



OPERATING SYSTEM

Multi-Processor & Real-time Scheduling

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Computer Science

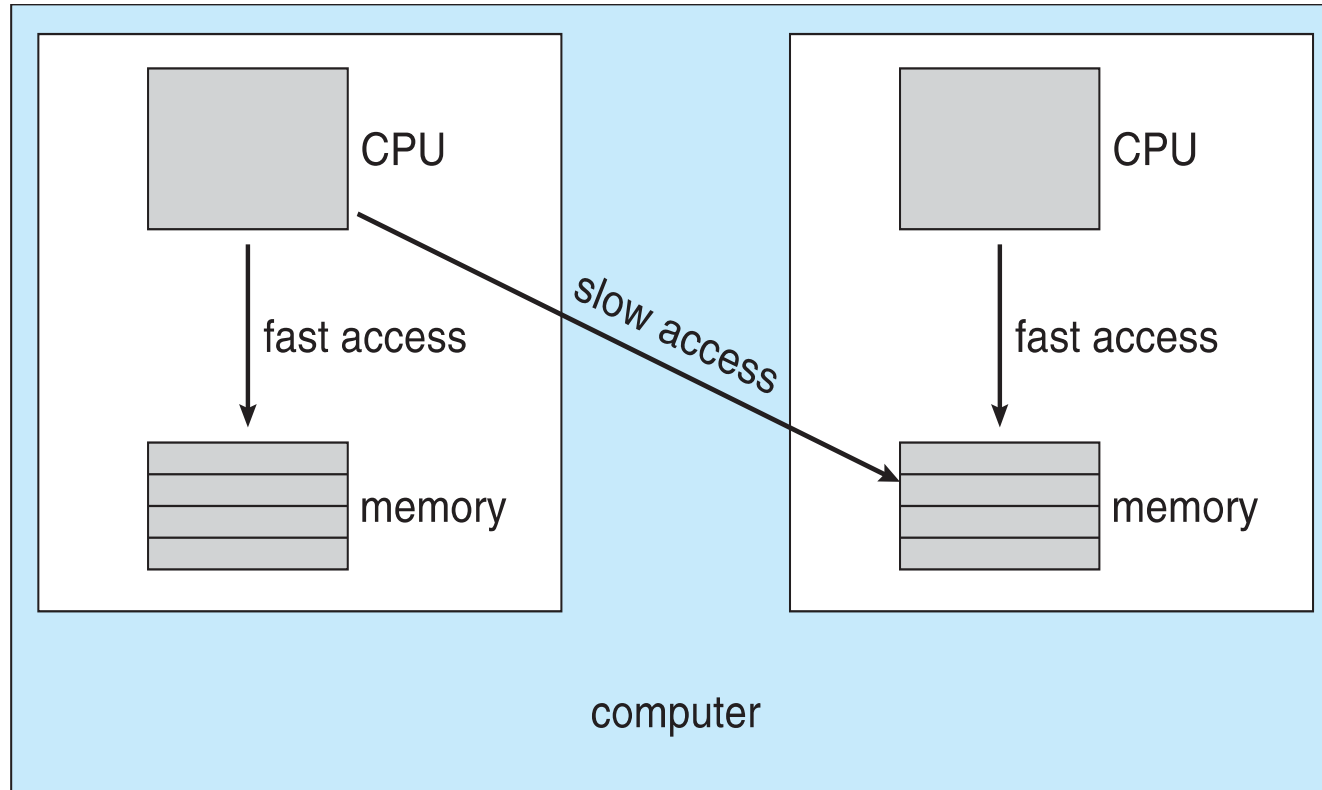
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- Multiprocessor Scheduling
- Real-time Scheduling
 - Priority Based Scheduling
 - Rate Monotonic Scheduling
 - Earliest Deadline First Scheduling
 - Proportional Share Scheduling

- CPU scheduling more complex when multiple CPUs are available
- **Homogeneous processors** within a multiprocessor
- **Asymmetric multiprocessing** – only one processor accesses the system data structures, alleviating the need for data sharing
- **Symmetric multiprocessing (SMP)** – each processor is self-scheduling, all processes in common ready queue, or each has its own private queue of ready processes
 - Currently, most common
- **Processor affinity** – process has affinity for processor on which it is currently running
 - **soft affinity, hard affinity**, Variations including **processor sets**



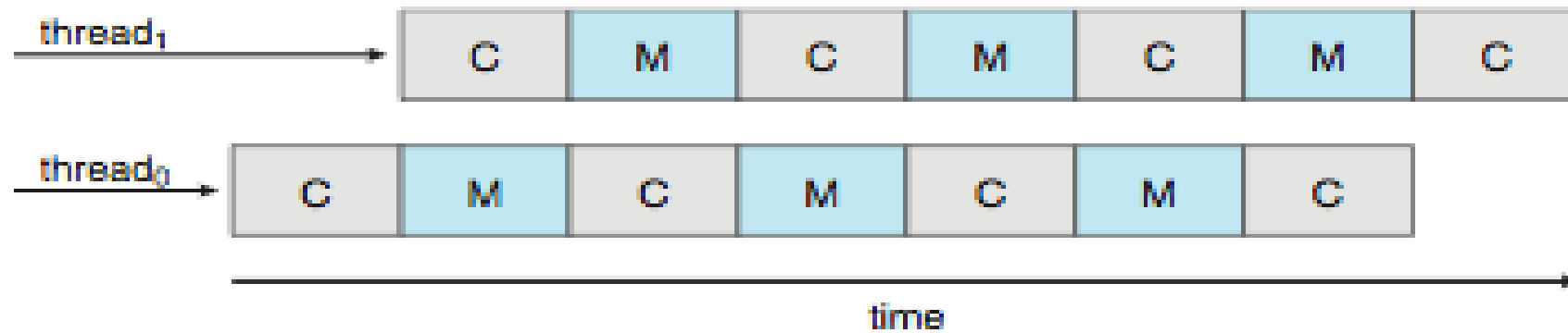
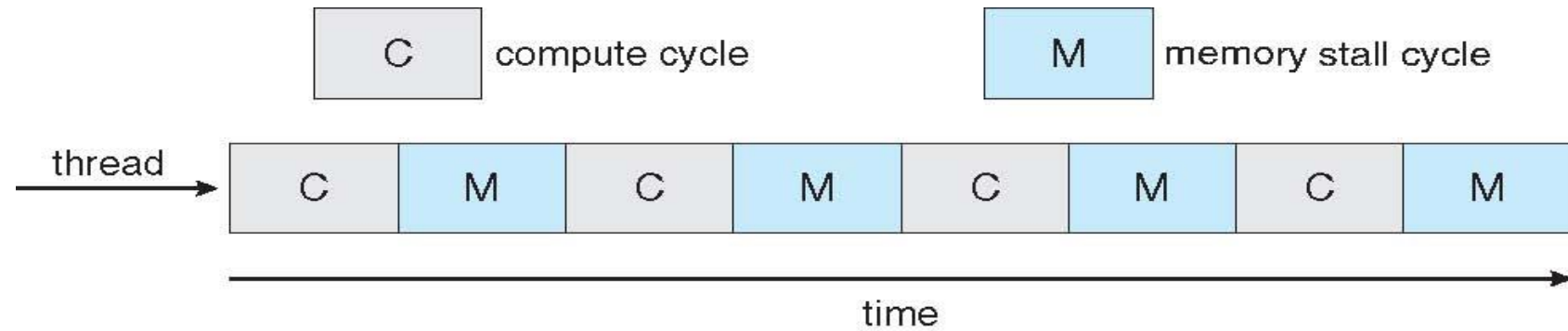
Note that memory-placement algorithms can also consider affinity

- If SMP, need to keep all CPUs loaded for efficiency
- **Load balancing** attempts to keep workload evenly distributed
- **Push migration** – periodic task checks load on each processor, and if found pushes task from overloaded CPU to other CPUs
- **Pull migration** – idle processors pulls waiting task from busy processor

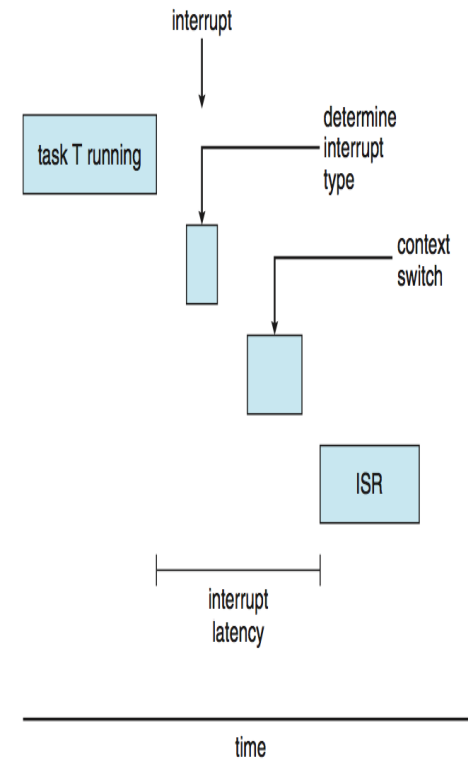
- Recent trend to place multiple processor cores on same physical chip
- Faster and consumes less power
- Multiple threads per core also growing
 - Takes advantage of memory stall to make progress on another thread while memory retrieve happens

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Multithreaded Multicore System



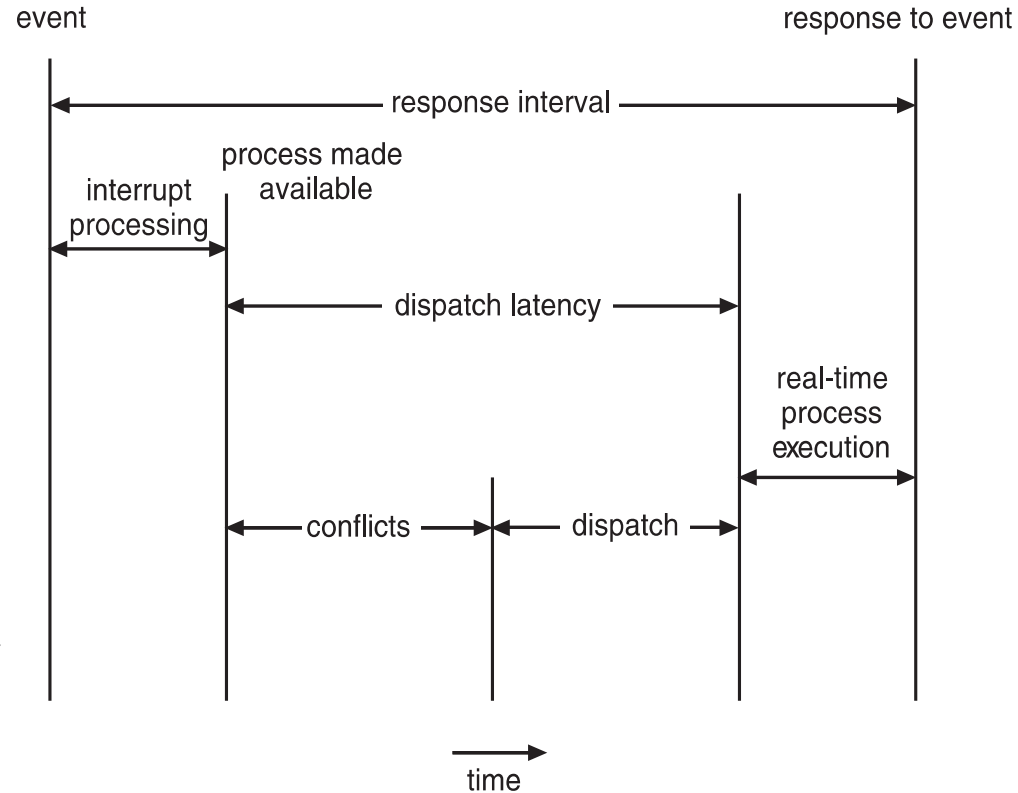
- Can present obvious challenges
- **Soft real-time systems** – no guarantee as to when critical real-time process will be scheduled
- **Hard real-time systems** – task must be serviced by its deadline
- Two types of latencies affect performance
 1. Interrupt latency – time from arrival of interrupt to start of routine that services interrupt
 2. Dispatch latency – time for schedule to take current process off CPU and switch to another



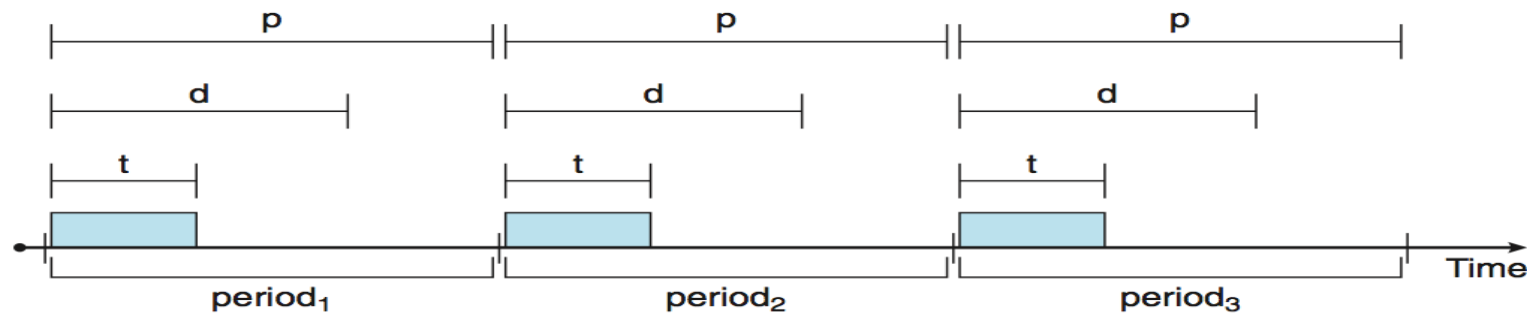
- Conflict phase of dispatch latency:

1. Preemption of any process running in kernel mode

2. Release by low-priority process of resources needed by high-priority processes



- For real-time scheduling, scheduler must support preemptive, priority-based scheduling
 - But only guarantees soft real-time
- For hard real-time must also provide ability to meet deadlines
- Processes have new characteristics: **periodic** ones require CPU at constant intervals
 - Has processing time t , deadline d , period p
 - $0 \leq t \leq d \leq p$
 - **Rate** of periodic task is $1/p$



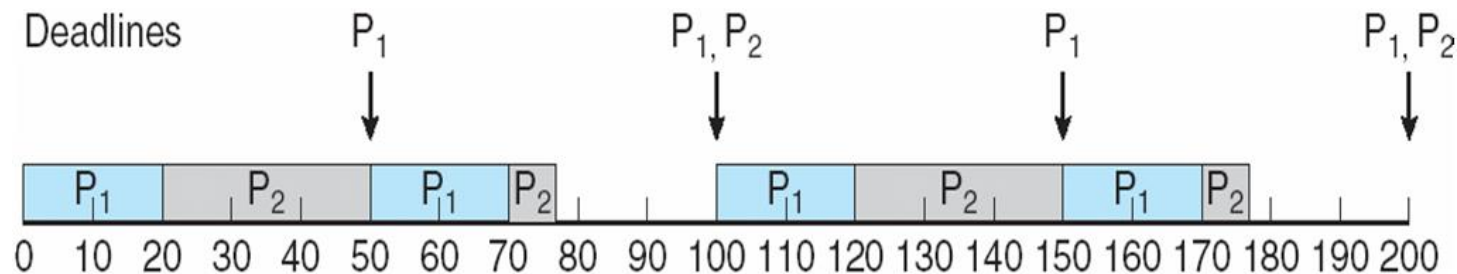
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Virtualization and Scheduling



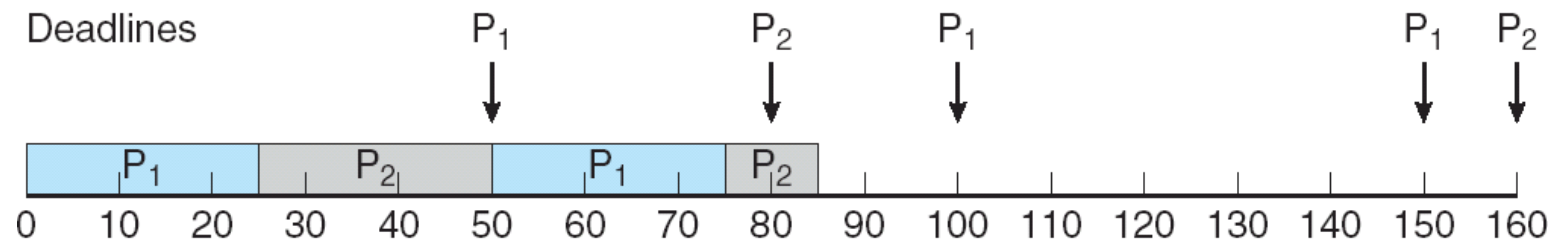
- Virtualization software schedules multiple guests onto CPU(s)
- Each guest doing its own scheduling
 - Not knowing it doesn't own the CPUs
 - Can result in poor response time
 - Can effect time-of-day clocks in guests
- Can undo good scheduling algorithm efforts of guests

- A priority is assigned based on the inverse of its period
- Shorter periods = higher priority;
- Longer periods = lower priority
- P_1 is assigned a higher priority than P_2 .



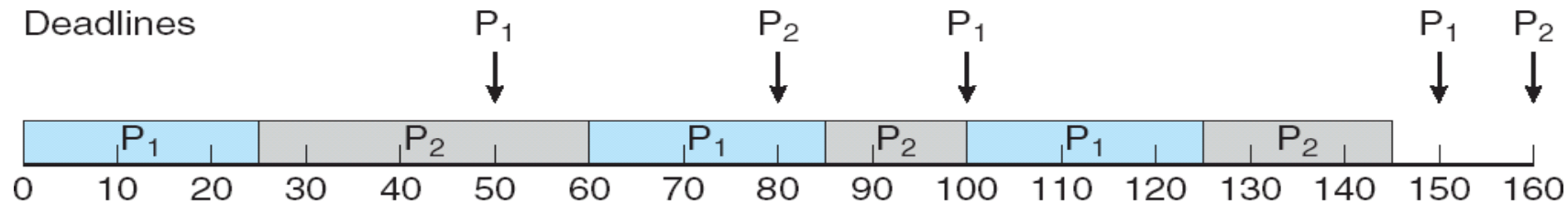
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Missed Deadlines with Rate Monotonic Scheduling



Earliest Deadline First Scheduling (EDF)

- Priorities are assigned according to deadlines
 - the earlier the deadline, the higher the priority;
 - the later the deadline, the lower the priority



- T shares are allocated among all processes in the system
- An application receives N shares where $N < T$
- This ensures each application will receive N / T of the total processor time



THANK YOU

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