

## DFA to Regular Expression-

The two popular methods for converting a DFA to its regular expression are-

### Converting DFA to Regular Expression (Methods)

Arden's Method

State Elimination Method

1. Arden's Method
2. State Elimination Method

### State Elimination Method-

This method involves the following steps in finding the regular expression for any given DFA-

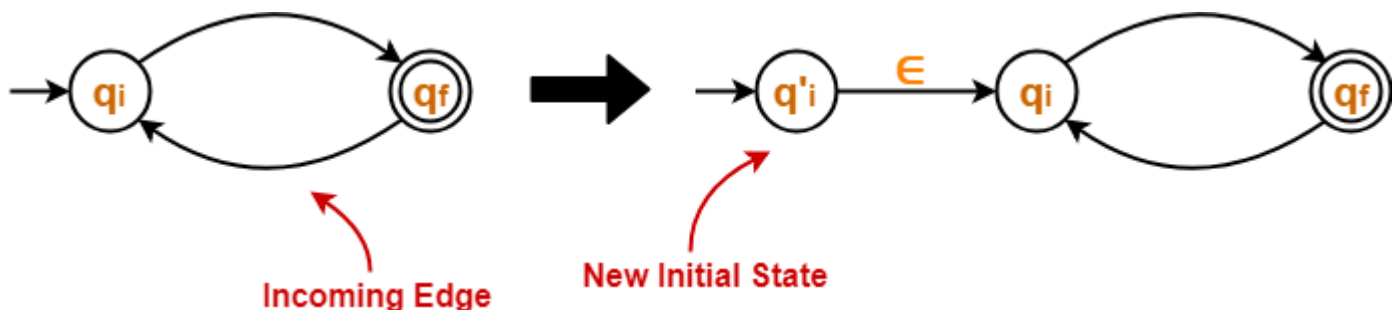
#### Step-01:

##### Thumb Rule

The initial state of the DFA must not have any incoming edge.

- If there exists any incoming edge to the initial state, then create a new initial state having no incoming edge to it.

#### Example-



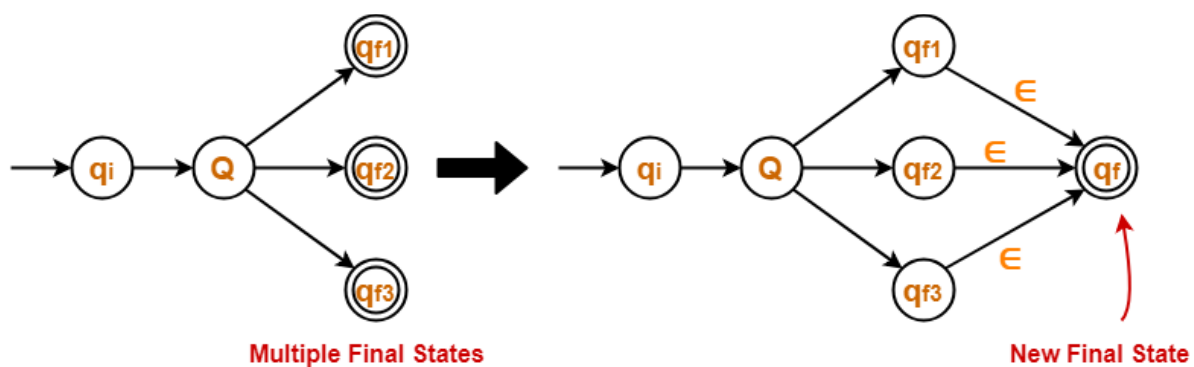
### Step-02:

#### Thumb Rule

There must exist only one final state in the DFA.

- If there exists multiple final states in the DFA, then convert all the final states into non-final states and create a new single final state.

#### Example-



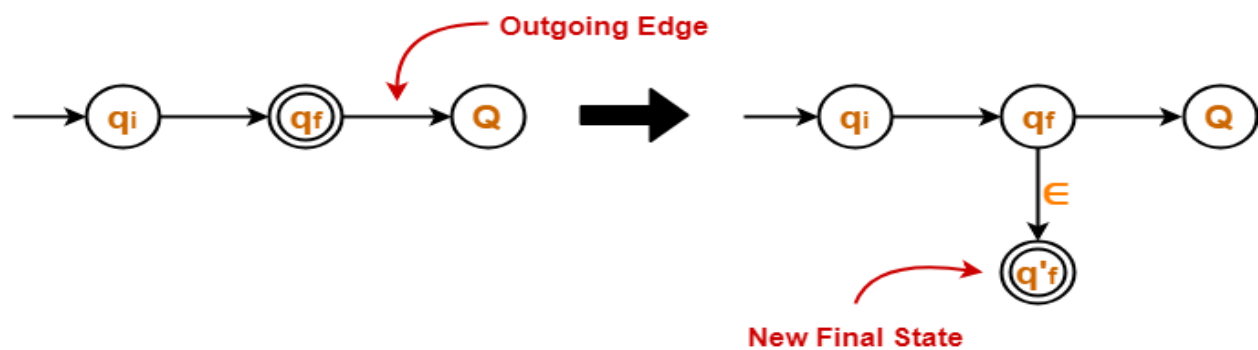
### Step-03:

#### Thumb Rule

The final state of the DFA must not have any outgoing edge.

- If there exists any outgoing edge from the final state, then create a new final state having no outgoing edge from it.

#### Example-



#### **Step-04:**

- Eliminate all the intermediate states one by one.
- These states may be eliminated in any order.

In the end,

- Only an initial state going to the final state will be left.
- The cost of this transition is the required regular expression.

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#### **NOTE**

The state elimination method can be applied to any finite automata.

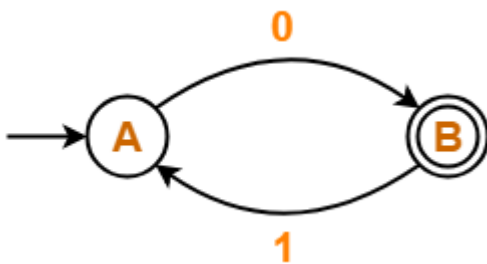
(NFA,  $\epsilon$ -NFA, DFA etc)

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### **PRACTICE PROBLEMS BASED ON CONVERTING DFA TO REGULAR EXPRESSION-**

#### **Problem-01:**

Find regular expression for the following DFA-

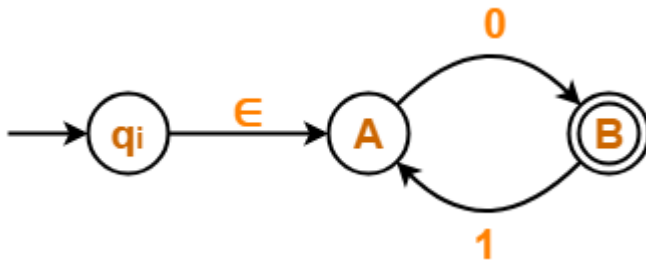


#### **Solution-**

##### **Step-01:**

- Initial state A has an incoming edge.
- So, we create a new initial state  $q_i$ .

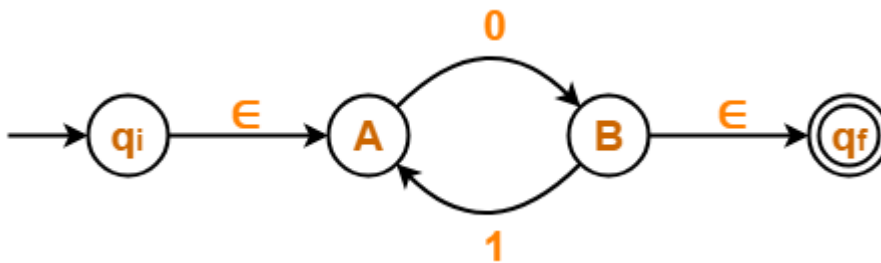
The resulting DFA is-



### Step-02:

- Final state B has an outgoing edge.
- So, we create a new final state  $q_f$ .

The resulting DFA is-



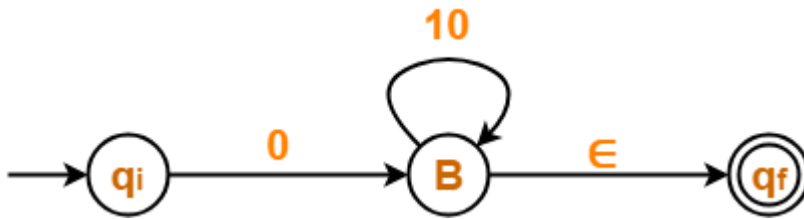
### Step-03:

Now, we start eliminating the intermediate states.

First, let us eliminate state A.

- There is a path going from state  $q_i$  to state B via state A.
- So, after eliminating state A, we put a direct path from state  $q_i$  to state B having cost  $\epsilon.0 = 0$
- There is a loop on state B using state A.
- So, after eliminating state A, we put a direct loop on state B having cost  $1.0 = 10$ .

Eliminating state A, we get-



#### Step-04:

Now, let us eliminate state B.

- There is a path going from state  $q_i$  to state  $q_f$  via state B.
- So, after eliminating state B, we put a direct path from state  $q_i$  to state  $q_f$  having cost 0.  
 $(10)^*.\epsilon = 0(10)^*$

Eliminating state B, we get-



From here,

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**Regular Expression =  $0(10)^*$**

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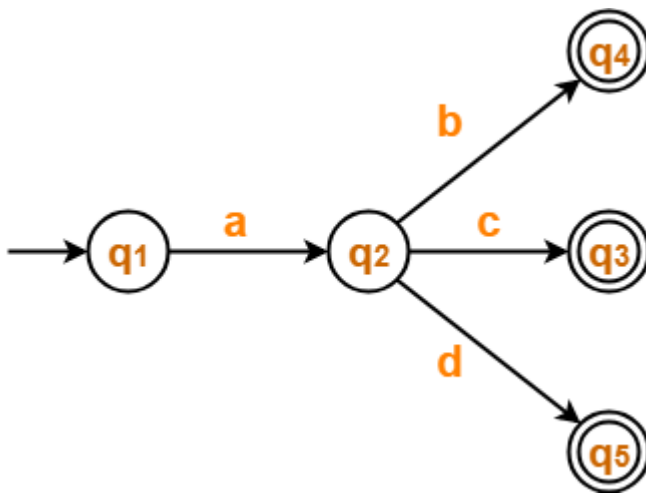
#### NOTE-

In the above question,

- If we first eliminate state B and then state A, then regular expression would be  $= (01)^*0$ .
- This is also the same and correct.

#### Problem-02:

Find regular expression for the following DFA-

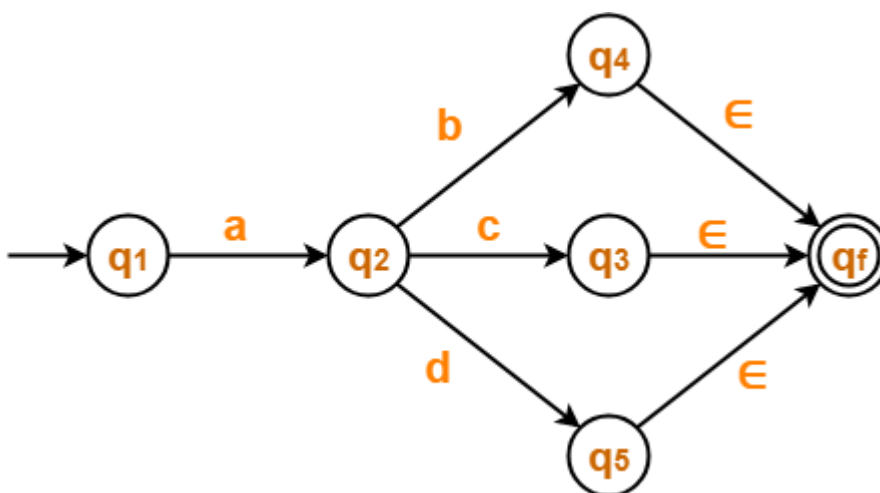


**Solution-**

**Step-01:**

- There exist multiple final states.
- So, we convert them into a single final state.

The resulting DFA is-

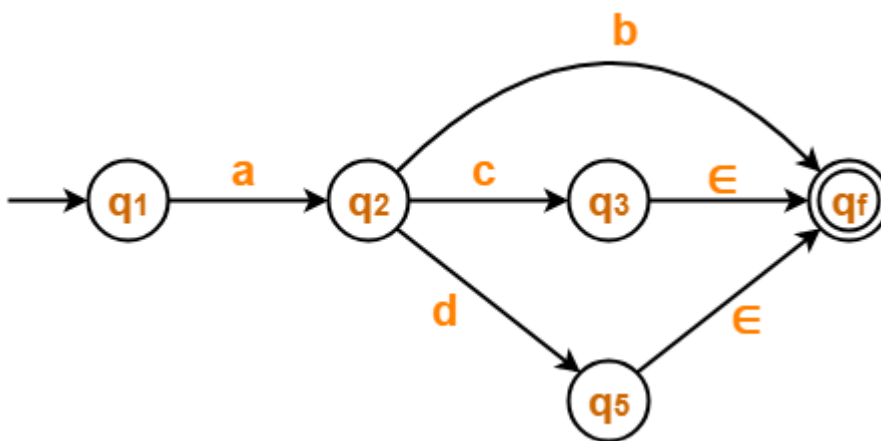


### Step-02:

Now, we start eliminating the intermediate states.

First, let us eliminate state  $q_4$ .

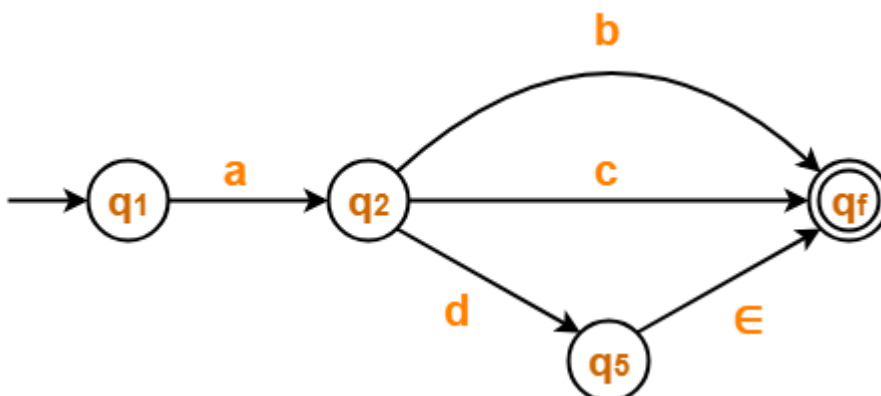
- There is a path going from state  $q_2$  to state  $q_f$  via state  $q_4$ .
- So, after eliminating state  $q_4$ , we put a direct path from state  $q_2$  to state  $q_f$  having cost  $b. \epsilon = b$ .



### Step-03:

Now, let us eliminate state  $q_3$ .

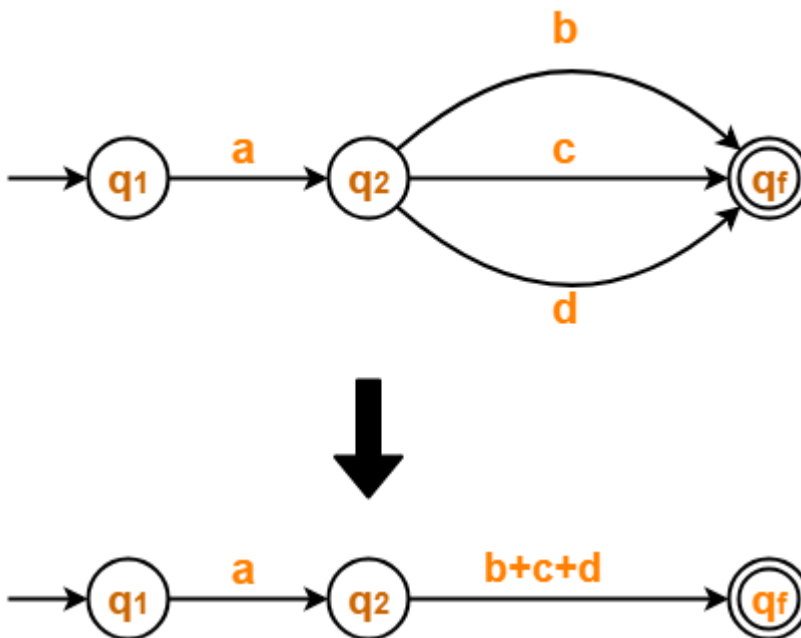
- There is a path going from state  $q_2$  to state  $q_f$  via state  $q_3$ .
- So, after eliminating state  $q_3$ , we put a direct path from state  $q_2$  to state  $q_f$  having cost  $c. \epsilon = c$ .



#### Step-04:

Now, let us eliminate state  $q_5$ .

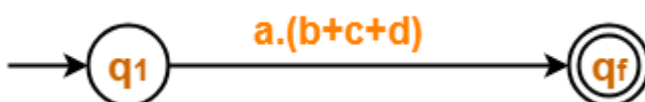
- There is a path going from state  $q_2$  to state  $q_f$  via state  $q_5$ .
- So, after eliminating state  $q_5$ , we put a direct path from state  $q_2$  to state  $q_f$  having cost  $d \in \mathbb{R}$ .



#### Step-05:

Now, let us eliminate state  $q_2$ .

- There is a path going from state  $q_1$  to state  $q_f$  via state  $q_2$ .
- So, after eliminating state  $q_2$ , we put a direct path from state  $q_1$  to state  $q_f$  having cost  $a \cdot (b+c+d)$ .



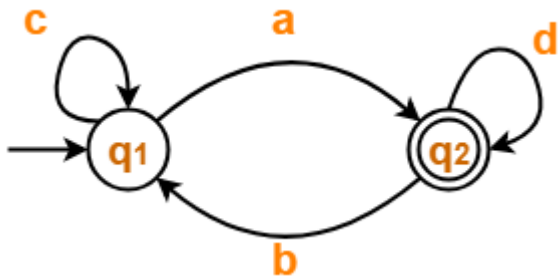
From here,



Regular Expression =  $a(b+c+d)$

**Problem-03:**

Find regular expression for the following DFA-

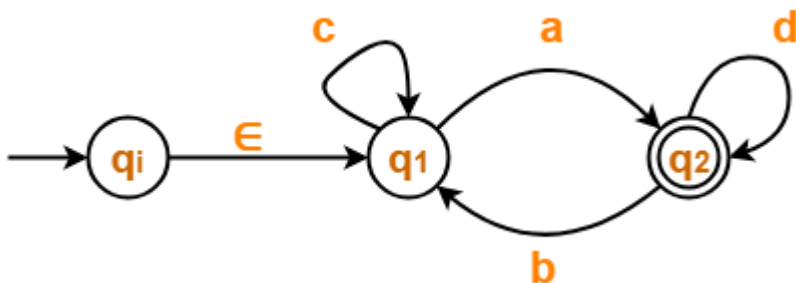


**Solution-**

**Step-01:**

- Initial state  $q_1$  has an incoming edge.
- So, we create a new initial state  $q_i$ .

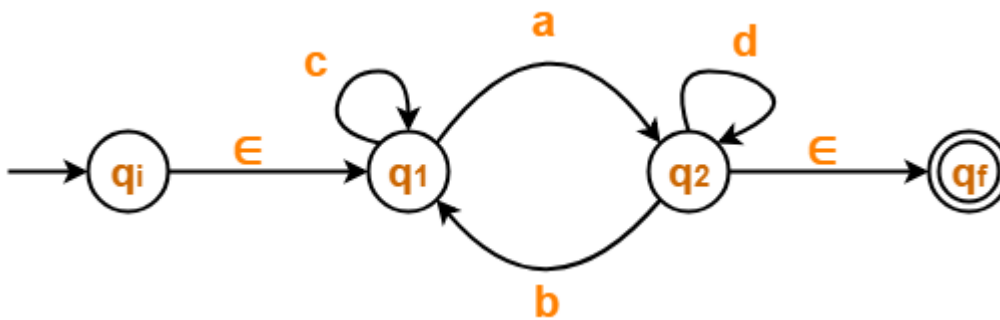
The resulting DFA is-



**Step-02:**

- Final state  $q_2$  has an outgoing edge.
- So, we create a new final state  $q_f$ .

The resulting DFA is-



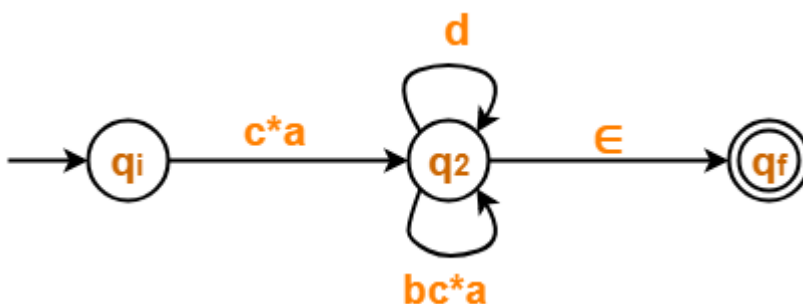
### Step-03:

Now, we start eliminating the intermediate states.

First, let us eliminate state  $q_1$ .

- There is a path going from state  $q_i$  to state  $q_2$  via state  $q_1$ .
- So, after eliminating state  $q_1$ , we put a direct path from state  $q_i$  to state  $q_2$  having cost  $\epsilon.c*.a = c*a$
- There is a loop on state  $q_2$  using state  $q_1$ .
- So, after eliminating state  $q_1$ , we put a direct loop on state  $q_2$  having cost  $b.c*.a = bc*a$

Eliminating state  $q_1$ , we get-

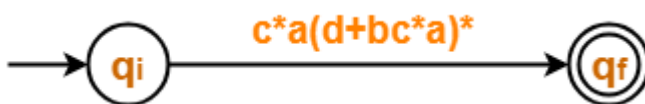


#### Step-04:

Now, let us eliminate state  $q_2$ .

- There is a path going from state  $q_i$  to state  $q_f$  via state  $q_2$ .
- So, after eliminating state  $q_2$ , we put a direct path from state  $q_i$  to state  $q_f$  having cost  $c*a(d+bc*a)^*\epsilon = c*a(d+bc*a)^*$

Eliminating state  $q_2$ , we get-



From here,

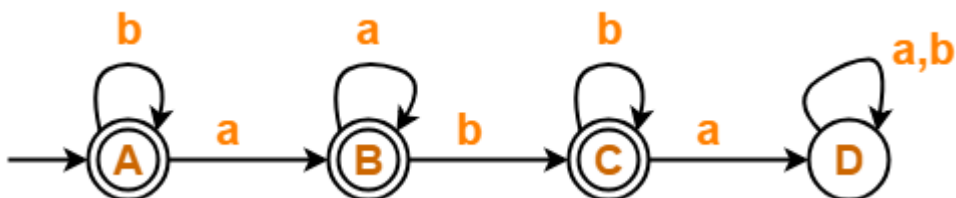
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**Regular Expression =  $c*a(d+bc*a)^*$**

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#### Problem-04:

Find regular expression for the following DFA-

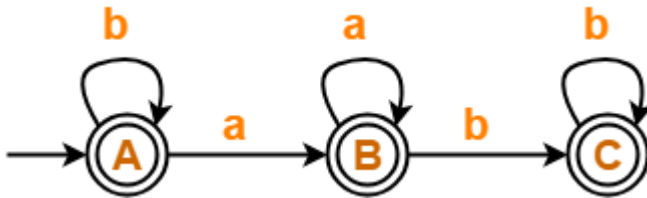


#### Solution-

##### Step-01:

- State D is a dead state as it does not reach to any final state.
- So, we eliminate state D and its associated edges.

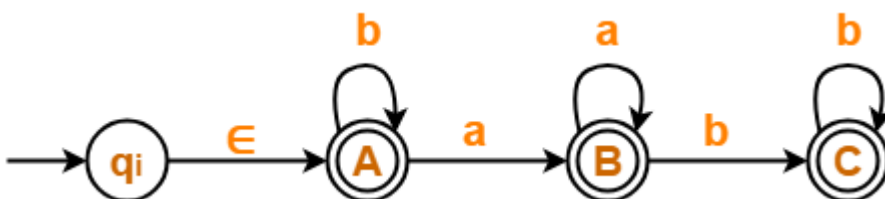
The resulting DFA is-



### Step-02:

- Initial state A has an incoming edge (self loop).
- So, we create a new initial state  $q_i$ .

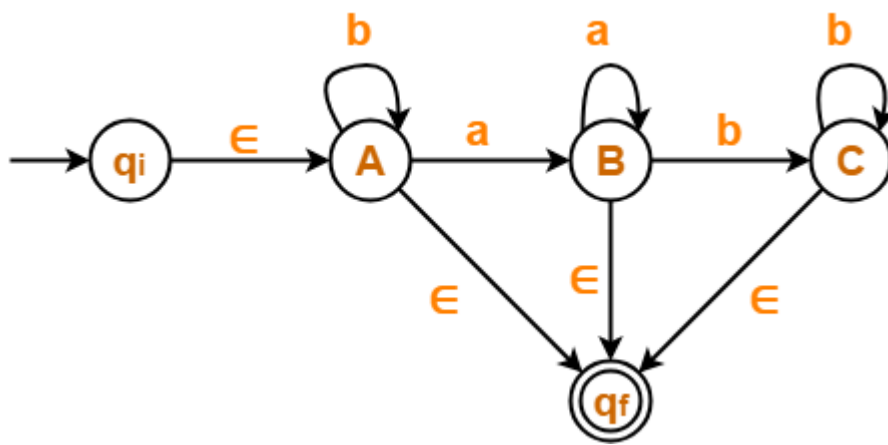
The resulting DFA is-



### Step-03:

- There exist multiple final states.
- So, we convert them into a single final state.

The resulting DFA is-



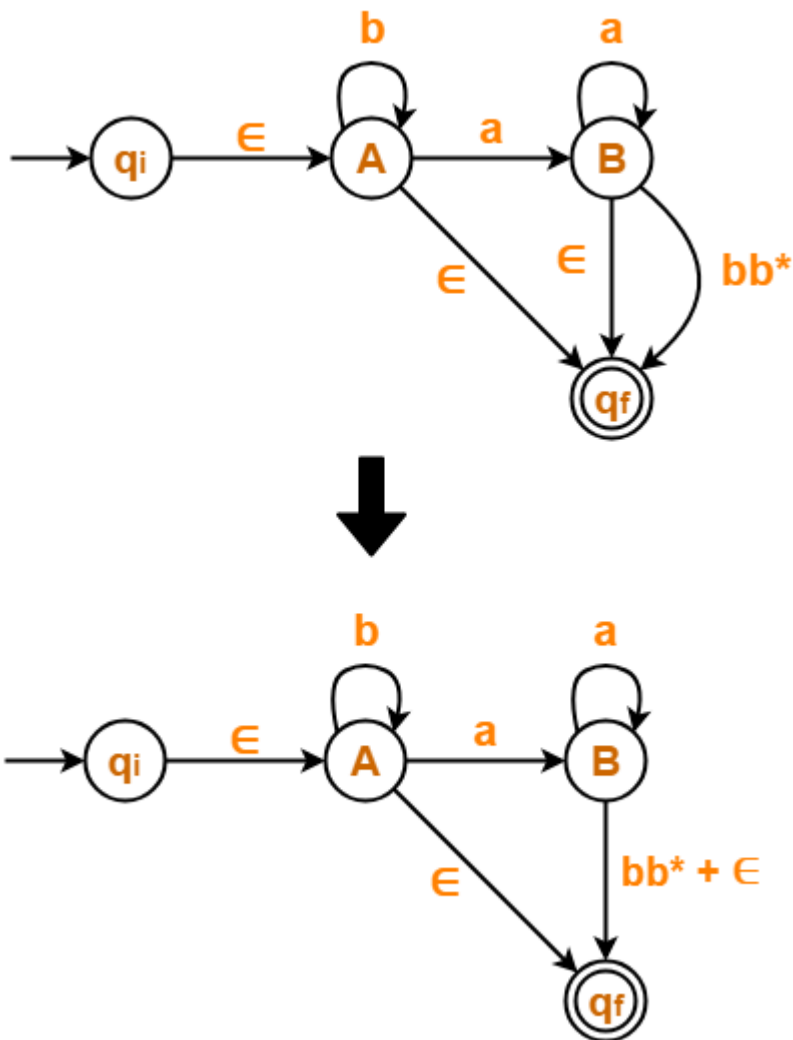
#### Step-04:

Now, we start eliminating the intermediate states.

First, let us eliminate state C.

- There is a path going from state B to state  $q_f$  via state C.
- So, after eliminating state C, we put a direct path from state B to state  $q_f$  having cost  $b.b^*.\epsilon$   
 $= bb^*$

Eliminating state C, we get-

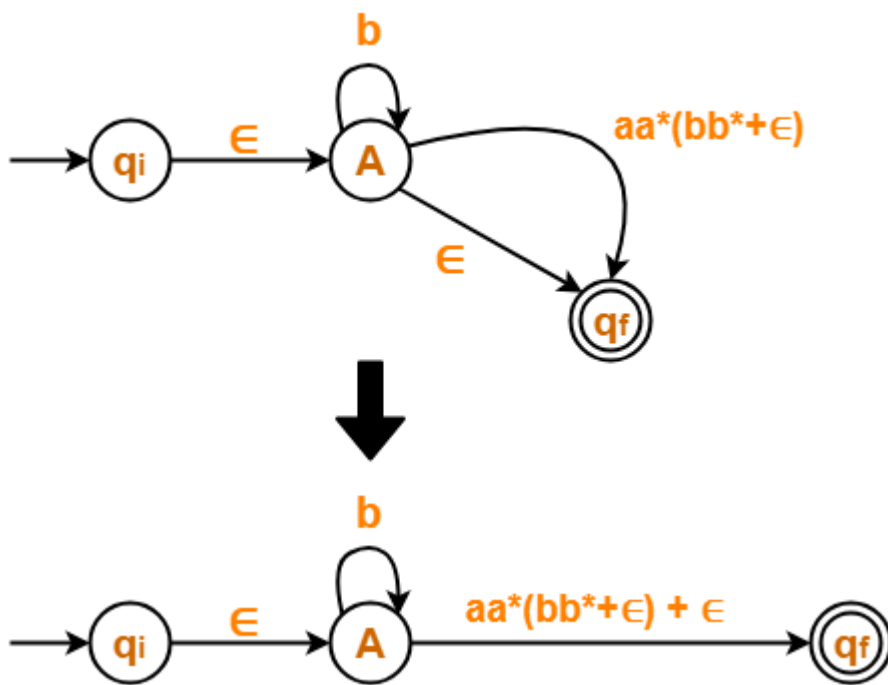


#### Step-05:

Now, let us eliminate state  $B$ .

- There is a path going from state  $A$  to state  $q_f$  via state  $B$ .
- So, after eliminating state  $B$ , we put a direct path from state  $A$  to state  $q_f$  having cost  $a.a^*.(bb^* + \epsilon) = aa^*(bb^* + \epsilon)$ .

Eliminating state  $B$ , we get-

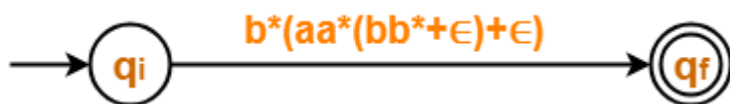


#### Step-06:

Now, let us eliminate state A.

- There is a path going from state  $q_i$  to state  $q_f$  via state A.
- So, after eliminating state A, we put a direct path from state  $q_i$  to state  $q_f$  having cost  $\epsilon.b^*$ .  
 $(aa^*(bb^*+\epsilon)+\epsilon) = b^*(aa^*(bb^*+\epsilon)+\epsilon)$

Eliminating state A, we get-



From here,

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**Regular Expression =  $b^*(aa^*(bb^*+\epsilon)+\epsilon)$**

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We know,  $bb^* + \epsilon = b^*$

So, we can also write-

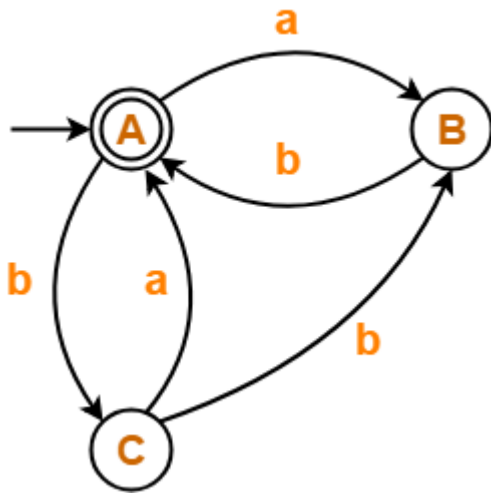
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**Regular Expression =  $b^*(aa^*b^*+\epsilon)$**

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**Problem-05:**

Find regular expression for the following DFA-



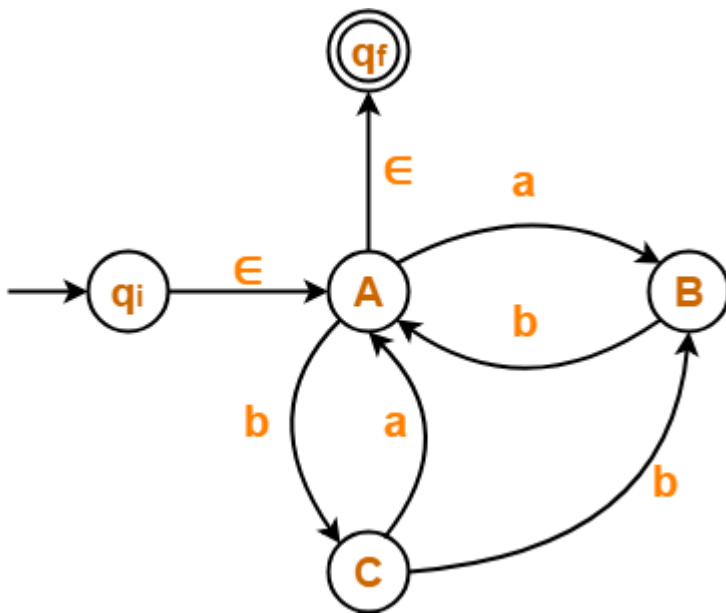
**Solution-**

**Step-01:**

- Since initial state A has an incoming edge, so we create a new initial state  $q_i$ .
- Since final state A has an outgoing edge, so we create a new final state  $q_f$ .

The resulting DFA is-





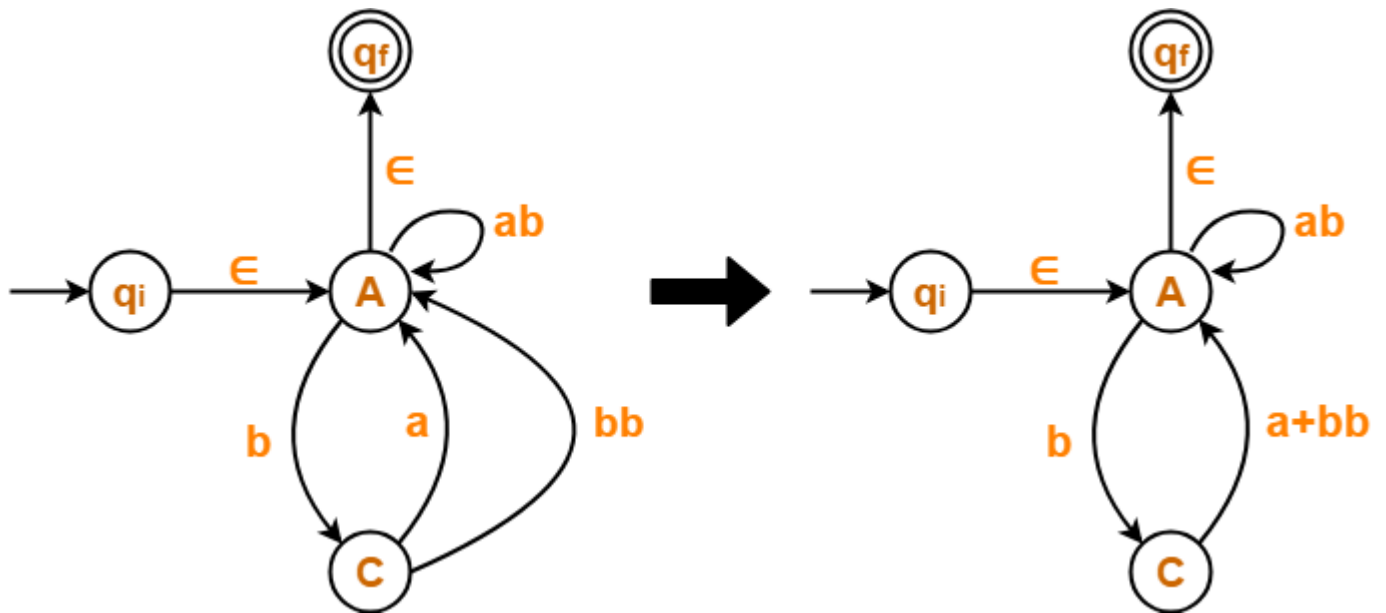
### Step-02:

Now, we start eliminating the intermediate states.

First, let us eliminate state B.

- There is a path going from state C to state A via state B.
- So, after eliminating state B, we put a direct path from state C to state A having cost  $b.b = bb$ .
- There is a loop on state A using state B.
- So, after eliminating state B, we put a direct loop on state A having cost  $a.b = ab$ .

Eliminating state B, we get-

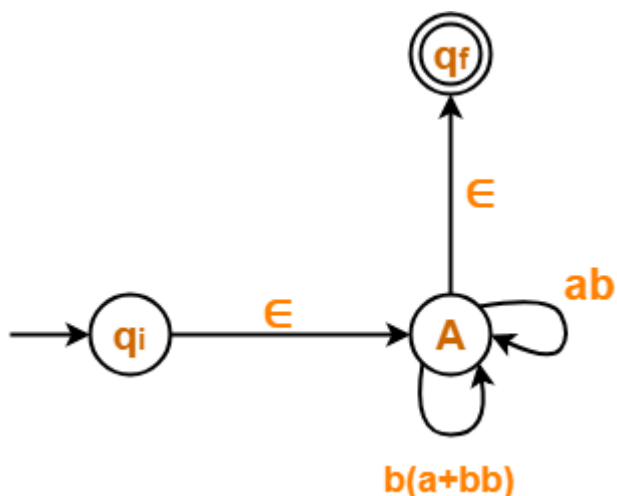


### Step-03:

Now, let us eliminate state C.

- There is a loop on state A using state C.
- So, after eliminating state C, we put a direct loop on state A having cost  $b.(a+bb) = b(a+bb)$

Eliminating state C, we get-

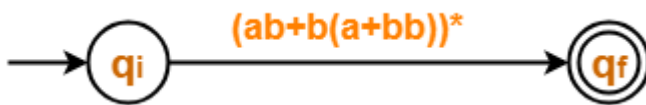


### Step-04:

Now, let us eliminate state A.

- There is a path going from state  $q_i$  to state  $q_f$  via state A.
- So, after eliminating state A, we put a direct path from state  $q_i$  to state  $q_f$  having cost  $\epsilon.(ab + b(a+bb))^*$

Eliminating state A, we get-

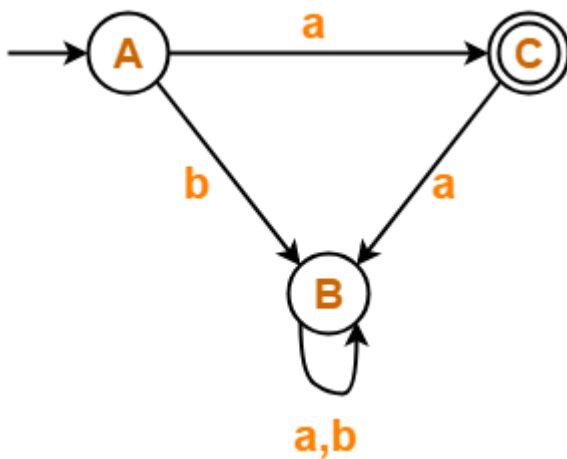


From here,

Regular Expression =  $(ab + b(a+bb))^*$

#### **Problem-06:**

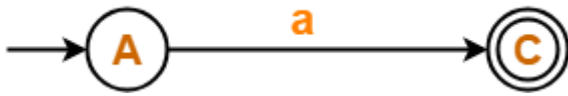
Find regular expression for the following DFA-



#### **Solution-**

- State B is a dead state as it does not reach to the final state.
- So, we eliminate state B and its associated edges.

The resulting DFA is-

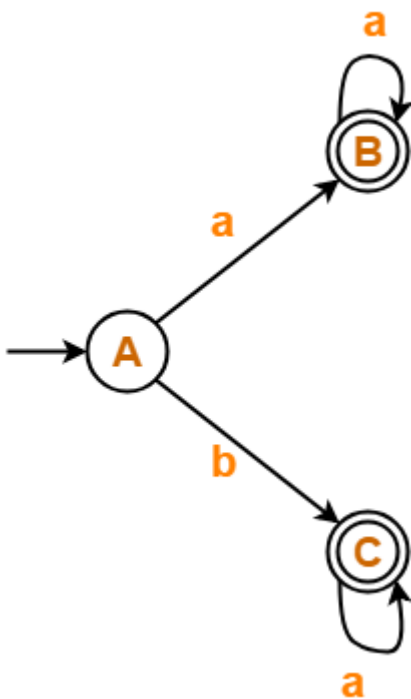


From here,

Regular Expression = a

**Problem-07:**

Find regular expression for the following DFA-



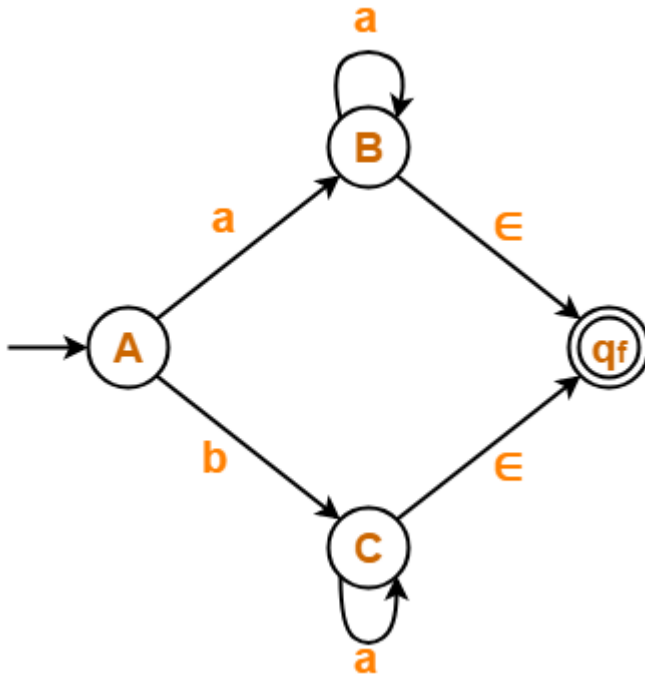
**Solution-**

**Step-01:**

- There exist multiple final states.

- So, we create a new single final state.

The resulting DFA is-



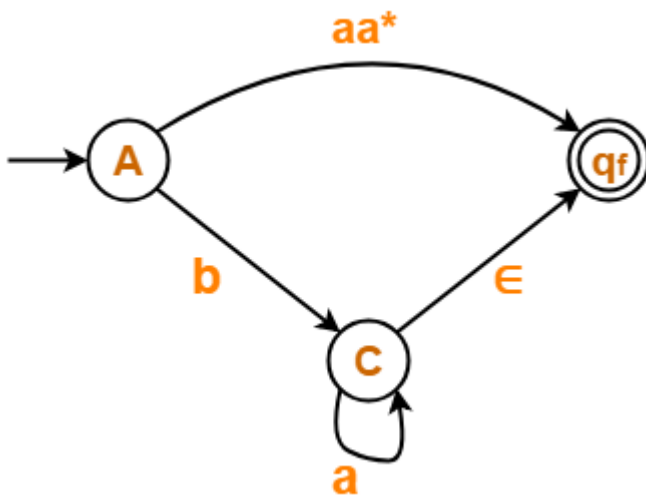
### Step-02:

Now, we start eliminating the intermediate states.

First, let us eliminate state B.

- There is a path going from state A to state  $q_f$  via state B.
- So, after eliminating state B, we put a direct path from state A to state  $q_f$  having cost  $a.a^*.ε$   
 $= aa^*.$

Eliminating state B, we get-

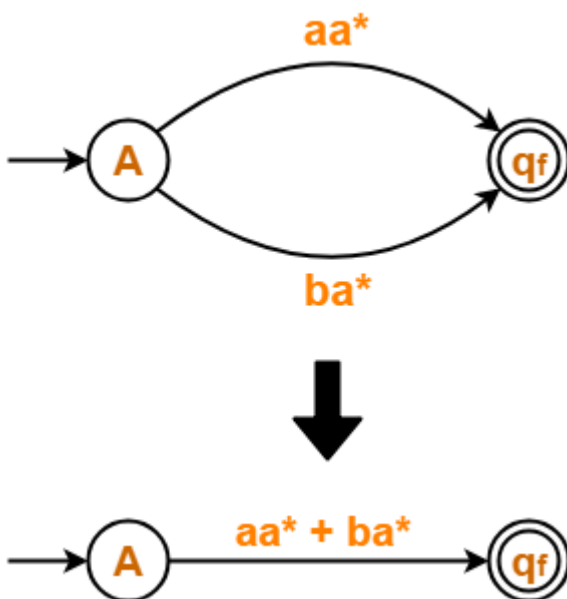


### Step-03:

Now, let us eliminate state C.

- There is a path going from state A to state  $q_f$  via state C.
- So, after eliminating state C, we put a direct path from state A to state  $q_f$  having cost  $b.a^*.ε = ba^*$ .

Eliminating state C, we get-



From here,

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**Regular Expression =  $aa^* + ba^*$**

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