



PRACTICAL FILE OF OPERATING SYSTEM

BTech: III Year

Department of Computer Science & Information Technology

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Branch & section : CSIT-1

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Year : III - YEAR

Department of Computer Science & Information Technology
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ACROPOLIS INSTITUTE OF TECHNOLOGY & RESEARCH, INDORE
Department of Computer Science & Information Technology

Certificate

This is to certify that the experimental work entered in this journal as per the Bachelor of Technology in Computer science and Technology III year syllabus prescribed by the RGPV was done by Mr. ANIKET RAI in Vth semester in the Laboratory of this institute during the academic year 2022- 2023

Signature of Head

Signature of the Faculty

ACROPOLIS INSTITUTE OF TECHNOLOGY & RESEARCH, INDORE
GENERAL INSTRUCTIONS FOR
LABORATORY CLASSES

DO'S

- Without Prior permission do not enter into the Laboratory.
- While entering into the LAB students should wear their ID cards.
- The Students should come with proper uniform.
- Students should maintain silence inside the laboratory.
- After completing the laboratory exercise, make sure to shutdown the system properly.

DONT'S

- Students bringing the bags inside the laboratory.
- Students using the computers in an improper way.
- Students scribbling on the desk and mishandling the chairs.
- Students using mobile phones inside the laboratory.
- Students making noise inside the laboratory

SYLLABUS

CS-502 – Operating System

Branch: Computer Science Information Technology V Semester

Course: CSIT 502 Operating System

Unit I

Introduction to System Programs & Operating Systems, Evolution of Operating System (mainframe, desktop, multiprocessor, Distributed, Network Operating System, Clustered & Handheld System), Operating system services, Operating system structure, System Call & System Boots, Operating system design & Implementations, System protection, Buffering & Spooling. Types of Operating System: Bare machine, Batch Processing, Real Time, Multitasking & Multiprogramming, time-sharing system.

Unit II

File: concepts, access methods, free space managements, allocation methods, and directory systems, protection, organization ,sharing & implementation issues, Disk & Drum Scheduling, I/O devices organization, I/O devices organization, I/O buffering, I/O Hardware, Kernel I/O subsystem, Transforming I/O request to hardware operations. Device Driver: Path managements, Sub module, Procedure, Scheduler, Handler, Interrupt Service Routine. File system in Linux & Windows

Unit III

Process: Concept, Process Control Blocks (PCB), Scheduling criteria Preemptive & non Preemptive process scheduling, Scheduling algorithms, algorithm evaluation, multiple processor scheduling, real time scheduling, operations on processes, threads; inter process communication, precedence graphs, critical section problem, semaphores, and classical problems of synchronization. Deadlock: Characterization, Methods for deadlock handling, deadlock prevention, deadlock avoidance, deadlock detection, recovery from deadlock, Process Management in Linux.

Unit IV

Memory Hierarchy, Concepts of memory management, MFT & MVT, logical and physical address space, swapping, contiguous and non-contiguous allocation, paging, segmentation, and paging combined with segmentation. Structure & implementation of Page table. Concepts of virtual memory, Cache Memory Organization, demand paging, page replacement algorithms, allocation of frames, thrashing, and demand segmentation.

Unit V

Distributed operating system:-Types, Design issues, File system, Remote file access, RPC, RMI, Distributed Shared Memory(DSM), Basic Concept of Parallel Processing & Concurrent Programming Security & threats protection: Security violation through Parameter, Computer Worms & Virus, Security Design Principle, Authentications, Protection Mechanisms. introduction to Sensor network and parallel operating system. Case study of Unix, Linux & Windows.

HARDWARE REQUIREMENTS:

Processors - 2.0 GHz or Higher

RAM - 256 MB or Higher

Hard Disk - 20 GB or Higher

SOFTWARE REQUIREMENTS:

Linux: Ubuntu / OpenSUSE / Fedora / Red Hat / Debian / Mint OS

WINDOWS: XP/7

Linux could be loaded in individual PCs.

RATIONALE:

The purpose of this subject is to cover the underlying concepts Operating System .This syllabus provides a comprehensive introduction of Operating System, Process Management, Memory Management, File Management and I/O management.

PREREQUISITE:

The students should have general idea about Operating System Concept, types of Operating System and their functionality.

Lab Plan

Operating System

CS-502

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ACROPOLIS INSTITUTE OF TECHNOLOGY & RESEARCH, INDORE

Name of Department –CI

Name of Laboratory – Operating System

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10.					
11.					

Experiment-1

FCFS SCHEDULING

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c++ program to implement the FCFS SCHEDULING.

FACILITIES REQUIRED

a) Facilities Required Doing The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	WINDOWS XP/7	

b) Concept of FCFS:

- Jobs are executed on first come, first serve basis.
- Easy to understand and implement.
- Poor in performance as average wait time is high.

Process	Arrival Time	Execute Time	Service Time
P0	0	5	0
P1	1	3	5
P2	2	8	8
P3	3	6	16



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Set the waiting of the first process as '0' and its burst time as its turn around time

Step 5: for each process in the Ready Q calculate

(a) Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)

(b) Turn around time for Process(n)= waiting time of Process(n)+ Burst time for process(n)

Step 6: Calculate

(a) Average waiting time = Total waiting Time / Number of process

(b) Average Turnaround time = Total Turnaround Time / Number of process
Step 7: Stop the process

d) Program:

➤ FCFS with 0 arrival time

```
#include<iostream>
using namespace std;
class FCFS
{
public:
void findCompletionTime(int process[],int n,int burst_time[],int comp_time[])
{
comp_time[0]=burst_time[0];
cout<<"Completion Time:"<<endl
cout<<comp_time[0]<<endl;
for(int i=1;i<n;i++)
{
comp_time[i]=comp_time[i-1]+burst_time[i];
cout<<comp_time[i]<<endl;
}
}
void findTurnAroundTime(int process[],int n,int ct[])
{
int tat[n],average;
cout<<"Turn around time:"<<endl;
for(int i=0;i<n;i++)
{
tat[i]=ct[i];
cout<<tat[i]<<endl;
average=average+tat[i];
}
cout<<"Average Turn around time:"<<average/n<<endl;
}
void findWaitingTime(int process[],int n,int ct[],int bt[],int wt[])
{
int tat[n],average=0;
cout<<"Waiting Time:"<<endl;
for(int i=0;i<n;i++)
{
tat[i]=ct[i];
wt[i]=tat[i]-bt[i];
average=average+wt[i];
}
```

```

cout<<wt[i]<<endl;
}
average=average/n;
cout<<"Average waiting time:"<<average;
}
};
int main()
{
FCFS f;
int n;
cout<<"Enter no of processes:";
cin>>n;
int process[n];
for(int i=1;i<=n;i++)
{
cin>>process[i];
}
int burst_time[]={2,3,4,5};
int comp_time[n],waiting_time[n],tat[n];
f.findCompletionTime(process,n,burst_time,comp_time);
f.findTurnAroundTime(process,n,comp_time);
f.findWaitingTime(process,n,comp_time,burst_time,waiting_time);
}

```

➤ FCFS with varying arrival time

```

#include<iostream>
#include<algorithm>
#include<string.h>
using namespace std;
void findCompletionTime(int process[],int n,int arrival_time[],int burst_time[],int ct[])
{
for(int i=0;i<n;i++)
{
if(i==0)
{
ct[0]=burst_time[0];
}
else
ct[i]=ct[i-1]+burst_time[i];
}
}

```

```

}
}
void findTurnAroundTime(int process[],int n,int ct[],int at[],int tat[])
{
for(int i=0;i<n;i++)
{
tat[i]=ct[i]-at[i];} }
void findWaitingTime(int process[],int n,int tat[],int bt[],int wt[])
{
for(int i=0;i<n;i++)
{
wt[i]=tat[i]-bt[i];} }
void display(int process[],int n,int at[],int bt[],int ct[],int tat[],int wt[])
{
cout<<"Process ID   Arrival Time   Burst Time       Completion Time   TurnAround Time
Waiting Time"<<endl;
for(int i=0;i<n;i++)
{
cout<<process[i]<<"           "<<at[i]<<"           "<<bt[i]<<"           "<<ct[i]<<"
"<<tat[i]<<"           "<<wt[i]<<endl;
}
}
void sortAccordingArrivalTime(int process[],int n,int at[], int bt[])
{
bool swapped;
int temp;
int stemp;
for (int i = 0; i < n; i++)
{
swapped = false;
for (int j = 0; j < n - i - 1; j++)
{
if (at[j] > at[j + 1])
{
temp = at[j];
at[j] = at[j + 1];
at[j + 1] = temp;
temp = bt[j];
bt[j] = bt[j + 1];
bt[j + 1] = temp;
temp = process[j];

```

```

process[j] = process[j + 1];
process[j + 1] = temp;
swapped = true;}}
if (swapped == false)
{
break;}}}}
void average(int process[],int n,int tat[],int wt[])
{
float avgtat=0,avgwt=0;
for(int i=0;i<n;i++)
{
avgtat=avgtat+tat[i];
avgwt=avgwt+wt[i];
}
cout<<"Average TurnAround Time:"<<avgtat/float(n)<<endl;
cout<<"Average Waiting Time:"<<avgwt/float(n)<<endl;
}
int main()
{
int n;
cout<<"Enter no. of processes:"<<endl;
cin>>n;
int process[n];
for(int i=0;i<n;i++)
{
cin>>process[i];
}
cout<<"Enter arrival time:"<<endl;
int arrival_time[n];
for(int i=0;i<n;i++)
{
cin>>arrival_time[i];
}
cout<<"Enter burst time:"<<endl;
int burst_time[n];
for(int i=0;i<n;i++)
{
cin>>burst_time[i];
}
sortAccordingArrivalTime(process,n,arrival_time,burst_time);
int ct[n],tat[n],wt[n];

```

```
findCompletionTime(process,n,arrival_time,burst_time,ct);
findTurnAroundTime(process,n,ct,arrival_time,tat);
findWaitingTime(process,n,tat,burst_time,wt);
display(process,n,arrival_time,burst_time,ct,tat,wt);
average(process,n,tat,wt);
return 0;
}
```

e) Output:

FCFS with 0 arrival time

```
Enter the number of process:4
Enter the process ID and burst time:1
3
Enter the process ID and burst time:2
0
Enter the process ID and burst time:3
4
Enter the process ID and burst time:4
1
p_ID completion TAT Waiting
1      3      3      0
2      3      3      3
3      7      7      3
4      8      8      7
```

FCFS with varying arrival time

```
Enter no. of processes:
4
1
2
3
4
Enter arrival time:
0
1
3
2
Enter burst time:
2
3
4
5
Process ID      Arrival Time      Burst Time      Completion Time      TurnAround Time      Waiting Time
1              0              2              2              2              0
2              1              3              5              4              1
4              2              5              10             8              3
3              3              4              14             11             7
Average TurnAround Time:6.25
Average Waiting Time:2.75
```

f) Result:

FCFS with 0 arrival time

Average Waiting Time – 3.25

Average Turnaround Time – 5.25

FCFS with varying arrival time

Average Waiting Time – 2.75

Average Turnaround Time -6.25

Experiment-2

SJF Scheduling

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c++ program to implement SJF CPU Scheduling Algorithm.

FACILITIES REQUIRED

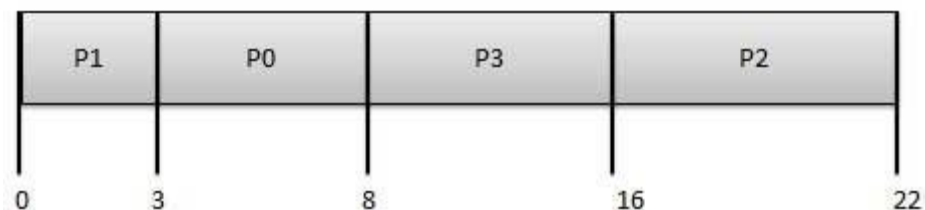
a) Facilities Required Doing The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	WINDOWS XP/7	

b) Concept of SJF:

- Best approach to minimize waiting time.
- Processer should know in advance how much time process will take.

Process	Arrival Time	Execute Time	Service Time
P0	0	5	0
P1	1	3	3
P2	2	8	8
P3	3	6	16



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time Step 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to highest burst time.

Step 5: Set the waiting time of the first process as '0' and its turnaround time as its burst time.

Step 6: For each process in the ready queue, calculate

Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n- 1)

Turnaround time for Process(n)= waiting time of Process(n)+ Burst time for process(n)

Step 7: Calculate

Average waiting time = Total waiting Time / Number of process

Average Turnaround time = Total Turnaround Time / Number of process Step 8: Stop the process

d) Program:

```
#include <bits/stdc++.h>
using namespace std;
struct data
{
    int pid;
    int bt;
    int at;
    int wt;
    int tat;
};
bool cmp( data x , data y)
{
    return ( x.bt<y.bt);
}

int main()
{
    data a[30];
    int i,k,n;
    float wtavg , tatavg;
    float avgwt,avgtat;
    cout<<"Enter the Number of process"<<endl;
    cin>>n;
    for(i=0;i<n;i++)
    {
        a[i].pid =i;
        //a[i].at=0;
        cout<<"Enter The Burst time for the process "<<i<<endl;
        cin>>a[i].bt;
        cout<<"Enter The Arrival time for the process "<<i<<endl;
        cin>>a[i].at;
    }

    sort( a , a+n , cmp);
    a[0].wt =0;
    a[0].tat = a[0].bt;
    for(i=1;i<n;i++)
    {
```

```

a[i].wt = a[i-1].wt + a[i-1].bt;
a[i].tat = a[i-1].tat + a[i].bt;
}
cout<<endl<<"P.id\tAt\tBt\tWt\tTat"<<endl;
for(i=0;i<n;i++)
{
avgwt=avgwt+a[i].wt;
avgtat=avgtat+a[i].tat;
cout<<a[i].pid<<"\t"<<a[i].at<<"\t"<<a[i].bt<<"\t"<<a[i].wt<<"\t"<<a[i].tat<<endl;
}
avgwt=avgwt/n;
avgtat=avgtat/n;
cout<<"Average waiting time:"<<avgwt<<endl;
cout<<"Average turn around time time:"<<avgtat<<endl;
};

```

e) Output:

```

Enter the Number of process
4
Enter The Burst time for the process 0
2
Enter The Arrival time for the process 0
4
Enter The Burst time for the process 1
5
Enter The Arrival time for the process 1
1
Enter The Burst time for the process 2
3
Enter The Arrival time for the process 2
1
Enter The Burst time for the process 3
5
Enter The Arrival time for the process 3
6

P.id    At    Bt    Wt    Tat
0       4     2     0     2
2       1     3     2     5
1       1     5     5     10
3       6     5     10    15
Average waiting time:4.25
Average turn around time time:8

```

f) Result:

Average Waiting Time : 4.25

Average Turnaround Time : 8

Experiment-3

SRTF Scheduling

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Date of Experiment
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement SRTF scheduling.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of SRTF Scheduling:

- Associate with each process the length of its next CPU burst. Use these lengths to schedule the process with the shortest time.
- Two schemes:
 1. non pre- emptive – once CPU given to the process it cannot be preempted until completes its CPU burst.
 2. Preemptive – if a new process arrives with CPU burst length less than remaining time of current executing process, preempt. This scheme is known as the Shortest-Remaining-Time-First (SRTF).

Example of Preemptive SJF

Process	Arrival Time	Burst Time
P1	0.0	7
P2	2.0	4
P3	4.0	1
P4	5.0	4
SJF (preemptive)		

P1	P2	P3	P2	P4	P1
0	2	4	5	7	11
					16

c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: For each process in the ready Q, Accept Arrival time

Step 4: Start the Ready Q according the shortest Burst time by sorting according to lowest to Highest burst time.

Step 5: Set the waiting time of the first process in Sorted Q as '0'.

Step 6: After every unit of time compare the remaining time of currently executing process (RT) and Burst time of newly arrived process (BT_n).

Step 7: If the burst time of newly arrived process (BT_n) is less than the currently executing process (RT) the processor will preempt the currently executing process and starts executing newly arrived process

Step 7: Calculate

(c) Average waiting time = Total waiting Time / Number of process

(d) Average Turnaround time = Total Turnaround Time / Number of process

Step 8: Stop the process

d) Program:

```
#include <iostream>
#include<algorithm>
using namespace std;

void findCompletionTime(int process[],int n,int burst_time[],int ct[])
{
    ct[0]=burst_time[0];
    for(int i=1;i<n;i++)
    {
        ct[i]=ct[i-1]+burst_time[i];
    }
}

void findWaitingTime(int process[],int n,int at[],int burst_time[],int wait[])
{
    wait[0]=0;
    for(int i=1;i<n;i++)
    {
        wait[i] = wait[i-1]+burst_time[i-1];
    }
}

void TurnaroundTime(int process[],int n,int at[],int burst_time[],int tat[],int wait[])
{
    tat[0]= burst_time[0];
    for(int i=1;i<n;i++)
    {
        tat[i] = tat[i-1]+burst_time[i];
    }
}

void display(int process[],int n,int at[],int bt[],int ct[],int tat[],int wt[])
{
    cout<<"Process ID    Arrival Time    Burst Time        Completion Time    TurnAround Time
Waiting Time"<<endl;
    for(int i=0;i<n;i++)
    {
        cout<<process[i]<<"          "<<at[i]<<"          "<<bt[i]<<"          "<<ct[i]<<"
"<<tat[i]<<"          "<<wt[i]<<endl;
    }
}

void average(int process[],int n,int tat[],int wt[])
{
    float avgtat=0,avgwt=0;
```



```

for(int i=0;i<n;i++)
{
    avgtat=avgtat+tat[i];
    avgwt=avgwt+wt[i];
}
cout<<"Average TurnAround Time:"<<avgtat/float(n)<<endl;
cout<<"Average Waiting Time:"<<avgwt/float(n)<<endl;

}
int main()
{
    int n;
    cout<<"Enter the number of process";
    cin>>n;

    int process[n],wait[n],tat[n],burst_time[n],comp[n],at[n]={0},bt[n];

    for(int i=0;i<n;i++)
    {
        cin>>process[i];
    }
    cout<<"Enter the burst time";
    for(int i=0;i<n;i++)
    {
        cin>>burst_time[i];
        bt[i]=burst_time[i];
    }
    sort(burst_time,burst_time+n);
    findWaitingTime(process,n,at,burst_time,wait);
    TurnaroundTime(process,n,at,burst_time,tat,wait);
    findCompletionTime(process,n,burst_time,comp);
    display(process,n,at,bt,comp,tat,wait);
    average(process,n,tat,wait);

    return 0;
}

```

e) Output:

```
Enter the number of process5
1
2
3
4
5
Enter the burst time1
0
2
4
3
Process ID    Arrival Time    Burst Time    Completion Time    TurnAround Time    Waiting Time
1              0                1              0                  0                  0
2              0                0              1                  1                  0
3              0                2              3                  3                  1
4              0                4              6                  6                  3
5              0                3              10                 10                 6
Average TurnAround Time:4
Average Waiting Time:2
```

f) Result:

Average Waiting Time – 2

Average Turnaround Time - 4

Experiment-4

ROUND ROBIN Scheduling

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Date of Experiment
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement Round Robin scheduling.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

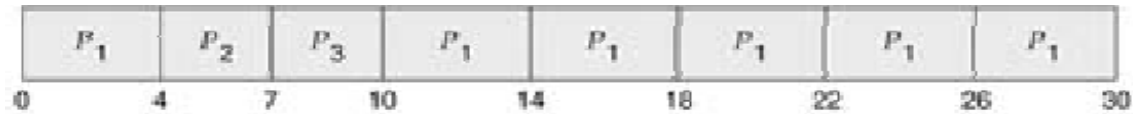
S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of Round Robin Scheduling:

This Algorithm is designed especially for time-sharing systems. A small unit of time, called time slices or **quantum** is defined. All runnable processes are kept in a circular queue. The CPU scheduler goes around this queue, allocating the CPU to each process for a time interval of one quantum. New processes are added to the tail of the queue. The CPU scheduler picks the first process from the queue, sets a timer to interrupt after one quantum, and dispatches the process. If the process is still running at the end of the quantum, the CPU is preempted and the process is added to the tail of the queue. If the process finishes before the end of the quantum, the process itself releases the CPU voluntarily Every time a process is granted the CPU, a **context switch** occurs, this adds overhead to the process execution time.

	Burst
Process	Time
P_1	24
P_2	3

P_2	3
Average	



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue and time quantum (or) time slice

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Calculate the no. of time slices for each process where

$$\text{No. of time slice for process}(n) = \text{burst time process}(n) / \text{time slice}$$

Step 5: If the burst time is less than the time slice then the no. of time slices =1.

Step 6: Consider the ready queue is a circular Q, calculate

(a) Waiting time for process(n) = waiting time of process(n-1) + burst time of process(n-1) + the time difference in getting the CPU from process(n-1)

(b) Turn around time for process(n) = waiting time of process(n) + burst time of process(n) + the time difference in getting CPU from process(n).

Step 7: Calculate

(g) Average waiting time = Total waiting Time / Number of process

(h) Average Turnaround time = Total Turnaround Time / Number of process

Step 8: Stop the process

d) Program:

```
#include<iostream>
using namespace std;
int main()
{
    int i,j,n,time,remain,flag=0,time_quantum;
    int wait_time=0,turnaround_time=0,process[10],at[10],bt[10],rt[10];
    cout<<"Enter Total Process:\t ";
    cin>>n;
    remain=n;
    for(i=0;i<n;i++)
    {
        cin>>process[i];
    }
    cout<<"Enter arrival time:";
    for(i=0;i<n;i++)
    {
        cin>>at[i];
    }
    cout<<"Enter burst time:";
    for(i=0;i<n;i++)
    {
        cin>>bt[i];
        rt[i]=bt[i];
    }
    cout<<"Enter Time Quantum:\t";
    cin>>time_quantum;
    cout<<"\n\nProcess\t|Turnaround Time|Waiting Time\n\n";
    for(time=0,i=0;remain!=0;)
    {
        if(rt[i]<=time_quantum && rt[i]>0)
        {
            time+=rt[i];
            rt[i]=0;
            flag=1;
        }
        else if(rt[i]>0)
        {
            rt[i]-=time_quantum;
            time+=time_quantum;
        }
        if(rt[i]==0 && flag==1)
        {
            remain--;
            cout<<process[i]<<"\t"<<time-at[i]<<"\t"<<time-at[i]-bt[i];
```

```
cout<<"\n";
wait_time+=time-at[i]-bt[i];
turnaround_time+=time-at[i];
flag=0;
}
if(i==n-1)
i=0;
else if(at[i+1]<=time)
i++;
else
i=0;
}
cout<<"\nAverage Waiting Time= \n"<<wait_time*1.0/n;
cout<<"\nAvg Turnaround Time =\n"<<turnaround_time*1.0/n;

return 0;
}
```

e) Output:

```
Enter Total Process:    5
1
2
3
4
5
Enter arrival time:1
0
2
3
5
Enter burst time:8
1
4
6
2
Enter Time Quantum:    2

Process |Turnaround Time|Waiting Time
2        3        2
5        4        2
3       11        7
4       16       10
1       20       12

Average Waiting Time=
6.6
Avg Turnaround Time =
10.8
```

f) Result:

Average Waiting Time – 6.6

Average Turnaround Time – 10.8

Experiment-5

PRIORITY SCHEDULING

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Date of Experiment
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To write c program to implement Priority scheduling.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of Priority Scheduling:

A priority is associated with each process, and the CPU is allocated to the process with the highest priority. Equal-priority processes are scheduled in FCFS order.

An SJF algorithm is simply a priority algorithm where the priority (p) is the inverse of the (predicted) next CPU burst. The larger the CPU burst, the lower the priority, and vice versa.

As an example, consider the following set of processes, assumed to have arrived at time 0, in the order

	Burst		Waiting	Turnaround
Process	Time	Priority	Time	Time
P_1	10	3	6	16
P_2	1	1	0	1
P_3	2	4	16	18
P_4	1	5	18	19

P_5	5	2	1	6
Average	-	-	8.2	12



c) Algorithm:

Step 1: Start the process

Step 2: Accept the number of processes in the ready Queue

Step 3: For each process in the ready Q, assign the process id and accept the CPU burst time

Step 4: Sort the ready queue according to the priority number.

Step 5: Set the waiting of the first process as '0' and its burst time as its turn around time

Step 6: For each process in the Ready Q calculate

(e) Waiting time for process(n)= waiting time of process (n-1) + Burst time of process(n-1)

(f) Turn around time for Process(n)= waiting time of Process(n)+ Burst time for process(n)

Step 7: Calculate

(i) Average waiting time = Total waiting Time / Number of process

(j) Average Turnaround time = Total Turnaround Time / Number of process

Step 8: Stop the process

d) Program:

```
#include<iostream>

using namespace std;

int main()
{
    int bt[20],p[20],wt[20],tat[20],pr[20],i,j,n,total=0,pos,temp,avg_wt,avg_tat;
    cout<<"Enter Total Number of Process:";
    cin>>n;

    cout<<"\nEnter Burst Time and Priority\n";
    for(i=0;i<n;i++)
    {
        cout<<"\nP["<<i+1<<"]\n";
        cout<<"Burst Time:";
        cin>>bt[i];
        cout<<"Priority:";
        cin>>pr[i];
        p[i]=i+1;
    }
    for(i=0;i<n;i++)

    {
        pos=i;
        for(j=i+1;j<n;j++)
        {
            if(pr[j]<pr[pos])
                pos=j;
        }

        temp=pr[i];
        pr[i]=pr[pos];
        pr[pos]=temp;

        temp=bt[i];
        bt[i]=bt[pos];
        bt[pos]=temp;

        temp=p[i];
        p[i]=p[pos];
        p[pos]=temp;
    }

    wt[0]=0;
    for(i=1;i<n;i++)
```

```

{
wt[i]=0;
for(j=0;j<i;j++)
wt[i]+=bt[j];

total+=wt[i];
}

avg_wt=total/n;    //average waiting time
total=0;

cout<<"\nProcess\t Burst Time \tWaiting Time\tTurnaround Time";
for(i=0;i<n;i++)
{
tat[i]=bt[i]+wt[i];    total+=tat[i];
cout<<"\nP["<<p[i]<<"]\t\t "<<bt[i]<<"\t\t "<<wt[i]<<"\t\t"<<tat[i];
}

avg_tat=total/n;
cout<<"\n\nAverage Waiting Time="<<avg_wt;
cout<<"\nAverage Turnaround Time="<<avg_tat;

return 0;

}

```

e) Output:

```
input
P[1]
Burst Time:1 2
Priority:
P[2]
Burst Time:1 4
Priority:
P[3]
Burst Time:4 2
Priority:
P[4]
Burst Time:1 4
Priority:
P[5]
Burst Time:1 4
Priority:
Process      Burst Time      Waiting Time      Turnaround Time
P[1]          1                0                 1
P[3]          4                1                 5
P[2]          1                5                 6
P[4]          1                6                 7
P[5]          1                7                 8

Average Waiting Time=3
Average Turnaround Time=5
```

f) Result:

Average Waiting Time - 3

Average Turnaround Time - 5

Experiment-6

BANKER ALGORITHM

Name of Student: ANIKET RAI		Class: CSIT-1	
Enrollment No: 0827CI201026		Batch: 2020-2024	
Date of Experiment	Date of Submission		Date of Experiment
Remarks by faculty:		Grade:	
Signature of student:		Signature of Faculty:	

OBJECTIVE OF THE EXPERIMENT

To write c program to implement deadlock avoidance & Prevention by using Banker's Algorithm.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Of BANKER'S Algorithm:

The Banker's algorithm is a resource allocation and deadlock avoidance algorithm that tests for safety by simulating the allocation of predetermined maximum possible amounts of all resources, and then makes an "s-state" check to test for possible deadlock conditions for all other pending activities, before deciding whether allocation should be allowed to continue.

- Always keep so many resources that satisfy the needs of at least one client
- Multiple instances.
- Each process must a priori claim maximum use.
- When a process requests a resource it may have to wait.
- When a process gets all its resources it must return them in a finite amount of time.

c) Algorithm:

1. Start the program.
2. Get the values of resources and processes.
3. Get the avail value.
4. After allocation find the need value.
5. Check whether it's possible to allocate.
6. If it is possible then the system is in safe state.
7. Else system is not in safety state.
8. If the new request comes then check that the system is in safety.
9. Or not if we allow the request.
10. Stop the program.

d) Program:

```
#include<iostream>
using namespace std;
class Bankers{
public:
int alloc[50][50];
int maxi[50][50];
int need[50][50];
int avail[50];
int check_safety(int j,int nr)
{
for(int i=0;i<nr;i++)
{
if(need[j][i]>avail[i])
return 0;
}
return 1;
}
int check(bool a[],int n)
{
for(int i=0;i<n;i++)
{
if(a[i]==false)
return 0;
}
return 1;
}
};
int main()
{
Bankers b;
int np=100, nr=100;
cout<<"\nEnter the no of processes : ";
cin>>np;
cout<<"\nEnter the no of resources : ";
cin>>nr;
cout<<"\nEnter the allocation data : \n";
for(int i=0;i<np;i++)
for(int j=0;j<nr;j++)
cin>>b.alloc[i][j];
cout<<"\nEnter the requirement data : \n";
for(int i=0;i<np;i++)
for(int j=0;j<nr;j++)
cin>>b.maxi[i][j];
for(int i=0;i<np;i++)
```

```

for(int j=0;j<nr;j++)
b.need[i][j]=b.maxi[i][j]-b.alloc[i][j];
cout<<"\nEnter the availability matrix : \n";
for(int i=0;i<nr;i++)
cin>>b.avail[i];
int ex_it=nr,flg;
bool completed[np];
while(10)
{
for(int i=0;i<np;i++)
{
if(!completed[i] && b.check_safety(i,nr))
{
for(int j=0;j<nr;j++)
b.avail[j]+=b.alloc[i][j];
}
completed[i]=true;
}
flg=b.check(completed,np);
ex_it--;
if(flg==1 || ex_it==0)
break;
}
cout<<"\nThe final availability matrix \n";
for(int i=0;i<nr;i++)
cout<<b.avail[i]<<" ";
if(flg==1)
cout<<"There is no deadlock";
else
cout<<"Sorry there is a possibility of deadlock";
return 0;
}

```


e) Output:

```
Enter the no of processes : 4
Enter the no of resources : 4
Enter the allocation data :
0 1 1
2 4 5
0 1 2
0 1 8
3 2 4
5 3 5

Enter the requirement data :
7 8 9
5 4 6
4 5 3
7 5 3
9 5 1

Enter the availability matrix :
3 2 5

The final availability matrix
1 3 2 5 There is no deadlock
```

f) Result:

The Sequence Is:
P1->P3>P2->P5

Experiment-7

FIFO PAGE REPLACEMENT

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Date of Experiment
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement page replacement algorithm FIFO.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept Fifo Page Replacement:

- Treats page frames allocated to a process as a circular buffer:
- When the buffer is full, the oldest page is replaced. Hence first-in, first-out: A frequently used page is often the oldest, so it will be repeatedly paged out by FIFO. Simple to implement: requires only a pointer that circles through the page frames of the process.

reference string

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1

7	7	7	2	2	2	4	4	4	0	0	0	7	7	7
	0	0	0	3	3	3	2	2	2	1	1	1	0	0
		1	1	1	0	0	0	3	3	3	2	2	2	1

page frames

- FIFO Replacement manifests Belady's Anomaly:

more frames \Rightarrow more page faults

- Reference string: 1, 2, 3, 4, 1, 2, 5, 1, 2, 3, 4, 5()

3 Frames:-9 page fault

4 Frames: - 10 page fault

c) Algorithm:

Step 1: Create a queue to hold all pages in memory

Step 2: When the page is required replace the page at the head of the queue

Step 3: Now the new page is inserted at the tail of the queue

d) Program:

```
#include <iostream>
#include<bits/stdc++.h>
using namespace std;

int pageFault(int page[],int n,int no_frame)
{
    int pagefault=0;

    vector<int> v1;
    int i;
    for(i=0;i<n;i++)
    {
        auto it=find(v1.begin(),v1.end(),page[i]);
        if(it==v1.end())
        {
            if(v1.size()==no_frame)
            {
                v1.erase(v1.begin());
            }
            v1.push_back(page[i]);
            pagefault++;
        }
    }
    return pagefault;
}

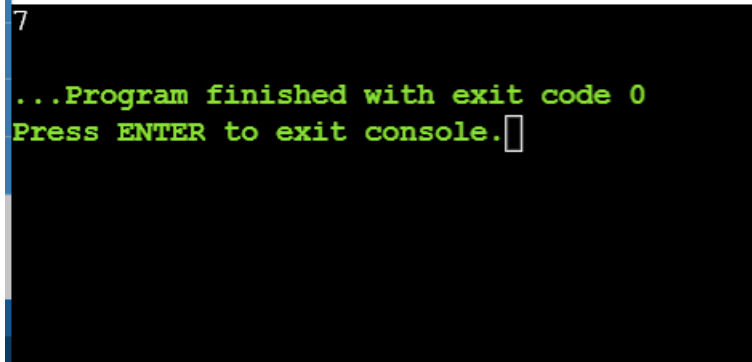
int main()
{
    int page[]={7,0,1,2,0,3,0,4,2,3,0,3,2};

    int n=13;

    int no_frame=4;
    cout<<pageFault(page,n,no_frame);

    return 0;
}
```

e) Output:

A screenshot of a terminal window with a black background and green text. The text displays the message "...Program finished with exit code 0" followed by "Press ENTER to exit console." and a cursor. Above the terminal window, the number "7" is visible, likely representing the number of page faults.

```
7  
...Program finished with exit code 0  
Press ENTER to exit console.
```

f) Result:

No. of page faults - 7

Experiment-8

LRU PAGE REPLACEMENT

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Date of Experiment
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement page replacement algorithm LRU.

FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept of LRU Algorithm:

Pages that have been heavily used in the last few instructions will probably be heavily used again in the next few. Conversely, pages that have not been used for ages will probably remain unused for a long time. when a page fault occurs, throw out the page that has been unused for the longest time. This strategy is called LRU (Least Recently Used) paging.

Page reference stream:

1 2 3 2 1 5 2 1 6 2 5 6 3 1 3 6 1 2 4 3

1	1	1	1	3	2	1	5	2	1	6	2	5	6	6	1	3	6	1	2
	2	2	3	2	1	5	2	1	6	2	5	6	3	1	3	6	1	2	4
		3	2	1	5	2	1	6	2	5	6	3	1	3	6	1	2	4	3
*	*	*			*			*		*		*	*			*	*	*	

LRU

Total 11 page faults

c) Algorithm:

Step 1: Create a queue to hold all pages in memory

Step 2: When the page is required replace the page at the head of the queue

Step 3: Now the new page is inserted at the tail of the queue

Step 4: Create a stack

Step 5: When the page fault occurs replace page present at the bottom of the stack

d) Program:

```
#include <iostream>
#include<bits/stdc++.h>

using namespace std;

int pageFault(int pages[],int n,int mem_capacity)
{
    int pagefault=0;
    vector<int> v1;

    for(int i=0;i<=n;i++)
    {
        auto it=find(v1.begin(),v1.end(),pages[i]);
        if(it==v1.end())
        {
            if(v1.size()==mem_capacity)
            {
                v1.erase(v1.begin());
            }
            v1.push_back(pages[i]);
            pagefault++;
        }
        Else
        {
            v1.erase(it);
            v1.push_back(pages[i]);
        }
    }
    return pagefault;
}

int main()
{
    int pages[]={ 7,0,1,2,0,3,0,4,2,3,0,3,2 };
    int n=sizeof(pages)/sizeof(pages[0]); //no of pages
    int mem_capacity;
    cout<<"Enter no. of frames";
    cin>>mem_capacity;

    cout<<pageFault(pages,n,mem_capacity);

    return 0;
```

}

d) OUTPUT:

```
Enter no. of frames:    4
7
...Program finished with exit code 0
Press ENTER to exit console.
```

e) Result:

No. of pages faults - 7

Experiment-9

FCFS Disk Scheduling Algorithm

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Submitted on:
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement FCFS Disk Scheduling Algorithm

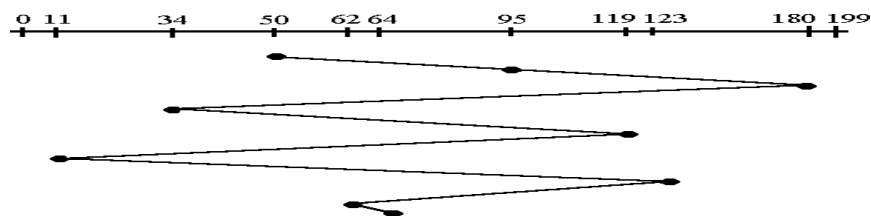
FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept of FCFS Disk Scheduling Algorithm:

All incoming requests are placed at the end of the queue. Whatever number that is next in the queue will be the next number served. Using this algorithm doesn't provide the best results. To determine the number of head movements you would simply find the number of tracks it took to move from one request to the next. For this case it went from 50 to 95 to 180 and so on. From 50 to 95 it moved 45 tracks. If you tally up the total number of tracks you will find how many tracks it had to go through before finishing the entire request. In this example, it had a total head movement of 640 tracks. The disadvantage of this algorithm is noted by the oscillation from track 50 to track 180 and then back to track 11 to 123 then to 64. As you will soon see, this is the worse algorithm that one can use.



c) Algorithm:

Step 1: Create a queue to hold all requests in disk

Step 2: Move the head to the request in FIFO order (Serve the request first that came first)

Step 3: Calculate the total head movement required to serve all request.

d) Program:

```
#include<bits/stdc++.h>
using namespace std;

class DiskSched{
public:
float fcfs(int arr[],int n,int init)
{
float answer=abs(init-arr[0]);
for(int i=1;i<n;i++)
{
answer=answer+abs(arr[i]-arr[i-1]);
}
return answer/n;
}
};

int main()
{
int disk[]={55,58,39,18,90,160,100,38,184};
DiskSched test;
cout<<test.fcfs(disk,9,100);
}
```

e) Output:

```
55.3333
...Program finished with exit code 0
Press ENTER to exit console.█
```

f) Result:

Total Head Movement Required Serving All Requests

Experiment-10

SSTF Disk Scheduling Algorithm

Name of Student: ANIKET RAI		Class: CSIT-1
Enrollment No: 0827CI201026		Batch: 2020-2024
Date of Experiment	Date of Submission	Date of Experiment
Remarks by faculty:		Grade:
Signature of student:		Signature of Faculty:

OBJECTIVE OF THE EXPERIMENT

To implement SSTF Disk Scheduling Algorithm

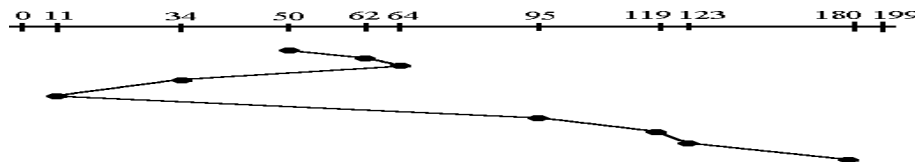
FACILITIES REQUIRED

a) Facilities Required To Do The Experiment:

S.NO	FACILITIES REQUIRED	QUANTITY
1	System	1
2	Windows XP/7	

b) Concept of SSTF Disk Scheduling Algorithm:

In this case request is serviced according to next shortest distance. Starting at 50, the next shortest distance would be 62 instead of 34 since it is only 12 tracks away from 62 and 16 tracks away from 34. The process would continue until all the process are taken care of. For example the next case would be to move from 62 to 64 instead of 34 since there are only 2 tracks between them and not 18 if it were to go the other way. Although this seems to be a better service being that it moved a total of 236 tracks, this is not an optimal one. There is a great chance that starvation would take place. The reason for this is if there were a lot of requests close to each other the other requests will never be handled since the distance will always be greater.



c) Algorithm:

Step 1: Create a queue to hold all requests in disk

Step 2: Calculate the shortest seek time every time before moving head from current head position

Step 3: Calculate the total head movement required to serve all request.

d) Program:

```
#include <bits/stdc++.h>
using namespace std;
// vector<int>:: iterator it;

int minDiff(int *req,int pos,int n)
{
    int newpos;
    int mini=INT_MAX;
    int diff;
    for(int i=0;i<n;i++)
    {
        if(req[i]!=-1)
        {
            diff=abs(pos-req[i]);
            if(mini>diff)
            {
                mini=diff;
                newpos=i;
            }
        }
    }
    // cout<<"request choosen : "<<req[newpos];
    return newpos;
}

float SSTF(int *req,int pos,int n)
{
    /* vector<int> :: iterator it;
    while(!req.empty())
    {
        it=req.find(req.begin(),req.end());
        int temp=req[it];
        if(abs(req[it]-req[it-1])<(req[it+1]-req[it]))
        {

            total+=abs(req[it]-req[it-1]);
            req.erase(req.begin()+it);

        }
        else
        {
```

```

total+=abs(req[it+1]-req[it]);
req.erase(req.begin()+it);
}
}*/
int posi=pos;
float total=0;

for(int i=0;i<n;i++)
{
int index=minDiff(req,posi,n);
// cout<<"diff: "<<abs(pos-req[index]);
total+=abs(posi-req[index]);
posi=req[index];
req[index]=-1;

}

float avg=total/n;
return avg;

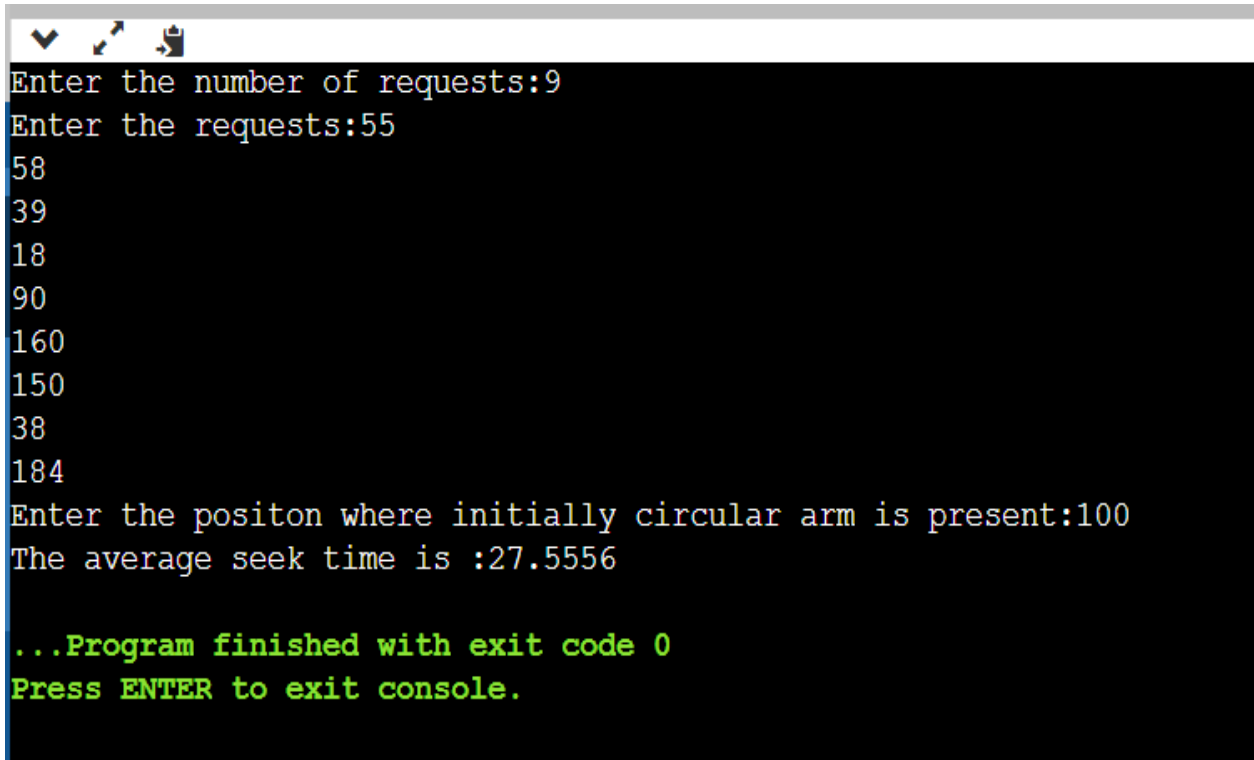
}

int main()
{
int n,positom;
cout<<"Enter the number of requests:";
cin>>n;

int req[n];
cout<<"Enter the requests:";
for(int i=0;i<n;i++)
{
// int r;
cin>>req[i];
// req.push_back(r);
}
cout<<"Enter the positon where initially circular arm is present:";
cin>>positom;
cout<<SSTF(req,positom,n);
return 0;

```

e) Output:

A screenshot of a console window with a black background and white text. The window has a title bar with standard Windows icons. The text in the console shows a program that takes user input for the number of requests and requests themselves, then outputs a list of values and the average seek time.

```
Enter the number of requests:9
Enter the requests:55
58
39
18
90
160
150
38
184
Enter the positon where initially circular arm is present:100
The average seek time is :27.5556

...Program finished with exit code 0
Press ENTER to exit console.
```

f) Result:

Total Head Movement Required Serving All Requests

FAQ's

1. What are different types of schedulers?
2. Explain types of Operating System?
3. Explain performance criteria for the selection of schedulers?
4. Explain priority based preemptive scheduling algorithm?
5. What is thread?
6. Explain different types of thread?
7. What is kernel level thread?
8. What is user level thread?
9. What is memory management?
10. Explain Belady's Anomaly.
11. What is a binary semaphore? What is its use?
12. What is thrashing?
13. List the Coffman's conditions that lead to a deadlock.
14. What are turnaround time and response time?
15. What is the Translation Lookaside Buffer (TLB)?
16. When is a system in safe state?
17. What is busy waiting?
18. Explain the popular multiprocessor thread-scheduling strategies.
19. What are local and global page replacements?
20. In the context of memory management, what are placement and replacement algorithms?
21. In loading programs into memory, what is the difference between load-time dynamic linking and run-time dynamic linking?
22. What are demand- and pre-paging?
23. Paging a memory management function, while multiprogramming a processor management functions, are the two interdependent?
24. What has triggered the need for multitasking in PCs?
25. What is SMP?
26. List out some reasons for process termination.
27. What are the reasons for process suspension?
28. What is process migration?
29. What is an idle thread?
30. What are the different operating systems?
31. What are the basic functions of an operating system?