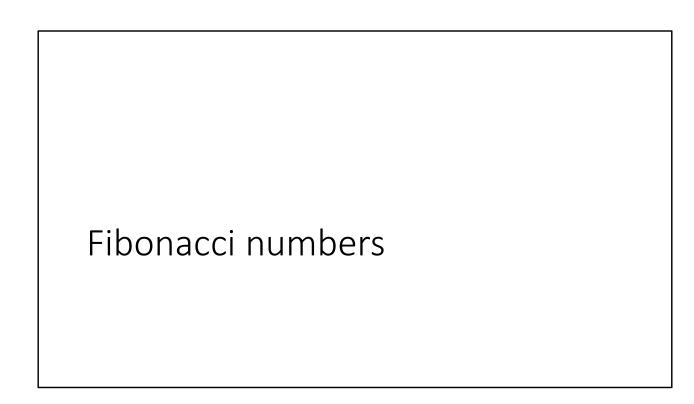


Today's class

- Fibonacci numbers
- Search algorithms
 - Linear, Binary
- Towers of Hanoi



• Recursive definition of Fibonacci numbers

$$f(n) = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ f(n-1) + f(n-2) & n \ge 2 \end{cases}$$

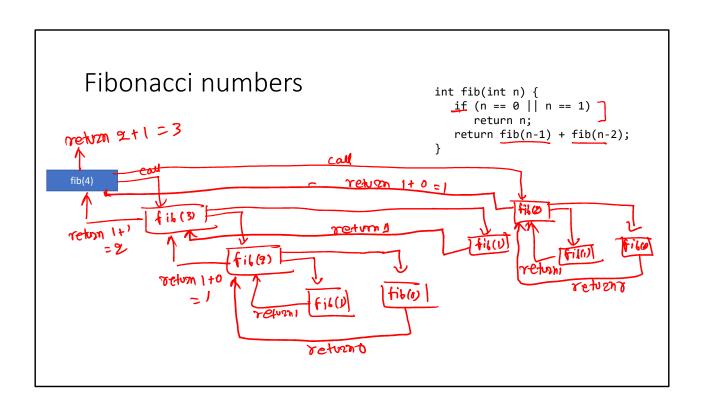
f(1) = 0 f(1) = 1 f(3) = 1 f(3) = 2 f(4) = 3

01173

• Recursive definition of Fibonacci numbers

$$f(n) = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ f(n-1) + f(n-2) & n \ge 2 \end{cases}$$

```
int fib(int n) {
   if (n == 0 || n == 1)
      return n;
   return fib(n-1) + fib(n-2);
}
```

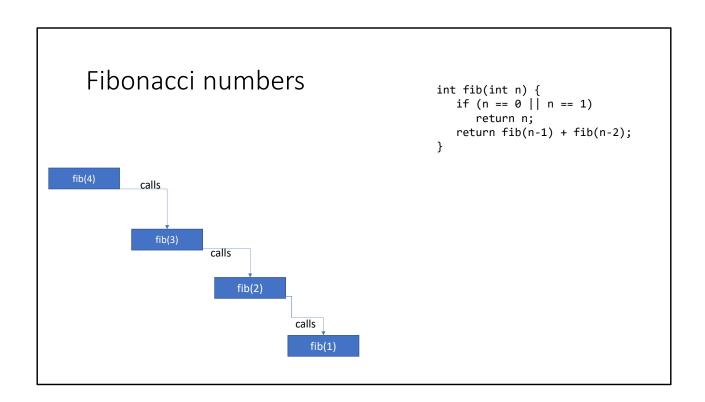


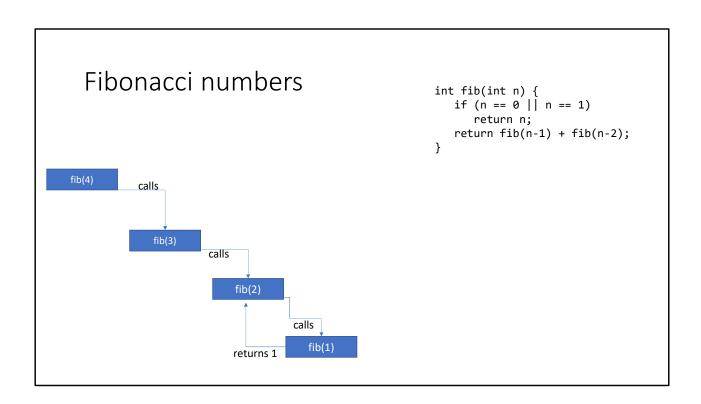
```
int fib(int n) {
   if (n == 0 || n == 1)
      return n;
   return fib(n-1) + fib(n-2);
}
```

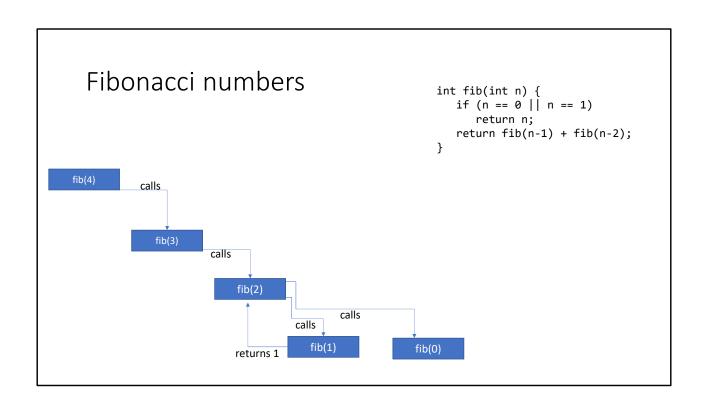
fib(4)

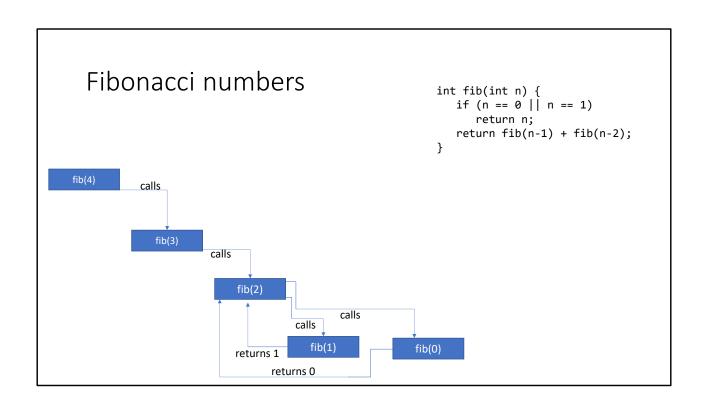
Fibonacci numbers int fib(int n) { if (n == 0 || n == 1) return n; return fib(n-1) + fib(n-2); } fib(4) calls

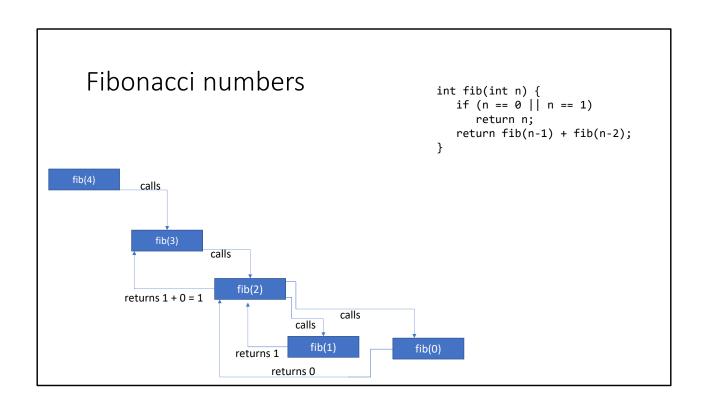
Fibonacci numbers int fib(int n) { if (n == 0 || n == 1) return n; return fib(n-1) + fib(n-2); } fib(4) calls fib(2)

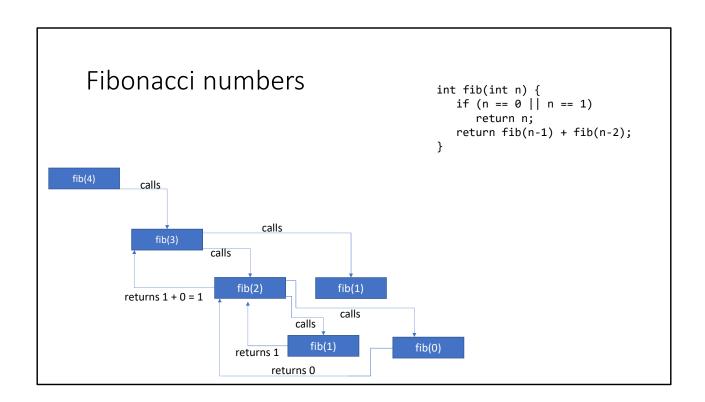


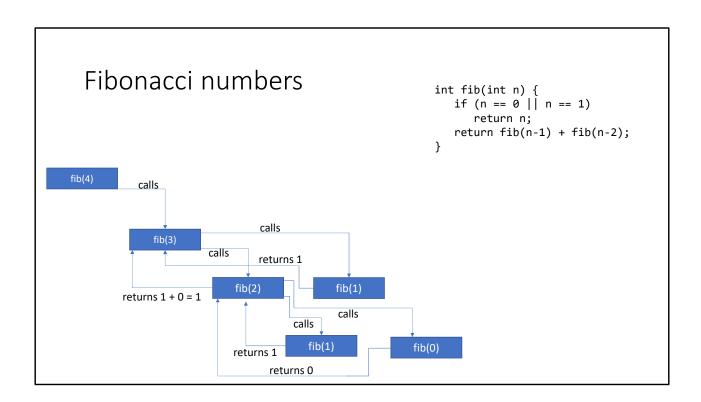


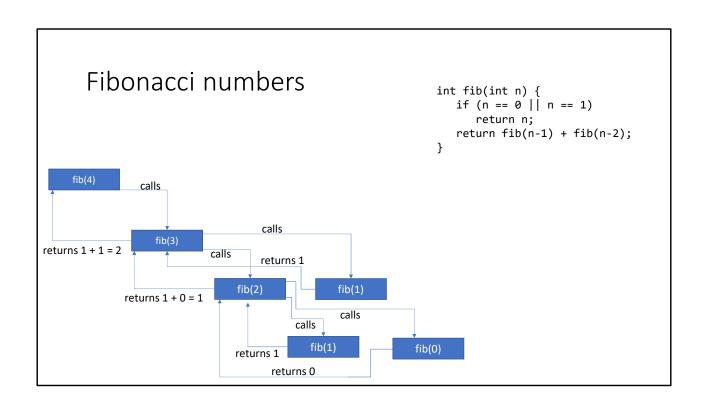


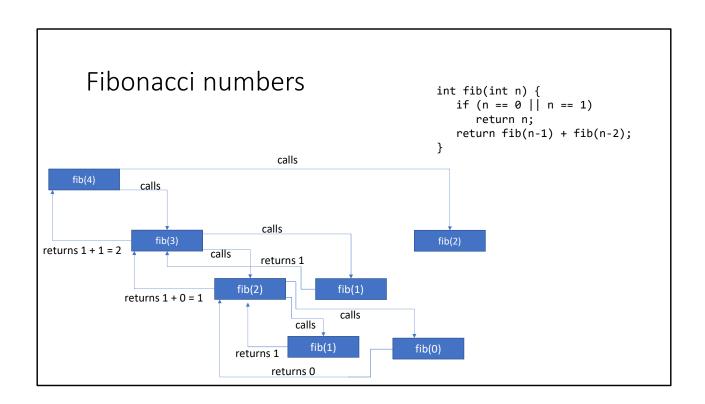


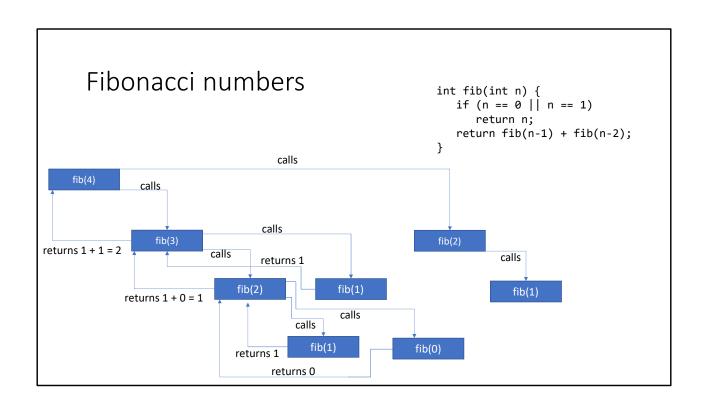


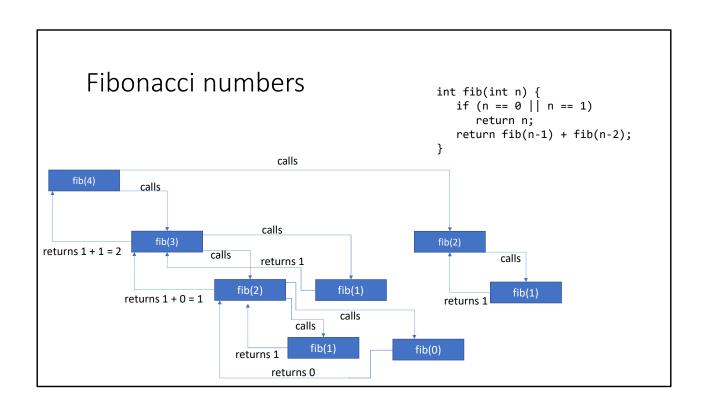


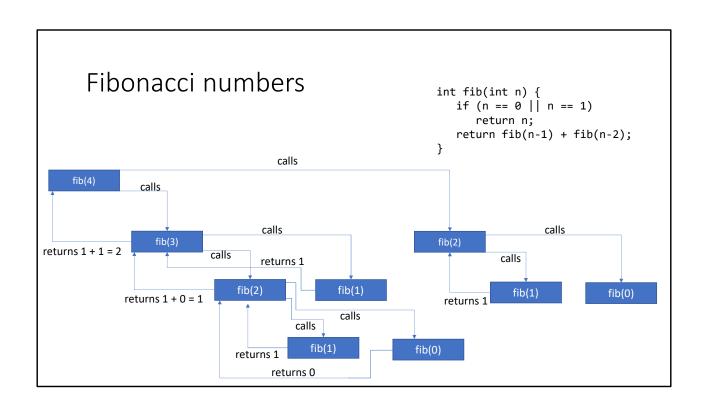


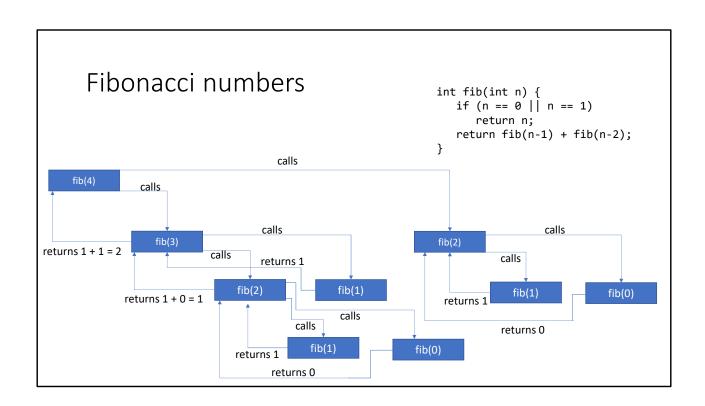


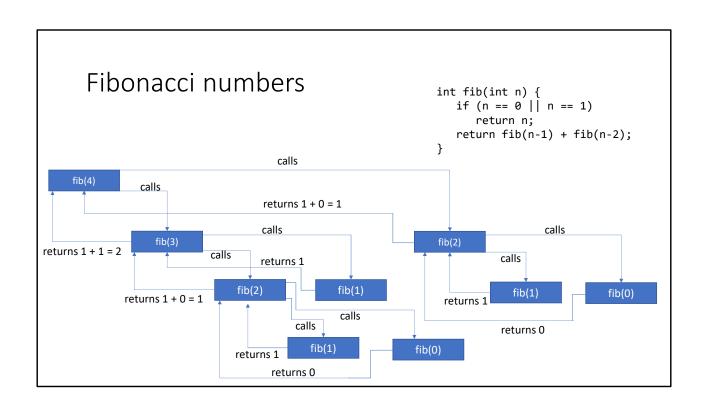


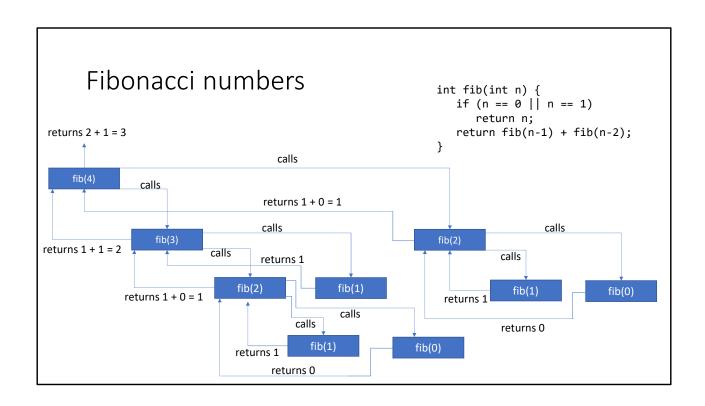


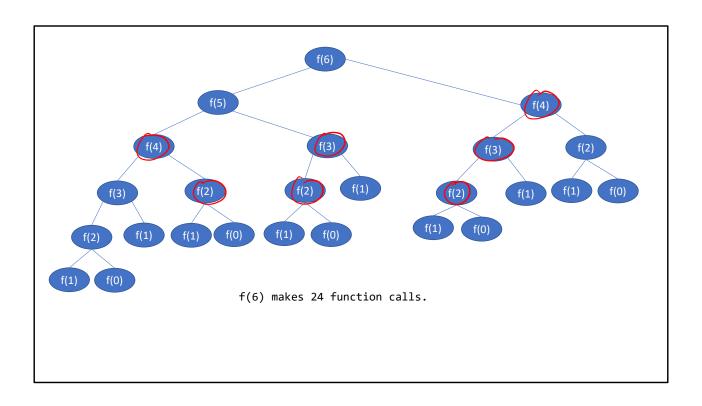


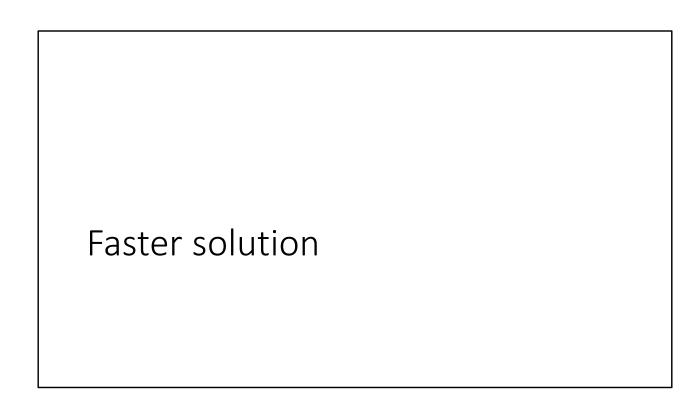


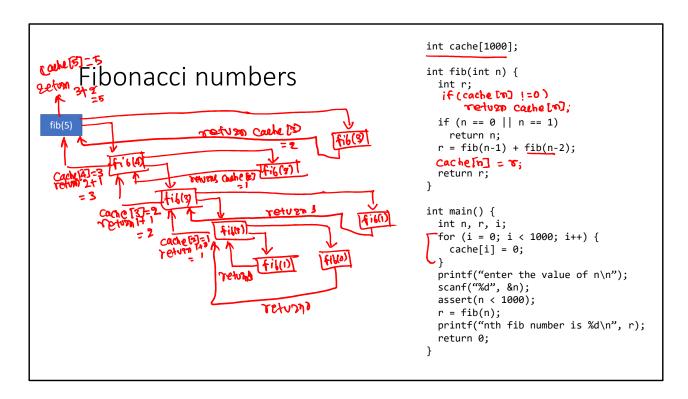








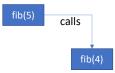




A cache can be used to reduce the number of repeating computations. We can save the result of a computation that might be used later in a cache (can be implemented using array / other data structures) and retrieve the result from the cache when needed, thus eliminating the need to recompute the result. However, such an approach may take more memory than the algorithm that doesn't use a cache.

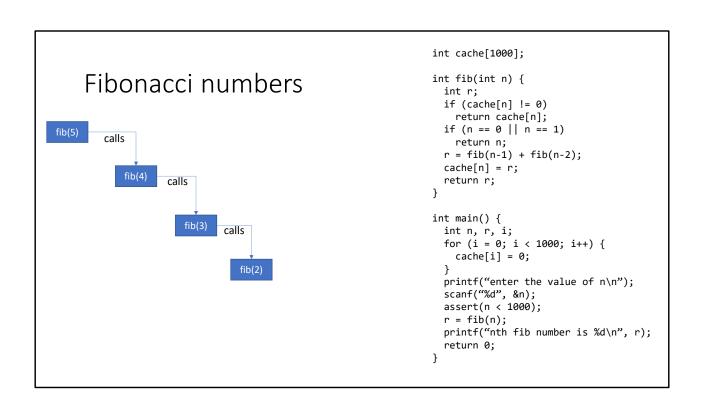
fib(5)

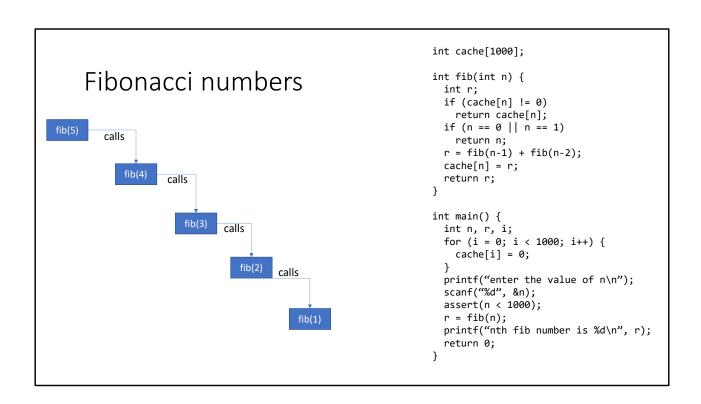
```
int cache[1000];
int fib(int n) {
  int r;
  if (cache[n] != 0)
    return cache[n];
  if (n == 0 || n == 1)
    return n;
  r = fib(n-1) + fib(n-2);
  cache[n] = r;
  return r;
int main() {
  int n, r, i;
for (i = 0; i < 1000; i++) {
  cache[i] = 0;
  printf("enter the value of n\n");
scanf("%d", &n);
assert(n < 1000);</pre>
  r = fib(n);
  printf("nth fib number is %d\n", r);
  return 0;
```

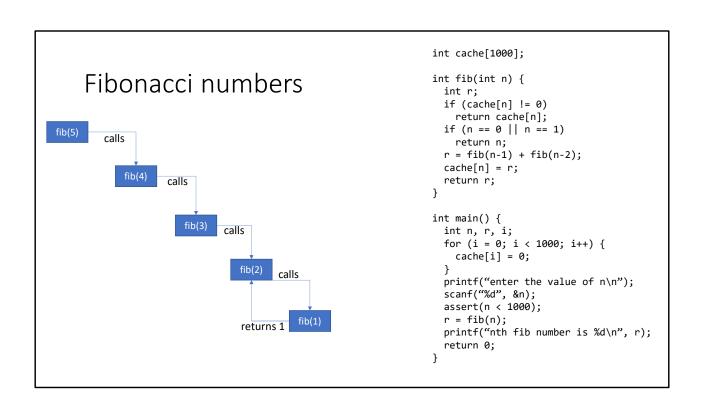


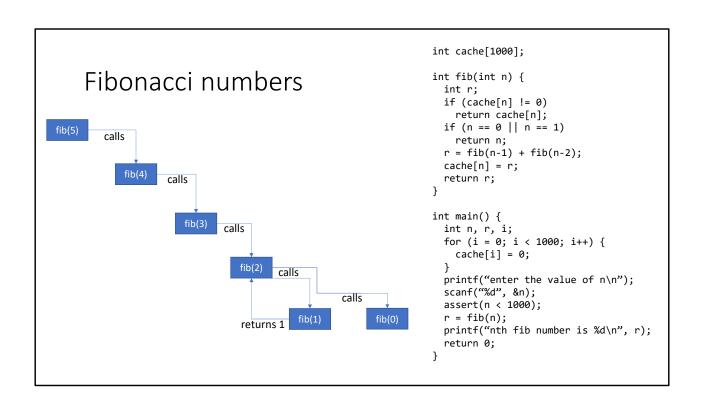
```
int cache[1000];
int fib(int n) {
  int r;
  if (cache[n] != 0)
    return cache[n];
  if (n == 0 || n == 1)
    return n;
  r = fib(n-1) + fib(n-2);
  cache[n] = r;
  return r;
int main() {
  int n, r, i;
for (i = 0; i < 1000; i++) {
    cache[i] = 0;
  printf("enter the value of n\n");
scanf("%d", &n);
assert(n < 1000);</pre>
  r = fib(n);
  printf("nth fib number is %d\n", r);
  return 0;
```

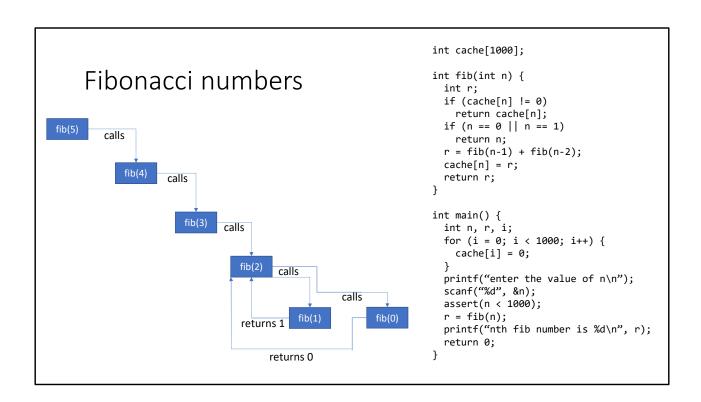
int cache[1000]; Fibonacci numbers int fib(int n) { int r; if (cache[n] != 0) return cache[n]; if (n == 0 || n == 1) calls return n; r = fib(n-1) + fib(n-2);cache[n] = r;fib(4) calls return r; int main() { fib(3) int n, r, i; for (i = 0; i < 1000; i++) { cache[i] = 0;printf("enter the value of n\n"); scanf("%d", &n); assert(n < 1000);</pre> r = fib(n);printf("nth fib number is %d\n", r); return 0;

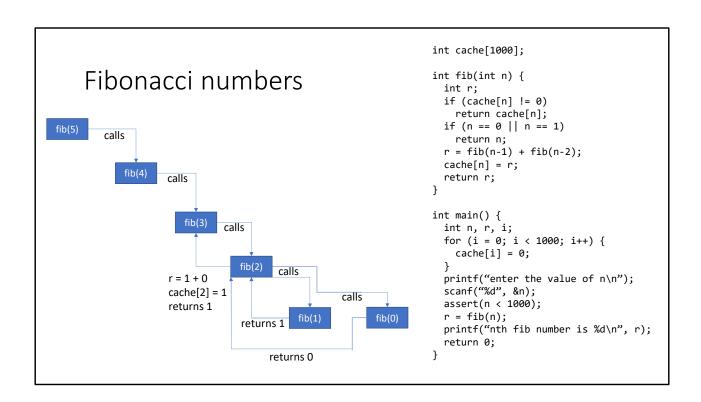


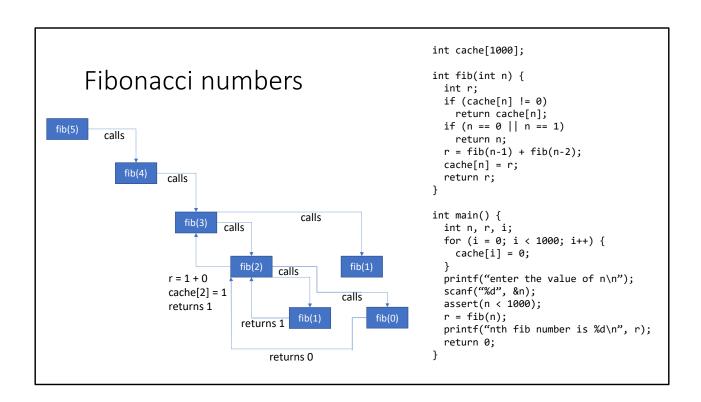


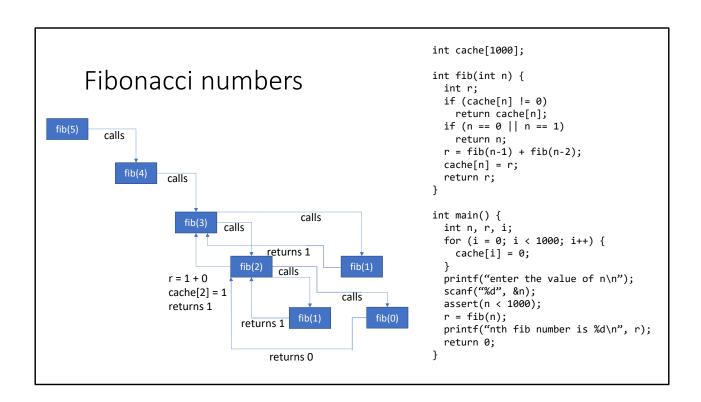


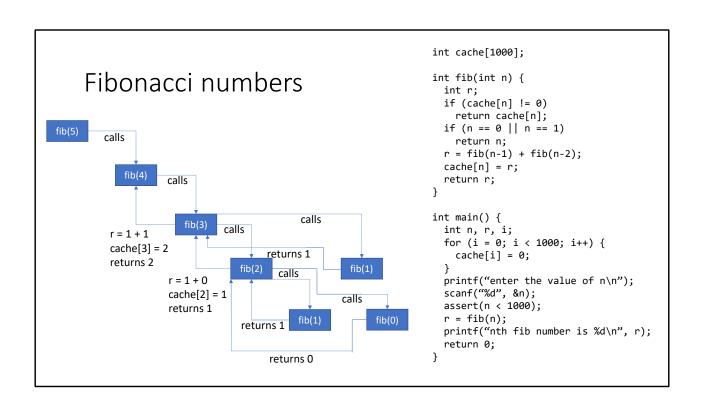


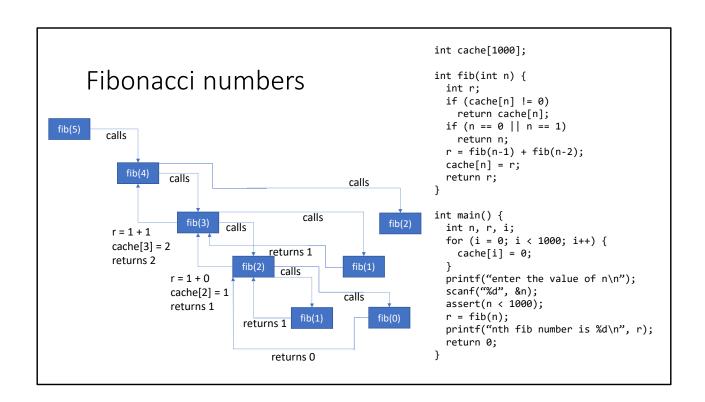


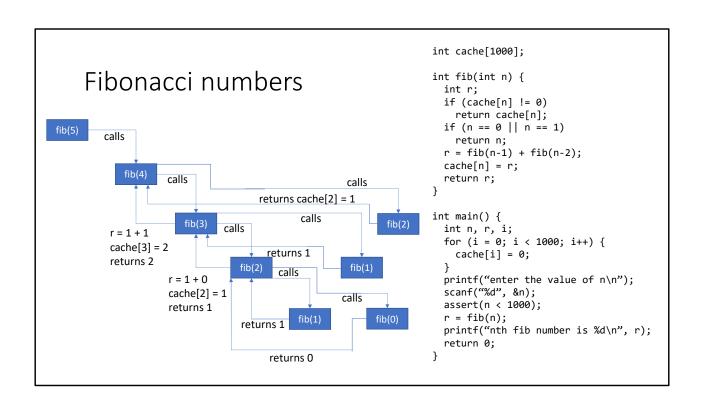


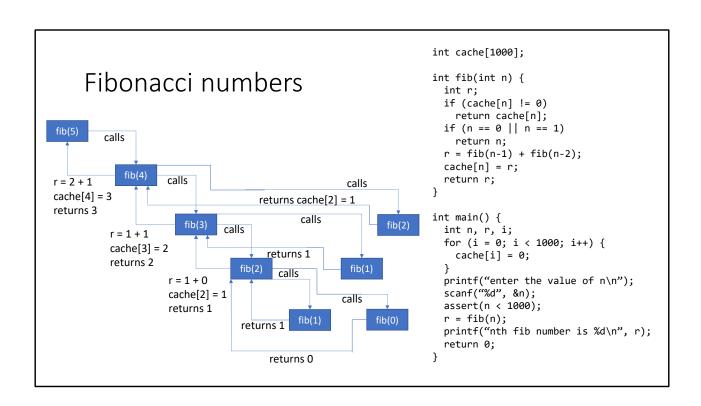


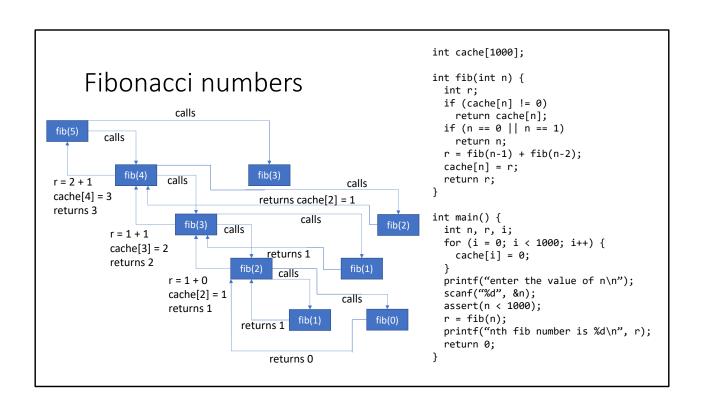


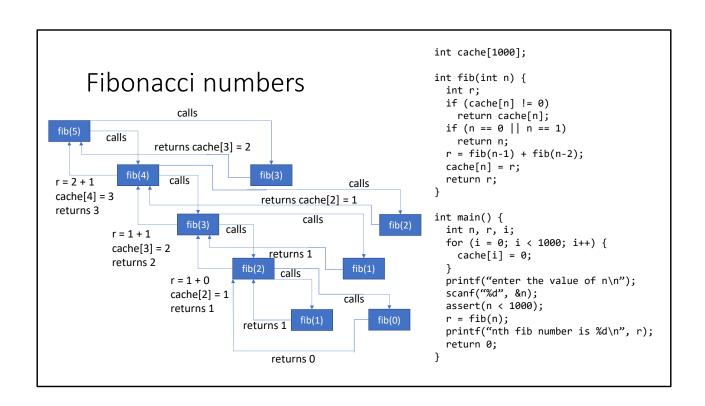


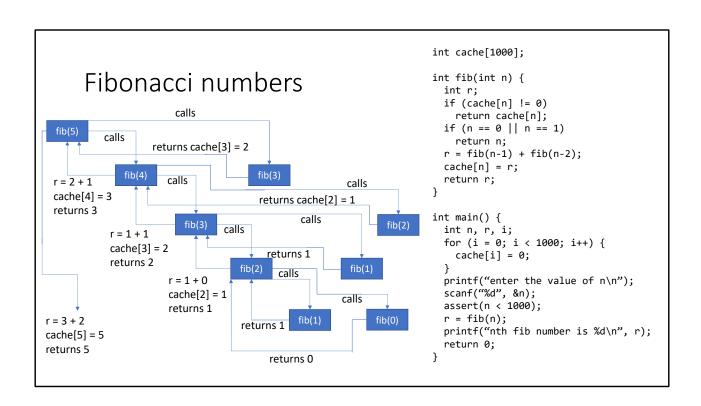








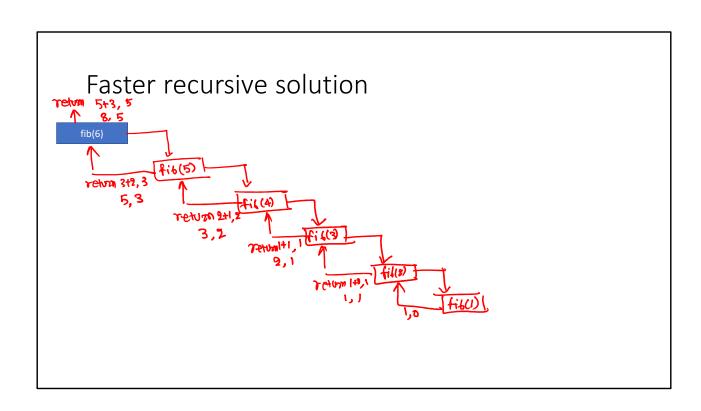


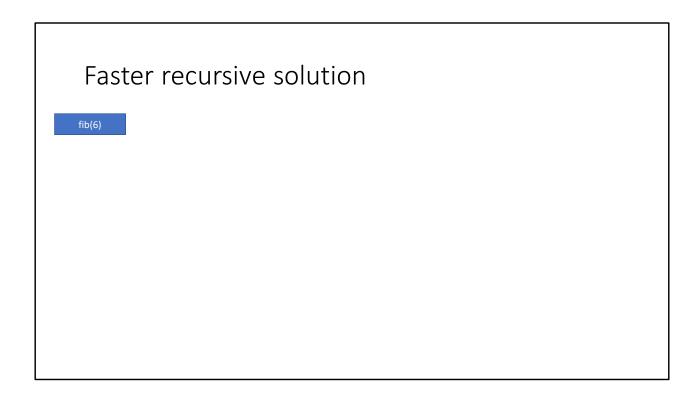


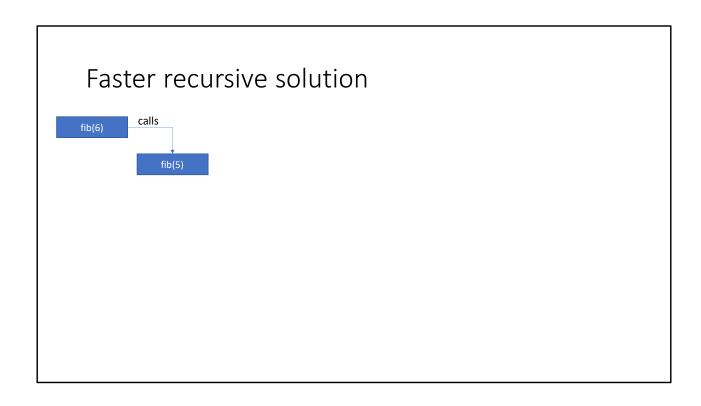
Faster recursive solution without caching

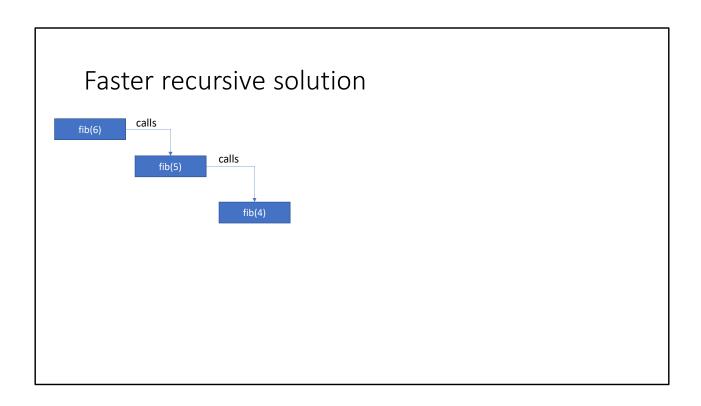
Faster recursive solution

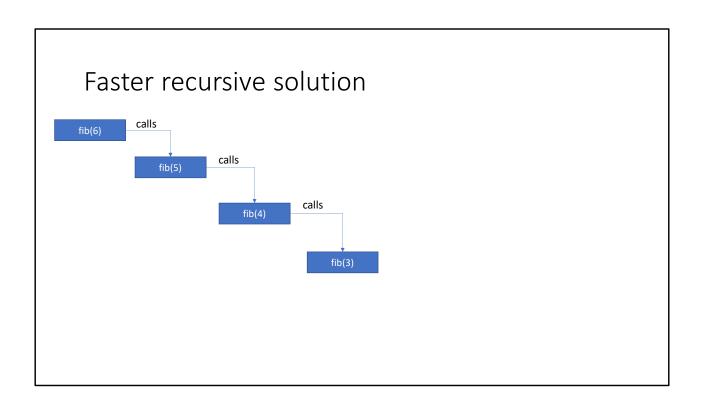
- Let's say fib(n) returns two values: fib(n) and fib(n-1), i.e.,
 - fib(5) returns fib(5) and fib(4)
 - fib(4) returns fib(4) and fib(3)
 - fib(3) returns fib(3) and fib(2)
 - fib(2) returns fib(2) and fib(1)
 - fib(1) returns 1 and 0
- How many function calls are needed to compute fib(6)?
 - The previous solution without caching requires 24 calls

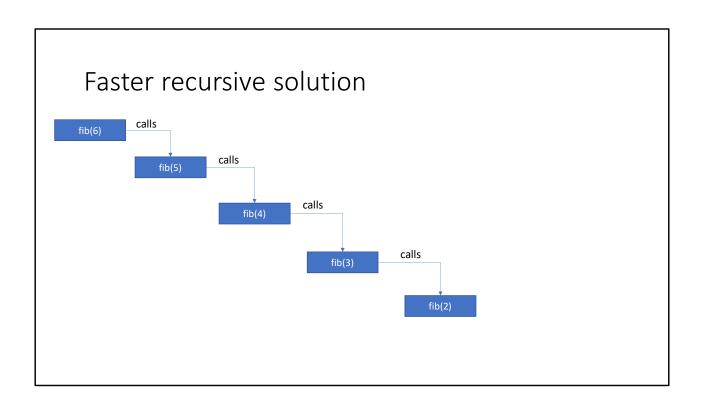


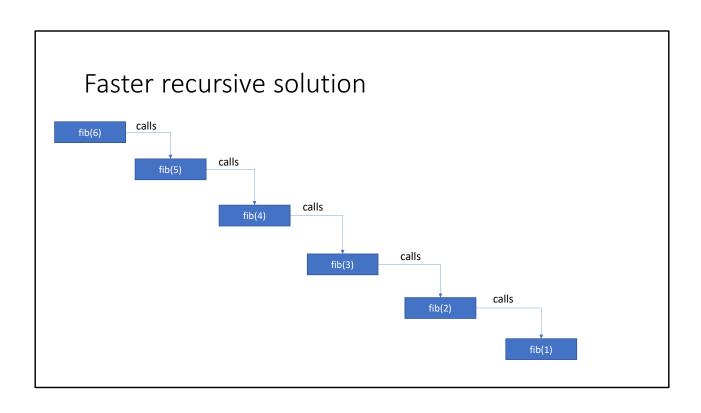


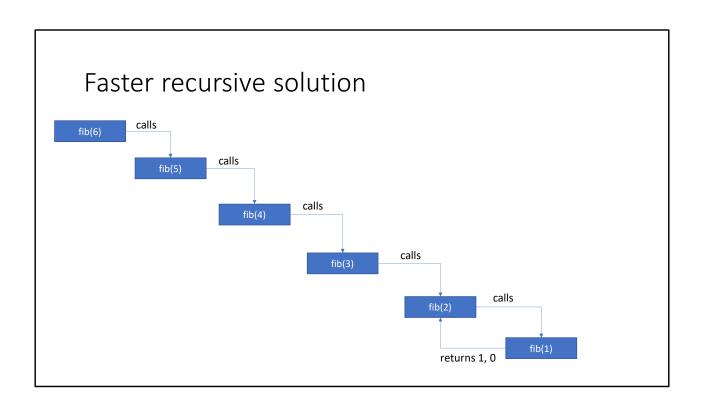


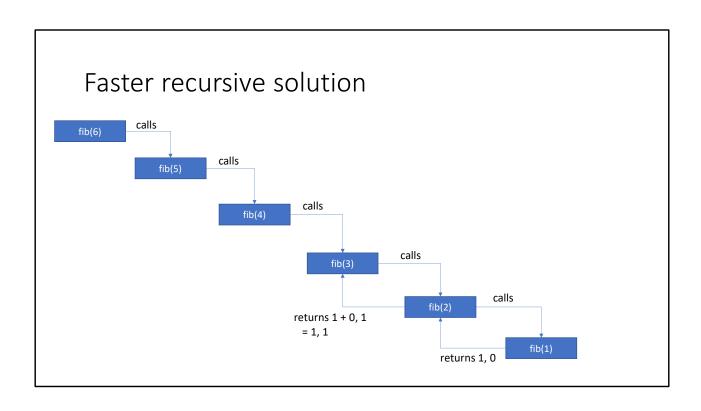


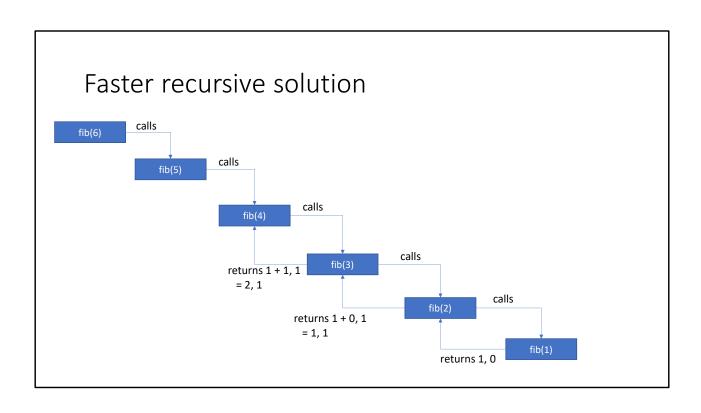


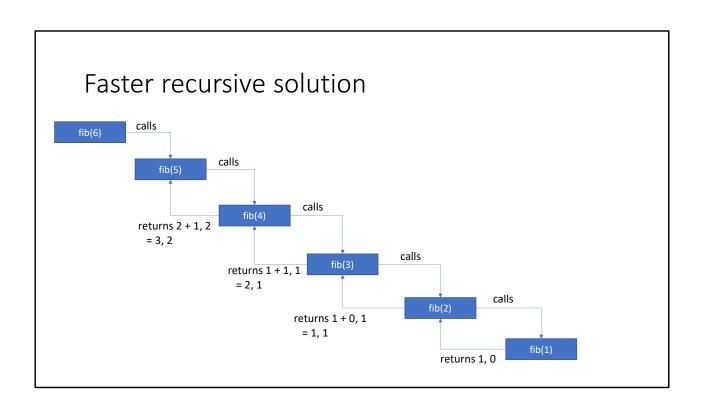


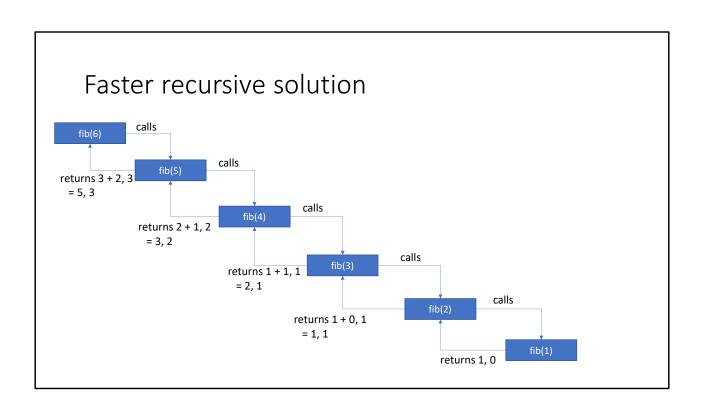


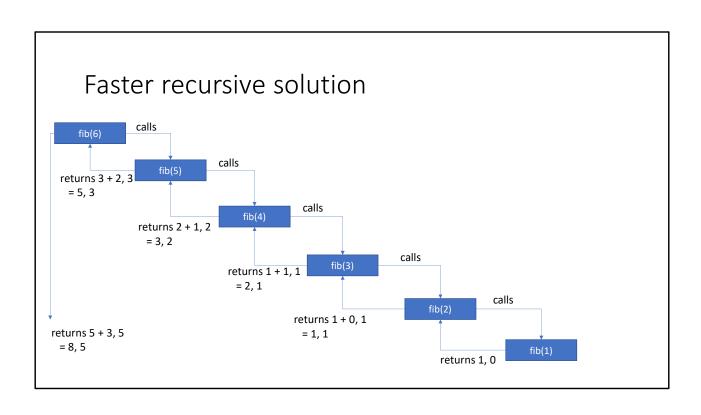












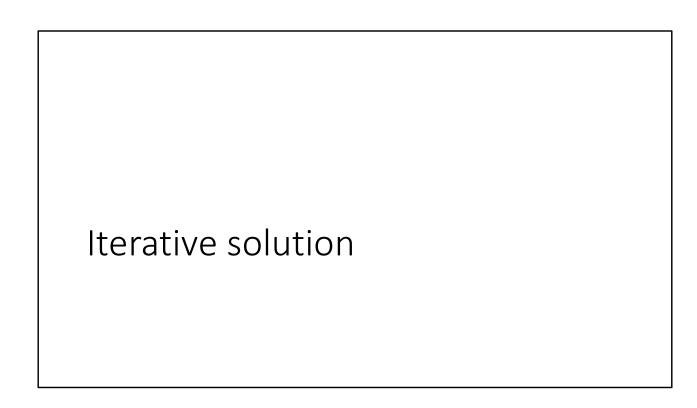
Faster recursive solution

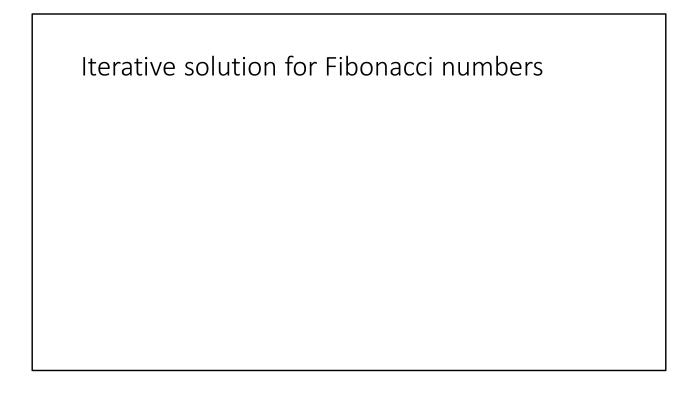
- Base case • if (n == 1) return (1, 0)
- Recursive step
 - Recursively call for n-1 to obtain (x, y) = (fib(n-1), fib(n-2))
 - return (x+y, x)

In this recursive algorithm, instead of just returning fib(n), the fib function returns two values: fib(n) and fib(n-1), resulting in an efficient algorithm that doesn't require multiple recursive calls or caches.

```
struct retual Ti
                                                                         into;
                                                 7.7 = 20;
Fibonacci numbers
 struct retval {
   int x;
   int y;
\};
 struct retval fib(int n) {
  struct retval ret;
   struct retval r;
   if (n == 1) { -
     r.x = 1;
     r.y = 0;
     return r;
   ret = fib(n - 1);_
   r.x = ret.x + ret.y;
   r.y = ret.x;
   return r;
 }
```

In C, "struct" is a way to define a new data type. A struct may contain multiple fields of possibly different types, including the struct type.





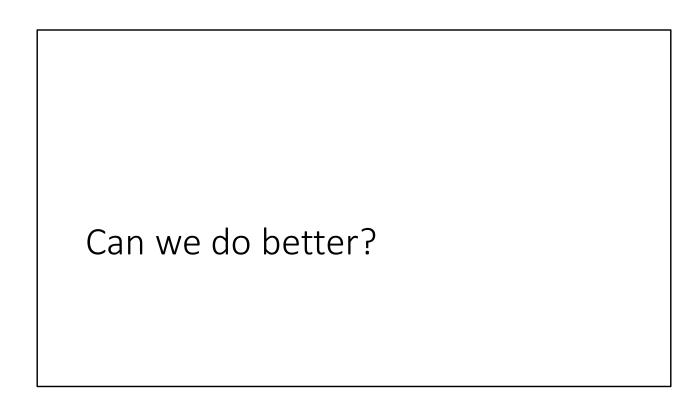
Iterative solution for Fibonacci numbers

```
Pprev=fib() prev=fib()
n= 4
                                           int fib(int n) {
1=2; 1<=4;
                                            if (n == 0 || n == 1) {
Tes = 1 +0 =1 fib(s)
                                              return n;
                          1=4; 144
 bbsen = fiff) = )
                          res = 1+2=3
 prev = file) = 1
                                            int prev = 1;
                           bbeen = 5
                                            int pprev = 0;
                            prev=3
                                            int res, i;
 1=3; i <=4
                             1:5
                                            for (i = 2; i <= n; i++) {
 rel= 1+1 =2
                                              res = prev + pprev;
                                              pprev = prev;
  ppren = 1
                                              prev = res;
 poer -2
                                            }
                                            return res;
  1-4
```

Iterative solution for Fibonacci numbers

```
iteration 3 : i = 4; i <= 5
compute fib(5):
                                                             int fib(int n) {
                           res = 2 + 1 = 3
                                                                if (n == 0 || n == 1) {
                          pprev = 2
prev = 1
                                                                  return n;
pprev = 0
                           prev = 3
                          i = 5
iteration 1:i=2;i \le 5
                                                                int prev = 1;
                          iteartion4: i = 5; i <= 5
res = 1 + 0 = 1
                                                                int pprev = 0;
pprev = 1
                          res = 3 + 2 = 5
                                                                int res, i;
                          pprev = 3
prev = 1
i = 3
                          prev = 5
                                                                for (i = 2; i <= n; i++) {
                          i = 6
                                                                  res = prev + pprev;
iteartion2: i = 3; i <= 5
                                                                  pprev = prev;
                          return res = 5
res = 1 + 1 = 2
                                                                  prev = res;
pprev = 1
                                                                }
prev = 2
                                                                return res;
i = 4
```

The iterative algorithm computes fib(2), fib(3), fib(4), ..., fib(n) in an iterative manner. It keeps track of the last two Fibonacci numbers needed to compute the next Fibonacci number.



• A better solution

$$\begin{bmatrix} f(n) \\ f(n-1) \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} f(n-1) \\ f(n-2) \end{bmatrix} = \begin{bmatrix} 1 & f(n-1) \\ 1 & f(n-1) \end{bmatrix} \begin{bmatrix} f(n-1) \\ 1 & f(n-2) \end{bmatrix}$$

$$= \begin{bmatrix} f(n-1) + f(n-2) \\ f(n-1) \end{bmatrix}$$

• A better solution

$$\begin{bmatrix} f(n) \\ f(n-1) \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} f(n-1) \\ f(n-2) \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^2 \begin{bmatrix} f(n-2) \\ f(n-3) \end{bmatrix}$$

• A better solution

$$\begin{bmatrix} f(n) \\ f(n-1) \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} f(n-1) \\ f(n-2) \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^2 \begin{bmatrix} f(n-2) \\ f(n-3) \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^3 \begin{bmatrix} f(n-3) \\ f(n-4) \end{bmatrix} = \dots = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1} \begin{bmatrix} f(1) \\ f(0) \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{n-1} \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

• Let
$$A = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} f(n) \\ f(n-1) \end{bmatrix} = A^{n-1} * \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

How to compute A^{n-1} fast?

```
Fibonacci numbers
// mul takes a matrix A and R as input
                                                               int fib(int n) {
// returns the result of A^n in R
void mul(int A[2][2], int R[2][2], int n) {
                                                                  if (n <= 1) {
 if (n == 1) {
                                                                    return n;
   R[0][0] = A[0][0]; R[0][1] = A[0][1]; R[1][0] = A[1][0]; R[1][1] = R[1][1];
                                                                  int A[2][2];
   return;
                                                                  int R[2][2];
if (n % 2 == 0) {
                                                                  A[0][0] = 1; A[0][1] = 1;
   mul(A, R, n/2);
    // mul2 takes two 2x2 matrices as input and
                                                                  A[1][0] = 1; A[1][1] = 0;
   // returns the multiplication in the first
   // matrix
                                                                  mul(A, R, n-1);
   mul2(R, R); // R \leftarrow R * R
                                                                  // R contains A<sup>n-1</sup>
                                                                  return R[0][0];
   mul(A, R, (n-1)/2);
   mul2(R, R); // R <- R * R
mul2(R, A); // R <- R * A
```

The fib routine initializes A (a 2x2 matrix) with the value discussed in the previous slide. The mul routine takes two 2x2 matrices, A and R; the value of n; and returns A^(n-1) in R. The mul2 routine takes two 2x2 matrices, P and Q; computes P x Q; and stores them in P before returning. The recursive algorithm to compute A^n is similar to the faster algorithm we discussed for computing x^n.

```
Fibonacci numbers
                                                       // mul takes a matrix A and R as input
                                                       // returns the result of A^n in R
                                                       void mul(int A[2][2], int R[2][2], int n) {
mul2 implementation?
                                                         if (n == 1) {
 mulz (int A [2] [2], int B[2] [2])
                                                           R[0][0] = A[0][0]; R[0][1] = A[0][1];
                                                           R[1][0] = A[1][0]; R[1][1] = R[1][1];
                                                           return;
 ind 七回到;
   +[0][0]- A[0][0]+ B[0][0]+ A[0][1], B[1][0] if (n % 2 == 0) {
mul(A, R, n/2);
                                                           // mul2 takes two 2x2 matrices as input and
                                                           // returns the multiplication in the first
                                                          // matrix
                                                          mul2(R, R); // R \leftarrow R * R
                                                         else {
                                                           mul(A, R, (n-1)/2);
                                                          mul2(R, R); // R <- R * R
mul2(R, A); // R <- R * A
                                                       }
```

Fibonacci numbers // mul takes a matrix A and R as input // returns the result of A^n in Rfib(31)—call void mul(int A[2][2], int R[2][2], int n) $\{$ if (n == 1) { R[0][0] = A[0][0]; R[0][1] = A[0][1];105) Som R[1][0] = A[1][0]; R[1][1] = R[1][1];return; mu((15) if (n % 2 == 0) { PERIR mul(A, R, n/2);// mul2 takes two 2x2 matrices as input and R=R+A JOHN BENE // returns the multiplication in the first mu (7) // matrix _mul2(R, R); // R <- R * R MU (3)else { mul(A, R, (n-1)/2); -mul2(R, R); // R <- R * R mul2(R, A); // R <- R * A R=RAA MUIN ceturn R=A

Fibonacci numbers

fib(31)

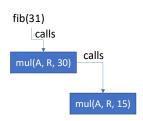
```
// mul takes a matrix A and R as input
// returns the result of A^n in R
void mul(int A[2][2], int R[2][2], int n) \{
  if (n == 1) {
    R[0][0] = A[0][0]; R[0][1] = A[0][1];
    R[1][0] = A[1][0]; R[1][1] = R[1][1];
    return;
  if (n % 2 == 0) {
    mul(A, R, n/2);
    // mul2 takes two 2x2 matrices as input and
    // returns the multiplication in the first
    // matrix
    mul2(R, R); // R \leftarrow R * R
  else {
    mul(A, R, (n-1)/2);
   mul2(R, R); // R <- R * R
mul2(R, A); // R <- R * A
}
```

Fibonacci numbers

```
fib(31)
calls
mul(A, R, 30)
```

```
// mul takes a matrix A and R as input
// returns the result of A^n in R
void mul(int A[2][2], int R[2][2], int n) \{
  if (n == 1) {
    R[0][0] = A[0][0]; R[0][1] = A[0][1];
    R[1][0] = A[1][0]; R[1][1] = R[1][1];
    return;
  if (n % 2 == 0) {
    mul(A, R, n/2);
    // mul2 takes two 2x2 matrices as input and
    // returns the multiplication in the first
    // matrix
    mul2(R, R); // R \leftarrow R * R
  else {
    mul(A, R, (n-1)/2);
   mul2(R, R); // R <- R * R
mul2(R, A); // R <- R * A
}
```

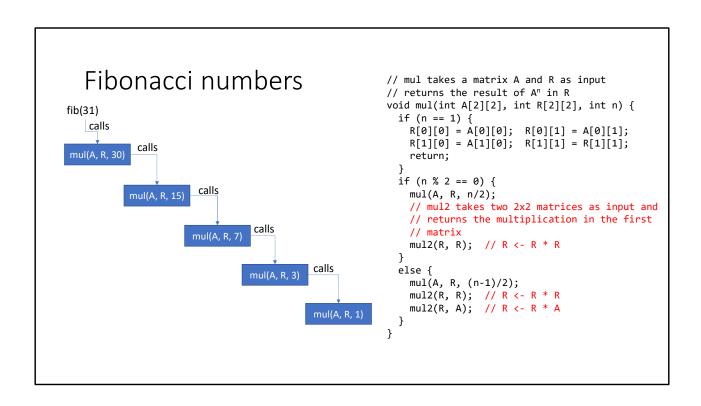
Fibonacci numbers

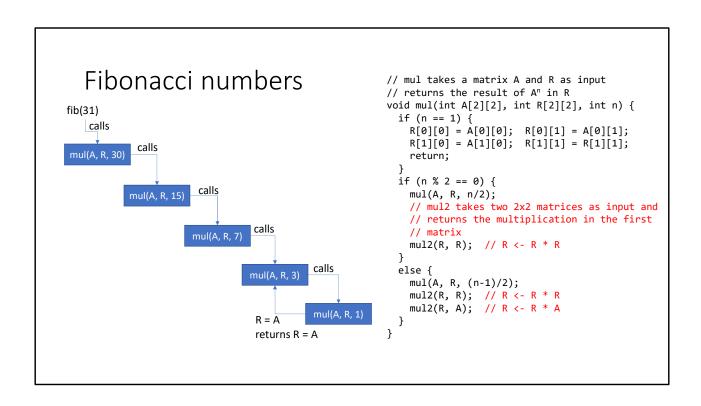


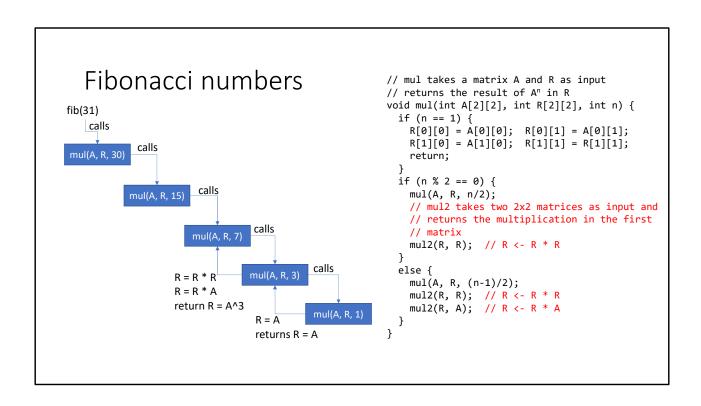
```
// mul takes a matrix A and R as input
// returns the result of A^n in R
void mul(int A[2][2], int R[2][2], int n) \{
  if (n == 1) {
    R[0][0] = A[0][0]; R[0][1] = A[0][1];
    R[1][0] = A[1][0]; R[1][1] = R[1][1];
    return;
  if (n % 2 == 0) {
    mul(A, R, n/2);
    // mul2 takes two 2x2 matrices as input and
    // returns the multiplication in the first
    // matrix
    mul2(R, R); // R \leftarrow R * R
  else {
    mul(A, R, (n-1)/2);
   mul2(R, R); // R <- R * R
mul2(R, A); // R <- R * A
}
```

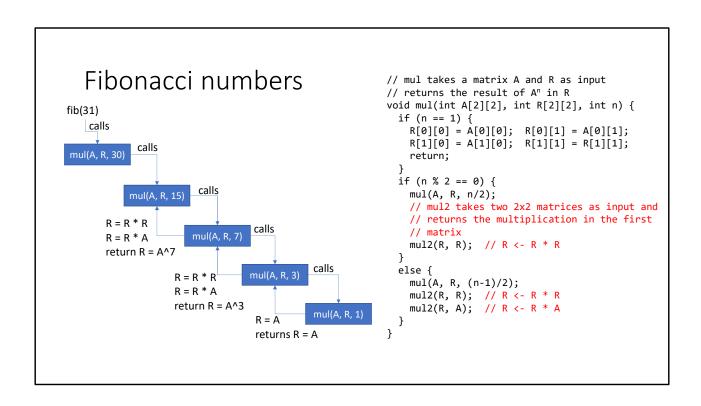
Fibonacci numbers // mul takes a matrix A and R as input // returns the result of A^n in Rvoid mul(int A[2][2], int R[2][2], int n) $\{$ fib(31) if (n == 1) { calls R[0][0] = A[0][0]; R[0][1] = A[0][1];R[1][0] = A[1][0]; R[1][1] = R[1][1];calls mul(A, R, 30) return; if (n % 2 == 0) { calls mul(A, R, 15) mul(A, R, n/2);// mul2 takes two 2x2 matrices as input and // returns the multiplication in the first // matrix mul(A, R, 7) $mul2(R, R); // R \leftarrow R * R$ else { mul(A, R, (n-1)/2);mul2(R, R); // R <- R * R mul2(R, A); // R <- R * A }

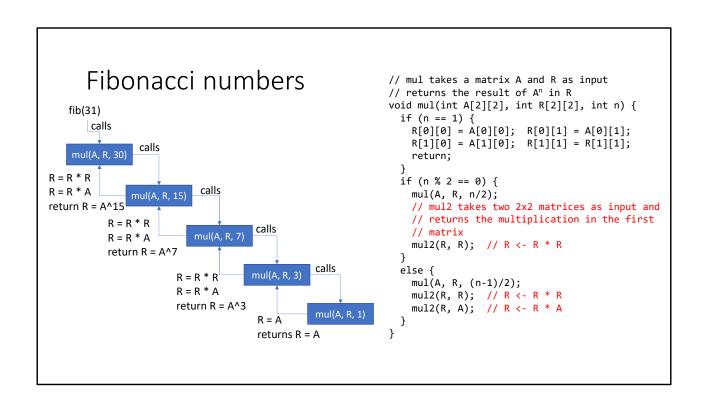
Fibonacci numbers // mul takes a matrix A and R as input // returns the result of A^n in Rvoid mul(int A[2][2], int R[2][2], int n) $\{$ fib(31) if (n == 1) { calls R[0][0] = A[0][0]; R[0][1] = A[0][1];R[1][0] = A[1][0]; R[1][1] = R[1][1];calls mul(A, R, 30) return; if (n % 2 == 0) { calls mul(A, R, 15) mul(A, R, n/2);// mul2 takes two 2x2 matrices as input and // returns the multiplication in the first calls // matrix mul(A, R, 7) $mul2(R, R); // R \leftarrow R * R$ else { mul(A, R, 3) mul(A, R, (n-1)/2);mul2(R, R); // R <- R * R mul2(R, A); // R <- R * A }

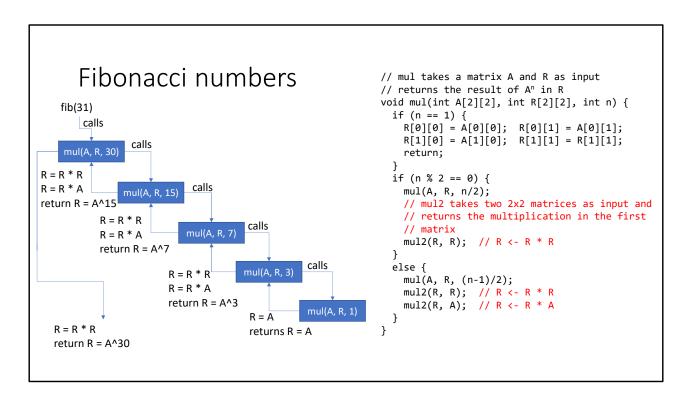








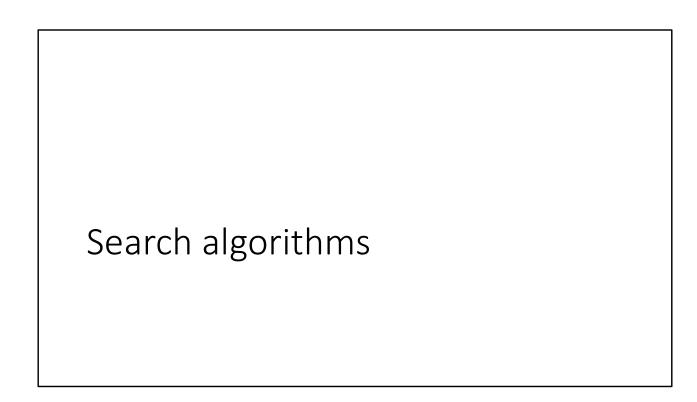


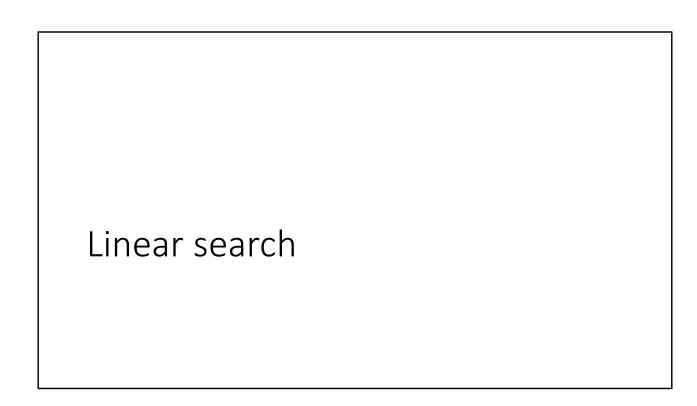


Notice that this algorithm only makes five recursive calls to compute f(31) in contrast to the previous recursive algorithm that makes around 30 calls.

Homework

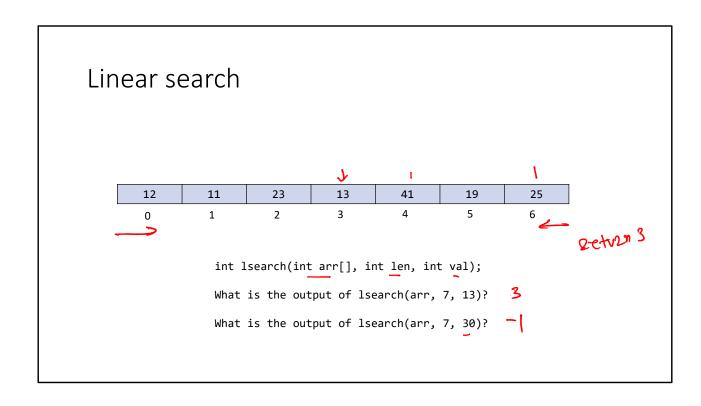
- Modify the fib(n) routines (matrix mul and iterative) discussed in the class to return "nth Fibonacci number % 1000" instead of "nth Fibonacci number"
- Compare the runtimes of matrix mul vs. iterative algorithms for various inputs





Linear search

• Let arr be an input array of length n. Given a value x, we want to find an index i ($\emptyset \le i \le n$), such that arr[i] == x. If no such index exists, then the algorithm returns -1.



Iterative linear search

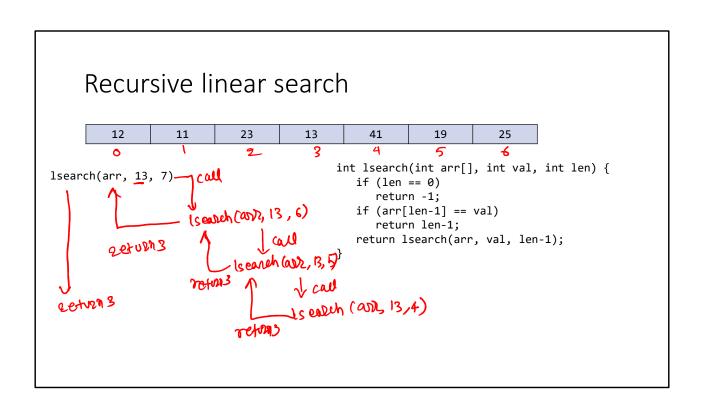
```
int lsearch(int arr[], int val, int len) {
   int i;
   for (i = len-1; i >= 0; i--) {
      if (arr[i] == val)
         return i;
   }
   return -1;
}
```

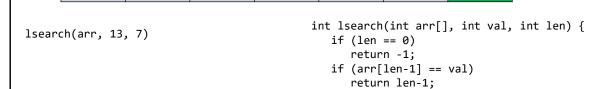
This linear search algorithm iterates all the elements in the array, starting from the last index to the start index. If the value of an array element is equals to val, it returns the corresponding index; otherwise, if val is not present in the array, Isearch returns - 1.

- •int lsearch(int arr[], int val, int len);
 - Initially, len is the length of the input arr
 - val is the value being searched

- •int lsearch(int arr[], int val, int len);
 - Initially, len is the length of the input arr
 - val is the value being searched
- Base cases
 - if (len == 0) return -1
 - if (arr[len-1] == val) return len-1
- Recursive step
 - Decrement the length of the array and recursively call search

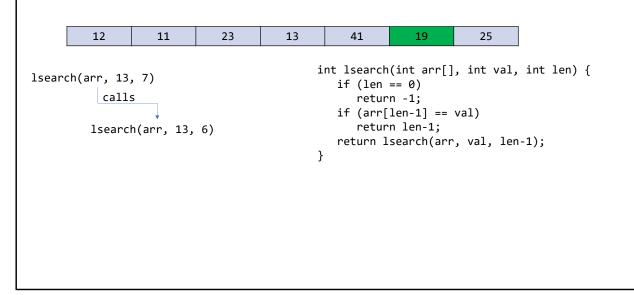
```
int lsearch(int arr[], int val, int len) {
   if (len == 0)
      return -1;
   if (arr[len-1] == val)
      return len-1;
   return lsearch(arr, val, len-1);
}
```

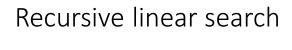


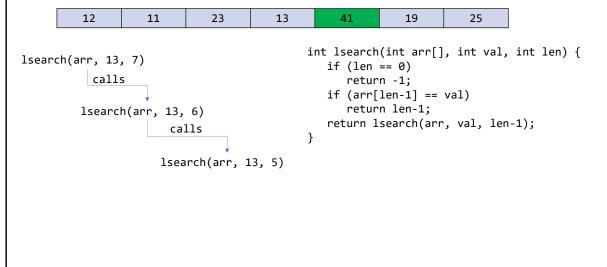


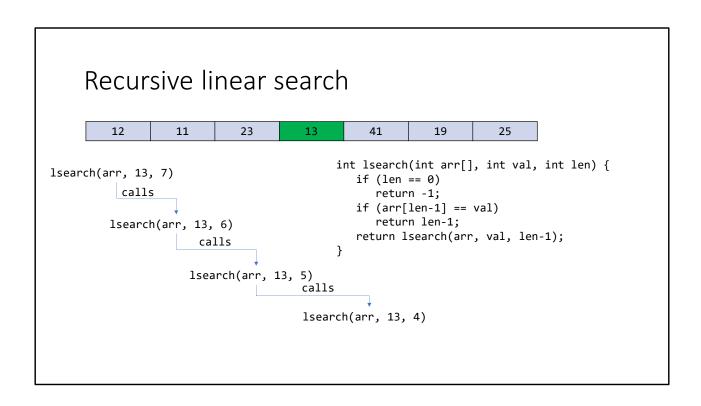
}

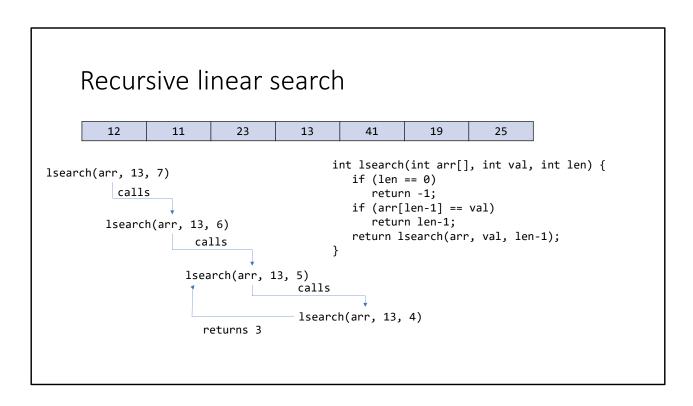
return lsearch(arr, val, len-1);



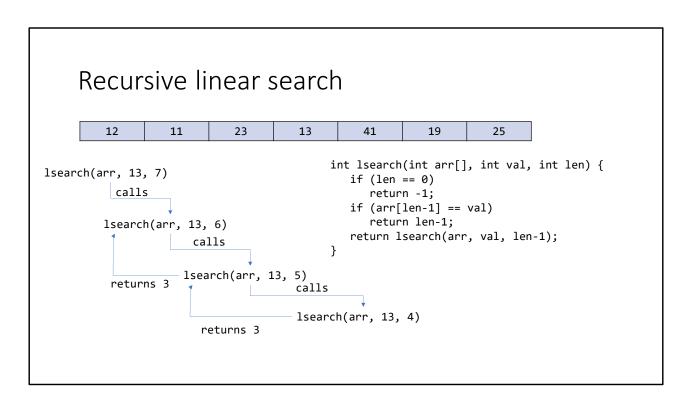




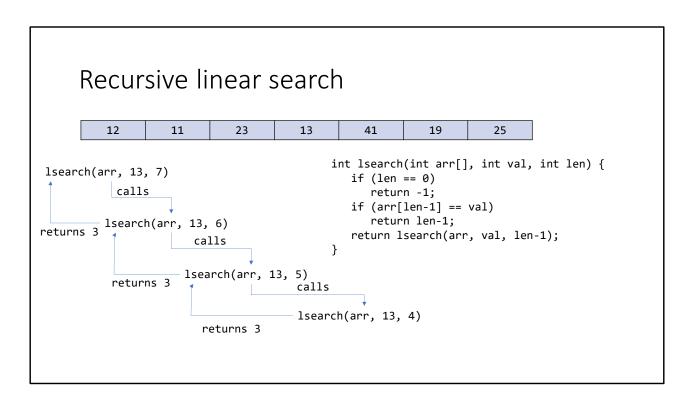




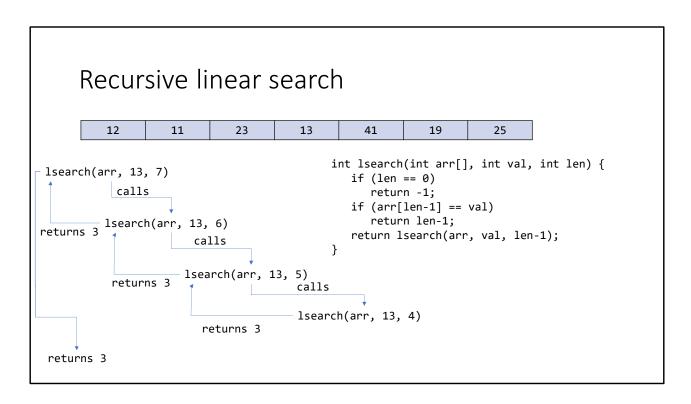
Isearch(arr, 13, 4) returns 3 because arr[len-1], i.e., arr[3] is equal to the value we are searching for, i.e., 13.



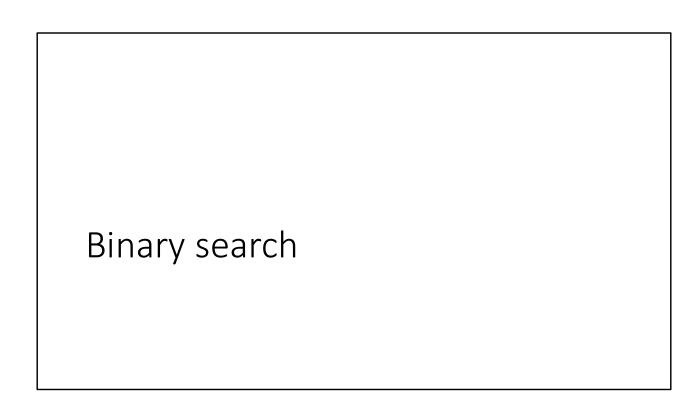
Isearch(arr, 13, 5) returns the return value of Isearch(arr, 13, 4), which is 3.



Isearch(arr, 13, 6) returns the return value of Isearch(arr, 13, 5), which is 3.

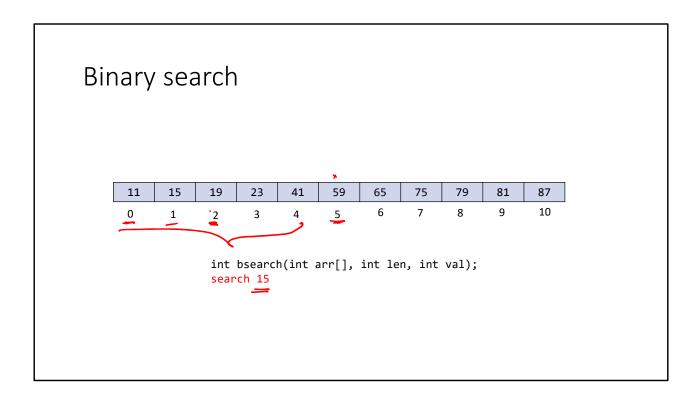


Isearch(arr, 13, 7) returns the return value of Isearch(arr, 13, 6), which is 3.



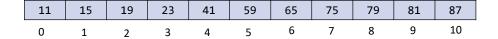
Binary search

• Let arr be an input sorted array (in ascending order) of length n. Given a value x, we want to find an index i (0 <= i < n), such that arr[i] == x. If no such index exists, then the algorithm returns -1.

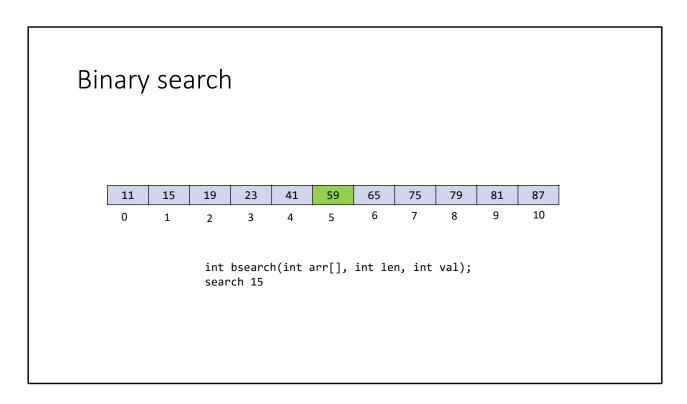


The binary search algorithm works on a sorted array. The algorithm first computes the middle index (mid), which is equal to (start+end)/2, where start and end are the first and last index of the array. In this case, initially, mid is 5. If the val == arr[mid], then the algorithm simply returns mid; otherwise, if val > arr[mid], then the searching is performed in the subarray starting from mid+1 to end, else we search the element in the subarray starting from start to mid-1. If the element is not present in the array, the start index will eventually become lesser than the end index, and in that case, the algorithm returns -1.

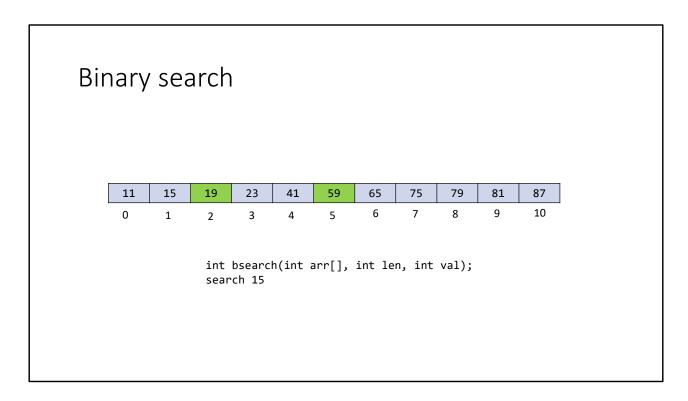
Binary search



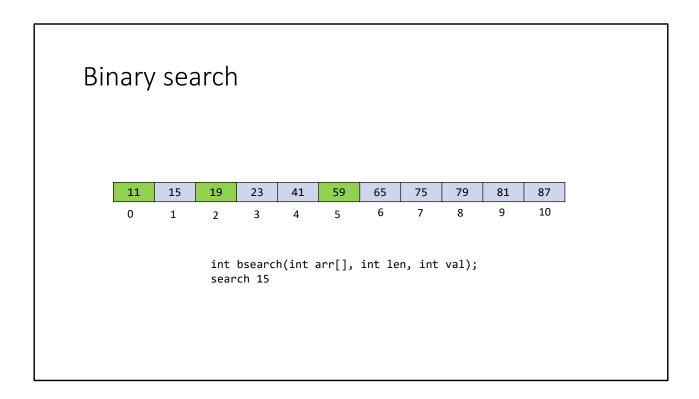
int bsearch(int arr[], int len, int val);
search 15



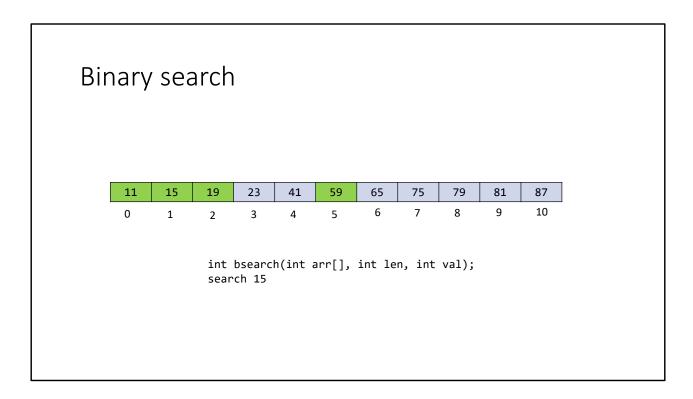
Searching 15 in the range (0,10). mid = (0+10)/2 = 5. 15 < arr[5], search in the range (0,4).



Searching 15 in the range (0,4). mid = (0+4)/2 = 2. 15 < arr[2], search in the range (0,1).

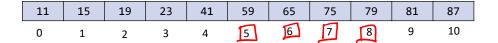


Searching 15 in the range (0,1). mid = (0+1)/2 = 0. 15 > arr[0], search in the range (1, 1).



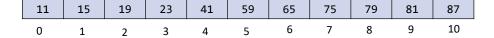
Searching 15 in the range (1,1). mid = (1+1)/2 = 1. 15 == arr[1], returns 1.

Binary search

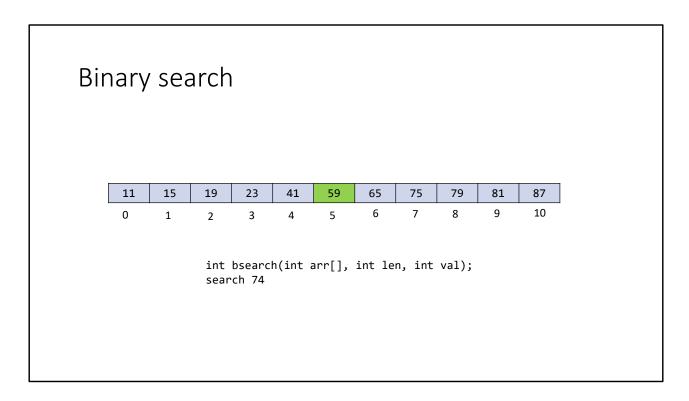


int bsearch(int arr[], int len, int val);
search 74

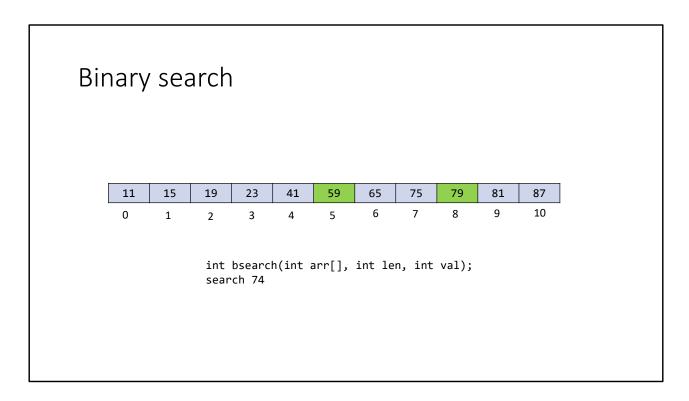
Binary search



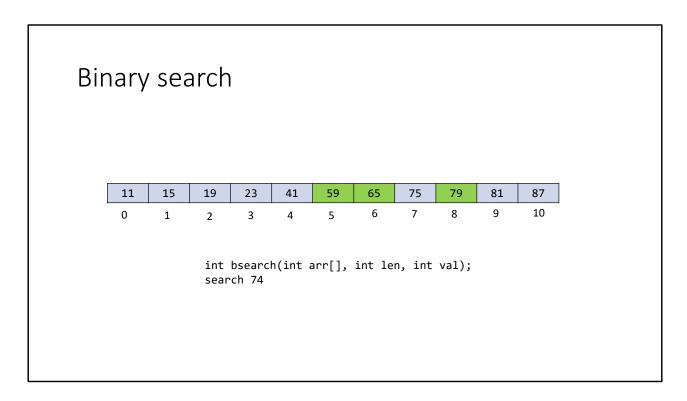
int bsearch(int arr[], int len, int val);
search 74



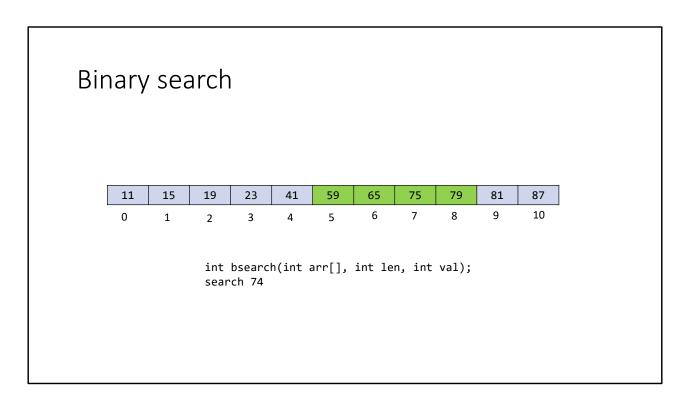
Searching 74 in the range (0,10). mid = (0+10)/2 = 5. 74 > arr[5], search in the range (6, 10).



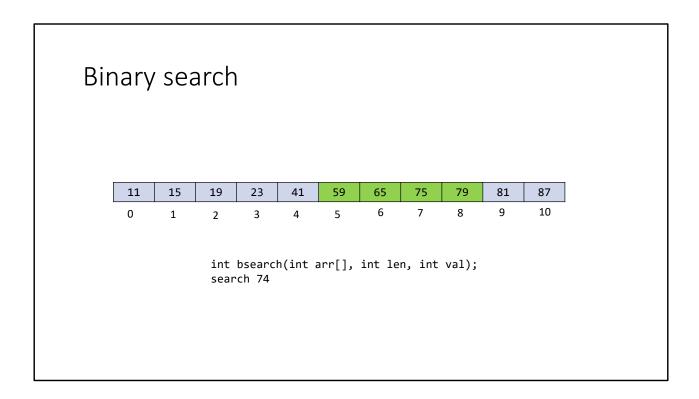
Searching 74 in the range (6,10). mid = (6+10)/2 = 8. 74 < arr[8], search in the range (6,7).



Searching 74 in the range (6,7). mid = (6+7)/2 = 6. 74 > arr[6], search in the range (7, 7).



Searching 74 in the range (7,7). mid = (7+7)/2 = 7. 74 < arr[7], search in the range (7,6).



Searching 74 in the range (7,6). start > end. The value is not present. Returns -1.

Recursive solution

```
int bsearch(int arr[], int val, int lo, int hi);
!nitially, lo = 0 and hi = length(arr) - 1
val is the value being searched
Base cases
if (lo > hi) return -1
mid = (lo + hi)/2
if (arr[mid] == val) return mid
Recursive step
!f val > arr[mid], search the second half (mid+1, lo), recursively
Otherwise, search the first half (lo, mid-1)
```

Recursive solution

```
int bsearch(int arr[], int val, int lo, int hi) {
   if (hi < lo)
      return -1;
   int mid = (lo + hi) / 2;
   if (arr[mid] == val)
      return mid;
   if (arr[mid] > val)
      return bsearch(arr, val, lo, mid-1);
   else
      return bsearch(arr, val, mid+1, hi);
}
```