



Vehicle Networks

Controller Area Network (CAN)

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Outline

- Introduction
 - Design requirements
 - History and application fields
 - Recapitulation: Communications basics
- CAN Physical Layer
- CAN Data Link Layer
- CAN Error Handling & Error Confinement

CAN





Introduction

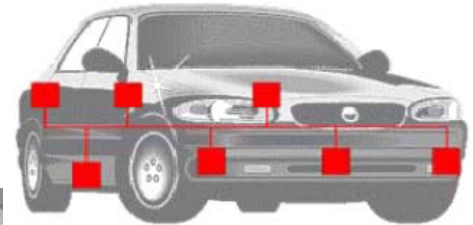
Requirements for automotive communications

- The development of CAN was mainly driven by Mercedes-Benz for networking of versatile *Electronic Control Units (ECUs)* with the following requirements:
 - **Error-resistance** to cope with strong electro-magnet interference
 - **Prioritized real-time** capabilities with short latency (e.g. for safety critical applications)
 - **Fast** data rate (Class C network: 125 kbit/s – 1 Mbit/s)
 - **Expandability** for versatile nodes
 - **Cost-effectiveness** for wires and nodes
- These requirements also hold for various other application fields in aviation and maritime industry, industrial and home automation, consumer electronics
 - ➔ Widespread distribution of CAN nowadays

Introduction

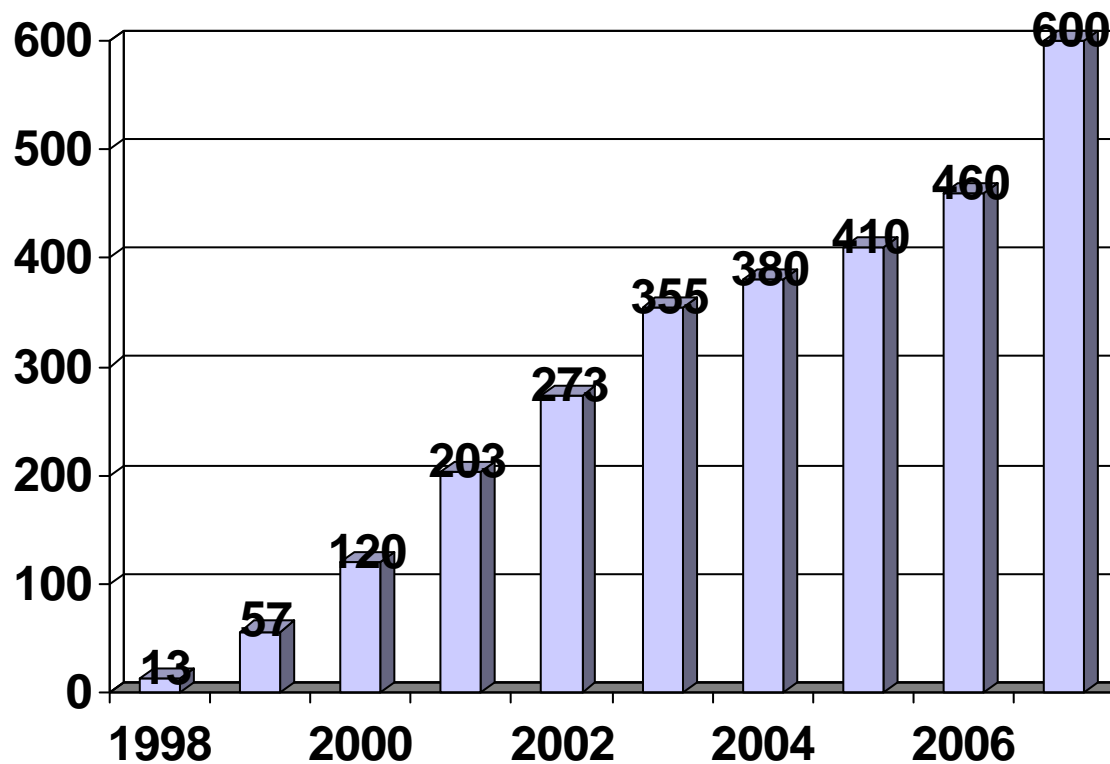
Application Fields

- Automotive, aviation, space, maritime industry
 - Cars, trucks, busses
 - Airplanes
 - Rockets, space shuttles
 - Ships
- Medical equipment
 - X-Ray, Electro-Cardiograms (ECG)
- Industrial and home automation
 - Production machines
 - Lifts and escalators
 - Shutter, heating, light control
- Household appliances
 - Washers, Dryers
 - Coffee machines ("Café au CANopen")
- Consumer electronics
 - Model railway



Introduction

Number of CAN nodes in millions



Produced by
Motorola,
Philips, Intel,
Infineon, etc.

Introduction

History

- 1985 Start of development of CAN at Robert Bosch GmbH
- 1986 V1.0 specification of CAN
- 1991 Specifications of the extended CAN2.0 protocol
 - Part 2.0A – 11-bit identifier
 - Part 2.0B – 29-bit identifier (extended frame format)
- 1992 CAN in Automation (CiA) established as the international users and manufacturers group
- 1993 First car, a Mercedes S-class, equipped with CAN
- 1994 First standardization at ISO is completed
- 1998 Development phase of time-triggered CAN (TTCAN) networks
- 1999 Explosion of CAN-linked equipment in all motor vehicle and industrial applications

Introduction

CAN in ISO/OSI Reference Model

No. of layer	ISO/OSI ref model	CAN protocol specification
7	<i>Application</i>	Application specific
6	<i>Presentation</i>	Optional: Higher Layer Protocols (HLP)
5	<i>Session</i>	
4	<i>Transport</i>	
3	<i>Network</i>	
2	<i>Data Link</i>	CAN protocol (with free choice of medium)
1	<i>Physical</i>	

Introduction

Standards

SAE J2284

High Speed CAN (HSC) For
Passenger Vehicle
Applications

SAE J1939

Recommended Practise for
Control and Communications
Network (Class C) on Truck
and Bus Applications

ISO 11898 Road vehicles – Controller Area Network (CAN)

Extensions to Data Link Layer

ISO 11898-4 Time-triggered CAN

ISO 11898-5 relates to high-speed CAN and low-power applications

Data Link Layer

ISO 11898-1 Data link layer and physical signalling

Physical Layer

ISO 11898-2 High-speed medium access unit

ISO 11898-3 Low-speed fault-tolerant medium-dependent interface

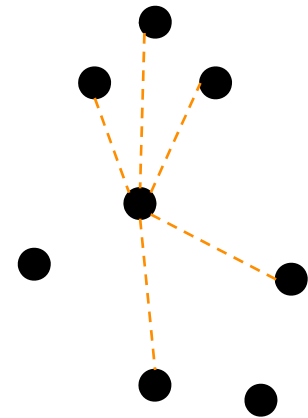
According to
original Bosch
specification
CAN2.0A and
CAN2.0B



Introduction

PHY: Overview

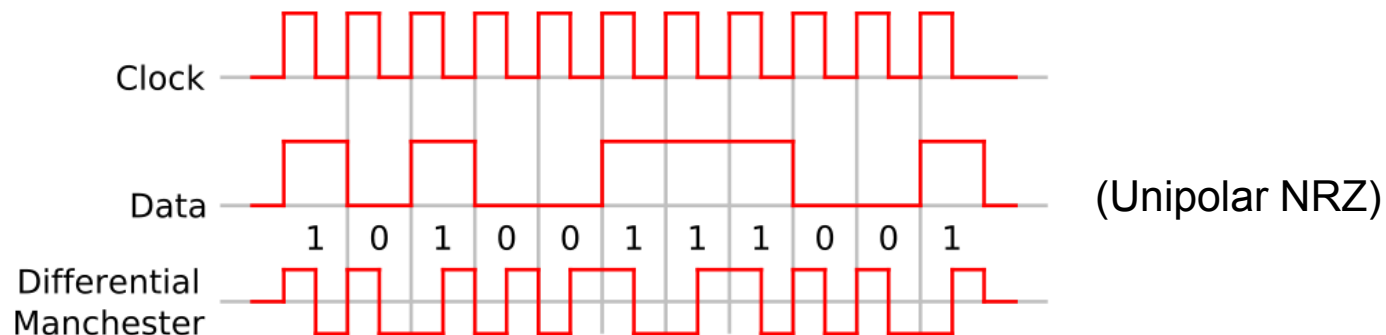
- Responsible for *bit-by-bit* translation of data into/from hardware-specific operations
- Defines electrical, mechanical, and procedural interface to the transmission medium (wire, air, optical fibre etc.)
- Main functions:
 - Line coding: (De-)modulation of the digital signal into/from the amplitude- and time-discrete signal
 - Channel coding: Insertion of redundant information for error detection/correction
 - Synchronisation: Determination of a common time basis



Introduction

PHY: Line Coding

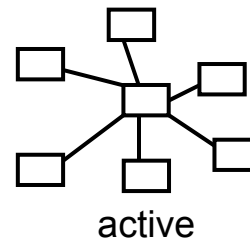
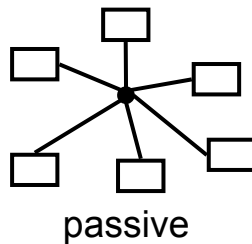
- Return-to-zero (RZ) coding (e.g. IrDA Serial Infrared)
- Non-Return-to-Zero (NRZ) coding (e.g. RS-232)
- Manchester coding (e.g. Ethernet)
- 4B3T (e.g. ISDN)



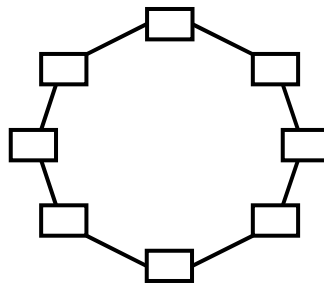
Introduction

PHY: Network topologies

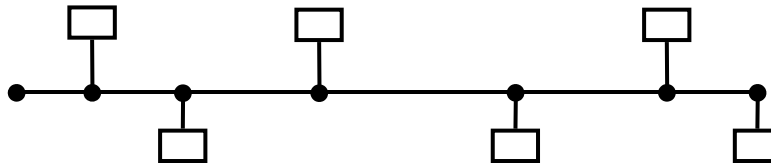
➤ Star topology



➤ Ring topology



➤ Bus topology



Introduction

PHY: Wiring

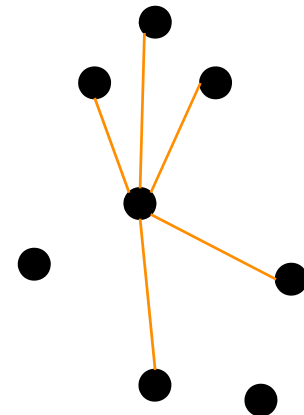
- One-wire: single feed line – return line over ground connection
- Two-wire:
 - Unshielded Twisted pair (UTP)
 - Shielded Twisted pair (STP)
- Coaxial:
Inner conductor shielded with fine woven wires or metallic foil
- Optical: Transmission of light through fiber
 - Single-mode fiber (SMF): only single *ray of light* can be carried
 - Multi-mode fiber (MMF): support of more than one propagation mode; reflection on fiber border



Introduction

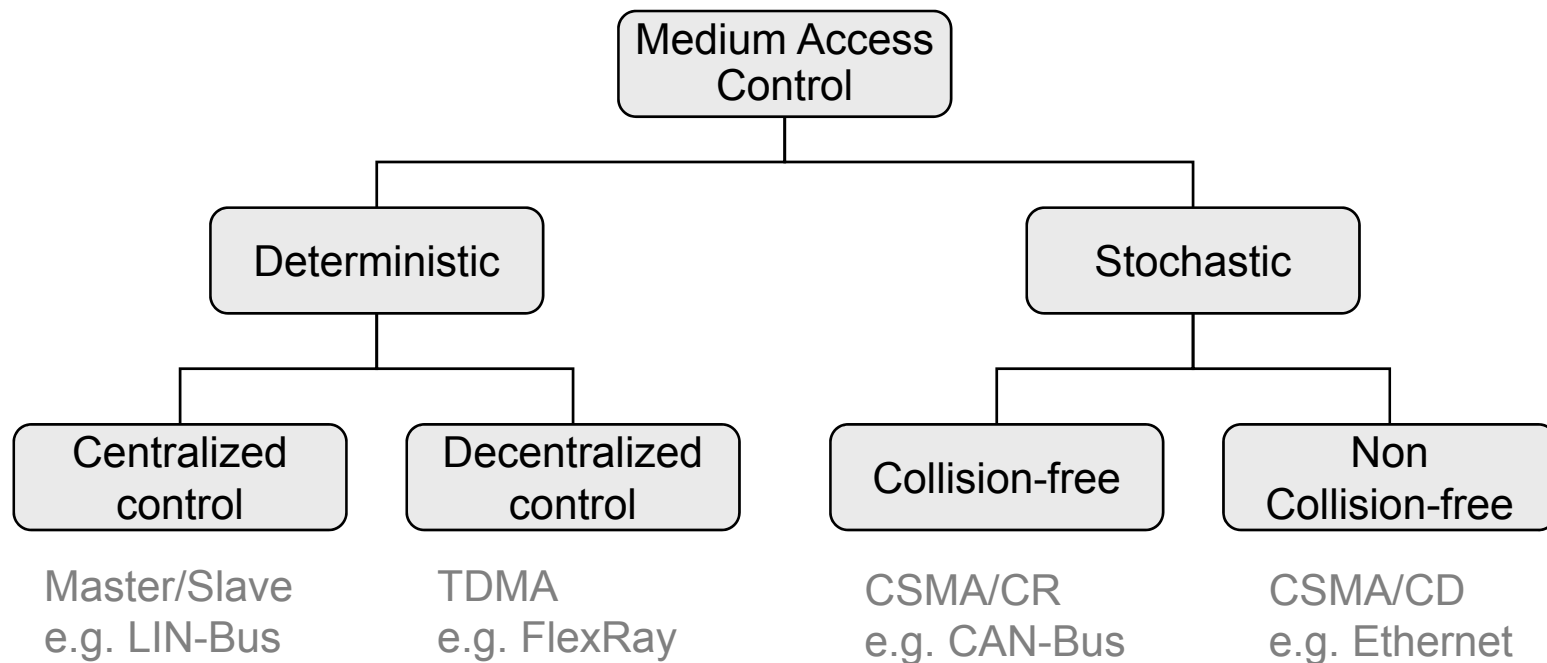
DL: Overview

- Responsible for *Point-to-Point / Point-to-Multipoint* transfer of data
- Main functions:
 - Framing (i.e. sequence of bits)
 - Addressing
 - Error protection/detection/correction
 - Flow control
 - Medium access
- Sublayers:
 - Medium Access Control (MAC)
e.g. CSMA/CD, CSMA/CA, ...
 - Logical link control (LLC)
e.g. IEEE 802.2

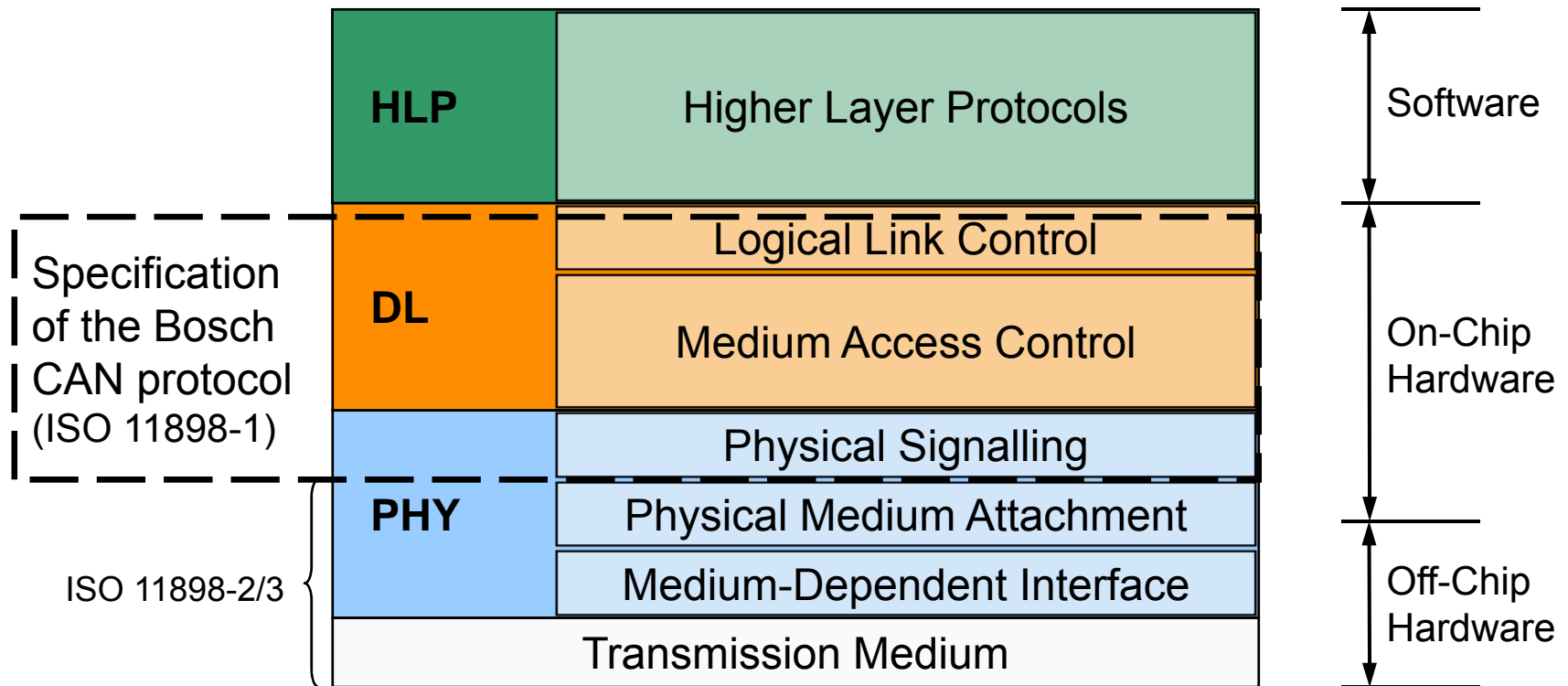



Introduction

DL: Medium Access Control



CAN Protocol Specification






Controller Area Network (CAN)

Physical Layer





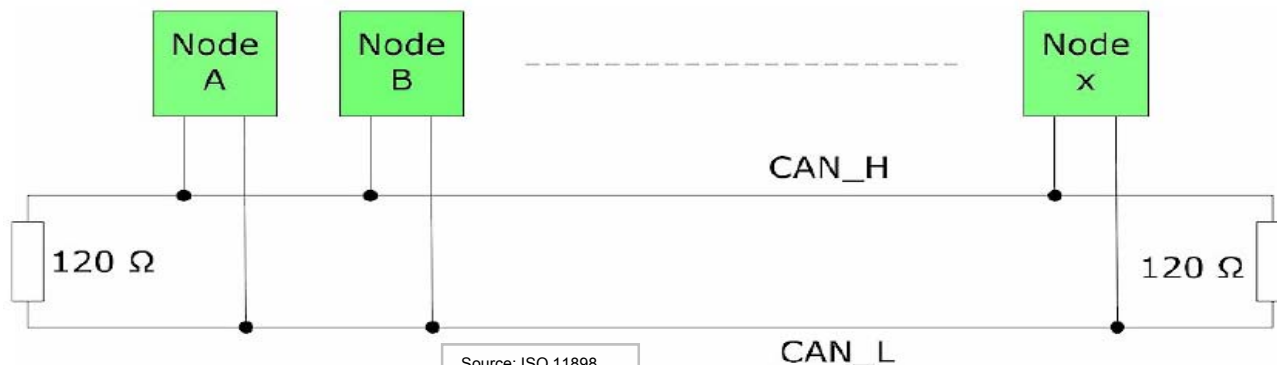
CAN Physical Layer

Overview

- Physical Line Signalling (PLS)
 - Bit Coding, Bit Timing
 - Synchronization
 - Error detection
- Physical Medium Attachment (PMA)
 - Transmitter/Receiver Specification, i.e.
Characteristics of the command stages and timing stages
- Medium-Dependent Interface (MDI)
 - Bus Cable, Bus Termination Networks, Connectors, etc.

CAN Physical Layer Hardware

- n nodes would require $O(n^2)$ wires for full meshed networking (up to 100 ECUs in modern cars)
- Bus topology only requires a single bus with n branches: $O(n)$
- Pure CAN does not specify transmission medium
- Automotive CAN according to ISO 11898-2/3 uses twisted pair with differential voltages on a bus topology (tolerant to single wire disturbance)
- Bus must be terminated with $120\ \Omega$ to:
 - remove signal reflections at the end of the bus
 - ensure the bus gets correct DC levels
- Max. 30 connected nodes



Source: ISO 11898

CAN Physical Layer

Data rates vs. bus length

- The signal has to propagate to the most remote node and back again (round trip) before the bit is sampled
 - ➔ Bus length and data rate are correlated
- Signal propagation of information is limited to approximately 0.4-0.7 times the speed of light ($\approx 0.1\text{-}0.2\text{ m/ns}$); depends on cable and node impedance
- A nominal bit time can be calculated at each node
- Further factors: transceiver delay, sample time tolerance

$$\text{max. bus length} < \frac{\text{signal velocity} * \text{nom. bittime}}{2}$$

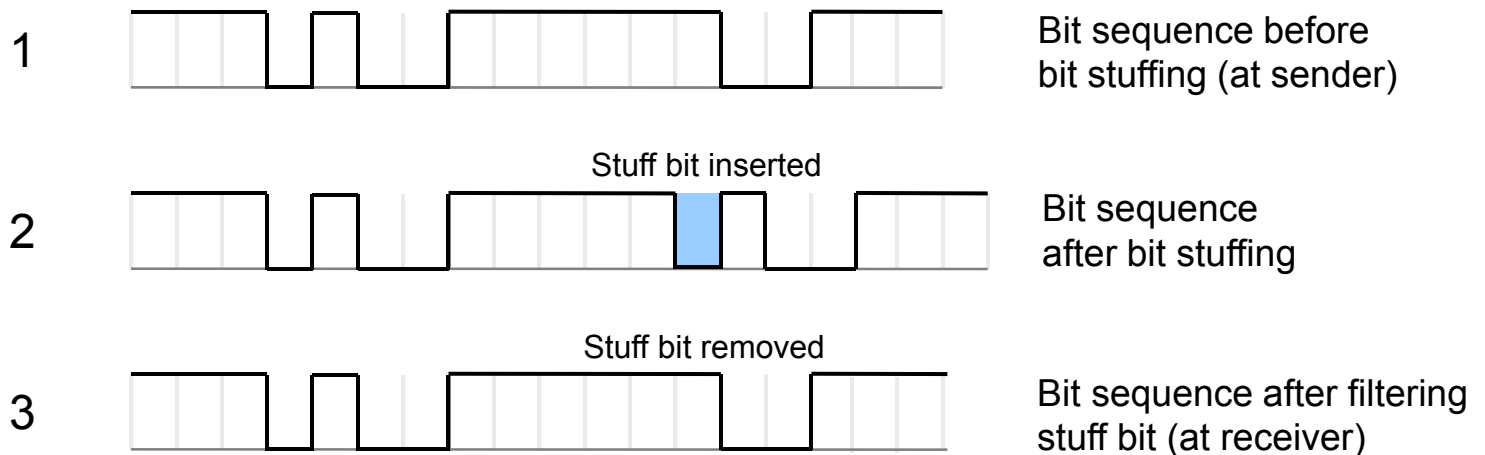
Data rate	Max. bus length*	Nominal Bit-Time
1000 kbit/s	40 m	1 μ s
500 kbit/s	130 m	2 μ s
250 kbit/s	270 m	4 μ s
125 kbit/s	530 m	8 μ s
50 kbit/s	1300 m	20 μ s
20 kbit/s	3300 m	50 μ s
10 kbit/s	6700 m	100 μ s

*approximation

CAN Physical Layer

Line Coding

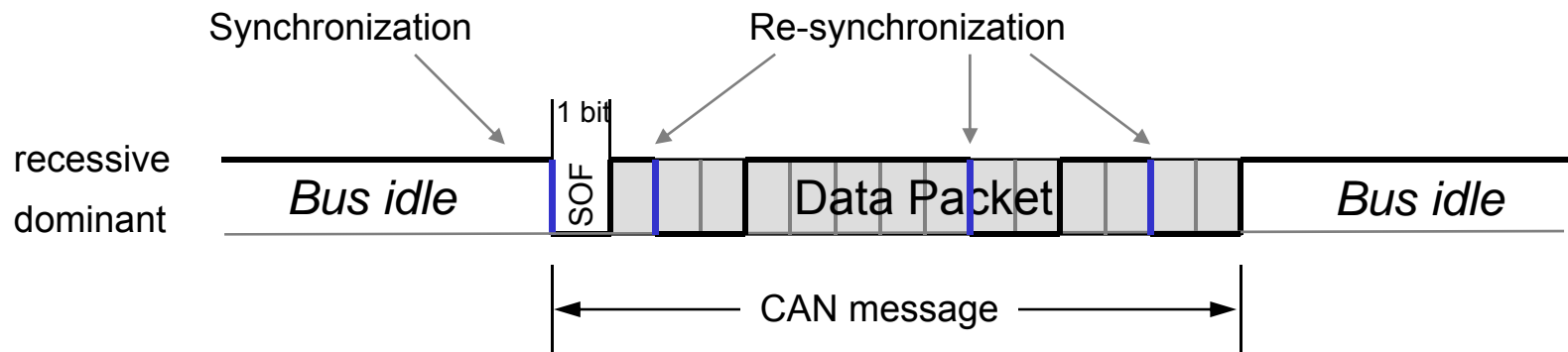
- Line coding: Non-Return-to-Zero (NRZ) with bit stuffing
 - Standard NRZ has not enough signal edges for synchronisation
 - ➔ Bit stuffing: after 5 consecutive bits with same polarity (dominant or recessive), a bit with complementary polarity will be inserted
 - Receiver filters the complementary bit ➔ not visible to higher layers



CAN Physical Layer

Synchronization

- No global time source, no dedicated clock signal
- Synchronization by edge detection in data signal
- Bit length known due to uniform clock rate for every node (e.g. 2 μ s for 500 kbit/s)
- Hard synchronization with first recessive-to-dominant edge (=dominant *Start Of Frame (SOF)* bit) after bus idle
- Continuous re-synchronization at every recessive-to-dominant edge transition



CAN Physical Layer

Error detection

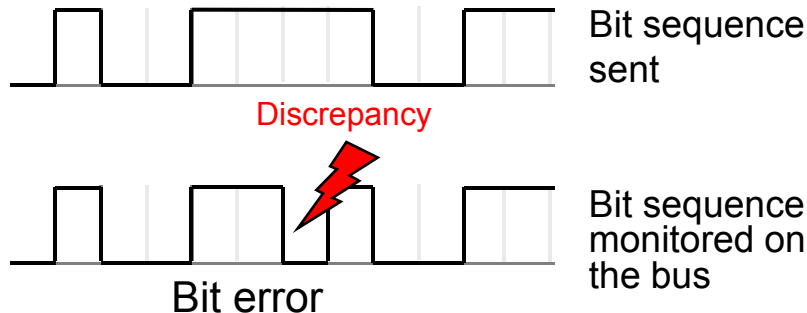
➤ At sender side:

- **Bit errors:** Senders monitor the bus during transmission. Discrepancies between sent and monitored bit are interpreted as transmission error

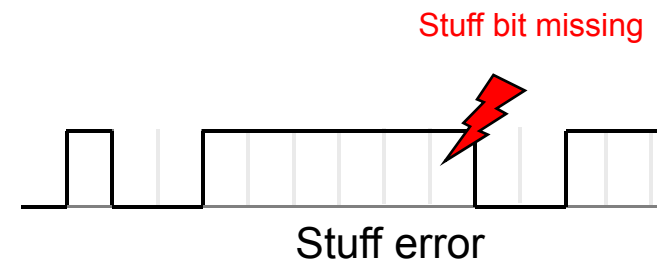
➤ At receiver side:


- **Stuff error:** Receiving nodes can detect errors by checking whether the bit sequence is in adherence with the bit stuffing rule

Sender side



Receiver side





Controller Area Network (CAN)

Data Link Layer





CAN Data Link Layer

Overview


- Sublayer: Medium Access Control (MAC)
 - Message Framing
 - Arbitration
 - Acknowledgement
 - Error detection, error handling

- Sublayer: Logical Link Control (LLC)
 - Flow control
 - Error recovery

CAN Data Link Layer

Characteristics

- Short messages assure **short latencies**
 - CAN2.0A (standard format): 11 bit identifier → 47-55 bit data frames (+stuff bits)
 - CAN2.0B (extended format): 29 bit identifier → 67-75 bit data frames (+ stuff bits)
- CAN identifiers do not specify the destination of the message but the content of the message
 - ➔ Every node on the bus receives all messages and filters as required
 - ➔ **Expandability**
- Contention-free medium access assures **real-time capabilities** (in contrast to Ethernet, WLAN, etc.)
- **Prioritisation** with message identifiers assured by non-destructive arbitration
- Sophisticated error detection, error handling and error confinement assures short- and long-term **error-resistance**



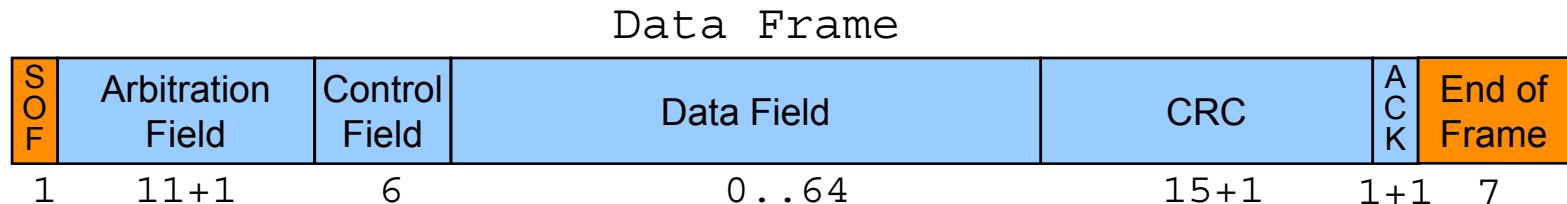
CAN Data Link Layer

Frame Types

- **Data frame:** Transport of data
 - Periodic or event-driven broadcast of data (e.g. temperature, yaw rate)
- **Remote frame:** Request of data
 - Same as Data frame with dominant RTR bit (empty payload)
 - Request for data according to specified identifier
- **Error frame:** Sent when a error is detected
- **Overload frame:** Request of supplementary time interval between preceding and following data frames
- **Interframes:** Separation of data/remote frames from preceding frames

CAN Data Link Layer

CAN2.0A data frame

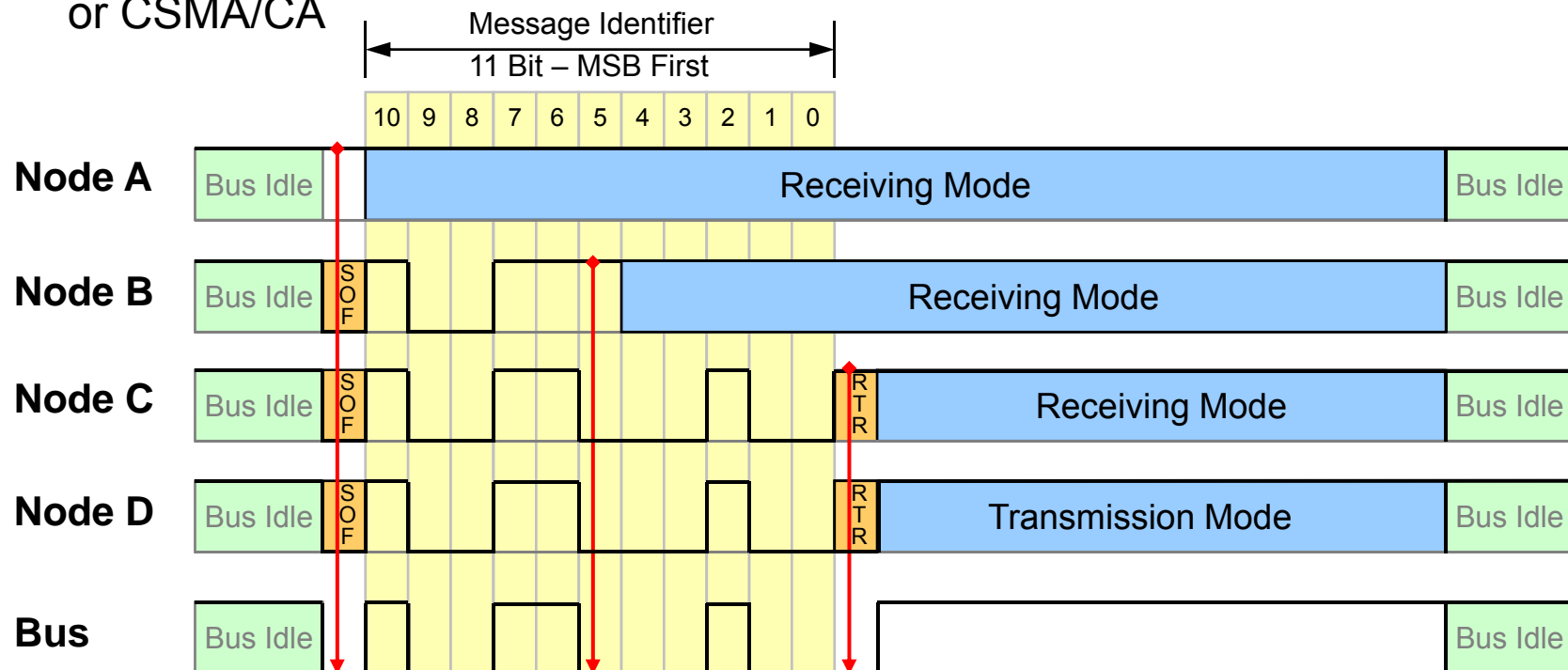


- SOF (Start of Frame): single dominant bit
- Arbitration Field:
 - 11 bit identifier
 - 1 bit RTR (Remote Transmission Request):
dominant for data frames, recessive for remote frames
- CRC (Cyclic Redundancy Check):
 - 15 bit CRC with generator polynomial $x^{15} + x^{14} + x^{10} + x^8 + x^7 + x^4 + x^3 + 1$
 - 1 bit CRC delimiter: single (always) recessive bit
- ACK (Acknowledgement)
 - 1 bit ACK slot: dominant overwriting
 - 1 bit ACK delimiter: single (always) recessive bit
- End of Frame: 7 recessive bits

CAN Data Link Layer

Arbitration

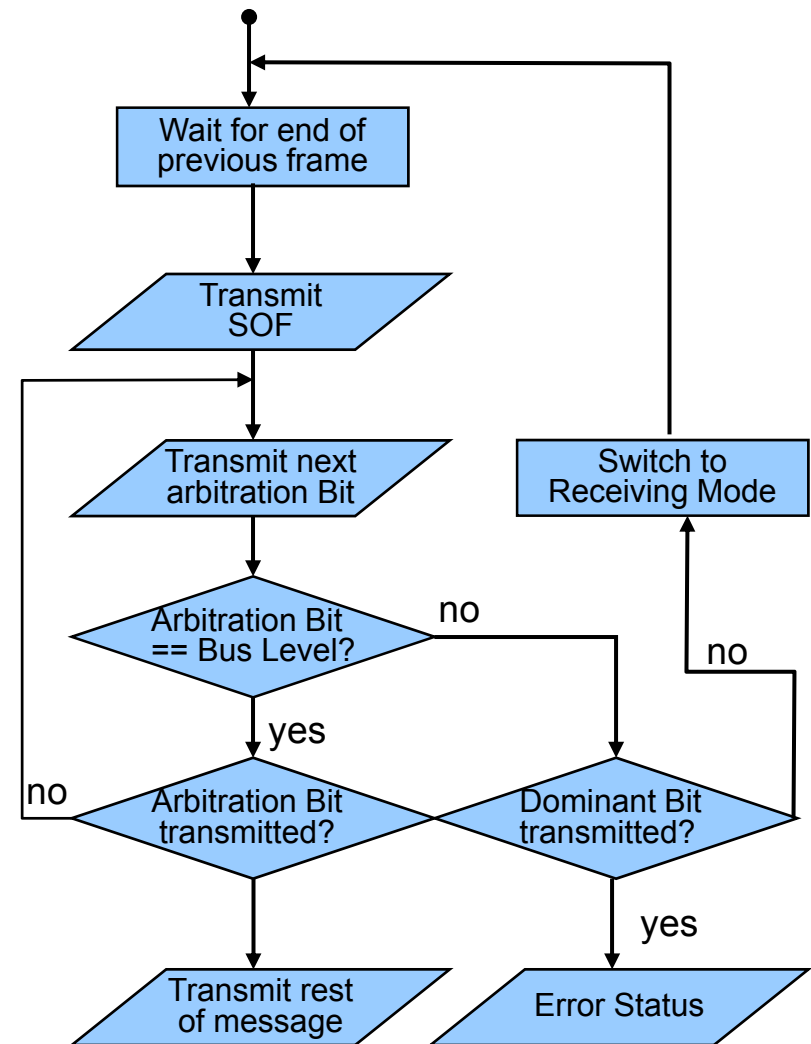
- Bit-wise (bits contained in the identifier)
- Non-destructive CSMA/CR: sender with lowest identifier (=highest priority) wins the arbitration, message is not destroyed as in CSMA/CD or CSMA/CA



CAN Data Link Layer

CSMA/CR

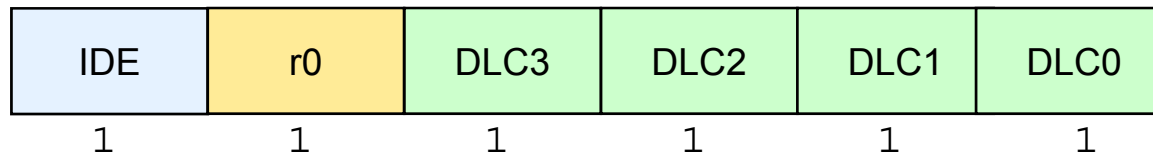
- Carrier Sense Multiple Access / Collision Resolution (CSMA/CR)
- Contention-free
- If collision happens, arbitration is used to resolve it
- High bandwidth utilisation (up to 100%)
- Allows prioritisation of messages



CAN Data Link Layer

CAN2.0A data frame (Control field)

Control Field



- IDE (IDentifier Extension): Distinction between standard format (CAN2.0A) and extended format (CAN2.0B)
- r0 (reserved)
- DLC3-0 (Data Length Code): size of the data field (0..8 bytes = 9 possibilities requires 4 bit field)

Number of Data Bytes	Data Length Code			
	DLC3	DLC2	DLC1	DLC0
0	dominant	dominant	dominant	dominant
1	dominant	dominant	dominant	recessive
2	dominant	dominant	recessive	dominant
3	dominant	dominant	recessive	recessive
4	dominant	recessive	dominant	dominant
5	dominant	recessive	dominant	recessive
6	dominant	recessive	recessive	dominant
7	dominant	recessive	recessive	recessive
8	recessive	dominant	dominant	dominant

CAN Data Link Layer

Cyclic Redundancy Check

Example message to transmit

Filler

G(x)

$$\begin{array}{r}
 \underbrace{0011000000101011011}_{\text{Example message to transmit}} \underbrace{0000000000000000}_{\text{Filler}} : \underbrace{1100010110011001}_{G(x)} = \text{something (not used)} \\
 \hline
 1100010110011001 \\
 000001010011010010000 \\
 \hline
 1100010110011001 \\
 01100011000010010 \\
 \hline
 1100010110011001 \\
 0000001110001011000000 \\
 \hline
 1100010110011001 \\
 001001110101100100 \\
 \hline
 1100010110011001 \\
 01011000111111010 \\
 \hline
 1100010110011001 \\
 01110100011000110 \\
 \hline
 1100010110011001 \\
 0010110101011111 \\
 \hline
 \underbrace{0010110101011111}_{\text{Remainder used as CRC checksum}}
 \end{array}$$

$$\begin{aligned}
 \text{CAN: } G(x) &= x^{15} + x^{14} + x^{10} + x^8 + x^7 + x^4 + x^3 + 1 \\
 &= 1100010110011001
 \end{aligned}$$

DIV mod 2 is
equal to XOR

CAN generator polynomial of degree 15
(requires 16 coefficients to be described,
delivers 15 bits of remainder)

Remainder used as CRC checksum

CAN Data Link Layer

Error Detection

- At receiver side:
 - **CRC error:** The calculated CRC by the receiver does not match the CRC contained in the frame
 - **Form error:** Received frame is not in adherence to the fixed frame fields
- At sender side:
 - **Acknowledgement error:** Message has not been acknowledged by a single node - caused by:
 - Transmission error
 - Corrupted ACK field
 - No receivers available
- At other receivers:
 - **Error signalling:** If an error has been detected all other nodes will immediately be informed by an error frame



Controller Area Network (CAN)

Error Handling & Error Confinement



Error Handling

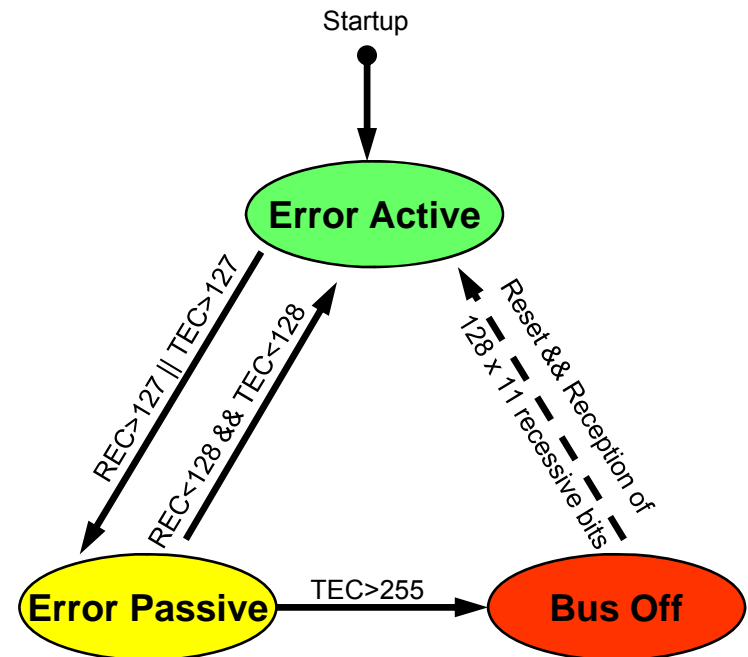
Error frames

- Error frames are sent immediately after an error has been detected
- Error flag:
 - Active error flag: 6 consecutive dominant bits (breaking the stuffing rule!)
 - Passive error flag: 6 recessive bits (can be squashed by error frames sent by other nodes!)
- Error delimiter: 8 recessive bits
- Majority vote to detect “perpetrator”:
 - Majority of nodes send error frame ➔ transmitter is perpetrator
 - Majority of nodes send no error frame ➔ receiver is perpetrator

Error Handling

Fault confinement

- Every node stores two kinds of errors:
 - Transmit error counter (TEC)
 - Receive error counter (REC)
- What a node does if the node is in one of the following states:
 - Error active: Transmission of *Active Error Flags* (dominant) if error is detected by this node
 - Error passive: Transmission of *Passive Error Flags* (recessive) if error is detected by this node
 - Bus off: No transmission on the bus

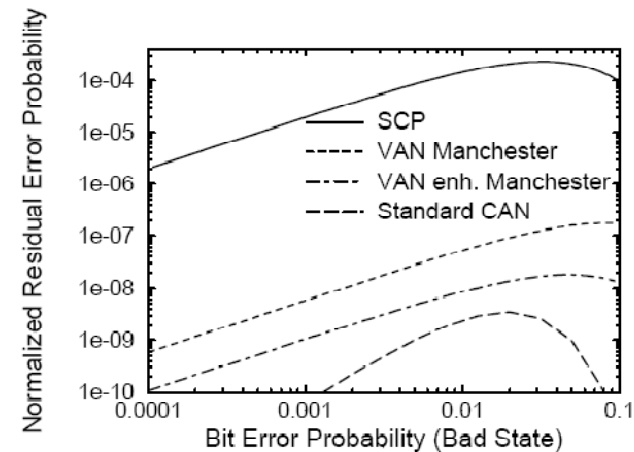
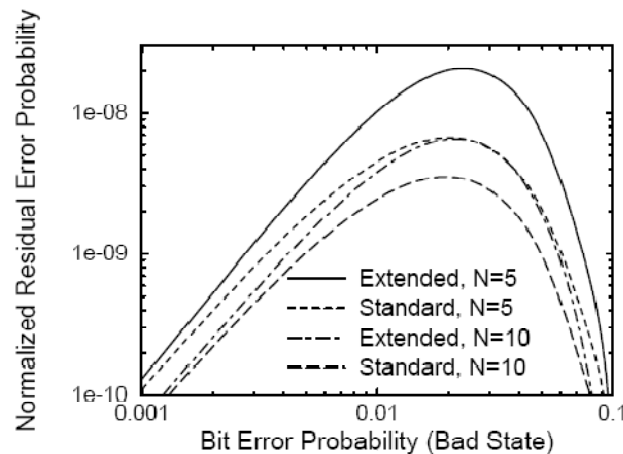


Error Handling

Residual error probability

➤ Example:

- 1 Bit error every 0.7s
- Bit rate: 500 kBit/s
- Operation of 8 hours/day and 365 days/year
- ➔ 1 undetected error in 1000 years



Source: CiA (1994): Performance of the Error Detection Mechanisms in CAN

Questions

- What were the main design criteria for CAN?
- What ISO OSI layers are defined by Bosch CAN specification (ISO-11898-1)?
- Why does the bus length depend on the data rate?
- How does CAN resolve simultaneous bus accesses?
- What types of errors can be detected by CAN?
- What is bit stuffing good for?
- Why are additional standards such as SAE J2284 and J1939 required for the practical application of CAN?