



OPERATING SYSTEM

Computer System Organization

Dr Rahul Nagpal
Computer Science

OPERATING SYSTEM

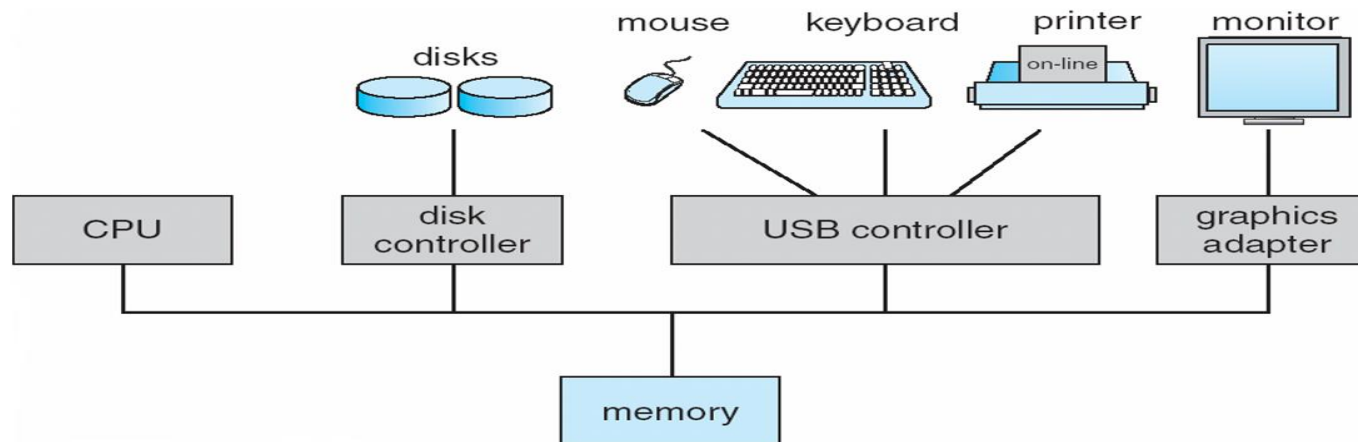
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OPERATING SYSTEMS

Computer System Organization

- Computer-system Organization
 - One or more CPUs, device controllers connect through common bus providing access to shared memory
 - Concurrent execution of CPUs and devices competing for memory cycles



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Computer System Operation



- I/O devices and the CPU can execute concurrently
- Each device controller is in charge of a particular device type
- Each device controller has a local buffer
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller
- Device controller interrupts CPU on completion

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Computer Startup



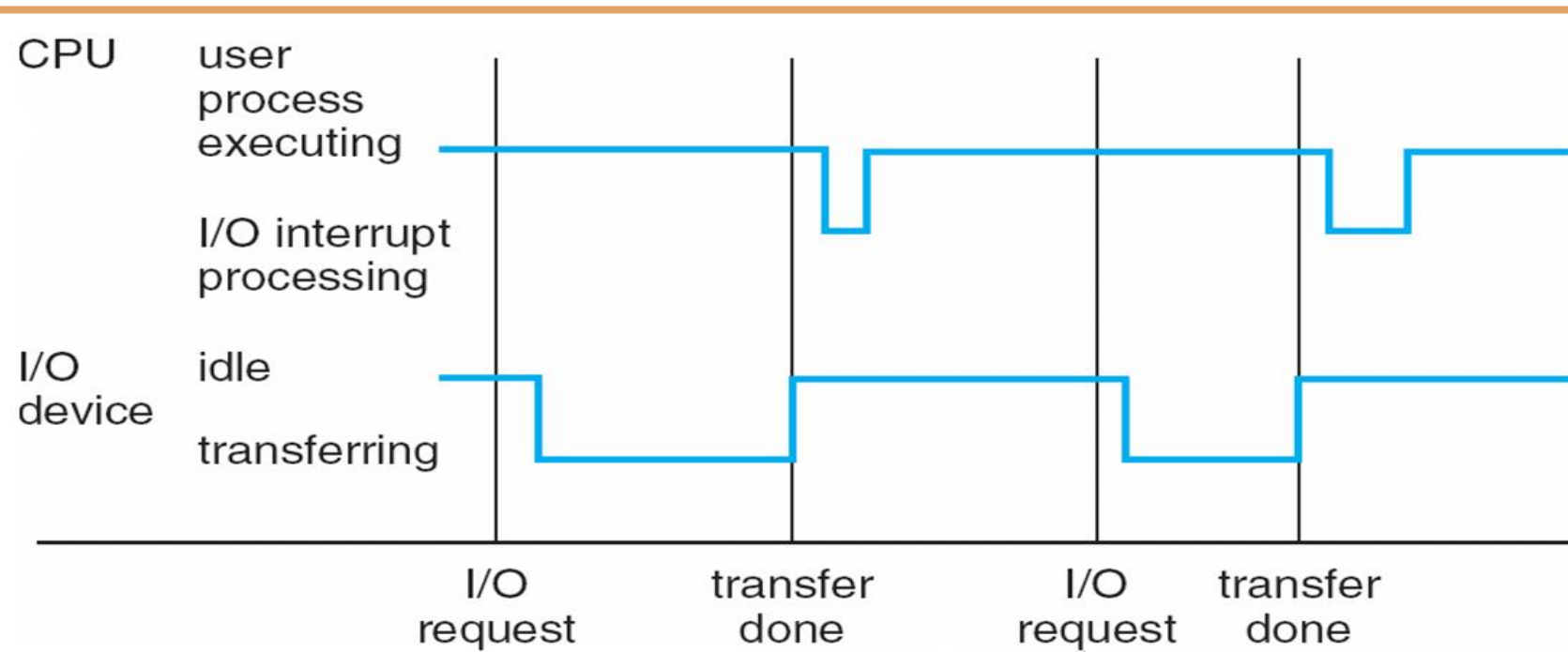
- **bootstrap program** is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, known as **firmware**
 - Initializes all aspects of system
 - Loads operating system kernel and starts execution

- An operating system is **interrupt driven**
- Interrupt transfers control to the interrupt service routine
- **interrupt vector** contains the addresses of all the service routines for transfer
- Interrupt architecture must save the address of the interrupted instruction
- A **trap** or **exception** is a software-generated interrupt caused either by an error or a user request

- The operating system preserves the state of the CPU by storing registers and the program counter
- Determines which type of interrupt has occurred:
 - **polling**
 - **vectored** interrupt system
- Separate segments of code determine what action should be taken for each type of interrupt

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Interrupt Timeline

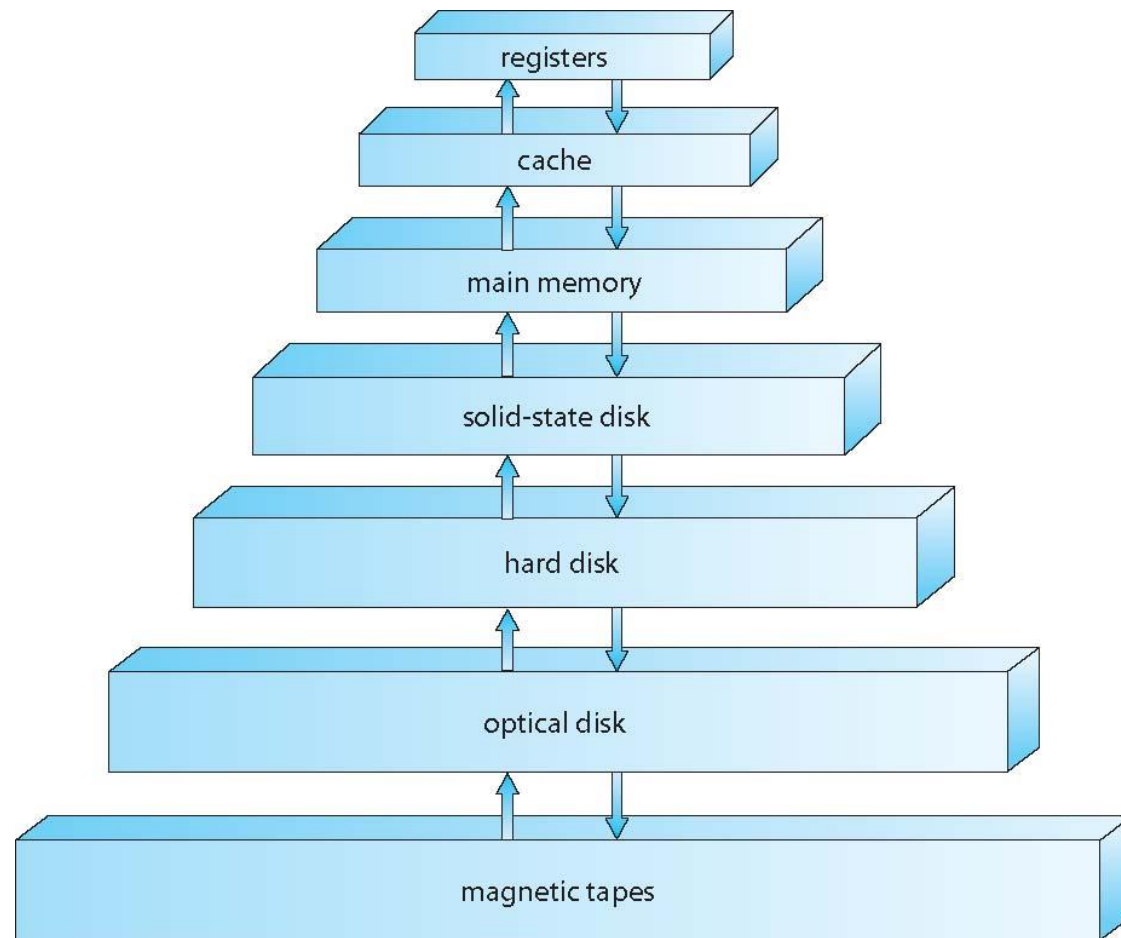


- Main memory – only large storage media directly access by CPU
 - **Random access and typically volatile**
- Secondary storage – extension of main memory that provides large **nonvolatile** storage capacity
- Hard disks – rigid metal or glass platters covered with magnetic recording material
 - Disk surface is logically divided into **tracks**, which are subdivided into **sectors**
 - The **disk controller** determines the logical interaction between the device and the computer
- **Solid-state disks** – faster than hard disks, nonvolatile
 - Various technologies
 - Becoming more popular

Storage Hierarchy



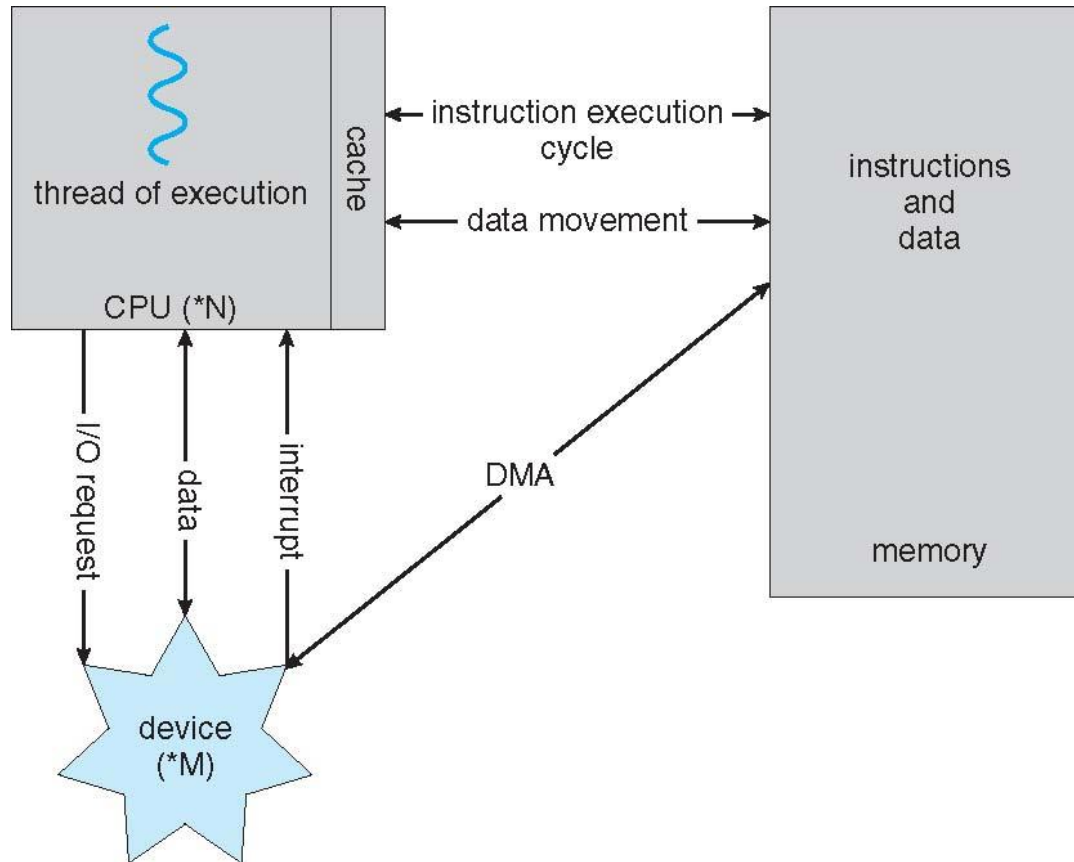
- Storage systems organized in hierarchy
 - Speed
 - Cost
 - Volatility
- **Caching** – copying information into faster storage system; main memory can be viewed as a cache for secondary storage
- **Device Driver** for each device controller to manage I/O
 - Provides uniform interface between controller and kernel



- Important principle, performed at many levels in a computer (in hardware, operating system, software)
- Information in use copied from slower to faster storage temporarily
- Faster storage (cache) checked first to determine if information is there
 - If it is, information used directly from the cache (fast)
 - If not, data copied to cache and used there
- Cache smaller than storage being cached
 - Cache management important design problem
 - Cache size and replacement policy

- ❑ After I/O starts, control returns to user program only upon I/O completion
 - ❑ Wait instruction idles the CPU until the next interrupt
 - ❑ Wait loop (contention for memory access)
 - ❑ At most one I/O request is outstanding at a time, no simultaneous I/O processing
- ❑ After I/O starts, control returns to user program without waiting for I/O completion
 - ❑ **System call** – request to the OS to allow user to wait for I/O completion
 - ❑ **Device-status table** contains entry for each I/O device indicating its type, address, and state
 - ❑ OS indexes into I/O device table to determine device status and to modify table entry to include interrupt

- Used for high-speed I/O devices able to transmit information at close to memory speeds
- Device controller transfers blocks of data from buffer storage directly to main memory without CPU intervention
- Only one interrupt is generated per block, rather than the one interrupt per byte



A von Neumann architecture



THANK YOU

Dr Rahul Nagpal

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rahulnagpal@pes.edu