



PESU Center for
Information Security,
Forensics and
Cyber Resilience



Welcome to
PES University
Ring Road Campus, Bengaluru



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Information Security,
Forensics and
Cyber Resilience



APPLIED CRYPTOGRAPHY

Private keys systems

Lecture 8

AES

Advanced Encryption Standard

AES animation

How to navigate through the animation:

- > press **Control + F** to get into full screen mode
- > use **Enter** key to advance
- > use **Slide controller** on bottom to navigate
- > press **c** to show/hide the slide controller

Origins

- A replacement for DES was needed
 - Key size is too small
- Can use Triple-DES – but slow, small block
- US NIST issued call for ciphers in 1997
- 15 candidates accepted in Jun 98
- 5 were shortlisted in Aug 99

AES Competition Requirements

- Private key symmetric block cipher
- 128-bit data, 128/192/256-bit keys
- Stronger & faster than Triple-DES
- Provide full specification & design details
- Both C & Java implementations

AES Shortlist

- After testing and evaluation, shortlist in Aug-99
 - MARS (IBM) - complex, fast, high security margin
 - RC6 (USA) - v. simple, v. fast, low security margin
 - Rijndael (Belgium) - clean, fast, good security margin
 - Serpent (Euro) - slow, clean, v. high security margin
 - Twofish (USA) - complex, v. fast, high security margin
- Found contrast between algorithms with
 - few complex rounds versus many simple rounds
 - Refined versions of existing ciphers versus new proposals

The AES Cipher - Rijndael

- Rijndael was selected as the AES in Oct-2000
 - Designed by Vincent Rijmen and Joan Daemen in Belgium
 - Issued as FIPS PUB 197 standard in Nov-2001
- An **iterative** rather than **Feistel** cipher
 - processes data as block of 4 columns of 4 bytes (128 bits)
 - operates on entire data block in every round
- Rijndael design:
 - simplicity
 - has 128/192/256 bit keys, 128 bits data
 - resistant against known attacks
 - speed and code compactness on many CPUs

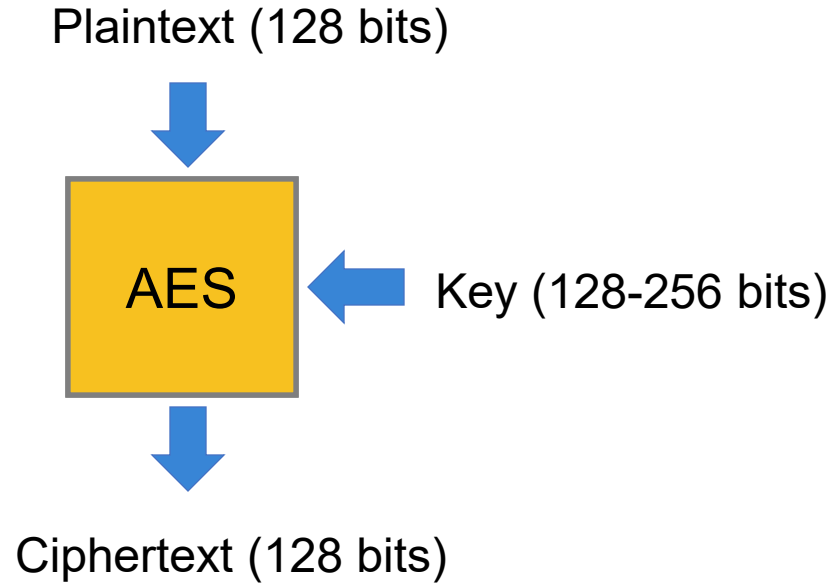


V. Rijmen

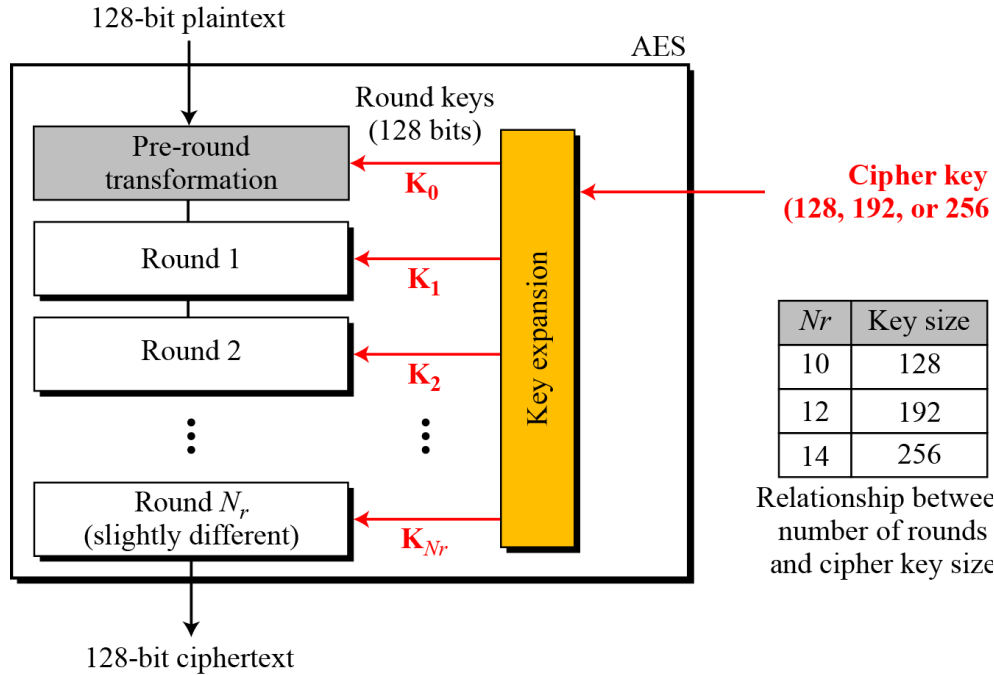


J. Daemen

AES Conceptual Scheme



Multiple rounds

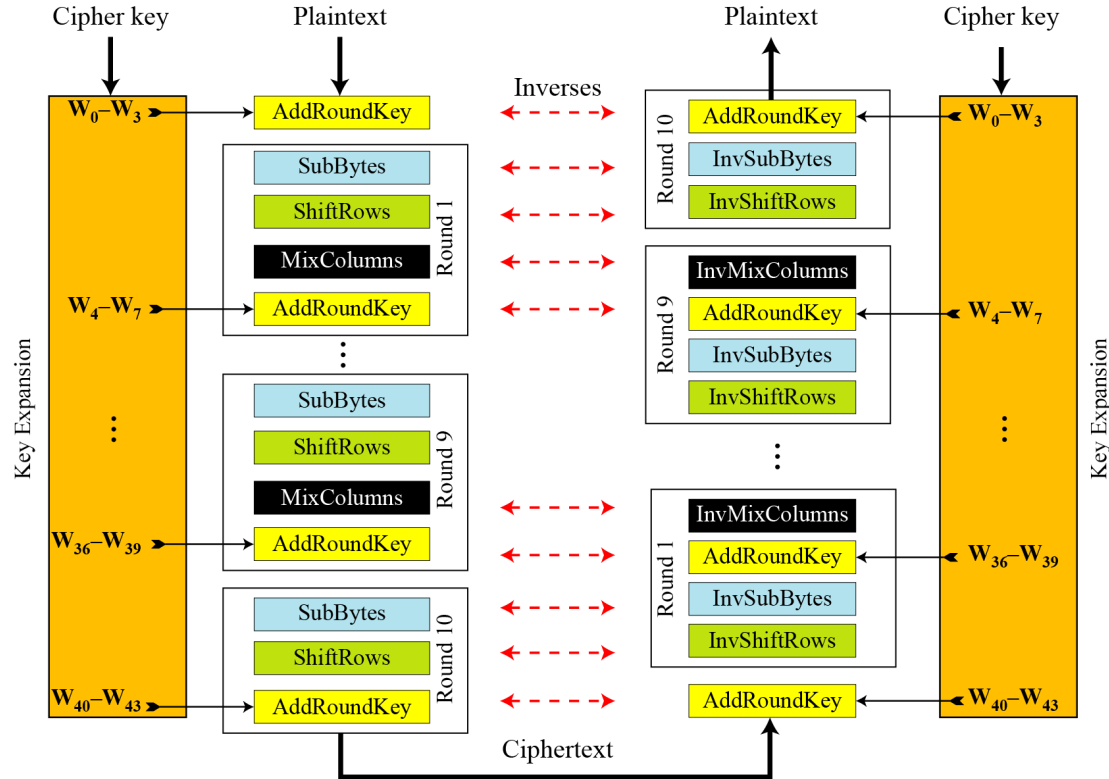


- Rounds are (almost) identical
- First and last round are a little different

High Level Description

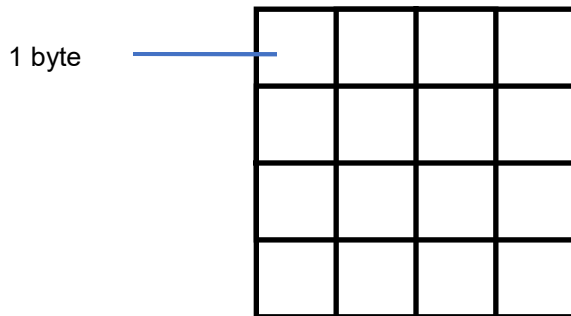
- **Key Expansion:** Round keys are derived from the cipher key using Rijndael's key schedule
- **Initial Round:** AddRoundKey : Each byte of the state is combined with the round key using bitwise xor
- **Rounds**
 - SubBytes : non-linear substitution step
 - ShiftRows : transposition step
 - MixColumns : mixing operation of each column.
 - AddRoundKey
- **Final Round:**
 - SubBytes
 - ShiftRows
 - AddRoundKey

Overall Structure

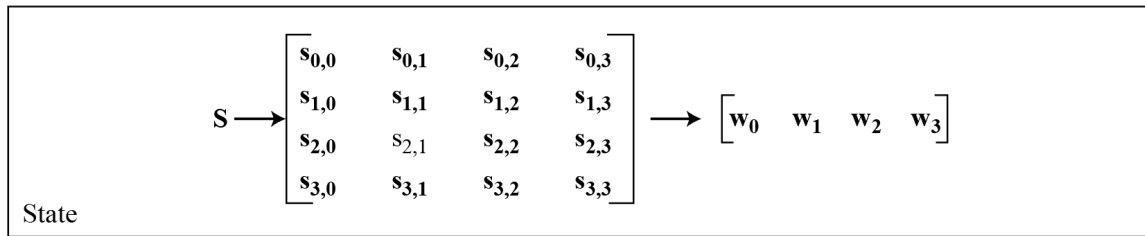
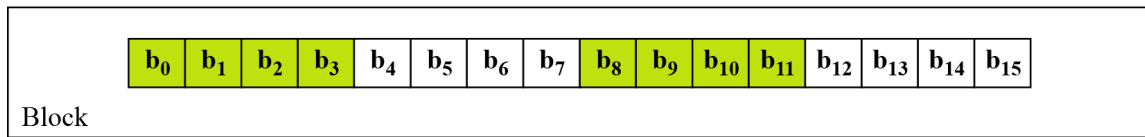
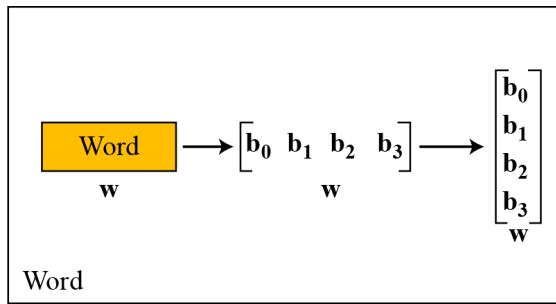
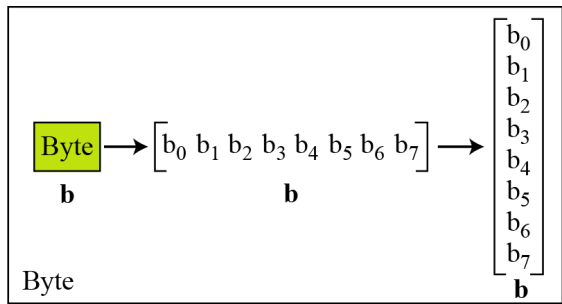


128-bit values

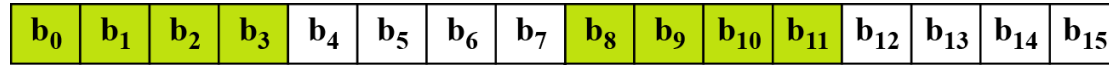
- Data block viewed as 4-by-4 table of bytes
- Represented as 4 by 4 matrix of 8-bit bytes.
- Key is expanded to array of 32 bits words



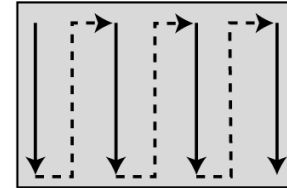
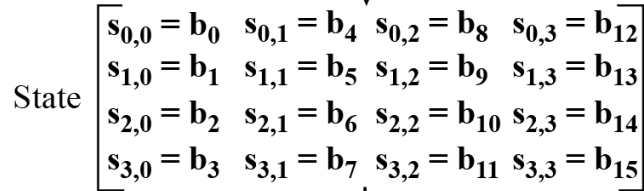
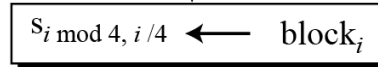
Data Unit



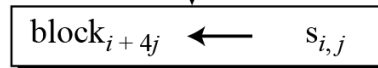
Unit Transformation



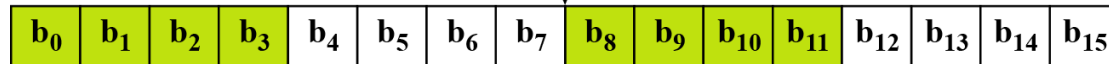
Block



Insertion and
extraction flow



Block



AES

Plaintext: AES USES MATRIX = 14 characters = $14 \times 8 = 112$

Converting plaintext to state which is 128 bits. $128 - 112 = 16$ bits needed

So add 2 extra character

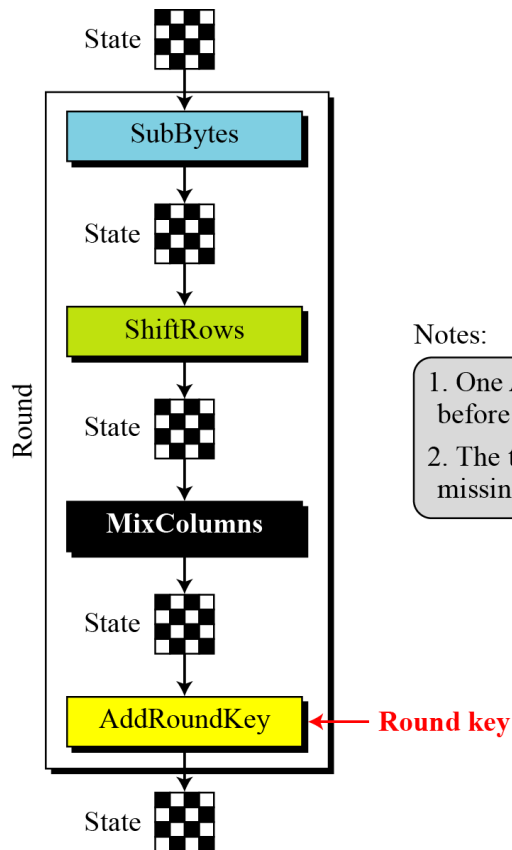
Plaintext →	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z
Value →	00	01	02	03	04	05	06	07	08	09	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25

Changing Plaintext to State

Text	A E S U S E S A M A T R I X Z Z															
Hexadecimal	00 04 12 14 12 04 12 00 0C 00 13 11 08 17 19 19															
<div><div><div>00120C08</div><div>04040017</div><div>12121319</div><div>14001119</div></div><div>State</div></div>																

Reference 2: example 7.1

Details of Each Round



Notes:

1. One AddRoundKey is applied before the first round.
2. The third transformation is missing in the last round.

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