# OPTIMIZATION (SI 416) – LECTURE 5

## Harsha Hutridurga

IIT Bombay

#### RECENT STORY

- $\clubsuit$  Take a continuously differentiable function  $f: \mathbb{R}^n \to \mathbb{R}$
- Gradient descent algorithm reads as follows:

$$\begin{cases} \text{Begin with a} & x^{(0)} \in \mathbb{R}^n \\ \text{Build iterates using} & x^{(n+1)} = x^{(n)} - \delta \nabla f(x^{(n)}) & \text{for } n = 0, 1, 2, \end{cases}$$

- $\clubsuit$  If f is strongly convex, then there is a unique global minimizer  $x_*$
- A If f is further assumed to be  $\beta$ -smooth, then picking  $\delta \in (0, \beta^{-1})$  yields a minimizing sequence, i.e.  $f(x^{(n+1)}) \leq f(x^{(n)})$
- Furthermore, we have the estimate:

$$||x^{(n)} - x_*|| \le \left(\frac{1}{1 + 2\delta\lambda}\right)^{\frac{n}{2}} ||x^{(0)} - x_*||$$
 for  $n = 0, 1, ...$ 

## RECENT STORY (CONTD.)

- \$\ \sim \text{Suppose we have a tolerance of } \varepsilon > 0\$, i.e we are looking for  $x^{(n)}$  which is at  $\varepsilon$  distance from  $x_*$
- ♣ Observe that

$$\left(\frac{1}{1+2\delta\lambda}\right)^{\frac{n}{2}} \left\| x^{(0)} - x_* \right\| \le \varepsilon \implies \left\| x^{(n)} - x_* \right\| \le \varepsilon$$

♣ That is

$$n \ge \frac{2}{\ln(1+2\delta\lambda)} \ln\left(\frac{\left\|x^{(0)} - x_*\right\|}{\varepsilon}\right)$$

 $\clubsuit$  Hence, for the  $n^{\text{th}}$  iterate to be  $\varepsilon$  close to  $x_*$ , we must have

$$n = \mathcal{O}(\ln(\varepsilon^{-1}))$$

#### AN IDEA OF NEWTON

- $\clubsuit$  Suppose  $f: \mathbb{R}^n \to \mathbb{R}$  is such that  $\nabla^2 f(x)$  is invertible for all x
- Consider the algorithm

$$\begin{cases} \text{ Begin with a } & x^{(0)} \in \mathbb{R}^n \\ \text{ Build iterates using } & x^{(n+1)} = x^{(n)} - \delta \left( \nabla^2 f(x^{(n)}) \right)^{-1} \nabla f(x^{(n)}) \end{cases}$$

- for  $n = 0, 1, 2, \dots$
- $\clubsuit$  The parameter  $\delta > 0$  to be chosen later
- Does this generate a minimizing sequence?
- & Employing Taylor's theorem, we get

$$f(x^{(n+1)}) = f(x^{(n)}) - \delta \left\langle \nabla f(x^{(n)}), \left( \nabla^2 f(x^{(n)}) \right)^{-1} \nabla f(x^{(n)}) \right\rangle + \frac{\delta^2}{2} \left\langle \nabla^2 f(y) \left( \nabla^2 f(x^{(n)}) \right)^{-1} \nabla f(x^{(n)}), \left( \nabla^2 f(x^{(n)}) \right)^{-1} \nabla f(x^{(n)}) \right\rangle$$

- $\bullet$  Observe that choosing  $\delta \ll 1$ , we can drop the term of  $\mathcal{O}(\delta^2)$
- $\clubsuit$  Note: Positive definite  $\nabla^2 f$  will generate minimizing sequence

 $\clubsuit$  Similar to the  $\beta$ -smoothness condition, we shall assume that

$$\left\| \nabla^2 f(x) v - \nabla^2 f(y) v \right\| \le \gamma \|x - y\| \|v\| \quad \text{ for all } x, y, v \in \mathbb{R}^n$$
 for some  $\gamma > 0$ 

- $\clubsuit$  Suppose f is strongly convex, i.e. a minimizer  $x_*$  exists
- $\clubsuit$  Take  $\delta = 1$  in the algorithm and note that

$$x^{(n+1)} - x_* = x^{(n)} - x_* - \left(\nabla^2 f(x^{(n)})\right)^{-1} \nabla f(x^{(n)})$$

$$= x^{(n)} - x_*$$

$$- \left(\nabla^2 f(x^{(n)})\right)^{-1} \int_0^1 \nabla^2 f(x_* + \alpha(x^{(n)} - x_*))(x^{(n)} - x_*) d\alpha$$

$$= \left(\nabla^2 f(x^{(n)})\right)^{-1} \left(\nabla^2 f(x^{(n)})\right) \left(x^{(n)} - x_*\right)$$

$$- \left(\nabla^2 f(x^{(n)})\right)^{-1} \int_0^1 \nabla^2 f(x_* + \alpha(x^{(n)} - x_*))(x^{(n)} - x_*) d\alpha$$

Recall that we had

$$x^{(n+1)} - x_* = \left(\nabla^2 f(x^{(n)})\right)^{-1}$$
$$\left(\int_0^1 \left(\nabla^2 f(x^{(n)}) - \nabla^2 f(x_* + \alpha(x^{(n)} - x_*))\right) d\alpha\right) \left(x^{(n)} - x_*\right)$$

♣ Hence we deduce that

$$\begin{aligned} & \left\| x^{(n+1)} - x_* \right\| \le \left\| \left( \nabla^2 f(x^{(n)}) \right)^{-1} \right\| \\ & \left\| \int_0^1 \left( \nabla^2 f(x^{(n)}) - \nabla^2 f(x_* + \alpha(x^{(n)} - x_*)) \right) \, d\alpha \right\| \left\| x^{(n)} - x_* \right\| \\ & \le \frac{\left\| x^{(n)} - x_* \right\|}{\lambda} \int_0^1 \left\| \nabla^2 f(x^{(n)}) - \nabla^2 f(x_* + \alpha(x^{(n)} - x_*)) \right\| \, d\alpha \end{aligned}$$

thanks to strong convexity of f

4 Using the smoothness assumption on the Hessian, we obtain

$$\left\| x^{(n+1)} - x_* \right\| \le \frac{\|x^{(n)} - x_*\|}{\lambda} \int_0^1 \gamma (1 - \alpha) \left\| x^{(n)} - x_* \right\| d\alpha$$
$$= \frac{\gamma}{2\lambda} \left\| x^{(n)} - x_* \right\|^2$$

A Hence we deduce that

$$||x^{(n)} - x_*|| \le \left(\frac{\gamma}{2\lambda}\right)^{2^n - 1} ||x^{(0)} - x_*||^{2^n}$$

♣ Observe that

$$\left\|x^{(0)} - x_*\right\| \le \frac{\lambda}{\gamma} \implies \left\|x^{(n)} - x_*\right\| \le \left(\frac{2\lambda}{\gamma}\right) 2^{-2^n}$$

- \$\iiiis\$ Suppose we have a tolerance of  $\varepsilon > 0$ , i.e we are looking for  $x^{(n)}$  which is at  $\varepsilon$  distance from  $x_*$
- ♣ Observe that

$$\left(\frac{2\lambda}{\gamma}\right)2^{-2^n} \le \varepsilon \implies \left\|x^{(n)} - x_*\right\| \le \varepsilon$$

A That is

$$n \ge \log_2\left(\log_2\left(\frac{2\lambda}{\gamma\varepsilon}\right)\right)$$

 $\clubsuit$  Hence, for the  $n^{\text{th}}$  iterate to be  $\varepsilon$  close to  $x_*$ , we must have

$$n = \mathcal{O}(\ln(\ln(\varepsilon^{-1})))$$

 $\clubsuit$  Recall that for gradient descent, we had  $n = \mathcal{O}(\ln(\varepsilon^{-1}))$ 

#### RATES OF CONVERGENCE

 $\Lambda$  If a sequence  $\{x^{(n)}\}\subset\mathbb{R}^n$  converging to a point  $x_*\in\mathbb{R}^n$ , then

$$\lim_{n \to \infty} \left\| x^{(n)} - x_* \right\| = 0$$

- \$\rightarrow\$ For a convergent sequence, we can talk about rate of convergence
  - ▶ The convergence is **linear** if there exists a  $\theta \in (0,1)$  such that

$$||x^{(n+1)} - x_*|| \le \theta ||x^{(n)} - x_*||$$

for all n sufficiently large

▶ The convergence is **superlinear** if

$$\lim_{n \to \infty} \frac{\|x^{(n+1)} - x_*\|}{\|x^{(n)} - x_*\|} = 0$$

▶ The convergence is **quadratic** if there exists a C > 0 such that

$$||x^{(n+1)} - x_*|| \le C ||x^{(n)} - x_*||^2$$

for all n sufficiently large

#### RATE OF CONVERGENCE (CONTD.)

 $\clubsuit$  Recall: For the gradient descent algorithm to minimize a strongly convex  $\beta$ -smooth function, we had

$$\left\| x^{(n+1)} - x_* \right\| \le \left( \frac{1}{1 + 2\delta\lambda} \right)^{\frac{1}{2}} \left\| x^{(n)} - x_* \right\|$$

- ♣ Hence the convergence here is linear
- A Recall: For the Newton's algorithm to minimize a smooth strongly convex function, we had

$$||x^{(n+1)} - x_*|| \le \frac{\gamma}{2\lambda} ||x^{(n)} - x_*||^2$$

♣ Hence the convergence here is quadratic

#### LINE SEARCH ALGORITHMS

- $\clubsuit$  Start with an initial vector  $x^{(0)} \in \mathbb{R}^n$  and a direction  $p^{(0)} \in \mathbb{R}^n$
- $\clubsuit$  Find the next iterate  $x^{(1)}$  along the line  $x^{(0)} + \alpha p^{(0)}$  with  $\alpha > 0$  s.t.

$$f(x^{(1)}) \le f(x^{(0)})$$

- $\clubsuit$  At the point  $x^{(1)}$ , pick a new direction  $p^{(1)} \in \mathbb{R}^n$
- $\clubsuit$  Find the next iterate  $x^{(2)}$  along the line  $x^{(1)} + \alpha p^{(1)}$  with  $\alpha > 0$  s.t.

$$f(x^{(2)}) \le f(x^{(1)})$$

- General principle of line search algorithms:
  - At the current iterate  $x^{(n)}$ , choose a direction  $p^{(n)}$
  - Pick the next iterate  $x^{(n+1)}$  along the line  $x^{(n)} + \alpha p^{(n)}$  with  $\alpha > 0$  such that

$$f(x^{(n+1)}) \le f(x^{(n)})$$

At each iteration step, we may perform a one-dimensional minimization problem:

$$\min_{\alpha>0} f(x^{(n)} + \alpha p^{(n)})$$

- ♣ But, in practice, we are content with finding a candidate that comes close to solving the above one-dimensional problem
- $\clubsuit$  The direction  $p^{(n)}$  is referred to as the SEARCH DIRECTION
- Recall the steepest descent algorithm:

$$x^{(n+1)} = x^{(n)} - \delta \nabla f(x^{(n)})$$

 $\clubsuit$  So, here the search direction at the  $n^{\text{th}}$  iteration step is

$$p^{(n)} = -\nabla f(x^{(n)})$$

 $\clubsuit$  At the iterate  $x^{(n)}$  and for any search direction  $p^{(n)}$ , we have

$$\begin{split} f(x^{(n)} + \alpha p^{(n)}) &= f(x^{(n)}) + \alpha \left\langle \nabla f(x^{(n)}), p^{(n)} \right\rangle \\ &+ \frac{\alpha^2}{2} \left\langle \nabla^2 f(x^{(n)} + s p^{(n)}) p^{(n)}, p^{(n)} \right\rangle \end{split}$$

for some  $s \in (0, \alpha)$ , thanks to Taylor's theorem.

 $\clubsuit$  Define a function  $g:[0,\infty)\to\mathbb{R}$  as follows:

$$g(\alpha) := f(x^{(n)} + \alpha p^{(n)}) \quad \text{for } \alpha \in [0, \infty).$$

♣ Observe that

$$g'(0) = \left\langle \nabla f(x^{(n)}), p^{(n)} \right\rangle$$

 $\clubsuit$  That is, the rate of change of f at the point  $x^{(n)}$  in the direction  $p^{(n)}$  is given by

$$\left\langle \nabla f(x^{(n)}), p^{(n)} \right\rangle$$

 $\clubsuit$  If we are interested in finding a unit direction of maximum decrease at the point  $x^{(n)}$ , we should understand

$$\min_{p \in \mathbb{R}^n, \|p\| = 1} \left\langle \nabla f(x^{(n)}), p \right\rangle$$

Recall that, if  $\theta_n$  denotes the angle between  $\nabla f(x^{(n)})$  and p, then  $\left\langle \nabla f(x^{(n)}), p \right\rangle = \|p\| \left\| \nabla f(x^{(n)}) \right\| \cos(\theta_n) = \left\| \nabla f(x^{(n)}) \right\| \cos(\theta_n)$ 

♣ So, the minimum possible value of  $\langle \nabla f(x^{(n)}), p \rangle$  is obtained when  $\cos(\theta_n) = -1$ 

 $\clubsuit$  Observe that the unit vector p which realises that is

$$p = -\frac{\nabla f(x^{(n)})}{\|\nabla f(x^{(n)})\|}$$

♣ We have seen that steepest descent is a line search algorithm where we take the search direction

$$p^{(n)} = -\nabla f(x^{(n)})$$

A Taylor's theorem says

$$f(x^{(n)} + \alpha p^{(n)}) = f(x^{(n)}) + \alpha \left\langle \nabla f(x^{(n)}), p^{(n)} \right\rangle + \frac{\alpha^2}{2} \left\langle \nabla^2 f(x^{(n)} + sp^{(n)}) p^{(n)}, p^{(n)} \right\rangle$$

 $\clubsuit$  Hence, if we take  $0 < \alpha \ll 1$ , and if we ensure that

$$\left\langle \nabla f(x^{(n)}), p^{(n)} \right\rangle < 0$$

then we find that  $f(x^{(n+1)}) < f(x^{(n)})$ 

 $\clubsuit$  Any such direction  $p^{(n)}$  is referred to as DESCENT DIRECTION

 $\clubsuit$  For any search direction p, we have by Taylor's theorem:

$$f(x^{(n)} + p) = f(x^{(n)}) + \left\langle \nabla f(x^{(n)}), p \right\rangle + \frac{1}{2} \left\langle \nabla^2 f(x^{(n)} + sp)p, p \right\rangle$$
 for some  $s \in (0, 1)$ .

- $\clubsuit$  Let us assume that  $\nabla^2 f(x^{(n)} + sp) \approx \nabla^2 f(x^{(n)})$
- ♣ Hence we obtain

$$f(x^{(n)} + p) \approx f(x^{(n)}) + \left\langle \nabla f(x^{(n)}), p \right\rangle + \frac{1}{2} \left\langle \nabla^2 f(x^{(n)})p, p \right\rangle =: F(p)$$

- $\clubsuit$  Observe that F is a quadratic function in p
- $\clubsuit$  If  $\nabla^2 f$  is positive definite, then F(p) has a unique global minimum
- $\clubsuit$  Recall: that global minimizer  $p_*$  is a critical point of F, i.e.

$$\nabla F(p_*) = 0 \implies p_* = -\left(\nabla^2 f(x^{(n)})\right)^{-1} \nabla f(x^{(n)})$$

♣ This is the search direction in Newton's algorithm

- A Newton's algorithm is also a line search algorithm
- ♣ The search direction in Newton's algorithm is

$$p^{(n)} = -\left(\nabla^2 f(x^{(n)})\right)^{-1} \nabla f(x^{(n)})$$

 $\clubsuit$  If  $\nabla^2 f$  is strictly positive definite, then

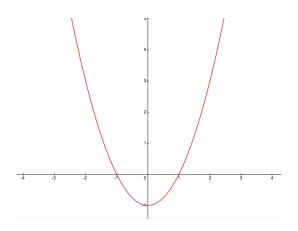
$$\left\langle \nabla f(x^{(n)}), p^{(n)} \right\rangle = -\left\langle \nabla f(x^{(n)}), \left( \nabla^2 f(x^{(n)}) \right)^{-1} \nabla f(x^{(n)}) \right\rangle < 0$$

 $\clubsuit$  Hence the above  $p^{(n)}$  is a descent direction

#### AN ILLUSTRATIVE EXAMPLE

 $\clubsuit$  Consider the function  $f: \mathbb{R} \to \mathbb{R}$  defined as follows:

$$f(x) := x^2 - 1$$
 for  $x \in \mathbb{R}$ 



 $\clubsuit$  The point x=0 is the minimizer and f(0)=-1

### AN ILLUSTRATIVE EXAMPLE (CONTD.)

$$Arr ext{Take } x^{(0)} = -2 \text{ and } p^{(0)} = 1$$

• Note that 
$$f(x^{(0)}) = 3$$

**.** Take 
$$\alpha_0 = 2 + \sqrt{3}$$
 so that  $x^{(1)} = x^{(0)} + \alpha_0 p^{(0)} = \sqrt{3}$ 

$$Arr$$
 Note that  $f(x^{(1)}) = 2$ 

♣ Take 
$$p^{(1)} = -1$$

**.** Take 
$$\alpha_1 = \sqrt{3} + \sqrt{2}$$
 so that  $x^{(2)} = x^{(1)} + \alpha_1 p^{(1)} = -\sqrt{2}$ 

$$Arr$$
 Note that  $f(x^{(2)}) = 1$ 

• Take 
$$p^{(2)} = 1$$

**A** Take 
$$\alpha_2 = \sqrt{2} + \sqrt{\frac{5}{3}}$$
 so that  $x^{(3)} = x^{(2)} + \alpha_2 p^{(2)} = \sqrt{\frac{5}{3}}$ 

• Note that 
$$f(x^{(3)}) = \frac{2}{3}$$

♣ Take 
$$p^{(3)} = -1$$

**.** Take 
$$\alpha_3 = \sqrt{\frac{5}{3}} + \sqrt{\frac{3}{2}}$$
 so that  $x^{(4)} = x^{(3)} + \alpha_3 p^{(3)} = -\sqrt{\frac{3}{2}}$ 

• Note that 
$$f(x^{(4)}) = \frac{1}{2}$$

♣ Observe that

$$f(x^{(0)}) > f(x^{(1)}) > f(x^{(2)}) > f(x^{(3)}) > f(x^{(4)})$$

#### AN ILLUSTRATIVE EXAMPLE (CONTD.)

 $\clubsuit$  We can thus build a minimizing sequence  $x^{(n)}$  such that

$$f(x^{(n)}) = \frac{2}{n}$$
 for  $n = 1, 2, ...$ 

- A But the limiting function value for this sequence is zero
- $\clubsuit$  Recall that the minimum value of the objective function is -1
- A This illustrates the possibility of a general line search algorithm
  - $\triangleright$  leading to insufficient reduction in f in each iteration
  - $\blacktriangleright$  failing to converge to the minimizer of f
- $\clubsuit$  The root cause for this behaviour stems from the choice of step lengths  $\alpha_n$  in each iteration step
- ♣ Here we encounter certain sufficient decrease conditions

## END OF LECTURE 5 THANK YOU FOR YOUR ATTENTION