

Syntactic Analysis

Experiment 03-01: Implementing context free grammar in C

$$S \rightarrow cAb$$
$$A \rightarrow ad \mid a$$

Experiment 03-02: Syntactic Analysis of simple arithmetic expressions.

A grammar for the concrete syntax of simple arithmetic expressions.

Production rules are:-

$$\langle \text{Exp} \rangle ::= \langle \text{Term} \rangle + \langle \text{Term} \rangle \mid \langle \text{Term} \rangle - \langle \text{Term} \rangle \mid \langle \text{Term} \rangle$$
$$\langle \text{Term} \rangle ::= \langle \text{Factor} \rangle * \langle \text{Factor} \rangle \mid \langle \text{Factor} \rangle / \langle \text{Factor} \rangle \mid \langle \text{Factor} \rangle$$
$$\langle \text{Factor} \rangle ::= (\langle \text{Exp} \rangle) \mid \text{ID} \mid \text{NUM}$$
$$\text{ID} ::= [\text{a-z}]$$
$$\text{NUM} ::= [0-9]$$

Non-terminal symbols:

$\langle \text{Exp} \rangle$, $\langle \text{Term} \rangle$, $\langle \text{Factor} \rangle$

Terminal symbols:

$+$, $-$, $*$, $/$, $($, $)$, a , b , c , 1 , 2 , 3 ...

Start symbol:

$\langle \text{Exp} \rangle$

Note the following are meta symbols:

$::=$ (read it as "can be replaced by")

\rightarrow is alternative notation for $::=$,

\mid (read it as "or by")

e.g. An $\langle \text{Exp} \rangle$ can be replaced by $\langle \text{Term} \rangle + \langle \text{Term} \rangle$ or by $\langle \text{Term} \rangle - \langle \text{Term} \rangle$.

Tasks

Task: 01 Consider the Context Free Grammar G. You have the set of rules as the following:

$$S \rightarrow S_1 \mid S_2$$

$$S_1 \rightarrow 0S_11 \mid \epsilon$$

$$S_2 \rightarrow 1S_20 \mid \epsilon$$

Implement the grammar that will accept strings accepted by this language $\{0^n1^n \mid n \geq 0\} \cup \{1^n0^n \mid n \geq 0\}$.

Task 02: Revise the grammar of simple arithmetic expressions to accept the strings like,

$$(abc - b2) * (-c * d3) - g$$

Note: You have to use the previous lab materials to implement it.

Task 03: Implement the following grammar in C.

$$\text{statement} \rightarrow \text{assign_stat} \mid \text{decision_stat} \mid \text{looping_stat}$$

$$\text{assign_stat} \rightarrow \text{id} = \text{expn}$$

$$\text{expn} \rightarrow \text{simple-expn} \text{ eprime}$$

$$\text{epime} \rightarrow \text{relop} \text{ simple-expn} \mid \epsilon$$

$$\text{decision-stat} \rightarrow \text{if (expn) stat dprime}$$

$$\text{dprime} \rightarrow \text{else stat} \mid \epsilon$$

$$\text{looping-stat} \rightarrow \text{while (expn) stat} \mid \text{for (assign_stat ; expn ; assign_stat) stat}$$

$$\text{relop} \rightarrow == \mid != \mid <= \mid >= \mid > \mid <$$

Note: The grammar related to simple expression is described in experiment 02.
