it = 3, jt = 2, 3, 6] : n = 6; [i = 2, j = 6] \rightarrow [it = 2, jt = 2, 2, 6] : n = 6; [i = 1, 1]
j = 6] -> [it = 2, jt = 2, 1, 6] : n = 6; [i = 6, j = 5] -> [it = 4, jt = 2, 6, 5] : n =
6; [i = 5, j = 5] \rightarrow [it = 4, jt = 2, 5, 5] : n = 6; [i = 4, j = 5] \rightarrow [it = 3, jt = 2, 5]
4, 5] : n = 6; [i = 3, j = 5] -> [it = 3, jt = 2, 3, 5] : <math>n = 6; [i = 2, j = 5] -> [it = 3, jt = 2, 3, 5]
2, jt = 2, 2, 5] : n = 6; [i = 1, j = 5] \rightarrow [it = 2, jt = 2, 1, 5] : <math>n = 6; [i = 6, j = 1]
4] -> [it = 3, jt = 1, 6, 4] : n = 6; [i = 5, j = 4] -> [it = 3, jt = 1, 5, 4] : n = 6;
[i = 4, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 3, j = 4] \rightarrow [it = 2, jt = 1, 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, j = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [i = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 2, jt = 1, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 2, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 3, jt = 4] \rightarrow [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n = 6; [it = 4, 4, 4] : n =
4] : n = 6; [i = 2, j = 4] -> [it = 1, jt = 1, 2, 4] : n = 6; [i = 1, j = 4] -> [it = 1,
jt = 1, 1, 4]: n = 6; [i = 6, j = 3] -> [it = 3, jt = 1, 6, 3]: n = 6; [i = 5, j = 3]
\rightarrow [it = 3, jt = 1, 5, 3] : n = 6; [i = 4, j = 3] \rightarrow [it = 2, jt = 1, 4, 3] : n = 6; [i =
3, j = 3] -> [it = 2, jt = 1, 3, 3] : n = 6; [i = 2, j = 3] -> [it = 1, jt = 1, 2, 3] : n
= 6; [i = 1, j = 3] \rightarrow [it = 1, jt = 1, 1, 3] : n = 6; [i = 6, j = 2] \rightarrow [it = 2, jt = 0,
6, 2] : n = 6; [i = 5, j = 2] -> [it = 2, jt = 0, 5, 2] : <math>n = 6; [i = 4, j = 2] -> [it = 2, jt = 0, 5, 2]
1, jt = 0, 4, 2] : n = 6; [i = 3, j = 2] \rightarrow [it = 1, jt = 0, 3, 2] : n = 6; [i = 2, j = 1]
2] -> [it = 0, jt = 0, 2, 2] : n = 6; [i = 1, j = 2] -> [it = 0, jt = 0, 1, 2] : n = 6;
[i = 6, j = 1] \rightarrow [it = 2, jt = 0, 6, 1] : n = 6; [i = 5, j = 1] \rightarrow [it = 2, jt = 0, 5, 1]
1] : n = 6; [i = 4, j = 1] -> [it = 1, jt = 0, 4, 1] : <math>n = 6; [i = 3, j = 1] -> [it = 1, jt = 1
jt = 0, 3, 1]: n = 6; [i = 2, j = 1] \rightarrow [it = 0, jt = 0, 2, 1]: n = 6; [i = 1, j = 1]
-> [it = 0, jt = 0, 1, 1] : n = 6 }
```



Zadanie 11.

Wygenerować pseudokod i równoległy kod kompilowalny implementujący kafelkowanie.

```
Wygenerowany pseudokod sekwencyjny:
for (int c0 = 0; c0 < floord(n + 1, 2); c0 += 1)
   for (int c1 = 0; c1 < (n + 1) / 2; c1 += 1)
     for (int c2 = 2 * c0 + 1; c2 <= min(n, 2 * c0 + 2); c2 += 1)
        for (int c3 = max(2, 2 * c1 + 1); c3 <= min(n, 2 * c1 + 2); c3 += 1)
        (c2, c3);

Kod kompilowany:
for (int c0 = 0; c0 < floord(n + 1, 2); c0 += 1)
     #pragma openmp parallel for
   for (int c1 = 0; c1 < (n + 1) / 2; c1 += 1)
     for (int c2 = 2 * c0 + 1; c2 <= min(n, 2 * c0 + 2); c2 += 1)
     for (int c3 = max(2, 2 * c1 + 1); c3 <= min(n, 2 * c1 + 2); c3 += 1)
     a[c2][c3] = a[c2][c3-2];</pre>
```

Zadanie 12.

Zastosować program porównujący wyniki obliczeń (zadanie 7, L2) do sprawdzania poprawności kodu docelowego w przestrzeni 6x6.

```
# gcc -fopenmp 12-joined.c -lm && ./a.out
Initial code result:
00 01 02 03 04 05 06
00 01 00 01 00 01 00
00 01 00 01 00 01 00
00 01 00 01 00 01 00
```

```
00 01 00 01 00 01 00 00 00 01 00 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00 01 00
```

Załączniki.

Skrypt implementujący zadania.

```
##Znalezienie zaleznosci
P := parse_file "1-my.c";
print "Loop domain:";
Domain := P[0];
LD := Domain;
print Domain;
Write := P[1] * Domain;
Read := P[3] * Domain;
Schedule := P[4];
Before := Schedule << Schedule;</pre>
RaW := (Write . (Read^-1)) * Before;
WaW := (Write . (Write^-1)) * Before;
WaR := (Read . (Write^-1)) * Before;
R := (RaW+WaW+WaR);
print "R";
print R;
print "scan (R*[n]->{:n=6})";
scan (R*[n]->{:n=6});
##krok 2
IslSchedule := schedule LD respecting R minimizing R;
print "IslSchedule";
IslSchedule;
##krok 3
SCHED:= map IslSchedule;
print "SCHED";
SCHED;
##krok 4
WAVE_FR:=[n] \rightarrow \{ S_0[i, j] \rightarrow [i+j, j] \}*LD;
print "WAVE_FR";
```

```
WAVE_FR;
##krok 5, kafelkowanie, generowanie kodu sekwencyjnego
CODE_TILING:= [n]->{
    [i,j]->[it, jt, i, j]:
    0 < i <= n and
    1 < j <= n and
    1<= i-2it <=2 and
    1<= j-2jt <=2
};
print "CODE_TILING";
CODE_TILING;
print "scan (CODE_TILING*[n]->{:n=6})";
scan (CODE_TILING*[n]->{:n=6});
##krok 7 generowanie kodu sekwencyjnego
print "codegen CODE_TILING";
codegen CODE_TILING;
##krok 9 tworzenie relacji pozwalajacej na implementacje wave-
#fronting na poziomie kafelkow:
CODE_TILING_PAR:= [n]->{
    [i,j]->[it, jt, i, j]:
        0 < i <= n and
        1 < j <= n and
        1<= i-2it <=2 and
        1<= j-2jt <=2
};
print "CODE_TILING_PAR";
CODE_TILING_PAR;
#krok 11 generowanie kodu rownoleglego
codegen CODE_TILING_PAR;
```