

There are a total of four problems. You have to solve **all** of them.

**Problem 1 (CO1): DFA and Regular Languages (10 points)**

Let  $\Sigma = \{0, 1\}$ . Consider the following languages over  $\Sigma$ .

$$L_1 = \{w : w \text{ starts with either } 01 \text{ or } 10\}$$

$$L_2 = \{w : w \text{ does not start with } 11\}$$

$$L_3 = \{w : \text{the length of } w \text{ is at least two}\}$$

Now solve the following problems.

- (a) Give the state diagram for a DFA that recognizes  $L_1$ . (3 points)
- (b) Give the state diagram for a DFA that recognizes  $L_2$ . (3 points)
- (c) Give the state diagram for a DFA that recognizes  $L_3$ . (2 points)
- (d) Give the state diagram for a DFA that recognizes  $\overline{L_1} \cap L_2 \cap L_3$  using only four states. Here  $\overline{L}$  denotes the complement of the language  $L$  i.e.,  $\overline{L} = \Sigma^* - L$ . (2 points)

**Problem 2 (CO1): Regular Expressions (10 points)**

Let  $\Sigma = \{0, 1\}$ . Consider the following pair of languages over  $\Sigma$ .

$$L_1 = \{w : w \text{ contains } 11 \text{ as a substring}\}$$

$$L_2 = \{w : w \text{ contains } 10 \text{ as a substring}\}$$

Now solve the following problems.

- (a) Write down a regular expression for the language  $L_1$ . (2 points)
- (b) Write down a regular expression for the language  $L_2$ . (2 points)
- (c) Your friend wants a regular expression for the language  $\overline{L_1 \cap L_2}$  where  $\overline{L}$  denotes the complement of the language  $L$  i.e.,  $\overline{L} = \Sigma^* - L$ . He wants your help. You tell him to make use of the fact  $\overline{L_1 \cap L_2} = \overline{L_1} \cup \overline{L_2}$ .
  - (i) Write down a regular expression for the language  $\overline{L_1}$ . (2 points)
  - (ii) Write down a regular expression for the language  $\overline{L_2}$ . (2 points)
  - (iii) Using the fact above, write down a regular expression for the language  $\overline{L_1 \cap L_2}$ . (2 points)

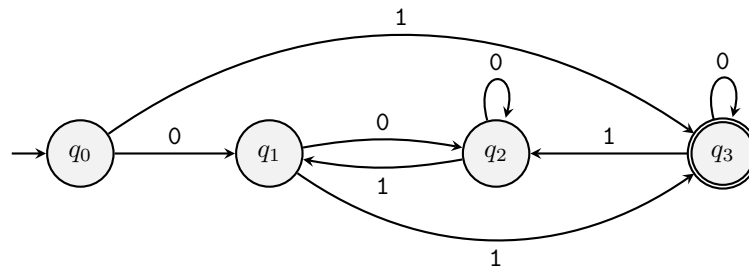
**Problem 3 (CO1): Converting Regular Expressions to NFAs (10 points)**

Convert the following regular expression into an equivalent NFA. Note that  $R_1 + R_2$  is the same as  $R_1 \cup R_2$ .

$$c + (ab^* + (bb^*c)^*)^*$$

**Problem 4 (CO1): Converting Finite Automata to Regular Expressions (10 points)**

Convert the following DFA into an equivalent regular expression using the state elimination method. First eliminate  $q_1$ , then  $q_2$  and finally  $q_3$ . You must show work.



After you are done with the test, please indicate where you stand on the smiley face spectrum.

