## Theory of Computation '23 Problem Set 3

**Notations.** Let  $\Sigma = \{a, b\}$ . For  $w \in \Sigma^*$ , let |w| denote the length of w. Let  $\#_a(w)$  denote the number of a's in w and let  $\#_b(w)$  denote the number of b's in w.

**Problem 1.** Show that regular languages are closed under the **repeat** operation, where **repeat** operation on a language L is given by

$$\mathbf{repeat}(L) = \{\ell_1 \ell_1 \ell_2 \ell_2 \dots \ell_k \ell_k \mid \ell_1 \ell_2 \dots \ell_k \in L\}$$

**Problem 2.** If A is any language, let  $A_{\frac{1}{2}}$  denote the set of all first halves of strings in A so that

$$A_{\frac{1}{2}-}=\{x\mid \text{for some }y,\, |x|=|y| \text{ and } xy\in A\}$$

Show that if A is regular, then so is  $A_{\frac{1}{2}-}$ .

**Problem 3.** For every string  $x \in \{0,1\}^+$  consider the number

$$0.x = x[1] \cdot \frac{1}{2} + x[2] \cdot \frac{1}{2^2} + \dots + x[|x|] \cdot \frac{1}{2^{|x|}}$$

where |x| is the length of x. For a real number  $\theta \in [0,1]$  let

$$L_{\theta} = \{x : 0.x \le \theta\}$$

Prove that  $L_{\theta}$  is regular if and only if  $\theta$  is rational.

**Problem 4** Prove or disprove that the following languages are regular.

- a) Twice =  $\{w \in \{a, b\}^* | \#_a(w) = 2\#_b(w)\}$
- b ) NEQ =  $\{0^i 1^j | i \neq j\}$  [This is a slightly hard problem.]

**Problem 5** Consider the language  $L = \{w \in \Sigma^* | \text{ 2nd letter from the end is } a\}.$ 

- a ) Draw an NFA for L.
- b) Using the ideas of subset construction draw a DFA for L.
- c ) Using the DFA minimization idea discussed in class, check whether the DFA thus constructed is minimal or not. If it not a minimal DFA then draw the correspoding minimal DFA for it.

d) Let  $L_k = \{w \in \Sigma^* | k \text{th letter from the end is } a \}$ . Prove using pigeon hole principle (or by any other method) that any DFA accepting  $L_k$  must have  $\Omega(k)$  states.

**Problem 6** Let L be a regular language. One of the following languages is regular and the other is not. Give a proof and provide a counterexample, respectively.

- a )  $\{w \in \{a,b\}^\star | \exists n \geq 0, \exists x \in L, x = w^n\}$
- b )  $\{w \in \{a, b\}^* | \exists n \ge 0, \exists x \in L, w = x^n\}$