

Student Information

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Answer 1

a)

$G_1 = (V, \Sigma, R, S)$ where

$V = \{S, a, b\}$,

$\Sigma = \{a, b\}$,

S is the start symbol ,

$R = \{S \rightarrow SaSbSbS | SbSaSbS | SbSbSaS | \epsilon\}$

such that $L(G_1) = L_1$.

b)

$G_2 = (V, \Sigma, R, S)$ where

$V = \{S, a, b\}$,

$\Sigma = \{a, b\}$,

S is the start symbol ,

$R = \{S \rightarrow aSb | aaSb | \epsilon\}$

such that $L(G_2) = L_2$.

c)

Let's introduce a Nondeterministic Pushdown Automaton M such that $L(M) = L_1$

$M = (K, \Sigma, V, \Delta, s, F)$ where ,

$K = \{p, q\}$

$\Sigma = \{a, b\}$

$V = \{a, b, S\}$

$s = p$

$F = \{q\}$

$\Delta = \{$
 $((p, \epsilon, \epsilon), (q, S)),$
 $((q, \epsilon, S), (q, SaSbSbS)),$
 $((q, \epsilon, S), (q, SbSaSbS)),$
 $((q, \epsilon, S), (q, SbSbSaS)),$
 $((q, \epsilon, S), (q, \epsilon)),$
 $((q, a, a), (q, \epsilon)),$
 $((q, b, b), (q, \epsilon))$
 $\}$

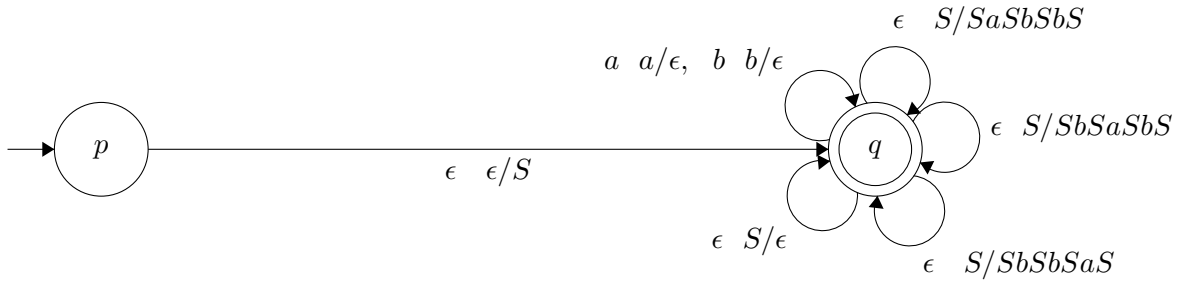


Figure: State Diagram of M

Be careful that $(a \quad a/\epsilon)$ and $(b \quad b/\epsilon)$ has been shown on the same loop for simplicity !

d)

$G_3 = (V, \Sigma, R, S)$ where

$V = \{S, S_1, S_2, a, b\}$,

$\Sigma = \{a, b\}$,

S is the start symbol ,

$R = \{$
 $S \rightarrow S_1 | S_2,$
 $S_1 \rightarrow S_1 a S_1 b S_1 b S_1 | S_1 b S_1 a S_1 b S_1 | S_1 b S_1 b S_1 a S_1 | \epsilon,$
 $S_2 \rightarrow a S_2 b | a a S_2 b | \epsilon$
 $\}$

such that $L(G_3) = L_3 = L_1 \cup L_2$.

Answer 2

a)

Consider the string $w = 00111 \in L(G_1)$:

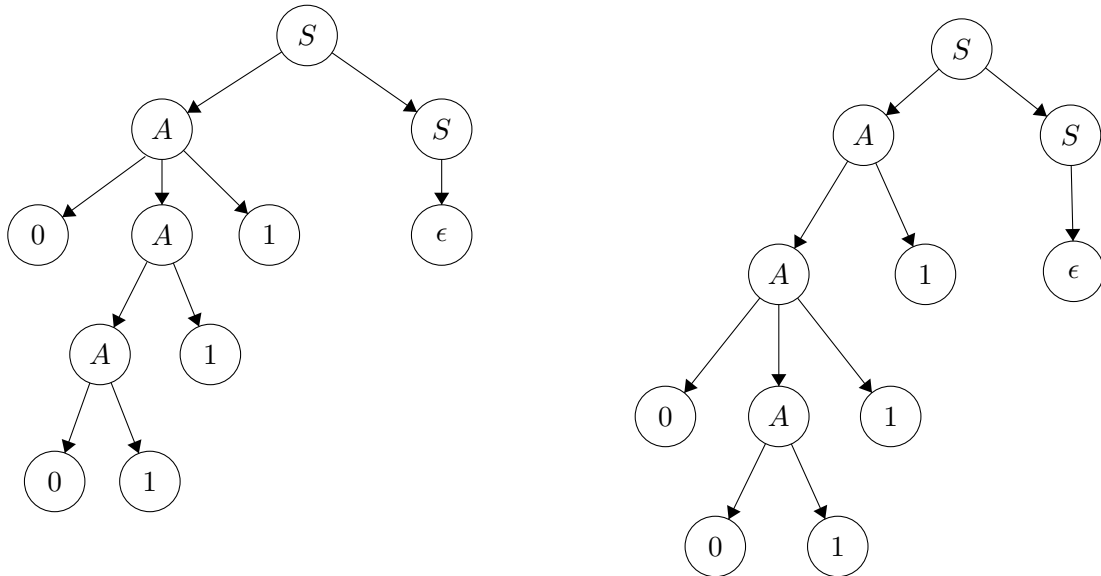


Figure 1 - 2: Two different parse trees for W

Since G_1 has two distinct parse trees, which are drawn above, for the string $w = 00111$. The grammar G_1 is ambiguous.

b)

Let's introduce an unambiguous grammar G_2 such that $L(G_1) = L(G_2)$:

$G_2 = (V, \Sigma, R, S)$ where

$V = \{S, A, B, 0, 1\}$,

$\Sigma = \{0, 1\}$,

S is the start symbol,

$R = \{$
 $S \rightarrow AS | \epsilon,$
 $A \rightarrow A1 | B,$
 $B \rightarrow 0B1 | 01$
 $\}$

c)

Left-most derivation :

$$S \Rightarrow AS \Rightarrow A1S \Rightarrow B1S \Rightarrow 0B11S \Rightarrow 00111S \Rightarrow 00111$$

Corrsponding parse tree:

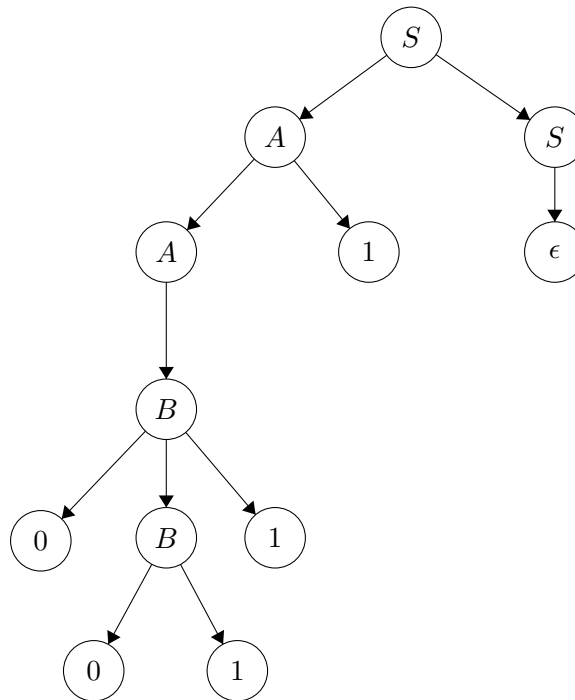


Figure: Corresponding parse tree