

# Data Structures and Algorithms

Lecture 04

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2026-01-19 Mon

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## 1 Agenda

### 1.1 Merge Sort

- Proof of Correctness
- Time Complexity

## 2 Merge Sort

### 2.1 Divide-Conquer-Combine Approach

- Divide the input into a number of subproblems.

- Conquer the subproblems by solving them recursively.
- Combine the solved subproblems to return the solution to the original problem.

### 3 Merge Sort

**3.1 Input:**  $A[1 : n]$

**3.2 Output:**  $A[1 : n]$  is sorted.

- Divide  $A[1 : n]$  into  $A[1 : \lfloor n/2 \rfloor]$  and  $A[\lfloor n/2 \rfloor + 1 : n]$
- Recursively run Merge Sort on  $A[1 : \lfloor n/2 \rfloor]$  and  $A[\lfloor n/2 \rfloor + 1 : n]$
- Merge the two sorted arrays to sort  $A[1 : n]$

**3.3 An Example:**

**3.4 Description of the Algorithm**

**3.5 Proof of Correctness**

- Inductive Proof
- Loop Invariant in the Merge function
- Draw the flow chart

**3.6 Time Complexity**

- Recurrence Relation

$$T(1) = \Theta(1)$$

$$T(n) = 2T(n/2) + \Theta(n)$$

- Prove using Recursion Tree

$$T(n) = \Theta(n \log n)$$